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MODULE CJupiter
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Model of our own *CJupiter* protocol.

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EXTENDS JupiterSerial, GraphsUtil

VARIABLES
    css      css[r]: the n-ary ordered state space at replica r ∈ Replica

vars ≜ ⟨intVars, ctxVars, serialVars, css⟩
```

A *css* is a directed graph (defined in moudule *GraphsUtil*) with labeled edges, Each node is characterized by its context, a set of oids. Each edge is labeled with an operation.

```
IsCSS(G) ≜
    ∧ G = [node ↦ G.node, edge ↦ G.edge]
    ∧ G.node ⊆ (SUBSET Oid)
    ∧ G.edge ⊆ [from : G.node, to : G.node, cop : Cop]

TypeOK ≜
    ∧   TypeOKInt
    ∧   TypeOKCtx
    ∧   TypeOKSerial
    ∧   Comm(Cop)!TypeOK
    ∧   ∀ r ∈ Replica : IsCSS(css[r])

Init ≜
    ∧ InitInt
    ∧ InitCtx
    ∧ InitSerial
    ∧ Comm(Cop)!Init
    ∧ css = [r ∈ Replica ↦ EmptyGraph]
```

Locate the node in *rcss* (the *css* at replica $r \in Replica$) that matches the context *ctx* of cop.

```
Locate(cop, rcss) ≜ CHOOSE n ∈ rcss.node : n = cop.ctx
```

xForm: Iteratively transform cop with a path through the *css* at replica $r \in Replica$, following the first edges.

```
xForm(cop, r) ≜
    LET rcss ≜ css[r]
        u ≜ Locate(cop, rcss)
        v ≜ u ∪ {cop.oid}
        RECURSIVE xFormHelper(_, _, _, _, _)
         'h' stands for "helper"; xcss: eXtra css created during transformation
        xFormHelper(uh, vh, coph, xcss, xcoph) ≜
            IF uh = ds[r]
             THEN ⟨xcss, xcoph⟩
             ELSE LET fedge ≜ CHOOSE e ∈ rcss.edge :
                                  ∧ e.from = uh
```

```
                                        ∧ ∀ uhe ∈ rcss.edge :
                                            (uhe.from = uh ∧ uhe ≠ e) ⇒ tb(e.cop.oid, uhe.cop.oid, serial[r])
                              uprime ≜ fedge.to
                              fcop ≜ fedge.cop
                              coph2fcop ≜ COT(coph, fcop)
                              fcop2coph ≜ COT(fcop, coph)
                               vprime ≜ vh ∪ {fcop.oid}
                         IN    xFormHelper(uprime, vprime, coph2fcop,
                                  [xcss EXCEPT !.node = @ ∪ {vprime},
                                    !.edge = @ ∪ {[from ↦ vh, to ↦ vprime, cop ↦ fcop2coph],
                                                  [from ↦ uprime, to ↦ vprime, cop ↦ coph2fcop]}],
                                    coph2fcop)
    IN    xFormHelper(u, v, cop, [node ↦ {v}, edge ↦ {[from ↦ u, to ↦ v, cop ↦ cop]}], cop)
```

Perform cop at replica $r \in Replica$.

```
Perform(cop, r) ≜
    LET xform ≜ xForm(cop, r)   xform: ⟨xcss, xcop⟩
        xcss ≜ xform[1]
        xcop ≜ xform[2]
    IN  ∧ css' = [css EXCEPT ![r] = @ ⊕ xcss]
        ∧ state' = [state EXCEPT ![r] = Apply(xcop.op, @)]
```

Client $c \in Client$ issues an operation op.

```
DoOp(c, op) ≜   op: the raw operation generated by the client c ∈ Client
    ∧ LET cop ≜ [op ↦ op, oid ↦ [c ↦ c, seq ↦ cseq'[c]], ctx ↦ ds[c]]
      IN  ∧ Perform(cop, c)
          ∧ UpdateDS(c, cop)
          ∧ Comm(Cop)!CSend(cop)

DoIns(c) ≜
    ∃ ins ∈ {op ∈ Ins : op.pos ∈ 1 .. (Len(state[c]) + 1) ∧ op.ch ∈ chins ∧ op.pr = Priority[c]} :
        ∧ DoOp(c, ins)
        ∧ chins' = chins \ {ins.ch}   We assume that all inserted elements are unique.

DoDel(c) ≜
    ∃ del ∈ {op ∈ Del : op.pos ∈ 1 .. Len(state[c])} :
        ∧ DoOp(c, del)
        ∧ UNCHANGED chins

Do(c) ≜
    ∧ DoCtx(c)
    ∧ DoSerial(c)
    ∧ ∨ DoIns(c)
      ∨ DoDel(c)
```

Client $c \in Client$ receives a message from the $Server$.

```
105 Rev(c) ≜
106       ∧ Comm(Cop)!CRev(c)
107       ∧ Perform(Head(cincoming[c]), c)
108       ∧ RevSerial(c)
109       ∧ RevCtx(c)
110       ∧ UNCHANGED chins
111 ----------------------------------------------------------------------
    The Server receives a message.
115 SRev ≜
116     ∧ Comm(Cop)!SRev
117     ∧ LET cop ≜ Head(sincoming)
118       IN    ∧ Perform(cop, Server)
119             ∧ Comm(Cop)!SSendSame(cop.oid.c, cop)   broadcast the original operation
120     ∧ SRevSerial
121     ∧ SRevCtx
122     ∧ UNCHANGED chins
123 ----------------------------------------------------------------------
124 Next ≜
125     ∨ ∃ c ∈ Client : Do(c) ∨ Rev(c)
126     ∨ SRev
    Fairness: There is no requirement that the clients ever generate operations.
130 Fairness ≜
131     WF_vars(SRev ∨ ∃ c ∈ Client : Rev(c))

133 Spec ≜ Init ∧ □[Next]_vars   ∧ Fairness (We care more about safety.)
134 ----------------------------------------------------------------------
    The compactness of CJupiter: the CSSes at all replicas are the same.
138 Compactness ≜
139     Comm(Cop)!EmptyChannel ⇒ Cardinality(Range(css)) = 1

141 THEOREM Spec ⇒ Compactness
142 ======================================================================
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