

Structures Presenting / Enabling Gaussian Concentration

Part 2!

Theorem: For all f of degree d ,
 $\forall \varepsilon, C, M$, there is $\tilde{f}: \mathbb{R}^N \rightarrow \mathbb{R}$
 of degree d s.t.

$$\|f - \tilde{f}\|_{L_Y} \lesssim_{C, M} \varepsilon^M \|f\|_{L_Y}$$

and \tilde{f} has a $(\varepsilon, \varepsilon^{-C})$ diffuse decomposition.

Theorem: If f is multilinear and has
 a (T, C, m, ε) decomposition
 then

$$|\mathbb{E}[\text{sgn}(f(x))] - \mathbb{E}[\text{sgn}(f(A))]|$$

$$\lesssim T^{1/5} C + \varepsilon^{1/d}$$

$$\bullet f: \mathbb{R}^N \rightarrow \mathbb{R}$$

$$\bullet f = h(q_1, \dots, q_m)$$

$$h: \mathbb{R}^m \rightarrow \mathbb{R}$$

$$q_i: \mathbb{R}^N \rightarrow \mathbb{R}$$

$$Q = (q_1, \dots, q_m)$$

$\bullet (\varepsilon, C)$ diffuse if

$$P(|Q(x) - a| \leq \varepsilon) \leq C \varepsilon^m$$

$\bullet f$ has a (T, C, m, ε)
 decomposition if there is
 \tilde{f} s.t.

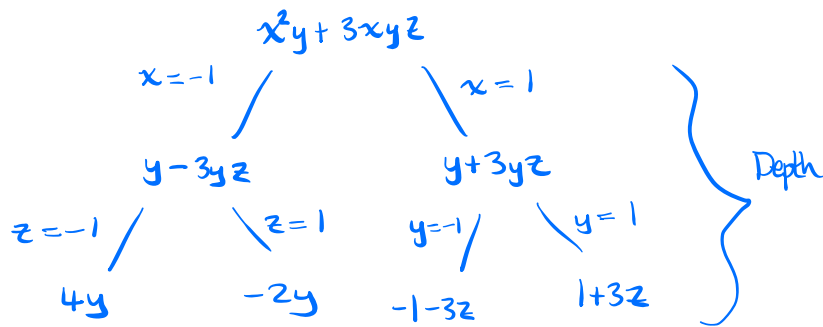
$$\|f - \tilde{f}\|_{L_Y} \leq \varepsilon^2 W(f(x))$$

and \tilde{f} has a $(T^{1/5}, C)$
 diffuse decomposition

(h, q_1, \dots, q_m) with

$$\inf_i (q_i) \lesssim T$$

Decision Trees



Theorem If f is a degree- d multilinear polynomial, then there is a decision tree of depth $O(\tau^{-1} \log(1/\tau)^{O(d)})$ s.t. all but a fraction τ of the leaf functions f_u either have $\mathbb{V}(f_u) < \tau^M \|f_u\|_{L_2}^2$ or f_u has a $(\tau, \tau^c, O(1), O(\tau^M))$.

$$\text{So } |\mathbb{E}[\text{sgn} f(x)] - \mathbb{E}[\text{sgn} f_u(A)]| \leq \tau^{1/5-c} + \tau^{M/d}$$

Invariance

$$\begin{aligned} \mathbb{E}[\text{sgn}(f(A))] &\approx \mathbb{E}[\text{sgn}(\hat{f}(A))] \\ &= \mathbb{E}[\text{sgn}(h(q_1(A), \dots, q_m(A)))] \\ &\approx \mathbb{E}(p(q_1(A), \dots, q_m(A))) \\ &\approx \mathbb{E}(p(q_1(x), \dots, q_m(x))) \\ &\approx \mathbb{E}(\text{sgn}(h(q_1(x), \dots, q_m(x)))) \\ &= \mathbb{E}[\text{sgn}(\hat{f}(x))] \end{aligned}$$

$$\approx \mathbb{E}[\text{sgn}(f(x))]$$

How do we define $p: \mathbb{R}^m \rightarrow \mathbb{R}$?

Theorem: If (h, q_1, \dots, q_m) is (ε, C) diffuse then there is $p: \mathbb{R}^m \rightarrow \mathbb{R}$ s.t.

