## 清华大学本科生考试试题专用纸

考试课程: 操作系统(A卷) 时间: 2011年06月22日下午2:30~4:30

任课教	<b>敛师:</b>	系别:	班级:	学号 <b>:</b>	姓名:	
		—————————————————————————————————————				
		2. 在答卷本上答题时, 要写明题号, 不必抄题。				
		3. 答题时, 要书写清楚和整洁。				
4. 请注意回答所有试题。本试卷有7个题目,共23页。						
	5. 考	5. 考试完毕, 必须将试题纸和答卷本一起交回。				
成从用户 需要补全I user/libs/	态函数re 的代码只 ′file.c	**	tsysfile_read(		所缺的代码,以正确完 这回过程。提示:每处	
int read(int fo }	l, void .(1)	*base, size_t l	en) {			
user/libs/	/syscal					
 #define MAX			5			
"act the MAN	L_ARGS		3			
va va ui in fo }	num, . t ret; _list a _start( nt32_t t i;	ap; ap, num); a[MAX_ARGS]; 0; i < MAX_ARGS a[i] = va_arg(				
as	m volat	:ile (				
		"int	%1;"			

```
: "i" (T_SYSCALL),
                          "a" (num),
                          "d" (a[0]),
                          "c" (a[1]),
                          "b" (a[2]),
                          "D" (a[3]),
                          "S" (a[4])
                         : "cc", "memory");
        return ret;
}
. . .
int
sys_read(int fd, void *base, size_t len) {
       ...(2)...
}
______
libs/stdarg.h
-----
typedef char * va_list;
#define
__va_size(type)
        ((sizeof(type) + (sizeof(long) - 1)) / sizeof(long) * sizeof(long))
#define va_start(ap, last)
        ((ap) = (va\_list)&(last) + \__va\_size(last))
#define va_arg(ap, type)
        (*(type *)((ap) += __va_size(type), (ap) - __va_size(type)))
#define va_end(ap)
                           ((void)0)
_____
libs/unistd.h
#define T_SYSCALL
                                       0x80
/* syscall number */
#define SYS_read
                                   102
#define SYS_write
                                   103
```

```
kern/syscall/syscall.c
struct trapframe {
         struct pushregs tf_regs;
         uint16_t tf_es;
        uint16_t tf_padding1;
        uint16_t tf_ds;
        uint16_t tf_padding2;
        uint32_t tf_trapno;
        /* below here defined by x86 hardware */
        uint32_t tf_err;
        uintptr_t tf_eip;
        uint16_t tf_cs;
        uint16_t tf_padding3;
        uint32_t tf_eflags;
         /* below here only when crossing rings, such as from user to kernel */
        uintptr_t tf_esp;
         uint16_t tf_ss;
        uint16_t tf_padding4;
};
kern/trap/trap.c
. . .
static void
trap_dispatch(struct trapframe *tf) {
         char c;
         int ret;
         switch (...(3)...) {
         case T_DEBUG:
         case T_BRKPT:
                   debug_monitor(tf);
                   break;
         case T_PGFLT:
                   if ((ret = pqfault_handler(tf)) != 0) {
                            print_trapframe(tf);
                            if (current == NULL) {
                                     panic("handle pgfault failed. %e\n", ret);
                            }
                            else {
                                     if (trap_in_kernel(tf)) {
```

```
panic("handle pgfault failed in
kernel mode. %e\n", ret);
                                      }
                                      cprintf("killed by kernel.\n");
                                      do_exit(-E_KILLED);
                            }
                   break;
         case T_SYSCALL:
                   ...(4)...
                   break;
         case IRQ_OFFSET + IRQ_TIMER:
                   ticks ++;
                   assert(current != NULL);
                   run_timer_list();
                   break;
         case IRQ_OFFSET + IRQ_COM1:
         case IRQ_OFFSET + IRQ_KBD:
                  if ((c = cons_getc()) == 13) {
                            debug_monitor(tf);
                   }
