Team notebook

${\bf HCMUS\text{-}IdentityImbalance}$

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Contents		5 Graphs	11	9.8 mod integer	
			5.1 bridges	11	9.9 mod inv
1	Algorithms	1	5.2 delete on dsu	12	9.10 mod mul
	1.1 Mo's algorithm on trees	1	5.3 euler path	13	9.11 mod pow
	1.2 Mo's algorithm	1	5.4 karp min mean cycle	13	9.12 number theoretic transform 2
	1.3 parallel binary search	1	5.5 konig's theorem	13	9.13 pollard rho factorize 2
			5.6 matching	13	9.14 primes
2	DP Optimizations	2	5.7 max flow min cost	14	9.15 totient sieve
	2.1 convex hull trick	2	5.8 minimum path cover in DAG	15	9.16 totient
	2.2 divide and conquer	2	5.9 planar graph (euler)	15	
			5.10 two sat (with kosaraju)	15	10 Strings 22
3	Data structures	3	6 Math	16	10.1 Incremental Aho Corasick 25
	3.1 dsu	3	6.1 Lucas theorem	16	10.2 minimal string rotation
	3.2 fake update	3	6.2 cumulative sum of divisors	16	10.3 suffix array
	3.3 fenwick tree	3	6.3 fft	16	10.4 suffix automaton
	3.4 hash table	4	6.4 fibonacci properties	17	10.5 z algorithm
	3.5 heavy light decomposition	4	6.5 others	17	
	3.6 persistent array	5	6.6 polynomials	17	1 Algorithms
	3.7 persistent seg tree	5	6.7 sigma function	18	
	3.8 persistent segment (v2)	6	6.8 system different constraints	18	1.1 Mo's algorithm on trees
	3.9 persistent trie	6	7 34 .	10	
	3.10 segment tree	7	7 Matrix	18 18	/*
	3.11 sparse table	8	7.1 matrix	18	
	3.12 trie	8	8 Misc	19	Problem:
	5.12 tile	O	8.1 dates	19	https://www.spoj.com/problems/COT2/
4	Geometry	8			Given a tree with N nodes and Q queries. Each node
-	4.1 center 2 points + radious	8	9 Number theory	19	has an integer weight.
	-	8	9.1 convolution	19	Each query provides two numbers u and v, ask for how many different integers weight of nodes
	1 1	_	9.2 crt	19	there are on path from u to v.
	4.3 convex diameter	9	9.3 diophantine equations	19	
	4.4 pick theorem	9	9.4 discrete logarithm	20	Modify DEC.
	4.5 squares	10	9.5 ext euclidean	20	Modify DFS:
	4.6 template	10	9.6 highest exponent factorial	20	For each node u, maintain the start and the end DFS
	4.7 triangles	11	9.7 miller rabin	20	time. Let's call them $ST(u)$ and $EN(u)$.

```
=> For each query, a node is considered if its
     occurrence count is one.
Query solving:
Let's query be (u, v). Assume that ST(u) \le ST(v).
    Denotes P as LCA(u, v).
Case 1: P = u
Our query would be in range [ST(u), ST(v)].
Case 2: P != u
Our query would be in range [EN(u), ST(v)] +
     [ST(q), ST(q)]
void update(int &L, int &R, int qL, int qR){
   while (L > qL) add(--L);
   while (R < qR) add(++R);
   while (L < qL) del(L++);</pre>
   while (R > qR) del(R--);
}
vector <int> MoQueries(int n, vector <query> Q){
   block_size = sqrt((int)nodes.size());
   sort(Q.begin(), Q.end(), [](const query &A,
        const query &B){
       return (ST[A.1]/block size !=
            ST[B.1]/block size)?
            (ST[A.1]/block size <
            ST[B.1]/block_size) : (ST[A.r] <
            ST[B.r]):
   }):
   vector <int> res:
   res.resize((int)0.size()):
   LCA lca:
   lca.initialize(n):
   int L = 1, R = 0;
   for(query q: Q){
       int u = q.1, v = q.r;
       if(ST[u] > ST[v]) swap(u, v); // assume that
            S[u] \leftarrow S[v]
       int parent = lca.get(u, v);
       if(parent == u){
           int qL = ST[u], qR = ST[v];
           update(L, R, qL, qR);
       }else{
           int aL = EN[u], aR = ST[v]:
           update(L, R, qL, qR);
           if(cnt val[a[parent]] == 0)
```

```
res[q.pos] += 1;
}
    res[q.pos] += cur_ans;
}
return res;
}
```

1.2 Mo's algorithm

```
https://www.spoj.com/problems/FREQ2/
vector <int> MoQueries(int n, vector <query> Q){
   block_size = sqrt(n);
   sort(Q.begin(), Q.end(), [](const query &A,
        const query &B){
       return (A.1/block_size != B.1/block_size)?
            (A.1/block size < B.1/block size) :
            (A.r < B.r):
   }):
   vector <int> res:
   res.resize((int)Q.size());
   int L = 1, R = 0:
   for(query q: Q){
       while (L > q.1) add(--L);
       while (R < q.r) add(++R);
       while (L < q.1) del(L++);</pre>
       while (R > q.r) del(R--);
       res[q.pos] = calc(1, R-L+1);
   }
   return res;
}
```

1.3 parallel binary search

```
int lo[N], mid[N], hi[N];
vector<int> vec[N];

void clear() //Reset
{
    memset(bit, 0, sizeof(bit));
}

void apply(int idx) //Apply ith update/query
{
```

```
if(ql[idx] <= qr[idx])</pre>
               update(ql[idx], qa[idx]),
                    update(gr[idx]+1, -ga[idx]);
        else
               update(1, qa[idx]);
               update(qr[idx]+1, -qa[idx]);
               update(ql[idx], qa[idx]);
       }
}
bool check(int idx) //Check if the condition is
     satisfied
        int req=reqd[idx];
       for(auto &it:owns[idx])
               req-=pref(it);
               if(req<0)
                       break;
        if(req<=0)</pre>
               return 1;
        return 0;
}
void work()
        for(int i=1:i<=q:i++)</pre>
               vec[i].clear():
        for(int i=1:i<=n:i++)</pre>
               if(mid[i]>0)
                       vec[mid[i]].push_back(i);
        clear():
       for(int i=1;i<=q;i++)</pre>
               applv(i):
               for(auto &it:vec[i]) //Add
                    appropriate check conditions
                       if(check(it))
                              hi[it]=i;
                       else
                              lo[it]=i+1;
               }
}
void parallel_binary()
        for(int i=1;i<=n;i++)</pre>
               lo[i]=1, hi[i]=q+1;
        bool changed = 1:
        while(changed)
```

2 DP Optimizations

2.1 convex hull trick

```
#define long long long
#define pll pair <long. long>
#define all(c) c.begin(), c.end()
#define fastio ios_base::sync_with_stdio(false);
     cin.tie(0)
struct line{
   long a, b;
   line() {};
   line(long a, long b) : a(a), b(b) {};
   bool operator < (const line &A) const {</pre>
               return pll(a,b) < pll(A.a,A.b);</pre>
};
bool bad(line A, line B, line C){
   return (C.b - B.b) * (A.a - B.a) <= (B.b - A.b)
        * (B.a - C.a):
void addLine(vector<line> &memo. line cur){
   int k = memo.size();
   while (k \ge 2 \&\& bad(memo[k - 2], memo[k - 1],
        cur)){
       memo.pop_back();
   memo.push_back(cur);
long Fn(line A, long x){
   return A.a * x + A.b;
```

```
}
long query(vector<line> &memo, long x){
   int lo = 0, hi = memo.size() - 1;
   while (lo != hi){
       int mi = (lo + hi) / 2;
       if (Fn(memo[mi], x) > Fn(memo[mi + 1], x)){
           lo = mi + 1:
       else hi = mi:
   return Fn(memo[lo], x):
const int N = 1e6 + 1:
long dp[N]:
int main()
   fastio;
   int n, c; cin >> n >> c;
   vector<line> memo;
   for (int i = 1; i <= n; i++){</pre>
       long val; cin >> val;
       addLine(memo, {-2 * val, val * val + dp[i -
       dp[i] = query(memo, val) + val * val + c;
   cout << dp[n] << '\n';
   return 0;
}
```

2.2 divide and conquer

B Data structures

3.1 dsu

```
class DSU{
public:
   vector <int> parent;
   void initialize(int n){
       parent.resize(n+1, -1);
   int findSet(int u){
       while(parent[u] > 0)
           u = parent[u];
       return u;
   void Union(int u, int v){
       int x = parent[u] + parent[v];
       if(parent[u] > parent[v]){
          parent[v] = x:
          parent[u] = v:
       }else{
           parent[u] = x:
           parent[v] = u:
};
```

3.2 fake update

```
vector <int> fake_bit[MAXN];
```

```
void fake_update(int x, int y, int limit_x){
   for(int i = x; i < limit_x; i += i&(-i))
       fake_bit[i].pb(y);
void fake_get(int x, int y){
   for(int i = x: i >= 1: i -= i&(-i))
       fake_bit[i].pb(y);
}
vector <int> bit[MAXN]:
void update(int x, int y, int limit_x, int val){
   for(int i = x: i < limit x: i += i&(-i)){
       for(int j = lower_bound(fake_bit[i].begin(),
            fake bit[i].end(), v) -
            fake_bit[i].begin(); j <</pre>
            fake_bit[i].size(); j += j&(-j))
           bit[i][j] = max(bit[i][j], val);
int get(int x, int y){
   int ans = 0;
   for(int i = x; i >= 1; i -= i&(-i)){
       for(int j = lower_bound(fake_bit[i].begin(),
            fake_bit[i].end(), y) -
            fake_bit[i].begin(); j >= 1; j -=
            i&(-i))
          ans = max(ans, bit[i][j]);
   }
   return ans:
int main(){
   io
   int n: cin >> n:
   vector <int> Sx, Sy;
   for(int i = 1: i <= n: i++){
       cin >> a[i].fi >> a[i].se:
       Sx.pb(a[i].fi);
       Sy.pb(a[i].se);
   unique_arr(Sx);
   unique_arr(Sy);
   // unique all value
   for(int i = 1; i <= n; i++){</pre>
       a[i].fi = lower_bound(Sx.begin(), Sx.end(),
            a[i].fi) - Sx.begin();
       a[i].se = lower_bound(Sy.begin(), Sy.end(),
            a[i].se) - Sy.begin();
   // do fake BIT update and get operator
   for(int i = 1; i \le n; i++){
```

```
fake_get(a[i].fi-1, a[i].se-1);
       fake_update(a[i].fi, a[i].se,
            (int)Sx.size());
   }
   for(int i = 0; i < Sx.size(); i++){</pre>
       fake bit[i].pb(INT MIN): // avoid zero
       sort(fake_bit[i].begin(), fake_bit[i].end());
       fake bit[i].resize(unique(fake bit[i].begin().
            fake bit[i].end()) -
            fake_bit[i].begin());
       bit[i].resize((int)fake bit[i].size(), 0);
   }
   // real update, get operator
   int res = 0:
   for(int i = 1: i <= n: i++){</pre>
       int maxCurLen = get(a[i].fi-1, a[i].se-1) +
       res = max(res, maxCurLen);
       update(a[i].fi, a[i].se, (int)Sx.size(),
            maxCurLen):
   }
}
```

3.3 fenwick tree

```
template <typename T>
class FenwickTree{
 vector <T> fenw;
 int n;
public:
 void initialize(int _n){
   this \rightarrow n = n:
   fenw.resize(n+1):
 void update(int id. T val) {
   while (id <= n) {
     fenw[id] += val:
     id += id&(-id):
     }
 T get(int id){
   T ans{};
   while(id >= 1){}
     ans += fenw[id]:
     id = id&(-id);
   return ans;
```

