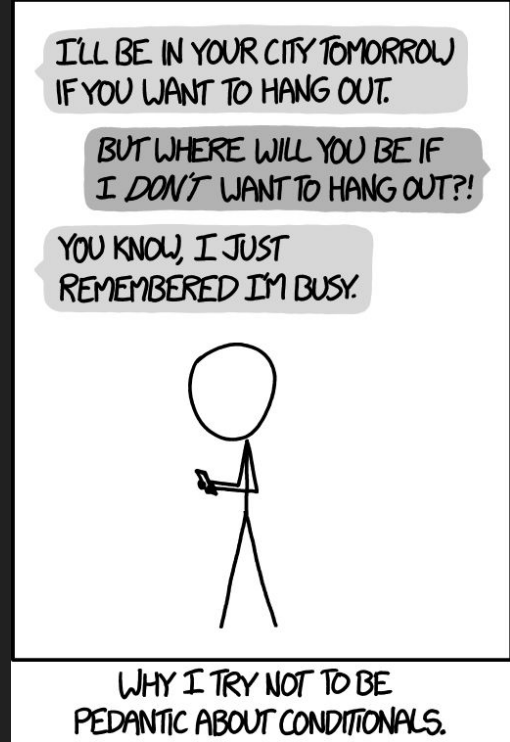


Propositional Logic and Inference

Marcus Hughes
CSCI 2824

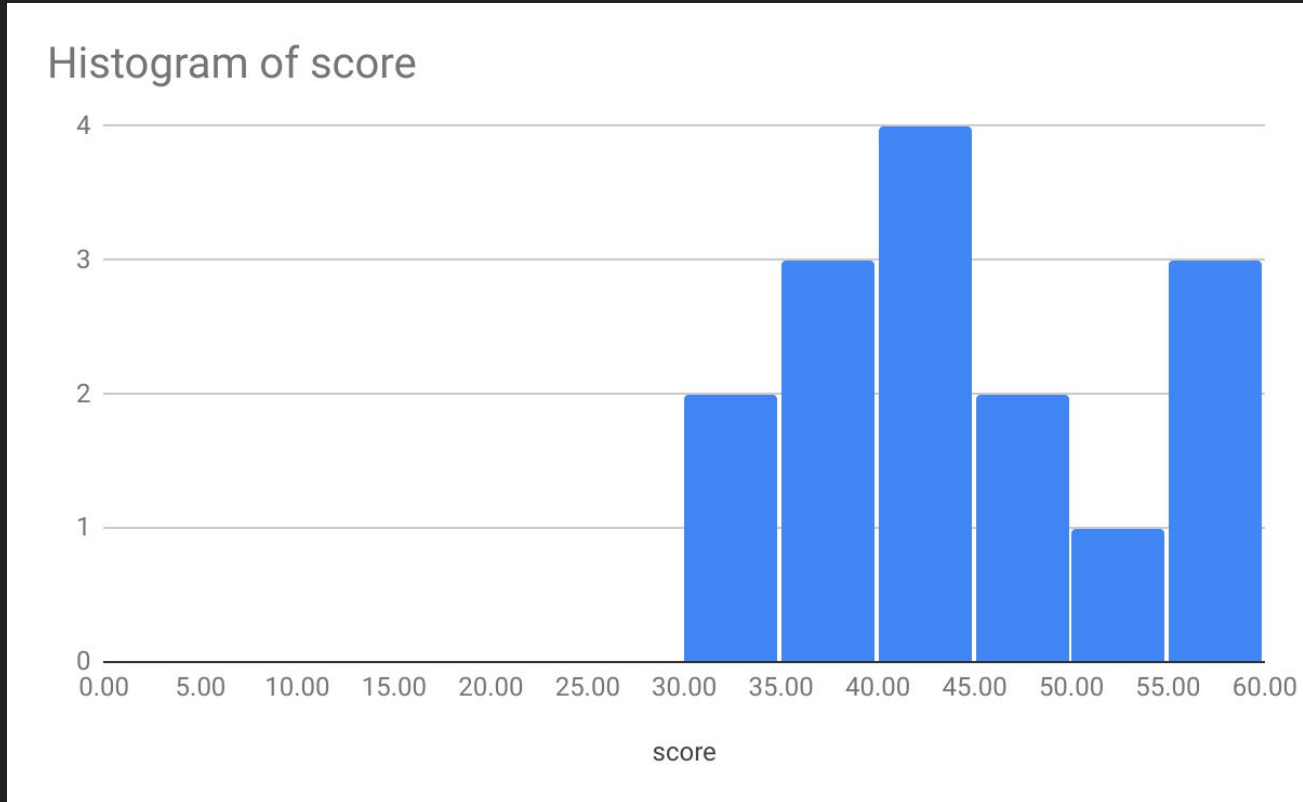
4 June 2019: Lecture 2



<https://xkcd.com/1652/>

<https://xkcd.com/1856/>

Mindset

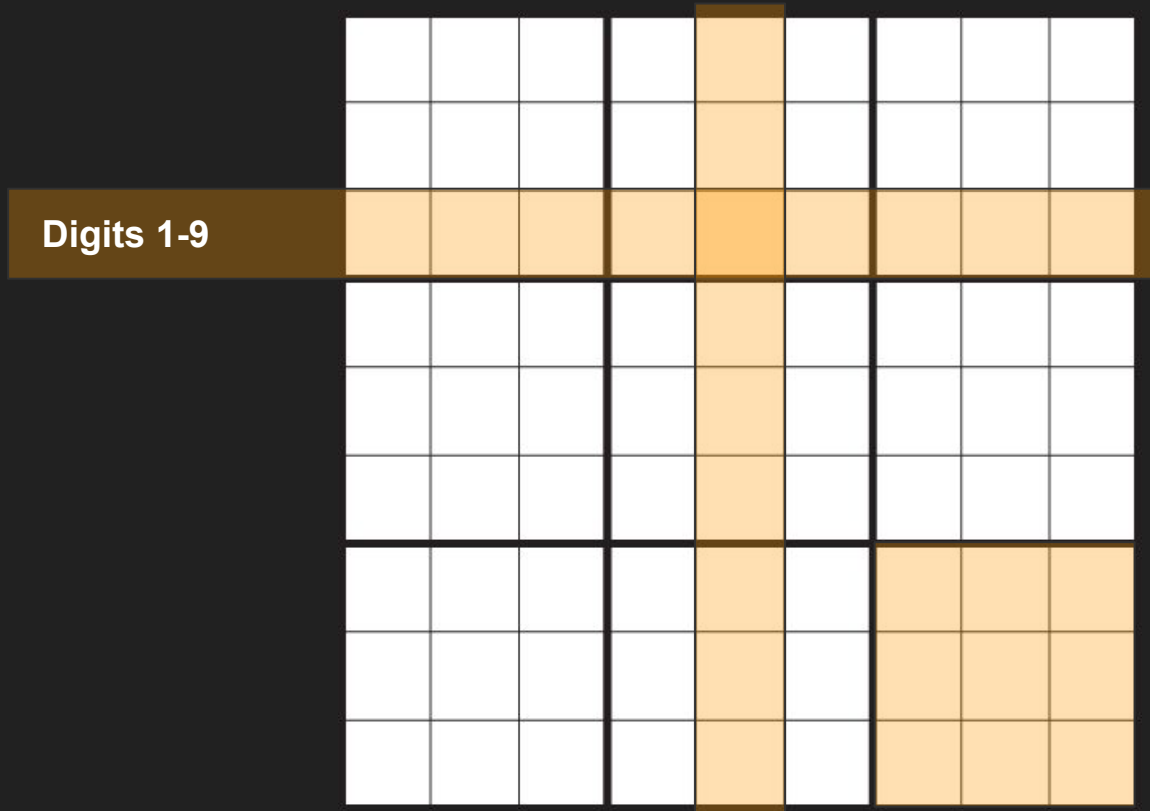


Sudoku as Logic!

Sudoku as Constraint Satisfaction

5	3			7				
6			1	9	5			
	9	8					6	
8				6				3
4			8		3			1
7				2				6
	6					2	8	
			4	1	9			5
				8			7	9

Sudoku as Constraint Satisfaction



Sudoku as Constraint Satisfaction

Digits 1-9	5	3			7			
	6			1	9	5		
		9	8				6	
	8				6			3
	4			8		3		1
	7				2			6
		6					2	8
				4	1	9		5
					8		7	9

Logical Equivalence

Logical Equivalence

p	q	$p \wedge q$	$q \wedge p$
T	T	T	T
T	F	F	F
F	T	F	F
F	F	F	F



$p \wedge q$ and $q \wedge p$ always have the same truth values, so they are logically equivalent

Are they equivalent?

$$\sim(p \wedge q)$$

$$\sim p \wedge \sim q$$

p	q	$\sim p$	$\sim q$	$p \wedge q$	$\sim(p \wedge q)$	$\sim p \wedge \sim q$
T	T	F	F	T	F	F
T	F	F	T	F	T	F
F	T	T	F	F	T	F
F	F	T	T	F	T	T



$\sim(p \wedge q)$ and $\sim p \wedge \sim q$ have different truth values in rows 2 and 3, so they are not logically equivalent

Valid?

$$p \rightarrow q \vee \sim r$$

