Team notebook

${\bf HCMUS\text{-}IdentityImbalance}$

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Contents			4 Geometry 1	۱6
			4.1 center 2 points + radious	16
1	Algorithms	1	4.2 closest pair problem	16
	1.1 Mo's algorithm on trees	1	4.3 squares	17
	1.2 Mo's algorithm	2	4.4 triangles	18
	1.3 sliding window	2		
				18
2	DP Optimizations	3	5.1 SCC kosaraju	
	2.1 convex hull trick	3	5.2 board	19
	2.2 divide and conquer	4	5.3 bridges	
	2.3 dp on trees	4	5.4 dijkstra	
			5.5 directed mst	21
3	Data structures	5	5.6 eulerian path	
	3.1 STL Treap	5	5.7 karp min mean cycle	22
	3.2 STL order statistics tree II	6	5.8 konig's theorem	
	3.3 STL order statistics tree	6	5.9 minimum path cover in DAG	
	3.4 binary index tree	7	5.10 planar graph (euler)	23
	3.5 dsu	7	5.11 query with lca	23
	3.6 hash table	8	5.12 tarjan scc	25
	3.7 heavy light decomposition	8	5.13 two sat (with kosaraju)	25
	3.8 persistent array	9	6 Math 2	27
	3.9 persistent seg tree	9	6.1 Lucas theorem	
	3.10 persistent trie		6.2 counting	
	3.11 segment tree		6.3 cumulative sum of divisors	
	3.12 sparse table		6.4 fft	
	3.13 splay tree			
	3.14 trie		6.5 fibonacci properties	
			F =	
	3.15 wavelet tree	19	6.7 sigma function	49

7	Mat 7.1	rix matrix	30
8	Mis	c	30
	8.1	Template Java	30
	8.2	dates	31
	8.3	fraction	32
	8.4	io	32
9	Nur	nber theory	32
J	9.1	convolution	32
	9.2	crt	33
	9.3	diophantine equations	33
	9.4	discrete logarithm	34
	$9.4 \\ 9.5$	ext euclidean	$\frac{34}{34}$
	9.6	highest exponent factorial	$\frac{34}{35}$
	$9.0 \\ 9.7$	miller rabin	- 55 - 35
	9.7	mod integer	აა 35
	9.9	mod inv	36
		mod mul	36
		mod pow	36
		number theoretic transform	36
		pollard rho factorize	37
		primes	37
		totient sieve	38
	9.16	totient	38
10	Stri	$_{ m ngs}$	38
	10.1	Incremental Aho Corasick	38
		minimal string rotation	40
		suffix array	40
		suffix automaton	41
	-	z algorithm	42
1	A	lgorithms	
1.	1 I	Mo's algorithm on trees	
<u></u>			
/*:	* oblem:		
PT.		tps://codeforces.com/gym/101161 problem E	

```
void flat(vector<vector<edge>> &g, vector<int> &a,
   vector<int> &le, vector<int> &ri, vector<int> &cost,
   int node, int pi, int &ts, int w) {
 cost[node] = w;
 le[node] = ts;
 a[ts] = node;
 ts++;
 for (auto e : g[node]) {
   if (e.to == pi) continue;
   flat(g, a, le, ri, cost, e.to, node, ts, e.w);
 ri[node] = ts;
 a[ts] = node;
 ts++;
* Case when the cost is in the edges.
void compute_queries(vector<vector<edge>> &g) {
 // g is undirected
 int n = g.size();
 lca_tree.init(g, 0);
 vector\langle int \rangle a(2 * n), le(n), ri(n), cost(n);
 // a: nodes in the flatten array
 // le: left id of the given node
 // ri: right id of the given node
 // cost: cost of the edge from the node to the parent
 int ts = 0; // timestamp
 flat(g, a, le, ri, cost, 0, -1, ts, 0);
 int q; cin >> q;
 vector<query> queries(q);
 for (int i = 0; i < q; i++) {</pre>
   int u, v;
   cin >> u >> v;
   u--; v--;
   int lca = lca_tree.query(u, v);
   if (le[u] > le[v])
     swap(u, v);
```

```
queries[i].id = i;
queries[i].lca = lca;
queries[i].u = u;
queries[i].v = v;
if (lca == u) {
   queries[i].a = le[u] + 1;
   queries[i].b = le[v];
} else {
   queries[i].a = ri[u];
   queries[i].b = le[v];
}
}
solve_mo(queries, a, le, cost); // this is the usal algorithm
```

1.2 Mo's algorithm

```
const int MN = 5 * 100000 + 1;
const int SN = 708;
struct Query {
 int a, b, id;
 Query() {}
 Query(int x, int y, int i) : a(x), b(y), id(i) {}
 bool operator<(const Query &o) const {</pre>
   if (a / SN != o.a / SN) return a < o.a;</pre>
   return a / SN & 1 ? b < o.b : b > o.b;
 }
};
struct DS {
 DS() : {}
 void Insert(int x) {}
 void Erase(int x) {}
 long long Query() {}
};
Query s[MN];
int ans[MN];
```

```
DS active;
int main() {
  int n;
  cin >> n;
  vector<int> a(n);
  for (auto &i : a) cin >> i;
  int q;
  cin >> q;
  for (int i = 0; i < q; ++i) {</pre>
   int b, e;
    cin >> b >> e;
    b--;
    e--;
    s[i] = Query(b, e, i);
  sort(s, s + q);
  int i = 0;
  int j = -1;
  for (int k = 0; k < (int)q; ++k) {
   int L = s[k].a;
    int R = s[k].b;
    while (j < R) active.Insert(a[++j]);</pre>
    while (j > R) active.Erase(a[j--]);
    while (i < L) active.Erase(a[i++]);</pre>
    while (i > L) active.Insert(a[--i]);
    ans[s[k].id] = active.Query();
  for (int i = 0; i < q; ++i) {</pre>
    cout << ans[i] << endl;</pre>
 return 0;
};
```

1.3 sliding window

```
/*
* Given an array ARR and an integer K, the problem boils down to
     computing for each index i: min(ARR[i], ARR[i-1], ..., ARR[i-K+1]).
* if mx == true, returns the maximun.
* http://people.cs.uct.ac.za/~ksmith/articles/sliding_window_minimum.html
* */
vector<int> sliding_window_minmax(vector<int> & ARR, int K, bool mx) {
 deque< pair<int, int> > window;
 vector<int> ans:
 for (int i = 0; i < ARR.size(); i++) {</pre>
   if (mx) {
     while (!window.empty() && window.back().first <= ARR[i])</pre>
       window.pop_back();
   } else {
     while (!window.empty() && window.back().first >= ARR[i])
       window.pop_back();
   window.push_back(make_pair(ARR[i], i));
   while(window.front().second <= i - K)</pre>
     window.pop_front();
   ans.push_back(window.front().first);
 }
 return ans:
```

2 DP Optimizations

2.1 convex hull trick

```
/**
 * Problems:
 * http://codeforces.com/problemset/problem/319/C
 * http://codeforces.com/contest/311/problem/B
 * https://csacademy.com/contest/archive/task/squared-ends
 * http://codeforces.com/contest/932/problem/F
 * */
struct line {
```

```
long long m, b;
 line (long long a, long long c) : m(a), b(c) {}
  long long eval(long long x) {
   return m * x + b;
};
long double inter(line a, line b) {
  long double den = a.m - b.m;
 long double num = b.b - a.b;
 return num / den:
 * min m_i * x_j + b_i, for all i.
      x_j \le x_{j+1}
      m_i >= m_{j+1}
struct ordered cht {
  vector<line> ch;
  int idx; // id of last "best" in query
  ordered_cht() {
   idx = 0;
  void insert_line(long long m, long long b) {
   line cur(m. b):
   // new line's slope is less than all the previous
   while (ch.size() > 1 &&
      (inter(cur, ch[ch.size() - 2]) >= inter(cur, ch[ch.size() - 1]))) {
       // f(x) is better in interval [inter(ch.back(), cur), inf)
       ch.pop_back();
   ch.push_back(cur);
  long long eval(long long x) { // minimum
   // current x is greater than all the previous x,
   // if that is not the case we can make binary search.
   idx = min<int>(idx, ch.size() - 1);
   while (idx + 1 < (int)ch.size() && ch[idx + 1].eval(x) <=
        ch[idx].eval(x))
     idx++;
```

```
return ch[idx].eval(x);
 }
};
 * Dynammic convex hull trick
typedef long long int64;
typedef long double float128;
const int64 is_query = -(1LL<<62), inf = 1e18;</pre>
struct Line {
  int64 m. b:
 mutable function<const Line*()> succ;
 bool operator<(const Line& rhs) const {</pre>
    if (rhs.b != is_query) return m < rhs.m;</pre>
    const Line* s = succ();
    if (!s) return 0;
    int64 x = rhs.m;
    return b - s -> b < (s -> m - m) * x;
 }
};
struct HullDynamic : public multiset<Line> { // will maintain upper hull
    for maximum
  bool bad(iterator y) {
    auto z = next(y);
    if (y == begin()) {
     if (z == end()) return 0;
     return y->m == z->m && y->b <= z->b;
    auto x = prev(y);
    if (z == end()) return y->m == x->m && y->b <= x->b;
    return (float128)(x->b - y->b)*(z->m - y->m) >= (float128)(y->b -
        z->b)*(y->m - x->m);
  void insert_line(int64 m, int64 b) {
    auto y = insert({ m, b });
    y->succ = [=] { return next(y) == end() ? 0 : &*next(y); };
    if (bad(y)) { erase(y); return; }
    while (next(y) != end() && bad(next(y))) erase(next(y));
    while (y != begin() && bad(prev(y))) erase(prev(y));
```

```
int64 eval(int64 x) {
    auto 1 = *lower_bound((Line) { x, is_query });
    return 1.m * x + 1.b;
};
```

2.2 divide and conquer

```
/**
* recurrence:
     dp[k][i] = min dp[k-1][j] + c[i][j-1], for all j > i;
  "comp" computes dp[k][i] for all i in O(n log n) (k is fixed)
* Problems:
* https://icpc.kattis.com/problems/branch
* http://codeforces.com/contest/321/problem/E
* */
void comp(int 1, int r, int le, int re) {
 if (1 > r) return;
 int mid = (1 + r) >> 1:
 int best = max(mid + 1, le);
 dp[cur][mid] = dp[cur ^ 1][best] + cost(mid, best - 1);
 for (int i = best; i <= re; i++) {</pre>
   if (dp[cur][mid] > dp[cur ^ 1][i] + cost(mid, i - 1)) {
     dp[cur][mid] = dp[cur ^ 1][i] + cost(mid, i - 1);
   }
 }
 comp(l, mid - 1, le, best);
 comp(mid + 1, r, best, re);
```

2.3 dp on trees

```
* This trick is very useful when doing DP on trees, basically, you can
 * the answer for each node as if it was the root of the tree. Partial
 * are also stored in order to query subtrees (taking the root and
     exclude some
 * child).
 * problems:
 * - http://codeforces.com/gym/101161, problem I : Sky tax
 * - http://codeforces.com/contest/791/problem/D
 * */
struct edge {
 int to, p_id;
 edge (int a, int b) : to(a), p_id(b) {}
};
struct state {
 bool seen;
 long long missing;
 long long total;
 vector<long long> partial;
  state() { clear(); }
  void clear() {
   seen = false;
   missing = 0;
   total = 0;
   partial.clear();
 }
};
void add_edge(int u, int v) {
 int id_u_v = g[u].size();
 int id_v_u = g[v].size();
 g[u].emplace_back(v, id_v_u); // id of the parent in the child's list
      (g[v][id] \rightarrow u)
 g[v].emplace_back(u, id_u_v); // id of the parent in the child's list
      (g[u][id] -> v)
}
```

```
int go(int node, int id_parent) {
  state &s = dp[node];
  if (!s.seen) {
   int ans = 1;
   s.partial.assign(g[node].size(), 0); // create the list of partial
        results.
   for (int i = 0; i < int(g[node].size()); i++) {</pre>
     int to = g[node][i].to;
     int pid = g[node][i].p_id;
     if (i != id_parent) {
       int tmp = go(to, pid);
       ans += tmp;
       s.partial[i] = tmp;
   }
   s.missing = id_parent;
   s.total = ans;
   s.seen = true;
   return ans:
 } else {
   if (s.missing == id_parent) { // the same id_parent than before, so
        we can not complete the results yet
     return s.total;
   if (s.missing !=-1) { // only one missing and is different of
        'id_parent'
     int tmp = go(g[node][s.missing].to, g[node][s.missing].p_id);
     s.partial[s.missing] = tmp;
     s.total += tmp;
     s.missing = -1;
   int extra = (id_parent == -1) ? 0 : s.partial[id_parent];
   return s.total - extra;
 }
}
```

