

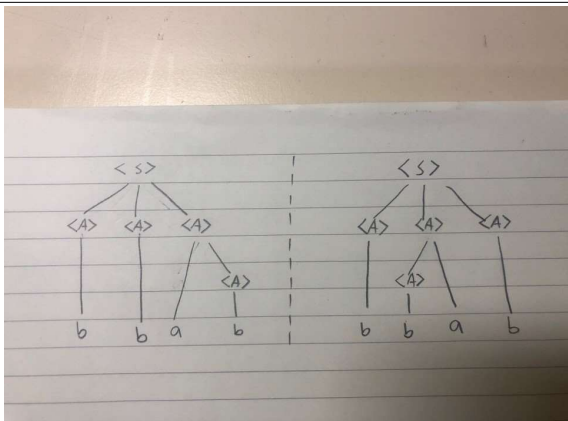
프로그래밍 언어론 HW1		
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문제	1.Alphabet set{0,1}상에서 다음 regular language를 정의하는 regular expression은 무엇인가.
해답a	(a)01로 끝나는 모든 string (0 1)*01은 001101, 010101 등등 모든 01로 끝나는 string을 표현할 수 있다.
해답b	(b)0이 짝수 개 포함되어 있는 string ((1)*0(1)*0(1)*)* 은 1001001,1111100,11010001 등 0이 짝수 개 포함되어 있는 string을 표현할 수 있다.

2. 다음 BNF에서 A가 start symbol일 때 $d=a+b*c^d$ 를 LeftMostDerivation과 RightMostDerivation을 이용하여 생성하시오.	
$\langle A \rangle \rightarrow \langle L \rangle = \langle E \rangle$ $\langle E \rangle \rightarrow \langle E \rangle + \langle T \rangle \mid \langle E \rangle - \langle T \rangle \mid \langle T \rangle$ $\langle T \rangle \rightarrow \langle T \rangle * \langle S \rangle \mid \langle T \rangle / \langle S \rangle \mid \langle S \rangle$ $\langle S \rangle \rightarrow \langle F \rangle ^ \langle S \rangle \mid \langle F \rangle$ $\langle F \rangle \rightarrow (\langle E \rangle) \mid \langle L \rangle$ $\langle L \rangle \rightarrow a \mid b \mid c \mid d$	
LeftMostDerivation $\langle A \rangle \rightarrow \langle L \rangle = \langle E \rangle$ $\langle A \rangle \rightarrow d = \langle E \rangle$ $\langle A \rangle \rightarrow d = \langle E \rangle + \langle T \rangle$ $\langle A \rangle \rightarrow d = \langle T \rangle + \langle T \rangle$ $\langle A \rangle \rightarrow d = \langle S \rangle + \langle T \rangle$ $\langle A \rangle \rightarrow d = \langle F \rangle + \langle T \rangle$ $\langle A \rangle \rightarrow d = \langle L \rangle + \langle T \rangle$ $\langle A \rangle \rightarrow d = a + \langle T \rangle$ $\langle A \rangle \rightarrow d = a + \langle T \rangle * \langle S \rangle$ $\langle A \rangle \rightarrow d = a + \langle S \rangle * \langle S \rangle$ $\langle A \rangle \rightarrow d = a + \langle F \rangle * \langle S \rangle$ $\langle A \rangle \rightarrow d = a + \langle L \rangle * \langle S \rangle$ $\langle A \rangle \rightarrow d = a + b * \langle S \rangle$ $\langle A \rangle \rightarrow d = a + b * \langle F \rangle ^ \langle S \rangle$ $\langle A \rangle \rightarrow d = a + b * \langle L \rangle ^ \langle S \rangle$	RightMostDerivation $\langle A \rangle \rightarrow \langle L \rangle = \langle E \rangle$ $\langle A \rangle \rightarrow \langle L \rangle = \langle E \rangle + \langle T \rangle$ $\langle A \rangle \rightarrow \langle L \rangle = \langle E \rangle + \langle T \rangle * \langle S \rangle$ $\langle A \rangle \rightarrow \langle L \rangle = \langle E \rangle + \langle T \rangle * \langle F \rangle ^ \langle S \rangle$ $\langle A \rangle \rightarrow \langle L \rangle = \langle E \rangle + \langle T \rangle * \langle F \rangle ^ \langle F \rangle$ $\langle A \rangle \rightarrow \langle L \rangle = \langle E \rangle + \langle T \rangle * \langle F \rangle ^ \langle L \rangle$ $\langle A \rangle \rightarrow \langle L \rangle = \langle E \rangle + \langle T \rangle * \langle F \rangle ^ d$ $\langle A \rangle \rightarrow \langle L \rangle = \langle E \rangle + \langle T \rangle * \langle L \rangle ^ d$ $\langle A \rangle \rightarrow \langle L \rangle = \langle E \rangle + \langle T \rangle * c ^ d$ $\langle A \rangle \rightarrow \langle L \rangle = \langle E \rangle + \langle S \rangle * c ^ d$ $\langle A \rangle \rightarrow \langle L \rangle = \langle E \rangle + \langle F \rangle * c ^ d$ $\langle A \rangle \rightarrow \langle L \rangle = \langle E \rangle + \langle L \rangle * c ^ d$ $\langle A \rangle \rightarrow \langle L \rangle = \langle E \rangle + b * c ^ d$ $\langle A \rangle \rightarrow \langle L \rangle = \langle T \rangle + b * c ^ d$ $\langle A \rangle \rightarrow \langle L \rangle = \langle S \rangle + b * c ^ d$

