My algorithms exercises

Evgeny Markin

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# Preface

Exercises for Introduction to Algorithms by Cormen et al., 4th ed. It has some exercises, that should be written down, mostly in math and whatnot.

## Part I

Appendix: Mathematical Background

### Chapter 1

## **Summations**

#### 1-1

Prove that  $\sum_{k=1}^{n} O(f_k(i)) = O(\sum_{k=1}^{n} f_k(i))$ 

Short answer:

$$\sum cg(x) = c\sum g(x)$$

Long answer:

Suppose that  $g \in O(f_k(i))$ . It follows that there exists  $n_i$  and  $c_i$  such that  $0 \le g(n) \le cf_i(n)$ . Thus we can pick  $n = \max\{n_0, n_1, ...\}$  and  $c = \max\{c_0, c_1, ...\}$ . We know that both n and c will work all of functions  $f_k$ . Therefore by linearity of summations

$$\sum_{k=1}^{n} O(f_k(i)) = \sum_{k=1}^{n} c f_k(i) == c \sum_{k=1}^{n} f_k(i) == O(\sum_{k=1}^{n} f_k(i))$$

(notation is a little abused and there is nothing is rigorously proven, but it'll do).

#### 1-2

Find a simple formula for  $\sum_{k=1}^{n} (2k-1)$ .

$$\sum_{k=1}^{n} (2k-1) = \sum_{k=1}^{n} (2k) - \sum_{k=1}^{n} (1) = 2\sum_{k=1}^{n} (k) - n = 2\frac{n(n+1)}{2} - n = n(n+1) - n = n^{2}$$

#### 1-3

Interpret the decimal number 111, 111, 111 in light of equation A.6

$$111, 111, 111 = \sum_{k=0}^{9} 10^k = \frac{10^{10} - 1}{10 - 1}$$

#### 1-4

Evaluate the infinite series  $1 - \frac{1}{2} + \frac{1}{4} - \frac{1}{8} + \frac{1}{16} - \dots$ The series converges absolutely to 2, so we are free to do anything with it.

$$1 - \frac{1}{2} + \frac{1}{4} - \frac{1}{8} + \frac{1}{16} - \dots = \sum_{k=0}^{\infty} \frac{1}{2}^{2k} - \sum_{k=0}^{\infty} \frac{1}{2}^{1+2k} = \sum_{k=0}^{\infty} \frac{1}{2}^{2k} - \frac{1}{2} \sum_{k=0}^{\infty} \frac{1}{2}^{2k} = \left(1 - \frac{1}{2}\right) \sum_{k=0}^{\infty} \frac{1}{2}^{2k} = \left(1 - \frac{1}{2}\right) \sum_{k=0}^{\infty} \frac{1}{4}^{2k} = \left(1 - \frac{1}{2}\right) \frac{1}{1 - \frac{1}{4}} = \frac{1}{2} * \frac{4}{3} = \frac{2}{3}$$

#### 1-5

Let  $c \ge 0$  be a constant. Show that  $\sum_{k=1}^{n} k^c = \Theta(n^{c+1})$ 

$$\sum_{k=1}^{n} k^{c} = \sum_{k=1}^{n-1} k^{c} + n^{c} = n^{c} \sum_{k=1}^{n} \frac{k^{c}}{n^{c}} =$$

Let  $f(n) = n^c$ . It can be seen from the graph that

$$\int_0^n f(x)dx \le \sum_{i=1}^n k^c \le \int_0^n f(x+1)dx$$

Thus

$$\int_0^n f(x)dx = \int_0^n x^c = \frac{n^{c+1}}{c+1} \in$$

$$\int_0^n f(x+1)dx = \int_0^n (x+1)^c = \frac{(n+1)^{c+1}}{c+1}$$

Thus we can state that  $\sum_{k=1}^{n} k^{c} = \Theta(n^{c+1})$  (I'm not good enough yet to show that  $\frac{(n+1)^{c+1}}{c+1} \in \Theta(n^{c+1})$ , but I'm pretty sure that it's true TODO).