};

3.4 hash table

```
/*
  * Micro hash table, can be used as a set.
  * Very efficient vs std::set
  *
  */
const int MN = 1001;
struct ht {
  int _s[(MN + 10) >> 5];
  int len;
  void set(int id) {
    len++;
    _s[id >> 5] |= (1LL << (id & 31));
}
bool is_set(int id) {
    return _s[id >> 5] & (1LL << (id & 31));
}
};</pre>
```

3.5 heavy light decomposition

```
const int N = 1e5 + 5;
const int LG = log2(N) + 1;
int n, tim = 0;
int a[N], level[N], tin[N], tout[N], rtin[N],
    nxt[N], subtree[N], parent[LG][N];
vector<int> g[N];
//Heavy Light Decomposition
void dfs(int u, int par, int lvl)
 parent[0][u] = par;
 level[u] = lvl:
 for (auto &it : g[u])
   if (it == par)
     continue:
   dfs(it, u, lvl + 1);
void dfs1(int u, int par)
 subtree[u] = 1;
```

```
for (auto &it : g[u])
   if (it == par)
     continue;
   dfs1(it, u);
   subtree[u] += subtree[it];
   if (subtree[it] > subtree[g[u][0]])
     swap(it, g[u][0]);
}
void dfs hld(int u. int par)
 tin[u] = ++tim:
 rtin[tim] = u;
 for (auto &v : g[u])
   if (v == par)
     continue;
   nxt[v] = (v == g[u][0] ? nxt[u] : v);
   dfs_hld(v, u);
  tout[u] = tim;
//LCA
int walk(int u, int h)
 for (int i = LG - 1: i >= 0: i--)
   if ((h >> i) & 1)
     u = parent[i][u];
 return u;
void precompute()
 for (int i = 1: i < LG: i++)</pre>
   for (int j = 1; j <= n; j++)</pre>
     if (parent[i - 1][j])
       parent[i][j] = parent[i - 1][parent[i -
            1][i]];
}
int LCA(int u, int v)
 if (level[u] < level[v])</pre>
   swap(u, v);
 int diff = level[u] - level[v];
 for (int i = LG - 1; i >= 0; i--)
   if ((1 << i) & diff)</pre>
   {
```

```
u = parent[i][u];
   }
 if (u == v)
   return u;
 for (int i = LG - 1; i >= 0; i--)
   if (parent[i][u] && parent[i][u] !=
        parent[i][v])
     u = parent[i][u];
     v = parent[i][v]:
 return parent[0][u];
int dist(int u. int v)
 return level[u] + level[v] - 2 * level[LCA(u, v)];
//Segment Tree
int st[4 * N], lazy[4 * N];
void build(int node, int L, int R)
 if (L == R)
   st[node] = a[rtin[L]]:
   return:
 int M = (L + R) / 2:
 build(node * 2, L, M);
 build(node * 2 + 1, M + 1, R);
 st[node] = min(st[node * 2], st[node * 2 + 1]):
}
void propagate(int node, int L, int R)
{
 if (L != R)
   lazy[node * 2] += lazy[node];
   lazy[node * 2 + 1] += lazy[node];
 st[node] += lazy[node];
 lazv[node] = 0;
int query(int node, int L, int R, int i, int j)
 if (lazv[node])
   propagate(node, L, R);
 if (i < L || i > R)
```

```
return 1e9;
 if (i <= L && R <= i)
   return st[node];
 int M = (L + R) / 2;
 int left = query(node * 2, L, M, i, j);
 int right = query(node * 2 + 1, M + 1, R, i, j);
 return min(left, right);
void update(int node, int L, int R, int i, int j,
    int val)
 if (lazv[node])
   propagate(node, L, R);
 if (j < L || i > R)
   return:
 if (i <= L && R <= j)</pre>
   lazv[node] += val;
   propagate(node, L, R);
   return;
 int M = (L + R) / 2;
 update(node * 2, L, M, i, j, val);
 update(node * 2 + 1, M + 1, R, i, j, val);
 st[node] = min(st[node * 2], st[node * 2 + 1]);
void upd(int 1, int r, int val)
 update(1, 1, n, l, r, val);
int get(int 1. int r)
 return query(1, 1, n, 1, r);
//Utility Functions
int query_up(int x, int y) //Assuming Y is an
     ancestor of X
 int res = 0;
 while (nxt[x] != nxt[v])
   res += get(tin[nxt[x]], tin[x]);
   x = parent[0][nxt[x]];
 res += get(tin[y] + 1, tin[x]); //use tin[y] to
      include Y
 return res;
int query_hld(int x, int y)
```

```
int lca = LCA(x, y);
 int res = query_up(x, lca) + query_up(y, lca);
 return res;
void update_up(int x, int y, int val) //Assuming Y
    is an ancestor of X
 while (nxt[x] != nxt[v])
   upd(tin[nxt[x]], tin[x], val);
   x = parent[0][nxt[x]];
 upd(tin[y] + 1, tin[x], val); //use tin[y] to
      include Y
void update_hld(int x, int y, int val)
 int lca = LCA(x, y);
 update_up(x, lca, val);
 update_up(y, lca, val);
void hld()
 dfs(1, 0, 1);
 dfs1(1, 0);
 dfs hld(1, 0):
 precompute();
 build(1, 1, n);
```

3.6 persistent array

```
if (1 == r) {
   ans-> val = what:
   return ans;
 int m = (1 + r) >> 1;
 if (at <= m) ans-> 1 = update(ans-> 1, 1, m, at,
      what):
 else ans-> r = update(ans-> r, m + 1, r, at,
      what):
 return ans:
int get(pnode cur, int 1, int r, int at) {
 if (cur == NULL) return 0:
 if (1 == r) return cur-> val:
 int m = (1 + r) >> 1:
 if (at <= m) return get(cur-> 1, 1, m, at);
            return get(cur-> r, m + 1, r, at);
 else
```

3.7 persistent seg tree

```
/* Problem: https://cses.fi/problemset/task/1737/
* Your task is to maintain a list of arrays which
     initially has a single array. You have to
     process the following types of queries:
* Query 1: Set the value a in array k to x.
* Query 2: Calculate the sum of values in range
     [a,b] in array k.
* Query 3: Create a copy of array k and add it to
     the end of the list.
* Idea to create a persistent segment tree to save
     all version of array.
vector <int> a;
struct Node{
   int val:
   Node *left, *right;
   Node(){
       left = right = NULL:
       val = 0:
   Node(Node* 1, Node *r, int v){
       left = 1;
       right = r;
       val = v;
};
void build(Node* &cur, int 1, int r){
```

```
if(1 == r){}
       cur->val = a[1];
       return;
   int mid = (l+r) >> 1;
   cur->left = new Node();
   cur->right = new Node();
   build(cur->left, 1, mid);
   build(cur->right, mid+1, r):
   cur->val = cur->left->val + cur->right->val:
}
void update(Node* prev, Node* &cur, int 1, int r,
    int i. int val){
   if(i < 1 || r < i)
       return:
   if(1 == r && 1 == i){
       cur->val = val:
       return;
   int mid = (l+r) >> 1;
   if(i \le mid){
       cur->right = prev->right;
       cur->left = new Node();
       update(prev->left, cur->left, 1, mid, i,
            val);
   }else{
       cur->left = prev->left;
       cur->right = new Node();
       update(prev->right, cur->right, mid+1, r, i,
            val):
   cur->val = cur->left->val + cur->right->val;
int get(Node* cur, int 1, int r, int u, int v){
   if(v < 1 \mid | r < u)
       return 0:
   if(u <= 1 && r <= v){
       return cur->val:
   int mid = (l+r) >> 1;
   int L = get(cur->left, l, mid, u, v);
   int R = get(cur->right, mid+1, r, u, v);
   return L + R;
}
Node* ver[MAXN];
```

3.8 persistent segment (v2)

/*

```
Find distinct numbers in a range (online query
        with persistent array)
struct Node{
   int lnode, rnode;
   int sum;
   Node(){
       lnode = rnode = sum = 0;
\ \text{Ver}[MAXN * 120]:
int sz = 0;
int build new node(int 1, int r){
   int next = ++sz:
   if(1 != r){
       int mid = (l+r) >> 1:
       ver[next].lnode = build new node(1, mid):
       ver[next].rnode = build new node(mid+1, r);
   return next;
int update(int cur, int 1, int r, int pos, int val){
   int next = ++sz;
   ver[next] = ver[cur];
   if(1 == r){}
       ver[next].sum = val;
       return next;
   else{
       int mid = (l+r) >> 1:
       if(pos <= mid)</pre>
           ver[next].lnode = update(ver[cur].lnode,
                1. mid. pos. val):
           ver[next].rnode = update(ver[cur].rnode
                . mid+1. r. pos. val):
   ver[next].sum = ver[ver[next].lnode].sum +
        ver[ver[next].rnode].sum;
   //cout << 1 << ', ' << r << ', ' << ver[next].sum
        << '\n';
   return next;
int get(int cur, int 1, int r, int u, int v){
   if(r < u \mid |v < 1)
       return 0:
   if(u <= 1 && r <= v){</pre>
       return ver[cur].sum:
   int mid = (1+r) \gg 1:
   return get(ver[cur].lnode, 1, mid, u, v) +
          get(ver[cur].rnode, mid+1, r, u, v);
```

3.9 persistent trie

```
// both tries can be tested with the problem:
    http://codeforces.com/problemset/problem/916/D
// Persistent binary trie (BST for integers)
const int MD = 31:
struct node_bin {
 node_bin *child[2];
 int val;
 node bin() : val(0) {
   child[0] = child[1] = NULL:
typedef node bin* pnode bin:
pnode_bin copy_node(pnode_bin cur) {
 pnode bin ans = new node bin():
 if (cur) *ans = *cur:
 return ans:
}
pnode_bin modify(pnode_bin cur, int key, int inc,
    int id = MD) {
 pnode_bin ans = copy_node(cur);
 ans->val += inc;
 if (id >= 0) {
   int to = (key >> id) & 1;
   ans->child[to] = modify(ans->child[to], key,
        inc, id - 1);
 return ans;
int sum_smaller(pnode_bin cur, int key, int id =
 if (cur == NULL) return 0:
 if (id < 0) return 0: // strictly smaller</pre>
 // if (id == - 1) return cur->val: // smaller or
      equal
 int ans = 0:
 int to = (key >> id) & 1;
   if (cur->child[0]) ans += cur->child[0]->val;
   ans += sum_smaller(cur->child[1], key, id - 1);
   ans = sum_smaller(cur->child[0], key, id - 1);
 return ans;
```

```
// Persistent trie for strings.
const int MAX_CHILD = 26;
struct node {
 node *child[MAX_CHILD];
 int val:
 node() : val(-1) {
   for (int i = 0: i < MAX CHILD: i++) {</pre>
     child[i] = NULL:
 }
}:
typedef node* pnode;
pnode copy_node(pnode cur) {
 pnode ans = new node():
 if (cur) *ans = *cur;
 return ans:
pnode set_val(pnode cur, string &key, int val, int
    id = 0) {
 pnode ans = copy_node(cur);
 if (id >= int(key.size())) {
   ans->val = val;
 } else {
   int t = kev[id] - 'a':
   ans->child[t] = set_val(ans->child[t], kev,
        val. id + 1):
 return ans;
pnode get(pnode cur, string &key, int id = 0) {
 if (id >= int(kev.size()) || !cur)
   return cur:
 int t = kev[id] - 'a':
 return get(cur->child[t], kev, id + 1):
```

