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set ts=4 sw=4 sta nu rnu sc stl+=%F cindent
set bg=dark ruler timeoutlen=1000
imap {<CR> {<CR>}<Esc>0
nmap <F2> 0V$%d
nmap <C-down> :m+1<CR>
nmap <C-up> :m-2<CR>
nmap <C-a> ggVG
nmap < S-up > :m-2 < CR >
nmap <S-down> :m+1<CR>
syntax on
vmap <C-c> "+y
set viminfo='20,\"1000
alias comp='q++ -std=c++17 -Wshadow -Wall -Wextra -Wformat=2
    →-Wconversion -fsanitize=address,undefined

→-fno-sanitize-recover -Wfatal-errors'

#include <bits/stdc++.h>
#define ff first
#define ss second
#define pb push_back
using namespace std;
using 11 = long long;
using ii = pair<int, int>;
const int N = 100005;
int main() {

  return 0;

1 Data Structures
```

1.1 Fenwick Tree 2D

```
vector<int> go[N];
vector<int> ft[N];
void prec_add(int x, int y) {
\triangleright for(; x < N; x += x & -x) {
▷ □ go[x].push_back(y);
⊳ }
void init() {
\triangleright for(int i = 1; i < N; i++) {
▷ ▷ sort(go[i].begin(), go[i].end());
▷ ▷ go[i].resize(unique(go[i].begin(), go[i].end()) -
     →go[i].begin());
b ft[i].assign(go[i].size() + 1, 0);
⊳ }
void add(int x, int y, int val) {
\triangleright for(; x < N; x += x & -x) {
▷ int id = int(upper_bound(go[x].begin(), go[x].end(), y) -
    \hookrightarrowgo[x].begin());
▷ b for(; id < (int)ft[x].size(); id += id & -id)</pre>
▷ ▷ ▷ ft[x][id] += val;
⊳ }
int sum(int x, int y) {
\triangleright int ans = 0;
\triangleright for(; x > 0; x -= x & -x) {
```

```
1.2 Wavelet Tree
template<typename T>
class wavelet { // 1-based!!
   T L, R;
   vector<int> 1;
▷ vector<T> sum: // <<</pre>
⊳ wavelet *lef, *rig;
b int r(int i) const{ return i - l[i]; }
public:

    template<typename ITER>

   wavelet(ITER bg, ITER en) { // it changes the argument array
▷ ▷ lef = rig = nullptr;
      L = *bq, R = *bq;
▷ b for(auto it = bg; it != en; it++)
         L = min(L, *it), R = max(R, *it);
▷ ▷ if(L == R) return;
      T \text{ mid} = L + (R - L)/2;
▷ ▷ l.reserve(std::distance(bg, en) + 1);
▷ ▷ sum.reserve(std::distance(bg, en) + 1);
▷ l.push_back(0), sum.push_back(0);
▷ ▷ sum.push_back(sum.back() + *it);
▷ ▷ ▷ return x <= mid;
▷ ▷ });
⊳ ~wavelet(){
⊳ ⊳ delete lef:

▷ b delete rig;

⊳ }
⊳ // 1 index, first is 1st
   T kth(int i, int j, int k) const{
      if(L >= R) return L;
      int c = l[j] - l[i-1];
      if(c >= k) return lef->kth(l[i-1]+1, l[j], k);
      else return rig->kth(r(i-1)+1, r(j), k - c);
\triangleright // # elements > x on [i, i]
p int cnt(int i, int j, T x) const{
\triangleright if(L > x) return j - i + 1;
\triangleright int ans = 0;
\triangleright if(lef) ans += lef->cnt(l[i-1]+1, l[j], x);
\triangleright if(rig) ans += rig->cnt(r(i-1)+1, r(j), x);
⊳ ⊳ return ans:
⊳ }
▷ // sum of elements <= k on [i, j]</pre>

    T sumk(int i, int j, T k){
      if(L == R) return R \le k ? L * (j - i + 1) : 0;
▷ if(R <= k) return sum[j] - sum[i-1];</pre>
```

```
▷ ▷ int ans = 0;
▷ ▷ if(lef) ans += lef->sumk(l[i-1]+1, l[j], k);
▷ ▷ if(rig) ans += rig->sumk(r(i-1)+1, r(j), k);
▷ ▷ return ans;
▷ }
▷ // swap (i, i+1) just need to update "array" l[i]
};
```

1.3 Order Set

1.4 Hash table

```
#include <ext/pb_ds/assoc_container.hpp>
using namespace and pbds:
struct custom_hash {
> static uint64_t splitmix64(uint64_t x) {
▷ ▷ // http://xorshift.di.unimi.it/splitmix64.c
\triangleright x += 0x9e3779b97f4a7c15;
\triangleright x = (x \hat{ } (x >> 30)) * 0xbf58476d1ce4e5b9;
\triangleright x = (x \hat{x} >> 27)) * 0x94d049bb133111eb;
\triangleright return x ^ (x >> 31):
⊳ }
> size_t operator()(uint64_t x) const {
▷ ▷ static const uint64_t FIXED_RANDOM =

chrono::steady_clock::now().time_since_epoch().count();

⊳ }
};
gp_hash_table<long long, int, custom_hash> table;
unordered_map<long long, int, custom_hash> uhash;
uhash.reserve(1 << 15);</pre>
uhash.max load factor(0.25):
```

1.5 Convex Hull Trick Simple

```
struct Line{
   11 m, b;
   inline 11 eval(11 x) const{
       return x * m + b;
};
// min => cht.back().m >= L.m
// max => cht.back().m <= L.m
void push_line(vector<Line> &cht, Line L){
 while((int)cht.size() >= 2){
   int sz = (int)cht.size();
   if((long double)(L.b-cht[sz-1].b)*(cht[sz-2].m-L.m)
  <= (long double)(L.b-cht[sz-2].b)*(cht[sz-1].m-L.m)){</pre>
     cht.pop_back();
   else break;
 cht.push_back(L);
// x increasing; pos = 0 in first call
11 linear_search(const vector<Line> &cht,ll x,int &pos){
```

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```
while(pos+1 < (int)cht.size()){</pre>
/*>>*/ if(cht[pos].eval(x) >= cht[pos+1].eval(x)) pos++;
       else break;
   return cht[pos].eval(x);
11 binary_search(const vector<Line> &cht, 11 x){
   int L = 0, R = (int)cht.size()-2;
   int bans = (int)cht.size()-1;
   while(L \le R)
       int mid = (L+R)/2:
       if(cht[mid].eval(x) >= cht[mid+1].eval(x)) // <<<</pre>
          L = mid + 1;
       else bans = mid, R = mid - 1;
   return cht[bans].eval(x);
```

1.6 Convex Hull Trick

```
const 11 is query = -(1LL<<62):
struct Line{
⊳ 11 m, b;

    mutable function < const Line*() > succ:
▶ bool operator<(const Line& rhs) const{</pre>
▷ ▷ if(rhs.b != is_query) return m < rhs.m;</pre>
▷ ▷ const Line* s = succ():
▷ ▷ if(!s) return 0;
\triangleright \triangleright 11 x = rhs.m:
\triangleright return b - s->b < (s->m - m) * x;
⊳ }
};
struct Cht : public multiset<Line>{ // maintain max
▶ bool bad(iterator v){
\triangleright \triangleright auto z = next(v):
\triangleright \triangleright \mathbf{if}(y == begin()){
\triangleright \triangleright \vdash if(z == end()) return 0;
\triangleright \triangleright return y->m == z->m && y->b <= z->b;
⊳ ⊳ }
\triangleright \triangleright auto x = prev(y);
\triangleright if(z == end()) return y->m == x->m && y->b <= x->b;
\Rightarrow return (long double)(x->b - v->b)*(z->m - v->m) >= (long
    \hookrightarrow double) (v->b-z->b)*(v->m-x->m):
⊳ }

    void insert line(ll m. ll b){
▷ ▷ auto y = insert({ m, b });
\triangleright \lor v->succ = [=]{ return next(y) == end() ? 0 : &*next(y); };
▷ if(bad(y)){ erase(y); return; }

    b while(y != begin() && bad(prev(y))) erase(prev(y));

⊳ ll eval(ll x){
▷ ▷ auto 1 = *lower_bound((Line) { x, is_query });
\triangleright return 1.m * x + 1.b;
⊳ }
};
```

1.7 Convex Hull Trick

```
* Author: Simon Lindholm
    →https://github.com/kth-competitive-programming/kactl/blob/maste
 * License: CC0
struct Line {

    mutable 11 m, b, p;
```

```
bool operator<(const Line& o) const { return m < o.m: }</pre>

    bool operator<(11 x) const { return p < x; }</pre>
struct LineContainer : multiset<Line, less<>> { // CPP14 only
\triangleright // (for doubles, use inf = 1/.0, div(a,b) = a/b)

    const 11 inf = LLONG_MAX;
\triangleright 11 div(11 a, 11 b) { // floored division

    return a / b - ((a ^ b) < 0 && a % b); }
</pre>
▶ bool isect(iterator x, iterator y) {
\triangleright if (y == end()) { x->p = inf; return false; }
\triangleright if (x->m == y->m) x->p = x->b > y->b ? inf : -inf;
\triangleright else x->p = div(y->b - x->b, x->m - y->m);
⊳ }

  void add(ll m, ll b) {
\triangleright auto z = insert({m, b, 0}), y = z++, x = y;
\triangleright while (isect(y, z)) z = erase(z);
\triangleright if (x != begin() && isect(--x, y)) isect(x, y = erase(y));
\Rightarrow while ((y = x) != begin() && (--x)->p >= y->p)
▷ ▷ isect(x, erase(y));
⊳ }
⊳ ll querv(ll x) {
▷ ▷ assert(!empty());
▷ ▷ auto 1 = *lower_bound(x);
\triangleright return 1.m * x + 1.b;
⊳ }
};
```

1.8 Min queue

```
template<typename T>
class minQ{

    deque<tuple<T, int, int> > p;
⊳ T delta;
⊳ int sz;
public:
p minQ() : delta(0), sz(0) {}

    inline int size() const{ return sz: }
p inline void add(T x){ delta += x: }

    inline void push(T x, int id){
▷ ▷ x -= delta, sz++;
\triangleright \triangleright int t = 1;
b while(p.size() > 0 && get<0>(p.back()) >= x)
▷ ▷ t += get<1>(p.back()), p.pop_back();
▷ p.emplace_back(x, t, id);
⊳ }

   inline void pop(){

    p get<1>(p.front())--, sz--;
b if(!get<1>(p.front())) p.pop_front();
p T getmin() const{ return get<0>(p.front())+delta; }
p int getid() const{ return get<2>(p.front()); }
```

1.9 Sparse Table

```
int fn(int i. int i){

  if(i == 0) return v[i]:
b if(~dn[i][j]) return dn[i][j];
\triangleright return dn[i][j] = min(fn(i, j-1), fn(i + (1 << (j-1)), j-1));
rint_getmn(int_ltrintur){/(ineconfliner.h
\Rightarrow int lz = lg(r - l + 1);
\triangleright return min(fn(1, lz), fn(r - (1 << lz) + 1, lz)):
1.10 Treap
```

```
// source:
   → https://qithub.com/victorsenam/caderno/blob/master/code/treap.cpp
//const int N = ; typedef int num;
num X[N]; int en = 1, Y[N], sz[N], L[N], R[N];
void calc (int u) { // update node given children info

    if(!u) return:
\triangleright sz[u] = sz[L[u]] + 1 + sz[R[u]];
⊳ // code here, no recursion
void unlaze (int u) {

    if(!u) return:
⊳ // code here, no recursion
void split_val(int u, num x, int &l, int &r) { // l gets <= x, r</pre>
\triangleright unlaze(u); if(!u) return (void) (1 = r = 0);
\triangleright if(X[u] <= x) { split_val(R[u], x, 1, r); R[u] = 1; 1 = u; }
else { split_val(L[u], x, 1, r); L[u] = r; r = u; }
⊳ calc(u):
void split sz(int u. int s. int &l. int &r) { // l gets first s. r
    ⇒gets remaining
\Rightarrow unlaze(u); if(!u) return (void) (1 = r = 0);
\triangleright if(sz[L[u]] < s) { split_sz(R[u], s - sz[L[u]] - 1, 1, r); R[u]
    \hookrightarrow= 1; 1 = u; }
else { split_sz(L[u], s, l, r); L[u] = r; r = u; }
⊳ calc(u):
int merge(int 1. int r) { // els on 1 <= els on r</pre>
\Rightarrow unlaze(1): unlaze(r): if(!1 || !r) return 1 + r: int u:
\triangleright if(Y[1] > Y[r]) { R[1] = merge(R[1], r); u = 1; }

    else { L[r] = merge(1, L[r]); u = r; }

⊳ calc(u); return u;
void init(int n=N-1) { // XXX call before using other funcs
\triangleright for(int i = en = 1; i <= n; i++) { Y[i] = i; sz[i] = 1; L[i] =
    \hookrightarrow R[i] = 0: 
random_shuffle(Y + 1, Y + n + 1);
void insert(int &u. int it){
□ unlaze(u);
\triangleright if(!u) u = it;

    else if(Y[it] > Y[u]) split val(u, X[it], L[it], R[it]), u = it:

    else insert(X[it] < X[u] ? L[u] : R[u], it);
</pre>
⊳ calc(u):
void erase(int &u, num key){
□ unlaze(u):

    if(!u) return;

\triangleright if(X[u] == key) u = merge(L[u], R[u]);

    else erase(key < X[u] ? L[u] : R[u], key);
</pre>
⊳ calc(u);
int create node(num kev){
\triangleright X[en] = key;
\triangleright sz[en] = 1;
\triangleright L[en] = R[en] = 0;

  return en++;
int query(int u, int 1, int r){//0 index
□ unlaze(u):

    if(u! or r < 0 or l >= sz[u]) return identity_element;

\triangleright if(1 <= 0 and r >= sz[u] - 1) return subt_data[u];

  int ans = query(L[u], 1, r);

\triangleright if(1 <= sz[L[u]] and sz[L[u]] <= r)
\triangleright \triangleright ans = max(ans, st[u]);
\rightarrow ans = max(ans, query(R[u], 1-sz[L[u]]-1, r-sz[L[u]]-1));
```

