

BLG435E

Artificial Intelligence

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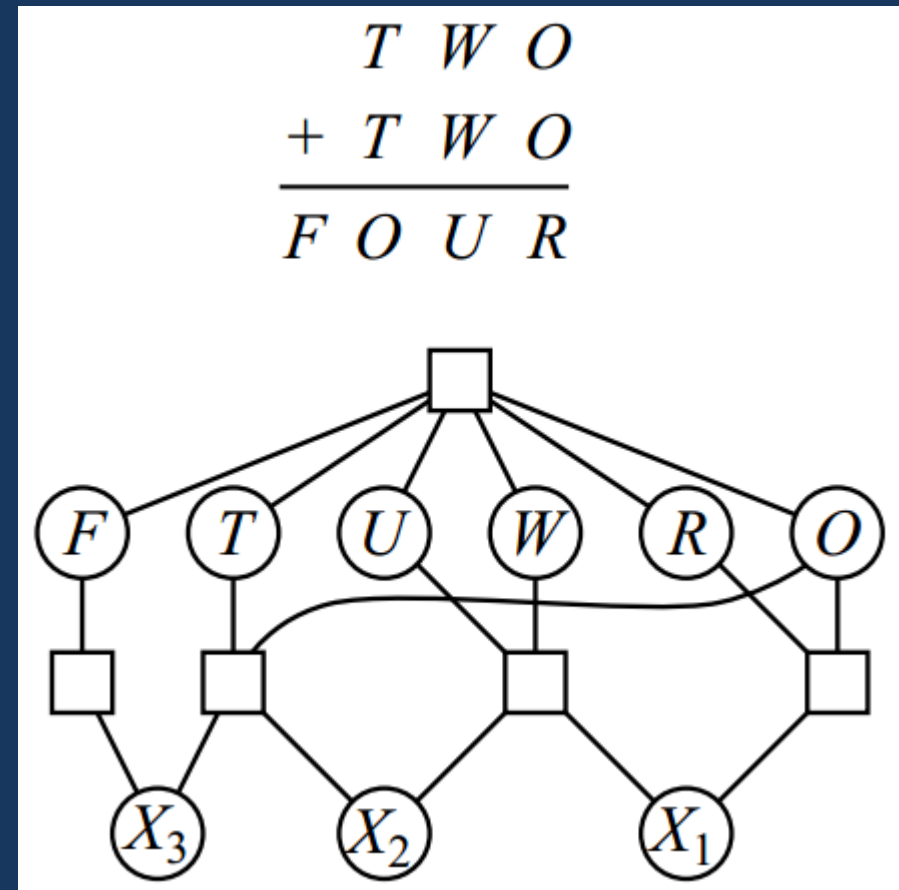


