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1  ┌────────────────── MODULE XJupiter ───────────────────┐
    Specification of the Jupiter protocol described in CSCW'2014 by Yi Xu, Chengzheng Sun, and
    Mo Li. We call it XJupiter, with 'X' for "Xu".
7  EXTENDS StateSpace
8  └──────────────────┐
9  VARIABLES
10     c2ss,      c2ss[c]: the 2D state space (2ss, for short) at client c ∈ Client
11     s2ss       s2ss[c]: the 2D state space maintained by the Server for client c ∈ Client
13     vars  $\triangleq$   $\langle \text{intVars}, \text{ctxVars}, \text{c2ss}, \text{s2ss} \rangle$ 
14 └──────────────────┐
15 TypeOK  $\triangleq$ 
16      $\wedge$  TypeOKInt
17      $\wedge$  TypeOKCtx
18      $\wedge \forall c \in \text{Client} : \text{IsSS}(\text{c2ss}[c]) \wedge \text{IsSS}(\text{s2ss}[c])$ 
19 └──────────────────┐
20 Init  $\triangleq$ 
21      $\wedge$  InitInt
22      $\wedge$  InitCtx
23      $\wedge \text{c2ss} = [c \in \text{Client} \mapsto \text{EmptySS}]$ 
24      $\wedge \text{s2ss} = [c \in \text{Client} \mapsto \text{EmptySS}]$ 
25 └──────────────────┐
26 xForm(cop, ss, cur)  $\triangleq$  Transform cop with an operation sequence in 2D state space ss.
27     LET u  $\triangleq$  Locate(cop, ss)
28     v  $\triangleq$  u  $\cup$  {cop.oid}
29     RECURSIVE xFormHelper(u, v, coph, xss)  $\triangleq$  xss: eXtra ss created during transformation
30     xFormHelper(uh, vh, coph, xss)  $\triangleq$ 
31         IF uh = cur THEN [xss  $\mapsto$  xss, xcop  $\mapsto$  coph]
32         ELSE LET e  $\triangleq$  CHOOSE e ∈ ss.edge : e.from = uh  $\wedge$  ClientOf(e.cop)  $\neq$  ClientOf(cop)
33             copprime  $\triangleq$  e.cop
34             uprime  $\triangleq$  e.to
35             vprime  $\triangleq$  vh  $\cup$  {copprime.oid}
36             coph2copprime  $\triangleq$  COT(coph, copprime)
37             copprime2coph  $\triangleq$  COT(copprime, coph)
38             IN xFormHelper(uprime, vprime, coph2copprime,
39                 xss  $\oplus$  [node  $\mapsto$  {vprime},
40                     edge  $\mapsto$  {[from  $\mapsto$  vh, to  $\mapsto$  vprime, cop  $\mapsto$  copprime2coph],
41                         [from  $\mapsto$  uprime, to  $\mapsto$  vprime, cop  $\mapsto$  coph2copprime]}])
42             IN xFormHelper(u, v, cop, [node  $\mapsto$  {v}, edge  $\mapsto$  {[from  $\mapsto$  u, to  $\mapsto$  v, cop  $\mapsto$  cop]}])
44 ClientPerform(c, cop)  $\triangleq$ 
45     LET xform  $\triangleq$  xForm(cop, c2ss[c], ds[c]) xform: [xss, xcop]
46     IN  $\wedge \text{c2ss}' = [\text{c2ss} \text{ EXCEPT } ![c] = @ \oplus \text{xform.xss}]$ 
47      $\wedge \text{SetNewAop}(c, \text{xform.xcop.op)$ 
49 ServerPerform(cop)  $\triangleq$ 

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50   LET  $c \triangleq ClientOf(cop)$ 
51    $scur \triangleq ds[Server]$ 
52    $xform \triangleq xForm(cop, s2ss[c], scur)$   $xform: [xss, xcop]$ 
53    $xcop \triangleq xform.xcop$ 
54    $xcur \triangleq scur \cup \{cop.oid\}$ 
55   IN  $\wedge s2ss' = [cl \in Client \mapsto$ 
56       IF  $cl = c$ 
57       THEN  $s2ss[cl] \oplus xform.xss$ 
58       ELSE  $s2ss[cl] \oplus [node \mapsto \{xcur\},$ 
59            $edge \mapsto \{[from \mapsto scur, to \mapsto xcur, cop \mapsto xcop]\}]$ 
60        $\wedge SetNewAop(Server, xcop.op)$ 
61        $\wedge Comm!SSendSame(c, xcop)$ 
62   |-----|
63    $DoOp(c, op) \triangleq$ 
64       LET  $cop \triangleq [op \mapsto op, oid \mapsto [c \mapsto c, seq \mapsto cseq[c]], ctx \mapsto ds[c]]$ 
65       IN  $\wedge ClientPerform(c, cop)$ 
66        $\wedge Comm!CSend(cop)$ 
67
68    $Do(c) \triangleq$ 
69        $\wedge DoInt(DoOp, c)$ 
70        $\wedge DoCtx(c)$ 
71        $\wedge UNCHANGED\ s2ss$ 
72
73    $Rev(c) \triangleq$ 
74        $\wedge RevInt(ClientPerform, c)$ 
75        $\wedge RevCtx(c)$ 
76        $\wedge UNCHANGED\ s2ss$ 
77
78    $SRev \triangleq$ 
79        $\wedge SRevInt(ServerPerform)$ 
80        $\wedge SRevCtx$ 
81        $\wedge UNCHANGED\ c2ss$ 
82   |-----|
83    $Next \triangleq$ 
84        $\vee \exists c \in Client : Do(c) \vee Rev(c)$ 
85        $\vee SRev$ 
86
87    $Fairness \triangleq$  There is no requirement that the clients ever generate operations.
88    $WF_{vars}(SRev \vee \exists c \in Client : Rev(c))$ 
89
90    $Spec \triangleq Init \wedge \Box[Next]_{vars} \wedge Fairness$ 
91   |-----|
92    $CSSync \triangleq$  Each client  $c \in Client$  is synchronized with the Server.
93    $\forall c \in Client : (ds[c] = ds[Server]) \Rightarrow c2ss[c] = s2ss[c]$ 
94
95   THEOREM  $Spec \Rightarrow \Box CSSync$ 
96   |-----|

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\\* Modification History  
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