Stochastic Simulation Generation of random variables Discrete sample space

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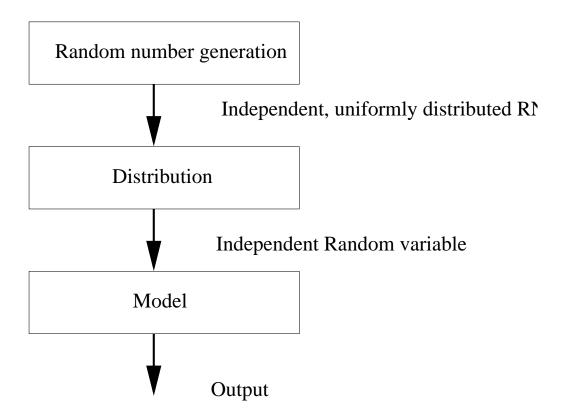
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Random variables



Aim

- The scope is the generation of **independent** random variables $X_1, X_2, ... X_n$ with a **given distribution**, $F_x(x)$, (or probability density function [pdf]).
- We assume we have access to a supply (U_i) of random numbers, independent samples from the uniform distribution on]0, 1[.
- Our task is to transform U_i into X_i .

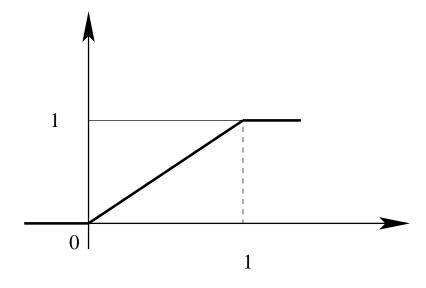
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Uniform distribution I



Our norm distribution or building brick, U(0,1)

$$f(x) = 1$$
 $F(x) = x$ for $0 \le x \le 1$

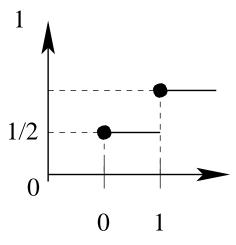


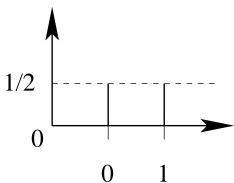
$$\mathbb{E}(X) = \frac{1}{2} \quad \mathbb{V}(X) = \frac{1}{12}$$

Coin



or uniform distribution







$$X = 0, 1$$

$$P(X=i) = \frac{1}{2}$$

$$X := \left(U > \frac{1}{2}\right) \quad X = \lfloor (2U) \rfloor$$

Bernoulli trial

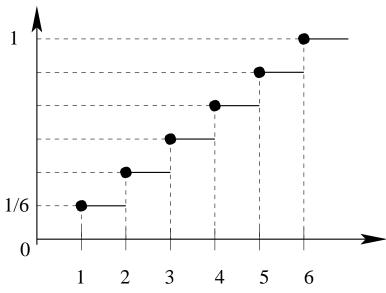


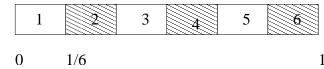
Toss a coin with P(X = 1) = p and P(X = 0) = 1 - p.

A fair die



or uniform distribution





$$X = 1, 2, \dots 6$$

$$P(X=i) = 1/6$$

$$X = \lfloor (6U) \rfloor + 1$$

Can be generalized $6 \rightarrow n$.

Discrete distribution - direct (crude) method

Suppose X can take n distinct values $x_1 < x_2 < \dots x_n$ with

$$p_i = P(X = x_i), \quad i = 1, 2, \dots, n$$

Then X takes the value x_i with probability p_i if U falls in an interval with length p_i . That is if

$$\sum_{j=1}^{i-1} p_j < U \le \sum_{j=1}^{i} p_j$$

or

$$X = x_i$$
 if $F(x_{i-1}) < U \le F(x_i)$

Geometric distribution, NB(1,p)

The discrete time version of waiting time. Memory-less.



$$f(n) = P(X = n) = (1 - p)^{n-1}p \quad n = 1, 2, \dots$$

$$F(n) = P(X \le n) = 1 - (1 - p)^n$$

$$X = n$$
 if $F(n-1) < U \le F(n)$ $1 - (1-p)^{n-1} < U \le 1 - (1-p)^n$

$$n - 1 < \frac{\log(1 - U)}{\log(1 - p)} \le n$$

$$X = \left\lfloor \left(\frac{\log(U)}{\log(1-p)} \right) \right\rfloor + 1$$

Discrete distribution II





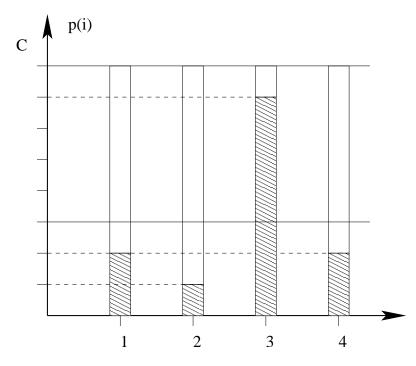
- 0 P1 P2
 - 1. Generate U
 - 2. Find the interval k which U belong to. $P_{k-1} < U \le P_k$
 - 3. output x_k
 - Linear search ($\mathbb{E}(X)$)
 - Rearrangement of intervals
 - Binary search
 - Indexed search

Rejection Method



Simple rejection More optimistic: acceptance method.

Assume $P(X = i) = p_i$ for $i = 1, 2, \dots n$.