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```
⊳ return ans:
1.11 ColorUpdate
   source:  \hookrightarrow \text{https://github.com/tfg50/Competitive-Programming/tree/master/Billioteca/Data%20Structures} // subtree of u => [in[u], out[u]) 
template <class Info = int>
class ColorUpdate {
⊳ set<Range> ranges;
public:
⊳ struct Range {
\triangleright \triangleright Range(int a = 0) : 1(a) {}
\triangleright Range(int a, int b, Info c) : 1(a), r(b), v(c) {}
⊳ ⊳ Info v;
▷ ▷ bool operator<(const Range &b) const { return 1 < b.1; }</pre>
vector<Range> upd(int 1, int r, Info v) {
▷ ▷ vector<Range> ans;
\triangleright if(1 >= r) return ans;
▷ auto it = ranges.lower bound(1):
▷ b if(it != ranges.begin()) {
⊳ ⊳ ⊳ it--;
\triangleright \triangleright if(it->r > 1) {
▷ ▷ ▷ ▷ auto cur = *it;
▷ ▷ ▷ ranges.erase(it):
▷ ▷ ▷ ranges.emplace(cur.1, 1, cur.v);
▷ ▷ ▷ ranges.emplace(1, cur.r, cur.v);
⊳ ⊳ }
▷ it = ranges.lower_bound(r);
▷ if(it != ranges.begin()) {
⊳ ⊳ ⊳ it--;
\triangleright \triangleright \triangleright if(it->r > r) {
▷ ▷ ▷ ▷ auto cur = *it:
▷ ▷ ▷ ▷ ranges.erase(it);
▷ ▷ ▷ ranges.emplace(cur.l, r, cur.v);
▷ ▷ ▷ ranges.emplace(r, cur.r, cur.v);
⊳ ⊳ }
▷ ▷ ▷ ans.push_back(*it);
▷ ranges.emplace(1, r, v);

⊳ return ans;

⊳ }
};
1.12 Heavy Light Decomposition
void dfs_sz(int u){
   sz[u] = 1;
   for(auto &v : g[u]) if(v == p[u]){
       swap(v, g[u].back()); g[u].pop_back();
   for(auto &v : g[u]){
```

```
p[v] = u; dfs_sz(v); sz[u] += sz[v];
       if(sz[v] > sz[g[u][0]])
          swap(v, g[u][0]);
   }
// nxt[u] = start of path with u
// set nxt[root] = root beforehand
void dfs_hld(int u){
   in[u] = t++;
```

```
rin[in[u]] = u;
   for(auto v : g[u]){
       nxt[v] = (v == g[u][0] ? nxt[u] : v); dfs_hld(v);
   out[u] = t;
// path from nxt[u] to u \Rightarrow [in[nxt[u]], in[u]]
1.13 Iterative Segtree
T query(int 1, int r){ // [1, r]
   T rl, rr;
   for(1 += n, r += n+1; 1 < r; 1 >>= 1, r >>= 1){
       if(l & 1) rl = merge(rl, st[l++]);
       if(r & 1) rr = merge(st[--r], rr);
   return merge(rl, rr);
// initially save v[i] in st[n+i] for all i in [0, n)
void build(){
   for(int p = n-1; p > 0; p--)
       st[p] = merge(st[2*p], st[2*p+1]);
void update(int p, T val){
   st[p += n] = val:
   while(p >>= 1) st[p] = merge(st[2*p], st[2*p+1]);
1.14 LiChao's Segtree
void add_line(line nw, int v = 1, int l = 0, int r = maxn) { //
    \hookrightarrow \lceil 1, r \rceil
   int m = (1 + r) / 2;
   bool lef = nw.eval(1) < st[v].eval(1);</pre>
   bool mid = nw.eval(m) < st[v].eval(m);</pre>
   if(mid) swap(st[v], nw);
   if(r - 1 == 1) {
   } else if(lef != mid) {
       add_line(nw, 2 * v, 1, m);
       add_line(nw, 2 * v + 1, m, r);
int get(int x, int v = 1, int l = 0, int r = maxn) {
   int m = (1 + r) / 2;
   if(r - 1 == 1) {
       return st[v].eval(x);
   \} else if(x < m) {
       return min(st[v].eval(x), get(x, 2*v, 1, m));
       return min(st[v].eval(x), get(x, 2*v+1, m, r));
1.15 Palindromic tree
```

```
#include <bits/stdc++.h>
using namespace std;
const int maxn = 3e5 + 1, sigma = 26;
int len[maxn], link[maxn], to[maxn][sigma];
int slink[maxn], diff[maxn], series_ans[maxn];
int sz, last, n;
char s[maxn];
```

```
void init()
   s[n++] = -1;
   link[0] = 1;
   len[1] = -1;
   sz = 2;
int get_link(int v)
   while(s[n - len[v] - 2] != s[n - 1]) v = link[v];
   return v;
void add_letter(char c)
   s[n++] = c -= 'a';
   last = get_link(last);
   if(!to[last][c])
      len[sz] = len[last] + 2;
      link[sz] = to[get_link(link[last])][c];
      diff[sz] = len[sz] - len[link[sz]];
      if(diff[sz] == diff[link[sz]])
          slink[sz] = slink[link[sz]];
          slink[sz] = link[sz];
      to[last][c] = sz++;
   last = to[last][c];
int main()
   ios::sync_with_stdio(0);
   cin.tie(0);
   init();
   string s;
   cin >> s;
   int n = s.size();
   int ans[n + 1];
   memset(ans, 63, sizeof(ans));
   ans[0] = 0:
   for(int i = 1; i \le n; i++)
      add_letter(s[i - 1]);
      for(int v = last; len[v] > 0; v = slink[v])
          series_ans[v] = ans[i - (len[slink[v]] + diff[v])];
          if(diff[v] == diff[link[v]])
             series_ans[v] = min(series_ans[v],
                ans[i] = min(ans[i], series_ans[v] + 1);
      cout << ans[i] << "\n";</pre>
   return 0;
```

2 Math

2.1 Extended Euclidean Algorithm

```
// a*x + b*y = gcd(a, b), < gcd, x, y>
tuple<int, int, int> gcd(int a, int b) {
b if(b == 0) return make_tuple(a, 1, 0);

    auto [q, w, e] = gcd(b, a % b);
```

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```
  return make_tuple(q, e, w - e * (a / b));
```

2.2 Chinese Remainder Theorem

```
// x = vet[i].first (mod vet[i].second)
11 crt(const vector<pair<11, 11>> &vet){
   11 \text{ ans} = 0, 1 \text{cm} = 1;
   ll a, b, g, x, y;
   for(const auto &p : vet) {
       tie(a, b) = p;
       tie(q. x. v) = qcd(lcm. b):
       if((a - ans) % g != 0) return -1; // no solution
       ans = ans + x * ((a - ans) / g) % (b / g) * lcm;
      lcm = lcm * (b / q);
       ans = (ans \% lcm + lcm) \% lcm;
   return ans:
```

2.3 Diophantine Solver

```
template<typename T>
T extqcd(T a, T b, T &x, T &y) {
 if (a == 0) {
   x = 0:
   y = 1;
   return b;
 T p = b / a;
 T g = extgcd(b - p * a, a, y, x);
 x -= p * y;
 return q;
template<typename T>
bool diophantine(T a, T b, T c, T &x, T &y, T &g) {
 if (a == 0 && b == 0) {
   if (c == 0) {
    x = y = g = 0;
     return true;
   return false:
 if (a == 0) {
   if (c % b == 0) {
    x = 0;
    y = c / b;
     g = abs(b);
     return true;
   return false;
 if (b == 0) {
   if (c % a == 0) {
    x = c / a:
    y = 0;
    q = abs(a);
    return true:
   return false;
 g = extgcd(a, b, x, y);
 if (c % g != 0) {
   return false:
 T dx = c / a:
 c -= dx * a;
 T dy = c / b;
```

```
c -= dv * b:
x = dx + mulmod(x, c / g, b);
y = dy + mulmod(y, c / g, a);
g = abs(g);
return true;
```

2.4 Preffix inverse

```
inv[1] = 1:
for(int i = 2; i < p; i++)
\triangleright inv[i] = (p - (p/i) * inv[p%i] % p) % p;
```

2.5 Pollard Rho

```
11 rho(11 n){
b if(n % 2 == 0) return 2;
⊳ 11 d, c, x, y, prod;
⊳ do{
\triangleright \triangleright c = 11rand(1, n - 1);
\triangleright \triangleright x = 11rand(1, n - 1);
\triangleright \triangleright y = x;
⊳ ⊳ prod = 1:
\triangleright for(int i = 0; i < 40; i++) {
\triangleright \triangleright \triangleright x = add(mul(x, x, n), c, n);
\triangleright \triangleright \lor y = add(mul(y, y, n), c, n);
\triangleright \triangleright \lor y = add(mul(y, y, n), c, n);
\triangleright \triangleright \mathsf{prod} = \mathsf{mul}(\mathsf{prod}, \mathsf{abs}(\mathsf{x} - \mathsf{y}), \mathsf{n}) ?: \mathsf{prod};
\triangleright \triangleright d = \_\_gcd(prod, n);
b } while(d == 1);
⊳ return d;
ll pollard_rho(ll n){
⊳ 11 x, c, y, d, k;
⊳ int i;
> do{
\triangleright \triangleright i = 1;
\triangleright x = 11rand(1, n-1), c = 11rand(1, n-1);
\triangleright \triangleright y = x, k = 4;
> do{
\triangleright \triangleright if(++i == k) y = x, k *= 2;
\triangleright \triangleright \triangleright x = add(mul(x, x, n), c, n);
\triangleright \triangleright \triangleright d = \_\gcd(abs(x - y), n);
▷ ▷ }while(d == 1);
> }while(d == n):
⊳ return d;
void factorize(ll val, map<ll, int> &fac){

    if(rabin(val)) fac[ val ]++;
⊳ else{
▷ ▷ 11 d = pollard_rho(val);
▷ ▷ factorize(d, fac);
⊳ }
map<ll. int> factor(ll val){

    map<11, int> fac;

    if(val > 1) factorize(val, fac);
⊳ return fac;
```

2.6 Miller Rabin

```
bool rabin(ll n){

    if(n <= 1) return 0;</pre>

    if(n <= 3) return 1:</pre>
> 11 s = 0, d = n - 1;
\triangleright while(d % 2 == 0) d /= 2, s++;
```

```
\triangleright for(int k = 0: k < 64: k++){
\triangleright 11 a = 11rand(2, n-2):
\triangleright 11 x = fexp(a, d, n);
\triangleright \mathbf{if}(x != 1 \&\& x != n-1) \{
\triangleright \triangleright \triangleright for(int r = 1; r < s; r++){
\triangleright \triangleright \triangleright \triangleright x = mul(x, x, n);
\triangleright \triangleright \triangleright if(x == 1) return 0;
\triangleright \triangleright \triangleright \triangleright  if(x == n-1) break;
\triangleright \triangleright if(x != n-1) return 0;
⊳ ⊳ }
⊳ }
⊳ return 1;
```

2.7 Primitive root

```
// a primitive root modulo n is any number g such that any c
   \hookrightarrowcoprime to n is congruent to a power of g modulo n.
bool exists root(ll n){
   if(n == 1 || n == 2 || n == 4) return true;
   if(n \% 2 == 0) n /= 2;
   if(n % 2 == 0) return false:
    // test if n is a power of only one prime
   for(11 i = 3: i * i \le n: i += 2) if(n % i == 0)
        while(n % i == 0) n /= i:
       return n == 1:
   return true;
11 primitive root(11 n){
   if(n == 1 || n == 2 || n == 4) return n - 1;
   if(not exists_root(n)) return -1;
   11 x = phi(n):
   auto pr = factorize(x);
   auto check = [x. n. pr](11 m){
        for(11 p : pr) if(fexp(m, x / p, n) == 1)
           return false:
       return true:
   for(11 m = 2; ; m++) if(\_gcd(m, n) == 1)
       if(check(m)) return m:
\ensuremath{/\!/} Let's denote R(n) as the set of primitive roots modulo n, p is
// g \in R(p) \Rightarrow (pow(g, p-1, p * p) == 1 ? g+p : g) \in R(pow(p, p-1, p * p) == 1 ? g+p : g)
    (\rightarrow k)), for all k > 1
// g \text{ in } R(pow(p, k)) \Rightarrow (g \% 2 == 1 ? g : g + pow(p, k)) \setminus in
    \hookrightarrow R(2*pow(p, k))
```

2.8 Mobius Function

```
memset(mu, 0, sizeof mu);
mu[1] = 1:
for(int i = 1; i < N; i++)
   for(int j = i + i; j < N; j += i)
       mu[j] -= mu[i];
// g(n) = sum\{f(d)\} \Rightarrow f(n) = sum\{mu(d)*g(n/d)\}
```

2.9 Mulmod TOP

```
constexpr uint64_t mod = (1ull<<61) - 1;</pre>
uint64_t modmul(uint64_t a, uint64_t b){
\rightarrow uint64_t 11 = (uint32_t)a, h1 = a>>32, 12 = (uint32_t)b, h2 =
\triangleright uint64_t l = 11*12, m = 11*h2 + 12*h1, h = h1*h2;

    uint64_t ret = (1&mod) + (1>>61) + (h << 3) + (m >> 29) + (m <<</pre>
   \hookrightarrow35 >> 3) + 1;
p ret = (ret & mod) + (ret>>61);
```

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```
> ret = (ret & mod) + (ret>>61);
> return ret-1;
}
```

2.10 Modular multiplication TOPPER

```
11 mulmod(11 a, 11 b, 11 mod) {
    11 q = 11((long double)a * (long double)b / (long double)mod);
    11 r = (a * b - mod * q) % mod;
    if(r < 0) r += mod;
    return r;
}</pre>
```

2.11 Division Trick

```
for(int l = 1, r; l <= n; l = r + 1) {
    r = n / (n / 1);
    // n / x yields the same value for l <= x <= r
}
for(int l, r = n; r > 0; r = l - 1) {
    int tmp = (n + r - 1) / r;
    l = (n + tmp - 1) / tmp;
    // (n+x-1) / x yields the same value for l <= x <= r
}</pre>
```

2.12 Matrix Determinant

```
long double a[n][n];
long double gauss(){
   long double det = 1;
   for(int i = 0; i < n; i++){
       int q = i;
       for(int j = i+1; j < n; j++){
          if(abs(a[j][i]) > abs(a[q][i]))
             q = j;
       if(abs(a[q][i]) < EPS){
          det = 0;
          break;
       if(i != q){
          for(int w = 0: w < n: w++)
              swap(a[i][w], a[q][w]);
          det = -det;
       det *= a[i][i]:
       for(int j = i+1; j < n; j++) a[i][j] /= a[i][i];</pre>
       for(int j = 0; j < n; j++) if(j != i){
          if(abs(a[i][i]) > EPS)
              for(int k = i+1; k < n; k++)
                 a[j][k] -= a[i][k] * a[j][i];
   }
   return det:
```

2.13 Simplex Method

