# Traf: a Graphical Proof Tree Viewer Cooperating with Coq through Proof General

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#### Readability of (tactic-based/procedural) formal proofs

Formal vs. Informal Proof --- Software Foundations

"A proof is an act of communication."

"Formal proofs are useful in many ways, but they are not very efficient ways of communicating ideas between human beings."

(You have to) "step through the tactics one after the other in your mind".

due to the lack of context and goal stack at each point

How about writing contexts and/or subgoals in the proof?

## Styles of formal proofs (and tools)

#### Declarative proofs

- informative, readable without tools
- laborious to write intermediate formulae

#### Tactic-based / procedural

- hard to read the proof scripts
- preferable for writing concise proofs interactively by making use of the theorem prover's automation facilities

(There are many combined approaches ... )

Do we have to change the style of formal proof?

Our approach:

provide support for writing and reading proof scripts

"interactive monitor"

for Proof General users

#### Visualizing a proof as a proof tree

```
Theorem pq_qp: forall P Q: Prop,
  P \/ Q -> Q \/ P.
Proof.
  intros P Q.
  intros H.
  destruct H as [HP | HQ].
  right. assumption.
  left. assumption.
Qed.
```

$$\frac{A}{A \vee B} \text{ ($\vee$-introl)}$$

$$\frac{A[y/x]}{\forall x.A} (\forall \text{-intro})$$

$$\frac{B}{A \vee B}$$
 ( $\vee$ -intror)

$$\begin{array}{c|ccc}
 & [A] & [B] \\
 & | & | \\
 \hline
 & A \lor B & C & C \\
\hline
 & C & & (\lor\text{-elim})
\end{array}$$

Gentzen-style natural deduction proof

#### Visualizing a proof as a proof tree

```
Theorem pq qp: forall P Q: Prop,
  P \/ O -> O \/ P.
Proof.
  intros P Q.
  intros H.
  destruct H as [HP | HQ].
  right. assumption.
  left. assumption.
Qed.
                                                      \frac{A}{A \vee B} (\vee-introl)
                                        [A]
                                           [B]
                             A \lor \underline{B} \quad C \quad C (\lor-elim)
```

Gentzen-style natural deduction proof

#### Visualizing a proof as a proof tree

```
Theorem pq_qp: forall P Q: Prop,
  P \/ Q -> Q \/ P.
Proof.
  intros P Q.
  intros H.
  destruct H as [HP | HQ].
  right. assumption.
  left. assumption.
Qed.
```

$$\frac{[P]^{2}}{Q \lor P} (\lor -intror) \qquad \frac{[Q]^{2}}{Q \lor P} (\lor -introl) \\
\frac{Q \lor P}{Q \lor P} (\lor -introl)^{2} \\
\frac{Q \lor P}{P \lor Q \to Q \lor P} (\to -introl)^{1} \\
\frac{P \lor Q \to Q \lor P}{\forall P \lor Q \to Q \lor P} (\forall -introl)$$

Gentzen-style natural deduction proof

```
Theorem pq_qp: forall P Q: Prop,
  P \/ Q -> Q \/ P.
Proof.
  intros P Q.
  intros H.
  destruct H as [HP | HQ].
  right. assumption.
  left. assumption.
Qed.
```

```
[HP:P]
                                          [HQ:Q]
                        -assumption.
                                                  -assumption.
                    Р
                     —right.
                                                 ⊣eft.
[H:PVQ]
                 QVP
                                           QVP
                                                 -destruct H as [HP I HQ].
                                 QVP
                                            intros H.
                             PVQ \rightarrow QVP
                                                   -intros P Q.
                      \forall PQ: Prop, PVQ \rightarrow QVP
```

Traf's tree

$$\frac{[P]^{2}}{[P \vee Q]^{1}} \frac{[Q]^{2}}{[Q \vee P]} (\forall -intror) \frac{[Q]^{2}}{[Q \vee P]} (\forall -introl) \frac{[Q \vee P]^{2}}{[Q \vee P]} (\forall -elim)^{2}$$

$$\frac{[Q \vee P]^{2}}{[Q \vee P]^{2}} (\forall -elim)^{2}$$

easy to read 👙

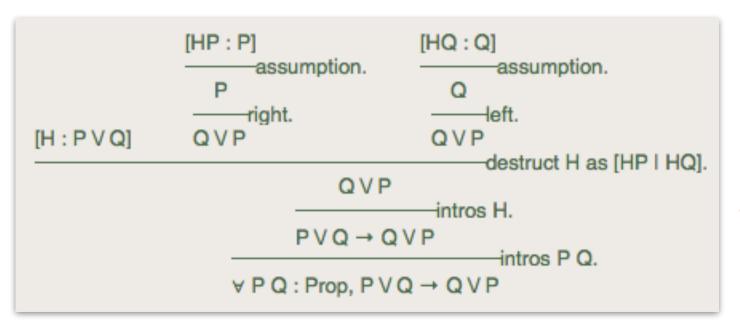


(if it is not too large)

