Team notebook

${\bf HCMUS\text{-}IdentityImbalance}$

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P as LCA(u, v).

8	Misc	28	Each query provides two numbers u and v, ask for how many different
	8.1 Template Java	28	integers weight of nodes
	8.2 dates	29	there are on path from u to v.
	8.3 fraction	29	
	8.4 io	30	Modify DFS:
9	Number theory	30	For each node u, maintain the start and the end DFS time. Let's call them
	9.1 convolution	30	ST(u) and EN(u).
	9.2 crt	31	=> For each query, a node is considered if its occurrence count is one.
	9.3 diophantine equations	31	
	9.4 discrete logarithm	32	
	9.5 ext euclidean		Query solving:
	9.6 highest exponent factorial		
	9.7 miller rabin		Let's query be (u, v) . Assume that $ST(u) \le ST(v)$. Denotes P as $LCA(u, v)$
	9.8 mod integer		Case 1: P = u
	9.9 mod inv		Our query would be in range [ST(u), ST(v)].
	9.10 mod mul	33	041 4412) 10414 20 111 14160 [21(4), 21(1)]
	9.11 mod pow		Case 2: P != u
	9.12 number theoretic transform		Our query would be in range [EN(u), ST(v)] + [ST(p), ST(p)]
	9.13 pollard rho factorize		*/
	9.14 primes		
	9.15 totient sieve		<pre>void update(int &L, int &R, int qL, int qR){</pre>
	9.16 totient	36	<pre>while (L > qL) add(L); while (R < qR) add(++R);</pre>
10	0.04	9.0	while (n \ qn) add(++n),
10	Strings	36	<pre>while (L < qL) del(L++);</pre>
	10.1 Incremental Aho Corasick		while (R > qR) del(R);
	10.2 minimal string rotation		}
	10.3 suffix array		
	10.4 suffix automaton		<pre>vector <int> MoQueries(int n, vector <query> Q){</query></int></pre>
	10.5 z algorithm	39	block_size = sqrt((int)nodes.size());
			sort(Q.begin(), Q.end(), [] (const query &A, const query &B){
1	Algorithms		<pre>return (ST[A.1]/block_size != ST[B.1]/block_size)? (ST[A.1]/block_size < ST[B.1]/block_size) : (ST[A.r] <</pre>
_	1116011111111		ST[B.r]);
1.	1 Mo's algorithm on trees		<pre>});</pre>
	2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2		<pre>vector <int> res;</int></pre>
/			<pre>res.resize((int)Q.size());</pre>
/* 			
Pr	oblem:		LCA lca;
			<pre>lca.initialize(n);</pre>
ht	tps://www.spoj.com/problems/COT2/		int L = 1, R = 0;
~ .			1,,

Given a tree with N nodes and Q queries. Each node has an integer weight.

```
for(query q: Q){
   int u = q.1, v = q.r;
   if(ST[u] > ST[v]) swap(u, v); // assume that S[u] \le S[v]
   int parent = lca.get(u, v);
   if(parent == u){
       int qL = ST[u], qR = ST[v];
       update(L, R, qL, qR);
   }else{
       int qL = EN[u], qR = ST[v];
       update(L, R, qL, qR);
       if(cnt_val[a[parent]] == 0)
          res[q.pos] += 1;
   }
   res[q.pos] += cur_ans;
}
return res;
```

1.2 Mo's algorithm

```
res[q.pos] = calc(1, R-L+1);
}
return res;
}
```

2 DP Optimizations

2.1 convex hull trick

```
/**
 * Problems:
     http://codeforces.com/problemset/problem/319/C
     http://codeforces.com/contest/311/problem/B
     https://csacademy.com/contest/archive/task/squared-ends
     http://codeforces.com/contest/932/problem/F
struct line {
 long long m, b;
 line (long long a, long long c) : m(a), b(c) {}
 long long eval(long long x) {
   return m * x + b;
};
long double inter(line a, line b) {
  long double den = a.m - b.m;
 long double num = b.b - a.b;
 return num / den;
 * min m_i * x_j + b_i, for all i.
      x_j \le x_{j+1}
      m_i >= m_{i} + 1
struct ordered_cht {
 vector<line> ch;
 int idx; // id of last "best" in query
  ordered cht() {
   idx = 0:
```

```
void insert_line(long long m, long long b) {
   line cur(m, b);
   // new line's slope is less than all the previous
   while (ch.size() > 1 &&
      (inter(cur, ch[ch.size() - 2]) >= inter(cur, ch[ch.size() - 1]))) {
       // f(x) is better in interval [inter(ch.back(), cur), inf)
       ch.pop_back();
   }
   ch.push_back(cur);
 }
 long long eval(long long x) { // minimum
   // current x is greater than all the previous x,
   // if that is not the case we can make binary search.
   idx = min<int>(idx, ch.size() - 1);
   while (idx + 1 < (int)ch.size() \&\& ch[idx + 1].eval(x) <=
        ch[idx].eval(x))
     idx++;
   return ch[idx].eval(x);
 }
};
 * Dynammic convex hull trick
 * */
typedef long long int64;
typedef long double float128;
const int64 is_query = -(1LL<<62), inf = 1e18;</pre>
struct Line {
 int64 m, b;
 mutable function<const Line*()> succ;
 bool operator<(const Line& rhs) const {</pre>
   if (rhs.b != is_query) return m < rhs.m;</pre>
   const Line* s = succ();
   if (!s) return 0:
   int64 x = rhs.m;
   return b - s->b < (s->m - m) * x:
 }
```

```
};
struct HullDynamic : public multiset<Line> { // will maintain upper hull
    for maximum
  bool bad(iterator y) {
   auto z = next(y);
   if (y == begin()) {
     if (z == end()) return 0;
     return y->m == z->m && y->b <= z->b;
   auto x = prev(y);
   if (z == end()) return y->m == x->m && y->b <= x->b;
   return (float128)(x-b - y-b)*(z-m - y-m) >= (float128)(y-b - y-m)
        z->b)*(y->m - x->m);
  void insert_line(int64 m, int64 b) {
   auto y = insert({ m, b });
   y->succ = [=] { return next(y) == end() ? 0 : &*next(y); };
   if (bad(y)) { erase(y); return; }
   while (next(y) != end() && bad(next(y))) erase(next(y));
   while (y != begin() && bad(prev(y))) erase(prev(y));
  int64 eval(int64 x) {
   auto 1 = *lower_bound((Line) { x, is_query });
   return 1.m * x + 1.b:
};
```

2.2 divide and conquer

```
/**
 * recurrence:
 * dp[k][i] = min dp[k-1][j] + c[i][j - 1], for all j > i;
 *
 * "comp" computes dp[k][i] for all i in O(n log n) (k is fixed)
 *
 * Problems:
 * https://icpc.kattis.com/problems/branch
 * http://codeforces.com/contest/321/problem/E
 * */
void comp(int l, int r, int le, int re) {
```

```
if (1 > r) return;
int mid = (1 + r) >> 1;

int best = max(mid + 1, le);
dp[cur][mid] = dp[cur ^ 1][best] + cost(mid, best - 1);
for (int i = best; i <= re; i++) {
   if (dp[cur][mid] > dp[cur ^ 1][i] + cost(mid, i - 1)) {
     best = i;
     dp[cur][mid] = dp[cur ^ 1][i] + cost(mid, i - 1);
   }
}

comp(1, mid - 1, le, best);
comp(mid + 1, r, best, re);
}
```

2.3 dp on trees

```
* This trick is very useful when doing DP on trees, basically, you can
 * the answer for each node as if it was the root of the tree. Partial
 * are also stored in order to query subtrees (taking the root and
     exclude some
 * child).
 * problems:
 * - http://codeforces.com/gym/101161, problem I : Sky tax
 * - http://codeforces.com/contest/791/problem/D
 * */
struct edge {
 int to, p_id;
 edge (int a, int b) : to(a), p_id(b) {}
};
struct state {
 bool seen:
 long long missing;
 long long total;
```

```
vector<long long> partial;
  state() { clear(); }
  void clear() {
   seen = false:
   missing = 0;
   total = 0;
   partial.clear();
}:
void add_edge(int u, int v) {
 int id_u_v = g[u].size();
 int id_v_u = g[v].size();
  g[u].emplace_back(v, id_v_u); // id of the parent in the child's list
      (g[v][id] \rightarrow u)
 g[v].emplace_back(u, id_u_v); // id of the parent in the child's list
      (g[u][id] -> v)
int go(int node, int id_parent) {
  state &s = dp[node];
  if (!s.seen) {
   int ans = 1;
   s.partial.assign(g[node].size(), 0); // create the list of partial
   for (int i = 0; i < int(g[node].size()); i++) {</pre>
     int to = g[node][i].to;
     int pid = g[node][i].p_id;
     if (i != id_parent) {
       int tmp = go(to, pid);
       ans += tmp;
       s.partial[i] = tmp;
   }
   s.missing = id_parent;
   s.total = ans;
   s.seen = true;
   return ans;
```

```
if (s.missing == id_parent) { // the same id_parent than before, so
    we can not complete the results yet
    return s.total;
}

if (s.missing != -1) { // only one missing and is different of
        'id_parent'
    int tmp = go(g[node][s.missing].to, g[node][s.missing].p_id);
    s.partial[s.missing] = tmp;
    s.total += tmp;
    s.missing = -1;
}

int extra = (id_parent == -1) ? 0 : s.partial[id_parent];
    return s.total - extra;
}
```

3 Data structures

3.1 STL Treap

```
#include <ext/rope> //header with rope
using namespace std;
using namespace __gnu_cxx; //namespace with rope and some additional stuff
int main()
   ios_base::sync_with_stdio(false);
   rope <int> v; //use as usual STL container
   int n, m;
   cin >> n >> m;
   for(int i = 1; i <= n; ++i)</pre>
       v.push_back(i); //initialization
   int 1, r;
   for(int i = 0; i < m; ++i)</pre>
   {
       cin >> 1 >> r;
       --1, --r;
       rope \langle int \rangle cur = v.substr(1, r - 1 + 1);
       v.erase(1, r - 1 + 1);
```

```
v.insert(v.mutable_begin(), cur);
}
for(rope <int>::iterator it = v.mutable_begin(); it !=
    v.mutable_end(); ++it)
    cout << *it << " ";
return 0;
}</pre>
```

3.2 STL order statistics tree II

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace std;
using namespace __gnu_pbds;
typedef tree<int,null_type,less<int>,rb_tree_tag,
tree_order_statistics_node_update> order_set;
order_set X;
int get(int y) {
  int l=0,r=1e9+1;
  while(l<r) {</pre>
   int m=l+((r-l)>>1);
   if (m-X.order_of_key(m+1)<y)</pre>
     1=m+1:
   else
     r=m;
 return 1;
main(){
  ios::sync_with_stdio(0);
  cin.tie(0);
  int n,m;
  cin>>n>>m;
  for(int i=0;i<m;i++) {</pre>
   char a:
   int b;
```

```
cin>>a>>b;
    if(a=='L')
      cout<<get(b)<<endl;</pre>
    else
      X.insert(get(b));
  }
}
/***
Input
20 7
L 5
D 5
L 4
L 5
D 5
L 4
L 5
Output
5
6
***/
```

3.3 STL order statistics tree

