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# University of Brasilia

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  set ts=4 sw=4 sta nu rnu sc stl+=%F cindent
set ba=dark ruler timeoutlen=1000
set viminfo='20,\"1000
imap {<CR> {<CR>}<Esc>0
nmap \langle F2 \rangle 0V$%d
nmap <C-down> :m+1<CR>
nmap <C-up> :m-2<CR>
nmap <C-a> ggVG
nmap \langle S-up \rangle : m-2 \langle CR \rangle
nmap <S-down> :m+1<CR>
vmap < C-c > "+v
svntax on
set timeoutlen=1000
alias comp='g++ -std=c++17 -Wshadow -Wall -Wextra -
  →Wformat=2 -Wconversion -fsanitize=address.

→undefined -fno-sanitize-recover -Wfatal-errors'

#include <bits/stdc++.h>
#define ff first
#define ss second
#define pb push_back
using namespace std;
using 11 = long long;
using ii = pair<int, int>;
const int N = 100005;
int main() {
 return 0;
1 Data Structures
1.1 Fenwick Tree 2D
vector<int> go[N];
vector<int> ft[N];
void prec_add(int x, int y) {
 for(; x < N; x += x & -x) {
```

```
go[x].push_back(y);
void init() {
  for(int i = 1; i < N; i++) {</pre>
   sort(go[i].begin(), go[i].end());
   ft[i].assign(go[i].size() + 1, 0);
void add(int x, int y, int val) {
 for(; x < N; x += x & -x) {
   int id = int(upper_bound(go[x].begin(), go[x].end
      \hookrightarrow(), y) - go[x].begin());
   for(; id < (int)ft[x].size(); id += id & -id)</pre>
     ft[x][id] += val;
 }
int sum(int x, int y) {
 int ans = 0;
 for(; x > 0; x -= x & -x) {
   int id = int(upper_bound(go[x].begin(), go[x].end
      \hookrightarrow(), y) - go[x].begin());
   for(; id > 0; id -= id & -id)
     ans += ft[x][id];
 return ans;
1.2 Wavelet Tree
template<typename T>
class wavelet { // 1-based!!
   T L, R;
   vector<int> 1;
 vector<T> sum; // <<</pre>
  wavelet *lef, *rig;
  int r(int i) const{ return i - l[i]; }
public:
  template<typename ITER>
   wavelet(ITER bg, ITER en) { // it changes the
      \hookrightarrowargument array
   lef = rig = nullptr;
       L = *bg, R = *bg;
```

for(auto it = bg; it != en; it++)

T mid = L + (R - L)/2;

if(L == R) return;

L = min(L, \*it), R = max(R, \*it);

```
1.reserve(std::distance(bg, en) + 1);
   sum.reserve(std::distance(bg, en) + 1);
   1.push_back(0), sum.push_back(0);
   for(auto it = bg; it != en; it++)
    l.push_back(l.back() + (*it <= mid)),</pre>
     sum.push_back(sum.back() + *it);
   auto tmp = stable_partition(bg, en, [mid](T x){
     return x <= mid:</pre>
   });
   if(bg != tmp) lef = new wavelet(bg, tmp);
   if(tmp != en) rig = new wavelet(tmp, en);
  ~wavelet(){
   delete lef;
   delete rig;
 // 1 index, first is 1st
   T kth(int i, int j, int k) const{
       if(L >= R) return L;
       int c = l[i] - l[i-1];
      if(c \ge k) return lef->kth(l[i-1]+1, l[j], k)
       else return rig->kth(r(i-1)+1, r(j), k - c);
 // # elements > x on [i, j]
 int cnt(int i, int j, T x) const{
   if(L > x) return i - i + 1;
   if(R <= x || L == R) return 0;
   int ans = 0;
   if(lef) ans += lef->cnt(l[i-1]+1, l[j], x);
   if(rig) ans += rig>cnt(r(i-1)+1, r(j), x);
   return ans;
 // sum of elements <= k on [i, i]
 T sumk(int i, int j, T k){
      if(L == R) return R <= k ? L * (j - i + 1) :
   if(R <= k) return sum[j] - sum[i-1];</pre>
   int ans = 0;
   if(lef) ans += lef->sumk(l[i-1]+1, l[j], k);
   if(rig) ans += rig -> sumk(r(i-1)+1, r(j), k);
   return ans:
 }
 // swap (i, i+1) just need to update "array" l[i]
};
1.3 Order Set
```

#include <ext/pb\_ds/assoc\_container.hpp>

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```
#include <ext/pb_ds/tree_policy.hpp>
#include <ext/pb_ds/detail/standard_policies.hpp>
using namespace __gnu_pbds; // or pb_ds;
template<typename T, typename B = null_type>
using oset = tree<T, B, less<T>, rb_tree_tag,
  // find_by_order / order_of_key
```

#### 1.4 Hash table

```
#include <ext/pb_ds/assoc_container.hpp>
using namespace __gnu_pbds;
struct custom_hash {
 static uint64_t splitmix64(uint64_t x) {
   // http://xorshift.di.unimi.it/splitmix64.c
   x += 0x9e3779b97f4a7c15:
   x = (x ^ (x >> 30)) * 0xbf58476d1ce4e5b9;
   x = (x ^(x >> 27)) * 0x94d049bb133111eb;
   return x (x >> 31):
 size_t operator()(uint64_t x) const {
   static const uint64_t FIXED_RANDOM = chrono::

    steady_clock::now().time_since_epoch().count

      \hookrightarrow();
   return splitmix64(x + FIXED_RANDOM);
}:
gp_hash_table<long long, int, custom_hash> table;
unordered_map<long long, int, custom_hash> uhash;
uhash.reserve(1 << 15):</pre>
uhash.max_load_factor(0.25);
```

#### 1.5 Convex Hull Trick Simple

```
struct Line{
   11 m. b:
   inline 11 eval(11 x) const{
      return x * m + b;
};
// min => cht.back().m >= L.m
// max => cht.back().m <= L.m
void push_line(vector<Line> &cht, Line L){
 while((int)cht.size() >= 2){
   int sz = (int)cht.size();
```

```
if((long double)(L.b-cht[sz-1].b)*(cht[sz-2].m-L.
   <= (long double)(L.b-cht[sz-2].b)*(cht[sz-1].m-L.m</pre>
     \hookrightarrow)){
     cht.pop_back();
    else break;
  cht.push_back(L);
// x increasing; pos = 0 in first call
11 linear_search(const vector<Line> &cht,ll x,int &
   \hookrightarrowpos){
   while(pos+1 < (int)cht.size()){</pre>
/*>>*/ if(cht[pos].eval(x) >= cht[pos+1].eval(x))
  ⇒pos++:
       else break;
   return cht[pos].eval(x);
}
11 binary_search(const vector<Line> &cht, ll x){
    int L = 0, R = (int)cht.size()-2;
   int bans = (int)cht.size()-1;
    while(L <= R){</pre>
       int mid = (L+R)/2;
       if(cht[mid].eval(x) >= cht[mid+1].eval(x)) //
          \hookrightarrow <<<
           L = mid + 1:
       else bans = mid, R = mid - 1;
   return cht[bans].eval(x);
}
```

#### 1.6 Convex Hull Trick

```
const ll is_query = -(1LL<<62);</pre>
struct Line{
 11 m, b;
 mutable function<const Line*()> succ;
  bool operator<(const Line& rhs) const{</pre>
   if(rhs.b != is_query) return m < rhs.m;</pre>
   const Line* s = succ();
   if(!s) return 0;
   11 x = rhs.m:
   return b - s -> b < (s -> m - m) * x:
};
struct Cht : public multiset<Line>{ // maintain max
 bool bad(iterator y){
```

```
auto z = next(y);
   if(y == begin()){
     if(z == end()) return 0;
     return y->m == z->m && y->b <= z->b;
    auto x = prev(y);
   if(z == end()) return y->m == x->m && y->b <= x->
   return (long double) (x->b - y->b)*(z->m - y->m)
      \Rightarrow = (long double)(y->b - z->b)*(y->m - x->m);
 void insert_line(ll m, ll b){
    auto y = insert({ m, b });
   y->succ = [=]{ return next(y) == end() ? 0 : &*
      \hookrightarrownext(y); };
   if(bad(y)){ erase(y); return; }
    while(next(y) != end() && bad(next(y))) erase(
      \hookrightarrownext(y));
   while(y != begin() && bad(prev(y))) erase(prev(y)
      \hookrightarrow):
 }
 ll eval(ll x){
   auto 1 = *lower_bound((Line) { x, is_query });
   return 1.m * x + 1.b;
 }
};
```

#### 1.7 Convex Hull Trick

```
* Author: Simon Lindholm
 * source: https://github.com/kth-competitive-
   →programming/kactl/blob/master/content/data-
   \hookrightarrowstructures/LineContainer.h
 * License: CC0
struct Line {
 mutable 11 m, b, p;
 bool operator<(const Line& o) const { return m < o</pre>
 bool operator<(11 x) const { return p < x; }</pre>
};
struct LineContainer : multiset<Line, less<>> { //
  \hookrightarrow CPP14 only
 // (for doubles, use inf = 1/.0, div(a,b) = a/b)
 const ll inf = LLONG_MAX;
 11 div(ll a, ll b) { // floored division
   return a / b - ((a ^ b) < 0 && a % b); }
 bool isect(iterator x, iterator y) {
```

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```
if (y == end()) { x->p = inf; return false; }
   if (x->m == y->m) x->p = x->b > y->b ? inf : -inf
     \hookrightarrow
   else x->p = div(y->b - x->b, x->m - y->m);
   return x->p >= y->p;
 void add(ll m, ll b) {
   auto z = insert(\{m, b, 0\}), y = z++, x = y;
   while (isect(y, z)) z = erase(z);
   if (x != begin() \&\& isect(--x, y)) isect(x, y =
      \hookrightarrowerase(y));
   while ((y = x) != begin() && (--x)->p >= y->p)
     isect(x, erase(y));
 11 query(ll x) {
   assert(!empty());
   auto 1 = *lower_bound(x);
   return 1.m * x + 1.b;
 }
};
1.8 Min queue
template<typename T>
class minQ{
 deque<tuple<T, int, int> > p;
 T delta:
 int sz:
public:
 minQ() : delta(0), sz(0) {}
 inline int size() const{ return sz; }
 inline void add(T x){ delta += x; }
 inline void push(T x, int id){
   x \rightarrow delta, sz++;
   int t = 1;
   while(p.size() > 0 && get<0>(p.back()) >= x)
     t += get<1>(p.back()), p.pop_back();
   p.emplace_back(x, t, id);
 inline void pop(){
   get<1>(p.front())--, sz--;
   if(!get<1>(p.front())) p.pop_front();
 T getmin() const{ return get<0>(p.front())+delta;
 int getid() const{ return get<2>(p.front()); }
};
1.9 Sparse Table
int fn(int i, int j){
```

if(j == 0) return v[i];

```
if(~dn[i][j]) return dn[i][j];
 \hookrightarrow-1)), j-1));
int getmn(int 1, int r){ // [1, r]
 int 1z = 1g(r - 1 + 1);
 return min(fn(1, lz), fn(r - (1 << lz) + 1, lz));
1.10 Treap
// source: https://github.com/victorsenam/caderno/
  ⇒blob/master/code/treap.cpp
//const int N = : typedef int num:
num X[N]; int en = 1, Y[N], sz[N], L[N], R[N];
void calc (int u) { // update node given children
  \hookrightarrow info
 if(!u) return;
  sz[u] = sz[L[u]] + 1 + sz[R[u]];
 // code here, no recursion
void unlaze (int u) {
 if(!u) return;
 // code here, no recursion
void split_val(int u, num x, int &l, int &r) { // 1
  \hookrightarrow gets <= x, r gets > x
  unlaze(u); if(!u) return (void) (1 = r = 0);
 if(X[u] \le x) \{ split_val(R[u], x, 1, r); R[u] = 1 \}
    \hookrightarrow; 1 = u; }
  else { split_val(L[u], x, l, r); L[u] = r; r = u;
    \hookrightarrow}
  calc(u);
void split_sz(int u, int s, int &l, int &r) { // 1
  \hookrightarrow gets first s, r gets remaining
  unlaze(u); if(!u) return (void) (1 = r = 0);
  if(sz[L[u]] < s)  { split_sz(R[u], s - sz[L[u]] -
    \hookrightarrow 1, 1, r); R[u] = 1; 1 = u; }
  else { split_sz(L[u], s, l, r); L[u] = r; r = u; }
  calc(u);
int merge(int 1, int r) { // els on 1 <= els on r</pre>
  unlaze(l); unlaze(r); if(!l || !r) return l + r;
    ⇒int u:
  if(Y[1] > Y[r]) { R[1] = merge(R[1], r); u = 1; }
  else { L[r] = merge(1, L[r]); u = r; }
  calc(u); return u;
```

```
void init(int n=N-1) { // XXX call before using
  →other funcs
 for(int i = en = 1; i \le n; i++) { Y[i] = i; sz[i]
    \hookrightarrow = 1; L[i] = R[i] = 0; }
 random_shuffle(Y + 1, Y + n + 1);
void insert(int &u, int it){
 unlaze(u);
 if(!u) u = it;
 else if(Y[it] > Y[u]) split_val(u, X[it], L[it], R
    \hookrightarrow[it]), u = it;
 else insert(X[it] < X[u] ? L[u] : R[u], it);
 calc(u);
void erase(int &u, num key){
 unlaze(u);
 if(!u) return:
 if(X[u] == key) u = merge(L[u], R[u]);
 else erase(key < X[u] ? L[u] : R[u], key);
 calc(u):
int create_node(num key){
 X[en] = key;
 sz[en] = 1;
 L[en] = R[en] = 0;
 return en++;
int query(int u, int 1, int r)\{//0 \text{ index}\}
 unlaze(u);
 if(u! or r < 0 or 1 >= sz[u]) return
    →identity_element;
 if(1 \le 0 \text{ and } r >= sz[u] - 1) \text{ return } subt_data[u];
 int ans = query(L[u], 1, r);
 if(1 \le sz[L[u]] and sz[L[u]] \le r)
   ans = max(ans, st[u]);
 ans = max(ans, query(R[u], 1-sz[L[u]]-1, r-sz[L[u]])
    \hookrightarrow]]-1));
 return ans:
1.11 ColorUpdate
// source: https://github.com/tfg50/Competitive-
  → Programming/tree/master/Biblioteca/Data%20
  \hookrightarrowStructures
#include <set>
#include <vector>
template <class Info = int>
class ColorUpdate {
```

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```
public:
 struct Range {
   Range(int l = 0) { this->l = 1; }
   Range(int 1, int r, Info v) {
     this->1 = 1;
     this->r = r;
     this->v = v;
   int 1, r;
   Info v;
   bool operator < (const Range &b) const { return 1</pre>
      \hookrightarrow < b.1; }
 };
 std::vector<Range> upd(int 1, int r, Info v) {
   std::vector<Range> ans;
   if(1 >= r) return ans;
   auto it = ranges.lower_bound(1);
   if(it != ranges.begin()) {
     it--;
     if(it->r > 1) {
       auto cur = *it;
       ranges.erase(it);
       ranges.insert(Range(cur.1, 1, cur.v));
       ranges.insert(Range(1, cur.r, cur.v));
     }
   it = ranges.lower_bound(r);
   if(it != ranges.begin()) {
     it--;
     if(it->r > r) {
       auto cur = *it:
       ranges.erase(it);
       ranges.insert(Range(cur.1, r, cur.v));
       ranges.insert(Range(r, cur.r, cur.v));
     }
   for(it = ranges.lower_bound(1); it != ranges.end
      \hookrightarrow() && it->l < r; it++) {
     ans.push_back(*it);
   ranges.erase(ranges.lower_bound(1), ranges.
      \hookrightarrowlower_bound(r));
   ranges.insert(Range(1, r, v));
   return ans;
private:
 std::set<Range> ranges;
```

