

Software Multicore Processors

Homework 3

In this report connected component algorithm has been parallelized using Galois system. First the sequential versions of the algorithm has been implemented using FIFO queue and Priority queue.

Pseudocode:

```
1: for v in graph.vertices():
2:   v.compId = v.id
3: worklist wl(graph.vertices())
4: while not wl.empty():
5:   v = wl.pop()
6:   for nv in v.neighbors():
7:     if v.compId < nv.compId:
8:       nv.compId = v.compId
9:       wl.add(nv)
```

Sequential Versions

1. Worklist implemented as FIFO queue

Worklist has been implemented as a FIFO queue using the queue. Following table shows the runtime for each graph.

| | | |
|----------|---------|--------|
| 16777216 | 13.7021 | Rmat |
| 23947347 | 245.968 | usa |
| 33554432 | 22.2624 | Random |

2. Worklist implemented as priority queue

Worklist has been implemented as a priority queue where nodes are prioritized based on the component id number. Following table shows the runtime for each graph.

| | | |
|----------|----------|--------|
| 16777216 | 13.7021 | Rmat |
| 23947347 | 16.90335 | usa |
| 33554432 | 22.2624 | Random |

Parallel Versions

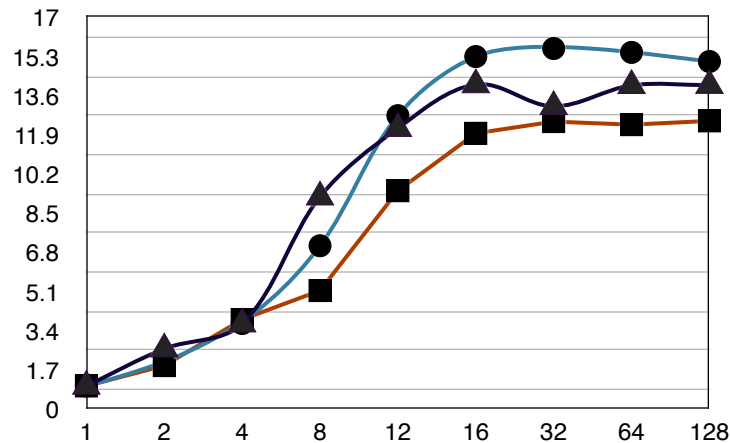
Following parallel versions have been implemented using Galois. Various scheduling policies were tried for worklist. Here are the results for each of them.

1. Worklist implemented as *dChunked FIFO queue Ordered by Integer Metric*.

Heuristic used for order by integer metric: Each node is assigned priority equal to $\text{component_id} \% \text{chunk_size}$. This creates lesser number of priority levels and nodes with lower component id are given higher priority.

Table below shows the runtime for each graph for different number of threads.

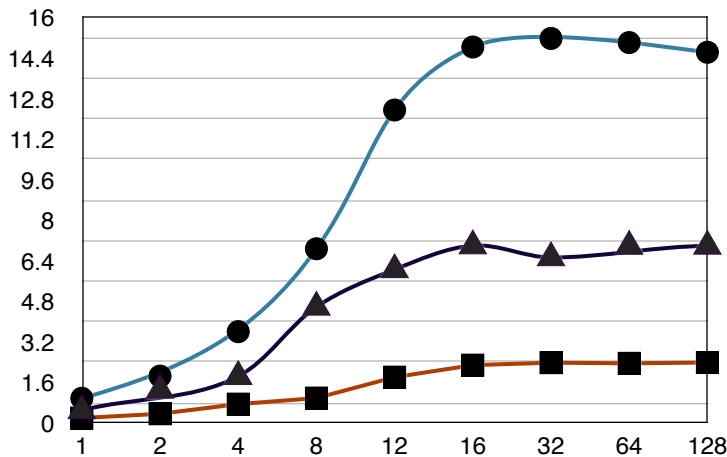
| n | 1 | 2 | 4 | 8 | 12 | 16 | 32 | 64 | 128 |
|---------------------|---------|---------|---------|---------|--------|---------|---------|---------|---------|
| $T_n(n = 16777210)$ | 27.5978 | 10.5478 | 7.45791 | 2.99623 | 2.2647 | 1.95979 | 2.10062 | 1.96454 | 1.96368 |
| $T_n(n = 23947347)$ | 253.779 | 133.296 | 68.3745 | 35.8171 | 19.931 | 16.5902 | 16.2203 | 16.3963 | 16.826 |
| $T_n(n = 33554432)$ | 116.992 | 62.123 | 30.113 | 22.7609 | 12.349 | 9.77683 | 9.37872 | 9.47971 | 9.35529 |



Scalability

Chunk size: 64.