$\langle A \rangle \rightarrow d = a + b * c^{\langle S \rangle}$	$\langle A \rangle \rightarrow \langle L \rangle = \langle F \rangle + b * c^d$
$\langle A \rangle \rightarrow d = a + b * c^{\langle F \rangle}$	$\langle A \rangle \rightarrow \langle L \rangle = \langle L \rangle + b * c^d$
$\langle A \rangle \rightarrow d = a + b * c^{\langle L \rangle}$	$\langle A \rangle \rightarrow \langle L \rangle = a + b * c^d$
$\langle A \rangle \rightarrow d = a + b * c^d$	$\langle A \rangle \rightarrow d = a + b * c^d$

문제	3. 다음 EBNF를 BNF로 변환하시오. $\langle S \rangle \rightarrow \langle A \rangle \{ b \langle A \rangle \}$ $\langle A \rangle \rightarrow a [b] \langle A \rangle$
해설	[] \rightarrow 있어도 되고 없어도 된다 {} \rightarrow 0번이상 반복된다 BNF로 변환하면 $\langle S \rangle \rightarrow \langle A \rangle \mid \langle A \rangle \langle s_bA \rangle$ $\langle s_bA \rangle \rightarrow b \langle A \rangle \mid b \langle A \rangle \langle s_bA \rangle$ $\langle A \rangle \rightarrow a \langle A \rangle \mid ab \langle A \rangle$

문제	4. 다음 grammar가 ambiguous하다는 것을 보이시오(startsymbol:S).
해설	 <p>위 그림과 같이 같은 sentence에 따라 서로 다른 parse tree가 존재하므로 ambiguous한 grammar입니다.</p>

문제	5. 다음과 같은 문법이 주어졌을 때, 각각의 non-terminal들에 대해 pairwise disjointness test를수행하시오(S=startsymbol). $\langle S \rangle \rightarrow a \langle S \rangle b \mid \langle B \rangle \langle A \rangle \langle A \rangle$ $\langle A \rangle \rightarrow b \{ a \langle B \rangle \} \mid a$ $\langle B \rangle \rightarrow a \langle B \rangle \mid c$
해설	S에 대한 pairwise disjointness test

	<p> $\text{First}(a \langle S \rangle b) = \{a\}$ $\text{First}(\langle B \rangle \langle A \rangle \langle A \rangle) = \{a, c\}$ 따라서 pairwise disjointness test를 통과하지 못한다 </p> <p> A에 대한 pairwise disjointness test $\text{First}(b\{a \langle B \rangle\}) = \{b\}$ $\text{First}(a) = \{a\}$ 따라서 pairwise disjointness test를 통과한다 </p> <p> B에 대한 pairwise disjointness test $\text{First}(a \langle B \rangle) = \{a\}$ $\text{First}(c) = \{c\}$ 따라서 pairwise disjointness test를 통과한다 </p>
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