burdensome 😥 to build



by using paper and pencil



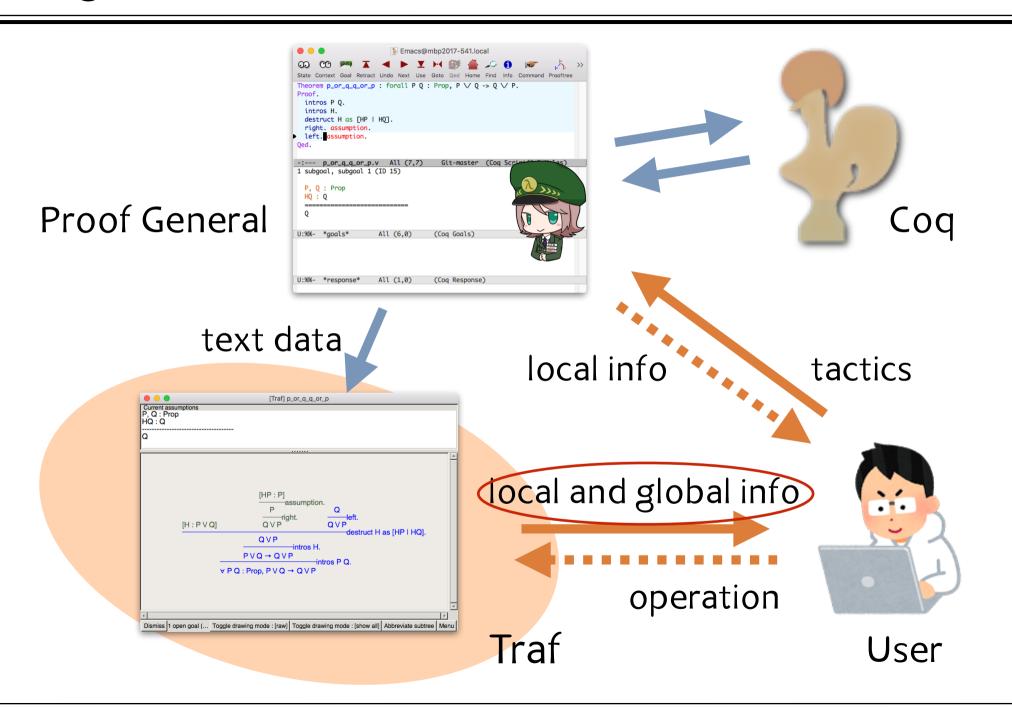
easy to read 👙



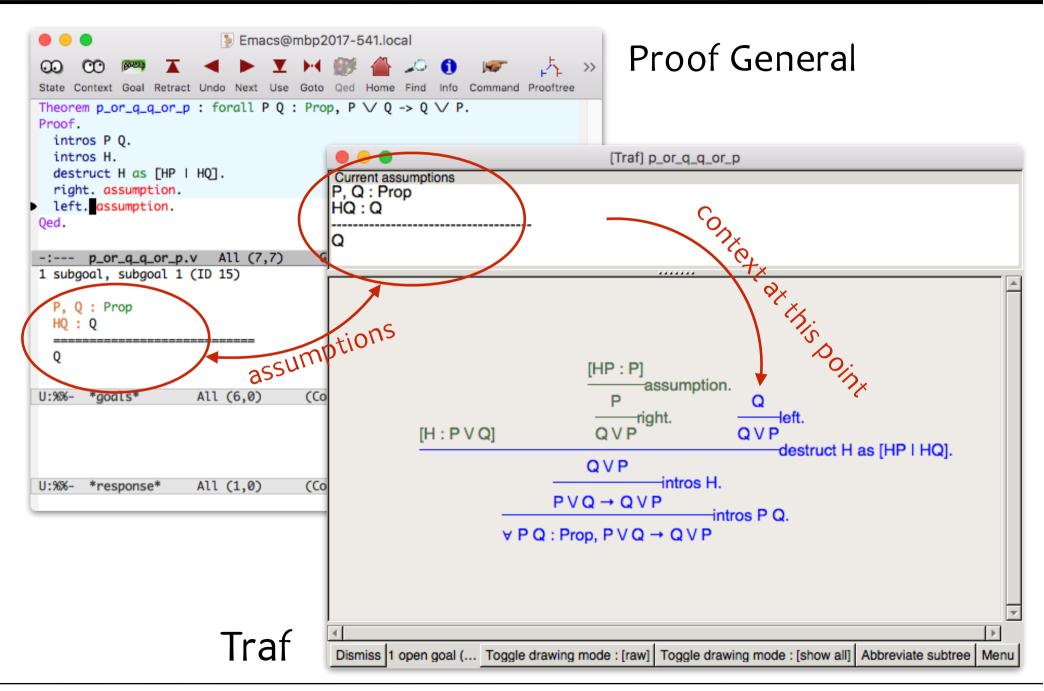
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no effort required easy to build a

