

Making all equalities equal

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Abstract

1 Equalities

There are multiple ways of defining equalities in a theorem prover. In the next sections, they will be defined.

1.1 Imports

First, it will be necessary to give some agda arguments:

```
{-# OPTIONS --cubical --cumulativity #-}  
module paper where
```

The cubical flag is necessary because we are using cubical equality, and the cumulativity flag is also necessary for level subtyping,

```
open import Agda.Primitive.Cubical using (I; i0; i1)
```

This library loads Cubical Agda Primitives as the equality interval.

1.2 Martin-Löf Equality

At the begin of Agda and in most theorems proves, equality is given by Martin-Löf's definition:

```
module Martin-Löf {ℓ} {A : Set ℓ} where  
  
data _≡_ (x : A) : A → Set ℓ where  
  refl : x ≡ x
```

This equality is very convenient in proof assistances like Agda because it is possible to pattern match using them:

```

private variable
  x y z : A

sym : x ≡ y → y ≡ x
sym refl = refl

trans : x ≡ y → y ≡ z → x ≡ z
trans refl refl = refl

```

But the problem of this equality is that it does not handle extensionality and other axioms very well.

```

module FunExt {ℓ ℓ'} {A : Set ℓ} {B : Set ℓ'} where
  open Martin-Löf

  funExt-Type = {f g : A → B}
    → ((x : A) → f x ≡ g x) → f ≡ g

```

1.3 Cubical Equality

To solve this problem, Agda adopted cubical type theory that equality is a function from the path to type:

```

module CubicalEquality {ℓ} {A : Set ℓ} where
  postulate
    PathP : (A : I → Set ℓ) → A i0 → A i1 → Set ℓ

  .≡_ : A → A → Set ℓ
  .≡_ = PathP λ _ → A

```

From this equality, I will define reflection, symmetry and extensionality:

```

module CubicalResults {ℓ ℓ'} {A : Set ℓ} {B : Set ℓ'} where
  open import Cubical.Core.Primitives

  private variable
    x y z : A

  refl : x ≡ x
  refl {x = x} = λ _ → x

  sym : x ≡ y → y ≡ x
  sym p i = p (~ i)

  funExt : {f g : A → B}
    → ((x : A) → f x ≡ g x) → f ≡ g
  funExt p i x = p x i

```

The operator \sim invert the interval. If the interval i goes from $i0$ to $i1$, the interval $\sim i$ goes from $i1$ to $i0$.

1.4 Leibniz equality

Leibniz equality is defined in this way: If a is equal to b , then for every propositional P , if $P a$, then $P b$. The main idea is that if both values are equal, then they are seen equal for every angle.

```
module LeibnizEquality {A : Set} where
  _≐_ : A → A → Set1
  a ≐ b = (P : A → Set) → P a → P b
```

2 Joining all equalities

All equalities have something in common. They are all equal to each other. So it will be defined as a common record that all equalities should have. In the next definition, all equalities are equal to cubical equality:

```
open import Cubical.Foundations.Prelude
open import Cubical.Foundations.Isomorphism
open import Cubical.Foundations.Equiv
open import Cubical.Foundations.Univalence
open import Cubical.Foundations.Function
open import Cubical.Data.Equality

module _ {a ℓ} {A : Set a} where
  ≐-Type = A → A → Set ℓ
  private
    ℓ1 = ℓ-max a ℓ

  record IsEquality (_≐_ : ≐-Type) : Set (ℓ-suc (ℓ-max a ℓ)) where
    constructor eq
    field
      ≐≡≡≡ : ∀ {x y} → let
        x≡y : Type ℓ1
        x≡y = x ≡ y

        x≐y : Type ℓ1
        x≐y = x ≐ y

        in _≡_ {ℓ-suc ℓ1}} x≐y x≡y

      ≡≡≡≡ : {x y : A} → let
        x≡y : Type (ℓ-max a ℓ)
        x≡y = x ≡ y
        x≐y = x ≐ y
        in x≡y ≡ x≐y
```

```

≡-≡-≡ = sym ≡-≡-≡

module _ { _≡_ : ≡-Type } where
  sym-Equality : (≡-≡-≡ : {x y : A} → let
    x≡y : Type (ℓ-max a ℓ)
    x≡y = x ≡ y
    x≡y = x ≡ y
    in x≡y ≡ x≡y)
    → IsEquality _≡_
  sym-Equality ≡-≡-≡ = eq (sym ≡-≡-≡)

record Equality : Set (ℓ-suc (ℓ-max a ℓ)) where
  constructor eqC
  field
    _≡_ : ≡-Type
    { isEquality } : IsEquality _≡_

EqFromInstance : { ≡ : ≡-Type } → IsEquality ≡ → Equality
EqFromInstance inst = eqC _ { inst }

eqsEqual : ( _≡_ 1 _≡_ 2 : ≡-Type )
  { ≡_1-eq : IsEquality _≡_ 1 }
  { ≡_2-eq : IsEquality _≡_ 2 }
  → ∀ {x y} → let
    x≡_1y : Type ℓ_1
    x≡_1y = x ≡_1 y

    x≡_2y : Type ℓ_1
    x≡_2y = x ≡_2 y

    in _≡_ {ℓ-suc ℓ_1} x≡_1y x≡_2y
  eqsEqual _ _ { eq ≡-≡-≡_1 } { eq ≡-≡-≡_2 } = ≡-≡-≡_1 • sym ≡-≡-≡_2