Practice Session 2: CSP

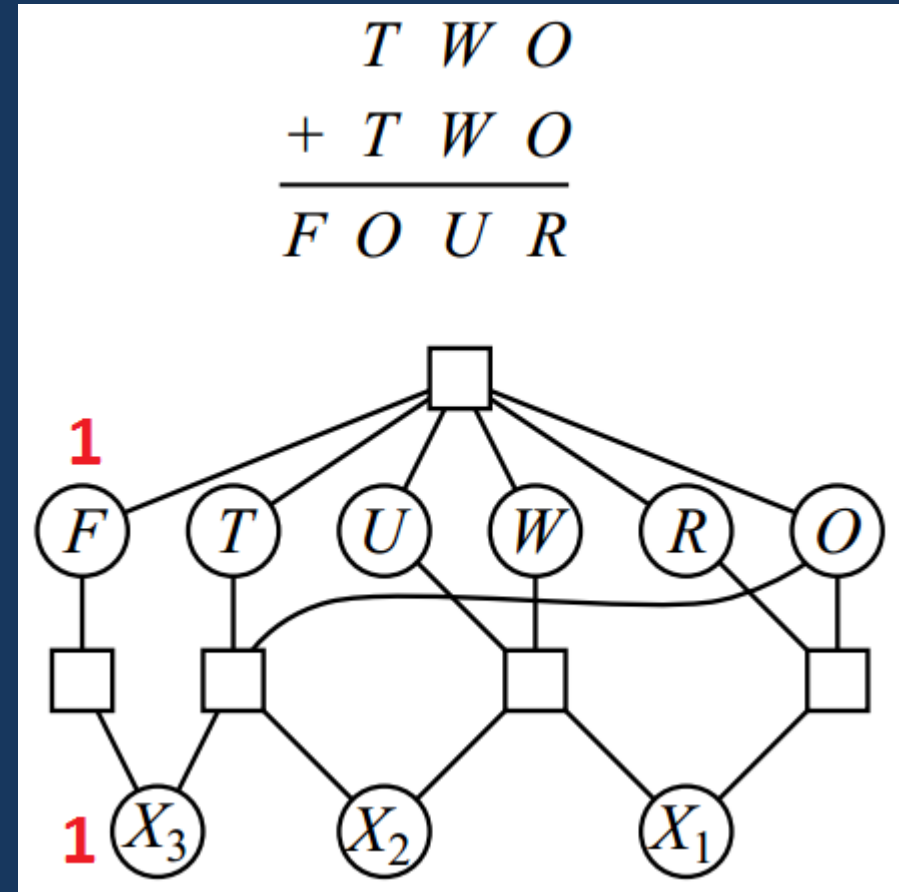


CSP 1

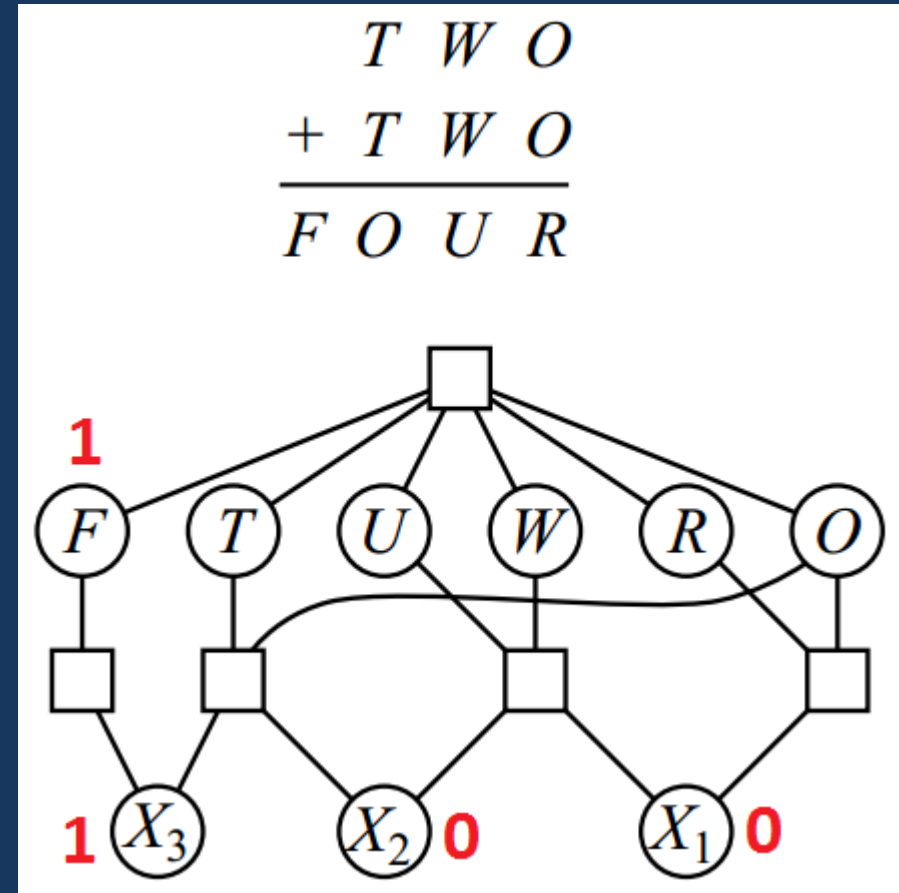
- Solve the cryptarithmic problem, using backtracking, forward checking, and the MRV and least-constraining-value heuristics.
- In this problem, the aim is to find a substitution of digits for letters such that the resulting sum is correct
 - where each letter stands for a distinct digit and
 - no leading zeroes are allowed



