

### Discrete Math, Homework 3, Chapter 3

1. (1 point) We often use pseudocode to state an algorithm because:
  - a. It can be more complex to state it in an actual programming language.
  - b. To state it in an actual programming language would favor users of that language over other students.
  - c. We can express the ideas of the algorithm more clearly in pseudocode.
  - \*d. All of the above.
  
2. (1 point) If  $f(x) = 148x$ , then  $f(x)$  is (*not restricting* ourselves to the tightest bound):
  - a.  $O(x)$
  - b.  $O(x^2)$
  - c.  $O(x^4)$
  - \*d. All of the above.
  
3. (1 point) If  $f(x) = 148x$ , then  $f(x)$  is (*restricting* ourselves to the tightest bound):
  - \*a.  $O(x)$
  - b.  $O(x^2)$
  - c.  $O(x^4)$
  - d. All of the above.
  
4. (1 point) Which of the following functions grows the most slowly?
  - a.  $f(x) = x \log x$
  - b.  $f(x) = 7x$
  - c.  $f(x) = \lg x$
  - \*d.  $f(x) = \sin x$
  
5. (1 point) Which of the following functions have the same big-O tight bound?
  - a.  $f(x) = 10x^2$  and  $f(x) = 10x$
  - b.  $f(x) = x^3$  and  $f(x) = 100x^2$
  - \*c.  $f(x) = 1x$  and  $f(x) = 1000000x + 1000000$
  - d.  $f(x) = 1000x$  and  $f(x) = x \lg x$
  
6. (1 point) What is the big-O of  $f(x) = x^2 * \log x + x^{2.2}$ ?
  - a.  $x^2$
  - b.  $x^2 \log x$
  - \*c.  $x^{2.2}$
  - d. All of the above.

7. (1 point) Consider the following function: what is its (tight bound) big-O run time?

f(n):

```
for i from 1 to n:
    for j from 1 to n:
        print i * j
        j = floor(j / 2)
    i = i + 1
```

a.  $n^2$

\*b.  $n \log n$

c.  $n$

d.  $n / 2$

8. (1 point) While big-O gives an upper bound for an algorithm's run-time complexity, if we want the lower bound, we use:

a. big-O also

\*b. big-Omega

c. big-Theta

d. little-O

9. (1 point) Determine the complexity of the function below:

```
SUM(s)
s = 0
for i = 1 to n
    s = s + i
return s
```

\*a.  $O(n)$

b.  $O(n^2)$

c.  $O(n \log n)$

d.  $O(e^n)$

10. (1 point) What is the Big-O runtime of the above function:

```
duplicates(A)
for i = 1 to A.length
    for j = i + 1 to A.length
        if A[i] == A[j]
            print(A[i])
```

a.  $O(n)$

\*b.  $O(n^2)$

c.  $O(n \log n)$

d.  $O(n^2 \log n)$

11. (1 point) What  $k$  could we choose as witness to make  $f = O(g)$ , if  $f(x) = x^{2.00001}$  and  $g(x) = x^2$ ?

a. 3

\*b. no such  $k$  exists

c. 42

d. 2.00001

12. (1 point) After what  $C$  will  $g$  rise above  $f$ , if  $g(x) = x^3$  and  $f(x) = 89x^2$ ?

a. 3

b. 2

c. 89

d. no such  $C$  exists