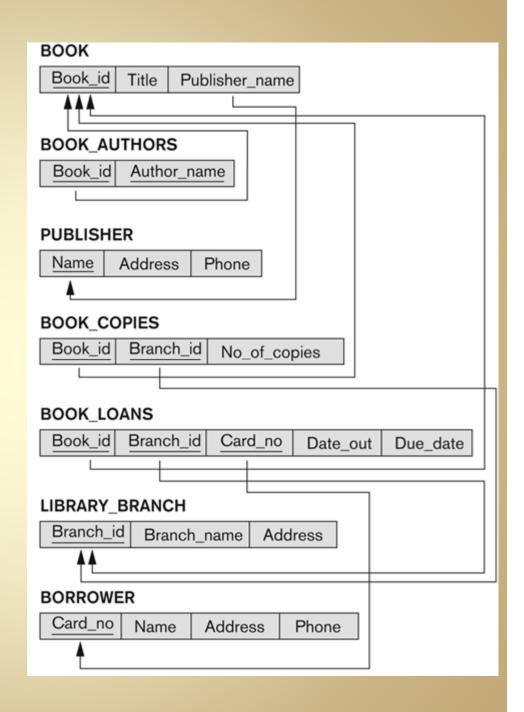
COMP163

Database Management Systems

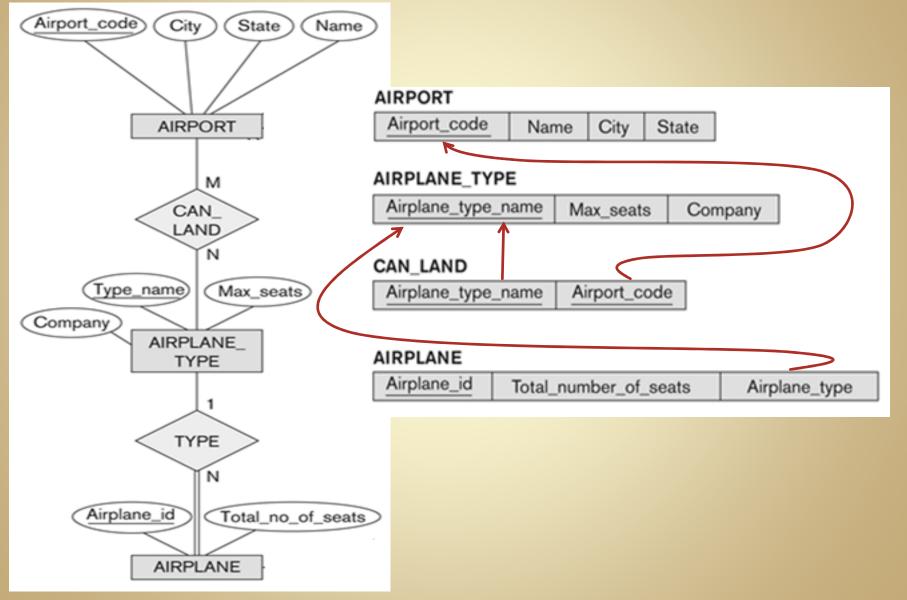
Lecture 5 – Chapter 9 ER to Relational Schema Translation

REVIEW

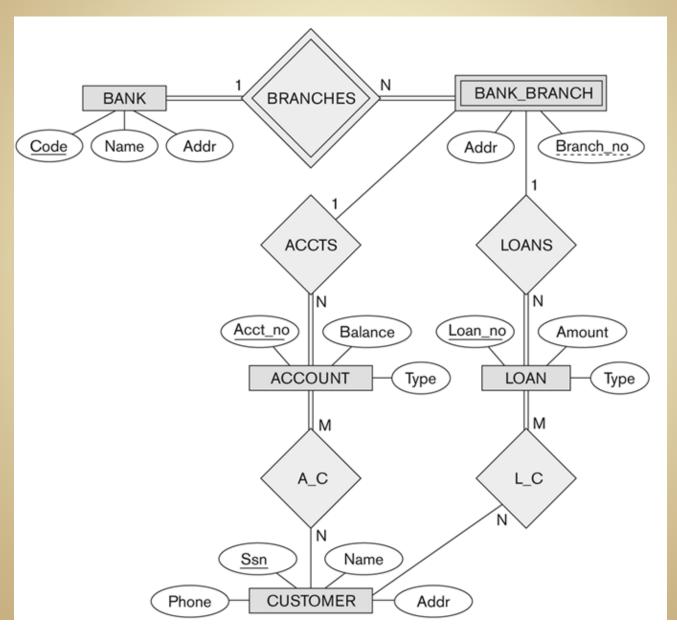
Reverse engineer this relational schema to find an equivalent ER schema.