Let
$$c > p_i$$
 (then $p_i/c < 1$).

1.
$$I = 1 + \lfloor (n * U_1) \rfloor$$

2. output I if $U_2 < p_I/c$.

frequency for
$$i$$
:
$$\frac{\frac{1}{n}\frac{p_i}{c}}{\sum \frac{1}{n}\frac{p_i}{c}} = p_i$$

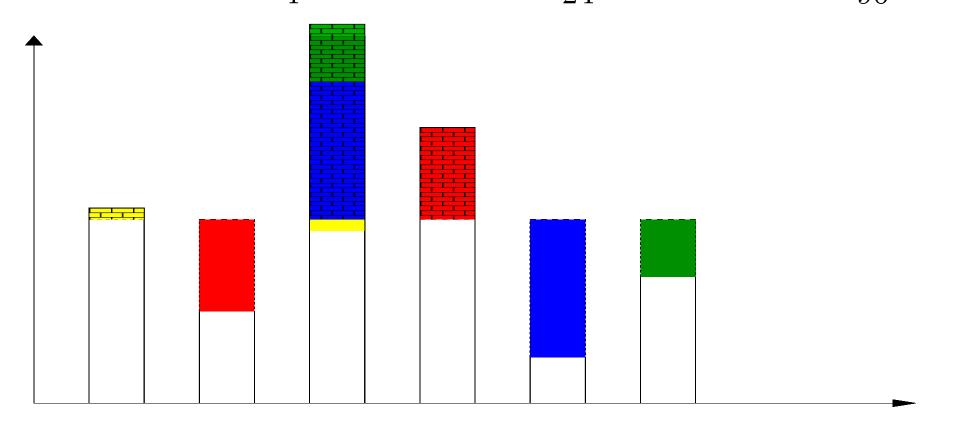
Alias method



- A method for generating discrete random variates of general type
- From discrete uniform to general discrete
- Generate one random number
- One comparison
- Potentially one table look-up
- Drawback: Complex set-up procedure

A six-point distribution

$$P\{X=1\} = \frac{17}{96} \qquad P\{X=2\} = \frac{1}{12} \qquad P\{X=3\} = \frac{1}{3} \implies P\{X=4\} = \frac{1}{4} \qquad P\{X=5\} = \frac{1}{24} \qquad P\{X=6\} = \frac{11}{96}$$



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Alias method

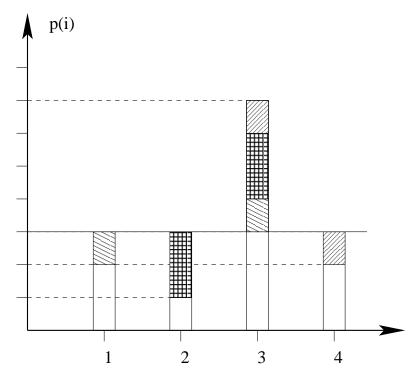


- Setup procedure
 - \diamond Generate the table of F(I)-values, (which part of the mass belongs to I itself.
 - \diamond Generate the table of L(I)-values, (the alias of class I
- Method at run time
 - \diamond Generate $I:I = |n * U_1|$
 - ♦ Test against F(I).If $U_2 \le F(I)$ then return X = I else return X = L(I). The L, F tables for the six-point distribution

$$F(1) = 1$$
 $F(2) = \frac{1}{2}$ $F(3) = \frac{15}{16}$ $F(4) = 1$ $F(5) = \frac{1}{4}$ $F(6) = \frac{11}{16}$ $L(1) = 1$ $L(2) = 4$ $L(3) = 1$ $L(4) = 4$ $L(5) = 3$ $L(6) = 3$

Alias Method





On setup: generate F and L.

- 1. $I = 1 + \lfloor (n * U_1) \rfloor$
- 2. output I if $U_2 < F(I)$ else E(I).

The Alias tables



Generate F and L.

Pseudo code. p is a vector containing the probabilities.

- 1. F=n*p (F=1 is equivalent for the uniform dist.)
- 2. G=find(F>=1) and S=find(F<=1)
- 3. while ~isempty(S),
 - (a) k=G(1) and j=S(1)
 - (b) L(j)=k and F(k)=F(k)-(1-F(j))
 - (c) if F(k)<1-eps then G(1)=[] and S=[S k]
 - (d) S(1) = []

Rejection Method

General method



Aim: We will generate X with probabilities $p_i = P\{X = i\}$.

Assume Y with probabilities $q_i = P\{Y = i\}$ is easily generated and $C > \frac{p_i}{q_i}$ for all $i = 1, \dots n$.

- 1. Generate Y with probability q_i and let j = Y.
- 2. Generate U_2 . Output X=Y if $U_2<\frac{p_j}{Cq_j}$ else reject.

Probability for X = i:

Frequency of output:

$$q_i \frac{p_i}{Cq_i} = \frac{p_i}{C}$$
$$\frac{p_i/C}{\sum p_i/C} = p_i$$

Excercise 2

Discrete random variables

In the excercise you can use a build in procedure for generating uniform random variables. Compare the results obtained in simulations with expected results. Use histograms (and tests).

- Choice a value for p in the geometric distribution and simulate 10,000 outcomes.
- Simulate the 6 point distribution with

X	1	2	3	4	5	6
p_i	7/48	5/48	1/8	1/16	1/4	5/16

- by applying a direct (crude) method
- by using the the rejction method
- by using the Alias mehtod