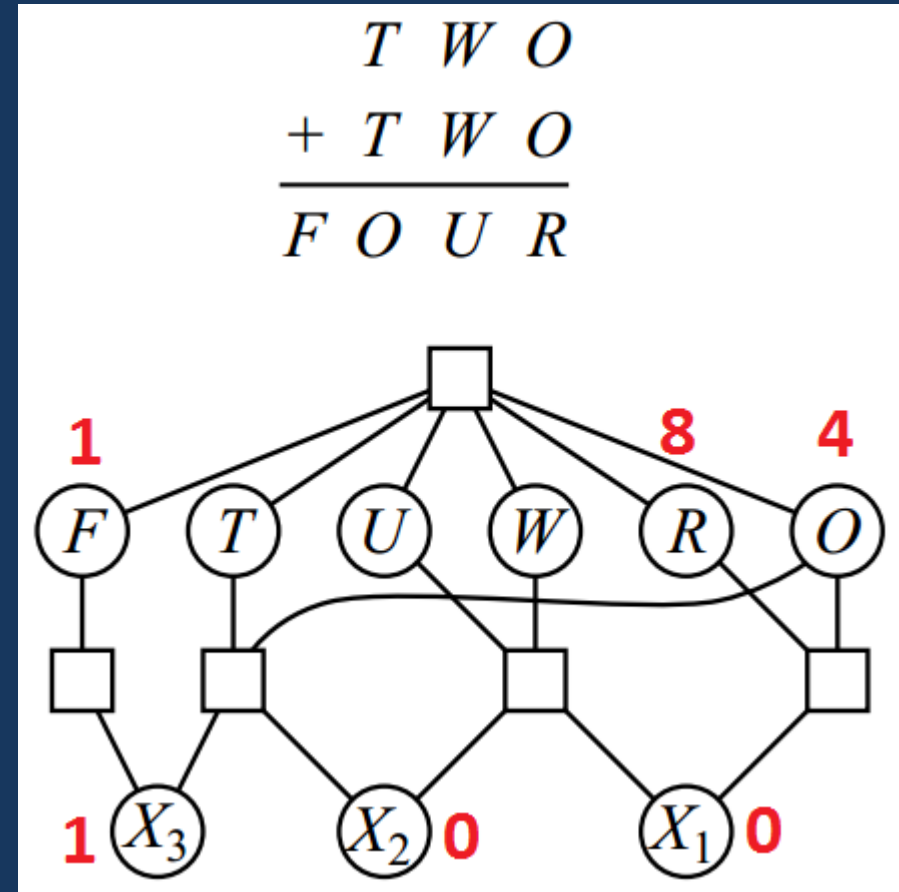
- The exact steps depend on certain choices you are free to make:
 - Choose the **X3** variable. Its domain is {0, 1}.
 - Choose the value 1 for **X3**. (We can't choose 0; it wouldn't survive forward checking, because it would force F to be 0, and the leading digit of the sum must be non-zero.)
 - Choose F, because it has only one remaining value.
 - Choose the value 1 for F.



- The exact steps depend on certain choices you are free to make:
 - Now **X2** and **X1** are tied for minimum remaining values at 2; let's choose **X2**.
 - Either value survives forward checking, let's choose 0 for **X2**.
 - Now **X1** has the minimum remaining values.
 - Again, arbitrarily choose 0 for the value of **X1**.



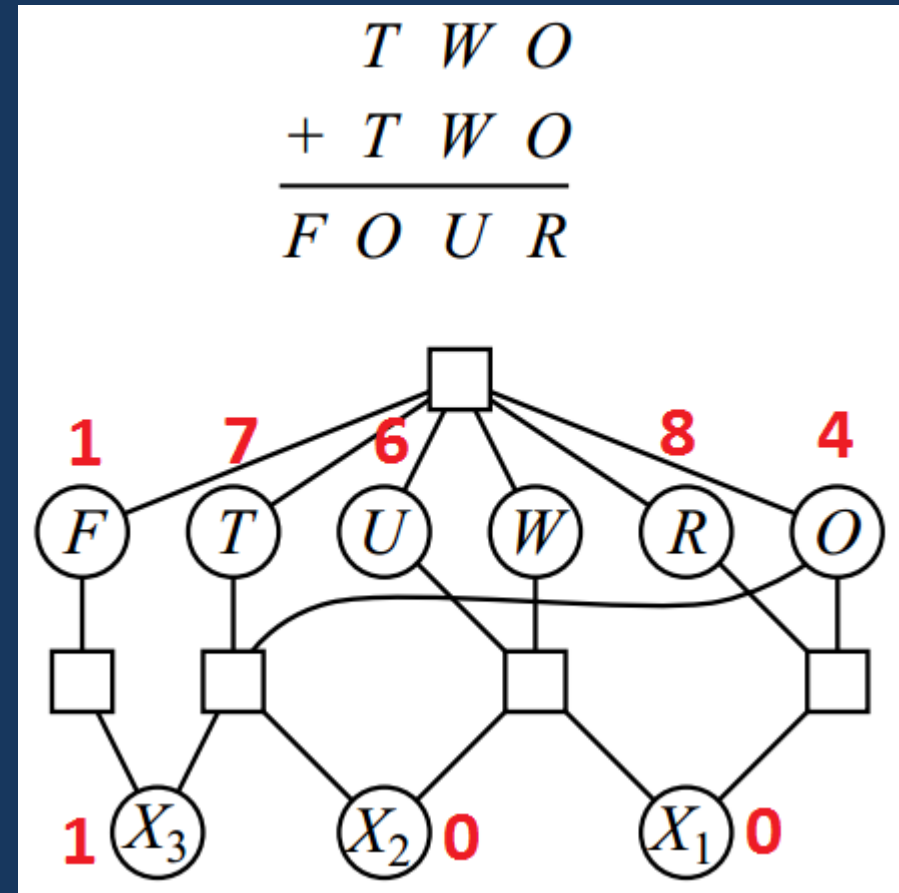
- The exact steps depend on certain choices you are free to make:
 - The variable **O** must be an even number (because it is the sum of **T** + **T**) less than 5 (because **O** + **O** = **R** + 10 x 0). That makes it most constrained.
 - Arbitrarily choose 4 as the value of **O**.
 - R** now has only 1 remaining value.
 - Choose the value 8 for **R**.