                   else {
                            extern void dev_stdin_write(char c);
                            dev_stdin_write(c);
                   }
                   break;
         case IRQ_OFFSET + IRQ_IDE1:
         case IRQ_OFFSET + IRQ_IDE2:
                   /* do nothing */
                   break;
         default:
                   print_trapframe(tf);
                   if (current != NULL) {
                            cprintf("unhandled trap.\n");
                            do_exit(-E_KILLED);
                   panic("unexpected trap in kernel.\n");
         }
}
void
trap(struct trapframe *tf) {
         // used for previous projects
         if (current == NULL) {
                   trap_dispatch(tf);
         }
         else {
```

```
// keep a trapframe chain in stack
                   struct trapframe *otf = current->tf;
                   current->tf = tf;
                   bool in_kernel = trap_in_kernel(tf);
                   ...(5)...
                   current->tf = otf;
                   if (!in_kernel) {
                            if (current->flags & PF_EXITING) {
                                     do_exit(-E_KILLED);
                            }
                            if (current->need_resched) {
                                     schedule();
                            }
                   }
         }
}
kern/syscall/syscall.c
static int
sys_read(uint32_t arg[]) {
        int fd = (int)arg[0];
         size_t len = (size_t)...(6)...;
        void *base = (\text{void *})...(7)...;
         ...(8a)...
}
static int (*syscalls[])(uint32_t arg[]) = {
         [SYS_read]
                                                   sys_read,
         [SYS_write]
                                                   sys_write,
. . .
         [SYS_mkfifo]
                                                  sys_mkfifo,
};
                             ((sizeof(syscalls)) / (sizeof(syscalls[0])))
#define NUM_SYSCALLS
void
syscall(void) {
         struct trapframe *tf = current->tf;
```

```
uint32_t arg[5];
         int num = tf \rightarrow ...(8b)...;
         if (num >= 0 && num < NUM_SYSCALLS) {
                   if (syscalls[num] != NULL) {
                            arg[0] = tf->tf_regs.reg_edx;
                            arg[1] = tf->tf_regs.reg_ecx;
                            ara[2] = tf->tf_reqs.req_ebx;
                            arg[3] = tf->tf_regs.reg_edi;
                            arg[4] = tf->tf_regs.reg_esi;
                            tf \rightarrow tf_regs.reg_eax = ...(9)...;
                            return ;
                   }
         print_trapframe(tf);
         panic("undefined syscall %d, pid = %d, name = %s.\n",
                            num, current->pid, current->name);
}
kern/fs/sysfile.c
-----
int
sysfile_read(int fd, void *base, size_t len) {
         struct mm_struct *mm = current->mm;
         if (len == 0) {
                   return 0;
         if (!file_testfd(fd, 1, 0)) {
                   return -E_INVAL;
         void *buffer;
        if ((buffer = kmalloc(IOBUF_SIZE)) == NULL) {
                   return -E_NO_MEM;
        }
         int ret = 0;
         size_t copied = 0, alen;
         while (len != 0) {
                   if ((alen = IOBUF_SIZE) > len) {
                            alen = len;
                   }
                   ret = ...(10)...;
                   if (alen != 0) {
                            lock_mm(mm);
```

```
if (copy_to_user(mm, base, buffer, alen)) {
                                                assert(len >= alen);
