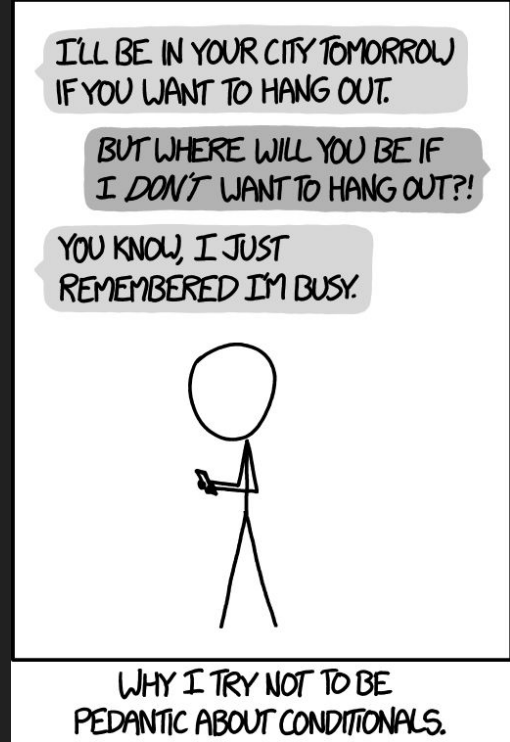


# Propositional Logic and Inference

Marcus Hughes  
CSCI 2824

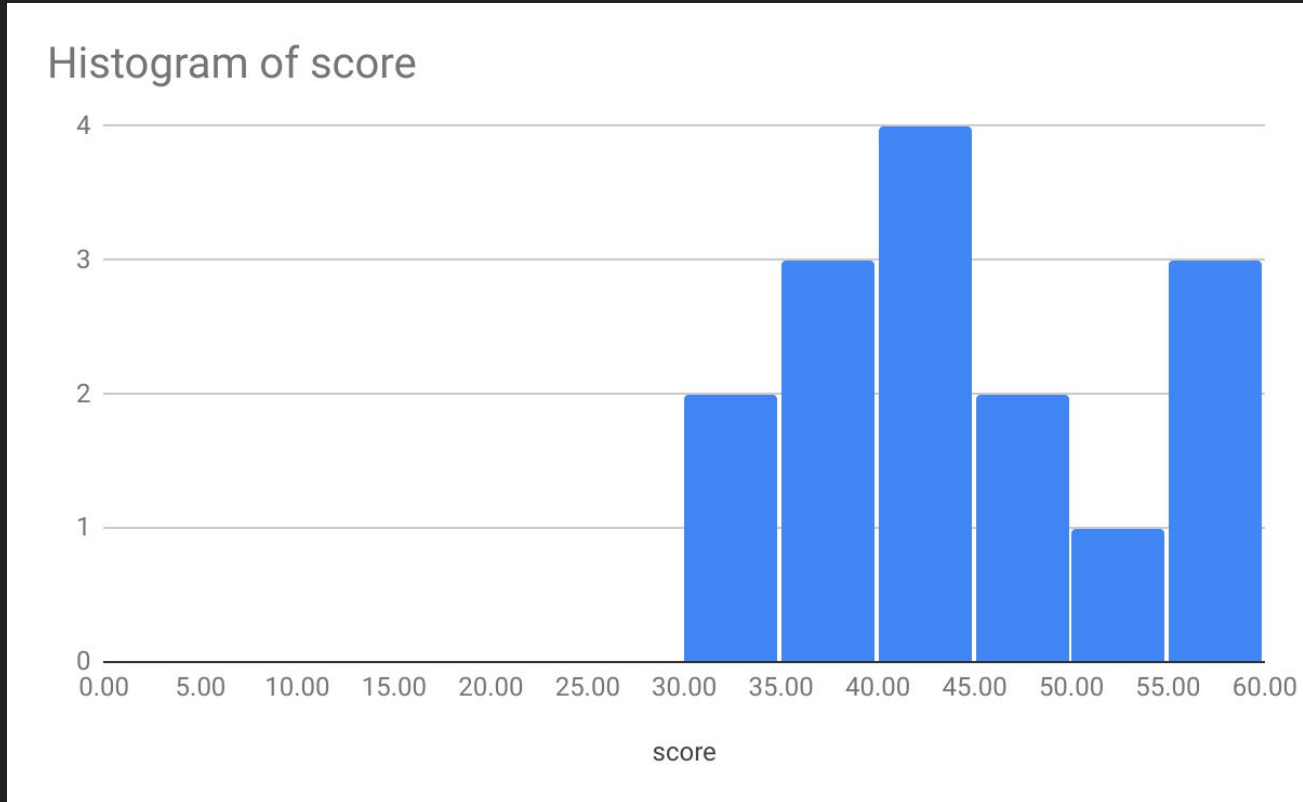
4 June 2019: Lecture 2



<https://xkcd.com/1652/>

<https://xkcd.com/1856/>

# Mindset

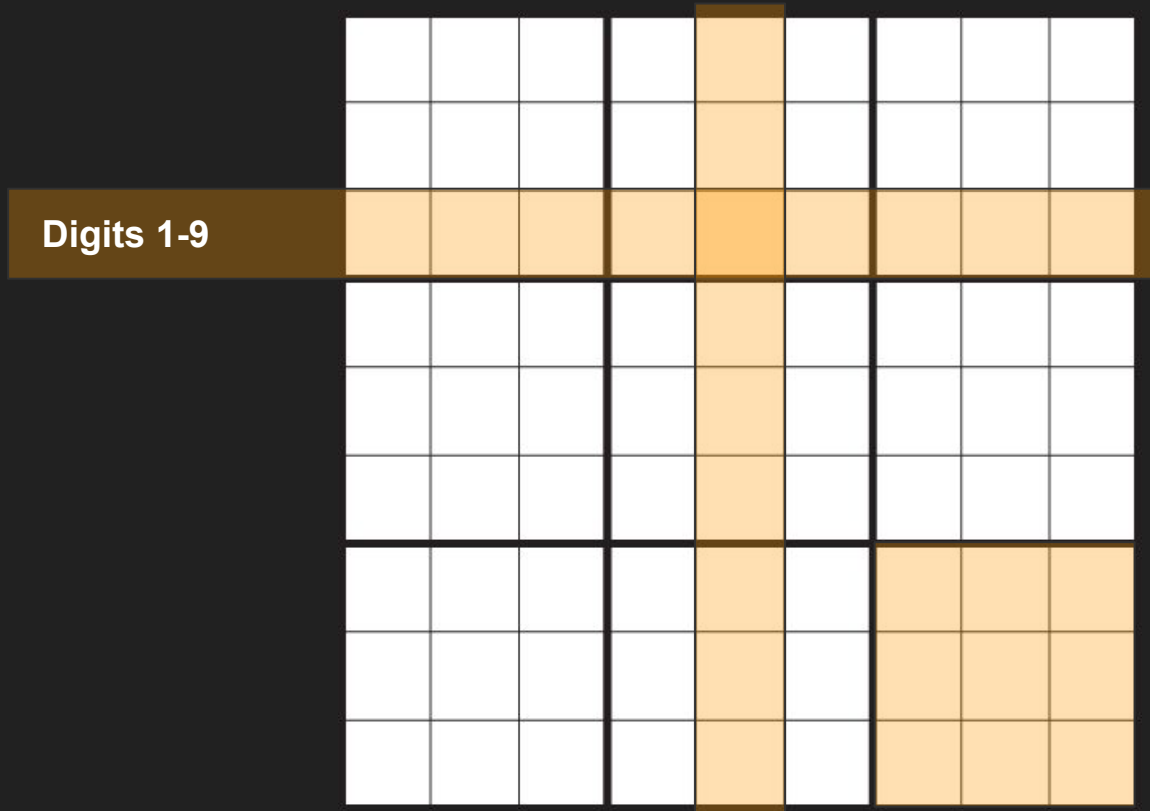


Sudoku as Logic!

# Sudoku as Constraint Satisfaction

5	3			7				
6			1	9	5			
	9	8					6	
8				6				3
4			8		3			1
7				2				6
	6					2	8	
			4	1	9			5
				8			7	9

# Sudoku as Constraint Satisfaction



# Sudoku as Constraint Satisfaction

Digits 1-9	5	3			7			
	6			1	9	5		
		9	8				6	
	8				6			3
	4			8		3		1
	7				2			6
		6					2	8
				4	1	9		5
					8		7	9

# Logical Equivalence

# Logical Equivalence

$p$	$q$	$p \wedge q$	$q \wedge p$
T	T	T	T
T	F	F	F
F	T	F	F
F	F	F	F



$p \wedge q$  and  $q \wedge p$  always  
