## Making all equalities equal

# Guilherme Horta Alvares da Silva

December 14, 2021

#### Abstract

### 1 Equalities

There are multiple ways of defining equalities in a theorem prover. In the next sections, they will be defined.

### 1.1 Imports

First, it will be necessary to give some agda arguments:

```
{-# OPTIONS -- cubical -- cumulativity #-} module paper where
```

The cubical flag is necessary because we are using cubical equality, and the cumulative flag is also necessary for level subtyping,

```
open import Agda. Primitive. Cubical using (I; i0; i1)
```

This library loads Cubical Agda Primitives as the equality interval.

#### 1.2 Martin-Löf Equality

At the beginning of Agda and in most theorems proves, equality is given by Martin-Löf's definition:

```
module Martin-Löf \{\ell\} \{A: \mathsf{Set}\ \ell\} where \mathsf{data} \ = \ (x:A): A \to \mathsf{Set}\ \ell \text{ where} \mathsf{refl}: x \equiv x
```

This equality is very convenient in proof assistances like Agda because it is possible to pattern match using them:

```
private variable x \ y \ z : A

sym: x \equiv y \rightarrow y \equiv x

sym refl = refl

trans: x \equiv y \rightarrow y \equiv z \rightarrow x \equiv z

trans refl refl = refl
```

But the problem of this equality is that it does not handle extensionality and other axioms very well.

```
module FunExt {\ell \ell '} {A : Set \ell} {B : Set \ell'} where open Martin-Löf funExt-Type = {f g : A \rightarrow B} \rightarrow ((x : A) \rightarrow f x \equiv g x) \rightarrow f \equiv g
```

#### 1.3 Cubical Equality

To solve this problem, Agda adopted cubical type theory that equality is a function from the path to type:

```
\begin{array}{l} \textbf{module CubicalEquality } \{\ell\} \ \{A: \mathsf{Set} \ \ell\} \ \textbf{where} \\ \textbf{postulate} \\ \textbf{PathP}: (A: \mathsf{I} \to \mathsf{Set} \ \ell) \to A \ \mathsf{i0} \to A \ \mathsf{i1} \to \mathsf{Set} \ \ell \\ \underline{=}_{-}: A \to A \to \mathsf{Set} \ \ell \\ \underline{=}_{-} = \mathsf{PathP} \ \lambda_{-} \to A \end{array}
```

From this equality, I will define reflection, symmetry and extensionality:

The operator  $\sim$  invert the interval. If the interval i goes from i0 to i1, the interval  $\sim i$  goes from i1 to i0.

#### 1.4 Leibniz equality

Leibniz equality is defined in this way: If a is equal to b, then for every propositional P, if P a, then P b. The main idea is that if both values are equal, then they are seen equal for every angle.

```
module LeibnizEquality \{A : Set\} where \underline{\dot{=}}_{-} : A \to A \to Set_{-1} a \doteq b = (P : A \to Set) \to P \ a \to P \ b
```