$$q \rightarrow p \wedge r$$

$$\therefore p \rightarrow r$$

			premises					conclusion
p	q	r	$\sim r$	$q \vee \sim r$	$p \wedge r$	$p \rightarrow q \vee \sim r$	$q \rightarrow p \wedge r$	$p \rightarrow r$
T	T	T	F	T	T	T	T	T
T	T	F	T	T	F	T	F	
T	F	T	F	F	T	F	T	
T	F	F	T	T	F	T	T	F
F	T	T	F	T	F	T	F	
F	T	F	T	T	F	T	F	
F	F	T	F	F	F	T	T	T
F	F	F	T	T	F	T	T	T

This row shows that an argument of this form can have true premises and a false conclusion. Hence this form of argument is invalid.



Inference

Modus Ponens	$p \rightarrow q$ p $\therefore q$	Elimination	a. $p \vee q$ $\sim q$ $\therefore p$	b. $p \vee q$ $\sim p$ $\therefore q$
Modus Tollens	$p \rightarrow q$ $\sim q$ $\therefore \sim p$	Transitivity	$p \rightarrow q$ $q \rightarrow r$ $\therefore p \rightarrow r$	
Generalization	a. p $\therefore p \vee q$	Proof by Division into Cases	$p \vee q$ $p \rightarrow r$ $q \rightarrow r$ $\therefore r$	
Specialization	b. q $\therefore p \vee q$			
	a. $p \wedge q$ $\therefore p$			
	b. $p \wedge q$ $\therefore q$			
Conjunction	p q $\therefore p \wedge q$	Contradiction Rule	$\sim p \rightarrow c$ $\therefore p$	

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What did you like? What did you not like?