University of Brasilia CONTENTS, 1

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# University of Brasilia

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University of Brasilia Data Structures, 2

```
set ts=4 sw=4 sta nu rnu sc stl+=%F cindent
set bg=dark ruler timeoutlen=1000
imap {<CR> {<CR>}<Esc>0
nmap <F2> 0V$%d
nmap <C-down> :m+1<CR>
nmap <C-up> :m-2<CR>
nmap <C-a> ggVG
nmap < S-up > :m-2 < CR >
nmap <S-down> :m+1<CR>
syntax on
vmap <C-c> "+y
set viminfo='20,\"1000
alias comp='q++ -std=c++17 -Wshadow -Wall -Wextra -Wformat=2 -
   →Wconversion -fsanitize=address,undefined -fno-sanitize-recover

— -Wfatal-errors'

#include <bits/stdc++.h>
#define ff first
#define ss second
#define pb push_back
using namespace std;
using 11 = long long;
using ii = pair<int, int>;
const int N = 100005;
int main() {

  return 0;

1 Data Structures
1.1 Fenwick Tree 2D
a82442781f9b43b07f5f5425c5308313, 31 lines
vector<int> go[N];
vector<int> ft[N];
void prec_add(int x, int y) {
\triangleright for(; x < N; x += x & -x) {
▷ □ go[x].push_back(y);
⊳ }
void init() {
\triangleright for(int i = 1; i < N; i++) {
▷ ▷ sort(go[i].begin(), go[i].end());
▷ ▷ go[i].resize(unique(go[i].begin(), go[i].end()) - go[i].begin())
b ft[i].assign(go[i].size() + 1, 0);
void add(int x, int y, int val) {
```

▷ int id = int(upper\_bound(go[x].begin(), go[x].end(), y) - go[x].

▷ b for(; id < (int)ft[x].size(); id += id & -id)</pre>

 $\triangleright$  for(; x < N; x += x & -x) {

▷ ▷ ▷ ft[x][id] += val;

int sum(int x, int y) {

 $\triangleright$  for(; x > 0; x -= x & -x) {

 $\triangleright$  int ans = 0;

⊳ }

```
\triangleright int id = int(upper_bound(go[x].begin(), go[x].end(), y) - go[x].
\triangleright \triangleright \triangleright ans += ft[x][id];
⊳ }

  return ans;
1.2 Wavelet Tree
9005683d2c15117fc322db509aa6299a, 65 lines
template<typename T>
class wavelet { // 1-based!!
   T L, R;
   vector<int> 1;
▷ vector<T> sum: // <<</pre>

    wavelet *lef, *rig;

b int r(int i) const{ return i - l[i]; }
public:

    template<typename ITER>

   wavelet(ITER bg, ITER en) { // it changes the argument array
▷ ▷ lef = rig = nullptr;
      L = *bq, R = *bq;
▷ for(auto it = bg; it != en; it++)
         L = min(L, *it), R = max(R, *it);
▷ ▷ if(L == R) return;
      T \text{ mid} = L + (R - L)/2;
▷ ▷ l.reserve(std::distance(bg, en) + 1);
▷ ▷ sum.reserve(std::distance(bg, en) + 1);
▷ ▷ l.push_back(0), sum.push_back(0);
▷ ▷ sum.push_back(sum.back() + *it);
▷ ▷ ▷ return x <= mid;
▷ ▷ });
⊳ ~wavelet(){
⊳ ⊳ delete lef:

▷ b delete rig;

⊳ }
⊳ // 1 index, first is 1st
   T kth(int i, int j, int k) const{
      if(L >= R) return L;
      int c = l[j] - l[i-1];
      if(c >= k) return lef > kth(l[i-1]+1, l[j], k);
      else return rig->kth(r(i-1)+1, r(j), k - c);
\triangleright // # elements > x on [i, i]
p int cnt(int i, int j, T x) const{
\triangleright if(L > x) return j - i + 1;
\triangleright int ans = 0;
\Rightarrow if(lef) ans += lef->cnt(l[i-1]+1, l[j], x);
\triangleright if(rig) ans += rig->cnt(r(i-1)+1, r(j), x);
⊳ ⊳ return ans:
⊳ }
▷ // sum of elements <= k on [i, j]</pre>

    T sumk(int i, int j, T k){
      if(L == R) return R <= k ? L * (j - i + 1) : 0;
```

```
\triangleright \triangleright int ans = 0:
\triangleright if(lef) ans += lef->sumk(l[i-1]+1, l[j], k);
\Rightarrow if(rig) ans += rig->sumk(r(i-1)+1, r(j), k);
⊳ ⊳ return ans;
⊳ }
▷ // swap (i, i+1) just need to update "array" 1[i]
};
1.3 Order Set
4e7ba1c597697798d1fad502cddf47cf, 8 lines
#include <bits/extc++.h>
using namespace __gnu_pbds; // or pb_ds;
template<typename T, typename B = null_type>
using oset = tree<T, B, less<T>, rb_tree_tag,

→tree_order_statistics_node_update>;
   find_by_order / order_of_key
1.4 Hash table
9507f2c486ea07a5bffb6f7635077f0d, 22 lines
#include <ext/pb_ds/assoc_container.hpp>
using namespace __gnu_pbds;
struct custom hash {
static uint64_t splitmix64(uint64_t x) {
▷ ▷ // http://xorshift.di.unimi.it/splitmix64.c
\Rightarrow x += 0x9e3779b97f4a7c15:
\triangleright x = (x \hat{x} >> 30)) * 0xbf58476d1ce4e5b9;
\triangleright x = (x \hat{x} > 27) * 0x94d049bb133111eb;
\triangleright return x ^ (x >> 31);
⊳ }
> size_t operator()(uint64_t x) const {
▷ static const uint64_t FIXED_RANDOM = chrono::steady_clock::now()
   →.time since epoch().count():
⊳ }
};
gp_hash_table<long long, int, custom_hash> table;
unordered_map<long long, int, custom_hash> uhash;
uhash.reserve(1 << 15);</pre>
uhash.max_load_factor(0.25);
1.5 Convex Hull Trick Simple
8014e282d7cdc23df5456b1927e26a3c, 42 lines
struct Line{
   11 m, b;
   inline 11 eval(11 x) const{
       return x * m + b;
};
// min => cht.back().m >= L.m
// max => cht.back().m <= L.m
void push_line(vector<Line> &cht, Line L){
  while((int)cht.size() >= 2){
   int sz = (int)cht.size();
   if((long double)(L.b-cht[sz-1].b)*(cht[sz-2].m-L.m)
   <= (long double)(L.b-cht[sz-2].b)*(cht[sz-1].m-L.m)){</pre>
     cht.pop_back();
   else break;
 cht.push_back(L);
 // x increasing; pos = 0 in first call
```

University of Brasilia Data Structures, 3

```
11 linear_search(const vector<Line> &cht,ll x,int &pos){
   while(pos+1 < (int)cht.size()){</pre>
/*>>*/ if(cht[pos].eval(x) >= cht[pos+1].eval(x)) pos++;
       else break:
   return cht[pos].eval(x);
11 binary_search(const vector<Line> &cht, 11 x){
   int L = 0. R = (int)cht.size()-2:
   int bans = (int)cht.size()-1;
   while(L <= R){</pre>
       int mid = (L+R)/2;
       if(cht[mid].eval(x) >= cht[mid+1].eval(x)) // <<<</pre>
          L = mid + 1:
       else bans = mid, R = mid - 1;
   return cht[bans].eval(x);
```

### 1.6 Convex Hull Trick

58f8d280705111c564412234b85d3c0b, 36 lines

```
const 11 is_query = -(1LL<<62);</pre>
struct Line{
⊳ 11 m. b:

    mutable function < const Line*() > succ;
▶ bool operator<(const Line& rhs) const{</pre>
▷ if(rhs.b != is_query) return m < rhs.m;</pre>
▷ ▷ const Line* s = succ();
▷ ▷ if(!s) return 0;
\triangleright \triangleright 11 x = rhs.m:
\triangleright return b - s->b < (s->m - m) * x;
⊳ }
};
struct Cht : public multiset<Line>{ // maintain max
▶ bool bad(iterator y){
\triangleright \triangleright auto z = next(v):
\triangleright \triangleright if(y == begin()){
\triangleright \triangleright \vdash if(z == end()) return 0;
\triangleright \triangleright return y->m == z->m && y->b <= z->b;
\triangleright \triangleright auto x = prev(y);
\triangleright if(z == end()) return y->m == x->m && y->b <= x->b;
\triangleright return (long double)(x->b - y->b)*(z->m - y->m) >= (long double)
    (y->b-z->b)*(y->m-x->m);
▶ void insert_line(ll m, ll b){
▷ ▷ auto y = insert({ m, b });
▷ if(bad(y)){ erase(y); return; }
⊳ }
⊳ ll eval(ll x){
▷ ▷ auto 1 = *lower_bound((Line) { x, is_query });
\triangleright return 1.m * x + 1.b;
⊳ }
};
```

### 1.7 Convex Hull Trick

39c065bc0c78205ce1cb7e5889083e17, 36 lines

```
* Author: Simon Lindholm
* source: https://github.com/kth-competitive-programming/kactl/
   ⇒blob/master/content/data-structures/LineContainer.h
* License: CC0
*/
```

```
struct Line {
⊳ mutable 11 m. b. p:
▷ bool operator<(const Line& o) const { return m < o.m; }</pre>

    bool operator<(11 x) const { return p < x; }</pre>
struct LineContainer : multiset<Line, less<>> { // CPP14 only
\triangleright // (for doubles, use inf = 1/.0, div(a,b) = a/b)

    const 11 inf = LLONG_MAX;
▶ bool isect(iterator x, iterator y) {
\triangleright if (y == end()) { x->p = inf; return false; }
\Rightarrow if (x->m == y->m) x->p = x->b > y->b ? inf : -inf;
\triangleright else x->p = div(y->b - x->b, x->m - y->m);
\triangleright return x->p >= y->p;
⊳ }

  void add(ll m, ll b) {
\triangleright auto z = insert({m, b, 0}), y = z++, x = y;
\triangleright while (isect(y, z)) z = erase(z);
\triangleright if (x != begin() && isect(--x, y)) isect(x, y = erase(y));
\Rightarrow while ((y = x) != begin() && (--x)->p >= y->p)
▷ ▷ isect(x, erase(y));
⊳ }
⊳ ll query(ll x) {
▷ ▷ assert(!empty());
\triangleright auto 1 = *lower_bound(x);
\triangleright return 1.m * x + 1.b;
⊳ }
};
1.8 Min queue
d27721ce5326145f04ded2c962063add, 23 lines
template<typename T>
class minQ{

    deque<tuple<T, int, int> > p;

→ T delta;

⊳ int sz;
public:
p minQ() : delta(0), sz(0) {}
p inline int size() const{ return sz; }

    inline void add(T x) { delta += x; }

    inline void push(T x, int id){
⊳ ⊳ x -= delta. sz++:
\triangleright \triangleright int t = 1;
b while(p.size() > 0 && get<0>(p.back()) >= x)
▷ ▷ t += get<1>(p.back()), p.pop_back();
▷ ▷ p.emplace_back(x, t, id);
⊳ }

   inline void pop(){

    p get<1>(p.front())--, sz--;
b if(!get<1>(p.front())) p.pop_front();
⊳ }
> T getmin() const{ return get<0>(p.front())+delta; }
p int getid() const{ return get<2>(p.front()); }
};
1.9 Sparse Table
e1916d98125f87032c8186bc10ca682e, 10 lines
int fn(int i, int j){
b if(j == 0) return v[i];
b if(~dn[i][j]) return dn[i][j];
\triangleright return dn[i][j] = min(fn(i, j-1), fn(i + (1 << (j-1)), j-1));
int getmn(int 1, int r) \{ // [1, r] \}
\triangleright int 1z = 1q(r - 1 + 1);
\triangleright return min(fn(1, lz), fn(r - (1 << lz) + 1, lz));
```

### 1.10 Treap

6299f3109fc13d0047223cc339e6bdb7, 64 lines

```
// source: https://github.com/victorsenam/caderno/blob/master/code/
   \hookrightarrowtreap.cpp
//const int N = ; typedef int num;
num X[N]; int en = 1, Y[N], sz[N], L[N], R[N];
void calc (int u) { // update node given children info

    if(!u) return:
\triangleright sz[u] = sz[L[u]] + 1 + sz[R[u]];
⊳ // code here, no recursion
void unlaze (int u) {

    if(!u) return;

⊳ // code here, no recursion
void split_val(int u, num x, int &l, int &r) { // l gets <= x, r</pre>
\triangleright unlaze(u); if(!u) return (void) (1 = r = 0);
b if(X[u] <= x) { split_val(R[u], x, 1, r); R[u] = 1; 1 = u; }

    else { split_val(L[u], x, 1, r); L[u] = r; r = u; }

void split_sz(int u, int s, int &l, int &r) { // l gets first s, r
    \hookrightarrow gets remaining
\triangleright unlaze(u); if(!u) return (void) (l = r = 0);
\Rightarrow if(sz[L[u]] < s) { split_sz(R[u], s - sz[L[u]] - 1, 1, r); R[u] =
   \hookrightarrow 1; 1 = u; }
else { split_sz(L[u], s, l, r); L[u] = r; r = u; }
⊳ calc(u):
int merge(int 1, int r) { // els on 1 <= els on r</pre>
\triangleright unlaze(1); unlaze(r); if(!1 || !r) return 1 + r; int u;
\triangleright if(Y[1] > Y[r]) { R[1] = merge(R[1], r); u = 1; }

    else { L[r] = merge(1, L[r]); u = r; }

⊳ calc(u); return u;
void init(int n=N-1) { // XXX call before using other funcs
\triangleright for(int i = en = 1; i <= n; i++) { Y[i] = i; sz[i] = 1; L[i] = R[
    \hookrightarrowi] = 0; }
\triangleright random_shuffle(Y + 1, Y + n + 1);
void insert(int &u. int it){
□ unlaze(u):
\triangleright if(!u) u = it;

    else if(Y[it] > Y[u]) split_val(u, X[it], L[it], R[it]), u = it;

    else insert(X[it] < X[u] ? L[u] : R[u], it);</pre>
⊳ calc(u);
void erase(int &u, num key){
□ unlaze(u):

    if(!u) return:
\triangleright if(X[u] == key) u = merge(L[u], R[u]);

    else erase(key < X[u] ? L[u] : R[u], key);
</pre>
⊳ calc(u);
int create_node(num key){
\triangleright X[en] = key;
\triangleright sz[en] = 1;
\triangleright L[en] = R[en] = 0;
⊳ return en++;
int query(int u, int 1, int r){//0 index

  unlaze(u);

    if(u! or r < 0 or l >= sz[u]) return identity_element;

\triangleright if(1 <= 0 and r >= sz[u] - 1) return subt_data[u];
\triangleright int ans = query(L[u], 1, r);
```

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```
\triangleright if(1 <= sz[ L[u] ] and sz[ L[u] ] <= r)
\triangleright ans = max(ans. st[u]):
\rightarrow ans = max(ans, query(R[u], l-sz[L[u]]-1, r-sz[L[u]]-1));
▶ return ans:
1.11 ColorUpdate
5b91213c31a8762d831ed083b04fad35, 43 lines
// source: https://github.com/tfg50/Competitive-Programming/tree/
   template <class Info = int>
class ColorUpdate {
⊳ set<Range> ranges;
public:
⊳ struct Range {
\triangleright Range(int a = 0) : 1(a) {}
\triangleright Range(int a, int b, Info c) : l(a), r(b), v(c) {}

    ▷ int 1. r:

⊳ ⊳ Info v;
▷ bool operator<(const Range &b) const { return 1 < b.1; }</pre>
⊳ };
▶ vector<Range> upd(int 1, int r, Info v) {
▷ ▷ vector<Range> ans:
\triangleright if(1 >= r) return ans;
▷ auto it = ranges.lower_bound(1);
▷ if(it != ranges.begin()) {
⊳ ⊳ ⊳ it--;
\triangleright \triangleright \vdash \mathbf{if}(\mathbf{it} - \mathbf{r} > 1)  {
▷ ▷ ▷ ▷ auto cur = *it;
▷ ▷ ▷ ranges.erase(it);
▷ ▷ ▷ ranges.emplace(cur.1, 1, cur.v);
▷ ▷ ▷ ranges.emplace(1, cur.r, cur.v);
⊳ ⊳ }
▷ it = ranges.lower_bound(r);
▷ if(it != ranges.begin()) {

