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1  |----- MODULE XJupiter -----|
   | Specification of the Jupiter protocol described in CSCW'2014 by Xu, Sun, and Li. |
5  | EXTENDS GraphStateSpace |
6  |-----|
7  VARIABLES
8      c2ss,      c2ss[c]: the 2D state space (2ss, for short) at client c ∈ Client
9      s2ss      s2ss[c]: the 2D state space maintained by the Server for client c ∈ Client
11 |
12 |
13 | TypeOK ≜
14 |   ∧ TypeOKInt
15 |   ∧ TypeOKCtx
16 |   ∧ ∀ c ∈ Client : IsSS(c2ss[c]) ∧ IsSS(s2ss[c])
17 |-----|
18 | Init ≜
19 |   ∧ InitInt
20 |   ∧ InitCtx
21 |   ∧ c2ss = [c ∈ Client ↦ EmptySS]
22 |   ∧ s2ss = [c ∈ Client ↦ EmptySS]
23 |-----|
24 | NextEdge(r, u, ss) ≜ Return the unique outgoing edge from u in 2D state space ss
25 |   CHOOSE e ∈ ss.edge : e.from = u before a transformation at u (r is not used).
26 |-----|
27 | ClientPerform(c, cop) ≜
28 |   LET xform ≜ xForm(NextEdge, c, cop, c2ss[c]) xform: [xcop, xss, lss]
29 |   IN   ∧ c2ss' = [c2ss EXCEPT ![c] = @ ⊕ xform.xss]
30 |       ∧ SetNewAop(c, xform.xcop.op)
31 |-----|
32 | ServerPerform(cop) ≜
33 |   LET c ≜ ClientOf(cop)
34 |   xform ≜ xForm(NextEdge, Server, cop, s2ss[c]) xform: [xcop, xss, lss]
35 |   xcop ≜ xform.xcop
36 |   IN   ∧ s2ss' = [cl ∈ Client ↦ IF cl = c
37 |                                     THEN s2ss[cl] ⊕ xform.xss
38 |                                     ELSE s2ss[cl] ⊕ xform.lss]
39 |       ∧ SetNewAop(Server, xcop.op)
40 |       ∧ Comm!SSendSame(c, xcop) broadcast the transformed xcop
41 |-----|
42 | DoOp(c, op) ≜
43 |   LET cop ≜ [op ↦ op, oid ↦ [c ↦ c, seq ↦ cseq[c], ctx ↦ ds[c]]
44 |   IN   ∧ ClientPerform(c, cop)
45 |       ∧ Comm!CSend(cop)
46 |-----|
47 | Do(c) ≜
48 |   ∧ DoInt(DoOp, c)

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49       $\wedge DoCtx(c)$ 
50       $\wedge \text{UNCHANGED } s2ss$ 

52   $Rev(c) \triangleq$ 
53       $\wedge RevInt(ClientPerform, c)$ 
54       $\wedge RevCtx(c)$ 
55       $\wedge \text{UNCHANGED } s2ss$ 

57   $SRev \triangleq$ 
58       $\wedge SRevInt(ServerPerform)$ 
59       $\wedge SRevCtx$ 
60       $\wedge \text{UNCHANGED } c2ss$ 

61  |-----|
62   $Next \triangleq$ 
63       $\vee \exists c \in Client : Do(c) \vee Rev(c)$ 
64       $\vee SRev$ 

66   $Fairness \triangleq$ 
67       $WF_{vars}(SRev \vee \exists c \in Client : Rev(c))$ 

69   $Spec \triangleq Init \wedge \Box[Next]_{vars} \wedge Fairness$ 
70  |-----|
71   $CSSync \triangleq$  Each client  $c \in Client$  is synchronized with the Server.
72       $\forall c \in Client : (ds[c] = ds[Server]) \Rightarrow c2ss[c] = s2ss[c]$ 

74  THEOREM  $Spec \Rightarrow \Box CSSync$ 
75  |-----|

  \ * Modification History
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