```
⊳ dbl ans:
⊳ int n. m:

    dbl sol[N];

    void pivot(int x, int y){
\triangleright \triangleright swap(X[y], Y[x]);
\triangleright \triangleright b[x] /= A[x][y];
\triangleright for(int i = 0; i < n; i++)
⊳ ⊳ ⊳ if(i != y)
\triangleright \triangleright \triangleright \triangleright A[x][i] /= A[x][y];
\triangleright \land A[x][y] = 1. / A[x][y];
\triangleright for(int i = 0; i < m; i++)
\triangleright \triangleright \models \mathbf{if}(i != x \&\& abs(A[i][y]) > eps) 
\triangleright \triangleright \triangleright \triangleright b[i] -= A[i][y] * b[x];
\triangleright \triangleright \triangleright \triangleright for(int j = 0; j < n; j++) if(j != y)
                           A[i][j] -= A[i][y] * A[x][j];
\triangleright \triangleright \triangleright A[i][y] = -A[i][y] * A[x][y];
▷ ▷ ▷ }
\triangleright ans += c[y] * b[x];
\triangleright for(int i = 0; i < n; i++)
▷ ▷ ▷ if(i != y)
\triangleright \triangleright \triangleright \triangleright c[i] -= c[y] * A[x][i];
\triangleright c[y] = -c[y] * A[x][y];
⊳ }
▷ // maximiza sum(x[i] * c[i])
⊳ // sujeito a
\triangleright // sum(a[i][j] * x[j]) <= b[i] para 0 <= i < m (Ax <= b)
\triangleright // x[i] >= 0 para 0 <= i < n (x >= 0)
⊳ // (n variaveis, m restricoes)
⊳ // guarda a resposta em ans e retorna o valor otimo

    dbl solve(int _n, int _m) {
\triangleright this->n = _n; this->m = _m;
           for(int i = 1; i < m; i++){
                int id = uniform_int_distribution<int>(0, i)(rng);
                swap(b[i], b[id]);
                for(int j = 0; j < n; j++)
                      swap(A[i][j], A[id][j]);
\triangleright \triangleright ans = 0.;
\triangleright for(int i = 0; i < n; i++) X[i] = i;
\triangleright for(int i = 0; i < m; i++) Y[i] = i + n;
▷ ▷ while(true) {
\triangleright \triangleright \triangleright int x = min_element(b, b + m) - b;
\triangleright \triangleright \triangleright \mathbf{if}(b[x] >= -eps)
▷ ▷ ▷ ▷ break;
\triangleright \triangleright int y = find_if(A[x], A[x] + n, [](dbl d) \{ return d < -eps; \}
    \hookrightarrow}) - A[x];
\triangleright \triangleright if(y == n) throw 1; // no solution
▷ ▷ ▷ pivot(x, y);
⊳ ⊳ }
▷ ▷ while(true) {
\triangleright \triangleright \triangleright int y = max_element(c, c + n) - c;
\triangleright \triangleright \triangleright if(c[y] \le eps) break;
\triangleright \triangleright int x = -1;
\triangleright \triangleright \triangleright dbl mn = 1. / 0.;
\triangleright \triangleright for(int i = 0; i < m; i++)
\triangleright \triangleright \triangleright \vdash if(A[i][y] > eps && b[i] / A[i][y] < mn)
\triangleright \triangleright \triangleright \triangleright \triangleright mn = b[i] / A[i][y], x = i;
▷ ▷ if(x == -1) throw 2; // unbounded
▷ ▷ ▷ pivot(x, y);
\triangleright for(int i = 0; i < m; i++)
\triangleright \triangleright \triangleright \mathbf{if}(Y[i] < n)
```

▷ ▷ ▷ ▷ sol[Y[i]] = b[i];

2.14 FFT

```
void fft(vector<base> &a, bool inv){
   int n = (int)a.size();
   for(int i = 1, j = 0; i < n; i++){
      int bit = n \gg 1:
      for(; j >= bit; bit >>= 1) j -= bit;
       j += bit;
       if(i < j) swap(a[i], a[j]);
   for(int sz = 2; sz <= n; sz <<= 1) {</pre>
       double ang = 2 * PI / sz * (inv ? -1 : 1);
      base wlen(cos(ang), sin(ang));
      for(int i = 0; i < n; i += sz){</pre>
          base w(1, 0);
          for(int j = 0; j < sz / 2; j++){
             base u = a[i+j], v = a[i+j + sz/2] * w;
             a[i+j] = u + v;
             a[i+i+sz/2] = u - v:
              w *= wlen;
      }
   if(inv) for(int i = 0; i < n; i++) a[i] /= 1.0 * n;
```

2.15 FFT Tourist

```
namespace fft {
 typedef double dbl;
 struct num {
   dbl x, y;
   num() \{ x = y = 0; \}
   num(dbl x, dbl y) : x(x), y(y) {}
 inline num operator+(num a, num b) { return num(a.x + b.x, a.y +
     \hookrightarrowb.v): }
 inline num operator-(num a, num b) { return num(a.x - b.x, a.y -
 inline num operator*(num a, num b) { return num(a.x * b.x - a.y
     \hookrightarrow * b.y, a.x * b.y + a.y * b.x); }
 inline num conj(num a) { return num(a.x, -a.y); }
 int base = 1;
 vector<num> roots = {{0, 0}, {1, 0}};
 vector<int> rev = {0, 1};
 const dbl PI = acosl(-1.0):
 void ensure_base(int nbase) {
   if(nbase <= base) return:</pre>
   rev.resize(1 << nbase);</pre>
   for(int i = 0; i < (1 << nbase); i++) {
    rev[i] = (rev[i >> 1] >> 1) + ((i & 1) << (nbase - 1));
   roots.resize(1 << nbase);</pre>
   while(base < nbase) {</pre>
     dbl \ angle = 2*PI / (1 << (base + 1));
     for(int i = 1 \ll (base - 1); i < (1 \ll base); i++) {
```

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```
roots[i << 1] = roots[i]:</pre>
     dbl angle i = angle * (2 * i + 1 - (1 << base)):
    roots[(i \ll 1) + 1] = num(cos(angle_i), sin(angle_i));
   base++:
 }
void fft(vector<num> &a, int n = -1) {
 if(n == -1) {
  n = a.size();
 }
 assert((n & (n-1)) == 0);
 int zeros = __builtin_ctz(n);
 ensure_base(zeros);
 int shift = base - zeros;
 for(int i = 0; i < n; i++) {
  if(i < (rev[i] >> shift)) {
    swap(a[i], a[rev[i] >> shift]);
   }
 for(int k = 1; k < n; k <<= 1) {
   for(int i = 0: i < n: i += 2 * k) {
     for(int j = 0; j < k; j++) {
      num z = a[i+j+k] * roots[j+k];
      a[i+j+k] = a[i+j] - z;
      a[i+j] = a[i+j] + z;
   }
vector<num> fa, fb;
vector<int> multiply(vector<int> &a, vector<int> &b) {
 int need = a.size() + b.size() - 1;
 int nbase = 0:
 while((1 << nbase) < need) nbase++;</pre>
 ensure_base(nbase);
 int sz = 1 << nbase;</pre>
 if(sz > (int) fa.size()) {
   fa.resize(sz);
 for(int i = 0; i < sz; i++) {
  int x = (i < (int) a.size() ? a[i] : 0);</pre>
   int y = (i < (int) b.size() ? b[i] : 0);</pre>
   fa[i] = num(x, y);
 fft(fa, sz);
 num r(0, -0.25 / sz);
 for(int i = 0; i \le (sz >> 1); i++) {
   int j = (sz - i) & (sz - 1);
   num z = (fa[j] * fa[j] - conj(fa[i] * fa[i])) * r;
   if(i != j) {
    fa[j] = (fa[i] * fa[i] - conj(fa[j] * fa[j])) * r;
   fa[i] = z;
 fft(fa, sz);
 vector<int> res(need);
 for(int i = 0; i < need; i++) {
  res[i] = fa[i].x + 0.5;
 return res;
vector<int> multiply_mod(vector<int> &a, vector<int> &b, int m,
   \hookrightarrowint eq = 0) {
 int need = a.size() + b.size() - 1;
```

```
int nbase = 0:
   while ((1 << nbase) < need) nbase++:</pre>
   ensure_base(nbase);
   int sz = 1 \ll nbase:
   if (sz > (int) fa.size()) {
     fa.resize(sz);
   for (int i = 0; i < (int) a.size(); i++) {</pre>
     int x = (a[i] \% m + m) \% m;
     fa[i] = num(x & ((1 << 15) - 1), x >> 15);
   fill(fa.begin() + a.size(), fa.begin() + sz, num {0, 0});
   fft(fa, sz);
   if (sz > (int) fb.size()) {
     fb.resize(sz);
   if (eq) {
     copy(fa.begin(), fa.begin() + sz, fb.begin());
   } else {
     for (int i = 0; i < (int) b.size(); i++) {</pre>
      int x = (b[i] \% m + m) \% m;
      fb[i] = num(x & ((1 << 15) - 1), x >> 15);
     fill(fb.begin() + b.size(), fb.begin() + sz, num {0, 0});
     fft(fb, sz);
   dbl ratio = 0.25 / sz;
   num r2(0, -1);
   num r3(ratio, 0);
   num r4(0, -ratio);
   num r5(0, 1);
   for (int i = 0; i \le (sz >> 1); i++) {
     int j = (sz - i) & (sz - 1);
     num a1 = (fa[i] + conj(fa[j]));
     num a2 = (fa[i] - conj(fa[j])) * r2;
     num b1 = (fb[i] + conj(fb[j])) * r3;
     num b2 = (fb[i] - conj(fb[j])) * r4;
     if (i != j) {
      num c1 = (fa[j] + conj(fa[i]));
      num c2 = (fa[j] - conj(fa[i])) * r2;
      num d1 = (fb[i] + conj(fb[i])) * r3;
      num d2 = (fb[j] - conj(fb[i])) * r4;
      fa[i] = c1 * d1 + c2 * d2 * r5;
       fb[i] = c1 * d2 + c2 * d1;
     fa[i] = a1 * b1 + a2 * b2 * r5;
     fb[j] = a1 * b2 + a2 * b1;
   fft(fa, sz);
   fft(fb, sz);
   vector<int> res(need);
   for (int i = 0; i < need; i++) {
     long long aa = fa[i].x + 0.5;
     long long bb = fb[i].x + 0.5;
     long long cc = fa[i].y + 0.5;
     res[i] = (aa + ((bb \% m) << 15) + ((cc \% m) << 30)) \% m;
   return res;
 vector<int> square_mod(vector<int> &a, int m) {
   return multiply_mod(a, a, m, 1);
2.16 NTT
```

const int mod = 7340033;

```
const int root = 5:
const int root 1 = 4404020:
const int root_pw = 1<<20;</pre>
void fft (vector<int> & a, bool invert) {
p int n = (int) a.size();
\triangleright for (int i=1, j=0; i<n; ++i) {
\triangleright int bit = n >> 1;
▷ for (; j>=bit; bit>>=1)
▷ ▷ ▷ j -= bit;
▷ ▷ j += bit;
▷ ▷ if (i < j)</pre>

▷ ▷ ▷ swap (a[i], a[j]);
⊳ }

    for (int len=2; len<=n; len<<=1) {</pre>
▷ int wlen = invert ? root_1 : root;
▷ for (int i=len; i<root_pw; i<<=1)</pre>
▷ ▷ ▷ wlen = int (wlen * 111 * wlen % mod);
\triangleright for (int i=0; i<n; i+=len) {
\triangleright \triangleright \triangleright int w = 1;
\triangleright \triangleright \triangleright for (int j=0; j<len/2; ++j) {
\triangleright \triangleright \triangleright \triangleright int u = a[i+j], v = int (a[i+j+len/2] * 111 * w % mod);
\triangleright \triangleright \triangleright \triangleright a[i+j] = u+v < mod ? u+v : u+v-mod;
\triangleright \triangleright \triangleright a[i+j+len/2] = u-v >= 0 ? u-v : u-v+mod;
\triangleright \triangleright \triangleright \triangleright w = int (w * 111 * wlen % mod);
▷ ▷ ▷ }
⊳ ⊳ }
⊳ }
▶ if (invert) {
▷ int nrev = reverse (n, mod);
\triangleright for (int i=0; i<n; ++i)
\triangleright \triangleright \vdash a[i] = int (a[i] * 111 * nrev % mod);
⊳ }
}
```

2.17 Gauss

```
// Solves systems of linear equations.
// To use, build a matrix of coefficients and call run(mat, R, C).
   \hookrightarrow If the i-th variable is free, row[i] will be -1, otherwise
   \hookrightarrowit's value will be ans[i].
namespace Gauss {
 const int MAXC = 1001;
 int row[MAXC]:
 double ans[MAXC];
 void run(double mat[][MAXC], int R, int C) {
   REP(i, C) row[i] = -1;
   int r = 0:
   REP(c, C) {
     int k = r:
     FOR(i, r, R) if(fabs(mat[i][c]) > fabs(mat[k][c])) k = i;
     if(fabs(mat[k][c]) < eps) continue;</pre>
     REP(j, C+1) swap(mat[r][j], mat[k][j]);
     REP(i, R) if (i != r) {
       double w = mat[i][c] / mat[r][c];
       REP(j, C+1) mat[i][j] -= mat[r][j] * w;
    row[c] = r++;
   REP(i, C) {
     int r = row[i];
```

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```
ans[i] = r == -1 ? 0 : mat[r][C] / mat[r][i];
}
}
```

2.18 Gauss Xor

```
const 11 MAX = 1e9:
const int LOG_MAX = 64 - __builtin_clzll((11)MAX);
struct Gauss {
  array<11, LOG_MAX> vet;
  int size;
  Gauss() : size(0) {
Gauss(vector<11> vals) : size(0) {
for(ll val : vals) add(val);
  bool add(11 val) {
      for(int i = 0; i < LOG_MAX; i++) if(val & (1LL << i)) {
        if(vet[i] == 0) {
           vet[i] = val;
           size++;
           return true:
        val ^= vetΓi]:
     return false;
```

2.19 Simpson

```
inline double simpson(double fl,double fr,double fmid,double
    \rightarrow1,double r) {

  return (fl + fr + 4.0 * fmid) * (r - 1) / 6.0;
double rsimpson(double slr, double fl, double fr, double fmid, double
    →1.double r) {
\triangleright double mid = (1+r)*0.5;

    double slm = simpson(fl, fmid, fml, l, mid);

    double smr = simpson(fmid, fr, fmr, mid, r);

    if(fabs(slr-slm-smr) < eps and r - l < delta) return slr;
</pre>

    return rsimpson(slm.fl.fmid.fml.l.mid) +
     →rsimpson(smr,fmid,fr,fmr,mid,r);
double integrate(double 1,double r) {
\triangleright double mid = (1+r)*0.5;
\triangleright double fl = f(1), fr = f(r), fmid = f(mid);

    return rsimpson(simpson(fl,fr,fmid,l,r),fl,fr,fmid,l,r);
```

2.20 Matrix

```
template <const size_t n, const size_t m, class T = modBase<>>
struct Matrix {
   T v[n][m];

Matrix(int d = 0) {
   for (int i = 0; i < n; i++) {
      for (int j = 0; j < m; j++) {
      v[i][j] = T(0);
    }
   if (i < m) {
      v[i][i] = T(d);
   }
}</pre>
```

```
template <size_t mm>
Matrix<n, mm, T> operator*(Matrix<m, mm, T> &o) {
    Matrix<n, mm, T> ans;
    for (int i = 0; i < n; i++) {
        for (int j = 0; j < mm; j++) {
            for (int k = 0; k < m; k++) {
                 ans.v[i][j] = ans.v[i][j] + v[i][k] * o.v[k][j];
            }
        }
    }
    return ans;
}</pre>
```

3 Graphs

3.1 Bipartite Matching

