

Making all equalities equal

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Abstract

1 Equalities

There are multiple ways of defining equalities in a theorem prover. In the next sections, they will be defined.

1.1 Imports

First, it will be necessary to give some agda arguments:

```
{-# OPTIONS --cubical --cumulativity #-}  
module paper where
```

The cubical flag is necessary because we are using cubical equality, and the cumulativity flag is also necessary for level subtyping,

```
open import Agda.Primitive.Cubical using (I; i0; i1)
```

This library loads Cubical Agda Primitives as the equality interval.

1.2 Martin-Löf Equality

At the begin of Agda and in most theorems proves, equality is given by Martin-Löf's definition:

```
module Martin-Löf {ℓ} {A : Set ℓ} where  
  
data _≡_ (x : A) : A → Set ℓ where  
  refl : x ≡ x
```

This equality is very convenient in proof assistances like Agda because it is possible to pattern match using them:

```

private variable
  x y z : A

sym : x ≡ y → y ≡ x
sym refl = refl

trans : x ≡ y → y ≡ z → x ≡ z
trans refl refl = refl

```

But the problem of this equality is that it does not handle extensionality and other axioms very well.

```

module FunExt {ℓ ℓ'} {A : Set ℓ} {B : Set ℓ'} where
  open Martin-Löf

  funExt-Type = {f g : A → B}
    → ((x : A) → f x ≡ g x) → f ≡ g

```

1.3 Cubical Equality

To solve this problem, Agda adopted cubical type theory that equality is a function from the path to type:

```

module CubicalEquality {ℓ} {A : Set ℓ} where
  postulate
    PathP : (A : I → Set ℓ) → A i0 → A i1 → Set ℓ

  .≡_ : A → A → Set ℓ
  .≡_ = PathP λ _ → A

```

From this equality, I will define reflection, symmetry and extensionality:

```

module CubicalResults {ℓ ℓ'} {A : Set ℓ} {B : Set ℓ'} where
  open import Cubical.Core.Primitives

  private variable
    x y z : A

  refl : x ≡ x
  refl {x = x} = λ _ → x

  sym : x ≡ y → y ≡ x
  sym p i = p (~ i)

  funExt : {f g : A → B}
    → ((x : A) → f x ≡ g x) → f ≡ g
  funExt p i x = p x i

```

The operator \sim invert the interval. If the interval i goes from $i0$ to $i1$, the interval $\sim i$ goes from $i1$ to $i0$.

1.4 Leibniz equality

Leibniz equality is defined in this way: If a is equal to b , then for every propositional P , if $P a$, then $P b$. The main idea is that if both values are equal, then they are seen equal for every angle.

```
module LeibnizEquality {A : Set} where
  _≐_ : A → A → Set
  a ≐ b = (P : A → Set) → P a → P b
```

