

# Program Verification

## Final Exam

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1 a) In short, if a segment is valid, then it can be proved using the segment calculus.  
And if a proof is done using the segment calculus then the segment is valid.

b) a state ~~in~~ at which all <sup>inputs</sup> ~~points~~ in a discrete flow eventually converge at is the convergence point.  
i.e. for  $D = \langle X, F \rangle$ ,  $\forall x \in X$  if  $F(x) = \ast$  a convergent pt.

e) The largest  $A$  is  $(7, 7)$

i.e.  $(u=7, v=7) \{F\} B$

2 a) Consider  $x, y \in \mathbb{N}^+$ ,  $i=0$ ,  $x \geq y$

~~Dir(x, y) = (q, r)~~  $\text{Dir}(x, y) = (q, r)$

where  $q$  is the quotient and  $r$  is the remainder.

~~Dir is computed by~~

it is true that  $x = q \times y + r$   
The problem is to write an algorithm that calculates Dir.

$$M = \langle P, F, \pi \rangle$$

$$b) P: N \times N^+ \rightarrow N \times N^+ \times N$$

$$X = N \times N^+ \times N$$

$$P(x, y) = (x, y, 0)$$

$$F: X \rightarrow X$$

$$F(x, y, i) = \begin{cases} (x-y, y, i+1) & \text{if } x \geq y \\ (x, y, i) & \text{otherwise} \end{cases}$$

$$\pi: X \rightarrow N \times N$$

$$\pi(x, y, i) = (x, i)$$

c) Consider a function

$$\lambda(x, y, i) = x$$

$$\lambda(F(x, y, i)) = x - y$$

$$\lambda(x, y, i) < \lambda(F(x, y, i))$$

it follows that  $\lambda(y, i) < \lambda(F(x, y, i))$   
 this is a well founded chain eventually

ends when  $x < y$ .

$x \in \mathbb{N}$  and  
 $M(\mathbb{N}, <)$  is a well founded  
 relation.  
 $\therefore M$  is an algorithm.

d) Consider the invariant function

$$\Theta(x, y, i) = y \times i + x \rightarrow \textcircled{1}$$

$$\begin{aligned} \Theta(F(x, y, i)) &= y \times (i+1) + (x-y) \\ &= y \times i + y + x - y \\ &= y \times i + x \\ &= \text{eq } \textcircled{1} \end{aligned}$$

$\therefore \Theta$  is an invariant function and  
 hence  $M$  is partially correct.

1. c) An invariant function is one such  
 that for a mapcode machine  $M = \langle P, F, \pi \rangle$   
 $\Theta(P(i)) = \Theta(F^j(P(i)))$

for any input  $i$  in  $I$  (input space)  
 and for any  $j \in \mathbb{N}$



3. Pseudocode for the repeated subtraction division algorithm is

1.  $(x, y, i) := (x, y, 0)$
2. while  $(x \geq y)$  do:
  - 2.1  $(x, y, i) := (x - y, y, i + 1)$

required specification is:

$$\text{TRUE} \{P\} \quad x' = \text{int } y + x$$

where  $x'$  is the initial  $x$ .