have the same truth  
values, so they are  
logically equivalent



Are they equivalent?

$$\sim(p \wedge q)$$

$$\sim p \wedge \sim q$$

$p$	$q$	$\sim p$	$\sim q$	$p \wedge q$	$\sim(p \wedge q)$	$\sim p \wedge \sim q$
T	T	F	F	T	F	F
T	F	F	T	F	T	F
F	T	T	F	F	T	F
F	F	T	T	F	T	T



$\sim(p \wedge q)$  and  $\sim p \wedge \sim q$  have different truth values in rows 2 and 3, so they are not logically equivalent

Valid?

$$p \rightarrow q \vee \sim r$$

$$q \rightarrow p \wedge r$$

$$\therefore p \rightarrow r$$

			premises					conclusion
$p$	$q$	$r$	$\sim r$	$q \vee \sim r$	$p \wedge r$	$p \rightarrow q \vee \sim r$	$q \rightarrow p \wedge r$	$p \rightarrow r$
T	T	T	F	T	T	T	T	T
T	T	F	T	T	F	T	F	
T	F	T	F	F	T	F	T	
T	F	F	T	T	F	T	T	F
F	T	T	F	T	F	T	F	
F	T	F	T	T	F	T	F	
F	F	T	F	F	F	T	T	T
F	F	F	T	T	F	T	T	T

