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MODULE *GJupiter*

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***** Google OT algorithm *****

class Msg
// type: Client → Server: [op, ack] client generated op and acked num
//   Server → Client: "ACKED" last msg is acked
//   op generated by other clients

class Client
  var outgoing // local generated operation buf: [op, stage]
  var ack      //init to 0

  synchronized procedure Do(op):
    Apply(op)
    Append(outgoing, [op, "READY"])
    Deliver()

  synchronized procedure Recv(msg):
    ack := ack + 1
    if msg = "ACKED"
      Remove(outgoing, 1)
      Deliver()
    else
      xop, outgoing := Xform(msg, outgoing)
      Apply(xop)

  procedure Deliver():
    if not Empty(outgoing) and outgoing[1].stage = "READY"
      Send(Server, [outgoing[1], ack])
      outgoing[1].stage := "SENT"

class Server // single server
  var outgoing // ops cannot be removed

  procedure SRecv(msg):
    xop, xops := Xform(msg.op, outgoing[msg.ack + 1 : Len(outgoing)])
    outgoing := outgoing[1 : msg.ack] + xops + [xop]
    Send(msg's sender, "ACKED")
    Send(other clients, xop)

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EXTENDS *JupiterInterface*, *OT*, *BufferStateSpace*

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VARIABLES

*outgoing*,    *outgoing*[*r*]: ops generated client or received by server.  
*stage*,        *stage*[*c*]: client msg sending stage.

$ack$        $ack[c]$ : client acked msg number.

CONSTANTS  $READY, SENT, ACKED$

$Stages \triangleq \{READY, SENT\}$

$vars \triangleq \langle intVars, outgoing, stage, ack \rangle$

$GMsg \triangleq$       messages exchanged by server and clients.  
 $[c : Client, op : Op \cup \{Nop\}, ack : Nat] \cup \{ACKED\} \cup Op \cup \{Nop\}$

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$TypeOK \triangleq$   
 $\wedge \quad TypeOKInt$   
 $\wedge \quad outgoing \in [Replica \rightarrow Seq(Op \cup \{Nop\})]$   
 $\wedge \quad stage \in [Client \rightarrow Seq(Stages)]$   
 $\wedge \quad ack \in [Client \rightarrow Nat]$

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$Init \triangleq$   
 $\wedge \quad InitInt$   
 $\wedge \quad outgoing = [r \in Replica \mapsto \langle \rangle]$   
 $\wedge \quad stage = [c \in Client \mapsto \langle \rangle]$   
 $\wedge \quad ack = [c \in Client \mapsto 0]$

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$Send(c) \triangleq$   
 IF  $Len(stage[c]) \neq 0$   
 THEN  $\wedge \quad stage[c][1] = READY$   
 $\wedge \quad stage' = [stage \text{ EXCEPT } ![c] = \langle SENT \rangle \circ Tail(@)]$   
 $\wedge \quad Comm!CSend([c \mapsto c, op \mapsto outgoing[c][1], ack \mapsto ack[c]])$   
 $\wedge \quad UNCHANGED \langle ack, aop, chins, outgoing, state \rangle$   
 ELSE FALSE

$ClientPerform(c, m) \triangleq$   
 IF  $m = ACKED$       last msg acked by server.  
 THEN  $\wedge \quad outgoing' = [outgoing \text{ EXCEPT } ![c] = Tail(@)]$   
 $\wedge \quad stage' = [stage \text{ EXCEPT } ![c] = Tail(@)]$   
 $\wedge \quad SetNewAop(c, Nop)$       a dummy operation.  
 ELSE       $op$  generated by other clients.  
 LET  $xform \triangleq xFormFull(OT, m, outgoing[c])$   
 IN  $\wedge \quad outgoing' = [outgoing \text{ EXCEPT } ![c] = xform.xops]$   
 $\wedge \quad UNCHANGED \quad stage$   
 $\wedge \quad SetNewAop(c, xform.xop)$

$ServerPerform(m) \triangleq$   
 LET  $xform \triangleq xFormAppend(OT, m.op, outgoing[Server], m.ack)$   
 IN  $\wedge \quad outgoing' = [outgoing \text{ EXCEPT } ![Server] = xform.xops]$

$$\begin{aligned}
& \wedge \text{SetNewAop}(\text{Server}, \text{xform.xop}) \\
& \wedge \text{Comm!SSendSameAck}(m.c, \text{ACKED}, \text{xform.xop}) \\
& \wedge \text{UNCHANGED } \langle \text{ack}, \text{stage} \rangle
\end{aligned}$$


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$$\begin{aligned}
\text{DoOp}(c, op) & \triangleq \\
& \wedge \text{SetNewAop}(c, op) \\
& \wedge \text{outgoing}' = [\text{outgoing} \text{ EXCEPT } ![c] = \text{Append}(@, op)] \\
& \wedge \text{stage}' = [\text{stage} \text{ EXCEPT } ![c] = \text{Append}(@, \text{READY})] \\
& \wedge \text{UNCHANGED } \langle \text{ack}, \text{cincoming}, \text{sincoming} \rangle \quad \text{DoOp will not send a msg.}
\end{aligned}$$


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$$\text{Do}(c) \triangleq \text{DoInt}(\text{DoOp}, c)$$

$$\begin{aligned}
\text{Rev}(c) & \triangleq \\
& \wedge \text{ack}' = [\text{ack} \text{ EXCEPT } ![c] = @ + 1] \\
& \wedge \text{RevInt}(\text{ClientPerform}, c)
\end{aligned}$$

$$\text{SRev} \triangleq \text{SRevInt}(\text{ServerPerform})$$


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$$\begin{aligned}
\text{Next} & \triangleq \\
& \vee \exists c \in \text{Client} : \text{Do}(c) \vee \text{Rev}(c) \vee \text{Send}(c) \\
& \vee \text{SRev}
\end{aligned}$$

$$\text{Fairness} \triangleq \text{WF}_{\text{vars}}(\text{SRev} \vee \exists c \in \text{Client} : \text{Rev}(c))$$

$$\text{Spec} \triangleq \text{Init} \wedge \Box[\text{Next}]_{\text{vars}} \quad \wedge \text{Fairness}$$


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$$\begin{aligned}
\text{EmptyOutgoing} & \triangleq \quad \text{all clients' outgoing is empty.} \\
& \forall c \in \text{Client} : \text{Len}(\text{stage}[c]) = 0
\end{aligned}$$

$$\begin{aligned}
\text{QC} & \triangleq \quad \text{Quiescent Consistency} \\
& \text{EmptyOutgoing} \wedge \text{Comm!EmptyChannel} \Rightarrow \text{Cardinality}(\text{Range}(\text{state})) = 1
\end{aligned}$$

$$\text{THEOREM } \text{Spec} \Rightarrow \Box \text{QC}$$


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\ \* Modification History  
\ \* Last modified Sat Apr 20 22:49:46 CST 2019 by tangruize  
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