# 1.12 Heavy Light Decomposition

```
void dfs_sz(int u){
   sz[u] = 1;
   for(auto &v : g[u]) if(v == p[u]){
       swap(v, g[u].back()); g[u].pop_back();
       break:
   for(auto &v : g[u]){
       p[v] = u; dfs_sz(v); sz[u] += sz[v];
       if(sz[v] > sz[g[u][0]])
          swap(v, g[u][0]);
   }
// nxt[u] = start of path with u
// set nxt[root] = root beforehand
void dfs_hld(int u){
   in[u] = t++;
   rin[in[u]] = u;
   for(auto v : g[u]){
       nxt[v] = (v == g[u][0] ? nxt[u] : v); dfs_hld
   out[u] = t;
// subtree of u => [ in[u], out[u] )
// path from nxt[u] to u \Rightarrow [in[nxt[u]], in[u]]
1.13 Iterative Segtree
T query(int 1, int r){ // [1, r]
   T rl, rr;
   for(1 += n, r += n+1; 1 < r; 1 >>= 1, r >>= 1){
       if(l & 1) rl = merge(rl, st[l++]);
       if(r & 1) rr = merge(st[--r], rr);
   return merge(rl, rr);
}
// initially save v[i] in st[n+i] for all i in [0, n
  \hookrightarrow )
void build(){
   for(int p = n-1; p > 0; p--)
       st[p] = merge(st[2*p], st[2*p+1]);
void update(int p, T val){
   st[p += n] = val;
   while(p >>= 1) st[p] = merge(st[2*p], st[2*p+1]);
```

#### 1.14 Recursive Segtree + lazy

```
class SegTree{
   vi st;
   vi lazy;
   int size:
   int el_neutro = -oo;
   inline int f(int a, int b){
      return max(a,b);
   inline int left(int i) {return 2 * i + 1;};
   inline int right(int i) {return 2 * i + 2;};
   void build(int sti, int stl, int str, vi& nums) {
      if(stl == str) {
          st[sti] = nums[stl];
          return:
       int mid = (stl + str) / 2;
       build(left(sti), stl, mid, nums);
      build(right(sti), mid + 1, str, nums);
       st[sti] = f(st[left(sti)], st[right(sti)]);
   }
   void propagate(int sti, int stl, int str){
      if(lazy[sti]){
          st[sti] += lazy[sti];
          if(stl != str)
             lazy[left(sti)] += lazy[sti];
             lazy[right(sti)] += lazy[sti];
          lazy[sti] = 0;
      }
   int query(int sti, int stl, int str, int l, int r
      ) {
      propagate(sti, stl, str);
      if(str < 1 || r < stl)
          return el_neutro;
      if(stl >= 1 and str <= r)
          return st[sti];
      int mid = (str+st1)/2;
      return f(query(left(sti),stl,mid,l,r),query(

→right(sti),mid+1,str,l,r));
```

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```
void update_range(int sti, int stl, int str, int
       \hookrightarrow1,int r, int amm){
       propagate(sti. stl. str):
       if(stl >= 1 and str <= r)
           lazy[sti] = amm:
           propagate(sti, stl, str);
           return;
       if(stl > r or str < 1)
           return:
       int mid = (stl + str)/2:
       update_range(left(sti),stl,mid,l,r,amm);
       update_range(right(sti),mid+1,str,l,r,amm);
       st[sti] = f(st[left(sti)],st[right(sti)]);
   void update(int sti, int stl, int str, int i, int
      \hookrightarrow amm){
       propagate(sti, stl, str);
       if(stl == i and str == i){
           st[sti] = amm;
           return;
       }
       if(stl > i or str < i)</pre>
           return:
       int mid = (stl + str)/2;
       update(left(sti),stl,mid,i,amm);
       update(right(sti),mid+1,str,i,amm);
       st[sti] = f(st[left(sti)],st[right(sti)]);
   public:
       SegTree(vi& v) : st(4*v.size(),0), lazy(4*v.
          \hookrightarrowsize(),0) {size = v.size(); build(0,0,
          \hookrightarrowsize - 1, v);}
       SegTree(int n) : st(4*n,0), lazy(4*n,0){size
          \hookrightarrow = n;
       int query(int 1, int r){return query(0,0,size
          \hookrightarrow-1.1.r):}
       void update_range(int 1, int r, int amm){
          \hookrightarrowupdate_range(0,0,size-1,1,r,amm);}
       void update(int i, int amm){update(0,0,size
          \hookrightarrow-1,i,amm);}
};
1.15 LiChao's Segtree
void add_line(line nw, int v = 1, int l = 0, int r = 1
   \hookrightarrow maxn) { // \lceil 1, r \rceil
   int m = (1 + r) / 2;
   bool lef = nw.eval(1) < st[v].eval(1);</pre>
   bool mid = nw.eval(m) < st[v].eval(m);</pre>
```

```
if(mid) swap(st[v], nw);
   if(r - 1 == 1) {
       return;
   } else if(lef != mid) {
       add_line(nw, 2 * v, 1, m);
   } else {
       add_line(nw, 2 * v + 1, m, r);
   }
}
int get(int x, int v = 1, int l = 0, int r = maxn) {
   int m = (1 + r) / 2;
   if(r - 1 == 1) {
       return st[v].eval(x);
   \} else if(x < m) {
       return min(st[v].eval(x), get(x, 2*v, 1, m));
       return min(st[v].eval(x), get(x, 2*v+1, m, r)
         \hookrightarrow);
   }
1.16 Palindromic tree
#include <bits/stdc++.h>
using namespace std;
const int maxn = 3e5 + 1, sigma = 26;
int len[maxn], link[maxn], to[maxn][sigma];
int slink[maxn], diff[maxn], series_ans[maxn];
int sz, last, n;
char s[maxn];
void init()
   s[n++] = -1;
   link[0] = 1;
   len[1] = -1;
   sz = 2;
int get_link(int v)
   while(s[n - len[v] - 2] != s[n - 1]) v = link[v];
   return v:
void add_letter(char c)
{
   s[n++] = c -= 'a';
```

```
last = get_link(last);
   if(!to[last][c])
       len[sz] = len[last] + 2;
       link[sz] = to[get_link(link[last])][c];
       diff[sz] = len[sz] - len[link[sz]];
       if(diff[sz] == diff[link[sz]])
           slink[sz] = slink[link[sz]];
       else
           slink[sz] = link[sz];
       to[last][c] = sz++;
   last = to[last][c];
int main()
   ios::sync_with_stdio(0);
   cin.tie(0);
   init():
   string s;
   cin >> s;
   int n = s.size();
   int ans[n + 1];
   memset(ans, 63, sizeof(ans));
   ans [0] = 0;
   for(int i = 1; i \le n; i++)
       add_letter(s[i - 1]);
       for(int v = last; len[v] > 0; v = slink[v])
           series_ans[v] = ans[i - (len[slink[v]] +
             \hookrightarrowdiff[v])];
           if(diff[v] == diff[link[v]])
              series_ans[v] = min(series_ans[v],
                 \hookrightarrow series_ans[link[v]]);
           ans[i] = min(ans[i], series_ans[v] + 1);
       cout << ans[i] << "\n";</pre>
   return 0;
2 Math
```

#### 2.1 Extended Euclidean Algorithm

```
// a*x + b*y = gcd(a, b), <gcd, x, y>
tuple<int, int, int> gcd(int a, int b) {
  if(b == 0) return make_tuple(a, 1, 0);
  int q, w, e;
```

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```
tie(q, w, e) = gcd(b, a \% b);
return make_tuple(q, e, w - e * (a / b));
```

#### 2.2 Chinese Remainder Theorem

```
// x = vet[i].first (mod vet[i].second)
11 crt(const vector<pair<11, 11>> &vet){
   11 \text{ ans} = 0, 1cm = 1;
   ll a, b, g, x, y;
   for(const auto &p : vet) {
       tie(a, b) = p;
       tie(g, x, y) = gcd(lcm, b);
       if((a - ans) \% q != 0) return -1; // no
          \hookrightarrowsolution
       ans = ans + x * ((a - ans) / g) % (b / g) *
          \hookrightarrow1cm;
       lcm = lcm * (b / g);
       ans = (ans \% lcm + lcm) \% lcm;
   return ans;
```

#### 2.3 Diophantine Solver

```
template<typename T>
T extgcd(T a, T b, T &x, T &y) {
 if (a == 0) {
   x = 0;
   y = 1;
   return b;
 T p = b / a;
 T g = extgcd(b - p * a, a, y, x);
 x -= p * y;
 return g;
template<typename T>
bool diophantine(T a, T b, T c, T &x, T &y, T &g) {
 if (a == 0 && b == 0) {
   if (c == 0) {
     x = y = g = 0;
     return true:
   return false;
 if (a == 0) {
   if (c % b == 0) {
     x = 0;
     y = c / b;
     g = abs(b);
```

```
return true:
 }
 return false:
if (b == 0) {
 if (c % a == 0) {
   x = c / a;
   y = 0;
   g = abs(a);
   return true;
 return false:
g = extgcd(a, b, x, y);
if (c % g != 0) {
 return false;
T dx = c / a;
c -= dx * a:
T dv = c / b:
c -= dy * b;
x = dx + mulmod(x, c / g, b);
y = dy + mulmod(y, c / g, a);
q = abs(q);
return true;
```

#### 2.4 Preffix inverse

```
inv[1] = 1;
for(int i = 2; i < p; i++)
 inv[i] = (p - (p/i) * inv[p%i] % p) % p;
```

#### 2.5 Pollard Rho

```
ll rho(ll n){
 if(n \% 2 == 0) return 2;
 11 d, c, x, y, prod;
   c = llrand(1, n - 1);
   x = 11rand(1, n - 1);
   y = x;
   prod = 1;
   for(int i = 0; i < 40; i++) {
     x = add(mul(x, x, n), c, n);
     y = add(mul(y, y, n), c, n);
     y = add(mul(y, y, n), c, n);
     prod = mul(prod, abs(x - y), n) ?: prod;
   d = \_gcd(prod, n);
 } while(d == 1);
 return d;
```

```
11 pollard_rho(ll n){
 11 x, c, y, d, k;
 int i;
 do{
   i = 1;
   x = 11rand(1, n-1), c = 11rand(1, n-1);
   y = x, k = 4;
   do{
    if(++i == k) y = x, k *= 2;
    x = add(mul(x, x, n), c, n);
    d = \_gcd(abs(x - y), n);
   }while(d == 1);
 }while(d == n);
 return d;
void factorize(ll val, map<ll, int> &fac){
 if(rabin(val)) fac[ val ]++;
 else{
   11 d = pollard_rho(val);
   factorize(d. fac):
   factorize(val / d, fac);
 }
map<ll, int> factor(ll val){
 map<ll, int> fac;
 if(val > 1) factorize(val, fac);
 return fac:
```

#### 2.6 Miller Rabin

return 1;

```
bool rabin(ll n){
 if(n \ll 1) return 0;
 if(n <= 3) return 1;
 11 s = 0, d = n - 1;
 while(d % 2 == 0) d /= 2, s++;
 for(int k = 0; k < 64; k++){
   11 a = 11rand(2, n-2);
   11 x = fexp(a, d, n);
   if(x != 1 \&\& x != n-1){
     for(int r = 1; r < s; r++){
      x = mul(x, x, n);
      if(x == 1) return 0;
      if(x == n-1) break:
     if(x != n-1) return 0;
 }
```

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```
2.7 Totiente
11 totiente(ll n){
 11 \text{ ans} = n:
 for(11 i = 2; i*i \le n; i++)
   if(n \% i == 0){
     ans = ans / i * (i - 1);
     while(n % i == 0) n /= i;
   }
 if(n > 1) ans = ans / n * (n - 1);
 return ans;
2.8 Primitive root
// a primitive root modulo n is any number g such
  \hookrightarrow of g modulo n.
bool exists root(ll n){
   if(n == 1 || n == 2 || n == 4) return true;
   if(n \% 2 == 0) n /= 2;
   if(n % 2 == 0) return false;
   // test if n is a power of only one prime
   for(11 i = 3; i * i <= n; i += 2) if(n % i == 0){
      while(n % i == 0) n /= i;
      return n == 1;
   }
   return true;
11 primitive_root(ll n){
   if(n == 1 || n == 2 || n == 4) return n - 1;
   if(not exists_root(n)) return -1;
   11 x = phi(n);
   auto pr = factorize(x);
   auto check = [x, n, pr](11 m){
      for(11 p : pr) if(fexp(m, x / p, n) == 1)
          return false;
      return true;
   for(11 m = 2; ; m++) if(\_gcd(m, n) == 1)
      if(check(m)) return m;
}
// Let's denote R(n) as the set of primitive roots
  \hookrightarrow modulo n, p is prime
// g \setminus in R(p) \Rightarrow (pow(g, p-1, p * p) == 1 ? g+p : g)
```

```
// g \text{ in } R(pow(p, k)) \Rightarrow (g \% 2 == 1 ? g : g + pow(p, k))
  \hookrightarrow k)) \in R(2*pow(p, k))
2.9 Mobius Function
memset(mu, 0, sizeof mu);
mu[1] = 1;
for(int i = 1; i < N; i++)
   for(int j = i + i; j < N; j += i)
       mu[i] -= mu[i];
// g(n) = sum\{f(d)\} \Rightarrow f(n) = sum\{mu(d)*g(n/d)\}
2.10 Mulmod TOP
constexpr uint64_t mod = (1ull<<61) - 1;</pre>
uint64_t modmul(uint64_t a, uint64_t b){
  uint64_t 11 = (uint32_t)a, h1 = a>>32, 12 = (
    \hookrightarrowuint32_t)b, h2 = b>>32;
  uint64_t l = 11*12, m = 11*h2 + 12*h1, h = h1*h2;
  uint64_t ret = (1\&mod) + (1>>61) + (h << 3) + (m
    \implies 29) + (m << 35 >> 3) + 1;
  ret = (ret \& mod) + (ret >> 61);
  ret = (ret \& mod) + (ret >> 61);
 return ret-1;
2.11 Modular multiplication TOPPER
11 mulmod(11 a, 11 b, 11 mod) {
   11 q = 11((long double)a * (long double)b / (long
      \hookrightarrow double)mod);
   11 r = (a * b - mod * q) \% mod;
   if(r < 0) r += mod;
   return r;
2.12 Division Trick
for(int l = 1, r; l \le n; l = r + 1) {
   r = n / (n / 1);
   // n / x yields the same value for 1 <= x <= r
for(int 1, r = n; r > 0; r = 1 - 1) {
   int tmp = (n + r - 1) / r;
   1 = (n + tmp - 1) / tmp;
   // (n+x-1) / x yields the same value for 1 <= x
      ∽<= r
2.13 Matrix Determinant
int n:
long double a[n][n];
long double gauss(){
   long double det = 1;
```

