
DEMO

Aside: Quantifiers

$\forall x \in \{3,4,5\}.\ P(x)$	$P(3) \wedge P(4) \wedge P(5)$
$\forall x.\ x \in \{3,4,5\} \Rightarrow P(x)$	$\text{true} \wedge \text{true} \wedge P(3) \wedge P(4) \wedge P(5) \wedge \text{true} \wedge \text{true} \wedge \ldots$
$\forall x.\ x \notin \{3,4,5\} \vee P(x)$	

$\exists x \in \{3,4,5\}.\ P(x)$	$P(3) \vee P(4) \vee P(5)$
$\exists x.\ x \in \{3,4,5\} \wedge P(x)$	$\text{false} \vee \text{false} \vee P(3) \vee P(4) \vee P(5) \vee \text{false} \vee \text{false} \vee \ldots$

Hardware & Software Verification

John Wickerson

Lecture 5: SAT and SMT solving

Automatic proof

- We often rely on automatic provers:
  - e.g. in Dafny, to show that **`invariant`** `P` is preserved,
  - e.g. in Isabelle methods like **`by`** `auto`.
- How do these automatic provers work?

SAT queries

- Simple case: proofs about Boolean statements.
  - $f = (A \wedge \neg B \wedge C) \Longrightarrow (C \vee (B \wedge A))$
- Formulas can be:

always true

VALID

sometimes true

SATISFIABLE

UNSATISFIABLE

always false

INVALID

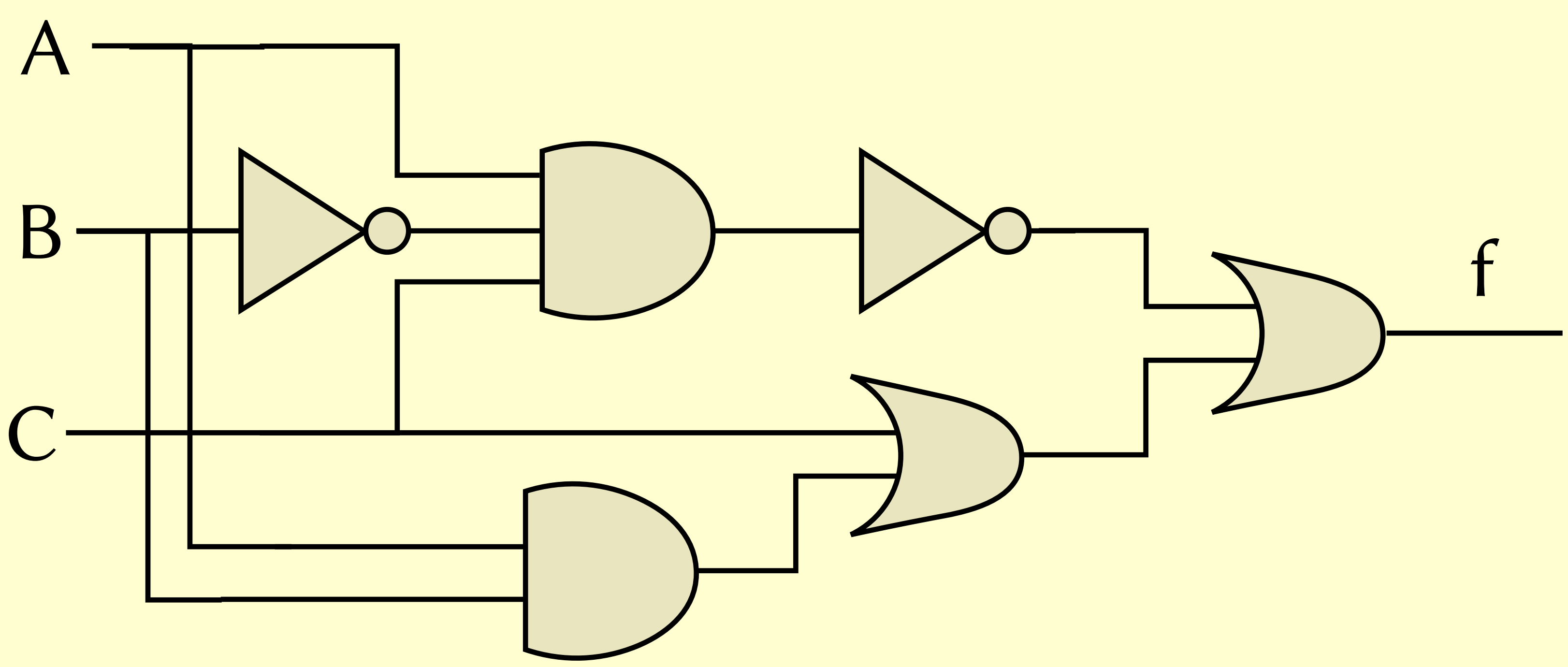
sometimes false

SAT queries

- Simple case: proofs about Boolean statements.
  - $f = (A \land \neg B \land C) \Longrightarrow (C \lor (B \land A))$

SAT queries

- Simple case: proofs about Boolean statements.
  - $f = \neg(A \wedge \neg B \wedge C) \vee (C \vee (B \wedge A))$



SAT queries

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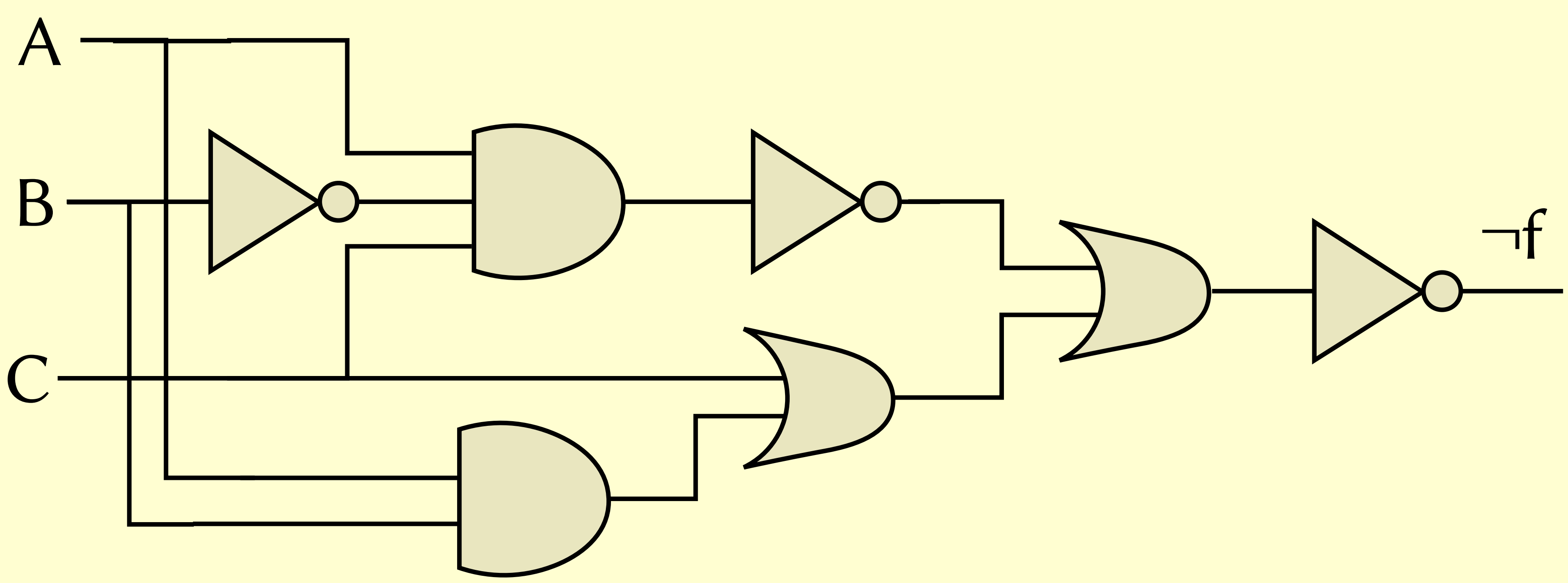


A	B	C	$\neg f$
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	0
1	1	0	0
1	1	1	0

SAT solving

- A simple algorithm:

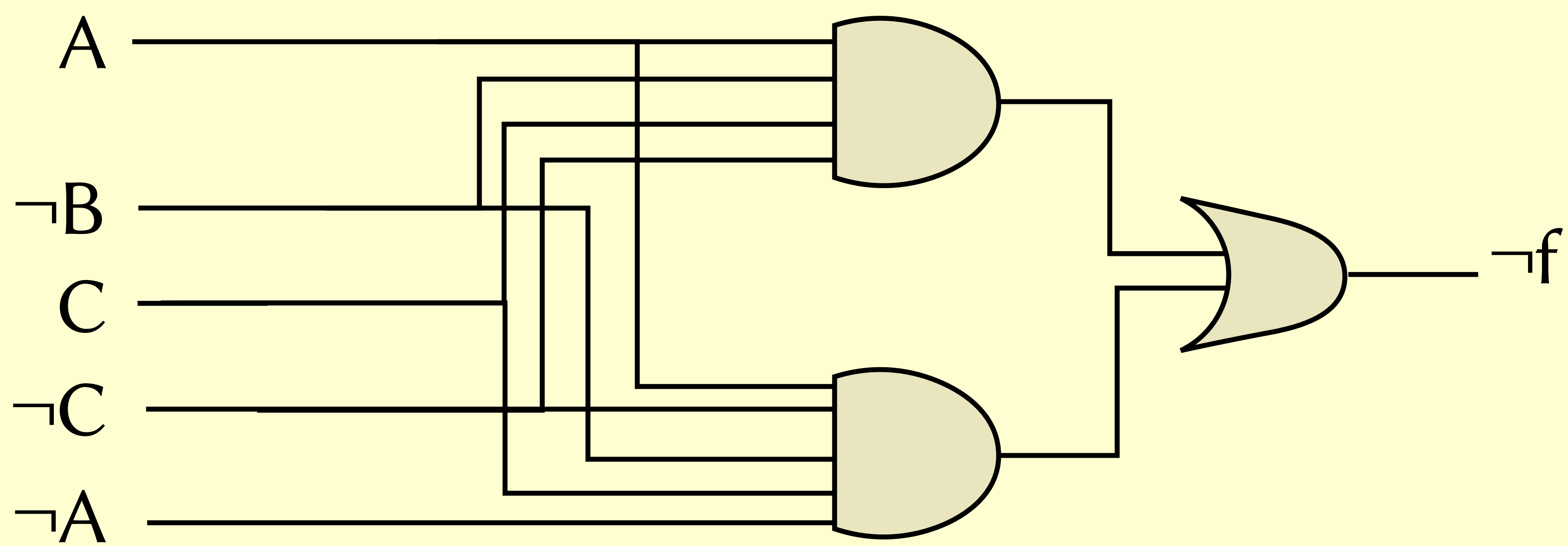
```
for A in {0, 1}:
  for B in {0, 1}:
    for C in {0, 1}:
      if ¬f(A,B,C) = 1:
        return ("SAT", [A,B,C])
return ("UNSAT")
```

