```
- Module XJupiter —
 1 [
    Specification of the Jupiter protocol described in CSCW'2014 by Yi Xu, Chengzheng Sun, and
    Mo Li. We call it XJupiter, with 'X' for "Xu".
    EXTENDS Integers, OT, TLC, AdditionalFunctionOperators, AdditionalSequenceOperators
 8 1
    CONSTANTS
         Client,
                         the set of client replicas
10
         Server,
11
                         the (unique) server replica
         Char.
12
                         set of characters allowed
         InitState
                         the initial state of each replica
13
     Replica \triangleq Client \cup \{Server\}
     List \stackrel{\triangle}{=} Seq(Char \cup Range(InitState)) all possible lists/strings
     MaxLen \stackrel{\Delta}{=} Cardinality(Char) + Len(InitState) the max length of lists in any states;
          We assume that all inserted elements are unique.
19
     ClientNum \triangleq Cardinality(Client)
21
     Priority \triangleq CHOOSE f \in [Client \rightarrow 1 .. ClientNum] : Injective(f)
     direction flags
24
     Local \triangleq 0
     Remote \stackrel{\Delta}{=} 1
26
          \land Range(InitState) \cap Char = \{\} due to the uniqueness requirement
29
30
          \land Priority \in [Client \rightarrow 1 .. ClientNum]
31
    The set of all operations. Note: The positions are indexed from 1.
    Rd \triangleq [type : \{ \text{``Rd''} \}]
    Del \stackrel{\Delta}{=} [type : \{ "Del" \}, pos : 1 .. MaxLen]
    Ins \stackrel{\triangle}{=} [type: \{ \text{"Ins"} \}, pos: 1... (MaxLen + 1), ch: Char, pr: 1... ClientNum] pr: priority
    Op \stackrel{\triangle}{=} Ins \cup Del
40
41 ⊦
    Cop: operation of type Op with context
   Oid \stackrel{\triangle}{=} [c:Client, seq:Nat] operation identifier
    Cop \triangleq [op : Op \cup \{Nop\}, oid : Oid, ctx : SUBSET Oid]
    OT of two operations of type Cop.
    COT(lcop, rcop) \stackrel{\Delta}{=} [lcop \ EXCEPT \ !.op = Xform(lcop.op, rcop.op), \ !.ctx = @ \cup \{rcop.oid\}]
52 L
   VARIABLES
        For the client replicas:
                    cseq[c]: local sequence number at client c \in Client
57
```

```
The 2D state spaces (ss, for short). Each client maintains one 2D state space. The server
         maintains n 2D state spaces, one for each client.
 63
          css.
                      css[c]: the 2D state space at client c \in Client
                      cur[c]: the current node of css[c]
 64
           ccur,
          sss,
                      sss[c]: the 2D state space maintained by the Server for client c \in Client
 65
                      scur[c]: the current node of sss[c]
 66
           scur,
        For all replicas
                      state[r]: state (the list content) of replica r \in Replica
 70
           state,
         For communication between the Server and the Clients:
           cincoming,
                             cincoming[c]: incoming channel at the client c \in Client
 74
           sincoming,
                             incoming channel at the Server
 75
         For model checking:
                     a set of chars to insert
           chins
 79
 80
      comm \stackrel{\triangle}{=} INSTANCE \ CSComm \ WITH \ Msg \leftarrow Cop
 81
 82
      eVars \triangleq
                   \langle chins \rangle
                               variables for the environment
      cVars \triangleq \langle cseq \rangle
                               variables for the clients
      cssVars \stackrel{\triangle}{=} \langle css, ccur \rangle
                                        variables for 2D state spaces at clients
      sssVars \stackrel{\triangle}{=} \langle sss, scur \rangle
                                         variables for 2D state spaces at the Server
      commVars \triangleq \langle cincoming, sincoming \rangle
                                                             variables for communication
      vars \triangleq \langle eVars, cVars, commVars, cssVars, sssVars, state \rangle all variables
 89
      A 2D state space is a directed graph with labeled edges. It is represented by a record with node
     field and edge field. Each node is characterized by its context, a set of operations. Each edge is
     labeled with an operation and a direction flag indicating whether this edge is LOCAL or REMOTE.
     For clarity, we denote edges by records instead of tuples.
