$Tools\ for\ primal\ degenerate\ linear\ programs$

 $Dual\text{-}guided\ pivot\ rules\ for\ LP$

Vector Space Decomposition for LP

J. Desrosiers¹

J.B. Gauthier¹ M.E. Lübbecke²

(1) HEC Montréal and GERAD

(2) RWTH Aachen University

Congrès de la ROADEF, Bordeaux 26 au 28 février 2014

• PS : Primal Simplex and degenerate solutions perturbation, pivot rules, etc. (Terlaky and Zhang 1993)

- PS: Primal Simplex and degenerate solutions perturbation, pivot rules, etc. (Terlaky and Zhang 1993)
- \circ CG : Column Generation for *huge* problems \circ CG \equiv PS perturbation of the master problem (1985-2000)

- PS: Primal Simplex and degenerate solutions perturbation, pivot rules, etc. (Terlaky and Zhang 1993)
- CG : Column Generation for *huge* problems $CG \equiv PS$ perturbation of the master problem (1985-2000)
- DVS: Dual Variable Stabilization (du Merle et al. 1999)
 Oukil et al (2007): cpu time reduction by a factor of 100

- PS: Primal Simplex and degenerate solutions perturbation, pivot rules, etc. (Terlaky and Zhang 1993)
- \bullet CG : Column Generation for *huge* problems \bullet CG \equiv PS perturbation of the master problem (1985-2000)
- DVS: Dual Variable Stabilization (du Merle et al. 1999)
 Oukil et al (2007): cpu time reduction by a factor of 100
- DCA: Dynamic Constraints Aggregation for Set Partitioning Problems (Elhallaoui, Metrane, Desaulniers, Soumis 2005-08)
 cpu time reduction by a factor of 100

- PS: Primal Simplex and degenerate solutions perturbation, pivot rules, etc. (Terlaky and Zhang 1993)
- CG : Column Generation for huge problems $CG \equiv PS$ perturbation of the master problem (1985-2000)
- DVS: Dual Variable Stabilization (du Merle et al. 1999)
 Oukil et al (2007): cpu time reduction by a factor of 100
- DCA: Dynamic Constraints Aggregation for Set Partitioning Problems (Elhallaoui, Metrane, Desaulniers, Soumis 2005-08)
 cpu time reduction by a factor of 100
- IPS : Improved Primal Simplex (Vincent, Soumis, Metrane, Orban 2008-11) generalizes DCA to LP; non-degenerate pivots; positive edge rule

- PS: Primal Simplex and degenerate solutions perturbation, pivot rules, etc. (Terlaky and Zhang 1993)
- \bullet CG : Column Generation for *huge* problems \bullet CG \equiv PS perturbation of the master problem (1985-2000)
- DVS: Dual Variable Stabilization (du Merle et al. 1999)
 Oukil et al (2007): cpu time reduction by a factor of 100
- DCA: Dynamic Constraints Aggregation for Set Partitioning Problems (Elhallaoui, Metrane, Desaulniers, Soumis 2005-08)
 cpu time reduction by a factor of 100
- IPS : Improved Primal Simplex (Vincent, Soumis, Metrane, Orban 2008-11) generalizes DCA to LP; non-degenerate pivots; positive edge rule
- MMCC: Minimum Mean Cost-cycle Canceling algorithm for networks 70%–90% of the PS pivots are degenerate; MMCC strongly polynomial.

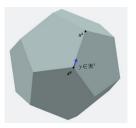
- PS: Primal Simplex and degenerate solutions perturbation, pivot rules, etc. (Terlaky and Zhang 1993)
- \bullet CG : Column Generation for *huge* problems \bullet CG \equiv PS perturbation of the master problem (1985-2000)
- DVS: Dual Variable Stabilization (du Merle et al. 1999)
 Oukil et al (2007): cpu time reduction by a factor of 100
- DCA: Dynamic Constraints Aggregation for Set Partitioning Problems (Elhallaoui, Metrane, Desaulniers, Soumis 2005-08)
 cpu time reduction by a factor of 100
- IPS : Improved Primal Simplex (Vincent, Soumis, Metrane, Orban 2008-11) generalizes DCA to LP; non-degenerate pivots; positive edge rule
- MMCC: Minimum Mean Cost-cycle Canceling algorithm for networks 70%–90% of the PS pivots are degenerate; MMCC strongly polynomial.

What are the links between PS, DCA, IPS, and MMCC?

Observation #1

LP IN STANDARD FORM

$$\begin{array}{rcl} \min & \mathbf{c}^\mathsf{T} \mathbf{x} \\ & \mathbf{A} \mathbf{x} &= \mathbf{b} \\ & \mathbf{x} & \geq \mathbf{0} \end{array} \quad [\pi]$$



From x^0 to x^1

- Find a *potential* improving direction $\mathbf{y}^0 \in \mathbb{R}^n$.
- **2** Determine step-size $\rho^0 \in \mathbb{R}$.
- **3** Compute $x^1 := x^0 + \rho^0 y^0$.

Primal/dual formulations of the pricing in Primal Simplex

$\overline{c_j} = 0, \forall j \in B$: Pricing for $j \in N$ (non-basic variables)

Selection of an entering variable into basis \mathbf{A}_B relies on the minimum reduced cost of non-basic variables

$$ar{c}_j = c_j - oldsymbol{\pi}^\intercal oldsymbol{a}_j, \ \ \forall j \in \mathcal{N}. \qquad oldsymbol{\pi}^\intercal = oldsymbol{c}_B^\intercal oldsymbol{A}_B^{-1}$$

PRIMAL/DUAL FORMULATIONS OF THE PRICING IN PRIMAL SIMPLEX

$\bar{c}_i = 0, \forall j \in B$: Pricing for $j \in N$ (non-basic variables)

Selection of an entering variable into basis \mathbf{A}_B relies on the minimum reduced cost of non-basic variables $\bar{c}_j = c_j - \boldsymbol{\pi}^\mathsf{T} \mathbf{a}_j, \ \forall j \in N. \qquad \boldsymbol{\pi}^\mathsf{T} = \mathbf{c}_B^\mathsf{T} \mathbf{A}_B^{-1}$

Find the minimum reduced cost value μ (Optimal if $\mu \geq 0$)

*** μ is the smallest reduced cost (given π). ***

PRIMAL/DUAL FORMULATIONS OF THE PRICING IN PRIMAL SIMPLEX

$\overline{c_i} = 0, \forall j \in B$: Pricing for $j \in N$ (non-basic variables)

Selection of an entering variable into basis $\mathbf{A}_{\mathcal{B}}$ relies on the minimum reduced cost of non-basic variables $\bar{c}_i = c_i - \boldsymbol{\pi}^\mathsf{T} \mathbf{a}_i, \ \forall j \in \mathcal{N}. \quad \boldsymbol{\pi}^\mathsf{T} = \mathbf{c}_R^\mathsf{T} \mathbf{A}_R^{-1}$

Find the minimum reduced cost value μ (Optimal if $\mu \geq 0$)

*** μ is the smallest reduced cost (given π). ***

Equivalent to finding a convex combination of non-basic variables

$$\begin{array}{ccccc} \mu = & \min & \sum_{j \in \mathcal{N}} \bar{c}_j y_j \\ & & \sum_{j \in \mathcal{N}} y_j & = & 1 & & [\mu] \\ & & & y_j & \geq & 0, & \forall j \in \mathcal{N} \end{array}$$

Direction $\mathbf{y}^0 \in \mathbb{R}^n$

 \mathbf{y}^{0}

Entering variable only impacts the current basic variables, at most m variables. All but one non-basic variables remain at zero.

 $\mathbf{y}^0 \in \mathbb{R}^n$ makes changes on at most m+1 components.

Direction $\mathbf{y}^0 \in \mathbb{R}^n$

\mathbf{y}^{0}

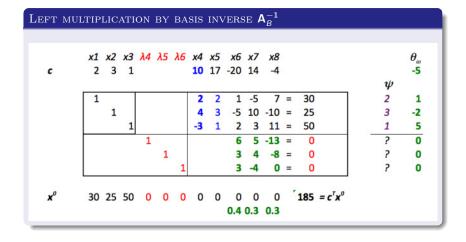
Entering variable only impacts the current basic variables, at most m variables. All but one non-basic variables remain at zero. $\mathbf{y}^0 \in \mathbb{R}^n$ makes changes on at most m+1 components.

y⁰ is formed by a combination of variables : the basic ones and the entering variable.

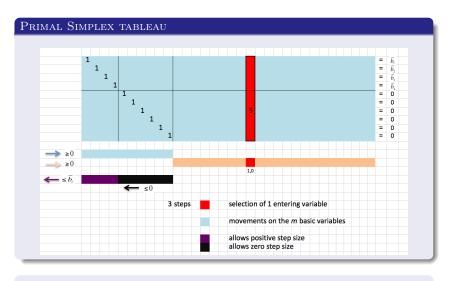
Step size computed such that $\mathbf{x}^1 := \mathbf{x}^0 + \rho^0 \mathbf{y}^0 \ge \mathbf{0}$.