3 Data structures

3.1 STL Treap

```
#include <ext/rope> //header with rope
using namespace std;
using namespace __gnu_cxx; //namespace with rope and some additional stuff
int main()
    ios_base::sync_with_stdio(false);
    rope <int> v; //use as usual STL container
    int n, m;
    cin >> n >> m;
    for(int i = 1; i <= n; ++i)</pre>
       v.push_back(i); //initialization
    int 1, r;
    for(int i = 0; i < m; ++i)</pre>
       cin >> 1 >> r;
       --1, --r;
       rope \langle int \rangle cur = v.substr(1, r - 1 + 1);
       v.erase(1, r - 1 + 1);
       v.insert(v.mutable_begin(), cur);
    for(rope <int>::iterator it = v.mutable_begin(); it !=
        v.mutable_end(); ++it)
       cout << *it << " ";
    return 0;
}
```

3.2 STL order statistics tree II

```
#include <bits/stdc++.h>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>

using namespace std;
using namespace __gnu_pbds;

typedef tree<int,null_type,less<int>,rb_tree_tag,
tree_order_statistics_node_update> order_set;
```

```
order_set X;
int get(int y) {
  int l=0,r=1e9+1;
  while(l<r) {</pre>
    int m=l+((r-l)>>1);
    if(m-X.order_of_key(m+1)<y)</pre>
    else
      r=m;
  return 1;
main(){
  ios::sync_with_stdio(0);
  cin.tie(0);
  int n,m;
  cin>>n>>m;
  for(int i=0;i<m;i++) {</pre>
    char a;
    int b;
    cin>>a>>b;
    if(a=='L')
      cout<<get(b)<<endl;</pre>
      X.insert(get(b));
  }
}
/***
Input
20 7
L 5
D 5
L 4
L 5
D 5
L 4
L 5
Output
5
4
```

```
6
4
7
***/
```

3.3 STL order statistics tree

```
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
#include <bits/stdc++.h>
using namespace __gnu_pbds;
using namespace std;
typedef
tree<
 pair<int,int>,
 null_type,
 less<pair<int,int>>,
 rb_tree_tag,
 tree_order_statistics_node_update>
ordered_set;
main()
    ios::sync_with_stdio(0);
    cin.tie(0);
    int n;
    int sz=0;
    cin>>n;
    vector<int> ans(n,0);
    ordered_set t;
    int x,y;
    for(int i=0;i<n;i++)</pre>
    {
       cin>>x>>y;
       ans[t.order_of_key(\{x,++sz\})]++;
       t.insert({x,sz});
    }
    for(int i=0;i<n;i++)</pre>
       cout<<ans[i]<<'\n';</pre>
```

```
/***
Input
5
1 1
5 1
7 1
3 3
5 5

Output
1
2
1
1
0
***/
```

3.4 binary index tree

```
struct binary_index_tree {
 int n;
 int t[2 * N];
 void add(int where, long long what){
   for (where++; where <= n; where += where & -where){</pre>
     t[where] += what;
   }
 }
 void add(int from, int to, long long what) {
   add(from, what);
   add(to + 1, -what);
 long long query(int where){
   long long sum = t[0];
   for (where++; where > 0; where -= where & -where){
     sum += t[where];
   return sum;
```

};

3.5 dsu

```
struct Dsu {
 vector<int> p;
 Dsu(int n) {
   p.resize(n);
   for (int i = 0; i < n; i++) {</pre>
     p[i] = i;
   }
 }
 int Find(int x) { return x == p[x] ? x : p[x] = Find(p[x]); }
 int Join(int x, int y) {
   int px = Find(x), py = Find(y);
   if (px == py) return 0;
   p[px] = py;
   return 1;
 }
};
```

3.6 hash table

```
/**
  * Micro hash table, can be used as a set.
  * Very efficient vs std::set
  * */

const int MN = 1001;
struct ht {
  int _s[(MN + 10) >> 5];
  int len;
  void set(int id) {
    len++;
    _s[id >> 5] |= (1LL << (id & 31));
}
bool is_set(int id) {
    return _s[id >> 5] & (1LL << (id & 31));
}</pre>
```

}

3.7 heavy light decomposition

```
// Heavy-Light Decomposition
struct TreeDecomposition {
 vector<int> g[MAXN], c[MAXN];
 int s[MAXN]; // subtree size
 int p[MAXN]; // parent id
 int r[MAXN]; // chain root id
 int t[MAXN]; // index used in segtree/bit/...
 int d[MAXN]; // depht
 int ts;
 void dfs(int v, int f) {
   p[v] = f;
   s[v] = 1;
   if (f != -1) d[v] = d[f] + 1;
   else d[v] = 0;
   for (int i = 0; i < g[v].size(); ++i) {</pre>
     int w = g[v][i];
     if (w != f) {
       dfs(w, v);
       s[v] += s[w];
   }
 void hld(int v, int f, int k) {
   t[v] = ts++;
   c[k].push_back(v);
   r[v] = k;
   int x = 0, y = -1;
   for (int i = 0; i < g[v].size(); ++i) {</pre>
     int w = g[v][i];
     if (w != f) {
       if (s[w] > x) {
        x = s[w]:
        y = w;
```

```
}
    }
    if (v != -1) {
     hld(y, v, k);
    for (int i = 0; i < g[v].size(); ++i) {</pre>
     int w = g[v][i];
     if (w != f && w != y) {
       hld(w, v, w);
     }
   }
 }
  void init(int n) {
   for (int i = 0; i < n; ++i) {</pre>
     g[i].clear();
   }
 }
  void add(int a, int b) {
    g[a].push_back(b);
    g[b].push_back(a);
  void build() {
    ts = 0:
   dfs(0, -1);
   hld(0, 0, 0);
 }
};
```

3.8 persistent array

```
struct node {
  node *1, *r;
  int val;

  node (int x) : 1(NULL), r(NULL), val(x) {}
  node () : 1(NULL), r(NULL), val(-1) {}
};

typedef node* pnode;
```

```
pnode update(pnode cur, int 1, int r, int at, int what) {
  pnode ans = new node();
  if (cur != NULL) {
   *ans = *cur:
  if (1 == r) {
   ans-> val = what;
   return ans;
 int m = (1 + r) >> 1;
 if (at <= m) ans-> 1 = update(ans-> 1, 1, m, at, what);
  else ans-> r = update(ans-> r, m + 1, r, at, what);
 return ans;
}
int get(pnode cur, int 1, int r, int at) {
 if (cur == NULL) return 0;
 if (1 == r) return cur-> val;
 int m = (1 + r) >> 1;
 if (at <= m) return get(cur-> 1, 1, m, at);
             return get(cur-> r, m + 1, r, at);
  else
```

3.9 persistent seg tree

```
/**
 * Problems:
 * http://codeforces.com/contest/813/problem/E

*
 * Important:
 * When using lazy propagation remembert to create new
 * versions for each push_down operation!!!
 * */

struct node {
 node *1, *r;
 long long acc;
 int flip;

node (int x) : l(NULL), r(NULL), acc(x), flip(0) {}
 node () : l(NULL), r(NULL), acc(0), flip(0) {}
```

```
};
typedef node* pnode;
pnode create(int 1, int r) {
 if (1 == r) return new node();
 pnode cur = new node();
 int m = (1 + r) >> 1;
  cur \rightarrow 1 = create(1, m);
  cur \rightarrow r = create(m + 1, r);
 return cur:
}
pnode copy_node(pnode cur) {
 pnode ans = new node();
  *ans = *cur;
 return ans;
}
void push_down(pnode cur, int 1, int r) {
  assert(cur);
  if (cur-> flip) {
   int len = r - 1 + 1;
    cur-> acc = len - cur-> acc;
    if (cur-> 1) {
     cur-> 1 = copy_node(cur-> 1);
     cur-> 1 -> flip ^= 1;
    if (cur-> r) {
     cur-> r = copy_node(cur-> r);
     cur-> r -> flip ^= 1;
    cur \rightarrow flip = 0;
}
int get_val(pnode cur) {
 assert(cur);
 assert((cur-> flip) == 0);
 if (cur) return cur-> acc;
 return 0;
}
pnode update(pnode cur, int 1, int r, int at, int what) {
 pnode ans = copy_node(cur);
```

```
if (1 == r) {
   assert(1 == at);
   ans-> acc = what;
   ans-> flip = 0;
   return ans;
  int m = (1 + r) >> 1;
 push_down(ans, 1, r);
 if (at <= m) ans-> 1 = update(ans-> 1, 1, m, at, what);
  else ans-> r = update(ans-> r, m + 1, r, at, what);
 push_down(ans-> 1, 1, m);
 push_down(ans-> r, m + 1, r);
  ans-> acc = get_val(ans-> 1) + get_val(ans-> r);
 return ans;
}
pnode flip(pnode cur, int 1, int r, int a, int b) {
 pnode ans = new node();
 if (cur != NULL) {
   *ans = *cur;
  if (1 > b | | r < a)
   return ans;
  if (1 >= a && r <= b) {</pre>
   ans-> flip ^= 1;
   push_down(ans, 1, r);
   return ans;
 int m = (1 + r) >> 1;
 ans-> 1 = flip(ans-> 1, 1, m, a, b);
  ans-> r = flip(ans-> r, m + 1, r, a, b);
 push_down(ans-> 1, 1, m);
 push_down(ans-> r, m + 1, r);
 ans-> acc = get_val(ans-> 1) + get_val(ans-> r);
 return ans;
}
long long get_all(pnode cur, int 1, int r) {
  assert(cur);
 push_down(cur, 1, r);
 return cur-> acc;
```

3.10 persistent trie

```
// both tries can be tested with the problem:
    http://codeforces.com/problemset/problem/916/D
// Persistent binary trie (BST for integers)
const int MD = 31;
struct node_bin {
 node_bin *child[2];
 int val;
 node_bin() : val(0) {
   child[0] = child[1] = NULL;
 }
};
typedef node_bin* pnode_bin;
pnode_bin copy_node(pnode_bin cur) {
 pnode_bin ans = new node_bin();
 if (cur) *ans = *cur;
 return ans;
}
pnode_bin modify(pnode_bin cur, int key, int inc, int id = MD) {
 pnode_bin ans = copy_node(cur);
 ans->val += inc;
 if (id >= 0) {
   int to = (key >> id) & 1;
   ans->child[to] = modify(ans->child[to], key, inc, id - 1);
```

```
}
 return ans;
int sum_smaller(pnode_bin cur, int key, int id = MD) {
  if (cur == NULL) return 0;
 if (id < 0) return 0; // strictly smaller</pre>
 // if (id == - 1) return cur->val; // smaller or equal
  int ans = 0;
  int to = (key >> id) & 1;
 if (to) {
   if (cur->child[0]) ans += cur->child[0]->val;
   ans += sum_smaller(cur->child[1], key, id - 1);
 } else {
   ans = sum_smaller(cur->child[0], key, id - 1);
 return ans;
// Persistent trie for strings.
const int MAX_CHILD = 26;
struct node {
 node *child[MAX_CHILD];
 int val:
 node() : val(-1) {
   for (int i = 0; i < MAX_CHILD; i++) {</pre>
     child[i] = NULL;
   }
 }
}:
typedef node* pnode;
pnode copy_node(pnode cur) {
 pnode ans = new node();
 if (cur) *ans = *cur;
 return ans;
}
pnode set_val(pnode cur, string &key, int val, int id = 0) {
  pnode ans = copy_node(cur);
 if (id >= int(key.size())) {
   ans->val = val;
```

```
} else {
   int t = key[id] - 'a';
   ans->child[t] = set_val(ans->child[t], key, val, id + 1);
}
return ans;
}

pnode get(pnode cur, string &key, int id = 0) {
   if (id >= int(key.size()) || !cur)
     return cur;
   int t = key[id] - 'a';
   return get(cur->child[t], key, id + 1);
}
```