     IsSS(G) \triangleq
 98
             \land G = [node \mapsto G.node, edge \mapsto G.edge]
 99
            \land G.node \subseteq (SUBSET\ Oid)
100
            \land G.edge \subseteq [from: G.node, to: G.node, cop: Cop, lr: \{Local, Remote\}]
101
      TypeOK \stackrel{\triangle}{=}
103
          For the client replicas:
           \land cseq \in [Client \rightarrow Nat]
107
          For the 2D state spaces:
           \land \forall c \in Client : IsSS(css[c]) \land IsSS(sss[c])
111
           \land ccur \in [Client \rightarrow SUBSET \ Oid]
112
           \land scur \in [Client \rightarrow SUBSET \ Oid]
113
           \land state \in [Replica \rightarrow List]
114
          For communication between the server and the clients:
           \land comm! TypeOK
118
          For model checking:
```

```
122
            \land chins \subseteq Char
123 ⊢
      The Init predicate.
     Init \triangleq
127
          For the client replicas:
            \land cseq = [c \in Client \mapsto 0]
131
          For the 2D state spaces:
            \land css = [c \in Client \mapsto [node \mapsto \{\{\}\}, edge \mapsto \{\}]]
135
136
            \land ccur = [c \in Client \mapsto \{\}]
            \land \mathit{sss} = [c \in \mathit{Client} \mapsto [\mathit{node} \mapsto \{\{\}\}\}, \, \mathit{edge} \mapsto \{\}]]
137
            \land scur = [c \in Client \mapsto \{\}]
138
           For all replicas:
            \land state = [r \in Replica \mapsto InitState]
142
           For communication between the server and the clients:
146
            \land comm!Init
          For model checking:
            \wedge chins = Char
150
151 F
      Locate the node in the 2D state space ss which matches the context ctx of cop.
     Locate(cop, ss) \stackrel{\Delta}{=} CHOOSE \ n \in (ss.node) : n = cop.ctx
155
      xForm: iteratively transform cop with a path through the 2D state space ss at some client,
      following the edges with the direction flag d.
      xForm(cop, ss, cur, d) \stackrel{\Delta}{=}
162
           LET u \triangleq Locate(cop, ss)
163
                 v \triangleq u \cup \{cop.oid\}
164
                 RECURSIVE xFormHelper(_, _, _, _)
165
                   'h' stands for "helper"; xss: eXtra \ ss created during transformation
166
                 xFormHelper(uh, vh, coph, xss) \stackrel{\Delta}{=}
167
                      If uh = cur
168
                        THEN xss
169
                        ELSE LET e \stackrel{\Delta}{=} \text{CHOOSE } e \in ss.edge : e.from = uh \land e.lr = d
170
                                       uprime \triangleq e.to
171
                                       copprime \triangleq e.cop
172
                                       coph2copprime \stackrel{\Delta}{=} COT(coph, copprime)
copprime2coph \stackrel{\Delta}{=} COT(copprime, coph)
173
174
                                        vprime \triangleq vh \cup \{copprime.oid\}
175
                                        xFormHelper(uprime, vprime, coph2copprime,
176
                                            [xss except !.node = @ \circ \langle vprime \rangle,
177
                                                                the order of recording edges here is important
178
                                                                so that the last one is labeled with the final transformed operation
179
                                                              !.edge = @ \circ \langle [from \mapsto vh, to \mapsto vprime, cop \mapsto copprime2coph, lr \mapsto vprime]
180
                                                                                  [from \mapsto uprime, to \mapsto vprime, cop \mapsto coph2copprime,
181
```

```
xFormHelper(u, v, cop, [node \mapsto \langle v \rangle,
182
                                                    edge \mapsto \langle [from \mapsto u, to \mapsto v, cop \mapsto cop, lr \mapsto 1 - d] \rangle])
183
184
      Client c \in Client perform operation cop guided by the direction flag d.
      ClientPerform(cop, c, d) \stackrel{\triangle}{=}
188
           LET xss \stackrel{\Delta}{=} xForm(cop, css[c], ccur[c], d)
189
                 xn \triangleq xss.node
190
                  xe \stackrel{\triangle}{=} xss.edge
191
                  xcur \triangleq Last(xn)
192
                 xcop \triangleq Last(xe).cop
193
                  \wedge css' = [css \ EXCEPT \ ![c].node = @ \cup Range(xn),
194
                                                 ![c].edge = @ \cup Range(xe)]
195
                  \wedge ccur' = [ccur \ EXCEPT \ ![c] = xcur]
196
                  \land state' = [state \ EXCEPT \ ![c] = Apply(xcop.op, @)]
197
      Client c \in Client issues an operation op.