▷ ▷ ▷ it--;

\triangleright \triangleright \triangleright if(it->r > r) {
▷ ▷ ▷ ▷ auto cur = *it;
▷ ▷ ▷ ranges.erase(it);
▷ ▷ ▷ ranges.emplace(cur.l, r, cur.v);
▷ ▷ ▷ ranges.emplace(r, cur.r, cur.v);
⊳ ⊳ }
\triangleright for(it = ranges.lower_bound(1); it != ranges.end() && it->1 < r;
    → it++) {
▷ ▷ ▷ ans.push_back(*it);
▷ ranges.erase(ranges.lower_bound(1), ranges.lower_bound(r));
▷ ranges.emplace(1, r, v);
▷ ▷ return ans;
⊳ }
};
1.12 Heavy Light Decomposition
7033165b33a01511eac564bfb0b02a0b, 24 lines
void dfs_sz(int u){
    sz[u] = 1:
    for(auto &v : g[u]) if(v == p[u]){
       swap(v, g[u].back()); g[u].pop_back();
       break:
   for(auto &v : g[u]){
       p[v] = u; dfs_sz(v); sz[u] += sz[v];
       if(sz[v] > sz[ g[u][0] ])
           swap(v, g[u][0]);
```

```
// nxt[u] = start of path with u
// set nxt[root] = root beforehand
void dfs_hld(int u){
   in[u] = t++;
    rin[in[u]] = u;
    for(auto v : g[u]){
       nxt[v] = (v == g[u][0] ? nxt[u] : v); dfs_hld(v);
   out[u] = t;
// subtree of u => [ in[u], out[u] )
// path from nxt[u] to u => [ in[ nxt[u] ], in[u] ]
1.13 Iterative Segtree
a185b7c8af7cdbd659d0a033b782d668, 20 lines
T query(int 1, int r){ // [1, r]
   T rl, rr;
    for(1 += n, r += n+1; 1 < r; 1 >>= 1, r >>= 1){
       if(1 & 1) rl = merge(rl. st[l++]):
       if(r \& 1) rr = merge(st[--r], rr);
   return merge(rl. rr):
// initially save v[i] in st[n+i] for all i in [0, n)
void build(){
   for(int p = n-1; p > 0; p--)
       st[p] = merge(st[2*p], st[2*p+1]);
void update(int p, T val){
    st[p += n] = val;
    while(p >>= 1) st[p] = merge(st[2*p], st[2*p+1]);
1.14 LiChao's Segtree
93d48bf6dfc3fa031e9c5e1f7c8a8bcc, 25 lines
void add_line(line nw, int v = 1, int l = 0, int r = maxn) { // [1,
    \hookrightarrow r)
    int m = (1 + r) / 2;
   bool lef = nw.eval(1) < st[v].eval(1);</pre>
   bool mid = nw.eval(m) < st[v].eval(m);</pre>
   if(mid) swap(st[v], nw);
   if(r - 1 == 1) {
       return;
    } else if(lef != mid) {
       add_line(nw, 2 * v, 1, m);
       add line(nw. 2 * v + 1. m. r):
}
int get(int x, int v = 1, int l = 0, int r = maxn) {
   int m = (1 + r) / 2;
    if(r - 1 == 1) {
       return st[v].eval(x);
    } else if(x < m) {
       return min(st[v].eval(x), get(x, 2*v, 1, m));
       return min(st[v].eval(x), get(x, 2*v+1, m, r));
1.15 Palindromic tree
be0430a857f84e03450f060e4edf3289, 68 lines
#include <bits/stdc++.h>
using namespace std;
```

```
const int maxn = 3e5 + 1. sigma = 26:
int len[maxn]. link[maxn]. to[maxn][sigma]:
int slink[maxn], diff[maxn], series_ans[maxn];
int sz. last. n:
char s[maxn];
void init()
   s[n++] = -1:
   link[0] = 1;
   len[1] = -1;
   sz = 2:
int get_link(int v)
   while(s[n - len[v] - 2] != s[n - 1]) v = link[v];
   return v:
void add_letter(char c)
   s[n++] = c -= 'a':
   last = get_link(last);
   if(!to[last][c])
       len[sz] = len[last] + 2;
       link[sz] = to[get_link(link[last])][c];
       diff[sz] = len[sz] - len[link[sz]];
       if(diff[sz] == diff[link[sz]])
          slink[sz] = slink[link[sz]];
          slink[sz] = link[sz];
       to[last][c] = sz++;
   last = to[last][c];
int main()
   ios::sync_with_stdio(0);
   cin.tie(0);
   init();
   string s:
   cin >> s:
   int n = s.size();
   int ans[n + 1];
   memset(ans, 63, sizeof(ans));
   ans[0] = 0;
   for(int i = 1; i \le n; i++)
       add_letter(s[i - 1]);
       for(int v = last; len[v] > 0; v = slink[v])
          series_ans[v] = ans[i - (len[slink[v]] + diff[v])];
          if(diff[v] == diff[link[v]])
             series_ans[v] = min(series_ans[v], series_ans[link[v]
          ans[i] = min(ans[i], series_ans[v] + 1);
       cout << ans[i] << "\n";</pre>
   return 0;
```

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### 2 Math

### 2.1 Extended Euclidean Algorithm

875e47dd763ddfd8c117cb89ff15f1c4, 6 lines

```
// a*x + b*v = gcd(a, b), < gcd, x, v>
tuple<int. int. int> qcd(int a. int b) {
b if(b == 0) return make_tuple(a, 1, 0);

    auto [q, w, e] = gcd(b, a % b);

p return make_tuple(q, e, w - e * (a / b));
```

### 2.2 Chinese Remainder Theorem

ab7a1345fe4768c32b7f483eb99ea2da, 14 lines

```
// x = vet[i].first (mod vet[i].second)
11 crt(const vector<pair<11, 11>> &vet){
   11 \text{ ans} = 0, 1cm = 1;
   11 a, b, g, x, y;
   for(const auto &p : vet) {
       tie(a, b) = p;
       tie(g, x, y) = gcd(lcm, b);
      if((a - ans) % g != 0) return -1; // no solution
       ans = ans + x * ((a - ans) / g) % (b / g) * lcm;
      lcm = lcm * (b / g);
       ans = (ans \% lcm + lcm) \% lcm;
   return ans;
```

## 2.3 Diophantine Solver

cf9cb1477bd7b69c0143b7673f8cf5d9, 54 lines

```
template<tvpename T>
T extgcd(T a, T b, T &x, T &y) {
 if (a == 0) {
   x = 0:
   y = 1;
   return b:
 T p = b / a;
 T g = extgcd(b - p * a, a, y, x);
 x -= p * y;
 return q;
template<tvpename T>
bool diophantine(T a, T b, T c, T &x, T &y, T &g) {
 if (a == 0 && b == 0) {
   if (c == 0) {
    x = y = g = 0;
     return true;
   }
   return false;
 if (a == 0) {
   if (c % b == 0) {
    x = 0:
    y = c / b;
    q = abs(b);
     return true:
   return false;
 if (b == 0) {
   if (c % a == 0) {
    x = c / a;
    y = 0:
     g = abs(a);
    return true;
```

```
return false:
  g = extgcd(a, b, x, y);
  if (c % g != 0) {
     return false;
  T dx = c / a:
   c -= dx * a;
  T dv = c / b:
   c -= dy * b;
  x = dx + mulmod(x, c / g, b);
  y = dy + mulmod(y, c / g, a);
  q = abs(q);
   return true;
2.4 Preffix inverse
b6de09916fe942ec868bf55fff5c4df9, 4 lines
inv[1] = 1;
for(int i = 2; i < p; i++)
\triangleright inv[i] = (p - (p/i) * inv[p%i] % p) % p;
2.5 Pollard Rho
8a037aeb5aca655be64856efe64261b6, 47 lines
ll rho(ll n){
b if(n % 2 == 0) return 2;
do{
\triangleright \triangleright c = 11rand(1, n - 1);
\triangleright \triangleright x = 1 | rand(1, n - 1);
\triangleright \triangleright y = x:
⊳ ⊳ prod = 1:
\triangleright for(int i = 0; i < 40; i++) {
\triangleright \triangleright \triangleright x = add(mul(x, x, n), c, n);
\triangleright \triangleright \lor y = add(mul(y, y, n), c, n);
\triangleright \triangleright \lor y = add(mul(y, y, n), c, n);
\triangleright \triangleright \mathsf{prod} = \mathsf{mul}(\mathsf{prod}, \mathsf{abs}(\mathsf{x} - \mathsf{y}), \mathsf{n}) ?: \mathsf{prod};
⊳ ⊳ }
\triangleright \triangleright d = \_gcd(prod, n);
b } while(d == 1);
⊳ return d;
ll pollard rho(ll n){
\triangleright 11 x, c, y, d, k;
⊳ int i;
do{
\triangleright \triangleright i = 1;
\triangleright x = 1 \text{lrand}(1, n-1), c = 1 \text{lrand}(1, n-1);
\triangleright \triangleright y = x, k = 4;
\triangleright \triangleright if(++i == k) y = x, k *= 2;
\triangleright \triangleright \triangleright x = add(mul(x, x, n), c, n);
\triangleright \triangleright d = \underline{gcd(abs(x - y), n)};
> }while(d == n);
⊳ return d;
void factorize(ll val, map<ll, int> &fac){

    if(rabin(val)) fac[ val ]++;
⊳ else{
▷ ▷ 11 d = pollard_rho(val);
⊳ b factorize(d, fac);
▷ b factorize(val / d, fac);
⊳ }
map<ll, int> factor(ll val){
```

map<11, int> fac;

```
    if(val > 1) factorize(val. fac):
⊳ return fac:
2.6 Miller Rabin
dbdaf062461f61fe2c4c5db55db1add3, 20 lines
bool rabin(ll n){
p if(n <= 1) return 0:</pre>

    if(n <= 3) return 1;</pre>
> 11 s = 0, d = n - 1;
\triangleright while(d % 2 == 0) d /= 2, s++;
\triangleright for(int k = 0; k < 64; k++){
\triangleright 11 a = 11rand(2, n-2);
\triangleright 11 x = fexp(a. d. n):
\triangleright \mathbf{if}(x != 1 \&\& x != n-1)
\triangleright \triangleright \triangleright \text{ for(int } r = 1; r < s; r++){}
\triangleright \triangleright \triangleright \triangleright x = mul(x, x, n);
\triangleright \triangleright \triangleright if(x == 1) return 0;
\triangleright \triangleright \triangleright if(x == n-1) break;
▷ ▷ ▷ }
\triangleright \triangleright \vdash \mathbf{if}(x != n-1) \mathbf{return 0};
⊳ ⊳ }
⊳ }
⊳ return 1;
2.7 Primitive root
8b6fd62a31a9ed12d4abbf643f3352b7, 30 lines
// a primitive root modulo n is any number g such that any c
    \hookrightarrowcoprime to n is congruent to a power of g modulo n.
bool exists root(ll n){
    if(n == 1 || n == 2 || n == 4) return true;
    if(n \% 2 == 0) n /= 2;
    if(n % 2 == 0) return false:
    // test if n is a power of only one prime
    for(11 i = 3; i * i <= n; i += 2) if(n % i == 0){
        while(n \% i == 0) n /= i;
        return n == 1:
    return true;
11 primitive root(11 n){
    if(n == 1 || n == 2 || n == 4) return n - 1;
    if(not exists_root(n)) return -1;
    11 x = phi(n);
    auto pr = factorize(x);
    auto check = [x, n, pr](11 m){
        for(11 p : pr) if(fexp(m, x / p, n) == 1)
            return false:
        return true:
    for(11 m = 2; ; m++) if(\_gcd(m, n) == 1)
        if(check(m)) return m;
// Let's denote R(n) as the set of primitive roots modulo n, p is
// g \in R(p) \Rightarrow (pow(g, p-1, p * p) == 1 ? g+p : g) \in R(pow(p, k))
    \hookrightarrow)), for all k > 1
// g \text{ in } R(pow(p, k)) \Rightarrow (g \% 2 == 1 ? g : g + pow(p, k)) \setminus in R(2*)
2.8 Mobius Function
```

02d66957b8c5648c0047a18fd1d912b3, 6 lines

for(int j = i + i; j < N; j += i)

memset(mu, 0, sizeof mu);

for(int i = 1: i < N: i++)

mu[j] -= mu[i];

mu[1] = 1;

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```
// g(n) = sum\{f(d)\} \Rightarrow f(n) = sum\{mu(d)*g(n/d)\}
2.9 Mulmod TOP
4219357bb28adeea0e8b9209479b86ba, 9 lines
```

```
constexpr uint64 t mod = (1ull<<61) - 1:</pre>
uint64_t modmul(uint64_t a, uint64_t b){
b uint64_t 11 = (uint32_t)a, h1 = a>>32, 12 = (uint32_t)b, h2 = b
   >>>32:
\triangleright uint64_t 1 = 11*12, m = 11*h2 + 12*h1, h = h1*h2;
\Rightarrow uint64 t ret = (1&mod) + (1>>61) + (h << 3) + (m >> 29) + (m <<
   \hookrightarrow35 >> 3) + 1:
p ret = (ret & mod) + (ret>>61);
p ret = (ret & mod) + (ret>>61):
▶ return ret-1;
2.10 Modular multiplication TOPPER
```

# 6cd65395928fcde0e3ddc7a7f8e82d1c, 6 lines

```
11 mulmod(11 a. 11 b. 11 mod) {
   11 q = 11((long double)a * (long double)b / (long double)mod);
   11 r = (a * b - mod * q) \% mod;
   if(r < 0) r += mod:
   return r;
```

### 2.11 Division Trick

ccbcf95b38d1ddc7b0411c86a88b6aa9, 9 lines

```
for(int l = 1, r; l \le n; l = r + 1) {
   r = n / (n / 1);
   // n / x yields the same value for 1 <= x <= r
for(int 1, r = n; r > 0; r = 1 - 1) {
   int tmp = (n + r - 1) / r;
   1 = (n + tmp - 1) / tmp;
   // (n+x-1) / x yields the same value for 1 <= x <= r
```

### 2.12 Matrix Determinant

60d2f1720b2577abfd897a9e194b060b, 32 lines

```
long double a[n][n];
long double gauss(){
   long double det = 1;
   for(int i = 0; i < n; i++){
       int a = i:
       for(int j = i+1; j < n; j++){
          if(abs(a[j][i]) > abs(a[q][i]))
             q = j;
       if(abs(a[q][i]) < EPS){
          det = 0;
          break:
       if(i != q){
          for(int w = 0: w < n: w++)
             swap(a[i][w], a[q][w]);
          det = -det;
       det *= a[i][i];
       for(int j = i+1; j < n; j++) a[i][j] /= a[i][i];</pre>
       for(int j = 0; j < n; j++) if(j != i){
          if(abs(a[j][i]) > EPS)
             for(int k = i+1; k < n; k++)
                 a[j][k] = a[i][k] * a[j][i];
   }
```

```
return det:
```

