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1  |----- MODULE GryadkaCasRegister -----|
2  | Written by Greg Rogers |
3  | TLA+ code: https://gist.github.com/grogers0/c7e87f9dfe58c6070b19db9d3c073b72 |
4  | Post (A TLA+ specification for Gryadka): https://medium.com/@grogepodge/tla-specification-for-gryadka-c80cd625944e |
5  |-----|
6  EXTENDS Integers, Sequences, FiniteSets
7  |-----|
8  | Timestamps is the set of possible timestamps for operations to choose from. |
9  | Each operation uses a unique timestamp. |
10 | Values is the set of possible values to set the register to. |
11 | Acceptors is the set of nodes which act as acceptors in the paxos sense. |
12 | Quorums is the set of all possible quorums, typically simple majorities. |
13 CONSTANTS Timestamps, Values, Acceptors, Quorums
14
15 ASSUME  $Timestamps \subseteq Nat$ 
16 ASSUME  $IsFiniteSet(Timestamps)$ 
17  $NoTS \triangleq -1$ 
18 ASSUME  $NoTS \notin Timestamps$ 
19
20 ASSUME  $Quorums \subseteq SUBSET Acceptors$ 
21 ASSUME  $\forall q1, q2 \in Quorums : q1 \cap q2 \neq \{\}$ 
22
23 | The initial value is chosen arbitrarily |
24  $InitVal \triangleq CHOOSE v \in Values : TRUE$ 
25
26 | msgs is the buffer of all messages. Messages can be delivered out of order or duplicated. |
27 | ops is the mapping from timestamp to  $CAS(old, new)$  for operations being proposed. |
28 | acceptorTS is the timestamp each acceptor is prepared for, only operations which match this value are accepted. |
29 | acceptorValTS is the timestamp of the last accepted value for each acceptor, or NoTS if none has been accepted yet. |
30 | acceptorValue is the last accepted value for each acceptor, or InitVal if none has been accepted yet. |
31 | history is the actual order of invoke/response actions for the operations identified by the timestamp. |
32 VARIABLES msgs, ops, acceptorTS, acceptorValTS, acceptorValue, history
33 |-----|
34  $Messages \triangleq [type : \{ "prepare" \}, acceptor : Acceptors, ts : Timestamps]$ 
35  $\cup [type : \{ "promise" \}, acceptor : Acceptors, ts : Timestamps,$ 
36  $prevTS : Timestamps \cup \{ NoTS \}, prevVal : Values]$ 
37  $\cup [type : \{ "accept" \}, acceptor : Acceptors, ts : Timestamps, val : Values]$ 
38  $\cup [type : \{ "accepted" \}, acceptor : Acceptors, ts : Timestamps, val : Values]$ 
39 | Each operation represents a CAS from an oldVal to a newVal. In Gryadka, |
40 | reads are treated the same as  $CAS(val, val)$  |
41  $Operations \triangleq [oldVal : Values, newVal : Values]$ 
42  $Events \triangleq [type : \{ "invoke", "response" \}, ts : Timestamps]$ 
43
44  $TypeOK \triangleq \wedge msgs \subseteq Messages$ 
45  $\wedge ops \in [Timestamps \rightarrow Operations]$ 
46  $\wedge acceptorTS \in [Acceptors \rightarrow Timestamps \cup \{ NoTS \}]$ 

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47 $\wedge \text{acceptorValTS} \in [\text{Acceptors} \rightarrow \text{Timestamps} \cup \{\text{NoTS}\}]$
 48 $\wedge \text{acceptorValue} \in [\text{Acceptors} \rightarrow \text{Values}]$
 49 $\wedge \text{history} \in \text{Seq}(\text{Events})$
 51 $\text{Init} \triangleq \wedge \text{msgs} = \{\}$
 52 $\wedge \text{ops} \in [\text{Timestamps} \rightarrow \text{Operations}]$
 53 $\wedge \text{acceptorTS} = [a \in \text{Acceptors} \mapsto \text{NoTS}]$
 54 $\wedge \text{acceptorValTS} = [a \in \text{Acceptors} \mapsto \text{NoTS}]$
 55 $\wedge \text{acceptorValue} = [a \in \text{Acceptors} \mapsto \text{InitVal}]$
 56 $\wedge \text{history} = \langle \rangle$
 58 $\text{SelectMessages}(\text{type}, \text{ts}) \triangleq \{m \in \text{msgs} : m.\text{type} = \text{type} \wedge m.\text{ts} = \text{ts}\}$
 59 $\text{PromisedValue}(\text{ts}) \triangleq \text{LET } \text{promiseMsgs} \triangleq \text{SelectMessages}(\text{"promise"}, \text{ts})$
 60 $\text{IN } (\text{CHOOSE } m \in \text{promiseMsgs} : \forall m2 \in \text{promiseMsgs} : m.\text{prevTS} \geq m2.\text{prevTS}).\text{prevVal}$
 62 $\text{Prepare}(\text{ts}) \triangleq \wedge \text{SelectMessages}(\text{"prepare"}, \text{ts}) = \{\}$ Each timestamp must be unique
 63 $\wedge \text{msgs}' = \text{msgs} \cup [\text{type} : \{\text{"prepare"}\}, \text{acceptor} : \text{Acceptors}, \text{ts} : \{\text{ts}\}]$
