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1  ┌────────────────── MODULE CAS Paxos ───────────────────┐
    This is a high-level specification of the CAS Paxos algorithm from the paper “CAS Paxos: Repl-
    icated State Machines without Logs” by Denis Rystsov.

    Please go to https://arxiv.org/abs/1802.07000 for the paper.

    This spec is adapted from that of Paxos consensus algorithm by Leslie Lamport, which can be
    found at https://github.com/tlaplus/Examples/blob/master/specifications/PaxosHowToWinATuringAward/Paxos.tla.

    Search “ $\langle + \rangle$ ” for the code added for CAS Paxos.

    TODO: It refines the spec in module Voting.
15 EXTENDS Integers
16 └──────────────────┐
17 CONSTANTS
18     Value,           the set of values to be proposed and chosen from
19     Acceptor,         the set acceptors
20     Quorum           the quorum system on acceptors

22 None  $\triangleq$  CHOOSE  $v : v \notin \text{Value}$ 

24 ASSUME  $\wedge \forall Q \in \text{Quorum} : Q \subseteq \text{Acceptor}$ 
25          $\wedge \forall Q1, Q2 \in \text{Quorum} : Q1 \cap Q2 \neq \{\}$ 
26 └──────────────────┐
27 Ballot  $\triangleq$  Nat

     $\langle + \rangle$  The set of all possible CAS operations. The CAS operations with cmpVal = None are
    initialization operations. We assume that the new values (i.e., swapVal) are not None.
34 CASOperation  $\triangleq$  [cmpVal : Value  $\cup$  {None}, swapVal : Value]

36 Message  $\triangleq$  the set of all possible messages that can be sent in the algorithm
37     [type : {“1a”}, bal : Ballot]
38      $\cup$  [type : {“1b”}, acc : Acceptor, bal : Ballot,
39         mbal : Ballot  $\cup$  { $-1$ }, mval : Value  $\cup$  {None}]
40      $\cup$  [type : {“2a”}, bal : Ballot, val : Value]
41      $\cup$  [type : {“2b”}, acc : Acceptor, bal : Ballot, val : Value]
42      $\cup$  [type : {“response”}, bal : Ballot]  $\langle + \rangle$  the messages sent to the user
43 └──────────────────┐
44 VARIABLES
45     maxBal[a]: the last ballot the acceptor  $a \in \text{Acceptor}$  has voted for
46     maxBal,
47      $\langle \text{maxVVal}[a], \text{maxVVal}[a] \rangle$  is the vote with the largest ballot cast by acceptor  $a \in \text{Acceptor}$ .
48     It equals  $\langle -1, \text{None} \rangle$  if  $a \in \text{Acceptor}$  has not cast any vote.
49     maxVVal, maxVVal,
50     msgs,           the set of all messages that have been sent
51     ops              $\langle + \rangle \text{ops}[b]$ : the CAS operation to be proposed at ballot  $b \in \text{Ballot}$ 

53 vars  $\triangleq$   $\langle \text{maxBal}, \text{maxVVal}, \text{maxVVal}, \text{msgs}, \text{ops} \rangle$ 

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55  $TypeOK \triangleq \wedge maxBal \in [Acceptor \rightarrow Ballot \cup \{-1\}]$

56  $\wedge maxVVal \in [Acceptor \rightarrow Value \cup \{None\}]$

57  $\wedge maxVVal \in [Acceptor \rightarrow Value \cup \{None\}]$

58  $\wedge msgs \subseteq Message$

59  $\wedge ops \in [Ballot \rightarrow CASOperation] \langle + \rangle$

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61  $Init \triangleq \wedge maxBal = [a \in Acceptor \mapsto -1]$

62  $\wedge maxVVal = [a \in Acceptor \mapsto -1]$

63  $\wedge maxVVal = [a \in Acceptor \mapsto None]$

64  $\wedge msgs = \{\}$

65  $\langle + \rangle ops$  remains unchanged; we utilize *TLC* to explore all possible *CAS* operations.

66  $\wedge ops \in [Ballot \rightarrow CASOperation]$

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67 |

68  $Send(m) \triangleq msgs' = msgs \cup \{m\}$

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69 |

The leader of ballot  $b \in Ballot$  sends a *Phase1a* message.

73  $Phase1a(b) \triangleq$

74  $\wedge Send([type \mapsto "1a", bal \mapsto b])$

75  $\wedge UNCHANGED \langle maxBal, maxVVal, maxVVal, ops \rangle$

The acceptor  $a \in Acceptor$  receives a *Phase1a* message and sends back a *Phase1b* message.

For refinement: This action implements the *IncreaseMaxBal*( $a, b$ ) action of the Voting algorithm for  $b = m.bal$ .

