

# 1. Levels and goals

## II. From Probability to Navigation

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# Objectives

- Understanding how probabilities can be used to model different game properties:
  - ① Player's skills, **PvP**: **Simplified Tennis (PONG)**
  - ② Player's chances at advancing the game: **Snakes and Ladders**.
- Using stochastic models for establishing the interconnection between probabilities and time:
  - ① Which is the probability of ending/winning the game in  $n$  turns?
  - ② Which is the average number of turns required for ending the game?
- Understanding how changes in the model and/or its probability distribution alters the previous questions, and their effect on gameplay.

# Game Modelling Goals

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Pacing is a fundamental analysis in **Game Balancing**:

- The aim of the game is to put obstacles and power-ups with trade-offs.
- Obstacles should never be insurmountable or unfair.

Statistical Analysis is often used to analyse games for balancing:

- We can identify faults (e.g., unbalanced areas) and make corrections
- The usual approach is to collect *logs* from real plays, and then infer statistics via *post-mortem analysis*.
- Still, is it possible to analyse the game play before sending the product to testers and/or sending it to production?

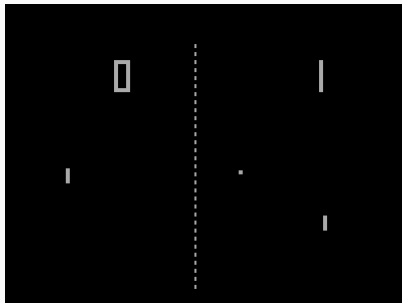
# Chances vs. Skills

- Chance allows a weaker player to beat a stronger one
- The outcome of the game should be mainly influenced by the skills
- Still, the skills of a player are hard to assess correctly.
- As a first approximation, we can model the player's skills as the chances of winning:
  - In a simplistic game, a player could only perform two actions: one, advancing the game towards the next turn, and the other making the player fail thus letting the opponent win.

# Player versus Player

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# Simplified Tennis (PONG) (1/3)



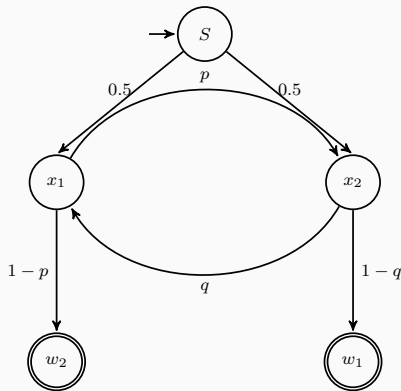
- One single graph represents one single (tennis) play (i.e., *round*).
- One of the two players, the *server*, starts the game (2 possible initial states). This player is randomly picked ( $S$ ).
- We have 2 states identifying one of the two players actively playing ( $x_1$  and  $x_2$ ), and 2 states ( $w_1$  and  $w_2$ ) declaring the winning for the given tennis game (e.g., *round*).
- The transition probabilities ( $p, q$ ) represent the player's expertise in performing the correct action

## Simplified Tennis (PONG) (2/3)

- The game is *memoryless*: the ability of responding to a service is completely unrelated to the previous move.
- We can represent such a game via a probabilistic process, namely a **Discrete Time Markov Chain (DTMC)**, represented as a transition matrix  $T$ .
- Such matrix can be graphically represented as a weighted directed graph.



# Simplified Tennis (PONG) (3/3)



$$T = \begin{matrix} & \begin{matrix} S & x_1 & x_2 & w_1 & w_2 \end{matrix} \\ \begin{matrix} S \\ x_1 \\ x_2 \\ w_1 \\ w_2 \end{matrix} & \begin{bmatrix} 0 & 0.5 & 0.5 & 0 & 0 \\ 0 & 0 & p & 0 & 1-p \\ 0 & q & 0 & 1-q & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \end{bmatrix} \end{matrix}$$

■ **State Space:** Determining the server ( $S$ ), Player1's turn ( $x_1$ ), Player2's turn ( $x_2$ ), Player 1 wins ( $w_1$ ), Player 2 wins ( $w_2$ ).

■ **Initial State:**  $S$ .

■ **Actions & Transitions:**

- ①  $P1$  serves,  $S \xrightarrow{0.5} x_1$
- ②  $P2$  serves,  $S \xrightarrow{0.5} x_2$
- ③  $P1$  responds,  $x_1 \xrightarrow{p} x_2$
- ④  $P2$  responds,  $x_2 \xrightarrow{q} x_1$
- ⑤  $P1$  does not catch,  $x_1 \xrightarrow{1-p} w_2$
- ⑥  $P2$  does not catch,  $x_2 \xrightarrow{1-q} w_1$

■ **Goal Test:** Either  $w_1$  or  $w_2$  are reached.