```
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
#include <bits/stdc++.h>

using namespace __gnu_pbds;
using namespace std;

typedef
tree<
  pair<int,int>,
  null_type,
  less<pair<int,int>>,
  rb_tree_tag,
  tree_order_statistics_node_update>
```

```
ordered_set;
main()
 {
    ios::sync_with_stdio(0);
    cin.tie(0);
    int n;
    int sz=0;
    cin>>n;
    vector<int> ans(n,0);
    ordered_set t;
    int x,y;
    for(int i=0;i<n;i++)</pre>
        cin>>x>>y;
        ans[t.order_of_key({x,++sz})]++;
        t.insert({x,sz});
    for(int i=0;i<n;i++)</pre>
        cout<<ans[i]<<'\n';</pre>
}
/***
Input
5
1 1
5 1
7 1
3 3
5 5
Output
1
2
1
0
***/
```

3.4 dsu

```
class DSU{
public:
   vector <int> parent;
   void initialize(int n){
       parent.resize(n+1, -1);
   }
   int findSet(int u){
       while(parent[u] > 0)
           u = parent[u];
       return u;
   }
   void Union(int u, int v){
       int x = parent[u] + parent[v];
       if(parent[u] > parent[v]){
           parent[v] = x;
           parent[u] = v;
       }else{
           parent[u] = x;
           parent[v] = u;
   }
};
```

3.5 fenwick tree

```
template <typename T>
class FenwickTree{
  vector <T> fenw;
  int n;
public:
  void initialize(int _n){
    this->n = _n;
    fenw.resize(n+1);
}

void update(int id, T val) {
  while (id <= n) {
    fenw[id] += val;
    id += id&(-id);
  }
}</pre>
```

```
T get(int id){
  T ans{};
  while(id >= 1){
     ans += fenw[id];
     id -= id&(-id);
  }
  return ans;
}
```

3.6 hash table

```
/*
 * Micro hash table, can be used as a set.
 * Very efficient vs std::set
 *
 */

const int MN = 1001;
struct ht {
  int _s[(MN + 10) >> 5];
  int len;
  void set(int id) {
    len++;
    _s[id >> 5] |= (1LL << (id & 31));
  }
 bool is_set(int id) {
    return _s[id >> 5] & (1LL << (id & 31));
  }
};</pre>
```

3.7 heavy light decomposition

```
// Heavy-Light Decomposition
struct TreeDecomposition {
  vector<int> g[MAXN], c[MAXN];
  int s[MAXN]; // subtree size
  int p[MAXN]; // parent id
  int r[MAXN]; // chain root id
  int t[MAXN]; // index used in segtree/bit/...
```

```
int d[MAXN]; // depht
int ts;
void dfs(int v, int f) {
 p[v] = f;
  s[v] = 1;
  if (f != -1) d[v] = d[f] + 1;
  else d[v] = 0;
  for (int i = 0; i < g[v].size(); ++i) {</pre>
   int w = g[v][i];
   if (w != f) {
     dfs(w, v);
     s[v] += s[w];
 }
}
void hld(int v, int f, int k) {
 t[v] = ts++;
  c[k].push_back(v);
 r[v] = k;
  int x = 0, y = -1;
  for (int i = 0; i < g[v].size(); ++i) {</pre>
   int w = g[v][i];
   if (w != f) {
     if (s[w] > x) {
       x = s[w];
       y = w;
     }
   }
  }
  if (y != -1) {
   hld(y, v, k);
  for (int i = 0; i < g[v].size(); ++i) {</pre>
   int w = g[v][i];
   if (w != f && w != y) {
     hld(w, v, w);
   }
 }
}
```

```
void init(int n) {
   for (int i = 0; i < n; ++i) {
     g[i].clear();
   }
}

void add(int a, int b) {
   g[a].push_back(b);
   g[b].push_back(a);
}

void build() {
   ts = 0;
   dfs(0, -1);
   hld(0, 0, 0);
}
};</pre>
```

3.8 persistent array

```
struct node {
 node *1, *r;
 int val;
 node (int x) : 1(NULL), r(NULL), val(x) {}
 node () : 1(NULL), r(NULL), val(-1) {}
};
typedef node* pnode;
pnode update(pnode cur, int 1, int r, int at, int what) {
 pnode ans = new node();
 if (cur != NULL) {
   *ans = *cur;
 if (1 == r) {
   ans-> val = what;
   return ans;
 int m = (1 + r) >> 1;
 if (at <= m) ans-> 1 = update(ans-> 1, 1, m, at, what);
  else ans-> r = update(ans-> r, m + 1, r, at, what);
```

3.9 persistent seg tree

```
/* Problem: https://cses.fi/problemset/task/1737/
 * Your task is to maintain a list of arrays which initially has a single
     array. You have to process the following types of queries:
 * Query 1: Set the value a in array k to x.
 * Query 2: Calculate the sum of values in range [a,b] in array k.
 * Query 3: Create a copy of array k and add it to the end of the list.
 * Idea to create a persistent segment tree to save all version of array.
*/
vector <int> a;
struct Node{
   int val;
   Node *left, *right;
   Node(){
       left = right = NULL;
       val = 0;
   Node(Node* 1, Node *r, int v){
       left = 1;
       right = r;
       val = v;
   }
};
void build(Node* &cur, int 1, int r){
   if(1 == r){
       cur->val = a[1];
       return;
   }
```

```
int mid = (l+r) >> 1;
   cur->left = new Node();
   cur->right = new Node();
   build(cur->left, 1, mid);
   build(cur->right, mid+1, r);
   cur->val = cur->left->val + cur->right->val;
void update(Node* prev, Node* &cur, int 1, int r, int i, int val){
   if(i < 1 || r < i)
       return:
   if(1 == r && 1 == i){
       cur->val = val;
       return;
   int mid = (l+r) >> 1;
   if(i <= mid){</pre>
       cur->right = prev->right;
       cur->left = new Node();
       update(prev->left, cur->left, 1, mid, i, val);
   }else{
       cur->left = prev->left;
       cur->right = new Node();
       update(prev->right, cur->right, mid+1, r, i, val);
   cur->val = cur->left->val + cur->right->val;
}
int get(Node* cur, int 1, int r, int u, int v){
   if(v < 1 \mid | r < u)
       return 0;
   if(u <= 1 && r <= v){
       return cur->val;
   int mid = (l+r) >> 1;
   int L = get(cur->left, l, mid, u, v);
   int R = get(cur->right, mid+1, r, u, v);
   return L + R;
}
Node* ver[MAXN];
```

3.10 persistent trie

```
// both tries can be tested with the problem:
    http://codeforces.com/problemset/problem/916/D
// Persistent binary trie (BST for integers)
const int MD = 31;
struct node_bin {
 node_bin *child[2];
 int val:
 node_bin() : val(0) {
   child[0] = child[1] = NULL;
 }
};
typedef node_bin* pnode_bin;
pnode_bin copy_node(pnode_bin cur) {
 pnode_bin ans = new node_bin();
 if (cur) *ans = *cur;
 return ans;
}
pnode_bin modify(pnode_bin cur, int key, int inc, int id = MD) {
 pnode_bin ans = copy_node(cur);
 ans->val += inc;
 if (id >= 0) {
   int to = (key >> id) & 1;
   ans->child[to] = modify(ans->child[to], key, inc, id - 1);
 }
 return ans;
}
int sum_smaller(pnode_bin cur, int key, int id = MD) {
 if (cur == NULL) return 0;
 if (id < 0) return 0; // strictly smaller</pre>
 // if (id == - 1) return cur->val; // smaller or equal
 int ans = 0;
 int to = (key >> id) & 1;
 if (to) {
   if (cur->child[0]) ans += cur->child[0]->val;
   ans += sum_smaller(cur->child[1], key, id - 1);
 } else {
```

```
ans = sum_smaller(cur->child[0], key, id - 1);
 return ans;
// Persistent trie for strings.
const int MAX_CHILD = 26;
struct node {
 node *child[MAX_CHILD];
 int val;
 node() : val(-1) {
   for (int i = 0; i < MAX_CHILD; i++) {</pre>
     child[i] = NULL;
   }
 }
};
typedef node* pnode;
pnode copy_node(pnode cur) {
 pnode ans = new node();
 if (cur) *ans = *cur;
 return ans:
pnode set_val(pnode cur, string &key, int val, int id = 0) {
 pnode ans = copy_node(cur);
 if (id >= int(key.size())) {
   ans->val = val;
 } else {
   int t = key[id] - 'a';
   ans->child[t] = set_val(ans->child[t], key, val, id + 1);
 return ans;
pnode get(pnode cur, string &key, int id = 0) {
 if (id >= int(key.size()) || !cur)
   return cur;
 int t = key[id] - 'a';
 return get(cur->child[t], key, id + 1);
```

3.11 segment tree

```
// Problem:
    https://codeforces.com/edu/course/2/lesson/4/1/practice/contest/273169/problem/B
struct SegmentTree {
#define m ((1 + r) \gg 1)
#define lc (i << 1)
#define rc (i << 1 | 1)
   vector<int> mn;
   int n:
   SegmentTree(int n = 0) : n(n){
       mn.resize(4 * n + 1, 0);
   }
   SegmentTree(const vector<int> &a) : n(a.size()) {
       mn.resize(4 * n + 1, 0);
       function<void(int, int, int)> build = [&](int i, int 1, int r){
           if (1 == r){
              mn[i] = a[l - 1];
              return:
           }
           build(lc, l, m); build(rc, m + 1, r);
           mn[i] = min(mn[lc], mn[rc]);
       };
       build(1, 1, n);
   }
   void update(int i, int l, int r, int p, long val){
       if (1 == r){
           mn[i] = val;
           return;
       if (p <= m) update(lc, l, m, p, val);</pre>
       else update(rc, m + 1, r, p, val);
       mn[i] = min(mn[lc], mn[rc]);
   }
   int get(int i, int 1, int r, int u, int v){
       if (v < 1 || r < u) return INF;</pre>
       if (u <= 1 && r <= v) return mn[i];</pre>
       return min(get(lc, 1, m, u, v), get(rc, m + 1, r, u, v));
   }
   void update(int p, long val){
```

```
update(1, 1, n, p, val);
   int get(int 1, int r){
       return get(1, 1, n, 1, r);
#undef m
#undef lc
#undef rc
};
// Problem: There are two operations:
// 1 l r val: add the value val to the segment from 1 to r
// 2 l v: calculate the minimum of elements from l to r
struct LazySegmentTree {
#define m ((1 + r) >> 1)
#define lc (i << 1)
#define rc (i << 1 | 1)
   vector<int> mn, lazy;
   int n;
   LazySegmentTree(int n = 0) : n(n){
       mn.resize(4 * n + 1, 0);
       lazy.resize(4 * n + 1, 0);
   }
   void push(int i, int 1, int r){
       if (lazy[i] == 0) return;
       mn[i] += lazy[i];
       if (1 != r){
           lazy[lc] += lazy[i];
           lazy[rc] += lazy[i];
       }
       lazy[i] = 0;
   void update(int i, int l, int r, int u, int v, int val){
       push(i, 1, r);
       if (v < 1 || r < u) return;</pre>
       if (u <= 1 && r <= v){</pre>
           lazy[i] += val;
           push(i, 1, r);
           return;
       }
```