PREVIEW: ER to Relational

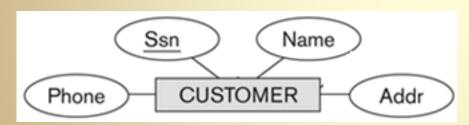


EER Bank Schema



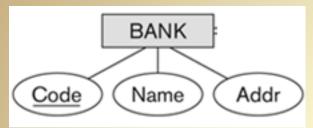
Step 1: Regular Entities

- Regular entity types become relations
 - include all simple attributes
 - include only components of compound attributes
 - keys become primary keys
 - if multiple keys (candidate keys) select a primary key

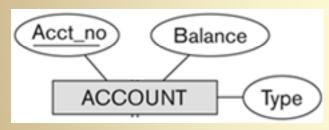


CUSTOMER(Ssn, Name, Addr, Phone)

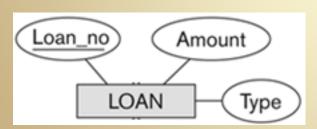
Step 1: Regular Entities



BANK(Code, Name, Addr)



ACCOUNT(Acct_no, Type, Balance)



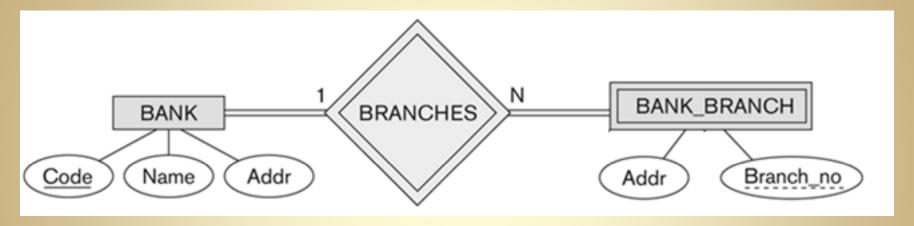
LOAN(Loan no, Type, Amount)

Step 2: Weak Entities

- Weak entity types become relations
 - include all simple attributes
 - include only components of compound attributes
 - create a primary key from partial key and key of owning entity type (through identifying relationship)
 - attributes acquired through identifying relationship become a foreign key*

Step 2: Weak Entities

Weak entity types become relations



BANK_BRANCH(Bank code, Branch No, Addr)

FK

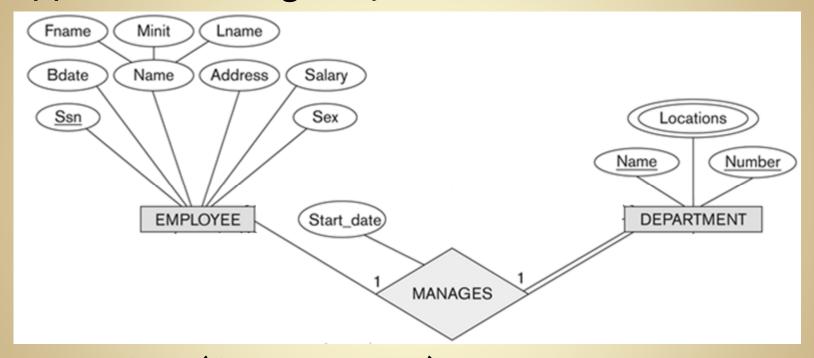
BANK(Code, Name, Addr)