Observation $\#\ 2$: Degenerate solution on simplex-tableau

Observation # 2: Degenerate solution on simplex-tableau

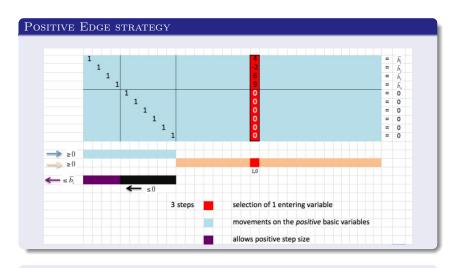


Observation # 3: Structures



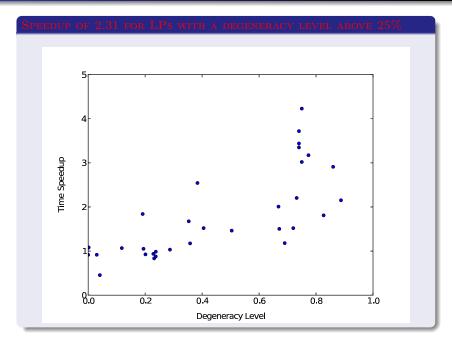
*** Changes on at most m+1 components. ***

Observation # 3: Structures



*** Changes on at most p+1 components. *** Non-degenerate pivot.

Positive Edge : identification in O(m)



Positive Edge

The Positive Edge is a pricing rule for the Primal Simplex : it identifies variables $x_j, j \notin B$ such that $\bar{\mathbf{a}}_{Zj} = \mathbf{0}$... without computing $\bar{\mathbf{a}}_{Zj}$.

Positive Edge

The Positive Edge is a pricing rule for the Primal Simplex: it identifies variables $x_j, j \notin B$ such that $\bar{\mathbf{a}}_{Zj} = \mathbf{0}$... without computing $\bar{\mathbf{a}}_{Zi}$.

If such a variable x_j has a negative reduced cost $\bar{c}_j < 0$, it strictly improves the objective function when entered into the basis.

COMPUTATIONAL COMPLEXITY OF THE POSITIVE EDGE

Let $\mathbf{v} \neq \mathbf{0}$ be a random vector of |Z| non-zero components.

COMPUTATIONAL COMPLEXITY OF THE POSITIVE EDGE

Let $\mathbf{v} \neq \mathbf{0}$ be a random vector of |Z| non-zero components.

$$\mathbf{v}^{\mathsf{T}}\mathbf{\bar{a}}_{Zj} = \mathbf{0} \quad \Longleftrightarrow \quad \mathbf{\bar{a}}_{Zj} = \mathbf{0} \quad \text{ or } \mathbf{v} \perp \mathbf{\bar{a}}_{Zj}$$

Let $\mathbf{v} \neq \mathbf{0}$ be a random vector of |Z| non-zero components.

$$\mathbf{v}^{\mathsf{T}}\mathbf{\bar{a}}_{Zj} = \mathbf{0} \quad \Longleftrightarrow \quad \mathbf{\bar{a}}_{Zj} = \mathbf{0} \quad \text{or} \quad \mathbf{v} \perp \mathbf{\bar{a}}_{Zj}$$

$$\mathbf{v}^{\mathsf{T}} \bar{\mathbf{a}}_{Zj} = \mathbf{v}^{\mathsf{T}} \ [\mathbf{A}_{B}^{-1}]_{Z} \ \mathbf{a}_{Zj} = \mathbf{v}^{\mathsf{T}} \ [\mathbf{A}_{B}^{-1}]_{Z} \ \mathbf{a}_{Zj}$$

Pre-compute once $\mathbf{w}^{\mathsf{T}} := \mathbf{v}^{\mathsf{T}} [\mathbf{A}_{B}^{-1}]_{Z}$, then compute $\mathbf{w}^{\mathsf{T}} \mathbf{a}_{Zj}$, on the **original data**.

Let $\mathbf{v} \neq \mathbf{0}$ be a random vector of |Z| non-zero components.

$$\mathbf{v}^{\mathsf{T}} \bar{\mathbf{a}}_{Zj} = \mathbf{0} \quad \Longleftrightarrow \quad \bar{\mathbf{a}}_{Zj} = \mathbf{0} \quad \text{ or } \quad \mathbf{v} \perp \bar{\mathbf{a}}_{Zj}$$

$$\mathbf{v}^{\mathsf{T}} \bar{\mathbf{a}}_{Zj} = \mathbf{v}^{\mathsf{T}} \ [\mathbf{A}_{B}^{-1}]_{Z} \ \mathbf{a}_{Zj} = \mathbf{v}^{\mathsf{T}} \ [\mathbf{A}_{B}^{-1}]_{Z} \ \mathbf{a}_{Zj}$$

Pre-compute once $\mathbf{w}^{\mathsf{T}} := \mathbf{v}^{\mathsf{T}} [\mathbf{A}_B^{-1}]_{\mathcal{Z}}$, then compute $\mathbf{w}^{\mathsf{T}} \mathbf{a}_{\mathcal{Z}_j}$, on the **original data**.

Determining if x_j satisfies the Positive Edge rule is done in O(m), as for $\bar{c}_j = c_j - \pi^{\mathsf{T}} \mathbf{a}_j, j \notin B$, where $\pi = \mathbf{c}_B^{\mathsf{T}} A_B^{-1}$.

Let $\mathbf{v} \neq \mathbf{0}$ be a random vector of |Z| non-zero components.

$$\mathbf{v}^{\mathsf{T}}\mathbf{\bar{a}}_{Zj} = \mathbf{0} \quad \Longleftrightarrow \quad \mathbf{\bar{a}}_{Zj} = \mathbf{0} \quad \text{or} \quad \mathbf{v} \perp \mathbf{\bar{a}}_{Zj}$$

$$\mathbf{v}^{\mathsf{T}}\mathbf{\bar{a}}_{Zj} = \mathbf{v}^{\mathsf{T}} \ [\mathbf{A}_{B}^{-1}]_{Z} \ \mathbf{a}_{Zj} = \mathbf{v}^{\mathsf{T}} \ [\mathbf{A}_{B}^{-1}]_{Z} \ \mathbf{a}_{Zj}$$

Pre-compute once $\mathbf{w}^{\mathsf{T}} := \mathbf{v}^{\mathsf{T}} [\mathbf{A}_B^{-1}]_{\mathcal{Z}}$, then compute $\mathbf{w}^{\mathsf{T}} \mathbf{a}_{\mathcal{Z}_j}$, on the **original data**.

Determining if x_j satisfies the Positive Edge rule is done in O(m), as for $\bar{c}_j = c_j - \pi^{\mathsf{T}} \mathbf{a}_j, j \notin B$, where $\pi = \mathbf{c}_B^{\mathsf{T}} A_B^{-1}$.

On large-scale instances ($m \approx 100~000$, $n \approx 450~000$), cpu time to identify all variables satisfying the PE rule : usual definition Positive Edge $O(m^2n)$ O(mn)

Let $\mathbf{v} \neq \mathbf{0}$ be a random vector of |Z| non-zero components.

$$\mathbf{v}^{\mathsf{T}} \bar{\mathbf{a}}_{Zj} = \mathbf{0} \quad \Longleftrightarrow \quad \bar{\mathbf{a}}_{Zj} = \mathbf{0} \quad \text{ or } \quad \mathbf{v} \perp \bar{\mathbf{a}}_{Zj}$$

$$\mathbf{v}^{\mathsf{T}}\mathbf{\bar{a}}_{Zj} = \mathbf{v}^{\mathsf{T}} \ [\mathbf{A}_{B}^{-1}]_{Z} \ \mathbf{a}_{Zj} = \mathbf{v}^{\mathsf{T}} \ [\mathbf{A}_{B}^{-1}]_{Z} \ \mathbf{a}_{Zj}$$

Pre-compute once $\mathbf{w}^{\mathsf{T}} := \mathbf{v}^{\mathsf{T}} [\mathbf{A}_B^{-1}]_Z$, then compute $\mathbf{w}^{\mathsf{T}} \mathbf{a}_{Zj}$, on the **original data**.

Determining if x_j satisfies the Positive Edge rule is done in O(m), as for $\bar{c}_j = c_j - \pi^{\mathsf{T}} \mathbf{a}_j, j \notin B$, where $\pi = \mathbf{c}_B^{\mathsf{T}} A_B^{-1}$.

