

# Lecture 6: Rational Number Proofs, Divisibility, and Induction

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## 1 Review

We learned about writing proofs. Let's practice with some ideas on rational numbers and divisibility.

## 2 Rational Number Proofs

What is a rational number? You may have heard it defined as a decimal that terminates or repeats in your early math days, but we are going to use a more formal definition.

**Definition:** *Rational Number*

A (real) number  $r$  is **rational** if, and only if, it can be expressed as a quotient of two integers with a nonzero denominator. A real number that is not rational is **irrational**. More formally, if  $r$  is a real number, then

$$r \text{ is rational} \leftrightarrow \exists a, b \in \mathbb{Z} \text{ such that } r = \frac{a}{b}, b \neq 0$$

What are some examples of rational numbers? Well integers are! Let's prove it.

**Theorem 2.1.** *Every integer is a rational number*

*Proof.* Let  $n$  be an arbitrary integer. Then  $n = \frac{n}{1}$  so  $n$  is a rational number. □

It's quite straight forward, but we are utilizing the idea of the generic particular. We do not specify what  $n$  is but show that for any arbitrary  $n$  we can make this argument. Seems pretty straight forward.

Now, try applying the same idea to show that any sum of rational numbers is rational.

**Theorem 2.2.** *The sum of two rational numbers is rational.*

*Proof.* Suppose  $r$  and  $s$  are rational numbers. Then, by the definition of rational,  $r = a/b$  and  $s = c/d$  for some  $a, b, c, d \in \mathbb{Z}$  where  $b \neq 0$  and  $d \neq 0$ . Thus,

$$\begin{aligned} r + s &= \frac{a}{b} + \frac{c}{d} \\ &= \frac{ad + bc}{bd} \end{aligned}$$

Let  $p = ad + bc$  and  $q = bd$ . Then,  $p$  and  $q$  are integers because products and sums of integers are integers and because  $a, b, c, d \in \mathbb{Z}$ . Also  $q \neq 0$  by the zero product property. Thus,  $r + s = \frac{p}{q}$  where  $p$  and  $q$  are integers and  $q \neq 0$ . Therefore,  $r + s$  is rational by the definition of a rational number.  $\square$

We can thus say that the rational numbers are closed under addition. Closed (in computer science lingo) means that performing that operation with a given input type always yields the same output type.

Can you show that any integer multiple of a rational number is rational? Is the product/quotient/difference of any two rational numbers also rational? What about the average?

### 3 Divisibility

Do you remember in elementary school when you talked about division? You knew that 12 wasn't divisible by 5 since it didn't go into it evenly. It may seem simplistic, but divisibility is a central idea in number theory; something so elementary is in advanced math! We can practice our proof writing skills in this domain too.

#### **Definition:** *Divisible*

If  $n$  and  $d$  are integers and  $d \neq 0$  then,  $n$  is **divisible by**  $d$  if, and only if,  $n$  equals  $d$  times some integer. The notation  $d|n$  is read " $d$  divides  $n$ ". Symbolically for  $n, d \in \mathbb{Z}$  and  $d \neq 0$ :

$$d|n \leftrightarrow \exists k \in \mathbb{Z} \text{ such that } n = dk$$

We call a number that divides another number a divisor of it. What do we know about divisors.

#### **Exercise**

For all integers  $a$  and  $b$ , if  $a$  and  $b$  are positive and  $a|b$ , then  $a \leq b$ .

Suppose  $a, b \in \mathbb{Z}^+$  and  $a|b$ . Then,  $\exists k \in \mathbb{Z}$  such that  $b = ak$ . From algebra (Property T25 of Appendix A),  $k$  must be positive because both  $a$  and  $b$  are positive. Then,  $1 \leq k$  because every positive integer is greater than or equal to 1. Multiplying both sides by  $a$  gives  $a \leq ka = b$  because multiplying both sides of an inequality by a positive number preserves the inequality property (T20 in Appendix A). Thus  $a \leq b$ .