```
for(int i = 0; i < n; i++){
       int q = i;
       for(int j = i+1; j < n; j++){
          if(abs(a[j][i]) > abs(a[q][i]))
       if(abs(a[q][i]) < EPS){
          det = 0;
          break;
      }
      if(i != q){
          for(int w = 0; w < n; w++)
              swap(a[i][w], a[q][w]);
          det = -det;
       det *= a[i][i];
       for(int j = i+1; j < n; j++) a[i][j] /= a[i][
         \hookrightarrowi]:
       for(int j = 0; j < n; j++) if(j != i){
          if(abs(a[j][i]) > EPS)
              for(int k = i+1; k < n; k++)
                 a[j][k] -= a[i][k] * a[j][i];
      }
   return det;
2.14 Simplex Method
typedef long double dbl;
const dbl eps = 1e-6;
const int N = , M = ;
mt19937 rng(chrono::steady_clock::now().
  →time_since_epoch().count());
struct simplex {
 int X[N], Y[M];
 dbl A[M][N], b[M], c[N];
 dbl ans:
 int n, m;
 dbl sol[N];
 void pivot(int x, int y){
   swap(X[y], Y[x]);
   b[x] /= A[x][y];
   for(int i = 0; i < n; i++)
    if(i != y)
      A[x][i] /= A[x][y];
   A[x][y] = 1. / A[x][y];
```

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```
for(int i = 0; i < m; i++)
   if(i != x \&\& abs(A[i][y]) > eps) {
     b[i] -= A[i][v] * b[x];
     for(int j = 0; j < n; j++) if(j != y)
                A[i][j] -= A[i][y] * A[x][j];
     A[i][y] = -A[i][y] * A[x][y];
   }
 ans += c[y] * b[x];
 for(int i = 0; i < n; i++)
   if(i != y)
     c[i] -= c[y] * A[x][i];
 c[y] = -c[y] * A[x][y];
// maximiza sum(x[i] * c[i])
// sujeito a
// sum(a[i][j] * x[j]) <= b[i] para 0 <= i < m (Ax)
  \hookrightarrow \langle = b \rangle
// x[i] >= 0 para 0 <= i < n (x >= 0)
// (n variaveis, m restricoes)
// guarda a resposta em ans e retorna o valor
  \hookrightarrowotimo
dbl solve(int _n, int _m) {
 this->n = _n; this->m = _m;
     for(int i = 1; i < m; i++){
         int id = uniform_int_distribution<int>(0,
            \hookrightarrowi)(rng);
         swap(b[i], b[id]);
         for(int j = 0; j < n; j++)
            swap(A[i][j], A[id][j]);
     }
 ans = 0.;
 for(int i = 0; i < n; i++) X[i] = i;
 for(int i = 0; i < m; i++) Y[i] = i + n;
 while(true) {
   int x = min element(b. b + m) - b:
   if(b[x] >= -eps)
     break;
   int y = find_if(A[x], A[x] + n, [](dbl d) {
      \hookrightarrow return d < -eps; }) - A[x];
   if(y == n) throw 1; // no solution
   pivot(x, y);
 while(true) {
   int y = max_element(c, c + n) - c;
   if(c[y] <= eps) break;</pre>
   int x = -1;
```

```
dbl mn = 1. / 0.;
     for(int i = 0; i < m; i++)
       \mathbf{if}(A[i][y] > \text{eps \&\& b[i]} / A[i][y] < \text{mn})
         mn = b[i] / A[i][y], x = i;
     if(x == -1) throw 2; // unbounded
     pivot(x, y);
   memset(sol, 0, sizeof(dbl) * n);
   for(int i = 0; i < m; i++)
     if(Y[i] < n)
       sol[Y[i]] = b[i];
   return ans:
 }
};
2.15 FFT
void fft(vector<base> &a, bool inv){
   int n = (int)a.size();
   for(int i = 1, j = 0; i < n; i++){
       int bit = n \gg 1;
       for(; j >= bit; bit >>= 1) j -= bit;
       j += bit;
       if(i < j) swap(a[i], a[i]);
   for(int sz = 2; sz <= n; sz <<= 1) {
       double ang = 2 * PI / sz * (inv ? -1 : 1);
       base wlen(cos(ang), sin(ang));
       for(int i = 0: i < n: i += sz){
           base w(1, 0);
           for(int j = 0; j < sz / 2; j++){
              base u = a[i+j], v = a[i+j + sz/2] * w
                \hookrightarrow ;
              a[i+j] = u + v;
              a[i+j+sz/2] = u - v;
              w *= wlen:
   if(inv) for(int i = 0; i < n; i++) a[i] /= 1.0 *
      \hookrightarrown;
2.16 FFT Tourist
namespace fft {
  typedef double dbl;
  struct num {
   dbl x, y;
```

```
num() \{ x = y = 0; \}
 num(dbl x, dbl y) : x(x), y(y) {}
inline num operator+(num a, num b) { return num(a.
  \hookrightarrow x + b.x, a.y + b.y; }
inline num operator-(num a, num b) { return num(a.
  \hookrightarrowx - b.x, a.y - b.y); }
inline num operator*(num a, num b) { return num(a.
  \rightarrowx * b.x - a.y * b.y, a.x * b.y + a.y * b.x); }
inline num conj(num a) { return num(a.x, -a.y); }
int base = 1;
vector<num> roots = \{\{0, 0\}, \{1, 0\}\};
vector<int> rev = {0, 1};
const dbl PI = acosl(-1.0);
void ensure_base(int nbase) {
 if(nbase <= base) return:</pre>
 rev.resize(1 << nbase);</pre>
  for(int i = 0; i < (1 << nbase); i++) {
   rev[i] = (rev[i >> 1] >> 1) + ((i & 1) << (
      \hookrightarrownbase - 1));
 roots.resize(1 << nbase);</pre>
  while(base < nbase) {</pre>
   dbl \ angle = 2*PI / (1 << (base + 1)):
    for(int i = 1 \ll (base - 1); i < (1 \ll base); i
      →++) {
     roots[i << 1] = roots[i];</pre>
     dbl \ angle_i = angle * (2 * i + 1 - (1 << base)
     roots[(i \ll 1) + 1] = num(cos(angle_i), sin(
        \rightarrowangle i)):
   }
   base++:
void fft(vector<num> &a, int n = -1) {
 if(n == -1) {
   n = a.size():
  assert((n & (n-1)) == 0);
 int zeros = __builtin_ctz(n);
  ensure_base(zeros);
```

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```
int shift = base - zeros:
 for(int i = 0; i < n; i++) {
   if(i < (rev[i] >> shift)) {
     swap(a[i], a[rev[i] >> shift]);
   }
 for(int k = 1; k < n; k <<= 1) {
   for(int i = 0; i < n; i += 2 * k) {
     for(int j = 0; j < k; j++) {
      num z = a[i+j+k] * roots[j+k];
       a[i+j+k] = a[i+j] - z;
      a[i+j] = a[i+j] + z;
     }
   }
vector<num> fa, fb;
vector<int> multiply(vector<int> &a, vector<int> &
 int need = a.size() + b.size() - 1;
 int nbase = 0;
 while((1 << nbase) < need) nbase++;</pre>
 ensure_base(nbase);
 int sz = 1 << nbase:</pre>
 if(sz > (int) fa.size()) {
   fa.resize(sz);
 for(int i = 0; i < sz; i++) {
   int x = (i < (int) a.size() ? a[i] : 0):
   int y = (i < (int) b.size() ? b[i] : 0);
   fa[i] = num(x, y);
 fft(fa, sz);
 num r(0, -0.25 / sz);
 for(int i = 0; i \le (sz >> 1); i++) {
   int j = (sz - i) & (sz - 1);
   num z = (fa[i] * fa[i] - conj(fa[i] * fa[i])) *
      \hookrightarrow r:
   if(i != j) {
     fa[j] = (fa[i] * fa[i] - conj(fa[j] * fa[j]))
       \hookrightarrow * r;
   fa[i] = z;
 }
 fft(fa, sz);
 vector<int> res(need);
 for(int i = 0; i < need; i++) {
   res[i] = fa[i].x + 0.5;
```

```
return res;
}
vector<int> multiply_mod(vector<int> &a, vector<</pre>
  \hookrightarrowint> &b, int m, int eq = 0) {
 int need = a.size() + b.size() - 1;
 int nbase = 0;
 while ((1 << nbase) < need) nbase++;</pre>
 ensure_base(nbase);
 int sz = 1 << nbase;
 if (sz > (int) fa.size()) {
   fa.resize(sz);
 for (int i = 0; i < (int) a.size(); i++) {
   int x = (a[i] \% m + m) \% m;
   fa[i] = num(x & ((1 << 15) - 1), x >> 15);
  fill(fa.begin() + a.size(), fa.begin() + sz, num
    \hookrightarrow{0. 0}):
  fft(fa, sz);
 if (sz > (int) fb.size()) {
   fb.resize(sz);
 }
 if (eq) {
   copy(fa.begin(), fa.begin() + sz, fb.begin());
 } else {
   for (int i = 0; i < (int) b.size(); i++) {
     int x = (b[i] \% m + m) \% m;
     fb[i] = num(x \& ((1 << 15) - 1), x >> 15);
   fill(fb.begin() + b.size(), fb.begin() + sz,
      \hookrightarrownum \{0, 0\});
   fft(fb, sz);
  dbl ratio = 0.25 / sz;
  num r2(0, -1);
 num r3(ratio. 0):
 num r4(0, -ratio);
 num r5(0, 1);
 for (int i = 0; i \le (sz >> 1); i++) {
   int j = (sz - i) & (sz - 1);
   num a1 = (fa[i] + conj(fa[j]));
   num a2 = (fa[i] - conj(fa[j])) * r2;
   num b1 = (fb[i] + conj(fb[j])) * r3;
   num b2 = (fb[i] - conj(fb[i])) * r4;
   if (i != j) {
     num c1 = (fa[j] + conj(fa[i]));
     num c2 = (fa[j] - conj(fa[i])) * r2;
```

```
num d1 = (fb[j] + conj(fb[i])) * r3;
      num d2 = (fb[i] - conj(fb[i])) * r4;
      fa[i] = c1 * d1 + c2 * d2 * r5;
       fb[i] = c1 * d2 + c2 * d1;
     fa[i] = a1 * b1 + a2 * b2 * r5;
     fb[j] = a1 * b2 + a2 * b1;
   fft(fa, sz);
   fft(fb, sz);
   vector<int> res(need);
   for (int i = 0; i < need; i++) {
     long long aa = fa[i].x + 0.5;
    long long bb = fb[i].x + 0.5;
     long long cc = fa[i].y + 0.5;
     res[i] = (aa + ((bb \% m) << 15) + ((cc \% m) <<
       →30)) % m:
   return res;
 vector<int> square_mod(vector<int> &a, int m) {
   return multiply_mod(a, a, m, 1);
 }
2.17 NTT
const int mod = 7340033;
const int root = 5:
const int root 1 = 4404020:
const int root_pw = 1<<20;</pre>
void fft (vector<int> & a, bool invert) {
 int n = (int) a.size();
 for (int i=1, j=0; i<n; ++i) {
   int bit = n \gg 1;
   for (; j>=bit; bit>>=1)
    i -= bit;
   j += bit;
   if (i < j)
     swap (a[i], a[j]);
 }
 for (int len=2: len<=n: len<<=1) {
   int wlen = invert ? root 1 : root:
   for (int i=len; i<root_pw; i<<=1)</pre>
     wlen = int (wlen * 111 * wlen % mod);
   for (int i=0; i<n; i+=len) {</pre>
     int w = 1;
```

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```
for (int j=0; j<len/2; ++j) {
       int u = a[i+j], v = int (a[i+j+len/2] * 111 *
          \hookrightarrow w % mod);
       a[i+j] = u+v < mod ? u+v : u+v-mod;
       a[i+j+len/2] = u-v >= 0 ? u-v : u-v+mod;
       w = int (w * 111 * wlen % mod);
     }
   }
 }
 if (invert) {
   int nrev = reverse (n, mod);
   for (int i=0: i<n: ++i)
     a[i] = int (a[i] * 1ll * nrev % mod);
2.18 Gauss
// Solves systems of linear equations.
// To use, build a matrix of coefficients and call
  ⇒run(mat. R. C). If the i-th variable is free.
  \hookrightarrowrow[i] will be -1. otherwise it's value will be
  \hookrightarrow ans[i].
namespace Gauss {
 const int MAXC = 1001;
 int row[MAXC];
 double ans[MAXC];
 void run(double mat[][MAXC], int R, int C) {
   REP(i, C) row[i] = -1;
   int r = 0;
   REP(c, C) {
     int k = r;
     FOR(i, r, R) if(fabs(mat[i][c]) > fabs(mat[k][c])
        \hookrightarrow])) k = i;
     if(fabs(mat[k][c]) < eps) continue;</pre>
     REP(j, C+1) swap(mat[r][j], mat[k][j]);
     REP(i, R) if (i != r) {
       double w = mat[i][c] / mat[r][c];
       REP(j, C+1) mat[i][j] -= mat[r][j] * w;
     }
     row[c] = r++;
   REP(i, C) {
     int r = row[i];
     ans[i] = r == -1 ? 0 : mat[r][C] / mat[r][i];
```

```
2.19 Gauss Xor
const 11 MAX = 1e9;
const int LOG_MAX = 64 - __builtin_clzll((11)MAX);
struct Gauss {
    array<11, LOG_MAX> vet;
   int size;
   Gauss() : size(0) {
   fill(vet.begin(), vet.end(), 0);
    Gauss(vector<ll> vals) : size(0) {
   fill(vet.begin(), vet.end(), 0):
       for(ll val : vals) add(val);
   bool add(ll val) {
       for(int i = 0; i < LOG_MAX; i++) if(val & (1
          \hookrightarrowLL \ll i)) {
           if(vet[i] == 0) {
               vet[i] = val;
               size++;
               return true;
           }
           val ^= vet[i];
       return false:
   }
};
2.20 Simpson
inline double simpson(double f1,double fr,double
   \hookrightarrowfmid,double 1,double r) {
  return (fl + fr + 4.0 * fmid) * (r - 1) / 6.0;
double rsimpson(double slr, double fl, double fr,
   \rightarrowdouble fmid,double 1,double r) {
  double mid = (1+r)*0.5;
  double fml = f((1+mid)*0.5), fmr = f((mid+r)*0.5);
  double slm = simpson(fl, fmid, fml, l, mid);
  double smr = simpson(fmid, fr, fmr, mid, r);
  if(fabs(slr-slm-smr) < eps and r - 1 < delta)
     \hookrightarrowreturn slr:
  return rsimpson(slm,fl,fmid,fml,l,mid) + rsimpson(
    \hookrightarrowsmr, fmid, fr, fmr, mid, r);
double integrate(double 1,double r) {
  double mid = (1+r)*0.5;
  double fl = f(1), fr = f(r), fmid = f(mid);
```

```
return rsimpson(simpson(fl,fr,fmid,l,r),fl,fr,fmid
    \hookrightarrow,1,r);
2.21 Modular Arithmetic
template <int mod = MOD>
struct modBase {
 modBase(int val = 0) : val(val) {}
 int val:
 modBase<mod> operator*(modBase<mod> o) {
   return (long long)val * o.val % mod;
 modBase<mod> operator+(modBase<mod> o) {
   return val + o.val > mod ? val + o.val - mod :
      \hookrightarrowval + o.val;
 }
};
template <class T>
T fexp(T x, long long e) {
 T ans(1):
 for (; e > 0; e /= 2) {
   if (e & 1) ans = ans *x;
   x = x * x:
 return ans;
2.22 Matrix
template <const size_t n, const size_t m, class T =</pre>
  \hookrightarrowmodBase<>>
struct Matrix {
 T v[n][m];
 Matrix(int d = 0) {
   for (int i = 0; i < n; i++) {
     for (int j = 0; j < m; j++) {
       v[i][i] = T(0);
     }
     if (i < m) {
       v[i][i] = T(d);
     }
  template <size_t mm>
  Matrix<n, mm, T> operator*(Matrix<m, mm, T> &o) {
   Matrix<n, mm, T> ans;
```

for (int i = 0; i < n; i++) {

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```
3.1 Bipartite Matching
// O(V * E)
int match[N];
int vis[N], pass;
vector<int> g[N];
bool dfs(int u) {
 vis[u] = pass;
 for(int v : g[u]) if(vis[v] != pass) {
   vis[v] = pass;
   if(match[v] == -1 or dfs(match[v])) {
     match[v] = u;
     match[u] = v;
     return true:
   }
 return false;
int max_maching() {
 memset(match, -1, sizeof match);
 int max_matching_size = 0;
 for(int u : vertices_on_side_A) {
   pass++;
   if(dfs(i)) max_matching_size++;
 return max_matching_size;
3.2 Dinic
const int N = 100005;
const int E = 2000006:
vector<int> g[N];
int ne;
struct Edge{
   int from, to; ll flow, cap;
} edge[E];
```

