# Issue XL: Modal Homotopy Type Theory

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#### Анотація

Formal definition of Cohesive Topos. **Keywords**: Differential Geometry, Topos Theory, Modal HoTT

### Зміст

1	Mod	dal Homotopy Type Theory	1
	1.1	Preliminaries	1
	1.2	Topos	4
	1.3	Geometric Morphism	4
	1.4	Cohesive Topos	٠
	1.5	Cohesive Adjunction Diagram and Modalities	
	1.6	Cohesive Modalities	4
	1.7	Differential Cohesion	,
		Graded Differential Cohesion	
	1.9	Adjoint String of Identity Modalities	,

## 1 Modal Homotopy Type Theory

#### 1.1 Preliminaries

A category C consists of:

- A class of **objects**, Ob(C),
- A class of **morphisms**,  $\operatorname{Hom}_{\mathcal{C}}(X,Y)$ , for each pair  $X,Y \in \operatorname{Ob}(\mathcal{C})$ ,
- Composition maps  $\circ$ : Hom $(Y, Z) \times \text{Hom}(X, Y) \to \text{Hom}(X, Z)$ ,
- Identity morphisms  $id_X \in Hom(X, X)$  for each X,

satisfying associativity and identity laws.

A functor  $F: \mathcal{C} \to \mathcal{D}$  assigns to each:

- Object  $X \in \mathcal{C}$  an object  $F(X) \in \mathcal{D}$ ,
- $\bullet \ \mathrm{Morphism} \ f:X\to Y \ \mathrm{a} \ \mathrm{morphism} \ F(f):F(X)\to F(Y),$

such that  $F(id_X) = id_{F(X)}$  and  $F(g \circ f) = F(g) \circ F(f)$ .

A natural transformation  $\eta: F \Rightarrow G$  between functors  $F, G: \mathcal{C} \to \mathcal{D}$  consists of morphisms  $\eta_X: F(X) \to G(X)$  such that for every  $f: X \to Y$  in  $\mathcal{C}$ ,

$$\begin{array}{ccc} F(X) & \xrightarrow{\eta_X} & G(X) \\ F(f) & & & \downarrow G(f) \\ F(Y) & \xrightarrow{\eta_Y} & G(Y) \end{array}$$

commutes.

An adjunction between categories  $\mathcal{C}$  and  $\mathcal{D}$  consists of functors

$$F:\mathfrak{C}\leftrightarrows\mathfrak{D}:G$$

and natural transformations (unit  $\eta$  and counit  $\epsilon$ )

$$\eta: \mathrm{Id}_{\mathfrak{C}} \Rightarrow \mathsf{G} \circ \mathsf{F}, \quad \varepsilon: \mathsf{F} \circ \mathsf{G} \Rightarrow \mathrm{Id}_{\mathfrak{D}}$$

satisfying the triangle identities.

### 1.2 Topos

A **topos**  $\mathcal{E}$  is a category that:

- Has all finite limits and colimits,
- Is Cartesian closed: has exponential objects [X, Y],
- Has a subobject classifier  $\Omega$ .

### 1.3 Geometric Morphism

A geometric morphism  $f: \mathcal{E} \to \mathcal{F}$  between topoi consists of an adjoint pair

$$f^*: \mathfrak{F} \leftrightarrows \mathcal{E}: f_*$$

with  $f^* \dashv f_*$ , where  $f^*$  preserves finite limits (i.e., is left exact).

### 1.4 Cohesive Topos

A **cohesive topos** is a topos  $\mathcal{E}$  equipped with a quadruple of adjoint functors:

$$\Pi \dashv \Delta \dashv \Gamma \dashv \nabla : \mathcal{E} \leftrightarrows \mathbf{Set}$$

such that:

- $\Gamma$  is the global sections functor,
- $\Delta$  is the constant sheaf functor,
- $\nabla$  sends a set to a codiscrete object,
- $\bullet~\Pi$  is the shape or fundamental groupoid functor,
- $\Delta$  and  $\nabla$  are fully faithful,
- $\Delta$  preserves finite limits,
- $\bullet~\Pi$  preserves finite products (in some variants).