3.10 segment tree

```
// Problem:
    https://codeforces.com/edu/course/2/lesson/4/1/practic
struct SegmentTree {
    #define m ((1 + r) >> 1)
    #define lc (i << 1)
    #define rc (i << 1 | 1)
        vector<int> mn;
    int n;
```

```
SegmentTree(int n = 0) : n(n){
       mn.resize(4 * n + 1, 0);
   SegmentTree(const vector<int> &a) : n(a.size())
       mn.resize(4 * n + 1, 0):
       function<void(int, int, int)> build =
            [%](int i, int 1, int r){
           if (1 == r){
              mn[i] = a[l - 1]:
              return:
           build(lc, l, m); build(rc, m + 1, r);
           mn[i] = min(mn[lc], mn[rc]);
       }:
       build(1, 1, n);
   void update(int i, int l, int r, int p, long
        val){
       if (1 == r){}
           mn[i] = val;
           return;
       }
       if (p <= m) update(lc, l, m, p, val);</pre>
       else update(rc, m + 1, r, p, val);
       mn[i] = min(mn[lc], mn[rc]):
   int get(int i, int l, int r, int u, int v){
       if (v < 1 \mid | r < u) return INF:
       if (u <= 1 && r <= v) return mn[i]:
       return min(get(lc, l, m, u, v), get(rc, m +
            1. r. u. v)):
   }
   void update(int p. long val){
       update(1, 1, n, p, val);
   int get(int 1, int r){
       return get(1, 1, n, 1, r);
#undef m
#undef lc
#undef rc
// Problem: There are two operations:
// 1 l r val: add the value val to the segment from
    1 to r
```

```
// 2 l v: calculate the minimum of elements from l
    to r
struct LazySegmentTree {
#define m ((1 + r) >> 1)
#define lc (i << 1)
#define rc (i << 1 | 1)
   vector<int> mn, lazy;
   int n:
   LazvSegmentTree(int n = 0) : n(n){
       mn.resize(4 * n + 1, 0):
       lazv.resize(4 * n + 1, 0):
   }
   void push(int i, int l, int r){
       if (lazv[i] == 0) return;
       mn[i] += lazv[i]:
       if (1 != r){
          lazv[lc] += lazv[i];
          lazy[rc] += lazy[i];
       lazy[i] = 0:
   }
   void update(int i, int l, int r, int u, int v,
        int val){
       push(i, 1, r);
       if (v < 1 || r < u) return;
       if (u \le 1 \&\& r \le v)
           lazv[i] += val:
           push(i, 1, r);
          return:
       update(lc. l. m. u. v. val): update(rc. m +
            1, r, u, v, val);
       mn[i] = min(mn[lc], mn[rc]):
   int get(int i, int l, int r, int u, int v){
       push(i, 1, r):
       if (v < 1 | | r < u) return INF:
       if (u <= 1 && r <= v) return mn[i];</pre>
       return min(get(lc, l, m, u, v), get(rc, m +
           1, r, u, v));
   }
   void update(int 1, int r, int val){
       update(1, 1, n, 1, r, val);
   int get(int 1, int r){
       return get(1, 1, n, 1, r);
#undef m
```

#undef lc

```
#undef rc
};
```

3.11 sparse table

```
template <typename T, typename func =</pre>
    function<T(const T, const T)>>
struct SparseTable {
   func calc;
   int n:
   vector<vector<T>> ans;
   SparseTable() {}
   SparseTable(const vector<T>& a, const func& f)
        : n(a.size()), calc(f) {
       int last = trunc(log2(n)) + 1;
       ans.resize(n);
       for (int i = 0: i < n: i++){</pre>
           ans[i].resize(last):
       for (int i = 0: i < n: i++){</pre>
           ans[i][0] = a[i]:
       for (int i = 1: i < last: i++){</pre>
           for (int i = 0; i \le n - (1 \le j); i++){
               ans[i][j] = calc(ans[i][j-1],
                    ans[i + (1 << (i - 1))][i - 1]):
       }
   T query(int 1, int r){
       assert(0 <= 1 && 1 <= r && r < n);
       int k = trunc(log2(r - 1 + 1)):
       return calc(ans[1][k], ans[r - (1 \ll k) +
            1][k]);
   }
};
```

3.12 trie

```
const int MN = 26; // size of alphabet
const int MS = 100010; // Number of states.

struct trie{
    struct node{
    int c;
    int a[MN];
```

```
};
 node tree[MS];
 int nodes:
 void clear(){
   tree[nodes].c = 0:
   memset(tree[nodes].a, -1, sizeof tree[nodes].a);
   nodes++:
 void init(){
   nodes = 0:
   clear():
 int add(const string &s, bool query = 0){
   int cur node = 0:
   for(int i = 0; i < s.size(); ++i){</pre>
     int id = gid(s[i]);
     if(tree[cur_node].a[id] == -1){
       if(query) return 0;
       tree[cur_node].a[id] = nodes;
       clear();
     cur_node = tree[cur_node].a[id];
   if(!query) tree[cur_node].c++;
   return tree[cur_node].c;
};
```

4 Geometry

4.1 center 2 points + radious

```
ans.push_back(e + x);
return ans;
}
```

4.2 closest pair problem

```
struct point {
 double x, y;
 int id;
 point() {}
 point (double a, double b) : x(a), y(b) {}
double dist(const point &o, const point &p) {
 double a = p.x - o.x, b = p.y - o.y;
 return sqrt(a * a + b * b):
double cp(vector<point> &p. vector<point> &x.
    vector<point> &y) {
 if (p.size() < 4) {</pre>
   double best = 1e100:
   for (int i = 0; i < p.size(); ++i)</pre>
    for (int j = i + 1; j < p.size(); ++j)</pre>
       best = min(best, dist(p[i], p[j]));
   return best;
 int ls = (p.size() + 1) >> 1;
 double l = (p[ls - 1].x + p[ls].x) * 0.5;
 vector<point> xl(ls), xr(p.size() - ls);
 unordered set<int> left:
 for (int i = 0; i < ls; ++i) {</pre>
   xl[i] = x[i]:
   left.insert(x[i].id):
 for (int i = ls; i < p.size(); ++i) {</pre>
   xr[i - ls] = x[i]:
 vector<point> yl, yr;
 vector<point> pl, pr;
 yl.reserve(ls); yr.reserve(p.size() - ls);
 pl.reserve(ls); pr.reserve(p.size() - ls);
 for (int i = 0; i < p.size(); ++i) {</pre>
   if (left.count(y[i].id))
     yl.push_back(y[i]);
     yr.push_back(y[i]);
   if (left.count(p[i].id))
     pl.push_back(p[i]);
```

```
pr.push_back(p[i]);
  double dl = cp(pl, xl, yl);
  double dr = cp(pr, xr, yr);
  double d = min(dl, dr);
  vector<point> yp; yp.reserve(p.size());
  for (int i = 0; i < p.size(); ++i) {</pre>
   if (fabs(y[i].x - 1) < d)
     yp.push_back(y[i]);
  for (int i = 0; i < yp.size(); ++i) {</pre>
   for (int j = i + 1; j < yp.size() && j < i + 7;</pre>
        ++i) {
     d = min(d, dist(vp[i], vp[i]));
  return d;
double closest_pair(vector<point> &p) {
  vector<point> x(p.begin(), p.end());
  sort(x.begin(), x.end(), [](const point &a, const
      point &b) {
   return a.x < b.x;</pre>
  vector<point> y(p.begin(), p.end());
  sort(y.begin(), y.end(), [](const point &a, const
      point &b) {
   return a.y < b.y;</pre>
  return cp(p, x, y);
```

4.3 convex diameter

```
struct point{
   int x, y;
};

struct vec{
   int x, y;
};

vec operator - (const point &A, const point &B){
   return vec{A.x - B.x, A.y - B.y};
}

int cross(vec A, vec B){
   return A.x*B.y - A.y*B.x;
}
```

```
int cross(point A, point B, point C){
   int val = A.x*(B.y - C.y) + B.x*(C.y - A.y) +
        C.x*(A.y - B.y);
   if(val == 0)
       return 0; // coline
   if(val < 0)
       return 1; // clockwise
   return -1: //counter clockwise
}
vector <point> findConvexHull(vector <point>
    points){
   vector <point> convex;
   sort(points.begin(), points.end(), [](const
        point &A. const point &B){
       return (A.x == B.x)? (A.y < B.y): (A.x <
            B.x);
   });
   vector <point> Up, Down;
   point A = points[0], B = points.back();
   Up.push_back(A);
   Down.push_back(A);
   for(int i = 0; i < points.size(); i++){</pre>
       if(i == points.size()-1 || cross(A,
            points[i], B) > 0){
           while(Up.size() > 2 &&
                cross(Up[Up.size()-2],
               Up[Up.size()-1], points[i]) <= 0)</pre>
              Up.pop_back();
           Up.push_back(points[i]);
       }
       if(i == points.size()-1 || cross(A,
            points[i], B) < 0){
           while(Down.size() > 2 &&
                cross(Down[Down.size()-2].
               Down[Down.size()-1], points[i]) >=
              Down.pop_back();
           Down.push_back(points[i]);
       }
   }
   for(int i = 0; i < Up.size(); i++)</pre>
        convex.push_back(Up[i]);
   for(int i = Down.size()-2; i > 0; i--)
        convex.push_back(Down[i]);
   return convex;
}
int dist(point A, point B){
   return (A.x - B.x)*(A.x - B.x) + (A.y -
        B.y)*(A.y - B.y);
}
```

```
double findConvexDiameter(vector <point>
    convexHull){
   int n = convexHull.size();
   int is = 0, js = 0;
   for(int i = 1; i < n; i++){</pre>
       if(convexHull[i].y > convexHull[is].y)
       if(convexHull[is].v > convexHull[i].v)
          is = i:
   }
   int maxd = dist(convexHull[is], convexHull[is]);
   int i, maxi, j, maxj;
   i = maxi = is:
   i = maxi = is:
   dof
       int ni = (i+1)%n, nj = (j+1)%n;
       if(cross(convexHull[ni] - convexHull[i],
            convexHull[nj] - convexHull[j]) <= 0){</pre>
          i = ni;
       }else{
          i = ni;
       int d = dist(convexHull[i], convexHull[j]);
       if(d > maxd){
          maxd = d;
          maxi = i;
          maxj = j;
   }while(i != is || j != js);
   return sqrt(maxd);
```

4.4 pick theorem

```
return abs(area/2.0);
}

11 boundary(vector <point> points){
    int n = (int)points.size();
    11 num_bound = 0;
    for(int i = 0; i < n; i++){
        11 dx = (points[i].x - points[(i+1)%n].x);
        11 dy = (points[i].y - points[(i+1)%n].y);
        num_bound += abs(__gcd(dx, dy)) - 1;
    }
    return num_bound;
}</pre>
```