### 2.13 Simplex Method

6c7be7a22b6ff4cee019c29e6ea12194, 79 lines

```
typedef long double dbl;
 const dbl eps = 1e-6;
const int N = , M = ;
mt19937 rng(chrono::steady_clock::now().time_since_epoch().count())
struct simplex {

  int X[N], Y[M];

    dbl A[M][N], b[M], c[N];

    dbl ans:

⊳ int n, m;

    dbl sol[N];

    void pivot(int x, int y){
\triangleright \triangleright swap(X[y], Y[x]);
\triangleright \triangleright b[x] /= A[x][y];
\triangleright for(int i = 0; i < n; i++)
▷ ▷ ▷ if(i != y)
\triangleright \triangleright \triangleright \triangleright A[x][i] /= A[x][y];
\triangleright \land A[x][y] = 1. / A[x][y];
\triangleright for(int i = 0; i < m; i++)
\triangleright \triangleright if(i != x \&\& abs(A[i][y]) > eps) {
\triangleright \triangleright \triangleright \triangleright b[i] -= A[i][y] * b[x];
\triangleright \triangleright \triangleright \triangleright for(int j = 0; j < n; j++) if(j != y)
                         A[i][j] -= A[i][y] * A[x][j];
\triangleright \triangleright \triangleright \triangleright A[i][y] = -A[i][y] * A[x][y];
⊳ ⊳ ⊳ }
\triangleright ans += c[y] * b[x];
\triangleright for(int i = 0; i < n; i++)
▷ ▷ ▷ if(i != y)
\triangleright \triangleright \triangleright \triangleright c[i] -= c[y] * A[x][i];
\triangleright c[y] = -c[y] * A[x][y];
⊳ }
▷ // maximiza sum(x[i] * c[i])
⊳ // sujeito a
▷ // sum(a[i][j] * x[j]) <= b[i] para 0 <= i < m (Ax <= b)</pre>
\triangleright // x[i] >= 0 para 0 <= i < n (x >= 0)
▷ // (n variaveis, m restricoes)
⊳ // guarda a resposta em ans e retorna o valor otimo
b dbl solve(int _n, int _m) {
\triangleright this->n = _n; this->m = _m;
          for(int i = 1; i < m; i++){
               int id = uniform_int_distribution<int>(0, i)(rng);
               swap(b[i], b[id]);
               for(int j = 0; j < n; j++)
                    swap(A[i][j], A[id][j]);
\triangleright \triangleright ans = 0.;
\triangleright for(int i = 0; i < n; i++) X[i] = i;
\triangleright for(int i = 0; i < m; i++) Y[i] = i + n;
▷ ▷ while(true) {
\triangleright \triangleright \triangleright  int x = min_element(b, b + m) - b;
\triangleright \triangleright \triangleright \mathbf{if}(b[x] >= -eps)
⊳ ⊳ ⊳ break:
\triangleright \triangleright int y = find_if(A[x], A[x] + n, [](dbl d) { return d < -eps;
     \hookrightarrow}) - A[x];
\triangleright \triangleright \triangleright if(y == n) throw 1; // no solution
▷ ▷ ▷ pivot(x, y);
⊳ ⊳ }
▷ ▷ while(true) {
```

```
\triangleright \triangleright int v = max element(c. c + n) - c:
▷ ▷ ▷ if(c[v] <= eps) break:
\triangleright \triangleright int x = -1;
\triangleright \triangleright dbl mn = 1. / 0.;
\triangleright \triangleright for(int i = 0; i < m; i++)
\triangleright \triangleright \triangleright if(A[i][y] > eps && b[i] / A[i][y] < mn)
\triangleright \triangleright \triangleright \triangleright \triangleright mn = b[i] / A[i][y], x = i;
\triangleright \triangleright if(x == -1) throw 2; // unbounded
▷ ▷ ▷ pivot(x, y);
⊳ ⊳ }
\triangleright \triangleright \text{ for(int } i = 0; i < m; i++)
\triangleright \triangleright \triangleright \mathbf{if}(Y[i] < n)
\triangleright \triangleright \triangleright sol[Y[i]] = b[i];
⊳ ⊳ return ans:
⊳ }
```

### 2.14 FFT

8f879ebf120408d0c2e84adea331b108, 25 lines

```
void fft(vector<base> &a, bool inv){
   int n = (int)a.size();
   for(int i = 1, j = 0; i < n; i++){
      int bit = n \gg 1;
      for(; j >= bit; bit >>= 1) j -= bit;
       j += bit;
       if(i < j) swap(a[i], a[j]);
   for(int sz = 2; sz <= n; sz <<= 1) {</pre>
       double ang = 2 * PI / sz * (inv ? -1 : 1);
       base wlen(cos(ang), sin(ang));
      for(int i = 0; i < n; i += sz){
          base w(1, 0);
          for(int j = 0; j < sz / 2; j++){
             base u = a[i+j], v = a[i+j + sz/2] * w;
             a[i+j] = u + v;
             a[i+j+sz/2] = u - v;
             w *= wlen;
      }
   if(inv) for(int i = 0: i < n: i++) a[i] /= 1.0 * n:
```

### 2.15 FFT Tourist

0a189fdba3448c9bb1d3acf98af394c0, 163 lines

```
namespace fft {
 typedef double dbl;
 struct num {
   dbl x, y;
   num() \{ x = y = 0; \}
   num(dbl x, dbl y) : x(x), y(y) {}
 inline num operator+(num a, num b) { return num(a.x + b.x, a.y +
 inline num operator-(num a, num b) { return num(a.x - b.x, a.y -
 inline num operator*(num a, num b) { return num(a.x * b.x - a.y *
     \hookrightarrow b.y, a.x * b.y + a.y * b.x); }
 inline num conj(num a) { return num(a.x, -a.y); }
 int base = 1:
 vector<num> roots = {{0, 0}, {1, 0}};
 vector < int > rev = \{0, 1\};
```

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```
const dbl PI = acosl(-1.0):
void ensure_base(int nbase) {
 if(nbase <= base) return;</pre>
 rev.resize(1 << nbase);
 for(int i = 0; i < (1 << nbase); i++) {
   rev[i] = (rev[i >> 1] >> 1) + ((i \& 1) << (nbase - 1));
 roots.resize(1 << nbase);</pre>
  while(base < nbase) {</pre>
   dbl \ angle = 2*PI / (1 << (base + 1));
   for(int i = 1 << (base - 1); i < (1 << base); i++) {</pre>
     roots[i << 1] = roots[i];</pre>
     dbl \ angle_i = angle * (2 * i + 1 - (1 << base));
     roots[(i << 1) + 1] = num(cos(angle_i), sin(angle_i));</pre>
   base++;
 }
void fft(vector<num> &a, int n = -1) {
 if(n == -1) {
   n = a.size();
 assert((n & (n-1)) == 0);
 int zeros = __builtin_ctz(n);
  ensure_base(zeros);
  int shift = base - zeros;
  for(int i = 0; i < n; i++) {
   if(i < (rev[i] >> shift)) {
     swap(a[i], a[rev[i] >> shift]);
 for(int k = 1; k < n; k <<= 1) {
   for(int i = 0; i < n; i += 2 * k) {
     for(int j = 0; j < k; j++) {
       num z = a[i+j+k] * roots[j+k];
       a[i+j+k] = a[i+j] - z;
       a[i+j] = a[i+j] + z;
     }
   }
vector<num> fa, fb;
vector<int> multiply(vector<int> &a, vector<int> &b) {
 int need = a.size() + b.size() - 1;
 int nbase = 0;
  while((1 << nbase) < need) nbase++;</pre>
  ensure_base(nbase);
  int sz = 1 << nbase;</pre>
  if(sz > (int) fa.size()) {
   fa.resize(sz);
 for(int i = 0; i < sz; i++) {
   int x = (i < (int) a.size() ? a[i] : 0);</pre>
   int y = (i < (int) b.size() ? b[i] : 0);</pre>
   fa[i] = num(x, y);
 fft(fa, sz);
 num r(0, -0.25 / sz);
  for(int i = 0; i \le (sz >> 1); i++) {
   int j = (sz - i) & (sz - 1);
   num z = (fa[j] * fa[j] - conj(fa[i] * fa[i])) * r;
   if(i != j) {
```

```
fa[j] = (fa[i] * fa[i] - conj(fa[j] * fa[j])) * r;
   fa[i] = z;
 fft(fa, sz);
 vector<int> res(need);
 for(int i = 0; i < need; i++) {
  res[i] = fa[i].x + 0.5;
 return res;
vector<int> multiply_mod(vector<int> &a, vector<int> &b, int m,
     \rightarrow int eq = 0) {
  int need = a.size() + b.size() - 1;
 int nbase = 0;
  while ((1 << nbase) < need) nbase++;</pre>
  ensure_base(nbase);
 int sz = 1 << nbase;</pre>
 if (sz > (int) fa.size()) {
   fa.resize(sz);
 for (int i = 0; i < (int) a.size(); i++) {</pre>
   int x = (a[i] \% m + m) \% m;
   fa[i] = num(x & ((1 << 15) - 1), x >> 15);
 fill(fa.begin() + a.size(), fa.begin() + sz, num {0, 0});
 fft(fa, sz);
 if (sz > (int) fb.size()) {
   fb.resize(sz);
 if (eq) {
   copy(fa.begin(), fa.begin() + sz, fb.begin());
   for (int i = 0; i < (int) b.size(); i++) {</pre>
     int x = (b[i] \% m + m) \% m;
     fb[i] = num(x & ((1 << 15) - 1), x >> 15);
   fill(fb.begin() + b.size(), fb.begin() + sz, num {0, 0});
   fft(fb, sz);
 dbl ratio = 0.25 / sz:
 num r2(0, -1):
 num r3(ratio. 0):
 num r4(0. -ratio):
 num r5(0, 1);
  for (int i = 0; i \le (sz >> 1); i++) {
   int j = (sz - i) & (sz - 1);
   num a1 = (fa[i] + conj(fa[j]));
   num a2 = (fa[i] - conj(fa[j])) * r2;
   num b1 = (fb[i] + conj(fb[j])) * r3;
   num b2 = (fb[i] - conj(fb[j])) * r4;
   if (i != j) {
     num c1 = (fa[j] + conj(fa[i]));
     num c2 = (fa[j] - conj(fa[i])) * r2;
     num d1 = (fb[j] + conj(fb[i])) * r3;
     num d2 = (fb[j] - conj(fb[i])) * r4;
     fa[i] = c1 * d1 + c2 * d2 * r5;
     fb[i] = c1 * d2 + c2 * d1;
   fa[j] = a1 * b1 + a2 * b2 * r5;
   fb[j] = a1 * b2 + a2 * b1;
 fft(fa, sz);
 fft(fb, sz);
 vector<int> res(need):
  for (int i = 0; i < need; i++) {
   long long aa = fa[i].x + 0.5;
```

```
long long bb = fb[i].x + 0.5;
      long long cc = fa[i].y + 0.5;
      res[i] = (aa + ((bb \% m) << 15) + ((cc \% m) << 30)) \% m;
    return res;
  vector<int> square_mod(vector<int> &a, int m) {
   return multiply_mod(a, a, m, 1);
2.16 NTT
6b8373484af527b5fc6f28a4b95d83d2, 37 lines
const int mod = 7340033:
const int root = 5;
const int root_1 = 4404020;
const int root_pw = 1<<20;</pre>
void fft (vector<int> & a, bool invert) {
p int n = (int) a.size();
\triangleright for (int i=1, j=0; i<n; ++i) {
\triangleright \triangleright int bit = n >> 1;
▷ for (; j>=bit; bit>>=1)
▷ ▷ ▷ j -= bit;
\triangleright \triangleright if (i < j)
\triangleright \triangleright \triangleright \text{ swap (a[i], a[j])};
⊳ }

    for (int len=2; len<=n; len<<=1) {</pre>
▷ int wlen = invert ? root_1 : root;
▷ ▷ ▷ wlen = int (wlen * 111 * wlen % mod);
▷ for (int i=0; i<n; i+=len) {</pre>
▷ ▷ ▷ int w = 1:
▷ ▷ for (int j=0; j<len/2; ++j) {</pre>
\triangleright \triangleright \triangleright  int u = a[i+j], v = int (a[i+j+len/2] * 111 * w % mod);
\triangleright \triangleright \triangleright a[i+j] = u+v < mod ? u+v : u+v-mod;
\triangleright \triangleright \triangleright a[i+j+len/2] = u-v >= 0 ? u-v : u-v+mod;
\triangleright \triangleright \triangleright \triangleright w = int (w * 111 * wlen % mod);
▷ ▷ ▷ }
⊳ ⊳ }
⊳ }

    if (invert) {
▷ int nrev = reverse (n, mod);
\triangleright for (int i=0; i<n; ++i)

▷ ▷ a[i] = int (a[i] * 1ll * nrev % mod);
⊳ }
2.17 Gauss
68a74ce2c6da5a785b5400546ab6d02e, 31 lines
// Solves systems of linear equations.
// To use, build a matrix of coefficients and call run(mat, R, C).
   → If the i-th variable is free, row[i] will be -1, otherwise it'
   \hookrightarrows value will be ans[i].
namespace Gauss {
  const int MAXC = 1001;
  int row[MAXC];
 double ans[MAXC];
  void run(double mat[][MAXC], int R, int C) {
   REP(i, C) row[i] = -1;
    int r = 0;
```

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```
REP(c. C) {
 int k = r:
 FOR(i, r, R) if(fabs(mat[i][c]) > fabs(mat[k][c])) k = i;
 if(fabs(mat[k][c]) < eps) continue;</pre>
 REP(j, C+1) swap(mat[r][j], mat[k][j]);
 REP(i, R) if (i != r) {
   double w = mat[i][c] / mat[r][c];
   REP(j, C+1) mat[i][j] -= mat[r][j] * w;
 row[c] = r++;
REP(i, C) {
 int r = row[i]:
 ans[i] = r == -1 ? 0 : mat[r][C] / mat[r][i];
```

### 2.18 Gauss Xor

95b10e530742fc319fba2c7759eb7096, 25 lines

```
const 11 MAX = 1e9:
const int LOG_MAX = 64 - __builtin_clzll((11)MAX);
struct Gauss {
   array<11, LOG_MAX> vet;
   int size;
   Gauss() : size(0) {
p fill(vet.begin(), vet.end(), 0);
   Gauss(vector<ll> vals) : size(0) {
▷ fill(vet.begin(), vet.end(), 0);
      for(ll val : vals) add(val);
   bool add(ll val) {
      for(int i = 0; i < LOG_MAX; i++) if(val & (1LL << i)) {
          if(vet[i] == 0) {
             vet[i] = val;
             size++;
             return true;
          val ^= vet[i];
      return false;
   }
```

## 2.19 Simpson

fc02c5a6437303ecda9da16cbb65150c, 16 lines

```
inline double simpson(double f1,double fr,double fmid,double 1,
    →double r) {

  return (fl + fr + 4.0 * fmid) * (r - 1) / 6.0;
double rsimpson(double slr.double fl.double fr.double fmid.double l
     →,double r) {
\triangleright double mid = (1+r)*0.5;

    double slm = simpson(fl, fmid, fml, l, mid);

    double smr = simpson(fmid, fr, fmr, mid, r);

    if(fabs(slr-slm-smr) < eps and r - l < delta) return slr;
</pre>

    return rsimpson(slm,fl,fmid,fml,1,mid) + rsimpson(smr,fmid,fr,fmr)

   \hookrightarrow.mid.r):
double integrate(double 1,double r) {
\triangleright double mid = (1+r)*0.5:
\triangleright double fl = f(1), fr = f(r), fmid = f(mid);

    return rsimpson(simpson(fl,fr,fmid,l,r),fl,fr,fmid,l,r);
```

### 2.20 Matrix

9b4a07b8f779901adbfec8348e388d1c, 28 lines

```
template <const size_t n, const size_t m, class T = modBase<>>
struct Matrix {
 T v[n][m];
 Matrix(int d = 0) {
   for (int i = 0; i < n; i++) {
     for (int j = 0; j < m; j++) {
      v[i][j] = T(0);
     if (i < m) {
      v[i][i] = T(d);
   }
 template <size_t mm>
 Matrix<n. mm. T> operator*(Matrix<m. mm. T> &o) {
   Matrix<n. mm. T> ans:
   for (int i = 0; i < n; i++) {
     for (int i = 0: i < mm: i++) {
      for (int k = 0; k < m; k++) {
        ans.v[i][j] = ans.v[i][j] + v[i][k] * o.v[k][j];
   }
   return ans;
};
    Graphs
```