```
update(lc, l, m, u, v, val); update(rc, m + 1, r, u, v, val);
       mn[i] = min(mn[lc], mn[rc]);
   }
   int get(int i, int 1, int r, int u, int v){
       push(i, 1, r);
       if (v < 1 || r < u) return INF;</pre>
       if (u <= 1 && r <= v) return mn[i];</pre>
       return min(get(lc, 1, m, u, v), get(rc, m + 1, r, u, v));
   }
   void update(int 1, int r, int val){
       update(1, 1, n, 1, r, val);
   }
   int get(int 1, int r){
       return get(1, 1, n, l, r);
   }
#undef m
#undef lc
#undef rc
};
```

3.12 sparse table

```
template <typename T, typename func = function<T(const T, const T)>>
struct SparseTable {
   func calc;
   int n;
   vector<vector<T>> ans;

   SparseTable() {}

   SparseTable(const vector<T>& a, const func& f) : n(a.size()), calc(f)
      {
      int last = trunc(log2(n)) + 1;
      ans.resize(n);
      for (int i = 0; i < n; i++){
            ans[i].resize(last);
      }
      for (int i = 0; i < n; i++){
            ans[i][0] = a[i];
      }
}</pre>
```

3.13 splay tree

```
using namespace std;
#include <bits/stdc++.h>
#define D(x) cout<<x<<endl;</pre>
typedef int T;
struct node{
 node *left, *right, *parent;
 T kev;
 node (T k) : key(k), left(0), right(0), parent(0) {}
struct splay_tree{
 node *root;
 void right_rot(node *x) {
   node *p = x->parent;
   if (x->parent = p->parent) {
     if (x->parent->left == p) x->parent->left = x;
     if (x->parent->right == p) x->parent->right = x;
   }
   if (p->left = x->right) p->left->parent = p;
   x->right = p;
   p->parent = x;
```

```
void left_rot(node *x) {
 node *p = x->parent;
 if (x->parent = p->parent) {
   if (x->parent->left == p) x->parent->left = x;
   if (x->parent->right == p) x->parent->right = x;
 if (p->right = x->left) p->right->parent = p;
 x \rightarrow left = p;
 p->parent = x;
void splay(node *x, node *fa = 0) {
 while( x->parent != fa and x->parent != 0) {
   node *p = x->parent;
   if (p->parent == fa)
     if (p->right == x)
       left_rot(x);
     else
       right_rot(x);
   else {
     node *gp = p->parent; //grand parent
     if (gp->left == p)
       if (p->left == x)
         right_rot(x), right_rot(x);
         left_rot(x),right_rot(x);
     else
       if (p->left == x)
         right_rot(x), left_rot(x);
         left_rot(x), left_rot(x);
   }
 }
 if (fa == 0) root = x;
void insert(T key) {
 node *cur = root;
 node *pcur = 0;
 while (cur) {
   pcur = cur;
   if (key > cur->key) cur = cur->right;
   else cur = cur->left;
```

```
cur = new node(key);
cur->parent = pcur;
if (!pcur) root = cur;
else if (key > pcur->key ) pcur->right = cur;
else pcur->left = cur;
splay(cur);
}

node *find(T key) {
   node *cur = root;
   while (cur) {
      if (key > cur->key) cur = cur->right;
      else if(key < cur->key) cur = cur->left;
      else return cur;
   }
   return 0;
}

splay_tree(){ root = 0;};
};
```

3.14 trie

```
const int MN = 26; // size of alphabet
const int MS = 100010; // Number of states.

struct trie{
    struct node{
        int c;
        int a[MN];
    };

    node tree[MS];
    int nodes;

void clear(){
    tree[nodes].c = 0;
    memset(tree[nodes].a, -1, sizeof tree[nodes].a);
    nodes++;
}

void init(){
```

```
nodes = 0:
   clear();
 }
 int add(const string &s, bool query = 0){
   int cur_node = 0;
   for(int i = 0; i < s.size(); ++i){</pre>
     int id = gid(s[i]);
     if(tree[cur_node].a[id] == -1){
       if(query) return 0;
       tree[cur_node].a[id] = nodes;
       clear();
     cur_node = tree[cur_node].a[id];
   if(!query) tree[cur_node].c++;
   return tree[cur_node].c;
 }
};
```

3.15 wavelet tree

```
// this can be tested in the problem:
    http://www.spoj.com/problems/ILKQUERY/
struct wavelet {
 vector<int> values, ori;
 vector<int> map_left, map_right;
 int 1, r, m;
 wavelet *left, *right;
  wavelet() : left(NULL), right(NULL) {}
 wavelet(int a, int b, int c) : 1(a), r(b), m(c), left(NULL),
      right(NULL) {}
};
wavelet *init(vector<int> &data, vector<int> &ind, int lo, int hi) {
 if (lo > hi || (data.size() == 0)) return NULL;
 int mid = ((long long)(lo) + hi) / 2;
 if (lo + 1 == hi) mid = lo; // handle negative values
  wavelet *node = new wavelet(lo, hi, mid);
```

```
vector<int> data_l, data_r, ind_l, ind_r;
  int ls = 0, rs = 0;
  for (int i = 0; i < int(data.size()); i++) {</pre>
   int value = data[i];
   if (value <= mid) {</pre>
     data_1.emplace_back(value);
     ind_l.emplace_back(ind[i]);
     ls++:
   } else {
     data_r.emplace_back(value);
     ind_r.emplace_back(ind[i]);
     rs++;
   node->map_left.emplace_back(ls);
   node->map_right.emplace_back(rs);
   node->values.emplace_back(value);
   node->ori.emplace_back(ind[i]);
  if (lo < hi) {</pre>
   node->left = init(data_1, ind_1, lo, mid);
   node->right = init(data_r, ind_r, mid + 1, hi);
 return node;
int kth(wavelet *node, int to, int k) {
 // returns the kth element in the sorted version of (a[0], ..., a[to])
 if (node->1 == node->r) return node->m;
  int c = node->map_left[to];
 if (k < c)
   return kth(node->left, c - 1, k);
 return kth(node->right, node->map_right[to] - 1, k - c);
}
int pos_kth_ocurrence(wavelet *node, int val, int k) {
  // returns the position on the original array of the kth ocurrence of
      the value "val"
  if (!node) return -1;
  if (node->1 == node->r) {
   if (int(node->ori.size()) <= k)</pre>
     return -1;
   return node->ori[k]:
```

```
if (val <= node->m)
    return pos_kth_ocurrence(node->left, val, k);
return pos_kth_ocurrence(node->right, val, k);
}
```

4 Geometry

4.1 center 2 points + radious

```
vector<point> find_center(point a, point b, long double r) {
 point d = (a - b) * 0.5;
 if (d.dot(d) > r * r) {
   return vector<point> ();
 }
 point e = b + d;
 long double fac = sqrt(r * r - d.dot(d));
 vector<point> ans;
 point x = point(-d.y, d.x);
 long double 1 = sqrt(x.dot(x));
 x = x * (fac / 1);
 ans.push_back(e + x);
 x = point(d.y, -d.x);
 x = x * (fac / 1);
 ans.push_back(e + x);
 return ans;
}
```

4.2 closest pair problem

```
struct point {
  double x, y;
  int id;
  point() {}
  point (double a, double b) : x(a), y(b) {}
};

double dist(const point &o, const point &p) {
  double a = p.x - o.x, b = p.y - o.y;
  return sqrt(a * a + b * b);
```

```
}
double cp(vector<point> &p, vector<point> &x, vector<point> &y) {
  if (p.size() < 4) {</pre>
   double best = 1e100;
    for (int i = 0; i < p.size(); ++i)</pre>
     for (int j = i + 1; j < p.size(); ++j)</pre>
       best = min(best, dist(p[i], p[j]));
   return best;
  int ls = (p.size() + 1) >> 1;
  double l = (p[ls - 1].x + p[ls].x) * 0.5;
  vector<point> xl(ls), xr(p.size() - ls);
  unordered_set<int> left;
  for (int i = 0; i < ls; ++i) {</pre>
    xl[i] = x[i];
   left.insert(x[i].id);
  for (int i = ls; i < p.size(); ++i) {</pre>
   xr[i - ls] = x[i];
  vector<point> yl, yr;
  vector<point> pl, pr;
  yl.reserve(ls); yr.reserve(p.size() - ls);
  pl.reserve(ls); pr.reserve(p.size() - ls);
  for (int i = 0; i < p.size(); ++i) {</pre>
    if (left.count(y[i].id))
     yl.push_back(y[i]);
    else
     yr.push_back(y[i]);
    if (left.count(p[i].id))
     pl.push_back(p[i]);
    else
     pr.push_back(p[i]);
  double dl = cp(pl, xl, yl);
  double dr = cp(pr, xr, yr);
  double d = min(dl, dr);
  vector<point> yp; yp.reserve(p.size());
  for (int i = 0; i < p.size(); ++i) {</pre>
    if (fabs(y[i].x - 1) < d)
```

```
yp.push_back(y[i]);
 }
 for (int i = 0; i < yp.size(); ++i) {</pre>
   for (int j = i + 1; j < yp.size() && j < i + 7; ++j) {
     d = min(d, dist(yp[i], yp[j]));
    }
 }
 return d:
}
double closest_pair(vector<point> &p) {
 vector<point> x(p.begin(), p.end());
  sort(x.begin(), x.end(), [](const point &a, const point &b) {
   return a.x < b.x;</pre>
 });
  vector<point> y(p.begin(), p.end());
  sort(y.begin(), y.end(), [](const point &a, const point &b) {
   return a.y < b.y;</pre>
 }):
  return cp(p, x, y);
```

4.3 squares

```
typedef long double ld;

const ld eps = 1e-12;
int cmp(ld x, ld y = 0, ld tol = eps) {
    return ( x <= y + tol) ? (x + tol < y) ? -1 : 0 : 1;
}

struct point{
    ld x, y;
    point(ld a, ld b) : x(a), y(b) {}
    point() {}
};

struct square{
    ld x1, x2, y1, y2,
        a, b, c;
    point edges[4];
    square(ld _a, ld _b, ld _c) {</pre>
```

```
a = _a, b = _b, c = _c;
   x1 = a - c * 0.5;
   x2 = a + c * 0.5;
   v1 = b - c * 0.5;
   v2 = b + c * 0.5;
   edges[0] = point(x1, y1);
   edges[1] = point(x2, y1);
   edges[2] = point(x2, y2);
   edges[3] = point(x1, y2);
 }
}:
ld min_dist(point &a, point &b) {
 1d x = a.x - b.x,
    y = a.y - b.y;
 return sqrt(x * x + y * y);
bool point_in_box(square s1, point p) {
 if (cmp(s1.x1, p.x) != 1 && cmp(s1.x2, p.x) != -1 &&
     cmp(s1.y1, p.y) != 1 && cmp(s1.y2, p.y) != -1)
   return true;
 return false;
bool inside(square &s1, square &s2) {
 for (int i = 0; i < 4; ++i)
   if (point_in_box(s2, s1.edges[i]))
     return true;
 return false;
bool inside_vert(square &s1, square &s2) {
 if ((cmp(s1.y1, s2.y1) != -1 && cmp(s1.y1, s2.y2) != 1) ||
     (cmp(s1.y2, s2.y1) != -1 \&\& cmp(s1.y2, s2.y2) != 1))
   return true;
return false;
bool inside_hori(square &s1, square &s2) {
 if ((cmp(s1.x1, s2.x1) != -1 && cmp(s1.x1, s2.x2) != 1) ||
     (cmp(s1.x2, s2.x1) != -1 \&\& cmp(s1.x2, s2.x2) != 1))
   return true;
 return false;
```