```
int lvl[N], vis[N], pass, start = N-2, target = N-1;
int qu[N], qt, px[N];
11 run(int s, int sink, ll minE){
   if(s == sink) return minE;
   11 \text{ ans} = 0:
   for(; px[s] < (int)g[s].size(); px[s]++){</pre>
       int e = g[s][ px[s] ];
       auto &v = edge[e], &rev = edge[e^1];
       if(lvl[v.to] != lvl[s]+1 || v.flow >= v.cap)
           continue; // v.cap - v.flow < lim</pre>
       11 tmp = run(v.to, sink,min(minE, v.cap-v.
          \hookrightarrowflow));
       v.flow += tmp, rev.flow -= tmp;
       ans += tmp, minE -= tmp;
       if(minE == 0) break;
   }
   return ans:
bool bfs(int source, int sink){
   qt = 0;
   qu[qt++] = source;
   lvl[source] = 1;
   vis[source] = ++pass;
   for(int i = 0; i < qt; i++){
       int u = qu[i]:
       px[u] = 0;
   if(u == sink) return true:
       for(auto& ed : q[u]) {
           auto v = edge[ed];
           if(v.flow >= v.cap || vis[v.to] == pass)
              continue; // v.cap - v.flow < lim</pre>
          vis[v.to] = pass;
           lvl[v.to] = lvl[u]+1;
           qu[qt++] = v.to;
       }
   }
   return false;
11 flow(int source = start, int sink = target){
   11 \text{ ans} = 0;
   //for(lim = (1LL << 62); lim >= 1; lim /= 2)
   while(bfs(source, sink))
   ans += run(source, sink, oo);
   return ans;
void addEdge(int u, int v, ll c = 1, ll rc = 0){
```

```
edge[ne] = {u, v, 0, c};
    g[u].push_back(ne++);
    edge[ne] = {v, u, 0, rc};
    g[v].push_back(ne++);
}
void reset_flow(){
    for(int i = 0; i < ne; i++)
        edge[i].flow = 0;
}
3.3     Push relabel
// Push relabel in O(V^2 E^0.5) w
// It's quite fast</pre>
```

```
// Push relabel in O(V^2 E^0.5) with gap heuristic
template<typename flow_t = long long>
struct PushRelabel {
   struct Edge { int to, rev; flow_t f, c; };
   vector<vector<Edge> > g;
   vector<flow_t> ec;
   vector<Edge*> cur;
   vector<vector<int> > hs;
   vector<int> H:
   PushRelabel(int n): g(n), ec(n), cur(n), hs(2*n)
      \hookrightarrow, H(n) \{\}
   void add_edge(int s, int t, flow_t cap, flow_t
      \hookrightarrowrcap=0) {
       if (s == t) return;
       Edge a = \{t, (int)g[t].size(), 0, cap\};
       Edge b = \{s, (int)g[s].size(), 0, rcap\};
       g[s].push_back(a);
       g[t].push_back(b);
   void add_flow(Edge& e, flow_t f) {
       Edge &back = g[e.to][e.rev];
       if (!ec[e.to] && f)
           hs[H[e.to]].push_back(e.to);
       e.f += f, ec[e.to] += f;
       back.f -= f, ec[back.to] -= f;
   flow_t max_flow(int s, int t) {
       int v = g.size();
       H[s] = v; ec[t] = 1;
       vector<int> co(2 * v);
       co[0] = v-1;
       for(int i = 0; i < v; ++i) cur[i] = g[i].data
         \hookrightarrow ():
       for(auto &e : g[s]) add_flow(e, e.c);
       if(hs[0].size())
       for (int hi = 0; hi >= 0;) {
           int u = hs[hi].back();
```

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```
hs[hi].pop_back();
           while (ec[u] > 0) // discharge u
                                                                   for(int e : g[u]) {
               if (cur[u] == g[u].data() + g[u].size
                                                                      auto v = edge[e];
                  \hookrightarrow()) {
                   H[u] = 1e9;
                                                                        \hookrightarrowtol:
                    for(auto &e:g[u])
                                                                      if(v.cap > 0 \text{ and } cand < d[v.to]) 
                       if (e.c - e.f && H[u] > H[e.to]
                                                                       p[v.to] = e;
                          \hookrightarrow]+1)
                                                                       d[v.to] = cand;
                            H[u] = H[e.to]+1, cur[u] = &
                                                                       q.emplace(-d[v.to], v.to);
                              ∽e;
                                                                     }
                   if (++co[H[u]], !--co[hi] \& hi < v
                                                                   }
                      \hookrightarrow)
                        for(int i = 0; i < v; ++i)
                                                                 return d[sink] < oo;</pre>
                            if (hi < H[i] && H[i] < v){
                               --co[H[i]];
                               H[i] = v + 1;
                                                                // <max flow, min cost>
                   hi = H[u];
                                                                   \hookrightarrowtarget){
               } else if (cur[u]->c - cur[u]->f && H[
                                                                   11 ans = 0, mf = 0;
                  \hookrightarrowu] == H[cur[u]->to]+1)
                                                                   while(dijkstra(source, sink)){
                   add_flow(*cur[u], min(ec[u], cur[u
                                                                       11 f = oo;
                      \hookrightarrow]->c - cur[u]->f));
               else ++cur[u];
                                                                          \hookrightarrow ].from)
           while (hi \geq 0 && hs[hi].empty()) --hi;
                                                                           f = min(f, edge[p[u]].cap);
       }
                                                                       mf += f:
       return -ec[s];
                                                                          \hookrightarrow1):
   }
};
                                                                           \hookrightarrow ].from){
3.4 Min Cost Max Flow
                                                                           edge[ p[u] ].cap -= f;
const 11 oo = 1e18;
                                                                           edge[p[u] ^1].cap += f;
const int N = 422, E = 2 * 10006;
vector<int> g[N];
                                                                      \hookrightarrow] + d[i]);
int ne;
```

# struct Edge{

```
int from, to; ll cap, cost;
} edge[E];
int start = N-1, target = N-2, p[N]; int inqueue[N];
ll d[N];
```

11 pot[N]; bool dijkstra(int source, int sink) { for(int i = 0; i < N; i++) d[i] = oo;

d[source] = 0;priority\_queue<pair<ll, int>> q; q.emplace(0, source);

if(dt > d[u]) continue;

ll dt; int u; while(!q.empty()) { tie(dt, u) = q.top(); q.pop(); dt = -dt; const int MAXN = 2020 + 1; // 1-based Vertex index int vis[MAXN], par[MAXN], orig[MAXN], match[MAXN],  $\hookrightarrow$ aux[MAXN], t, N;

```
if(u == sink) return true;
     const 11 cand = d[u] + v.cost + pot[u] - pot[v.
pair<11, 11> mincost(int source = start, int sink =
       for(int u = sink; u != source; u = edge[ p[u]
       ans += f * (d[sink] - pot[source] + pot[sink
       for(int u = sink; u != source; u = edge[ p[u]
   for(int i = 0; i < N; i++) pot[i] = min(oo, pot[i
   return {mf, ans};
void addEdge(int u, int v, ll c, ll cost){
 assert(cost >= 0);
   edge[ne] = \{u, v, c, cost\};
   g[u].push_back(ne++);
   edge[ne] = \{v, u, 0, -cost\};
   g[v].push_back(ne++);
```

## 3.5 Blossom Algorithm for General Matching

```
vector<int> conn[MAXN];
queue<int> Q;
void addEdge(int u, int v) {
 conn[u].push_back(v); conn[v].push_back(u);
void init(int n) {
 N = n; t = 0;
 for(int i=0; i<=n; ++i)
   conn[i].clear(), match[i] = aux[i] = par[i] = 0;
void augment(int u, int v) {
 int pv = v, nv;
 do {
   pv = par[v]; nv = match[pv];
   match[v] = pv; match[pv] = v;
   v = nv;
 } while(u != pv);
int lca(int v, int w) {
 ++t:
 while(true) {
   if(v) {
     if(aux[v] == t) return v; aux[v] = t;
     v = orig[par[match[v]]];
   swap(v, w);
 }
void blossom(int v, int w, int a) {
 while(orig[v] != a) {
   par[v] = w; w = match[v];
   if(vis[w] == 1) Q.push(w), vis[w] = 0;
   orig[v] = orig[w] = a; v = par[w];
 }
bool bfs(int u) {
 fill(vis+1, vis+1+N, -1); iota(orig + 1, orig + N
    \hookrightarrow+ 1. 1):
 Q = queue < int > (); Q.push(u); vis[u] = 0;
 while(!Q.empty()) {
   int v = Q.front(); Q.pop();
   for(int x: conn[v]) {
     if(vis[x] == -1) {
       par[x] = v; vis[x] = 1;
       if(!match[x]) return augment(u, x), true;
       Q.push(match[x]); vis[match[x]] = 0;
     else if(vis[x] == 0 && orig[v] != orig[x]) {
       int a = lca(orig[v], orig[x]);
```

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```
blossom(x, v, a); blossom(v, x, a);
     }
   }
 }
 return false;
int Match() {
 int ans = 0;
 // find random matching (not necessary, constant
    \hookrightarrow improvement)
 vector<int> V(N-1); iota(V.begin(), V.end(), 1);
  shuffle(V.begin(), V.end(), mt19937(0x94949));
 for(auto x: V) if(!match[x]){
   for(auto y: conn[x]) if(!match[y]) {
     match[x] = y, match[y] = x;
     ++ans; break;
   }
 for(int i=1; i<=N; ++i) if(!match[i] && bfs(i)) ++</pre>
    ⇒ans:
 return ans:
```

# 3.6 Blossom Algorithm for Weighted General Matching

```
// N^3 (but fast in practice)
static const int INF = INT_MAX;
static const int N = 514:
struct edge{
 int u,v,w; edge(){}
 edge(int ui,int vi,int wi)
    :u(ui),v(vi),w(wi){}
};
int n,n_x;
edge g[N*2][N*2];
int lab[N*2];
int match[N*2],slack[N*2],st[N*2],pa[N*2];
int flo_from[N*2][N+1],S[N*2],vis[N*2];
vector<int> flo[N*2];
queue<int> q;
int e_delta(const edge &e){
 return lab[e.u]+lab[e.v]-g[e.u][e.v].w*2;
void update_slack(int u,int x){
 if(!slack[x]||e_delta(g[u][x])<e_delta(g[slack[x</pre>
    \hookrightarrow]][x]))slack[x]=u;
void set_slack(int x){
 slack[x]=0;
```

```
for(int u=1;u<=n;++u)
   if(g[u][x].w>0&&st[u]!=x&&S[st[u]]==0)
     update_slack(u,x);
void q_push(int x){
 if(x \le n)q.push(x);
  else for(size_t i=0;i<flo[x].size();i++)</pre>
   q_push(flo[x][i]);
void set_st(int x,int b){
  st[x]=b:
 if(x>n)for(size_t i=0;i<flo[x].size();++i)
   set_st(flo[x][i],b);
int get_pr(int b,int xr){
  int pr=find(flo[b].begin(),flo[b].end(),xr)-flo[b
    \hookrightarrow].begin();
  if(pr%2==1){
   reverse(flo[b].begin()+1,flo[b].end());
   return (int)flo[b].size()-pr;
  }else return pr;
void set_match(int u,int v){
  match[u]=g[u][v].v;
  if(u<=n) return;</pre>
  edge e=g[u][v];
  int xr=flo_from[u][e.u],pr=get_pr(u,xr);
  for(int i=0;i<pr;++i)set_match(flo[u][i],flo[u][i</pre>
    \hookrightarrow 1]);
  set match(xr.v):
  rotate(flo[u].begin(),flo[u].begin()+pr,flo[u].end
    \hookrightarrow());
void augment(int u,int v){
 for(;;){
   int xnv=st[match[u]];
   set_match(u,v);
   if(!xnv)return:
   set_match(xnv,st[pa[xnv]]);
   u=st[pa[xnv]],v=xnv;
 }
int get_lca(int u,int v){
  static int t=0;
  for(++t;u|v;swap(u,v)){
   if(u==0)continue;
   if(vis[u]==t)return u;
   vis[u]=t;
   u=st[match[u]];
```