On large-scale instances ($m \approx 100~000$, $n \approx 450~000$), cpu time to identify all variables satisfying the PE rule : usual definition Positive Edge $O(m^2n)$ O(mn) 2500 seconds

Let $\mathbf{v} \neq \mathbf{0}$ be a random vector of |Z| non-zero components.

$$\mathbf{v}^{\mathsf{T}}\mathbf{\bar{a}}_{Zj} = \mathbf{0} \quad \Longleftrightarrow \quad \mathbf{\bar{a}}_{Zj} = \mathbf{0} \quad \text{ or } \ \mathbf{v} \perp \mathbf{\bar{a}}_{Zj}$$

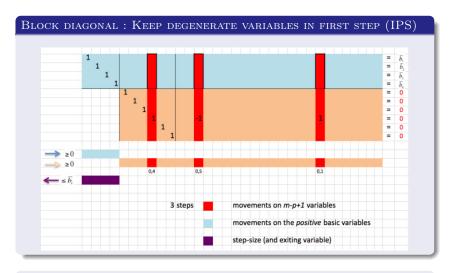
$$\mathbf{v}^{\mathsf{T}}\mathbf{\bar{a}}_{Zj} = \mathbf{v}^{\mathsf{T}} \ [\mathbf{A}_{B}^{-1}]_{Z} \ \mathbf{a}_{Zj} = \mathbf{v}^{\mathsf{T}} \ [\mathbf{A}_{B}^{-1}]_{Z} \ \mathbf{a}_{Zj}$$

Pre-compute once $\mathbf{w}^{\mathsf{T}} := \mathbf{v}^{\mathsf{T}} [\mathbf{A}_B^{-1}]_{Z}$, then compute $\mathbf{w}^{\mathsf{T}} \mathbf{a}_{Zj}$, on the **original data**.

Determining if x_j satisfies the Positive Edge rule is done in O(m), as for $\bar{c}_j = c_j - \pi^{\mathsf{T}} \mathbf{a}_j, j \notin B$, where $\pi = \mathbf{c}_B^{\mathsf{T}} A_B^{-1}$.

On large-scale instances ($m \approx 100~000$, $n \approx 450~000$), cpu time to identify all variables satisfying the PE rule : usual definition Positive Edge $O(m^2n)$ O(mn) 2500 seconds 0.5 seconds

Observation #3: Structures



*** Changes on at most m+1 components. ***

NOTATION

Vectors and matrices are written in **bold face**.

 I_{ℓ} : the $\ell \times \ell$ identity matrix.

0 (1): a vector/matrix with all zeros (ones) entries of appropriate dimensions.

 \mathbf{A}_{RC} : sub-matrix of \mathbf{A} containing the rows and columns indexed by R and C.

Standard notation : basis A_B , inverse A_B^{-1} , $c_B^T x_B$, $A_B x_B$, $\pi = c_B A_B^{-1}$...

$$I_F < x_F < u_F, \ x_L = I_L, \ x_U = u_U$$

Useful decomposition of $x \in \mathbb{R}^n$ in Ax = b, I < x < u

$$\mathbf{x} = \begin{bmatrix} \mathbf{x}_F \\ \mathbf{x}_L \\ \mathbf{x}_U \end{bmatrix} \qquad \mathbf{x} = \begin{bmatrix} \mathbf{x}_B \\ \mathbf{x}_N \end{bmatrix} = \begin{bmatrix} \mathbf{x}_F \\ \mathbf{x}_{B_L} \\ \mathbf{x}_{B_U} \\ \mathbf{x}_{N_L} \\ \mathbf{x}_{N_U} \end{bmatrix}$$

For
$$\emptyset \subseteq S \subseteq B : \mathbf{x} = \begin{bmatrix} \mathbf{x}_S \\ \mathbf{x}_{\bar{S}} \end{bmatrix} = \begin{bmatrix} \mathbf{x}_{S_F} \\ \mathbf{x}_{S_L} \\ \mathbf{x}_{S_F} \\ \mathbf{x}_{\bar{S}_U} \\ \mathbf{x}_{\bar{S}_L} \\ \mathbf{x}_{\bar{c}_U} \end{bmatrix} = \begin{bmatrix} \mathbf{x}_F \\ \mathbf{x}_L \\ \mathbf{x}_U \end{bmatrix}$$

$$\mathbf{x} = \left[\begin{array}{c} \mathbf{x}_{\emptyset} \\ \mathbf{x} \end{array} \right] = \left[\begin{array}{c} \mathbf{x}_{F} \\ \mathbf{x}_{L} \\ \mathbf{x}_{U} \end{array} \right]$$

Dual Guided Pivot Rules for LPs

LINEAR PROGRAM LP

$$z^* := \min c^T x$$

st. $Ax = b$
 $I \le x \le u$

Dual Guided Pivot Rules for LPs

Linear program LP

$$\begin{aligned} z^{\star} &:= & & \text{min} & & c^{\mathsf{T}} x \\ \text{st.} & & & \mathsf{A} x &= b \\ & & & \mathsf{I} &\leq x &\leq u \end{aligned}$$

Generic Algorithm with single parameter set S, $\emptyset \subseteq S \subseteq B$

- Let k = 0 and assume a feasible basic solution \mathbf{x}^k to LP.
- **2** For $\emptyset \subseteq S \subseteq B$, construct the residual problem $LP_S(\mathbf{x}^k)$.
- **3** Fix a subset of the dual variables in row-set $R: \pi_R^{\mathsf{T}} = \mathbf{c}_S^{\mathsf{T}} \mathbf{A}_{RS}^{-1}$.
- Determine the value of the smallest reduced cost μ_s^k . If $\mu_s^k \ge 0$, STOP. Current solution \mathbf{x}^k is optimal for LP.
- **3** Retrieve direction $\mathbf{y}_S^k \in \mathbb{R}^n$ and compute its maximum step-size ρ_S^k .
- Update $x^{k+1} := x^k + \rho_S^k y_S^k;$ $z^{k+1} := z^k + \rho_S^k \mu_S^k;$ k := k + 1.
- Goto Step 2.

Step 2 Construct residual problem $LP(\mathbf{x}^k)$

Assume a feasible solution
$$\mathbf{x}^k = \begin{bmatrix} \mathbf{x}_F^k \\ \mathbf{x}_L^k \\ \mathbf{x}_U^k \end{bmatrix}$$

$$\mathbf{x} := \mathbf{x}^k + \mathbf{y}, \quad \mathbf{y} \in \mathbb{R}^n$$

Step 2 Construct residual problem $LP(x^k)$

Assume a feasible solution
$$\mathbf{x}^k = \begin{bmatrix} \mathbf{x}_F^2 \\ \mathbf{x}_L^k \\ \mathbf{x}_U^k \end{bmatrix}$$

$$\begin{split} \mathbf{x} &:= \mathbf{x}^k + \mathbf{y}, \quad \mathbf{y} \in \mathbb{R}^n \\ &= \mathbf{x}^k + (\vec{\mathbf{y}} - \mathbf{\ddot{y}}), & \quad \vec{\mathbf{y}}, \mathbf{\ddot{y}} \geq \mathbf{0}, \quad \vec{\mathbf{y}}^\mathsf{T} \mathbf{\ddot{y}} = \mathbf{0}, \quad \vec{\mathbf{y}} \leq \vec{\mathbf{r}}^k, \quad \mathbf{\ddot{y}} \leq \mathbf{\ddot{r}}^k \end{split}$$

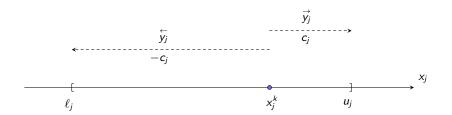
STEP 2 CONSTRUCT RESIDUAL PROBLEM $LP(\mathbf{x}^k)$

Assume a feasible solution
$$\mathbf{x}^k = \left[egin{array}{c} \mathbf{x}_F^k \\ \mathbf{x}_L^k \\ \mathbf{x}_U^k \end{array} \right]$$

 $\mathbf{x} := \mathbf{x}^k + \mathbf{y}, \quad \mathbf{y} \in \mathbb{R}^n$

$$= \mathbf{x}^{k} + (\vec{\mathbf{y}} - \mathbf{\bar{y}}), \qquad \qquad \vec{\mathbf{y}}, \mathbf{\bar{y}} \ge \mathbf{0}, \quad \vec{\mathbf{y}}^{\mathsf{T}} \mathbf{\bar{y}} = 0, \quad \vec{\mathbf{y}} \le \mathbf{\bar{r}}^{k}, \quad \mathbf{\bar{y}} \le \mathbf{\bar{r}}^{k}$$

$$= \mathbf{x}^{k} + \left(\begin{bmatrix} \vec{\mathbf{y}}_{F} \\ \vec{\mathbf{y}}_{L} \\ \vec{\mathbf{y}}_{U} \end{bmatrix} - \begin{bmatrix} \mathbf{\bar{y}}_{F} \\ \mathbf{\bar{y}}_{L} \\ \mathbf{\bar{y}}_{U} \end{bmatrix} \right); \quad \vec{\mathbf{y}}, \mathbf{\bar{y}} \ge \mathbf{0}, \quad \vec{\mathbf{y}}^{\mathsf{T}} \mathbf{\bar{y}} = 0, \quad \vec{\mathbf{y}} \le \mathbf{\bar{r}}^{k}, \quad \mathbf{\bar{y}} \le \mathbf{\bar{r}}^{k}$$



