μKanren

A Minimal Functional Core for Relational Programming

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- Motivation
- Beispiele
- Konzepte
- Implementierung
- Zusammenfassung

Motivation

- relationale Programmierung
- Scheme (39 Zeilen Code)
- Einbettung in anderen Sprachen
- Erweiterbarkeit
- Eleganz

Scheme

```
(append '(1 2 3) '(4 5 6)); (1 2 3 4 5 6)
```

Beispiele

```
(run* (q)
  (=== q 5))
; (5)
```

```
(run* (q)
 (=== q 5))
; (5)
```

```
(run* (q)
  (=== q q))
; (_.0)
```

Konzepte

```
(((#(0) . 5) (#(1) . #(0)) (#(2) . (1 2 3))) . 3); 
Substitution Counter ^
```

```
(((#(0) . 5) (#(1) . #(0)) (#(2) . (1 2 3))) . 3); 
Substitution Counter ^
```

```
(((#(0) . 5) (#(1) . #(0)) (#(2) . (1 2 3))) . 3); 
Substitution Counter ^
```

```
(((#(0) . 5) (#(1) . #(0)) (#(2) . (1 2 3))) . 3); \tag{Counter \}
```

Goal

Funktion mit State als Argument

```
(=== (var 0) 5); <goal>
```

Stream

Liste von States

```
((=== (var 0) 5) '(() . 1)); ((((#(0) . 5)) . 1))
```

Implementierung

```
((call/fresh
  (lambda (q)
        (=== q 5)))
empty-state)
```

```
((call/fresh
  (lambda (q)
        (=== q 5)))
empty-state)
; (((#(0) . 5)) . 1))
```

```
((=== (var 0) 5) '(() . 1)); ((((#(0) . 5)) . 1))
```

$$(===)(x)$$
 2 3)
 (1) y 3))
 $(x = 1, y = 2)$

```
((=== (var 0) 6) `((,(var 0) . 5) . 1)); ()
```

```
((=== (var 0) 6) `((,(var 0) . 5) . 1));
```

```
(define (=== u v)
   (lambda (s/c)
      (let ((s (unify u v (car s/c))))
       (if s
            (unit `(,s . ,(cdr s/c)))
            mzero))))
```

```
(define (unify u v s)
 (let ((u (walk u s))
        (v (walk v s)))
    (cond
     ((and (var? u) (var? v) (var=? u v)) s)
     ((var? u) (ext-s u v s))
    ((var? v) (ext-s v u s))
     ((and (pair? u) (pair? v))
      (let ((s (unify (car u) (car v) s)))
        (and s (unify (cdr u) (cdr v) s)))
     (else (and (eqv? u v) s))))
```

```
(define (unify u v s)
 (let ((u (walk u s))
        (v (walk v s)))
    (cond
     ((and (var? u) (var? v) (var=? u v)) s)
     ((var? u) (ext-s u v s))
     ((var? v) (ext-s v u s))
     ((and (pair? u) (pair? v))
      (let ((s (unify (car u) (car v) s)))
        (and s (unify (cdr u) (cdr v) s)))
     (else (and (eqv? u v) s))))
```

```
(walk #(1) ((#(0) . 5)
(#(1) . 6)))
```

```
(walk #(1) ((#(0) . 5)
(#(1) . 6)))
```

```
(walk #(1) ((#(0) . 5)
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```

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(#(1) . 6)))
```

```
(walk #(1) ((#(0) . 5)
(#(1) . 6)))
```

```
(walk #(1) ((#(0) . 5))
            (#(1) . #(2))
            (#(2).7)
(walk #(2) ((#(0) . 5)
            (#(1) . #(2))
            (#(2).7)
```

```
(walk #(1) ((#(0) . 5))
            (#(1) . #(2))
            (#(2).7)
(walk #(2) ((#(0) . 5)
            (#(1) . #(2))
            (#(2).7)
```

```
(walk #(1) ((#(0) . 5))
            (#(1) . #(2))
            (#(2).7)
(walk #(2) ((#(0) . 5)
            (#(1) . #(2))
            (#(2).7)
```