4.5 squares

```
typedef long double ld;
const ld eps = 1e-12;
int cmp(ld x, ld y = 0, ld tol = eps) {
   return ( x \le v + tol) ? (x + tol \le v) ? -1 : 0
        : 1:
}
struct point{
 ld x, v;
 point(ld a, ld b) : x(a), y(b) {}
 point() {}
};
struct square{
 ld x1, x2, y1, y2,
    a, b, c;
  point edges[4];
  square(ld _a, ld _b, ld _c) {
   a = _a, b = _b, c = _c;
   x1 = a - c * 0.5:
   x2 = a + c * 0.5:
   v1 = b - c * 0.5:
   v2 = b + c * 0.5:
   edges[0] = point(x1, v1):
   edges[1] = point(x2, y1);
   edges[2] = point(x2, y2);
   edges[3] = point(x1, y2);
};
ld min_dist(point &a, point &b) {
 1d x = a.x - b.x,
    y = a.y - b.y;
  return sqrt(x * x + y * y);
```

```
}
bool point_in_box(square s1, point p) {
 if (cmp(s1.x1, p.x) != 1 && cmp(s1.x2, p.x) != -1
     cmp(s1.v1, p.v) != 1 && cmp(s1.v2, p.v) != -1)
   return true:
 return false:
bool inside(square &s1, square &s2) {
 for (int i = 0: i < 4: ++i)
   if (point in box(s2, s1.edges[i]))
     return true:
 return false;
}
bool inside_vert(square &s1, square &s2) {
 if ((cmp(s1.v1, s2.v1) != -1 && cmp(s1.v1, s2.v2)
      != 1) ||
     (cmp(s1.y2, s2.y1) != -1 \&\& cmp(s1.y2, s2.y2)
         != 1))
   return true;
return false;
bool inside_hori(square &s1, square &s2) {
 if ((cmp(s1.x1, s2.x1) != -1 \&\& cmp(s1.x1, s2.x2))
     (cmp(s1.x2, s2.x1) != -1 \&\& cmp(s1.x2, s2.x2)
          != 1))
   return true:
return false:
ld min dist(square &s1, square &s2) {
 if (inside(s1, s2) || inside(s2, s1))
   return 0:
 ld ans = 1e100:
 for (int i = 0; i < 4; ++i)
   for (int j = 0; j < 4; ++j)
     ans = min(ans, min_dist(s1.edges[i],
          s2.edges[i]));
 if (inside_hori(s1, s2) || inside_hori(s2, s1)) {
   if (cmp(s1.v1, s2.v2) != -1)
     ans = min(ans, s1.v1 - s2.v2);
   if (cmp(s2.v1, s1.v2) != -1)
     ans = min(ans. s2.v1 - s1.v2):
 }
```

```
if (inside_vert(s1, s2) || inside_vert(s2, s1)) {
   if (cmp(s1.x1, s2.x2) != -1)
      ans = min(ans, s1.x1 - s2.x2);
   else
   if (cmp(s2.x1, s1.x2) != -1)
      ans = min(ans, s2.x1 - s1.x2);
}
return ans;
}
```

4.6 template

```
#define EPS 1e-6
const double PI = acos(-1.0):
double DEG TO RAD(double d) { return d * PI /
    180.0: }
double RAD TO DEG(double r) { return r * 180.0 /
inline int cmp(double a, double b) {
   return (a < b - EPS) ? -1 : ((a > b + EPS) ? 1
       : 0):
struct Point{
   double x, y;
   Point(){
      x = y = 0.0;
   Point(double x, double y): x(x), y(y) {}
   Point operator + (const Point& a) const {
        return Point(x+a.x, v+a.v): }
   Point operator - (const Point& a) const {
       return Point(x-a.x. v-a.v): }
   Point operator * (double k) const { return
        Point(x*k, y*k); }
   Point operator / (double k) const { return
        Point(x/k, v/k): }
   double dot(const Point& a) const { return x*a.x
        + v*a.v: } // dot product
   double cross(const Point& a) const { return
       x*a.y - y*a.x; } // cross product
   int cmp(const Point& q) const {
      if (x != q.x) return ::cmp(x, q.x);
      <u>return</u> ::cmp(y, q.y);
```

```
#define Comp(x) bool operator x (Point q) const
        { return cmp(q) x 0; }
   Comp(>) Comp(<) Comp(==) Comp(>=) Comp(<=)</pre>
        Comp(!=)
   #undef Comp
   double norm() { return x*x + y*y; }
   double len() { return sqrt(norm()); }
   // Rotate vector
   Point rotate(double alpha) {
       double cosa = cos(alpha). sina = sin(alpha):
       return Point(x * cosa - v * sina. x * sina +
            v * cosa):
   }
}:
istream& operator >> (istream& cin, Point& p) {
   cin >> p.x >> p.y;
   return cin;
ostream& operator << (ostream& cout, Point& p) {
   cout << p.x << ' ' << p.y;
   return cout;
struct Line{
   double a, b, c;
   Point A. B:
   Line(double a, double b, double c): a(a), b(b),
        c(c) {}
   Line(Point A. Point B): A(A), B(B) {
       a = B.v - A.v:
       b = A.x - B.x:
       c = -(a * A.x + b * A.v):
   // initialize a line with slope k
   Line(Point P. double k) {
       a = -k;
       b = 1;
       c = k * P.x - P.y;
   double f(Point A){
       return a * A.x + b * A.y + c;
};
bool areParallel(Line 11, Line 12) {
   return cmp(l1.a*l2.b, l1.b*l2.a) == 0:
```

```
bool areSame(Line 11, Line 12) {
   return areParallel(11, 12) && cmp(11.c*12.a,
        12.c*11.a) == 0
          && cmp(11.c*12.b, 11.b*12.c) == 0;
}
bool areIntersect(Line 11, Line 12, Point &p) {
   if (areParallel(11, 12))
       return false:
   double dx = 11.b*12.c - 12.b*11.c:
   double dy = 11.c*12.a - 12.c*11.a;
   double d = 11.a*12.b - 12.a*11.b:
   p = Point(dx / d, dv / d):
   return true:
}
// distance from p to line ab
double distToLine(Point p, Point a, Point b, Point
   Point ap = p - a, ab = b - a;
   double k = ap.dot(ab) / ab.norm();
   c = a + (ab * k):
   return (p - c).len();
// closest point from p in line 1.
void closestPoint(Line 1, Point p, Point &ans) {
   if (fabs(1.b) < EPS) {</pre>
       ans.x = -(1.c) / 1.a; ans.y = p.y;
       return:
   }
   if (fabs(l.a) < EPS) {</pre>
       ans.x = p.x; ans.y = -(1.c) / 1.b;
       return:
   Line perp(1.b, -1.a, - (1.b*p.x - 1.a*p.y));
   areIntersect(1, perp, ans);
}
// reflect point p over line 1
void reflectionPoint(Line 1, Point p, Point &ans) {
   closestPoint(1, p, b);
   ans = p + (b - p) * 2;
}
```

4.7 triangles

Let a, b, c be length of the three sides of a triangle.

$$p = (a+b+c)*0.5$$

The inradius is defined by:

$$iR = \sqrt{\frac{(p-a)(p-b)(p-c)}{p}}$$

The radius of its circumcircle is given by the formula:

$$cR = \frac{abc}{\sqrt{(a+b+c)(a+b-c)(a+c-b)(b+c-a)}}$$

5 Graphs

5.1 bridges

```
struct Graph {
 vector<vector<Edge>> g;
 vector<int> vi, low, d, pi, is_b;
 int bridges computed:
 int ticks, edges;
 Graph(int n, int m) {
   g.assign(n, vector<Edge>());
   is_b.assign(m, 0);
   vi.resize(n);
   low.resize(n);
   d.resize(n);
   pi.resize(n);
   edges = 0;
   bridges_computed = 0;
 void AddEdge(int u, int v) {
   g[u].push back(Edge(v. edges)):
   g[v].push_back(Edge(u, edges));
   edges++:
 void Dfs(int u) {
   vi[u] = true:
   d[u] = low[u] = ticks++:
   for (int i = 0: i < (int)g[u].size(): ++i) {</pre>
     int v = g[u][i].to;
     if (v == pi[u]) continue;
     if (!vi[v]) {
      pi[v] = u;
       Dfs(v):
       if (d[u] < low[v]) is_b[g[u][i].id] = true;</pre>
      low[u] = min(low[u], low[v]);
     } else {
```

```
low[u] = min(low[u], d[v]);
 // Multiple edges from a to b are not allowed.
 // (they could be detected as a bridge).
 // If you need to handle this, just count
 // how many edges there are from a to b.
 void CompBridges() {
   fill(pi.begin(), pi.end(), -1);
   fill(vi.begin(), vi.end(), 0):
   fill(low.begin(), low.end(), 0):
   fill(d.begin(), d.end(), 0);
   ticks = 0:
   for (int i = 0: i < (int)g.size(): ++i)</pre>
     if (!vi[i]) Dfs(i):
   bridges_computed = true;
 map<int, vector<Edge>> BridgesTree() {
   if (!bridges_computed) CompBridges();
   int n = g.size();
   Dsu dsu(g.size());
   for (int i = 0; i < n; i++)
     for (auto e : g[i])
       if (!is_b[e.id]) dsu.Join(i, e.to);
   map<int, vector<Edge>> tree;
   for (int i = 0; i < n; i++)
     for (auto e : g[i])
       if (is b[e.id])
        tree[dsu.Find(i)].emplace_back(dsu.Find(e.to),
              e.id):
   return tree;
};
```

5.2 delete on dsu

```
class dsu_rollback{
   public:
       vector <int> parent;
       int comps;
       stack <dsu_save> st_op;
       dsu_rollback(){};
       dsu rollback(int n){
           parent.resize(n+1, -1);
           comps = n;
       }
       int find set(int u){
           while(parent[u] > 0)
              u = parent[u]:
           return u;
       }
       bool Union(int u, int v){
           int U = find_set(u);
           int V = find_set(v);
           if(U == V)
              return false;
           comps--;
           st_op.push(dsu_save(U, parent[U], V,
               parent[V]));
           int x = parent[U] + parent[V];
           if(parent[U] > parent[V]){
              parent[U] = V;
              parent[V] = x;
           }else{
              parent[U] = x;
              parent[V] = U:
           return true;
       }
       void rollback(){
           if(st_op.empty())
              return;
           dsu_save x = st_op.top();
           st_op.pop();
           comps++;
           parent[x.u] = x.par_u;
           parent[x.v] = x.par_v;
       }
};
struct query{
   int u, v;
   bool united;
}:
class QuervTree{
```

```
vector <vector <query>> t;
   dsu_rollback dsu;
   int T;
   public:
       QueryTree(int _T, int n){
           this \rightarrow T = _T;
           this->dsu = dsu_rollback(n);
           t.resize(4*T + 4);
       void add_to_tree(int id, int 1, int r, int
            u. int v. querv a){
           if(v < 1 | | r < u | | u > v)
              return:
           if(u <= 1 && r <= v){</pre>
              t[id].push back(q):
              return;
           int mid = (l+r) >> 1;
           add_to_tree(2*id, 1, mid, u, v, q);
           add_to_tree(2*id+1, mid+1, r, u, v, q);
       void add_query(query q, int 1, int r){
           add_to_tree(1, 0, T-1, 1, r, q);
       void DFS(int id, int 1, int r, vector <int>
            &ans){
           for(query &q: t[id])
              q.united = dsu.Union(q.u, q.v);
           if(1 == r){
               ans[1] = dsu.comps:
              int mid = (l+r) >> 1:
              DFS(2*id, 1, mid, ans):
              DFS(2*id+1, mid+1, r, ans);
           for(query &q: t[id])
               if (q.united)
                  dsu.rollback();
       vector <int> compute(){
           vector <int> ans(T); // T query
           DFS(1, 0, T-1, ans);
           return ans;
};
```

5.3 euler path

```
struct DirectedEulerPath
       int n;
       vector<vector<int> > g;
       vector<int> path;
       void init(int _n){
              n = _n;
              g = vector < vector < int > (n + 1,
                   vector<int> ());
              path.clear();
       }
       void add_edge(int u, int v){
              g[u].push_back(v);
       void dfs(int u)
               while(g[u].size())
                      int v = g[u].back();
                      g[u].pop_back();
                      dfs(v);
              path.push_back(u);
       }
       bool getPath(){
               int ctEdges = 0;
               vector<int> outDeg, inDeg;
              outDeg = inDeg = vector<int> (n + 1,
              for(int i = 1: i <= n: i++)
                      ctEdges += g[i].size();
                      outDeg[i] += g[i].size();
                      for(auto &u:g[i])
                             inDeg[u]++:
               int ctMiddle = 0. src = 1:
              for(int i = 1: i <= n: i++)</pre>
                      if(abs(inDeg[i] - outDeg[i])
                           > 1)
                             return 0;
                      if(inDeg[i] == outDeg[i])
                             ctMiddle++;
                      if(outDeg[i] > inDeg[i])
                             src = i;
              }
```

