

System Protection

- Goals of Protection
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- Access Matrix
- Implementation of Access Matrix
- Revocation of Access Rights
- Capability-Based Systems
- Language-Based Protection

Protection

- Protection must ensure that only those processes that have gained proper authorization from the OS can operate on memory segments, the CPU, and other resources.
- Operating system consists of a collection of objects, hardware or software
- Each object has a unique name and can be accessed through a well-defined set of operations.
- Protection problem - ensure that each object is accessed correctly and only by those processes that are allowed to do so.
- Enforcement of the policies governing resource usage.
- A protection system must have the flexibility to enforce a variety of policies that can be declared to it.

Goals of Protection

- Provide mechanisms for enforcement of policies.
 - ✦ Mechanisms determine how some thing will be done. Policies decide what will be done.
- Policies may change over time and can be decided by application programmer or system programmer.
- In this chapter we discuss the protection mechanism the OS should provide so that the application designers can design their own protection software.

Guiding principle

■ Principle of least privilege

- ✦ Dictates that programs, users and even systems be given just enough privileges to perform their task.

■ Uses fine grained access controls

■ Beneficial to create a audit trial

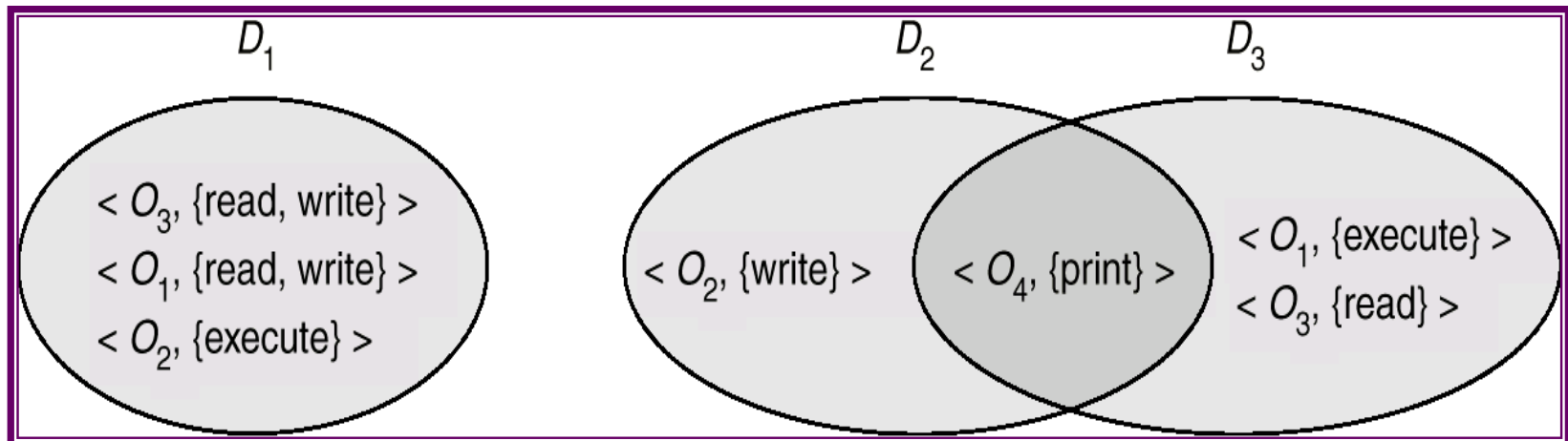
■ Audit trial allows tracing of violations.

Domain Structure

- Computer system is a collection of processes and objects
- Objects are abstract data types
 - ✦ Hardware objects (cpu, memory segments, printers..) and software objects (files, programs, and semaphores).
- Each object is accessed by a well-defined and meaningful operations.
- Operations depend on the object
 - ✦ A CPU may only be executed on. Memory systems may read or written. Files can be open, close, read, write, executed, deleted.
- The process should access only those resources which it requires or allowed to access.
 - ✦ Need to know principle

Domain Structure

- A process operates within a **protection domain**.
- Domain specifies the resources a process may access
- Access right= The ability to execute an operation on an object
- Domain is a collection of access rights.
- Access-right = $\langle \text{object-name}, \text{rights-set} \rangle$
where *rights-set* is a subset of all valid operations that can be performed on the object.



Ways of realizing domains

■ Each user may be a domain

- ◆ Set of objects that can be accessed depends up on the identity of the user.

