Feedback — Week 5 - Problem Set

Help

You submitted this homework on **Thu 1 May 2014 5:56 PM PDT**. You got a score of **13.00** out of **15.00**. You can attempt again in 10 minutes.

Question 1

Consider the toy key exchange protocol using an online trusted 3rd party (TTP) discussed in Lecture 9.1. Suppose Alice, Bob, and Carol are three users of this system (among many others) and each have a secret key with the TTP denoted k_a, k_b, k_c respectively. They wish to generate a group session key k_{ABC} that will be known to Alice, Bob, and Carol but unknown to an eavesdropper. How would you modify the protocol in the lecture to accomodate a group key exchange of this type? (note that all these protocols are insecure against active attacks)

Your Answer		Score	Explanation
○Alice contacts the TTP. TTP generates a random k_{ABC} and sends to Alice $E(k_a, k_{ABC}), \quad \text{ticket}_1 \leftarrow E_(k_c, E(k_b, k_{ABC})), \quad \text{ticket}_2 \leftarrow E(k_b, k$. Alice sends k_{ABC} to Bob and k_{ABC} to Carol.			
●Bob contacts the TTP. TTP generates random k_{ABC} and sends to Bob $E(k_b, k_{ABC}), \quad \text{ticket}_1 \leftarrow E(k_a, k_{ABC}), \quad \text{ticket}_2 \leftarrow E(k_c, k_{ABC})$ Bob sends ticket_1 to Alice and ticket_2 to Carol.	✓	1.00	The protocol works because it lets Alice, Bob, and Carol obtain k_{ABC} but an eaesdropper only sees encryptions of k_{ABC} under keys he does not

have.

○ Bob contacts the TTP. TTP generates a random k_{AB} and a random k_{BC}. It sends to Bob

$$E(k_a, k_{AB}), \quad \text{ticket}_1 \leftarrow E(k_a, k_{AB}), \quad \text{ticket}_2 \leftarrow E(k_c, k_{BC})$$

Bob sends ticket_1 to Alice and ticket_2 to Carol.

○ Alice contacts the TTP. TTP generates a random k_{ABC} and sends to Alice

$$E(k_a, k_{ABC}), \quad \text{ticket}_1 \leftarrow k_{ABC}, \quad \text{ticket}_2 \leftarrow k_{ABC}$$

Alice sends ticket_1 to Bob and ticket_2 to Carol.

Total 1.00 / 1.00

Question 2

Let G be a finite cyclic group (e.g. $G = \mathbb{Z}_p^*$) with generator g. Suppose the Diffie-Hellman function $\text{DH}_g(g^x, g^y) = g^{xy}$ is difficult to compute in G. Which of the following functions is also difficult to compute:

As usual, identify the f below for which the contra-positive holds: if $f(\cdot,\cdot)$ is easy to compute then so is $\text{DH}_g(\cdot,\cdot)$. If you can show that then it will follow that if DH_g is hard to compute in G then so must be f.

Your Answer		Score	Explanation
□ $f(g^x, g^y) = g^{x+y}$	✓	0.25	It is easy to compute f as $f(g^x, g^y) = g^x \cdot g^y$.
☑ $f(g^x, g^y) = g^{xy+x+y+1}$	✓	0.25	an algorithm for calculating $f(g^x, g^y)$ can easily be converted into an algorithm for calculating $\text{DH}(\cdot,\cdot)$. Therefore, if f were easy to compute then so would DH, contrading the assumption.
☑ $f(g^x, g^y) = g^{x(y+1)}$	✓	0.25	an algorithm for calculating $f(g^x, g^y)$ can easily be converted into an algorithm for calculating $\text{DH}(\cdot,\cdot)$. Therefore, if f were easy to compute then so would DH, contrading the assumption.
□ $f(g^x, g^y) = (\sqrt{g})^{x+y}$	✓	0.25	It is easy to compute f as $f(g^x, g^y) = \sqrt{g^x \cdot g^y}$.

Total		1.00 / 1.00	

Question 3

Suppose we modify the Diffie-Hellman protocol so that Alice operates as usual, namely chooses a random a in $\{1,\ldots,p-1\}$and sends to Bob $A \leftarrow g^{a}$. Bob, however, chooses a random b in $\{1,\ldots,p-1\}$and sends to Alice $B \leftarrow g^{1/b}$. What shared secret can they generate and how would they do it?

Your Answer		Score	Explanation
○ $\text{secret} = g^{b/a}$. Alice computes the secret as B^a and Bob computes $A^{1/b}$.			
○ $\text{secret} = g^{a/b}$. Alice computes the secret as $B^{1/b}$ and Bob computes A^a.			
● $\text{secret} = g^{a/b}$. Alice computes the secret as B^a and Bob computes $A^{1/b}$.	✓	1.00	This is correct since it is not difficult to see that both will obtain $g^{a/b}$
○ $\text{secret} = g^{a/b}$. Alice computes the secret as $B^{1/a}$ and Bob computes A^b.			
Total		1.00 / 1.00	

Question 4

Consider the toy key exchange protocol using public key encryption described in Lecture 9.4. Suppose that when sending his reply $c \leftarrow E(pk,x)$to Alice, Bob appends a MAC $t := S(x,c)$ to the ciphertext so that what is sent to Alice is the pair (c,t). Alice verifies the tag t and rejects the message from Bob if the tag does not verify. Will this additional step prevent the man in the middle attack described in the lecture?