```
// O(V * E)
int match[N];
int vis[N], pass;
vector<int> g[N];
bool dfs(int u) {
vis[u] = pass;
p for(int v : g[u]) if(vis[v] != pass) {
▷ ▷ vis[v] = pass;
\triangleright \triangleright \triangleright match[v] = u;
\triangleright \triangleright \triangleright match[u] = v;
▷ ▷ ▷ return true;
⊳ ⊳ }
⊳ }

    return false;

int max_maching() {

    memset(match, -1, sizeof match);

    int max_matching_size = 0;

    for(int u : vertices_on_side_A) {
▷ if(dfs(i)) max_matching_size++;
⊳ }

    return max_matching_size;
}
```

3.2 Dinic

```
const int N = 100005;
const int E = 2000006;
vector<int> g[N];
int ne;
struct Edge{
    int from, to; ll flow, cap;
} edge[E];
int lv1[N], vis[N], pass, start = N-2, target = N-1;
int qu[N], qt, px[N];

ll run(int s, int sink, ll minE){
    if(s == sink) return minE;

    ll ans = 0;

    for(; px[s] < (int)g[s].size(); px[s]++){
        int e = g[s][ px[s] ];
        auto &v = edge[e], &rev = edge[e^1];</pre>
```

```
if(|v| | v.to] != |v| | s| +1 | v.flow >= v.cap)
           continue: // v.cap - v.flow < lim
       11 tmp = run(v.to, sink,min(minE, v.cap-v.flow));
       v.flow += tmp, rev.flow -= tmp;
       ans += tmp, minE -= tmp;
       if(minE == 0) break;
   return ans;
bool bfs(int source, int sink){
   qt = 0;
   qu[qt++] = source;
   lvl[source] = 1;
   vis[source] = ++pass;
   for(int i = 0; i < qt; i++){
       int u = qu[i];
       px[u] = 0;
▷ ▷ if(u == sink) return true;
       for(auto& ed : g[u]) {
           auto v = edge[ed];
          if(v.flow >= v.cap || vis[v.to] == pass)
              continue; // v.cap - v.flow < lim</pre>
           vis[v.to] = pass;
          lvl[v.to] = lvl[u]+1;
           qu[qt++] = v.to;
      }
   return false;
11 flow(int source = start, int sink = target){
   11 \text{ ans} = 0:
   //for(lim = (1LL << 62); lim >= 1; lim /= 2)
   while(bfs(source, sink))
▷ ▷ ans += run(source, sink, oo);
   return ans;
void addEdge(int u, int v, ll c = 1, ll rc = 0){
   edge[ne] = {u, v, 0, c};
   g[u].push_back(ne++);
   edge[ne] = {v, u, 0, rc};
   g[v].push_back(ne++);
void reset_flow(){
\triangleright for(int i = 0: i < ne: i++)
▷ ▷ edge[i].flow = 0;
```

3.3 Push relabel

```
// Push relabel in O(V^2 E^0.5) with gap heuristic
// It's quite fast
template<typename flow_t = long long>
struct PushRelabel {
   struct Edge { int to, rev; flow_t f, c; };
   vector<vector<Edge> > g;
   vector<flow_t> ec;
   vector<Edge*> cur;
   vector<vector<int> > hs;
   vector<int> H;
   PushRelabel(int n) : g(n), ec(n), cur(n), hs(2*n), H(n) {}
   void add_edge(int s, int t, flow_t cap, flow_t rcap=0) {
      if (s == t) return;
       Edge a = \{t, (int)g[t].size(), 0, cap\};
       Edge b = {s, (int)g[s].size(), 0, rcap};
       g[s].push_back(a);
       g[t].push_back(b);
   void add_flow(Edge& e, flow_t f) {
```

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```
Edge &back = g[e.to][e.rev];
       if (!ec[e.to] && f)
          hs[H[e.to]].push_back(e.to);
       e.f += f, ec[e.to] += f;
       back.f -= f, ec[back.to] -= f;
   flow_t max_flow(int s, int t) {
       int v = g.size();
       H[s] = v; ec[t] = 1;
       vector<int> co(2 * v);
       co[0] = v-1;
       for(int i = 0; i < v; ++i) cur[i] = g[i].data();</pre>
       for(auto &e : g[s]) add_flow(e, e.c);
       if(hs[0].size())
       for (int hi = 0; hi >= 0;) {
          int u = hs[hi].back();
          hs[hi].pop_back();
          while (ec[u] > 0) // discharge u
              if (cur[u] == g[u].data() + g[u].size()) {
                  H[u] = 1e9;
                  for(auto &e:g[u])
                     if (e.c - e.f && H[u] > H[e.to]+1)
                         H[u] = H[e.to]+1, cur[u] = &e;
                  if (++co[H[u]], !--co[hi] && hi < v)</pre>
                     for(int i = 0; i < v; ++i)
                         if (hi < H[i] && H[i] < v){</pre>
                             --co[H[i]];
                             H[i] = v + 1;
                  hi = H[u];
              } else if (cur[u]->c - cur[u]->f && H[u] ==
                  \hookrightarrowH[cur[u]->to]+1)
                  add_flow(*cur[u], min(ec[u], cur[u]->c -
                     \hookrightarrow cur[u]->f));
              else ++cur[u];
           while (hi >= 0 && hs[hi].empty()) --hi;
       return -ec[s];
};
```

3.4 Min Cost Max Flow

```
const 11 oo = 1e18:
const int N = 422, E = 2 * 10006;
vector<int> g[N];
int ne;
struct Edge{
    int from, to; ll cap, cost;
int start = N-1, target = N-2, p[N]; int inqueue[N];
11 d[N];
11 pot[N]:
bool dijkstra(int source, int sink) {
p for(int i = 0; i < N; i++) d[i] = oo;</pre>
▷ d[source] = 0;
priority_queue<pair<ll, int>> q;

    q.emplace(0, source);
⊳ 11 dt; int u;
▶ while(!q.empty()) {
\triangleright tie(dt, u) = q.top(); q.pop(); dt = -dt;
▷ if(dt > d[u]) continue;
▷ if(u == sink) return true;
▷ b for(int e : g[u]) {
\triangleright \triangleright \triangleright auto v = edge[e];
\triangleright \triangleright \triangleright const 11 cand = d[u] + v.cost + pot[u] - pot[v.to];
```

```
\triangleright \triangleright \vdash \mathbf{if}(v.cap > 0 \text{ and } cand < d[v.to]) 
▷ ▷ ▷ ▷ p[v.to] = e;
\triangleright \triangleright \triangleright \triangleright d[v.to] = cand;
▷ ▷ ▷ □ q.emplace(-d[v.to], v.to);
⊳ ⊳ ⊳ }
⊳ ⊳ }
⊳ }

  return d[sink] < oo;</pre>
// <max flow, min cost>
pair<11, 11> mincost(int source = start, int sink = target){
    11 ans = 0, mf = 0;
    while(dijkstra(source, sink)){
        11 f = oo:
        for(int u = sink; u != source; u = edge[ p[u] ].from)
            f = min(f, edge[ p[u] ].cap);
        mf += f;
        ans += f * (d[sink] - pot[source] + pot[sink]);
        for(int u = sink; u != source; u = edge[ p[u] ].from){
            edge[ p[u] ].cap -= f;
            edge[ p[u] ^ 1 ].cap += f;
\triangleright for(int i = 0; i < N; i++) pot[i] = min(oo, pot[i] + d[i]);
    return {mf, ans};
void addEdge(int u, int v, ll c, ll cost){
▷ assert(cost >= 0);
    edge[ne] = {u, v, c, cost};
    g[u].push_back(ne++);
    edge[ne] = \{v, u, 0, -cost\};
    g[v].push_back(ne++);
```

3.5 Blossom Algorithm for General Matching

```
const int MAXN = 2020 + 1;
// 1-based Vertex index
int vis[MAXN], par[MAXN], orig[MAXN], match[MAXN], aux[MAXN], t, N;
vector<int> conn[MAXN];
queue<int> Q;
void addEdge(int u, int v) {
p conn[u].push_back(v); conn[v].push_back(u);
void init(int n) {
\triangleright N = n; t = 0;

    for(int i=0; i<=n; ++i)</pre>
b conn[i].clear(), match[i] = aux[i] = par[i] = 0;
void augment(int u, int v) {
\triangleright int pv = v, nv;
⊳ do {
▷ pv = par[v]; nv = match[pv];

▷ match[v] = pv; match[pv] = v;
\triangleright \triangleright v = nv;
int lca(int v, int w) {
→ ++t;
▶ while(true) {
▷ ▷ if(v) {
▷ ▷ if(aux[v] == t) return v; aux[v] = t;
▷ ▷ ▷ v = orig[par[match[v]]];
⊳ ⊳ }
\triangleright \triangleright swap(v, w);
⊳ }
```

```
void blossom(int v, int w, int a) {

  while(orig[v] != a) {
\triangleright \triangleright par[v] = w; w = match[v];
\triangleright if(vis[w] == 1) Q.push(w), vis[w] = 0;
▷ orig[v] = orig[w] = a; v = par[w];
⊳ }
bool bfs(int u) {
\triangleright fill(vis+1, vis+1+N, -1); iota(orig + 1, orig + N + 1, 1);

    Q = queue<int>(); Q.push(u); vis[u] = 0;
▶ while(!Q.empty()) {
▷ int v = Q.front(); Q.pop();
▷ b for(int x: conn[v]) {
\triangleright \triangleright \triangleright \triangleright par[x] = v; vis[x] = 1;
▷ ▷ ▷ if(!match[x]) return augment(u, x), true;
▷ ▷ ▷ Q.push(match[x]); vis[match[x]] = 0;
▷ ▷ ▷ }

▷ ▷ ▷ else if(vis[x] == 0 && orig[v] != orig[x]) {
\triangleright \triangleright \triangleright \triangleright int a = lca(orig[v], orig[x]);
▷ ▷ ▷ blossom(x, v, a); blossom(v, x, a);
▷ ▷ ▷ }
⊳ ⊳ }
⊳ }
⊳ return false:
int Match() {
\triangleright int ans = 0;
▷ // find random matching (not necessary, constant improvement)

  vector<int> V(N-1); iota(V.begin(), V.end(), 1);
shuffle(V.begin(), V.end(), mt19937(0x94949));
p for(auto x: V) if(!match[x]){
b for(auto y: conn[x]) if(!match[y]) {
\triangleright \triangleright match[x] = y, match[y] = x;
▷ ▷ ▷ ++ans; break;
⊳ ⊳ }

    for(int i=1; i<=N; ++i) if(!match[i] && bfs(i)) ++ans;
</pre>

  return ans;
```

3.6 Blossom Algorithm for Weighted General Match-

ing

```
// N^3 (but fast in practice)
static const int INF = INT_MAX;
static const int N = 514;
struct edge{
p int u,v,w; edge(){}

    edge(int ui,int vi,int wi)

▷ ▷ :u(ui),v(vi),w(wi){}
};
int n,n_x;
edge g[N*2][N*2];
int lab[N*2];
int match[N*2],slack[N*2],st[N*2],pa[N*2];
int flo_from[N*2][N+1],S[N*2],vis[N*2];
vector<int> flo[N*2];
queue<int> q;
int e_delta(const edge &e){

  return lab[e.u]+lab[e.v]-g[e.u][e.v].w*2;

void update_slack(int u,int x){
   \hookrightarrow if(!slack[x]||e_delta(g[u][x])<e_delta(g[slack[x]][x]))slack[x]=u;
void set_slack(int x){
```

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```
    slack[x]=0:

    for(int u=1;u<=n;++u)
</pre>
▷ ▷ if(g[u][x].w>0&&st[u]!=x&&S[st[u]]==0)
▷ ▷ update_slack(u,x);
void q_push(int x){

    if(x<=n)q.push(x);</pre>
b else for(size_t i=0;i<flo[x].size();i++)</pre>
▷ □ q_push(flo[x][i]);
}
void set_st(int x,int b){
b if(x>n)for(size_t i=0;i<flo[x].size();++i)</pre>
▷ ▷ set_st(flo[x][i],b);
int get_pr(int b,int xr){
b int pr=find(flo[b].begin(),flo[b].end(),xr)-flo[b].begin();

    if(pr%2==1){
p reverse(flo[b].begin()+1,flo[b].end());
▷ return (int)flo[b].size()-pr;
⊳ }else return pr;
void set_match(int u,int v){

  match[u]=q[u][v].v;

    if(u<=n) return;</pre>
⊳ edge e=g[u][v];
b int xr=flo_from[u][e.u],pr=get_pr(u,xr);
p for(int i=0;i<pr;++i)set_match(flo[u][i],flo[u][i^1]);</pre>
⊳ set_match(xr,v);
p rotate(flo[u].begin(),flo[u].begin()+pr,flo[u].end());
void augment(int u,int v){
p for(;;){
▷ int xnv=st[match[u]];
▷ ▷ set_match(u,v);
▷ ▷ if(!xnv)return;
▷ ▷ set_match(xnv,st[pa[xnv]]);
▷ ▷ u=st[pa[xnv]],v=xnv;
⊳ }
int get_lca(int u,int v){
⊳ static int t=0;
\triangleright for(++t;u||v;swap(u,v)){
▷ ▷ if(u==0)continue:
▷ ▷ if(vis[u]==t)return u:
▷ ▷ vis[u]=t;
▷ ▷ u=st[match[u]];
▷ b if(u)u=st[pa[u]];
⊳ }

  return 0;
void add_blossom(int u,int lca,int v){
▶ int b=n+1:

    while(b<=n_x&&st[b])++b;</pre>
\triangleright if(b>n_x)++n_x;
□ lab[b]=0,S[b]=0;

    match[b]=match[lca];
b flo[b].clear();
p flo[b].push_back(lca);
p for(int x=u,y;x!=lca;x=st[pa[y]])
p reverse(flo[b].begin()+1,flo[b].end());
p for(int x=v,y;x!=lca;x=st[pa[y]])
b flo[b].push_back(x),flo[b].push_back(y=st[match[x]]),q_push(y);
⊳ set_st(b,b);
p for(int x=1;x<=n_x;++x)g[b][x].w=g[x][b].w=0;</pre>
\triangleright for(int x=1;x<=n;++x)flo_from[b][x]=0;
p for(size_t i=0;i<flo[b].size();++i){</pre>
```