■ Each process may be a domain

- ◆ Set of objects that can be accessed depends on the identity of the process.

■ Each procedure may be a domain

- ◆ Set of objects that can be accessed corresponds to the local variables defined within the procedure.

Domain Implementation (UNIX)

■ System consists of 2 domains:

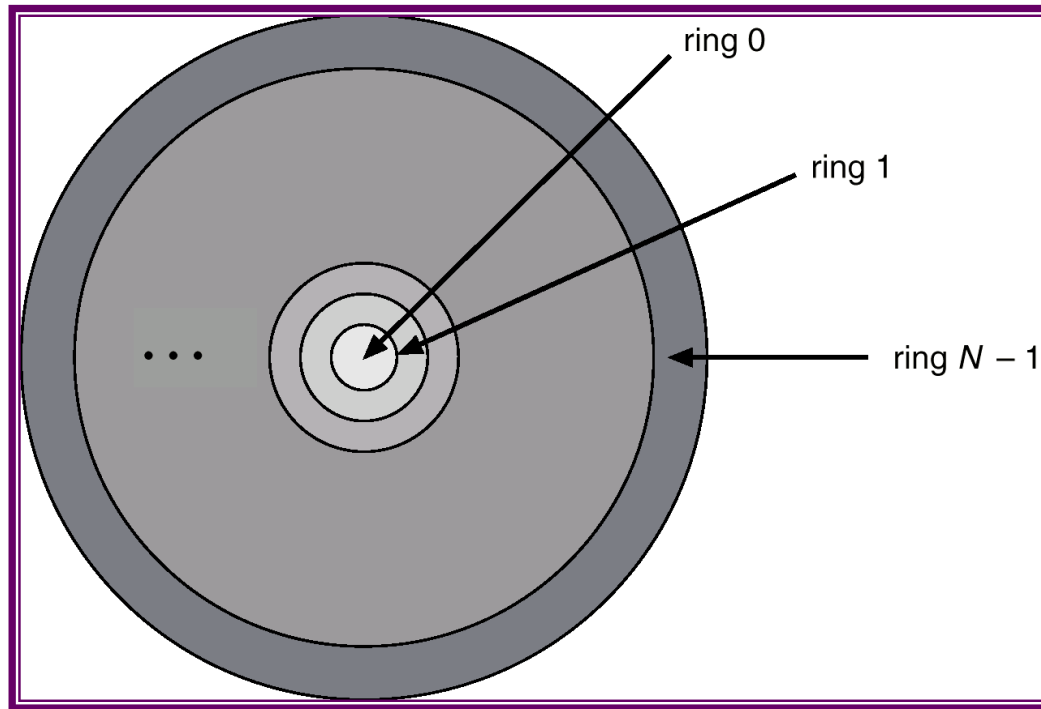
- ✦ User
- ✦ Supervisor

■ UNIX

- ✦ Domain = user-id
- ✦ Domain switch accomplished via file system.
 - ✓ Each file has associated with it a domain bit (setuid bit).
 - ✓ When file is executed and setuid = on, then user-id is set to owner of the file being executed. When execution completes user-id is reset.

Domain Implementation (Multics)

- Protection domains are organized hierarchically into a ring structure.
- Let D_i and D_j be any two domain rings.
- If $j < i \Rightarrow D_i \subseteq D_j$
- Protection system is more complex and less efficient.



Multics Rings

Access Matrix Method

- View protection as a matrix (*access matrix*)
- Rows represent domains
- Columns represent objects
- Each entry in the matrix consists of set of access rights.
- $\text{Access}(i, j)$ is the set of operations that a process executing in Domain_i can invoke on Object_j
- Access Matrix gives flexibility to implement various policies.

Access Matrix

domain \ object				
	F_1	F_2	F_3	printer
D_1	read		read	
D_2				print
D_3		read	execute	
D_4	read write		read write	

Figure A

Use of Access Matrix

- If a process in Domain D_i tries to do “op” on object O_j , then “op” must be in the access matrix.
- Can be expanded to dynamic protection.
 - ✦ Operations to add, delete access rights.
 - ✦ Special access rights:
 - ✓ *owner of O_i*
 - ✓ *copy op from O_i to O_j*
 - ✓ *control – D_i can modify D_j access rights*
 - ✓ *transfer – switch from domain D_i to D_j*

Use of Access Matrix (Cont.)

