CS 475 Parallel Computing Instructor: Matthew Meyn Student: **Khuong Luu (luukh)**

Project 6

1. What machine did you run the programs on

Due heavy loads on the rabbit server, eventually I decided to install everything on my local computer and run it by myself.

- Processor: Intel(R) Core(TM) i7-7700HQ CPU @ 2.80GHz

- Memory: 16GB, 2400Mhz

- Operating System: Ubuntu 16.04

Kernel: 4.13.0-37-genericNumber of Cores: 4Cache line size: 64 bytes

- Compiler: gcc

\$ clinfo

Number of platforms 1

Platform Name NVIDIA CUDA

Platform Vendor NVIDIA Corporation

Platform Version OpenCL 1.2 CUDA 9.0.368

Platform Profile FULL PROFILE

Platform Extensions cl_khr_global_int32_base_atomics cl_khr_global_int32_extended_atomics cl_khr_local_int32_base_atomics

 $cl_khr_local_int32_extended_atomics\ cl_khr_fp64\ cl_khr_byte_addressable_store\ cl_khr_icd\ cl_khr_gl_sharing\ cl_nv_compiler_options\ cl_nv_device_attribute_query\ cl_nv_pragma_unroll\ cl_nv$

cl_nv_copy_opts cl_nv_create_buffer

Platform Extensions function suffix NV

Platform Name NVIDIA CUDA

Number of devices 1

Device Name GeForce GTX 1060

Device Vendor NVIDIA Corporation

Device Vendor ID 0x10de

Device Version OpenCL 1.2 CUDA

Driver Version 384.130

Device OpenCL C Version OpenCL C 1.2

Device Type GPU

I use OpenCL version 1.1 from Kronos

```
#define CL_TARGET_OPENCL_VERSION 120
#include <CL/cl.h>
```

2. Multiply and Multiply-Add

2.1. Tables

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LOCAL_SIZE	1024	10240	102400	1024000	10240000	40960000	81920000
1	0.035	0.172	0.283	0.275	0.3	0.302	0.369
2	0.04	0.276	0.696	0.66	0.819	0.835	0.838
4	0.039	0.328	1.199	1.088	1.543	1.656	1.357
8	0.037	0.356	1.73	1.526	2.683	2.846	2.885
16	0.043	0.374	2.508	1.941	5.005	5.794	5.938
32	0.041	0.393	3.147	2.444	8.308	10.066	10.646
64	0.042	0.391	2.292	2.436	8.981	11.377	11.845
128	0.039	0.409	2.264	2.434	9.105	11.312	11.839
256	0.041	0.417	3.604	2.6	8.973	11.172	11.883
512	0.042	0.409	3.419	2.64	8.951	11.348	11.872
1024	0.042	0.403	3.342	2.533	8.965	11.297	11.814

Multiply-Add program

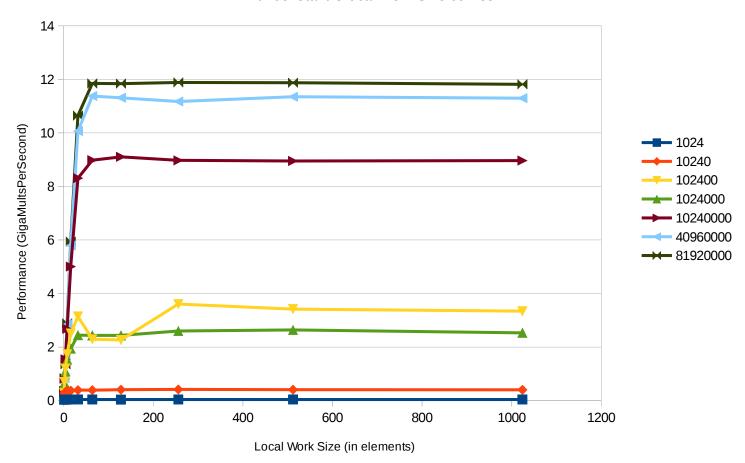
P 8							
LOCAL_SIZE	1024	10240	102400	1024000	10240000	40960000	81920000
1	0.037	0.207	0.38	0.367	0.414	0.414	0.415
2	0.039	0.274	0.699	0.666	0.818	0.83	0.837
4	0.039	0.328	1.188	1.102	1.584	1.651	1.662
8	0.039	0.361	1.836	1.508	2.795	3.042	3.076
16	0.041	0.367	2.537	2.058	4.987	5.588	5.761
32	0.04	0.262	2.996	2.395	6.903	8.262	8.492
64	0.041	0.401	3.254	2.437	6.811	8.639	9.166
128	0.039	0.402	3.303	2.433	7.269	8.785	9.141
256	0.036	0.4	3.445	2.364	7.272	8.69	9.12
512	0.042	0.414	3.372	2.272	7.347	8.779	9.131
1024	0.042	0.413	3.306	2.312	7.27	8.733	9.109

