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- Module Agreement -
EXTENDS Naturals, FiniteSets, Commons
CONSTANT NPROCESSES
CONSTANT NMESSAGES
CONSTANT CONFLICTR(_, _)
Since this algorithm is for failure-free environments, the set of all processes is the same as the
correct ones.
LOCAL Processes \triangleq \{i : i \in 1 ... NPROCESSES\}
LOCAL ChooseProcess \stackrel{\Delta}{=} CHOOSE x \in Processes : TRUE
LOCAL Create(id) \triangleq [id \mapsto id, d \mapsto Processes, o \mapsto ChooseProcess]
LOCAL AllMessages \stackrel{\triangle}{=} \{Create(id) : id \in 1 .. NMESSAGES\}
VARIABLES
    K,
    Pending,
    Delivering,
    Delivered,
    PreviousMsqs,
    Votes,
    QuasiReliable Channel
Initialize the instance for the Generic Multicast 0. The INITIAL_MESSAGES is a set with
NMESSAGES, unordered, a tuple with the starting state S0 and the message.
Algorithm \stackrel{\Delta}{=} INSTANCE Generic Multicast 0 WITH
    INITIAL\_MESSAGES \leftarrow \{\langle \text{``SO''}, m \rangle : m \in AllMessages \}
 Weak fairness is necessary.
Spec \triangleq Algorithm! SpecFair
If a correct process deliver a message m, then all correct processes in m.d eventually delivers m.
We verify that all messages in AllMessages, for all the processes that delivered a message, even-
tually, all the correct members in the destination will deliver.
Agreement \triangleq
    \forall m \in AllMessages:
       \forall p \in Processes:
          Algorithm! WasDelivered(p, m)
               \rightsquigarrow \forall q \in m.d:
                    q \in Processes \land Algorithm! WasDelivered(q, m)
```