### 3.1 Bipartite Matching

0fa7a518a130e597ac64ec4e1f1505fb, 28 lines

003c2ff033ac00d7616028972bdff695, 65 lines

```
int match[N];
int vis[N], pass;
vector<int> q[N];
bool dfs(int u) {
vis[u] = pass;
p for(int v : g[u]) if(vis[v] != pass) {
▷ ▷ vis[v] = pass;
\triangleright \triangleright \triangleright match[v] = u;
\triangleright \triangleright \triangleright match[u] = v;
⊳ ⊳ ⊳ return true;
⊳ ⊳ }
⊳ }
⊳ return false:
int max_maching() {

    memset(match, -1, sizeof match);

    int max_matching_size = 0;

    for(int u : vertices_on_side_A) {
▷ if(dfs(i)) max_matching_size++;
⊳ }
p return max_matching_size;
3.2 Dinic
```

```
const int N = 100005:
const int E = 2000006:
vector<int> q[N];
int ne:
struct Edge{
   int from, to; ll flow, cap;
int lvl[N], vis[N], pass, start = N-2, target = N-1;
int qu[N], qt, px[N];
ll run(int s, int sink, ll minE){
   if(s == sink) return minE:
   11 \text{ ans} = 0;
   for(; px[s] < (int)g[s].size(); px[s]++){</pre>
       int e = g[s][ px[s] ];
       auto &v = edge[e], &rev = edge[e^1];
       if(lvl[v.to] != lvl[s]+1 || v.flow >= v.cap)
          continue; // v.cap - v.flow < lim</pre>
       11 tmp = run(v.to, sink,min(minE, v.cap-v.flow));
       v.flow += tmp, rev.flow -= tmp;
       ans += tmp, minE -= tmp;
       if(minE == 0) break;
   return ans:
bool bfs(int source, int sink){
   qt = 0;
   qu[qt++] = source;
   lvl[source] = 1:
   vis[source] = ++pass;
   for(int i = 0; i < qt; i++){
       int u = qu[i];
       px[u] = 0;
▷ ▷ if(u == sink) return true;
       for(auto& ed : g[u]) {
           auto v = edge[ed];
           if(v.flow >= v.cap || vis[v.to] == pass)
              continue; // v.cap - v.flow < lim</pre>
           vis[v.to] = pass;
          lvl[v.to] = lvl[u]+1;
           qu[qt++] = v.to;
      }
   }
   return false;
11 flow(int source = start, int sink = target){
   11 \text{ ans} = 0;
   //for(lim = (1LL << 62); lim >= 1; lim /= 2)
   while(bfs(source, sink))
▷ ▷ ans += run(source, sink, oo);
   return ans:
void addEdge(int u, int v, ll c = 1, ll rc = 0){
   edge[ne] = \{u, v, 0, c\};
   g[u].push_back(ne++);
   edge[ne] = {v, u, 0, rc};
   g[v].push_back(ne++);
void reset_flow(){
\triangleright for(int i = 0; i < ne; i++)
▷ ▷ edge[i].flow = 0;
3.3 Push relabel
7f5c584b894a5295a1c15aebc9928f06, 58 lines
```

// Push relabel in O(V^2 E^0.5) with gap heuristic

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```
// It's quite fast
template<typename flow_t = long long>
struct PushRelabel {
   struct Edge { int to, rev; flow_t f, c; };
   vector<vector<Edge> > g;
   vector<flow_t> ec;
   vector<Edge*> cur;
   vector<vector<int> > hs;
   vector<int> H:
   PushRelabel(int n) : g(n), ec(n), cur(n), hs(2*n), H(n) {}
   void add_edge(int s, int t, flow_t cap, flow_t rcap=0) {
       if (s == t) return:
       Edge a = {t, (int)g[t].size(), 0, cap};
       Edge b = \{s, (int)g[s].size(), 0, rcap\};
       g[s].push_back(a);
       g[t].push_back(b);
   void add_flow(Edge& e, flow_t f) {
       Edge &back = g[e.to][e.rev];
       if (!ec[e.to] && f)
          hs[H[e.to]].push_back(e.to);
       e.f += f, ec[e.to] += f;
       back.f -= f, ec[back.to] -= f;
   flow_t max_flow(int s, int t) {
       int v = g.size();
      H[s] = v; ec[t] = 1;
      vector<int> co(2 * v);
       co[0] = v-1;
       for(int i = 0; i < v; ++i) cur[i] = g[i].data();</pre>
       for(auto &e : g[s]) add_flow(e, e.c);
       if(hs[0].size())
       for (int hi = 0; hi >= 0;) {
          int u = hs[hi].back();
          hs[hi].pop_back();
          while (ec[u] > 0) // discharge u
             if (cur[u] == g[u].data() + g[u].size()) {
                 H[u] = 1e9;
                 for(auto &e:g[u])
                     if (e.c - e.f && H[u] > H[e.to]+1)
                        H[u] = H[e.to]+1, cur[u] = &e;
                 if (++co[H[u]], !--co[hi] && hi < v)</pre>
                     for(int i = 0; i < v; ++i)
                        if (hi < H[i] && H[i] < v){</pre>
                            --co[H[i]];
                            H[i] = v + 1;
                 hi = H[u];
              } else if (cur[u]->c - cur[u]->f && H[u] == H[cur[u
                 \hookrightarrow]->to]+1)
                 add_flow(*cur[u], min(ec[u], cur[u]->c - cur[u]->f
                     →)):
              else ++cur[u];
          while (hi >= 0 \& hs[hi].empty()) --hi;
       return -ec[s];
3.4 Min Cost Max Flow
546b8e1a0c2cc055a8e6747620146f31, 59 lines
const 11 oo = 1e18:
const int N = 422, E = 2 * 10006;
vector<int> g[N];
```

int ne;

struct Edge{

```
int from, to; 11 cap, cost;
} edae[E]:
int start = N-1, target = N-2, p[N]; int inqueue[N];
11 d[N];
11 pot[N];
bool dijkstra(int source, int sink) {
\triangleright for(int i = 0; i < N; i++) d[i] = oo;

    d[source] = 0;
priority_queue<pair<ll, int>> q;

    q.emplace(0, source);
⊳ 11 dt; int u;
▶ while(!q.empty()) {
▷ ▷ if(dt > d[u]) continue;
▷ ▷ if(u == sink) return true;
▷ b for(int e : g[u]) {
▷ ▷ ▷ auto v = edge[e];
\triangleright \triangleright \triangleright const 11 cand = d[u] + v.cost + pot[u] - pot[v.to];
\triangleright \triangleright \mathbf{if}(v.cap > 0 \text{ and } cand < d[v.to]) 
▷ ▷ ▷ ▷ p[v.to] = e;
\triangleright \triangleright \triangleright \triangleright d[v.to] = cand;
▷ ▷ ▷ □ q.emplace(-d[v.to], v.to);
⊳ ⊳ ⊳ }
⊳ ⊳ }
⊳ }
▶ return d[sink] < oo;</pre>
// <max flow, min cost>
pair<11, 11> mincost(int source = start, int sink = target){
    11 ans = 0, mf = 0;
    while(dijkstra(source, sink)){
        11 f = oo;
        for(int u = sink; u != source; u = edge[ p[u] ].from)
            f = min(f, edge[ p[u] ].cap);
        ans += f * (d[sink] - pot[source] + pot[sink]);
        for(int u = sink; u != source; u = edge[ p[u] ].from){
            edge[ p[u] ].cap -= f;
            edge[ p[u] ^ 1 ].cap += f;
\triangleright \mathsf{for}(\mathsf{int}\ i=0;\ i< N;\ i++)\ \mathsf{pot}[i]=\mathsf{min}(\mathsf{oo},\ \mathsf{pot}[i]+\mathsf{d}[i]);
   return {mf, ans};
void addEdge(int u, int v, ll c, ll cost){

    assert(cost >= 0); //IF not, pot[i]=short.path source

    edge[ne] = {u, v, c, cost};
    g[u].push_back(ne++);
    edge[ne] = \{v, u, 0, -cost\};
    g[v].push_back(ne++);
3.5 Blossom Algorithm for General Matching
```

eeb18cdbaa4a43f8574d250ba3d8fbe1, 71 lines

```
const int MAXN = 2020 + 1;
// 1-based Vertex index
int vis[MAXN], par[MAXN], orig[MAXN], match[MAXN], aux[MAXN], t, N;
vector<int> conn[MAXN];
queue<int> 0;
void addEdge(int u, int v) {
        conn[u].push_back(v); conn[v].push_back(u);
    }
void init(int n) {
        N = n; t = 0;
        b for(int i=0; i<=n; ++i)
        b > conn[i].clear(), match[i] = aux[i] = par[i] = 0;
```

```
void augment(int u, int v) {
\triangleright int pv = v. nv:
▷ ▷ pv = par[v]; nv = match[pv];
▷ b match[v] = pv; match[pv] = v;
\triangleright \triangleright v = nv;
int lca(int v, int w) {
→ ++t;
▶ while(true) {
▷ ▷ if(v) {
▷ ▷ if(aux[v] == t) return v; aux[v] = t;
▷ ▷ ▷ v = orig[par[match[v]]];
⊳ ⊳ }
▷ ▷ swap(v, w);
⊳ }
void blossom(int v, int w, int a) {
b while(orig[v] != a) {
▷ ▷ par[v] = w; w = match[v];
\triangleright if(vis[w] == 1) Q.push(w), vis[w] = 0;
▷ orig[v] = orig[w] = a; v = par[w];
⊳ }
bool bfs(int u) {

    fill(vis+1, vis+1+N, -1); iota(orig + 1, orig + N + 1, 1);

\triangleright Q = queue<int>(); Q.push(u); vis[u] = 0;
▶ while(!Q.empty()) {
▷ int v = Q.front(); Q.pop();
▷ b for(int x: conn[v]) {
\triangleright \triangleright \triangleright \triangleright par[x] = v; vis[x] = 1;
▷ ▷ ▷ if(!match[x]) return augment(u, x), true;
▷ ▷ ▷ Q.push(match[x]); vis[match[x]] = 0;
\triangleright \triangleright \triangleright \mathsf{else} \; \mathsf{if}(\mathsf{vis}[\mathsf{x}] == 0 \; \&\& \; \mathsf{orig}[\mathsf{v}] \; != \mathsf{orig}[\mathsf{x}]) \; \{
▷ ▷ ▷ int a = lca(orig[v], orig[x]);
\triangleright \triangleright \triangleright \triangleright blossom(x, v, a); blossom(v, x, a);
▷ ▷ ▷ }
⊳ ⊳ }
⊳ }

    return false;

int Match() {
\triangleright int ans = 0;

  vector<int> V(N-1); iota(V.begin(), V.end(), 1);
> shuffle(V.begin(), V.end(), mt19937(0x94949));
For(auto x: V) if(!match[x]){
b for(auto y: conn[x]) if(!match[y]) {
\triangleright \triangleright \triangleright match[x] = y, match[y] = x;
▷ ▷ ▷ ++ans; break;
⊳ ⊳ }
⊳ }

    for(int i=1; i<=N; ++i) if(!match[i] && bfs(i)) ++ans;
</pre>
▶ return ans;
```

# 3.6 Blossom Algorithm for Weighted General Matching

cdfc6f3b970dfec881deb021a9ce429b, 205 lines

```
// N^3 (but fast in practice)
static const int INF = INT_MAX;
static const int N = 514;
struct edge{
> int u,v,w; edge(){}
```

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```
    edge(int ui.int vi.int wi)

▷ ▷ :u(ui),v(vi),w(wi){}
};
int n,n_x;
edge g[N*2][N*2];
int lab[N*2];
int match[N*2],slack[N*2],st[N*2],pa[N*2];
int flo_from[N*2][N+1],S[N*2],vis[N*2];
vector<int> flo[N*2];
queue<int> q;
int e_delta(const edge &e){

  return lab[e.u]+lab[e.v]-g[e.u][e.v].w*2;
void update_slack(int u,int x){
b if(!slack[x]||e_delta(g[u][x])<e_delta(g[slack[x]][x]))slack[x]=u</pre>
void set_slack(int x){
⊳ slack[x]=0;

    for(int u=1;u<=n;++u)</pre>
▷ ▷ if(g[u][x].w>0&&st[u]!=x&&S[st[u]]==0)
▷ ▷ □ update_slack(u,x);
void q_push(int x){

    if(x<=n)q.push(x);</pre>
b else for(size_t i=0;i<flo[x].size();i++)</pre>
▷ ▷ q_push(flo[x][i]);
void set_st(int x,int b){
⊳ st[x]=b;
b if(x>n)for(size_t i=0;i<flo[x].size();++i)</pre>
▷ ▷ set_st(flo[x][i],b);
int get_pr(int b,int xr){
p int pr=find(flo[b].begin(),flo[b].end(),xr)-flo[b].begin();
b if(pr%2==1){
▷ reverse(flo[b].begin()+1,flo[b].end());
▷ return (int)flo[b].size()-pr;
▶ }else return pr;
void set_match(int u,int v){
p match[u]=g[u][v].v;

    if(u<=n) return;</pre>
⊳ edge e=g[u][v];
p int xr=flo_from[u][e.u],pr=get_pr(u,xr);
b for(int i=0;i<pr;++i)set_match(flo[u][i],flo[u][i^1]);</pre>
⊳ set_match(xr,v);
p rotate(flo[u].begin(),flo[u].begin()+pr,flo[u].end());
void augment(int u,int v){
▶ for(;;){
▷ int xnv=st[match[u]];
▷ ▷ set_match(u,v);
▷ ▷ if(!xnv)return;
▷ ▷ set_match(xnv,st[pa[xnv]]);
▷ ▷ u=st[pa[xnv]],v=xnv;
⊳ }
int get_lca(int u,int v){
⊳ static int t=0;
\triangleright for(++t;u||v;swap(u,v)){
▷ ▷ if(u==0)continue;
▷ if(vis[u]==t)return u;
▷ ▷ vis[u]=t;
▷ ▷ u=st[match[u]];
▷ ▷ if(u)u=st[pa[u]];
⊳ }

  return 0;
```

```
void add_blossom(int u,int lca,int v){

   int b=n+1;

    while(b<=n_x&&st[b])++b;</pre>

    if(b>n_x)++n_x;

  lab[b]=0,S[b]=0;
p match[b]=match[lca];
b flo[b].clear();
p flo[b].push_back(lca);
p for(int x=u,y;x!=lca;x=st[pa[y]])
b flo[b].push_back(x),flo[b].push_back(y=st[match[x]]),q_push(y);
p reverse(flo[b].begin()+1,flo[b].end());
p for(int x=v,y;x!=lca;x=st[pa[y]])
b flo[b].push_back(x),flo[b].push_back(y=st[match[x]]),q_push(y);
⊳ set st(b.b):
p for(int x=1;x<=n_x;++x)g[b][x].w=g[x][b].w=0;</pre>
p for(int x=1;x<=n;++x)flo_from[b][x]=0;</pre>

    for(size_t i=0;i<flo[b].size();++i){</pre>
▷ ▷ int xs=flo[b][i];
\triangleright \triangleright if(g[b][x].w==0||e_delta(g[xs][x])<e_delta(g[b][x]))
\triangleright \ \triangleright \ | \ g[b][x]=g[xs][x],g[x][b]=g[x][xs];
▷ ▷ if(flo_from[xs][x])flo_from[b][x]=xs;
⊳ }
⊳ set_slack(b);
void expand_blossom(int b){
p for(size_t i=0;i<flo[b].size();++i)</pre>
> set_st(flo[b][i],flo[b][i]);
p int xr=flo_from[b][g[b][pa[b]].u],pr=get_pr(b,xr);

    for(int i=0;i<pr;i+=2){</pre>
▷ int xs=flo[b][i],xns=flo[b][i+1];
▷ ▷ pa[xs]=g[xns][xs].u;
\triangleright \triangleright S[xs]=1,S[xns]=0;
▷ □ q_push(xns);
> S[xr]=1,pa[xr]=pa[b];
p for(size_t i=pr+1;i<flo[b].size();++i){</pre>
▷ ▷ int xs=flo[b][i];
▷ ▷ S[xs]=-1,set_slack(xs);
⊳ }
⊳ st[b]=0;
bool on_found_edge(const edge &e){

  int u=st[e.u],v=st[e.v];

    if(S[v]==-1){
▷ ▷ pa[v]=e.u,S[v]=1;
▷ ▷ int nu=st[match[v]];
▷ ▷ slack[v]=slack[nu]=0;
▷ ▷ S[nu]=0,q_push(nu);

   }else if(S[v]==0){
▷ ▷ int lca=get_lca(u,v);
⊳ }

  return false;
bool matching(){

    memset(S+1,-1,sizeof(int)*n_x);
p memset(slack+1,0,sizeof(int)*n_x);

p q=queue<int>();

    for(int x=1; x<=n_x;++x)</pre>

    if(q.empty())return false;
p for(;;){
▷ ▷ while(q.size()){
```

```
▷ ▷ ▷ int u=q.front();q.pop();
▷ ▷ ▷ if(S[st[u]]==1)continue:
▷ ▷ for(int v=1;v<=n;++v)</pre>
\triangleright \triangleright \triangleright \vdash \mathbf{if}(g[u][v].w>0\&st[u]!=st[v])
\triangleright \triangleright \triangleright \triangleright  if (e_delta(g[u][v])==0) {
▷ ▷ ▷ ▷ ▷ if(on_found_edge(g[u][v]))return true;
▷ ▷ ▷ ▷ ▷ }else update_slack(u,st[v]);
▷ ▷ ▷ ▷ }
⊳ ⊳ }
▷ ▷ int d=INF;
\triangleright for(int b=n+1;b<=n_x;++b)
▷ ▷ if(st[b]==b&&S[b]==1)d=min(d,lab[b]/2);
\triangleright for(int x=1;x<=n_x;++x)

▷ ▷ □ if(st[x]==x&&slack[x]){
▷ ▷ ▷ if(S[x]==-1)d=min(d,e_delta(g[slack[x]][x]));

▷ ▷ ▷ else if(S[x]==0)d=min(d,e_delta(g[slack[x]][x])/2);
▷ ▷ ▷ }
\triangleright for(int u=1;u<=n;++u){
▷ ▷ ▷ if(S[st[u]]==0){
▷ ▷ ▷ if(lab[u]<=d)return 0;
▷ ▷ ▷ ▷ lab[u]-=d;
▷ ▷ } else if(S[st[u]]==1)lab[u]+=d;
⊳ ⊳ }
▷ ▷ ▷ if(st[b]==b){
▷ ▷ ▷ if(S[st[b]]==0)lab[b]+=d*2;
▷ ▷ ▷ else if(S[st[b]]==1)lab[b]-=d*2;
▷ ▷ ▷ }
▷ for(int x=1;x<=n_x;++x)</pre>
\triangleright \triangleright if(st[x]==x&&slack[x]&&st[slack[x]]!=x&&e_delta(g[slack[x]][x)
    →1)==0)
▷ ▷ ▷ if(on_found_edge(g[slack[x]][x]))return true;
\triangleright for(int b=n+1:b<=n x:++b)
\triangleright \triangleright if(st[b]==b\&\&S[b]==1\&\&lab[b]==0) expand_blossom(b);
⊳ }

    return false;

pair<long long,int> solve(){
p memset(match+1,0,sizeof(int)*n);
⊳ n x=n:

    int n_matches=0;

    long long tot_weight=0;

    for(int u=0;u<=n;++u)st[u]=u,flo[u].clear();</pre>

  int w max=0:
p for(int u=1;u<=n;++u)</pre>
\triangleright \triangleright flo_from[u][v]=(u==v?u:0);
▷ ▷ W_max=max(W_max,g[u][v].w);
⊳ ⊳ }

    for(int u=1;u<=n;++u)lab[u]=w_max;</pre>

    while(matching())++n_matches;

    for(int u=1;u<=n;++u)
</pre>
▷ if(match[u]&&match[u]<u)</pre>
▷ ▷ tot_weight+=g[u][match[u]].w;
preturn make_pair(tot_weight,n_matches);
void add_edge( int ui , int vi , int wi ){

    g[ui][vi].w = g[vi][ui].w = wi;

void init( int _n ){
\triangleright n = n:

    for(int u=1;u<=n;++u)</pre>
▷ for(int v=1; v<=n; ++v)</pre>
\triangleright \triangleright g[u][v]=edge(u,v,0);
3.7 Small to Large
75d792e0726aaba92088f4b4819687be, 24 lines
```