### 2 Joining all equalities

All equalities have something in common. They are all equal to each other. So it will be defined as a common record that all equalities should have. In the next definition, all equalities are equal to cubical equality:

```
open import Cubical.Foundations.Prelude
open import Cubical.Foundations.Isomorphism
open import Cubical.Foundations.Equiv
open import Cubical.Foundations.Univalence
open import Cubical.Foundations.Function
open import Cubical. Data. Equality
module _{-}\{a \ \ell\} \ \{A : \mathsf{Set} \ a\} \ \mathsf{where}
  \triangleq-Type = A \rightarrow A \rightarrow Set \ell
  private
     \ell_1 = \ell-max a \ell
  private variable
     xyz:A
  record IsEquality (\_\triangleq\_: \triangleq-Type) : Set (\ell-suc (\ell-max a \ell)) where
     constructor eq
     field
         ≜-≡-≡ : let
           x\equiv y: Type \ell_1
           x \equiv y = x \equiv y
           x \triangleq y: Type \ell_1
           x \triangleq y = x \triangleq y
           in _{=} _{\ell} {\ell-suc \ell<sub>1</sub>} x \triangleq y x \equiv y
     =-=-≜ : let
         x\equiv y: Type (\ell-max a \ell)
```

```
x \equiv y = x \equiv y
        x \triangleq y = x \triangleq y
        in x \equiv y \equiv x \triangleq y
   ≡-≡-≜ = sym ≜-≡-≡
module _ {_≜_ : ≜-Type} where
   sym-Equality : (\equiv -\equiv -\triangleq : \{x \ y : A\} \rightarrow let
        x\equiv y: Type (\ell-max a \ell)
        x \equiv y = x \equiv y
        x \triangleq y = x \triangleq y
        in x \equiv y \equiv x \triangleq y
        → IsEquality _≜_
   sym-Equality \equiv - \equiv - \triangleq = eq (sym \equiv - \equiv - \triangleq)
record Equality : Set (\ell-suc (\ell-max a \ell)) where
   constructor eqC
   field
         _≜_ : ≜-Type
        { isEquality } : IsEquality _≜_
EgFromInstance : \{ \triangleq : \triangleq \text{-Type} \} \rightarrow \text{IsEquality} \triangleq \rightarrow \text{Equality}
EqFromInstance inst = eqC _ { inst }
\{ \triangleq_1 - eq : IsEquality \_ \triangleq_{1-} \}
   \{ \triangleq 2 - eq : IsEquality _ \triangleq 2 - \}
   \rightarrow \forall \{x \ y\} \rightarrow let
       x \triangleq y : \text{Type } \ell_1
        x \triangleq {}_{1}y = x \triangleq {}_{1}y
        x \triangleq 2y: Type \ell_1
        x \triangleq 2y = x \triangleq 2y
        in = \{\ell-suc \ell_1\} x \triangleq y x \triangleq y
eqsEqual \_ { eq \triangleq-\equiv-\equiv _1 } { eq \triangleq-\equiv-\equiv _2 } = \triangleq-\equiv-\equiv _1 • sym \triangleq-\equiv-\equiv _2
```

It will be defined for each equality, its instance:

#### 2.1 Cubical Equality

The simplest example is cubical equality hence this equality is already equal in itself.

```
module _{-}\{a\} {A : Set a} where instance = -IsEquality : IsEquality {A = A} _{-}=_{-} = -IsEquality = eq refl = -Equality : Equality {\ell = a} = -Equality = eqC _{-}=_{-}
```

#### 2.2 Martin-Löf equality

The proof of Martin-Löf equality is more difficult, but it is already in the Cubical library as p-c.

```
instance
\equivp-IsEquality : IsEquality \{A = A\} \_\equivp\_
\equivp-IsEquality = sym-Equality p-c
\equivp-Equality : Equality \{\ell = a\}
\equivp-Equality = eqC \_\equivp\_
```

#### 2.3 Isomorphism

Isomorphism is equality between types.

```
module _{-}\{\ell\} where univalencePath' : \{A \ B : \text{Type } \ell\} \rightarrow (A \equiv B) \equiv (A \simeq B) univalencePath' \{A\} \ \{B\} = ua \{\ell\text{-suc } \ell\} \ \{A \equiv B\} \ \{A \simeq B\} \ (\text{compEquiv (univalence } \{\ell\} \ \{A\} \ \{B\}) \ (\text{isoToEquiv (iso } \{\ell\} \ \{\ell\text{-suc } \ell\} \ (\lambda \ x \rightarrow x) \ (\lambda \ x \rightarrow x) \ (\lambda \ b \ i \rightarrow b) \ \lambda \ a \ i \rightarrow a)))
```

univalencePath is already defined in Agda library, but with  $A \simeq B$  instead of *Lifted*  $(A \simeq B)$ . This change can be done because of the cumulative flag.

```
instance \simeq-IsEquality: IsEquality \{A = \text{Type } \ell\} \_\simeq-\simeq-IsEquality = sym-Equality univalencePath' \simeq-Equality: Equality \{\ell = \ell\} \simeq-Equality = eqC \_\simeq-
```

