# **FOLLOW Sets**

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# Objectives

▶ Compute the FOLLOW sets for the nonterminal symbols of a grammar.

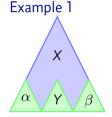
### **FOLLOW Sets**

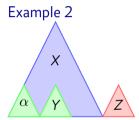
► Given a non terminal symbol *S*, what terminal symbols could come after strings that are derived from *S*?

#### The algorithm:

- Put \$ in FOLLOW(S), where S is the start symbol.
   \$ represents the "end of input."
- 2. If there is a production  $X \to \alpha Y \beta$ , then add  $FIRST(\beta)$  (but not  $\epsilon$ ) to FOLLOW(Y).
- 3. If there is a production  $X \to \alpha Y$ , or if there is a production  $X \to \alpha Y \beta$ , where  $\epsilon \in \mathit{FIRST}(\beta)$  then add  $\mathit{FOLLOW}(X)$  to  $\mathit{FOLLOW}(Y)$ .

# Diagram





- 1. If there is a production  $X \to \alpha Y \beta$ , then add  $FIRST(\beta)$  (but not  $\epsilon$ ) to FOLLOW(Y).
- 2. If there is a production  $X \to \alpha Y$ , or if there is a production  $X \to \alpha Y \beta$ , where  $\epsilon \in FIRST(\beta)$  then add FOLLOW(X) to FOLLOW(Y).

# **Small Examples**

### Example 1

 $S \rightarrow x A y$ FOLLOW set of A is  $\{y\}$ .

## Example 3

 $B \rightarrow CED$   $FIRST(D) = \{a, b\}$ FOLLOW set of D is  $\{y\}$ . FOLLOW set of E is  $\{a,b\}$ .

## Example 2

 $A \rightarrow q B$  FOLLOW set of B is also  $\{y\}$ .

## Example 4

 $B \rightarrow C E D$   $FIRST(D) = \{a, b, \epsilon\}$ FOLLOW set of D is  $\{y\}$ . FOLLOW set of E is  $\{a,b,v\}$ .

#### Grammar

```
S \rightarrow \text{if } E \text{ then } S;

S \rightarrow \text{print } E;

E \rightarrow E + E

E \rightarrow P \text{ id } P

P \rightarrow * P

P \rightarrow \epsilon
```

### Result

```
S={ $}
E={}
P={}
```

## Action

Make a chart, add \$ to S.

#### Grammar

```
S \rightarrow \text{if } E \text{ then } S ; \Leftarrow S \rightarrow \text{print } E; \\ E \rightarrow E + E \\ E \rightarrow P \text{ id } P \\ P \rightarrow * P \\ P \rightarrow \epsilon
```

## Result

```
S={$, ;}
E={ then}
P={}
```

## Action

Check productions: add then to FOLLOW(E), and; to FOLLOW(S).

#### Grammar

```
S \rightarrow \text{if } E \text{ then } S;

S \rightarrow \text{print } E; \Leftarrow

E \rightarrow E + E \Leftarrow

E \rightarrow P \text{ id } P

P \rightarrow *P

P \rightarrow \epsilon
```

## Result

```
S={$,;}
E={then,;,+}
P={}
```

## Action

Check productions: add; and + to FOLLOW(E).

#### Grammar

```
S \rightarrow \text{if } E \text{ then } S;

S \rightarrow \text{print } E;

E \rightarrow E + E

E \rightarrow P \text{ id } P \Leftarrow

P \rightarrow * P

P \rightarrow \epsilon
```

## Result

```
S={$,;}
E={then,;,+}
P={ id}
```

### Action

Check productions: add id to FOLLOW(P).

#### Grammar

```
S \rightarrow \text{if } E \text{ then } S;

S \rightarrow \text{print } E;

E \rightarrow E + E

E \rightarrow P \text{ id } P \Leftarrow

P \rightarrow * P

P \rightarrow \epsilon
```

## Result

```
S={$,;}
E={then,;,+}
P={id, then,;,+}
```

## Action

Check endings: P ends this rule, so add FOLLOW(E) to FOLLOW(P).

#### Grammar

```
S \rightarrow \text{if } E \text{ then } S;

S \rightarrow \text{print } E;

E \rightarrow E + E

E \rightarrow P \text{ id } P

P \rightarrow * P

P \rightarrow \epsilon
```

### Result

```
S={$,;}
E={then,;,+}
P={id,then,;,+}
```

### Action

Done.