STEP 2 CONSTRUCT RESIDUAL PROBLEM $LP(\mathbf{x}^k)$

Residual problem
$$LP(\mathbf{x}^k)$$
; $\mathbf{x} := \mathbf{x}^k + (\vec{\mathbf{y}} - \mathbf{\ddot{y}})$ (change of variables)

$$z^* := \min \quad \mathbf{c}^{\mathsf{T}} \mathbf{x}^k \quad + \quad \mathbf{c}^{\mathsf{T}} (\vec{\mathbf{y}} - \mathbf{\bar{y}})$$

$$\text{st.} \quad \mathbf{A} \mathbf{x}^k \quad + \quad \mathbf{A} (\vec{\mathbf{y}} - \mathbf{\bar{y}}) \quad = \quad \mathbf{b}$$

$$\mathbf{0} \le \vec{\mathbf{y}} \le \vec{\mathbf{r}}^k$$

$$\mathbf{0} \le \mathbf{\bar{y}} \le \mathbf{\bar{r}}^k$$

STEP 2 CONSTRUCT RESIDUAL PROBLEM $LP(\mathbf{x}^k)$

Residual problem
$$LP(\mathbf{x}^k)$$
; $\mathbf{x} := \mathbf{x}^k + (\mathbf{y} - \mathbf{\bar{y}})$ (change of variables)

$$\begin{split} z^{\star} &:= \min \quad \mathbf{c}^{\mathsf{T}} \mathbf{x}^{k} \quad + \quad \mathbf{c}^{\mathsf{T}} (\vec{\mathbf{y}} - \mathbf{\tilde{y}}) \\ &\text{st.} \quad \mathbf{A} \mathbf{x}^{k} \quad + \quad \mathbf{A} (\vec{\mathbf{y}} - \mathbf{\tilde{y}}) \quad = \quad \mathbf{b} \\ &\mathbf{0} \leq \vec{\mathbf{y}} \leq \mathbf{\tilde{r}}^{k} \\ &\mathbf{0} \leq \mathbf{\tilde{y}} \leq \mathbf{\tilde{r}}^{k} \end{split}$$

RESIDUAL PROBLEM $LP(\mathbf{x}^k)$

$$z^\star := \mathbf{c}^\intercal \mathbf{x}^k + \min$$
 $\mathbf{c}^\intercal (\vec{\mathbf{y}} - \mathbf{\bar{y}})$ st. $\mathbf{A} (\vec{\mathbf{y}} - \mathbf{\bar{y}}) = \mathbf{0}$ $\mathbf{0} \le \vec{\mathbf{y}} \le \vec{\mathbf{r}}^k$ $\mathbf{0} < \mathbf{\bar{y}} < \mathbf{\bar{r}}^k$

STEP 2 CONSTRUCT RESIDUAL PROBLEM $LP(x^k)$

RESIDUAL PROBLEM WITH
$$S = F$$
, $\bar{S} = L \cup U$

$$z^{\star} := \mathbf{c}^{\mathsf{T}} \mathbf{x}^{k} +$$

$$\min \quad \mathbf{c}_{F}^{\mathsf{T}} (\vec{\mathbf{y}}_{F} - \mathbf{\bar{y}}_{F}) + \mathbf{c}_{L}^{\mathsf{T}} (\vec{\mathbf{y}}_{L}) - \mathbf{c}_{U}^{\mathsf{T}} (\mathbf{\bar{y}}_{U})$$

$$\text{st.} \quad \mathbf{A}_{F} (\vec{\mathbf{y}}_{F} - \mathbf{\bar{y}}_{F}) + \mathbf{A}_{L} (\vec{\mathbf{y}}_{L}) - \mathbf{A}_{U} (\mathbf{\bar{y}}_{U}) = \mathbf{0}$$

$$\vec{\mathbf{y}}_{F} \geq \mathbf{0}, \ \vec{\mathbf{y}}_{F} \geq \mathbf{0}, \qquad \vec{\mathbf{y}}_{L} \geq \mathbf{0}, \qquad \mathbf{\bar{y}}_{U} \geq \mathbf{0}$$

$$\vec{\mathbf{y}}_{F} \leq \vec{\mathbf{r}}_{F}, \ \vec{\mathbf{y}}_{F} \leq \mathbf{\bar{r}}_{F}, \qquad \vec{\mathbf{y}}_{L} \leq \vec{\mathbf{r}}_{L}, \qquad \mathbf{\bar{y}}_{U} \leq \mathbf{\bar{r}}_{U}$$

STEP 2 CONSTRUCT RESIDUAL PROBLEM $LP(x^k)$

Residual problem – General case : $\emptyset \subseteq S \subseteq B$

$$\begin{array}{lll} \mathbf{c}^\intercal \mathbf{x}^k + & & \min \ \mathbf{c}_S^\intercal (\vec{y}_S - \ddot{\mathbf{y}}_S) \ & & \\ & \text{st.} & \mathbf{A}_S (\vec{y}_S - \ddot{\mathbf{y}}_S) \ & & \\ & & \vec{y}_S \geq \mathbf{0}, \ \ddot{\mathbf{y}}_S \geq \mathbf{0} \\ & & & \ddot{\mathbf{y}}_S \geq \mathbf{0}, \ \ddot{\mathbf{y}}_S \geq \mathbf{0} \\ & & & \ddot{\mathbf{y}}_{\bar{S}_L} \leq \mathbf{0}, \ \ddot{\mathbf{y}}_{\bar{S}_U} \leq \mathbf{0} \\ & & & \ddot{\mathbf{y}}_{\bar{S}_L} \leq \mathbf{0}, \ \ddot{\mathbf{y}}_{\bar{S}_U} \leq \mathbf{0} \\ & & & \ddot{\mathbf{y}}_{\bar{S}_L} \leq \mathbf{0}, \ \ddot{\mathbf{y}}_{\bar{S}_U} \leq \mathbf{0} \\ & & \ddot{\mathbf{y}}_{\bar{S}_L} \leq \mathbf{0}, \ \ddot{\mathbf{y}}_{\bar{S}_U} \leq \mathbf{0} \\ & & \ddot{\mathbf{y}}_{\bar{S}_L} \leq \mathbf{0}, \ \ddot{\mathbf{y}}_{\bar{S}_U} \leq \mathbf{0} \\ & & \ddot{\mathbf{y}}_{\bar{S}_L} \leq \mathbf{0}, \ \ddot{\mathbf{y}}_{\bar{S}_U} \leq \mathbf{0} \\ & & \ddot{\mathbf{y}}_{\bar{S}_L} \leq \mathbf{0}, \ \ddot{\mathbf{y}}_{\bar{S}_U} \leq \mathbf{0} \\ \end{array}$$

$$S = S_F \cup S_L \cup S_U; \qquad \bar{S} = \bar{S}_F \cup \bar{S}_L \cup \bar{S}_U$$

STEP 2 CONSTRUCT RESIDUAL PROBLEM $LP(\mathbf{x}^k)$

RESIDUAL PROBLEM WITH
$$S = B$$
 (BASIC), $\bar{S} = N$ (NON-BASIC)

$$c^{\mathsf{T}}x^k + \min c^{\mathsf{T}}_B(\vec{y}_B - \ddot{y}_B) + c^{\mathsf{T}}_N(\vec{y}_N - \ddot{y}_N) \\
\text{st.} \quad A_B(\vec{y}_B - \ddot{y}_B) + A_N(\vec{y}_N - \ddot{y}_N) = 0 \\
\vec{y}_B \ge 0, \ \ddot{y}_B \ge 0 \qquad \vec{y}_N \ge 0, \ \ddot{y}_N \ge 0 \\
\vec{y}_{N_L} \le 0, \ \vec{y}_{N_U} \le 0$$

$$\vec{y}_F \le \vec{r}_F, \ \ddot{y}_F \le \ddot{r}_F \\
\ddot{y}_{B_L} \le 0, \ \vec{y}_{B_U} \le 0 \qquad \vec{y}_{N_L} \le \vec{r}_{N_L}, \ \ddot{y}_{N_U} \le \ddot{r}_{N_U}$$

 $B = F \cup B_I \cup B_{II}$; $N = N_I \cup N_{II}$

STEP 3 FIX SOME DUAL VARIABLES IN ROW-SET R

For
$$\emptyset \subseteq S \subseteq B$$
, find $\mathbf{y}^k = \begin{bmatrix} \mathbf{y}_S^k \\ \mathbf{y}_S^k \end{bmatrix} = \begin{bmatrix} (\mathbf{\vec{y}}_S^k - \mathbf{\bar{y}}_S^k) \\ (\mathbf{\vec{y}}_S^k - \mathbf{\bar{y}}_S^k) \end{bmatrix}$ of min reduced cost μ^k .