      DoOp(c, op) \stackrel{\Delta}{=} op: the raw operation generated by the client c \in Client
201
               \land cseq' = [cseq \ \texttt{EXCEPT} \ ![c] = @+1]
202
               \wedge LET cop \stackrel{\triangle}{=} [op \mapsto op, oid \mapsto [c \mapsto c, seq \mapsto cseq'[c]], ctx \mapsto ccur[c]]
203
                          \land ClientPerform(cop, c, Remote)
204
                          \land comm! CSend(cop)
205
      DoIns(c) \stackrel{\Delta}{=}
207
           \exists ins \in Ins :
208
               \land ins.pos \in 1 \dots (Len(state[c]) + 1)
209
               \land ins.ch \in chins
210
               \wedge ins.pr = Priority[c]
211
               \wedge chins' = chins \setminus \{ins.ch\} We assume that all inserted elements are unique.
212
213
               \wedge DoOp(c, ins)
               \land UNCHANGED \langle sssVars \rangle
214
      DoDel(c) \triangleq
216
           \exists del \in Del:
217
               \land del.pos \in 1 \dots Len(state[c])
218
               \wedge DoOp(c, del)
219
               \land Unchanged \langle sssVars, eVars \rangle
220
      Do(c) \triangleq
222
              \vee DoIns(c)
223
             \vee DoDel(c)
224
      Client c \in Client receives a message from the Server.
      Rev(c) \triangleq
228
             \land comm! CRev(c)
229
             \wedge LET cop \stackrel{\triangle}{=} Head(cincoming[c]) the received (transformed) operation
230
                 IN ClientPerform(cop, c, Local)
231
             \land UNCHANGED \langle eVars, cVars, sssVars \rangle
232
```

```
233 |
             The Server performs operation cop.
             ServerPerform(cop) \triangleq
237
                         Let c \triangleq cop.oid.c
238
                                             \stackrel{\Delta}{=} xForm(cop, sss[c], scur[c], Remote)
239
                                              \stackrel{\Delta}{=} \textit{ xss.node}
                                 xn
240
                                              \triangleq xss.edge
241
                            xcur \stackrel{\triangle}{=} Last(xn)
242
                           xcop \triangleq Last(xe).cop
243
                                       \wedge sss' = [cl \in Client \mapsto
244
                                                                         If cl = c
245
                                                                            THEN [sss[cl]] EXCEPT !.node = @ \cup Range(xn),
246
                                                                                                                                                !.edge = @ \cup Range(xe)]
247
                                                                            ELSE LET scurcl \triangleq scur[cl]
248
                                                                                                             scurclprime \triangleq scurcl \cup \{cop.oid\}
249
                                                                                                            [sss[cl] \ EXCEPT \ !.node = @ \cup \{scurclprime\},
250
                                                                                                                                                              !.edge = @ \cup \{[from \mapsto scurcl, to \mapsto scurclprime, \}\}
251
                                                                                                                                                                                                               cop \mapsto xcop, lr \mapsto Remote]\}]
252
                                        \begin{tabular}{l} \begin{tabu
253
254
                                                                          IF cl = c THEN xcur ELSE scur[cl] \cup \{cop.oid\}]
255
                                        \land state' = [state \ EXCEPT \ ! [Server] = Apply(xcop.op, @)]
256
                                        \land comm! SSendSame(c, xcop) broadcast the transformed operation
257
             The Server receives a message.
             SRev \triangleq
261
                          \land \ comm \, ! \, SRev
262
                          \wedge \text{ LET } cop \stackrel{\triangle}{=} Head(sincoming)
263
                                   IN ServerPerform(cop)
264
                          \land UNCHANGED \langle eVars, cVars, cssVars \rangle
265
266
             The next-state relation.
             Next \triangleq
270
                          \vee \exists c \in Client : Do(c) \vee Rev(c)
271
                          \vee SRev
272
             The Spec. (TODO: Check the fairness condition.)
             Spec \stackrel{\triangle}{=} Init \wedge \Box [Next]_{vars} \wedge WF_{vars}(Next)
276
277 L
              \* Modification History
              \* Last modified Wed Oct 10 15:00:09 CST 2018 by hengxin
              \^* Created Tue Oct 09 16:33:18 CST 2018 by hengxin
```