```
if(u)u=st[pa[u]];
 }
 return 0;
void add_blossom(int u,int lca,int v){
 int b=n+1;
 while(b<=n_x&&st[b])++b;
 if(b>n_x)++n_x;
 lab[b]=0,S[b]=0;
 match[b]=match[lca];
 flo[b].clear();
  flo[b].push_back(lca);
  for(int x=u,y;x!=lca;x=st[pa[y]])
   flo[b].push_back(x),flo[b].push_back(y=st[match[x
      \hookrightarrow]]),q_push(y);
 reverse(flo[b].begin()+1,flo[b].end());
  for(int x=v,y;x!=lca;x=st[pa[y]])
   flo[b].push_back(x),flo[b].push_back(y=st[match[x
      \hookrightarrow]]),q_push(y);
  set st(b.b):
  for(int x=1; x \le n_x; ++x)g[b][x].w=g[x][b].w=0;
  for(int x=1;x<=n;++x)flo_from[b][x]=0;</pre>
  for(size_t i=0;i<flo[b].size();++i){</pre>
   int xs=flo[b][i];
   for(int x=1; x<=n_x;++x)
     if(g[b][x].w==0||e_delta(g[xs][x])< e_delta(g[b]
        \hookrightarrow][x]))
       g[b][x]=g[xs][x],g[x][b]=g[x][xs];
   for(int x=1;x\leq n;++x)
     if(flo_from[xs][x])flo_from[b][x]=xs;
 set_slack(b);
void expand_blossom(int b){
 for(size_t i=0;i<flo[b].size();++i)</pre>
   set_st(flo[b][i],flo[b][i]);
 int xr=flo_from[b][g[b][pa[b]].u],pr=get_pr(b,xr);
 for(int i=0;i<pr;i+=2){
   int xs=flo[b][i],xns=flo[b][i+1];
   pa[xs]=g[xns][xs].u;
   S[xs]=1,S[xns]=0;
   slack[xs]=0, set_slack(xns);
   q_push(xns);
 S[xr]=1,pa[xr]=pa[b];
 for(size_t i=pr+1;i<flo[b].size();++i){</pre>
   int xs=flo[b][i];
   S[xs]=-1,set_slack(xs);
```

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```
st[b]=0:
bool on_found_edge(const edge &e){
 int u=st[e.u],v=st[e.v];
 if(S[v]==-1){
   pa[v]=e.u.S[v]=1;
   int nu=st[match[v]];
   slack[v]=slack[nu]=0;
   S[nu]=0,q_push(nu);
 }else if(S[v]==0){
   int lca=get_lca(u,v);
   if(!lca)return augment(u,v),augment(v,u),true;
   else add_blossom(u,lca,v);
 return false;
bool matching(){
 memset(S+1,-1,sizeof(int)*n_x);
 memset(slack+1,0,sizeof(int)*n_x);
 q=queue<int>():
 for(int x=1;x\leq n_x;++x)
   if(st[x]==x\&\&!match[x])pa[x]=0,S[x]=0,q_push(x);
 if(q.empty())return false;
 for(;;){
   while(q.size()){
     int u=q.front();q.pop();
     if(S[st[u]]==1)continue;
     for(int v=1; v<=n; ++v)
      if(g[u][v].w>0&&st[u]!=st[v]){
        if(e_delta(g[u][v])==0){
          if(on_found_edge(g[u][v]))return true;
        }else update_slack(u,st[v]);
      }
   int d=INF:
   for(int b=n+1;b<=n_x;++b)
     if(st[b]==b\&\&S[b]==1)d=min(d,lab[b]/2);
   for(int x=1:x \le n x:++x)
     if(st[x]==x&&slack[x]){
      if(S[x]==-1)d=min(d,e_delta(g[slack[x]][x]));
      else if(S[x]==0)d=min(d,e_delta(g[slack[x]][x
         \hookrightarrow])/2);
   for(int u=1; u<=n;++u){
     if(S[st[u]]==0){
      if(lab[u]<=d)return 0;</pre>
      lab[u]-=d;
     }else if(S[st[u]]==1)lab[u]+=d;
```

```
for(int b=n+1;b<=n_x;++b)
     if(st[b]==b){
       if(S[st[b]]==0)lab[b]+=d*2;
       else if(S[st[b]]==1)lab[b]-=d*2;
   q=queue<int>();
   for(int x=1; x<=n_x;++x)
     if(st[x]==x&&slack[x]&&st[slack[x]]!=x&&e_delta
        \hookrightarrow (g[slack[x]][x])==0)
       if(on_found_edge(g[slack[x]][x]))return true;
   for(int b=n+1;b<=n_x;++b)
     if(st[b]==b&&S[b]==1&&lab[b]==0)expand_blossom(
        →b);
 }
 return false;
pair<long long,int> solve(){
 memset(match+1,0,sizeof(int)*n);
  n x=n:
  int n matches=0:
  long long tot_weight=0;
  for(int u=0;u<=n;++u)st[u]=u,flo[u].clear();</pre>
  int w_max=0;
  for(int u=1;u<=n;++u)
   for(int v=1; v<=n;++v){
     flo_from[u][v]=(u==v?u:0);
     w_{max}=max(w_{max},g[u][v].w);
  for(int u=1;u<=n;++u)lab[u]=w_max;
  while(matching())++n matches:
  for(int u=1;u<=n;++u)
   if(match[u]&&match[u]<u)</pre>
     tot_weight+=g[u][match[u]].w;
 return make_pair(tot_weight,n_matches);
void add_edge( int ui , int vi , int wi ){
 g[ui][vi].w = g[vi][ui].w = wi;
void init( int _n ){
 n = _n;
  for(int u=1;u<=n;++u)
   for(int v=1; v<=n;++v)
     g[u][v]=edge(u,v,0);
}
3.7 Small to Large
void cnt_sz(int u, int p = -1){
   sz[u] = 1;
   for(int v : q[u]) if(v != p)
```

 $cnt_sz(v, u), sz[u] += sz[v];$ 

```
void add(int u, int p, int big = -1){
   // Update info about this vx in global answer
   for(int v : g[u]) if(v != p && v != big)
       add(v, u);
void dfs(int u, int p, int keep){
   int big = -1, mmx = -1;
   for(int v : g[u]) if(v != p \&\& sz[v] > mmx)
       mmx = sz[v], big = v;
   for(int v : g[u]) if(v != p && v != big)
       dfs(v. u. 0):
   if(big != -1) dfs(big, u, 1);
   add(u, p, big);
   for(auto x : q[u]){
      // answer all gueries for this vx
   if(!keep){ /*Remove data from this subtree*/ }
```

#### 3.8 Centroid Decomposition

```
void decomp(int v, int p){
 int treesize = calc_sz(v, v);
 if(treesize < k) return;</pre>
 int cent = centroid(v, v, treesize);
 erased[cent] = 1;
 for(int i = 1; i <= treesize; i++) dist[i] = 1e18;
   for(pair<int,int> x : G[cent]) if(!erased[x.ff]){
       procurar_ans(x.ff, cent, 1, x.ss); // linear
       atualiza_dist(x.ff, cent, 1, x.ss); // linear
 for(pair<int,int> x : G[cent]) if(!erased[x.ff])
       decomp(x.ff, cent);
```

#### 3.9 Kosaraju

```
vector<int> g[N], gt[N], S; int vis[N], cor[N];
void dfs(int u){
 vis[u] = 1; for(int v : g[u]) if(!vis[v]) dfs(v);
 S.push_back(u);
void dfst(int u, int e){
 cor[u] = e:
 for(int v : gt[u]) if(!cor[v]) dfst(v, e);
void kosaraju(){
 for(int i = 1; i <= n; i++) if(!vis[i]) dfs(i);
```

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```
for(int i = 1; i \le n; i \leftrightarrow j for(int j : g[i])
      gt[j].push_back(i);
 int e = 0; reverse(S.begin(), S.end());
 for(int u : S) if(!cor[u]) dfst(u, ++e);
3.10 Tarjan
int cnt = 0, root;
void dfs(int u, int p = -1){
 low[u] = num[u] = ++t;
 for(int v : q[u]){
   if(!num[v]){
     dfs(v, u);
            if(u == root) cnt++;
     if(low[v] >= num[u]) u PONTO DE ARTICULAÇÃO;
     if(low[v] > num[u]) ARESTA u->v PONTE;
     low[u] = min(low[u], low[v]);
   else if(v != p) low[u] = min(low[u], num[v]);
root PONTO DE ARTICULAÇÃO <=> cnt > 1
void tarjanSCC(int u){
 low[u] = num[u] = ++cnt;
 vis[u] = 1:
 S.push_back(u);
 for(int v : q[u]){
   if(!num[v]) tarjanSCC(v);
   if(vis[v]) low[u] = min(low[u], low[v]);
 if(low[u] == num[u]){
   ssc[u] = ++ssc_cnt; int v;
     v = S.back(); S.pop_back(); vis[v] = 0;
     ssc[v] = ssc_cnt;
   }while(u != v):
}
3.11 Max Clique
long long adj[N], dp[N];
for(int i = 0; i < n; i++){
 for(int j = 0; j < n; j++){
   int x;
   scanf("%d",&x);
   if(x \mid | i == j)
     adj[i] |= 1LL << j;
```

```
int resto = n - n/2:
int C = n/2;
for(int i = 1; i < (1 << resto); i++){}
 int x = i;
 for(int j = 0; j < resto; j++)
   if(i & (1 << j))
     x \&= adj[j + C] >> C;
 if(x == i)
   dp[i] = __builtin_popcount(i);
 }
}
for(int i = 1; i < (1 << resto); i++)</pre>
 for(int j = 0; j < resto; j++)
   if(i & (1 << j))
     dp[i] = max(dp[i], dp[i ^ (1 << j)]);
int maxCliq = 0;
for(int i = 0; i < (1 << C); i++){
 int x = i, y = (1 << resto) - 1;
 for(int i = 0; i < C; i++)
   if(i & (1 << j))
     x \&= adj[j] \& ((1 << C) - 1), y \&= adj[j] >> C;
 if(x != i) continue:
 maxCliq = max(maxCliq, __builtin_popcount(i) + dp[
    →y]);
3.12 Dominator Tree
vector<int> g[N], gt[N], T[N];
vector<int> S;
int dsu[N], label[N];
int sdom[N], idom[N], dfs_time, id[N];
vector<int> bucket[N]:
vector<int> down[N];
void prep(int u){
 S.push_back(u);
 id[u] = ++dfs_time;
 label[u] = sdom[u] = dsu[u] = u;
 for(int v : g[u]){
   if(!id[v])
     prep(v), down[u].push_back(v);
   gt[v].push_back(u);
```

```
int fnd(int u, int flag = 0){
 if(u == dsu[u]) return u;
 int v = fnd(dsu[u], 1), b = label[ dsu[u] ];
 if(id[ sdom[b] ] < id[ sdom[ label[u] ] ])</pre>
  label[u] = b;
 dsu[u] = v;
 return flag ? v : label[u];
void build_dominator_tree(int root, int sz){
 // memset(id, 0, sizeof(int) * (sz + 1));
 // for(int i = 0; i <= sz; i++) T[i].clear();
 prep(root);
 reverse(S.begin(), S.end());
 int w;
 for(int u : S){
   for(int v : gt[u]){
    w = fnd(v);
    if(id[ sdom[w] ] < id[ sdom[u] ])
      sdom[u] = sdom[w];
   gt[u].clear();
   if(u != root) bucket[ sdom[u] ].push_back(u);
   for(int v : bucket[u]){
     w = fnd(v):
     if(sdom[w] == sdom[v]) idom[v] = sdom[v];
     else idom[v] = w;
   bucket[u].clear();
   for(int v : down[u]) dsu[v] = u;
   down[u].clear();
 reverse(S.begin(), S.end());
 for(int u : S) if(u != root){
   if(idom[u] != sdom[u]) idom[u] = idom[ idom[u] ];
   T[ idom[u] ].push_back(u);
 S.clear();
3.13 Min Cost Matching
// Min cost matching
// O(n^2 * m)
```

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```
// n == nro de linhas
// m == nro de colunas
// n \ll m / flow == n
// a[i][j] = custo pra conectar i a j
vector<int> u(n + 1), v(m + 1), p(m + 1), way(m + 1)
for(int i = 1; i \le n; ++i){
   p[0] = i;
   int j0 = 0;
   vector<int> minv(m + 1 , oo);
   vector<char> used(m + 1 , false);
   do{
      used[i0] = true;
      int i0 = p[j0] , delta = oo, j1;
      for(int j = 1; j \le m; ++j)
          if(! used[i]){
              int cur = a[i0][j] - u[i0] - v[j];
             if(cur < minv[j])</pre>
                 minv[j] = cur, way[j] = j0;
             if(minv[j] < delta)</pre>
                 delta = minv[j] , j1 = j;
       for(int j = 0; j \le m; ++j)
          if(used[i])
              u[p[j]] += delta, v[j] -= delta;
          else
              minv[j] -= delta;
      i0 = i1;
   }while(p[j0] != 0);
   do{
      int j1 = way[j0];
      p[j0] = p[j1];
      j0 = j1;
   }while(j0);
// match[i] = coluna escolhida para linha i
vector<int> match(n + 1);
for(int j = 1; j \le m; ++j)
   match[p[j]] = j;
int cost = -v[0];
4 Strings
4.1 Aho Corasick
int to[N][A];
int ne = 2, fail[N], term[N];
```

void add\_string(const char \*str, int id){

```
int p = 1;
   for(int i = 0; str[i]; i++){
       int ch = str[i] - 'a';
       if(!to[p][ch]) to[p][ch] = ne++;
       p = to[p][ch];
   term[p]++;
void init(){
   for(int i = 0; i < ne; i++) fail[i] = 1;</pre>
   queue<int> q; q.push(1);
   int u, v; char c;
   while(!q.empty()){
       u = q.front(); q.pop();
       for(int i = 0; i < A; i++){
          if(to[u][i]){
              v = to[u][i]; q.push(v);
              if(u != 1){
                  fail[v] = to[ fail[u] ][i];
                  term[v] += term[ fail[v] ];
              }
          else if(u != 1) to[u][i] = to[ fail[u] ][i
             \hookrightarrow
          else to[u][i] = 1;
       }
   }
void clean() {
   memset(to, 0, ne * sizeof(to[0]));
   memset(fail, 0, ne * sizeof(fail[0]));
   memset(term, 0, ne * sizeof(term[0]));
   memset(to, 0, ne * sizeof(to[0]));
   ne = 2;
4.2 Suffix Array
int lcp[N], c[N];
// Caractere final da string '\0' esta sendo
  \hookrightarrow considerado parte da string s
void build_sa(char s[], int n, int a[]){
   const int A = 300; // Tamanho do alfabeto
   int c1[n], a1[n], h[n + A];
   memset(h, 0, sizeof h);
   for(int i = 0; i < n; i++) {
       c[i] = s[i];
       h[c[i] + 1]++;
```

```
partial_sum(h, h + A, h);
               for(int i = 0; i < n; i++)
                             a[h[c[i]]++] = i;
               for(int i = 0; i < n; i++)
                            h[c[i]]--;
               for(int L = 1; L < n; L <<= 1) {
                             for(int i = 0; i < n; i++) {
                                           int j = (a[i] - L + n) \% n;
                                           a1[h[c[j]]++] = j;
                            }
                            int cc = -1;
                            for(int i = 0; i < n; i++) {
                                         if(i == 0 \mid \mid c[a1[i]] \mid = c[a1[i-1]] \mid \mid c[(a1[i-1]] \mid (a1[i-1]) \mid c[(a1[i-1]] \mid a[(a1[i-1]) \mid c[(a1[i-1]] \mid a[(a1[i-1]) \mid a[(a1[
                                                    \hookrightarrowa1[i] + L) % n] != c[(a1[i-1] + L) %
                                                    \hookrightarrown])
                                                        h[++cc] = i;
                                           c1[a1[i]] = cc;
                            }
                             memcpy(a, a1, sizeof a1);
                            memcpy(c, c1, sizeof c1);
                            if(cc == n-1) break;
}
void build_lcp(char s[], int n, int a[]){ // lcp[i]
           \hookrightarrow = lcp(s[:i], s[:i+1])
              int k = 0;
              //memset(lcp, 0, sizeof lcp);
               for(int i = 0; i < n; i++){
                            if(c[i] == n-1) continue;
                            int j = a[c[i]+1];
                            while(i+k < n \& j+k < n \& s[i+k] == s[j+k])
                                     \hookrightarrow k++;
                           lcp[c[i]] = k;
                             if(k) k--;
             }
}
int comp_lcp(int i, int j){
              if(i == j) return n - i;
              if(c[i] > c[j]) swap(i, j);
              return min(lcp[k]  for k in [c[i], c[j]-1]);
}
```