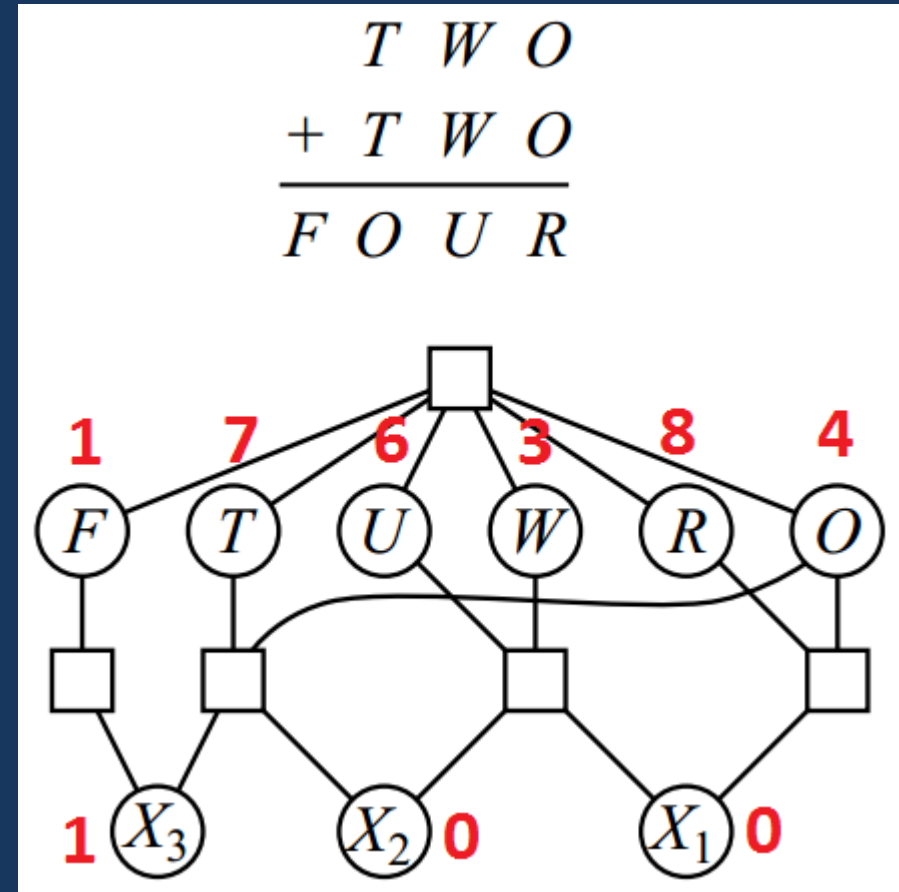
CSP 1

- The exact steps depend on certain choices you are free to make:

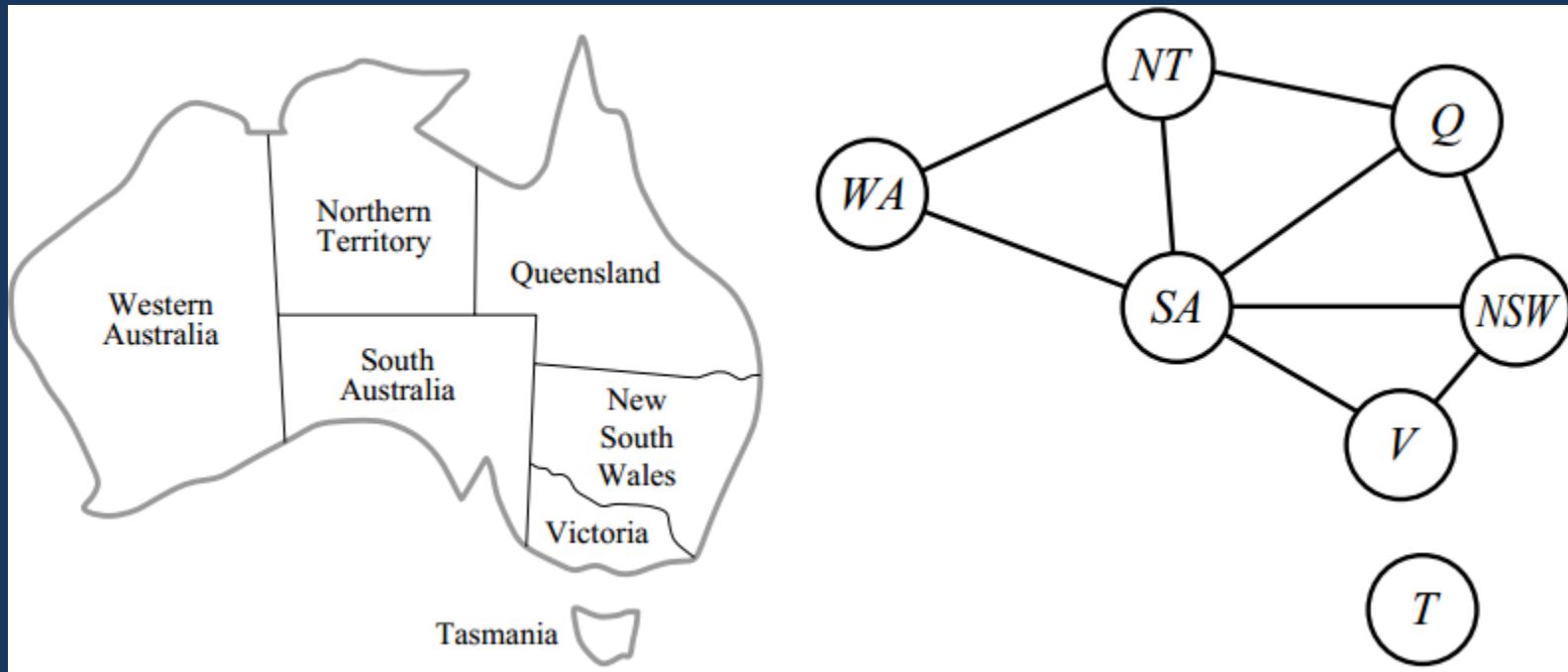
- m) T now has only 1 remaining value.
- n) Choose the value 7 for T.
- o) U must be an even number less than 9; choose U.
- p) The only value for U that survives forward checking is 6