5.4 karp min mean cycle

```
* Finds the min mean cycle, if you need the max
     mean cvcle
 * just add all the edges with negative cost and
     print
 * ans * -1
 * test: uva. 11090 - Going in Cycle!!
const int MN = 1000:
struct edge{
 int v;
 long long w;
 edge()\{\} edge(int v, int w) : v(v), w(w) \{\}
};
long long d[MN][MN];
// This is a copy of g because increments the size
// pass as reference if this does not matter.
int karp(vector<vector<edge> > g) {
 int n = g.size();
 g.resize(n + 1): // this is important
 for (int i = 0; i < n; ++i)
   if (!g[i].emptv())
     g[n].push_back(edge(i,0));
 for(int i = 0:i<n:++i)</pre>
   fill(d[i],d[i]+(n+1),INT_MAX);
 d[n - 1][0] = 0:
 for (int k = 1; k \le n; ++k) for (int u = 0; u \le n
      n: ++u) {
   if (d[u][k - 1] == INT_MAX) continue;
   for (int i = g[u].size() - 1; i >= 0; --i)
     d[g[u][i].v][k] = min(d[g[u][i].v][k], d[u][k]
          - 1] + g[u][i].w);
```

```
bool flag = true;
for (int i = 0; i < n && flag; ++i)</pre>
 if (d[i][n] != INT_MAX)
   flag = false;
if (flag) {
 return true; // return true if there is no a
double ans = 1e15:
for (int u = 0: u + 1 < n: ++u) {
 if (d[u][n] == INT_MAX) continue;
 double W = -1e15:
 for (int k = 0; k < n; ++k)
   if (d[u][k] != INT_MAX)
     W = max(W, (double)(d[u][n] - d[u][k]) / (n
          - k)):
 ans = min(ans, W);
// printf("%.21f\n", ans);
cout << fixed << setprecision(2) << ans << endl;</pre>
return false:
```

5.5 konig's theorem

In any bipartite graph, the number of edges in a maximum matching equals the number of vertices in a minimum vertex cover

5.6 matching

```
struct Hopcroft_Karp
{
    static const int inf = 1e9;
    int n;
    vector<int> matchL, matchR, dist;
    vector<vector<int> > g;
    Hopcroft_Karp(int n) :
```

```
n(n), matchL(n+1), matchR(n+1),
            dist(n+1), g(n+1) {}
void addEdge(int u, int v)
       g[u].push_back(v);
}
bool bfs()
       queue<int> q;
       for(int u=1:u<=n:u++)
              if(!matchL[u])
                      dist[u]=0:
                      q.push(u);
              }
              else
                      dist[u]=inf;
       dist[0]=inf;
       while(!q.empty())
              int u=q.front();
              q.pop();
              for(auto v:g[u])
                      if(dist[matchR[v]] ==
                          inf)
                             dist[matchR[v]]
                                  = dist[u]
                                  + 1:
                             q.push(matchR[v]);
                     }
              }
       }
       return (dist[0]!=inf);
}
bool dfs(int u)
       if(!u)
              return true;
       for(auto v:g[u])
              if(dist[matchR[v]] ==
                   dist[u]+1
                   &&dfs(matchR[v]))
                      matchL[u]=v;
                      matchR[v]=u:
```

```
return true;
                       }
               }
               dist[u]=inf;
               return false;
       }
       int max_matching()
               int matching=0:
               while(bfs())
                       for(int u=1:u<=n:u++)</pre>
                       {
                               if(!matchL[u])
                                      if(dfs(u))
                                              matching++:
               return matching;
       }
};
```

5.7 max flow min cost

```
struct edge
{
       long long x, y, cap, flow, cost;
};
struct MinCostMaxFlow
{
       long long n, S, T;
       vector < vector <long long> > a;
       vector <long long> dist, prev, done, pot;
       vector <edge> e;
       MinCostMaxFlow() {}
       MinCostMaxFlow(long long _n, long long _S,
            long long _T)
       {
               n = _n; S = _S; T = _T;
               a = vector < vector <long long> >(n
                   + 1):
               dist = vector <long long>(n + 1);
               prev = vector <long long>(n + 1);
               done = vector <long long>(n + 1);
               pot = vector \langle long long \rangle (n + 1, 0);
       }
       void addEdge(long long x, long long y, long
            long _cap, long long _cost)
```

```
{
       edge e1 = \{x, y, \_cap, 0, \_cost\};
       edge e2 = \{y, x, 0, 0, -\_cost\};
       a[x].push_back(e.size());
            e.push_back(e1);
       a[v].push_back(e.size());
            e.push_back(e2);
}
pair <long long,long long> dijkstra()
       long long flow = 0, cost = 0:
       for (long long i = 1: i <= n: i++)
            done[i] = 0. dist[i] = oo:
       priority_queue < pair<long long,long</pre>
            long> > a:
       dist[S] = 0; prev[S] = -1;
       q.push(make_pair(0, S));
       while (!q.empty())
               long long x = q.top().second;
                    q.pop();
               if (done[x]) continue;
               done[x] = 1;
               for (int i = 0; i <</pre>
                    int(a[x].size()); i++)
                       long long id =
                           a[x][i], y =
                           e[id].y;
                       if (e[id].flow <</pre>
                            e[id].cap)
                              long long D =
                                   dist[x] +
                                   e[id].cost
                                   + pot[x] -
                                   pot[y];
                              if (!done[y] &&
                                   D <
                                   dist[y])
                                      dist[y]
                                          D;
                                           prev[v]
                                           id;
                                          y));
                              }
                      }
              }
       }
```

```
for (long long i = 1; i <= n; i++)
                   pot[i] += dist[i];
              if (done[T])
                      flow = oo;
                      for (long long id = prev[T];
                          id >= 0; id =
                          prev[e[id].x])
                             flow = min(flow.
                                  e[id].cap -
                                  e[id].flow):
                      for (long long id = prev[T]:
                           id >= 0: id =
                          prev[e[id].x])
                             cost += e[id].cost *
                                  flow:
                             e[id].flow += flow;
                             e[id ^ 1].flow -= flow;
                      }
              }
              return make_pair(flow, cost);
       }
       pair <long long,long long> minCostMaxFlow()
              long long totalFlow = 0, totalCost =
              while (1)
                      pair <long long,long long> u
                          = diikstra():
                      if (!done[T]) break;
                      totalFlow += u.first:
                          totalCost += u.second:
              return make_pair(totalFlow,
                   totalCost):
       }
};
```

5.8 minimum path cover in DAG

q.push(make_pair(Gisehya, directed acyclic graph G = (V, E), we are to find the minimum number of vertex-disjoint paths to cover each vertex in V.

We can construct a bipartite graph $G' = (Vout \cup Vin, E')$ from G, where :

```
Vout = \{v \in V : v \text{ has positive out } - degree\}
Vin = \{v \in V : v \text{ has positive } in - degree\}
E' = \{(u, v) \in Vout \times Vin : (u, v) \in E\}
```

Then it can be shown, via König's theorem, that G has a matching of size m if and only if there exists n-m vertex-disjoint paths that cover each vertex in G, where n is the number of vertices in G and m is the maximum cardinality bipartite mathching in G.

Therefore, the problem can be solved by finding the maximum cardinality matching in G' instead.

NOTE: If the paths are note necessarily disjoints, find the transitive closure and solve the problem for disjoint paths.

5.9 planar graph (euler)

Euler's formula states that if a finite, connected, planar graph is drawn in the plane without any edge intersections, and v is the number of vertices, e is the number of edges and f is the number of faces (regions bounded by edges, including the outer, infinitely large region), then:

$$f + v = e + 2$$

It can be extended to non connected planar graphs with c connected components:

$$f + v = e + c + 1$$

5.10 two sat (with kosaraju)

```
/**

* Given a set of clauses (a1 v a2)^(a2 v a3)....

* this algorithm find a solution to it set of clauses.

* test:
http://lightoj.com/volume_showproblem.php?problem=1251

**/

#include<bits/stdc++.h>
using namespace std;
#define MAX 100000
```

```
#define endl '\n'
vector<int> G[MAX];
vector<int> GT[MAX];
vector<int> Ftime;
vector<vector<int> > SCC;
bool visited[MAX];
int n:
void dfs1(int n){
 visited[n] = 1:
 for (int i = 0: i < G[n].size(): ++i) {</pre>
   int curr = G[n][i]:
   if (visited[curr]) continue:
   dfs1(curr):
 Ftime.push_back(n);
void dfs2(int n, vector<int> &scc) {
 visited[n] = 1;
 scc.push_back(n);
 for (int i = 0;i < GT[n].size(); ++i) {</pre>
   int curr = GT[n][i];
   if (visited[curr]) continue:
   dfs2(curr. scc):
}
void kosaraju() {
 memset(visited, 0, sizeof visited);
 for (int i = 0; i < 2 * n; ++i) {</pre>
   if (!visited[i]) dfs1(i):
 memset(visited, 0, sizeof visited);
 for (int i = Ftime.size() - 1; i >= 0; i--) {
   if (visited[Ftime[i]]) continue;
   vector<int> _scc;
   dfs2(Ftime[i],_scc);
   SCC.push_back(_scc);
* After having the SCC, we must traverse each scc,
     if in one SCC are -b v b, there is not a
     solution.
```

```
* Otherwise we build a solution, making the first
      "node" that we find truth and its complement
bool two_sat(vector<int> &val) {
 kosaraju();
  for (int i = 0: i < SCC.size(): ++i) {</pre>
   vector<bool> tmpvisited(2 * n, false);
   for (int j = 0; j < SCC[i].size(); ++j) {</pre>
     if (tmpvisited[SCC[i][j] ^ 1]) return 0;
     if (val[SCC[i][i]] != -1) continue:
       val[SCC[i][j]] = 0;
       val[SCC[i][i] ^ 1] = 1:
     tmpvisited[SCC[i][j]] = 1;
  return 1;
// Example of use
int main() {
  int m, u, v, nc = 0, t; cin >> t;
 // n = "nodes" number. m = clauses number
  while (t--) {
   cin >> m >> n:
   Ftime.clear():
   SCC.clear():
   for (int i = 0; i < 2 * n; ++i) {</pre>
     G[i].clear():
     GT[i].clear():
   // (a1 v a2) = (a1 -> a2) = (a2 -> a1)
   for (int i = 0; i < m ; ++i) {</pre>
     cin >> u >> v;
     int t1 = abs(u) - 1;
     int t2 = abs(v) - 1;
     int p = t1 * 2 + ((u < 0)? 1 : 0);
     int q = t2 * 2 + ((v < 0)? 1 : 0);
     G[p ^ 1].push_back(q);
     G[q ^ 1].push_back(p);
     GT[p].push_back(q ^ 1);
     GT[q].push_back(p ^ 1);
   vector < int > val(2 * n. -1):
   cout << "Case " << ++nc <<": ";
   if (two sat(val)) {
```

```
cout << "Yes" << endl;
vector<int> sol;
for (int i = 0; i < 2 * n; ++i)
   if (i % 2 == 0 and val[i] == 1)
      sol.push_back(i / 2 + 1);
cout << sol.size();

for (int i = 0; i < sol.size(); ++i) {
   cout << " " << sol[i];
   }
   cout << endl;
} else {
   cout << "No" << endl;
}
return 0;</pre>
```

6 Math

6.1 Lucas theorem

For non-negative integers m and n and a prime p, the following congruence relation holds: :

$$\binom{m}{n} \equiv \prod_{i=0}^{k} \binom{m_i}{n_i} \pmod{p},$$

where :

$$m = m_k p^k + m_{k-1} p^{k-1} + \dots + m_1 p + m_0,$$

and:

$$n = n_k p^k + n_{k-1} p^{k-1} + \dots + n_1 p + n_0$$

are the base p expansions of m and n respectively. This uses the convention that $\binom{m}{n} = 0$ if $m \le n$.