                                                base += alen, len -= alen, copied
+= alen;
                                     else if (ret == 0) {
                                                ret = -E_INVAL;
                                     }
                            }
                            unlock_mm(mm);
                   if (ret != 0 || alen == 0) {
                            goto out;
                   }
        }
out:
         kfree(buffer);
         if (copied != 0) {
                   return copied;
         return ret;
}
kern/fs/file.c
int
file_read(int fd, void *base, size_t len, size_t *copied_store) {
         int ret;
         struct file *file;
         *copied_store = 0;
        if ((ret = fd2file(fd, &file)) != 0) {
                   return ret;
         }
         if (!file->readable) {
                   return -E_INVAL;
         }
         filemap_acquire(file);
         struct iobuf __iob, *iob = iobuf_init(&__iob, base, len, file->pos);
         ret = vop_read(file->node, iob);
         size_t copied = iobuf_used(iob);
         if (file->status == FD_OPENED) {
```

```
file->pos += copied;
       *copied_store = copied;
       filemap_release(file);
       return ret;
}
      (10分)给出程序fork.c的输出结果。注: 1) getpid()和getppid()是两个系统调
  用,分别返回本进程标识和父进程标识。2)你可以假定每次新进程创建时生成的进程标识是顺
  序加1得到的;在进程标识为1000的命令解释程序shell中启动该程序的执行。
#include <sys/types.h>
#include <unistd.h>
/* getpid() and fork() are system calls declared in unistd.h. They return */
/* values of type pid_t. This pid_t is a special type for process ids. */
/* It's equivalent to int. */
int main(void)
  pid_t childpid;
   int x = 5;
     int i;
   childpid = fork();
  for (i = 0; i < 2; i++) {
     printf("This is process %d; childpid = %d; The parent of this process has
id %d; i = %d; x = %d\n", qetpid(), childpid, qetppid(), i, x);
           sleep(1);
     X++;
  }
   return 0;
}
      (18分)调度器是操作系统内核中依据调度算法进行进程切换选择的模块。1)试描述时间
  片轮转算法(Round Robin)的基本原理。2)下面代码是ucore中调度器和时间片轮转算法
  的实现代码。请补全其中所缺代码,以实现调度器和调度算法的功能。提示:每处需要补全的
  代码只需要一行,一共有7个空要填。
sched.h
struct proc_struct;
typedef struct {
   unsigned int expires;
```

```
struct proc_struct *proc;
   list_entry_t timer_link;
} timer_t;
#define le2timer(le, member)
   to_struct((le), timer_t, member)
static inline timer_t *
timer_init(timer_t *timer, struct proc_struct *proc, int expires) {
   timer->expires = expires;
   timer->proc = proc;
   list_init(&(timer->timer_link));
   return timer;
}
struct run_queue;
struct sched_class {
   const char *name;
   void (*init)(struct run_queue *rq);
   void (*enqueue)(struct run_queue *rq, struct proc_struct *proc);
   void (*dequeue)(struct run_queue *rq, struct proc_struct *proc);
   struct proc_struct *(*pick_next)(struct run_queue *rq);
   void (*proc_tick)(struct run_queue *rq, struct proc_struct *proc);
};
struct run_queue {
   list_entry_t run_list;
   unsigned int proc_num;
   int max_time_slice;
   list_entry_t rq_link;
};
#define le2rq(le, member)
   to_struct((le), struct run_queue, member)
void sched_init(void);
void wakeup_proc(struct proc_struct *proc);
void schedule(void);
void add_timer(timer_t *timer);
void del_timer(timer_t *timer);
void run_timer_list(void);
extern struct proc_struct *idleproc, *initproc, *current;
extern struct proc_struct *kswapd;
```

```
sched.c
static list_entry_t timer_list;
static struct sched_class *sched_class;
static struct run_queue *rq;
static inline void
sched_class_enqueue(struct proc_struct *proc) {
   if (proc != idleproc) {