```
▷ ▷ int xs=flo[b][i]:
\triangleright \triangleright \mathbf{if}(g[b][x].w==0||e_delta(g[xs][x]) < e_delta(g[b][x]))
\triangleright \triangleright \triangleright \triangleright g[b][x]=g[xs][x],g[x][b]=g[x][xs];
▷ ▷ if(flo_from[xs][x])flo_from[b][x]=xs;
⊳ }
⊳ set_slack(b);
void expand_blossom(int b){
p for(size_t i=0;i<flo[b].size();++i)</pre>
> set_st(flo[b][i],flo[b][i]);
p int xr=flo_from[b][g[b][pa[b]].u],pr=get_pr(b,xr);

    for(int i=0;i<pr;i+=2){</pre>
▷ int xs=flo[b][i],xns=flo[b][i+1];
▷ ▷ pa[xs]=g[xns][xs].u;
▷ ▷ S[xs]=1,S[xns]=0;
▷ ▷ slack[xs]=0,set_slack(xns);
▷ □ q_push(xns);
⊳ }
> S[xr]=1,pa[xr]=pa[b];
p for(size_t i=pr+1;i<flo[b].size();++i){</pre>
▷ ▷ int xs=flo[b][i];
▷ ▷ S[xs]=-1,set_slack(xs);
⊳ }
⊳ st[b]=0;
bool on_found_edge(const edge &e){

    int u=st[e.u],v=st[e.v];

    if(S[v]==-1){
▷ ▷ pa[v]=e.u,S[v]=1;
▷ ▷ int nu=st[match[v]];
▷ ▷ slack[v]=slack[nu]=0;
▷ ▷ S[nu]=0,q_push(nu);
> }else if(S[v]==0){
▷ ▷ int lca=get_lca(u,v);
▷ ▷ else add_blossom(u,lca,v);
⊳ }

    return false;
bool matching(){

    memset(S+1,-1,sizeof(int)*n_x);
p memset(slack+1,0,sizeof(int)*n_x);

p q=queue<int>();

    for(int x=1; x<=n_x; ++x)</pre>
\Rightarrow if(st[x]==x&&!match[x])pa[x]=0,S[x]=0,q_push(x);

    if(q.empty())return false;

    for(;;){
▷ ▷ while(q.size()){
▷ ▷ int u=q.front();q.pop();
▷ ▷ if(S[st[u]]==1)continue;
▷ ▷ ▷ if(g[u][v].w>0&&st[u]!=st[v]){
\triangleright \triangleright \triangleright \triangleright \vdash \mathbf{if}(e_{delta}(g[u][v])==0){
▷ ▷ ▷ ▷ ▷ if(on_found_edge(g[u][v]))return true;
▷ ▷ ▷ ▷ ▷ }else update_slack(u,st[v]);
▷ ▷ ▷ ▷ }
⊳ ⊳ }
▷ ▷ int d=INF;
▷ ▷ if(st[b]==b&&S[b]==1)d=min(d,lab[b]/2);
▷ ▷ if(st[x]==x&&slack[x]){
▷ ▷ ▷ if(S[x]==-1)d=min(d,e_delta(g[slack[x]][x]));

▷ ▷ ▷ ▷ else if(S[x]==0)d=min(d,e_delta(g[slack[x]][x])/2);
▷ ▷ ▷ }
▷ for(int u=1;u<=n;++u){</pre>
```

```
▷ ▷ ▷ if(S[st[u]]==0){
▷ ▷ ▷ if(lab[u]<=d)return 0;
▷ ▷ ▷ ▷ lab[u]-=d;
▷ ▷ } else if(S[st[u]]==1)lab[u]+=d;
⊳ ⊳ }
▷ ▷ if(st[b]==b){
▷ ▷ ▷ if(S[st[b]]==0)lab[b]+=d*2;
▷ ▷ ▷ else if(S[st[b]]==1)lab[b]-=d*2;
⊳ ⊳ ⊳ }
\triangleright for(int x=1:x<=n x:++x)
\triangleright \triangleright \triangleright
    \hookrightarrow if(st[x]==x&&slack[x]&&st[slack[x]]!=x&&e_delta(g[slack[x]][x])==0
▷ ▷ ▷ if(on_found_edge(g[slack[x]][x]))return true;
\triangleright for(int b=n+1;b<=n_x;++b)

▷ ▷ if(st[b]==b&&S[b]==1&&lab[b]==0)expand_blossom(b);

⊳ }

    return false;

pair<long long,int> solve(){
p memset(match+1,0,sizeof(int)*n);
⊳ n x=n:
p int n_matches=0;

    long long tot_weight=0;
p for(int u=0;u<=n;++u)st[u]=u,flo[u].clear();</pre>

  int w_max=0;

    for(int u=1;u<=n;++u)</pre>
▷ for(int v=1; v<=n; ++v){</pre>
\triangleright \triangleright \vdash flo\_from[u][v]=(u==v?u:0);
\triangleright \triangleright \lor w_{max=max}(w_{max},g[u][v].w);
⊳ ⊳ }

    for(int u=1;u<=n;++u)lab[u]=w_max;</pre>
b while(matching())++n_matches;
p for(int u=1;u<=n;++u)</pre>
▷ if(match[u]&&match[u]<u)</pre>
▷ ▷ tot_weight+=g[u][match[u]].w;
preturn make_pair(tot_weight,n_matches);
void add_edge( int ui , int vi , int wi ){

    g[ui][vi].w = g[vi][ui].w = wi;

void init( int _n ){
p n = n:

    for(int u=1;u<=n;++u)</pre>
\triangleright \triangleright \triangleright g[u][v] = edge(u,v,0);
3.7 Small to Large
void cnt_sz(int u, int p = -1){
    sz[u] = 1;
    for(int v : g[u]) if(v != p)
        cnt_sz(v, u), sz[u] += sz[v];
void add(int u, int p, int big = -1){
     // Update info about this vx in global answer
    for(int v : g[u]) if(v != p && v != big)
        add(v, u);
void dfs(int u, int p, int keep){
    int big = -1, mmx = -1;
    for(int v : g[u]) if(v != p \&\& sz[v] > mmx)
        mmx = sz[v], big = v;
    for(int v : g[u]) if(v != p && v != big)
        dfs(v, u, 0);
    if(big != -1) dfs(big, u, 1);
```

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```
add(u, p, big);
for(auto x : q[u]){
   // answer all queries for this vx
if(!keep){ /*Remove data from this subtree*/ }
```

3.8 Kosaraju

```
vector<int> g[N], gt[N], S; int vis[N], cor[N];
void dfs(int u){
vis[u] = 1; for(int v : g[u]) if(!vis[v]) dfs(v);
▷ S.push_back(u);
void dfst(int u. int e){
⊳ cor[u] = e;

    for(int v : gt[u]) if(!cor[v]) dfst(v, e);
void kosaraju(){
p for(int i = 1; i <= n; i++) if(!vis[i]) dfs(i);</pre>
   for(int i = 1; i \le n; i++) for(int j : g[i])
       gt[j].push_back(i);
p int e = 0; reverse(S.begin(), S.end());
p for(int u : S) if(!cor[u]) dfst(u, ++e);
```

3.9 Tarjan

```
int cnt = 0, root;
void dfs(int u, int p = -1){
\triangleright low[u] = num[u] = ++t:

    for(int v : g[u]){
▷ ▷ if(!num[v]){
\triangleright \triangleright \triangleright dfs(v, u);
⊳⊳ if(u == root) cnt++;
▷ ▷ ▷ if(low[v] >= num[u]) u PONTO DE ARTICULAÇÃO;
▷ ▷ ▷ if(low[v] > num[u]) ARESTA u->v PONTE;
\triangleright \triangleright \log[u] = \min(\log[u], \log[v]);
⊳ ⊳ }
▷ ▷ else if(v != p) low[u] = min(low[u], num[v]);
⊳ }
root PONTO DE ARTICULAÇÃO <=> cnt > 1
void tarjanSCC(int u){
⊳ low[u] = num[u] = ++cnt:
▷ vis[u] = 1;
p for(int v : q[u]){
▷ if(!num[v]) tarjanSCC(v);
▷ ▷ if(vis[v]) low[u] = min(low[u], low[v]);
⊳ }
\triangleright if(low[u] == num[u]){
▷ ▷ ssc[u] = ++ssc cnt: int v:
\triangleright \triangleright \lor v = S.back(); S.pop_back(); vis[v] = 0;
▷ ▷ ▷ ssc[v] = ssc_cnt;
\triangleright \triangleright }while(u != v);
⊳ }
```

3.10 Max Clique

```
long long adj[N], dp[N];
for(int i = 0; i < n; i++){
\triangleright \mathbf{for}(\mathbf{int} \ j = 0; \ j < n; \ j++) \{
⊳ ⊳ int x;
▷ ▷ scanf("%d",&x);
```

```
\triangleright \triangleright \mathbf{if}(x \mid | i == i)
▷ ▷ ▷ adj[i] |= 1LL << j;</pre>
⊳ }
}
int resto = n - n/2;
int C = n/2;
for(int i = 1; i < (1 << resto); i++){</pre>
▷ int x = i:
\triangleright for(int j = 0; j < resto; j++)
▷ ▷ if(i & (1 << j))</pre>
\triangleright \triangleright x \&= adj[j + C] >> C;
\triangleright \mathbf{if}(\mathbf{x} == \mathbf{i})
▷ b dp[i] = __builtin_popcount(i);
⊳ }
}
for(int i = 1; i < (1 << resto); i++)</pre>
\triangleright for(int j = 0; j < resto; j++)
\triangleright \ \mathbf{if}(i \& (1 << j))
\triangleright \triangleright dp[i] = max(dp[i], dp[i ^ (1 << j)]);
int maxClig = 0:
for(int i = 0; i < (1 << C); i++){
\triangleright int x = i, y = (1 << resto) - 1;
\triangleright for(int j = 0; j < C; j++)
▷ ▷ if(i & (1 << j))</pre>
\triangleright \triangleright x \& = adj[j] \& ((1 << C) - 1), y \& = adj[j] >> C;

    if(x != i) continue;
b maxCliq = max(maxCliq, __builtin_popcount(i) + dp[y]);
```

3.11 Dominator Tree

```
vector<int> q[N], qt[N], T[N];
vector<int> S;
int dsu[N], label[N];
int sdom[N], idom[N], dfs_time, id[N];
vector<int> bucket[N];
vector<int> down[N];
void prep(int u){

    S.push_back(u);

  id[u] = ++dfs_time;
\triangleright label[u] = sdom[u] = dsu[u] = u;

    for(int v : g[u]){
▷ ▷ if(!id[v])
▷ ▷ prep(v), down[u].push_back(v);
▷ ▷ gt[v].push_back(u);
⊳ }
}
int fnd(int u, int flag = 0){

    if(u == dsu[u]) return u;
p int v = fnd(dsu[u], 1), b = label[ dsu[u] ];
b if(id[ sdom[b] ] < id[ sdom[ label[u] ] ])</pre>
\triangleright \triangleright label[u] = b;
\triangleright dsu[u] = v;

  return flag ? v : label[u];
void build_dominator_tree(int root, int sz){
▷ // memset(id, 0, sizeof(int) * (sz + 1));
▷ // for(int i = 0; i <= sz; i++) T[i].clear();</pre>

   prep(root);

    reverse(S.begin(), S.end());
```

```
⊳ int w:

    for(int u : S){
▷ ▷ for(int v : gt[u]){
\triangleright \triangleright \triangleright w = fnd(v);
\triangleright \triangleright \mathsf{if}(\mathsf{id}[\mathsf{sdom}[\mathsf{w}]] < \mathsf{id}[\mathsf{sdom}[\mathsf{u}]])
\triangleright \triangleright \triangleright \triangleright sdom[u] = sdom[w];
⊳ ⊳ }
▷ ▷ gt[u].clear();
▷ b if(u != root) bucket[ sdom[u] ].push_back(u);
▷ b for(int v : bucket[u]){
\triangleright \triangleright \triangleright w = fnd(v);
▷ ▷ if(sdom[w] == sdom[v]) idom[v] = sdom[v];
▷ ▷ ▷ else idom[v] = w;
⊳ ⊳ }
▷ bucket[u].clear();
▷ for(int v : down[u]) dsu[v] = u;
▷ ▷ down[u].clear();
⊳ }

    reverse(S.begin(), S.end());

    for(int u : S) if(u != root){
▷ b if(idom[u] != sdom[u]) idom[u] = idom[ idom[u] ];
▷ ▷ T[ idom[u] ].push_back(u);
⊳ }
▷ S.clear();
}
```

3.12 Min Cost Matching

```
// Min cost matching
// O(n^2 * m)
// n == nro de linhas
// m == nro de colunas
// n <= m | flow == n
// a[i][j] = custo pra conectar i a j
vector<int> u(n + 1), v(m + 1), p(m + 1), way(m + 1);
for(int i = 1; i \le n; ++i){
   p[0] = i;
   int j0 = 0;
   vector<int> minv(m + 1 . oo):
   vector<char> used(m + 1 , false);
      used[j0] = true;
       int i0 = p[j0] , delta = oo, j1;
       for(int j = 1; j \ll m; ++j)
          if(! used[j]){
             int cur = a[i0][j] - u[i0] - v[j];
             if(cur < minv[j])</pre>
                 minv[j] = cur, way[j] = j0;
             if(minv[j] < delta)</pre>
                 delta = minv[j] , j1 = j;
       for(int j = 0; j \le m; ++j)
          if(used[j])
             u[p[j]] += delta, v[j] -= delta;
          else
             minv[j] -= delta;
       j0 = j1;
   }while(p[j0] != 0);
      int j1 = way[j0];
      p[j0] = p[j1];
       j0 = j1;
```

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```
}while(j0);
}

// match[i] = coluna escolhida para linha i
vector<int> match(n + 1);
for(int j = 1; j <= m; ++j)
    match[p[j]] = j;

int cost = -v[0];</pre>
```

4 Strings

4.1 Aho Corasick

```
int to[N][A];
int ne = 2, fail[N], term[N];
void add_string(const char *str, int id) {
   int p = 1;
   for(int i = 0; str[i]; i++) {
       int ch = str[i] - 'a';
      if(!to[p][ch]) to[p][ch] = ne++;
      p = to[p][ch];
   term[p]++;
void init() {
   for(int i = 0; i < ne; i++) fail[i] = 1;
   queue<int> q; q.push(1);
   while(!q.empty()){
       int u = q.front(); q.pop();
       for(int i = 0; i < A; i++){
          if(to[u][i]) {
             int v = to[u][i]; q.push(v);
             if(u != 1) {
                 fail[v] = to[ fail[u] ][i];
                 term[v] += term[ fail[v] ];
             }
          else if(u != 1) to[u][i] = to[ fail[u] ][i];
          else to[u][i] = 1;
   }
void clean() {
   memset(to, 0, ne * sizeof(to[0]));
   memset(fail, 0, ne * sizeof(fail[0]));
   memset(term, 0, ne * sizeof(term[0]));
   ne = 2:
```

```
for(int i = 0: i < n: i++)
       h[c[i]]--:
    for(int L = 1; L < n; L <<= 1) {
       for(int i = 0; i < n; i++) {
           int j = (a[i] - L + n) \% n;
           a1[h[c[j]]++] = j;
       int cc = -1;
       for(int i = 0; i < n; i++) {</pre>
           if(i == 0 || c[a1[i]] != c[a1[i-1]] || c[(a1[i] + L) %
               \rightarrown] != c[(a1[i-1] + L) % n])
              h[++cc] = i;
           c1[a1[i]] = cc;
       memcpy(a, a1, sizeof a1);
       memcpy(c, c1, sizeof c1);
       if(cc == n-1) break;
   }
void build_lcp(char s[], int n, int a[]){ // lcp[i] =
    \hookrightarrow lcp(s[:a[i]], s[:a[i+1]])
   int k = 0;
    //memset(lcp, 0, sizeof lcp);
    for(int i = 0; i < n; i++){
       if(c[i] == n-1) continue;
       int j = a[c[i]+1];
       while(i+k < n \&\& j+k < n \&\& s[i+k] == s[j+k]) k++;
       lcp[c[i]] = k;
       if(k) k--;
   }
}
int comp_lcp(int i, int j){
   if(i == j) return n - i;
   if(c[i] > c[j]) swap(i, j);
   return min(lcp[k] for k in [c[i], c[j]-1]);
```

4.3 Adamant Suffix Tree