Step 3: Binary 1:1 Relationships

- Approach 1: Foreign Key
 - Chose one of the related entity types to hold the relationship (chose one with total participation, if possible)
 - add FK to other relation
 - move all relationship attributes to this relation
 - this approach is preferable, except as noted below
- Approach 2: Merged Relation
 - combine the relations for the related entities into a single relation
 - use only when both participations are total
- Approach 3: Separate Relation
 - same as binary M:N relationship (see step 5)
 - not generally a good option

Step 3: Binary 1:1 Relationships

Approach 1: Foreign Key



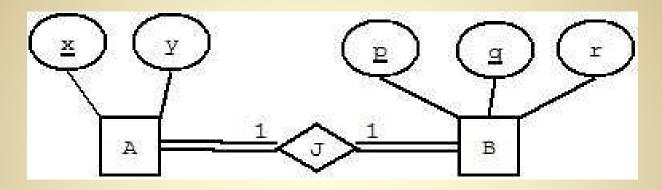
EMPLOYEE(Ssn, Name, ...)

FK

DEPARTMENT(Name, Number, Mgr, Mgr_start_date)

Step 3: Binary 1:1 Relationships

Approach 2: Merged Relation

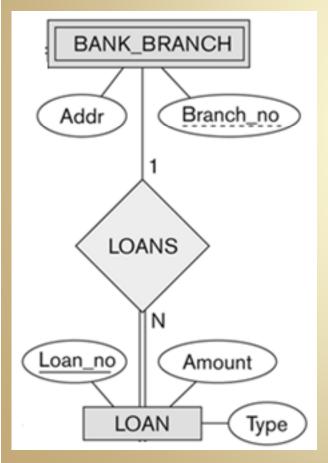


or

AJB(x, y, p,
$$q$$
, r)

Step 4: Binary 1:N Relationships

- 1:N Relationships become foreign key at N side
 - any relationship attributes also go to N side

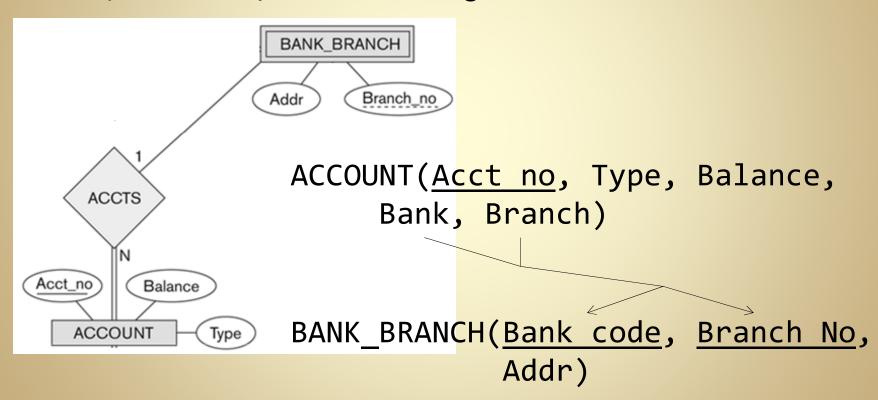


LOAN(Loan no, Type, Amount,
Bank, Branch)

BANK_BRANCH(Bank code, Branch No,
Addr)

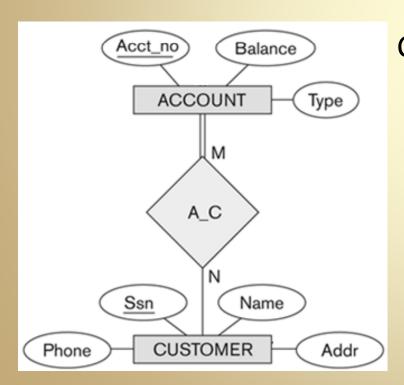
Step 4: Binary 1:N Relationships

- 1:N Relationships become foreign key at N side
 - any relationship attributes also go to N side