3.11 segment tree

```
* Taken from: http://codeforces.com/blog/entry/18051
* */
const int MN = 1e5; // limit for array size
struct seg_tree {
 int n; // array size
 int t[2 * MN];
 seg_tree(int _n) : n(_n) {}
 void clear() {
   memset(t, 0, sizeof t);
 void build() { // build the tree
   for (int i = n - 1; i > 0; --i) t[i] = t[i << 1] + t[i << 1|1];
 }
 // Single modification, range query.
 void modify(int p, int value) { // set value at position p
   for (t[p += n] = value; p > 1; p >>= 1) t[p>>1] = t[p] + t[p^1];
 }
 int query(int 1, int r) { // sum on interval [1, r)
   int res = 0;
```

```
for (1 += n, r += n; 1 < r; 1 >>= 1, r >>= 1) {
     if (1&1) res += t[1++];
     if (r&1) res += t[--r]:
   return res;
};
// Range modification, single query.
void modify(int 1, int r, int value) {
 for (1 += n, r += n; 1 < r; 1 >>= 1, r >>= 1) {
   if (1&1) t[1++] += value;
   if (r&1) t[--r] += value;
}
int query(int p) {
 int res = 0;
 for (p += n; p > 0; p >>= 1) res += t[p];
 return res;
}
* If at some point after modifications we need to inspect all the
 * elements in the array, we can push all the modifications to the
 * leaves using the following code. After that we can just traverse
 * elements starting with index n. This way we reduce the complexity
 * from O(n \log(n)) to O(n) similarly to using build instead of n
     modifications.
 * */
void push() {
 for (int i = 1; i < n; ++i) {</pre>
   t[i<<1] += t[i];
   t[i<<1|1] += t[i];
   t[i] = 0;
}
// Non commutative combiner functions.
void modify(int p, const S& value) {
 for (t[p += n] = value; p >>= 1;) t[p] = combine(t[p<<1], t[p<<1|1]);
```

```
S query(int 1, int r) {
 S resl, resr;
 for (1 += n, r += n; 1 < r; 1 >>= 1, r >>= 1) {
   if (l\&1) resl = combine(resl, t[l++]);
    if (r&1) resr = combine(t[--r], resr);
 return combine(resl, resr);
/**
 * segment tree for intervals
 * */
const int MN = 100000 + 100;
struct seg_tree {
 int val[MN * 4 + 4];
 int pending [MN * 4 + 4];
  seg_tree() {
    memset(val, -1, sizeof val);
   memset(pending, -1, sizeof pending);
 }
  void propagate(int node, int b, int e) {
    if (pending[node] != -1) {
     val[node] = pending[node];
     if (b < e) {</pre>
       pending[node << 1] = pending[node];</pre>
       pending[node << 1 | 1] = pending[node];</pre>
     pending[node] = -1;
 }
  void set(int node, int b, int e, int from, int to, int v) {
    if (b > to || e < from) return;</pre>
    if (b >= from && e <= to) {</pre>
     pending[node] = v;
     propagate(node, b, e);
     return;
    }
```

```
int mid = (b + e) >> 1;
    set(node << 1, b, mid, from, to, v);
   set(node << 1 | 1, mid + 1, e, from, to, v);
  int query(int node, int b, int e, int pos) {
   propagate(node, b, e);
   if (b == e && b == pos) {
     return val[node];
   int mid = (b + e) \gg 1;
   if (pos <= mid)</pre>
     return query(node << 1, b, mid, pos);</pre>
    return query(node << 1 | 1, mid + 1, e, pos);</pre>
  void set(int from, int to, int v) {
    return set(1, 0, MN - 1, from, to, v);
  int query(int pos) {
    return query(1, 0, MN - 1, pos);
 }
};
```

3.12 sparse table

```
// RMQ.
const int MN = 100000 + 10; // Max number of elements
const int ML = 18; // ceil(log2(MN));

struct st {
  int data[MN];
  int M[MN][ML];
  int n;

void init(const vector<int> &d) {
```

```
n = d.size();
   for (int i = 0; i < n; ++i)</pre>
     data[i] = d[i];
   build();
 }
 void build() {
   for (int i = 0; i < n; ++i)
     M[i][0] = data[i];
   for (int j = 1, p = 2, q = 1; p \le n; ++j, p \le 1, q \le 1)
     for (int i = 0; i + p - 1 < n; ++i)
       M[i][j] = max(M[i][j-1], M[i+q][j-1]);
 }
 int query(int b, int e) {
   int k = log2(e - b + 1);
   return max(M[b][k], M[e + 1 - (1<<k)][k]);</pre>
 }
};
```

3.13 splay tree

```
using namespace std;
#include<bits/stdc++.h>
#define D(x) cout<<x<<endl;

typedef int T;

struct node{
  node *left, *right, *parent;
  T key;
  node (T k) : key(k), left(0), right(0), parent(0) {}
};

struct splay_tree{

  node *root;

  void right_rot(node *x) {
    node *p = x->parent;
    if (x->parent = p->parent) {
        if (x->parent->left = x;
```

```
if (x->parent->right == p) x->parent->right = x;
 if (p->left = x->right) p->left->parent = p;
 x->right = p;
 p->parent = x;
void left_rot(node *x) {
 node *p = x->parent;
 if (x->parent = p->parent) {
   if (x->parent->left == p) x->parent->left = x;
   if (x->parent->right == p) x->parent->right = x;
 if (p->right = x->left) p->right->parent = p;
 x->left = p;
 p->parent = x;
void splay(node *x, node *fa = 0) {
 while( x->parent != fa and x->parent != 0) {
   node *p = x->parent;
   if (p->parent == fa)
     if (p->right == x)
      left_rot(x);
     else
      right_rot(x);
   else {
     node *gp = p->parent; //grand parent
     if (gp->left == p)
       if (p->left == x)
         right_rot(x), right_rot(x);
        left_rot(x),right_rot(x);
     else
       if (p->left == x)
        right_rot(x), left_rot(x);
        left_rot(x), left_rot(x);
 if (fa == 0) root = x;
void insert(T key) {
```

```
node *cur = root;
   node *pcur = 0;
   while (cur) {
     pcur = cur;
     if (key > cur->key) cur = cur->right;
     else cur = cur->left;
   cur = new node(key);
   cur->parent = pcur;
   if (!pcur) root = cur;
   else if (key > pcur->key ) pcur->right = cur;
   else pcur->left = cur;
   splay(cur);
 }
 node *find(T key) {
   node *cur = root;
   while (cur) {
     if (key > cur->key) cur = cur->right;
     else if(key < cur->key) cur = cur->left;
     else return cur;
   }
   return 0;
 }
  splay_tree(){ root = 0;};
};
```

3.14 trie

```
const int MN = 26; // size of alphabet
const int MS = 100010; // Number of states.

struct trie{
    struct node{
      int c;
      int a[MN];
    };

    node tree[MS];
    int nodes;

    void clear(){
```

```
tree[nodes].c = 0;
   memset(tree[nodes].a, -1, sizeof tree[nodes].a);
   nodes++;
 void init(){
   nodes = 0;
   clear();
  int add(const string &s, bool query = 0){
   int cur_node = 0;
   for(int i = 0; i < s.size(); ++i){</pre>
     int id = gid(s[i]);
     if(tree[cur_node].a[id] == -1){
       if(query) return 0;
       tree[cur_node].a[id] = nodes;
       clear();
     cur_node = tree[cur_node].a[id];
   if(!query) tree[cur_node].c++;
   return tree[cur_node].c;
};
```

3.15 wavelet tree

```
// this can be tested in the problem:
    http://www.spoj.com/problems/ILKQUERY/

struct wavelet {
    vector<int> values, ori;
    vector<int> map_left, map_right;
    int 1, r, m;
    wavelet *left, *right;
    wavelet() : left(NULL), right(NULL) {}
    wavelet(int a, int b, int c) : l(a), r(b), m(c), left(NULL),
        right(NULL) {}
};

wavelet *init(vector<int> &data, vector<int> &int lo, int lo, int hi) {
```

```
if (lo > hi || (data.size() == 0)) return NULL;
 int mid = ((long long)(lo) + hi) / 2;
 if (lo + 1 == hi) mid = lo: // handle negative values
  wavelet *node = new wavelet(lo, hi, mid);
 vector<int> data_1, data_r, ind_1, ind_r;
  int ls = 0, rs = 0;
 for (int i = 0; i < int(data.size()); i++) {</pre>
   int value = data[i];
   if (value <= mid) {</pre>
     data_l.emplace_back(value);
     ind_l.emplace_back(ind[i]);
     ls++:
   } else {
     data_r.emplace_back(value);
     ind_r.emplace_back(ind[i]);
     rs++;
   node->map_left.emplace_back(ls);
   node->map_right.emplace_back(rs);
   node->values.emplace_back(value);
   node->ori.emplace_back(ind[i]);
 if (lo < hi) {</pre>
   node->left = init(data_1, ind_1, lo, mid);
   node->right = init(data_r, ind_r, mid + 1, hi);
 }
 return node;
}
int kth(wavelet *node, int to, int k) {
 // returns the kth element in the sorted version of (a[0], ..., a[to])
 if (node->1 == node->r) return node->m;
 int c = node->map_left[to];
 if (k < c)
   return kth(node->left, c - 1, k);
 return kth(node->right, node->map_right[to] - 1, k - c);
}
int pos_kth_ocurrence(wavelet *node, int val, int k) {
 // returns the position on the original array of the kth ocurrence of
      the value "val"
 if (!node) return -1;
```

```
if (node->l == node->r) {
   if (int(node->ori.size()) <= k)
     return -1;
   return node->ori[k];
}

if (val <= node->m)
   return pos_kth_ocurrence(node->left, val, k);
   return pos_kth_ocurrence(node->right, val, k);
}
```

4 Geometry

4.1 center 2 points + radious

```
vector<point> find_center(point a, point b, long double r) {
  point d = (a - b) * 0.5;
  if (d.dot(d) > r * r) {
    return vector<point> ();
  }
  point e = b + d;
  long double fac = sqrt(r * r - d.dot(d));
  vector<point> ans;
  point x = point(-d.y, d.x);
  long double l = sqrt(x.dot(x));
  x = x * (fac / l);
  ans.push_back(e + x);
  x = point(d.y, -d.x);
  x = x * (fac / l);
  ans.push_back(e + x);
  return ans;
}
```