6.2 cumulative sum of divisors

n, is defined as below:

```
/*
The function SOD(n) (sum of divisors) is defined as the summation of all the actual divisors of an integer number n. For example,

SOD(24) = 2+3+4+6+8+12 = 35.

The function CSOD(n) (cumulative SOD) of an integer
```

```
csod(n) = \sum_{i = 1}^{n} sod(i)

It can be computed in O(sqrt(n)):
*/

long long csod(long long n) {
  long long ans = 0;
  for (long long i = 2; i * i <= n; ++i) {
    long long j = n / i;
    ans += (i + j) * (j - i + 1) / 2;
    ans += i * (j - i);
  }
  return ans;
}</pre>
```

6.3 fft

```
* Fast Fourier Transform.
* Useful to compute convolutions.
* \hat{C}(f \operatorname{star} g)[n] = \operatorname{sum}_m(f[m] * g[n - m])
* for all n.
* test: icpc live archive, 6886 - Golf Bot
using namespace std;
#include <bits/stdc++.h>
#define D(x) cout << #x " = " << (x) << endl
#define endl '\n'
const int MN = 262144 << 1:</pre>
int d[MN + 10], d2[MN + 10];
const double PI = acos(-1.0):
struct cpx {
 double real, image;
 cpx(double _real, double _image) {
   real = real:
   image = _image;
 cpx(){}
cpx operator + (const cpx &c1, const cpx &c2) {
 return cpx(c1.real + c2.real, c1.image +
      c2.image);
```

```
cpx operator - (const cpx &c1, const cpx &c2) {
 return cpx(c1.real - c2.real, c1.image -
      c2.image);
cpx operator * (const cpx &c1, const cpx &c2) {
 return cpx(c1.real*c2.real - c1.image*c2.image,
      c1.real*c2.image + c1.image*c2.real);
int rev(int id, int len) {
 int ret = 0:
  for (int i = 0; (1 << i) < len; i++) {</pre>
   ret <<= 1:
   if (id & (1 << i)) ret |= 1:
 return ret;
cpx A[1 << 20];
void FFT(cpx *a, int len, int DFT) {
  for (int i = 0; i < len; i++)</pre>
   A[rev(i, len)] = a[i];
  for (int s = 1; (1 << s) <= len; s++) {
   int m = (1 << s);
   cpx wm = cpx(cos(DFT * 2 * PI / m), sin(DFT *
        2 * PI / m)):
   for(int k = 0; k < len; k += m) {</pre>
     cpx w = cpx(1, 0);
     for(int j = 0; j < (m >> 1); j++) {
       cpx t = w * A[k + j + (m >> 1)];
       cpx u = A[k + i]:
       A[k + j] = u + t;
       A[k + j + (m >> 1)] = u - t;
  if (DFT == -1) for (int i = 0; i < len; i++)</pre>
      A[i].real /= len, A[i].image /= len;
  for (int i = 0; i < len; i++) a[i] = A[i];</pre>
 return;
cpx in[1 << 20];
void solve(int n) {
 memset(d, 0, sizeof d);
  for (int i = 0; i < n; ++i) {</pre>
   cin >> t;
   d[t] = true;
```

```
int m;
 cin >> m;
 vector<int> q(m);
 for (int i = 0; i < m; ++i)</pre>
    cin >> q[i];
 for (int i = 0; i < MN; ++i) {</pre>
    if (d[i])
     in[i] = cpx(1, 0);
    else
      in[i] = cpx(0, 0);
 FFT(in, MN, 1):
 for (int i = 0; i < MN; ++i) {</pre>
   in[i] = in[i] * in[i]:
 FFT(in, MN, -1);
 int ans = 0;
 for (int i = 0; i < q.size(); ++i) {</pre>
    if (in[q[i]].real > 0.5 || d[q[i]]) {
     ans++;
  cout << ans << endl;</pre>
int main() {
 ios_base::sync_with_stdio(false);cin.tie(NULL);
 while (cin >> n)
    solve(n);
 return 0:
}
```

6.4 fibonacci properties

Let A, B and n be integer numbers.

$$k = A - B \tag{1}$$

$$F_A F_B = F_{k+1} F_A^2 + F_k F_A F_{A-1} \tag{2}$$

$$\sum_{i=0}^{n} F_i^2 = F_{n+1} F_n \tag{3}$$

ev(n) = returns 1 if n is even.

$$\sum_{i=0}^{n} F_i F_{i+1} = F_{n+1}^2 - ev(n) \tag{4}$$

$$\sum_{i=0}^{n} F_i F_{i-1} = \sum_{i=0}^{n-1} F_i F_{i+1}$$
 (5)

6.5 others

Approximate factorial

$$n! = \sqrt{2.\pi \cdot n} \cdot \left(\frac{n}{e}\right)^n \tag{6}$$

6.6 polynomials

```
// TODO: what's this ?
const double pi = acos(-1);
struct poly {
 deque <double> coef;
 double x_lo, x_hi;
 double evaluate(double x) {
   double ans = 0:
   for (auto it : coef)
     ans = (ans * x + it):
   return ans:
 }
 double volume(double x, double dx=1e-6) {
   dx = (x_hi - x_lo) / 1000000.0;
   double ans = 0;
   for (double ix = x lo: ix <= x: ix += dx) {
     double rad = evaluate(ix);
     ans += pi * rad * rad * dx;
   }
   return ans;
};
```

6.7 sigma function

the sigma function is defined as:

$$\sigma_x(n) = \sum_{d|n} d^x$$

when x = 0 is called the divisor function, that counts the number of positive divisors of n.

Now, we are interested in find

$$\sum_{d|n} \sigma_0(d)$$

if n is written as prime factorization:

$$n = \prod_{i=1}^{k} P_i^{e_k}$$

we can demonstrate that:

$$\sum_{d|n} \sigma_0(d) = \prod_{i=1}^k g(e_k + 1)$$

where g(x) is the sum of the first x positive numbers:

$$q(x) = (x * (x + 1))/2$$

6.8 system different constraints

```
/* http://poj.org/problem?id=2983 */
/*
Given a system of inequations of the form x_j - x_i
Find any solution x_1, x_2, \ldots, x_n or show that
     the system has no solution.
We construct a n-vertex graph (vertext i represents
     variable x_i). For each inequation x_j - x_i
we add an edge from i to j with weight w_ij.
If the graph has negative cycle, there's no
Else, create a virtual vertex s, add edge with
     weight 0 from s to every x_i,
the solution is the shortest path from s to n
     vertices.
typedef long long 11;
struct edge{
   int u. v. c:
   check if negative cycle
bool bellman_ford(int n, vector <edge> edges){
   int m = (int)edges.size();
   vector <1l> dist(n+1);
   for(int i = 1; i < n; i++)</pre>
       for(int j = 0; j < m; j++){
```

int u = edges[j].u;

```
int v = edges[j].v;
           int c = edges[j].c;
           if(dist[v] > dist[u] + c)
              dist[v] = dist[u] + c:
       }
   for(int j = 0; j < m; j++){
       int u = edges[j].u;
       int v = edges[j].v;
       int c = edges[j].c;
       if(dist[v] > dist[u] + c)
           return true:
   }
   return false:
void solve(int n. int m){
   vector <edge> edges;
   while(m--){
       char t;
       cin >> t;
       if(t == 'P'){
          int u, v, c;
           cin >> u >> v >> c;
           edges.push_back({u, v, c});
           edges.push_back({v, u, -c});
       }else{
           int u, v; cin >> u >> v;
           edges.push_back({v, u, -1});
       }
   }
   if(bellman ford(n, edges))
       cout << "Unreliable" << '\n':</pre>
   else cout << "Reliable" << '\n':</pre>
```

7 Matrix

7.1 matrix

```
const int MN = 111;
const int mod = 10000;

struct matrix {
  int r, c;
  int m[MN] [MN];

matrix (int _r, int _c) : r (_r), c (_c) {
   memset(m, 0, sizeof m);
  }

void print() {
```

```
for (int i = 0; i < r; ++i) {</pre>
     for (int j = 0; j < c; ++j)
       cout << m[i][j] << " ";
     cout << endl:
 }
 int x[MN][MN];
  matrix & operator *= (const matrix &o) {
   memset(x, 0, sizeof x);
   for (int i = 0; i < r; ++i)</pre>
     for (int k = 0: k < c: ++k)
       if (m[i][k] != 0)
         for (int j = 0; j < c; ++j) {</pre>
           x[i][j] = (x[i][j] + ((m[i][k] *
                o.m[k][i]) % mod) ) % mod:
   memcpy(m, x, sizeof(m));
   return *this;
};
void matrix_pow(matrix b, long long e, matrix &res)
  memset(res.m, 0, sizeof res.m);
 for (int i = 0; i < b.r; ++i)</pre>
   res.m[i][i] = 1;
  if (e == 0) return:
  while (true) {
   if (e & 1) res *= b:
   if ((e >>= 1) == 0) break:
   b *= b:
 }
}
```

8 Misc

8.1 dates

```
//
// Time - Leap years
//

// A[i] has the accumulated number of days from months previous to i

const int A[13] = { 0, 0, 31, 59, 90, 120, 151, 181, 212, 243, 273, 304, 334 };

// same as A, but for a leap year

const int B[13] = { 0, 0, 31, 60, 91, 121, 152, 182, 213, 244, 274, 305, 335 };
```

```
// returns number of leap years up to, and
    including, v
int leap_years(int y) { return y / 4 - y / 100 + y
    / 400: }
bool is_leap(int y) { return y % 400 == 0 || (y % 4
    == 0 && v % 100 != 0); }
// number of days in blocks of years
const int p400 = 400*365 + leap_years(400);
const int p100 = 100*365 + leap_years(100);
const int p4 = 4*365 + 1:
const int p1 = 365;
int date to days(int d. int m. int v)
 return (y - 1) * 365 + leap_years(y - 1) +
      (is_{pap}(y) ? B[m] : A[m]) + d;
void days_to_date(int days, int &d, int &m, int &v)
 bool top100; // are we in the top 100 years of a
      400 block?
 bool top4; // are we in the top 4 years of a
      100 block?
 bool top1; // are we in the top year of a 4
      block?
 y = 1:
 top100 = top4 = top1 = false;
 y += ((days-1) / p400) * 400;
 d = (days-1) \% p400 + 1;
 if (d > p100*3) top100 = true, d = 3*p100, v +=
 else v += ((d-1) / p100) * 100, d = (d-1) % p100
      + 1:
 if (d > p4*24) top4 = true, d = 24*p4, v += 24*4:
 else y += ((d-1) / p4) * 4, d = (d-1) % p4 + 1;
 if (d > p1*3) top1 = true, d -= p1*3, y += 3;
 else y += (d-1) / p1, d = (d-1) % p1 + 1;
 const int *ac = top1 && (!top4 || top100) ? B : A;
 for (m = 1; m < 12; ++m) if (d \le ac[m + 1])
      break;
 d = ac[m];
```

