

Phase 1

- First, check what are the functions of the bomb using: *info functions*
- Create a breakpoint at phase_1 function: *break phase_1*
- run < input* and then, *disas*
 - `0x000000000400ee4 <+4>: mov $0x402400,%esi`
 - `0x000000000400ee9 <+9>: callq 0x401338 <strings_not_equal>`
 - `0x000000000400eee <+14>: test %eax,%eax`
 - `0x000000000400ef0 <+16>: je 0x400ef7 <phase_1+23>`
 - `0x000000000400ef2 <+18>: callq 0x40143a <explode_bomb>`
- We can see that it moves \$0x402400 to %esi, which is a register used to store arguments for function calls. Since it calls a function called "string_not_equal", let's see what is this constant. To do this we step into instruction phase_1<+9> and see what's in %esi
- p (char *) \$esi*
\$3 = 0x402400 "Border relations with Canada have never been better."
- If we check now what is the other argument, in %rsi: *p (char *) \$rsi*
\$4 = 0x402400 "The input I typed."
- Now, since the code for phase_1 only checks if the strings are equal and if they are the bomb is defused, the password for phase 1 is:

"Border relations with Canada have never been better."

Phase 2

- Create a breakpoint at phase_1 function: *break phase_2*
- run < input* and then *disas*
 - `0x000000000400f02 <+6>: mov %rsp,%rsi`
 - `0x000000000400f05 <+9>: callq 0x40145c <read_six_numbers>`
 - `0x000000000400f0a <+14>: cmpl $0x1,(%rsp)`
 - `0x000000000400f0e <+18>: je 0x400f30 <phase_2+52>`
 - `0x000000000400f10 <+20>: callq 0x40143a <explode_bomb>`
 - `0x000000000400f15 <+25>: jmp 0x400f30 <phase_2+52>`
 - `0x000000000400f17 <+27>: mov -0x4(%rbx),%eax`
 - `0x000000000400f1a <+30>: add %eax,%eax`
 - `0x000000000400f1c <+32>: cmp %eax,(%rbx)`
 - `0x000000000400f1e <+34>: je 0x400f25 <phase_2+41>`
 - `0x000000000400f20 <+36>: callq 0x40143a <explode_bomb>`
 - `0x000000000400f25 <+41>: add $0x4,%rbx`
 - `0x000000000400f29 <+45>: cmp %rbp,%rbx`
 - `0x000000000400f2c <+48>: jne 0x400f17 <phase_2+27>`
 - `0x000000000400f2e <+50>: jmp 0x400f3c <phase_2+64>`
 - `0x000000000400f30 <+52>: lea 0x4(%rsp),%rbx`
 - `0x000000000400f35 <+57>: lea 0x18(%rsp),%rbp`
 - `0x000000000400f3a <+62>: jmp 0x400f17 <phase_2+27>`
 - `0x000000000400f3c <+64>: add $0x28,%rsp`
 - `0x000000000400f40 <+68>: pop %rbx`
 - `0x000000000400f41 <+69>: pop %rbp`
 - `0x000000000400f42 <+70>: retq`
- In the first lines that the function read_six_numbers receives the stack pointer as an argument. I typed 6 numbers to see what happened after that to the stack, and:
 - (gdb) p *(int *) (\$rsp)
 - \$5 = <first typed number>
 - (gdb) p *(int *) (\$rsp + 0x4)
 - \$7 = <second typed number>

- (gdb) p *(int *) (\$rsp + 0x8)
- \$8 = <third typed number>
- (gdb) p *(int *) (\$rsp + 0xc)
- \$9 = <4th typed number>
- (gdb) p *(int *) (\$rsp + 0x10)
- \$11 = <5th typed number>
- (gdb) p *(int *) (\$rsp + 0x14)
- \$12 = <6th typed number>
- The function read six integers and put them in the stack. After that there is a comparison between \$0x1 and the first read number. If they are not equal the bomb explodes. Then the 1st number must be 1.
- After that, there is a jump to <phase_2 + 52>, where:
 - 0x0000000000400f30 <+52>: lea 0x4(%rsp),%rbx
 - 0x0000000000400f35 <+57>: lea 0x18(%rsp),%rbp
 - 0x0000000000400f3a <+62>: jmp 0x400f17 <phase_2+27>

We copy the address of the second argument to %rbx and the 6th older element in the stack. Then we jump to <phase_2+27>

- In <phase_2+27>:
 - 0x0000000000400f17 <+27>: mov -0x4(%rbx),%eax
 - 0x0000000000400f1a <+30>: add %eax,%eax
 - 0x0000000000400f1c <+32>: cmp %eax,(%rbx)
 - 0x0000000000400f1e <+34>: je 0x400f25 <phase_2+41>
 - 0x0000000000400f20 <+36>: callq 0x40143a <explode_bomb>

In this part we verify if the element that came after the current element pointed by %rbx is exactly twice the first one. If not the bomb explodes. If they are we jump to <phase_2+41>.

- In <phase_2+41>:
 - 0x0000000000400f25 <+41>: add \$0x4,%rbx
 - 0x0000000000400f29 <+45>: cmp %rbp,%rbx
 - 0x0000000000400f2c <+48>: jne 0x400f17 <phase_2+27>
 - 0x0000000000400f2e <+50>: jmp 0x400f3c <phase_2+64>

In this part, we update %rbx to point to the next entered number and we compare if %rbp and %rbx points to the same place. Remember that %rbx saves the address right after the 6th entered number, so this will be true when we add 0x4 to %rbx 6 times. If they are equal we end the function, otherwise we jump back to line 27 to verify the same thing, if the next number is also twice the last one. Then, the answer to this phase is:

1 2 4 8 16 32

Other Phases

I also defused phase 3 and phase 4, which has, respectively, the passwords **3 256** and **1 0**.