4.2 closest pair problem

```
struct point {
  double x, y;
  int id;
  point() {}
```

```
point (double a, double b) : x(a), y(b) {}
};
double dist(const point &o, const point &p) {
 double a = p.x - o.x, b = p.y - o.y;
 return sqrt(a * a + b * b);
}
double cp(vector<point> &p, vector<point> &x, vector<point> &y) {
 if (p.size() < 4) {</pre>
    double best = 1e100:
   for (int i = 0; i < p.size(); ++i)</pre>
     for (int j = i + 1; j < p.size(); ++j)</pre>
       best = min(best, dist(p[i], p[j]));
   return best;
 }
  int ls = (p.size() + 1) >> 1;
  double l = (p[ls - 1].x + p[ls].x) * 0.5;
  vector<point> xl(ls), xr(p.size() - ls);
  unordered_set<int> left;
 for (int i = 0; i < ls; ++i) {</pre>
   xl[i] = x[i];
    left.insert(x[i].id);
 for (int i = ls; i < p.size(); ++i) {</pre>
   xr[i - ls] = x[i];
 }
  vector<point> yl, yr;
  vector<point> pl, pr;
  yl.reserve(ls); yr.reserve(p.size() - ls);
 pl.reserve(ls); pr.reserve(p.size() - ls);
 for (int i = 0; i < p.size(); ++i) {</pre>
   if (left.count(v[i].id))
     yl.push_back(y[i]);
    else
     yr.push_back(y[i]);
    if (left.count(p[i].id))
     pl.push_back(p[i]);
    else
     pr.push_back(p[i]);
  }
```

```
double dl = cp(pl, xl, yl);
  double dr = cp(pr, xr, yr);
  double d = min(dl, dr);
  vector<point> yp; yp.reserve(p.size());
  for (int i = 0; i < p.size(); ++i) {</pre>
   if (fabs(y[i].x - 1) < d)
     yp.push_back(y[i]);
  for (int i = 0; i < yp.size(); ++i) {</pre>
   for (int j = i + 1; j < yp.size() && j < i + 7; ++j) {
     d = min(d, dist(yp[i], yp[j]));
   }
 }
  return d;
}
double closest_pair(vector<point> &p) {
  vector<point> x(p.begin(), p.end());
  sort(x.begin(), x.end(), [](const point &a, const point &b) {
   return a.x < b.x;</pre>
 });
  vector<point> y(p.begin(), p.end());
  sort(y.begin(), y.end(), [](const point &a, const point &b) {
   return a.y < b.y;</pre>
 });
  return cp(p, x, y);
```

4.3 squares

```
typedef long double ld;

const ld eps = 1e-12;
int cmp(ld x, ld y = 0, ld tol = eps) {
    return ( x <= y + tol) ? (x + tol < y) ? -1 : 0 : 1;
}

struct point{
    ld x, y;
    point(ld a, ld b) : x(a), y(b) {}
    point() {}
};</pre>
```

```
struct square{
 ld x1, x2, y1, y2,
    a, b, c;
 point edges[4];
  square(ld _a, ld _b, ld _c) {
   a = _a, b = _b, c = _c;
   x1 = a - c * 0.5:
   x2 = a + c * 0.5:
   v1 = b - c * 0.5;
   v2 = b + c * 0.5:
   edges[0] = point(x1, y1);
   edges[1] = point(x2, y1);
   edges[2] = point(x2, y2);
   edges[3] = point(x1, y2);
 }
};
ld min_dist(point &a, point &b) {
 1d x = a.x - b.x
    y = a.y - b.y;
 return sqrt(x * x + y * y);
}
bool point_in_box(square s1, point p) {
 if (cmp(s1.x1, p.x) != 1 && cmp(s1.x2, p.x) != -1 &&
     cmp(s1.y1, p.y) != 1 && cmp(s1.y2, p.y) != -1)
   return true;
 return false;
}
bool inside(square &s1, square &s2) {
 for (int i = 0; i < 4; ++i)</pre>
   if (point_in_box(s2, s1.edges[i]))
     return true;
 return false;
}
bool inside_vert(square &s1, square &s2) {
 if ((cmp(s1.y1, s2.y1) != -1 && cmp(s1.y1, s2.y2) != 1) ||
     (cmp(s1.y2, s2.y1) != -1 \&\& cmp(s1.y2, s2.y2) != 1))
   return true;
return false;
}
```

```
bool inside_hori(square &s1, square &s2) {
 if ((cmp(s1.x1, s2.x1) != -1 && cmp(s1.x1, s2.x2) != 1) ||
     (cmp(s1.x2, s2.x1) != -1 \&\& cmp(s1.x2, s2.x2) != 1))
   return true;
return false;
ld min_dist(square &s1, square &s2) {
 if (inside(s1, s2) || inside(s2, s1))
   return 0:
 1d ans = 1e100;
 for (int i = 0; i < 4; ++i)
   for (int j = 0; j < 4; ++j)
     ans = min(ans, min_dist(s1.edges[i], s2.edges[j]));
 if (inside_hori(s1, s2) || inside_hori(s2, s1)) {
   if (cmp(s1.v1, s2.v2) != -1)
     ans = min(ans, s1.y1 - s2.y2);
   else
   if (cmp(s2.y1, s1.y2) != -1)
     ans = min(ans, s2.y1 - s1.y2);
 if (inside_vert(s1, s2) || inside_vert(s2, s1)) {
   if (cmp(s1.x1, s2.x2) != -1)
     ans = min(ans, s1.x1 - s2.x2);
   if (cmp(s2.x1, s1.x2) != -1)
     ans = min(ans, s2.x1 - s1.x2);
 return ans;
```

4.4 triangles

Let a, b, c be length of the three sides of a triangle.

$$p = (a + b + c) * 0.5$$

The inradius is defined by:

$$iR = \sqrt{\frac{(p-a)(p-b)(p-c)}{p}}$$

The radius of its circumcircle is given by the formula:

$$cR = \frac{abc}{\sqrt{(a+b+c)(a+b-c)(a+c-b)(b+c-a)}}$$

5 Graphs

5.1 SCC kosaraju

```
struct SCC {
 vector<vector<int> > g, gr;
 vector<bool> used;
 vector<int> order, component;
 int total_components;
 SCC(vector<vector<int> > &adj) {
   g = adi;
   int n = g.size();
   gr.resize(n);
   for (int i = 0; i < n; i++)</pre>
     for (auto to : g[i])
       gr[to].push_back(i);
   used.assign(n, false);
   for (int i = 0; i < n; i++)</pre>
     if (!used[i])
       GenTime(i);
   used.assign(n, false);
   component.assign(n, -1);
   total_components = 0;
   for (int i = n - 1; i >= 0; i--) {
     int v = order[i];
     if (!used[v]) {
       vector<int> cur_component;
       Dfs(cur_component, v);
       for (auto node : cur_component)
         component[node] = total_components;
       total_components++;
```

```
}
 }
  void GenTime(int node) {
   used[node] = true;
   for (auto to : g[node])
     if (!used[to])
       GenTime(to):
   order.push_back(node);
 void Dfs(vector<int> &cur, int node) {
   used[node] = true;
   cur.push_back(node);
   for (auto to : gr[node])
     if (!used[to])
       Dfs(cur, to);
 }
  vector<vector<int>> CondensedGraph() {
   vector<vector<int>> ans(total_components);
   for (int i = 0; i < int(g.size()); i++) {</pre>
     for (int to : g[i]) {
       int u = component[i], v = component[to];
       if (u != v)
         ans[u].push_back(v);
     }
   return ans;
};
```

5.2 board

```
struct board {
  int n, m, r;
  board(int a, int b, int c = 1) : n(a), m(b), r(c) {}

long long frec(int x, int y) {
   // returns how many squares of r x r contain the cell (x, y)
  long long a = min(x, n - r) - max(x - r + 1, 0) + 1;
  long long b = min(y, m - r) - max(y - r + 1, 0) + 1;
  return a * b;
```

```
}
bool valid(int x, int y) {
   return x >= 0 && x < n && y >= 0 && y < m;
}
};</pre>
```

5.3 bridges

```
struct Graph {
 vector<vector<Edge>> g;
 vector<int> vi, low, d, pi, is_b;
 int bridges_computed;
 int ticks, edges;
 Graph(int n, int m) {
   g.assign(n, vector<Edge>());
   is_b.assign(m, 0);
   vi.resize(n);
   low.resize(n);
   d.resize(n);
   pi.resize(n);
   edges = 0;
   bridges_computed = 0;
 }
 void AddEdge(int u, int v) {
   g[u].push_back(Edge(v, edges));
   g[v].push_back(Edge(u, edges));
   edges++;
 }
 void Dfs(int u) {
   vi[u] = true;
   d[u] = low[u] = ticks++;
   for (int i = 0; i < (int)g[u].size(); ++i) {</pre>
     int v = g[u][i].to;
     if (v == pi[u]) continue;
     if (!vi[v]) {
       pi[v] = u;
       Dfs(v);
       if (d[u] < low[v]) is_b[g[u][i].id] = true;</pre>
```

```
low[u] = min(low[u], low[v]);
     } else {
       low[u] = min(low[u], d[v]);
   }
 }
 // Multiple edges from a to b are not allowed.
 // (they could be detected as a bridge).
 // If you need to handle this, just count
 // how many edges there are from a to b.
 void CompBridges() {
   fill(pi.begin(), pi.end(), -1);
   fill(vi.begin(), vi.end(), 0);
   fill(low.begin(), low.end(), 0);
   fill(d.begin(), d.end(), 0);
   ticks = 0;
   for (int i = 0; i < (int)g.size(); ++i)</pre>
     if (!vi[i]) Dfs(i);
   bridges_computed = true;
 }
 map<int, vector<Edge>> BridgesTree() {
   if (!bridges_computed) CompBridges();
   int n = g.size();
   Dsu dsu(g.size());
   for (int i = 0; i < n; i++)</pre>
     for (auto e : g[i])
       if (!is_b[e.id]) dsu.Join(i, e.to);
   map<int, vector<Edge>> tree;
   for (int i = 0; i < n; i++)</pre>
     for (auto e : g[i])
       if (is_b[e.id])
         tree[dsu.Find(i)].emplace_back(dsu.Find(e.to), e.id);
   return tree;
 }
};
```

5.4 dijkstra

```
struct edge {
 int to;
 long long w;
 edge () {}
  edge (int a, long long b) : to(a), w(b) {}
 bool operator < (const edge &o) const {</pre>
   return w > o.w;
 }
};
typedef vector<vector<edge>> graph;
const long long inf = 1000000LL * 10000000LL;
pair<vector<int>, vector<long long>> dijkstra(graph &g, int start) {
 int n = g.size();
 vector<long long> d(n, inf);
 vector<int> p(n, -1);
 d[start] = 0;
 priority_queue<edge> q;
 q.push(edge(start, 0));
  while (!q.empty()) {
   int node = q.top().to;
   long long dist = q.top().w;
   q.pop();
   if (dist > d[node]) continue;
   for (int i = 0; i < (int)g[node].size(); i++) {</pre>
     int to = g[node][i].to;
     long long w_extra = g[node][i].w;
     if (dist + w_extra < d[to]) {</pre>
       p[to] = node;
       d[to] = dist + w_extra;
       q.push(edge(to, d[to]));
   }
 }
 return {p, d};
```

5.5 directed mst

```
const int inf = 1000000 + 10;
struct edge {
 int u, v, w;
 edge() {}
 edge(int a,int b,int c): u(a), v(b), w(c) {}
* Computes the minimum spanning tree for a directed graph
* - edges : Graph description in the form of list of edges.
     each edge is: From node u to node v with cost w
* - root : Id of the node to start the DMST.
* - n : Number of nodes in the graph.
* */
int dmst(vector<edge> &edges, int root, int n) {
 int ans = 0;
 int cur_nodes = n;
 while (true) {
   vector<int> lo(cur_nodes, inf), pi(cur_nodes, inf);
   for (int i = 0; i < edges.size(); ++i) {</pre>
     int u = edges[i].u, v = edges[i].v, w = edges[i].w;
     if (w < lo[v] and u != v) {
       lo[v] = w;
       pi[v] = u;
   }
   lo[root] = 0;
   for (int i = 0; i < lo.size(); ++i) {</pre>
     if (i == root) continue;
     if (lo[i] == inf) return -1;
   int cur_id = 0;
   vector<int> id(cur_nodes, -1), mark(cur_nodes, -1);
   for (int i = 0; i < cur_nodes; ++i) {</pre>
     ans += lo[i];
     int u = i:
     while (u != root and id[u] < 0 and mark[u] != i) {</pre>
       mark[u] = i;
       u = pi[u];
```

```
if (u != root and id[u] < 0) { // Cycle}
      for (int v = pi[u]; v != u; v = pi[v])
        id[v] = cur_id;
      id[u] = cur_id++;
   }
  }
  if (cur id == 0)
   break;
  for (int i = 0; i < cur_nodes; ++i)</pre>
    if (id[i] < 0) id[i] = cur_id++;</pre>
  for (int i = 0; i < edges.size(); ++i) {</pre>
    int u = edges[i].u, v = edges[i].v, w = edges[i].w;
    edges[i].u = id[u];
    edges[i].v = id[v];
    if (id[u] != id[v])
     edges[i].w -= lo[v];
  }
  cur_nodes = cur_id;
  root = id[root];
}
return ans;
```

