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1  |----- MODULE GryadkaCasRegister -----|
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   | TLA+ code: https://gist.github.com/grogers0/c7e87f9dfe58c6070b19db9d3c073b72
   | Post (A TLA+ specification for Gryadka): https://medium.com/@grogepodge/tla-specification-for-gryadka-c80cd625944e
10 EXTENDS Integers, Sequences, FiniteSets
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12 | Timestamps is the set of possible timestamps for operations to choose from.
13 |   Each operation uses a unique timestamp.
14 | Values is the set of possible values to set the register to.
15 | Acceptors is the set of nodes which act as acceptors in the paxos sense.
16 | Quorums is the set of all possible quorums, typically simple majorities.
17 CONSTANTS Timestamps, Values, Acceptors, Quorums
   |
19 ASSUME Timestamps  $\subseteq$  Nat
20 ASSUME IsFiniteSet(Timestamps)
21 NoTS  $\triangleq$  - 1
22 ASSUME NoTS  $\notin$  Timestamps
   |
24 ASSUME Quorums  $\subseteq$  SUBSET Acceptors
25 ASSUME  $\forall q1, q2 \in \textit{Quorums} : q1 \cap q2 \neq \{\}$ 
   |
27 | The initial value is chosen arbitrarily
28 InitVal  $\triangleq$  CHOOSE v  $\in$  Values : TRUE
   |
30 | msgs is the buffer of all messages. Messages can be delivered out of order or duplicated.
31 | ops is the mapping from timestamp to CAS(old, new) for operations being proposed.
32 | acceptorTS is the timestamp each acceptor is prepared for, only operations which match this value are accepted.
33 | acceptorValTS is the timestamp of the last accepted value for each acceptor, or NoTS if none has been accepted yet.
34 | acceptorValue is the last accepted value for each acceptor, or InitVal if none has been accepted yet.
35 | history is the actual order of invoke/response actions for the operations identified by the timestamp.
36 VARIABLES msgs, ops, acceptorTS, acceptorValTS, acceptorValue, history
   |-----|
38 Messages  $\triangleq$  [type : { "prepare" }, acceptor : Acceptors, ts : Timestamps]
39    $\cup$  [type : { "promise" }, acceptor : Acceptors, ts : Timestamps,
40     prevTS : Timestamps  $\cup$  { NoTS }, prevVal : Values]
41    $\cup$  [type : { "accept" }, acceptor : Acceptors, ts : Timestamps, val : Values]
42    $\cup$  [type : { "accepted" }, acceptor : Acceptors, ts : Timestamps, val : Values]
43 | Each operation represents a CAS from an oldVal to a newVal. In Gryadka,
44 | reads are treated the same as CAS(val, val)
45 Operations  $\triangleq$  [oldVal : Values, newVal : Values]
46 Events  $\triangleq$  [type : { "invoke", "response" }, ts : Timestamps]
   |
48 TypeOK  $\triangleq$   $\wedge$  msgs  $\subseteq$  Messages
49    $\wedge$  ops  $\in$  [Timestamps  $\rightarrow$  Operations]
50    $\wedge$  acceptorTS  $\in$  [Acceptors  $\rightarrow$  Timestamps  $\cup$  { NoTS }]