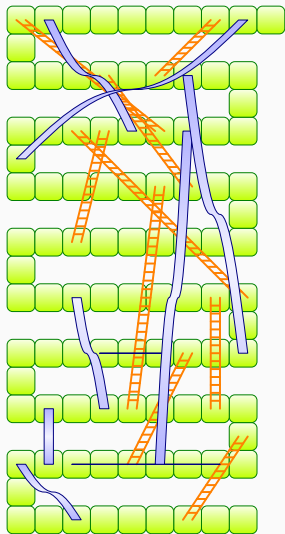
# Player versus Environment

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# Player versus Environment (PvE)

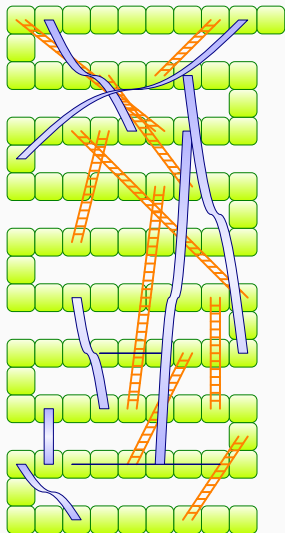
- They are also referred as *One-person games*, or *Games against nature*.
- Before tackling the problem of an interaction with another agent, we should first ask ourselves whether we are able to model agent that is capable of taking **meaningful decisions**:
  - No decision is with no effect,
  - but, one among them, will lead to the optimal solution.
- This person may take the wrong decision, but there does not exist a conscious opponent.
- One-person games are greatly investigated in AI (robotics) and in Planning Theory.

# Snakes and Ladders (1/3)



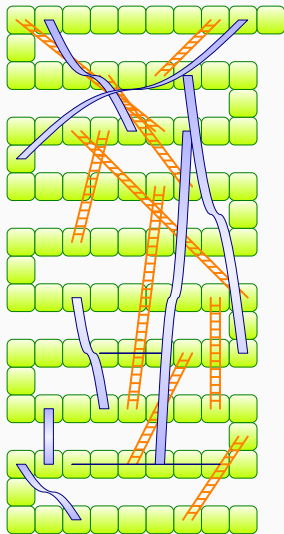
- The game contains 100 cells:
- The player starts the game from the first cell.
- **Actions:**
  - ① At each turn, the player rolls a die:
  - ② if they step into the lower end of a ladder, the player progresses towards its upper end;
  - ③ if they step into the upper end of a snake, the player backtracks towards its lower end;
  - ④ otherwise, the player stays in the reached state.
- The game ends when the last cell is reached.

## Snakes and Ladders (2/3)



- The game is *memoryless*: at a given point in the game, the player's progression from the current square is independent of how they arrived at that square.
- Each edge weight is now independent from each player's skills.

# Snakes and Ladders (3/3)



Without Snakes and Ladders:

- Each non-rigged die generates numbers in  $I = \{ 1, 2, 3, 4, 5, 6 \}$  with probability  $\mathbb{P}(I = i) = 1/6$  for each  $i \in I$ .
- Given a initial cell  $i$ , I can only move from  $i$  to one of the following states:  $\{ i + 1, \dots, i + 6 \} \setminus \{ n \in \mathbb{N} \mid n > 100 \}$ .

With Snakes and Ladders:

- arriving to a starting point of a snake/ladder from cell  $i$  has 0 probability,
- while the probability of reaching the end of the snake/ladder from cell  $i$  is increased by  $1/6$ .

## **Measuring easiness through probability and time**

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# Average Hitting Time

The *average hitting time* determines the mean number of turns  $x_i$  required to reach one of the goals for the first time from state  $i$ . We can prove that this reduces to solve the following *linear (equation) system* for all the states  $i$ :

$$\begin{cases} x_i = 0 & i \text{ goal} \\ -x_i + \sum_{j \neq i} T_{ij} x_j = -1 & \text{otherwise} \end{cases}$$

This reduces to solve a linear system  $A\vec{x} = \vec{b}$ , where:

$$\vec{b}_i = \begin{cases} 0 & i \text{ goal} \\ -1 & \text{otherwise} \end{cases} \quad A_{ij} = \begin{cases} 1 & i = j, i \text{ goal} \\ -1 & i = j, i \text{ not goal} \\ 0 & i \neq j, i \text{ goal} \\ T_{ij} & \text{otherwise} \end{cases}$$



## Average Hitting Time: Snakes and Ladders (1/2)

Given that our problem has only one initial state,  $\text{cell}=1$ , we are interested in  $x_1$ .