### 1.5 Cohesive Adjunction Diagram and Modalities

$$\varepsilon \xrightarrow{\begin{picture}(50,0) \put(0,0){\line(1,0){100}} \put(0,0){\line(1,0)$$



### 1.6 Cohesive Modalities

The above adjoint quadruple canonically induces a triple of endofunctors on  $\mathcal{E}$ :

$$(\int \dashv \flat \dashv \sharp) : \mathcal{E} \to \mathcal{E}$$

defined as follows:

$$\int := \Delta \circ \Pi$$
$$\flat := \Delta \circ \Gamma$$
$$\sharp := \nabla \circ \Gamma$$

This yields an **adjoint triple** of endofunctors on  $\mathcal{E}$ :

$$\int -|b| + |\pm|$$

These are:

- $\int$  the **shape modality**: captures the fundamental shape or homotopy type,
- b the **flat modality**: forgets cohesive structure while remembering discrete shape,
- # the **sharp modality**: codiscretizes the structure, reflecting the full cohesion.

Each of these is an **idempotent** (co)monad, hence a *modality* in the internal language (type theory) of  $\mathcal{E}$ .

### 1.7 Differential Cohesion

A differential cohesive topos is a cohesive topos  $\mathcal{E}$  equipped with an additional adjoint triple of endofunctors:

$$(\mathfrak{R}\dashv\mathfrak{I}\dashv\mathfrak{L}):\mathcal{E}\to\mathcal{E}$$

These are:

- $\Re$ : the **reduction modality** forgets nilpotents,
- $\Im$ : the **infinitesimal shape modality** retains infinitesimal data,
- &: the infinitesimal flat modality reflects formally smooth structure.

Important object classes:

- An object X is **reduced** if  $\Re(X) \cong X$ .
- It is **coreduced** if  $\&(X) \cong X$ .
- It is **formally smooth** if the unit map  $X \to \& X$  is an effective epimorphism.

Formally étale maps are those morphisms  $f: X \to Y$  such that the square

$$\begin{array}{ccc} X & \longrightarrow \mathfrak{I}X \\ \downarrow^{\mathfrak{I}} & & \downarrow^{\mathfrak{I}(f)} \\ Y & \longrightarrow \mathfrak{I}Y \end{array}$$

is a pullback.

### 1.8 Graded Differential Cohesion

In **graded differential cohesion**, such as used in synthetic supergeometry, one introduces an adjoint triple:

$$10) \Rightarrow \dashv \rightsquigarrow \dashv Rh$$

$$(\Rightarrow \dashv \rightsquigarrow \dashv Rh) : \mathcal{E} \to \mathcal{E}$$

These are:

- ullet  $\rightrightarrows$ : the **fermionic modality** captures anti-commuting directions,
- ullet  $\leadsto$ : the **bosonic modality** filters out fermionic directions,
- Rh: the **rheonomic modality** encodes constraint structures.

These modal operators form part of the internal logic of supergeometric or supersymmetric type theories.

### 1.9 Adjoint String of Identity Modalities

In Homotopy Type Theory (HoTT), identity systems (Contractible, Strict Id, Quotient, Isomorphism, Path = Equivalence) are modeled as modalities in the  $\infty$ -topos  $\mathcal{E} = \infty$ Grp. We construct an adjoint quadruple extending the Jacobs-Lawvere triple C  $\dashv$  Id<sub>A</sub>  $\dashv$  Q(-/  $\sim$ ), incorporating Isomorphism and Path = Equivalence. The modalities are ordered by adjointness: Contractible  $\leq$  Strict Id  $\leq$  Quotient  $\leq$  Isomorphism  $\leq$  Path = Equivalence, reflecting their structure in HoTT, where Strict Id, Quotient, and Path = Equivalence are mere propositions for h-sets, while Isomorphism is not.

Homotopy Type Theory (HoTT) provides a framework for reasoning about equality via the identity type  $\mathrm{Id}_A(x,y)$ . In the  $\infty$ -topos  $\mathcal{E}=\infty\mathrm{Grp}$ , identity systems are modalities (monads), ordered by adjointness. The classical Jacobs-Lawvere adjunction triple  $C\dashv \mathrm{Id}_A\dashv Q(-/\sim)$  captures Contractible, Strict, and Quotient. We extend this to a quadruple, including Isomorphism and Path = Equiv, respecting the HoTT equivalence of Path and Equivalence and the propositional nature of Strict Id, Quotient, and Path = Equivalence for h-sets.