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(walk #(1) ((#(0) . 5)
            (#(1) . #(2))
            (#(2).7)
(walk #(2) ((#(0) . 5))
            (#(1) . #(2))
            (#(2).7)
```

```
(walk #(1) ((#(0) . 5)
            (#(1) . #(2))
            (#(2).7)
(walk #(2) ((#(0) . 5)
            (#(1) . #(2))
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```

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(walk #(1) ((#(0) . 5)
            (#(1) . #(2))
            (#(2).7)
(walk #(2) ((#(0) . 5)
            (#(1) . #(2))
            (#(2).7)
```

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(walk #(1) ((#(0) . 5)
            (#(1) . #(2))
            (#(2).7)
(walk #(2) ((#(0) . 5)
            (#(1) . #(2))
            (#(2).7)
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            (#(2).7)
(walk #(2) ((#(0) . 5)
            (#(1) . #(2))
            (#(2).7)
```

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(walk #(1) ((#(0) . 5))
            (#(1) . #(2))
            (#(2).7)
(walk #(2) ((#(0) . 5)
            (#(1) . #(2))
            (#(2).7)
```

```
(walk #(1) ((#(0) . 5))
            (#(1) . #(2))
            (#(2).7)
(walk #(2) ((#(0) . 5)
            (#(1) . #(2))
            (#(2).7)
```

```
(define (unify u v s)
 (let ((u (walk u s))
        (v (walk v s)))
    (cond
     ((and (var? u) (var? v) (var=? u v)) s)
     ((var? u) (ext-s u v s))
    ((var? v) (ext-s v u s))
     ((and (pair? u) (pair? v))
      (let ((s (unify (car u) (car v) s)))
        (and s (unify (cdr u) (cdr v) s)))
     (else (and (eqv? u v) s))))
```

```
(define (unify u v s)
  (let ((u (walk u s))
        (v (walk v s)))
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     ((and (var? u) (var? v) (var=? u v)) s)
     ((var? u) (ext-s u v s))
     ((var? v) (ext-s v u s))
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     ((var? u) (ext-s u v s))
    ((var? v) (ext-s v u s))
     ((and (pair? u) (pair? v))
      (let ((s (unify (car u) (car v) s)))
        (and s (unify (cdr u) (cdr v) s)))
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     ((var? u) (ext-s u v s))
     ((var? v) (ext-s v u s))
     ((and (pair? u) (pair? v))
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        (and s (unify (cdr u) (cdr v) s)))
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     ((and (var? u) (var? v) (var=? u v)) s)
     ((var? u) (ext-s u v s))
     ((var? v) (ext-s v u s))
     ((and (pair? u) (pair? v))
      (let ((s (unify (car u) (car v) s)))
        (and s (unify (cdr u) (cdr v) s)))
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    (cond
     ((and (var? u) (var? v) (var=? u v)) s)
     ((var? u) (ext-s u v s))
     ((var? v) (ext-s v u s))
     ((and (pair? u) (pair? v))
      (let ((s (unify (car u) (car v) s)))
        (and s (unify (cdr u) (cdr v) s)))
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     ((and (var? u) (var? v) (var=? u v)) s)
     ((var? u) (ext-s u v s))
     ((var? v) (ext-s v u s))
     ((and (pair? u) (pair? v))
      (let ((s (unify (car u) (car v) s)))
        (and s (unify (cdr u) (cdr v) s)))
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```
(define (unify u v s)
  (let ((u (walk u s))
        (v (walk v s)))
    (cond
     ((and (var? u) (var? v) (var=? u v)) s)
     ((var? u) (ext-s u v s))
    ((var? v) (ext-s v u s))
     ((and (pair? u) (pair? v))
      (let ((s (unify (car u) (car v) s)))
        (and s (unify (cdr u) (cdr v) s)))
     (else (and (eqv? u v) s))))
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```
(define (unify u v s)
  (let ((u (walk u s))
        (v (walk v s)))
    (cond
     ((and (var? u) (var? v) (var=? u v)) s)
     ((var? u) (ext-s u v s))
     ((var? v) (ext-s v u s))
     ((and (pair? u) (pair? v))
      (let ((s (unify (car u) (car v) s)))
        (and s (unify (cdr u) (cdr v) s)))
     (else (and (eqv? u v) s))))
```