The above is a nice universal proof to know.

**Exercise**

Prove: The only divisors of 1 are 1 and -1.

Since  $1 \times 1 = 1$  and  $(-1)(-1) = 1$  both 1 and  $-1$  are divisors of 1. Now suppose  $m$  is any integer that divides 1. Then  $\exists n \in \mathbb{Z}$  such that  $1 = mn$ . By Theorem T25 in Appendix A, either both  $m$  and  $n$  are positive or both  $m$  and  $n$  are negative. If both  $m$  and  $n$  are positive, then  $m$  is a positive integer divisor of 1. By the previous theorem,  $m \leq 1$ , and since the only positive integer that is less than or equal to 1 is 1 itself, it follows that  $m = 1$ . On the other hand, if both  $m$  and  $n$  are negative, then by Theorem T12 in Appendix A,  $(-m)(-n) = mn = 1$ . In this case,  $-m$  is a positive integer divisor of 1, and so by the same reasoning  $-m = 1$  and thus  $m = -1$ . Therefore, there are only two possibilities: either  $m = 1$  or  $m = -1$ . So the only divisors of 1 are 1 and -1.

**Exercise**

Problem: For all integers  $a$ ,  $b$ , and  $c$  if  $a$  divides  $b$  and  $b$  divides  $c$ , then  $a$  divides  $c$ .

Suppose  $a$ ,  $b$ , and  $c$  are integers such that  $a$  divides  $b$  and  $b$  divides  $c$ . By definition of divisibility  $\exists r, s \in \mathbb{Z}$  such that  $b = ar$  and  $c = bs$ . Then,

$$\begin{aligned} c &= bs \\ &= (ar)s \\ &= a(rs) \end{aligned}$$

Let  $k = rs$ . Then,  $k$  is an integer since it is a product of integers and therefore  $c = ak$ . Thus,  $a$  divides  $c$  by definition of divisibility.

**Exercise**

Prove: Any integer  $n > 1$  is divisible by a prime number.

Suppose  $n$  is an integer that is greater than 1. If  $n$  is prime, then  $n$  is divisible by itself, a prime number, and we are done. If  $n$  is not prime, then  $n = r_0 s_0$  where  $r_0$  and  $s_0$  are integers and  $1 < r_0 < n$  and  $1 < s_0 < n$ . It follows from the definition of divisibility that  $r_0 | n$ . If  $r_0$  is prime, then  $r_0$  is a prime number that divides  $n$  and we are done. If  $r_0$  is not prime, then,  $r_0 = r_1 s_1$  where  $r_1, s_1 \in \mathbb{Z}$  and  $1 < r_1 < r_0$  and  $1 < s_1 < r_0$ . It follows that  $r_1 | r_0$ . But we already know that  $r_0 | n$ . So by the previous proof  $r_1 | n$ .

We may continue factoring in this way, until we find a prime factor. We must succeed in a finite number of steps because each new factor is both less than the previous one and greater than 1, and there are fewer than  $n$  integers strictly between 1 and  $n$ . Thus, we obtain a sequence  $r_0, r_1, r_2, \dots, r_k$  where  $k \geq 0$ ,  $1 < r_k < r_{k-1} < \dots < r_2 < r_1 < r_0 < n$  and  $r_i | n$  for each  $i = 0, 1, \dots, k$ . The condition for termination is that  $r_k$  should be prime. Hence,  $r_k$  is a prime number that divides  $n$ .

Is the following statement true? For all integers  $a$  and  $b$  if  $a|b$  and  $b|a$  then  $a = b$ . No. find a counterexample with  $a = 2$  and  $b = -2$ .