For basic columns A_S , select a $s \times s$ square matrix A_{RS} , a set of s independent rows.

$$\mathbf{T} = \begin{bmatrix} \mathbf{A}_{RS} & \mathbf{0} \\ \mathbf{A}_{ZS} & \mathbf{I}_{m-s} \end{bmatrix} \quad \mathbf{T}^{-1} = \begin{bmatrix} \mathbf{A}_{RS}^{-1} & \mathbf{0} \\ -\mathbf{A}_{ZS}\mathbf{A}_{RS}^{-1} & \mathbf{I}_{m-s} \end{bmatrix}$$

STEP 3 FIX SOME DUAL VARIABLES IN ROW-SET R

$$\emptyset \subset S \subset B$$

$$\begin{array}{lll} \text{Perform } \mathbf{T}^{-1} \mathbf{A}(\vec{\mathbf{y}} - \vec{\mathbf{y}}) = \mathbf{0}. & \text{General case} & \boldsymbol{\pi}^{\mathsf{T}} = \boldsymbol{\psi}^{\mathsf{T}} \mathbf{T}^{-1} \\ \mathbf{c}^{\mathsf{T}} \mathbf{x}^k + \min \ \mathbf{c}_{S}^{\mathsf{T}} (\vec{\mathbf{y}}_S - \vec{\mathbf{y}}_S) \ & \\ \text{st.} & (\vec{\mathbf{y}}_S - \vec{\mathbf{y}}_S) \ & \\ \mathbf{x}_S \cdot (\vec{\mathbf{y}}_S -$$

Observe $\bar{c}_S = 0$ (basic variables).

STEP 3 FIX SOME DUAL VARIABLES IN ROW-SET R

Perform
$$\mathbf{T}^{-1}\mathbf{A}(\vec{\mathbf{y}} - \mathbf{\ddot{\mathbf{y}}}) = \mathbf{0}$$
. General case : details $\boldsymbol{\pi}^{\mathsf{T}} = \boldsymbol{\psi}^{\mathsf{T}}\mathbf{T}^{-1}$

$$\mathbf{c}^{\mathsf{T}}\mathbf{x}^k + \min \ \mathbf{c}^{\mathsf{T}}_{S}(\vec{\mathbf{y}}_{S} - \mathbf{\ddot{\mathbf{y}}}_{S}) + \mathbf{c}^{\mathsf{T}}_{S}(\vec{\mathbf{y}}_{\bar{S}} - \mathbf{\ddot{\mathbf{y}}}_{\bar{S}})$$
st. $(\vec{\mathbf{y}}_{S} - \mathbf{\ddot{\mathbf{y}}}_{S}) + \mathbf{A}_{R\bar{S}}(\vec{\mathbf{y}}_{\bar{S}} - \mathbf{\ddot{\mathbf{y}}}_{\bar{S}}) = \mathbf{0} \ [\boldsymbol{\psi}_{R} = \mathbf{c}_{S}]$

$$\mathbf{\ddot{\mathbf{y}}}_{S} \geq \mathbf{0}, \ \mathbf{\ddot{\mathbf{y}}}_{S} \geq \mathbf{0}$$

$$\mathbf{\ddot{\mathbf{A}}}_{Z\bar{S}}(\vec{\mathbf{y}}_{\bar{S}} - \mathbf{\ddot{\mathbf{y}}}_{\bar{S}}) = \mathbf{0} \ [\boldsymbol{\psi}_{Z} \ unknown]$$

$$\mathbf{\ddot{\mathbf{y}}}_{\bar{S}} \geq \mathbf{0}, \ \mathbf{\ddot{\mathbf{y}}}_{\bar{S}} \geq \mathbf{0}$$

$$\mathbf{\ddot{\mathbf{y}}}_{\bar{S}_{L}} \leq \mathbf{0}, \ \mathbf{\ddot{\mathbf{y}}}_{\bar{S}_{U}} \leq \mathbf{0}$$

$$\mathbf{\ddot{\mathbf{y}}}_{S_{F}} \leq \mathbf{\ddot{\mathbf{r}}}_{S_{F}}, \ \mathbf{\ddot{\mathbf{y}}}_{S_{F}} \leq \mathbf{\ddot{\mathbf{r}}}_{S_{F}}$$

$$\mathbf{\ddot{\mathbf{y}}}_{S_{L}} \leq \mathbf{0}, \ \mathbf{\ddot{\mathbf{y}}}_{S_{U}} \leq \mathbf{0}$$

$$\mathbf{\ddot{\mathbf{y}}}_{\bar{S}_{L}} \leq \mathbf{\ddot{\mathbf{r}}}_{\bar{S}_{L}}, \ \mathbf{\ddot{\mathbf{y}}}_{\bar{S}_{U}} \leq \mathbf{\ddot{\mathbf{r}}}_{\bar{S}_{U}}$$

 $S = S_F \cup S_L \cup S_U$: $\bar{S} = \bar{S}_F \cup \bar{S}_U \cup \bar{S}_U$

Step 4 Find minimum reduced cost μ_S^k

Pricing of the variables

$$S \subseteq B$$
, hence $\overline{c}_S = 0$.

Therefore pricing of $\vec{y}_{\bar{S}}$ and $\vec{y}_{\bar{S}}$ needed to get partial direction $(\vec{y}_{\bar{S}}^k - \vec{y}_{\bar{S}}^k)$

followed by impact on $(\vec{\mathbf{y}}_{S}^{k} - \vec{\mathbf{y}}_{S}^{k})$ to complete direction $\mathbf{y}^{k} = \begin{bmatrix} \mathbf{y}_{S}^{k} \\ \mathbf{y}_{S}^{k} \end{bmatrix}$.

PRIMAL/DUAL FORMULATIONS OF THE PRICING

$$\boldsymbol{\psi}_{\mathit{R}}^{\intercal} = \mathbf{c}_{\mathit{S}}^{\intercal} \mathbf{A}_{\mathit{RS}}^{-1}$$

$$\begin{array}{lll} \max \; \mu & \text{st. } \mu \mathbf{1}^\intercal & \leq & \mathbf{c}_{\bar{\mathsf{S}}}^\intercal - \boldsymbol{\psi}_R^\intercal \mathsf{A}_{R\bar{\mathsf{S}}} - \boldsymbol{\psi}_Z^\intercal \bar{\mathsf{A}}_{Z\bar{\mathsf{S}}} & \left[\vec{\boldsymbol{\mathsf{y}}}_{\bar{\mathsf{S}}} \right] \\ & \mu \mathbf{1}^\intercal & \leq - (\mathbf{c}_{\bar{\mathsf{S}}}^\intercal - \boldsymbol{\psi}_R^\intercal \mathsf{A}_{R\bar{\mathsf{S}}} - \boldsymbol{\psi}_Z^\intercal \bar{\mathsf{A}}_{Z\bar{\mathsf{S}}}) & \left[\mathbf{\bar{\mathsf{y}}}_{\bar{\mathsf{S}}} \right] \end{array}$$

A convex combination of the variables $\vec{y}_{\bar{\varsigma}}$ and $\dot{\bar{y}}_{\bar{\varsigma}}$

*** Optimal solution : μ_S^k , $\vec{\mathbf{y}}_{\bar{5}}^k$ and $\dot{\mathbf{y}}_{\bar{5}}^k$. ***

Step 5 Retrieve direction $\mathbf{y}_{\mathsf{S}}^k$ of minimum reduced cost μ_{S}^k

... Complete direction \mathbf{y}_{5}^{k}

Given $\mu^k < 0$ and $(\vec{\mathbf{y}}_{\bar{S}}^k - \vec{\mathbf{y}}_{\bar{S}}^k)$, find impacts on other variables $(\vec{\mathbf{y}}_S^k - \vec{\mathbf{y}}_S^k)$.

$$\begin{aligned} & (\vec{y}_S - \vec{y}_S) \quad + \quad \bar{\mathbf{A}}_{R\bar{S}} (\vec{\mathbf{y}_S^k} - \vec{\mathbf{y}_S^k}) \quad = \quad \mathbf{0} \\ \vec{\mathbf{y}}_S \geq \mathbf{0}, \ \, \mathbf{\tilde{y}}_S \geq \mathbf{0} \qquad \qquad \qquad \vec{y}_j \, \mathbf{\tilde{y}}_j = \mathbf{0}, \quad \forall j \in S_F \end{aligned}$$

DIRECTION \mathbf{y}_{s}^{k}

$$\mathbf{y}_{S}^{k} = \begin{bmatrix} (\ddot{\mathbf{y}}_{S}^{k} - \ddot{\mathbf{y}}_{S}^{k}) \\ (\ddot{\mathbf{y}}_{S}^{k} - \ddot{\mathbf{y}}_{S}^{k}) \end{bmatrix} = \begin{bmatrix} -\bar{\mathbf{A}}_{R\bar{S}}(\ddot{\mathbf{y}}_{S}^{k} - \ddot{\mathbf{y}}_{S}^{k}) \\ (\ddot{\mathbf{y}}_{S}^{k} - \ddot{\mathbf{y}}_{S}^{k}) \end{bmatrix} = \begin{bmatrix} Impact \ level \ Pricing \ level \end{bmatrix}$$

Positive part; Negative part.