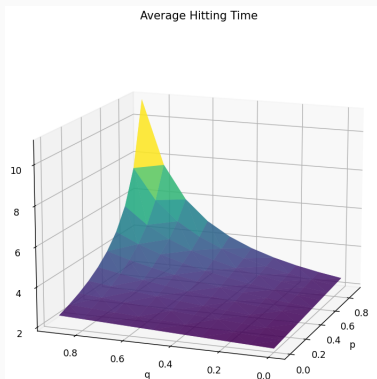
The following result confirm the intuitive results in length of game-play from performing *random walks* over the stochastic process:

- Without Snakes and Ladders:  $\sim 29$  turns.
- With only Ladders:  $\sim 18$  turns.
- With only Snakes:  $\sim 46$  turns.
- With both Snakes and Ladders:  $\sim 25$  turns.

## Average Hitting Time: PONG

By solving the system, we obtain the following *Average Hitting Time*:

$$x_s = \frac{4 + p + q - 2pq}{2 - 2pq}$$



We now exploit *average hitting time* for determining how long one single round of the PONG game will take dependently from the success probabilities  $p$  and  $q$  for both players.

- If only one player is strong, it will take as little as  $\sim 2$  turns after tossing the coin.
- The more experienced the two players are, the longer it will take!

# Probability of winning in $n$ turns

Initial Probability Distribution  $\mu$ :

- Given that in our game we have only one initial state,  $S$ , we will have an **initial probability distribution**  $\mu$  as an empty vector except for  $\mu_S = 1$ .

Accepting States/Goal Tests:

- Given that the two states declaring the winner,  $w_1$  and  $w_2$ , have no outgoing edges and given that all the intermediate states can reach such a state, we want to assess the probability to reach such states.

Reachability Probability at step  $n$ :

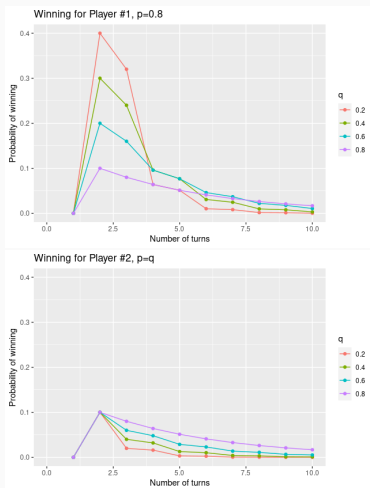
- $\mu T$  returns a vector  $\mu^{(1)}$ , where  $\mu_i^{(1)}$  indicates the probability of reaching a state  $i$  in one step.
- The probability of Player1 winning in 1 turn is  $\mu_{w_1}^{(1)} = 0$ .
- The probability of reaching any game configuration in two steps is  $\mu^{(2)} = \mu^{(1)}T = \mu TT = \mu T^2$ :
- so, the probability of Player1 winning in  $n$  turns is given by  $[\mu T^n]_{w_1}$

## Probability of winning in $n$ turns: PONG (1/2)

In the previous slide we discussed how long should it take to end the game when  $p$  and  $q$  vary:

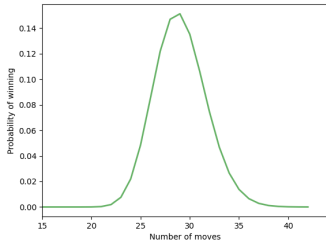
- This gives no information of the chances of winning of one player in a given amount of turns.
- To simplify the analysis, we assume that Player1 is an expert,  $p = 0.8$ , while we test varying opponent (Player2) skills.
- Furthermore, we want to analyse the probability of Player1 and Player2 of winning after  $n$  turns in the game:
- This can be assessed by assessing the Reachability Probability at step  $n$ .

# Probability of winning in $n$ turns: PONG (2/2)

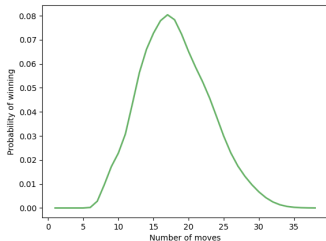


- We need to first toss the coin before assessing a winning probability.
- The two players have higher chances of winning at the first stages of the game.
- When both opponents have the same skills ( $q = 0.8$ ), they both have the same distribution of chances of winning.
- The lower the skills for Player2, the higher are the changes of winning for Player1.