```
namespace sf {
const int inf = 1e9;
const int maxn = 200005;
char s[maxn];
map<int, int> to[maxn];
int len[maxn], fpos[maxn], link[maxn];
int node, pos;
int sz = 1. n = 0:
int make_node(int _pos, int _len) {
 fpos[sz] = _pos;
 len[sz] = _len;
 return sz++;
void go_edge() {
 while (pos > len[to[node][s[n - pos]]]) {
   node = to[node][s[n - pos]];
   pos -= len[node];
void add_letter(int c) {
```

```
s[n++] = (char)c:
 pos++:
 int last = 0;
  while (pos > 0) {
   go_edge();
   int edge = s[n - pos];
   int &v = to[node][edge];
   int t = s[fpos[v] + pos - 1];
   if (v == 0) {
    v = make_node(n - pos, inf);
     link[last] = node;
     last = 0:
   } else if (t == c) {
     link[last] = node;
     return:
   } else {
     int u = make_node(fpos[v], pos - 1);
     to[u][c] = make\_node(n - 1, inf);
     to[u][t] = v;
     fpos[v] += pos - 1;
     len[v] -= pos - 1;
     v = u:
     link[last] = u:
     last = u;
   if (node == 0)
    pos--;
   else
    node = link[node];
 }
void add_string(char *str) {
 for (int i = 0; str[i]; i++) add_letter(str[i]);
 add letter('$'):
bool is_leaf(int u) { return len[u] > n; }
int get_len(int u) {
 if (!u) return 0;
 if (is_leaf(u)) return n - fpos[u];
 return len[u];
int leafs[maxn];
int calc_leafs(int u = 0) {
 leafs[u] = is_leaf(u);
 for (const auto &c : to[u]) leafs[u] += calc_leafs(c.second);
 return leafs[u];
}; // namespace sf
int main() { sf::len[0] = sf::inf; }
```

4.4 Z Algorithm

```
vector<int> z_algo(const string &s) {
    int n = s.size(), L = 0, R = 0;
    vector<int> z(n, 0);
    for(int i = 1; i < n; i++){
        b if(i <= R) z[i] = min(z[i-L], R - i + 1);
        b while(z[i]+i < n && s[ z[i]+i ] == s[ z[i] ])
        b > b z[i]++;
        b if(i+z[i]-1 > R) L = i, R = i + z[i] - 1;
        }
    }
    return z;
}
```

4.5 Prefix function/KMP

```
vector<int> preffix_function(const string &s){
> int n = s.size(); vector<int> b(n+1);
```

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```
b [0] = -1; int i = 0, j = -1;
\triangleright while(i < n){
b \in while(j >= 0 \&\& s[i] != s[j]) j = b[j];
\triangleright \triangleright b[++i] = ++j;
⊳ }
⊳ return b;
void kmp(const string &t, const string &p){
vector<int> b = preffix_function(p);
p int n = t.size(), m = p.size();
b int i = 0:
\triangleright for(int i = 0: i < n: i++){
b while(j >= 0 && t[i] != p[j]) j = b[j];
\triangleright \triangleright if(j == m){
▷ ▷ ▷ //patern of p found on t
\triangleright \triangleright \triangleright j = b[j];
⊳ ⊳ }
⊳ }
```

4.6 Min rotation

```
// remember std::rotate
int min rotation(int *s. int N) {
 REP(i. N) s[N+i] = s[i]:
 int a = 0:
 REP(b, N) REP(i, N) {
   if (a+i == b \mid | s[a+i] < s[b+i]) { b += max(0, i-1); break; }
   if (s[a+i] > s[b+i]) { a = b; break; }
 return a;
```

4.7 Manacher

```
// rad[2 * i] = largest palindrome cetered at char i
// rad[2 * i + 1] = largest palindrome cetered between chars i and
void manacher(char *s, int n, int *rad) {
⊳ static char t[2*MAX];
\triangleright int m = 2 * n - 1;
\triangleright \mathbf{for}(\mathbf{int} \ i = 0; \ i < m; \ i++) \ t\lceil i\rceil = -1;
\triangleright for(int i = 0; i < n; i++) t[2 * i] = s[i];
\triangleright int x = 0:
p rad[0] = 0; // <</pre>
\triangleright for(int i = 1; i < m; i++) {
\triangleright int &r = rad[i] = 0;
\triangleright if(i <= x+rad[x]) r = min(rad[x+x-i],x+rad[x]-i);
\triangleright while(i - r - 1 >= 0 and i + r + 1 < m and
\triangleright \triangleright \mathsf{t}[\mathsf{i} - \mathsf{r} - \mathsf{1}] == \mathsf{t}[\mathsf{i} + \mathsf{r} + \mathsf{1}]) ++\mathsf{r};
\triangleright if(i + r >= x + rad[x]) x = i:
⊳ }
\triangleright for(int i = 0: i < m: i++) {
\triangleright // for(int i = 0; i < m; i++) rad[i] /= 2;
```

4.8 Suffix Automaton

```
map<char, int> to[2*N];
int link[2*N], len[2*N], last = 0, sz = 1;
void add_letter(char c){
   int p = last;
```

```
last = sz++:
len[last] = len[p] + 1:
for(; !to[p][c]; p = link[p]) to[p][c] = last;
if(to[p][c] == last){
   link[last] = 0;
   return;
int u = to[p][c];
if(len[u] == len[p]+1){
   link[last] = u;
   return;
int c1 = sz++;
to[c1] = to[u];
link[c1] = link[u];
len[c1] = len[p]+1;
link[last] = link[u] = c1;
for(; to[p][c] == u; p = link[p]) to[p][c] = c1;
```

Geometry

5.1 2D basics

```
typedef double cod;
double eps = 1e-7;
bool eq(cod a, cod b){ return abs(a - b) <= eps; }</pre>
struct vec{
⊳ cod x, y; int id;
\triangleright vec(cod a = 0, cod b = 0) : x(a), y(b) {}

    vec operator+(const vec &o) const{
\triangleright return {x + o.x, y + o.y};

    vec operator-(const vec &o) const{
\triangleright return {x - o.x, y - o.y};

    vec operator*(cod t) const{
▷ return {x * t, y * t};

  vec operator/(cod t) const{
▷ return {x / t, y / t};

    cod operator*(const vec &o) const{ // cos
▷ return x * o.x + y * o.y;

    cod operator^(const vec &o) const{ // sin
▷ return x * o.y - y * o.x;

    bool operator==(const vec &o) const{
\triangleright return eq(x, o.x) && eq(y, o.y);
⊳ }

    bool operator<(const vec &o) const{</pre>
\triangleright if(!eq(x, o.x)) return x < o.x;
⊳ ⊳ return y < o.y;
⊳ cod cross(const vec &a, const vec &b) const{
⊳ }
   int ccw(const vec &a, const vec &b) const{
       cod tmp = cross(a, b);
       return (tmp > eps) - (tmp < -eps);</pre>
⊳ cod dot(const vec &a, const vec &b) const{

    return (a-(*this)) * (b-(*this));
⊳ }
▷ cod len() const{
```

```
    return atan2(cross(a, b), dot(a, b));

▷ return cross(a, b) / dot(a, b);
⊳ }
▶ vec unit() const{
▷ return operator/(len());
⊳ }

   int quad() const{
\Rightarrow if(x > 0 && y >=0) return 0;
\triangleright if(x <=0 && y > 0) return 1;
\triangleright if(x < 0 && y <=0) return 2;
⊳ ⊳ return 3:

    bool comp(const vec &a, const vec &b) const{

    return (a - *this).comp(b - *this);
bool comp(vec b){

    return (*this) * (*this) < b * b:
</pre>
⊳ }

    template<class T>

▶ void sort_by_angle(T first, T last) const{
▷ ▷ std::sort(first, last, [=](const vec &a, const vec &b){
▷ ▷ ▷ return comp(a, b);
▷ ▷ });
⊳ }
vec rot90() const{ return {-y, x}; }

  vec rot(double a) const{
\triangleright return {cos(a)*x -sin(a)*y, sin(a)*x +cos(a)*y};
   vec proj(const vec &b) const{ // proj of *this onto b
       cod k = operator*(b) / (b * b);
       return b * k:
   // proj of (*this) onto the plane orthogonal to b
   vec rejection(vec b) const{
       return (*this) - proj(b);
};
struct line{
⊳ cod a, b, c; vec n;
\triangleright line(vec q, vec w){ // q.cross(w, (x, y)) = 0
\triangleright \triangleright a = -(w.y-q.y);
\triangleright \triangleright b = w.x-q.x;
\triangleright c = -(a * q.x + b * q.y);
\triangleright \triangleright n = \{a, b\};
⊳ }
⊳ cod dist(const vec &o) const{
▷ return abs(eval(o)) / n.len();
▶ bool contains(const vec &o) const{
\triangleright return eq(a * o.x + b * o.y + c, 0);

    cod dist(const line &o) const{
▷ if(!parallel(o)) return 0;
\triangleright if(!eq(o.a * b, o.b * a)) return 0;
▷ ▷ if(!eq(a, 0))
▷ ▷ return abs(c - o.c * a / o.a) / n.len();
▷ ▷ if(!eq(b, 0))
▷ ▷ return abs(c - o.c * b / o.b) / n.len();
▷ ▷ return abs(c - o.c):
▶ bool parallel(const line &o) const{
```

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```
▷ return eq(n ^ o.n, 0);
> }

    bool operator==(const line &o) const{
▷ if(!eq(a*o.b, b*o.a)) return false;
▷ if(!eq(a*o.c, c*o.a)) return false;
▷ if(!eq(c*o.b, b*o.c)) return false;
▷ ▷ return true;
⊳ }
▶ bool intersect(const line &o) const{
▷ return !parallel(o) || *this == o;
⊳ }

  vec inter(const line &o) const{
▷ ▷ if(parallel(o)){
▷ ▷ ▷ if(*this == 0){ }
▷ ▷ ▷ else{ /* dont intersect */ }
\triangleright auto tmp = n o.n;

    return {(o.c*b -c*o.b)/tmp, (o.a*c -a*o.c)/tmp};
⊳ }
▷ vec at_x(cod x) const{
\triangleright return {x, (-c-a*x)/b};
> }

  vec at_y(cod y) const{
\triangleright return {(-c-b*y)/a, y};
⊳ }

    cod eval(const vec &o) const{

    return a * o.x + b * o.y + c;

⊳ }
};
struct segment{
⊳ vec p, q;

    segment(vec a = vec(), vec b = vec()): p(a), q(b) {}
▶ bool onstrip(const vec &o) const{ // onstrip strip

    return p.dot(o, q) >= -eps && q.dot(o, p) >= -eps;

    cod len() const{
▷ return (p-q).len();
⊳ }

    cod dist(const vec &o) const{
▷ if(onstrip(o)) return line(p, q).dist(o);
p return min((o-q).len(), (o-p).len());
⊳ }
▶ bool contains(const vec &o) const{

    return eq(p.cross(q, o), 0) && onstrip(o);

▶ bool intersect(const segment &o) const{
▷ if(contains(o.p)) return true;
▷ if(contains(o.q)) return true;
▷ if(o.contains(q)) return true;
▷ if(o.contains(p)) return true;
\triangleright return p.ccw(q, o.p) * p.ccw(q, o.q) == -1
       && o.p.ccw(o.q, q) * o.p.ccw(o.q, p) == -1;
▶ bool intersect(const line &o) const{
▷ return o.eval(p) * o.eval(q) <= 0;</pre>
⊳ }

    cod dist(const segment &o) const{
▷ ▷ if(onstrip(o.p) || onstrip(o.q)
▷ ▷ ▷ || o.onstrip(p) || o.onstrip(q))
▷ ▷ ▷ return line(p, q).dist(line(o.p, o.q));
▷ ▷ else if(intersect(o)) return 0;

    return min(min(dist(o.p), dist(o.q)),
▷ ▷ ▷ min(o.dist(p), o.dist(q)));
⊳ }
```

```
⊳ cod dist(const line &o) const{
▷ if(line(p, q).parallel(o))
▷ ▷ return line(p, q).dist(o);
▷ ▷ else if(intersect(o)) return 0;
▷ return min(o.dist(p), o.dist(q));
⊳ }
};
struct hrav{
⊳ vec p, q;
\triangleright hray(vec a = vec(), vec b = vec()): p(a), q(b){}
▶ bool onstrip(const vec &o) const{ // onstrip strip
▷ return p.dot(q, o) >= -eps;
⊳ }
⊳ cod dist(const vec &o) const{
▷ if(onstrip(o)) return line(p, q).dist(o);
▷ ▷ return (o-p).len();
⊳ }
▶ bool intersect(const segment &o) const{
▷ if(line(o.p, o.q).parallel(line(p,q)))
▷ ▷ return contains(o.p) || contains(o.q);

    return contains(line(p,q).inter(line(o.p,o.q)));
▶ bool contains(const vec &o) const{

    return eq(line(p, q).eval(o), 0) && onstrip(o);

    cod dist(const segment &o) const{
▷ if(line(p, q).parallel(line(o.p, o.q))){
▷ ▷ if(onstrip(o.p) || onstrip(o.q))
▷ ▷ ▷ return line(p, q).dist(line(o.p, o.q));
▷ ▷ ▷ return o.dist(p);
⊳ ⊳ }
▷ ▷ else if(intersect(o)) return 0;
▷ return min(min(dist(o.p), dist(o.q)),
▷ ▷ ▷ ▷ o.dist(p));
⊳ }
▶ bool intersect(const hray &o) const{
▷ if(!line(p, q).parallel(line(o.p, o.q)))
▷ ▷ ▷ return false;
p auto pt = line(p, q).inter(line(o.p, o.q));
▶ bool intersect(const line &o) const{
▷ if(o.contains(p) || o.contains(q)) return true;
p return (o.eval(p) >= -eps)^(o.eval(p)<o.eval(q));</pre>
▷ return contains(o.inter(line(p, q)));
⊳ }

    cod dist(const line &o) const{
▷ if(line(p,q).parallel(o))
▷ ▷ return line(p,q).dist(o);
▷ ▷ else if(intersect(o)) return 0;
▷ ▷ return o.dist(p);
⊳ }

    cod dist(const hray &o) const{
▷ ▷ if(onstrip(o.p) || o.onstrip(p))
▷ ▷ ▷ return line(p,q).dist(line(o.p, o.q));
▷ ▷ return (p-o.p).len();
⊳ ⊳ }
▷ ▷ else if(intersect(o)) return 0;
▷ return min(dist(o.p), o.dist(p));
⊳ }
};
double heron(cod a, cod b, cod c){
> cod s = (a + b + c) / 2;
```