```

It will be defined for each equality, its instance:

2.1 Cubical Equality

The simplest example is the cubical equality hence this equality is already equal to itself.

```

module _ {a} {A : Set a} where
  instance
    ≡-IsEquality : IsEquality {A = A} _≡_
    ≡-IsEquality = eq refl
    ≡-Equality : Equality {ℓ = a}
    ≡-Equality = eqC _≡_

```

2.2 Martin-Löf equality

The proof of Martin-Löf equality is more difficult, but it is already in Cubical library as `p-c`.

```
instance
  ≡p-IsEquality : IsEquality {A = A} ≡p_
  ≡p-IsEquality = sym-Equality p-c
  ≡p-Equality : Equality {ℓ = a}
  ≡p-Equality = eqC ≡p_
```

2.3 Isomorphism

The isomorphism is an equality between types.

```
module _ {ℓ} where
  univalencePath' : {A B : Type ℓ} → (A ≡ B) ≡ (A ≃ B)
  univalencePath' {A} {B} =
    ua {ℓ-suc ℓ} {A ≡ B} {A ≃ B} (compEquiv (univalence {ℓ} {A} {B}))
    (isoToEquiv (iso {ℓ} {ℓ-suc ℓ}
      (λ x → x) (λ x → x) (λ b i → b) λ a i → a)))
```

`univalencePath` is already defined in Agda library, but with $A \simeq B$ instead of *Lifted* ($A \simeq B$). This change can be done because of the cumulativity flag.

```
instance
  ≃-IsEquality : IsEquality
  {A = Type ℓ} ≃-_
  ≃-IsEquality = sym-Equality univalencePath'
  ≃-Equality : Equality {ℓ = ℓ}
  ≃-Equality = eqC ≃-_
```

2.4 Leibniz Equality

The hardest equality to proof that is equality is the Leibniz Equality.

```
liftIso : ∀ {a b} {A : Type a} {B : Type b}
  → Iso {a} {b} A B → Iso {ℓ-max a b} {ℓ-max a b} A B
liftIso {a} f = iso fun inv
  (λ x i → rightInv x i) (λ x i → leftInv x i)
```

This `liftIso` will be used to lift the Isomorphism to types of the same maximum level of both.

```
where open Iso f
```

```
open import leibniz
open Leibniz
```

It is importing the definition of Leibniz equality made by [?]. In this work, there is already a proof of the isomorphism between the Leibniz and the Martin-Löf equality.

```
module FinalEquality {A : Set} where
  open MainResult A

  ≐≐≐ : ∀ {a b} → Iso (a ≐ b) (a ≐p b)
  ≐≐≐ = iso j i (ptoc ∘ ji) (ptoc ∘ ij)
```

In Cubical Library, the definition of isomorphism uses cubical equality instead of Martin-Löf equality when we have to proof that $\forall x \rightarrow from (to x) \equiv x$ and $\forall x \rightarrow to (from x) \equiv x$. `ptoc` is necessary to do this conversion from these equalities.

```
≐≐≐ : ∀ {a b} → (a ≐ b) ≐c (a ≐p b)
≐≐≐ = let lifted = liftIso ≐≐≐ in isoToPath lifted
```

Using the univalence and `liftIso` defined previously, it is possible to transform the isomorphism into an equality.

```
open IsEquality

instance
  ≐-IsEquality : IsEquality {A = A} _≐_
  ≐-IsEquality = eq λ {x} {y} → ≐≐≐ •
    λ i → ≐p-IsEquality {ℓ-zero} . ≐≐≐≐ {x} {y} i

  ≐-Equality : Equality {ℓ = ℓ-suc ℓ-zero}
  ≐-Equality = eqC _≐_
```

The last pass is to join the three equalities between equalities: Leibniz to Martin-Löf to cubical equality.