```
}
ld min_dist(square &s1, square &s2) {
 if (inside(s1, s2) || inside(s2, s1))
   return 0;
 ld ans = 1e100;
 for (int i = 0; i < 4; ++i)
   for (int j = 0; j < 4; ++j)
     ans = min(ans, min_dist(s1.edges[i], s2.edges[j]));
 if (inside_hori(s1, s2) || inside_hori(s2, s1)) {
   if (cmp(s1.y1, s2.y2) != -1)
     ans = min(ans, s1.y1 - s2.y2);
   else
   if (cmp(s2.y1, s1.y2) != -1)
     ans = min(ans, s2.y1 - s1.y2);
 }
 if (inside_vert(s1, s2) || inside_vert(s2, s1)) {
   if (cmp(s1.x1, s2.x2) != -1)
     ans = min(ans, s1.x1 - s2.x2);
   else
   if (cmp(s2.x1, s1.x2) != -1)
     ans = min(ans, s2.x1 - s1.x2);
 }
 return ans;
```

4.4 triangles

Let a, b, c be length of the three sides of a triangle.

$$p = (a+b+c)*0.5$$

The inradius is defined by:

$$iR = \sqrt{\frac{(p-a)(p-b)(p-c)}{p}}$$

The radius of its circumcircle is given by the formula:

$$cR = \frac{abc}{\sqrt{(a+b+c)(a+b-c)(a+c-b)(b+c-a)}}$$

5 Graphs

5.1 SCC kosaraju

```
struct SCC {
 vector<vector<int> > g, gr;
 vector<bool> used;
 vector<int> order, component;
 int total_components;
 SCC(vector<vector<int> > &adj) {
   g = adj;
   int n = g.size();
   gr.resize(n);
   for (int i = 0; i < n; i++)</pre>
     for (auto to : g[i])
       gr[to].push_back(i);
   used.assign(n, false);
   for (int i = 0; i < n; i++)</pre>
     if (!used[i])
       GenTime(i);
   used.assign(n, false);
   component.assign(n, -1);
   total_components = 0;
   for (int i = n - 1; i >= 0; i--) {
     int v = order[i]:
     if (!used[v]) {
       vector<int> cur_component;
       Dfs(cur_component, v);
       for (auto node : cur_component)
         component[node] = total_components;
       total_components++;
 void GenTime(int node) {
   used[node] = true;
```

```
for (auto to : g[node])
     if (!used[to])
       GenTime(to);
   order.push_back(node);
  void Dfs(vector<int> &cur, int node) {
   used[node] = true;
   cur.push_back(node);
   for (auto to : gr[node])
     if (!used[to])
       Dfs(cur, to);
 }
 vector<vector<int>> CondensedGraph() {
   vector<vector<int>> ans(total_components);
   for (int i = 0; i < int(g.size()); i++) {</pre>
     for (int to : g[i]) {
       int u = component[i], v = component[to];
       if (u != v)
         ans[u].push_back(v);
     }
   }
   return ans;
 }
};
```

5.2 board

```
struct board {
  int n, m, r;
  board(int a, int b, int c = 1) : n(a), m(b), r(c) {}

long long frec(int x, int y) {
    // returns how many squares of r x r contain the cell (x, y)
    long long a = min(x, n - r) - max(x - r + 1, 0) + 1;
    long long b = min(y, m - r) - max(y - r + 1, 0) + 1;
    return a * b;
}

bool valid(int x, int y) {
    return x >= 0 && x < n && y >= 0 && y < m;
}</pre>
```

};

5.3 bridges

```
struct Graph {
 vector<vector<Edge>> g;
 vector<int> vi, low, d, pi, is_b;
 int bridges_computed;
 int ticks, edges;
 Graph(int n, int m) {
   g.assign(n, vector<Edge>());
   is_b.assign(m, 0);
   vi.resize(n);
   low.resize(n):
   d.resize(n);
   pi.resize(n);
   edges = 0;
   bridges_computed = 0;
 void AddEdge(int u, int v) {
   g[u].push_back(Edge(v, edges));
   g[v].push_back(Edge(u, edges));
   edges++;
 void Dfs(int u) {
   vi[u] = true;
   d[u] = low[u] = ticks++;
   for (int i = 0; i < (int)g[u].size(); ++i) {</pre>
     int v = g[u][i].to;
     if (v == pi[u]) continue;
     if (!vi[v]) {
       pi[v] = u;
       Dfs(v);
       if (d[u] < low[v]) is_b[g[u][i].id] = true;</pre>
       low[u] = min(low[u], low[v]);
     } else {
       low[u] = min(low[u], d[v]);
```

```
}
 }
 // Multiple edges from a to b are not allowed.
 // (they could be detected as a bridge).
 // If you need to handle this, just count
 // how many edges there are from a to b.
 void CompBridges() {
   fill(pi.begin(), pi.end(), -1);
   fill(vi.begin(), vi.end(), 0);
   fill(low.begin(), low.end(), 0);
   fill(d.begin(), d.end(), 0);
   ticks = 0;
   for (int i = 0; i < (int)g.size(); ++i)</pre>
     if (!vi[i]) Dfs(i);
   bridges_computed = true;
 map<int, vector<Edge>> BridgesTree() {
   if (!bridges_computed) CompBridges();
   int n = g.size();
   Dsu dsu(g.size());
   for (int i = 0; i < n; i++)</pre>
     for (auto e : g[i])
       if (!is_b[e.id]) dsu.Join(i, e.to);
   map<int, vector<Edge>> tree;
   for (int i = 0; i < n; i++)</pre>
     for (auto e : g[i])
       if (is_b[e.id])
         tree[dsu.Find(i)].emplace_back(dsu.Find(e.to), e.id);
   return tree;
 }
};
```

5.4 directed mst

```
const int inf = 1000000 + 10;
struct edge {
  int u, v, w;
  edge() {}
```

```
edge(int a,int b,int c) : u(a), v(b), w(c) {}
};
 * Computes the minimum spanning tree for a directed graph
 * - edges : Graph description in the form of list of edges.
 * each edge is: From node u to node v with cost w
 * - root : Id of the node to start the DMST.
 * - n : Number of nodes in the graph.
* */
int dmst(vector<edge> &edges, int root, int n) {
  int ans = 0;
  int cur_nodes = n;
  while (true) {
   vector<int> lo(cur_nodes, inf), pi(cur_nodes, inf);
   for (int i = 0; i < edges.size(); ++i) {</pre>
     int u = edges[i].u, v = edges[i].v, w = edges[i].w;
     if (w < lo[v] and u != v) {
       lo[v] = w;
       pi[v] = u;
   }
   lo[root] = 0;
   for (int i = 0; i < lo.size(); ++i) {</pre>
     if (i == root) continue:
     if (lo[i] == inf) return -1;
   }
   int cur_id = 0;
   vector<int> id(cur_nodes, -1), mark(cur_nodes, -1);
   for (int i = 0; i < cur_nodes; ++i) {</pre>
     ans += lo[i];
     int u = i:
     while (u != root and id[u] < 0 and mark[u] != i) {</pre>
       mark[u] = i;
       u = pi[u];
     if (u != root and id[u] < 0) { // Cycle}
        for (int v = pi[u]; v != u; v = pi[v])
          id[v] = cur_id;
        id[u] = cur_id++;
```

```
if (cur_id == 0)
    break;

for (int i = 0; i < cur_nodes; ++i)
    if (id[i] < 0) id[i] = cur_id++;

for (int i = 0; i < edges.size(); ++i) {
    int u = edges[i].u, v = edges[i].v, w = edges[i].w;
    edges[i].u = id[u];
    edges[i].v = id[v];
    if (id[u] != id[v])
        edges[i].w -= lo[v];
}

cur_nodes = cur_id;
root = id[root];
}

return ans;</pre>
```

5.5 eulerian path

}

```
// Taken from https://github.com/lbv/pc-code/blob/master/code/graph.cpp
// Eulerian Trail
struct Euler {
 ELV adj; IV t;
 Euler(ELV Adj) : adj(Adj) {}
 void build(int u) {
   while(! adj[u].empty()) {
     int v = adj[u].front().v;
     adj[u].erase(adj[u].begin());
     build(v);
   t.push_back(u);
 }
};
bool eulerian_trail(IV &trail) {
 Euler e(adj);
 int odd = 0, s = 0;
    for (int v = 0; v < n; v++) {
    int diff = abs(in[v] - out[v]);
```

```
if (diff > 1) return false;
if (diff == 1) {
   if (++odd > 2) return false;
   if (out[v] > in[v]) start = v;
   }
   }
   */
e.build(s);
reverse(e.t.begin(), e.t.end());
trail = e.t;
return true;
}
```

5.6 karp min mean cycle

```
/**
 * Finds the min mean cycle, if you need the max mean cycle
 * just add all the edges with negative cost and print
 * ans * -1
 * test: uva, 11090 - Going in Cycle!!
 * */
const int MN = 1000;
struct edge{
 int v;
 long long w;
  edge(){} edge(int v, int w) : v(v), w(w) {}
};
long long d[MN][MN];
// This is a copy of g because increments the size
// pass as reference if this does not matter.
int karp(vector<vector<edge> > g) {
 int n = g.size();
  g.resize(n + 1); // this is important
 for (int i = 0; i < n; ++i)
   if (!g[i].empty())
     g[n].push_back(edge(i,0));
  ++n:
```

```
for(int i = 0:i < n:++i)
 fill(d[i],d[i]+(n+1),INT_MAX);
d[n - 1][0] = 0;
for (int k = 1: k \le n: ++k) for (int u = 0: u \le n: ++u) {
  if (d[u][k - 1] == INT_MAX) continue;
 for (int i = g[u].size() - 1; i >= 0; --i)
   d[g[u][i].v][k] = min(d[g[u][i].v][k], d[u][k-1] + g[u][i].w);
}
bool flag = true;
for (int i = 0; i < n && flag; ++i)</pre>
  if (d[i][n] != INT_MAX)
   flag = false;
if (flag) {
  return true: // return true if there is no a cycle.
double ans = 1e15;
for (int u = 0: u + 1 < n: ++u) {
  if (d[u][n] == INT_MAX) continue;
  double W = -1e15:
  for (int k = 0; k < n; ++k)
   if (d[u][k] != INT_MAX)
     W = max(W, (double)(d[u][n] - d[u][k]) / (n - k));
  ans = min(ans, W):
}
// printf("%.21f\n", ans);
cout << fixed << setprecision(2) << ans << endl;</pre>
return false;
```

5.7 konig's theorem

}

In any bipartite graph, the number of edges in a maximum matching equals the number of vertices in a minimum vertex cover

5.8 minimum path cover in DAG

Given a directed acyclic graph G = (V, E), we are to find the minimum number of vertex-disjoint paths to cover each vertex in V.

We can construct a bipartite graph $G' = (Vout \cup Vin, E')$ from G, where :

$$Vout = \{v \in V : v \text{ has positive out} - degree\}$$

$$Vin = \{v \in V : v \text{ has positive in} - degree\}$$

$$E' = \{(u, v) \in Vout \times Vin : (u, v) \in E\}$$

Then it can be shown, via König's theorem, that G' has a matching of size m if and only if there exists n-m vertex-disjoint paths that cover each vertex in G, where n is the number of vertices in G and m is the maximum cardinality bipartite mathching in G'.

Therefore, the problem can be solved by finding the maximum cardinality matching in G' instead.

NOTE: If the paths are note necessarily disjoints, find the transitive closure and solve the problem for disjoint paths.

5.9 planar graph (euler)

Euler's formula states that if a finite, connected, planar graph is drawn in the plane without any edge intersections, and v is the number of vertices, e is the number of edges and f is the number of faces (regions bounded by edges, including the outer, infinitely large region), then:

$$f + v = e + 2$$

It can be extended to non connected planar graphs with c connected components:

$$f + v = e + c + 1$$

5.10 query with lca