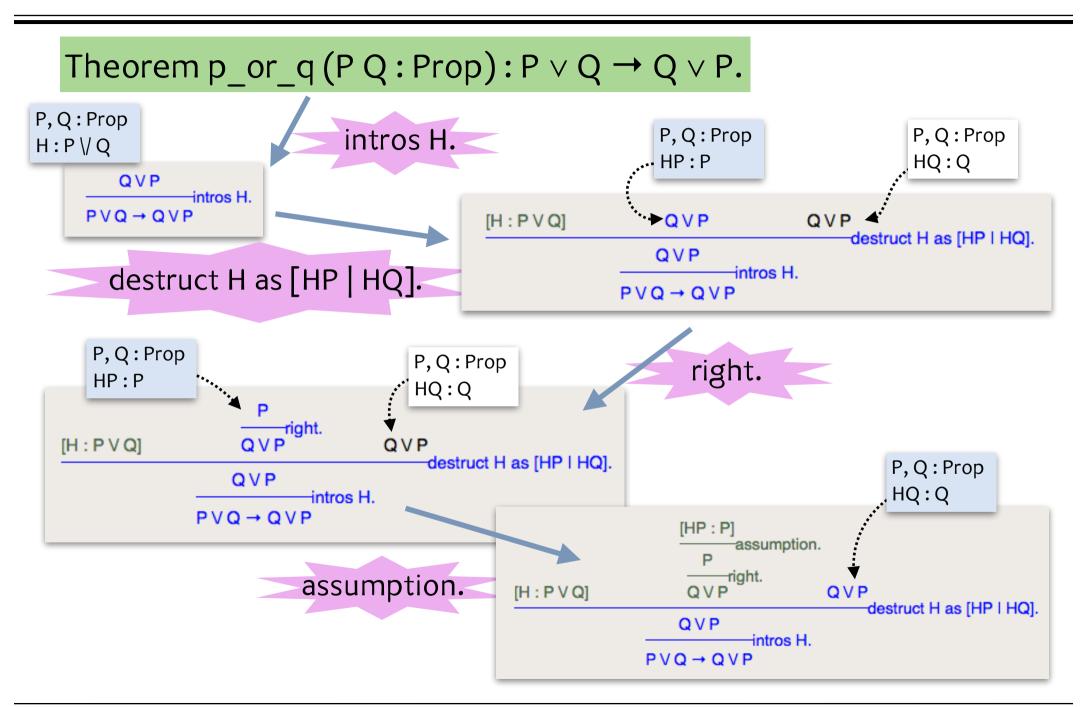
#### Using Traf: General structure



#### Using Traf: General structure



## Interactive construction of proof trees by Traf



```
Theorem pq_qp: forall P Q: Prop,
  P \/ Q -> Q \/ P.
Proof.
  intros P Q.
  intros H.
  destruct H as [HP | HQ].
  right. assumption.
  left. assumption.
Qed.
```

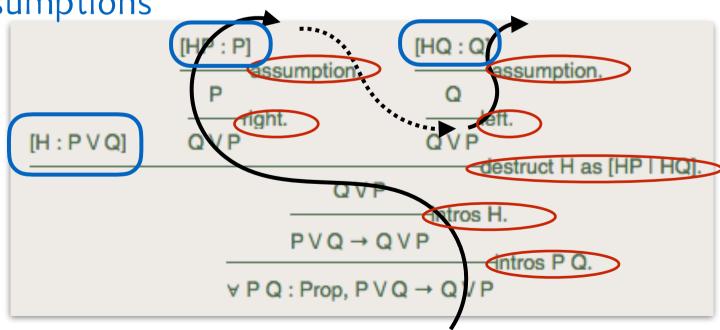
grows as if it is performing pre-order traversal (like depth-first search)