A significant amount of scalability and speedup is achieved using chunked FIFO queue ordered by integer metric.

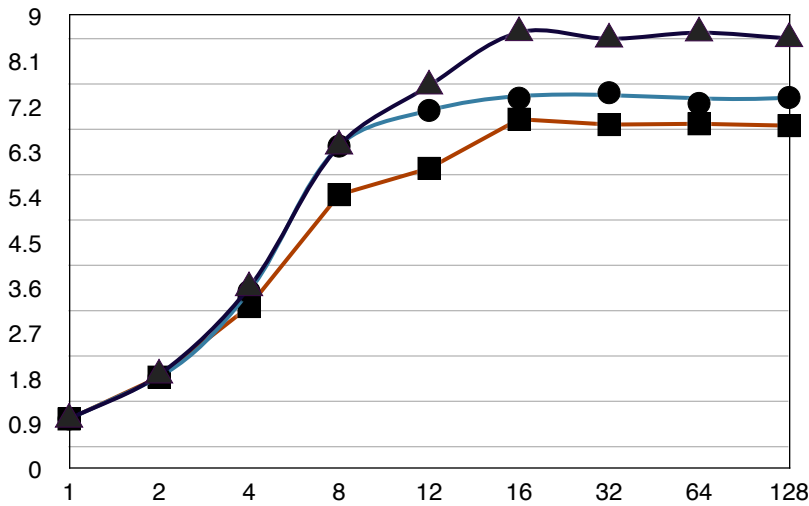


Speedup

Significant amount of speedup can be seen.

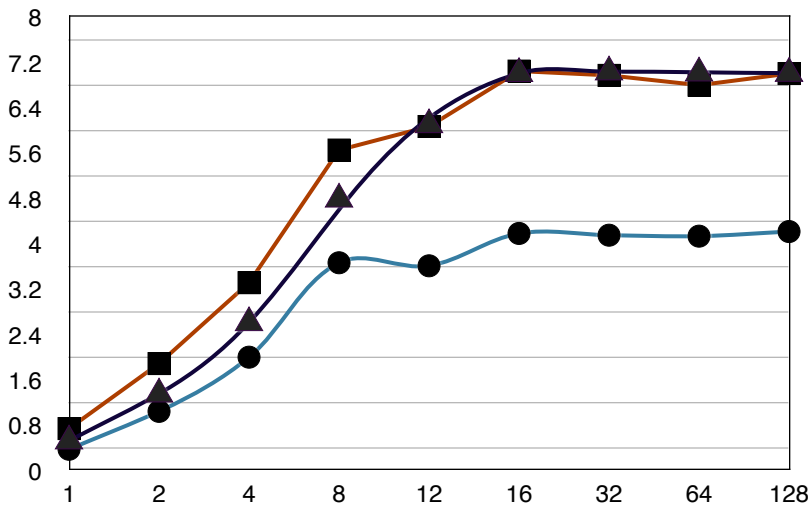
2. Worklist implemented as *dChunkedFIFO* queue without the OBIM.

| n | 1 | 2 | 4 | 8 | 12 | 16 | 32 | 64 | 128 |
|---------------------|---------|---------|---------|---------|--------|---------|---------|---------|---------|
| $T_n(n = 16777210)$ | 15.4453 | 8.23198 | 4.27823 | 2.40482 | 2.0283 | 1.78075 | 1.80713 | 1.78184 | 1.80541 |
| $T_n(n = 23947347)$ | 432.281 | 235.787 | 122.379 | 67.3642 | 60.666 | 58.6669 | 57.8133 | 59.505 | 58.5562 |
| $T_n(n = 33554432)$ | 21.712 | 11.8588 | 6.71972 | 3.98159 | 3.6338 | 3.12376 | 3.17235 | 3.16506 | 3.18215 |



Scalability

Again the algorithm scales well due to parallelism.