83  $Phase1b(a) \triangleq$

84  $\wedge \exists m \in msgs :$

85  $\wedge m.type = "1a"$

86  $\wedge m.bal > maxBal[a]$

87  $\wedge maxBal' = [maxBal \text{ EXCEPT } ![a] = m.bal]$

88  $\wedge Send([type \mapsto "1b", acc \mapsto a, bal \mapsto m.bal,$

89  $mbal \mapsto maxVVal[a], mval \mapsto maxVVal[a]])$

90  $\wedge UNCHANGED \langle maxVVal, maxVVal, ops \rangle$

In the *Phase2a*( $b, v$ ) action, the ballot  $b$  leader sends a type "2a" message asking the acceptors to vote for some value computed based on  $v$  in ballot number  $b$ .

For refinement: the enabling conditions of the action—its first two conjuncts—ensure that the second through fourth conjuncts of the four enabling conditions of action *VoteFor*( $a, b, v$ ) in module Voting will be true when acceptor  $a$  receives that message.

100  $Phase2a(b, v) \triangleq$

101  $\wedge \neg \exists m \in msgs : m.type = "2a" \wedge m.bal = b$

102  $\wedge \exists Q \in Quorum :$

103  $LET Q1b \triangleq \{m \in msgs : \wedge m.type = "1b"$

104  $\wedge m.acc \in Q$

105  $\wedge m.bal = b\}$

106  $Q1bv \triangleq \{m \in Q1b : m.mbal \geq 0\}$

107  $IN \wedge \forall a \in Q : \exists m \in Q1b : m.acc = a$

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108       $\wedge \vee \wedge Q1bv = \{ \} \quad \langle + \rangle CAS(None, v)$  as an initialization operation
109       $\wedge ops[b].cmpVal = None \quad \langle + \rangle$ 
110       $\vee \exists m \in Q1bv : \quad \langle + \rangle CAS(v, ops[b].swapVal)$  as an atomic compare-and-swap operation
111       $\wedge m.mval = v$ 
112       $\wedge \forall mm \in Q1bv : m.mbal \geq mm.mbal$ 
113       $\wedge ops[b].cmpVal = v \quad \langle + \rangle$ 
114       $\wedge Send([type \mapsto "2a", bal \mapsto b, val \mapsto ops[b].swapVal]) \quad \langle + \rangle val \mapsto ops[b].swapVal$ 
115       $\wedge UNCHANGED \langle maxBal, maxVVal, maxVVal, ops \rangle$ 

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The *Phase2b(a)* action describes what  $a \in Acceptor$  does when it receives a phase 2a message  $m \in msgs$ , which is sent by the leader of ballot  $m.bal$  asking acceptors to vote for  $m.val$  in that ballot.

For refinement: The enabling condition of the *Phase2b(a)* action together with the receipt of the phase 2a message  $m$  implies that the *VoteFor(a, m.bal, m.val)* action of module Voting is enabled and can be executed.

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127 Phase2b(a)  $\triangleq$ 
128    $\wedge \exists m \in msgs :$ 
129      $\wedge m.type = "2a"$ 
130      $\wedge m.bal \geq maxBal[a]$ 
131      $\wedge maxBal' = [maxBal \text{ EXCEPT } ![a] = m.bal]$ 
132      $\wedge maxVVal' = [maxVVal \text{ EXCEPT } ![a] = m.val]$ 
133      $\wedge maxVVal' = [maxVVal \text{ EXCEPT } ![a] = m.val]$ 
134      $\wedge Send([type \mapsto "2b", acc \mapsto a, bal \mapsto m.bal, val \mapsto m.val])$ 
135      $\wedge UNCHANGED \langle ops \rangle$ 

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$\langle + \rangle$  The leader of ballot  $b \in Ballot$  responds to the user.

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139 Respond(b)  $\triangleq$ 
140    $\wedge \neg \exists m \in msgs : m.type = "response" \wedge m.bal = b$ 
141    $\wedge \exists Q \in Quorum :$ 
142     LET  $Q2b \triangleq \{ m \in msgs : \wedge m.type = "2b"$ 
143        $\wedge m.acc \in Q$ 
144        $\wedge m.bal = b \}$ 
145     IN  $\forall a \in Q : \exists m \in Q2b : m.acc = a$ 
146    $\wedge Send([type \mapsto "response", bal \mapsto b])$ 
147    $\wedge UNCHANGED \langle maxBal, maxVVal, maxVVal, ops \rangle$ 

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149 Next  $\triangleq$ 
150    $\vee \exists b \in Ballot :$ 
151      $\vee Phase1a(b)$ 
152      $\vee \exists v \in Value : Phase2a(b, v)$ 
153      $\vee Respond(b) \quad \langle + \rangle$ 
154    $\vee \exists a \in Acceptor :$ 
155      $\vee Phase1b(a)$ 
156      $\vee Phase2b(a)$ 

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158 Spec  $\triangleq Init \wedge \Box [Next]_{vars}$ 
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\\* Modification History  
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