all tactics are shown

Traf puts used assumptions

with names

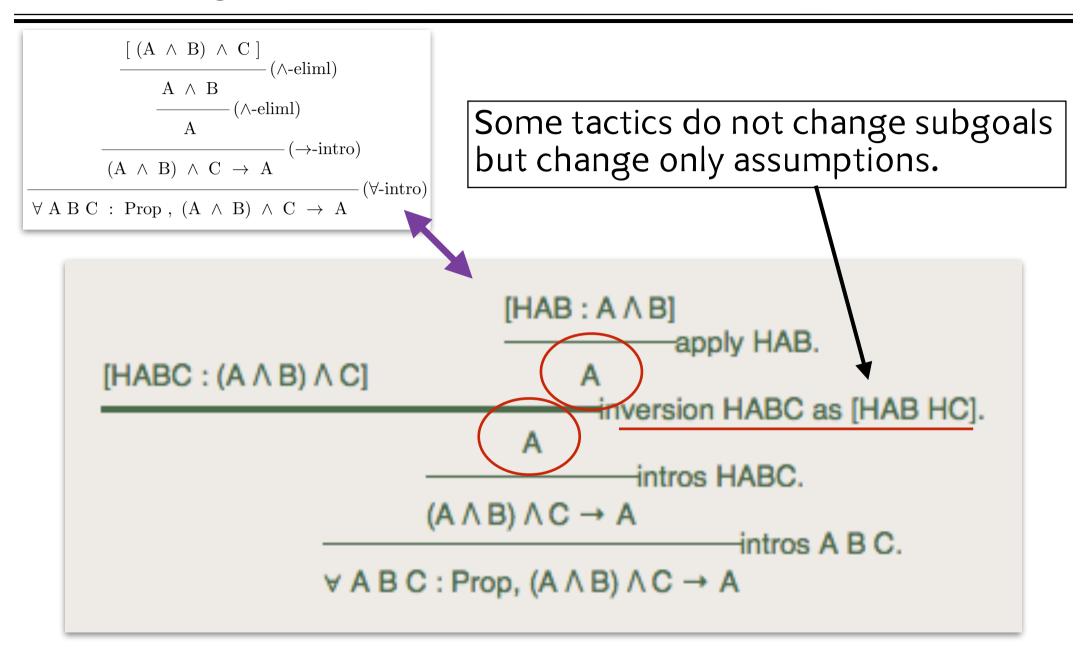
makes proof trees "readable"