Your Answer	Score	Explanation

○ it depends on what public key encryption system is used.			
○ it depends on what MAC system is used.			
● no	✓	1.00	an active attacker can still decrypt $E(pk', x)$ to recover x and then replace (c, t) by (c', t') where $c' \leftarrow E(pk, x)$ and $t \leftarrow S(x, c')$.
○ yes			
Total		1.00 / 1.00	

Question 5

The numbers 7 and 23 are relatively prime and therefore there must exist integers a and b such that $7a + 23b = 1$ Find such a pair of integers (a, b) with the smallest possible $a > 0$. Given this pair, can you determine the inverse of 7 in $\mathbb{Z}_{23}$?

Enter below comma separated values for $a,\ b$, and for 7^{-1} in $\mathbb{Z}_{23}$.

You entered:

```
10, -3, 10
```

Your Answer		Score	Explanation
10, -3, 10	✓	1.00	
Total		1.00 / 1.00	

Question 6

Solve the equation $3x + 2 = 7$ in $\mathbb{Z}_{19}$.

You entered:

8

Your Answer		Score	Explanation
8	✓	1.00	
Total		1.00 / 1.00	

Question 7

How many elements are there in $\mathbb{Z}_{35}^*$?

You entered:

24

Your Answer		Score	Explanation
24	✓	1.00	
Total		1.00 / 1.00	

Question 8

How much is $2^{10001} \bmod 11$? (please do not use a calculator for this)

Hint: use Fermat's theorem.

You entered:

2

Your Answer		Score	Explanation
2	✓	1.00	
Total		1.00 / 1.00	

Question 9

While we are at it, how much is $2^{245} \bmod 35$?

Hint: use Euler's theorem (you should not need a calculator)

You entered:

32

Your Answer		Score	Explanation
32	✓	1.00	
Total		1.00 / 1.00	

Question 10

What is the order of 2 in $\mathbb{Z}_{35}^*$?

You entered:

4

Your Answer		Score	Explanation
4	✗	0.00	
Total		0.00 / 1.00	

Question 11

Which of the following numbers is a generator of $\mathbb{Z}_{13}^*$?

Your Answer		Score	Explanation
□ $5, \quad \langle 5 \rangle = \{1, 5, 12, 8\}$	✓	0.20	No, 5 only generates four

			elements in $\mathbb{Z}_{13}^*$.
☑ $7, \quad \langle 7 \rangle = \{1, 7, 10, 5, 9, 11, 12, 6, 3, 8, 4, 2$	✓	0.20	correct, 7 generates the entire group $\mathbb{Z}_{13}^*$
□ $9, \quad \langle 9 \rangle = \{1, 9, 3\}$	✓	0.20	No, 9 only generates three elements in $\mathbb{Z}_{13}^*$.
□ $10, \quad \langle 10 \rangle = \{1, 10, 9, 12, 3, 4\}$	✓	0.20	No, 10 only generates six elements in $\mathbb{Z}_{13}^*$.
☑ $2, \quad \langle 2 \rangle = \{1, 2, 4, 8, 3, 6, 12, 11, 9, 5, 10, 7$	✓	0.20	correct, 2 generates the entire group $\mathbb{Z}_{13}^*$
Total		1.00 / 1.00	

Question 12

Solve the equation $x^2 + 4x + 1 = 0$in $\mathbb{Z}_{23}$. Use the method described in lecture 9.3 using the quadratic formula.

You entered:

8 and 7

Your Answer		Score	Explanation
8 and 7	✗	0.00	
Total		0.00 / 1.00	

Question 13

What is the 11th root of 2 in $\mathbb{Z}_{19}$? (i.e. what is $2^{1/11}$ in $\mathbb{Z}_{19}$)

Hint: observe that $11^{-1} = 5$ in $\mathbb{Z}_{18}$.

You entered:

13

Your Answer		Score	Explanation
13	✓	1.00	
Total		1.00 / 1.00	

Question 14

What is the discete log of 5 base 2 in $\mathbb{Z}_{13}$? (i.e. what is $\mathrm{Dlog}_2(5)$)

Recall that the powers of 2 in $\mathbb{Z}_{13}$ are $\langle 2 \rangle = \{1, 2, 4, 8, 3, 6, 12, 11, 9, 5, 10, 7\}$

You entered:

9

Your Answer		Score	Explanation
9	✓	1.00	
Total		1.00 / 1.00	

Question 15

If p is a prime, how many generators are there in $\mathbb{Z}_p^*$?

Your Answer		Score	Explanation
○ $\sqrt{p}$			
● $\varphi(p-1)$	✓	1.00	The answer is $\varphi(p-1)$. Here is why. Let g be some generator of $\mathbb{Z}_p^*$ and let $h = g^x$ for some x. It is not difficult to see that h is a generator exactly when we can write g as $g = h^y$ for some integer y (h is a generator because if $g = h^y$ then any power of g can also be written as a power of h). Since $y = x^{-1} \mod p - 1$this y exists exactly when x is relatively prime to $p - 1$. The number of such x is the size of $\mathbb{Z}_{p-1}$ which is precisely $\varphi(p-1)$.

○ $(p+1)/2$

○ $p-1$

Total	1.00 / 1.00