- The exact steps depend on certain choices you are free to make:
 - The only variable left is **W**.
 - The only value left for **W** is 3.
 - This is a solution.

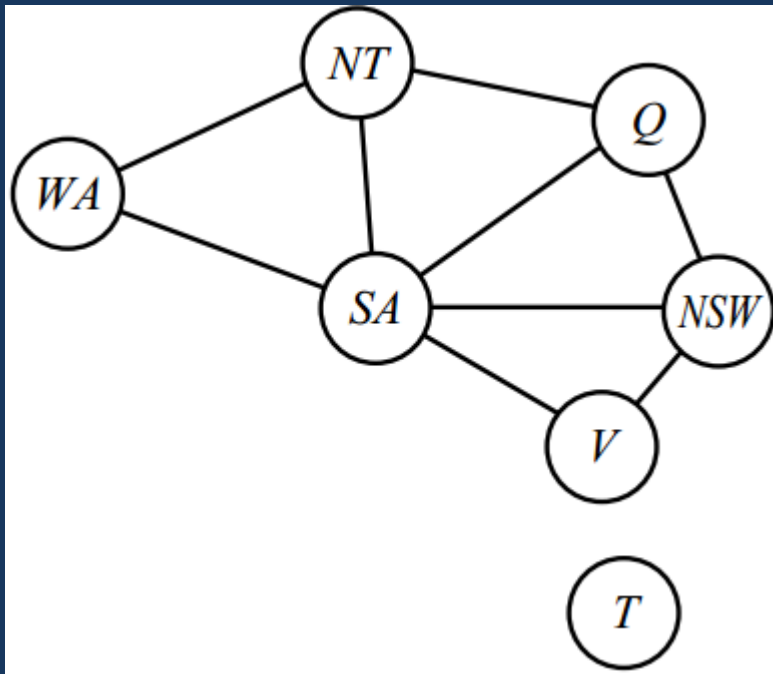


- Use the AC-3 algorithm to show that arc consistency is able to detect the inconsistency of the partial assignment $\{WA = \text{red}, V = \text{blue}\}$ for the problem shown on the right where the goal is to assign colors to each region so that no neighboring regions have the same color.