```
struct lowest_ca {
  int T[MN], L[MN], W[MN];
  int P[MN][ML], MI[MN][ML], MA[MN][ML];

void dfs(vector<vector<edge> > &g, int root, int pi = -1) {
  if (pi == -1) {
```

```
L[root] = W[root] = 0;
   T[root] = -1;
  }
  for (int i = 0; i < (int)g[root].size(); ++i) {</pre>
   int to = g[root][i].v;
   if (to != pi) {
     T[to] = root;
     W[to] = g[root][i].w;
     L[to] = L[root] + 1;
     dfs(g, to, root);
   }
 }
}
void init(vector<vector<edge> > &g, int root) {
  // g is undirected
  dfs(g, root);
  int N = g.size(), i, j;
  for (i = 0; i < N; i++) {</pre>
   for (j = 0; 1 << j < N; j++) {
     P[i][j] = -1;
     MI[i][j] = inf;
   }
  }
  for (i = 0; i < N; i++) {</pre>
   P[i][0] = T[i];
   MI[i][0] = W[i];
  }
  for (j = 1; 1 << j < N; j++)
   for (i = 0; i < N; i++)</pre>
     if (P[i][i - 1] != -1) {
       P[i][j] = P[P[i][j-1]][j-1];
       MI[i][j] = min(MI[i][j-1], MI[P[i][j-1]][j-1]);
     }
}
int query(int p, int q) {
  int tmp, log, i;
  int mmin = inf;
  if (L[p] < L[q])
    tmp = p, p = q, q = tmp;
```

```
for (log = 1; 1 << log <= L[p]; log++);</pre>
 log--;
 for (i = log; i >= 0; i--)
   if (L[p] - (1 << i) >= L[q]) {
     mmin = min(mmin, MI[p][i]);
     p = P[p][i];
   }
 if (p == q)
   // return p;
   return mmin;
 for (i = log; i >= 0; i--)
   if (P[p][i] != -1 && P[p][i] != P[q][i]) {
     mmin = min(mmin, min(MI[p][i], MI[q][i]));
     p = P[p][i], q = P[q][i];
 // return T[p];
 return min(mmin, min(MI[p][0], MI[q][0]));
int get_child(int p, int q) { // p is ancestor of q
 if (p == q) return -1;
 int i, log;
 for (log = 1; 1 << log <= L[q]; log++) {}
 log--;
 for (i = log; i >= 0; i--)
   if (L[q] - (1 << i) > L[p]) {
     q = P[q][i];
 assert(P[q][0] == p);
 return q;
int is_ancestor(int p, int q) {
 if (L[p] >= L[q])
   return false;
 int dist = L[q] - L[p];
```

```
int cur = q;
int step = 0;
while (dist) {
   if (dist & 1)
      cur = P[cur][step];
   step++;
   dist >>= 1;
}
return cur == p;
}
};
```

5.11 tarjan scc

```
const int MN = 20002;
struct tarjan_scc {
 int scc[MN], low[MN], d[MN], stacked[MN];
 int ticks, current_scc;
 deque<int> s; // used as stack.
 tarjan_scc() {}
 void init () {
   memset(scc, -1, sizeof scc);
   memset(d, -1, sizeof d);
   memset(stacked, 0, sizeof stacked);
   s.clear();
   ticks = current_scc = 0;
 }
 void compute(vector<vector<int> > &g, int u) {
   d[u] = low[u] = ticks++;
   s.push_back(u);
   stacked[u] = true;
   for (int i = 0; i < g[u].size(); ++i) {</pre>
     int v = g[u][i];
     if (d[v] == -1)
       compute(g, v);
     if (stacked[v]) {
       low[u] = min(low[u], low[v]);
```

```
}
}

if (d[u] == low[u]) { // root
    int v;
    do {
       v = s.back();s.pop_back();
       stacked[v] = false;
       scc[v] = current_scc;
    } while (u != v);
    current_scc++;
}
}
```

5.12 two sat (with kosaraju)

```
* Given a set of clauses (a1 v a2)^(a2 v a3)....
* this algorithm find a solution to it set of clauses.
* test: http://lightoj.com/volume_showproblem.php?problem=1251
#include<bits/stdc++.h>
using namespace std;
#define MAX 100000
#define endl '\n'
vector<int> G[MAX];
vector<int> GT[MAX];
vector<int> Ftime;
vector<vector<int> > SCC;
bool visited[MAX];
int n;
void dfs1(int n){
 visited[n] = 1;
 for (int i = 0; i < G[n].size(); ++i) {</pre>
   int curr = G[n][i]:
   if (visited[curr]) continue;
   dfs1(curr);
```

```
}
 Ftime.push back(n):
}
void dfs2(int n, vector<int> &scc) {
 visited[n] = 1;
 scc.push_back(n);
 for (int i = 0;i < GT[n].size(); ++i) {</pre>
   int curr = GT[n][i]:
   if (visited[curr]) continue;
   dfs2(curr, scc);
 }
}
void kosaraju() {
 memset(visited, 0, sizeof visited);
 for (int i = 0; i < 2 * n; ++i) {</pre>
   if (!visited[i]) dfs1(i);
 }
 memset(visited, 0, sizeof visited);
 for (int i = Ftime.size() - 1; i >= 0; i--) {
   if (visited[Ftime[i]]) continue;
   vector<int> _scc;
   dfs2(Ftime[i],_scc);
   SCC.push_back(_scc);
 }
}
* After having the SCC, we must traverse each scc, if in one SCC are -b
     y b, there is not a solution.
 * Otherwise we build a solution, making the first "node" that we find
     truth and its complement false.
 **/
bool two_sat(vector<int> &val) {
 kosaraju();
 for (int i = 0; i < SCC.size(); ++i) {</pre>
   vector<bool> tmpvisited(2 * n, false);
```

```
for (int j = 0; j < SCC[i].size(); ++j) {</pre>
     if (tmpvisited[SCC[i][j] ^ 1]) return 0;
     if (val[SCC[i][j]] != -1) continue;
     else {
       val[SCC[i][j]] = 0;
       val[SCC[i][j] ^ 1] = 1;
     tmpvisited[SCC[i][j]] = 1;
   }
 }
 return 1:
// Example of use
int main() {
 int m, u, v, nc = 0, t; cin >> t;
 // n = "nodes" number. m = clauses number
 while (t--) {
   cin >> m >> n;
   Ftime.clear();
   SCC.clear():
   for (int i = 0; i < 2 * n; ++i) {</pre>
     G[i].clear();
     GT[i].clear();
   // (a1 v a2) = (a1 -> a2) = (a2 -> a1)
   for (int i = 0; i < m ; ++i) {</pre>
     cin >> u >> v:
     int t1 = abs(u) - 1;
     int t2 = abs(v) - 1;
     int p = t1 * 2 + ((u < 0)? 1 : 0);
     int q = t2 * 2 + ((v < 0)? 1 : 0);
     G[p ^ 1].push_back(q);
     G[q ^ 1].push_back(p);
     GT[p].push_back(q ^ 1);
     GT[q].push_back(p ^ 1);
   vector < int > val(2 * n, -1);
   cout << "Case " << ++nc <<": ";
   if (two_sat(val)) {
```

```
cout << "Yes" << endl;
vector<int> sol;
for (int i = 0; i < 2 * n; ++i)
    if (i % 2 == 0 and val[i] == 1)
        sol.push_back(i / 2 + 1);
cout << sol.size();

for (int i = 0; i < sol.size(); ++i) {
    cout << " " << sol[i];
}
    cout << endl;
} else {
    cout << "No" << endl;
}
return 0;</pre>
```

6 Math

6.1 Lucas theorem

For non-negative integers m and n and a prime p, the following congruence relation holds: :

$$\binom{m}{n} \equiv \prod_{i=0}^{k} \binom{m_i}{n_i} \pmod{p},$$

where:

$$m = m_k p^k + m_{k-1} p^{k-1} + \dots + m_1 p + m_0,$$

and:

$$n = n_k p^k + n_{k-1} p^{k-1} + \dots + n_1 p + n_0$$

are the base p expansions of m and n respectively. This uses the convention that $\binom{m}{n} = 0$ if $m \le n$.

6.2 cumulative sum of divisors

```
/*
The function SOD(n) (sum of divisors) is defined
as the summation of all the actual divisors of
an integer number n. For example,
```

```
SOD(24) = 2+3+4+6+8+12 = 35.
The function CSOD(n) (cumulative SOD) of an integer n, is defined as below:
    csod(n) = \sum_{{i = 1}^{n}} sod(i)

It can be computed in O(sqrt(n)):
*/
long long csod(long long n) {
    long long ans = 0;
    for (long long i = 2; i * i <= n; ++i) {
        long long j = n / i;
        ans += (i + j) * (j - i + 1) / 2;
        ans += i * (j - i);
    }
    return ans;
}</pre>
```

6.3 fft

```
/**
 * Fast Fourier Transform.
 * Useful to compute convolutions.
 * computes:
 * C(f star g)[n] = sum_m(f[m] * g[n - m])
 * for all n.
 * test: icpc live archive, 6886 - Golf Bot
 * */

using namespace std;
#include <bits/stdc++.h>
#define D(x) cout << #x " = " << (x) << endl
#define endl '\n'

const int MN = 262144 << 1;
int d[MN + 10], d2[MN + 10];

const double PI = acos(-1.0);</pre>
```

```
struct cpx {
 double real, image;
 cpx(double _real, double _image) {
   real = _real;
   image = _image;
 }
 cpx(){}
}:
cpx operator + (const cpx &c1, const cpx &c2) {
 return cpx(c1.real + c2.real, c1.image + c2.image);
}
cpx operator - (const cpx &c1, const cpx &c2) {
 return cpx(c1.real - c2.real, c1.image - c2.image);
}
cpx operator * (const cpx &c1, const cpx &c2) {
 return cpx(c1.real*c2.real - c1.image*c2.image, c1.real*c2.image +
      c1.image*c2.real);
}
int rev(int id, int len) {
 int ret = 0:
 for (int i = 0; (1 << i) < len; i++) {
   ret <<= 1:
   if (id & (1 << i)) ret |= 1;</pre>
 }
 return ret;
}
cpx A[1 << 20];
void FFT(cpx *a, int len, int DFT) {
 for (int i = 0; i < len; i++)</pre>
   A[rev(i, len)] = a[i];
 for (int s = 1; (1 << s) <= len; s++) {
   int m = (1 << s);
   cpx wm = cpx(cos(DFT * 2 * PI / m), sin(DFT * 2 * PI / m));
   for(int k = 0; k < len; k += m) {
     cpx w = cpx(1, 0);
     for(int j = 0; j < (m >> 1); j++) {
       cpx t = w * A[k + j + (m >> 1)];
       cpx u = A[k + j];
       A[k + j] = u + t;
```