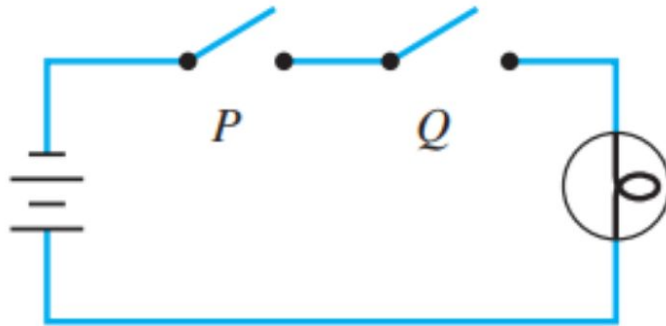
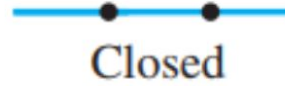
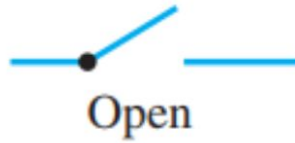
This row shows that an argument of this form can have true premises and a false conclusion. Hence this form of argument is invalid.



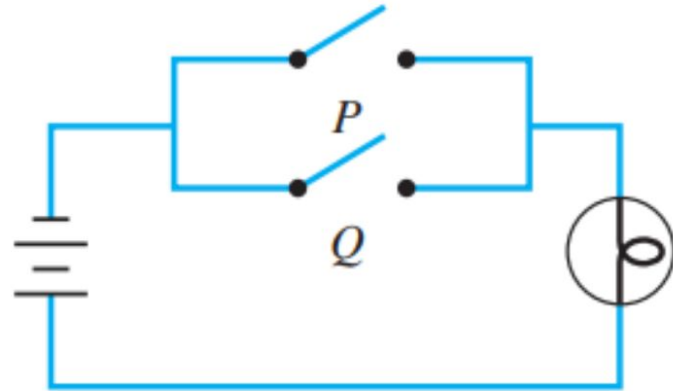
# Inference

<b>Modus Ponens</b>	$p \rightarrow q$ $p$ $\therefore q$	<b>Elimination</b>	<b>a.</b> $p \vee q$ $\sim q$ $\therefore p$	<b>b.</b> $p \vee q$ $\sim p$ $\therefore q$
<b>Modus Tollens</b>	$p \rightarrow q$ $\sim q$ $\therefore \sim p$	<b>Transitivity</b>	$p \rightarrow q$ $q \rightarrow r$ $\therefore p \rightarrow r$	
<b>Generalization</b>	<b>a.</b> $p$ $\therefore p \vee q$	<b>Proof by Division into Cases</b>	$p \vee q$ $p \rightarrow r$ $q \rightarrow r$ $\therefore r$	
<b>Specialization</b>	<b>b.</b> $q$ $\therefore p \vee q$			
	<b>a.</b> $p \wedge q$ $\therefore p$			
	<b>b.</b> $p \wedge q$ $\therefore q$			
<b>Conjunction</b>	$p$ $q$ $\therefore p \wedge q$	<b>Contradiction Rule</b>	$\sim p \rightarrow c$ $\therefore p$	

# Digital Circuits



Switches “in series”



Switches “in parallel”

# Circuits and Logic?

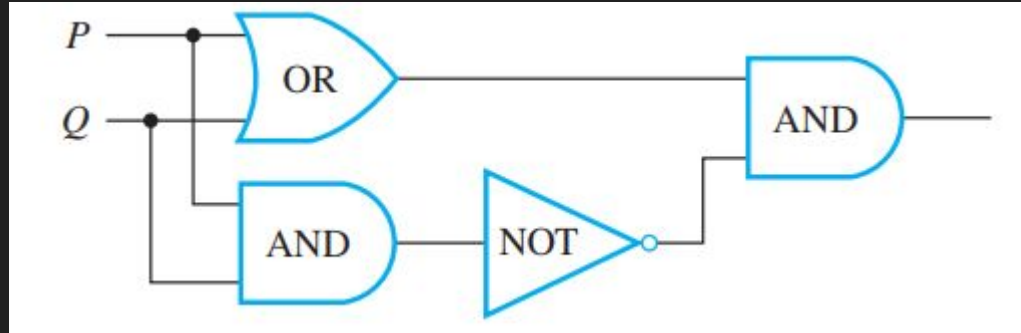
(a) Switches in Series

Switches		Light Bulb
$P$	$Q$	State
closed	closed	on
closed	open	off
open	closed	off
open	open	off

(b) Switches in Parallel

Switches		Light Bulb
$P$	$Q$	State
closed	closed	on
closed	open	on
open	closed	on
open	open	off

What is the formula?





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*What did you like? What did you not like?*