2 Joining all equalities

All equalities have something in common. They are all equal to each other. So it will be defined as a common record that all equalities should have. In the next definition, all equalities are equal to cubical equality:

```
open import Cubical.Foundations.Prelude
open import Cubical.Foundations.Isomorphism
open import Cubical.Foundations.Equiv
open import Cubical.Foundations.Univalence
open import Cubical.Foundations.Function
open import Cubical.Data.Equality

module _ {a ℓ} {A : Set a} where
  ≐-Type = A → A → Set ℓ
  private
    ℓ1 = ℓ-max a ℓ

  private variable
    x y z : A

  record IsEquality (≐- : ≐-Type) : Set (ℓ-suc (ℓ-max a ℓ)) where
    constructor eq
    field
      ≐-≡≡ : let
        x≡y : Type ℓ1
        x≡y = x ≡ y

        x≐y : Type ℓ1
        x≐y = x ≐ y

        in ≐- {ℓ-suc ℓ1} x≐y x≡y

      ≡≡≐ : let
        x≡y : Type (ℓ-max a ℓ)
```

```

 $x \equiv y = x \equiv y$ 
 $x \triangle y = x \triangle y$ 
in  $x \equiv y \equiv x \triangle y$ 
 $\equiv \equiv \triangle = \text{sym } \triangle \equiv \equiv$ 

module _ {  $\triangle$  :  $\triangle$ -Type } where
  sym-Equality : ( $\equiv \equiv \triangle$  : {  $x y : A$  }  $\rightarrow$ ) let
     $x \equiv y$  : Type ( $\ell$ -max  $a \ell$ )
     $x \equiv y = x \equiv y$ 
     $x \triangle y = x \triangle y$ 
    in  $x \equiv y \equiv x \triangle y$ 
     $\rightarrow$  IsEquality  $\triangle$ 
  sym-Equality  $\equiv \equiv \triangle = \text{eq (sym } \equiv \equiv \triangle)$ 

record Equality : Set ( $\ell$ -suc ( $\ell$ -max  $a \ell$ )) where
  constructor eqC
  field
     $\triangle$  :  $\triangle$ -Type
    { isEquality } : IsEquality  $\triangle$ 

EqFromInstance : {  $\triangle$  :  $\triangle$ -Type }  $\rightarrow$  IsEquality  $\triangle$   $\rightarrow$  Equality
EqFromInstance inst = eqC _ { inst }

eqsEqual : ( $\triangle$   $\triangle_1$   $\triangle$   $\triangle_2$  :  $\triangle$ -Type)
  {  $\triangle$   $\triangle_1$ -eq : IsEquality  $\triangle$   $\triangle_1$  }
  {  $\triangle$   $\triangle_2$ -eq : IsEquality  $\triangle$   $\triangle_2$  }
 $\rightarrow \forall \{x y\} \rightarrow$  let
   $x \triangle_1 y$  : Type  $\ell_1$ 
   $x \triangle_1 y = x \triangle_1 y$ 

   $x \triangle_2 y$  : Type  $\ell_1$ 
   $x \triangle_2 y = x \triangle_2 y$ 

  in  $\triangle \triangle_1 \{ \ell$ -suc  $\ell_1 \} x \triangle_1 y x \triangle_2 y$ 
eqsEqual _ _ { eq  $\triangle \equiv \equiv_1$  } { eq  $\triangle \equiv \equiv_2$  } =  $\triangle \equiv \equiv_1 \bullet \text{sym } \triangle \equiv \equiv_2$ 

```

It will be defined for each equality, its instance:

2.1 Cubical Equality

The simplest example is the cubical equality hence this equality is already equal to itself.

```

module _ {  $a$  } {  $A : \text{Set } a$  } where
  instance
     $\equiv$ -IsEquality : IsEquality {  $A = A$  }  $\equiv$ 
     $\equiv$ -IsEquality = eq refl
     $\equiv$ -Equality : Equality {  $\ell = a$  }
     $\equiv$ -Equality = eqC  $\equiv$ 

```

2.2 Martin-Löf equality

The proof of Martin-Löf equality is more difficult, but it is already in Cubical library as `p-c`.

```
instance
  ≡p-IsEquality : IsEquality {A = A} ≡p_
  ≡p-IsEquality = sym-Equality p-c
  ≡p-Equality : Equality {ℓ = a}
  ≡p-Equality = eqC ≡p_
```

2.3 Isomorphism

The isomorphism is an equality between types.

```
module _ {ℓ} where
  univalencePath' : {A B : Type ℓ} → (A ≡ B) ≡ (A ≃ B)
  univalencePath' {A} {B} =
    ua {ℓ-suc ℓ} {A ≡ B} {A ≃ B} (compEquiv (univalence {ℓ} {A} {B}))
    (isoToEquiv (iso {ℓ} {ℓ-suc ℓ}
      (λ x → x) (λ x → x) (λ b i → b) λ a i → a)))
```

`univalencePath` is already defined in Agda library, but with $A \simeq B$ instead of *Lifted* ($A \simeq B$). This change can be done because of the cumulativity flag.

```
instance
  ≃-IsEquality : IsEquality
  {A = Type ℓ} ≃-_
  ≃-IsEquality = sym-Equality univalencePath'
  ≃-Equality : Equality {ℓ = ℓ}
  ≃-Equality = eqC ≃_
```

2.4 Leibniz Equality

The hardest equality to prove that is equality is the Leibniz Equality.

```
liftIso : ∀ {a b} {A : Type a} {B : Type b}
  → Iso {a} {b} A B → Iso {ℓ-max a b} {ℓ-max a b} A B
liftIso {a} f = iso fun inv
  (λ x i → rightInv x i) (λ x i → leftInv x i)
```

This `liftIso` will be used to lift the Isomorphism to types of the same maximum level of both.

```
where open Iso f
```

```
open import leibniz
open Leibniz
```