- **Problem:** if the formula has N variables, this algorithm has exponential time-complexity, $O(2^N)$. ☹️

SAT solving

- **Idea:** Use de Morgan's laws to convert formula into *disjunctive normal form*.

$$\neg f \;=\; (A \land \neg B \land C \land \neg C) \lor (A \land \neg C \land \neg B \land C \land \neg A)$$



- **Hooray:** checking satisfiability becomes trivial!

SAT solving

- **Problem:** converting into disjunctive normal form has exponential time-complexity.

$$A \land \neg B \land C \land D \land \neg E \land F \land (G \lor H)$$

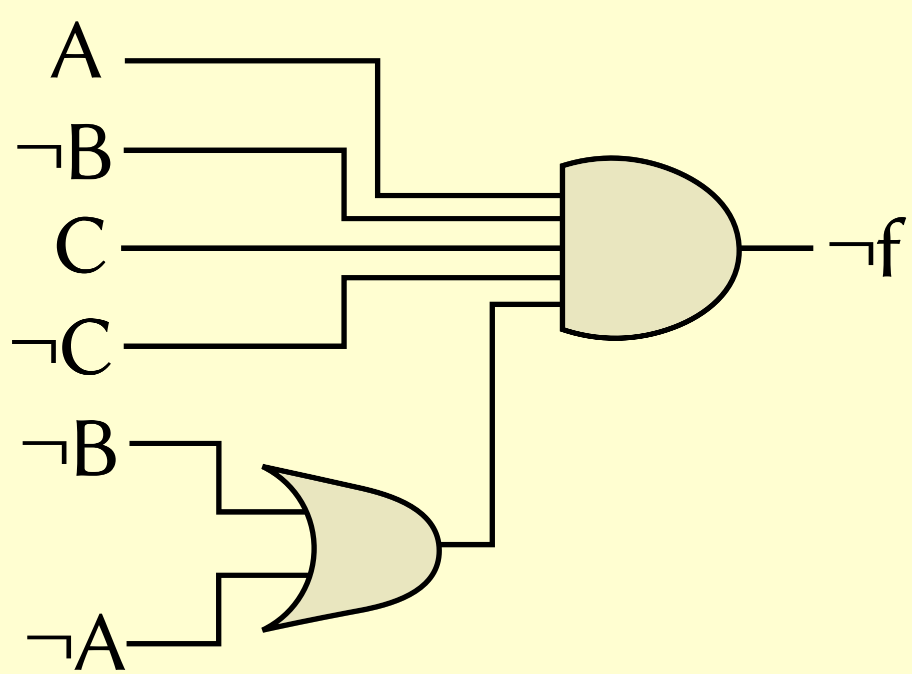
$\Downarrow$

$$(A \land \neg B \land C \land D \land \neg E \land F \land G) \lor (A \land \neg B \land C \land D \land \neg E \land F \land H)$$

SAT solving

- **Idea:** Use de Morgan's laws to convert formula into *conjunctive normal form.*

$$\neg f \;=\; A \wedge \neg B \wedge C \wedge \neg C \wedge (\neg B \vee \neg A)$$



SAT solving

- **Problem:** converting into conjunctive normal form still has exponential time-complexity.

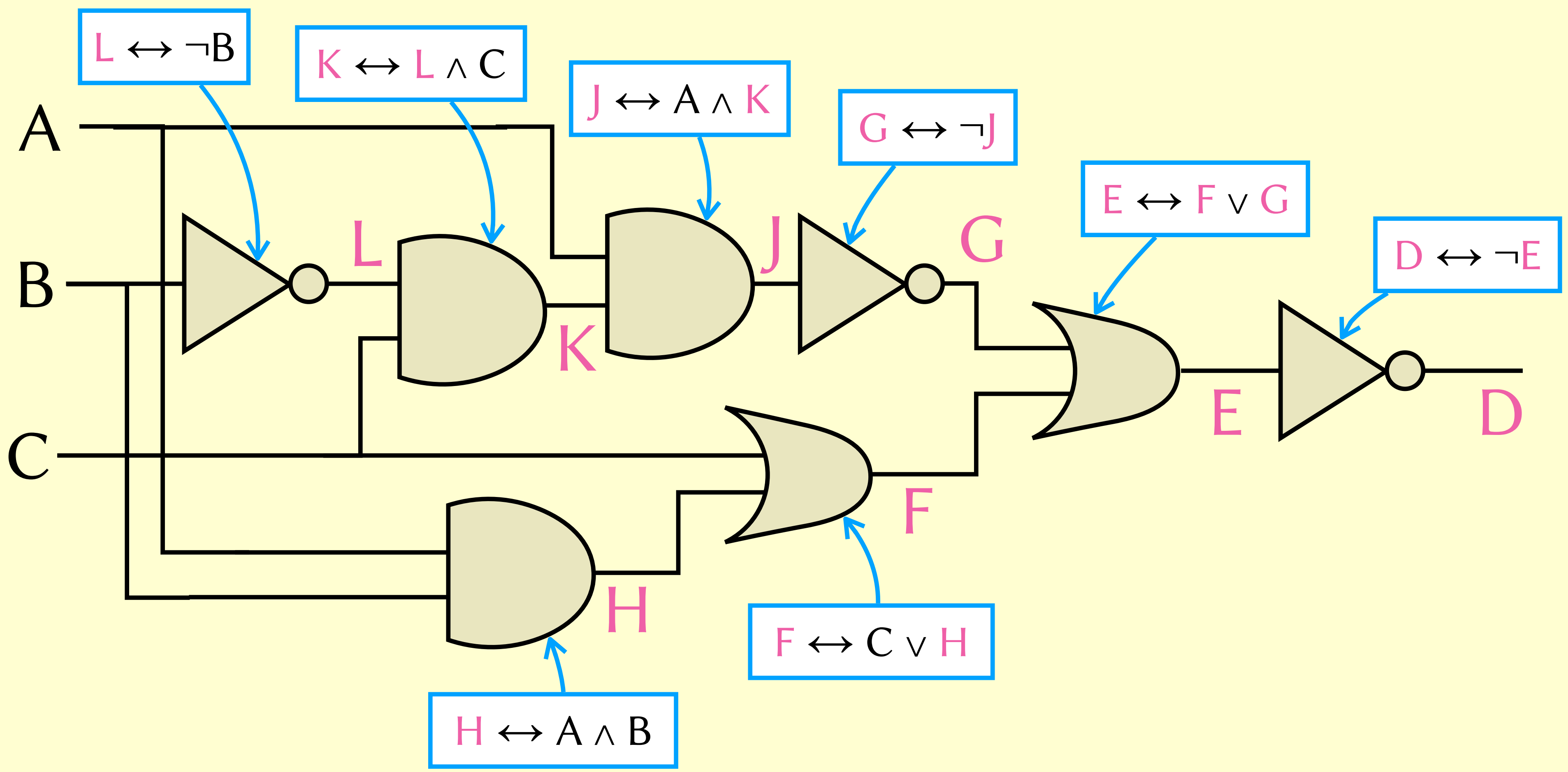
$$A \lor \neg B \lor C \lor D \lor \neg E \lor F \lor (G \land H)$$

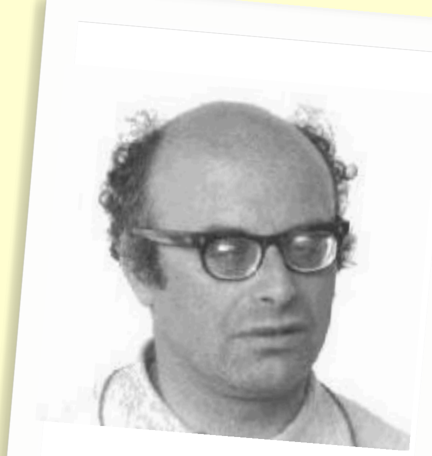
$$\Downarrow$$

$$(A \lor \neg B \lor C \lor D \lor \neg E \lor F \lor G) \land (A \lor \neg B \lor C \lor D \lor \neg E \lor F \lor H)$$

SAT solving

- **Solution:** the Tseitin transform. Make a fresh variable for each wire. Associate each gate with an equality.

