# Interpreters, Part 2

Dr. Mattox Beckman

University of Illinois at Urbana-Champaign Department of Computer Science

## Define the Types - Types.hs

```
idata Exp = IntExp Integer
          IntOpExp String Exp Exp
          RelOpExp String Exp Exp
          BoolOpExp String Exp Exp
          | BoolExp Bool
    deriving (Show, Eq)
8 data Val = IntVal Integer
          BoolVal Bool
    deriving (Show, Eq)
10
```

#### Eval – Bools, and, or

```
1 boolOps = [ ("&&", (&&))
           , ("||",(||))]
3
4 liftBoolOp f (BoolVal i1) (BoolVal i2) = BoolVal (f i1 i2)
                                           = BoolVal False
5liftBoolOp f _
6
reval (BoolExp b) _ = BoolVal b
8
9 eval (BoolOpExp op e1 e2) env =
   let v1 = eval e1 env
       v2 = eval e2 env
11
       Just f = lookup op boolOps
12
    in liftBoolOp f v1 v2
13
```

#### Adding Comparisons

```
relOps = [ ("<", (<)) , ("<=", (<=)) . (">". (>))
          . (">=", (<=)), ("==", (<=)), ("/=", (/=))]
4 liftRelOp f (IntVal i1) (IntVal i2) = BoolVal (f i1 i2)
                                     = BoolVal False
5 liftRelOp f
7 eval (RelOpExp op e1 e2) env =
   let v1 = eval e1 env
      v2 = eval e2 env
       Just f = lookup op relOps
10
    in liftRelOp f v1 v2
```

### A Simple Let Expression

We want to define local variables:

```
1 i4> 3 + let x = 2 + 3 in x * x end
2 IntVal 28
```

▶ Need two new Exp constructors.

# Coding Eval for Variables

▶ For variables, we look them up in the environment.

```
1 eval (VarExp var) env =
2   case lookup var env of
3     Just val -> val
4     Nothing -> IntVal 0
```

## Coding Eval for Let

```
leval (LetExp var e1 e2) env =
let v1 = eval e1 env
in eval e2 (insert var v1 env)
```

▶ The insert var v1 env call acts like pushing a value onto a stack!

#### **Next Time**

- ▶ You now have some interesting things for your interpreter.
- ► The reference implementation is in i4.
- ▶ We've also added a IfExp to the types if you want to try adding this.