



# An Empirical Study of Partial Deduction for MINIKANREN

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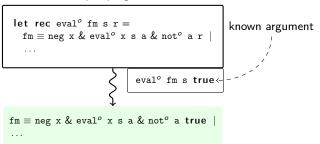
#### input program

```
let rec eval° fm s r = fm \equiv neg \ x \ \& \ eval^\circ \ x \ s \ a \ \& \ not^\circ \ a \ r \ | \dots
```

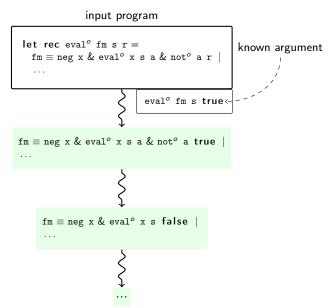
#### input program

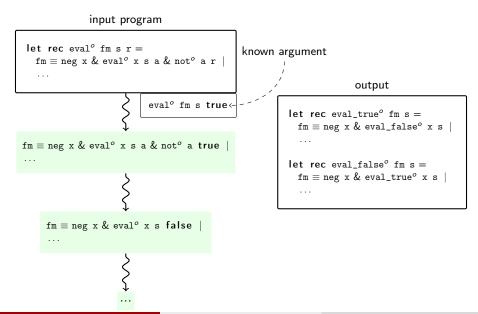
```
let rec eval° fm s r =  fm \equiv neg \ x \ \& \ eval° \ x \ s \ a \ \& \ not° \ a \ r \ |  ...  eval° \ fm \ s \ true \leftarrow
```

#### input program



# input program let rec eval<sup>o</sup> fm s r = known argument $fm \equiv neg x \& eval^o x s a \& not^o a r$ evalo fm s true < $fm \equiv neg x \& eval^o x s a \& not^o a true$ $fm \equiv neg x \& eval^o x s false$





## Partial Deduction: Specialization for Logic Programming

#### input

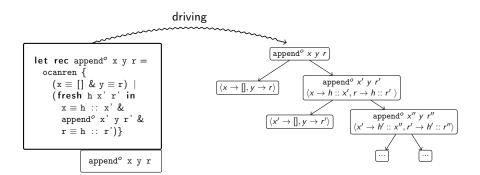
```
let double_appendo x y z r =
  ocanren {
    fresh t in
      appendo x y t &
      appendo t z r}
let rec appendo x y r =
  ocanren {
    (x \equiv [] \& y \equiv r) \mid
    (fresh h x' r' in
      x = h \cdot \cdot \cdot x' \&
      append° x' y r' &
      r \equiv h :: r')
```

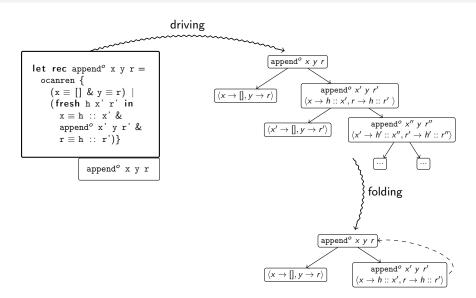
double\_append° x y z r

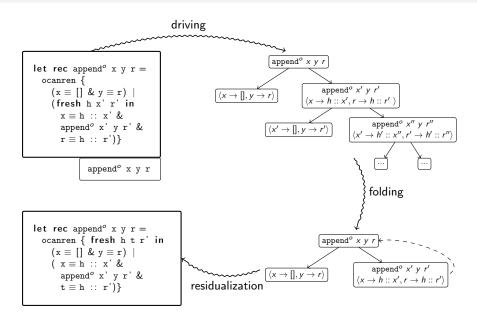
#### output

```
let double_appendo x y z r =
  ocanren {
    (x \equiv [] \& append^o y z r) |
    (fresh h x' r' in
      x \equiv h :: x' \&
      double_append° x' y z r' &
      r \equiv h :: r')
let rec appendo x y r =
  ocanren {
    (x \equiv [] \& y \equiv r) \mid
    (fresh h x' r' in
      x \equiv h :: x' \&
      appendo x'y r'&
      r \equiv h :: r')
```

```
let rec append° x y r =
  ocanren {
    (x = [] & y = r) |
    (fresh h x' r' in
      x = h :: x' &
      append° x' y r' &
    r = h :: r')}
```



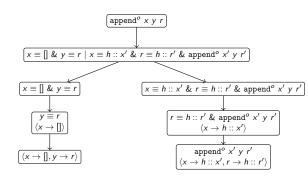




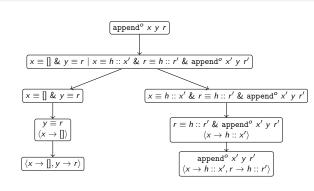
```
let rec append° x y r =
ocanren {
  (x = [] & y = r) |
  (fresh h x' r' in
  x = h :: x' &
  append° x' y r' &
  r = h :: r')}

append° x y r
```

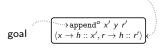
```
let rec append° x y r =
    ocarren {
        (x = [] & y = r) |
        (fresh h x' r' in
        x = h :: r' &
        append° x' y r' &
        r = h :: r')}
```



