

Dependent Types in GHC

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References

- [Pointer to this talk on github](#)
- [Pointer to references](#)

What are Dependent Types?

```
data  $\mathbb{N}$  : Set where
```

```
  zero :  $\mathbb{N}$ 
```

```
  suc  :  $\mathbb{N} \rightarrow \mathbb{N}$ 
```

```
data Vec (A : Set) :  $\mathbb{N} \rightarrow$  Set where
```

```
  [] : Vec A zero
```

```
  ::_ : {n :  $\mathbb{N}$ }  $\rightarrow$  A  $\rightarrow$  Vec A n  $\rightarrow$  Vec A (suc n)
```

Vector Append

$_+__ : \mathbb{N} \rightarrow \mathbb{N} \rightarrow \mathbb{N}$

$\text{zero} + n = n$

$\text{suc } m + n = \text{suc } (m + n)$

$_++_ : \{A : \text{Set}\} \rightarrow \{m\ n : \mathbb{N}\} \rightarrow$
 $\text{Vec } A\ m \rightarrow \text{Vec } A\ n \rightarrow \text{Vec } A\ (m + n)$

$[] ++ y = y$

$(x :: xs) ++ y = x :: (xs ++ y)$

Vector Lookup

```
data Fin :  $\mathbb{N}$   $\rightarrow$  Set where  
  zero : {n :  $\mathbb{N}$ }  $\rightarrow$  Fin (suc n)  
  suc  : {n :  $\mathbb{N}$ }  $\rightarrow$  (m : Fin n)  $\rightarrow$  Fin (suc n)
```

```
lookup : {A : Set}  $\rightarrow$  {n :  $\mathbb{N}$ }  $\rightarrow$  Fin n  $\rightarrow$  Vec A n  $\rightarrow$  A  
lookup zero (x :: _) = x  
lookup (suc n) (_ :: xs) = lookup n xs
```

Some Basic Types

```
data Bool : Set where  
  true  : Bool  
  false : Bool
```

```
data _≡_ {A : Set} : A → A → Set where  
  refl : {a : A} → a ≡ a
```

Vector Lookup 2

```
_j_ :  $\mathbb{N} \rightarrow \mathbb{N} \rightarrow \text{Bool}$   
 $m \text{ j zero} = \text{false}$   
 $\text{zero j suc } n = \text{true}$   
 $\text{suc } m \text{ j suc } n = m \text{ j } n$ 
```

```
lookup' :  $\{A : \text{Set}\} \rightarrow \{n : \mathbb{N}\} \rightarrow$   
           $(m : \mathbb{N}) \rightarrow m \text{ j } n \equiv \text{true} \rightarrow \text{Vec } A \ n \rightarrow A$   
lookup' _ () []  
lookup' zero refl (x :: _ ) = x  
lookup' (suc m) p ( _ :: xs ) = lookup' m p xs
```

Dependent Types

Advantages over other types

Applications

- Better correctness guarantees for software.
- Mechanical verification of mathematics.

Dependent Types are Not New

Timeline

The Golden Age is Now

- DeepSpec, etc.
- BigProof, etc.

Robert Harper Quote

Dependent Types in Haskell

Richard Eisenberg's PhD Thesis

- 1 Introduction
- 2 Preliminaries
- 3 Motivation
- 4 Dependent Haskell
- 5 PICO: The Intermediate Language
- 6 Type Inference and Elaboration, or how to BAKE a PICO
- 7 Implementation
- 8 Related and Future Work