9 Number theory

9.1 convolution

```
typedef long long int LL;
typedef pair<LL, LL> PLL;
inline bool is_pow2(LL x) {
 return (x & (x-1)) == 0;
inline int ceil_log2(LL x) {
 int ans = 0:
 --x:
 while (x != 0) {
   x >>= 1:
   ans++:
 return ans;
/* Returns the convolution of the two given vectors
    in time proportional to n*log(n).
* The number of roots of unity to use nroots_unity
     must be set so that the product of the first
 * nroots_unity primes of the vector
     nth_roots_unity is greater than the maximum
     value of the
 * convolution. Never use sizes of vectors bigger
     than 2^24, if you need to change the values of
 * the nth roots of unity to appropriate primes for
     those sizes.
vector<LL> convolve(const vector<LL> &a. const
    vector<LL> &b. int nroots unity = 2) {
 int N = 1 << ceil_log2(a.size() + b.size());</pre>
 vector<LL> ans(N,0), fA(N), fB(N), fC(N);
 LL modulo = 1:
 for (int times = 0: times < nroots unity:</pre>
      times++) {
   fill(fA.begin(), fA.end(), 0):
   fill(fB.begin(), fB.end(), 0);
   for (int i = 0: i < a.size(): i++) fA[i] = a[i]:</pre>
   for (int i = 0: i < b.size(): i++) fB[i] = b[i]:</pre>
   LL prime = nth_roots_unity[times].first;
   LL inv_modulo = mod_inv(modulo % prime, prime);
   LL normalize = mod_inv(N, prime);
   ntfft(fA, 1, nth_roots_unity[times]);
   ntfft(fB, 1, nth_roots_unity[times]);
   for (int i = 0; i < N; i++) fC[i] = (fA[i] *</pre>
        fB[i]) % prime;
   ntfft(fC, -1, nth_roots_unity[times]);
   for (int i = 0; i < N; i++) {</pre>
     LL curr = (fC[i] * normalize) % prime;
     LL k = (curr - (ans[i] % prime) + prime) %
          prime:
     k = (k * inv modulo) % prime:
     ans[i] += modulo * k:
```

```
modulo *= prime;
}
return ans;
}
```

9.2 crt

```
/**

* Chinese remainder theorem.

* Find z such that z % x[i] = a[i] for all i.

* */
long long crt(vector<long long> &a, vector<long long> &x) {
  long long z = 0;
  long long n = 1;
  for (int i = 0; i < x.size(); ++i)
    n *= x[i];

for (int i = 0; i < a.size(); ++i) {
  long long tmp = (a[i] * (n / x[i])) % n;
  tmp = (tmp * mod_inv(n / x[i], x[i])) % n;
  z = (z + tmp) % n;
}

return (z + n) % n;
}
```

9.3 diophantine equations

```
long long gcd(long long a, long long b, long long
    &x, long long &y) {
 if (a == 0) {
   x = 0:
   y = 1;
   return b:
 long long x1, y1;
 long long d = gcd(b \% a, a, x1, y1);
 x = v1 - (b / a) * x1:
 y = x1;
 return d:
bool find_any_solution(long long a, long long b,
    long long c, long long &x0,
   long long &v0, long long &g) {
 g = gcd(abs(a), abs(b), x0, y0);
 if (c % g) {
   return false;
```

```
x0 *= c / g;
 v0 *= c / g;
 if (a < 0) x0 = -x0;
 if (b < 0) v0 = -v0;
 return true;
void shift solution(long long &x. long long &v.
    long long a, long long b,
   long long cnt) {
 x += cnt * b:
 y -= cnt * a;
long long find all solutions(long long a, long long
     b. long long c.
   long long minx, long long maxx, long long miny,
   long long maxy) {
 long long x, y, g;
 if (!find_any_solution(a, b, c, x, y, g)) return
 a /= g;
 b /= g;
 long long sign_a = a > 0 ? +1 : -1;
 long long sign_b = b > 0 ? +1 : -1;
 shift_solution(x, y, a, b, (minx - x) / b);
 if (x < minx) shift_solution(x, y, a, b, sign_b);</pre>
 if (x > maxx) return 0:
 long long lx1 = x:
 shift solution(x, v, a, b, (maxx - x) / b):
 if (x > maxx) shift_solution(x, y, a, b, -sign_b);
 long long rx1 = x;
 shift_solution(x, y, a, b, -(miny - y) / a);
 if (y < miny) shift_solution(x, y, a, b, -sign_a);</pre>
 if (v > maxv) return 0:
 long long 1x2 = x;
 shift_solution(x, y, a, b, -(maxy - y) / a);
 if (y > maxy) shift_solution(x, y, a, b, sign_a);
 long long rx2 = x;
 if (1x2 > rx2) swap(1x2, rx2);
 long long lx = max(lx1, lx2);
 long long rx = min(rx1, rx2);
 if (lx > rx) return 0;
 return (rx - lx) / abs(b) + 1:
```

9.4 discrete logarithm

```
// Computes x which a \hat{x} = b \mod n.
long long d_log(long long a, long long b, long long
 long long m = ceil(sqrt(n));
 long long aj = 1;
 map<long long, long long> M;
 for (int i = 0; i < m; ++i) {</pre>
   if (!M.count(aj))
     M[ai] = i;
   aj = (aj * a) % n;
 long long coef = mod_pow(a, n - 2, n);
 coef = mod_pow(coef, m, n);
 // coef = a^{-} (-m)
 long long gamma = b;
 for (int i = 0: i < m: ++i) {
   if (M.count(gamma)) {
     return i * m + M[gamma]:
   } else {
     gamma = (gamma * coef) % n:
 return -1:
```

9.5 ext euclidean

9.6 highest exponent factorial

```
int highest_exponent(int p, const int &n){
  int ans = 0;
  int t = p;
```

```
while(t <= n){
    ans += n/t;
    t*=p;
}
return ans;
}</pre>
```

9.7 miller rabin

```
const int rounds = 20;
// checks whether a is a witness that n is not
    prime, 1 < a < n
bool witness(long long a, long long n) {
 // check as in Miller Rabin Primality Test
      described
 long long u = n - 1;
 int t = 0:
 while (u % 2 == 0) {
   t.++:
   u >>= 1:
 long long next = mod_pow(a, u, n);
 if (next == 1) return false;
 long long last;
 for (int i = 0; i < t; ++i) {</pre>
   last = next:
   next = mod_mul(last, last, n);
   if (next == 1) {
     return last != n - 1;
 return next != 1;
// Checks if a number is prime with prob 1 - 1 / (2
// D(miller_rabin(999999999999997LL) == 1);
// D(miller rabin(999999999971LL) == 1):
// D(miller rabin(7907) == 1):
bool miller_rabin(long long n, int it = rounds) {
 if (n <= 1) return false:
 if (n == 2) return true;
 if (n % 2 == 0) return false;
 for (int i = 0; i < it; ++i) {</pre>
   long long a = rand() \% (n - 1) + 1;
   if (witness(a, n)) {
     return false;
 return true;
```

.

9.8 mod integer

```
template<class T, T mod>
struct mint_t {
   T val;
   mint_t() : val(0) {}
   mint_t(T v) : val(v % mod) {}

mint_t operator + (const mint_t& o) const {
   return (val + o.val) % mod;
}

mint_t operator - (const mint_t& o) const {
   return (val - o.val) % mod;
}

mint_t operator * (const mint_t& o) const {
   return (val * o.val) % mod;
}

stypedef mint_t<long long, 998244353> mint;
```

9.9 mod inv

```
long long mod_inv(long long n, long long m) {
  long long x, y, gcd;
  ext_euclid(n, m, x, y, gcd);
  if (gcd != 1)
   return 0;
  return (x + m) % m;
}
```

9.10 mod mul

```
// Computes (a * b) % mod
long long mod_mul(long long a, long long b, long
    long mod) {
    long long x = 0, y = a % mod;
    while (b > 0) {
        if (b & 1)
            x = (x + y) % mod;
        y = (y * 2) % mod;
        b /= 2;
    }
    return x % mod;
}
```

$9.11 \mod pow$

```
// Computes ( a ^ exp ) % mod.
long long mod_pow(long long a, long long exp, long
    long mod) {
 long long ans = 1;
 while (exp > 0) {
   if (exp & 1)
     ans = mod_mul(ans, a, mod);
   a = mod_mul(a, a, mod);
   exp >>= 1;
 return ans;
```

number theoretic transform

```
typedef long long int LL;
typedef pair<LL, LL> PLL:
/* The following vector of pairs contains pairs
     (prime, generator)
 * where the prime has an Nth root of unity for N
     being a power of two.
 * The generator is a number g s.t g^(p-1)=1 (mod p)
 * but is different from 1 for all smaller powers */
vector<PLL> nth_roots_unity {
 {1224736769,330732430},{1711276033,927759239},{167772161,1674893220 pollard_rho(long long n) {
  \{469762049, 343261969\}, \{754974721, 643797295\}, \{1107296297, 883865065\}  x, y, i = 1, k = 2, d;
PLL ext_euclid(LL a, LL b) {
 if (b == 0)
   return make_pair(1,0);
 pair<LL,LL> rc = ext_euclid(b, a % b);
 return make pair(rc.second, rc.first - (a / b) *
       rc.second):
}
//returns -1 if there is no unique modular inverse
LL mod inv(LL x. LL modulo) {
 PLL p = ext_euclid(x, modulo);
 if ( (p.first * x + p.second * modulo) != 1 )
   return -1:
 return (p.first+modulo) % modulo;
//Number theory fft. The size of a must be a power
    of 2
void ntfft(vector<LL> &a, int dir, const PLL
    &root_unity) {
 int n = a.size();
```

```
LL prime = root_unity.first;
 LL basew = mod_pow(root_unity.second, (prime-1) /
      n, prime);
 if (dir < 0) basew = mod_inv(basew, prime);</pre>
 for (int m = n; m >= 2; m >>= 1) {
   int mh = m >> 1;
   LL w = 1:
   for (int i = 0; i < mh; i++) {</pre>
     for (int i = i: i < n: i += m) {
       int k = j + mh;
       LL x = (a[j] - a[k] + prime) \% prime;
       a[j] = (a[j] + a[k]) \% prime;
       a[k] = (w * x) \% prime:
     w = (w * basew) % prime;
   basew = (basew * basew) % prime;
 for (int j = 1; j < n - 1; j++) {
   for (int k = n >> 1; k > (i ^= k); k >>= 1);
   if (j < i) swap(a[i], a[j]);</pre>
 }
}
```

9.13 pollard rho factorize

```
x = y = rand() \% n;
 while (1) {
   x = mod_mul(x, x, n);
   if (x \ge n) x = n:
   if (x == y) return 1;
   d = \_gcd(abs(x - y), n);
   if (d != 1) return d:
   if (i == k) {
    y = x;
     k *= 2:
 return 1;
// Returns a list with the prime divisors of n
vector<long long> factorize(long long n) {
 vector<long long> ans;
 if (n == 1)
   return ans;
```

```
if (miller_rabin(n)) {
 ans.push_back(n);
} else {
 long long d = 1;
 while (d == 1)
   d = pollard_rho(n);
 vector<long long> dd = factorize(d);
 ans = factorize(n / d);
 for (int i = 0: i < dd.size(): ++i)</pre>
   ans.push back(dd[i]):
return ans:
```

