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MODULE *PartialOrder*

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EXTENDS *Naturals, FiniteSets, Commons*

CONSTANT *NGROUPS, NPROCESSES, NMESSAGES, CONFLICTR*(-, -)

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This algorithm works in an environment with crash-stop failures, but we do not model processes failing. The set of all processes contains all correct ones.

LOCAL *Processes*  $\triangleq 1 \dots NPROCESSES$

LOCAL *Groups*  $\triangleq 1 \dots NGROUPS$

LOCAL *ProcessesInGroup*  $\triangleq [g \in Groups \mapsto Processes]$

LOCAL *AllMessages*  $\triangleq CreateMessages(NMESSAGES, Groups, Processes)$

LOCAL *MessagesCombinations*  $\triangleq CreatePossibleMessages(AllMessages)$

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VARIABLES *K, PreviousMsgs, Delivered, Votes, MemoryBuffer, QuasiReliableChannel, AtomicBroadcastBuffer*

Initialize the instance for the Generic Multicast 2. The *INITIAL\_MESSAGES* is a sequence, partially ordered. The sequence elements are sets of messages, messages that commute can share a set.

*Algorithm*  $\triangleq$  INSTANCE *GenericMulticast2* WITH

*INITIAL\_MESSAGES*  $\leftarrow [g \in Groups \mapsto$

*PartiallyOrdered*(

*MessagesCombinations*[(*g*%*NMESSAGES*) + 1], *CONFLICTR*)]

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*Spec*  $\triangleq$  *Algorithm*!*Spec*

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LOCAL *BothDelivered*(*g*, *p1*, *p2*, *m1*, *m2*)  $\triangleq$

$\wedge$  *Algorithm*!*WasDelivered*(*g*, *p1*, *m1*)  $\wedge$  *Algorithm*!*WasDelivered*(*g*, *p1*, *m2*)

$\wedge$  *Algorithm*!*WasDelivered*(*g*, *p2*, *m1*)  $\wedge$  *Algorithm*!*WasDelivered*(*g*, *p2*, *m2*)

LOCAL *LHS*(*g*, *p1*, *p2*, *m1*, *m2*)  $\triangleq$

$\wedge \{p1, p2\} \subseteq (m1.d \cap m2.d)$

$\wedge CONFLICTR(m1, m2)$

$\wedge BothDelivered(g, p1, p2, m1, m2)$

LOCAL *RHS*(*g*, *p1*, *p2*, *m1*, *m2*)  $\triangleq$

(*Algorithm*!*DeliveredInstant*(*g*, *p1*, *m1*) <

*Algorithm*!*DeliveredInstant*(*g*, *p1*, *m2*))

$\equiv$  (*Algorithm*!*DeliveredInstant*(*g*, *p2*, *m1*) <

*Algorithm*!*DeliveredInstant*(*g*, *p2*, *m2*))

For every two messages, if they conflict, given a pair of processes, they are in the messages' destination, then both must deliver in the same order.

*PartialOrder*  $\triangleq$

$\square \forall g \in Groups :$

$\forall p1, p2 \in ProcessesInGroup[g] :$

$\forall m1, m2 \in AllMessages :$

$LHS(g, p1, p2, m1, m2) \Rightarrow RHS(g, p1, p2, m1, m2)$

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