5.6 eulerian path

```
// Taken from https://github.com/lbv/pc-code/blob/master/code/graph.cpp
// Eulerian Trail

struct Euler {
    ELV adj; IV t;
    Euler(ELV Adj) : adj(Adj) {}
    void build(int u) {
        while(! adj[u].empty()) {
            int v = adj[u].front().v;
            adj[u].erase(adj[u].begin());
            build(v);
        }
        t.push_back(u);
    }
}
```

```
};
bool eulerian_trail(IV &trail) {
 Euler e(adj);
 int odd = 0, s = 0;
    for (int v = 0; v < n; v++) {
    int diff = abs(in[v] - out[v]);
    if (diff > 1) return false:
    if (diff == 1) {
    if (++odd > 2) return false;
    if (out[v] > in[v]) start = v;
    }
    */
 e.build(s);
 reverse(e.t.begin(), e.t.end());
 trail = e.t;
 return true;
```

5.7 karp min mean cycle

```
/**
* Finds the min mean cycle, if you need the max mean cycle
 * just add all the edges with negative cost and print
 * ans * -1
 * test: uva, 11090 - Going in Cycle!!
const int MN = 1000;
struct edge{
 int v;
 long long w;
  edge()\{\} edge(int v, int w) : v(v), w(w) \{\}
};
long long d[MN][MN];
// This is a copy of g because increments the size
// pass as reference if this does not matter.
int karp(vector<vector<edge> > g) {
  int n = g.size();
```

```
g.resize(n + 1); // this is important
for (int i = 0; i < n; ++i)
  if (!g[i].empty())
   g[n].push_back(edge(i,0));
++n;
for(int i = 0:i < n:++i)
  fill(d[i],d[i]+(n+1),INT_MAX);
d[n - 1][0] = 0:
for (int k = 1; k \le n; ++k) for (int u = 0; u \le n; ++u) {
  if (d[u][k - 1] == INT_MAX) continue;
  for (int i = g[u].size() - 1; i >= 0; --i)
   d[g[u][i].v][k] = min(d[g[u][i].v][k], d[u][k-1] + g[u][i].w);
}
bool flag = true;
for (int i = 0; i < n && flag; ++i)</pre>
  if (d[i][n] != INT_MAX)
   flag = false;
if (flag) {
  return true; // return true if there is no a cycle.
double ans = 1e15;
for (int u = 0; u + 1 < n; ++u) {
 if (d[u][n] == INT_MAX) continue;
  double W = -1e15;
  for (int k = 0; k < n; ++k)
   if (d[u][k] != INT_MAX)
     W = max(W, (double)(d[u][n] - d[u][k]) / (n - k));
  ans = min(ans, W);
}
// printf("%.21f\n", ans);
cout << fixed << setprecision(2) << ans << endl;</pre>
return false;
```

5.8 konig's theorem

In any bipartite graph, the number of edges in a maximum matching equals the number of vertices in a minimum vertex cover

5.9 minimum path cover in DAG

Given a directed acyclic graph G = (V, E), we are to find the minimum number of vertex-disjoint paths to cover each vertex in V.

We can construct a bipartite graph $G' = (Vout \cup Vin, E')$ from G, where :

$$Vout = \{v \in V : v \text{ has positive out} - degree\}$$

$$Vin = \{v \in V : v \text{ has positive in} - degree\}$$

$$E' = \{(u, v) \in Vout \times Vin : (u, v) \in E\}$$

Then it can be shown, via König's theorem, that G' has a matching of size m if and only if there exists n-m vertex-disjoint paths that cover each vertex in G, where n is the number of vertices in G and m is the maximum cardinality bipartite mathching in G'.

Therefore, the problem can be solved by finding the maximum cardinality matching in G' instead.

NOTE: If the paths are note necessarily disjoints, find the transitive closure and solve the problem for disjoint paths.

5.10 planar graph (euler)

Euler's formula states that if a finite, connected, planar graph is drawn in the plane without any edge intersections, and v is the number of vertices, e is the number of edges and f is the number of faces (regions bounded by edges, including the outer, infinitely large region), then:

$$f + v = e + 2$$

It can be extended to non connected planar graphs with \boldsymbol{c} connected components:

$$f + v = e + c + 1$$

5.11 query with lca

```
struct lowest_ca {
 int T[MN], L[MN], W[MN];
 int P[MN][ML], MI[MN][ML], MA[MN][ML];
 void dfs(vector<vector<edge> > &g, int root, int pi = -1) {
   if (pi == -1) {
     L[root] = W[root] = 0;
     T[root] = -1;
   for (int i = 0; i < (int)g[root].size(); ++i) {</pre>
     int to = g[root][i].v;
     if (to != pi) {
       T[to] = root;
       W[to] = g[root][i].w;
      L[to] = L[root] + 1;
       dfs(g, to, root);
    }
   }
 }
 void init(vector<vector<edge> > &g, int root) {
   // g is undirected
   dfs(g, root);
   int N = g.size(), i, j;
   for (i = 0; i < N; i++) {</pre>
     for (j = 0; 1 << j < N; j++) {
       P[i][j] = -1;
      MI[i][j] = inf;
     }
   }
   for (i = 0; i < N; i++) {</pre>
     P[i][0] = T[i];
     MI[i][0] = W[i];
   }
   for (j = 1; 1 << j < N; j++)
     for (i = 0; i < N; i++)</pre>
       if (P[i][j - 1] != -1) {
        P[i][j] = P[P[i][j - 1]][j - 1];
        MI[i][j] = min(MI[i][j-1], MI[P[i][j-1]][j-1]);
       }
```

```
}
int query(int p, int q) {
  int tmp, log, i;
  int mmin = inf;
  if (L[p] < L[q])
   tmp = p, p = q, q = tmp;
  for (log = 1; 1 << log <= L[p]; log++);</pre>
  log--;
  for (i = log; i >= 0; i--)
   if (L[p] - (1 << i) >= L[q]) {
     mmin = min(mmin, MI[p][i]);
     p = P[p][i];
  if (p == q)
   // return p;
   return mmin;
  for (i = log; i >= 0; i--)
   if (P[p][i] != -1 && P[p][i] != P[q][i]) {
     mmin = min(mmin, min(MI[p][i], MI[q][i]));
     p = P[p][i], q = P[q][i];
   }
  // return T[p];
  return min(mmin, min(MI[p][0], MI[q][0]));
int get_child(int p, int q) { // p is ancestor of q
  if (p == q) return -1;
 int i, log;
  for (log = 1; 1 << log <= L[q]; log++) {}
  log--;
 for (i = log; i >= 0; i--)
   if (L[q] - (1 << i) > L[p]) {
     q = P[q][i];
  assert(P[q][0] == p);
```

```
return q;
 }
  int is_ancestor(int p, int q) {
   if (L[p] >= L[q])
     return false;
   int dist = L[q] - L[p];
   int cur = q;
   int step = 0;
   while (dist) {
     if (dist & 1)
       cur = P[cur][step];
     step++;
     dist >>= 1;
   return cur == p;
 }
};
```

5.12 tarjan scc

```
const int MN = 20002;

struct tarjan_scc {
  int scc[MN], low[MN], d[MN], stacked[MN];
  int ticks, current_scc;
  deque<int> s; // used as stack.

tarjan_scc() {}

void init () {
  memset(scc, -1, sizeof scc);
  memset(d, -1, sizeof d);
  memset(stacked, 0, sizeof stacked);
  s.clear();
  ticks = current_scc = 0;
}

void compute(vector<vector<int> > &g, int u) {
  d[u] = low[u] = ticks++;
```

```
s.push_back(u);
   stacked[u] = true;
   for (int i = 0; i < g[u].size(); ++i) {</pre>
     int v = g[u][i];
     if (d[v] == -1)
       compute(g, v);
     if (stacked[v]) {
       low[u] = min(low[u], low[v]);
     }
   }
   if (d[u] == low[u]) { // root
     int v;
     do {
       v = s.back();s.pop_back();
       stacked[v] = false;
       scc[v] = current_scc;
     } while (u != v);
     current_scc++;
 }
};
```

5.13 two sat (with kosaraju)

```
/**
 * Given a set of clauses (a1 v a2)^(a2 v a3)....
 * this algorithm find a solution to it set of clauses.
 * test: http://lightoj.com/volume_showproblem.php?problem=1251
 **/

#include<bits/stdc++.h>
using namespace std;
#define MAX 100000
#define endl '\n'

vector<int> G[MAX];
vector<int> GT[MAX];
vector<int> Ftime;
vector<vector<int> > SCC;
bool visited[MAX];
int n;
```

```
void dfs1(int n){
 visited[n] = 1;
 for (int i = 0; i < G[n].size(); ++i) {</pre>
   int curr = G[n][i]:
   if (visited[curr]) continue;
   dfs1(curr);
 }
 Ftime.push_back(n);
}
void dfs2(int n, vector<int> &scc) {
 visited[n] = 1;
 scc.push_back(n);
 for (int i = 0;i < GT[n].size(); ++i) {</pre>
   int curr = GT[n][i]:
   if (visited[curr]) continue;
   dfs2(curr, scc);
 }
}
void kosaraju() {
 memset(visited, 0, sizeof visited);
 for (int i = 0; i < 2 * n; ++i) {</pre>
   if (!visited[i]) dfs1(i);
 }
 memset(visited, 0, sizeof visited);
 for (int i = Ftime.size() - 1; i >= 0; i--) {
   if (visited[Ftime[i]]) continue;
   vector<int> _scc;
   dfs2(Ftime[i],_scc);
   SCC.push_back(_scc);
 }
}
 * After having the SCC, we must traverse each scc, if in one SCC are -b
     y b, there is not a solution.
```

```
* Otherwise we build a solution, making the first "node" that we find
     truth and its complement false.
bool two sat(vector<int> &val) {
 kosaraju();
 for (int i = 0; i < SCC.size(); ++i) {</pre>
   vector<bool> tmpvisited(2 * n, false);
   for (int j = 0; j < SCC[i].size(); ++j) {</pre>
     if (tmpvisited[SCC[i][j] ^ 1]) return 0;
     if (val[SCC[i][j]] != -1) continue;
     else {
       val[SCC[i][j]] = 0;
       val[SCC[i][j] ^ 1] = 1;
     tmpvisited[SCC[i][j]] = 1;
 }
 return 1;
// Example of use
int main() {
 int m, u, v, nc = 0, t; cin >> t;
 // n = "nodes" number, m = clauses number
 while (t--) {
   cin >> m >> n;
   Ftime.clear():
   SCC.clear();
   for (int i = 0; i < 2 * n; ++i) {
     G[i].clear();
     GT[i].clear();
   // (a1 v a2) = (a1 -> a2) = (a2 -> a1)
   for (int i = 0; i < m ; ++i) {</pre>
     cin >> u >> v;
     int t1 = abs(u) - 1;
     int t2 = abs(v) - 1;
     int p = t1 * 2 + ((u < 0)? 1 : 0);
     int q = t2 * 2 + ((v < 0)? 1 : 0);
```

```
G[p ^ 1].push_back(q);
      G[q ^ 1].push_back(p);
      GT[p].push_back(q ^ 1);
      GT[q].push_back(p ^ 1);
    vector < int > val(2 * n, -1);
    cout << "Case " << ++nc <<": ":
    if (two_sat(val)) {
      cout << "Yes" << endl;</pre>
      vector<int> sol:
     for (int i = 0; i < 2 * n; ++i)
        if (i % 2 == 0 and val[i] == 1)
          sol.push_back(i / 2 + 1);
      cout << sol.size() ;</pre>
      for (int i = 0; i < sol.size(); ++i) {</pre>
        cout << " " << sol[i];
      cout << endl;</pre>
    } else {
      cout << "No" << endl;</pre>
    }
 }
 return 0;
}
```

6 Math

6.1 Lucas theorem

For non-negative integers m and n and a prime p, the following congruence relation holds: :

$$\binom{m}{n} \equiv \prod_{i=0}^{k} \binom{m_i}{n_i} \pmod{p},$$

where:

$$m = m_k p^k + m_{k-1} p^{k-1} + \dots + m_1 p + m_0,$$

and:

$$n = n_k p^k + n_{k-1} p^{k-1} + \dots + n_1 p + n_0$$

are the base p expansions of m and n respectively. This uses the convention that $\binom{m}{n} = 0$ if $m \le n$.