9.14 primes

```
namespace primes {
 const int MP = 100001:
 bool sieve[MP]:
 long long primes[MP];
 int num_p;
 void fill sieve() {
   num_p = 0;
   sieve[0] = sieve[1] = true;
   for (long long i = 2; i < MP; ++i) {</pre>
     if (!sieve[i]) {
       primes[num_p++] = i;
       for (long long j = i * i; j < MP; j += i)
         sieve[i] = true;
   }
 // Finds prime numbers between a and b, using
      basic primes up to sqrt(b)
 // a must be greater than 1.
 vector<long long> seg_sieve(long long a, long
      long b) {
   long long ant = a;
   a = max(a, 3LL):
   vector<bool> pmap(b - a + 1);
   long long sqrt_b = sqrt(b);
   for (int i = 0: i < num p: ++i) {
     long long p = primes[i];
     if (p > sqrt_b) break;
     long long j = (a + p - 1) / p;
     for (long long v = (j == 1) ? p + p : j * p;
          v \le b; v += p) {
       pmap[v - a] = true;
   vector<long long> ans;
```

```
if (ant == 2) ans.push_back(2);
   int start = a % 2 ? 0 : 1:
   for (int i = start, I = b - a + 1; i < I; i +=</pre>
     if (pmap[i] == false)
       ans.push_back(a + i);
   return ans;
 vector<pair<int, int>> factor(int n) {
   vector<pair<int, int>> ans;
   if (n == 0) return ans:
   for (int i = 0: primes[i] * primes[i] <= n:</pre>
        ++i) {
     if ((n % primes[i]) == 0) {
       int expo = 0:
       while ((n % primes[i]) == 0) {
         expo++:
         n /= primes[i];
       ans.emplace_back(primes[i], expo);
   if (n > 1) {
     ans.emplace_back(n, 1);
   return ans;
}
```

9.15 totient sieve

```
for (int i = 1; i < MN; i++)
  phi[i] = i;

for (int i = 1; i < MN; i++)
  if (!sieve[i]) // is prime
  for (int j = i; j < MN; j += i)
    phi[j] -= phi[j] / i;</pre>
```

9.16 totient

```
long long totient(long long n) {
  if (n == 1) return 0;
  long long ans = n;
  for (int i = 0; primes[i] * primes[i] <= n; ++i) {
    if ((n % primes[i]) == 0) {
      while ((n % primes[i]) == 0) n /= primes[i];
      ans -= ans / primes[i];
  }
}</pre>
```

```
}
if (n > 1) {
  ans -= ans / n;
}
return ans;
}
```

10 Strings

10.1 Incremental Aho Corasick

```
class IncrementalAhoCorasic {
 static const int Alphabets = 26:
 static const int AlphabetBase = 'a':
 struct Node {
   Node *fail:
   Node *next[Alphabets];
   int sum:
   Node(): fail(NULL), next{}, sum(0) { }
 struct String {
   string str;
   int sign;
 };
public:
 //totalLen = sum of (len + 1)
 void init(int totalLen) {
   nodes.resize(totalLen);
   nNodes = 0:
   strings.clear();
   roots.clear():
   sizes.clear():
   que.resize(totalLen);
 void insert(const string &str, int sign) {
   strings.push back(String{ str. sign }):
   roots.push back(nodes.data() + nNodes):
   sizes.push back(1):
   nNodes += (int)str.size() + 1:
   auto check = [&]() { return sizes.size() > 1 &&
        sizes.end()[-1] == sizes.end()[-2]: }:
   if(!check())
     makePMA(strings.end() - 1, strings.end(),
         roots.back(), que);
   while(check()) {
     int m = sizes.back();
     roots.pop_back();
     sizes.pop_back();
```

```
sizes.back() += m;
     if(!check())
       makePMA(strings.end() - m * 2,
            strings.end(), roots.back(), que);
 }
  int match(const string &str) const {
   int res = 0:
   for(const Node *t : roots)
     res += matchPMA(t, str):
   return res:
 }
private:
  static void
      makePMA(vector<String>::const iterator
      begin, vector<String>::const_iterator end,
      Node *nodes, vector<Node*> &que) {
   int nNodes = 0:
   Node *root = new(&nodes[nNodes ++]) Node();
   for(auto it = begin; it != end; ++ it) {
     Node *t = root:
     for(char c : it->str) {
       Node *&n = t->next[c - AlphabetBase];
       if(n == nullptr)
         n = new(&nodes[nNodes ++]) Node();
       t = n;
     t->sum += it->sign;
   int at = 0:
   for(Node *&n : root->next) {
     if(n != nullptr) {
       n->fail = root:
       que[qt ++] = n;
     } else {
       n = root:
   for(int qh = 0; qh != qt; ++ qh) {
     Node *t = que[qh];
     int a = 0:
     for(Node *n : t->next) {
       if(n != nullptr) {
         que[qt ++] = n;
         Node *r = t->fail:
         while(r->next[a] == nullptr)
          r = r->fail;
         n\rightarrow fail = r\rightarrow next[a];
         n->sum += r->next[a]->sum;
       ++ a:
```

```
}
 static int matchPMA(const Node *t, const string
   int res = 0;
   for(char c : str) {
     int a = c - AlphabetBase;
     while(t->next[a] == nullptr)
       t = t->fail:
     t = t- next[a]:
     res += t->sum:
   return res;
 vector<Node> nodes:
 int nNodes:
 vector<String> strings;
 vector<Node*> roots;
 vector<int> sizes;
 vector<Node*> que;
};
int main() {
 int m;
 while(~scanf("%d", &m)) {
   IncrementalAhoCorasic iac;
   iac.init(600000):
   rep(i, m) {
     int ty;
     char s[300001];
     scanf("%d%s", &ty, s);
     if(tv == 1) {
       iac.insert(s, +1);
     } else if(ty == 2) {
       iac.insert(s, -1):
     } else if(ty == 3) {
       int ans = iac.match(s):
       printf("%d\n", ans);
       fflush(stdout);
     } else {
       abort();
 return 0;
```

10.2 minimal string rotation

```
// Lexicographically minimal string rotation
```

```
int lmsr() {
 string s;
 cin >> s;
 int n = s.size();
 s += s;
 vector<int> f(s.size(), -1);
 int k = 0:
 for (int j = 1; j < 2 * n; ++j) {
   int i = f[j - k - 1];
   while (i != -1 && s[j] != s[k + i + 1]) {
     if (s[j] < s[k + i + 1])
      k = i - i - 1:
     i = f[i]:
   if (i == -1 && s[j] != s[k + i + 1]) {
     if (s[i] < s[k + i + 1]) {
      k = j;
    f[i - k] = -1;
   } else {
     f[i - k] = i + 1;
 }
 return k;
```

10.3 suffix array

```
const int MAXN = 200005;
const int MAX_DIGIT = 256;
void countingSort(vector<int>& SA, vector<int>& RA,
    int k = 0) {
   int n = SA.size():
   vector<int> cnt(max(MAX DIGIT, n), 0);
   for (int i = 0; i < n; i++)</pre>
       if (i + k < n)
           cnt[RA[i + k]]++:
       else
           cnt[0]++:
   for (int i = 1: i < cnt.size(): i++)</pre>
       cnt[i] += cnt[i - 1]:
   vector<int> tempSA(n);
   for (int i = n - 1; i \ge 0; i--)
       if (SA[i] + k < n)
           tempSA[--cnt[RA[SA[i] + k]]] = SA[i];
           tempSA[--cnt[0]] = SA[i];
   SA = tempSA;
vector <int> constructSA(string s) {
```

```
int n = s.length();
   vector <int> SA(n);
   vector <int> RA(n);
   vector <int> tempRA(n);
   for (int i = 0; i < n; i++) {</pre>
       RA[i] = s[i];
       SA[i] = i;
   }
   for (int step = 1; step < n; step <<= 1) {</pre>
       countingSort(SA, RA, step);
       countingSort(SA, RA, 0);
       int c = 0:
       tempRA[SA[O]] = c:
       for (int i = 1; i < n; i++) {</pre>
           if (RA[SA[i]] == RA[SA[i - 1]] &&
                RA[SA[i] + step] == RA[SA[i - 1] +
                stepl)
                  tempRA[SA[i]] = tempRA[SA[i - 1]];
               tempRA[SA[i]] = tempRA[SA[i - 1]] +
       }
       RA = tempRA;
       if (RA[SA[n-1]] == n-1) break;
   return SA;
}
vector<int> computeLCP(const string& s, const
     vector<int>& SA) {
   int n = SA.size():
   vector<int> LCP(n), PLCP(n), c(n, 0);
   for (int i = 0; i < n; i++)</pre>
       c[SA[i]] = i:
   int k = 0:
   for (int j, i = 0; i < n-1; i++) {</pre>
       if(c[i] - 1 < 0)
           continue;
       i = SA[c[i] - 1]:
       k = \max(k - 1, 0);
       while (i+k < n \&\& j+k < n \&\& s[i+k] == s[j
           k++;
       PLCP[i] = k;
   for (int i = 0; i < n; i++)
       LCP[i] = PLCP[SA[i]];
   return LCP;
```

10.4 suffix automaton

```
* Suffix automaton:
 * This implementation was extended to maintain
      (online) the
 * number of different substrings. This is
     equivalent to compute
 * the number of paths from the initial state to
     all the other
 * states.
 * The overall complexity is O(n)
 * can be tested here:
     https://www.urionlinejudge.com.br/judge/en/problems/view/1580t g = sa[p].next[c]:
struct state {
 int len. link:
 long long num_paths;
 map<int, int> next;
};
const int MN = 200011;
state sa[MN << 1];
int sz, last;
long long tot_paths;
void sa_init() {
 sz = 1;
 last = 0:
 sa[0].len = 0:
 sa[0].link = -1:
 sa[0].next.clear();
 sa[0].num_paths = 1;
 tot paths = 0:
void sa extend(int c) {
 int cur = sz++;
 sa[cur].len = sa[last].len + 1:
 sa[cur].next.clear():
```

```
sa[cur].num_paths = 0;
for (p = last; p != -1 && !sa[p].next.count(c); p
     = sa[p].link) {
  sa[p].next[c] = cur;
 sa[cur].num_paths += sa[p].num_paths;
 tot_paths += sa[p].num_paths;
if (p == -1) {
  sa[cur].link = 0;
} else {
  if (sa[p].len + 1 == sa[q].len) {
   sa[cur].link = q;
 } else {
   int clone = sz++;
   sa[clone].len = sa[p].len + 1;
   sa[clone].next = sa[q].next;
   sa[clone].num_paths = 0;
   sa[clone].link = sa[q].link;
   for (; p!= -1 && sa[p].next[c] == q; p =
        sa[p].link) {
     sa[p].next[c] = clone;
     sa[q].num_paths -= sa[p].num_paths;
     sa[clone].num_paths += sa[p].num_paths;
    sa[q].link = sa[cur].link = clone;
last = cur;
```

10.5 z algorithm

```
using namespace std:
#include<bits/stdc++.h>
```

```
vector<int> compute_z(const string &s){
 int n = s.size():
 vector<int> z(n,0);
 int 1.r:
 r = 1 = 0;
 for(int i = 1; i < n; ++i){
   if(i > r) {
     1 = r = i:
     while(r < n and s[r - 1] == s[r])r++:
     z[i] = r - 1:r--:
   }else{
     int k = i-1:
     if(z[k] < r - i +1) z[i] = z[k]:
     else {
      1 = i:
       while(r < n and s[r - 1] == s[r])r++:
       z[i] = r - 1:r--:
 return z;
int main(){
 //string line;cin>>line;
 string line = "alfalfa";
 vector<int> z = compute_z(line);
 for(int i = 0: i < z.size(): ++i ){</pre>
   if(i)cout<<" ":
   cout<<z[i]:
 cout << endl:
 // must print "0 0 0 4 0 0 1"
 return 0;
```