3 New Equalities types

The equalities used previously were defined using the cubical equality. Now I will define them using other equalities.

```
module Equalities {a ℓ} {A : Set a} where
  private
    ≐-Type' = ≐-Type {a} {ℓ} {A}
    ℓ1 = ℓ-max a ℓ
```

Loaded the modules using the levels to be more generic.

```
module _
  (Eq1 : Equality {-} {ℓ} {A})
```

where

open Equality Eq_1 renaming (\triangleq to \equiv_1 ; isEquality to eq₁)

I am importing a generic equality to use it to define a more generic equality.

```
record IsEquality2 ( $\triangleq$  :  $\triangleq$ -Type') : Set ( $\ell$ -suc  $\ell_1$ ) where
  constructor eq
  field
     $\triangleq$ - $\equiv$ - $\equiv$  :  $\forall \{x\ y\} \rightarrow$  let
       $x \equiv y$  : Type  $\ell_1$ 
       $x \equiv y = x \equiv_1 y$ 
```

Different from previously definition of IsEquality, the cubical equality defined in the line above was substituted by the more generic equality \equiv_1 .

```
 $x \triangleq y$  : Type  $\ell_1$ 
 $x \triangleq y = x \triangleq y$ 

in  $\equiv_1$  { $\ell$ -suc  $\ell_1$ }  $x \triangleq y\ x \equiv y$ 

 $\equiv$ - $\equiv$ - $\triangleq$  :  $\{x\ y : A\} \rightarrow$  let
   $x \equiv y$  : Type  $\ell_1$ 
   $x \equiv y = x \equiv_1 y$ 
   $x \triangleq y = x \triangleq y$ 
in  $x \equiv y \equiv x \triangleq y$ 
 $\equiv$ - $\equiv$ - $\triangleq$  = sym  $\triangleq$ - $\equiv$ - $\equiv$ 
```

The rest of the definition is the same.

```
instance
   $\triangleq$ -isEquality : IsEquality  $\triangleq$ -
   $\triangleq$ -isEquality = eq ( $\triangleq$ - $\equiv$ - $\equiv$  • IsEquality. $\triangleq$ - $\equiv$ - $\equiv$  eq1)
```

From a more generic definition of equality, it is easily possible to return to the less generic definition.

```
module _
  ( $Eq_2$  : Equality { $\ell$ } { $\ell$ } {A})
  where

    open Equality  $Eq_2$  renaming ( $\triangleq$  to  $\equiv_2$ ; isEquality to eq2)
```

I am defining a new generic equality to prove that it is an equality of type 2:

```
eqsEqual2 :  $\forall \{x\ y\} \rightarrow$  let
   $x \triangleq_1 y$  : Type  $\ell_1$ 
```

```

 $x \dot{=}_1 y = x \equiv_1 y$ 

 $x \dot{=}_2 y : \text{Type } \ell_1$ 
 $x \dot{=}_2 y = x \equiv_2 y$ 

in  $\_ \equiv\_ \{ \ell\text{-suc } \ell_1 \} x \dot{=}_1 y x \dot{=}_2 y$ 
eqsEqual2 = eqsEqual  $\_ \equiv_1 \_ \equiv_2 \_$ 

instance
   $\equiv_2\text{-Equality}_2 : \text{IsEquality}_2 \_ \equiv_2 \_$ 
   $\equiv_2\text{-Equality}_2 = \text{eq} (\text{sym eqsEqual}_2)$ 
  where open IsEquality

module  $\_ \{ \_ \dot{=} \_ : \dot{=}\text{-Type} \}$  where
  sym-Equality2 :  $(\equiv \equiv \dot{=} : \{x y : A\} \rightarrow \text{let}$ 
     $x \equiv y : \text{Type } \ell_1$ 
     $x \equiv y = x \equiv_1 y$ 
     $x \dot{=} y = x \dot{=} y$ 
    in  $x \equiv y \equiv x \dot{=} y$ )
     $\rightarrow \text{IsEquality}_2 \_ \dot{=} \_$ 
  sym-Equality2  $\equiv \equiv \dot{=} = \text{eq} (\text{sym } \equiv \equiv \dot{=})$ 

```

Given a symmetric definition of the previous equality, it is easy to prove that it is also an equality of type 2.

```

module  $\_$ 
  ( $Eq_3 : \text{Equality } \{A = \text{Set } \ell_1\}$ )
  where

    open Equality  $Eq_3$  renaming ( $\_ \dot{=} \_$  to  $\_ \equiv_3 \_$ ; IsEquality to eq3)

    record IsEquality3 ( $\_ \dot{=} \_ : \dot{=}\text{-Type}'$ ) : Set ( $\ell\text{-suc } \ell_1$ ) where
      constructor eq
      field
         $\dot{=} \equiv \equiv : \forall \{x y\} \rightarrow \text{let}$ 
           $x \equiv y : \text{Type } \ell_1$ 
           $x \equiv y = x \equiv_1 y$ 

           $x \dot{=} y : \text{Type } \ell_1$ 
           $x \dot{=} y = x \dot{=} y$ 

          in  $x \dot{=} y \equiv_3 x \equiv y$ 

    instance
       $\dot{=}\text{-IsEquality}_2 : \text{IsEquality}_2 \_ \dot{=} \_$ 
       $\dot{=}\text{-IsEquality}_2 = \text{eq} (\text{transport } (\text{IsEquality}.\dot{=} \equiv \equiv \equiv \equiv_3) \dot{=} \equiv \equiv \equiv)$ 

