Homework 1

Jiří Klepl

(10 points) Let G = (V, E) be an undirected graph. Suggest a propositional formula that is satisfiable if and only if G contains a Hamiltonian cycle.

We will describe a formula φ such that the assignment $\alpha \models \varphi$ will give us the Hamiltonian cycle. This formula will be satisfiable iff G has a Hamiltonian cycle.

Let a variable $x_{u,v}$ denote the fact the cycle takes edge $\{u,v\} \in E$ (with a direction from u to v) iff $\alpha(x_{u,v}) = 1$. The formula then has to describe the fact we visit each vertex exactly once and that we exit the vertex exactly once and that there is only one path:

• At least once:

$$\varphi_1 \equiv \bigwedge_{u \in V} \left(\bigvee_{\{u,v\} \in E} x_{u,v} \land \bigvee_{\{u,v\} \in E} x_{v,u} \land \bigwedge_{\{u,v\} \in E} (\neg x_{u,v} \lor \neg x_{v,u}) \right)$$

• At most once:

$$\varphi_2 \equiv \bigwedge_{(u,v),\{u,w\}\in E \land w \neq v} ([\neg x_{u,v} \lor \neg x_{u,w}] \land [\neg x_{v,u} \lor \neg x_{w,u}])$$

- One path: Let a variable $y_{u,v,w}$ denote that v is accessible by a path from u through w that is a part of the Hamiltonian cycle iff $\alpha(y_{u,v,w}) = 1$, where $u, v, w \in V$. Then:
 - Extensibility:

$$\varphi_3 \equiv \bigwedge_{u,v,w,x \in V \land \{v,x\} \in E \land |\{u,v,x\}| = 3} ([y_{u,v,w} \land x_{v,x}] \to y_{u,x,w})$$

- Asymmetry (non-triviality):

$$\varphi_4 \equiv \bigwedge_{u,v,w \in V \land |\{u,v,w\}|=3} ([y_{u,v,w} \to \neg y_{v,u,w}] \land [y_{u,v,w} \to \neg y_{u,w,v}])$$

- Universality:

$$\varphi_5 \equiv \bigwedge_{u,v \in V \land u \neq v} (y_{u,v,u} \land y_{u,v,v})$$

The formula has the form: $\varphi \equiv \varphi_1 \wedge \varphi_2 \wedge \varphi_3 \wedge \varphi_4 \wedge \varphi_5$.

 φ_1, φ_2 ensure that every vertex will be visited and exited exactly once.

 $\varphi_3, \varphi_4, \varphi_5$ ensure that the result will not be broken into separate cycles.