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#### 4.3 Adamant Suffix Tree

```
namespace sf {
const int inf = 1e9:
const int maxn = 200005:
char s[maxn];
map<int, int> to[maxn];
int len[maxn], fpos[maxn], link[maxn];
int node. pos:
int sz = 1, n = 0;
int make_node(int _pos, int _len) {
 fpos[sz] = _pos;
 len[sz] = _len;
 return sz++;
void go_edge() {
 while (pos > len[to[node][s[n - pos]]]) {
   node = to[node][s[n - pos]];
   pos -= len[node];
void add_letter(int c) {
 s[n++] = (char)c;
 pos++;
 int last = 0;
 while (pos > 0) {
   go_edge();
   int edge = s[n - pos];
   int &v = to[node][edge];
   int t = s[fpos[v] + pos - 1];
   if (v == 0) {
     v = make_node(n - pos, inf);
     link[last] = node;
     last = 0:
   } else if (t == c) {
     link[last] = node;
     return;
   } else {
     int u = make_node(fpos[v], pos - 1);
     to[u][c] = make\_node(n - 1, inf);
     to[u][t] = v;
     fpos[v] += pos - 1;
     len[v] -= pos - 1;
     v = u;
     link[last] = u;
     last = u;
   if (node == 0)
```

```
pos--;
   else
     node = link[node];
 }
}
void add_string(char *str) {
 for (int i = 0; str[i]; i++) add_letter(str[i]);
 add_letter('$');
bool is_leaf(int u) { return len[u] > n; }
int get_len(int u) {
 if (!u) return 0:
 if (is_leaf(u)) return n - fpos[u];
 return len[u];
int leafs[maxn];
int calc_leafs(int u = 0) {
 leafs[u] = is_leaf(u);
 for (const auto &c : to[u]) leafs[u] += calc_leafs
    \hookrightarrow (c.second):
 return leafs[u];
}; // namespace sf
int main() { sf::len[0] = sf::inf; }
4.4 Z Algorithm
vector<int> z_algo(const string &s) {
 int n = s.size(), L = 0, R = 0;
 vector<int> z(n. 0):
 for(int i = 1; i < n; i++){
   if(i \le R) z[i] = min(z[i-L], R - i + 1);
   while(z[i]+i < n \& s[z[i]+i] == s[z[i]])
     z[i]++;
   if(i+z[i]-1 > R) L = i, R = i + z[i] - 1;
 return z;
4.5 Prefix function/KMP
vector<int> preffix_function(const string &s){
 int n = s.size(); vector<int> b(n+1);
 b[0] = -1; int i = 0, j = -1;
 while(i < n)
   while(j \ge 0 \& s[i] != s[j]) j = b[j];
   b[++i] = ++j;
 return b;
```

void kmp(const string &t, const string &p){

```
vector<int> b = preffix_function(p);
 int n = t.size(), m = p.size();
 int i = 0;
 for(int i = 0; i < n; i++){
   while(j \ge 0 \& t[i] != p[j]) j = b[j];
   i++;
   if(j == m)
    //patern of p found on t
    i = b[i];
 }
4.6 Min rotation
int min_rotation(int *s, int N) {
 REP(i, N) s[N+i] = s[i];
 int a = 0;
 REP(b, N) REP(i, N) {
   if (a+i == b \mid | s[a+i] < s[b+i]) { b += max(0, i)}
     \hookrightarrow-1); break; }
   if (s[a+i] > s[b+i])  { a = b: break: }
 }
 return a:
4.7 Manacher
// rad[2 * i] = largest palindrome cetered at char i
// rad[2 * i + 1] = largest palindrome cetered
  ⇒between chars i and i+i
void manacher(char *s, int n, int *rad) {
 static char t[2*MAX];
 int m = 2 * n - 1;
 for(int i = 0; i < m; i++) t[i] = -1;
 for(int i = 0; i < n; i++) t[2 * i] = s[i];
 int x = 0:
 rad[0] = 0; // <
 for(int i = 1; i < m; i++) {
   int &r = rad[i] = 0;
   if(i \le x+rad[x]) r = min(rad[x+x-i],x+rad[x]-i);
   while(i - r - 1 >= 0 and i + r + 1 < m and
      t[i - r - 1] == t[i + r + 1]) ++r;
   if(i + r >= x + rad[x]) x = i;
 }
 for(int i = 0; i < m; i++) {
   if(i-rad[i] == 0 || i+rad[i] == m-1) ++rad[i];
```

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```
// for(int i = 0; i < m; i++) rad[i] /= 2;
4.8 Suffix Automaton
map<char, int> to[2*N];
int link[2*N], len[2*N], last = 0, sz = 1;
void add_letter(char c){
   int p = last;
   last = sz++;
   len[last] = len[p] + 1;
   for(; !to[p][c]; p = link[p]) to[p][c] = last;
   if(to[p][c] == last){
      link[last] = 0;
      return;
   int u = to[p][c];
   if(len[u] == len[p]+1){
      link[last] = u;
      return;
```

#### 4.9 Suffix Tree

int c1 = sz++;

to[c1] = to[u];

link[c1] = link[u];

len[c1] = len[p]+1;

link[last] = link[u] = c1;

```
namespace sf {
// const int NS = ; const int N = * 2;
int cn, cd, ns, en = 1, lst;
string S[NS]; int si = -1;
vector<int> sufn[N]; // sufn[si][i] no do sufixo S[
  ∽si][i...]
struct node {
 int 1, r, si, p, suf;
 map<char, int> adj;
 node() : 1(0), r(-1), suf(0), p(0) {}
 node(int L, int R, int S, int P) : l(L), r(R), si(
    \hookrightarrowS), p(P) {}
 inline int len() { return r - 1 + 1; }
 inline int operator[](int i) { return S[si][l + i
    \hookrightarrow]; }
 inline int& operator()(char c) { return adj[c]; }
inline int new_node(int L, int R, int S, int P) { t[
  \hookrightarrowen] = node(L, R, S, P); return en++; }
void add_string(string s) {
```

for(; to[p][c] == u; p = link[p]) to[p][c] = c1;

```
s += '; S[++si] = s; sufn[si].resize(s.size() +
    \hookrightarrow1); cn = cd = 0;
  int i = 0; const int n = s.size();
  for(int j = 0; j < n; j++)
   for(; i <= j; i++) {
     if(cd == t[cn].len() \&\& t[cn](s[j])) { cn = t[}
        \hookrightarrow cn](s[j]); cd = 0; }
     if(cd < t[cn].len() \&\& t[cn][cd] == s[j]) {
       cd++:
       if(j < s.size() - 1) break;
       else {
         if(i) t[lst].suf = cn;
         for(; i <= j; i++) { sufn[si][i] = cn; cn =
            \hookrightarrowt[cn].suf; }
     } else if(cd == t[cn].len()) {
       sufn[si][i] = en;
       if(i) t[lst].suf = en; lst = en;
       t[cn](s[j]) = new_node(j, n - 1, si, cn);
       cn = t[cn].suf; cd = t[cn].len();
     } else {
       int mid = new_node(t[cn].l, t[cn].l + cd - 1,
          \hookrightarrow t[cn].si, t[cn].p);
       t[t[cn].p](t[cn][0]) = mid;
       if(ns) t[ns].suf = mid;
       if(i) t[lst].suf = en; lst = en;
       sufn[si][i] = en;
       t[mid](s[j]) = new_node(j, n - 1, si, mid);
       t[mid](t[cn][cd]) = cn;
       t[cn].p = mid; t[cn].l += cd; cn = t[mid].p;
       int g = cn? j - cd : i + 1; cn = t[cn].suf;
       while(g < j \&\& g + t[t[cn](S[si][g])].len()
          = j) {
         cn = t[cn](S[si][g]); g += t[cn].len();
       if(g == j) \{ ns = 0; t[mid].suf = cn; cd = t[
          \hookrightarrowcnl.len(): }
       else { ns = mid; cn = t[cn](S[si][g]); cd = j
          \hookrightarrow - q; }
 }
};
   Geometry
5.1 2D basics
```

```
typedef double cod;
double eps = 1e-7;
bool eq(cod a, cod b){ return abs(a - b) <= eps; }</pre>
```

```
struct vec{
 cod x, y; int id;
 vec(cod a = 0, cod b = 0) : x(a), y(b) {}
 vec operator+(const vec &o) const{
   return \{x + o.x, y + o.y\};
 vec operator-(const vec &o) const{
   return \{x - o.x, y - o.y\};
 vec operator*(cod t) const{
   return {x * t, y * t};
 vec operator/(cod t) const{
   return {x / t, y / t};
 cod operator*(const vec &o) const{ // cos
   return x * o.x + y * o.y;
 cod operator^(const vec &o) const{ // sin
   return x * o.y - y * o.x;
 bool operator==(const vec &o) const{
   return eq(x, o.x) && eq(y, o.y);
 bool operator<(const vec &o) const{</pre>
   if(!eq(x, o.x)) return x < o.x;
   return y < o.y;</pre>
 cod cross(const vec &a. const vec &b) const{
   return (a-(*this)) ^ (b-(*this));
   int ccw(const vec &a, const vec &b) const{
       cod tmp = cross(a, b);
       return (tmp > eps) - (tmp < -eps);</pre>
 cod dot(const vec &a, const vec &b) const{
   return (a-(*this)) * (b-(*this));
 cod len() const{
   return sqrt(x * x + y * y); // <
 double angle(const vec &a, const vec &b) const{
   return atan2(cross(a, b), dot(a, b));
 }
 double tan(const vec &a, const vec &b) const{
   return cross(a, b) / dot(a, b);
 }
 vec unit() const{
```

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```
return operator/(len());
 int quad() const{
   if(x > 0 \&\& y >= 0) return 0;
   if(x \le 0 \& y > 0) return 1;
   if(x < 0 \&\& y <=0) return 2;
   return 3;
 bool comp(const vec &a, const vec &b) const{
   return (a - *this).comp(b - *this);
 bool comp(vec b){
   if(quad() != b.quad()) return quad() < b.quad();</pre>
   if(!eq(operator^(b), 0)) return operator^(b) > 0;
   return (*this) * (*this) < b * b;
  template<class T>
 void sort_by_angle(T first, T last) const{
   std::sort(first, last, [=](const vec &a, const
      →vec &b){
     return comp(a, b);
   });
 vec rot90() const{ return {-y, x}; }
 vec rot(double a) const{
   return \{\cos(a)*x - \sin(a)*y, \sin(a)*x + \cos(a)*y\};
   vec proj(const vec &b) const{ // proj of *this
      \hookrightarrowonto b
       cod k = operator*(b) / (b * b):
       return b * k;
   // proj of (*this) onto the plane orthogonal to b
   vec rejection(vec b) const{
       return (*this) - proj(b);
};
struct line{
 cod a, b, c; vec n;
 line(vec q, vec w){ // q.cross(w, (x, y)) = 0
   a = -(w.y-q.y);
   b = w.x-q.x;
   c = -(a * q.x + b * q.y);
   n = \{a, b\};
 cod dist(const vec &o) const{
   return abs(eval(o)) / n.len();
```

```
bool contains(const vec &o) const{
   return eq(a * o.x + b * o.y + c, \emptyset);
  cod dist(const line &o) const{
   if(!parallel(o)) return 0;
   if(!eq(o.a * b, o.b * a)) return 0;
   if(!eq(a, 0))
     return abs(c - o.c * a / o.a) / n.len();
   if(!eq(b, 0))
     return abs(c - o.c * b / o.b) / n.len();
   return abs(c - o.c);
  bool parallel(const line &o) const{
   return eq(n ^ o.n, 0);
  bool operator==(const line &o) const{
   if(!eq(a*o.b, b*o.a)) return false;
   if(!eq(a*o.c, c*o.a)) return false;
   if(!eq(c*o.b, b*o.c)) return false;
   return true:
 }
 bool intersect(const line &o) const{
   return !parallel(o) || *this == o;
 vec inter(const line &o) const{
   if(parallel(o)){
     if(*this == 0){ }
     else{ /* dont intersect */ }
   auto tmp = n \hat{o}.n;
   return \{(o.c*b -c*o.b)/tmp, (o.a*c -a*o.c)/tmp\};
 vec at_x(cod x) const{
   return \{x, (-c-a*x)/b\};
 vec at_y(cod y) const{
   return \{(-c-b*y)/a, y\};
 cod eval(const vec &o) const{
   return a * o.x + b * o.y + c;
 }
};
struct segment{
 vec p, q;
 segment(vec a = vec(), vec b = vec()): p(a), q(b)
 bool onstrip(const vec &o) const{ // onstrip strip
```

```
return p.dot(o, q) \geq -eps && q.dot(o, p) \geq -eps
 cod len() const{
   return (p-q).len();
 cod dist(const vec &o) const{
   if(onstrip(o)) return line(p, q).dist(o);
   return min((o-q).len(), (o-p).len());
 bool contains(const vec &o) const{
   return eq(p.cross(q, o), 0) && onstrip(o);
 bool intersect(const segment &o) const{
   if(contains(o.p)) return true;
   if(contains(o.g)) return true;
   if(o.contains(q)) return true;
   if(o.contains(p)) return true;
   return p.ccw(q, o.p) * p.ccw(q, o.q) == -1
       && o.p.ccw(o.q, q) * o.p.ccw(o.q, p) == -1;
 bool intersect(const line &o) const{
   return o.eval(p) * o.eval(q) <= 0;</pre>
  cod dist(const segment &o) const{
   if(line(p, q).parallel(line(o.p, o.q))){
     if(onstrip(o.p) || onstrip(o.q)
     || o.onstrip(p) || o.onstrip(q))
       return line(p, q).dist(line(o.p, o.q));
   else if(intersect(o)) return 0;
   return min(min(dist(o.p), dist(o.q)),
         min(o.dist(p), o.dist(q)));
  cod dist(const line &o) const{
   if(line(p, q).parallel(o))
     return line(p, q).dist(o);
   else if(intersect(o)) return 0:
   return min(o.dist(p), o.dist(q));
 }
};
struct hray{
 vec p, q;
 hray(vec a = vec(), vec b = vec()): p(a), q(b){}
 bool onstrip(const vec &o) const{ // onstrip strip
   return p.dot(q, o) >= -eps;
 cod dist(const vec &o) const{
```

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```
if(onstrip(o)) return line(p, q).dist(o);
 return (o-p).len();
bool intersect(const segment &o) const{
 if(!o.intersect(line(p,q))) return false;
 if(line(o.p, o.q).parallel(line(p,q)))
   return contains(o.p) || contains(o.q);
 return contains(line(p,q).inter(line(o.p,o.q)));
bool contains(const vec &o) const{
 return eq(line(p, q).eval(o), 0) && onstrip(o);
cod dist(const segment &o) const{
 if(line(p, q).parallel(line(o.p, o.q))){
   if(onstrip(o.p) || onstrip(o.q))
     return line(p, q).dist(line(o.p, o.q));
   return o.dist(p);
 else if(intersect(o)) return 0;
 return min(min(dist(o.p), dist(o.q)),
       o.dist(p));
bool intersect(const hray &o) const{
 if(!line(p, q).parallel(line(o.p, o.q)))
   return false:
 auto pt = line(p, q).inter(line(o.p, o.q));
 return contains(pt) && o.contains(pt); // <<</pre>
bool intersect(const line &o) const{
 if(line(p, q).parallel(o)) return line(p, q)== o;
 if(o.contains(p) || o.contains(q)) return true;
 return (o.eval(p) >= -eps)^(o.eval(p)<o.eval(q));</pre>
 return contains(o.inter(line(p, q)));
cod dist(const line &o) const{
 if(line(p,q).parallel(o))
   return line(p,q).dist(o);
 else if(intersect(o)) return 0:
 return o.dist(p);
cod dist(const hray &o) const{
 if(line(p, q).parallel(line(o.p, o.q))){
   if(onstrip(o.p) || o.onstrip(p))
     return line(p,q).dist(line(o.p, o.q));
   return (p-o.p).len();
 else if(intersect(o)) return 0;
 return min(dist(o.p), o.dist(p));
```