#### **Definition 1.** In HoTT, the identity systems are:

- Contractible: (-1)-truncated types, mere propositions.
- Strict:  $Id_A(x, y)$  for h-sets (0-truncated), a mere proposition.
- Quotient: Set-quotients  $A/\sim$ , 0-truncated, equivalent to Strict Id.
- Isomorphism:  $Iso_A(x, y)$ , a triple (f, g, p), not a mere proposition.
- Path = Equiv: Path<sub>A</sub>(x, y)  $\simeq$  ( $x \simeq y$ ), equivalent in HoTT.

In  $\mathcal{E} = \infty$ Grp, we define categories:

- $\mathcal{E}_{contr} = \mathcal{E}_{<-1}$ : Mere propositions.
- $\mathcal{E}_{\text{strict}} = \mathcal{E}_{\leq 0} \cong \text{Set: h-sets (Strict Id)}.$
- $\mathcal{E}_{quot} = \mathcal{E}_{<0} \cong Set: h\text{-sets (Quotient)}.$
- $\mathcal{E}_{iso} \cong \mathcal{E}$ :  $\infty$ -groupoids with isomorphisms.
- $\mathcal{E}_{path/equiv} \cong \mathcal{E}$ :  $\infty$ -groupoids with paths/equivalences.

The Jacobs-Lawvere triple  $C\dashv \mathrm{Id}_A\dashv Q(-/\sim)$  is extended to an adjoint quadruple:

$$\mathcal{E}_{\mathrm{contr}} \xrightarrow{F_4} \mathcal{E}_{\mathrm{strict}} \xrightarrow{F_3} \mathcal{E}_{\mathrm{quot}} \xrightarrow{F_2} \mathcal{E}_{\mathrm{iso}} \xrightarrow{F_1} \mathcal{E}_{\mathrm{path/equiv}}$$

**Theorem 1.** The functors form an adjoint quadruple with adjunctions:

$$F_4 \dashv U_4$$
,  $F_3 \dashv U_3$ ,  $F_2 \dashv U_2$ ,  $F_1 \dashv U_1$ 

- $F_4: \mathcal{E}_{\mathbf{contr}} \to \mathcal{E}_{\mathbf{strict}}$ : Inclusion of (-1)-truncated objects into 0-truncated objects. Right adjoint  $U_4$ : (-1)-truncation,  $U_4(X) = ||X||_{-1}$ .
- F<sub>3</sub>: ε<sub>strict</sub> → ε<sub>quot</sub>: Canonical map to quotient structure, viewing h-sets as quotiented by trivial relations. Right adjoint U<sub>3</sub>: Inverse map preserving h-set structure.
- $F_2: \mathcal{E}_{\mathbf{quot}} \to \mathcal{E}_{\mathbf{iso}}$ : Inclusion of h-sets into  $\mathcal{E}$ ,  $\operatorname{core}(X) \cong X$ . Right adjoint  $U_2: 0$ -truncation,  $U_2(X) = ||X||_0$ .
- $F_1: \mathcal{E}_{\mathbf{iso}} \to \mathcal{E}_{\mathbf{path/equiv}}$ : Canonical inclusion of  $\infty$ -groupoids with isomorphisms into full  $\infty$ -groupoids with paths/equivalences. Right adjoint  $U_1$ : Core map, preserving isomorphism structure.

The adjunctions induce the ordering:

 $Contractible \leq Strict\ Id \leq Quotient \leq Isomorphism \leq Path = Equivalence$ 

- Contractible: Coarsest, mere propositions ((-1)-truncated).
- Strict: h-sets,  $Id_A(x, y)$  is a mere proposition.
- Quotient: Equivalent to Strict Id, 0-truncated set-quotients.
- **Isomorphism**:  $Iso_A(x, y)$  is not a mere proposition for general types.
- Path = Equivalence: Finest, full ∞-groupoid structure, equivalent via univalence.

The adjoint quadruple extends the Jacobs-Lawvere triple, capturing the structure of identity systems in HoTT. The ordering reflects their increasing complexity, with Strict Id, Quotient, and Path = Equivalence collapsing to mere propositions for h-sets, while Isomorphism retains higher structure. Future work could explore these adjunctions in other  $\infty$ -topoi or specific CTT models.