6.2 counting

```
const int MN = 1e5 + 100;
long long fact[MN];

void fill_fact() {
  fact[0] = 1;
  for (int i = 1; i < MN; i++) {
    fact[i] = mult(fact[i - 1], i);
  }
}

long long perm_rep(vector<int> &frec) {
  int total = 0;
  long long den = 1;
  for (int i = 0; i < (int)frec.size(); i++) {
    den = mult(den, mod_inv(fact[frec[i]]));
    total += frec[i];
  }
  return mult(fact[total], den);
}</pre>
```

6.3 cumulative sum of divisors

```
/**
The function SOD(n) (sum of divisors) is defined
as the summation of all the actual divisors of
an integer number n. For example,

SOD(24) = 2+3+4+6+8+12 = 35.

The function CSOD(n) (cumulative SOD) of an integer n, is defined as
below:

csod(n) = \sum_{{i = 1}^{n}} sod(i)

It can be computed in O(sqrt(n)):
*/

long long csod(long long n) {
 long long ans = 0;
 for (long long i = 2; i * i <= n; ++i) {
  long long i = n / i;
}</pre>
```

```
ans += (i + j) * (j - i + 1) / 2;
ans += i * (j - i);
}
return ans;
}
```

6.4 fft.

```
* Fast Fourier Transform.
 * Useful to compute convolutions.
 * computes:
 * C(f \operatorname{star} g)[n] = \operatorname{sum}_m(f[m] * g[n - m])
 * for all n.
 * test: icpc live archive, 6886 - Golf Bot
 * */
using namespace std;
#include<bits/stdc++.h>
#define D(x) cout << #x " = " << (x) << endl
#define endl '\n'
const int MN = 262144 << 1:</pre>
int d[MN + 10], d2[MN + 10];
const double PI = acos(-1.0);
struct cpx {
 double real, image;
  cpx(double _real, double _image) {
   real = _real;
    image = _image;
 }
  cpx(){}
};
cpx operator + (const cpx &c1, const cpx &c2) {
 return cpx(c1.real + c2.real, c1.image + c2.image);
}
cpx operator - (const cpx &c1, const cpx &c2) {
```

```
return cpx(c1.real - c2.real, c1.image - c2.image);
cpx operator * (const cpx &c1, const cpx &c2) {
 return cpx(c1.real*c2.real - c1.image*c2.image, c1.real*c2.image +
      c1.image*c2.real);
}
int rev(int id, int len) {
 int ret = 0;
 for (int i = 0; (1 << i) < len; i++) {
   ret <<= 1;
   if (id & (1 << i)) ret |= 1;</pre>
 return ret;
}
cpx A[1 << 20];
void FFT(cpx *a, int len, int DFT) {
 for (int i = 0; i < len; i++)</pre>
   A[rev(i, len)] = a[i];
 for (int s = 1; (1 << s) <= len; s++) {
   int m = (1 << s);
   cpx wm = cpx(cos(DFT * 2 * PI / m), sin(DFT * 2 * PI / m));
   for(int k = 0; k < len; k += m) {</pre>
     cpx w = cpx(1, 0);
     for(int j = 0; j < (m >> 1); j++) {
       cpx t = w * A[k + j + (m >> 1)];
       cpx u = A[k + j];
       A[k + j] = u + t;
       A[k + j + (m >> 1)] = u - t;
       w = w * wm;
     }
   }
 }
 if (DFT == -1) for (int i = 0; i < len; i++) A[i].real /= len,</pre>
      A[i].image /= len;
 for (int i = 0; i < len; i++) a[i] = A[i];</pre>
 return:
cpx in[1 << 20];
void solve(int n) {
```

```
memset(d, 0, sizeof d);
  int t;
 for (int i = 0; i < n; ++i) {</pre>
    cin >> t;
    d[t] = true;
 }
  int m;
  cin >> m;
  vector<int> q(m);
 for (int i = 0; i < m; ++i)</pre>
    cin >> q[i];
 for (int i = 0; i < MN; ++i) {</pre>
    if (d[i])
     in[i] = cpx(1, 0);
    else
      in[i] = cpx(0, 0);
 }
 FFT(in, MN, 1);
 for (int i = 0; i < MN; ++i) {</pre>
   in[i] = in[i] * in[i];
 FFT(in, MN, -1);
 int ans = 0:
 for (int i = 0; i < q.size(); ++i) {</pre>
   if (in[q[i]].real > 0.5 || d[q[i]]) {
      ans++;
   }
 }
  cout << ans << endl;</pre>
}
int main() {
  ios_base::sync_with_stdio(false);cin.tie(NULL);
  int n;
  while (cin >> n)
    solve(n);
  return 0;
```

6.5 fibonacci properties

Let A, B and n be integer numbers.

$$k = A - B \tag{1}$$

$$F_A F_B = F_{k+1} F_A^2 + F_k F_A F_{A-1} \tag{2}$$

$$\sum_{i=0}^{n} F_i^2 = F_{n+1} F_n \tag{3}$$

ev(n) = returns 1 if n is even.

$$\sum_{i=0}^{n} F_i F_{i+1} = F_{n+1}^2 - ev(n) \tag{4}$$

$$\sum_{i=0}^{n} F_i F_{i-1} = \sum_{i=0}^{n-1} F_i F_{i+1}$$
 (5)

6.6 polynomials

```
const double pi = acos(-1);
struct poly {
 deque<double> coef;
 double x_lo, x_hi;
 double evaluate(double x) {
   double ans = 0;
   for (auto it : coef)
     ans = (ans * x + it);
   return ans;
 double volume(double x, double dx=1e-6) {
   dx = (x_hi - x_lo) / 1000000.0;
   double ans = 0;
   for (double ix = x_lo; ix <= x; ix += dx) {
     double rad = evaluate(ix);
     ans += pi * rad * rad * dx;
   return ans;
```

}:

6.7 sigma function

the sigma function is defined as:

$$\sigma_x(n) = \sum_{d|n} d^x$$

when x = 0 is called the divisor function, that counts the number of positive divisors of n.

Now, we are interested in find

$$\sum_{d|n} \sigma_0(d)$$

if n is written as prime factorization:

$$n = \prod_{i=1}^{k} P_i^{e_k}$$

we can demonstrate that:

$$\sum_{d|n} \sigma_0(d) = \prod_{i=1}^k g(e_k + 1)$$

where g(x) is the sum of the first x positive numbers:

$$g(x) = (x * (x+1))/2$$

7 Matrix

7.1 matrix

```
const int MN = 111;
const int mod = 10000;

struct matrix {
  int r, c;
  int m[MN] [MN];

matrix (int _r, int _c) : r (_r), c (_c) {
```

```
memset(m, 0, sizeof m);
   for (int i = 0; i < r; ++i) {</pre>
     for (int j = 0; j < c; ++j)
       cout << m[i][j] << " ";
     cout << endl:</pre>
   }
 }
  int x[MN][MN];
  matrix & operator *= (const matrix &o) {
   memset(x, 0, sizeof x);
   for (int i = 0; i < r; ++i)
     for (int k = 0; k < c; ++k)
       if (m[i][k] != 0)
         for (int j = 0; j < c; ++j) {
           x[i][j] = (x[i][j] + ((m[i][k] * o.m[k][j]) % mod)) % mod;
   memcpy(m, x, sizeof(m));
   return *this;
};
void matrix_pow(matrix b, long long e, matrix &res) {
 memset(res.m, 0, sizeof res.m);
 for (int i = 0; i < b.r; ++i)</pre>
   res.m[i][i] = 1;
  if (e == 0) return;
  while (true) {
   if (e & 1) res *= b;
   if ((e >>= 1) == 0) break;
```

8 Misc

8.1 Template Java

```
import java.io.*;
import java.util.StringTokenizer;
public class Template {
   public static void main(String []args) throws IOException {
       Scanner in = new Scanner(System.in);
       OutputWriter out = new OutputWriter(System.out);
       Task solver = new Task():
       solver.solve(in, out);
       out.close();
   }
}
class Task{
   public void solve(Scanner in, OutputWriter out){
class Scanner{
   public BufferedReader reader;
   public StringTokenizer st;
   public Scanner(InputStream stream){
       reader = new BufferedReader(new InputStreamReader(stream));
       st = null;
   }
   public String next(){
       while(st == null || !st.hasMoreTokens()){
           try{
              String line = reader.readLine();
              if(line == null) return null;
              st = new StringTokenizer(line);
           }catch (Exception e){
              throw (new RuntimeException());
           }
       return st.nextToken();
   }
   public int nextInt(){
```

```
return Integer.parseInt(next());
   public long nextLong(){
       return Long.parseLong(next());
   public double nextDouble(){
       return Double.parseDouble(next());
}
class OutputWriter{
   BufferedWriter writer;
   public OutputWriter(OutputStream stream){
       writer = new BufferedWriter(new OutputStreamWriter(stream));
   public void print(int i) throws IOException {
       writer.write(i):
   public void print(String s) throws IOException {
       writer.write(s);
   public void print(char []c) throws IOException {
       writer.write(c);
   public void close() throws IOException {
       writer.close();
```

8.2 dates

```
// same as A, but for a leap year
const int B[13] = { 0, 0, 31, 60, 91, 121, 152, 182, 213, 244, 274, 305,
    335 }:
// returns number of leap years up to, and including, y
int leap_years(int y) { return y / 4 - y / 100 + y / 400; }
bool is_leap(int y) { return y % 400 == 0 || (y % 4 == 0 && y % 100 !=
    0); }
// number of days in blocks of years
const int p400 = 400*365 + leap_years(400);
const int p100 = 100*365 + leap_years(100);
const int p4 = 4*365 + 1;
const int p1 = 365;
int date_to_days(int d, int m, int y)
 return (y - 1) * 365 + leap_years(y - 1) + (is_leap(y) ? B[m] : A[m]) +
      d:
void days_to_date(int days, int &d, int &m, int &y)
 bool top100; // are we in the top 100 years of a 400 block?
 bool top4; // are we in the top 4 years of a 100 block?
 bool top1; // are we in the top year of a 4 block?
 y = 1;
 top100 = top4 = top1 = false;
 y += ((days-1) / p400) * 400;
 d = (days-1) \% p400 + 1;
 if (d > p100*3) top100 = true, d -= 3*p100, y += 300;
 else y += ((d-1) / p100) * 100, d = (d-1) % p100 + 1;
 if (d > p4*24) top4 = true, d = 24*p4, y += 24*4;
 else y += ((d-1) / p4) * 4, d = (d-1) % p4 + 1;
 if (d > p1*3) top1 = true, d -= p1*3, y += 3;
 else y += (d-1) / p1, d = (d-1) % p1 + 1;
 const int *ac = top1 && (!top4 || top100) ? B : A;
 for (m = 1; m < 12; ++m) if (d \le ac[m + 1]) break;
 d = ac[m];
}
```

8.3 fraction

```
struct frac{
  long long x, y;
  frac(long long a, long long b) {
    long long g = __gcd(a, b);
    x = a / g;
    y = b / g;
  }
  bool operator < (const frac &o) const {
    return (x * o.y < y * o.x);
  }
};</pre>
```