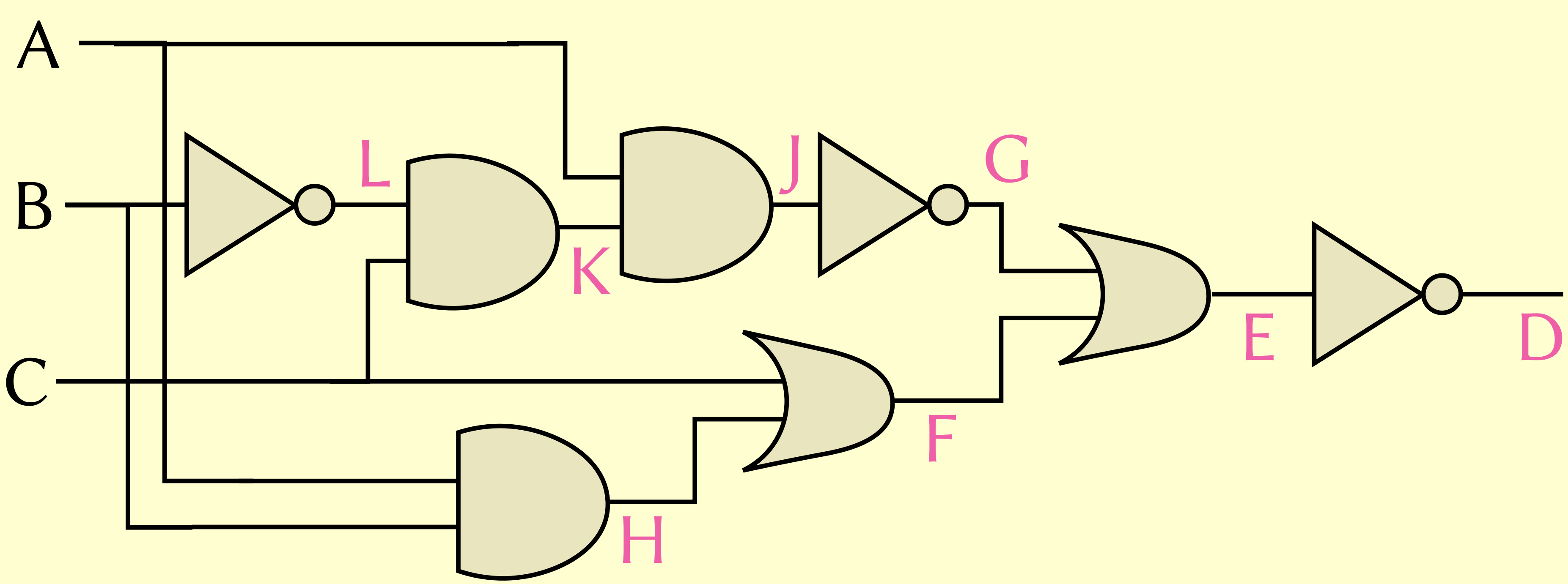




Gregory Tseitin
1936–2022

SAT solving

- **Solution:** the Tseitin transform. Make a fresh variable for each wire. Associate each gate with an equality.



$\boxed{L \leftrightarrow \neg B} \;\wedge$

$\boxed{K \leftrightarrow L \wedge C} \;\wedge$

$\boxed{J \leftrightarrow A \wedge K} \;\wedge$

$\boxed{G \leftrightarrow \neg J} \;\wedge$

$\boxed{E \leftrightarrow F \vee G} \;\wedge$

$\boxed{D \leftrightarrow \neg E} \;\wedge$

$\boxed{F \leftrightarrow C \vee H} \;\wedge$

$\boxed{H \leftrightarrow A \wedge B} \;\wedge D$

SAT solving

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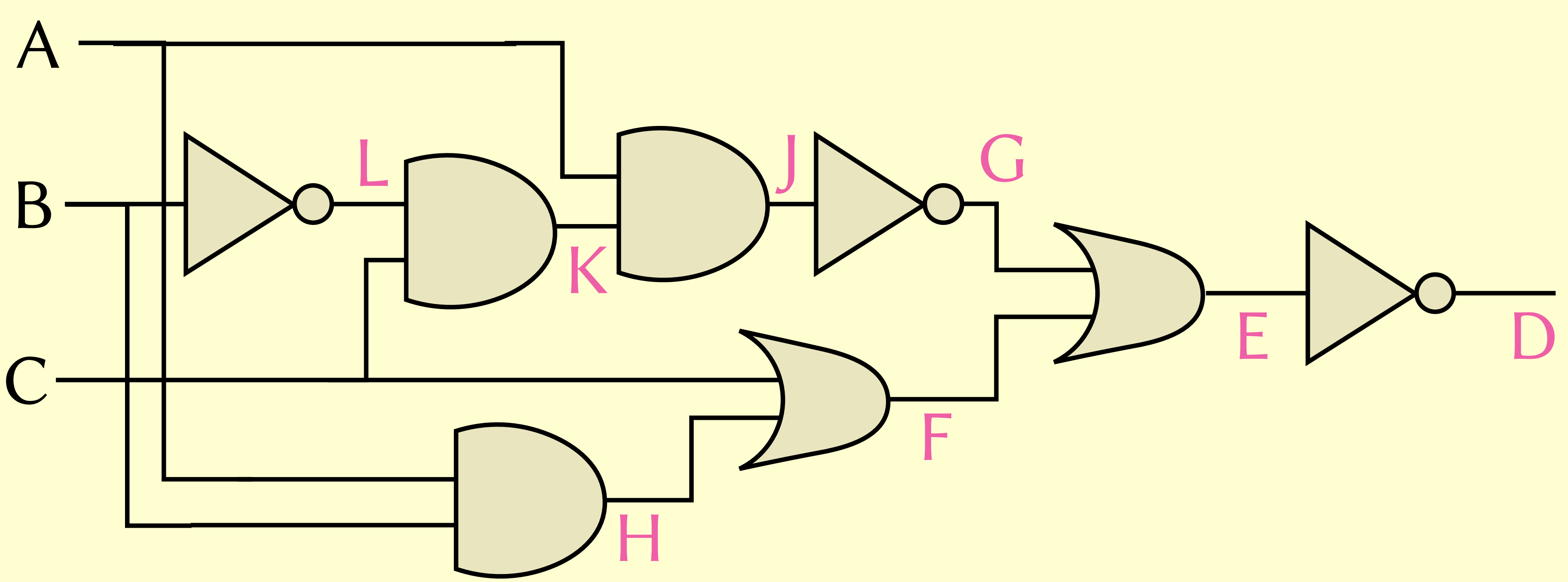

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$\boxed{E \leftrightarrow F \vee G} \wedge$

$\boxed{D \leftrightarrow \neg E} \wedge$

$\boxed{F \leftrightarrow C \vee H} \wedge$

$\boxed{H \rightarrow A \wedge B} \wedge \boxed{A \wedge B \rightarrow H} \wedge D$

SAT solving

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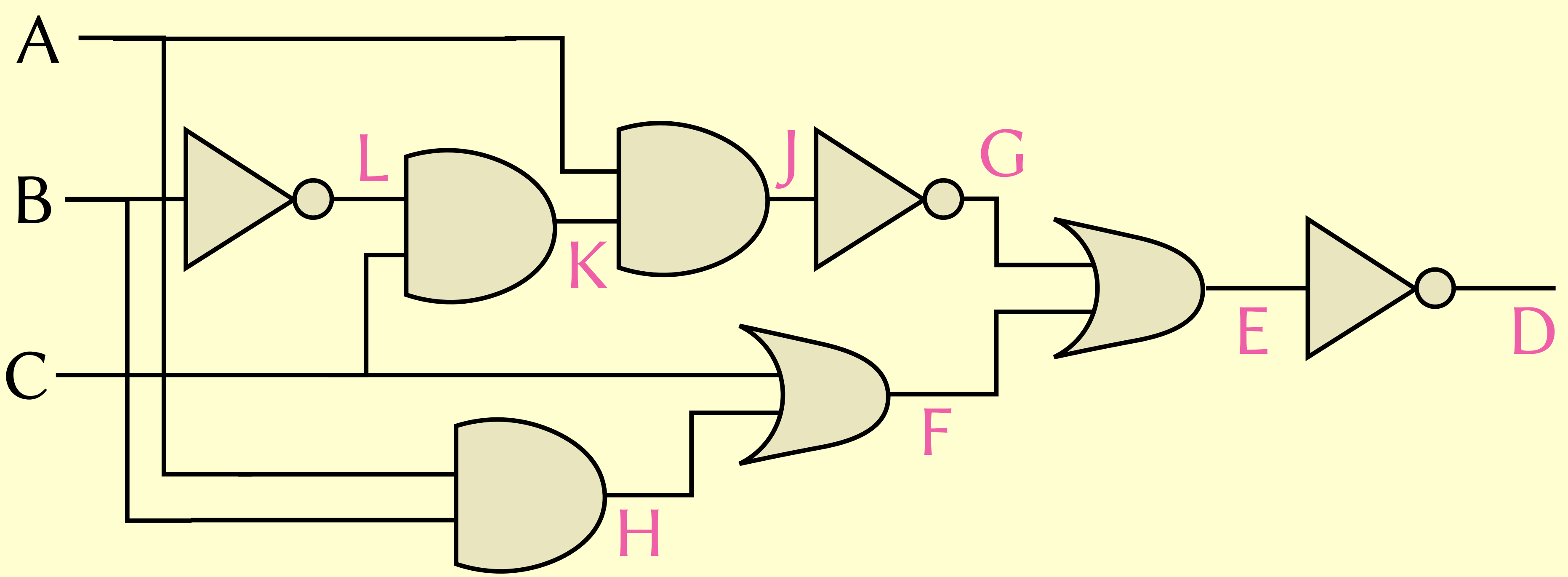

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$\boxed{D \leftrightarrow \neg E} \wedge$

$\boxed{F \leftrightarrow C \vee H} \wedge$

$\boxed{\neg H \vee (A \wedge B)} \wedge \boxed{\neg(A \wedge B) \vee H} \wedge D$