```
(define (unify u v s)
  (let ((u (walk u s))
        (v (walk v s)))
    (cond
     ((and (var? u) (var? v) (var=? u v)) s)
     ((var? u) (ext-s u v s))
     ((var? v) (ext-s v u s))
     ((and (pair? u) (pair? v))
      (let ((s (unify (car u) (car v) s)))
        (and s (unify (cdr u) (cdr v) s)))
     (else (and (eqv? u v) s))))
```

```
(define (unify u v s)
  (let ((u (walk u s))
        (v (walk v s)))
    (cond
     ((and (var? u) (var? v) (var=? u v)) s)
     ((var? u) (ext-s u v s))
     ((var? v) (ext-s v u s))
     ((and (pair? u) (pair? v))
      (let ((s (unify (car u) (car v) s)))
        (and s (unify (cdr u) (cdr v) s)))
     (else (and (eqv? u v) s))))
```

```
(define (unify u v s)
  (let ((u (walk u s))
        (v (walk v s)))
    (cond
     ((and (var? u) (var? v) (var=? u v)) s)
     ((var? u) (ext-s u v s))
     ((var? v) (ext-s v u s))
     ((and (pair? u) (pair? v))
      (let ((s (unify (car u) (car v) s)))
        (and s (unify (cdr u) (cdr v) s)))
     (else (and (eqv? u v) s))))
```

```
((=== `(,(var 0) 2 3) `(1 ,(var 1) 3)) '(() . 2)); (((#(0) . 1) (#(1) . 2)) . 2))
```

```
(run* (q)
  (disj
     (=== q 5)
     (=== q 6)))
```

```
(run* (q)
  (conj
     (=== q 5)
     (=== q 6)))
```

```
(define (conj g1 g2)
  (lambda (s/c)
        (bind (g1 s/c) g2)))
```

```
(define (conj g1 g2)
  (lambda (s/c)
      (bind (g1 s/c) g2)))
```

```
(run* (q)
  (conj
     (=== q 5)
     (=== q 6)))
```

```
(define (fives x)
  (disj
    (=== x 5)
    (fives x)))
```

```
(define (fives x)
 (disj
    (=== x 5)
    (fives x)))
; (while (not end-of-the-universe)
; ...)
```

```
(define (fives x)
  (disj
    (=== \times 5)
    (lambda (s/c)
      (lambda ()
        ((fives x) s/c))))
```

```
(define (mplus $1 $2)
  (cond
   ((null? $1) $2)
   (else (cons (car $1) (mplus (cdr $1) $2))))
```

```
(define (mplus $1 $2)
  (cond
  ((null? $1) $2)
   ((procedure? $1) (lambda () (mplus ($1) $2)))
   (else (cons (car $1) (mplus (cdr $1) $2))))
```

```
(define (bind $ g)
  (cond
   ((null? $) mzero)
   (else (mplus (g (car $)) (bind (cdr $) g)))))
```

```
(define (bind $ g)
  (cond
  ((null? $) mzero)
   ((procedure? $) (lambda () (bind ($) g)))
   (else (mplus (g (car $)) (bind (cdr $) g))))
```

```
(define (fives x))
  (disj
    (=== \times 5)
    (lambda (s/c)
      (lambda ()
        ((fives x) s/c))))
```

```
(define (sixes x)
  (disj
    (=== x 6)
    (lambda (s/c)
      (lambda ()
        ((sixes x) s/c))))
```