 64 $\wedge \text{history}' = \text{Append}(\text{history}, [\text{type} \mapsto \text{"invoke"}, \text{ts} \mapsto \text{ts}])$
 65 $\wedge \text{UNCHANGED } \langle \text{ops}, \text{acceptorTS}, \text{acceptorValTS}, \text{acceptorValue} \rangle$
 68 $\text{RecvPrepare}(a, \text{ts}) \triangleq \wedge \text{acceptorTS}[a] = \text{NoTS} \vee \text{acceptorTS}[a] < \text{ts}$
 69 $\wedge \text{acceptorTS}' = [\text{acceptorTS} \text{ EXCEPT } ![a] = \text{ts}]$
 70 $\wedge \text{msgs}' = \text{msgs} \cup \{[\text{type} \mapsto \text{"promise"}, \text{acceptor} \mapsto a, \text{ts} \mapsto \text{ts},$
 71 $\text{prevTS} \mapsto \text{acceptorValTS}[a], \text{prevVal} \mapsto \text{acceptorValue}[a]]\}$
 72 $\wedge \text{UNCHANGED } \langle \text{ops}, \text{acceptorValTS}, \text{acceptorValue}, \text{history} \rangle$
 74 $\text{Accept}(\text{ts}) \triangleq \wedge \{m.\text{acceptor} : m \in \text{SelectMessages}(\text{"promise"}, \text{ts})\} \in \text{Quorums}$
 75 $\wedge \text{ops}[\text{ts}].\text{oldVal} = \text{PromisedValue}(\text{ts})$
 76 $\wedge \text{msgs}' = \text{msgs} \cup [\text{type} : \{\text{"accept"}\}, \text{acceptor} : \text{Acceptors}, \text{ts} : \{\text{ts}\}, \text{val} : \{\text{ops}[\text{ts}].\text{newVal}\}]$
 77 $\wedge \text{UNCHANGED } \langle \text{ops}, \text{acceptorTS}, \text{acceptorValTS}, \text{acceptorValue}, \text{history} \rangle$
 79 $\text{RecvAccept}(a, \text{ts}, v) \triangleq \wedge \text{acceptorTS}[a] = \text{ts}$
 80 $\wedge \text{acceptorValTS}' = [\text{acceptorValTS} \text{ EXCEPT } ![a] = \text{ts}]$
 81 $\wedge \text{acceptorValue}' = [\text{acceptorValue} \text{ EXCEPT } ![a] = v]$
 82 $\wedge \text{msgs}' = \text{msgs} \cup [\text{type} : \{\text{"accepted"}\}, \text{acceptor} : \{a\}, \text{ts} : \{\text{ts}\}, \text{val} : \{v\}]$
 83 $\wedge \text{UNCHANGED } \langle \text{ops}, \text{acceptorTS}, \text{history} \rangle$
 85 $\text{Accepted}(\text{ts}) \triangleq \wedge \{m.\text{acceptor} : m \in \text{SelectMessages}(\text{"accepted"}, \text{ts})\} \in \text{Quorums}$
 86 $\wedge \{hpos \in \text{DOMAIN } \text{history} : \text{history}[hpos] = [\text{type} \mapsto \text{"response"}, \text{ts} \mapsto \text{ts}]\} = \{\}$
 87 $\wedge \text{history}' = \text{Append}(\text{history}, [\text{type} \mapsto \text{"response"}, \text{ts} \mapsto \text{ts}])$
 88 $\wedge \text{UNCHANGED } \langle \text{msgs}, \text{ops}, \text{acceptorTS}, \text{acceptorValTS}, \text{acceptorValue} \rangle$
 90 $\text{Next} \triangleq \vee \exists \text{ts} \in \text{Timestamps} : \vee \text{Prepare}(\text{ts})$
 91 $\vee \text{Accept}(\text{ts})$
 92 $\vee \text{Accepted}(\text{ts})$
 93 $\vee \exists m \in \text{msgs} : \vee m.\text{type} = \text{"prepare"} \wedge \text{RecvPrepare}(m.\text{acceptor}, m.\text{ts})$

94 $\vee m.type = \text{"accept"} \wedge RecvAccept(m.acceptor, m.ts, m.val)$
 96 $Spec \triangleq Init \wedge \Box[Next]_{\langle msgs, ops, acceptorTS, acceptorValTS, acceptorValue, history \rangle}$
 97 \vdash
 98 $FiniteSeq(S) \triangleq \text{UNION } \{[1 \dots n \rightarrow S] : n \in 1 \dots Cardinality(S)\}$
 99 $SeqAsSet(S) \triangleq \{S[i] : i \in \text{DOMAIN } S\}$
 101 $HistoryIsLinearizable \triangleq \exists order \in \{\langle \rangle\} \cup FiniteSeq(Timestamps) :$
 102 $\quad \wedge \forall H \in SeqAsSet(history) : H.type = \text{"response"} \Rightarrow H.ts \in SeqAsSet(order)$
 103 $\quad \wedge \forall H1_i, H2_i \in \text{DOMAIN } history :$
 104 $\quad (history[H1_i].type = \text{"response"} \wedge history[H2_i].type = \text{"invoke"} \wedge H1_i < H2_i) \Rightarrow$
 105 $\quad ($
 106 $\quad \quad history[H2_i].ts \in SeqAsSet(order) \Rightarrow$
 107 $\quad \quad \exists i1, i2 \in \text{DOMAIN } order :$
 108 $\quad \quad \quad \wedge order[i1] = history[H1_i].ts$
 109 $\quad \quad \quad \wedge order[i2] = history[H2_i].ts$
 110 $\quad \quad \quad \wedge i1 < i2$
 111 $\quad \quad)$
 112 $\quad \wedge \forall i1, i2 \in \text{DOMAIN } order :$
 113 $\quad \quad i2 = i1 + 1 \Rightarrow ops[order[i1]].newVal = ops[order[i2]].oldVal$
 114 $\quad \wedge order \neq \langle \rangle \Rightarrow InitVal = ops[order[1]].oldVal$
 116 $Inv \triangleq \wedge TypeOK$
 117 $\quad \wedge HistoryIsLinearizable$
 119 \vdash