Step 5: Binary M:N Relationships

- M:N Relationships must become a new relation
 - contains FKs to both related entities
 - combined FKs become PK for new relations
 - relationship attributes go in new relation



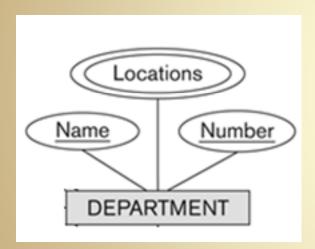
CUSTOMER(<u>Ssn</u>, Name, Addr, Phone)

A_C(<u>Acct</u>, <u>Cust</u>)

ACCOUNT(<u>Acct no</u>, Type, Balance, Bank, Branch)

Step 6: Multivalued Attributes

- Multivalued attributes must become new relations
 - FK to associated entity type
 - PK is whole relation

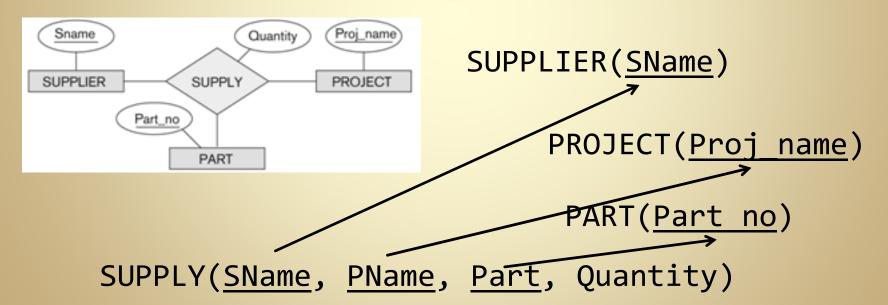


DEPARTMENT(Name, Number, Mgr, Mgr_start_date)

DEPT_LOCATIONS(DName, Dno, Location)

Step 7: N-ary Relationships

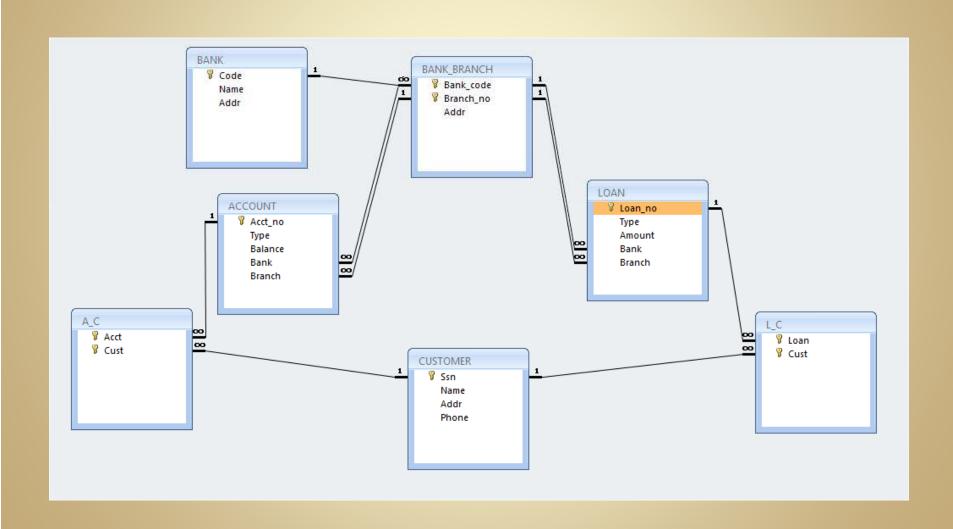
- Non-Binary Relationships become new relations
 - FKs to all participating entity types
 - Combine FKs to make a PK (exclude entities with max participation of 1)
 - Include any relationship attributes