       sched_class->enqueue(rq, proc);
   }
}
static inline void
sched_class_dequeue(struct proc_struct *proc) {
   sched_class->dequeue(rq, proc);
}
static inline struct proc_struct *
sched_class_pick_next(void) {
   return sched_class->pick_next(rq);
}
static void
sched_class_proc_tick(struct proc_struct *proc) {
   if (proc != idleproc) {
       sched_class->proc_tick(rq, proc);
   }
   else {
       proc->need_resched = 1;
   }
}
static struct run_queue __rq[4];
void
sched_init(void) {
   list_init(&timer_list);
   rq = \_\_rq;
   list_init(&(rq->rq_link));
   rq->max_time_slice = 8;
```

```
int i;
   for (i = 1; i < sizeof(\_rq) / sizeof(\_rq[0]); i ++) {
       list_add_before(&(rq->rq_link), &(__rq[i].rq_link));
       __rq[i].max_time_slice = rq->max_time_slice * (1 << i);</pre>
   }
   sched_class = &MLFQ_sched_class;
   sched_class->init(rq);
   cprintf("sched class: %s\n", sched_class->name);
}
void
wakeup_proc(struct proc_struct *proc) {
   assert(proc->state != PROC_ZOMBIE);
   bool intr_flag;
   local_intr_save(intr_flag);
       if (proc->state != PROC_RUNNABLE) {
          proc->state = PROC_RUNNABLE;
          proc->wait_state = 0;
          sched_class_enqueue(proc);
       }
       else {
          warn("wakeup runnable process.\n");
       }
   local_intr_restore(intr_flag);
}
void
schedule(void) {
   bool intr_flag;
   struct proc_struct *next;
   local_intr_save(intr_flag);
   {
       current->need_resched = 0;
       if (current->state == PROC_RUNNABLE) {
          ...(1)...
       if ((next = sched_class_pick_next()) != NULL) {
          ...(2)...
       }
   local_intr_restore(intr_flag);
```

```
if (next == NULL) {
       next = ...(3)...;
   }
   next->runs ++;
   if (next != current) {
       ...(4)...
}
sched_RR.c
static void
RR_init(struct run_queue *rq) {
   list_init(&(rq->run_list));
   rq - proc_num = 0;
}
static void
RR_enqueue(struct run_queue *rq, struct proc_struct *proc) {
   assert(list_empty(&(proc->run_link)));
   list_add_before(&(rq->run_list), &(proc->run_link));
   if (proc->time_slice == 0 || proc->time_slice > rq->max_time_slice) {
       ...(5)...
   }
   proc->rq = rq;
   rq->proc_num ++;
}
static void
RR_dequeue(struct run_queue *rq, struct proc_struct *proc) {
   assert(!list_empty(&(proc->run_link)) && proc->rq == rq);
   list_del_init(&(proc->run_link));
   rq->proc_num --;
}
static struct proc_struct *
RR_pick_next(struct run_queue *rq) {
   list_entry_t *le = list_next(&(rq->run_list));
   if (le != &(rq->run_list)) {
       return le2proc(le, run_link);
   }
   return NULL;
}
```

```
static void
RR_proc_tick(struct run_queue *rq, struct proc_struct *proc) {
   if (proc->time_slice > 0) {
      ...(6)...
   if (proc->time_slice == 0) {
      ...(7)...
}
struct sched_class RR_sched_class = {
   .name = "RR_scheduler",
   .init = RR_init,
   .enqueue = RR_enqueue,
   .dequeue = RR_dequeue,
   .pick_next = RR_pick_next,
   .proc_tick = RR_proc_tick,
};
proc.c
// proc_run - make process "proc" running on cpu
// NOTE: before call switch_to, should load base addr of "proc"'s new PDT
proc_run(struct proc_struct *proc) {
   if (proc != current) {
      bool intr_flag;
      struct proc_struct *prev = current, *next = proc;
      local_intr_save(intr_flag);
      {
         current = proc;
         load_esp0(next->kstack + KSTACKSIZE);
         lcr3(next->cr3);
         switch_to(&(prev->context), &(next->context));
      local_intr_restore(intr_flag);
   }
}
       (22分) 管程是操作系统提供的一种进程同步机制,利用管程可解决进程间通信时遇到的
四、
  同步互斥问题。读者-写者问题(Reader-writer problem)是一个经典的同步问题。写者
```