2.2. Graphs

Multiply program

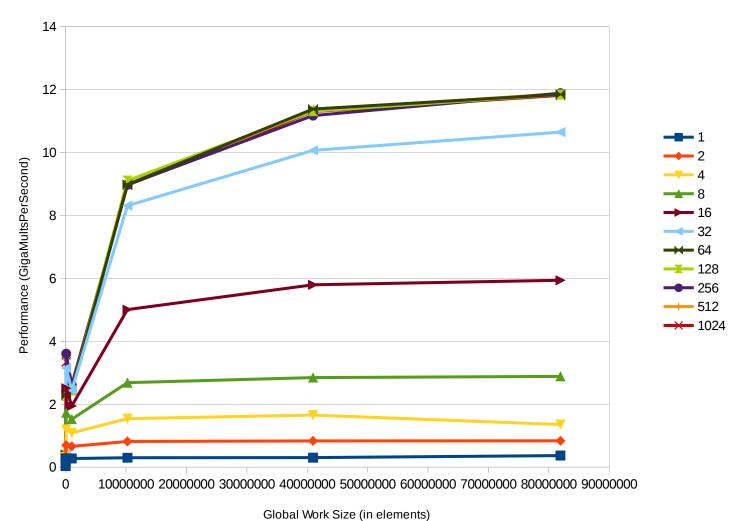
Performance vs. Local Work Size of the Multiply program

with constant Global Work Size curves



Performance vs. Global Work Size of the Multiply program

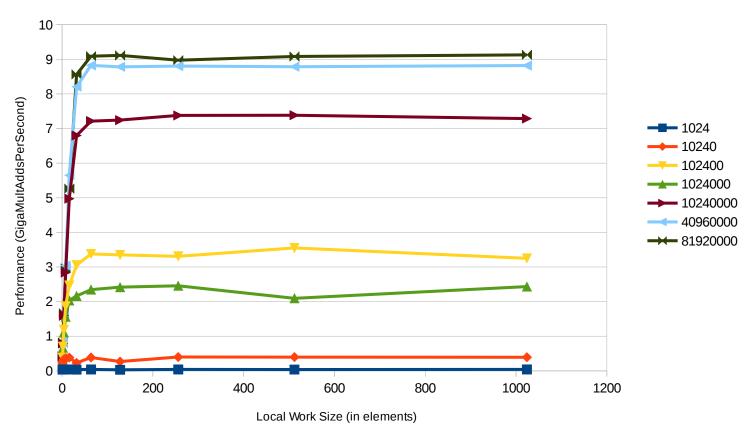
with constant Local Work Size curves



Multiply-Add Program

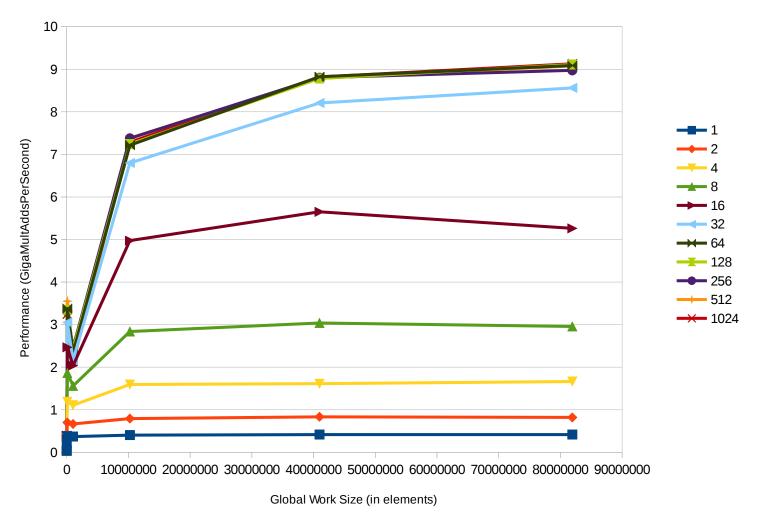
Performance vs. Local Work Size of the Multiply-Add program

with constant Global Work Size curves



Performance vs. Global Work Size of the Multiply-Add program

with constant Local Work Size curves



3. Pattern and Explanation for Multiply and Multiply-Add program

3.1. Patterns in each of the 2 programs

Pattern 1: Given the constant local sizes, the bigger the global size, the significantly better the performance until the performance converges