```
return sqrt(s * (s - a) * (s - b) * (s - c));
line mediatrix(const vec &a, const vec &b) {
\triangleright auto tmp = (b - a) * 2;

   return line(tmp.x, tmp.y, a * a - b * b);
struct circle {
⊳ vec c; cod r;

    circle() : c(0, 0), r(0) {}
\triangleright circle(const vec o) : c(o), r(0) {}
⊳ circle(const vec &a, const vec &b) {
\triangleright c = (a + b) * 0.5; r = (a - c).len();
⊳ }
⊳ circle(const vec &a, const vec &b, const vec &cc) {
▷ ▷ c = mediatrix(a, b).inter(mediatrix(b, cc));
\triangleright \triangleright r = (a - c).len();
⊳ }
▶ bool inside(const vec &a) const {
▷ return (a - c).len() <= r;</pre>
⊳ }
};
circle min_circle_cover(vector<vec> v) {
random_shuffle(v.begin(), v.end());
⊳ circle ans;
p int n = (int)v.size();
\triangleright for(int i = 0; i < n; i++) if(!ans.inside(v[i])) {
▷ ▷ ans = circle(v[i]);
\triangleright \mathsf{for}(\mathsf{int}\ j = 0;\ j < i;\ j++)\ \mathsf{if}(!\mathsf{ans.inside}(\mathsf{v}[j]))

▷ ▷ ans = circle(v[i], v[j]);
\triangleright \triangleright  for(int k=0; k<j; k++)if(!ans.inside(v[k])){
\triangleright \triangleright \triangleright \triangleright ans = circle(v[i], v[j], v[k]);
▷ ▷ ▷ }
⊳ ⊳ }
⊳ }
▶ return ans;
```

5.2 Circle line intersection

```
// intersection of line a * x + b * y + c = 0
// and circle centered at the origin with radius r
double r, a, b, c; // given as input
double x0 = -a*c/(a*a+b*b), y0 = -b*c/(a*a+b*b);
if(c*c > r*r*(a*a+b*b)+EPS)
   puts("no points");
else if(abs(c*c - r*r*(a*a+b*b)) < EPS){
   puts("1 point");
   cout << x0 << ' ' << y0 << '\n';
else {
   double d = r*r - c*c/(a*a+b*b);
   double mult = sqrt (d / (a*a+b*b));
   double ax, ay, bx, by;
   ax = x0 + b * mult;
   bx = x0 - b * mult;
   ay = y0 - a * mult;
   bv = v0 + a * mult:
   puts ("2 points");
   cout<<ax<<' '<<ay<<'\n'<<bx<<' '<<by<<'\n';
```

5.3 Half plane intersection

```
const double eps = 1e-8;
typedef pair<long double, long double> pi;
bool z(long double x){ return fabs(x) < eps; }
struct line{
   long double a, b, c;
   bool operator<(const line &l)const{</pre>
```

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```
\triangleright bool flag1 = pi(a, b) > pi(0, 0);
\triangleright bool flag2 = pi(1.a, 1.b) > pi(0, 0);
▷ ▷ if(flag1 != flag2) return flag1 > flag2;
\triangleright long double t = ccw(pi(0, 0), pi(a, b), pi(l.a, l.b));
\triangleright return z(t) ? c * hypot(1.a, 1.b) < 1.c * hypot(a, b) : t > 0;
p pi slope(){ return pi(a, b); }
};
pi cross(line a, line b){

    long double det = a.a * b.b - b.a * a.b;

    return pi((a.c * b.b - a.b * b.c) / det, (a.a * b.c - a.c * b.a)

    →/ det):
bool bad(line a, line b, line c){

    if(ccw(pi(0, 0), a.slope(), b.slope()) <= 0) return false;</pre>
p pi crs = cross(a, b);
▶ return crs.first * c.a + crs.second * c.b >= c.c;
bool solve(vector<line> v, vector<pi> &solution){ // ax + by <= c;</pre>
⊳ sort(v.begin(), v.end());

    dequedq;

    for(auto &i : v){
\triangleright if(!dq.empty() && z(ccw(pi(0, 0), dq.back().slope(),
    \rightarrowi.slope()))) continue;
b while(dq.size() >= 2 && bad(dq[dq.size()-2], dq.back(), i))
    \hookrightarrowdq.pop_back();
\triangleright while(dq.size() >= 2 && bad(i, dq[0], dq[1])) dq.pop_front();
▷ b dq.push_back(i);
⊳ }
b while(dq.size() > 2 && bad(dq[dq.size()-2], dq.back(), dq[0]))
    →dq.pop_back();
while(dq.size() > 2 && bad(dq.back(), dq[0], dq[1]))
    →dq.pop_front();

  vector<pi> tmp;
p for(int i=0; i<dq.size(); i++){</pre>
\triangleright line cur = dq[i], nxt = dq[(i+1)%dq.size()];

    if(ccw(pi(0, 0), cur.slope(), nxt.slope()) <= eps) return false;</pre>
b tmp.push_back(cross(cur, nxt));
⊳ }
⊳ solution = tmp;
▷ return true;
```

5.4 Detect empty Half plane intersection

```
// abs(point a) = absolute value of a
// ccw(a, b, c) = a.ccw(b, c)
pair<bool, point> half_inter(vector<pair<point,point> > &vet){
   random_shuffle(all(vet));
   rep(i,0,sz(vet)) if(ccw(vet[i].x,vet[i].y,p) != 1){
       point dir = (vet[i].y - vet[i].x) / abs(vet[i].y -
          \hookrightarrow vet[i].x);
       point 1 = vet[i].x - dir*1e15;
       point r = vet[i].x + dir*1e15;
       if(r < 1) swap(1, r);
       rep(j, 0, i){
          if(ccw(point(), vet[i].x-vet[i].y, vet[j].x-vet[j].y) ==
              →0){
              if(ccw(vet[j].x, vet[j].y, p) == 1)
                 continue;
              return mp(false, point());
          if(ccw(vet[j].x, vet[j].y, 1) != 1)
                  →line_intersect(vet[i].x,vet[i].y,vet[j].x,vet[j].y));
          if(ccw(vet[j].x, vet[j].y, r) != 1)
```

5.5 Circle Circle intersection

Assume that the first circle is centered at the origin and second at (x2, y2). Find circle line intersection of first circle and line Ax + By + C = 0, where $A = -2x_2$, $B = -2y_2$, $C = x_2^2 + y_2^2 + r_1^2 - r_2^2$.

Be aware of corner case with two circles centered at the same point.

5.6 Tangents of two circles

```
// solve first for same circle(and infinitely many tangents)
// Find up to four tangents of two circles
void tangents(pt c, double r1, double r2, vector<line> & ans){
   double r = r2 - r1:
   double z = c.x * c.x + c.y * c.y;
   double d = z - r * r;
   if(d < -EPS) return;</pre>
   d = sqrt(abs(d));
   line 1:
   1.a = (c.x * r + c.y * d) / z;
   1.b = (c.y * r - c.x * d) / z;
   1.c = r1:
   ans.push_back (1);
vector<line> tangents(circle a, circle b){
   vector<line> ans:
   pt aux = a.center - b.center;
   for(int i = -1: i <= 1: i += 2)
       for(int j = -1; j \le 1; j += 2)
          tangents(aux, a.r * i, b.r * j, ans);
    for(size_t i = 0; i < ans.size(); ++i)</pre>
       ans[i].c -= ans[i].a * a.x + ans[i].b * a.y;
   return ans;
```

5.7 Convex Hull

```
U.pop_back();

U.push_back(p);
}

L.pop_back(), U.pop_back();

L.reserve(L.size() + U.size());
L.insert(L.end(), U.begin(), U.end());

return L;
}
```

5.8 Check point inside polygon

5.9 Check point inside polygon without lower/upper hull

```
/// borders included
// must not have 3 colinear consecutive points
bool inside_poly(const vector<vec> &v, vec p){
    if(v[0].ccw(v[1], p) < 0) return false;
    if(v[0].ccw(v.back(), p) > 0) return 0;
    if(v[0].ccw(v.back(), p) == 0)
        return v[0].dot(p, v.back()) >= 0
        && v.back().dot(p, v[0]) >= 0;

int L = 1, R = (int)v.size() - 1, ans = 1;

while(L <= R){
    int mid = (L+R)/2;
    if(v[0].ccw(v[mid], p) >= 0) ans = mid, L = mid+1;
    else R = mid-1;
    }

return v[ans].ccw(v[(ans+1)%v.size()], p) >= 0;
}
```

5.10 Minkowski sum

```
vector<vec> msum(vector<vec>& a, vector<vec>& b) {
    int i = 0, j = 0;
    b for(int k = 0; k < (int)a.size(); k++)
    b if(a[k] < a[i]) i = k;
    b for(int k = 0; k < (int)b.size(); k++)
    b if(b[k] < b[j]) j = k;

    vector<vec> c;
    c.reserve(a.size() + b.size());
```

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```
\triangleright for(int k = 0; k < int(a.size()+b.size()); k++){
▷ ▷ vec pt{a[i] + b[j]};
▷ ▷ if((int)c.size() >= 2
\triangleright \triangleright \triangleright && c[c.size()-2].ccw(c.back(), pt) == 0)
▷ ▷ ▷ c.pop_back();
▷ ▷ c.push_back(pt);
\triangleright int q = i+1, w = j+1;
\triangleright if(q == int(a.size())) q = 0;
\triangleright if(w == int(b.size())) w = 0;
\triangleright if(c.back().ccw(a[i]+b[w], a[q]+b[j]) < 0) i = q;
\triangleright \triangleright else j = w;
⊳ }
▷ c.shrink_to_fit();
▶ return c;
```

5.11 Geo Notes

5.11.1 Center of mass

System of points(2D/3D): Mass weighted average of

Frame(2D/3D): Get middle point of each segment solve as previously.

Triangle: Average of vertices.

2D Polygon: Compute **signed** area and center of mass of triangle $((0,0), p_i, p_{i+1})$. Then solve as system of points.

Polyhedron surface: Solve each face as a 2D polygon(be aware of (0, 0) then replace each face with its center of mass and solve as system of points.

Tetrahedron(Triangular pyramid): As triangles, its the average of points.

Polyhedron: Can be done as 2D polygon, but with tetrahedralization intead of triangulation.

5.11.2 Pick's Theorem

Given a polygon without self-intersections and all its vertices on integer coordinates in some 2D grid. Let A be its area, *I* the number of points with integer coordinates stricly inside the polygon and B the number of points with integer coordinates in the border of the polygon. The following formula holds: $A = I + \frac{B}{2} - 1$.

Miscellaneous

6.1 Cute LIS

```
multiset<int> S;
for(int i = 0; i < n; i++){

    auto it = S.upper_bound(a[i]); // low for inc

    if(it != S.end()) S.erase(it);
▷ S.insert(a[i]);
ans = S.size();
```

6.2 Cute LIS

```
multiset<int> S;
for(int i = 0; i < n; i++){

    auto it = S.upper_bound(a[i]); // low for inc

b if(it != S.end()) S.erase(it);
▷ S.insert(a[i]);
ans = S.size();
6.3 Efficient recursive lambda
```

```
template<class Fun>
class y_combinator_result {
 Fun fun_;
public:
 template<class T>
 explicit y_combinator_result(T &&fun):
      →fun_(std::forward<T>(fun)) {}
 template<class ...Args>
 decltype(auto) operator()(Args &&...args) {
   return fun_(std::ref(*this), std::forward<Args>(args)...);
template<class Fun>
decltype(auto) y_combinator(Fun &&fun) {
      →y_combinator_result<std::decay_t<Fun>>(std::forward<Fun>(fun))
```

// auto gcd = y_combinator([](auto gcd, int a, int b) -> int {

6.4 Bitsets

// return b == 0 ? a : gcd(b, a % b);

```
#define private public
#include <bitset>
#undef private
#include <bits/stdc++.h>
using namespace std;
#define tab _M_w
using biti = typename
   -remove_reference<decltype(bitset<404>().tab[0])>::type;
const int SIZE = 8 * sizeof(biti);
const int LOG = __builtin_ctz(SIZE);
template<size_t Nw>
int find_prev(const bitset<Nw> &x, int v) {

   int start = v >> LOG;

    int first_bits = v & (SIZE - 1);

▶ if(first_bits) {
▷ ▷ biti curr = x.tab[start];
▷ ▷ if(curr)
▷ ▷ ▶ return start << LOG | (SIZE - __builtin_clzl(curr) - 1);
\triangleright for(int i = start - 1; i >= 0; i--) {
▷ ▷ biti curr = x.tab[i];
▷ ▷ if(curr) {
▷ ▷ ▷ return (i << LOG) | (SIZE - __builtin_clzl(curr) - 1);
⊳ ⊳ }
⊳ }
⊳ return -1;
```

```
// s._Find_first(); s._Find_next(k); find_prev(s, k+1);
// _Unchecked_set/_Unchecked_reset/_Unchecked_flip
```

6.5 Buildings

```
// count the number of circular arrays of size m, with elements on
   \hookrightarrow range [1, c^{**}(n^*n)]
int n, m, c; cin >> n >> m >> c;
int x = f_{exp}(c, n * n); int ans = f_{exp}(x, m);
for(int i = 1; i \le m; i++) if(m % i == 0) {
 int y = f_exp(x, i);
 for(int j = 1; j < i; j++) if(i % j == 0)
     y = sub(y, mult(j, dp[j]));
 dp[i] = mult(y, inv(i));
 ans = sub(ans, mult(i - 1, dp[i]));
cout << ans << '\n';</pre>
6.6 Rand
```

```
#include <random>
#include <chrono>
cout << RAND_MAX << endl;</pre>
     >rng(chrono::steady_clock::now().time_since_epoch().count());
shuffle(p.begin(), p.end(), rng);
uniform_int_distribution<int>(a,b)(rng);
```

6.7 Klondike

```
// minimum number of moves to make
// all elements equal
// move: change a segment of equal value
// elements to any value
int v[305], dp[305][305], rec[305][305];
int f(int 1, int r){
 if(r == 1) return 1;
 if(r < 1) return 0;</pre>
 if(dp[l][r] != -1) return dp[l][r];
 int ans = f(1+1, r) + 1;
 for(int i = 1+1; i <= r; i++)
   if(v[i] == v[1])
     ans = min(ans, f(1, i - 1) + f(i+1, r));
 return dp[l][r] = ans;
```

6.8 Hilbert Order

```
// maybe use B = n / sqrt(q) before this
inline int64_t hilbertOrder(int x, int y, int pow = 21, int rotate
   ⇒= 0) {
b if(pow == 0) return 0;
\triangleright int hpow = 1 << (pow-1);
\triangleright int seg = (x < hpow) ? (
\triangleright \triangleright (y < hpow) ? 0 : 3
⊳):(
\triangleright \triangleright (y < hpow) ? 1 : 2
▷ );
⊳ seg = (seg + rotate) & 3;

    const int rotateDelta[4] = {3, 0, 0, 1};