SAT solving

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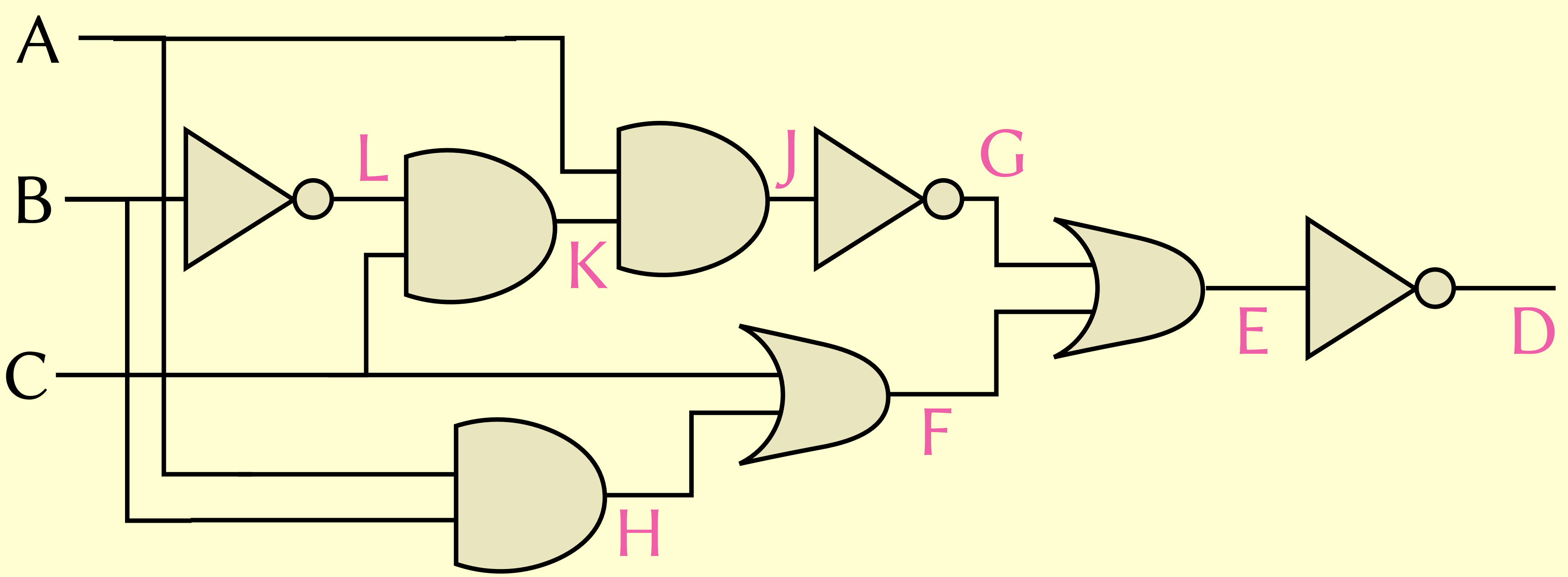

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$\boxed{D \leftrightarrow \neg E} \;\wedge$

$\boxed{F \leftrightarrow C \vee H} \;\wedge$

$\boxed{\neg H \vee A} \;\wedge\; \boxed{\neg H \vee B} \;\wedge\; \boxed{\neg(A \wedge B) \vee H} \;\wedge\; D$

SAT solving

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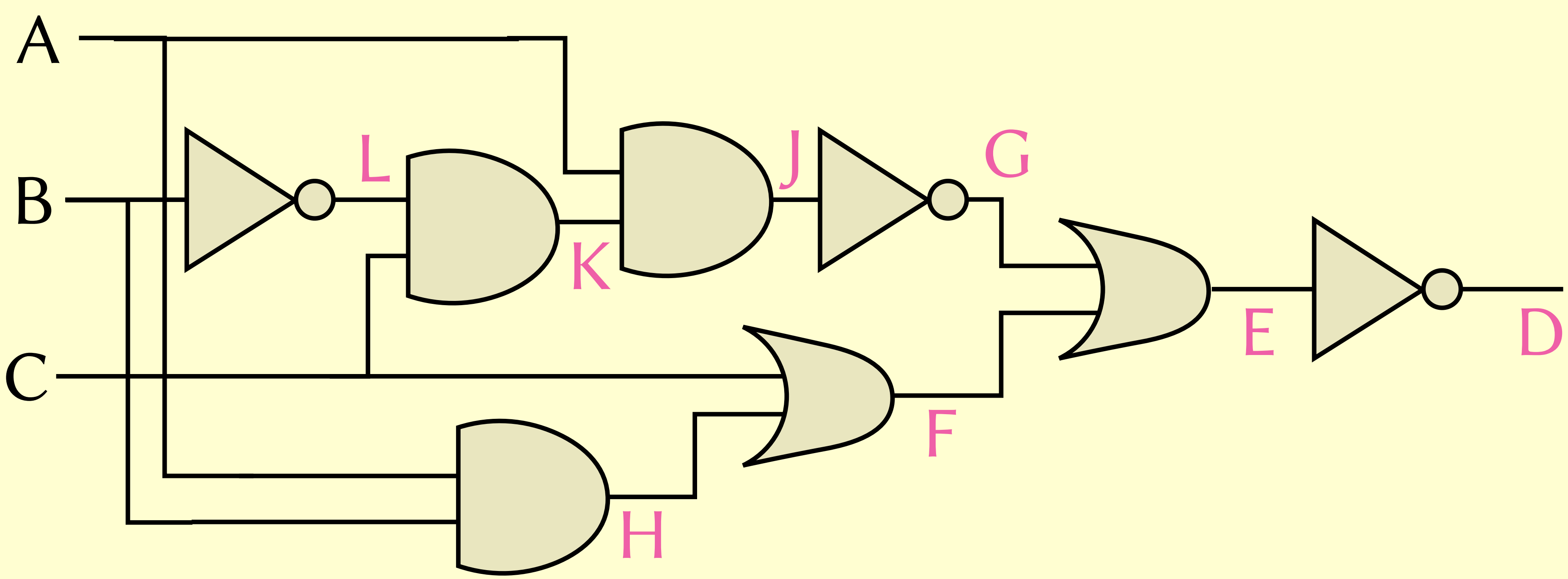

$$\boxed{L \leftrightarrow \neg B} \land$$
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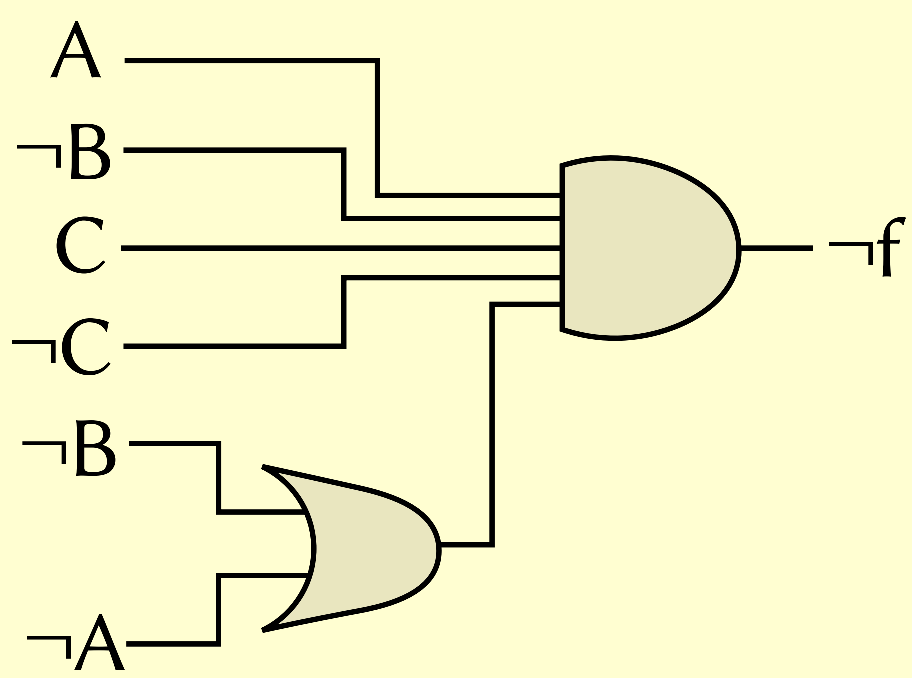
$$\boxed{D \leftrightarrow \neg E} \land$$
$$\boxed{F \leftrightarrow C \lor H} \land$$
$$\boxed{\neg H \lor A} \land \boxed{\neg H \lor B} \land \boxed{\neg A \lor \neg B \lor H} \land D$$

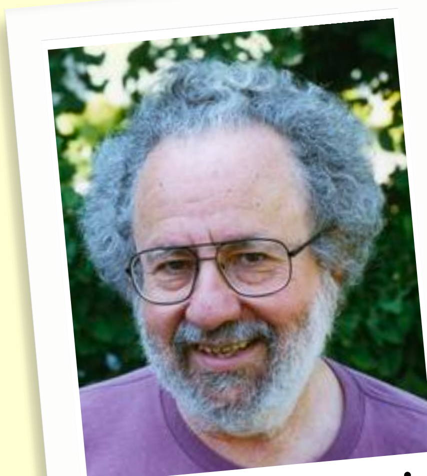
SAT solving

- **Problem:** the satisfiability problem for CNF is difficult (unlike for DNF).

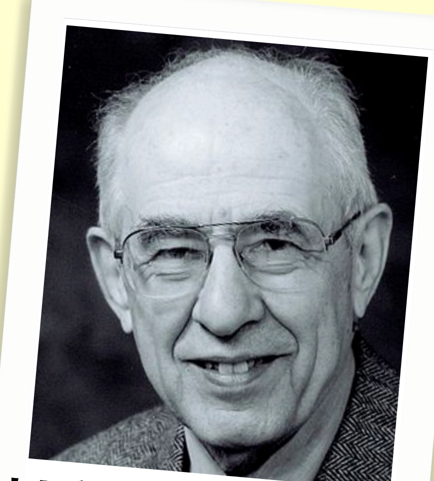


SAT solving

- **Problem:** the satisfiability problem for CNF is difficult (unlike for DNF).
- **Solution:** Davis and Putnam to the rescue!