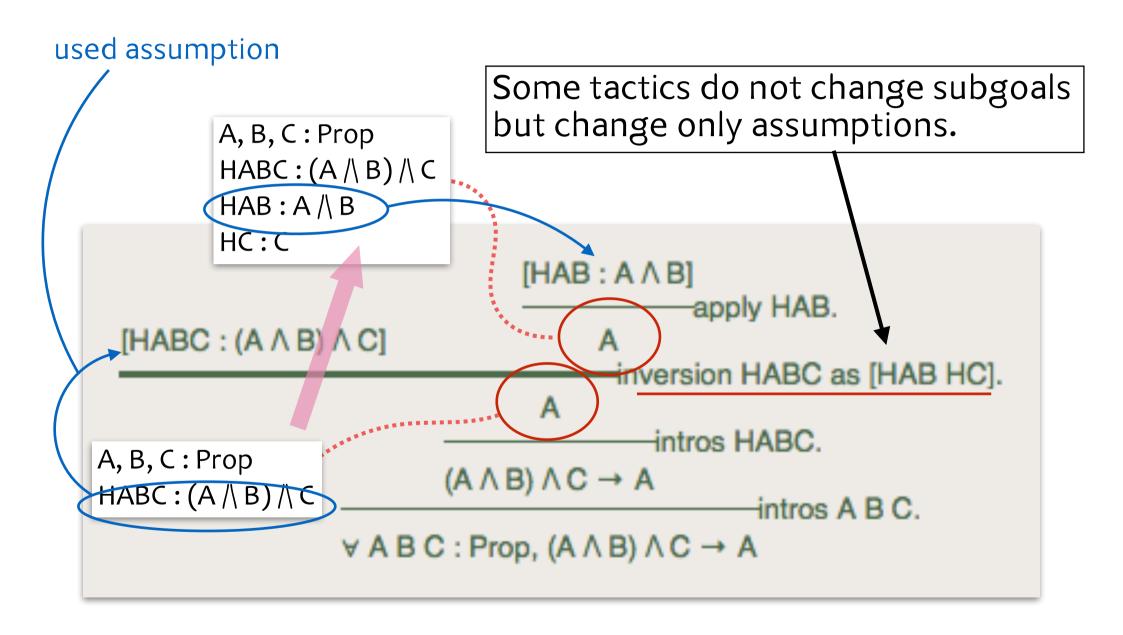


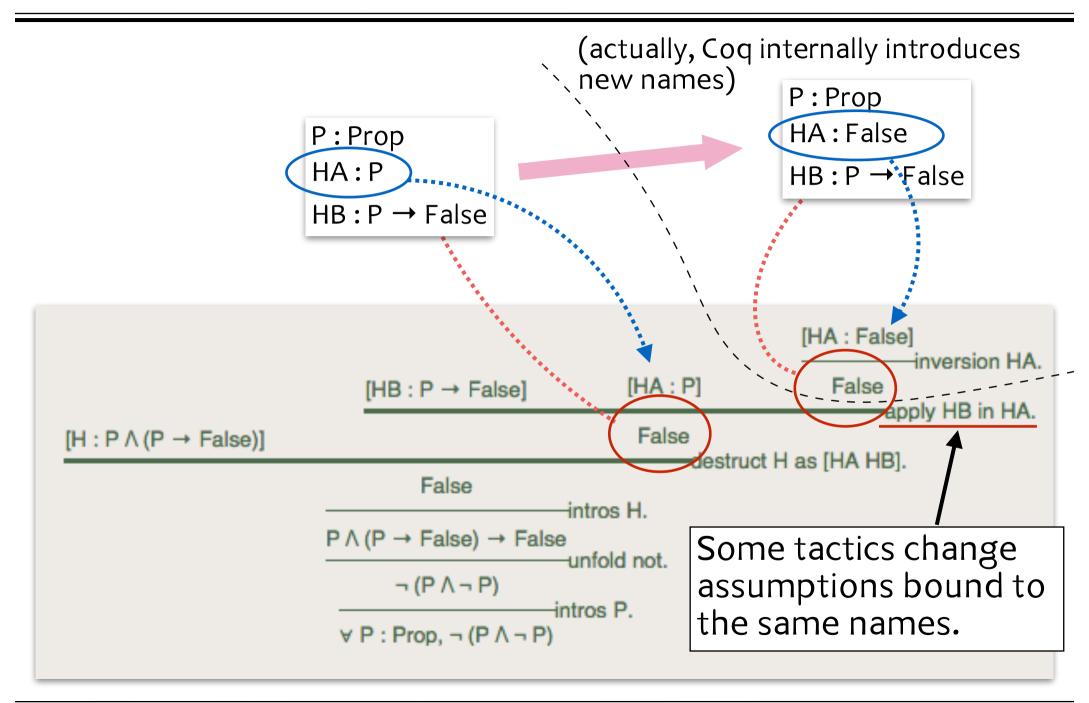
#### Another style of proof tree

```
\frac{\text{n: nat, IHn: } n + 0 = n + S \text{ n} = S \text{ n}}{\text{n: nat, IHn: } n + 0 = n + S \text{ (n + 0)} = S \text{ n}} \text{rewrite} \rightarrow \text{IHn.}
\frac{\text{n: nat, IHn: } n + 0 = n + S \text{ (n + 0)} = S \text{ n}}{\text{n: nat, IHn: } n + 0 = n + S \text{ n} + 0 = S \text{ n}} \text{simpl.}
\frac{\text{n: nat} + \text{n + 0} = n}{\text{n: nat, n + 0} = n} \text{induction n.}
```

expressing everything could make the tree less readable







Use of tacticals and/or automation

```
intros H; destruct H; auto.
PVQ → QVP
```

sometimes they could make proof trees needless ...

```
\forall P Q R : Prop, (P \rightarrow Q) \rightarrow (Q \rightarrow R) \rightarrow P \rightarrow R
```



```
(fun (P Q : Prop) (H : P \/ Q) =>
match H with
| or_introl HO => or_intror HO
| or_intror HO => or_introl HO
end)
```

In that case, corresponding proof object might be informative. But, ...

 $(fun (P Q R : Prop) (H : P \rightarrow Q) (HO : Q \rightarrow R) (H1 : P) \Rightarrow HO (H H1))$ 

```
Theorem ceval_deterministic': forall c st st1 st2,
      c / st \\ st1 ->
      c / st \\ st2 ->
                                                         in Auto.v, Software Foundations
      st1 = st2.
Proof.
  intros c st st1 st2 E1 E2;
  generalize dependent st2;
   induction E1:
               intros st2 E2; inv E2; repeat find eqn; try find rwinv; auto.
                                            (fun (c:com) (st st1 st2: state) (E1: c / st \\ st1) (E2: c / st \\ st2)
Oed.
       One-line proof script
                                             ceval ind
                                              (fun (st0: state) (c0: com) (st3: state) =>
                                              forall st4: state, c0 / st0 \ st4 -> st3 = st4)
                                              (fun (stO st3 : state) (E3 : SKIP / stO \\ st3) =>
  It depends.
                                              let H: st0 = st0 \rightarrow SKIP = SKIP \rightarrow st3 = st3 \rightarrow st0 = st3 :=
                                               match
                                                E3 in (c0 / s \\ s0)
  (Traf currently does not
                                                return (s = st0 \rightarrow c0 = SKIP \rightarrow s0 = st3 \rightarrow st0 = st3)
                                               with
  look at proof objects)
                                               | E Skip st4 =>
                                                 fun (H · c+4 - c+0) (HO · SKIP - SKIP) (H1 · c+4 - c+3) ->
                                                     Corresponding proof object
                                                     (> 2500 lines)
```

$$[pq:P \rightarrow Q] \qquad [qr:Q \rightarrow R] \\ by move /pq /qr. \\ \hline P \rightarrow R \\ \hline \hline move => P Q R pq qr. \\ \forall P Q R : Prop, (P \rightarrow Q) \rightarrow (Q \rightarrow R) \rightarrow P \rightarrow R$$

Traf recognizes used assumptions by analyzing commands syntactically.

