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1  ┌────────────────── MODULE CAS Paxos ───────────────────┐
    This is a high-level specification of the CAS Paxos algorithm from the paper “CAS Paxos: Repl-
    icated State Machines without Logs” by Denis Rystsov.

    Please go to https://arxiv.org/abs/1802.07000 for the paper.

    This spec is adapted from that of Paxos consensus algorithm by Leslie Lamport, which can be
    found at https://github.com/tlaplus/Examples/blob/master/specifications/PaxosHowToWinATuringAward/Paxos.tla.

    Search “ $\langle + \rangle$ ” for the code added for CAS Paxos.

    TODO: It refines the spec in module Voting.
15 EXTENDS Integers
16 └──────────────────┐
17 CONSTANTS
18     Value,           the set of values to be proposed and chosen from
19     Acceptor,        the set acceptors
20     Quorum           the quorum system on acceptors

22 None  $\triangleq$  CHOOSE  $v : v \notin \textit{Value}$ 

24 ASSUME  $\bigwedge \forall Q \in \textit{Quorum} : Q \subseteq \textit{Acceptor}$ 
25          $\bigwedge \forall Q1, Q2 \in \textit{Quorum} : Q1 \cap Q2 \neq \{\}$ 
26 └──────────────────┐
27 Ballot  $\triangleq$  Nat

     $\langle + \rangle$  The set of all possible CAS operations. The CAS operations with cmpVal = None are
    initialization operations. We assume that the new values (i.e., swapVal) are not None.
34 CASOperation  $\triangleq$  [cmpVal : Value  $\cup$  {None}, swapVal : Value]

36 Message  $\triangleq$  the set of all possible messages that can be sent in the algorithm
37     [type : {“1a”}, bal : Ballot]
38      $\cup$  [type : {“1b”}, acc : Acceptor, bal : Ballot,
39         mbal : Ballot  $\cup$  { $-1$ }, mval : Value  $\cup$  {None}]
40      $\cup$  [type : {“2a”}, bal : Ballot, val : Value]
41      $\cup$  [type : {“2b”}, acc : Acceptor, bal : Ballot, val : Value]
42      $\cup$  [type : {“response”}, bal : Ballot]  $\langle + \rangle$  the messages sent to the user
43 └──────────────────┐
44 VARIABLES
45     maxBal[a]: the last ballot the acceptor  $a \in \textit{Acceptor}$  has voted for
46     maxBal,
47      $\langle \textit{maxVVal}[a], \textit{maxVVal}[a] \rangle$  is the vote with the largest ballot cast by acceptor  $a \in \textit{Acceptor}$ .
48     It equals  $\langle -1, \textit{None} \rangle$  if  $a \in \textit{Acceptor}$  has not cast any vote.
49     maxVVal, maxVVal,
50     msgs,           the set of all messages that have been sent
51     ops              $\langle + \rangle \textit{ops}[b]$ : the CAS operation to be proposed at ballot  $b \in \textit{Ballot}$ 

53 vars  $\triangleq$   $\langle \textit{maxBal}, \textit{maxVVal}, \textit{maxVVal}, \textit{msgs}, \textit{ops} \rangle$ 

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54 |

55 $TypeOK \triangleq \wedge maxBal \in [Acceptor \rightarrow Ballot \cup \{-1\}]$

56 $\wedge maxVVal \in [Acceptor \rightarrow Value \cup \{None\}]$

57 $\wedge maxVVal \in [Acceptor \rightarrow Value \cup \{None\}]$

58 $\wedge msgs \subseteq Message$

59 $\wedge ops \in [Ballot \rightarrow CASOperation] \langle + \rangle$

60 |

61 $Init \triangleq \wedge maxBal = [a \in Acceptor \mapsto -1]$

62 $\wedge maxVVal = [a \in Acceptor \mapsto -1]$

63 $\wedge maxVVal = [a \in Acceptor \mapsto None]$

64 $\wedge msgs = \{\}$

65 $\langle + \rangle ops$ remains unchanged; we utilize *TLC* to explore all possible *CAS* operations.

66 $\wedge ops \in [Ballot \rightarrow CASOperation]$

67 |

68 $Send(m) \triangleq msgs' = msgs \cup \{m\}$

69 |

The leader of ballot $b \in Ballot$ sends a *Phase1a* message.

73 $Phase1a(b) \triangleq$

74 $\wedge Send([type \mapsto "1a", bal \mapsto b])$

75 $\wedge UNCHANGED \langle maxBal, maxVVal, maxVVal, ops \rangle$

The acceptor $a \in Acceptor$ receives a *Phase1a* message and sends back a *Phase1b* message.

For refinement: This action implements the *IncreaseMaxBal*(a, b) action of the Voting algorithm for $b = m.bal$.