```



```

 $\hat{=}$ -isEquality : IsEquality  $\_ \hat{=}$ 
 $\hat{=}$ -isEquality = eq (IsEquality  $\_ \hat{=}$   $\hat{=}$ -isEquality  $\_ \bullet$  IsEquality  $\_ \hat{=}$  eq  $\_$ )

 $\equiv$ - $\hat{=}$ - $\hat{=}$  : {x y : A} → let
  x $\equiv$ y : Type  $\ell$   $\_$ 
  x $\equiv$ y = x  $\equiv$   $\_$  y
  x $\hat{=}$ y = x  $\hat{=}$  y
  in x $\equiv$ y  $\equiv$   $\_$  x $\hat{=}$ y
 $\equiv$ - $\hat{=}$ - $\hat{=}$  = let
   $\alpha$   $\_$  = IsEquality. $\equiv$ - $\hat{=}$ - $\hat{=}$  eq  $\_$ 
   $\alpha$   $\_$  = IsEquality. $\hat{=}$ - $\hat{=}$ - $\hat{=}$  eq  $\_$ 
   $\alpha$   $\_$  = IsEquality. $\equiv$ - $\hat{=}$ - $\hat{=}$   $\hat{=}$ -isEquality
  in transport  $\alpha$   $\_$  ( $\alpha$   $\_ \bullet$   $\alpha$   $\_$ )

module  $\_$ 
  (Eq  $\_$  : Equality { $\_$ } { $\ell$ } {A})
  where

    open Equality Eq  $\_$  renaming ( $\_ \hat{=}$  to  $\_ \equiv$   $\_$ ; isEquality to eq  $\_$ )

    instance
       $\equiv$   $\_$ -Equality  $\_$  : IsEquality  $\_ \equiv$   $\_$ 
       $\equiv$   $\_$ -Equality  $\_$  = eq  $\alpha$ 
      where
        open IsEquality eq  $\_$ 
         $\alpha$  :  $\forall$  {x y} → (x  $\equiv$   $\_$  y)  $\equiv$   $\_$  (x  $\equiv$   $\_$  y)
         $\alpha$  = transport  $\equiv$ - $\hat{=}$ - $\hat{=}$  (IsEquality  $\_ \hat{=}$   $\hat{=}$ - $\hat{=}$  ( $\equiv$   $\_$ -Equality  $\_$  (eqC  $\_ \equiv$   $\_$ )))

    module  $\_$  { $\_ \hat{=}$  :  $\hat{=}$ -Type'} where
      sym-Equality  $\_$  : ( $\equiv$ - $\hat{=}$ - $\hat{=}$  : {x y : A} → let
        x $\equiv$ y : Type ( $\ell$ -max a  $\ell$ )
        x $\equiv$ y = x  $\equiv$   $\_$  y
        x $\hat{=}$ y = x  $\hat{=}$  y
        in x $\equiv$ y  $\equiv$   $\_$  x $\hat{=}$ y)
        → IsEquality  $\_ \hat{=}$   $\_$ 
      sym-Equality  $\_ \equiv$ - $\hat{=}$ - $\hat{=}$  = eq (let
         $\alpha$   $\_$  = IsEquality. $\equiv$ - $\hat{=}$ - $\hat{=}$  eq  $\_$ 
         $\alpha$   $\_$  = transport (sym  $\alpha$   $\_$ )  $\equiv$ - $\hat{=}$ - $\hat{=}$ 
        in transport  $\alpha$   $\_$  (sym  $\alpha$   $\_$ ))

    module LeibnizFromPEquality {A : Set} where
      open Equalities { $\ell$ -zero} { $\ell$ -suc  $\ell$ -zero}

       $\_ \equiv$   $\_$  : A → A → Set  $\_$ 
      x  $\equiv$   $\_$  y = x  $\equiv$   $\_$  y

      instance

```

```

≡p1-isEquality : IsEquality _≡p1_
≡p1-isEquality = eq λ {x y} → (sym λ i → let
  α : Type1
  α = p-c {ℓ-zero} {x = x} {y = y} i
  in α)

leibniz : IsEquality {A = A} _≡_
leibniz =
  IsEquality2.≡-isEquality {Eq1 = eqC _≡p_} (eq FinalEquality.≡≡≡)

```

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