Adapting Traf for another tactic library only requires an additional parser for commands.

#### Use of assert

```
(0+n) * m = n * m
\forall n m : nat, (0+n) * m = n * m
just an addition of a new subgoal
```

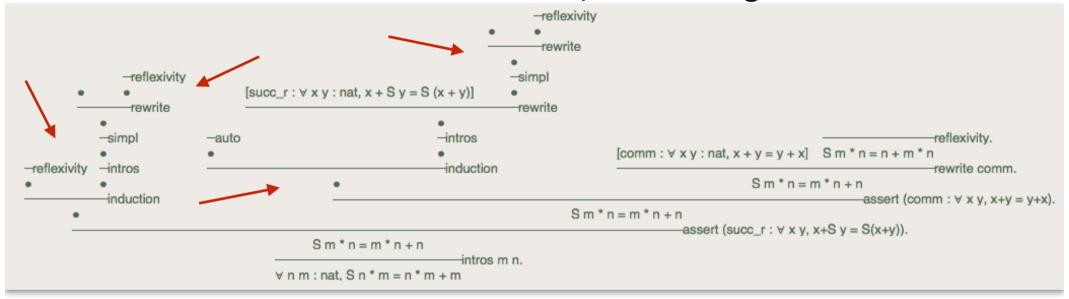
Use of special notations

Traf (currently) does not parse each formula

```
[HP:Pst]
                                                                                                                                 assumption.
                                                                               [Himp: P->> P']
                                                                                                                        P st
                                            [Hc:c/st\\st']
                                                               assumption.
                                                                                                                             apply Himp.
                                                c / st \\ st'
                                                                                                       P'st
[Hhoare : {{P'}} c {{Q}}}]
                                                                                                            -apply (Hhoare st st').
                                                         Q st<sup>t</sup>
                                                                   intros st st' Hc HP.
                                                    {{P}} c {{Q}}
                                                                                                        intros P P' Q c Hhoare Himp.
               \forall (P P'Q : Assertion) (c : com), \{\{P'\}\} c \{\{Q\}\}\} \rightarrow P ->> P' \rightarrow \{\{P\}\}\} c \{\{Q\}\}\}
```

#### What if the tree becomes large?

- You can shrink nodes that are not very interesting.



- Or you would like to shorten the proof itself by using tacticals.

#### What if the tree becomes very large?



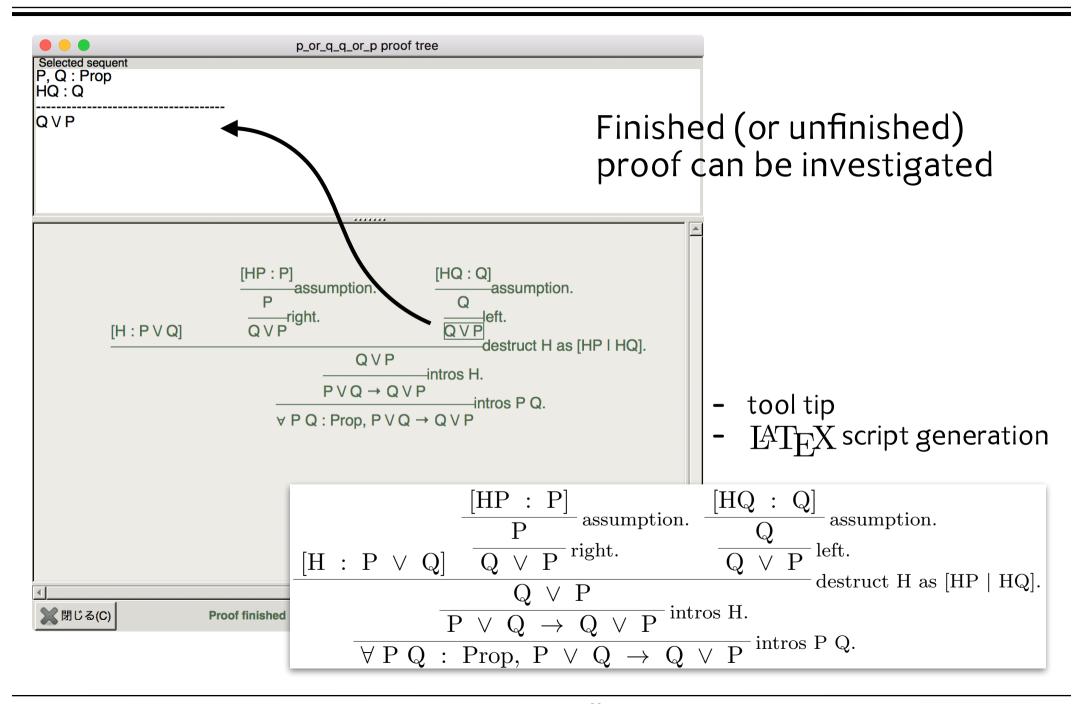
(scroll bars are available)

```
\forall (d:dcom) (P:Assertion), verification conditions P d \rightarrow {{P}} extract d {{post d}}.
 induction d; intros P H; simple in *.
|-- {{P}} SKIP {{a}}
     eapply hoare_consequence_pre.
   |-- {{?P'}} SKIP {{a}}
     apply hoare skip.
  \-- P ->> ?P'
      assumption.
|-- {{P}} extract d1;; extract d2 {{post d2}}
     inversion H as [H1 H2].
  |-- [H: verification conditions P d1 /\ verification conditions (post d1) d2]
  \-{{P}} extract d1;; extract d2{{post d2}}
       clear H.
     \--{{P}} extract d1;; extract d2 {{post d2}}
         eapply hoare seq.
         -- {{?Q}} extract d2 {{post d2}}
           apply IHd2.
           |---[[IHd2: ∀P: Assertion, verification conditions P d2 → {{P}}] extract d2 {{post d2}}]
           \-- verification conditions ?Q d2
             apply H2.
             \---- [H2: verification conditions (post d1) d2]
        \-- \{\{P\}\} extract d1 \{\{?Q\}\}
            apply IHd1.
           |---- [[IHd1: ∀ P: Assertion< verification conditions P d1 → {{P}} extract d1 {{post d1}}]
           \-- verification conditions P d1
```