\triangleright int nx = x & (x ^ hpow), ny = y & (y ^ hpow);
b int nrot = (rotate + rotateDelta[seg]) & 3;

    int64_t subSquareSize = int64_t(1) << (2*pow - 2);
</pre>
b int64_t ans = seg * subSquareSize;

    int64_t add = hilbert0rder(nx, ny, pow-1, nrot);
\triangleright ans += (seg == 1 || seg == 2) ? add : (subSquareSize - add - 1);
▶ return ans;
```

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6.9 Modular Factorial

```
// Compute (1*2*...*(p-1)*1*(p+1)*(p+2)*..*n) % p
// in O(p*lg(n))
int factmod(int n, int p){
   int ans = 1:
   while (n > 1) {
       for(int i = 2; i <= n % p; i++)</pre>
          ans = (ans * i) % p:
      n /= p:
      if(n \% 2) ans = p - ans;
   return ans % p;
int fac_pow(int n, int p){
   int ans = 0;
   while(n) n /= p, ans += n;
   return ans;
int C(int n. int k. int p){
   if(fac_pow(n, p) > fac_pow(n-k, p) + fac_pow(k, p))
   int tmp = factmod(k, p) * factmod(n-k, p) % p;
   return (f_{exp}(tmp, p - 2, p) * factmod(n, p)) % p;
```

6.10 Enumeration all submasks of a bitmask

```
// loop through all submask of a given bitmask
// it does not include mask 0
for(int sub = mask; sub; sub = (sub - 1) & mask){
// loop through all supermasks of a given bitmask
for(int super = mask; super < (1 << n); super = (super + 1) |</pre>
   ⇒mask) {
```

6.11 Knapsack Bounded with Cost

```
// menor custo para conseguir peso ate M usando N tipos diferentes

→ de elementos, sendo que o i-esimo elemento pode ser usado

    \hookrightarrow b[i] vezes, tem peso w[i] e custo c[i]
// O(N * M)
int b[N], w[N], c[N];
MinOueue O[M]
int d[M] //d[i] = custo minimo para conseguir peso i
for(int i = 0; i \leftarrow M; i++) d[i] = i ? oo : 0;
for(int i = 0; i < N; i++){
\triangleright for(int j = 0; j < w[i]; j++)
▷ ▷ Q[j].clear();
\triangleright for(int j = 0; j \leftarrow M; j++){
\triangleright \ \mathsf{q} = \mathsf{Q[j \% w[i]]};
▷ if(q.size() >= q) q.pop();
▷ □ q.add(c[i]);
▷ ▷ q.push(d[j]);
▷ ▷ d[j] = q.getmin();
⊳ }
```

6.12 LCA < O(nlgn), O(1) >

```
int start[N], dfs_time;
int tour[2*N], id[2*N];
void dfs(int u){
```

```
start[u] = dfs time:
   id[dfs time] = u:
   tour[dfs_time++] = start[u];
   for(int v : g[u]){
       dfs(v);
      id[dfs_time] = u;
      tour[dfs_time++] = start[u];
int LCA(int u, int v){
   if(start[u] > start[v]) swap(u, v);
   return id[min(tour[k]for k in [start[u],start[v]])];
```

6.13 Buffered reader

```
→https://github.com/ngthanhtrung23/ACM_Notebook_new/blob/master/w6.15 d_Edge coloring CPP
int INP.AM.REACHEOF:
#define BUFSIZE (1<<12)</pre>
char BUF[BUFSIZE+1]. *inp=BUF:
#define GETCHAR(INP) { \
   if(!*inp && !REACHEOF) { \
      memset(BUF.0.sizeof BUF):\
      int inpzzz = fread(BUF,1,BUFSIZE,stdin);\
      if (inpzzz != BUFSIZE) REACHEOF = true:\
      inp=BUF: \
   } \
   INP=*inp++; \
#define DIG(a) (((a)>='0')&&((a)<='9'))
#define GN(j) { \
   GETCHAR(INP): while(!DIG(INP) && INP!='-') GETCHAR(INP):\
   if (INP=='-') {AM=1:GETCHAR(INP):} \
   j=INP-'0'; GETCHAR(INP); \
   while(DIG(INP)){j=10*j+(INP-'0');GETCHAR(INP);} \
   if (AM) j=-j;\
```

6.14 Modular summation

```
//calcula (sum(0 \ll i \ll n) P(i)) \% mod,
//onde P(i) eh uma PA modular (com outro modulo)
namespace sum_pa_mod{
▷ ▷ assert(a&&b);
▷ b if(a >= b){
\triangleright \triangleright 11 \text{ ret } = ((n*(n+1)/2) \text{mod})*(a/b);
\triangleright \triangleright if(a\%b) \text{ ret } = (\text{ret } + \text{calc}(a\%b,b,n,mod))\%\text{mod};
▷ ▷ else ret = (ret+n+1)%mod;
▷ ▷ ▷ return ret;
⊳ ⊳ }
\triangleright return ((n+1)*(((n*a)/b+1)%mod) - calc(b,a,(n*a)/b,mod) + mod +
    \hookrightarrown/b + 1)%mod:
\triangleright //P(i) = a*i mod m
\triangleright \triangleright a = (a\%m + m)\%m;
▷ ▷ if(!a) return 0;
\triangleright 11 ret = (n*(n+1)/2)%mod;
▷ ▷ ret = (ret*a)%mod;
\triangleright \triangleright 11 g = \_gcd(a,m);
\triangleright ret -= m*(calc(a/g,m/g,n,mod)-n-1);
▷ return (ret%mod + mod)%mod;
⊳ }
```

```
P(i) = a + r*i \mod m
\triangleright \triangleright a = (a\%m + m)\%m;
\triangleright r = (r\%m + m)\%m;
▷ if(!a) return solve(r, n, m, mod);
\triangleright \triangleright 11 g, x, y;
\triangleright g = gcdExtended(r, m, x, y);
\triangleright \quad x = (x\%m + m)\%m;
\triangleright \ 11 d = a - (a/g)*g;
⊳ ⊳ a -= d;
\triangleright \triangleright x = (x*(a/q))%m:
\triangleright return (solve(r, n+x, m, mod) - solve(r, x-1, m, mod) + mod +
    \hookrightarrowd*(n+1))%mod;
⊳ }
};
```

```
const int MX = 300:
int C[MX][MX] = \{\}, G[MX][MX] = \{\};
void solve(vector<pii> &E, int N){
   int X[MX] = \{\}, a, b;
   auto update = [&](int u){ for(X[u] = 1; C[u][X[u]]; X[u]++); };
   auto color = [&](int u, int v, int c){
      int p = G[u][v];
       G[u][v] = G[v][u] = c;
       C[u][c] = v; C[v][c] = u;
      C[u][p] = C[v][p] = 0;
       if(p) X[u] = X[v] = p;
       else update(u), update(v);
       return p; };
   auto flip = [\&] (int u, int c1, int c2){
       int p = C[u][c1], q = C[u][c2];
       swap(C[u][c1], C[u][c2]);
      if( p ) G[u][p] = G[p][u] = c2;
      if( !C[u][c1] ) X[u] = c1;
      if( !C[u][c2] ) X[u] = c2;
      return p; };
   for(int i = 1; i \le N; i++) X[i] = 1;
   for(int t = 0; t < E.size(); t++){</pre>
      int u = E[t].first, v0 = E[t].second, v = v0, c0 = X[u], c
          \Rightarrow= c0. d:
       vector<pii> L;
       int vst[MX] = {};
       while(!G[u][v0]){
          L.emplace_back(v, d = X[v]);
          if(!C[v][c]) for(a = (int)L.size()-1; a >= 0; a--) c =
              \hookrightarrow if(!C[u][d])for(a=(int)L.size()-1;a>=0;a--)color(u,L[a]
          else if( vst[d] ) break;
          else vst[d] = 1, v = C[u][d];
      if( !G[u][v0] ){
          for(;v; v = flip(v, c, d), swap(c, d));
          if(C[u][c0]){
             for(a = (int)L.size()-2; a >= 0 && L[a].second != c;
                 →a--):
             for(; a >= 0; a--) color(u, L[a].first, L[a].second);
         } else t--;
      }
```

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6.16 Burnside's Lemma

Let (G, \oplus) be a finite group that acts on a set X. It should hold that $e_g * x = x$ and $g_1 * (g_2 * x) = (g_1 \oplus g_2) * x$, $\forall x \in X$, $g_1, g_2 \in G$. For each $g \in G$ let $X^g = \{x \in X \mid g * x = x\}$. The number of orbits its given by:

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|$$

6.17 Wilson's Theorem

 $(n-1)! = -1 \mod n \iff n \text{ is prime}$

6.18 Fibonacci

- $F_{n-1}F_{n+1} F_n^2 = (-1)^n$
- $F_{n+k} = F_k F_{n+1} + F_{k-1} F_n$
- $GCD(F_n, F_m) = F_{GCD(n,m)}$
- $F_n = \frac{(\frac{1+\sqrt{5}}{2})^n (\frac{1-\sqrt{5}}{2})^n}{\sqrt{5}}$

6.19 Lucas's Theorem

For non-negative integers m and n and a prime p, the following congruence holds:

$$\binom{m}{n} \equiv \prod_{i=0}^{k} \binom{m_i}{n_i} \pmod{p}$$

where m_i is the i-th digit of m in base p. $\binom{a}{b} = 0$ if a < b.

6.20 Kirchhoff's Theorem

Laplacian matrix is L = D - A, where D is a diagonal matrix with vertex degrees on the diagonals and A is adjacency matrix.

The number of spanning trees is any cofactor of L. i-th cofactor is determinant of the matrix gotten by removing i-th row and column of L.

6.20.1 Multigraphs

In D[i][i] all loops are excluded. A[i][j] = number of edges from i to j.

6.20.2 Directed multigraphs

D[i][i] = indegree of i minus the number of loops at i. A[i][j] = number of edges from i to j.

The number of oriented spanning trees rooted at a vertex i is the determinant of the matrix gotten by removing the ith row and column of L.

6.21 Matroid

Let *X* set of objects, $I \subseteq 2^X$ set of independents sets such that:

- 1. $\emptyset \in I$
- 2. $A \in I.B \subseteq A \implies B \in I$
- 3. Exchange axiom, $A \in I, B \in I, |B| > |A| \implies \exists x \in B \setminus A : A \cup \{x\} \in I$
- 4. $A \subseteq X$ and I and I' are maximal independent subsets of A then |I| = |I'|

Then (X, I) is a matroid. The combinatorial optimization problem associated with it is: Given a weight $w(e) \ge 0 \ \forall e \in X$, find an independet subset that has the largest possible total weight.

6.22 Matroid intersection

```
// Input two matroids (X, I_a) and (X, I_b)
// output set I of maximum size, I \setminusin I_a and I \setminusin I_b
set<> I;
while(1){
   for(e_i : X \setminus I)
       if(I + e_i \in I_a \text{ and } I + e_i \in I_b)
          I = I + e_i;
    set<> A, T; queue<> Q;
    for(x : X) label[x] = MARK1;
    for(e_i : X \setminus I){
       if(I + e_i \in I_a)
           Q.push(e_i), label[e_i] = MARK2;
           for(x such that I - x + e_i \setminus in I_a)
              A[x].push(e_i);
       if(I + e_i \setminus in I_b)
           T = T + \{e_i\}
           for(x such that I - x + e_i \in I_b)
               A[e_i].push(x);
   if(T.empty()) break;
   bool found = false;
   while(!Q.empty() and !found){
       auto e = Q.front(); Q.pop();
       for(x : A[e]) if(label[x] == MARK1){
           label[x] = e; Q.push(x);
           if(x \in T)
               found = true; put = 1;
               while(label[x] != MARK2){
                  I = put ? (I + x) : (I - x);
                  put = 1 - put;
              I = I + x;
               break;
```

Where path(e) = [e] if label[e] = MARK2, path(label[e]) + [e] otherwise.

6.22.1 Matroid Union

Given k matroids over the same set of objects (X, I_1) , (X, I_2) , ..., (X, I_k) find $A_1 \in I_1$, $A_2 \in I_2$, ..., $A_k \in I_k$ such that $i \neq j$, $A_i \cap A_j = \emptyset$ and $\bigcup_{i=1}^k A_i |$ is maximum. Matroid union can be reduced to matroid intersection as follows.

Let $X' = X \times \{1, 2, ..., k\}$, ie, k copies of each element of X with different colors. M1 = (X', Q) where $B \in Q \iff \forall 1 \le i \le k$, $\{x \mid (x, i) \in B\} \in I_i$, ie, for each color, B is independent. M2 = (X', W) where $B \in W \iff i \ne j \implies \neg((x, i) \in B \land (x, j) \in B)$, ie, each element is picked by at most one color.

Intersection of *M*1 and *M*2 is the answer for the combinatorial problem of matroid union.

6.23 Notes

When we repeat something and each time we have probability p to succeed then the expected number or tries is $\frac{1}{n}$, till we succeed.

Small to large

Trick in statement If k sets are given you should note that the amount of different set sizes is $O(\sqrt{s})$ where s is total size of those sets. And no more than \sqrt{s} sets have size greater than \sqrt{s} . For example, a path to the root in Aho-Corasick through suffix links will have at most $O(\sqrt{s})$ vertices.

gcd on subsegment, we have at most $log(a_i)$ different values in { $gcd(a_i, a_{i+1}, ..., a_i)$ for i < i }.

From static set to expandable. To insert, create a new set with the new element. While there are two sets with same size, merge them. There will be at most log(n) disjoints sets.

Matrix exponentiation optimization. Save binary power of A_{nxn} and answer q queries $b = A^m x$ in $O((n^3 + n^2))$

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 $qn^2)log(m)$).

Ternary search on integers into binary search, comparing f(mid) and f(mid+1), binary search on derivative

Dynamic offline set For each element we will wind segment of time [a, b] such that element is present in the set during this whole segment. Now we can come up with recursive procedure which handles [l, r] time segment considering that all elements such that $[l, r] \subset [a, b]$ are already included into the set. Now, keeping this invariant

we recursively go into [l, m] and [m + 1, r] subsegments. Finally when we come into segment of length 1.

$$a > b \implies a \mod b < \frac{a}{2}$$

Convex Hull. The expected number of points in the convex hull of a random set of points is O(log(n)). The number of points in a convex hull with points coordinates limited by *L* is $O(L^{2/3})$.

enough. Just do adamant's hld sorting subtrees by their 264241152, 382205952, 530841600

size and remap vertices indexes.

Range query offline can be solved by a sweep, ordering queries by R.

Maximal number of divisors of any n-digit number. 7 4, 12, 32, 64, 128, 240, 448, 768, 1344, 2304, 4032, 6720, 10752, 17280, 26880, 41472, 64512, 103680, 161280, 245760, 368640, 552960, 860160, 1290240, 1966080, 2764800, 4128768, 6193152, 8957952, 13271040, 19660800, 28311552, Tree path query. Sometimes the linear query is fast | 41287680, 59719680, 88473600, 127401984, 181665792,