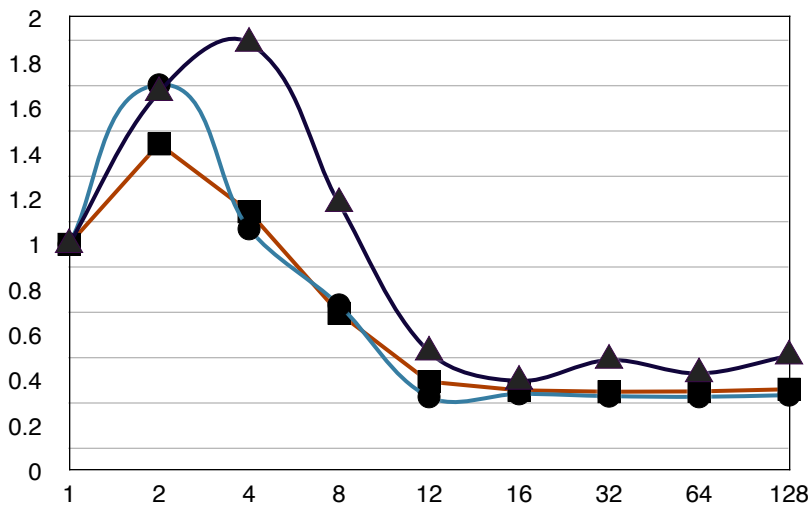


Speedup

A significant amount of speedup can be seen when ran for larger number of threads.

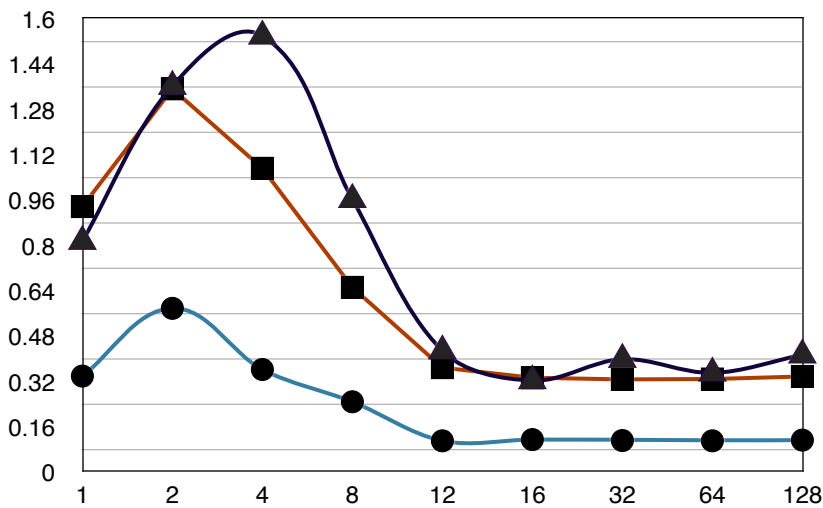
3. Worklist implemented as FIFO queue

| $n(\text{FIFO})$ | 1 | 2 | 4 | 8 | 12 | 16 | 32 | 64 | 128 |
|----------------------|---------|---------|---------|---------|--------|---------|---------|---------|---------|
| $T_n (n = 16777210)$ | 16.7943 | 10.0526 | 8.89564 | 14.2192 | 31.93 | 42.3024 | 34.307 | 39.044 | 33.036 |
| $T_n (n = 23947347)$ | 724.091 | 425.015 | 677.792 | 988.054 | 2212.2 | 2127.74 | 2192.54 | 2212.54 | 2162.54 |
| $T_n (n = 33554432)$ | 23.7347 | 16.4475 | 20.7483 | 34.1036 | 60.261 | 66.6905 | 67.8724 | 67.5958 | 65.8084 |



Scalability:

Scalability is not achieved due to overhead associated. Since all the threads are competing for the same queue.