第 13 页/共 23 页

优先的读者-写者问题是指,假定有多个并发的读进程和写进程都要访问一个共享的数据结构,要求: (1)读写互斥; (2)写写互斥; (3)允许多个读进程同时访问; (4)只要有写进程提出申

请,其后提出申请的读进程就必须等待该写进程完成访问。下面是ucore中管程机制和写者优先的读者-写者问题的实现代码。请尝试补全其中所缺的代码,以正确实现管程机制和读者-写者间的读写操作协调。提示:文件"cdt\_wf.c"中的补全代码可能需要在一处加多行代码,其他需要补全的代码只需要一行,一共有11个空要填。

```
condition.h
typedef struct {
   int numWaiting;
   int valid;
   wait_queue_t wait_queue;
} condition_t;
#define cdtid2cdt(cdt_id)
   ((condition_t *)((uintptr_t)(cdt_id) + KERNBASE))
#define cdt2cdtid(cdt)
   ((cdt_t)((uintptr_t)(cdt) - KERNBASE))
void
condition_value_init(condition_t *cdt) {
   ...(1)...
   cdt->valid=1;
       wait_queue_init(&(cdt->wait_queue));
}
int
condition_init(){
   condition_t *cdt;
   if ((cdt = kmalloc(sizeof(condition_t))) != NULL) {
       condition_value_init( cdt );
   if (cdt != NULL) {
          return cdt2cdtid(cdt);
   return -E_INVAL;
}
int
condition_free(cdt_t cdt_id) {
   condition_t *cdt = cdtid2cdt(cdt_id);
   int ret = -E_INVAL;
   if (cdt != NULL) {
       bool intr_flag;
```

```
local_intr_save(intr_flag);
       {
              cdt->valid = 0, ret = 0;
              wakeup_queue(&(cdt->wait_queue), WT_INTERRUPTED, 1);
              kfree(cdt);
       }
       local_intr_restore(intr_flag);
   }
   return ret;
}
int
condition_wait(cdt_t cdt_id, klock_t kl_id){
   condition_t *cdt = cdtid2cdt(cdt_id);
   bool intr_flag;
   local_intr_save(intr_flag);
   ...(2)...
   wait_t __wait, *wait = &__wait;
   ...(3)...
   local_intr_restore(intr_flag);
   sys_unlock(kl_id);
   schedule();
   sys_lock(kl_id);
   //local_intr_save(intr_flag);
   //wait_current_del(&(cdt->wait_queue), wait);
   //local_intr_restore(intr_flag);
   if (wait->wakeup_flags != WT_UCONDITION) {
       return wait->wakeup_flags;
   }
   return 0;
}
int
condition_signal(cdt_t cdt_id){
   condition_t *cdt = cdtid2cdt(cdt_id);
   if (cdt == NULL) {
       return -E_INVAL;
   }
   bool intr_flag;
   local_intr_save(intr_flag);
   if (cdt->numWaiting > 0) {
```

```
wait_t *wait;
       if ((wait = wait_queue_first(&(cdt->wait_queue))) != NULL) {
              assert(wait->proc->wait_state == WT_UCONDITION);
              ...(4)...
       }
       ...(5)...
   local_intr_restore(intr_flag);
   return 0;
}
ulib.c
cdt_t
cdt_init(){
   return sys_cdt_init();
}
int
cdt_signal(cdt_t cdt_id){
   return sys_cdt_signal(cdt_id);
}
int
cdt_wait(cdt_t cdt_id ,klock_t klock_id){
   return sys_cdt_wait(cdt_id ,klock_id);
}
int
cdt_free(cdt_t cdt_id){
   return sys_cdt_free(cdt_id);
}
klock_t
klock_init(){
   return sys_klock_init();
}
int
klock_aquire(klock_t klock_id){
   return sys_klock_aquire(klock_id);
}
int
```

```
klock_release(klock_t klock_id){
   return sys_klock_release(klock_id);
}
int
klock_free(klock_t klock_id){
   return sys_klock_free(klock_id);
}
cdt_wf.c
int *active_reader : // # count of active readers
int *active_writer ; // # count of active writers
int *waiting_reader ; // # count of waiting readers
int *waiting_writer; // # count of waiting writers
cdt_t cdt_okToRead;
cdt_t cdt_okToWrite;
klock_t lock;
void
failed(void) {
   cprintf("FAIL: T.T\n");
   exit(-1);
}
void
init(void) {
   if ((cdt_okToRead = cdt_init()) < 0 || (cdt_okToWrite = cdt_init()) < 0) {</pre>
      failed();
   if ((lock = klock_init()) < 0) {</pre>
      failed();
   }
   if ((active_reader = shmem_malloc(sizeof(int))) == NULL || (active_writer =
shmem_malloc(sizeof(int))) == NULL
   | (waiting_reader = shmem_malloc(sizeof(int))) == NULL | (waiting_writer =
shmem_malloc(sizeof(int))) == NULL) {
      failed();
   *active_reader = *active_writer = *waiting_reader = *waiting_writer = 0;
}
void
check_init_value(void) {
```

```
if (cdt_okToRead < 0 || cdt_okToWrite < 0 ) {</pre>
   failed();
   }
   if (lock < 0 ) {
       failed();
   }
   if (*active_reader != 0 || *active_writer != 0 || *waiting_reader != 0 ||
*waiting_writer != 0) {
       failed();
   }
}
void
free_wf(void){
   if ( cdt_free(cdt_okToRead) < 0 || cdt_free(cdt_okToWrite) < 0 ){</pre>
       scprintf(" conditon free failed! \n");
       exit(-1);
   if ( klock_free(lock) < 0 ){</pre>
       scprintf(" kernal lock free failed! \n");
       exit(-1);
   }
}
void
start_read(void) {
   klock_aquire(lock);
       ...(6)...
       klock_release(lock);
}
void
done_read(void) {