#### Grammar

 $S \rightarrow Ax$ 

 $S \to By$ 

 $S \rightarrow z$  $A \rightarrow 1CB$ 

 $A \rightarrow 2B$ 

 $B \rightarrow 3B$ 

 $B \rightarrow C$ 

 $C \rightarrow 4$ 

 $C o \epsilon$ 

### Result

S = { \$}

A={}

 $B = {}$ 

C={}

#### Grammar

 $S \rightarrow Ax \Leftarrow$ 

S o By

 $S \rightarrow z$  $A \rightarrow 1CB$ 

 $A \rightarrow 2B$ 

 $B \rightarrow 3B$ 

 $B \rightarrow C$ 

 $C \rightarrow 4$ 

 $C o \epsilon$ 

## Result

S ={ \$}

 $A=\{ \frac{x}{x} \}$ 

B={}

C={}

#### Grammar

 $S \rightarrow Ax$ 

 $S \rightarrow By \Leftarrow$ 

 $S \rightarrow z$ 

A o 1CB

A 
ightarrow 2B

B o 3B

 $B \rightarrow C$ 

 $C \rightarrow 4$ 

 $C o \epsilon$ 

### Result

S ={ \$}

 $A=\{x\}$ 

 $B = {y}$ 

C={}

#### Grammar

 $S \rightarrow Ax$ 

 $\mathsf{S} o \mathsf{B} \mathsf{y}$ 

 $S \rightarrow z \Leftarrow$ 

A o 1CB

 $A \rightarrow 2B$  $B \rightarrow 3B \Leftarrow$ 

 $B \rightarrow C$ 

 $C \rightarrow 4 \Leftarrow$ 

 $C \rightarrow \epsilon \Leftarrow$ 

### Result

S ={ \$}

 $A=\{x\}$ 

 $B = {y}$ 

C={}

#### Grammar

 $S \rightarrow Ax$ 

 $S \to By$ 

 $S \to z$ 

 $A \rightarrow 1CB \Leftarrow$ 

 $A \rightarrow 2B$ 

 $B \rightarrow 3B$ 

 $B \rightarrow C$ 

 $C \rightarrow 4$ 

 $C 
ightarrow \epsilon$ 

### Result

 $S = { \$ }$ 

 $A=\{x\}$ 

 $B = \{y\}$ 

 $C=\{3, 4\}$ 

#### Grammar

 $S \rightarrow Ax$ 

 $S \rightarrow Bv$ 

 $S \rightarrow z$ 

 $A \rightarrow 1CB \Leftarrow$ 

 $A \rightarrow 2B \Leftarrow$ 

 $B \rightarrow 3B$ 

 $B \rightarrow C$ 

 $C \rightarrow 4$  $C 
ightarrow \epsilon$ 

## Result

 $S = \{ \$ \}$ 

 $A=\{x\}$ 

FOLLOW Sets

 $B = \{ x, y \}$ 

 $C={3, 4}$ 

#### Grammar

 $S \rightarrow Ax$ 

 $\mathsf{S}\to \mathsf{B}\mathtt{y}$ 

 $S \rightarrow z$ 

 $A \rightarrow 1CB \Leftarrow$ 

 $A \rightarrow 2B$ 

 $B \rightarrow 3B$ 

 $B \rightarrow C$ 

 $C \rightarrow 4$ 

 $C 
ightarrow \epsilon$ 

### Result

S ={ \$}

 $A=\{x\}$ 

 $B = \{x, y\}$ 

 $C = \{ x, 3, 4 \}$ 

#### Grammar

 $S \rightarrow Ax$ 

 $S \to By$ 

S o z

 $A \rightarrow 1CB$ 

 $A \rightarrow 2B$  $B \rightarrow 3B$ 

 $B \rightarrow C \Leftarrow$ 

 $C \rightarrow 4$ 

 $extstyle C 
ightarrow \epsilon$ 

#### Result

S ={ \$}

 $A=\{x\}$   $B=\{x, y\}$ 

 $C=\{x, y, 3, 4\}$ 

## Action

Add FOLLOW(B) to FOLLOW(C). Now we're done.