#### 1-6

Show that  $\sum_{k=0}^{\infty}k^2x^k=x(1+x)/(1-x)^3$  for |x|<1 We know that for |x|<1

$$\sum_{k=0}^{\infty} kx^k = \frac{x}{(1-x)^2}$$

thus if we differentiate both sides we get

$$\sum_{k=0}^{\infty} k^2 x^{k-1} = \frac{2x}{(1-x)^3} + \frac{1}{(1-x)^2}$$

and then if we multiply all of it by x we'll get

$$\sum_{k=0}^{\infty} k^2 x^k = \frac{2x^2}{(1-x)^3} + \frac{x}{(1-x)^2}$$

thus if we factor all of this jazz we'll get

$$\sum_{k=0}^{\infty} k^2 x^k = -\frac{x(x+1)}{(x-1)^3}$$

and if we tuck this minus into denominator we'll get (which we can do because the power is odd)

$$\sum_{k=0}^{\infty} k^2 x^k = \frac{x(x+1)}{(1-x)^3}$$

as desired.

#### 1-7

Prove that  $\sum_{k=1}^{n} \sqrt{k \lg k} = \Theta(n^{3/2} \lg^{1/2} n)$ 

$$\int \sqrt{k \lg k} =$$

#### 1-9

Show that

$$\sum_{k=0}^{\infty} (k-1)/2^k = 0$$

$$\sum_{k=0}^{\infty} (k-1)/2^k = \sum_{k=0}^{\infty} k/2^k - \sum_{k=0}^{\infty} 1/2^k = \sum_{k=0}^{\infty} k/2^k - 2 = 0$$

$$\sum_{k=0}^{\infty} k/2^k - 2 = 0$$

$$\sum_{k=0}^{\infty} k/2^k = 2$$

$$\sum_{k=0}^{\infty} k/2^k - 2 = 0$$

$$\sum_{k=0}^{\infty} k/2^k - 2 = 0$$

$$\sum_{k=0}^{\infty} k - 2^{k+1} = 0$$

### Chapter 2

## Sets, Etc.

#### 1-1

Draw Venn diagrams that illustrate the first of the distributive laws (B.1) TODO, add picture here

#### 1-2

Prove the generalization of DeMorgan's laws to any finite collection of sets Copy from real analysis exercises

Suppose that  $x \in (\bigcup_{\lambda \in \Lambda} E_{\lambda})^c$ . It follows, that x is not in the union of given sets. Therefore there is no set  $E_n$  such that  $x \in E_n$  (because if there would be such a set, then x wouldn't be in  $(\bigcup_{\lambda \in \Lambda} E_{\lambda})^c$ ). Therefore  $x \in \bigcap_{\lambda \in \Lambda} E_{\lambda}^c$ . Therefore

$$(\cup_{\lambda \in \Lambda} E_{\lambda})^{c} \subseteq \cap_{\lambda \in \Lambda} E_{\lambda}^{c}$$

The proof of reverse inclusion is the same as with the forward, but in reverse order.

 $x \in (\cap_{\lambda \in \Lambda} E_{\lambda})^c$  implies that x is not in every  $E_n$ . Therefore there exists  $x \in E_n^c$  for some  $E_n$ . therefore it is in  $\bigcup_{\lambda \in \Lambda} E_{\lambda}^c$ . The proof of reverse inclusion uses the same argument, but in other direction.

#### 1-3

TODO

#### 1-4

Show that the set of odd natural numbers is countable.

Let us set a function  $f: A \to N$ , where A denotes the set of odd natural numbers

$$f(n) = (n+1)/2$$

for this function we've got

$$f^{-1}(n) = 2n - 1$$

Both functions are injective and therefore f is bijective. Therefore we've got a bijective function between A and N, therefore  $A \sim N$ , therefore it is conuntable, as desired.

#### 1-5

Show that for any finite set S, the power set  $2^{S}$  has  $2^{|S|}$  elements (that is, there are  $2^{|S|}$  distinct subsets of S).

Another copy from real analysis

This proof is dumb, but intuitive:

Every subset is corresponding to a number in binary system: 0 for excluded, 1 for included. Therefore there exist  $2^n$  possible combinations.

For a more concrete proof let's resort to induction.

Base case(s): subsets of  $\emptyset$  are  $\emptyset$  itseft ( $2^0 = 1$  in total). Subsets of set with one element are  $\emptyset$  and set itself ( $2^1 = 1$  in total).

Proposition is that set with n elements has  $2^n$  subsets.

Inductive step is that for set with n+1 elements can either have or hot have the n+1'th element. Therefore there exist  $2^n + 2^n = 2 * 2^n = 2^{n+1}$  subsets, as desired.

#### 1-6

Give an inductive definition for an n-tuple by extending the set-theoretic definition for an ordered pair.

The tuple is actually just a re-writing of particular set

$$(a_1, a_2, ..., a_n) = \{\{a_1\}, \{a_1, a_2\}, \{a_1, a_2, a_3\}, ..., \{a_1, a_2, a_3, ..., a_n\}\}$$