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```
void dfs(int u. int p. int keep){
    int big = -1, mmx = -1;
   for(int v : g[u]) if(v != p \&\& sz[v] > mmx)
       mmx = sz[v], big = v;
    for(int v : g[u]) if(v != p && v != big)
       dfs(v, u, 0);
    if(big != -1) dfs(big, u, 1);
    add(u, p, big);
    for(auto x : q[u]){
       // answer all queries for this vx
   if(!keep){ /*Remove data from this subtree*/ }
3.8 Kosaraiu
7e83f291d6d74e576ffc597c1a7b6e09, 16 lines
vector<int> g[N], gt[N], S; int vis[N], cor[N];
void dfs(int u){
vis[u] = 1; for(int v : g[u]) if(!vis[v]) dfs(v);
▷ S.push_back(u);
void dfst(int u, int e){
\triangleright cor[u] = e;

    for(int v : gt[u]) if(!cor[v]) dfst(v, e);

void kosaraju(){
p for(int i = 1; i <= n; i++) if(!vis[i]) dfs(i);</pre>
   for(int i = 1; i \le n; i++) for(int j : g[i])
        qt[i].push_back(i);
p int e = 0; reverse(S.begin(), S.end());
p for(int u : S) if(!cor[u]) dfst(u, ++e);
3.9 Tarjan
289a63a44243b11663e0693e6b4c90e8, 33 lines
int cnt = 0, root;
void dfs(int u, int p = -1){
\triangleright low[u] = num[u] = ++t;

    for(int v : g[u]){
▷ if(!num[v]){
\triangleright \triangleright \land dfs(v, u);
\triangleright \triangleright if(u == root) cnt++;
▷ ▷ ▷ if(low[v] >= num[u]) u PONTO DE ARTICULAÇÃO;
▷ ▷ ▷ if(low[v] > num[u]) ARESTA u->v PONTE;
\triangleright \triangleright \log[u] = \min(\log[u], \log[v]);
⊳ ⊳ }
▷ ▷ else if(v != p) low[u] = min(low[u], num[v]);
⊳ }
root PONTO DE ARTICULAÇÃO <=> cnt > 1
void tarjanSCC(int u){
\triangleright low[u] = num[u] = ++cnt;
\triangleright vis[u] = 1;
▷ S.push_back(u);

    for(int v : g[u]){
▷ if(!num[v]) tarjanSCC(v);
```

void cnt\_sz(int u, int p = -1){

for(int v : g[u]) if(v != p)

void add(int u, int p, int big = -1){

 $cnt_sz(v, u), sz[u] += sz[v];$ 

// Update info about this vx in global answer

for(int v : g[u]) if(v != p && v != big)

sz[u] = 1;

add(v, u);

```
▷ if(vis[v]) low[u] = min(low[u], low[v]);
⊳ }
\triangleright \mathbf{if}(\log[u] == num[u]) \{
\triangleright \triangleright ssc[u] = ++ssc\_cnt; int v;
> do{
\triangleright \triangleright \lor v = S.back(); S.pop_back(); vis[v] = 0;
▷ ▷ ▷ ssc[v] = ssc_cnt;
▷ ▷ }while(u != v);
⊳ }
3.10 Max Clique
a38a2529d78304f7da330a35b85d81e3, 37 lines
long long adj[N], dp[N];
for(int i = 0; i < n; i++){
\triangleright \text{ for(int } j = 0; j < n; j++)\{
⊳ ⊳ int x;
▷ ▷ scanf("%d",&x);
\triangleright \triangleright \mathbf{if}(\mathbf{x} \mid \mid \mathbf{i} == \mathbf{j})
▷ ▷ ▷ adj[i] |= 1LL << j;</pre>
⊳ }
}
int resto = n - n/2;
int C = n/2;
for(int i = 1; i < (1 << resto); i++){</pre>
\triangleright int x = i;
\triangleright for(int j = 0; j < resto; j++)
▷ ▷ if(i & (1 << j))</pre>
\triangleright \triangleright \triangleright x \&= adj[j + C] >> C;
\triangleright if(x == i){
▷ ▷ dp[i] = __builtin_popcount(i);
⊳ }
}
for(int i = 1; i < (1 << resto); i++)
\triangleright for(int j = 0; j < resto; j++)
\triangleright \ \mathbf{if}(i \& (1 << j))
\triangleright \triangleright dp[i] = max(dp[i], dp[i ^ (1 << j)]);
int maxCliq = 0;
for(int i = 0; i < (1 << C); i++){
\triangleright int x = i, y = (1 << resto) - 1;
\triangleright for(int j = 0; j < C; j++)
▷ ▷ if(i & (1 << j))</pre>
\triangleright \triangleright x \& = adj[j] \& ((1 << C) - 1), y \& = adj[j] >> C;

    if(x != i) continue;
b maxCliq = max(maxCliq, __builtin_popcount(i) + dp[y]);
3.11 Dominator Tree
df0303f87faa2c0f68eeb291d10a5d3e, 64 lines
vector<int> g[N], gt[N], T[N];
vector<int> S:
int dsu[N], label[N];
int sdom[N], idom[N], dfs_time, id[N];
vector<int> bucket[N];
vector<int> down[N];
void prep(int u){

  id[u] = ++dfs_time;
\triangleright label[u] = sdom[u] = dsu[u] = u;

    for(int v : g[u]){
▷ ▷ if(!id[v])
```

```
▷ ▷ prep(v), down[u].push_back(v);
▷ □ gt[v].push_back(u);
⊳ }
int fnd(int u, int flag = 0){

    if(u == dsu[u]) return u;

    int v = fnd(dsu[u], 1), b = label[ dsu[u] ];

b if(id[ sdom[b] ] < id[ sdom[ label[u] ] ])</pre>
\triangleright \triangleright label[u] = b;
\triangleright dsu[u] = v;
▶ return flag ? v : label[u];
void build_dominator_tree(int root, int sz){
▷ // memset(id, 0, sizeof(int) * (sz + 1));
▷ // for(int i = 0; i <= sz; i++) T[i].clear();</pre>
prep(root);

    reverse(S.begin(), S.end());
⊳ int w;

    for(int u : S){
▷ b for(int v : gt[u]){
\triangleright \triangleright \triangleright w = fnd(v);
\triangleright \triangleright \triangleright \triangleright sdom[u] = sdom[w];
⊳ ⊳ }
▷ ▷ gt[u].clear();
▷ b if(u != root) bucket[ sdom[u] ].push_back(u);
▷ br(int v : bucket[u]){
\triangleright \triangleright \triangleright w = fnd(v);
▷ ▷ if(sdom[w] == sdom[v]) idom[v] = sdom[v];
\triangleright \triangleright \triangleright \mathsf{else} \mathsf{idom}[v] = w;
▷ bucket[u].clear();
▷ for(int v : down[u]) dsu[v] = u;
▷ b down[u].clear();
⊳ }

    reverse(S.begin(), S.end());

    for(int u : S) if(u != root){
▷ ▷ if(idom[u] != sdom[u]) idom[u] = idom[ idom[u] ];
▷ ▷ T[ idom[u] ].push_back(u);
⊳ }
▷ S.clear();
3.12 Min Cost Matching
```

3f80580181f29cdafa11d5c83eb5988e, 44 lines

```
// Min cost matching
// O(n^2 * m)
// n == nro de linhas
// m == nro de colunas
// n <= m | flow == n
// a[i][j] = custo pra conectar i a j
vector<int> u(n + 1), v(m + 1), p(m + 1), way(m + 1);
for(int i = 1; i <= n; ++i){
    p[0] = i;
    int j0 = 0;
    vector<int> minv(m + 1 , oo);
    vector<char> used(m + 1 , false);
    do{
        used[j0] = true;
        int i0 = p[j0] , delta = oo, j1;
        for(int j = 1; j <= m; ++j)</pre>
```

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```
if(! used[i]){
              int cur = a[i0][j] - u[i0] - v[j];
              if(cur < minv[j])</pre>
                 minv[j] = cur, way[j] = j0;
              if(minv[j] < delta)</pre>
                 delta = minv[j] , j1 = j;
       for(int j = 0; j \le m; ++j)
          if(used[j])
              u[p[j]] += delta, v[j] -= delta;
          else
              minv[j] -= delta;
       i0 = i1;
   }while(p[j0] != 0);
       int j1 = way[j0];
       p[j0] = p[j1];
       j0 = j1;
   }while(j0);
// match[i] = coluna escolhida para linha i
vector<int> match(n + 1);
for(int j = 1; j \le m; ++j)
   match[p[j]] = j;
int cost = -v[0];
```

# 4 Strings

### 4.1 Aho Corasick

b4babc9c71e8a8481608f1b5fd5f3cfb, 35 lines

```
int to[N][A]:
int ne = 2, fail[N], term[N];
void add_string(const char *str, int id) {
   int p = 1;
   for(int i = 0; str[i]; i++) {
       int ch = str[i] - 'a';
      if(!to[p][ch]) to[p][ch] = ne++;
      p = to[p][ch];
   term[p]++;
void init() {
   for(int i = 0; i < ne; i++) fail[i] = 1;</pre>
   queue<int> q; q.push(1);
   while(!q.empty()){
       int u = q.front(); q.pop();
       for(int i = 0; i < A; i++){
          if(to[u][i]) {
             int v = to[u][i]; q.push(v);
             if(u != 1) {
                 fail[v] = to[ fail[u] ][i];
                 term[v] += term[ fail[v] ];
          else if(u != 1) to[u][i] = to[ fail[u] ][i];
          else to[u][i] = 1;
  }
void clean() {
   memset(to, 0, ne * sizeof(to[0]));
   memset(fail, 0, ne * sizeof(fail[0]));
   memset(term, 0, ne * sizeof(term[0]));
   ne = 2;
```

### 4.2 Suffix Array

namespace sf {

f013040501ada8671938368eb37d6eab, 57 lines

```
int lcp[N], c[N];
// Caractere final da string '\0' esta sendo considerado parte da
   \hookrightarrowstring s
void build_sa(char s[], int n, int a[]){
   const int A = 300: // Tamanho do alfabeto
   int c1[n], a1[n], h[n + A];
   memset(h, 0, sizeof h);
    for(int i = 0; i < n; i++) {
       c[i] = s[i];
       h[c[i] + 1]++;
   partial_sum(h, h + A, h);
    for(int i = 0; i < n; i++)
       a[h[c[i]]++] = i:
    for(int i = 0; i < n; i++)
       h[c[i]]--;
    for(int L = 1; L < n; L <<= 1) {
       for(int i = 0: i < n: i++) {
          int j = (a[i] - L + n) \% n;
          a1[h[c[j]]++] = j;
       int cc = -1;
       for(int i = 0; i < n; i++) {
          if(i == 0 || c[a1[i]] != c[a1[i-1]] || c[(a1[i] + L) % n]
              \hookrightarrow != c[(a1[i-1] + L) \% n])
              h[++cc] = i;
          c1[a1[i]] = cc:
       memcpy(a, a1, sizeof a1);
       memcpy(c, c1, sizeof c1);
       if(cc == n-1) break:
   }
}
void build_lcp(char s[], int n, int a[]){ // lcp[i] = lcp(s[:a[i]],
   \hookrightarrow s[:a[i+1]])
   int k = 0;
    //memset(lcp, 0, sizeof lcp);
    for(int i = 0; i < n; i++){
       if(c[i] == n-1) continue;
       int i = a[c[i]+1]:
       while(i+k < n && j+k < n && s[i+k] == s[j+k]) k++;
       lcp[c[i]] = k;
       if(k) k--;
}
int comp_lcp(int i, int j){
   if(i == i) return n - i:
   if(c[i] > c[j]) swap(i, j);
   return min(lcp[k] for k in [c[i], c[j]-1]);
4.3 Adamant Suffix Tree
5b7e1ff58ad2765a604e07a57b74afd5, 72 lines
```

```
const int inf = 1e9:
const int maxn = 200005:
char s[maxn];
map<int, int> to[maxn];
int len[maxn], fpos[maxn], link[maxn];
int node, pos;
int sz = 1, n = 0;
int make_node(int _pos, int _len) {
 fpos[sz] = _pos;
 len[sz] = _len;
 return sz++:
void go_edge() {
 while (pos > len[to[node][s[n - pos]]]) {
   node = to[node][s[n - pos]];
   pos -= len[node];
void add_letter(int c) {
 s[n++] = (char)c;
 pos++:
  int last = 0:
  while (pos > 0) {
   go_edge();
   int edge = s[n - pos];
   int &v = to[node][edge];
   int t = s[fpos[v] + pos - 1];
   if (v == 0) {
     v = make_node(n - pos, inf);
     link[last] = node:
     last = 0;
   } else if (t == c) {
     link[last] = node;
     return;
   } else {
     int u = make_node(fpos[v], pos - 1);
     to[u][c] = make\_node(n - 1, inf);
     to[u][t] = v;
     fpos[v] += pos - 1;
     len[v] -= pos - 1;
     v = 11:
     link[last] = u;
     last = u:
   if (node == 0)
     pos--;
   else
     node = link[node];
void add_string(char *str) {
 for (int i = 0; str[i]; i++) add_letter(str[i]);
 add_letter('$');
bool is_leaf(int u) { return len[u] > n; }
int get_len(int u) {
 if (!u) return 0;
 if (is_leaf(u)) return n - fpos[u];
 return len[u];
int leafs[maxn];
int calc_leafs(int u = 0) {
 leafs[u] = is_leaf(u);
 for (const auto &c : to[u]) leafs[u] += calc_leafs(c.second);
 return leafs[u]:
}; // namespace sf
```