   klock_aquire(lock);
       ...(7)...
   klock_release(lock);
}
void
start_write(void) {
   klock_aquire(lock);
       ...(8)...
   klock_release(lock);
}
void
```

```
done_write(void) {
   klock_aquire(lock);
       ...(9)...
       if ((*waiting_writer) > 0) {
           ...(10)...
       }
       else if ((*waiting_reader) > 0) {
       int wakecount=0;
       while(...(11)...){
          cdt_signal(cdt_okToRead);
          wakecount++;
       }
   klock_release(lock);
}
void
writer(int id, int time) {
   scprintf("writer %d: (pid:%d) arrive \n", id, getpid());
       start_write();
                   writer_wf %d: (pid:%d) start %d\n", id, getpid(), time);
       scprintf("
   sleep(time);
   scprintf(" writer_wf %d: (pid:%d) end %d\n", id, getpid(), time);
       done_write();
}
void
reader(int id, int time) {
   scprintf("reader %d: (pid:%d) arrive\n", id, getpid());
   start_read();
   scprintf("
                reader_wf %d: (pid:%d) start %d\n", id, getpid(), time);
   sleep(time);
   scprintf(" reader_wf %d: (pid:%d) end %d\n", id, getpid(), time);
   done_read();
}
void
read_test_wf(void) {
}
void
write_test_wf(void) {
}
```

```
void
read_write_test_wf(void) {
}
int
main(void) {
   init();
   read_test_wf();
   write_test_wf();
   read_write_test_wf();
   free_wf();
   cprintf("condition reader_writer_wf_test pass..\n");
   return 0;
}
void
wait_init(wait_t *wait, struct proc_struct *proc) {
   wait->proc = proc;
   wait->wakeup_flags = WT_INTERRUPTED;
   list_init(&(wait->wait_link));
}
void
wait_queue_init(wait_queue_t *queue) {
   list_init(&(queue->wait_head));
}
void
wait_queue_add(wait_queue_t *queue, wait_t *wait) {
   assert(list_empty(&(wait->wait_link)) && wait->proc != NULL);
   wait->wait_queue = queue;
   list_add_before(&(queue->wait_head), &(wait->wait_link));
}
void
wait_queue_del(wait_queue_t *queue, wait_t *wait) {
   assert(!list_empty(&(wait->wait_link)) && wait->wait_queue == queue);
   list_del_init(&(wait->wait_link));
}
wait_t *
```

```
wait_queue_next(wait_queue_t *queue, wait_t *wait) {
   assert(!list_empty(&(wait->wait_link)) && wait->wait_queue == queue);
   list_entry_t *le = list_next(&(wait->wait_link));
   if (le != &(queue->wait_head)) {
      return le2wait(le, wait_link);
   }
   return NULL;
}
wait_t *
wait_queue_prev(wait_queue_t *queue, wait_t *wait) {
   assert(!list_empty(&(wait->wait_link)) && wait->wait_queue == queue);
   list_entry_t *le = list_prev(&(wait->wait_link));
   if (le != &(queue->wait_head)) {
      return le2wait(le, wait_link);
   }
   return NULL;
}
wait_t *
wait_queue_first(wait_queue_t *queue) {
   list_entry_t *le = list_next(&(queue->wait_head));
   if (le != &(queue->wait_head)) {
      return le2wait(le, wait_link);
   }
   return NULL;
}
wait_t *
wait_queue_last(wait_queue_t *queue) {
   list_entry_t *le = list_prev(&(queue->wait_head));
   if (le != &(queue->wait_head)) {
      return le2wait(le, wait_link);
   }
   return NULL;
}
bool
wait_queue_empty(wait_queue_t *queue) {
   return list_empty(&(queue->wait_head));
}
bool
wait_in_queue(wait_t *wait) {
   return !list_empty(&(wait->wait_link));
}
```

```
void
wakeup_wait(wait_queue_t *queue, wait_t *wait, uint32_t wakeup_flags, bool del)
{
   if (del) {
      wait_queue_del(queue, wait);
   }
   wait->wakeup_flags = wakeup_flags;
   wakeup_proc(wait->proc);
}
void
wakeup_first(wait_queue_t *queue, uint32_t wakeup_flags, bool del) {
   wait_t *wait;
   if ((wait = wait_queue_first(queue)) != NULL) {
      wakeup_wait(queue, wait, wakeup_flags, del);
   }
}
void
wakeup_queue(wait_queue_t *queue, uint32_t wakeup_flags, bool del) {
   wait_t *wait;
   if ((wait = wait_queue_first(queue)) != NULL) {
      if (del) {
          do {
             wakeup_wait(queue, wait, wakeup_flags, 1);
          } while ((wait = wait_queue_first(queue)) != NULL);
      else {
          do {
             wakeup_wait(queue, wait, wakeup_flags, 0);
          } while ((wait = wait_queue_next(queue, wait)) != NULL);
      }
   }
}
void
wait_current_set(wait_queue_t *queue, wait_t *wait, uint32_t wait_state) {
   assert(current != NULL);
   wait_init(wait, current);
   current->state = PROC_SLEEPING;
   current->wait_state = wait_state;
   wait_queue_add(queue, wait);
proc.h
```

