MODULE *XJupiterImplCJupiter*

In this module, we show that *XJupiter* implements *CJupiter*. To this end, we first extends *XJupiter* by replacing its *Cop* with that used in *CJupiter*.

EXTENDS *XJupiterExtended*

VARIABLES
- $cincomingCJ$, *cincoming* for *CJupiter* which contains original operations instead of transformed ones in *XJupiter*
- $op2ss$, a function from an operation (represented by its *Oid*) to the part of $2D$ state space produced while the operation is transformed
- $c2ssX$ $c2ssX[c]$: redundant (*eXtra*) $2D$ state space maintained for client $c \in Client$

$$varsImpl \triangleq \langle vars, cincomingCJ, op2ss, c2ssX \rangle$$

$$\begin{aligned}
&TypeOKImpl \triangleq \\
&\quad \land TypeOK \\
&\quad \land cincomingCJ \in [Client \to Seq(Cop)] \\
&\quad \land \forall\, oid \in \text{DOMAIN } op2ss : oid \in Oid \land IsSS(op2ss[oid]) \\
&\quad \land \forall\, c \in Client : IsSS(c2ssX[c])
\end{aligned}$$

The *Init* predicate.

$$\begin{aligned}
&InitImpl \triangleq \\
&\quad \land Init \\
&\quad \land cincomingCJ = [c \in Client \mapsto \langle\rangle] \\
&\quad \land op2ss = \langle\rangle \\
&\quad \land c2ssX = [c \in Client \mapsto [node \mapsto \{\{\}\}, edge \mapsto \{\}]]
\end{aligned}$$

Client $c \in Client$ generates an operation and performs it locally.

$$\begin{aligned}
&DoImpl(c) \triangleq \\
&\quad \land Do(c) \\
&\quad \land \text{UNCHANGED } \langle cincomingCJ, op2ss, c2ssX \rangle
\end{aligned}$$

$$\begin{aligned}
&ss1 \oplus ss2 \triangleq \\
&\quad [ss1 \text{ EXCEPT } !.node = @ \cup ss2.node, \\
&\qquad\qquad !.edge = @ \cup ss2.edge]
\end{aligned}$$

Client $c \in Client$ receives a message and processes it.

$$\begin{aligned}
&RevImpl(c) \triangleq \\
&\quad \land \quad Rev(c) \\
&\quad \land \quad cincomingCJ[c] \neq \langle\rangle \quad \text{there are (original) operations to handle with} \\
&\quad \land \quad cincomingCJ' = [cincomingCJ \text{ EXCEPT } ![c] = Tail(@)] \quad \text{also consume a message} \\
&\quad \land \quad \text{LET } cop \triangleq Head(cincoming[c]) \\
&\qquad\quad \text{IN} \quad c2ssX' = [c2ssX \text{ EXCEPT } ![c] = @ \oplus op2ss[cop.oid]] \\
&\quad \land \quad \text{UNCHANGED } \langle op2ss \rangle
\end{aligned}$$

Also broadcast the original operation to clients (using the *cincomingCJ* channels)

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SRevImpl ≜
    ∧ SRev
    ∧ LET cop ≜ [Head(sincoming) EXCEPT !.sctx = soids]
          c ≜ cop.oid.c
          ss ≜ xForm(cop, s2ss[c], scur[c], Remote)
      IN  ∧ cincomingCJ' = [cl ∈ Client ↦
                              IF cl = c
                              THEN cincomingCJ[cl]
                              ELSE Append(cincomingCJ[cl], cop)]
          ∧ op2ss' = op2ss @@ (cop.oid :> [node ↦ Range(ss.node), edge ↦ Range(ss.edge)])
    ∧ UNCHANGED ⟨c2ssX⟩

The next-state relation.
NextImpl ≜
    ∨ ∃ c ∈ Client : DoImpl(c) ∨ RevImpl(c)
    ∨ SRevImpl

The specification.
SpecImpl ≜ InitImpl ∧ □[NextImpl]_varsImpl ∧ WF_varsImpl(SRevImpl ∨ ∃ c ∈ Client : RevImpl(c))

Ignore the lr field in edges of 2D state space.
IgnoreDir(ss) ≜
    [ss EXCEPT !.edge =
        {[field ∈ (DOMAIN e \ {"lr"}) ↦ e.field] : e ∈ @}]
      {[from ↦ e.from, to ↦ e.to, cop ↦ e.cop] : e ∈ @}]

CJ ≜ INSTANCE CJupiter
        WITH cincoming ← cincomingCJ,
             css ← [r ∈ Replica ↦
                      IF r = Server
                      THEN IgnoreDir(SetReduce(⊕, Range(s2ss), [node ↦ {{}}, edge ↦ {}]))
                      ELSE IgnoreDir(c2ss[r] ⊕ c2ssX[r])],
             cur ← [r ∈ Replica ↦
                      IF r = Server
                      It SHOULD be that Cardinality(Range(scur)) = 1
                      THEN CHOOSE n ∈ Range(scur) : TRUE
                      ELSE ccur[r]]

THEOREM SpecImpl ⇒ CJ!Spec

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