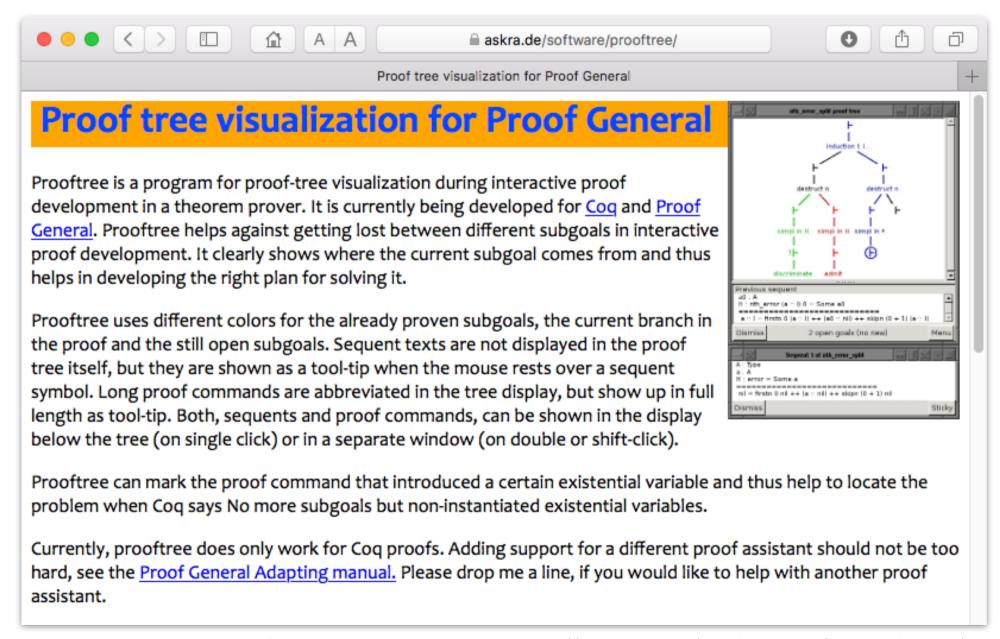
Turning the tree clockwise by 90 degrees would be preferable. (not supported yet)

#### Traf as a Proving Situation Monitor Emacs@mbp2017-541.local else S (double k). "Current Goal" mode intros n. induction n as [In' IH]. - exists 0. reflexivity. inversion IH as [k' IHk]. automatically shrinks irrelevant nodes rewrite -> evenb\_S. + simpl. exists k'. rewrite -> IHk. reflexivity. - keeps the tree's size modest + simpl. exists (S k'). rewrite -> IHk. reflexivity. 1 subgoal (ID 398) evenb\_double\_conv proof tree Current assumptions IH : exists k : nat, n' = S (double k) $IH : \exists k : nat, n' = S (double k)$ IHk : n' = S (double k') k' : nat S(S(double k')) = double(Sk')IHk : n' = S (double k')S(S(double k)) = double(S(k))U:%%- \*goals\* All (8,0) (Coa Goals) (Cog Response) -reflexivity S(S(double k')) = double(S(k'))-rewrite rewrite → IHk. S n' = double (S k')–∃ (S k'). --∃ k' ∃ k : nat. S n' = double k -simpl $\exists$ k : nat, S n' = (if negb false then double k else S (double k)) -destruct (evenb n'). $\exists$ k : nat, S n' = (if negb (evenb n') then double k else S (double k)) -reflexivity rewrite → evenb\_S. $\exists$ k : nat, S n' = (if evenb (S n') then double k else S (double k)) **-∃** 0 inversion IH as [k' IHk]. ∃ k : nat, S n' = (if evenb (S n') then double k else S (double k)) induction n as [ln' lH]. $\exists$ k : nat, n = (if evenb n then double k else S (double k)) intros n. $\forall$ n : nat, $\exists$ k : nat, n = (if evenb n then double k else S (double k)) 🤾 閉じる(C) Toggle drawing mode : [current goal] | Abbreviate subtree | Menu Retract to 635