Completed Bank Schema

```
CUSTOMER(Ssn, Name, Addr, Phone)
BANK(Code, Name, Addr)
ACCOUNT(Acct_no, Type, Balance, Bank, Branch)
LOAN(Loan_no, Type, Amount, Bank, Branch)
BANK_BRANCH(Bank_code, Branch_No, Addr)
A_C(Acct, Cust)
L_C(Loan, Cust)
BANK_BRANCH(Bank_code) refers to BANK
LOAN(Bank, Branch) refers to BANK_BRANCH
ACCOUNT(Bank, Branch) refers to BANK BRANCH
A_C(Acct) refers to ACCOUNT
A_C(Cust) refers to CUSTOMER
L_C(Loan) refers to LOAN
L C(Cust) refers to CUSTOMER
```

Bank Schema: MS Access



- Option a: Each entity type becomes a relation
 - all have same PK (from superclass)
 - PKs in subclasses are FKs to superclass
 - most general solution

PERSON(<u>ID</u>, Name)

STUDENT(<u>ID</u>, Major, Class)

PROFESSOR(<u>ID</u>, Dept, Office)

Name

PERSON

PERSON

PROFESSOR(ID) refers to PERSON

STUDENT

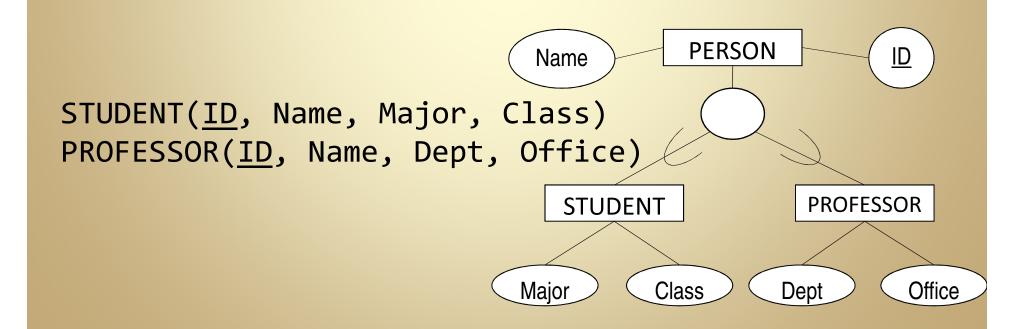
PROFESSOR

STUDENT

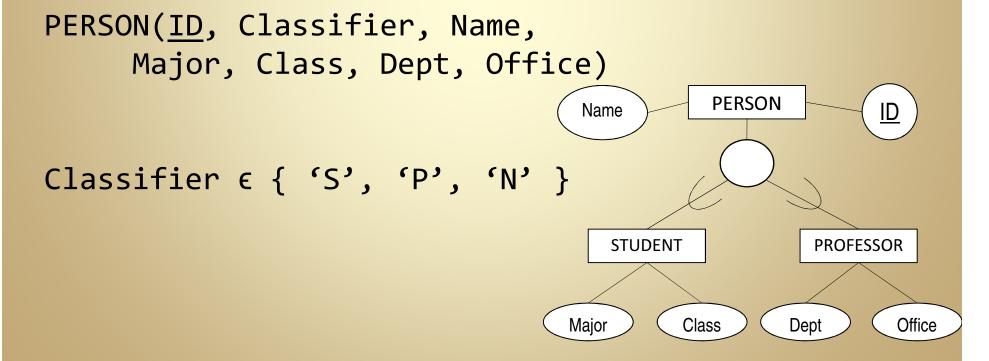
PROFESSOR

Office

- Option b: Each subclass becomes a relation
 - all have same PK (from superclass)
 - each relation gets all superclass attributes
 - restriction: only works for covering inheritance
 - problem: need to join tables to find all PERSONs



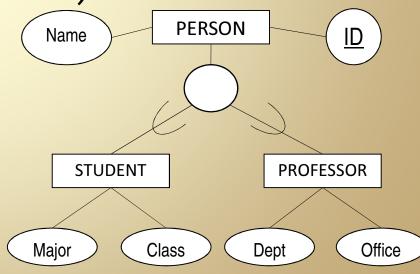
- Option c: Single relation with a type discriminator
 - PK from superclass
 - all attributes from all classes
 - restriction: only works for disjoint inheritance
 - problem: lots of NULL values