# Snakes and Ladders: Probability of winning in $n$ turns (1/2)



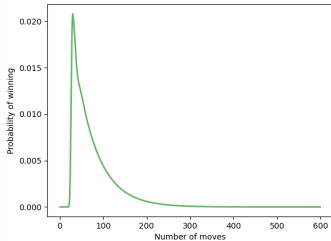
*Without Snakes and Ladders.*



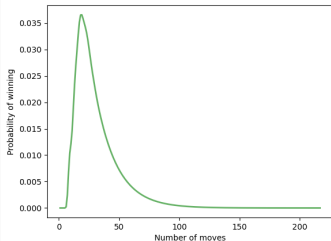
*With Ladders.*

- We need at least 17 turns to win the game.
- It is more likely to win the game with probability 0.15 after 29 turns.
- The game fastens up, we now need at least 6 turns to win the game.

# Snakes and Ladders: Probability of winning in $n$ turns (2/2)



*With Snakes.*



*With Snakes and Ladders.*

- The game slows down, as we might get stuck in a loop.
  - Furthermore, the overall chances of winning decreases.
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- Adding back the ladders, the chances of winning increase as well as the length of loops decrease.

## Wrap Up: Snakes and Ladders

This simple game allows us to determine how **power-ups** (*ladders*) and **penalties** (*snakes*) might affect the duration and the probability of reaching the end of the game in a given amount of steps:

- By only adding ladders, the time for the game-play is considerably reduced:
  - ① We generate fewer and shorter paths, but there might be multiple possible way to get paths of the same length:
  - ② The probability of reaching the final state in a given amount of steps is reduced.
- By only adding snakes, the time of the game-play is considerably increased:
  - ① Even if it is extremely improbable, players might get stuck in loops, thus increasing the length of the possible paths;
  - ② Therefore, the probability of reaching the end of the game in a given amount of steps if considerably decreased.
- Intuitively, adding both snakes and ladder will provide a trade-off situation between possible length of game-play and probability of reaching the end of the game.



## Lessons Learned

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# Creating a game: Balancing Process

While assessing the game balance, we should proceed as follows:

- ❶ First, we should create a balanced game:
  - *E.g., we should generate a board with neither snakes nor ladders.*
- ❷ Next, we might consider the average running time (*average hitting time*) and the probability of winning the game in a given number of steps.
  - *Collect such values as reasonable outcomes for the game.*
- ❸ Introduce minimal changes, so to always determine how single changes affect the overall system.
- ❹ Compare the values obtained with the previous configuration, which can be thought as reasonable if returns similar or better values than the previous configuration.
- ❺ If the game configuration is reasonable, go back to the 3<sup>rd</sup> item and continue to refine the game.

## More complex games

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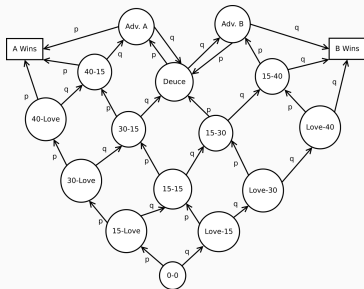
# Monopoly

The proposed methodology can be also applied to board games such as **Monopoly**:

- The real Monopoly is not a stochastic model:
  - if a pawn is in jail, the person in control might choose whether to stay in jail if doubles are not thrown.
  - if a pawn has been in jail two turns, it must be moved out of jail.
- This is because the current models are *memoryless*. There are two possible directions to be taken:
  - Force the process to be memoryless, and change the jail rule.
  - Simulate the notion of memory by exploding the number of possible states.
- Still, simplification of the model will give you a first intuition of how the real game might actually work.



# Semi-Realistic Tennis



- Each PONG process could be one single nested process within the diagram on the left.
- Please observe: from theory, nested process are equivalent to un-nested ones, so:
  - 1 Replace each circle state  $c$  from the left with a whole PONG game instance, where
  - 2 the ingoing edges to  $w_1$  (and  $w_2$ ) state have now the same target of  $c \xrightarrow{p} c'$  (and  $x \xrightarrow{q} c''$ ), and
  - 3 all the ingoing edges to a circle are redirected to each node  $S$ .
  - 4 Remove the unconnected  $w_1$  and  $w_2$ .
- **Exercise:** compare the results of PONG with the ones with Semi-Realistic Tennis.

This section introduced the key concepts to understand the relevance of determining the *average hitting times* and the *probability of terminating a game in a given amount of steps*. We also showed how changes in probability affect the overall game duration.