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51  $\wedge \text{acceptorValTS} \in [\text{Acceptors} \rightarrow \text{Timestamps} \cup \{\text{NoTS}\}]$   
 52  $\wedge \text{acceptorValue} \in [\text{Acceptors} \rightarrow \text{Values}]$   
 53  $\wedge \text{history} \in \text{Seq}(\text{Events})$   
 55  $\text{Init} \triangleq \wedge \text{msgs} = \{\}$   
 56  $\wedge \text{ops} \in [\text{Timestamps} \rightarrow \text{Operations}]$   
 57  $\wedge \text{acceptorTS} = [a \in \text{Acceptors} \mapsto \text{NoTS}]$   
 58  $\wedge \text{acceptorValTS} = [a \in \text{Acceptors} \mapsto \text{NoTS}]$   
 59  $\wedge \text{acceptorValue} = [a \in \text{Acceptors} \mapsto \text{InitVal}]$   
 60  $\wedge \text{history} = \langle \rangle$   
 62  $\vdash$   
 63  $\text{SelectMessages}(\text{type}, \text{ts}) \triangleq \{m \in \text{msgs} : m.\text{type} = \text{type} \wedge m.\text{ts} = \text{ts}\}$   
 64  $\text{PromisedValue}(\text{ts}) \triangleq \text{LET } \text{promiseMsgs} \triangleq \text{SelectMessages}(\text{"promise"}, \text{ts})$   
 65  $\text{IN } (\text{CHOOSE } m \in \text{promiseMsgs} : \forall m2 \in \text{promiseMsgs} : m.\text{prevTS} \geq m2.\text{prevTS}).\text{prevVal}$   
 66  $\vdash$   
 67  $\text{Prepare}(\text{ts}) \triangleq \wedge \text{SelectMessages}(\text{"prepare"}, \text{ts}) = \{\}$  Each timestamp must be unique  
 68  $\wedge \text{msgs}' = \text{msgs} \cup [\text{type} : \{\text{"prepare"}\}, \text{acceptor} : \text{Acceptors}, \text{ts} : \{\text{ts}\}]$   
 69  $\wedge \text{history}' = \text{Append}(\text{history}, [\text{type} \mapsto \text{"invoke"}, \text{ts} \mapsto \text{ts}])$   
 70  $\wedge \text{UNCHANGED } \langle \text{ops}, \text{acceptorTS}, \text{acceptorValTS}, \text{acceptorValue} \rangle$   
 72  $\text{RecvPrepare}(a, \text{ts}) \triangleq \wedge \text{acceptorTS}[a] = \text{NoTS} \vee \text{acceptorTS}[a] < \text{ts}$   
 73  $\wedge \text{acceptorTS}' = [\text{acceptorTS} \text{ EXCEPT } ![a] = \text{ts}]$   
 74  $\wedge \text{msgs}' = \text{msgs} \cup \{[\text{type} \mapsto \text{"promise"}, \text{acceptor} \mapsto a, \text{ts} \mapsto \text{ts},$   
 75  $\text{prevTS} \mapsto \text{acceptorValTS}[a], \text{prevVal} \mapsto \text{acceptorValue}[a]]\}$   
 76  $\wedge \text{UNCHANGED } \langle \text{ops}, \text{acceptorValTS}, \text{acceptorValue}, \text{history} \rangle$   
 78  $\text{Accept}(\text{ts}) \triangleq \wedge \{m.\text{acceptor} : m \in \text{SelectMessages}(\text{"promise"}, \text{ts})\} \in \text{Quorums}$   
 79  $\wedge \text{ops}[\text{ts}].\text{oldVal} = \text{PromisedValue}(\text{ts})$   
 80  $\wedge \text{msgs}' = \text{msgs} \cup [\text{type} : \{\text{"accept"}\}, \text{acceptor} : \text{Acceptors}, \text{ts} : \{\text{ts}\}, \text{val} : \{\text{ops}[\text{ts}].\text{newVal}\}]$   
 81  $\wedge \text{UNCHANGED } \langle \text{ops}, \text{acceptorTS}, \text{acceptorValTS}, \text{acceptorValue}, \text{history} \rangle$   
 83  $\text{RecvAccept}(a, \text{ts}, v) \triangleq \wedge \text{acceptorTS}[a] = \text{ts}$   
 84  $\wedge \text{acceptorValTS}' = [\text{acceptorValTS} \text{ EXCEPT } ![a] = \text{ts}]$   
 85  $\wedge \text{acceptorValue}' = [\text{acceptorValue} \text{ EXCEPT } ![a] = v]$   
 86  $\wedge \text{msgs}' = \text{msgs} \cup [\text{type} : \{\text{"accepted"}\}, \text{acceptor} : \{a\}, \text{ts} : \{\text{ts}\}, \text{val} : \{v\}]$   
 87  $\wedge \text{UNCHANGED } \langle \text{ops}, \text{acceptorTS}, \text{history} \rangle$   
 89  $\text{Accepted}(\text{ts}) \triangleq \wedge \{m.\text{acceptor} : m \in \text{SelectMessages}(\text{"accepted"}, \text{ts})\} \in \text{Quorums}$   
 90  $\wedge \{hpos \in \text{DOMAIN } \text{history} : \text{history}[hpos] = [\text{type} \mapsto \text{"response"}, \text{ts} \mapsto \text{ts}]\} = \{\}$   
 91  $\wedge \text{history}' = \text{Append}(\text{history}, [\text{type} \mapsto \text{"response"}, \text{ts} \mapsto \text{ts}])$   
 92  $\wedge \text{UNCHANGED } \langle \text{msgs}, \text{ops}, \text{acceptorTS}, \text{acceptorValTS}, \text{acceptorValue} \rangle$   
 94  $\text{Next} \triangleq \vee \exists \text{ts} \in \text{Timestamps} : \vee \text{Prepare}(\text{ts})$   
 95  $\vee \text{Accept}(\text{ts})$   
 96  $\vee \text{Accepted}(\text{ts})$   
 97  $\vee \exists m \in \text{msgs} : \vee m.\text{type} = \text{"prepare"} \wedge \text{RecvPrepare}(m.\text{acceptor}, m.\text{ts})$

98  $\vee m.type = \text{"accept"} \wedge RecvAccept(m.acceptor, m.ts, m.val)$

100  $Spec \triangleq Init \wedge \Box[Next]_{\langle msgs, ops, acceptorTS, acceptorValTS, acceptorValue, history \rangle}$

101  $\frac{}{}$

102  $FiniteSeq(S) \triangleq \text{UNION } \{[1 \dots n \rightarrow S] : n \in 1 \dots Cardinality(S)\}$

103  $SeqAsSet(S) \triangleq \{S[i] : i \in \text{DOMAIN } S\}$

105  $HistoryIsLinearizable \triangleq \exists order \in \{\langle \rangle\} \cup FiniteSeq(Timestamps) :$

106  $\quad \wedge \forall H \in SeqAsSet(history) : H.type = \text{"response"} \Rightarrow H.ts \in SeqAsSet(order)$

107  $\quad \wedge \forall H1\_i, H2\_i \in \text{DOMAIN } history :$

108  $\quad (history[H1\_i].type = \text{"response"} \wedge history[H2\_i].type = \text{"invoke"} \wedge H1\_i < H2\_i) \Rightarrow$

109  $\quad ($

110  $\quad \quad history[H2\_i].ts \in SeqAsSet(order) \Rightarrow$

111  $\quad \quad \exists i1, i2 \in \text{DOMAIN } order :$

112  $\quad \quad \quad \wedge order[i1] = history[H1\_i].ts$

113  $\quad \quad \quad \wedge order[i2] = history[H2\_i].ts$

114  $\quad \quad \quad \wedge i1 < i2$

115  $\quad )$

116  $\quad \wedge \forall i1, i2 \in \text{DOMAIN } order :$

117  $\quad \quad i2 = i1 + 1 \Rightarrow ops[order[i1]].newVal = ops[order[i2]].oldVal$

118  $\quad \wedge order \neq \langle \rangle \Rightarrow InitVal = ops[order[1]].oldVal$

120  $Inv \triangleq \wedge TypeOK$

121  $\quad \wedge HistoryIsLinearizable$

123  $\frac{}{}$