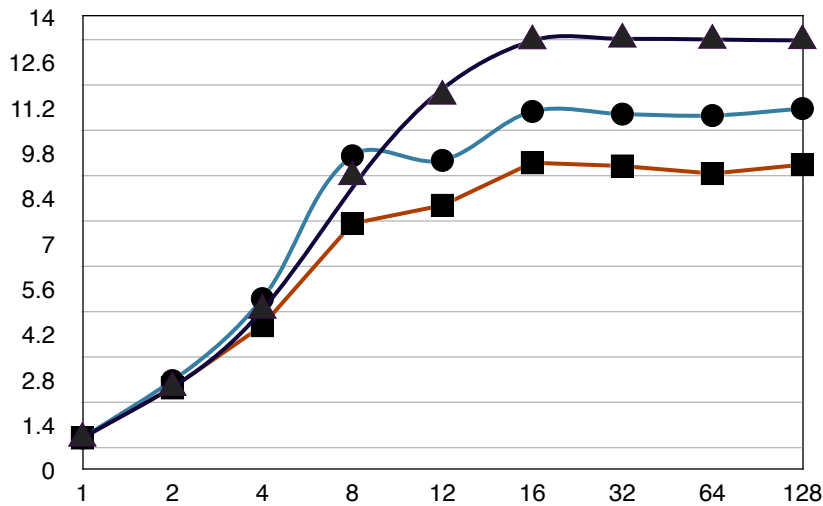


Speedup:

Due to similar reasons as above speedup is not achieved.

4. Worklist implemented as ChunkedFIFO queue

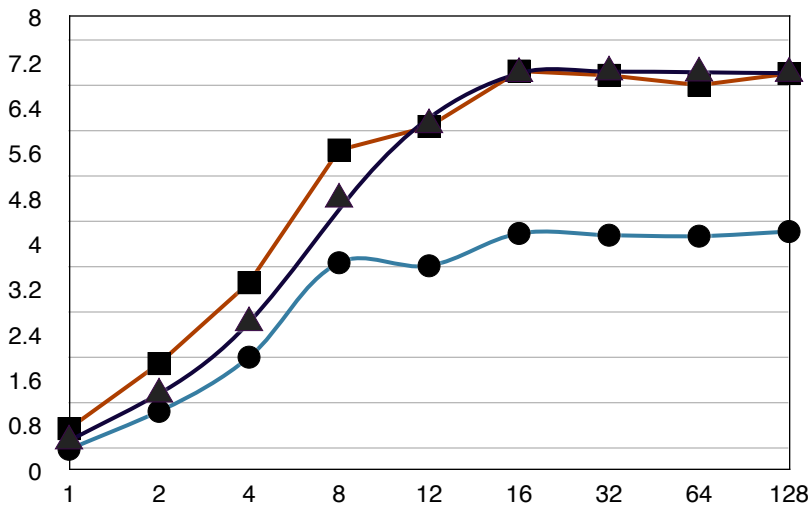
| $n(\text{chunked fifo})$ | 1 | 2 | 4 | 8 | 12 | 16 | 32 | 64 | 128 |
|--------------------------|---------|---------|---------|---------|--------|---------|--------|--------|--------|
| $T_n(n = 16777210)$ | 25.9309 | 10.0981 | 5.2382 | 2.85677 | 2.2428 | 1.95398 | 1.9461 | 1.9514 | 1.9531 |
| $T_n(n = 23947347)$ | 651.217 | 235.1 | 122.813 | 67.0412 | 68.039 | 58.7574 | 59.192 | 59.436 | 58.290 |
| $T_n(n = 33554432)$ | 29.9922 | 11.781 | 6.71209 | 3.93461 | 3.6653 | 3.15713 | 3.1913 | 3.2723 | 3.1770 |



Scalability

Again the algorithm scales well due to parallelism.

Chunk size: 64



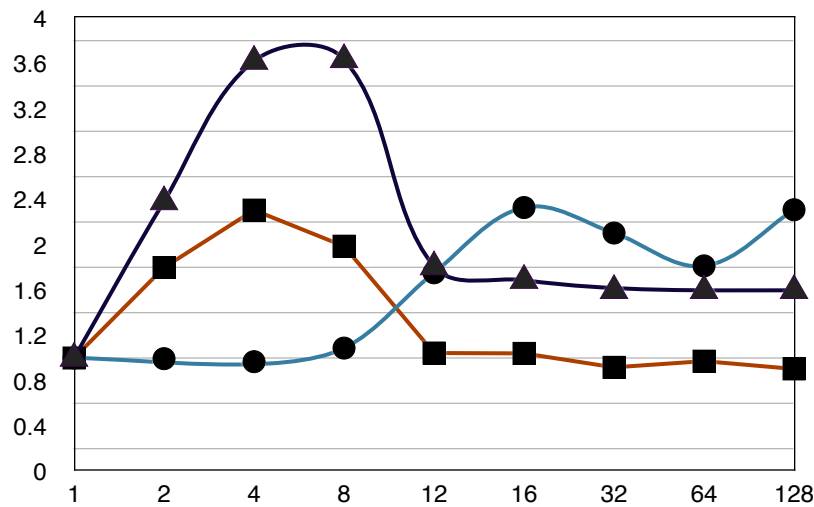
Speedup

Again the algorithm scales well due to parallelism.