```
};
double heron(cod a, cod b, cod c){
 cod s = (a + b + c) / 2:
 return sgrt(s * (s - a) * (s - b) * (s - c));
line mediatrix(const vec &a, const vec &b) {
 auto tmp = (b - a) * 2;
 return line(tmp.x, tmp.y, a * a - b * b);
struct circle {
 vec c: cod r:
  circle() : c(0, 0), r(0) {}
  circle(const vec o) : c(o), r(0) {}
  circle(const vec &a, const vec &b) {
   c = (a + b) * 0.5; r = (a - c).len();
  circle(const vec &a, const vec &b, const vec &cc)
   c = mediatrix(a, b).inter(mediatrix(b, cc));
   r = (a - c).len();
 bool inside(const vec &a) const {
   return (a - c).len() \ll r;
 }
};
circle min_circle_cover(vector<vec> v) {
 random_shuffle(v.begin(), v.end());
  circle ans;
  int n = (int)v.size():
  for(int i = 0; i < n; i++) if(!ans.inside(v[i])) {
   ans = circle(v[i]);
   for(int j = 0; j < i; j++) if(!ans.inside(v[j])){
     ans = circle(v[i], v[j]);
     for(int k=0; k<j; k++)if(!ans.inside(v[k])){
       ans = circle(v[i], v[i], v[k]);
     }
   }
 }
 return ans;
5.2 Circle line intersection
```

```
// intersection of line a * x + b * y + c = 0
// and circle centered at the origin with radius r
double r, a, b, c; // given as input
double x0 = -a*c/(a*a+b*b), y0 = -b*c/(a*a+b*b);
if(c*c > r*r*(a*a+b*b)+EPS)
puts("no points");
else if(abs(c*c - r*r*(a*a+b*b)) < EPS){
```

```
puts("1 point");
  cout << x0 << ' ' << y0 << '\n';
}
else {
    double d = r*r - c*c/(a*a+b*b);
    double mult = sqrt (d / (a*a+b*b));
    double ax, ay, bx, by;
    ax = x0 + b * mult;
    bx = x0 - b * mult;
    ay = y0 - a * mult;
    by = y0 + a * mult;
    puts ("2 points");
    cout<<ax<<' '<<ay<<'\n'<<bx<<' '<<by<<'\n';
}</pre>
```

```
5.3 Half plane intersection
const double eps = 1e-8;
typedef pair<long double. long double> pi;
bool z(long double x){ return fabs(x) < eps; }</pre>
struct line{
 long double a, b, c;
 bool operator<(const line &l)const{</pre>
   bool flag1 = pi(a, b) > pi(0, 0);
   bool flag2 = pi(1.a, 1.b) > pi(0, 0);
   if(flag1 != flag2) return flag1 > flag2;
   long double t = ccw(pi(0, 0), pi(a, b), pi(1.a, 1)
      \hookrightarrow.b)):
   return z(t) ? c * hypot(1.a, 1.b) < 1.c * hypot(a
      \hookrightarrow, b) : t > 0;
 pi slope(){ return pi(a, b); }
pi cross(line a, line b){
 long double det = a.a * b.b - b.a * a.b;
 return pi((a.c * b.b - a.b * b.c) / det, (a.a * b.
    \hookrightarrowc - a.c * b.a) / det);
bool bad(line a, line b, line c){
 if(ccw(pi(0, 0), a.slope(), b.slope()) \le 0)
    →return false;
 pi crs = cross(a, b);
 return crs.first * c.a + crs.second * c.b >= c.c;
bool solve(vector<line> v, vector<pi> &solution){ //
  \hookrightarrow ax + bv <= c:
 sort(v.begin(), v.end());
 deque<line> dq;
 for(auto &i : v){
   if(!dq.empty() && z(ccw(pi(0, 0), dq.back().slope
      \hookrightarrow(), i.slope()))) continue;
```

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```
while(dq.size() >= 2 && bad(dq[dq.size()-2], dq.
     \hookrightarrowback(), i)) dq.pop_back();
 while(dq.size() \geq 2 && bad(i, dq[0], dq[1])) dq.
     \hookrightarrowpop_front();
 dq.push_back(i);
while(dq.size() > 2 && bad(dq[dq.size()-2], dq.
   \hookrightarrowback(), dq[0])) dq.pop_back();
while(dq.size() > 2 && bad(dq.back(), dq[0], dq
  \hookrightarrow[1])) dq.pop_front();
vector<pi> tmp;
for(int i=0; i<dq.size(); i++){</pre>
 line cur = dq[i], nxt = dq[(i+1)%dq.size()];
 if(ccw(pi(0, 0), cur.slope(), nxt.slope()) <= eps</pre>

→) return false;

 tmp.push_back(cross(cur, nxt));
solution = tmp;
return true;
```

#### 5.4 Detect empty Half plane intersection

```
// abs(point a) = absolute value of a
// ccw(a, b, c) = a.ccw(b, c)
pair<bool, point> half_inter(vector<pair<point,point</pre>
   \Rightarrow >  &vet){
   random_shuffle(all(vet));
   point p;
   rep(i,0,sz(vet)) if(ccw(vet[i].x,vet[i].y,p) !=
       \hookrightarrow1){
       point dir = (vet[i].y - vet[i].x) / abs(vet[i
          \hookrightarrow].y - vet[i].x);
       point l = vet[i].x - dir*1e15;
       point r = vet[i].x + dir*1e15;
       if(r < 1) swap(1, r);
       rep(j, 0, i){
           if(ccw(point(), vet[i].x-vet[i].y, vet[j].
              \hookrightarrow x\text{-vet[j].y}) == 0){
               if(ccw(vet[j].x, vet[j].y, p) == 1)
                   continue;
               return mp(false, point());
           if(ccw(vet[j].x, vet[j].y, 1) != 1)
               1 = max(1, line_intersect(vet[i].x,vet
                  \hookrightarrow[i].y,vet[j].x,vet[j].y));
           if(ccw(vet[j].x, vet[j].y, r) != 1)
               r = min(r, line_intersect(vet[i].x,vet
                  \hookrightarrow[i].y,vet[j].x,vet[j].y));
           if(!(1 < r)) return mp(false, point());</pre>
       }
```

```
p = r;
return mp(true, p);
```

#### 5.5 Circle Circle intersection

Assume that the first circle is centered at the origin and second at (x2, y2). Find circle line intersection of first circle and line Ax + By + C = 0, where  $A = -2x_2$ ,  $B = -2y_2$ ,  $C = x_2^2 + y_2^2 + r_1^2 - r_2^2$ .

Be aware of corner case with two circles centered at the same point.

#### 5.6 Tangents of two circles

```
// solve first for same circle(and infinitely many
   \hookrightarrow tangents)
// Find up to four tangents of two circles
void tangents(pt c, double r1, double r2, vector<</pre>
   \hookrightarrowline> & ans){
   double r = r2 - r1;
   double z = c.x * c.x + c.y * c.y;
   double d = z - r * r;
   if(d < -EPS) return;</pre>
   d = sqrt(abs(d));
   line 1;
   1.a = (c.x * r + c.y * d) / z;
   1.b = (c.y * r - c.x * d) / z;
   1.c = r1;
   ans.push_back (1);
vector<line> tangents(circle a, circle b){
   vector<line> ans;
   pt aux = a.center - b.center;
   for(int i = -1; i \le 1; i += 2)
       for(int j = -1; j \le 1; j += 2)
           tangents(aux, a.r * i, b.r * j, ans);
   for(size_t i = 0; i < ans.size(); ++i)
       ans[i].c = ans[i].a * a.x + ans[i].b * a.y;
   return ans;
5.7 Convex Hull
vector<vec> monotone_chain_ch(vector<vec> P){
```

sort(P.begin(), P.end());

// BE CAREFUL WITH OVERFLOW!

vector<vec> L, U;

for(auto p : P){

```
// MAX VALUE (2*A)^2, where 0 \le abs(p.x),
          \hookrightarrow abs(p.y) <= A
       while(L.size() >= 2 && L[L.size() - 2].cross(
         \hookrightarrowL.back(), p) <= 0)
          L.pop_back();
       L.push_back(p);
   reverse(P.begin(), P.end());
   for(auto p : P){
       while(U.size() >= 2 && U[U.size() - 2].cross(
          \hookrightarrow U.back(), p) <= 0)
          U.pop_back();
       U.push_back(p);
   L.pop_back(), U.pop_back();
   L.reserve(L.size() + U.size());
   L.insert(L.end(), U.begin(), U.end());
   return L;
5.8 Check point inside polygon
bool below(const vector<vec> &vet, vec p){
 auto it = lower_bound(vet.begin(), vet.end(), p);
   if(it == vet.end()) return false;
 if(it == vet.begin()) return *it == p;
 return prev(it)->cross(*it, p) <= 0;</pre>
bool above(const vector<vec> &vet, vec p){
 auto it = lower_bound(vet.begin(), vet.end(), p);
   if(it == vet.end()) return false;
 if(it == vet.begin()) return *it == p;
 return prev(it)->cross(*it, p) >= 0;
// lowerhull, upperhull and point, borders included
bool inside_poly(const vector<vec> &lo, const vector
  \hookrightarrow<vec> &hi, vec p){
 return below(hi, p) && above(lo, p);
5.9 Check point inside polygon without lower/upper
```

hull

// borders included

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```
// must not have 3 colinear consecutive points
bool inside_poly(const vector<vec> &v, vec p){
   if(v[0].ccw(v[1], p) < 0) return false;
   if(v[0].ccw(v.back(), p) > 0) return 0;
   if(v[0].ccw(v.back(), p) == 0)
       return v[0].dot(p, v.back()) >= 0
          && v.back().dot(p, v[0]) >= 0;
   int L = 1, R = (int)v.size() - 1, ans = 1;
   while(L <= R){</pre>
       int mid = (L+R)/2:
       if(v[0].ccw(v[mid], p) >= 0) ans = mid, L =
          \hookrightarrowmid+1;
       else R = mid-1;
   }
   return v[ans].ccw(v[(ans+1)%v.size()], p) >= 0;
}
```

#### 5.10 Minkowski sum

```
vector<vec> mk(const vector<vec>&a,const vector<vec</pre>
  ⇒>&b){
   int i = 0, j = 0;
   for(int k = 0; k < (int)a.size(); k++)if(a[k] < a[
      →i])
      i = k:
   for(int k = 0; k < (int)b.size(); k++)if(b[k] < b[
      j])
      j = k;
   vector<vec> c;
   c.reserve(a.size() + b.size());
   for(int k = 0; k < int(a.size()+b.size()); k++){
      vec pt{a[i] + b[j]};
      if((int)c.size() >= 2
       && c[c.size()-2].ccw(c.back(), pt) == 0)
          c.pop_back();
      c.push_back(pt);
      int q = i+1, w = j+1;
      if(q == int(a.size())) q = 0;
      if(w == int(b.size())) w = 0;
      if(c.back().ccw(a[i]+b[w], a[q]+b[j]) < 0) i
         \hookrightarrow = q;
       else j = w;
   c.shrink_to_fit();
   return c;
```

#### 5.11 Geo Notes

#### 5.11.1 Center of mass

System of points(2D/3D): Mass weighted average of points.

Frame(2D/3D): Get middle point of each segment solve as previously.

**Triangle:** Average of vertices.

**2D Polygon:** Compute **signed** area and center of mass of triangle  $((0,0), p_i, p_{i+1})$ . Then solve as system of points.

**Polyhedron surface:** Solve each face as a 2D polygon(be aware of (0, 0)) then replace each face with its center of mass and solve as system of points.

Tetrahedron(Triangular pyramid): As triangles, its the average of points.

**Polyhedron:** Can be done as 2D polygon, but with tetrahedralization intead of triangulation.

#### 5.11.2 Pick's Theorem

Given a polygon without self-intersections and all its vertices on integer coordinates in some 2D grid. Let A be its area, I the number of points with integer coordinates stricly inside the polygon and B the number of points with integer coordinates in the border of the polygon. The following formula holds:  $A = I + \frac{B}{2} - 1$ .

#### 6 Miscellaneous

#### 6.1 LIS

```
multiset<int> S;
for(int i = 0; i < n; i++){
 auto it = S.upper_bound(a[i]); // low for inc
 if(it != S.end()) S.erase(it);
 S.insert(a[i]);
ans = S.size();
```

#### 6.2 DSU rollback

```
struct DSU{
 vector<int> sz, p, change;
 vector<tuple<int, int, int>> modifications;
 vector<size_t> saves;
 bool bipartite;
 DSU(int n): sz(n+1, 1), p(n+1), change(n+1),
    ⇒bipartite(true){
   iota(p.begin(), p.end(), 0);
 }
```

```
void add_edge(int u, int v){
  if(!bipartite) return;
  int must_change = get_colour(u) == get_colour(v);
  int a = rep(u), b = rep(v);
  if(sz[a] < sz[b]) swap(a, b);
 if(a != b){
   p[b] = a;
   modifications.emplace_back(b, change[b],
      \hookrightarrowbipartite);
   change[b] ^= must_change;
   sz[a] += sz[b];
  else if(must_change){
   modifications.emplace_back(0, change[0],
      ⇒bipartite);
   bipartite = false;
int rep(int u){
 return p[u] == u ? u : rep(p[u]);
int get_colour(int u){
 if(p[u] == u) return change[u];
 return change[u] ^ get_colour(p[u]);
void reset(){
  modifications.clear();
  saves.clear();
  iota(p.begin(), p.end(), 0);
  fill(sz.begin(), sz.end(), 1);
  fill(change.begin(), change.end(), 0);
  bipartite = true;
}
void rollback(){
  int u = get<0>(modifications.back());
  tie(ignore, change[u], bipartite) = modifications
    \hookrightarrow.back();
  sz[ p[u] ] -= sz[u];
 p[u] = u;
  modifications.pop_back();
void reload(){
  while(modifications.size() > saves.back())
   rollback();
```

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```
saves.pop_back();
 void save(){
   saves.push_back(modifications.size());
};
6.3 Buildings
// count the number of circular arrays of size m,
  \hookrightarrow with elements on range [1, c^{**}(n^*n)]
int n, m, c; cin >> n >> m >> c;
int x = f_{exp}(c, n * n); int ans = f_{exp}(x, m);
for(int i = 1; i \le m; i++) if(m % i == 0) {
 int y = f_{exp}(x, i);
 for(int j = 1; j < i; j++) if(i % j == 0)
     y = sub(y, mult(j, dp[j]));
 dp[i] = mult(y, inv(i));
 ans = sub(ans, mult(i - 1, dp[i]));
cout << ans << '\n';
6.4 Rand
#include <random>
#include <chrono>
cout << RAND MAX << endl:</pre>
mt19937 rng(chrono::steady_clock::now().
  →time_since_epoch().count());
shuffle(p.begin(), p.end(), rng);
uniform_int_distribution<int>(a,b)(rng);
6.5 Klondike
// minimum number of moves to make
// all elements equal
// move: change a segment of equal value
// elements to any value
int v[305], dp[305][305], rec[305][305];
int f(int 1, int r){
 if(r == 1) return 1;
 if(r < 1) return 0;
 if(dp[l][r] != -1) return dp[l][r];
 int ans = f(1+1, r) + 1;
 for(int i = l+1; i \le r; i++)
   if(v[i] == v[l])
     ans = min(ans, f(1, i - 1) + f(i+1, r));
 return dp[l][r] = ans;
```