A very powerful theorem is:

**Theorem 3.1.** *Given any integer  $n > 1$ , there exists a positive integer  $k$ , distinct prime numbers  $p_1, p_2, \dots, p_k$  and positive integers  $e_1, e_2, \dots, e_k$  such that  $n = p_1^{e_1} p_2^{e_2} \dots p_k^{e_k}$  and any other expression for  $n$  as a product of prime numbers is identical to this, except, perhaps, for the order in which the factors are written. We call this the standard factored form of  $n$  when  $p_1 < p_2 < \dots < p_k$ .*

We will come back to proving this later.

## 4 Mathematical Induction

We introduced the idea of mathematical induction yesterday. We will define two forms of

### **Definition:** *Weak Induction*

Let  $S_n$  denote a statement regarding an integer  $n$  and let  $k \in \mathbb{Z}$  be fixed. If

1.  $S_k$  holds and
2. for every  $m \geq k$ ,  $S_m \rightarrow S_{m+1}$

then for every  $n \geq k$ , the statement  $S_n$  holds.

### **Definition:** *Strong Induction*

Let  $S_n$  denote a statement regarding an integer  $n$ . If

1.  $S_k$  holds and
2. For every  $m \geq k$ ,  $[S_k \wedge S_{k+1} \wedge \dots \wedge S_m] \rightarrow S_{m+1}$

then for every  $n \geq k$ , the statement  $S_n$  is true.

### **Exercise**

Show that strong and weak induction are equivalent.

Let  $S_n$  be our statement we wish to prove.

First suppose that strong induction holds for our statement. This means that  $S_1$  holds and whenever  $n \leq k$ , it must hold for  $n = k + 1$ . This implies that  $n = k$  holds and  $n = k + 1$  holds, the criteria for weak induction. Therefore, weak induction follows from strong induction.

Now, suppose that weak induction works, also specifically for our statement. This means that  $S_1$  holds and for  $n = k$ , it must be that  $n = k + 1$  holds. Let  $Q_k$  be the statement " $S_n$  is true for all  $n \leq k$ ." We will prove  $Q_n$  is true for all positive integers  $n$  by weak induction. Since we have  $S_1$  then we also have the base case of  $Q_1$ . Suppose the inductive hypothesis of  $Q_k$ , i.e.  $S_k$  is true for all  $n \leq k$ . This tells us that  $S_{k+1}$  is also true by the outer induction proof. This then implies that  $Q_{k+1}$  holds since  $S_{k+1}$  and  $Q_k$  hold. This establishes the inductive step for statement  $Q_n$ . Therefore, this holds for all positive integers  $n$ . This is the same conclusion as strong induction makes for  $S_n$  meaning it is equivalent.

**Exercise**

Prove that  $1 + 2 + \dots + n = \frac{n(n+1)}{2}$ .

We could prove this directly by using Euler's trick of folding the sequence  $1, 2, 3, \dots, n$  on itself and creating pairs  $(1, n), (2, n-1), (3, n-2), \dots, (n/2 - 1, n/2 + 1)$  which each sum to  $n + 1$ . There are  $n/2$  such pairs so the total sum is  $\frac{n(n+1)}{2}$ .

However, you might not have that insight and thus rely on mathematical induction. Let  $S_n$  be the statement that  $1 + 2 + \dots + n = \frac{n(n+1)}{2}$ . Then,  $S_1$  says  $1 = \frac{1(1+1)}{2} = 1$  which holds. Now assume the inductive hypothesis that statement  $S_k$  holds for some  $k \geq 1$ , i.e.  $1 + 2 + \dots + k = \frac{k(k+1)}{2}$ . Then, we will show the inductive step that  $S_{k+1}$  follows.

$$\begin{aligned}
 1 + 2 + \dots + (k-1) + k + (k+1) &= [1 + \dots + (k-1) + k] + (k+1) && \text{regrouping} \\
 &= \frac{k(k+1)}{2} + k + 1 && \text{by inductive hypothesis} \\
 &= \frac{k(k+1) + 2k + 2}{2} && \text{algebra} \\
 &= \frac{k^2 + 3k + 2}{2} && \text{algebra} \\
 &= \frac{(k+1)(k+2)}{2} && \text{algebra}
 \end{aligned}$$

Therefore, we have proven the inductive step and by mathematical induction the theorem is proven.