Step 5 Compute maximum step-size ρ_{S}^k

$$\rho \left[\begin{array}{c} \vec{\mathbf{y}}_{S_F}^k \\ \vec{\mathbf{y}}_{S_L}^k \\ \vec{\mathbf{y}}_{S_U} \end{array} \right] \leq \left[\begin{array}{c} \vec{\mathbf{r}}_{S_F}^k \\ \vec{\mathbf{r}}_{S_L}^k \\ \mathbf{0} \end{array} \right]; \quad \rho \left[\begin{array}{c} \vec{\mathbf{y}}_{S_F}^k \\ \vec{\mathbf{y}}_{S_L}^k \\ \vec{\mathbf{y}}_{S_U}^k \end{array} \right] \leq \left[\begin{array}{c} \vec{\mathbf{r}}_{S_F}^k \\ \mathbf{0} \\ \vec{\mathbf{r}}_{S_U}^k \end{array} \right]$$

$$\rho \left[\begin{array}{c} \vec{\mathbf{y}}_{\bar{S}_F}^k \\ \vec{\mathbf{y}}_{\bar{S}_L}^k \\ \mathbf{0} \end{array} \right] \leq \left[\begin{array}{c} \vec{\mathbf{r}}_{\bar{S}_F}^k \\ \vec{\mathbf{r}}_{\bar{S}_L}^k \\ \mathbf{0} \end{array} \right]; \quad \rho \left[\begin{array}{c} \vec{\mathbf{y}}_{\bar{S}_F}^k \\ \mathbf{0} \\ \vec{\mathbf{y}}_{\bar{S}_U}^k \end{array} \right] \leq \left[\begin{array}{c} \vec{\mathbf{r}}_{\bar{S}_F}^k \\ \mathbf{0} \\ \vec{\mathbf{r}}_{\bar{S}_U}^k \end{array} \right]$$

10 out of 12 types of residual upper bonds to verify. $\emptyset \subseteq S \subseteq B$

Step 5 Compute maximum step-size ρ_{S}^k

$$\rho \left[\begin{array}{c} \vec{\mathbf{y}}_{F}^{k} \\ \vec{\mathbf{y}}_{B_{L}}^{k} \\ \vec{\mathbf{y}}_{B_{U}} \end{array} \right] \leq \left[\begin{array}{c} \vec{\mathbf{r}}_{F}^{k} \\ \vec{\mathbf{r}}_{B_{L}}^{k} \\ \mathbf{0} \end{array} \right]; \quad \rho \left[\begin{array}{c} \vec{\mathbf{y}}_{F}^{k} \\ \vec{\mathbf{y}}_{B_{L}}^{k} \\ \vec{\mathbf{y}}_{B_{U}}^{k} \end{array} \right] \leq \left[\begin{array}{c} \vec{\mathbf{r}}_{F}^{k} \\ \mathbf{0} \\ \vec{\mathbf{r}}_{B_{U}}^{k} \end{array} \right]$$

$$\rho \left[\begin{array}{c} \vec{\mathbf{y}}_{N_L}^k \end{array} \right] \leq \left[\begin{array}{c} \vec{\mathbf{r}}_{N_L}^k \end{array} \right]; \quad \rho \left[\begin{array}{c} \overline{\mathbf{y}}_{N_U}^k \end{array} \right] \leq \left[\begin{array}{c} \overline{\mathbf{r}}_{N_U}^k \end{array} \right]$$

8 types of residual upper bonds to verify for Primal Simplex. S =

Step 5 Compute maximum step-size ρ_S^k

$$\rho \left[\begin{array}{c} \mathbf{\ddot{y}}_P^k \\ \mathbf{\ddot{y}}_Z^k \end{array} \right] \leq \left[\begin{array}{c} \mathbf{\ddot{r}}_P^k \\ \mathbf{0} \end{array} \right]$$

Only 2 types of residual upper bonds to verify for PS in standard form. $S = B = P \cup Z$ (positive and zero variables)

Step 5 Compute maximum step-size ρ_{S}^k

$$\rho \left[\begin{array}{c} \vec{\mathbf{y}}_{F}^{k} \\ \vec{\mathbf{y}}_{L}^{k} \end{array} \right] \leq \left[\begin{array}{c} \vec{\mathbf{r}}_{F}^{k} \\ \vec{\mathbf{r}}_{L}^{k} \end{array} \right]; \quad \rho \left[\begin{array}{c} \mathbf{\tilde{y}}_{F}^{k} \\ \mathbf{\tilde{y}}_{U}^{k} \end{array} \right] \leq \left[\begin{array}{c} \mathbf{\tilde{r}}_{F}^{k} \\ \mathbf{\tilde{r}}_{U}^{k} \end{array} \right]$$

4 types of strictly positive residual upper bonds to verify in MMCC. $S = \emptyset$

Step 5 Compute maximum step-size ρ_{S}^k

$$\rho \left[\begin{array}{c} \vec{\mathbf{y}}_{F}^{k} \end{array} \right] \leq \left[\begin{array}{c} \vec{\mathbf{r}}_{F}^{k} \end{array} \right]; \quad \rho \left[\begin{array}{c} \mathbf{\tilde{y}}_{F}^{k} \end{array} \right] \leq \left[\begin{array}{c} \vec{\mathbf{r}}_{F}^{k} \end{array} \right]$$

$$\rho \left[\begin{array}{c} \vec{\mathbf{y}}_{L}^{k} \end{array}\right] \leq \left[\begin{array}{c} \vec{\mathbf{r}}_{L}^{k} \end{array}\right]; \quad \rho \left[\begin{array}{c} \overline{\mathbf{y}}_{U}^{k} \end{array}\right] \leq \left[\begin{array}{c} \overline{\mathbf{r}}_{U}^{k} \end{array}\right]$$

4 types of strictly positive residual upper bonds to verify in IPS. S = F

Special case #1:S=B

$$T = [A_B \quad \emptyset], \quad T^{-1} = \begin{bmatrix} A_B^{-1} \\ \emptyset \end{bmatrix}$$

 $B = F \cup B_L \cup B_U; \quad N = N_L \cup N_U$

Special case #1: S = B

$$\mathbf{T} = [\mathbf{A}_B \quad {}^{\bullet}], \quad \mathbf{T}^{-1} = \begin{bmatrix} \mathbf{A}_B^{-1} \\ \emptyset \end{bmatrix}$$
$$B = F \cup B_I \cup B_{II}: \quad N = N_I \cup N_{II}$$

Primal simplex method (Dantzig 1945)

$$\begin{array}{lll} \mathbf{c}^{\mathsf{T}}\mathbf{x}^k + \min \ \mathbf{c}_B^{\mathsf{T}}(\vec{\mathbf{y}}_B - \mathbf{\bar{y}}_B) & + & \mathbf{c}_N^{\mathsf{T}}(\vec{\mathbf{y}}_N - \mathbf{\bar{y}}_N) \\ & \mathrm{st.} & (\vec{\mathbf{y}}_B - \mathbf{\bar{y}}_B) & + & \bar{\mathbf{A}}_N(\vec{\mathbf{y}}_N - \mathbf{\bar{y}}_N) & = & \mathbf{0} \quad [\boldsymbol{\psi}^{\mathsf{T}} = \mathbf{c}_B] \\ & \vec{\mathbf{y}}_B \geq \mathbf{0}, \ \mathbf{\bar{y}}_B \geq \mathbf{0} & \\ & & \vec{\mathbf{y}}_N \geq \mathbf{0}, \ \mathbf{\bar{y}}_N \geq \mathbf{0} \\ & & \mathbf{\bar{y}}_N \geq \mathbf{0}, \ \mathbf{\bar{y}}_N \geq \mathbf{0} \\ & & \mathbf{\bar{y}}_N \geq \mathbf{0}, \ \mathbf{\bar{y}}_N \geq \mathbf{0} \\ & & \mathbf{\bar{y}}_F \leq \mathbf{\bar{r}}_F, \ \mathbf{\bar{y}}_F \leq \mathbf{\bar{r}}_F \\ & \mathbf{\bar{y}}_{B_I} \leq \mathbf{0}, \ \mathbf{\bar{y}}_{B_U} \leq \mathbf{0} & & \mathbf{\bar{y}}_{N_L} \leq \mathbf{\bar{r}}_{N_L}, \mathbf{\bar{y}}_{N_U} \leq \mathbf{\bar{r}}_{N_U} & *\rho_B \geq \mathbf{0}* \end{array}$$

Properties of PS

No equality constraints in the pricing problem.