CSP 2

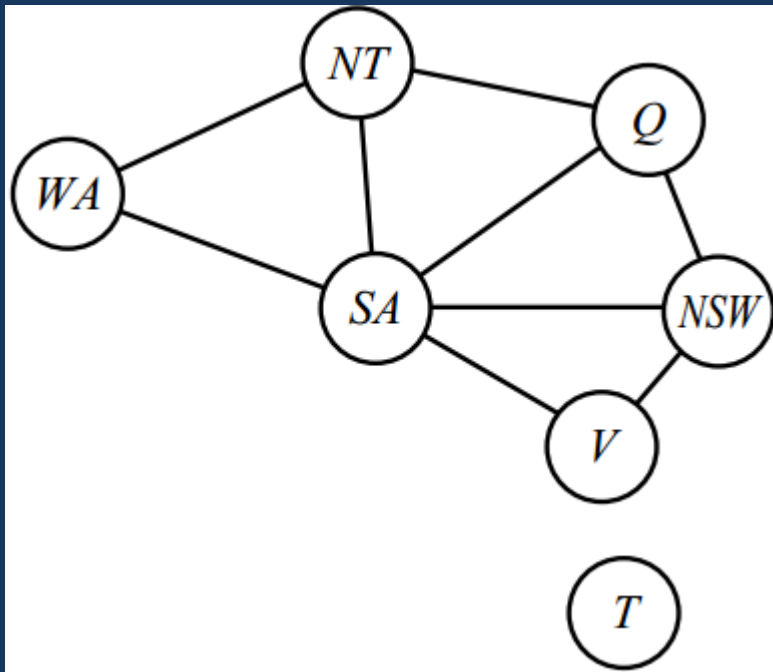
- Remove SA-WA, delete R from SA because there is no value from the domain of WA being consistent with $SA = R$. Insert NT-SA, Q-SA, NSW-SA, V-SA into the queue (other neighbours of SA).



WA	NT	Q	NSW	V	SA	T

WA	NT	Q	NSW	V	SA	T

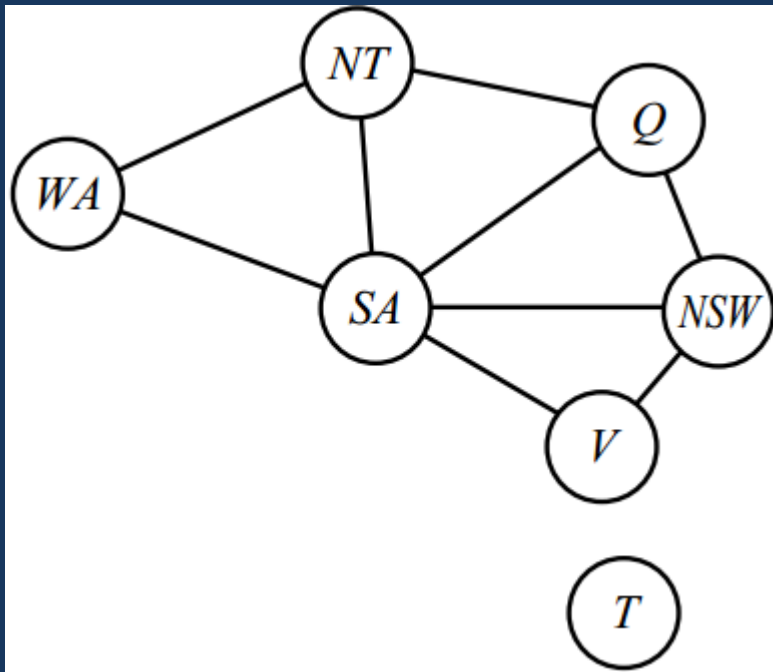
- Remove SA-V, delete B from SA, leaving only G. (For following processes, we will skip the inserting part.)



WA	NT	Q	NSW	V	SA	T

WA	NT	Q	NSW	V	SA	T

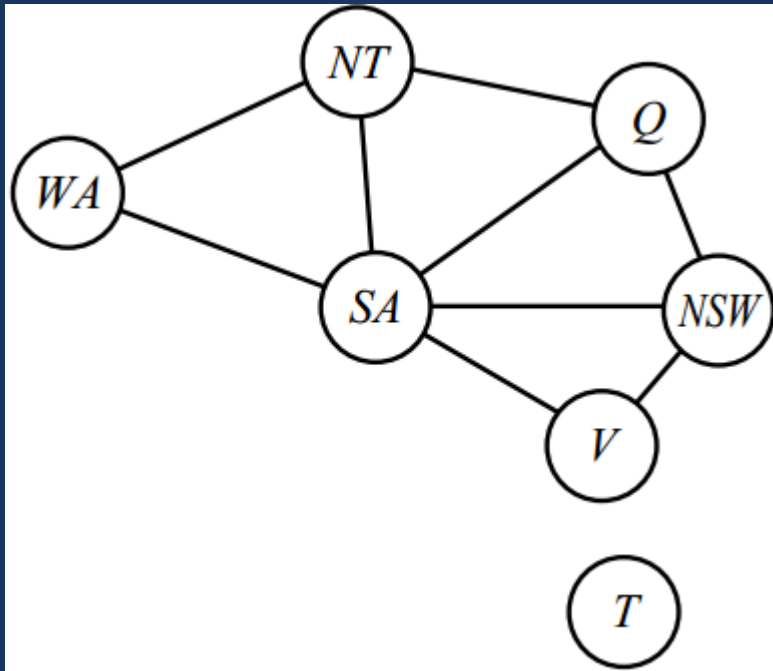
- Remove NT-WA, delete R from NT.