6.6 Hilbert Order

```
// maybe use B = n / sqrt(q)
inline int64_t hilbertOrder(int x, int y, int pow =
  \hookrightarrow21, int rotate = 0) {
 if(pow == 0) return 0;
 int hpow = 1 \ll (pow-1);
 int seq = (x < hpow) ? (
   (y < hpow) ? 0 : 3
 ):(
   (y < hpow) ? 1 : 2
 ):
  seg = (seg + rotate) & 3;
  const int rotateDelta[4] = {3, 0, 0, 1};
  int nx = x & (x ^ hpow), ny = y & (y ^ hpow);
  int nrot = (rotate + rotateDelta[seq]) & 3;
 int64_t subSquareSize = int64_t(1) << (2*pow - 2);</pre>
 int64_t ans = seg * subSquareSize;
 int64_t add = hilbertOrder(nx, ny, pow-1, nrot);
  ans += (seg == 1 || seg == 2) ? add : (
    ⇒subSquareSize - add - 1);
 return ans:
```

#### 6.7 Modular Factorial

```
// Compute (1*2*...*(p-1)*1*(p+1)*(p+2)*..*n) % p
// in O(p*lg(n))
int factmod(int n, int p){
   int ans = 1;
   while (n > 1)
       for(int i = 2; i <= n % p; i++)
           ans = (ans * i) % p;
       n /= p;
       if(n \% 2) ans = p - ans;
   }
   return ans % p;
int fac_pow(int n, int p){
   int ans = 0;
   while(n) n \neq p, ans += n;
   return ans;
int C(int n, int k, int p){
   if(fac_pow(n, p) > fac_pow(n-k, p) + fac_pow(k, p)
      \hookrightarrow))
       return 0;
   int tmp = factmod(k, p) * factmod(n-k, p) % p;
   return (f_{exp}(tmp, p - 2, p) * factmod(n, p)) % p
```

#### 6.8 Enumeration all submasks of a bitmask

#### 6.9 Slope Trick

#### 6.10 Knapsack Bounded with Cost

```
// menor custo para conseguir peso ate M usando N
  ⇒tipos diferentes de elementos, sendo que o i-
  ⇒esimo elemento pode ser usado b[i] vezes, tem
  \hookrightarrow peso w[i] e custo c[i]
// O(N * M)
int b[N], w[N], c[N];
MinQueue Q[M]
int d[M] //d[i] = custo minimo para conseguir peso i
for(int i = 0; i \le M; i++) d[i] = i ? oo : 0;
for(int i = 0; i < N; i++){
 for(int j = 0; j < w[i]; j++)
   Q[j].clear();
 for(int j = 0; j \le M; j++){
   q = Q[i \% w[i]];
   if(q.size() >= q) q.pop();
   q.add(c[i]);
   q.push(d[j]);
```

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```
6.11 LCA < O(nlgn), O(1)>
int start[N], dfs_time;
int tour[2*N], id[2*N];
void dfs(int u){
   start[u] = dfs_time;
   id[dfs_time] = u;
   tour[dfs_time++] = start[u];
   for(int v : g[u]){
      dfs(v);
      id[dfs_time] = u;
      tour[dfs_time++] = start[u];
   }
}
int LCA(int u, int v){
   if(start[u] > start[v]) swap(u, v);
   return id[min(tour[k] for k in [start[u], start[v
      \hookrightarrow]])];
6.12 Buffered reader
// source: https://github.com/ngthanhtrung23/
  → ACM_Notebook_new/blob/master/buffered_reader.h
int INP,AM,REACHEOF;
#define BUFSIZE (1<<12)</pre>
char BUF[BUFSIZE+1], *inp=BUF;
#define GETCHAR(INP) { \
   if(!*inp && !REACHEOF) { \
      memset(BUF,0,sizeof BUF);\
      int inpzzz = fread(BUF,1,BUFSIZE,stdin);\
      if (inpzzz != BUFSIZE) REACHEOF = true;\
      inp=BUF; \
   } \
   INP=*inp++; \
#define DIG(a) (((a)>='0')\&\&((a)<='9'))
```

GETCHAR(INP); while(!DIG(INP) && INP!='-')

while(DIG(INP)){j=10\*j+(INP-'0');GETCHAR(INP);} \

if (INP=='-') {AM=1;GETCHAR(INP);} \

d[j] = q.getmin();

#define GN(j) { \

**if** (AM)  $j=-j;\$ 

 $\hookrightarrow$  GETCHAR(INP);\

j=INP-'0'; GETCHAR(INP); \

 $AM=0; \$ 

}

#### 6.13 Modular summation

```
//calcula (sum(0 \ll i \ll n) P(i)) \% mod.
//onde P(i) eh uma PA modular (com outro modulo)
namespace sum_pa_mod{
    11 calc(ll a, ll b, ll n, ll mod){
          assert(a&&b);
          if(a >= b){
               11 ret = ((n*(n+1)/2) \mod)*(a/b);
               if(a\%b) ret = (ret + calc(a\%b,b,n,mod))\%mod;
               else ret = (ret+n+1) \mod;
               return ret:
          return ((n+1)*(((n*a)/b+1))*(nod) - calc(b,a,(n*a)/b+1)*(nod) - calc(b,a,(n*a)/b+1
                  \hookrightarrowb,mod) + mod + n/b + 1)%mod;
     }
     //P(i) = a*i \mod m
     11 solve(l1 a, l1 n, l1 m, l1 mod){
          a = (a\%m + m)\%m;
          if(!a) return 0;
          11 ret = (n*(n+1)/2)%mod;
          ret = (ret*a)%mod;
          11 g = \_gcd(a,m);
          ret -= m*(calc(a/g,m/g,n,mod)-n-1);
          return (ret%mod + mod)%mod;
     //P(i) = a + r*i \mod m
     11 solve(ll a, ll r, ll n, ll m, ll mod){
          a = (a\%m + m)\%m;
          r = (r\%m + m)\%m;
          if(!r) return (a*(n+1))%mod;
          if(!a) return solve(r, n, m, mod);
          11 g, x, y;
           g = gcdExtended(r, m, x, y);
          x = (x\%m + m)\%m;
          11 d = a - (a/q)*q;
           a -= d:
          x = (x*(a/g))%m;
          return (solve(r, n+x, m, mod) - solve(r, x-1, m,
                  \hookrightarrow mod) + mod + d*(n+1))%mod;
    }
};
6.14 Edge coloring CPP
const int MX = 300;
int C[MX][MX] = {}, G[MX][MX] = {};
```

void solve(vector<pii> &E, int N){

```
int X[MX] = \{\}, a, b;
auto update = [\&](int u)\{for(X[u] = 1; C[u][X[u])\}
  \hookrightarrow]]; X[u]++); };
auto color = [&](int u, int v, int c){
   int p = G[u][v]:
   G[u][v] = G[v][u] = c;
   C[u][c] = v; C[v][c] = u;
   C[u][p] = C[v][p] = 0;
   if(p) X[u] = X[v] = p;
   else update(u), update(v);
   return p: }:
auto flip = [\&] (int u, int c1, int c2){
   int p = C[u][c1], q = C[u][c2];
   swap(C[u][c1], C[u][c2]);
   if(p) G[u][p] = G[p][u] = c2;
   if( !C[u][c1] ) X[u] = c1;
   if( !C[u][c2] ) X[u] = c2;
   return p; };
for(int i = 1; i \le N; i++) X[i] = 1;
for(int t = 0; t < E.size(); t++){</pre>
   int u = E[t].first, v0 = E[t].second, v = v0,
      \hookrightarrow c0 = X[u], c = c0, d;
   vector<pii> L;
   int vst[MX] = {};
   while(!G[u][v0]){
       L.emplace_back(v, d = X[v]);
       if(!C[v][c]) for(a = (int)L.size()-1; a >=
          \hookrightarrow 0; a--) c = color(u, L[a].first, c);
       else if(!C[u][d])for(a=(int)L.size()-1;a
          \Rightarrow =0; a--)color(u,L[a].first,L[a].
          ⇒second):
       else if( vst[d] ) break;
       else vst[d] = 1, v = C[u][d];
   if( !G[u][v0] ){
       for(;v; v = flip(v, c, d), swap(c, d));
       if(C[u][c0]){
           for(a = (int)L.size()-2; a >= 0 \&\& L[a
             \hookrightarrow].second != c; a--);
           for(; a >= 0; a--) color(u, L[a].first
             \hookrightarrow, L[a].second);
       } else t--;
   }
```

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#### 6.15 Burnside's Lemma

Let  $(G, \oplus)$  be a finite group that acts on a set X. It should hold that  $e_g * x = x$  and  $g_1 * (g_2 * x) = (g_1 \oplus g_2) * x$ ,  $\forall x \in X$ ,  $g_1, g_2 \in G$ . For each  $g \in G$  let  $X^g = \{x \in X \mid g * x = x\}$ . The number of orbits its given by:

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|$$

#### 6.16 Wilson's Theorem

 $(n-1)! = -1 \mod n \iff n \text{ is prime}$ 

#### 6.17 Fibonacci

- $F_{n-1}F_{n+1} F_n^2 = (-1)^n$
- $F_{n+k} = F_k F_{n+1} + F_{k-1} F_n$
- $GCD(F_n, F_m) = F_{GCD(n,m)}$
- $F_n = \frac{(\frac{1+\sqrt{5}}{2})^n (\frac{1-\sqrt{5}}{2})^n}{\sqrt{5}}$

#### 6.18 Lucas's Theorem

For non-negative integers m and n and a prime p, the following congruence holds:

$$\binom{m}{n} \equiv \prod_{i=0}^{k} \binom{m_i}{n_i} \pmod{p}$$

where  $m_i$  is the i-th digit of m in base p.  $\binom{a}{b} = 0$  if a < b.

#### 6.19 Kirchhoff's Theorem

Laplacian matrix is L = D - A, where D is a diagonal matrix with vertex degrees on the diagonals and A is adjacency matrix.

The number of spanning trees is any cofactor of L. i-th cofactor is determinant of the matrix gotten by removing i-th row and column of L.

### 6.19.1 Multigraphs

In D[i][i] all loops are excluded. A[i][j] = number of edges from i to j.

### 6.19.2 Directed multigraphs

D[i][i] = indegree of i minus the number of loops at i. A[i][j] = number of edges from i to j.

The number of oriented spanning trees rooted at a vertex i is the determinant of the matrix gotten by removing the ith row and column of L.

#### 6.20 Matroid

Let *X* set of objects,  $I \subseteq 2^X$  set of independents sets such that:

- 1.  $\emptyset \in I$
- 2.  $A \in I.B \subseteq A \implies B \in I$
- 3. Exchange axiom,  $A \in I, B \in I, |B| > |A| \implies \exists x \in B \setminus A : A \cup \{x\} \in I$
- 4.  $A \subseteq X$  and I and I' are maximal independent subsets of A then |I| = |I'|

Then (X, I) is a matroid. The combinatorial optimization problem associated with it is: Given a weight  $w(e) \ge 0 \ \forall e \in X$ , find an independet subset that has the largest possible total weight.

#### 6.21 Matroid intersection

```
// Input two matroids (X, I_a) and (X, I_b)
// output set I of maximum size, I \in I_a and I \in
  \hookrightarrow I_b
set<> I;
while(1){
   for(e_i : X \setminus I)
       if(I + e_i \in I_a \text{ and } I + e_i \in I_b)
           I = I + e_i;
   set<> A, T; queue<> Q;
   for(x : X) label[x] = MARK1;
   for(e_i : X \setminus I){
       if(I + e_i \in I_a)
           Q.push(e_i), label[e_i] = MARK2;
       else{
           for(x such that I - x + e_i \setminus I_a)
               A[x].push(e_i);
       if(I + e_i \in I_b)
           T = T + \{e_i\}
           for (x \text{ such that } I - x + e_i \in I_b)
               A[e_i].push(x);
       }
   if(T.empty()) break;
   bool found = false;
   while(!Q.empty() and !found){
```

```
auto e = Q.front(); Q.pop();
for(x : A[e]) if(label[x] == MARK1){
    label[x] = e; Q.push(x);
    if(x \in T){
        found = true; put = 1;
        while(label[x] != MARK2){
            I = put ? (I + x) : (I - x);
            put = 1 - put;
        }
        I = I + x;
        break;
    }
}
if(!found) break;
}
return I;
```

Where path(e) = [e] if label[e] = MARK2, path(label[e]) + [e] otherwise.

#### 6.21.1 Matroid Union

Given k matroids over the same set of objects  $(X, I_1)$ ,  $(X, I_2)$ , ...,  $(X, I_k)$  find  $A_1 \in I_1$ ,  $A_2 \in I_2$ , ...,  $A_k \in I_k$  such that  $i \neq j$ ,  $A_i \cap A_j = \emptyset$  and  $\bigcup_{i=1}^k A_i$  is maximum. Matroid union can be reduced to matroid intersection as follows.

Let  $X' = X \times \{1, 2, ..., k\}$ , ie, k copies of each element of X with different colors. M1 = (X', Q) where  $B \in Q \iff \forall 1 \le i \le k$ ,  $\{x \mid (x, i) \in B\} \in I_i$ , ie, for each color, B is independent. M2 = (X', W) where  $B \in W \iff i \ne j \implies \neg((x, i) \in B \land (x, j) \in B)$ , ie, each element is picked by at most one color.

Intersection of *M*1 and *M*2 is the answer for the combinatorial problem of matroid union.

#### 6.22 Notes

When we repeat something and each time we have probability p to succeed then the expected number or tries is  $\frac{1}{n}$ , till we succeed.

## Small to large

**Trick in statement** If k sets are given you should note that the amount of different set sizes is  $O(\sqrt{s})$  where s is total size of those sets. And no more than  $\sqrt{s}$  sets have size greater than  $\sqrt{s}$ . For example, a path to the root

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in Aho-Corasick through suffix links will have at most  $O(\sqrt{s})$  vertices.

**gcd on subsegment**, we have at most  $\log(a_i)$  different values in  $\{\gcd(a_j, a_{j+1}, ..., a_i) \text{ for } j < i\}$ .

From static set to expandable. To insert, create a new set with the new element. While there are two sets with same size, merge them. There will be at most  $\log(n)$  disjoints sets.

**Matrix exponentiation optimization**. Save binary power of  $A_{nxn}$  and answer q queries  $b = A^m x$  in  $O((n^3 + qn^2)log(m))$ .

**Ternary search on integers into binary search**, comparing f(mid) and f(mid+1), binary search on deriva-

tive

**Dynamic offline set** For each element we will wind segment of time [a, b] such that element is present in the set during this whole segment. Now we can come up with recursive procedure which handles [l, r] time segment considering that all elements such that  $[l, r] \subset [a, b]$  are already included into the set. Now, keeping this invariant we recursively go into [l, m] and [m + 1, r] subsegments. Finally when we come into segment of length 1.

 $a > b \implies a \mod b < \frac{a}{2}$ 

**Convex Hull**. The expected number of points in the convex hull of a random set of points is O(log(n)). The number of points in a convex hull with points coordinates 264241152, 382205952, 530841600

limited by *L* is  $O(L^{2/3})$ .

**Tree path query**. Sometimes the linear query is fast enough. Just do adamant's hld sorting subtrees by their size and remap vertices indexes.

**Range query** offline can be solved by a sweep, ordering queries by R.

Maximal number of divisors of any n-digit number. 7 4, 12, 32, 64, 128, 240, 448, 768, 1344, 2304, 4032, 6720, 10752, 17280, 26880, 41472, 64512, 103680, 161280, 245760, 368640, 552960, 860160, 1290240, 1966080, 2764800, 4128768, 6193152, 8957952, 13271040, 19660800, 28311552, 41287680, 59719680, 88473600, 127401984, 181665792, 264241152, 382205952, 530841600