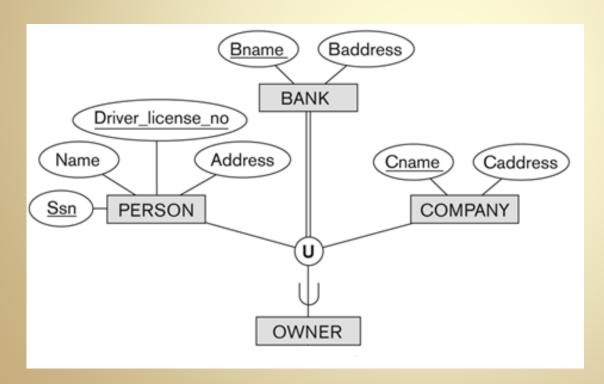
- Option d: Single relation with multiple discriminators
 - PK from superclass
 - all attributes from all classes
 - works for overlapping inheritance
 - problem: lots of NULL values

PERSON(ID, isStudent, isProfessor,
Name, Major, Class, Dept, Office)

dom(isStudent) = Boolean
dom(isProfessor) = Boolean



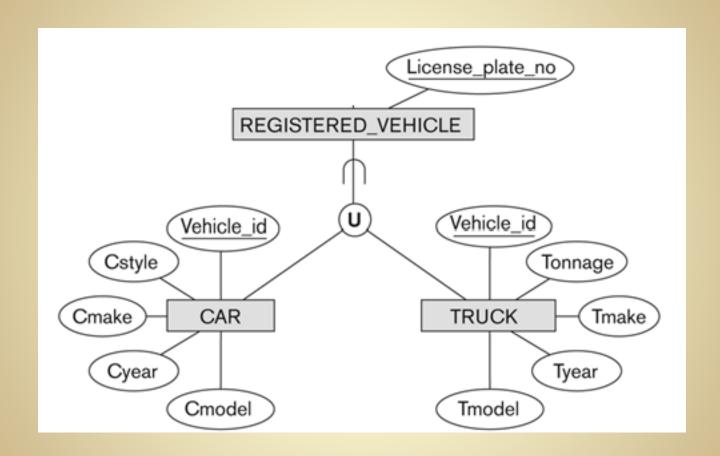
- Union types become a new relation of surrogate keys
 - surrogate keys are added to all defining classes
 - attributes of the union type go in the new relation



add surrogate key to OWNER

```
PERSON(<u>Driver license no</u>, Ssn, Name,
Address, Owner_id)
BANK(<u>Bname</u>, Baddress, Owner_id)
COMPANY(<u>Cname</u>, Caddress, Owner_id)
OWNER(<u>ID</u>)
```

PERSON(Owner_id) refers to OWNER
BANK(Owner_id) refers to OWNER
COMPANY(Owner_id) refers to OWNER

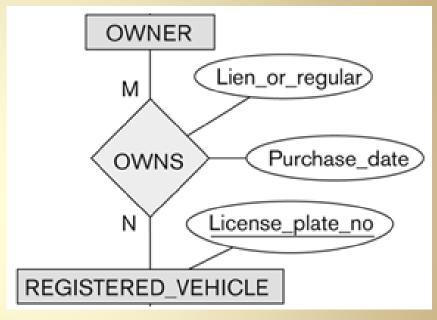


add surrogate key to REGISTERED_VEHICLE

CAR(Vehicle_id) refers to REGISTERED_VEHICLE
TRUCK(Vehicle_id) refers to REGISTERED_VEHICLE

in this case, we don't need to invent a surrogate key, since the domains of CAR keys and TRUCK keys are the same (and non-overlapping)

OWNS relation uses the surrogate keys



OWNS(<u>Owner_id</u>, <u>Vehicle_id</u>, Purchase_date, Lien_or_regular)

OWNS(Owner_id) refers to OWNER
OWNS(Vehicle_id) refers to REGISTERED_VEHICLE

EXERCISES

Create relational schema from the following:

