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- MODULE CJupiter
 1 [
    Model of our own CJupiter protocol.
   EXTENDS JupiterSerial
 6 |
 7
    VARIABLES
                    cseq[c]: local sequence number at client c \in Client
         cseq,
         For all replicas: the n-ary ordered state space
                    css[r]: the n-ary ordered state space at replica r \in Replica
12
                    cur[r]: the current node of css at replica r \in Replica
13
         cur
    vars \triangleq \langle chins, cseq, css, cur, state, cincoming, sincoming, serialVars \rangle
15
16 |
    A css is a directed graph with labeled edges, represented by a record with node field and edge field.
    Each node is characterized by its context, a set of oids. Each edge is labeled with an operation.
    IsCSS(G) \triangleq
23
          \land G = [node \mapsto G.node, edge \mapsto G.edge]
24
          \land G.node \subseteq (SUBSET\ Oid)
25
          \land G.edge \subseteq [from : G.node, to : G.node, cop : Cop]
26
     EmptySS \stackrel{\Delta}{=} [node \mapsto \{\{\}\}, edge \mapsto \{\}]
28
     TypeOK \triangleq
30
                TypeOKInt
          Λ
31
                TypeOKSerial
32
                Comm(Cop)! TypeOK
33
                cseq \in [Client \rightarrow Nat]
34
         For all replicas: the n-ary ordered state space
          \land \forall r \in Replica : IsCSS(css[r])
38
          \land cur \in [Replica \rightarrow SUBSET\ Oid]
39
40
    Init \triangleq
41
          \land \mathit{InitInt}
42
          \land \mathit{InitSerial}
43
          \land Comm(Cop)!Init
44
          \land cseq = [c \in Client \mapsto 0]
45
         For all replicas: the n-ary ordered state space
          \land css = [r \in Replica \mapsto EmptySS]
49
          \land cur = [r \in Replica \mapsto \{\}]
50
51
    Locate the node in rcss (the css at replica r \in Replica) that matches the context ctx of cop.
    Locate(cop, rcss) \stackrel{\Delta}{=} CHOOSE \ n \in rcss.node : n = cop.ctx
    Take union of two state spaces ss1 and ss2.
   ss1 \oplus ss2 \triangleq [node \mapsto ss1.node \cup ss2.node, edge \mapsto ss1.edge \cup ss2.edge]
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xForm: Iteratively transform cop with a path through the css at replica r \in Replica, following
      the first edges.
      xForm(cop, r) \triangleq
 64
           Let rcss \stackrel{\triangle}{=} css[r]
 65
                 u \triangleq Locate(cop, rcss)
 66
                 v \ \stackrel{\scriptscriptstyle \Delta}{=} \ u \cup \{cop.oid\}
 67
                 RECURSIVE xFormHelper(\_, \_, \_, \_, \_, \_)
 68
                   'h' stands for "helper"; xcss: eXtra css created during transformation
                 xFormHelper(uh, vh, coph, xcss, xcoph, xcurh) \stackrel{\Delta}{=}
 70
                      IF uh = cur[r]
                       THEN \langle xcss, xcoph, xcurh \rangle
 72
                       ELSE LET fedge \stackrel{\Delta}{=} \text{CHOOSE } e \in rcss.edge :
 73
                                                      \wedge e.from = uh
 74
 75
                                                      \land \forall uhe \in rcss.edge:
                                                          (uhe.from = uh \land uhe \neq e) \Rightarrow tb(e.cop.oid, uhe.cop.oid, serial[r])
 76
                                       uprime \triangleq fedge.to
 77
                                      fcop \triangleq fedge.cop
 78
                                       coph2fcop \stackrel{\Delta}{=} COT(coph, fcop)
 79
                                      fcop2coph \triangleq COT(fcop, coph)
 80
                                        vprime \triangleq vh \cup \{fcop.oid\}
 81
                                       xFormHelper(uprime, vprime, coph2fcop,
                                IN
 82
                                           [xcss \ EXCEPT \ !.node = @ \cup \{vprime\},
 83
                                              !.edge = @ \cup \{[from \mapsto vh, to \mapsto vprime, cop \mapsto fcop2coph],
                                                                  [from \mapsto uprime, to \mapsto vprime, cop \mapsto coph2fcop]\}],
 85
                                            coph2fcop, vprime
 86
                 xFormHelper(u, v, cop, [node \mapsto \{v\}, edge \mapsto \{[from \mapsto u, to \mapsto v, cop \mapsto cop]\}], cop, v)
 87
      Perform cop at replica r \in Replica.
      Perform(cop, r) \triangleq
 91
           LET xform \stackrel{\triangle}{=} xForm(cop, r) xform: \langle xcss, xcop, xcur \rangle
 92
                 xcss \triangleq xform[1]
 93
                 xcop \triangleq xform[2]
 94
                 xcur \triangleq xform[3]
                  \wedge css' = [css \text{ except } ![r] = @ \oplus xcss]
 96
                  \wedge cur' = [cur \ EXCEPT \ ![r] = xcur]
 97
                  \wedge state' = [state \ EXCEPT \ ![r] = Apply(xcop.op, @)]
 98
 99
      Client c \in Client issues an operation op.
     DoOp(c, op) \stackrel{\Delta}{=} op: the raw operation generated by the client c \in Client
103
              \land cseq' = [cseq \ EXCEPT \ ![c] = @ + 1]
104
              \wedge LET cop \stackrel{\Delta}{=} [op \mapsto op, oid \mapsto [c \mapsto c, seq \mapsto cseq'[c]], ctx \mapsto cur[c]]
105
                        \land Perform(cop, c)
106
                        \land Comm(Cop)! CSend(cop)
107
      DoIns(c) \triangleq
109
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110

 $\exists ins \in \{op \in Ins : op.pos \in 1 .. (Len(state[c]) + 1) \land op.ch \in chins \land op.pr = Priority[c]\}:$

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\wedge DoOp(c, ins)
111
              \wedge chins' = chins \setminus \{ins.ch\} We assume that all inserted elements are unique.
112
     DoDel(c) \triangleq
114
          \exists del \in \{op \in Del : op.pos \in 1 .. Len(state[c])\}:
115
              \wedge DoOp(c, del)
116
              \land UNCHANGED chins
117
      Do(c) \triangleq
119
             \land DoSerial(c)
120
            \land \lor DoIns(c)
121
                \vee DoDel(c)
122
     Client c \in Client receives a message from the Server.
     Rev(c) \triangleq
126
            \land Comm(Cop)! CRev(c)
127
            \land Perform(Head(cincoming[c]), c)
128
            \land RevSerial(c)
129
            \land UNCHANGED \langle chins, cseq \rangle
130
131
     The Server receives a message.
     SRev \triangleq
135
           \land Comm(Cop)!SRev
136
           \wedge LET cop \stackrel{\triangle}{=} Head(sincoming)
137
                    \land Perform(cop, Server)
138
                    \land Comm(Cop)!SSendSame(cop.oid.c, cop) broadcast the original operation
139
           \land \ SRevSerial
140
           \land UNCHANGED \langle chins, cseq \rangle
141
142 |
     Next \triangleq
143
           \vee \exists c \in Client : Do(c) \vee Rev(c)
144
145
     Fairness: There is no requirement that the clients ever generate operations.
     Fairness \triangleq
149
          WF_{vars}(SRev \vee \exists c \in Client : Rev(c))
150
     Spec \stackrel{\Delta}{=} Init \wedge \Box [Next]_{vars} \wedge Fairness (We care more about safety.)
152
153 F
     The compactness of CJupiter: the CSSes at all replicas are the same.
     Compactness \triangleq
157
          Comm(Cop)!EmptyChannel \Rightarrow Cardinality(Range(css)) = 1
158
    Theorem Spec \Rightarrow Compactness
161 L
      \ * Modification History
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