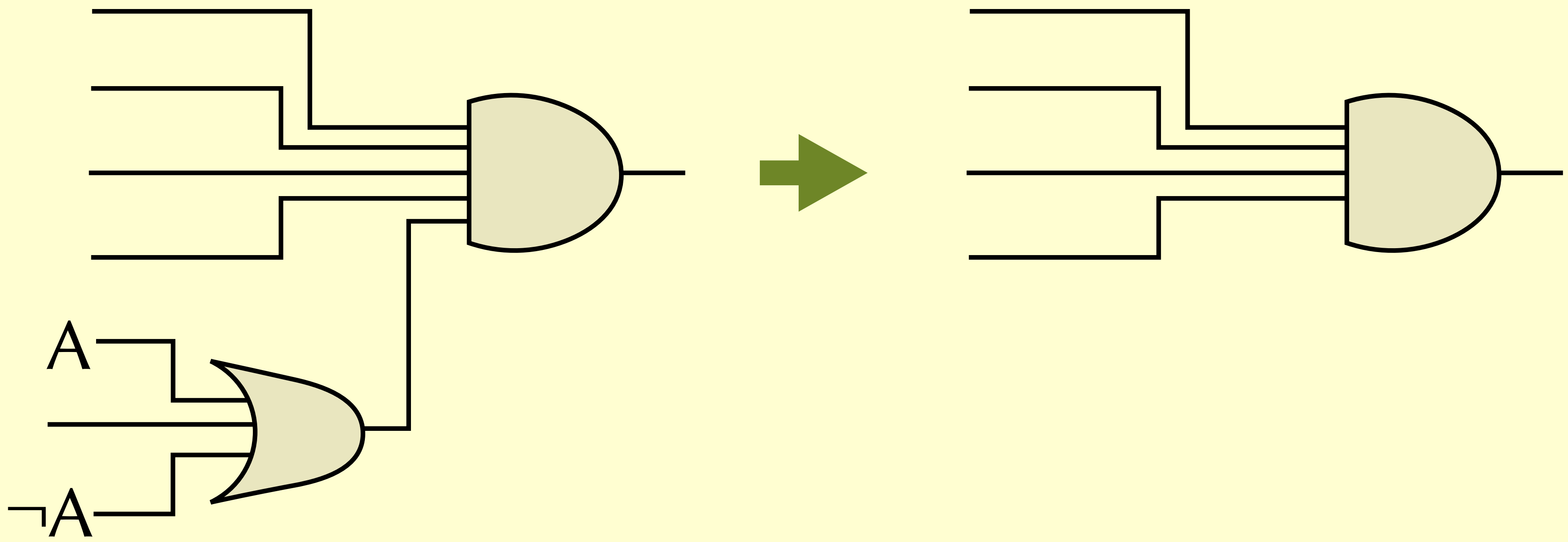
Martin Davis
1928–2023



Hilary Putnam
1926–2016

The DP method

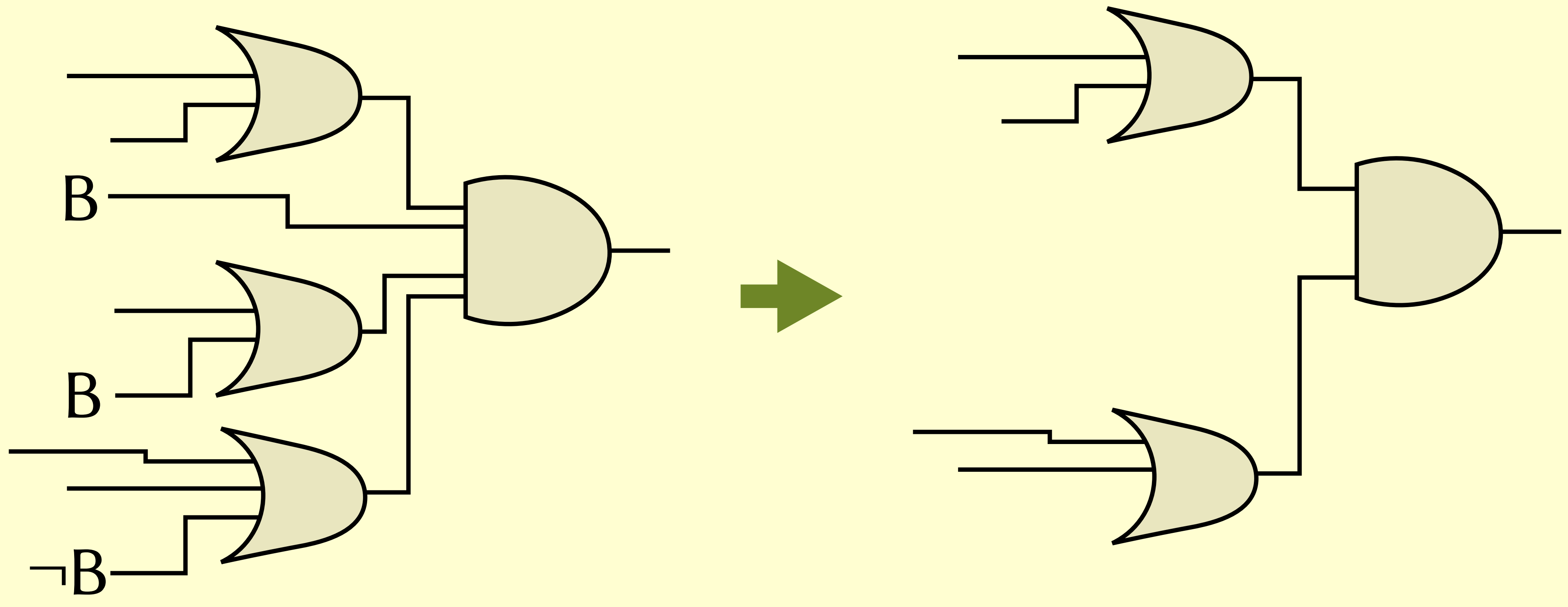
1. If an OR-gate takes both L and ¬L, delete it.



The DP method

1. If an OR-gate takes both L and ¬L, delete it.

2. If L is connected directly to the AND-gate, delete it, delete all OR-gates that take L, and delete any connections to ¬L.
   (The solution, if it exists, will surely involve setting L=1.)



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3. If L is unused, delete all OR-gates that take ¬L.
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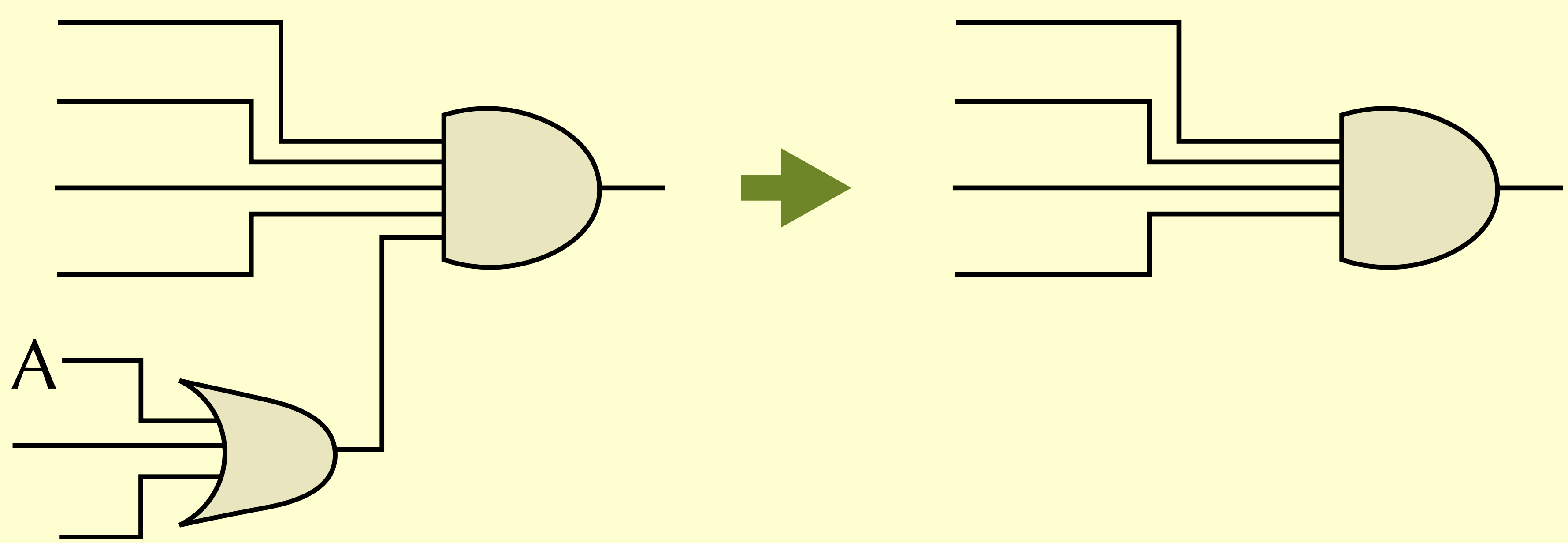