Explanation: Given the constant local size, the bigger the global size leads to significantly more work groups being used, which in turn increase the parallelization of the program, which in turn certainly increase the performance significantly unlike in pattern 2 which not necessary increase the parallelization. until we reach the optimal number of work groups of the hardware. This also explain why increase in global work size is more effective than increase in local work size.

Pattern 2: Given the constant global size, increase in local size leads to increase in performance but only for the increases of smalls local sizes (1, 2, 4, 8, 16, 32, 64). After that, the performance doesn't increase much if not to say not at all since the local size of 64.

Explanation: This is hard to explain but rather it tell us a meaningful and useful conclusion that the optimal value for local size for this kind of task and hardware is 64 (times the size of a float)

3.2. Different between the Multiply program and Mutiply-Add program

Difference: The performance of the Multiply program is higher than that of the Multiply-Add program. Given the same local size and global size, the performance of the Multiply program is of 3 Giga-Multiplies per Second bigger than the performance of the Multiply-Add program.

Explanation: because the kernel of the later has to do (one) more task compared to the former: Adding.

3.3. What does that mean for the proper use of GPU parallel computing?

From the patterns and their explanation above, we have learned that increasing the global work size is more effective than increasing the local work size; and there is a optimal value for local work size that we need to find out in order to allocate the local size that use the GPU the most effective way.

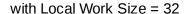
4. Benchmark and Analyze Multiply Reduce program

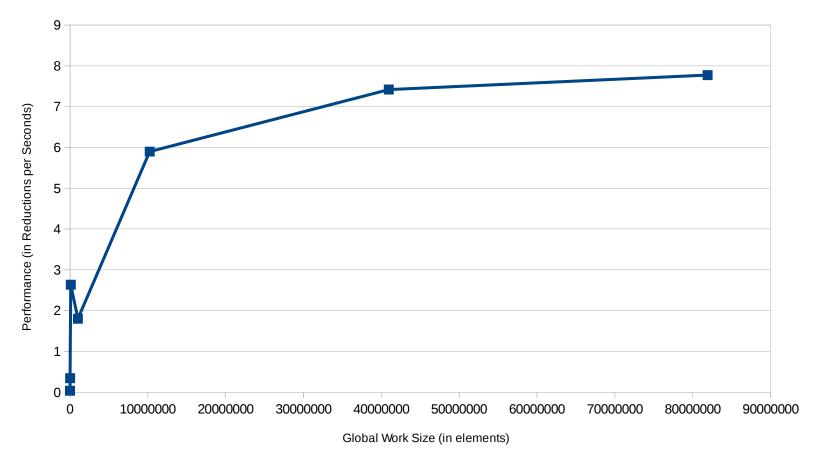
4.1. Table

Global Work Size	1024	10240	102400	1024000	10240000	40960000	81920000
Performance	0.036	0.344	2.635	1.802	5.897	7.418	7.772

4.2. Graph

Performance vs. Global Work Size of the Multiply-Reduce program





4.3. Patterns and Explanations

Pattern 1: Increase in global work size leads to increase in performance, until the performance start converging at global work size of 80M.

Explanation: This is the same explanation as with the Multiply program: The increase in Global Work Size leads to more number of work groups being use and thus increase the parallelization of the program, which certainly leads to better performance until we reach the optimal number of work groups.

Pattern 2: The performance of this program is less than the performance of the Multiply and the Multiply-Add programs

Explanation: This is because the kernel of this program has to perform more work than the kernel in the other two programs.

4.4. What does this mean for proper use of GPU computing?

The fact that more busy kernel code has the same performance pattern as the less busy kernel code gives us a lesson that we should do as much work in each work-groups as possible before passing data to/from the host. Data passing is expensive.