It is importing the definition of Leibniz equality made by [?]. In this work, there is already a proof of the isomorphism between the Leibniz and the Martin-Löf equality.

```
module FinalEquality {A : Set} where
  open MainResult A

  ≐≐≐ : ∀ {a b} → Iso (a ≐ b) (a ≐p b)
  ≐≐≐ = iso j i (ptoc ∘ ji) (ptoc ∘ ij)
```

In Cubical Library, the definition of isomorphism uses cubical equality instead of Martin-Löf equality when we have to proof that $\forall x \rightarrow from (to x) \equiv x$ and $\forall x \rightarrow to (from x) \equiv x$. `ptoc` is necessary to do this conversion from these equalities.

```
≐≐≐ : ∀ {a b} → (a ≐ b) ≐c (a ≐p b)
≐≐≐ = let lifted = liftIso ≐≐≐ in isoToPath lifted
```

Using the univalence and `liftIso` defined previously, it is possible to transform the isomorphism into an equality.

```
open IsEquality

instance
  ≐-IsEquality : IsEquality {A = A} _≐_
  ≐-IsEquality = eq λ {x} {y} → ≐≐≐ •
    λ i → ≐p-IsEquality {ℓ-zero} . ≐≐≐ {x} {y} i

  ≐-Equality : Equality {ℓ = ℓ-suc ℓ-zero}
  ≐-Equality = eqC _≐_
```

The last pass is to join the three equalities between equalities: Leibniz to Martin-Löf to cubical equality.

3 New Equalities types

The equalities used previously were defined using the cubical equality. Now I will define them using other equalities.

```
module Equalities {a ℓ} {A : Set a} where
  private
    ≐-Type' = ≐-Type {a} {ℓ} {A}
    ℓ1 = ℓ-max a ℓ

  private variable
    x y z : A
```

Loaded the modules using the levels to be more generic.

```
module _
  (Eq1 : Equality {·} {ℓ} {A})
  where

  open Equality Eq1 renaming (·≐ to ·≐1; isEquality to eq1)
```

I am importing a generic equality to use it to define a more generic equality.

```
record IsEquality2 (·≐ : ≐-Type') : Set (ℓ-suc ℓ1) where
  constructor eq
  field
    ≐≐≐≐ : let
      x≐y : Type ℓ1
      x≐y = x ≐1 y
```

Different from previously definition of `IsEquality`, the cubical equality defined in the line above was substituted by the more generic equality `≐1`.

```
x≐y : Type ℓ1
x≐y = x ≐1 y

in ≐≐≐≐ {ℓ-suc ℓ1} x≐y x≐y

≐≐≐≐ : let
  x≐y : Type ℓ1
  x≐y = x ≐1 y
  x≐y = x ≐1 y
  in x≐y ≐ x≐y
≐≐≐≐ = sym ≐≐≐≐
```

The rest of the definition is the same.

```
instance
  ≐≐≐≐ : IsEquality2 ·≐
  ≐≐≐≐ = eq (≐≐≐≐ • IsEquality.≐≐≐≐ eq1)
```

From a more generic definition of equality, it is easily possible to return to the less generic definition.

```
module _
  (Eq2 : Equality {·} {ℓ} {A})
  where

  open Equality Eq2 renaming (·≐ to ·≐2; isEquality to eq2)
```

I am defining a new generic equality to prove that it is an equality of type 2:

```

eqsEqual2 : let
  xΔ1y : Type ℓ1
  xΔ1y = x ≡1 y

  xΔ2y : Type ℓ1
  xΔ2y = x ≡2 y

  in  $\equiv$  {ℓ-suc ℓ1} xΔ1y xΔ2y
eqsEqual2 = eqsEqual  $\equiv$ 1  $\equiv$ 2

instance
   $\equiv$ 2-Equality2 : IsEquality2  $\equiv$ 2
   $\equiv$ 2-Equality2 = eq (sym eqsEqual2)
  where open IsEquality

module  $\equiv$  { $\Delta$  : Δ-Type} where
  sym-Equality2 : (≡-≡- $\Delta$  : {x y : A} → let
    x≡y : Type ℓ1
    x≡y = x ≡1 y
    xΔy = x Δ y
    in x≡y ≡ xΔy)
    → IsEquality2  $\Delta$ 
  sym-Equality2 ≡-≡- $\Delta$  = eq (sym ≡-≡- $\Delta$ )