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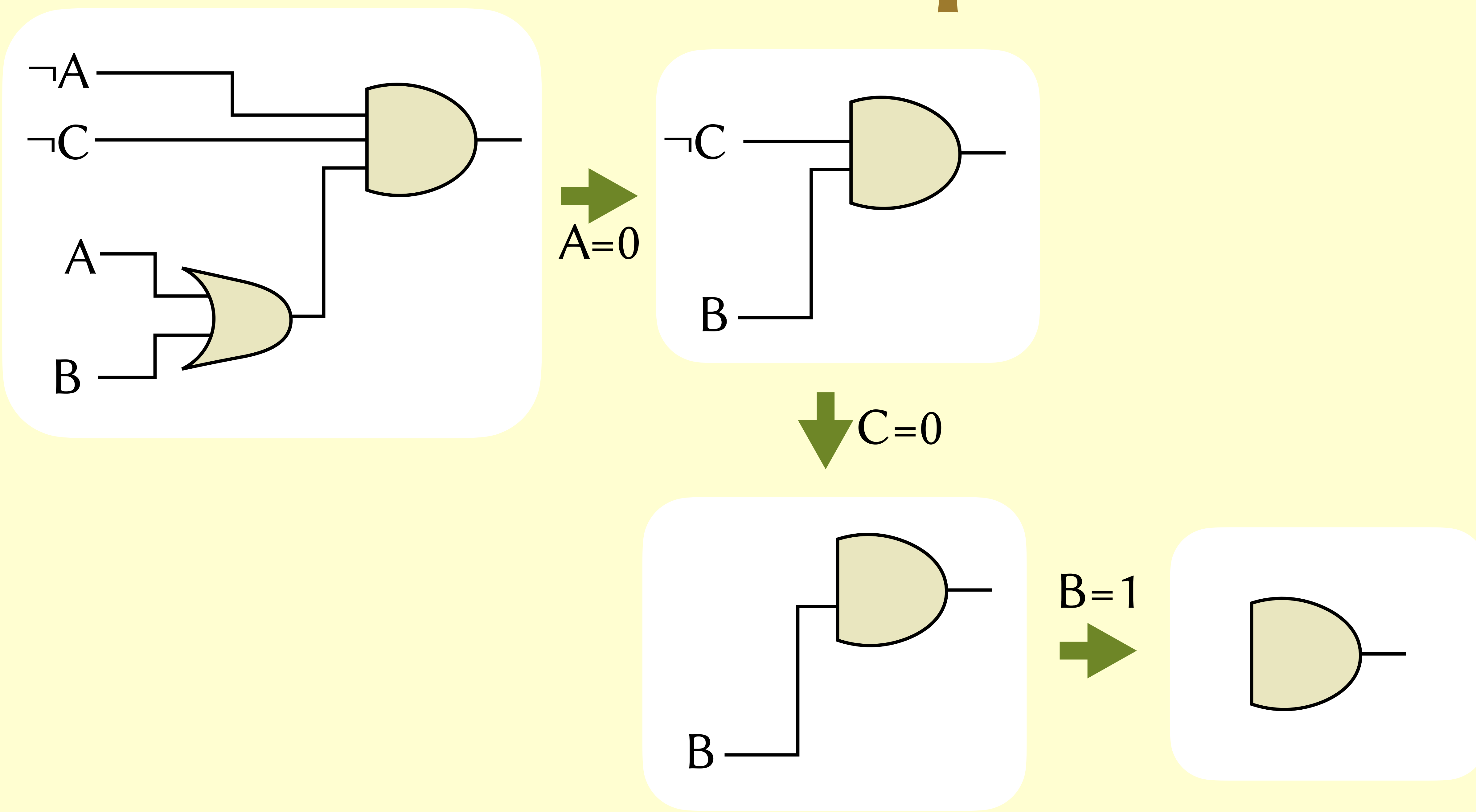
3. If L is unused, delete all OR-gates that take ¬L.
   (The solution, if it exists, will surely involve setting L=0.)

4. If any OR-gate has no inputs, the formula is false.

5. If the AND-gate has no inputs, the formula is true.

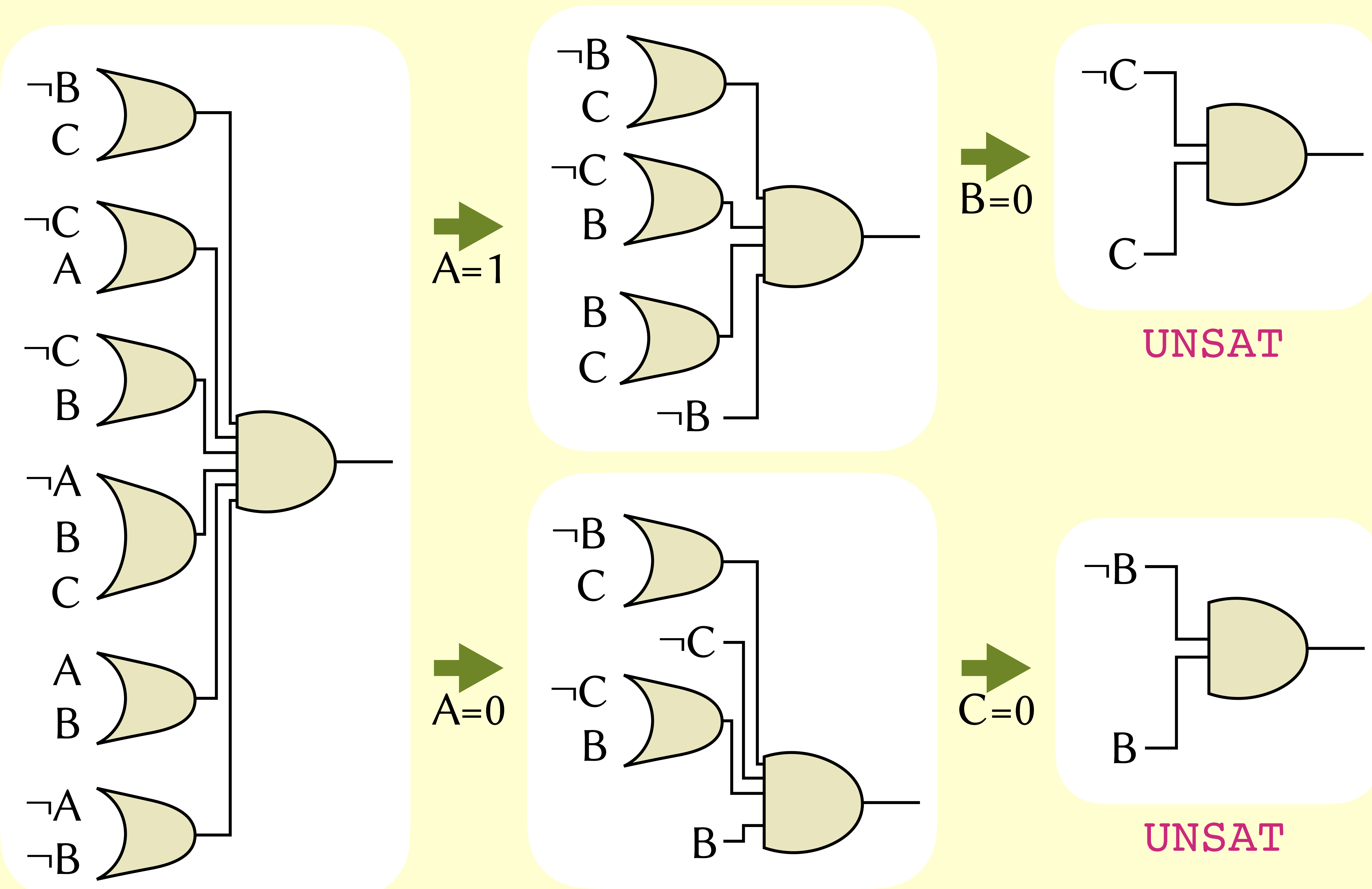
6. Pick a literal L and repeat the above for the cases L=0 and L=1.

DP example 1



SAT, A=0, B=1, C=0

DP example 2

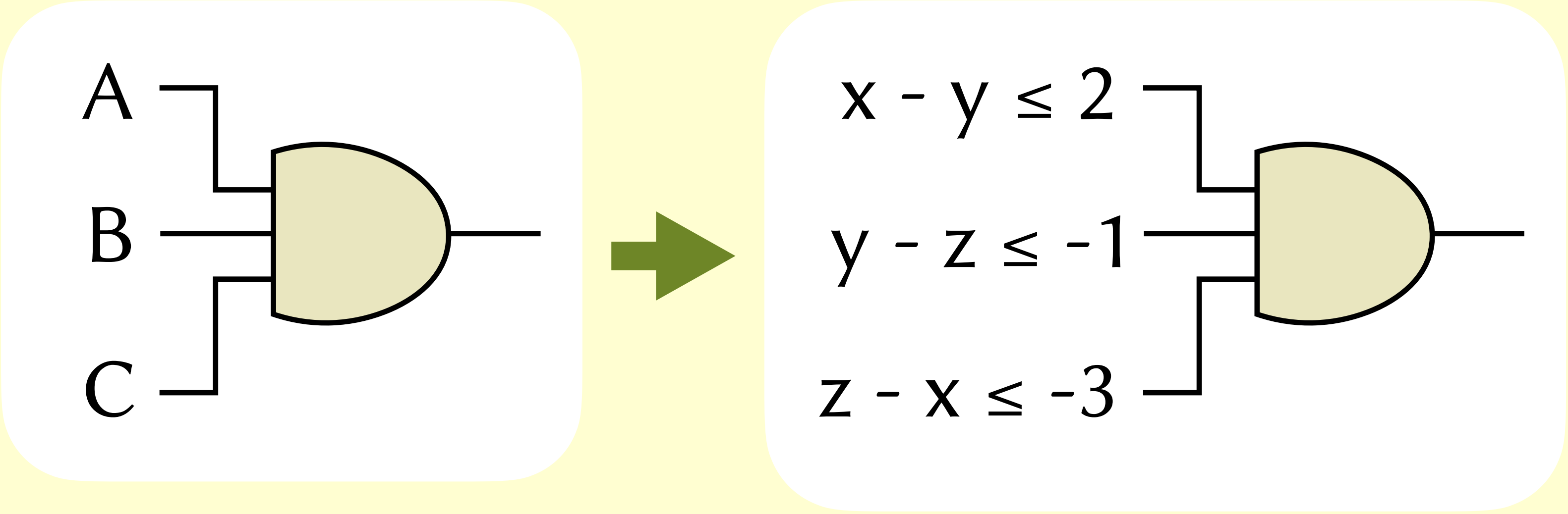


Towards SMT solving

- We can now prove basic Boolean formulas. But what about proving something like $A \times (B + C) = A \times B + A \times C$?

- If these are 32-bit integers, we could make this a SAT problem by treating each variable as 32 Boolean variables and encoding the rules of Boolean arithmetic.