```
A[k + j + (m >> 1)] = u - t;
   }
 if (DFT == -1) for (int i = 0; i < len; i++) A[i].real /= len,
      A[i].image /= len;
 for (int i = 0; i < len; i++) a[i] = A[i];</pre>
 return:
cpx in[1 << 20];
void solve(int n) {
 memset(d, 0, sizeof d);
 for (int i = 0; i < n; ++i) {
   cin >> t;
   d[t] = true:
 int m;
  cin >> m;
  vector<int> q(m);
 for (int i = 0; i < m; ++i)</pre>
   cin >> q[i];
 for (int i = 0; i < MN; ++i) {</pre>
   if (d[i])
     in[i] = cpx(1, 0);
   else
     in[i] = cpx(0, 0);
 FFT(in, MN, 1);
 for (int i = 0; i < MN; ++i) {</pre>
   in[i] = in[i] * in[i];
 FFT(in, MN, -1);
 int ans = 0;
 for (int i = 0; i < q.size(); ++i) {</pre>
   if (in[q[i]].real > 0.5 || d[q[i]]) {
     ans++;
   }
 }
```

```
cout << ans << endl;
}
int main() {
  ios_base::sync_with_stdio(false);cin.tie(NULL);
  int n;
  while (cin >> n)
    solve(n);
  return 0;
}
```

6.4 fibonacci properties

Let A, B and n be integer numbers.

$$k = A - B \tag{1}$$

$$F_A F_B = F_{k+1} F_A^2 + F_k F_A F_{A-1} \tag{2}$$

$$\sum_{i=0}^{n} F_i^2 = F_{n+1} F_n \tag{3}$$

ev(n) = returns 1 if n is even.

$$\sum_{i=0}^{n} F_i F_{i+1} = F_{n+1}^2 - ev(n) \tag{4}$$

$$\sum_{i=0}^{n} F_i F_{i-1} = \sum_{i=0}^{n-1} F_i F_{i+1}$$
 (5)

6.5 others

Approximate factorial $n! = \sqrt{2.\pi \cdot n} \cdot (\frac{n}{e})^n$

6.6 polynomials

```
// TODO: what's this ?
const double pi = acos(-1);
struct poly {
```

```
deque <double> coef;
double x_lo, x_hi;

double evaluate(double x) {
   double ans = 0;
   for (auto it : coef)
      ans = (ans * x + it);
   return ans;
}

double volume(double x, double dx=1e-6) {
   dx = (x_hi - x_lo) / 1000000.0;
   double ans = 0;
   for (double ix = x_lo; ix <= x; ix += dx) {
      double rad = evaluate(ix);
      ans += pi * rad * rad * dx;
   }
   return ans;
}
</pre>
```

6.7 sigma function

the sigma function is defined as:

$$\sigma_x(n) = \sum_{d|n} d^x$$

when x = 0 is called the divisor function, that counts the number of positive divisors of n.

Now, we are interested in find

$$\sum_{d|n} \sigma_0(d)$$

if n is written as prime factorization:

$$n = \prod_{i=1}^{k} P_i^{e_k}$$

we can demonstrate that:

$$\sum_{d|n} \sigma_0(d) = \prod_{i=1}^k g(e_k + 1)$$

where g(x) is the sum of the first x positive numbers:

$$q(x) = (x * (x + 1))/2$$

7 Matrix

7.1 matrix

```
const int MN = 111;
const int mod = 10000;
struct matrix {
 int r, c;
 int m[MN][MN];
 matrix (int _r, int _c) : r (_r), c (_c) {
   memset(m, 0, sizeof m);
 }
 void print() {
   for (int i = 0; i < r; ++i) {</pre>
     for (int j = 0; j < c; ++j)
       cout << m[i][j] << " ";
     cout << endl:</pre>
 }
  int x[MN][MN];
 matrix & operator *= (const matrix &o) {
   memset(x, 0, sizeof x);
   for (int i = 0; i < r; ++i)
     for (int k = 0; k < c; ++k)
       if (m[i][k] != 0)
         for (int j = 0; j < c; ++j) {
           x[i][j] = (x[i][j] + ((m[i][k] * o.m[k][j]) % mod)) % mod;
   memcpy(m, x, sizeof(m));
   return *this;
 }
};
void matrix_pow(matrix b, long long e, matrix &res) {
 memset(res.m, 0, sizeof res.m);
```

```
for (int i = 0; i < b.r; ++i)
  res.m[i][i] = 1;

if (e == 0) return;
while (true) {
  if (e & 1) res *= b;
  if ((e >>= 1) == 0) break;
  b *= b;
}
```

8 Misc

8.1 Template Java

```
import java.io.*;
import java.util.StringTokenizer;
public class Template {
   public static void main(String []args) throws IOException {
       Scanner in = new Scanner(System.in);
       OutputWriter out = new OutputWriter(System.out);
       Task solver = new Task();
       solver.solve(in. out):
       out.close();
}
class Task{
   public void solve(Scanner in, OutputWriter out){
class Scanner{
   public BufferedReader reader;
   public StringTokenizer st;
   public Scanner(InputStream stream){
       reader = new BufferedReader(new InputStreamReader(stream));
```

```
st = null:
   }
   public String next(){
       while(st == null || !st.hasMoreTokens()){
           try{
              String line = reader.readLine();
              if(line == null) return null;
              st = new StringTokenizer(line);
           }catch (Exception e){
              throw (new RuntimeException());
           }
       }
       return st.nextToken();
   }
   public int nextInt(){
       return Integer.parseInt(next());
   public long nextLong(){
       return Long.parseLong(next());
   public double nextDouble(){
       return Double.parseDouble(next());
}
class OutputWriter{
   BufferedWriter writer;
   public OutputWriter(OutputStream stream){
       writer = new BufferedWriter(new OutputStreamWriter(stream));
   }
   public void print(int i) throws IOException {
       writer.write(i);
   }
   public void print(String s) throws IOException {
       writer.write(s);
   }
   public void print(char []c) throws IOException {
       writer.write(c);
   }
```

```
public void close() throws IOException {
    writer.close();
}
```

8.2 dates

```
// Time - Leap years
// A[i] has the accumulated number of days from months previous to i
const int A[13] = { 0, 0, 31, 59, 90, 120, 151, 181, 212, 243, 273, 304,
    334 };
// same as A, but for a leap year
const int B[13] = { 0, 0, 31, 60, 91, 121, 152, 182, 213, 244, 274, 305,
    335 };
// returns number of leap years up to, and including, y
int leap_vears(int y) { return y / 4 - y / 100 + y / 400; }
bool is_leap(int y) { return y % 400 == 0 || (y % 4 == 0 && y % 100 !=
    0); }
// number of days in blocks of years
const int p400 = 400*365 + leap_years(400);
const int p100 = 100*365 + leap_years(100);
const int p4 = 4*365 + 1;
const int p1 = 365;
int date_to_days(int d, int m, int y)
 return (y - 1) * 365 + leap_years(y - 1) + (is_leap(y) ? B[m] : A[m]) +
void days_to_date(int days, int &d, int &m, int &y)
 bool top100; // are we in the top 100 years of a 400 block?
 bool top4; // are we in the top 4 years of a 100 block?
 bool top1; // are we in the top year of a 4 block?
 y = 1;
 top100 = top4 = top1 = false;
 y += ((days-1) / p400) * 400;
  d = (days-1) \% p400 + 1;
```

```
if (d > p100*3) top100 = true, d -= 3*p100, y += 300;
else y += ((d-1) / p100) * 100, d = (d-1) % p100 + 1;

if (d > p4*24) top4 = true, d -= 24*p4, y += 24*4;
else y += ((d-1) / p4) * 4, d = (d-1) % p4 + 1;

if (d > p1*3) top1 = true, d -= p1*3, y += 3;
else y += (d-1) / p1, d = (d-1) % p1 + 1;

const int *ac = top1 && (!top4 || top100) ? B : A;
for (m = 1; m < 12; ++m) if (d <= ac[m + 1]) break;
d -= ac[m];</pre>
```

8.3 fraction

```
struct frac{
  long long x, y;
  frac(long long a, long long b) {
    long long g = __gcd(a, b);
    x = a / g;
    y = b / g;
  }
  bool operator < (const frac &o) const {
    return (x * o.y < y * o.x);
  }
};</pre>
```

8.4 io

```
// taken from :
    https://github.com/lbv/pc-code/blob/master/solved/c-e/diablo/diablo.cpp
// this is very fast as well :
    https://github.com/lbv/pc-code/blob/master/code/input.cpp

typedef unsigned int u32;
#define BUF 524288
struct Reader {
    char buf[BUF]; char b; int bi, bz;
    Reader() { bi=bz=0; read(); }
```

```
void read() {
   if (bi==bz) { bi=0; bz = fread(buf, 1, BUF, stdin); }
   b = bz ? buf[bi++] : 0; }
void skip() { while (b > 0 && b <= 32) read(); }
u32 next_u32() {
   u32 v = 0; for (skip(); b > 32; read()) v = v*10 + b-48; return v; }
int next_int() {
   int v = 0; bool s = false;
   skip(); if (b == '-') { s = true; read(); }
   for (; 48<=b&&b<=57; read()) v = v*10 + b-48; return s ? -v : v; }
char next_char() { skip(); char c = b; read(); return c; }
};</pre>
```

9 Number theory

9.1 convolution

```
typedef long long int LL;
typedef pair<LL, LL> PLL;
inline bool is_pow2(LL x) {
 return (x & (x-1)) == 0;
inline int ceil_log2(LL x) {
 int ans = 0;
 --x:
 while (x != 0) {
   x >>= 1;
   ans++:
 return ans;
/* Returns the convolution of the two given vectors in time proportional
    to n*log(n).
* The number of roots of unity to use nroots_unity must be set so that
     the product of the first
* nroots_unity primes of the vector nth_roots_unity is greater than the
     maximum value of the
* convolution. Never use sizes of vectors bigger than 2^24, if you need
     to change the values of
```