#### Traf as a Proof Script Browser

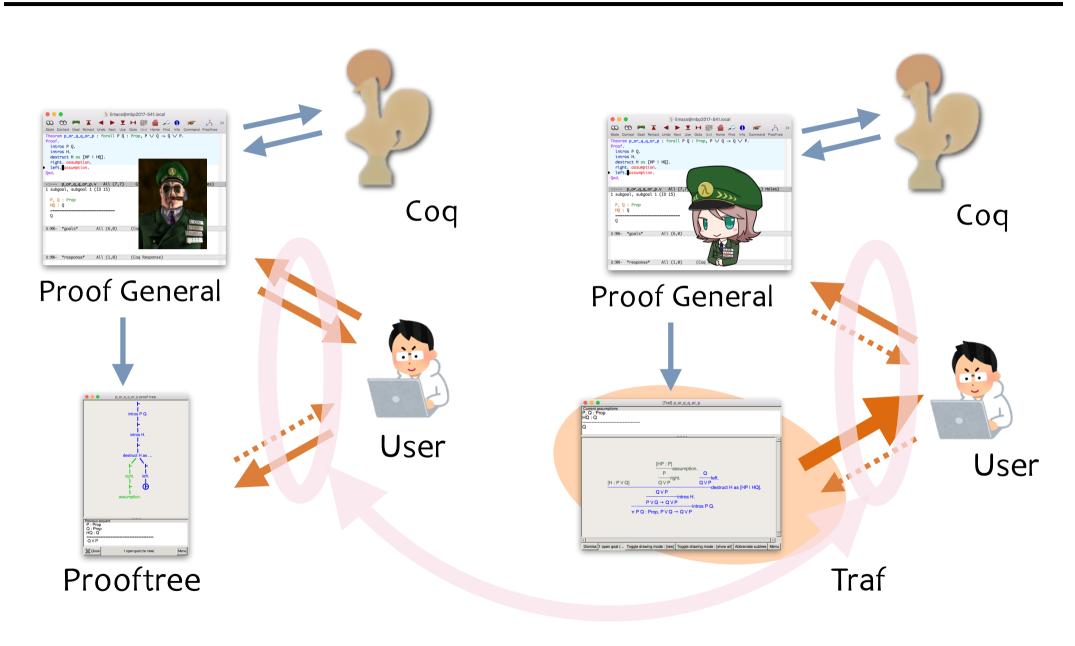


#### Traf is based on Prooftree



Hendrik Tews: Prooftree homepage https://askra.de/software/prooftree/

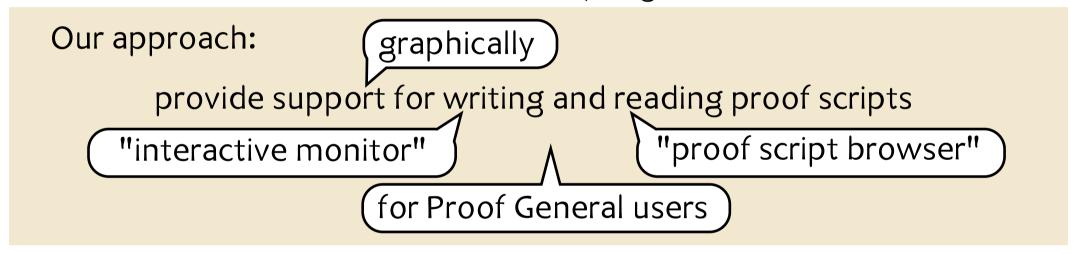
#### Comparison: Prooftree and Traf



Both are implemented in OCaml with LablGtk library.

# Traf: a Graphical Proof Tree Viewer Cooperating with Coq through Proof General

Available at https://github.com/hide-kawabata/traf



#### Future work

- support for large proof, long commands and/or tacticals, automation, ...
- support for referencing external lemmas, ...
- (evaluation based on user study)