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```
int main() { sf::len[0] = sf::inf: }
4.4 Z Algorithm
21483733f3e48570db9f780a714b967e, 11 lines
vector<int> z_algo(const string &s) {
\triangleright int n = s.size(), L = 0, R = 0;
▷ vector<int> z(n, 0);
\triangleright for(int i = 1; i < n; i++){
\triangleright if(i \le R) z[i] = min(z[i-L], R - i + 1);
\triangleright while(z[i]+i < n && s[ z[i]+i ] == s[ z[i] ])

▷ ▷ ▷ Z[i]++:
\triangleright if(i+z[i]-1 > R) L = i, R = i + z[i] - 1:
⊳ }
⊳ return z;
4.5 Prefix function/KMP
e940251be9782350ed6cd7ea97be7f94, 22 lines
vector<int> preffix_function(const string &s){
```

```
b int n = s.size(); vector<int> b(n+1);
b[0] = -1: int i = 0. i = -1:
\triangleright while(i < n){
b \in \mathbf{while}(j >= 0 \& s[i] != s[j]) j = b[j];
\triangleright \triangleright b\lceil ++i\rceil = ++i:
⊳ }
⊳ return b:
void kmp(const string &t, const string &p){
vector<int> b = preffix_function(p);
p int n = t.size(), m = p.size();
b int i = 0:
\triangleright for(int i = 0: i < n: i++){
b \in while(j >= 0 \&\& t[i] != p[j]) j = b[j];
\triangleright \triangleright if(i == m)
▷ ▷ ▷ //patern of p found on t
\triangleright \triangleright \triangleright j = b[j];
⊳ ⊳ }
⊳ }
```

# 4.6 Min rotation

e8ff5e094a8ae5d9ff7f169744e09963, 11 lines

```
// remember std::rotate
int min_rotation(int *s, int N) {
 REP(i, N) s[N+i] = s[i];
 int a = 0;
 REP(b, N) REP(i, N) {
  if (a+i == b \mid | s[a+i] < s[b+i]) { b += max(0, i-1); break; }
  if (s[a+i] > s[b+i]) \{ a = b; break; \}
 return a;
4.7 Manacher
```

7aacc600318d9535cfb0b8ec0371139b, 24 lines

```
// rad[2 * i] = largest palindrome cetered at char i
// rad[2 * i + 1] = largest palindrome cetered between chars i and
    \hookrightarrow i+i
void manacher(char *s, int n, int *rad) {
⊳ static char t[2*MAX];
\triangleright int m = 2 * n - 1:
\triangleright for(int i = 0; i < m; i++) t[i] = -1;
\triangleright for(int i = 0; i < n; i++) t[2 * i] = s[i];
\triangleright int x = 0:
```

```
p rad[0] = 0: // <</pre>
\triangleright for(int i = 1: i < m: i++) {
\triangleright \triangleright int \&r = rad[i] = 0;
\Rightarrow if(i <= x+rad[x]) r = min(rad[x+x-i],x+rad[x]-i);
\triangleright while(i - r - 1 >= 0 and i + r + 1 < m and
\triangleright \triangleright t[i - r - 1] == t[i + r + 1]) ++r;
\triangleright \vdash if(i + r >= x + rad[x]) x = i;
⊳ }
\triangleright for(int i = 0; i < m; i++) {
b if(i-rad[i] == 0 || i+rad[i] == m-1) ++rad[i];
\triangleright // for(int i = 0; i < m; i++) rad[i] /= 2;
4.8 Suffix Automaton
3b484c75f789b7207d1755bdcabbecc6, 25 lines
map<char, int> to[2*N];
int link[2*N], len[2*N], last = 0, sz = 1;
void add_letter(char c){
    int p = last;
    last = sz++;
    len[last] = len[p] + 1;
    for(; !to[p][c]; p = link[p]) to[p][c] = last;
    if(to[p][c] == last){}
        link[last] = 0;
        return:
    int u = to[p][c];
    if(len[u] == len[p]+1){
        link[last] = u;
        return:
    int c1 = sz++;
    to[c1] = to[u]:
    link[c1] = link[u];
    len[c1] = len[p]+1;
   link[last] = link[u] = c1;
    for(; to[p][c] == u; p = link[p]) to[p][c] = c1;
5 Geometry
5.1 2D basics
0fa022b65a23394df67bb6b2176b7a55, 284 lines
typedef double cod:
double eps = 1e-7;
bool eq(cod a, cod b){ return abs(a - b) <= eps; }</pre>
struct vec{
⊳ cod x, y; int id;
\triangleright vec(cod a = 0, cod b = 0) : x(a), y(b) {}

    vec operator+(const vec &o) const{
\triangleright return {x + o.x, y + o.y};
⊳ }

    vec operator-(const vec &o) const{
\triangleright return {x - o.x, y - o.y};

  vec operator*(cod t) const{
▷ return {x * t, y * t};
⊳ }

  vec operator/(cod t) const{
▷ return {x / t, y / t};
⊳ }

    cod operator*(const vec &o) const{ // cos
```

▷ return x \* o.x + y \* o.y;

```
    cod operator^(const vec &o) const{ // sin

    bool operator==(const vec &o) const{
\triangleright return eq(x, o.x) && eq(y, o.y);
▶ bool operator<(const vec &o) const{</pre>
\triangleright if(!eq(x, o.x)) return x < o.x;
⊳ ⊳ return y < o.y;
⊳ }
⊳ cod cross(const vec &a, const vec &b) const{
⊳ }
   int ccw(const vec &a, const vec &b) const{
      cod tmp = cross(a, b);
      return (tmp > eps) - (tmp < -eps);</pre>

    cod dot(const vec &a, const vec &b) const{

    return (a-(*this)) * (b-(*this));

⊳ }
▷ cod len() const{

    double angle(const vec &a, const vec &b) const{

    return atan2(cross(a, b), dot(a, b));

▷ return cross(a, b) / dot(a, b);
⊳ }
▶ vec unit() const{
▷ return operator/(len());
⊳ }

   int quad() const{
\triangleright if(x > 0 && y >=0) return 0;
\triangleright if(x <=0 && y > 0) return 1;
\triangleright if(x < 0 && y <=0) return 2;
⊳ ⊳ return 3:

    bool comp(const vec &a, const vec &b) const{
p return (a - *this).comp(b - *this);
bool comp(vec b){

    return (*this) * (*this) < b * b:
</pre>
⊳ }

    template<class T>

▶ void sort_by_angle(T first, T last) const{
▷ ▷ std::sort(first, last, [=](const vec &a, const vec &b){
▷ ▷ ▷ return comp(a, b);
▷ ▷ });
⊳ }
vec rot90() const{ return {-y, x}; }

  vec rot(double a) const{
\triangleright return {cos(a)*x -sin(a)*y, sin(a)*x +cos(a)*y};
   vec proj(const vec &b) const{ // proj of *this onto b
      cod k = operator*(b) / (b * b);
      return b * k;
   // proj of (*this) onto the plane orthogonal to b
   vec rejection(vec b) const{
      return (*this) - proj(b);
};
struct line{
⊳ cod a, b, c; vec n;
```

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```
\triangleright line(vec q, vec w){ // q.cross(w, (x, y)) = 0
\triangleright \triangleright a = -(w.y-q.y);
\triangleright \triangleright b = w.x-q.x;
\triangleright c = -(a * q.x + b * q.y);
\triangleright \triangleright n = \{a, b\};
⊳ }
⊳ cod dist(const vec &o) const{
▷ return abs(eval(o)) / n.len();
⊳ }
▶ bool contains(const vec &o) const{
\triangleright return eq(a * o.x + b * o.y + c, 0);
⊳ }

    cod dist(const line &o) const{
▷ if(!parallel(o)) return 0;
▷ ▷ if(!eq(a, 0))
▷ ▷ if(!eq(b, 0))
▷ ▷ ▷ return abs(c - o.c * b / o.b) / n.len();
▷ return abs(c - o.c);
⊳ }
▶ bool parallel(const line &o) const{
\triangleright return eq(n ^ o.n, 0);
⊳ }
▶ bool operator==(const line &o) const{
▷ if(!eq(a*o.b, b*o.a)) return false;
▷ if(!eq(a*o.c, c*o.a)) return false;
▷ if(!eq(c*o.b, b*o.c)) return false;
▷ ▷ return true;
▶ bool intersect(const line &o) const{
▷ return !parallel(o) || *this == o;

  vec inter(const line &o) const{
▷ ▷ if(parallel(o)){
▷ ▷ if(*this == o){ }
▷ ▷ ▷ else{ /* dont intersect */ }
> > }
\triangleright \triangleright auto tmp = n \hat{} o.n;

    return {(o.c*b -c*o.b)/tmp, (o.a*c -a*o.c)/tmp};
⊳ }
▷ vec at_x(cod x) const{
\triangleright return {x, (-c-a*x)/b};
⊳ }

  vec at_y(cod y) const{
\triangleright return {(-c-b*y)/a, y};
⊳ }

    cod eval(const vec &o) const{

    return a * o.x + b * o.y + c;

⊳ }
};
struct segment{
⊳ vec p, q;
⊳ segment(vec a = vec(), vec b = vec()): p(a), q(b) {}
▶ bool onstrip(const vec &o) const{ // onstrip strip
▷ return p.dot(o, q) >= -eps && q.dot(o, p) >= -eps;
⊳ }
⊳ cod len() const{
▷ return (p-q).len();
⊳ }

    cod dist(const vec &o) const{
▷ b if(onstrip(o)) return line(p, q).dist(o);
▷ return min((o-q).len(), (o-p).len());
⊳ }
▶ bool contains(const vec &o) const{

    return eq(p.cross(q, o), 0) && onstrip(o);
```

```
▶ bool intersect(const seament &o) const{
▷ if(contains(o.p)) return true;
▷ if(contains(o.q)) return true;
▷ if(o.contains(q)) return true;
▷ if(o.contains(p)) return true;
\triangleright return p.ccw(q, o.p) * p.ccw(q, o.q) == -1
      && o.p.ccw(o.q, q) * o.p.ccw(o.q, p) == -1;
▶ bool intersect(const line &o) const{
▷ return o.eval(p) * o.eval(q) <= 0;</pre>
⊳ }

    cod dist(const segment &o) const{
▷ ▷ if(onstrip(o.p) || onstrip(o.q)
▷ ▷ ▷ || o.onstrip(p) || o.onstrip(q))
▷ ▷ ▷ return line(p, q).dist(line(o.p, o.q));
⊳ ⊳ }
▷ ▷ else if(intersect(o)) return 0;
▷ return min(min(dist(o.p), dist(o.q)),
▷ ▷ ▷ min(o.dist(p), o.dist(q)));
⊳ }

    cod dist(const line &o) const{
▷ b if(line(p, q).parallel(o))
▷ ▷ return line(p, q).dist(o);
▷ ▷ else if(intersect(o)) return 0;
▷ return min(o.dist(p), o.dist(q));
⊳ }
};
struct hray{
⊳ vec p, q;
\triangleright hray(vec a = vec(), vec b = vec()): p(a), q(b){}
▶ bool onstrip(const vec &o) const{ // onstrip strip
▷ ▷ return p.dot(q, o) >= -eps;
⊳ cod dist(const vec &o) const{
▷ if(onstrip(o)) return line(p, q).dist(o);
▷ ▷ return (o-p).len();
▶ bool intersect(const segment &o) const{
▷ if(line(o.p, o.q).parallel(line(p,q)))
▷ ▷ return contains(o.p) || contains(o.q);

    return contains(line(p,q).inter(line(o.p,o.q)));
▶ bool contains(const vec &o) const{

    return eq(line(p, q).eval(o), 0) && onstrip(o);

    cod dist(const segment &o) const{
▷ if(line(p, q).parallel(line(o.p, o.q))){
▷ ▷ if(onstrip(o.p) || onstrip(o.q))
▷ ▷ ▷ return line(p, q).dist(line(o.p, o.q));
▷ ▷ ▷ return o.dist(p);
⊳ ⊳ }
▷ ▷ else if(intersect(o)) return 0;

    return min(min(dist(o.p), dist(o.q)),
▷ ▷ ▷ ▷ o.dist(p));
⊳ }
▶ bool intersect(const hray &o) const{
▷ if(!line(p, q).parallel(line(o.p, o.q)))

⊳ ⊳ return false;

▷ return contains(pt) && o.contains(pt); // <<</pre>
▶ bool intersect(const line &o) const{
```

```
p return (o.eval(p) >= -eps)^(o.eval(p)<o.eval(q));</pre>

    return contains(o.inter(line(p, q)));
⊳ }

    cod dist(const line &o) const{
▷ if(line(p,q).parallel(o))
▷ ▷ return line(p,q).dist(o);
▷ ▷ else if(intersect(o)) return 0;
▷ ▷ return o.dist(p);
⊳ }

    cod dist(const hray &o) const{
▷ ▷ if(onstrip(o.p) || o.onstrip(p))
▷ ▷ ▷ return line(p,q).dist(line(o.p, o.q));
▷ ▷ ▷ return (p-o.p).len();
⊳ ⊳ }
▷ ▷ else if(intersect(o)) return 0;
▷ return min(dist(o.p), o.dist(p));
⊳ }
};
double heron(cod a, cod b, cod c){
\triangleright cod s = (a + b + c) / 2;

  return sqrt(s * (s - a) * (s - b) * (s - c));
line mediatrix(const vec &a, const vec &b) {
\triangleright auto tmp = (b - a) * 2;
p return line(tmp.x, tmp.y, a * a - b * b);
struct circle {
⊳ vec c; cod r;

    circle() : c(0, 0), r(0) {}
\triangleright circle(const vec o) : c(o), r(0) {}
⊳ circle(const vec &a, const vec &b) {
\triangleright c = (a + b) * 0.5; r = (a - c).len();
⊳ circle(const vec &a, const vec &b, const vec &cc) {
▷ ▷ c = mediatrix(a, b).inter(mediatrix(b, cc));
\triangleright \triangleright r = (a - c).len();
▶ bool inside(const vec &a) const {
\triangleright return (a - c).len() <= r;
⊳ }
};
circle min_circle_cover(vector<vec> v) {

    random_shuffle(v.begin(), v.end());
⊳ circle ans;
p int n = (int)v.size();
\triangleright for(int i = 0; i < n; i++) if(!ans.inside(v[i])) {
▷ ▷ ans = circle(v[i]);
\triangleright \triangleright \mathbf{for}(\mathbf{int} \ j = 0; \ j < i; \ j++) \ \mathbf{if}(!ans.inside(v[j]))
▷ ▷ ▷ ans = circle(v[i], v[j]);
\triangleright \triangleright  for(int k=0; k<j; k++)if(!ans.inside(v[k])){
\triangleright \triangleright \triangleright \triangleright ans = circle(v[i], v[j], v[k]);
▷ ▷ ▷ }
⊳ ⊳ }
⊳ }
▶ return ans;
5.2 Circle line intersection
```