8.4 io

```
// taken from :
    https://github.com/lbv/pc-code/blob/master/solved/c-e/diablo/diablo.cpp
// this is very fast as well :
    https://github.com/lbv/pc-code/blob/master/code/input.cpp
typedef unsigned int u32;
#define BUF 524288
struct Reader {
 char buf[BUF]; char b; int bi, bz;
 Reader() { bi=bz=0; read(); }
 void read() {
   if (bi==bz) { bi=0; bz = fread(buf, 1, BUF, stdin); }
   b = bz ? buf[bi++] : 0; }
 void skip() { while (b > 0 && b <= 32) read(); }</pre>
  u32 next u32() {
   u32 v = 0; for (skip(); b > 32; read()) v = v*10 + b-48; return v; }
 int next_int() {
   int v = 0; bool s = false;
   skip(); if (b == '-') { s = true; read(); }
   for (; 48<=b&b<=57; read()) v = v*10 + b-48; return s ? -v : v; }
  char next_char() { skip(); char c = b; read(); return c; }
};
```

9 Number theory

9.1 convolution

```
typedef long long int LL;
typedef pair<LL, LL> PLL;
inline bool is_pow2(LL x) {
 return (x & (x-1)) == 0:
inline int ceil_log2(LL x) {
 int ans = 0;
 --x;
 while (x != 0) {
   x >>= 1;
   ans++:
 }
 return ans;
/* Returns the convolution of the two given vectors in time proportional
* The number of roots of unity to use nroots_unity must be set so that
     the product of the first
* nroots_unity primes of the vector nth_roots_unity is greater than the
     maximum value of the
* convolution. Never use sizes of vectors bigger than 2^24, if you need
     to change the values of
* the nth roots of unity to appropriate primes for those sizes.
vector<LL> convolve(const vector<LL> &a, const vector<LL> &b, int
    nroots_unity = 2) {
 int N = 1 << ceil_log2(a.size() + b.size());</pre>
 vector<LL> ans(N,0), fA(N), fB(N), fC(N);
 LL modulo = 1;
 for (int times = 0; times < nroots_unity; times++) {</pre>
   fill(fA.begin(), fA.end(), 0);
   fill(fB.begin(), fB.end(), 0);
   for (int i = 0; i < a.size(); i++) fA[i] = a[i];</pre>
   for (int i = 0; i < b.size(); i++) fB[i] = b[i];</pre>
   LL prime = nth_roots_unity[times].first;
   LL inv_modulo = mod_inv(modulo % prime, prime);
   LL normalize = mod_inv(N, prime);
```

```
ntfft(fA, 1, nth_roots_unity[times]);
ntfft(fB, 1, nth_roots_unity[times]);
for (int i = 0; i < N; i++) fC[i] = (fA[i] * fB[i]) % prime;
ntfft(fC, -1, nth_roots_unity[times]);
for (int i = 0; i < N; i++) {
   LL curr = (fC[i] * normalize) % prime;
   LL k = (curr - (ans[i] % prime) + prime) % prime;
   k = (k * inv_modulo) % prime;
   ans[i] += modulo * k;
}
modulo *= prime;
}
return ans;
}</pre>
```

9.2 crt

```
/**
  * Chinese remainder theorem.
  * Find z such that z % x[i] = a[i] for all i.
  * */
long long crt(vector<long long> &a, vector<long long> &x) {
  long long z = 0;
  long long n = 1;
  for (int i = 0; i < x.size(); ++i)
    n *= x[i];

  for (int i = 0; i < a.size(); ++i) {
    long long tmp = (a[i] * (n / x[i])) % n;
    tmp = (tmp * mod_inv(n / x[i], x[i])) % n;
    z = (z + tmp) % n;
}

  return (z + n) % n;
}</pre>
```

9.3 diophantine equations

```
long long gcd(long long a, long long b, long long &x, long long &y) {
  if (a == 0) {
    x = 0;
```

```
y = 1;
   return b;
 }
 long long x1, y1;
 long long d = gcd(b \% a, a, x1, y1);
 x = y1 - (b / a) * x1;
 y = x1;
 return d:
}
bool find_any_solution(long long a, long long b, long long c, long long
    &x0.
   long long &y0, long long &g) {
 g = gcd(abs(a), abs(b), x0, y0);
 if (c % g) {
   return false:
 x0 *= c / g;
 v0 *= c / g;
 if (a < 0) x0 = -x0;
 if (b < 0) y0 = -y0;
 return true;
}
void shift_solution(long long &x, long long &y, long long a, long long b,
   long long cnt) {
 x += cnt * b;
 y -= cnt * a;
}
long long find_all_solutions(long long a, long long b, long long c,
   long long minx, long long maxx, long long miny,
   long long maxy) {
 long long x, y, g;
 if (!find_any_solution(a, b, c, x, y, g)) return 0;
 a /= g;
 b /= g;
 long long sign_a = a > 0 ? +1 : -1;
 long long sign_b = b > 0 ? +1 : -1;
  shift_solution(x, y, a, b, (minx - x) / b);
 if (x < minx) shift_solution(x, y, a, b, sign_b);</pre>
 if (x > maxx) return 0;
```

```
long long lx1 = x;
shift_solution(x, y, a, b, (maxx - x) / b);
if (x > maxx) shift_solution(x, y, a, b, -sign_b);
long long rx1 = x;
shift_solution(x, y, a, b, -(miny - y) / a);
if (y < miny) shift_solution(x, y, a, b, -sign_a);</pre>
if (y > maxy) return 0;
long long 1x2 = x;
shift_solution(x, y, a, b, -(maxy - y) / a);
if (y > maxy) shift_solution(x, y, a, b, sign_a);
long long rx2 = x;
if (1x2 > rx2) swap(1x2, rx2);
long long lx = max(lx1, lx2);
long long rx = min(rx1, rx2);
if (1x > rx) return 0;
return (rx - lx) / abs(b) + 1;
```

9.4 discrete logarithm

```
// Computes x which a ^ x = b mod n.
long long d_log(long long a, long long b, long long n) {
  long long m = ceil(sqrt(n));
  long long aj = 1;
  map<long long, long long> M;
  for (int i = 0; i < m; ++i) {
    if (!M.count(aj))
        M[aj] = i;
    aj = (aj * a) % n;
}

long long coef = mod_pow(a, n - 2, n);
  coef = mod_pow(coef, m, n);
// coef = a ^ (-m)
  long long gamma = b;
  for (int i = 0; i < m; ++i) {</pre>
```

```
if (M.count(gamma)) {
    return i * m + M[gamma];
} else {
    gamma = (gamma * coef) % n;
}
return -1;
}
```

9.5 ext euclidean

```
void ext_euclid(long long a, long long b, long long &x, long long &y,
    long long &g) {
    x = 0, y = 1, g = b;
    long long m, n, q, r;
    for (long long u = 1, v = 0; a != 0; g = a, a = r) {
        q = g / a, r = g % a;
        m = x - u * q, n = y - v * q;
        x = u, y = v, u = m, v = n;
    }
}
```

9.6 highest exponent factorial

```
int highest_exponent(int p, const int &n){
  int ans = 0;
  int t = p;
  while(t <= n){
    ans += n/t;
    t*=p;
  }
  return ans;
}</pre>
```

9.7 miller rabin

```
const int rounds = 20;
// checks whether a is a witness that n is not prime, 1 < a < n</pre>
```

```
bool witness(long long a, long long n) {
 // check as in Miller Rabin Primality Test described
 long long u = n - 1;
  int t = 0;
  while (u % 2 == 0) {
   t++;
   u >>= 1;
 long long next = mod_pow(a, u, n);
  if (next == 1) return false;
  long long last;
  for (int i = 0; i < t; ++i) {</pre>
   last = next;
   next = mod_mul(last, last, n);
   if (next == 1) {
     return last != n - 1;
 }
 return next != 1;
// Checks if a number is prime with prob 1 - 1 / (2 ^it)
// D(miller_rabin(999999999999997LL) == 1);
// D(miller_rabin(99999999971LL) == 1);
// D(miller_rabin(7907) == 1);
bool miller_rabin(long long n, int it = rounds) {
 if (n <= 1) return false;</pre>
 if (n == 2) return true;
 if (n % 2 == 0) return false;
  for (int i = 0; i < it; ++i) {</pre>
   long long a = rand() \% (n - 1) + 1;
   if (witness(a, n)) {
     return false;
 }
 return true;
```

9.8 mod integer

```
template<class T, T mod>
struct mint_t {
```

```
T val;
mint_t() : val(0) {}
mint_t(T v) : val(v % mod) {}

mint_t operator + (const mint_t& o) const {
    return (val + o.val) % mod;
}

mint_t operator - (const mint_t& o) const {
    return (val - o.val) % mod;
}

mint_t operator * (const mint_t& o) const {
    return (val * o.val) % mod;
}

stypedef mint_t<long long, 998244353> mint;
```

9.9 mod inv

```
long long mod_inv(long long n, long long m) {
  long long x, y, gcd;
  ext_euclid(n, m, x, y, gcd);
  if (gcd != 1)
   return 0;
  return (x + m) % m;
}
```

9.10 mod mul

```
// Computes (a * b) % mod
long long mod_mul(long long a, long long b, long long mod) {
  long long x = 0, y = a % mod;
  while (b > 0) {
    if (b & 1)
        x = (x + y) % mod;
        y = (y * 2) % mod;
        b /= 2;
    }
  return x % mod;
}
```

9.11 mod pow

```
// Computes ( a ^ exp ) % mod.
long long mod_pow(long long a, long long exp, long long mod) {
  long long ans = 1;
  while (exp > 0) {
    if (exp & 1)
      ans = mod_mul(ans, a, mod);
    a = mod_mul(a, a, mod);
    exp >>= 1;
  }
  return ans;
}
```

9.12 number theoretic transform

```
typedef long long int LL;
typedef pair<LL, LL> PLL;
/* The following vector of pairs contains pairs (prime, generator)
 * where the prime has an Nth root of unity for N being a power of two.
 * The generator is a number g s.t g^(p-1)=1 \pmod{p}
 * but is different from 1 for all smaller powers */
vector<PLL> nth_roots_unity {
  {1224736769,330732430},{1711276033,927759239},{167772161,167489322},
  {469762049,343261969},{754974721,643797295},{1107296257,883865065}};
PLL ext_euclid(LL a, LL b) {
 if (b == 0)
   return make_pair(1,0);
 pair<LL,LL> rc = ext_euclid(b, a % b);
 return make_pair(rc.second, rc.first - (a / b) * rc.second);
//returns -1 if there is no unique modular inverse
LL mod_inv(LL x, LL modulo) {
 PLL p = ext_euclid(x, modulo);
 if ( (p.first * x + p.second * modulo) != 1 )
   return -1:
 return (p.first+modulo) % modulo;
}
```

```
//Number theory fft. The size of a must be a power of 2
void ntfft(vector<LL> &a, int dir, const PLL &root_unity) {
 int n = a.size():
 LL prime = root_unity.first;
 LL basew = mod_pow(root_unity.second, (prime-1) / n, prime);
 if (dir < 0) basew = mod_inv(basew, prime);</pre>
 for (int m = n; m >= 2; m >>= 1) {
    int mh = m >> 1:
   LL w = 1:
    for (int i = 0; i < mh; i++) {</pre>
     for (int j = i; j < n; j += m) {</pre>
       int k = j + mh;
       LL x = (a[j] - a[k] + prime) \% prime;
       a[j] = (a[j] + a[k]) \% prime;
       a[k] = (w * x) \% prime;
     w = (w * basew) % prime;
    basew = (basew * basew) % prime;
 int i = 0;
 for (int j = 1; j < n - 1; j++) {
   for (int k = n >> 1; k > (i ^= k); k >>= 1);
   if (j < i) swap(a[i], a[j]);</pre>
 }
}
```