```
* the nth roots of unity to appropriate primes for those sizes.
vector<LL> convolve(const vector<LL> &a, const vector<LL> &b, int
    nroots_unity = 2) {
 int N = 1 << ceil_log2(a.size() + b.size());</pre>
 vector<LL> ans(N,0), fA(N), fB(N), fC(N);
 LL modulo = 1;
 for (int times = 0; times < nroots_unity; times++) {</pre>
   fill(fA.begin(), fA.end(), 0);
   fill(fB.begin(), fB.end(), 0);
   for (int i = 0; i < a.size(); i++) fA[i] = a[i];</pre>
   for (int i = 0; i < b.size(); i++) fB[i] = b[i];</pre>
   LL prime = nth_roots_unity[times].first;
   LL inv_modulo = mod_inv(modulo % prime, prime);
   LL normalize = mod_inv(N, prime);
   ntfft(fA, 1, nth_roots_unity[times]);
   ntfft(fB, 1, nth_roots_unity[times]);
   for (int i = 0; i < N; i++) fC[i] = (fA[i] * fB[i]) % prime;</pre>
   ntfft(fC, -1, nth_roots_unity[times]);
   for (int i = 0; i < N; i++) {</pre>
     LL curr = (fC[i] * normalize) % prime;
     LL k = (curr - (ans[i] % prime) + prime) % prime;
     k = (k * inv_modulo) % prime;
     ans[i] += modulo * k;
   modulo *= prime;
 }
 return ans;
```

9.2 crt

```
/**
 * Chinese remainder theorem.
 * Find z such that z % x[i] = a[i] for all i.
 * */
long long crt(vector<long long> &a, vector<long long> &x) {
 long long z = 0;
 long long n = 1;
 for (int i = 0; i < x.size(); ++i)
    n *= x[i];

for (int i = 0; i < a.size(); ++i) {</pre>
```

```
long long tmp = (a[i] * (n / x[i])) % n;
tmp = (tmp * mod_inv(n / x[i], x[i])) % n;
z = (z + tmp) % n;
}
return (z + n) % n;
}
```

9.3 diophantine equations

```
long long gcd(long long a, long long b, long long &x, long long &y) {
 if (a == 0) {
   x = 0:
   y = 1;
   return b;
 long long x1, y1;
 long long d = gcd(b \% a, a, x1, y1);
 x = y1 - (b / a) * x1;
 v = x1;
 return d;
bool find_any_solution(long long a, long long b, long long c, long long
    &x0,
   long long &y0, long long &g) {
 g = gcd(abs(a), abs(b), x0, y0);
 if (c % g) {
   return false;
 x0 *= c / g;
 v0 *= c / g;
 if (a < 0) x0 = -x0;
 if (b < 0) v0 = -v0;
 return true;
void shift_solution(long long &x, long long &y, long long a, long long b,
   long long cnt) {
 x += cnt * b:
 y -= cnt * a;
```

```
long long find_all_solutions(long long a, long long b, long long c,
   long long minx, long long maxx, long long miny,
   long long maxy) {
 long long x, y, g;
 if (!find_any_solution(a, b, c, x, y, g)) return 0;
 a /= g;
 b /= g;
 long long sign_a = a > 0 ? +1 : -1;
 long long sign_b = b > 0 ? +1 : -1;
 shift_solution(x, y, a, b, (minx - x) / b);
 if (x < minx) shift_solution(x, y, a, b, sign_b);</pre>
 if (x > maxx) return 0;
 long long lx1 = x;
 shift_solution(x, y, a, b, (maxx - x) / b);
 if (x > maxx) shift_solution(x, y, a, b, -sign_b);
 long long rx1 = x;
 shift_solution(x, y, a, b, -(miny - y) / a);
 if (y < miny) shift_solution(x, y, a, b, -sign_a);</pre>
 if (y > maxy) return 0;
 long long 1x2 = x;
 shift_solution(x, y, a, b, -(maxy - y) / a);
 if (y > maxy) shift_solution(x, y, a, b, sign_a);
 long long rx2 = x;
 if (1x2 > rx2) swap(1x2, rx2);
 long long lx = max(lx1, lx2);
 long long rx = min(rx1, rx2);
 if (lx > rx) return 0;
 return (rx - lx) / abs(b) + 1;
}
```

9.4 discrete logarithm

```
// Computes x which a \hat{x} = b \mod n.
```

```
long long d_log(long long a, long long b, long long n) {
 long long m = ceil(sqrt(n));
 long long aj = 1;
 map<long long, long long> M;
 for (int i = 0; i < m; ++i) {</pre>
   if (!M.count(aj))
     M[aj] = i;
   aj = (aj * a) % n;
 long long coef = mod_pow(a, n - 2, n);
  coef = mod_pow(coef, m, n);
 // coef = a ^ (-m)
 long long gamma = b;
 for (int i = 0; i < m; ++i) {</pre>
   if (M.count(gamma)) {
     return i * m + M[gamma];
   } else {
     gamma = (gamma * coef) % n;
 }
 return -1;
```

9.5 ext euclidean

```
void ext_euclid(long long a, long long b, long long &x, long long &y,
    long long &g) {
    x = 0, y = 1, g = b;
    long long m, n, q, r;
    for (long long u = 1, v = 0; a != 0; g = a, a = r) {
        q = g / a, r = g % a;
        m = x - u * q, n = y - v * q;
        x = u, y = v, u = m, v = n;
    }
}
```

9.6 highest exponent factorial

```
int highest_exponent(int p, const int &n){
  int ans = 0;
```

```
int t = p;
while(t <= n){
   ans += n/t;
   t*=p;
}
return ans;
}</pre>
```

9.7 miller rabin

```
const int rounds = 20;
// checks whether a is a witness that n is not prime, 1 < a < n
bool witness(long long a, long long n) {
 // check as in Miller Rabin Primality Test described
 long long u = n - 1;
 int t = 0;
 while (u % 2 == 0) {
   t++;
   u >>= 1;
 long long next = mod_pow(a, u, n);
 if (next == 1) return false;
 long long last;
 for (int i = 0; i < t; ++i) {</pre>
   last = next;
   next = mod_mul(last, last, n);
   if (next == 1) {
     return last != n - 1;
   }
 return next != 1;
// Checks if a number is prime with prob 1 - 1 / (2 ^ it)
// D(miller_rabin(999999999999997LL) == 1);
// D(miller_rabin(999999999971LL) == 1);
// D(miller_rabin(7907) == 1);
bool miller_rabin(long long n, int it = rounds) {
 if (n <= 1) return false;</pre>
 if (n == 2) return true;
 if (n % 2 == 0) return false;
```

```
for (int i = 0; i < it; ++i) {
  long long a = rand() % (n - 1) + 1;
  if (witness(a, n)) {
    return false;
  }
}
return true;
}</pre>
```

9.8 mod integer

```
template<class T, T mod>
struct mint_t {
   T val;
   mint_t() : val(0) {}
   mint_t(T v) : val(v % mod) {}

mint_t operator + (const mint_t& o) const {
    return (val + o.val) % mod;
   }
   mint_t operator - (const mint_t& o) const {
    return (val - o.val) % mod;
   }
   mint_t operator * (const mint_t& o) const {
    return (val * o.val) % mod;
   }
}

mint_t operator * (const mint_t& o) const {
   return (val * o.val) % mod;
   }
};

typedef mint_t<long long, 998244353> mint;
```

9.9 mod inv

```
long long mod_inv(long long n, long long m) {
  long long x, y, gcd;
  ext_euclid(n, m, x, y, gcd);
  if (gcd != 1)
    return 0;
  return (x + m) % m;
}
```

9.10 mod mul

```
// Computes (a * b) % mod
long long mod_mul(long long a, long long b, long long mod) {
  long long x = 0, y = a % mod;
  while (b > 0) {
    if (b & 1)
        x = (x + y) % mod;
        y = (y * 2) % mod;
        b /= 2;
    }
    return x % mod;
}
```

9.11 mod pow

```
// Computes ( a ^ exp ) % mod.
long long mod_pow(long long a, long long exp, long long mod) {
  long long ans = 1;
  while (exp > 0) {
    if (exp & 1)
      ans = mod_mul(ans, a, mod);
    a = mod_mul(a, a, mod);
    exp >>= 1;
  }
  return ans;
}
```

9.12 number theoretic transform

```
typedef long long int LL;
typedef pair<LL, LL> PLL;

/* The following vector of pairs contains pairs (prime, generator)
 * where the prime has an Nth root of unity for N being a power of two.
 * The generator is a number g s.t g^(p-1)=1 (mod p)
 * but is different from 1 for all smaller powers */
vector<PLL> nth_roots_unity {
    {1224736769,330732430},{1711276033,927759239},{167772161,167489322},
    {469762049,343261969},{754974721,643797295},{1107296257,883865065}};
```

```
PLL ext_euclid(LL a, LL b) {
 if (b == 0)
   return make_pair(1,0);
 pair<LL,LL> rc = ext_euclid(b, a % b);
 return make_pair(rc.second, rc.first - (a / b) * rc.second);
//returns -1 if there is no unique modular inverse
LL mod inv(LL x. LL modulo) {
 PLL p = ext_euclid(x, modulo);
 if ( (p.first * x + p.second * modulo) != 1 )
   return -1;
 return (p.first+modulo) % modulo;
//Number theory fft. The size of a must be a power of 2
void ntfft(vector<LL> &a, int dir, const PLL &root_unity) {
 int n = a.size():
 LL prime = root_unity.first;
 LL basew = mod_pow(root_unity.second, (prime-1) / n, prime);
  if (dir < 0) basew = mod_inv(basew, prime);</pre>
 for (int m = n; m >= 2; m >>= 1) {
   int mh = m >> 1:
   LL w = 1;
   for (int i = 0; i < mh; i++) {</pre>
     for (int j = i; j < n; j += m) {</pre>
       int k = j + mh;
       LL x = (a[j] - a[k] + prime) % prime;
       a[j] = (a[j] + a[k]) \% prime;
       a[k] = (w * x) % prime;
     w = (w * basew) % prime;
   basew = (basew * basew) % prime;
  int i = 0;
  for (int j = 1; j < n - 1; j++) {
   for (int k = n >> 1; k > (i ^= k); k >>= 1);
   if (j < i) swap(a[i], a[j]);</pre>
 }
}
```

9.13 pollard rho factorize

```
long long pollard_rho(long long n) {
 long long x, y, i = 1, k = 2, d;
 x = y = rand() \% n;
 while (1) {
   ++i;
   x = mod_mul(x, x, n);
   x += 2:
   if (x \ge n) x = n;
   if (x == y) return 1;
   d = \_gcd(abs(x - y), n);
   if (d != 1) return d;
   if (i == k) {
     y = x;
     k *= 2;
 }
 return 1;
}
// Returns a list with the prime divisors of n
vector<long long> factorize(long long n) {
 vector<long long> ans;
 if (n == 1)
   return ans;
 if (miller_rabin(n)) {
   ans.push_back(n);
 } else {
   long long d = 1;
   while (d == 1)
     d = pollard_rho(n);
   vector<long long> dd = factorize(d);
   ans = factorize(n / d);
   for (int i = 0; i < dd.size(); ++i)</pre>
     ans.push_back(dd[i]);
 }
 return ans;
}
```

9.14 primes