93f656aefa8c75b101df761fbb414d98, 21 lines

```
// intersection of line a * x + b * y + c = 0

// and circle centered at the origin with radius r

double r, a, b, c; // given as input

double x0 = -a^*c/(a^*a+b^*b), y0 = -b^*c/(a^*a+b^*b);

if(c^*c > r^*r^*(a^*a+b^*b)+EPS)

puts("no points");

else if(abs(c^*c - r^*r^*(a^*a+b^*b)) < EPS){
```

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```
puts("1 point");
  cout << x0 << ' ' ' << y0 << '\n';
}
else {
  double d = r*r - c*c/(a*a+b*b);
  double mult = sqrt (d / (a*a+b*b));
  double ax, ay, bx, by;
  ax = x0 + b * mult;
  bx = x0 - b * mult;
  ay = y0 - a * mult;
  by = y0 + a * mult;
  puts ("2 points");
  cout<<ax<' '<<ay<'\n'<<bx<' '<<by<'\n';
}</pre>
```

### 5.3 Half plane intersection

6eb846164d5affc44d8780095397a027, 43 lines

1ff067095f462e5c38bf1bce952d1368, 26 lines

```
const double eps = 1e-8;
typedef pair<long double, long double> pi;
bool z(long double x){ return fabs(x) < eps; }</pre>
struct line{

   long double a, b, c;

▶ bool operator<(const line &l)const{</pre>
\triangleright bool flag1 = pi(a, b) > pi(0, 0);
\triangleright bool flag2 = pi(1.a, 1.b) > pi(0, 0);
▷ ▷ if(flag1 != flag2) return flag1 > flag2;
\triangleright long double t = ccw(pi(0, 0), pi(a, b), pi(1.a, 1.b));
\triangleright return z(t) ? c * hypot(l.a, l.b) < l.c * hypot(a, b) : t > 0;
pi slope(){ return pi(a, b); }
pi cross(line a, line b){

    long double det = a.a * b.b - b.a * a.b;

    return pi((a.c * b.b - a.b * b.c) / det, (a.a * b.c - a.c * b.a)

   \rightarrow/ det):
bool bad(line a, line b, line c){
b if(ccw(pi(0, 0), a.slope(), b.slope()) <= 0) return false;</pre>
p pi crs = cross(a, b);
p return crs.first * c.a + crs.second * c.b >= c.c;
bool solve(vector<line> v, vector<pi> &solution){ // ax + by <= c;</pre>
⊳ sort(v.begin(), v.end());

    deque<line> dq;

    for(auto &i : v){
\triangleright if(!dq.empty() && z(ccw(pi(0, 0), dq.back().slope(), i.slope()))
→pop_back();
\triangleright while(dq.size() >= 2 && bad(i, dq[0], dq[1])) dq.pop_front();
▷ b dq.push_back(i);
⊳ }
b while(dq.size() > 2 && bad(dq[dq.size()-2], dq.back(), dq[0])) dq
▶ while(dq.size() > 2 && bad(dq.back(), dq[0], dq[1])) dq.pop_front

  vector<pi> tmp;
p for(int i=0; i<dq.size(); i++){</pre>
\triangleright line cur = dq[i], nxt = dq[(i+1)%dq.size()];
▷ if(ccw(pi(0, 0), cur.slope(), nxt.slope()) <= eps) return false;</pre>

    b tmp.push_back(cross(cur, nxt));

⊳ }
⊳ solution = tmp;
⊳ return true;
5.4 Detect empty Half plane intersection
```

```
// ccw(a, b, c) = a.ccw(b, c)
pair<bool, point> half_inter(vector<pair<point,point> > &vet){
   random_shuffle(all(vet));
   rep(i,0,sz(vet)) if(ccw(vet[i].x,vet[i].y,p) != 1){
       point dir = (vet[i].y - vet[i].x) / abs(vet[i].y - vet[i].x)
       point 1 = vet[i].x - dir*1e15;
       point r = vet[i].x + dir*1e15;
       if(r < 1) swap(1, r);
       rep(j, 0, i){
          if(ccw(point(), vet[i].x-vet[i].y, vet[j].x-vet[j].y) ==
              if(ccw(vet[j].x, vet[j].y, p) == 1)
                  continue;
              return mp(false, point());
          if(ccw(vet[j].x, vet[j].y, 1) != 1)
              l = max(l, line_intersect(vet[i].x,vet[i].y,vet[j].x,
                  \hookrightarrowvet[j].y));
          if(ccw(vet[j].x, vet[j].y, r) != 1)
              r = min(r, line_intersect(vet[i].x,vet[i].y,vet[j].x,
                  \hookrightarrowvet[i].y));
          if(!(1 < r)) return mp(false, point());</pre>
       }
       p = r;
   return mp(true, p);
```

### 5.5 Circle Circle intersection

// abs(point a) = absolute value of a

Assume that the first circle is centered at the origin and second at (x2, y2). Find circle line intersection of first circle and line Ax + By + C = 0, where  $A = -2x_2$ ,  $B = -2y_2$ ,  $C = x_2^2 + y_2^2 + r_1^2 - r_2^2$ .

Be aware of corner case with two circles centered at the same point.

## 5.6 Tangents of two circles

236b49cb5fffc14ce471c189d5568dbe, 25 lines

```
// solve first for same circle(and infinitely many tangents)
// Find up to four tangents of two circles
void tangents(pt c, double r1, double r2, vector<line> & ans){
   double r = r2 - r1;
   double z = c.x * c.x + c.y * c.y;
   double d = z - r * r:
   if(d < -EPS) return;</pre>
   d = sqrt(abs(d));
   line 1:
   1.a = (c.x * r + c.y * d) / z;
   1.b = (c.y * r - c.x * d) / z;
   1.c = r1:
   ans.push_back (1);
vector<line> tangents(circle a, circle b){
   vector<line> ans;
   pt aux = a.center - b.center;
   for(int i = -1; i \le 1; i += 2)
       for(int j = -1; j <= 1; j += 2)
          tangents(aux, a.r * i, b.r * j, ans);
   for(size_t i = 0; i < ans.size(); ++i)</pre>
```

```
ans[i].c = ans[i].a * a.x + ans[i].b * a.y;
   return ans:
5.7 Convex Hull
7988fcc13274e6fd81a2b039e434efa9, 29 lines
vector<vec> monotone_chain_ch(vector<vec> P){
   sort(P.begin(), P.end());
   vector<vec> L, U;
   for(auto p : P){
       // BE CAREFUL WITH OVERFLOW!
       // MAX VALUE (2*A)^2, where 0 \le abs(p.x), abs(p.y) \le A
       while(L.size() >= 2 && L[L.size() - 2].cross(L.back(), p) <=</pre>
          L.pop_back();
       L.push_back(p);
   reverse(P.begin(), P.end());
   for(auto p : P){
       while(U.size() >= 2 && U[U.size() - 2].cross(U.back(), p) <=</pre>
          U.pop_back();
      U.push_back(p);
   L.pop_back(), U.pop_back();
   L.reserve(L.size() + U.size());
   L.insert(L.end(), U.begin(), U.end());
   return L;
5.8 Check point inside polygon
c72c1e4a59d54003ea32f314dbd5dec1, 18 lines
bool below(const vector<vec> &vet, vec p){
b auto it = lower_bound(vet.begin(), vet.end(), p);
   if(it == vet.end()) return false:
b if(it == vet.begin()) return *it == p;
return prev(it)->cross(*it, p) <= 0;</pre>
bool above(const vector<vec> &vet, vec p){
p auto it = lower_bound(vet.begin(), vet.end(), p);
   if(it == vet.end()) return false;
b if(it == vet.begin()) return *it == p;
return prev(it)->cross(*it, p) >= 0;
// lowerhull, upperhull and point, borders included
bool inside_poly(const vector<vec> &lo, const vector<vec> &hi, vec
return below(hi, p) && above(lo, p);
5.9 Check point inside polygon without lower/upper
      hull
9c7fe9fd47da401b280b35249c70709f, 19 lines
// borders included
// must not have 3 colinear consecutive points
bool inside_poly(const vector<vec> &v, vec p){
   if(v[0].ccw(v[1], p) < 0) return false;
```

if(v[0].ccw(v.back(), p) > 0) return 0;

if(v[0].ccw(v.back(), p) == 0)

```
return v[0].dot(p, v.back()) >= 0
    && v.back().dot(p, v[0]) >= 0;

int L = 1, R = (int)v.size() - 1, ans = 1;

while(L <= R){
    int mid = (L+R)/2;
    if(v[0].ccw(v[mid], p) >= 0) ans = mid, L = mid+1;
    else R = mid-1;
}

return v[ans].ccw(v[(ans+1)%v.size()], p) >= 0;
```

### 5.10 Minkowski sum

b667e546607ccd322f49c3d48e684d53, 25 lines

```
vector<vec> msum(vector<vec>& a, vector<vec>& b) {
\triangleright int i = 0, j = 0;
\triangleright for(int k = 0; k < (int)a.size(); k++)
\triangleright if(a[k] < a[i]) i = k;
\triangleright for(int k = 0; k < (int)b.size(); k++)
\triangleright \triangleright \mathbf{if}(b[k] < b[j]) j = k;
▶ vector<vec> c:
p c.reserve(a.size() + b.size());
p for(int k = 0; k < int(a.size()+b.size()); k++){</pre>
▷ ▷ if((int)c.size() >= 2
\triangleright \triangleright \triangleright && c[c.size()-2].ccw(c.back(), pt) == 0)
▷ ▷ ▷ c.pop_back();
▷ ▷ c.push_back(pt);
\triangleright int q = i+1, w = j+1;
\triangleright if(q == int(a.size())) q = 0;
b if(w == int(b.size())) w = 0;
b \in \mathbf{if}(c.back().ccw(a[i]+b[w], a[q]+b[j]) < 0) i = q;
\triangleright \triangleright else j = w;
⊳ }
▷ c.shrink_to_fit();
⊳ return c:
```

### 5.11 Geo Notes

### 5.11.1 Center of mass

**System of points(2D/3D):** Mass weighted average of points.

**Frame(2D/3D):** Get middle point of each segment solve as previously.

**Triangle:** Average of vertices.

**2D Polygon:** Compute **signed** area and center of mass of triangle  $((0,0), p_i, p_{i+1})$ . Then solve as system of points.

**Polyhedron surface:** Solve each face as a 2D polygon(be aware of (0, 0)) then replace each face with its center of mass and solve as system of points.

**Tetrahedron(Triangular pyramid):** As triangles, its the average of points.

**Polyhedron:** Can be done as 2D polygon, but with tetrahedralization intead of triangulation.

### 5.11.2 Pick's Theorem

Given a polygon without self-intersections and all its vertices on integer coordinates in some 2D grid. Let A be its area, I the number of points with integer coordinates stricly inside the polygon and B the number of points with integer coordinates in the border of the polygon. The following formula holds:  $A = I + \frac{B}{2} - 1$ .

### 6 Miscellaneous

p int first\_bits = v & (SIZE - 1);

### 6.1 Cute LIS

```
ed1e47c2a22e616f141bb36cd05e167f, 7 lines
multiset<int> S:
for(int i = 0; i < n; i++){

    auto it = S.upper_bound(a[i]); // low for inc

p if(it != S.end()) S.erase(it);
▷ S.insert(a[i]);
ans = S.size();
6.2 Efficient recursive lambda
c675838f4f976169cec6d509c5bd1bf5, 21 lines
template<class Fun>
class y_combinator_result {
 Fun fun_;
public:
 template<class T>
 explicit y_combinator_result(T &&fun): fun_(std::forward<T>(fun))
 template<class ...Args>
 decltype(auto) operator()(Args &&...args) {
   return fun_(std::ref(*this), std::forward<Args>(args)...);
};
template<class Fun>
decltype(auto) y_combinator(Fun &&fun) {
 return y_combinator_result<std::decay_t<Fun>>(std::forward<Fun>(
// auto gcd = y_combinator([](auto gcd, int a, int b) -> int {
// return b == 0 ? a : gcd(b, a % b);
6.3 Bitsets
022047d007596ba7d8fee4d0057c9c71, 33 lines
#define private public
#include <bitset>
#undef private
#include <bits/stdc++.h>
using namespace std;
#define tab _M_w
using biti = typename remove_reference<decltype(bitset<404>().tab
    \hookrightarrow [0])>::type;
const int SIZE = 8 * sizeof(biti);
const int LOG = __builtin_ctz(SIZE);
template<size_t Nw>
int find_prev(const bitset<Nw> &x, int v) {
p int start = v >> LOG;
```

```
▷ if(first bits) {
▷ ▷ biti curr = x.tab[start];
▷ ▷ ▷ return start << LOG | (SIZE - __builtin_clzl(curr) - 1);
\triangleright for(int i = start - 1; i >= 0; i--) {
▷ ▷ biti curr = x.tab[i];
▷ ▷ if(curr) {
\triangleright \triangleright return (i << LOG) | (SIZE - __builtin_clzl(curr) - 1);
⊳ }

  return -1;

// s._Find_first(); s._Find_next(k); find_prev(s, k+1);
// _Unchecked_set/_Unchecked_reset/_Unchecked_flip
6.4 Buildings
8290ed09c45bdfce47ffca1ce6ae5c8f, 11 lines
// count the number of circular arrays of size m. with elements on
   \hookrightarrow range [1, c^{**}(n^*n)]
int n, m, c; cin >> n >> m >> c;
int x = f_{exp}(c, n * n); int ans = f_{exp}(x, m);
for(int i = 1; i \le m; i++) if(m \% i == 0) {
  int y = f_exp(x, i);
  for(int j = 1; j < i; j++) if(i % j == 0)
     y = sub(y, mult(j, dp[j]));
  dp[i] = mult(y, inv(i));
 ans = sub(ans, mult(i - 1, dp[i]));
cout << ans << '\n';</pre>
6.5 Rand
2c44688813ebbf7967f35d07e6553580, 6 lines
#include <random>
#include <chrono>
cout << RAND MAX << endl:</pre>
mt19937 rng(chrono::steady_clock::now().time_since_epoch().count())
shuffle(p.begin(), p.end(), rng);
uniform_int_distribution<int>(a,b)(rng);
6.6 Klondike
88a355600bcefd0e5b7607c18edbec23, 17 lines
// minimum number of moves to make
// all elements equal
// move: change a segment of equal value
// elements to any value
int v[305], dp[305][305], rec[305][305];
int f(int 1, int r){
 if(r == 1) return 1;
 if(r < 1) return 0;</pre>
 if(dp[l][r] != -1) return dp[l][r];
  int ans = f(1+1, r) + 1;
  for(int i = l+1; i \le r; i++)
   \mathbf{if}(v[i] == v[1])
     ans = min(ans, f(1, i - 1) + f(i+1, r));
  return dp[l][r] = ans;
6.7 Hilbert Order
940598d402c0c52731fbd959a4207b9e, 19 lines
// maybe use B = n / sqrt(q) before this
inline int64_t hilbertOrder(int x, int y, int pow = 21, int rotate
```

**⇒= 0)** {

if(pow == 0) return 0;