Pricing contains • convex combination of the non-basic variables

• non-negativity restrictions (a cone).

Due to the step-size $*\rho_B \ge 0*$, possible degenerate pivots.

Oscillation of μ_B ; it may even not converge towards 0.

Special case $\#2: S = F \Rightarrow \bar{S} = L \cup U$

$$\mathbf{T} = \left[egin{array}{ccc} \mathbf{A}_{RF} & \mathbf{0} & & \\ \mathbf{A}_{ZF} & \mathbf{I}_{m-r} & \\ \end{array}
ight], \quad \mathbf{T}^{-1} = \left[egin{array}{ccc} \mathbf{A}_{RF}^{-1} & \mathbf{0} & & \\ -\mathbf{A}_{ZS}\mathbf{A}_{RF}^{-1} & \mathbf{I}_{m-r} & \\ \end{array}
ight].$$

Special case $\#2: S = F \Rightarrow \bar{S} = L \cup U$

$$\mathbf{T} = \left[\begin{array}{cc} \mathbf{A}_{RF} & \mathbf{0} \\ \mathbf{A}_{ZF} & \mathbf{I}_{m-r} \end{array} \right], \quad \mathbf{T}^{-1} = \left[\begin{array}{cc} \mathbf{A}_{RF}^{-1} & \mathbf{0} \\ -\mathbf{A}_{ZS} \mathbf{A}_{RF}^{-1} & \mathbf{I}_{m-r} \end{array} \right].$$

Improved Primal Simplex method (Elhallaoui et al. 2011)

$$\begin{array}{lllll} \mathbf{c}^{\intercal}\mathbf{x}^{k} + \min & \mathbf{c}_{F}^{\intercal}(\vec{\mathbf{y}}_{F} - \mathbf{\bar{y}}_{F}) & + & \mathbf{c}_{L}^{\intercal}\vec{\mathbf{y}}_{L} - \mathbf{c}_{U}^{\intercal}\mathbf{\bar{y}}_{U} \\ & \text{st.} & (\vec{\mathbf{y}}_{F} - \mathbf{\bar{y}}_{F}) & + & \mathbf{\bar{A}}_{RL}\vec{\mathbf{y}}_{L} - \mathbf{\bar{A}}_{RU}\mathbf{\bar{y}}_{U} & = & \mathbf{0} & [\boldsymbol{\psi}_{R} = \mathbf{c}_{F}] \\ & & \vec{\mathbf{y}}_{F} \geq \mathbf{0}, \ \mathbf{\bar{y}}_{F} \geq \mathbf{0} & \\ & & & \mathbf{\bar{A}}_{ZL}\vec{\mathbf{y}}_{L} - \mathbf{\bar{A}}_{ZU}\mathbf{\bar{y}}_{U} & = & \mathbf{0} & [\boldsymbol{\psi}_{Z}] \\ & & & \vec{\mathbf{y}}_{L} \geq \mathbf{0}, \ \mathbf{\bar{y}}_{U} \geq \mathbf{0} & \\ & & & \vec{\mathbf{y}}_{F} \leq \vec{\mathbf{r}}_{F}, \ \mathbf{\bar{y}}_{F} \leq \mathbf{\bar{r}}_{F}, & \vec{\mathbf{y}}_{L} \leq \vec{\mathbf{r}}_{L}, \mathbf{\bar{y}}_{U} \leq \mathbf{\bar{r}}_{U} & *\rho > 0* \end{array}$$

Properties of IPS

f equality constraints in the master problem, m - f in the pricing problem.

Non-degenerate pivots only ($\rho_F > 0$).

 $z^0>z^1>z^2>\cdots=z^*$ cost strictly decreasing at each iteration $(\rho_F>0)$.

Oscillation of μ_F but converging towards 0.

Special case $\#3:S=\emptyset$

Select
$$S = \emptyset$$
. $T = \begin{bmatrix} \emptyset & I_m \end{bmatrix}$, $T^{-1} = \begin{bmatrix} \emptyset & I_m \end{bmatrix}$

Special case $\#3: S = \emptyset$

Select
$$S = \emptyset$$
. $\mathbf{T} = \begin{bmatrix} \emptyset & \mathbf{I_m} \end{bmatrix}, \mathbf{T}^{-1} = \begin{bmatrix} \emptyset \\ \mathbf{I_m} \end{bmatrix}$

MINIMUM MEAN CYCLE-CANCELING ALGORITHM ADAPTED FOR LP

$$\begin{array}{lll} \mathbf{c}^{\mathsf{T}}\mathbf{x}^k + & \min & \mathbf{c}^{\mathsf{T}}(\vec{\mathbf{y}} - \overline{\mathbf{y}}) & \\ & \mathrm{st.} & \mathbf{A}(\vec{\mathbf{y}} - \overline{\mathbf{y}}) &= \mathbf{0} & \textit{Directions} \\ & & \vec{\mathbf{y}}, \ \overline{\mathbf{y}} &\geq \mathbf{0} & \textit{in the cone} \\ & & & \overline{\mathbf{y}}_L, \ \overline{\mathbf{y}}_U &\leq \mathbf{0} & \textit{at vertex } \mathbf{x}^k \\ \\ & & & \vec{\mathbf{y}}_F, \overline{\mathbf{y}}_F, \ \overline{\mathbf{y}}_L, \ \overline{\mathbf{y}}_U &\leq \vec{\mathbf{r}}_F^k, \ \overline{\mathbf{r}}_F^k, \ \overline{\mathbf{r}}_L^k, \ \overline{\mathbf{r}}_U^k & *\textit{Step size } \rho_\emptyset > 0 * \end{array}$$

Properties of MMCC

All equality constraints in the pricing problem.

Upper bounds in the master problem.

$$\begin{split} z^0 > z^1 > z^2 > \cdots = z^\star & \text{cost strictly decreasing at each iteration } (\rho_\emptyset > 0). \\ \mu^0 \le \mu^1 \le \mu^2 \le \cdots = 0 & \text{smallest reduced cost non decreasing}. \end{split}$$

MMCC is strongly polynomial for network flow problems in O(mn) phases.

Goldberg and Tarjan (1989), Radzick and Goldberg (1994)

Illustration of z (Network with n=1025, m=91,220)

Computational results

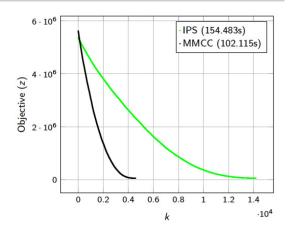


Illustration of μ (Network with n=1025, m=91,220)

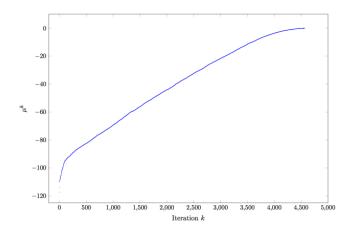
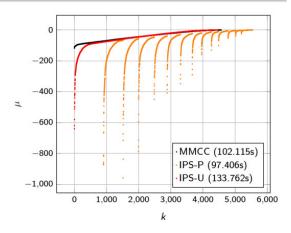


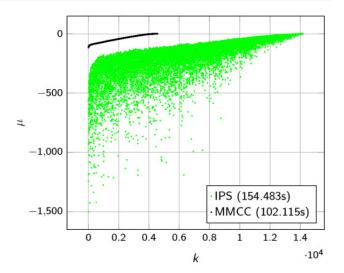
Illustration of the phases on μ (n=1025, m=91,220)

COMPUTATIONAL RESULTS



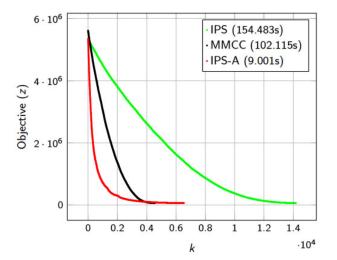
μ for MMCC and IPS on a network (n=1025, m=91,220)

Computational results



z for MMCC, IPS and IPS-Accelerated (n=1025, m=91,220)

