



## An Empirical Study of Partial Deduction for MINIKANREN

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#### input program

```
let rec eval° fm s r = fm \equiv neg x & eval° x s a & not° a r | ...
```

#### input program

#### input program

```
let rec eval° fm s r = fm \equiv neg x & eval° x s a & not° a r | eval° fm s true \langle --- \rangle known argument fm \equiv neg x & eval° x s a & not° a true | ...
```

## input program let rec eval<sup>o</sup> fm s r = known argument $fm \equiv neg x \& eval^o x s a \& not^o a r$ evalo fm s true < $fm \equiv neg x \& eval^o x s a \& not^o a true$ $fm \equiv neg x \& eval^o x s false$

# input program let rec eval<sup>o</sup> fm s r = known argument $fm \equiv neg x \& eval^o x s a \& not^o a r$ evalo fm s true < $fm \equiv neg x \& eval^o x s a \& not^o a true$ $fm \equiv neg x \& eval^o x s false$

```
input program
 let rec eval<sup>o</sup> fm s r =
                                                  known argument
   fm \equiv neg x \& eval^o x s a \& not^o a r
                             evalo fm s true <
                                                                            output
                                                             let rec eval true fm s =
                                                               fm \equiv neg x \& eval\_false^o x s
fm \equiv neg x \& eval^o x s a \& not^o a true
                                                             let rec eval_false<sup>o</sup> fm s =
                                                               fm \equiv neg x \& eval\_true^o x s
     fm \equiv neg x \& eval^o x s false
```

## Partial Deduction: Specialization for Logic Programming

#### input

```
let double_appendo x y z r =
  ocanren {
    fresh t in
      appendo x y t &
      appendo t z r}
let rec appendo x y r =
  ocanren {
    (x \equiv [] \& y \equiv r) \mid
    (fresh h x' r' in
      x = h \cdot \cdot \cdot x' \&
      append° x' y r' &
      r \equiv h :: r')
```

double\_appendo x y z r

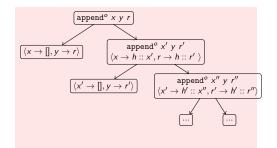
#### output

```
let double_appendo x y z r =
  ocanren {
    (x \equiv [] \& append^o y z r) |
    (fresh h x' r' in
      x \equiv h :: x' \&
      double_append° x' y z r' &
      r \equiv h :: r')
let rec appendo x y r =
  ocanren {
    (x \equiv [] \& y \equiv r) \mid
    (fresh h x' r' in
      x \equiv h :: x' \&
      appendo x' y r' &
      r \equiv h :: r')
```

### Partial Deduction for MINIKANREN: Bird's-eye View

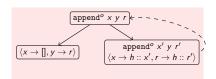
#### input

```
let rec appendo x y r =
  ocanren {
    (x = [] & y = r) |
    (fresh h x' r' in
        x = h :: x' &
        appendo x' y r' &
        r = h :: r')}
```



#### output

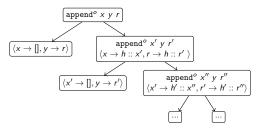
```
let rec append° x y r = ocanren { fresh h t r' in (x \equiv [] \& y \equiv r) \mid (x \equiv h :: x' \& append° x' y r' \& t \equiv h :: r')}
```



#### Partial Deduction: Bird's-eye View

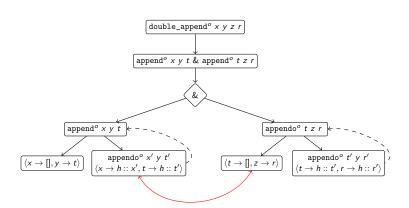
```
let rec append° x y r = ocanren { (x \equiv [] \& y \equiv r) | (fresh h x' r' in x \equiv h :: x' \& append° x' y r' & r \equiv h :: r')}
```