- Or we can move up to SMT: *satisfiability modulo theories.*



Some theories

- **Equality and uninterpreted functions**, which knows that x=y and y=z implies x=z, and that x=y implies f(x)=f(y).
- **Difference logic**, where statements take the form $x - y \leq c$.
- **Presburger arithmetic**, which allows statements about naturals containing +, 0, 1, and =. For instance, n is a McNugget number if $\exists x\ y\ z.\ n = 6x + 9y + 20z$.

Mojżesz Presburger
1904–c.1943

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- **Presburger arithmetic**, which allows statements about naturals containing +, 0, 1, and =. For instance, n is a McNugget number if $\exists x\ y\ z.\ n = 6x + 9y + 20z$.
- **Non-linear arithmetic**, which allows queries like:

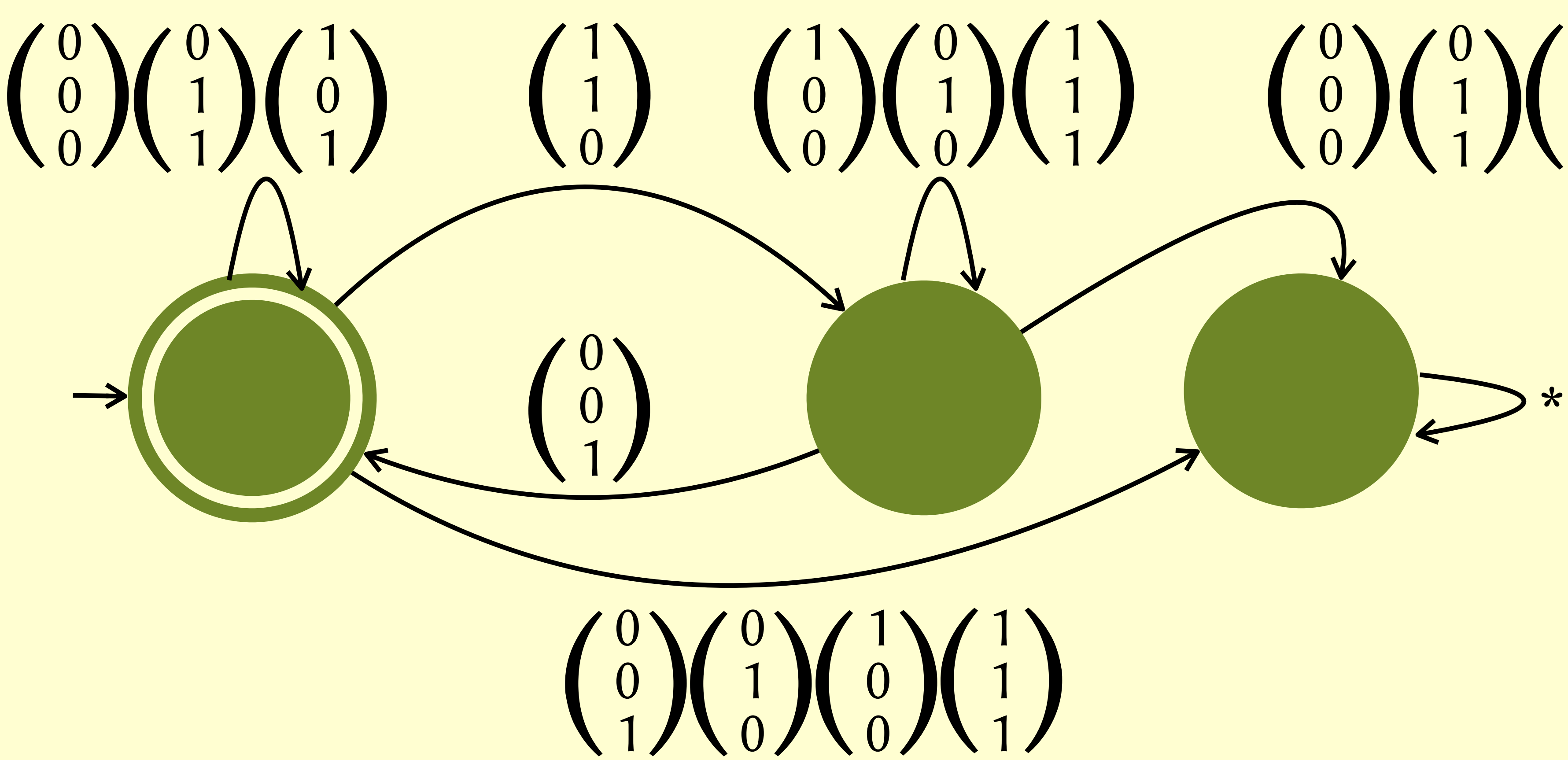
$$(\sin(x)^3 = \cos(\log(y) \cdot x) \lor b \lor -x^2 \geq 2.3y) \land \left(\neg b \lor y < -34.4 \lor \exp(x) > \frac{y}{x}\right)$$

- **Theory of arrays**, **theory of bit-vectors**, etc.