■ Access matrix design separates mechanism from policy.

✦ Mechanism (how something will be done)

- ✓ Operating system provides access-matrix + rules.

- ✓ If ensures that the matrix is only manipulated by authorized agents and that rules are strictly enforced.

✦ Policy (what will be done)

- ✓ User dictates policy.

- ✓ Who can access what object and in what mode.

Implementation of Access Matrix

- Each column = Access-control list for one object
Defines who can perform what operation.

Domain 1 = Read, Write

Domain 2 = Read

Domain 3 = Read

⋮

- Each Row = Capability List (like a key)
Fore each domain, what operations allowed on what objects.

Object 1 – Read

Object 4 – Read, Write, Execute

Object 5 – Read, Write, Delete, Copy

Access Matrix of Figure A With Domains as Objects

A process in domain D4 can switch to D1, and one in domain D1 can switch to D2

object \ domain	F_1	F_2	F_3	laser printer	D_1	D_2	D_3	D_4
D_1	read		read			switch		
D_2				print			switch	switch
D_3		read	execute					
D_4	read write		read write		switch			

Figure B

Access Matrix with Copy Rights

- A process executing in domain D_2 can copy the read operation in to any entry associated with F_2 . Propagation may be limited.

object \ domain	F_1	F_2	F_3
D_1	execute		write*
D_2	execute	read*	execute
D_3	execute		

(a)

object \ domain	F_1	F_2	F_3
D_1	execute		write*
D_2	execute	read*	execute
D_3	execute	read	

(b)

Access Matrix With *Owner* Rights

- Domain D_1 is owner of F_1 and can add or delete any valid right in F_1 column. Similarly D_2 is owner of F_2 and F_3 .

object \ domain	F_1	F_2	F_3
D_1	owner execute		write
D_2		read* owner	read* owner write*
D_3	execute		

(a)

object \ domain	F_1	F_2	F_3
D_1	owner execute		
D_2		owner read* write*	read* owner write*
D_3		write	write

(b)

Access Matrix: Switch control

A process executing in D_2 could modify domain D_4 .

object domain	F_1	F_2	F_3	laser printer	D_1	D_2	D_3	D_4
D_1	read		read			switch		
D_2				print			switch	switch control
D_3		read	execute					
D_4	write		write		switch			

Confinement problem

- *The copy and owner rights provide us the mechanism to limit the propagation of access rights.*
- *However, they do not give appropriate tools for preventing the propagation of information.*
- The problem of guaranteeing that no information initially held in an object can migrate outside of its execution environment is called **the confinement problem**.
- **Confinement problem is unsolvable.**

Implementation of the Access Matrix

- *Global table of <domain, object, rights-set>*
 - ✦ *Table becomes large and additional I/O is needed*
- *Access list for every object*
 - ✦ *Each column can be implemented as a access list for the object. The llist for each object consists of <Domain, rights-set>*
- *Capability List* for domains
 - ✦ *A capability list a domain is a list objects together with the operations allowed on those objects*

Capability-Based Systems

■ Hydra

- ✦ Fixed set of access rights known to and interpreted by the system.
- ✦ Interpretation of user-defined rights performed solely by user's program; system provides access protection for use of these rights.

■ Cambridge CAP System

- ✦ Data capability - provides standard read, write, execute of individual storage segments associated with object.
- ✦ Software capability -interpretation left to the subsystem, through its protected procedures.

Language-Based Protection

- Specification of protection in a programming language allows the high-level description of policies for the allocation and use of resources.
- Language implementation can provide software for protection enforcement when automatic hardware-supported checking is unavailable.
- Interpret protection specifications to generate calls on whatever protection system is provided by the hardware and the operating system.

Protection in Java 2

- Protection is handled by the Java Virtual Machine (JVM)
- A class is assigned a protection domain when it is loaded by the JVM.
- The protection domain indicates what operations the class can (and cannot) perform.
- If a library method is invoked that performs a privileged operation, the stack is inspected to ensure the operation can be performed by the library.

Stack Inspection

protection domain:	untrusted applet	URL loader	networking
socket permission:	none	*.lucent.com:80, connect	any
class:	gui: ... get(url); open(addr); ...	get(URL u): ... doPrivileged { open('proxy.lucent.com:80'); } <request u from proxy> ...	open(Addr a): ... checkPermission(a, connect); connect (a); ...