WA	NT	Q	NSW	V	SA	T

WA	NT	Q	NSW	V	SA	T

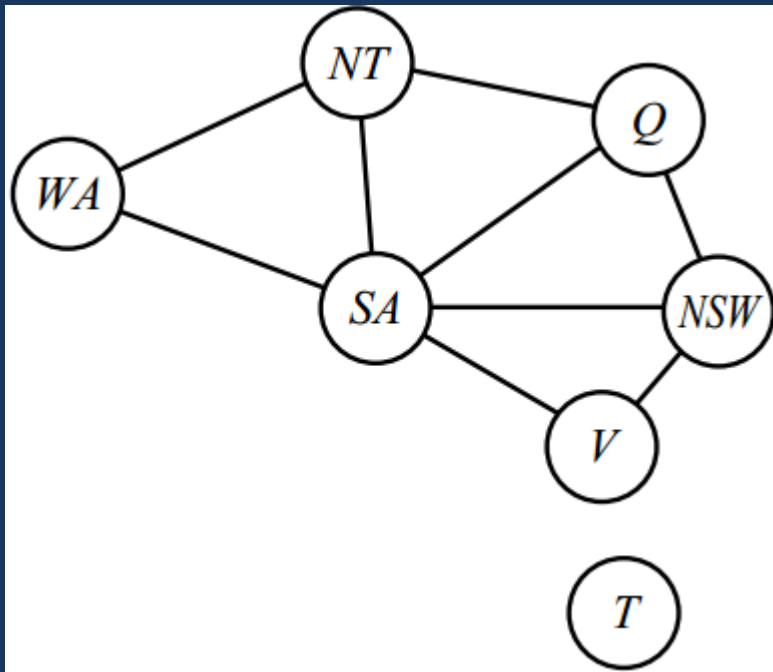
- Remove NT-SA, delete G from NT, leaving only B.



WA	NT	Q	NSW	V	SA	T

WA	NT	Q	NSW	V	SA	T

- Remove NSW-SA, delete G from NSW.

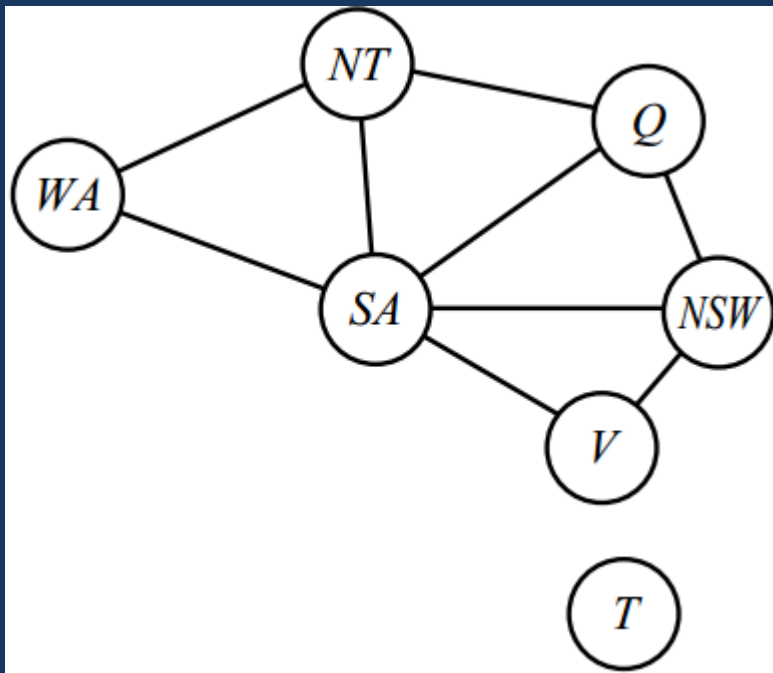


WA	NT	Q	NSW	V	SA	T

WA	NT	Q	NSW	V	SA	T

CSP 2

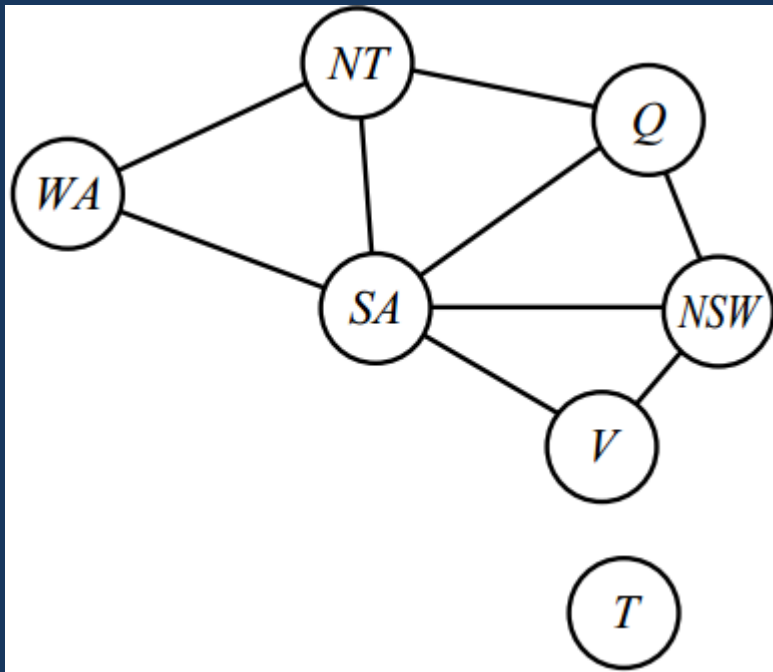
- Remove NSW-V, delete B from NSW, leaving only R.