```

Given a symmetric definition of the previous equality, it is easy to prove that it is also an equality of type 2.

```

module  $\equiv$ 
  (Eq3 : Equality {A = Set ℓ1})
  where
    open Equality Eq3 renaming ( $\Delta$  to  $\equiv$ 3; isEquality to eq3)

    record IsEquality3 ( $\Delta$  : Δ-Type') : Set (ℓ-suc ℓ1) where
      constructor eq
      field
         $\Delta$ -≡-≡ : let
          x≡y : Type ℓ1
          x≡y = x ≡1 y

          xΔy : Type ℓ1
          xΔy = x Δ y

          in xΔy ≡3 x≡y

    instance

```



```

 $\hat{=}$ -isEquality  $_2$  : IsEquality  $_2$   $\hat{=}$ -
 $\hat{=}$ -isEquality  $_2$  = eq (transport (IsEquality. $\hat{=}$ -≡ eq  $_3$ )  $\hat{=}$ -≡)

 $\hat{=}$ -isEquality : IsEquality  $\hat{=}$ -
 $\hat{=}$ -isEquality = eq (IsEquality  $_2$ . $\hat{=}$ -≡  $\hat{=}$ -isEquality  $_2$  • IsEquality. $\hat{=}$ -≡ eq  $_1$ )

≡-≡- $\hat{=}$  : let
  x≡y : Type  $\ell_1$ 
  x≡y = x ≡  $_1$  y
  x $\hat{=}$ y = x  $\hat{=}$  y
  in x≡y ≡  $_3$  x $\hat{=}$ y
≡-≡- $\hat{=}$  = let
   $\alpha_1$  = IsEquality.≡-≡- $\hat{=}$  eq  $_3$ 
   $\alpha_2$  = IsEquality. $\hat{=}$ -≡ eq  $_1$ 
   $\alpha_3$  = IsEquality.≡-≡- $\hat{=}$   $\hat{=}$ -isEquality
  in transport  $\alpha_1$  ( $\alpha_2$  •  $\alpha_3$ )

module _
  (Eq  $_2$  : Equality {·} {·} {A})
  where

  open Equality Eq  $_2$  renaming ( $\hat{=}$ - to  $\hat{=}$ - $_2$ ; isEquality to eq  $_2$ )

  instance
     $\hat{=}$ - $_2$ -Equality  $_3$  : IsEquality  $_3$   $\hat{=}$ - $_2$ -
     $\hat{=}$ - $_2$ -Equality  $_3$  = eq  $\alpha$ 
    where
      open IsEquality eq  $_3$ 
       $\alpha$  : (x  $\hat{=}$ - $_2$  y) ≡  $_3$  (x ≡  $_1$  y)
       $\alpha$  = transport ≡-≡- $\hat{=}$  (IsEquality  $_2$ . $\hat{=}$ -≡ (≡  $_2$ -Equality  $_2$  (eqC  $\hat{=}$ - $_2$ )))

  module _ { $\hat{=}$ - :  $\hat{=}$ -Type'} where
    sym-Equality  $_3$  : (≡-≡- $\hat{=}$  : {x y : A}) → let
      x≡y : Type ( $\ell$ -max a  $\ell$ )
      x≡y = x ≡  $_1$  y
      x $\hat{=}$ y = x  $\hat{=}$  y
      in x≡y ≡  $_3$  x $\hat{=}$ y)
      → IsEquality  $_3$   $\hat{=}$ -
    sym-Equality  $_3$  ≡-≡- $\hat{=}$  = eq (let
       $\alpha_1$  = IsEquality.≡-≡- $\hat{=}$  eq  $_3$ 
       $\alpha_2$  = transport (sym  $\alpha_1$ ) ≡-≡- $\hat{=}$ 
      in transport  $\alpha_1$  (sym  $\alpha_2$ ))

  module LeibnizFromPEquality {A : Set} where
    open Equalities { $\ell$ -zero} { $\ell$ -suc  $\ell$ -zero}

     $\hat{=}$ -p  $_1$ - : A → A → Set  $_1$ 
    x  $\hat{=}$ -p  $_1$  y = x  $\hat{=}$ -p y

```

```

instance
  ≡p1-isEquality : IsEquality _≡p1_
  ≡p1-isEquality = eq λ {x y} → (sym λ i → let
    α : Type1
    α = p-c {ℓ-zero} {x = x} {y = y} i
    in α)

leibniz : IsEquality {A = A} _≡_
leibniz =
  IsEquality2.≡-isEquality {Eq1 = eqC _≡p_} (eq FinalEquality.≡≡≡)

```

Acknowledgements