```
namespace primes {
  const int MP = 100001;
  bool sieve[MP];
 long long primes[MP];
  int num_p;
  void fill_sieve() {
   num_p = 0;
   sieve[0] = sieve[1] = true;
   for (long long i = 2; i < MP; ++i) {</pre>
     if (!sieve[i]) {
       primes[num_p++] = i;
       for (long long j = i * i; j < MP; j += i)
         sieve[j] = true;
     }
   }
 }
  // Finds prime numbers between a and b, using basic primes up to sqrt(b)
  // a must be greater than 1.
  vector<long long> seg_sieve(long long a, long long b) {
   long long ant = a;
   a = max(a, 3LL);
   vector<bool> pmap(b - a + 1);
   long long sqrt_b = sqrt(b);
   for (int i = 0; i < num_p; ++i) {</pre>
     long long p = primes[i];
     if (p > sqrt_b) break;
     long long j = (a + p - 1) / p;
     for (long long v = (j == 1) ? p + p : j * p; v <= b; v += p) {
       pmap[v - a] = true;
   }
   vector<long long> ans;
   if (ant == 2) ans.push_back(2);
   int start = a % 2 ? 0 : 1;
   for (int i = start, I = b - a + 1; i < I; i += 2)</pre>
     if (pmap[i] == false)
       ans.push_back(a + i);
   return ans;
  vector<pair<int, int>> factor(int n) {
   vector<pair<int, int>> ans;
   if (n == 0) return ans;
```

```
for (int i = 0; primes[i] * primes[i] <= n; ++i) {
   if ((n % primes[i]) == 0) {
     int expo = 0;
     while ((n % primes[i]) == 0) {
        expo++;
        n /= primes[i];
     }
     ans.emplace_back(primes[i], expo);
   }
}

if (n > 1) {
   ans.emplace_back(n, 1);
}
return ans;
}
```

9.15 totient sieve

```
for (int i = 1; i < MN; i++)
  phi[i] = i;

for (int i = 1; i < MN; i++)
  if (!sieve[i]) // is prime
   for (int j = i; j < MN; j += i)
     phi[j] -= phi[j] / i;</pre>
```

9.16 totient

```
long long totient(long long n) {
  if (n == 1) return 0;
  long long ans = n;
  for (int i = 0; primes[i] * primes[i] <= n; ++i) {
    if ((n % primes[i]) == 0) {
      while ((n % primes[i]) == 0) n /= primes[i];
      ans -= ans / primes[i];
    }
  }
  if (n > 1) {
    ans -= ans / n;
}
```

```
}
return ans;
}
```

10 Strings

10.1 Incremental Aho Corasick

```
class IncrementalAhoCorasic {
 static const int Alphabets = 26;
 static const int AlphabetBase = 'a';
 struct Node {
   Node *fail:
   Node *next[Alphabets];
   int sum;
   Node() : fail(NULL), next{}, sum(0) { }
 };
 struct String {
   string str;
   int sign;
 };
public:
 //totalLen = sum of (len + 1)
 void init(int totalLen) {
   nodes.resize(totalLen);
   nNodes = 0;
   strings.clear();
   roots.clear();
   sizes.clear();
   que.resize(totalLen);
 void insert(const string &str, int sign) {
   strings.push_back(String{ str, sign });
   roots.push_back(nodes.data() + nNodes);
   sizes.push_back(1);
   nNodes += (int)str.size() + 1;
   auto check = [&]() { return sizes.size() > 1 && sizes.end()[-1] ==
        sizes.end()[-2]; };
   if(!check())
```

```
makePMA(strings.end() - 1, strings.end(), roots.back(), que);
   while(check()) {
     int m = sizes.back();
     roots.pop_back();
     sizes.pop_back();
     sizes.back() += m;
     if(!check())
       makePMA(strings.end() - m * 2, strings.end(), roots.back(), que);
   }
 }
  int match(const string &str) const {
   int res = 0;
   for(const Node *t : roots)
     res += matchPMA(t, str);
   return res;
 }
private:
  static void makePMA(vector<String>::const_iterator begin,
      vector<String>::const_iterator end, Node *nodes, vector<Node*>
      &que) {
   int nNodes = 0;
   Node *root = new(&nodes[nNodes ++]) Node();
   for(auto it = begin; it != end; ++ it) {
     Node *t = root:
     for(char c : it->str) {
       Node *&n = t->next[c - AlphabetBase];
       if(n == nullptr)
         n = new(&nodes[nNodes ++]) Node();
       t = n;
     t->sum += it->sign;
   int qt = 0;
   for(Node *&n : root->next) {
     if(n != nullptr) {
       n->fail = root;
       que[qt ++] = n;
     } else {
       n = root;
     }
   for(int qh = 0; qh != qt; ++ qh) {
     Node *t = que[qh];
```

```
int a = 0;
     for(Node *n : t->next) {
       if(n != nullptr) {
         que[qt ++] = n;
         Node *r = t->fail;
         while(r->next[a] == nullptr)
           r = r \rightarrow fail;
         n->fail = r->next[a];
         n->sum += r->next[a]->sum;
       }
       ++ a:
  static int matchPMA(const Node *t, const string &str) {
   int res = 0;
   for(char c : str) {
     int a = c - AlphabetBase;
     while(t->next[a] == nullptr)
       t = t->fail;
     t = t-next[a];
     res += t->sum;
   return res;
  vector<Node> nodes;
  int nNodes;
  vector<String> strings;
  vector<Node*> roots:
 vector<int> sizes;
 vector<Node*> que;
};
int main() {
  int m;
  while(~scanf("%d", &m)) {
   IncrementalAhoCorasic iac;
   iac.init(600000);
   rep(i, m) {
     int ty;
     char s[300001];
     scanf("%d%s", &ty, s);
```

```
if(ty == 1) {
    iac.insert(s, +1);
} else if(ty == 2) {
    iac.insert(s, -1);
} else if(ty == 3) {
    int ans = iac.match(s);
    printf("%d\n", ans);
    fflush(stdout);
} else {
    abort();
}
}
return 0;
```

10.2 minimal string rotation

```
// Lexicographically minimal string rotation
int lmsr() {
 string s;
 cin >> s;
 int n = s.size();
 s += s;
 vector<int> f(s.size(), -1);
 int k = 0;
 for (int j = 1; j < 2 * n; ++j) {
   int i = f[j - k - 1];
   while (i != -1 && s[j] != s[k + i + 1]) {
     if (s[j] < s[k + i + 1])
      k = j - i - 1;
     i = f[i];
   }
   if (i == -1 \&\& s[j] != s[k + i + 1]) {
     if (s[j] < s[k + i + 1]) {
       k = j;
     f[j - k] = -1;
   } else {
     f[j - k] = i + 1;
   }
 }
 return k;
```

}

10.3 suffix array

```
/**
 * 0 (n log^2 (n))
 * See http://web.stanford.edu/class/cs97si/suffix-array.pdf for reference
 * */
struct entry{
 int a, b, p;
  entry(){}
  entry(int x, int y, int z): a(x), b(y), p(z){}
  bool operator < (const entry &o) const {</pre>
   return (a == o.a)? (b == o.b)? (p < o.p): (b < o.b): (a < o.a);
 }
};
struct SuffixArray{
  const int N;
  string s;
  vector<vector<int> > P;
  vector<entry> M;
  SuffixArray(const string &s) : N(s.length()), s(s), P(1, vector<int>
      (N, O)), M(N) {
    for (int i = 0; i < N; ++i)</pre>
     P[0][i] = (int) s[i];
    for (int skip = 1, level = 1; skip < N; skip *= 2, level++) {</pre>
     P.push_back(vector<int>(N, 0));
     for (int i = 0 ; i < N; ++i) {</pre>
       int next = ((i + skip) < N) ? P[level - 1][i + skip] : -10000;
       M[i] = entry(P[level - 1][i], next, i);
     sort(M.begin(), M.end());
     for (int i = 0; i < N; ++i)
       P[level][M[i].p] = (i > 0 \text{ and } M[i].a == M[i - 1].a \text{ and } M[i].b ==
            M[i - 1].b) ? P[level][M[i - 1].p] : i;
   }
  }
  vector<int> getSuffixArray(){
```

```
vector<int> &rank = P.back();
   vector<pair<int, int> > inv(rank.size());
   for (int i = 0; i < rank.size(); ++i)</pre>
     inv[i] = make_pair(rank[i], i);
   sort(inv.begin(), inv.end());
   vector<int> sa(rank.size());
   for (int i = 0; i < rank.size(); ++i)</pre>
     sa[i] = inv[i].second;
   return sa:
 }
 // returns the length of the longest common prefix of s[i...L-1] and
      s[j...L-1]
  int lcp(int i, int j) {
   int len = 0;
   if (i == j) return N - i;
   for (int k = P.size() - 1; k \ge 0 && i < N && j < N; --k) {
     if (P[k][i] == P[k][j]) {
       i += 1 << k:
       i += 1 << k;
       len += 1 << k;
   return len;
 }
};
```

10.4 suffix automaton

```
/*
 * Suffix automaton:
 * This implementation was extended to maintain (online) the
 * number of different substrings. This is equivalent to compute
 * the number of paths from the initial state to all the other
 * states.
 *
 * The overall complexity is O(n)
 * can be tested here:
    https://www.urionlinejudge.com.br/judge/en/problems/view/1530
 * */

struct state {
    int len, link;
```

```
long long num_paths;
 map<int, int> next;
};
const int MN = 200011;
state sa[MN << 1];</pre>
int sz, last;
long long tot_paths;
void sa_init() {
  sz = 1:
 last = 0;
  sa[0].len = 0;
  sa[0].link = -1;
  sa[0].next.clear();
  sa[0].num_paths = 1;
 tot_paths = 0;
void sa_extend(int c) {
  int cur = sz++;
  sa[cur].len = sa[last].len + 1;
  sa[cur].next.clear();
  sa[cur].num_paths = 0;
  for (p = last; p != -1 && !sa[p].next.count(c); p = sa[p].link) {
   sa[p].next[c] = cur;
   sa[cur].num_paths += sa[p].num_paths;
   tot_paths += sa[p].num_paths;
  if (p == -1) {
   sa[cur].link = 0;
 } else {
   int q = sa[p].next[c];
   if (sa[p].len + 1 == sa[q].len) {
     sa[cur].link = q;
   } else {
     int clone = sz++;
     sa[clone].len = sa[p].len + 1;
     sa[clone].next = sa[q].next;
     sa[clone].num_paths = 0;
     sa[clone].link = sa[q].link;
     for (; p!= -1 && sa[p].next[c] == q; p = sa[p].link) {
       sa[p].next[c] = clone;
```

```
sa[q].num_paths -= sa[p].num_paths;
sa[clone].num_paths += sa[p].num_paths;
}
sa[q].link = sa[cur].link = clone;
}
last = cur;
}
```

10.5 z algorithm

```
using namespace std;
#include<bits/stdc++.h>

vector<int> compute_z(const string &s){
   int n = s.size();
   vector<int> z(n,0);
   int l,r;
   r = l = 0;
   for(int i = 1; i < n; ++i){
      if(i > r) {
        l = r = i;
      while(r < n and s[r - l] == s[r])r++;
      z[i] = r - l;r--;
   }else{
      int k = i-l;</pre>
```

```
if(z[k] < r - i +1) z[i] = z[k];
     else {
       1 = i;
       while (r < n \text{ and } s[r - 1] == s[r])r++;
       z[i] = r - 1;r--;
     }
   }
 return z;
int main(){
 //string line;cin>>line;
  string line = "alfalfa";
 vector<int> z = compute_z(line);
 for(int i = 0; i < z.size(); ++i ){</pre>
   if(i)cout<<" ";
   cout<<z[i];
  cout<<endl;
 // must print "0 0 0 4 0 0 1"
 return 0;
```