```
(run*(x)
  (disj
    (fives x)
    (sixes x)))
```

```
(run*(x)
  (disj
    (fives x)
    (sixes x)))
; only fives
```

```
(define (mplus $1 $2)
  (cond
  ((null? $1) $2)
   ((procedure? $1) (lambda () (mplus ($1) $2)))
   (else (cons (car $1) (mplus (cdr $1) $2))))
```

```
(define (mplus $1 $2)
  (cond
  ((null? $1) $2)
   ((procedure? $1) (lambda () (mplus ($1) $2)))
   (else (cons (car $1) (mplus (cdr $1) $2))))
```

```
(define (mplus $1 $2)
  (cond
  ((null? $1) $2)
   ((procedure? $1) (lambda () (mplus $2 ($1))))
   (else (cons (car $1) (mplus (cdr $1) $2))))
```

```
(run* (x)
  (disj
        (fives x)
        (sixes x)))
; fives and sixes!
```

```
(run* (x)
  (disj
      (fives x) ; (5 ...)
      (sixes x)))
```

```
(run* (x)
  (disj
        (fives x)
        (sixes x))); (6 ...)
```

```
(run* (x)
  (disj
        (fives x)
        (sixes x)))
```

Grenzen

- Negation (=/= in cKanren)
- Zahlen (miniKanren, cKanren)
- "Interface" (run* und andere mit Macros)
- Performance (u.a. Tabling in miniKanren)

Zusammenfassung

- Eleganz
- Erweiterbarkeit (cKanren, αKanren, rKanren)
- Einbettung in anderen Sprachen (> 20)
- relationale Programmierung

```
(define (var x) (vector x))
(define (var? x) (vector? x))
(define (var=? x1 x2) (= (vector-ref x1 0) (vector-ref x2 0)))
(define (walk u s)
 (let ((pr (and (var? u) (assp (lambda (v) (var=? u v)) s))))
  (if pr (walk (cdr pr) s) u)))
(define (ext-s x v s) `((,x .,v) .,s))
(define (=== u v)
 (lambda (s/c)
    (let ((s (unify u v (car s/c))))
      (if s (unit `(,s . ,(cdr s/c))) mzero))))
(define (unit s/c) (cons s/c mzero))
(define mzero ,())
(define (unify uv vv s)
  (let ((u (walk uv s)) (v (walk vv s)))
  (cond
     ((and (var? u) (var? v) (var=? u v)) s)
     ((var? u) (ext-s u v s))
     ((var? v) (ext-s v u s))
     ((and (pair? u) (pair? v))
     (let ((s (unify (car u) (car v) s)))
        (and s (unify (cdr u) (cdr v) s))))
     (else (and (eqv? u v) s)))))
(define (call/fresh f)
 (lambda (s/c)
   (let ((c (cdr s/c)))
     ((f (var c)) (,(car s/c) . ,(+ c 1)))))
(define (disj g1 g2) (lambda (s/c) (mplus (g1 s/c) (g2 s/c))))
(define (conj q1 q2) (lambda (s/c) (bind (q1 s/c) q2)))
(define (mplus $1 $2)
  (cond
  ((null? $1) $2)
  ((procedure? $1) (lambda () (mplus $2 ($1))))
   (else (cons (car $1) (mplus (cdr $1) $2)))))
(define (bind $ q)
  (cond
  ((null? $) mzero)
  ((procedure? $) (lambda () (bind ($) g)))
   (else (mplus (q (car $)) (bind (cdr $) q)))))
```

(define + -)

```
(define + -)
; have fun!
```

```
(define + -)
; have fun!
; any questions?
```

Quellen

Jason Hemann und Daniel P. Friedman, <u>µKanren: A Minimal Functional Core for Relational Programming</u>, 2013.

Daniel P. Friedman, William E. Byrd und Oleg Kiselyov, <u>The Reasoned Schemer</u>, Oktober 2005.

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