```
for(int i = 0: i < N: i++){
\triangleright int hpow = 1 << (pow-1):
\triangleright int seq = (x < hpow) ? (
                                                                         \triangleright for(int j = 0; j < w[i]; j++)
\triangleright \triangleright (y < hpow) ? 0 : 3
                                                                        ▷ ▷ Q[i].clear();
                                                                        \triangleright \mathbf{for}(\mathbf{int} \ j = 0; \ j \leftarrow M; \ j++)\{
▷): (
\triangleright \triangleright (y < hpow) ? 1 : 2
                                                                        \triangleright p = Q[j \% w[i]];
                                                                        ⊳);
                                                                        ▷ □ q.add(c[i]);
⊳ seg = (seg + rotate) & 3;

    const int rotateDelta[4] = {3, 0, 0, 1};

                                                                        ▷ push(d[j]);
\triangleright int nx = x & (x ^ hpow), ny = y & (y ^ hpow);
                                                                        ▷ ▷ d[j] = q.getmin();

    int nrot = (rotate + rotateDelta[seg]) & 3;

                                                                         ⊳ }
b int64_t subSquareSize = int64_t(1) << (2*pow - 2);</pre>
                                                                         6.11 LCA < O(nlgn), O(1)>
p int64_t ans = seg * subSquareSize;
                                                                         57ee488ab554896ae9c3fd0e9fc8d6f0, 19 lines

    int64_t add = hilbert0rder(nx, ny, pow-1, nrot);
\triangleright ans += (seg == 1 || seg == 2) ? add : (subSquareSize - add - 1);
                                                                         int start[N], dfs_time;
▶ return ans:
                                                                         int tour[2*N], id[2*N];
6.8 Modular Factorial
                                                                         void dfs(int u){
ce57a0d9307f6be1554bca7f60816caa, 23 lines
                                                                             start[u] = dfs_time;
// Compute (1*2*...*(p-1)*1*(p+1)*(p+2)*..*n) % p
                                                                            id[dfs\_time] = u;
// in O(p*lg(n))
                                                                            tour[dfs_time++] = start[u];
int factmod(int n. int p){
                                                                            for(int v : a[u]){
   int ans = 1;
                                                                                dfs(v);
   while (n > 1) {
                                                                                id[dfs_time] = u;
       for(int i = 2; i \le n \% p; i++)
                                                                                tour[dfs_time++] = start[u];
          ans = (ans * i) % p;
                                                                        }
       n /= p;
      if(n \% 2) ans = p - ans;
                                                                         int LCA(int u. int v){
   return ans % p:
                                                                            if(start[u] > start[v]) swap(u, v):
                                                                            return id[min(tour[k]for k in [start[u],start[v]])];
int fac_pow(int n, int p){
                                                                         6.12 Buffered reader
   int ans = 0;
                                                                         d0e88e8e88f781ee107ad1ff4442d439, 22 lines
   while(n) n /= p, ans += n;
   return ans;
                                                                         // source: https://github.com/ngthanhtrung23/ACM Notebook new/blob/
                                                                            int C(int n, int k, int p){
                                                                         int INP.AM.REACHEOF:
   if(fac_pow(n, p) > fac_pow(n-k, p) + fac_pow(k, p))
                                                                         #define BUFSIZE (1<<12)</pre>
       return 0;
                                                                         char BUF[BUFSIZE+1], *inp=BUF;
   int tmp = factmod(k, p) * factmod(n-k, p) % p;
                                                                         #define GETCHAR(INP) { \
   return (f_{exp}(tmp, p - 2, p) * factmod(n, p)) % p;
                                                                            if(!*inp && !REACHEOF) { \
                                                                                memset(BUF,0,sizeof BUF);\
6.9 Iterate over submasks
                                                                                int inpzzz = fread(BUF,1,BUFSIZE,stdin);\
2e979d7fe5b96f7bc29b625ab42eeeef, 10 lines
                                                                                if (inpzzz != BUFSIZE) REACHEOF = true;\
// loop through all submask of a given bitmask
                                                                                inp=BUF: \
// it does not include mask 0
for(int sub = mask; sub; sub = (sub - 1) & mask){
                                                                            INP=*inp++; \
                                                                         #define DIG(a) (((a)>='0')&&((a)<='9'))
                                                                         #define GN(j) { \
// loop through all supermasks of a given bitmask
                                                                            AM=0:\
for(int super = mask; super < (1 << n); super = (super + 1) | mask)</pre>
                                                                            GETCHAR(INP); while(!DIG(INP) && INP!='-') GETCHAR(INP);\
                                                                            if (INP=='-') {AM=1;GETCHAR(INP);} \
   \hookrightarrow {
                                                                            j=INP-'0'; GETCHAR(INP); \
                                                                            while(DIG(INP)){j=10*j+(INP-'0');GETCHAR(INP);} \
6.10 Knapsack Bounded with Cost
                                                                            if (AM) j=-j;\
2a7be64fb7e82706b487228a4dc2c2b3, 19 lines
                                                                         6.13 Modular summation
// menor custo para conseguir peso ate M usando N tipos diferentes
                                                                         6ce73a9b66343d3281c83808039a686e, 40 lines

→ de elementos, sendo que o i-esimo elemento pode ser usado b[i]

   \hookrightarrow vezes, tem peso w[i] e custo c[i]
                                                                         //calcula (sum(0 \ll i \ll n) P(i)) \% mod,
// O(N * M)
                                                                         //onde P(i) eh uma PA modular (com outro modulo)
                                                                         namespace sum_pa_mod{
                                                                        int b[N], w[N], c[N];
```

▷ ▷ assert(a&&b);

 $\triangleright \triangleright 11 \text{ ret} = ((n*(n+1)/2)\%\text{mod})*(a/b);$ 

 $\triangleright \triangleright if(a\%b) \text{ ret} = (\text{ret} + \text{calc}(a\%b,b,n,mod))\%\text{mod};$ 

▷ b if(a >= b){

MinQueue Q[M]

int d[M] //d[i] = custo minimo para conseguir peso i

for(int i = 0;  $i \leftarrow M$ ; i++) d[i] = i ? oo : 0;

```
▷ ▷ ▷ return ret;
⊳ ⊳ }
\triangleright return ((n+1)*(((n*a)/b+1)%mod) - calc(b,a,(n*a)/b,mod) + mod +
   \hookrightarrown/b + 1)%mod;
⊳ }
\triangleright //P(i) = a*i mod m
\triangleright \triangleright a = (a\%m + m)\%m;
▷ ▷ if(!a) return 0:
> 11 \text{ ret} = (n*(n+1)/2) \text{ mod}:
▷ ▷ ret = (ret*a)%mod;
\triangleright \triangleright 11 q = \__qcd(a,m);
\triangleright ret -= m*(calc(a/g,m/g,n,mod)-n-1);
▷ ▷ return (ret%mod + mod)%mod;
\triangleright //P(i) = a + r*i \mod m
\triangleright \triangleright a = (a\%m + m)\%m:
\triangleright r = (r\%m + m)\%m;
\triangleright if(!r) return (a*(n+1))%mod:
▷ if(!a) return solve(r, n, m, mod);
\triangleright \triangleright 11 g, x, y;
\triangleright g = gcdExtended(r, m, x, y);
\triangleright \triangleright x = (x\%m + m)\%m;
\triangleright \ \triangleright \ 11 \ d = a - (a/g)*g;
\triangleright \triangleright a = d;
\triangleright \triangleright x = (x*(a/q))\%m;
▷ return (solve(r, n+x, m, mod) - solve(r, x-1, m, mod) + mod + d
    \hookrightarrow*(n+1))%mod;
⊳ }
};
6.14 Edge coloring CPP
bd934a3806a9c263b93d393cb773525e, 44 lines
const int MX = 300;
int C[MX][MX] = \{\}, G[MX][MX] = \{\};
void solve(vector<pii> &E, int N){
    int X[MX] = \{\}, a, b;
    auto update = [&](int u){ for(X[u] = 1; C[u][X[u]]; X[u]++); };
    auto color = [&](int u, int v, int c){
        int p = G[u][v]:
        G[u][v] = G[v][u] = c;
        C[u][c] = v; C[v][c] = u;
        C[u][p] = C[v][p] = 0;
        if(p) X[u] = X[v] = p;
        else update(u), update(v);
        return p; };
    auto flip = [&](int u, int c1, int c2){
        int p = C[u][c1], q = C[u][c2];
        swap(C[u][c1], C[u][c2]);
        if( p ) G[u][p] = G[p][u] = c2;
        if( !C[u][c1] ) X[u] = c1;
        if( !C[u][c2] ) X[u] = c2;
        return p; };
    for(int i = 1; i <= N; i++) X[i] = 1;</pre>
    for(int t = 0; t < E.size(); t++){</pre>
        int u = E[t].first, v0 = E[t].second, v = v0, c0 = X[u], c =
            \hookrightarrow c0. d:
        vector<pii> L;
        int vst[MX] = {};
        while(!G[u][v0]){
            L.emplace_back(v, d = X[v]);
```

▷ ▷ **else** ret = (ret+n+1)%mod:

### 6.15 Burnside's Lemma

Let  $(G, \oplus)$  be a finite group that acts on a set X. It should hold that  $e_g * x = x$  and  $g_1 * (g_2 * x) = (g_1 \oplus g_2) * x$ ,  $\forall x \in X, g_1, g_2 \in G$ . For each  $g \in G$  let  $X^g = \{x \in X \mid g * x = x\}$ . The number of orbits its given by:

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|$$

### 6.16 Wilson's Theorem

 $(n-1)! = -1 \mod n \iff n \text{ is prime}$ 

### 6.17 Fibonacci

- $F_{n-1}F_{n+1} F_n^2 = (-1)^n$
- $F_{n+k} = F_k F_{n+1} + F_{k-1} F_n$
- $GCD(F_n, F_m) = F_{GCD(n,m)}$
- $F_n = \frac{(\frac{1+\sqrt{5}}{2})^n (\frac{1-\sqrt{5}}{2})^n}{\sqrt{5}}$

### 6.18 Lucas's Theorem

For non-negative integers m and n and a prime p, the following congruence holds:

$$\binom{m}{n} \equiv \prod_{i=0}^{k} \binom{m_i}{n_i} \pmod{p}$$

where  $m_i$  is the i-th digit of m in base p.  $\binom{a}{b} = 0$  if a < b.

### 6.19 Kirchhoff's Theorem

Laplacian matrix is L = D - A, where D is a diagonal matrix with vertex degrees on the diagonals and A is adjacency matrix.

The number of spanning trees is any cofactor of L. i-th cofactor is determinant of the matrix gotten by removing i-th row and column of L.

### 6.19.1 Multigraphs

In D[i][i] all loops are excluded. A[i][j] = number of edges from i to j.

### 6.19.2 Directed multigraphs

D[i][i] = indegree of i minus the number of loops at i. A[i][j] = number of edges from i to j.

The number of oriented spanning trees rooted at a vertex i is the determinant of the matrix gotten by removing the ith row and column of L.

### 6.20 Matroid

Let *X* set of objects,  $I \subseteq 2^X$  set of independents sets such that:

- 1.  $\emptyset \in I$
- 2.  $A \in I, B \subseteq A \implies B \in I$
- 3. Exchange axiom,  $A \in I, B \in I, |B| > |A| \implies \exists x \in B \setminus A : A \cup \{x\} \in I$
- 4.  $A \subseteq X$  and I and I' are maximal independent subsets of A then |I| = |I'|

Then (X, I) is a matroid. The combinatorial optimization problem associated with it is: Given a weight  $w(e) \ge 0 \ \forall e \in X$ , find an independet subset that has the largest possible total weight.

### 6.21 Matroid intersection

55f1a2840ddf17787899e0809856b6bd, 43 lines

```
// Input two matroids (X, I_a) and (X, I_b)
// output set I of maximum size, I \setminusin I_a and I \setminusin I_b
set<> I;
while(1){
    for(e_i : X \setminus I)
        if(I + e_i \in I_a \text{ and } I + e_i \in I_b)
            I = I + e_i;
    set<> A, T; queue<> Q;
    for(x : X) label[x] = MARK1;
    for(e_i : X \setminus I){
        if(I + e_i \in I_a)
            Q.push(e_i), label[e_i] = MARK2;
        else{
            for (x \text{ such that } I - x + e_i \setminus in I_a)
                A[x].push(e_i);
        if(I + e_i \setminus in I_b)
           T = T + \{e_i\}
```

```
for(x such that I - x + e_i \setminus in I_b)
              A[e_i].push(x);
      }
   if(T.empty()) break;
   bool found = false;
   while(!Q.empty() and !found){
       auto e = Q.front(); Q.pop();
       for(x : A[e]) if(label[x] == MARK1){
          label[x] = e; Q.push(x);
          if(x \in T)
              found = true; put = 1;
              while(label[x] != MARK2){
                 I = put ? (I + x) : (I - x);
                 put = 1 - put;
              I = I + x;
              break;
       }
   if(!found) break;
return I;
```

Where path(e) = [e] if label[e] = MARK2, path(label[e]) + [e] otherwise.

### 6.21.1 Matroid Union

Given k matroids over the same set of objects  $(X, I_1)$ ,  $(X, I_2)$ , ...,  $(X, I_k)$  find  $A_1 \in I_1$ ,  $A_2 \in I_2$ , ...,  $A_k \in I_k$  such that  $i \neq j$ ,  $A_i \cap A_j = \emptyset$  and  $\bigcup_{i=1}^k A_i |$  is maximum. Matroid union can be reduced to matroid intersection as follows.

Let  $X' = X \times \{1, 2, ..., k\}$ , ie, k copies of each element of X with different colors. M1 = (X', Q) where  $B \in Q \iff \forall 1 \le i \le k$ ,  $\{x \mid (x, i) \in B\} \in I_i$ , ie, for each color, B is independent. M2 = (X', W) where  $B \in W \iff i \ne j \implies \neg((x, i) \in B \land (x, j) \in B)$ , ie, each element is picked by at most one color.

Intersection of *M*1 and *M*2 is the answer for the combinatorial problem of matroid union.

### 6.22 Notes

When we repeat something and each time we have probability p to succeed then the expected number or tries is  $\frac{1}{n}$ , till we succeed.

### Small to large

**Trick in statement** If k sets are given you should note that the amount of different set sizes is  $O(\sqrt{s})$  where s is total size of those sets. And no more than  $\sqrt{s}$  sets have size greater than  $\sqrt{s}$ . For example, a path to the root

in Aho-Corasick through suffix links will have at most  $O(\sqrt{s})$  vertices.

**gcd on subsegment**, we have at most  $\log(a_i)$  different values in  $\{\gcd(a_j, a_{j+1}, ..., a_i) \text{ for } j < i\}$ .

**From static set to expandable.** To insert, create a new set with the new element. While there are two sets with same size, merge them. There will be at most log(n) disjoints sets.

**Matrix exponentiation optimization**. Save binary power of  $A_{nxn}$  and answer q queries  $b = A^m x$  in  $O((n^3 + qn^2)log(m))$ .

**Ternary search on integers into binary search**, comparing f(mid) and f(mid+1), binary search on deriva-

tive

**Dynamic offline set** For each element we will wind segment of time [a, b] such that element is present in the set during this whole segment. Now we can come up with recursive procedure which handles [l, r] time segment considering that all elements such that  $[l, r] \subset [a, b]$  are already included into the set. Now, keeping this invariant we recursively go into [l, m] and [m + 1, r] subsegments. Finally when we come into segment of length 1.

 $a > b \implies a \mod b < \frac{a}{2}$ 

**Convex Hull**. The expected number of points in the convex hull of a random set of points is O(log(n)). The number of points in a convex hull with points coordinates 264241152, 382205952, 530841600

limited by *L* is  $O(L^{2/3})$ .

**Tree path query**. Sometimes the linear query is fast enough. Just do adamant's hld sorting subtrees by their size and remap vertices indexes.

**Range query** offline can be solved by a sweep, ordering queries by R.

Maximal number of divisors of any n-digit number. 7 4, 12, 32, 64, 128, 240, 448, 768, 1344, 2304, 4032, 6720, 10752, 17280, 26880, 41472, 64512, 103680, 161280, 245760, 368640, 552960, 860160, 1290240, 1966080, 2764800, 4128768, 6193152, 8957952, 13271040, 19660800, 28311552, 41287680, 59719680, 88473600, 127401984, 181665792, 264241152, 382205952, 530841600