Chunk size: 64

5. Worklist implemented as LocalQueue queue

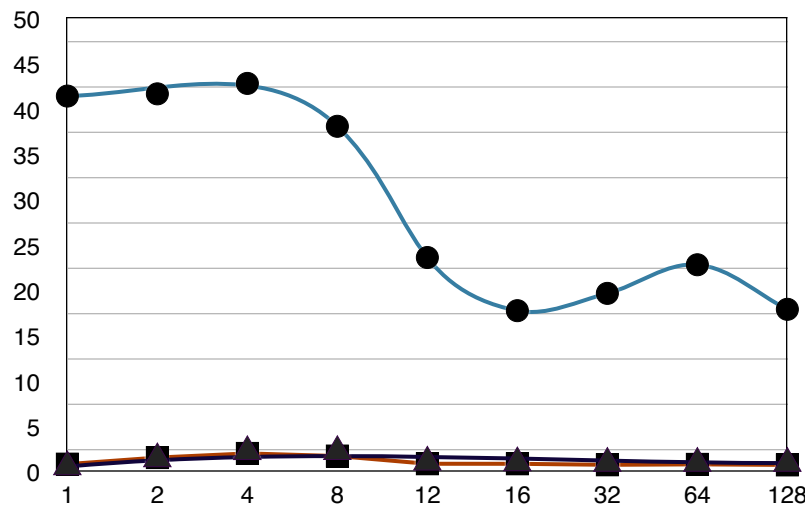
| $n(\text{rmat local queue chunk fifo})$ | 1 | 2 | 4 | 8 | 12 | 16 | 32 | 64 | 128 |
|---|---------|---------|---------|---------|--------|---------|---------|---------|---------|
| $T_n(n = 16777210)$ | 21.4701 | 9.01762 | 5.93116 | 5.9072 | 11.884 | 12.7283 | 13.441 | 13.441 | 13.441 |
| $T_n(n = 23947347)$ | 5.92695 | 5.89223 | 5.73414 | 6.4447 | 10.382 | 13.7954 | 12.4623 | 10.7439 | 13.6801 |
| $T_n(n = 33554432)$ | 24.6082 | 13.6828 | 10.6759 | 12.3865 | 23.571 | 23.6487 | 26.818 | 25.3101 | 27.2372 |



Scalability

It scales quite well for lesser number of threads in case of random and rmat graph.

However it scales well for US graph since there are clusters in graph and nodes in the local queue are updated more frequently.

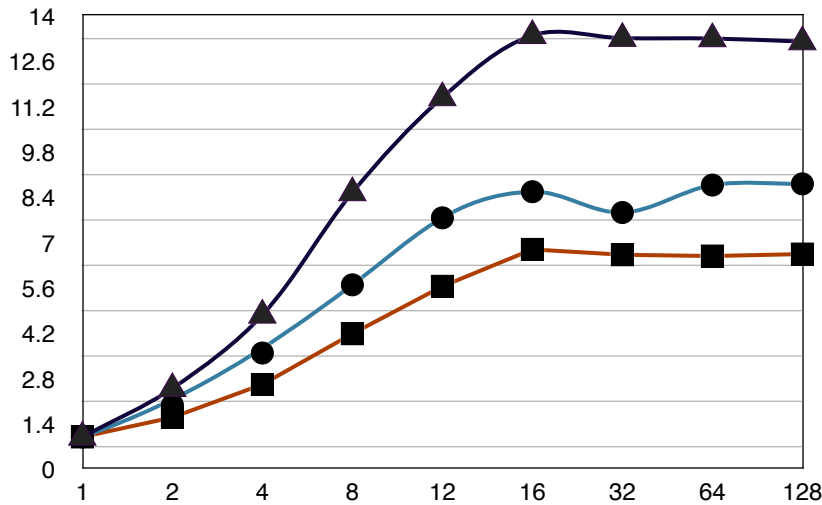


Speedup

Speed is obtained but the overhead takes over in later stages when as there are more threads competing for the same queue.