- $p(q_1, \dots, q_m) \geq \text{sgn}(h(q_1, \dots, q_m))$
- $\mathbb{E}[p(q_1(x), \dots, q_m(x))] - \mathbb{E}[\text{sgn}(h(q_1, \dots, q_m))] \leq C\varepsilon \log(1/\varepsilon)^{d_m/2+1}$
- $\|D^k p\|_{\infty} \lesssim \varepsilon^{-k}$

Construction Let

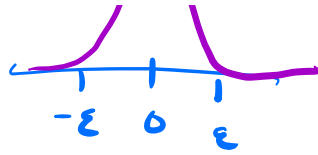
$$g(x) = \begin{cases} 1 & : \exists y: |x-y| < \varepsilon \text{ and } h(y) \geq 0 \\ -1 & : \text{otherwise} \end{cases}$$

Then $g \geq \text{sgn}(h(q_1, \dots, q_m))$

even after mollifying in an ε neighborhood

Let $P = g * \eta_\varepsilon$





To show

$$\mathbb{E}[p(q_1(x), \dots, q_m(x))] - \mathbb{E}[\text{sgn}(h(q_1^{(x)}, \dots, q_m^{(x)}))] \lesssim C \varepsilon \log(1/\varepsilon)^{dm/2+1}$$

$p(q_1, \dots, q_m) = \text{sgn}(h(q_1, \dots, q_m))$
 except if $d(Q(x), \bar{h}'(0)) \leq 2\varepsilon$

Bound this by covering
 argument + diffuseness

Algorithm for Thm 1

Theorem: For all f of degree d ,
 $\forall \varepsilon, c, M$, there is $\tilde{f}: \mathbb{R}^n \rightarrow \mathbb{R}$
of degree d s.t.

$$\|f - \tilde{f}\|_{L_Y} \lesssim_{c, M} \varepsilon^M \|f\|_{L_Y}$$

and \tilde{f} has a $(\varepsilon, \varepsilon^{-c})$ diffuse decomposition.

Associate with each (h, q_1, \dots, q_m)

a d -tuple $K = (K_d, \dots, K_1) \in \mathbb{N}^d$

$$K_i = \sum_{\deg(q_j)=i} 3^{n-j} \times \text{'complexity of } q_j'$$

Define ordering $K \succ K'$ if $K_i = K'_i \quad i > i_0$
and $K_{i_0} > K'_{i_0}$.

Then \mathbb{N}^d has no infinite decreasing subsequence

Goal: If (h, q_1, \dots, q_m) is not
diffuse, can write $h(q_1, \dots, q_m) = h'(q'_1, \dots, q'_m)$

s.t. if $K, K' \in \mathbb{N}^d$, then $K \succ K'$

Eventually we end up with something diffuse,
or where q_i are all linear + simple (and thus
diffuse).

Goal: Reduce Degrees of q_i

Technique: Strong Anticoncentration

$$|f(x) - a| \geq |\nabla f(x)|$$

with high probability

Simplest Case (Step 1)

$$h(x) = x \quad q(x) = f(x)$$

Theorem: $f: \mathbb{R}^n \rightarrow \mathbb{R}$ has degree d ,

$$\mathbb{P}(|f(x) - a| < \varepsilon \mid \nabla f(x)) \lesssim \varepsilon 2^{O(d)} \log 1/\varepsilon$$

Thus either:

- 1) $|\nabla f(x)|$ is small with non-negligible probability
- 2) $f(x)$ is already diffuse.

If 2) we're done.

If 1) following result applies

Theorem: If $f: \mathbb{R}^n \rightarrow \mathbb{R}$ has degree d ,

$$\|f\|_{L^2_Y} = 1, \text{ and } \mathbb{P}(|\nabla f(x)| < \varepsilon) > \varepsilon^M$$

Then there are a_i, b_i with $0 \leq \deg a_i, \deg b_i < d$
 s.t.

$$\epsilon \lesssim \epsilon^{\epsilon^{o(1)}} \text{ fixed length}$$

$$\| (f(x) - \sum_{i=1}^{\ell} a_i(x) b_i(x))^{[d]} \|_{L_Y^2} \lesssim \epsilon^{1-c}$$

So replace f with $h(q_1, \dots, q_{m-1}) + q_m$

$$\{q_1, \dots, q_{m-1}\} = \underbrace{\{a_i\} \cup \{b_i\}}_{\deg < d} \quad \underbrace{\quad}_{\substack{\text{degree } d \\ \text{but} \\ O(\epsilon^{1-c})}}$$

General Step

Given (h, q_1, \dots, q_m)

Assume $\deg q_1 \geq \dots \geq \deg q_m$

Strong anticoncentration

$$\mathbb{P}\left(\prod_{i=1}^m |q_i(x) - a_i| < \varepsilon \mid \bigwedge_{i=1}^m \nabla q_i(x)\right) \\ \leq \varepsilon 2^{O(\sum \deg q_i)} O(\sqrt{m})^{m+1} \log(1/\varepsilon)^m$$

So either diffuse, or

$$\left| \bigwedge \nabla q_i(x) \right| \lesssim \varepsilon^{c/2}$$

with non-negligible probability.

If true, can choose i s.t.

$$\left| \nabla q_i - \sum_{j>i} a_j \nabla q_j \right| \lesssim \varepsilon^{c/3^i}$$

Now replace q_i with $q_i - \sum a_j q_j$

Then $|\nabla q_i| \lesssim \varepsilon^{c/3^i}$ so can

decompose $\|q_i - \sum a_j q_j\|_{L_Y} \lesssim \varepsilon^{1-c}$