MODULE *CJupiter*

Model of our own *CJupiter* protocol.

EXTENDS *Integers*, *OT*, *TLC*, *AdditionalFunctionOperators*

CONSTANTS
- *Client*, the set of client replicas
- *Server*, the (unique) server replica
- *Char*, set of characters allowed
- *InitState* the initial state of each replica

$Replica \triangleq Client \cup \{Server\}$

$List \triangleq Seq(Char \cup Range(InitState))$ all possible lists/strings

$MaxLen \triangleq Cardinality(Char) + Len(InitState)$ the max length of lists in any states; We assume that all inserted elements are unique.

$ClientNum \triangleq Cardinality(Client)$

$Priority \triangleq$ CHOOSE $f \in [Client \to 1 \,..\, ClientNum] : Injective(f)$

ASSUME
- $\land Range(InitState) \cap Char = \{\}$
- $\land Priority \in [Client \to 1 \,..\, ClientNum]$

The set of all operations. Note: The positions are indexed from 1.

$Rd \triangleq [type : \{\text{“Rd”}\}]$

$Del \triangleq [type : \{\text{“Del”}\}, pos : 1 \,..\, MaxLen]$

$Ins \triangleq [type : \{\text{“Ins”}\}, pos : 1 \,..\, (MaxLen + 1), ch : Char, pr : 1 \,..\, ClientNum]$ *pr*: priority

$Op \triangleq Ins \cup Del$ Now we don't consider *Rd* operations.

$Oid \triangleq [c : Client, seq : Nat]$ operation identifier

$Cop \triangleq [op : Op, oid : Oid, ctx :$ SUBSET $Oid, sctx :$ SUBSET $Oid]$ operation with context

VARIABLES

For the client replicas:
- *cseq*, *cseq*[*c*]: local sequence number at client $c \in Client$
- *cstate*, *cstate*[*c*]: state (the list content) of the client $c \in Client$

For the server replica:
- *sstate*, *sstate*: state (the list content) of the server *Server*

For all replicas: the n-ary ordered state space
- *css*, *css*[*r*]: the n-ary ordered state space at replica r
- *cur*, *cur*[*r*]: the current node of *css* at replica r

For communication between the *Server* and the Clients:

62 $cincoming$, *cincoming*[c]: incoming channel at the client $c \in Client$
63 $sincoming$, incoming channel at the *Server*

For model checking:

67 $chins$ a set of chars to insert

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70 $comm \triangleq$ INSTANCE $CSComm$

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72 $eVars \triangleq \langle chins \rangle$ variables for the environment
73 $cVars \triangleq \langle cseq, cstate \rangle$ variables for the clients
74 $ecVars \triangleq \langle eVars, cVars \rangle$ variables for the clients and the environment
75 $sVars \triangleq \langle sstate \rangle$ variables for the server
76 $commVars \triangleq \langle cincoming, sincoming \rangle$ variables for communication
77 $vars \triangleq \langle eVars, cVars, sVars, commVars, css, cur \rangle$ all variables

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An *css* is a directed graph with labeled edges.

It is represented by a record with node field and edge field.

Each node is characterized by its context, a set of operations.

Each edge is labeled with an operation. For clarity, we denote edges by records instead of tuples.

89 $IsCSS(G) \triangleq$
90 $\land\ G = [node \mapsto G.node,\ edge \mapsto G.edge]$
91 $\land\ G.node \subseteq (\text{SUBSET } Oid)$
92 $\land\ G.edge \subseteq [from : G.node,\ to : G.node,\ cop : Cop]$

94 $TypeOK \triangleq$

For the client replicas:

98 $\land\ cseq \in [Client \to Nat]$
99 $\land\ cstate \in [Client \to List]$

For the server replica:

103 $\land\ sstate \in List$

For all replicas: the n-ary ordered state space

107 $\land\ \forall\, r \in Replica : IsCSS(r)$
108 $\land\ cur \in [Client \to \text{SUBSET } Oid]$

For communication between the server and the clients:

112 $\land\ comm\,!\,TypeOK$

For model checking:

116 $\land\ chins \subseteq Char$

117 ----

The *Init* predicate.