WA	NT	Q	NSW	V	SA	T

WA	NT	Q	NSW	V	SA	T

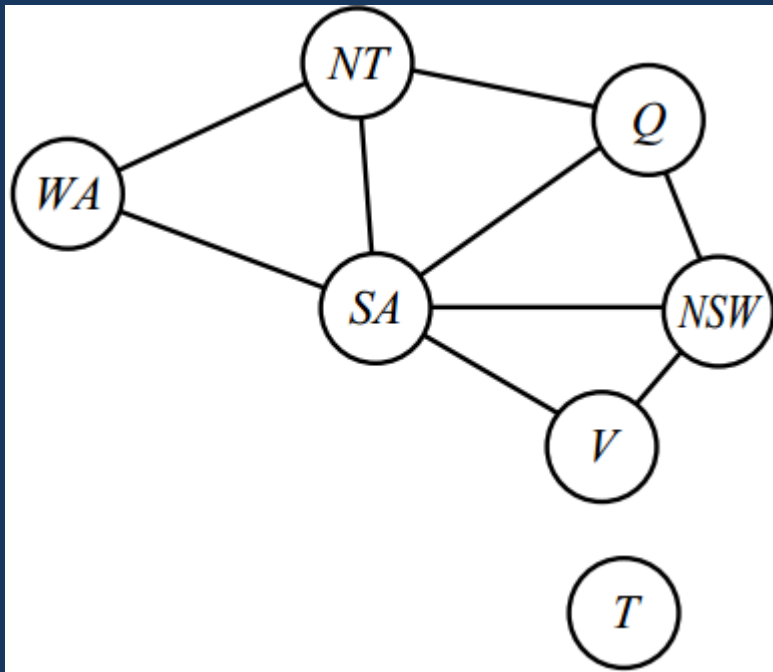
- Remove Q-NT, delete B from Q.



WA	NT	Q	NSW	V	SA	T

WA	NT	Q	NSW	V	SA	T

- Remove Q-SA, delete G from Q.

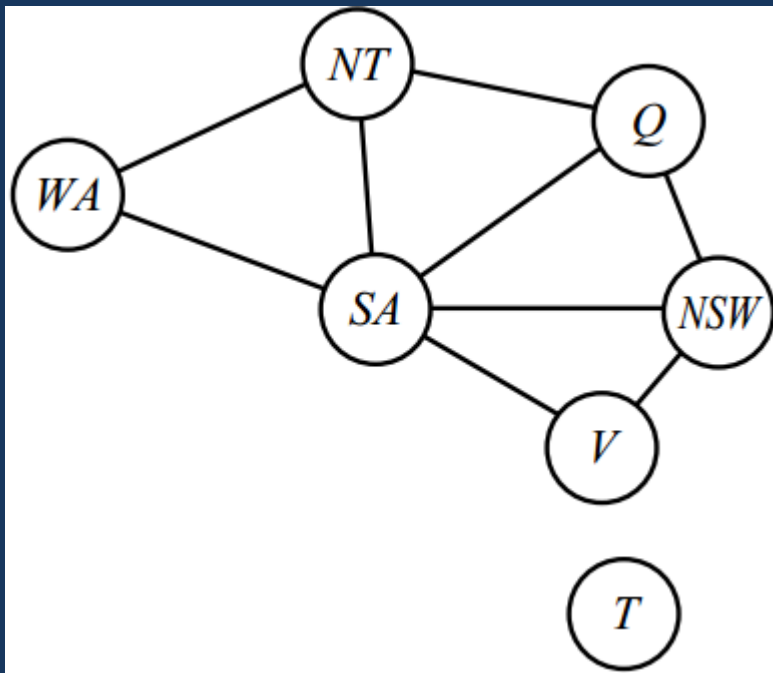


WA	NT	Q	NSW	V	SA	T

WA	NT	Q	NSW	V	SA	T

CSP 2

- Remove Q-NSW, delete R from Q, leaving no domain for Q.



WA	NT	Q	NSW	V	SA	T

WA	NT	Q	NSW	V	SA	T