9.13 pollard rho factorize

```
long long pollard_rho(long long n) {
  long long x, y, i = 1, k = 2, d;
  x = y = rand() % n;
  while (1) {
    ++i;
    x = mod_mul(x, x, n);
    x += 2;
    if (x >= n) x -= n;
    if (x == y) return 1;
    d = __gcd(abs(x - y), n);
    if (d != 1) return d;
    if (i == k) {
        y = x;
        k *= 2;
    }
}
```

```
}
 return 1;
// Returns a list with the prime divisors of n
vector<long long> factorize(long long n) {
 vector<long long> ans;
 if (n == 1)
   return ans:
 if (miller_rabin(n)) {
   ans.push_back(n);
 } else {
   long long d = 1;
   while (d == 1)
     d = pollard_rho(n);
   vector<long long> dd = factorize(d);
   ans = factorize(n / d);
   for (int i = 0; i < dd.size(); ++i)</pre>
     ans.push_back(dd[i]);
 }
 return ans;
```

9.14 primes

```
// Finds prime numbers between a and b, using basic primes up to sqrt(b)
 // a must be greater than 1.
  vector<long long> seg_sieve(long long a, long long b) {
   long long ant = a;
   a = max(a, 3LL);
   vector<bool> pmap(b - a + 1);
   long long sqrt_b = sqrt(b);
   for (int i = 0; i < num_p; ++i) {</pre>
     long long p = primes[i];
     if (p > sqrt_b) break;
     long long j = (a + p - 1) / p;
     for (long long v = (j == 1) ? p + p : j * p; v <= b; v += p) {
       pmap[v - a] = true;
     }
   }
   vector<long long> ans;
   if (ant == 2) ans.push_back(2);
   int start = a % 2 ? 0 : 1;
   for (int i = start, I = b - a + 1; i < I; i += 2)</pre>
     if (pmap[i] == false)
       ans.push_back(a + i);
   return ans;
 }
 vector<pair<int, int>> factor(int n) {
   vector<pair<int, int>> ans;
   if (n == 0) return ans;
   for (int i = 0; primes[i] * primes[i] <= n; ++i) {</pre>
     if ((n % primes[i]) == 0) {
       int expo = 0;
       while ((n % primes[i]) == 0) {
         expo++;
         n /= primes[i];
       ans.emplace_back(primes[i], expo);
   }
   if (n > 1) {
     ans.emplace_back(n, 1);
   }
   return ans;
 }
}
```

9.15 totient sieve

```
for (int i = 1; i < MN; i++)
  phi[i] = i;

for (int i = 1; i < MN; i++)
  if (!sieve[i]) // is prime
  for (int j = i; j < MN; j += i)
    phi[j] -= phi[j] / i;</pre>
```

9.16 totient

```
long long totient(long long n) {
   if (n == 1) return 0;
   long long ans = n;
   for (int i = 0; primes[i] * primes[i] <= n; ++i) {
      if ((n % primes[i]) == 0) {
        while ((n % primes[i]) == 0) n /= primes[i];
        ans -= ans / primes[i];
    }
   if (n > 1) {
      ans -= ans / n;
   }
   return ans;
}
```

10 Strings

10.1 Incremental Aho Corasick

```
class IncrementalAhoCorasic {
  static const int Alphabets = 26;
  static const int AlphabetBase = 'a';
  struct Node {
   Node *fail;
   Node *next[Alphabets];
   int sum;
   Node() : fail(NULL), next{}, sum(0) { }
};
```

```
struct String {
   string str;
   int sign;
 };
public:
 //totalLen = sum of (len + 1)
 void init(int totalLen) {
   nodes.resize(totalLen);
   nNodes = 0:
   strings.clear();
   roots.clear();
   sizes.clear();
   que.resize(totalLen);
 }
 void insert(const string &str, int sign) {
   strings.push_back(String{ str, sign });
   roots.push_back(nodes.data() + nNodes);
   sizes.push_back(1);
   nNodes += (int)str.size() + 1;
   auto check = [&]() { return sizes.size() > 1 && sizes.end()[-1] ==
        sizes.end()[-2]; };
   if(!check())
     makePMA(strings.end() - 1, strings.end(), roots.back(), que);
   while(check()) {
     int m = sizes.back();
     roots.pop_back();
     sizes.pop_back();
     sizes.back() += m;
     if(!check())
       makePMA(strings.end() - m * 2, strings.end(), roots.back(), que);
   }
 }
 int match(const string &str) const {
   int res = 0;
   for(const Node *t : roots)
     res += matchPMA(t, str);
   return res;
 }
private:
```

```
static void makePMA(vector<String>::const_iterator begin,
    vector<String>::const_iterator end, Node *nodes, vector<Node*>
    &que) {
 int nNodes = 0;
 Node *root = new(&nodes[nNodes ++]) Node();
 for(auto it = begin; it != end; ++ it) {
   Node *t = root;
   for(char c : it->str) {
     Node *&n = t->next[c - AlphabetBase];
     if(n == nullptr)
      n = new(&nodes[nNodes ++]) Node():
     t = n;
   t->sum += it->sign;
 int qt = 0;
 for(Node *&n : root->next) {
   if(n != nullptr) {
     n->fail = root:
     que[qt ++] = n;
   } else {
     n = root;
   }
 }
 for(int qh = 0; qh != qt; ++ qh) {
   Node *t = que[qh];
   int a = 0:
   for(Node *n : t->next) {
     if(n != nullptr) {
       que[qt ++] = n;
       Node *r = t->fail;
       while(r->next[a] == nullptr)
        r = r->fail;
      n-fail = r-next[a]:
       n->sum += r->next[a]->sum;
     }
static int matchPMA(const Node *t, const string &str) {
 int res = 0;
 for(char c : str) {
   int a = c - AlphabetBase;
```

```
while(t->next[a] == nullptr)
       t = t->fail;
     t = t- next[a]:
     res += t->sum;
   return res;
 }
 vector<Node> nodes;
 int nNodes:
 vector<String> strings;
 vector<Node*> roots;
 vector<int> sizes;
 vector<Node*> que;
};
int main() {
 int m:
 while(~scanf("%d", &m)) {
   IncrementalAhoCorasic iac;
   iac.init(600000);
   rep(i, m) {
     int ty;
     char s[300001];
     scanf("%d%s", &ty, s);
     if(ty == 1) {
       iac.insert(s, +1);
     } else if(ty == 2) {
       iac.insert(s, -1);
     } else if(ty == 3) {
       int ans = iac.match(s);
       printf("%d\n", ans);
       fflush(stdout);
     } else {
       abort();
   }
 }
 return 0;
```

10.2 minimal string rotation

```
// Lexicographically minimal string rotation
int lmsr() {
 string s;
 cin >> s;
 int n = s.size();
 s += s;
 vector<int> f(s.size(), -1);
 int k = 0;
 for (int j = 1; j < 2 * n; ++j) {
   int i = f[j - k - 1];
   while (i != -1 && s[j] != s[k + i + 1]) {
     if (s[i] < s[k + i + 1])
      k = j - i - 1;
     i = f[i];
   if (i == -1 \&\& s[j] != s[k + i + 1]) {
     if (s[i] < s[k + i + 1]) {
       k = j;
     f[j - k] = -1;
   } else {
     f[j - k] = i + 1;
 }
 return k;
```

10.3 suffix array

```
/**
 * 0 (n log^2 (n))
 * See http://web.stanford.edu/class/cs97si/suffix-array.pdf for reference
 * */

struct entry{
  int a, b, p;
  entry(){}
  entry(int x, int y, int z): a(x), b(y), p(z){}
  bool operator < (const entry &o) const {
    return (a == o.a) ? (b == o.b) ? ( p < o.p) : (b < o.b) : (a < o.a);
  }
};</pre>
```

```
struct SuffixArray{
 const int N;
 string s;
 vector<vector<int> > P;
 vector<entry> M;
 SuffixArray(const string &s) : N(s.length()), s(s), P(1, vector<int>
      (N, O)), M(N) {
   for (int i = 0; i < N; ++i)</pre>
     P[0][i] = (int) s[i];
   for (int skip = 1, level = 1; skip < N; skip *= 2, level++) {</pre>
     P.push_back(vector<int>(N, 0));
     for (int i = 0 ; i < N; ++i) {</pre>
       int next = ((i + skip) < N) ? P[level - 1][i + skip] : -10000;
       M[i] = entry(P[level - 1][i], next, i);
     sort(M.begin(), M.end());
     for (int i = 0; i < N; ++i)</pre>
       P[level][M[i].p] = (i > 0 \text{ and } M[i].a == M[i - 1].a \text{ and } M[i].b ==
           M[i - 1].b) ? P[level][M[i - 1].p] : i;
   }
 }
 vector<int> getSuffixArray(){
   vector<int> &rank = P.back();
   vector<pair<int, int> > inv(rank.size());
   for (int i = 0; i < rank.size(); ++i)</pre>
     inv[i] = make_pair(rank[i], i);
   sort(inv.begin(), inv.end());
   vector<int> sa(rank.size());
   for (int i = 0; i < rank.size(); ++i)</pre>
     sa[i] = inv[i].second;
   return sa;
 }
 // returns the length of the longest common prefix of s[i...L-1] and
      s[j...L-1]
 int lcp(int i, int j) {
   int len = 0;
   if (i == j) return N - i;
   for (int k = P.size() - 1; k \ge 0 && i < N && j < N; --k) {
     if (P[k][i] == P[k][j]) {
       i += 1 << k;
```

```
j += 1 << k;
len += 1 << k;
}

return len;
}
</pre>
```

10.4 suffix automaton

```
/*
 * Suffix automaton:
 * This implementation was extended to maintain (online) the
 * number of different substrings. This is equivalent to compute
 * the number of paths from the initial state to all the other
 * states.
 * The overall complexity is O(n)
 * can be tested here:
     https://www.urionlinejudge.com.br/judge/en/problems/view/1530
 * */
struct state {
 int len, link;
 long long num_paths;
 map<int, int> next;
};
const int MN = 200011;
state sa[MN << 1];</pre>
int sz, last;
long long tot_paths;
void sa init() {
  sz = 1;
 last = 0;
  sa[0].len = 0;
  sa[0].link = -1;
  sa[0].next.clear();
  sa[0].num_paths = 1;
  tot_paths = 0;
}
```

```
void sa_extend(int c) {
 int cur = sz++;
 sa[cur].len = sa[last].len + 1;
 sa[cur].next.clear();
 sa[cur].num_paths = 0;
 int p;
 for (p = last; p != -1 && !sa[p].next.count(c); p = sa[p].link) {
   sa[p].next[c] = cur;
   sa[cur].num_paths += sa[p].num_paths;
   tot_paths += sa[p].num_paths;
 if (p == -1) {
   sa[cur].link = 0;
 } else {
   int q = sa[p].next[c];
   if (sa[p].len + 1 == sa[q].len) {
     sa[cur].link = q;
   } else {
     int clone = sz++;
     sa[clone].len = sa[p].len + 1;
     sa[clone].next = sa[q].next;
     sa[clone].num_paths = 0;
     sa[clone].link = sa[q].link;
     for (; p!= -1 && sa[p].next[c] == q; p = sa[p].link) {
       sa[p].next[c] = clone;
       sa[q].num_paths -= sa[p].num_paths;
       sa[clone].num_paths += sa[p].num_paths;
     sa[q].link = sa[cur].link = clone;
 }
 last = cur;
```

10.5 z algorithm

```
using namespace std;
#include<bits/stdc++.h>
```

```
vector<int> compute_z(const string &s){
  int n = s.size();
  vector<int> z(n,0);
  int 1,r;
 r = 1 = 0;
  for(int i = 1; i < n; ++i){</pre>
   if(i > r) {
     l = r = i:
     while (r < n \text{ and } s[r - 1] == s[r])r++;
     z[i] = r - 1:r--:
   }else{
      int k = i-1;
      if(z[k] < r - i +1) z[i] = z[k];
      else {
       1 = i:
       while (r < n \text{ and } s[r - 1] == s[r])r++;
       z[i] = r - 1;r--;
   }
 return z;
}
int main(){
  //string line;cin>>line;
  string line = "alfalfa";
  vector<int> z = compute_z(line);
  for(int i = 0; i < z.size(); ++i ){</pre>
   if(i)cout<<" ";</pre>
   cout<<z[i];
  cout << endl;
  // must print "0 0 0 4 0 0 1"
 return 0;
```