Decidability of Presburger

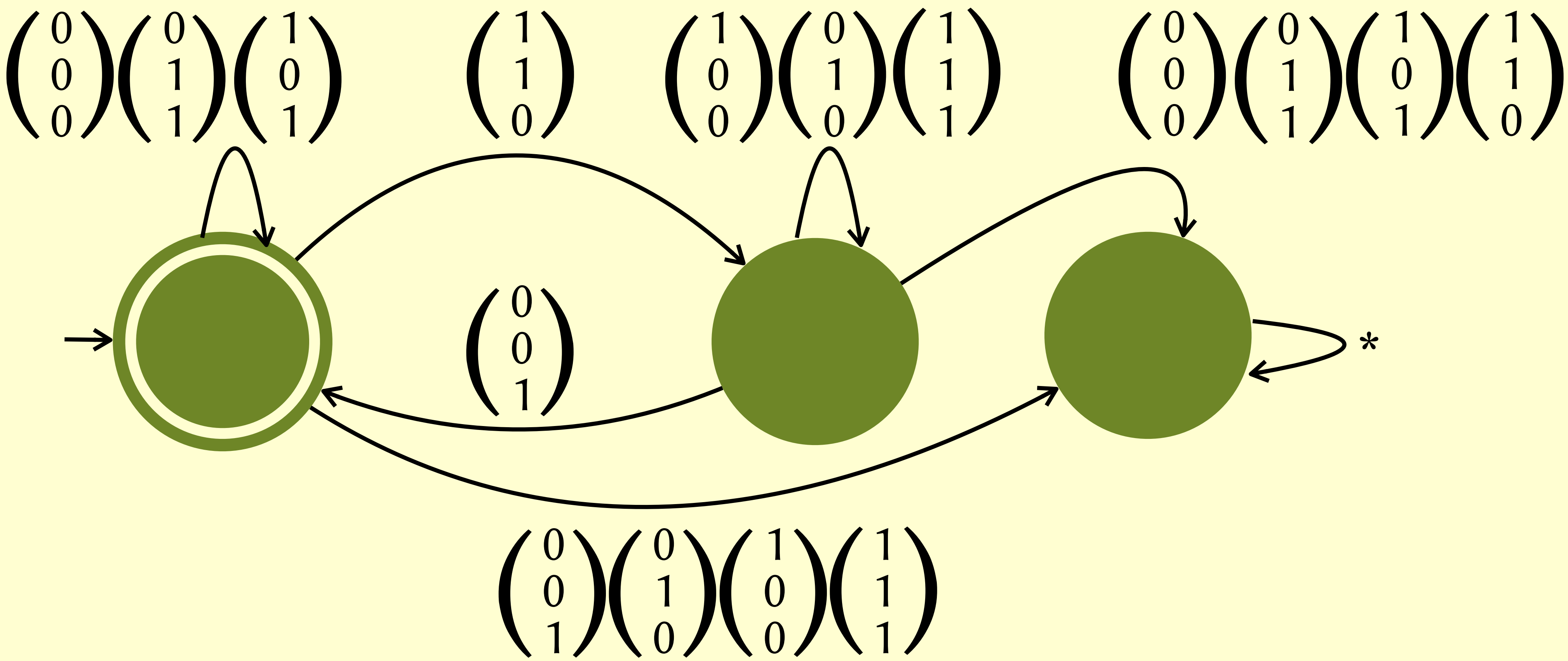
$x + y = z$

	1	2	4	8	16	32	64
x =	0	1	0	0	1	0	0
y =	0	1	0	1	0	1	0
z =	0	0	1	1	1	1	0



Julius Richard Büchi
1924–1984

x + y = z



$\begin{pmatrix}0\\0\\0\end{pmatrix}\begin{pmatrix}0\\1\\1\end{pmatrix}\begin{pmatrix}1\\0\\1\end{pmatrix}$
$\begin{pmatrix}1\\1\\0\end{pmatrix}$
$\begin{pmatrix}1\\0\\0\end{pmatrix}\begin{pmatrix}0\\1\\0\end{pmatrix}\begin{pmatrix}1\\1\\1\end{pmatrix}$
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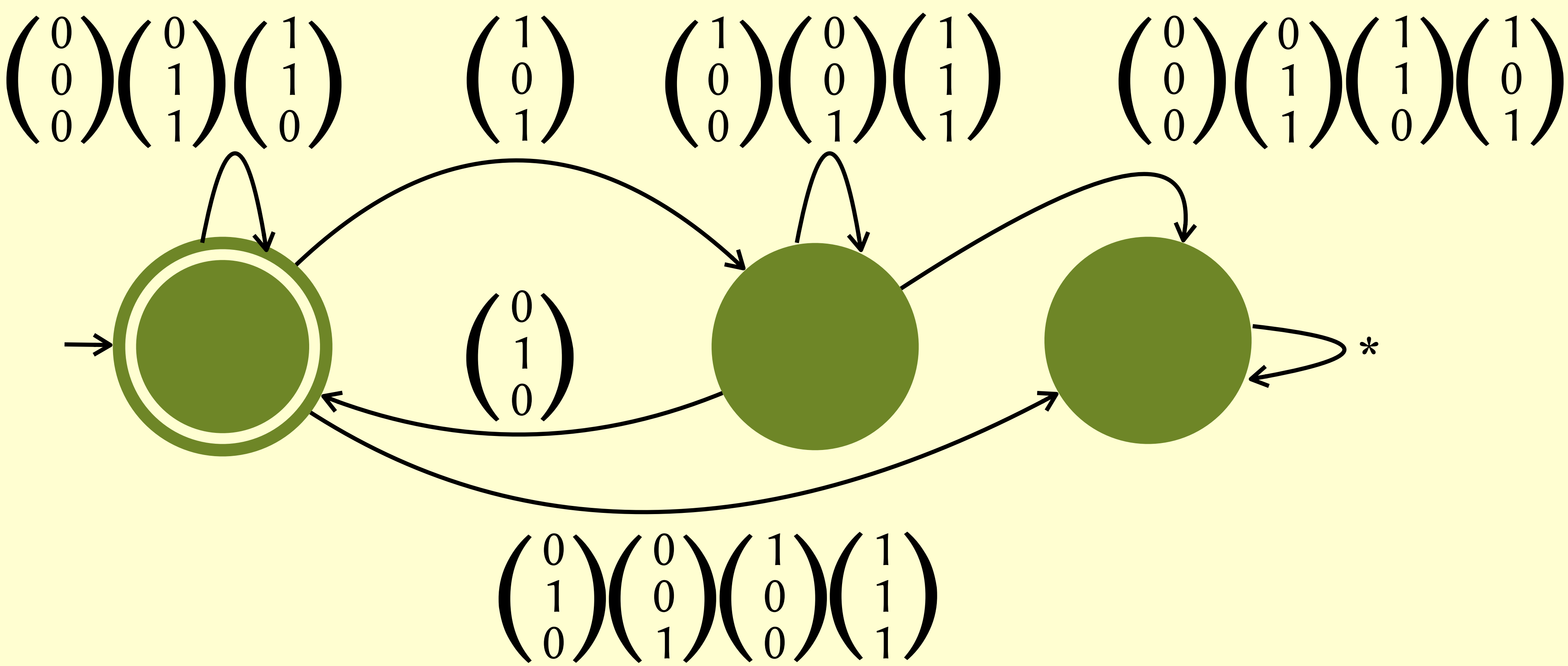
*
$\begin{pmatrix}0\\0\\1\end{pmatrix}\begin{pmatrix}0\\1\\0\end{pmatrix}\begin{pmatrix}1\\0\\0\end{pmatrix}\begin{pmatrix}1\\1\\1\end{pmatrix}$

x + z = y



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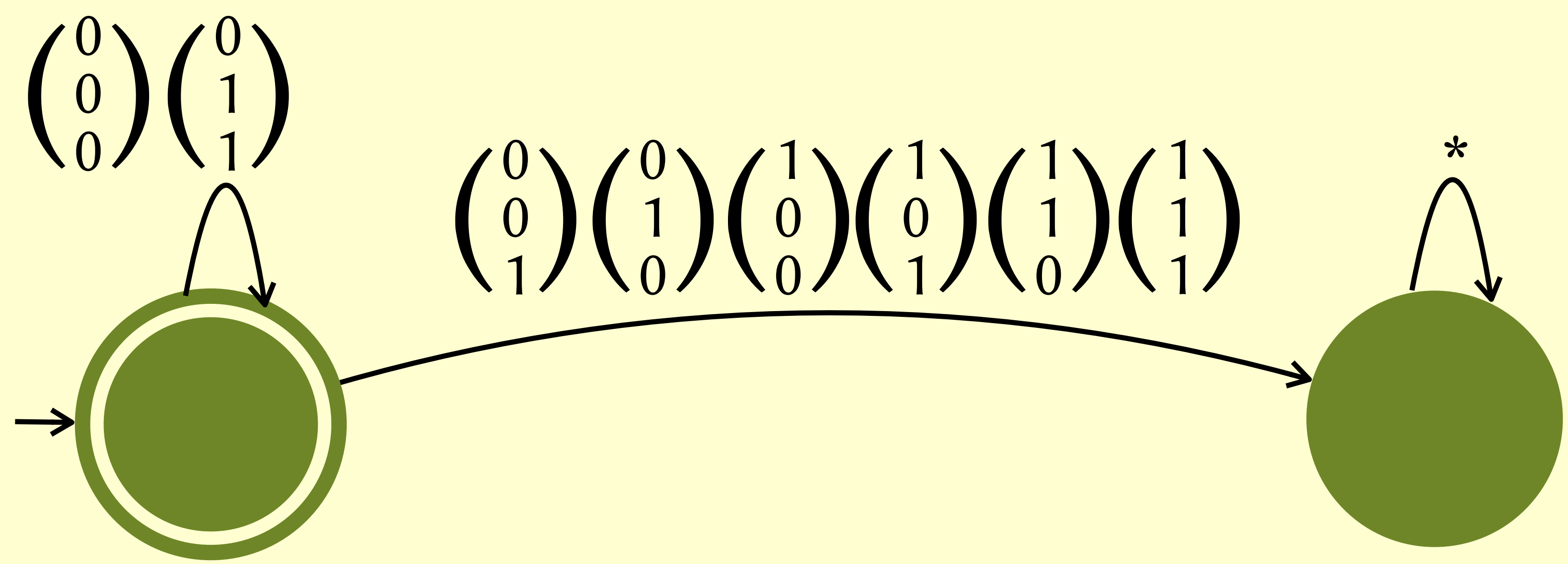
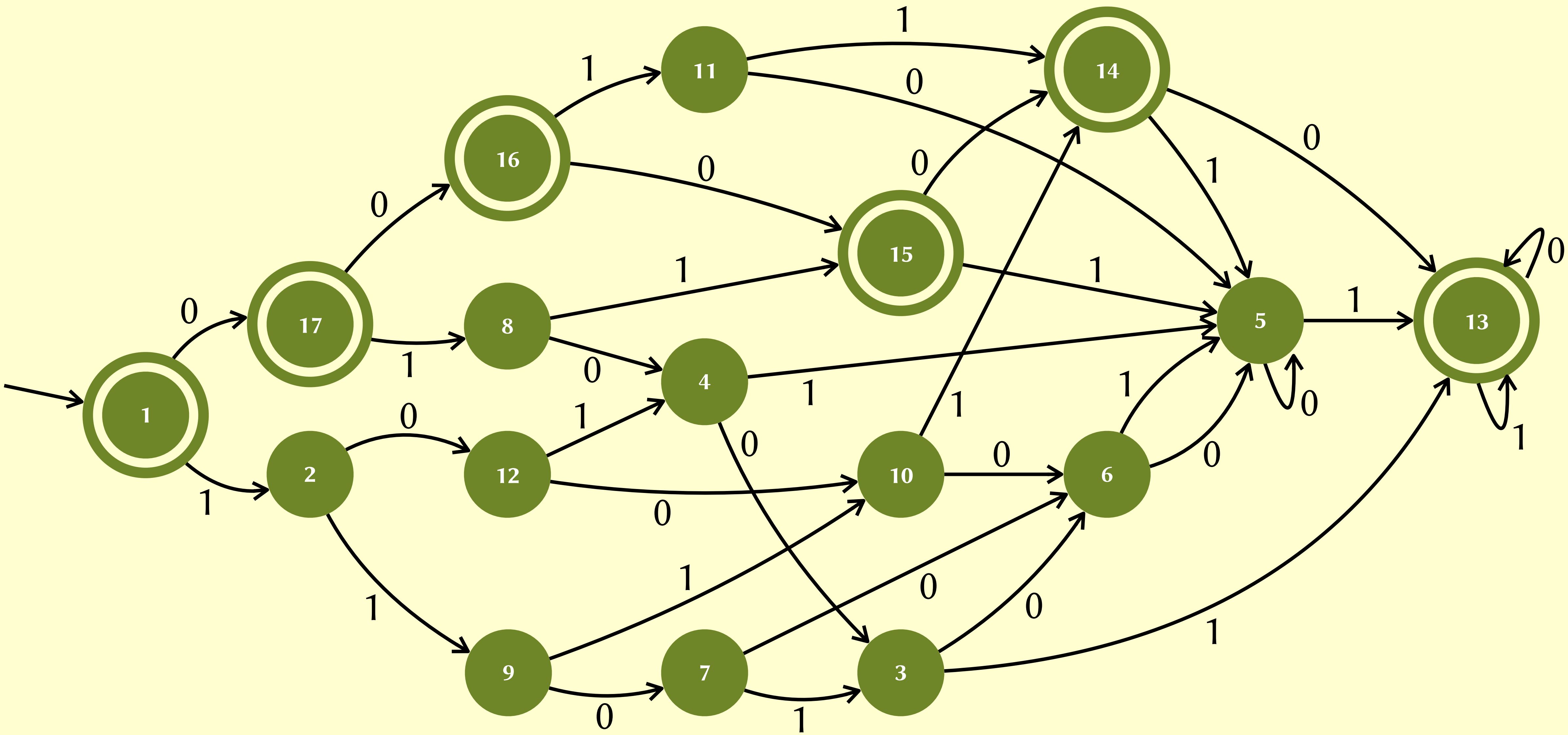
Decidability of Presburger

$$x + y = z \\ \wedge\ x + z = y$$



Decidability of Presburger

$\exists x\ y\ z.\ n = 6x + 9y + 20z$



Adding multiplication

- If we add multiplication, we can write a statement representing the Collatz conjecture: does there exist an infinite sequence of positive integers x_0, x_1, x_2, ... such that

$$
\begin{array}{ll}
2 \times x_{i+1} = x_i & \text{if } x_i \text{ is even} \\
x_{i+1} = 3 \times x_i + 1 & \text{if } x_i \text{ is odd}
\end{array}
$$

- So **if** arithmetic with multiplication were decidable, we could solve the Collatz conjecture automatically!

Automatic proof

- We often rely on automatic provers:
  - e.g. in Dafny, to show that `invariant P` is preserved,
  - e.g. in Isabelle methods like `by auto`.
- How do these automatic provers work?