Monads

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Objectives

- Describe the problem that monads attempt to solve.
- Know the three monad laws.
- ► Know the syntax for declaring monadic operations.
- Be able to give examples using the Maybe and List monads.

Introducing Monads

- Monads are a way of defining computation.
- A monad is a container type m along with two functions:
 - ▶ return :: a -> m a
 - ▶ bind :: m a -> (a -> m b) -> m b
 - In HASKELL, bind is written as >>=
- These functions must obey three laws:

Left identity return a >>= f is the same as f a.

Right identity m >>= return is the same as m.

Associativity (m >>= f) >>= g is the same as m >>= (x -> f x >>= g).

Understanding Return

- ▶ return :: a -> m a
- ▶ The return keyword takes an element and puts it into a monad.
- ► This is a one-way trip!
- Very much like pure in the applicative type class.

```
1 instance Monad Maybe where
2  return a = Just a
3 instance Monad [] where
4  return a = [a]
5 instance Monad (Either a) where
6  return a = Right a
```

Understanding Bind

- All the magic happens in bind.
- ▶ bind :: m a -> (a -> m b) -> m b
 - ► The first argument is a monad.
 - The second argument takes a monad, unpacks it, and repackages it with the help of the function argument.
 - Exactly how it does that is the magic part.

Bind for Maybe

```
1 Nothing >>= f = Nothing
2 (Just a) >>= f = f a
3   -- Remember that f returns a monad
```

Motivation

- ▶ They are similar to continuations, but even more powerful.
- ▶ They are also related to the applicative functors from last time.
- Consider this program:

```
inc1 a = a + 1
2r1 = inc1 <$> Just 10 -- result: Just 11
3r2 = inc1 <$> Nothing -- result: Nothing
```

But what if we have functions like this?

```
linc2 a = Just (a+1)
recip a | a =/ 0 = Just (1/a)
lotherwise = Nothing
```

How can we pass a Nothing to it? How can we use what we get from it?



Notice the Pattern

- ► Applicatives take the values out of the parameters, run them through a function, and then repackage the result for us.
- ▶ The functions have no control: the applicative makes all the decisions.
- Monads let the functions themselves decide what should happen.

A Calculator, with Monads

- Okay, the above code works, but here's a better way.
- First define functions lift to convert a function to monadic form for us!

These are part of Control. Monad:

Continued

Lifting

```
iminc = liftM inc
imadd = liftM2 add
imsub = liftM2 sub
implies a b = a >>= (\aa ->
b >>= (\bb ->
if bb == 0 then fail "/0"
else return (aa `div` bb)))
```

- ▶ fail is another useful monadic function, defined in the MonadFail typeclass.
- ► Here it's defined as Nothing.

The Maybe Monad

▶ Here is the complete monad definition for Maybe.

Maybe Monad

```
1 instance Monad Maybe where
2  return = Just
3
4  (>>=) Nothing f = Nothing
5  (>>=) (Just a) f = f a
6
7  fail s = Nothing
```

Example with Maybe

```
Prelude > minc (Just 10)
Just 11
Prelude > madd (minc (Just 10)) (Just 20)
Just 31
Prelude > mdiv (minc (Just 10)) (minc (Just 2))
Just 3
Prelude > minc (mdiv (minc (Just 10)) (minc (Just 2)))
Just 4
Prelude > minc (mdiv (minc (Just 10)) (Just 0))
Nothing
```

The List Monad

Lists can be monads too. The trick is deciding what bind should do.

List Monad

```
instance Monad [] where
return a = [a]

(>>=) [] f = []
(>>=) xs f = concatMap f xs

fail s = []
```

▶ Note that we do not have to change *anything* in our lifted calculator example!

Example with List

```
Prelude> minc [1,2,3]
[2,3,4]

Prelude> madd [1,2,3] [10,200]
[11,201,12,202,13,203]

Prelude> minc (mdiv [10] [0])
[]

Prelude> minc (mdiv [10] [0,2,5])
[5,2]
```