//the wait state (0x00000001 | WT\_INTERRUPTED) // wait child process #define WT\_CHILD #define WT\_TIMER (0x00000002 | WT\_INTERRUPTED) // wait timer #define WT KSWAPD 0x00000003 // wait kswapd to free page #define WT\_KSEM 0x00000100 // wait kernel semaphore #define WT USEM (0x00000101 | WT INTERRUPTED) // wait user semaphore #define WT\_EVENT\_SEND (0x00000110 | WT\_INTERRUPTED) // wait the sending event #define WT\_EVENT\_RECV (0x00000111 | WT\_INTERRUPTED) // wait the recving event #define WT\_MBOX\_SEND (0x00000120 | WT\_INTERRUPTED) // wait the sending mbox #define WT\_MBOX\_RECV (0x00000121 | WT\_INTERRUPTED) // wait the recving mbox #define WT\_UCONDITION (0x00000130 | WT\_INTERRUPTED) // wait user condition --liuruilin #define WT\_INTERRUPTED 0x80000000 // the wait state could be interrupted

- 五、 (10分)银行家算法(Banker's Algorithm)是一种在资源分配过程中避免出现死锁的 算法,资源管理者可以有进程申请资源时,使用银行家算法来判断分配相应资源后是否可能出 现死锁。试回答下列问题。
  - 1) 形成死锁的条件是什么?
  - 2) 试用伪代码描述银行家算法。
  - 3)假设系统中有A、B、C和D这四类资源,有P1、P2和P3这三个进程正在使用这些资源。下面某次资源申请后的资源占用情况。请问这个状态是否安全?如果是安全的,请给出一个可能的资源分配和回收序列。

当前的可用资源情况:

ABCD

3 1 1 2

当前各进程的已分配资源情况:

ABCD

P1 1 0 3 3

P2 1 2 2 1

P3 1 2 1 0

各进程声称的最大资源申请情况:

ABCD

P1 1 2 3 4

P2 3 3 2 2

P3 1 3 5 0

- 六、 (10分)基本的文件组织方式有哪几种?请用图示方式描述UNIX文件系统UFS的文件组织方式。
- 七、 (10分) 磁盘缓存置换算法的作用是什么? 试简要描述访问频率置换算法 (Frequency-based Replacement)的基本原理。磁盘调度算法的作用是什么? 试简要描述 扫描算法 (SCAN) 的基本原理。