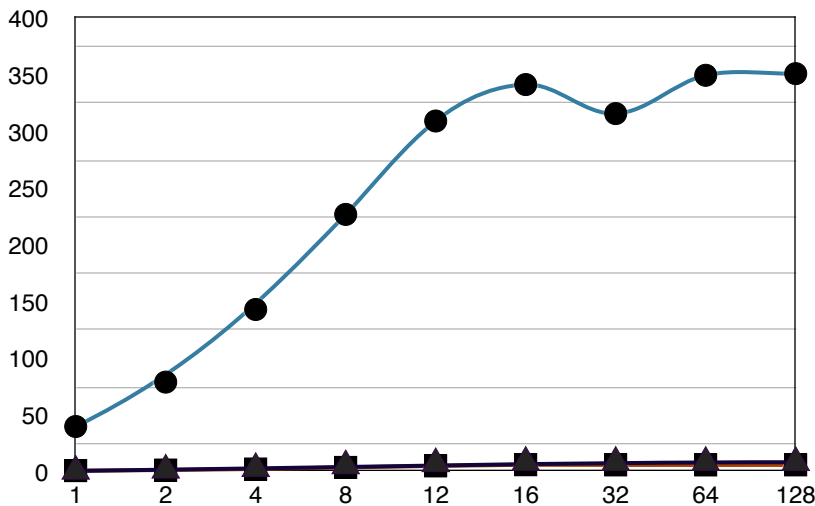
6. Worklist implemented as LocalQueue with dChunkFIFO<64>

| $n(\text{rmat local queue dchunk fifo})$ | 1 | 2 | 4 | 8 | 12 | 16 | 32 | 64 | 128 |
|--|---------|---------|---------|---------|--------|---------|---------|---------|---------|
| $T_n (n = 16777210)$ | 21.7539 | 8.70306 | 4.56197 | 2.54503 | 1.8945 | 1.62324 | 1.6337 | 1.6351 | 1.6468 |
| $T_n (n = 23947347)$ | 6.17394 | 3.11272 | 1.71899 | 1.08367 | 0.7949 | 0.72011 | 0.77847 | 0.70328 | 0.70034 |
| $T_n (n = 33554432)$ | 24.8104 | 15.4735 | 9.46749 | 5.92441 | 4.3896 | 3.6532 | 3.74414 | 3.76669 | 3.73196 |



Scalability:

Huge scalability can be seen with the LocalQueue implementation with dChunkedFIFO scheduling.



Speedup:

Speedup for USA road network is exceptionally high since it has many clusters and using this scheduling policy is the optimal.

Observations

The data structure used for parallel implementation of the algorithm in Galois.

Graph was stored as LC_CSR graph, since we only have local computations and the component id was stored as a node data. Various scheduling policies that were experimented with are:

1. dChunkedFIFO with OBIM
2. dChunkedFIFO
3. FIFO
4. ChunkedFIFO
5. LocalQueue with dChunkedFIFO as scheduling policy
6. LocalQueue with default scheduling policy.

It can be seen from the above results that the scheduling policy is input and algorithm dependent. For some inputs like USA-road a huge amount of speedup can be seen in case of LocalQueues with scheduling policy as dChunkedFIFO as compared to the other algorithms.

It can be seen that the scheduling policy of priority queue for USA-road network graph has a dramatically low runtime than the FIFO counterpart. A similar result can be seen in case of parallel implementations also. This is due to the fact that there are very few connected components in the graph(means large clusters) and this means that there is lot of work that needs to be done and since we are initializing the wordlist to be all the nodes then the nodes with updated component id which is smaller are added to the end in case of non-priority scheduling and thus more work is generated. Hence a scheduling policy of priority queue is optimal for the given algorithm. In addition due to lower overhead of threads, LocalQueue with dChunkedFIFO performed quite well since the tasks were ran in parallel.

Overall Galois provides a very easy way to try out different scheduling policies and actually focussing on the algorithm in hand rather than the implementation details and parallelization part. However, it is important to choose the right amount of scheduling policy for a given input.