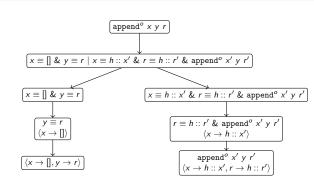
```
let rec append° x y r =
    ocanren {
        (x ≡ [] & y ≡ r) |
        (fresh h x' r' in
        x ≡ h :: x' &
        append° x' y r' &
        r ≡ h :: r')}
```

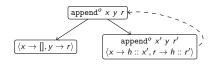


#### substitution



```
let rec append° x y r =
    ocanren {
        (x ≡ [] & y ≡ r) |
        (fresh h x' r' in
        x ≡ h :: x' &
        append° x' y r' &
        r ≡ h :: r')}
```

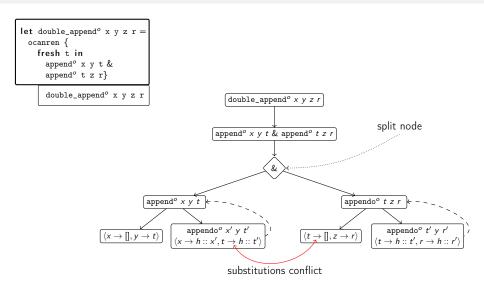




substitution

goal  $(x \to h :: x', r \to h :: r')$ 

#### Partial Deduction



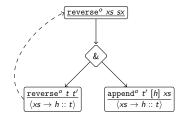
### Conjunctive Partial Deduction

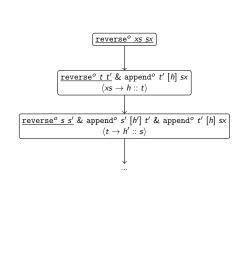
```
let double_appendo x y z r =
  ocanren {
      fresh t in
         appendo x v t &
         appendo t z r}
                                                                                   double_appendo x y z r
                                                  call to unfold
         double_appendo x y z r
                                                                             \sqrt{\text{append}^o \times y \ t \ \& \ \text{append}^o \ t \ z \ r}
                                                           appendo tzr
                                                                                                         append^{o} x' y t' & append^{o} (h :: t') z t'
                                                            (x \to [], y \to t)
                                                                                                                    \langle x \rightarrow h :: x', t \rightarrow h :: t' \rangle
                                                                          appendo t'zr'
                                                                                                             appendo x' v t' & appendo t' z r'
                                           \langle t \to [], z \to r \rangle
                                                                     \langle t \rightarrow h :: t', r \rightarrow h :: r' \rangle
                                                                                                                            \langle r \rightarrow h :: r' \rangle
```

```
let double_append° x y z r =
    ocanren {
        (x ≡ [] & append° y z r) |
        (fresh h x' r' in
        x ≡ h :: x' &
        double_append° x' y z r' &
        r ≡ h :: r')}
```

# CPD: Split is Necessary

```
let rec reverse<sup>o</sup> xs sx =
    ocanren {
    (xs = [] & sx = []) |
    (fresh h t t' in
            xs = h :: t &
            reverse<sup>o</sup> t t' &
            append<sup>o</sup> t' [h] sx}
```

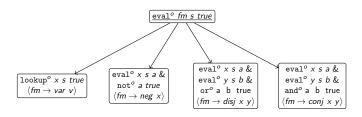




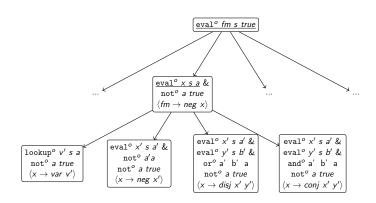
#### Decisions in Partial Deduction

- What to unfold: which calls, how many calls?
  - CPD: the leftmost call, which does not have a predecessor embedded into it
- How to unfold: to what depth a call should be unfolded?
  - CPD: unfold once
- When to stop driving?
  - When a goal is an instance of some goal in the process tree
- When to split?
  - When there is a predecessor embedded into the goal

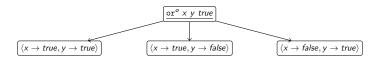
## Evaluator of Logic Formulas: Unfolding Step 1



## Evaluator of Logic Formulas: Unfolding Step 2

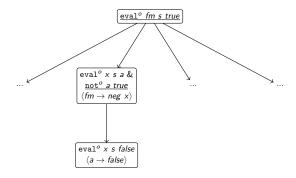


#### Unfolding of Boolean Connectives

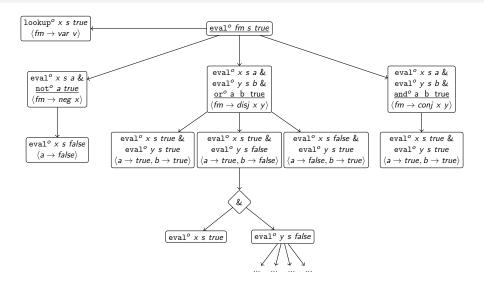




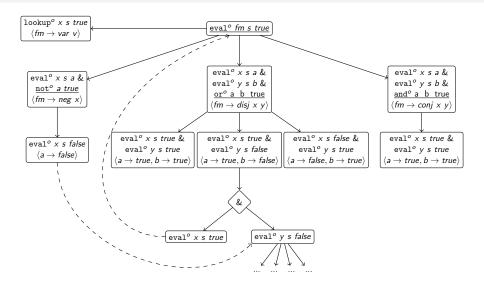
#### Unfolding Boolean Connectives First