#### Process tree construction

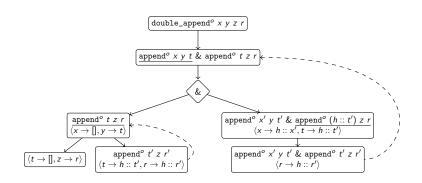


#### Residualization

#### Partial Deduction



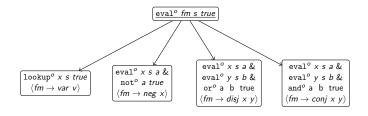
#### Conjunctive Partial Deduction



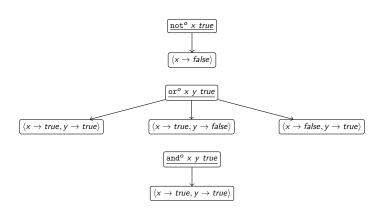
#### Evaluator of Logic Formulas

```
let rec eval° fm s r =
  ocanren { fresh v x y a b in
    (fm = var v & lookup° v s r) |
    (fm = neg x & eval° x s a & not° a r) |
    (fm = conj x y & eval° x s a & eval° y s b & and° a b r) |
    (fm = disj x y & eval° x s a & eval° y s b & oro° a b r)
```

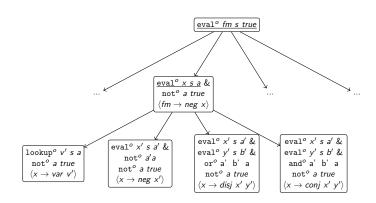
## Evaluator of Logic Formulas: Unfolding



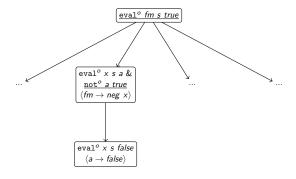
#### Boolean Connectives



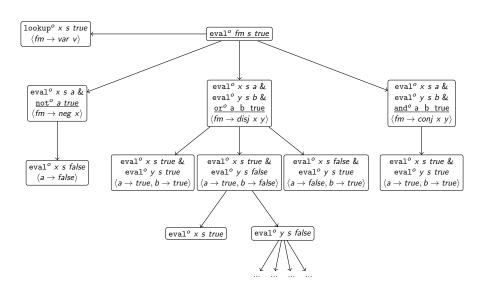
## Evaluator of Logic Formulas: Unfolding 2



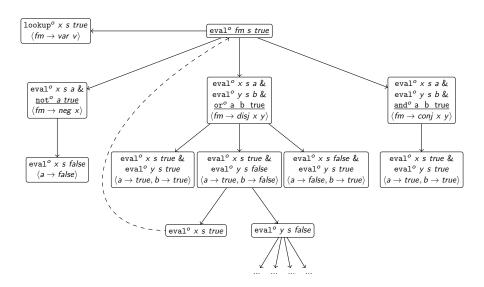
## Evaluator of Logic Formulas: Unfolding 3



## Evaluator of Logic Formulas: ConsPD

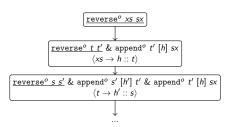


#### Evaluator of Logic Formulas: ConsPD

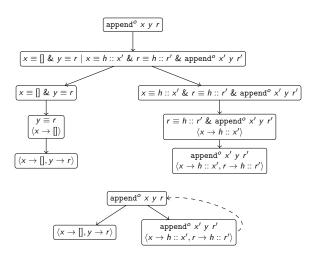


#### reverse<sup>o</sup>

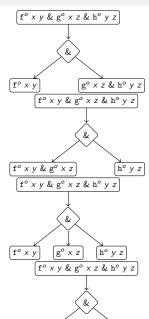
```
let rec reverse° xs sx =
  ocanren {
    (xs = [] & sx = []) |
    (fresh h t t' in
        xs = h :: t &
        reverse° t t' &
        append° t' [h] sx}
```



## Unfolding

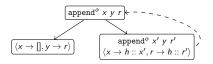


### Split



#### Conservative Partial Deduction

#### Branching Heuristics



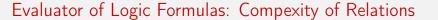


#### **Evaluation**

## Evaluator of Logic Formulas

## Evaluator of Logic Formulas: Order of Calls

:



## Evaluator of Logic Formulas: Results

#### Unification

#### Path Search

#### **Evaluation Results**

	last	plain	unify	isPath
Original	1.06s	1.84s	_	_
CPD	_	1.13s	14.12s	3.62s
ConsPD	0.93s	0.99s	0.96s	2.51s
Branching	3.11s	7.53s	3.53s	0.54s

Table: Evaluation results

#### Conclusion

- Conservative Partial Deduction
  - Less-branching heuristics
- Evaluation shows some improvement, but not for every query
- Models to predict performance can help