#### 2.4 Leibniz Equality

The hardest equality to prove that is equality is the Leibniz Equality.

```
\begin{array}{l} \text{liftIso}: \forall \ \{a\ b\} \ \{A: \mathsf{Type}\ a\} \ \{B: \mathsf{Type}\ b\} \\ \rightarrow \mathsf{Iso} \ \{a\} \ \{b\} \ A\ B \rightarrow \mathsf{Iso} \ \{\ell\text{-max}\ a\ b\} \ \{\ell\text{-max}\ a\ b\} \ A\ B \\ \text{liftIso} \ \{a\} \ f = \mathsf{iso} \ \mathsf{fun} \ \mathsf{inv} \\ (\lambda \ x\ i \rightarrow \mathsf{rightInv}\ x\ i) \ (\lambda \ x\ i \rightarrow \mathsf{leftInv}\ x\ i) \end{array}
```

This liftlso will be used to lift the Isomorphism to types of the same maximum level of both.

```
where open lso f
```

```
open import leibniz open Leibniz
```

It is importing the definition of Leibniz equality made by [?]. In this work, there is already proof of the isomorphism between Leibniz and Martin-Löf equality.

```
module FinalEquality \{A : Set\} where open MainResult A
\stackrel{\cong}{=} : \forall \{a \ b\} \rightarrow \mathsf{Iso} \ (a \doteq b) \ (a \equiv p \ b)
\stackrel{\cong}{=} = \mathsf{iso} \ \mathsf{j} \ \mathsf{i} \ (\mathsf{ptoc} \circ \mathsf{ji}) \ (\mathsf{ptoc} \circ \mathsf{ij})
```

In Cubical Library, the definition of isomorphism uses cubical equality instead of Martin-Löf equality when we have to prove that  $\forall x \to from \ (to \ x) \equiv x$  and  $\forall x \to to \ (from \ x) \equiv x$ . ptoc is necessary to do this conversion from these equalities.

Using the univalence and liftlso defined previously, it is possible to transform the isomorphism into equality.

The last pass is to join the three equalities between equalities: Leibniz to Martin-Löf to cubical equality.

### 3 New Equalities types

The equalities used previously were defined using cubical equality. Now I will define them using other equalities.

Loaded the modules using the levels to be more generic.

I am importing generic equality to use it to define more generic equality.

```
record IsEquality _2 (_=: =-Type') : Set (\ell-suc \ell _1) where constructor eq field =-=-=: let x=y: Type \ell _1 x=y = x = _1 y
```

Different from previously definition of IsEquality, the cubical equality defined in the line above was substituted by the more generic equality  $\equiv 1$ .

```
x \triangleq y: Type \ell_1

x \triangleq y = x \triangleq y

in _= _= \{\ell\text{-suc }\ell_1\} x \triangleq y x \equiv y

\equiv _= = _= \triangleq: let

x \equiv y: Type \ell_1

x \equiv y = x \equiv _1 y

x \triangleq y = x \triangleq y

in x \equiv y \equiv x \triangleq y

\equiv _= = _= \triangleq sym \triangleq _= = _=
```

The rest of the definition is the same.

From a more generic definition of equality, it is easily possible to return to the less generic definition.

```
module \_ (Eq_2: Equality \{\_\} \{\ell\} \{A\}) where open Equality Eq_2 renaming (\_\triangleq\_ to \_\equiv_2\_; is Equality to eq _2)
```

I am defining a new generic equality to prove that it is an equality of type 2:

```
eqsEqual 2: let
   x \triangleq y : \text{Type } \ell_1
   x \triangleq y = x \equiv y
  x \triangleq 2y: Type \ell_1
  x \triangleq 2y = x \equiv 2y
   in = \{\ell-suc \ell_1\} x \triangleq 1y x \triangleq 2y
eqsEqual _2 = eqsEqual _{-}\equiv _{1-} _{-}\equiv _{2-}
instance
   \equiv 2-Equality 2: IsEquality 2 \equiv 2-
   \equiv 2-Equality 2 = eq (sym eqsEqual 2)
      where open IsEquality
module _ {_≜_: ≜-Type} where
   sym-Equality _2: (\equiv -\equiv -\triangleq : \{x \ y : A\} \rightarrow let
      x\equiv y: Type \ell_1
      x \equiv y = x \equiv {}_1 y
      x \triangleq y = x \triangleq y
      in x \equiv y \equiv x \triangleq y
      \rightarrow IsEquality <sub>2</sub> \triangleq _
   sym-Equality _2 \equiv - \equiv - \triangleq = eq (sym \equiv - \equiv - \triangleq)
```

Given a symmetric definition of the previous equality, it is easy to prove that it is also equality of type 2.

#### 3.1 Everything is an equality

In this part, a relation is equality when it is equal (using general equality) to cubical equality.

```
module _ (Eq_3 : \text{Equality } \{A = \text{Set } \ell_1\}) where open Equality Eq_3 renaming (_\(\delta_-\) to _\(\exists _3_-\); is Equality to eq_3) record Is Equality _3 (_\(\delta_- : \delta_-\) Type') : Set (\(\ell\)-suc (\ell_1)) where constructor eq field _{-\infty} = -\infty : (x \triangleq y) \equiv_3 (x \equiv_1 y)
```

With this definition of equality, it is possible to prove that if equality is equal to cubical equality, so it is equal (using the general or cubical equality) to the cubical equality.

This is proof that the symmetric definition of equality is also valid.

```
≡-≡-≜: (x ≡ _1 y) ≡ _3 (x ≜ y)

≡-≡-≜ = let

\alpha _1 = IsEquality.=-=-≜ eq _3

\alpha _2 = IsEquality.≜-=-= eq _1

\alpha _3 = IsEquality.=-=-≜ ≜-isEquality

in transport \alpha _1 (\alpha _2 • \alpha _3)
```

It is possible to prove that a general equality is equality from this definition:

```
module _ (Eq_2 : Equality {_} {\ell} {A}) where open Equality Eq_2 renaming (_\triangleq_ to _\equiv_2_; isEquality to eq_2) instance =_2-Equality _3 : IsEquality _3 _\equiv_2_ =_2-Equality _3 = eq \alpha where open IsEquality eq_3 \alpha : (x \equiv _2 y) \equiv _3 (x \equiv _1 y) \alpha = transport \equiv_=\triangleq (IsEquality _2.\triangleq-\equiv-\equiv (\equiv_2-Equality _2 (eqC _\equiv_2_)))
```

If there is proof of symmetrical equality, so it is also equality from this definition:

```
module _{-} {_{-} \triangleq_{-} : \triangleq_{-} Type'} where sym-Equality _{3} : 
 (\equiv_{-} = \triangleq : \forall {x y} \rightarrow (x \equiv_{1} y) \equiv_{3} (x \triangleq y)) \rightarrow IsEquality _{3} =_{-} sym-Equality _{3} =_{-} \triangleq_{-} eq (let \alpha _{1} = IsEquality.\equiv_{-} \equiv_{-} \triangleq_{-} eq _{3} \alpha _{2} = transport (sym \alpha _{1}) \equiv_{-} \equiv_{-} \triangleq_{-} in transport \alpha _{1} (sym \alpha _{2}))
```

### 4 Using the definitions

The best part of defining all of this stuff is that it is now easy to prove that Leibniz equality is equality.

```
module LeibnizFromPEquality {A : Set}} where open Equalities {\ell-zero} {\ell-suc \ell-zero}
= p_{1-} : A \to A \to Set_1
x \equiv p_1 \ y = x \equiv p \ y
```

I redefined this equality because it must be a set of universe one. And because of that, I have to prove again that this is an equality:

With just one line of code, it is possible now to prove that Leibniz equality is an equality from Martin-Löf Equality.

```
leibniz : IsEquality \{A=A\} \_\doteq_ leibniz = IsEquality \{Eq\ _1=eqC\ _\equiv p_-\} (eq FinalEquality.\doteq\equiv\equiv)
```

### Acknowledgements