**Exercise**

Conjecture a formula for the sum of the first  $n$  positive odd integers. Then prove your conjecture using mathematical induction.

First, we think about the conjecture. We might make a table like follows:

$n$	Integers	Sum
1	1	$1 = 1^2$
2	1, 3	$4 = 2^2$
3	1, 3, 5	$9 = 3^2$
4	1, 3, 5, 7	$16 = 4^2$

Thus, it seems we want to prove the sum of the first  $n$  positive odd integers is  $n^2$ .

Let  $S_n$  be the statement that the sum of the first  $n$  positive odd integers is  $n^2$ . The base case is then  $S_1$ , which holds since  $1 = 1^2$ .

Suppose the inductive hypothesis that  $S_k$  holds, i.e. the sum of the first  $k$  odd integers is  $k^2$ . Another way to state this is that  $1 + 3 + 5 + \dots + (2k-1) = k^2$ . Now we show that  $S_{k+1}$  holds.

$$\begin{aligned}
 1 + 3 + 5 + \dots + (2k-1) + (2k+1) &= k^2 + 2k + 1 \\
 &= (k+1)(k+1) = (k+1)^2
 \end{aligned}$$

Therefore, the inductive step holds and we have proven the claim.

**Exercise**

Use mathematical induction to show that  $1 + 2 + 2^2 + \dots + 2^n = 2^{n+1} - 1$ .

Let  $S_n$  be the statement that  $2^0 + 2^1 + 2^2 + \dots + 2^n = 2^{n+1} - 1$ . The base case,  $S_0$ , holds since  $2^0 = 1 = 2^{0+1} - 1$ .

Suppose the inductive hypothesis that  $S_k$  holds, i.e.  $1 + 2 + 2^2 + \dots + 2^k = 2^{k+1} - 1$ . Then, we wish to show  $S_{k+1}$  holds:

$$\begin{aligned} 1 + 2 + 2^2 + \dots + 2^k + 2^{k+1} &= 2^{k+1} - 1 + 2^{k+1} \\ &= 2(2^{k+1}) - 1 \\ &= 2^{k+2} - 1 \end{aligned}$$

Thus, the inductive step holds and we have proven the claim.

**Exercise**

Prove that for all  $n \in \mathbb{Z}^+$  that  $3|n^3 - n$ .

We could try the clever direct proof by noting that  $3|n^3 - n$  is the same as saying  $3|(n-1)n(n+1)$  if we factor the polynomial. Since  $n-1$ ,  $n$ , and  $n+1$  are consecutive one of them must be divisible by three meaning that we have a factor of three multiplied into the polynomial and thus we have our proof.

We could alternatively write an inductive proof. Then,  $S_n$  is the statement  $3|n^3 - n$ . The base case is  $S_1$ , i.e.  $3|1^3 - 1$ . We can let  $k = 0$  to satisfy  $0 = 3k$  and prove the base case holds.

Now suppose the inductive hypothesis of  $3|k^3 - k$ . We wish to prove the inductive step  $S_{k+1}$ . Note:

$$\begin{aligned} (k+1)^3 - (k+1) &= (k^3 + 3k^2 + 3k + 1) - (k+1) \\ &= (k^3 - k) + 3(k^2 + k) \end{aligned}$$

We know that  $3|k^3 - k$  by the inductive hypothesis and thus  $\exists m \in \mathbb{Z}$  such that  $3m = k^3 - k$ . So by substitution we get  $3m + 3(k^2 + k)$  which factors to  $3(m + k^2 + k)$ . Since integers are closed under multiplication and addition  $m + k^2 + k$  is an integer and we have shown that  $(k+1)^3 - (k+1)$  is indeed divisible by 3.