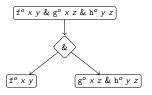
#### Evaluator of Logic Formulas: Conservative PD

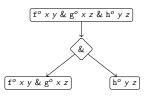


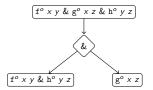
#### Evaluator of Logic Formulas: Conservative PD

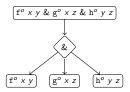


# Split: Which Way is the Right Way?









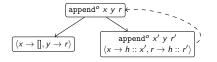
#### Conservative Partial Deduction

- Split conjunction into individual calls
- Unfold each call in isolation
- Unfold until embedding is encountered
- Find a call which narrows the search state (less-branching heuristics)
- Join the result of unfolding the selected call with the other calls not unfolded
- Continue driving the constucted conjunction

#### Less-branching Heuristics

Less-branching heuristics is used to select a call to unfold

If the call has less branches in the process tree than the relation can possible have, unfold the call





#### **Evaluation**

We implemented the Conservative Partial Deduction and compared it with CPD for  $\mbox{\scriptsize MINIKANREN}$  and CPD with branching heuristics on the following relations

- Two implementations of an evaluator of logic formulas
- A program to compute a unifier of two terms
- A program to search for paths of a specific length in a graph

## Evaluator of Logic Formulas

```
let rec eval° fm s r =
  ocanren { fresh v x y a b in
    (fm = var v & lookup° v s r) |
    (fm = neg x & eval° x s a & not° a r) |
    (fm = conj x y & eval° x s a & eval° y s b & and° a b r) |
    (fm = disj x y & eval° x s a & eval° y s b & oro° a b r)
```

## Evaluator of Logic Formulas: Order of Calls

```
let rec eval<sup>o</sup> fm s r =
  ocanren { fresh v x y a b in
    (fm \equiv var v \& lookup^o v s r) \mid
    (fm \equiv neg x \& not^o a r \& eval^o x s a)
    (fm \equiv conj x y \& and^o a b r \& eval^o x s a \& eval^o y s b)
    (fm \equiv disj \times y \& oro^o a b r \& eval^o \times s a \& eval^o y s b)
let rec evalo fm s r =
  ocanren { fresh v x y a b in
    (fm \equiv var v \& lookup^o v s r)
    (fm \equiv neg x \& eval^o x s a \& not^o a r)
    (fm \equiv conj x y \& eval^o x s a \& eval^o y s b \& and^o a b r)
    (fm \equiv disj \times y \& eval^o \times s a \& eval^o y s b \& oro^o a b r)
```

#### Evaluator of Logic Formulas: Compexity of Relations

```
let rec or ^{o} x y r =
  ocanren {
     (x \equiv true \& y \equiv true \& r \equiv true)
     (x \equiv true \& y \equiv false \& r \equiv true)
     (x \equiv false \& y \equiv true \& r \equiv true)
     (x \equiv false \& y \equiv false \& r \equiv false) 
let or o x y r =
  ocanren {
     fresh a b in
       (nand^o \times \times a \& nand^o \vee \vee b \& nand^o a b r)
let rec nando x v r =
  ocanren {
     (x \equiv true \& y \equiv true \& r \equiv false)
     (x \equiv true \& y \equiv false \& r \equiv true)
     (x \equiv false \& y \equiv true \& r \equiv true)
     (x \equiv false \& y \equiv false \& r \equiv false) }
```

### Evaluator of Logic Formulas: Evaluation

#### Implementations:

- last: boolean connectives last, implemented via nand<sup>o</sup>
- plain: boolean connectives first, straightforward implementation

	last	plain
Original	1.06s	1.84s
CPD	_	1.13s
ConsPD	0.93s	0.99s
Branching	3.11s	7.53s

Table: Evaluation results

#### Unification

Relation to find a unifier of two terms

Query: unification of terms f(X, X, g(Z, t)) and f(g(p, L), Y, Y).

#### Path Search

Relation to search for paths in a graph

Query: find 5 paths in a graph with 20 vertices and 30 edges

#### **Evaluation Results**

	last	plain	unify	isPath
Original	1.06s	1.84s	_	_
CPD	_	1.13s	14.12s	3.62s
ConsPD	0.93s	0.99s	0.96s	2.51s
Branching	3.11s	7.53s	3.53s	0.54s

Table: Evaluation results

#### Conclusion

- We developed and implemented Conservative Partial Deduction
  - Less-branching heuristics
- Evaluation shows some improvement, but not for every query
- Future work:
  - Develop models to predict execution time
  - Develop specialization which is more predictable, stable and well-behaved