83 $Phase1b(a) \triangleq$

84 $\wedge \exists m \in msgs :$

85 $\wedge m.type = "1a"$

86 $\wedge m.bal > maxBal[a]$

87 $\wedge maxBal' = [maxBal \text{ EXCEPT } ![a] = m.bal]$

88 $\wedge Send([type \mapsto "1b", acc \mapsto a, bal \mapsto m.bal,$

89 $mbal \mapsto maxVVal[a], mval \mapsto maxVVal[a]])$

90 $\wedge UNCHANGED \langle maxVVal, maxVVal, ops \rangle$

In the *Phase2a*(b, v) action, the ballot b leader sends a type "2a" message asking the acceptors to vote for some value computed based on v in ballot number b .

For refinement: the enabling conditions of the action—its first two conjuncts—ensure that the second through fourth conjuncts of the four enabling conditions of action *VoteFor*(a, b, v) in module Voting will be true when acceptor a receives that message.

100 $Phase2a(b, v) \triangleq$

101 $\wedge \neg \exists m \in msgs : m.type = "2a" \wedge m.bal = b$

102 $\wedge \exists Q \in Quorum :$

103 $LET Q1b \triangleq \{m \in msgs : \wedge m.type = "1b"$

104 $\wedge m.acc \in Q$

105 $\wedge m.bal = b\}$

106 $Q1bv \triangleq \{m \in Q1b : m.mbal \geq 0\}$

107 $IN \wedge \forall a \in Q : \exists m \in Q1b : m.acc = a$

108 $\wedge \vee \wedge Q1bv = \{ \} \quad \langle + \rangle CAS(None, v)$ as an initialization operation
 109 $\wedge ops[b].cmpVal = None \quad \langle + \rangle$
 110 $\vee \exists m \in Q1bv : \quad \langle + \rangle CAS(v, ops[b].swapVal)$ as an atomic compare-and-swap operation
 111 $\wedge m.mval = v$
 112 $\wedge \forall mm \in Q1bv : m.mbal \geq mm.mbal$
 113 $\wedge ops[b].cmpVal = v \quad \langle + \rangle$
 114 $\wedge Send([type \mapsto "2a", bal \mapsto b, val \mapsto ops[b].swapVal]) \quad \langle + \rangle val \mapsto ops[b].swapVal$
 115 $\wedge UNCHANGED \langle maxBal, maxVVal, maxVVal, ops \rangle$

The *Phase2b(a)* action describes what $a \in Acceptor$ does when it receives a phase 2a message $m \in msgs$, which is sent by the leader of ballot $m.bal$ asking acceptors to vote for $m.val$ in that ballot.

For refinement: The enabling condition of the *Phase2b(a)* action together with the receipt of the phase 2a message m implies that the *VoteFor(a, m.bal, m.val)* action of module Voting is enabled and can be executed.

127 $Phase2b(a) \triangleq$
 128 $\wedge \exists m \in msgs :$
 129 $\wedge m.type = "2a"$
 130 $\wedge m.bal \geq maxBal[a]$
 131 $\wedge maxBal' = [maxBal \text{ EXCEPT } ![a] = m.bal]$
 132 $\wedge maxVVal' = [maxVVal \text{ EXCEPT } ![a] = m.bal]$
 133 $\wedge maxVVal' = [maxVVal \text{ EXCEPT } ![a] = m.val]$
 134 $\wedge Send([type \mapsto "2b", acc \mapsto a, bal \mapsto m.bal, val \mapsto m.val])$
 135 $\wedge UNCHANGED \langle ops \rangle$

$\langle + \rangle$ The leader of ballot $b \in Ballot$ responds to the user.

139 $Respond(b) \triangleq$
 140 $\wedge \neg \exists m \in msgs : m.type = "response" \wedge m.bal = b$
 141 $\wedge \exists Q \in Quorum :$
 142 $\text{LET } Q2b \triangleq \{ m \in msgs : \wedge m.type = "2b"$
 143 $\wedge m.acc \in Q$
 144 $\wedge m.bal = b \}$
 145 $\text{IN } \forall a \in Q : \exists m \in Q2b : m.acc = a$
 146 $\wedge Send([type \mapsto "response", bal \mapsto b])$
 147 $\wedge UNCHANGED \langle maxBal, maxVVal, maxVVal, ops \rangle$

149 $Next \triangleq \vee \exists b \in Ballot : \vee Phase1a(b)$
 150 $\vee \exists v \in Value : Phase2a(b, v)$
 151 $\vee Respond(b) \quad \langle + \rangle$
 152 $\vee \exists a \in Acceptor : \vee Phase1b(a)$
 153 $\vee Phase2b(a)$

155 $Spec \triangleq Init \wedge \Box [Next]_{vars}$
 156