121 $Init \triangleq$
122 $\land\ chins = Char$

For the client replicas:

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126     ∧ cseq = [c ∈ Client ↦ 0]
127     ∧ cstate = [c ∈ Client ↦ InitState]
        For the server replica:
131     ∧ sstate = InitState
        For all replicas: the n-ary ordered state space
135     ∧ css  = [c ∈ Client ↦ [node ↦ {}, edge ↦ {}]]
136     ∧ cur = {}
        For communication between the server and the clients:
140     ∧ comm ! Init
141 ├──────────────────────────────────────────────────────────────────────────┤
    Client c ∈ Client issues an operation op.
145 DoOp(c, op) ≜
146         ∧ cstate′ = [cstate EXCEPT ![c] = Apply(op, @)]
147         ∧ cseq′ = [cseq EXCEPT ![c] = @ + 1]
148         ∧ LET cop ≜ [op ↦ op, oid ↦ [c ↦ c, seq ↦ cseq′[c]],
149               ctx ↦ cur[c], sctx ↦ {}]
150                   v ≜ cur ∪ {cop.oid}
151               IN    ∧ css′ = [css EXCEPT ![c].node = @ ∪ {v},
152                                          ![c].edge = @ ∪ {[from ↦ cur, to ↦ v, cop ↦ cop]}]
153                     ∧ cur′ = v
154                     ∧ comm ! CSend([c ↦ c, op ↦ cop])

156 DoIns(c) ≜
157     ∃ ins ∈ Ins :
158         ∧ ins.pos ∈ 1 .. (Len(cstate[c]) + 1)
159         ∧ ins.ch ∈ chins
160         ∧ ins.pr = Priority[c]
161         ∧ chins′ = chins \ {ins.ch}   We assume that all inserted elements are unique.
162         ∧ DoOp(c, ins)
163         ∧ UNCHANGED sVars

165 DoDel(c) ≜
166     ∃ del ∈ Del :
167         ∧ del.pos ∈ 1 .. Len(cstate[c])
168         ∧ DoOp(c, del)
169         ∧ UNCHANGED ⟨sVars, eVars⟩

171 Do(c) ≜
172       ∨ DoIns(c)
173       ∨ DoDel(c)
    Locate the node in rcss which matches the context ctx of cop.
    rcss: the css at replica r ∈ Replica
179 Locate(cop, rcss) ≜ CHOOSE n ∈ (rcss.node) : n = cop.ctx
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xForm: .

$xForm(cop,\ rcss) \triangleq$ TRUE *TODO*

Client $c \in Client$ receives a message from the *Server*.

$Rev(c) \triangleq$
  $\land\ comm!CRev(c)$
  $\land$ LET $m \triangleq Head(cincoming[c])$
    IN $\land$ TRUE
      $\land\ cstate' = [cstate$ EXCEPT $![c] = Apply(xop, @)]$ * apply the transformed operation *xop*
  $\land$ UNCHANGED $\langle sVars,\ eVars\rangle$

The *Server* receives a message.

$SRev \triangleq$
  $\land\ comm!SRev$
  $\land$ LET $m \triangleq Head(sincoming)$ the message to handle with
    IN $\land$ TRUE
      $\land\ sstate' = Apply(xop,\ sstate)$ * apply the transformed operation
      $\land\ comm!SSend(c,\ srec,\ xop)$
  $\land$ UNCHANGED $ecVars$

The next-state relation.

$Next \triangleq$
  $\lor \exists\, c \in Client : Do(c) \lor Rev(c)$
  $\lor\ SRev$

The *Spec*. (*TODO*: Check the fairness condition.)

$Spec \triangleq Init \land \Box[Next]_{vars} \land \mathrm{WF}_{vars}(Next)$

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