COMPUTATIONAL RESULTS



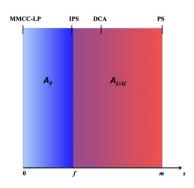
Dual Guided Pivot Rules for LPs

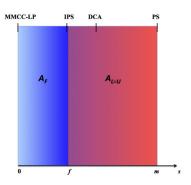
Linear program LP

$$\begin{aligned} z^{\star} &:= & & \text{min} & & c^{\mathsf{T}} x \\ \text{st.} & & & \mathbf{A} x &= \mathbf{b} \\ & & & \mathbf{I} \, \leq \, x & \leq \mathbf{u} \end{aligned}$$

Generic Algorithm with single parameter set S, $\emptyset \subseteq S \subseteq B$

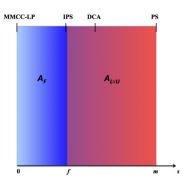
- **1** Let k = 0 and assume a feasible basic solution \mathbf{x}^k to LP.
- **2** For $\emptyset \subseteq S \subseteq B$, construct the residual problem $LP_S(\mathbf{x}^k)$.
- **3** Fix a subset of the dual variables in row-set $R: \pi_R^T = \mathbf{c}_S^T \mathbf{A}_{RS}^{-1}$.
- Determine the value of the smallest reduced cost μ_s^k . If $\mu_s^k \ge 0$, STOP. Current solution \mathbf{x}^k is optimal for LP.
- **3** Retrieve direction $\mathbf{y}_S^k \in \mathbb{R}^n$ and compute its maximum step-size ρ_S^k .
- Update $x^{k+1} := x^k + \rho_S^k y_S^k;$ $z^{k+1} := z^k + \rho_S^k \mu_S^k;$ k := k + 1.
- Goto Step 2.





- $S \cap \{L \cup U\} \neq \emptyset$: It may come up with degenerate pivots and not converge.
 - Primal Simplex method (PS). S =

Properties

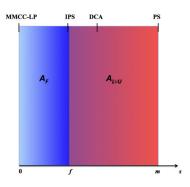


- $S \cap \{L \cup U\} \neq \emptyset$: It may come up with degenerate pivots and not converge.
 - Primal Simplex method (PS).

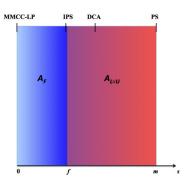
$$S = B$$
, * $ρ$ _B ≥ 0 *

- $\emptyset \subseteq S \subseteq F$: It ensures a non-degenerate pivot at every iteration.
 - Improved Primal Simplex algorithm (IPS). $S = F_1 * \rho_F > 0 *$

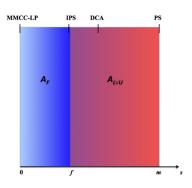
$$S = F$$
, * $\rho_F > 0$



- $S \cap \{L \cup U\} \neq \emptyset$: It may come up with degenerate pivots and not converge.
 - Primal Simplex method (PS). S = B, * $\rho_B \ge 0$ *
- $\emptyset \subseteq S \subseteq F$: It ensures a non-degenerate pivot at every iteration.
 - Improved Primal Simplex algorithm (IPS). S = F, * $\rho_F > 0$ *
 - Minimum mean cycle-canceling algorithm (MMCC) $S = \emptyset$, * $\rho_{\emptyset} > 0$ *



- $S \cap \{L \cup U\} \neq \emptyset$: It may come up with degenerate pivots and not converge.
 - Primal Simplex method (PS). S = B, * $\rho_B \ge 0$ *
- $\emptyset \subseteq S \subseteq F$: It ensures a non-degenerate pivot at every iteration.
 - Improved Primal Simplex algorithm (IPS). S = F, * $\rho_F > 0$ *
 - Minimum mean cycle-canceling algorithm (MMCC) $S = \emptyset$, * $\rho_{\emptyset} > 0$ * Strongly polynomial for network flow problems.



- $S \cap \{L \cup U\} \neq \emptyset$: It may come up with degenerate pivots and not converge.
- Primal Simplex method (PS).

$$S = B, * \rho_B \ge 0 *$$

- $\emptyset \subseteq S \subseteq F$: It ensures a non-degenerate pivot at every iteration.
 - Improved Primal Simplex algorithm (IPS). $S = F_1 * \rho_F > 0 *$
 - Minimum mean cycle-canceling algorithm (MMCC) $S = \emptyset$, * $\rho_{\emptyset} > 0$ * Strongly polynomial for network flow problems.
 - * $S \subset F$: Optimal direction \mathbf{y}_S^k can be an <u>interior</u> ray.

VECTOR SPACE DECOMPOSITION

- *** Imagine the same transformation is kept for a while...
 - \bullet T := $[\Lambda, \Lambda^{\perp}]$
 - $\Lambda = \begin{bmatrix} \Lambda_R \\ \Lambda_Z \end{bmatrix}$, where Λ_R is a set of s independent rows.

$$\bullet \ \mathsf{T} = \left[\begin{array}{cc} \mathsf{\Lambda}_R & \mathsf{0} \\ \mathsf{\Lambda}_Z & \mathsf{I}_{m-s} \end{array} \right]$$

$$\mathsf{T}^{-1} = \left[\begin{array}{cc} \mathsf{\Lambda}_R^{-1} & \mathsf{0} \\ \mathsf{\Lambda}_Z \mathsf{\Lambda}_R^{-1} & \mathsf{I}_{m-s} \end{array} \right]$$

- T^{-1} splits row-space \mathbb{R}^m of $LP(\mathbf{x}^k)$ into two vector subspaces \mathbf{V} and \mathbf{V}^{\perp} .
- Vector subspace basis Λ spans V of dimension $0 \le s \le m$.
- Vector $\mathbf{a} \in \mathbf{V}$ if and only if $\bar{\mathbf{a}}_Z = \mathbf{0}$, where $\bar{\mathbf{a}} = \mathbf{T}^{-1}\mathbf{a} = \begin{bmatrix} \bar{\mathbf{a}}_R \\ \mathbf{0} \end{bmatrix}$.
- $\emptyset \subseteq S \subseteq B$: index subset of basic columns spanned by Λ .
- Algorithmic properties derived according to subset S.

At iteration k, a dynamic Dantzig-Wolfe decomposition

$$\begin{array}{lll} \mathbf{c}^{\intercal}\mathbf{x}^{k} + & \min \ \mathbf{c}_{S}^{\intercal}(\vec{y}_{S} - \mathbf{\bar{y}}_{S}) \ + & \mathbf{c}_{\bar{S}}^{\intercal}(\vec{y}_{\bar{S}} - \mathbf{\bar{y}}_{\bar{S}}) \\ & \text{st.} & \mathbf{A}_{RS}(\vec{y}_{S} - \mathbf{\bar{y}}_{S}) \ + & \mathbf{A}_{R\bar{S}}(\vec{y}_{\bar{S}} - \mathbf{\bar{y}}_{\bar{S}}) \ & = \mathbf{0} \ [\pi_{R}] \\ & & \bar{\mathbf{A}}_{Z\bar{S}}(\mathbf{\bar{y}}_{\bar{S}} - \mathbf{\bar{y}}_{\bar{S}}) \ & = \mathbf{0} \\ & & \bar{\mathbf{y}}_{S} \geq \mathbf{0}, \ \mathbf{\bar{y}}_{S} \geq \mathbf{0} \\ & & \mathbf{\bar{y}}_{\bar{S}} \geq \mathbf{0}, \ \mathbf{\bar{y}}_{\bar{S}} \geq \mathbf{0} \\ & & \bar{\mathbf{y}}_{\bar{S}_{L}} \leq \mathbf{0}, \ \mathbf{\bar{y}}_{\bar{S}_{U}} \leq \mathbf{0} \\ & & \bar{\mathbf{y}}_{S_{F}} \leq \mathbf{\bar{r}}_{S_{F}}, \ \mathbf{\bar{y}}_{S_{F}} \leq \mathbf{\bar{r}}_{\bar{S}_{F}} \\ & \bar{\mathbf{y}}_{S_{I}} \leq \mathbf{0}, \ \mathbf{\bar{y}}_{S_{U}} \leq \mathbf{0} \\ & & \bar{\mathbf{y}}_{\bar{S}_{I}} \leq \mathbf{\bar{r}}_{\bar{S}_{F}}, \ \mathbf{\bar{y}}_{\bar{S}_{F}} \leq \mathbf{\bar{r}}_{\bar{S}_{F}} \\ & \bar{\mathbf{y}}_{S_{I}} \leq \mathbf{0}, \ \mathbf{\bar{y}}_{S_{U}} \leq \mathbf{0} \\ & & \bar{\mathbf{y}}_{\bar{S}_{I}} \leq \mathbf{\bar{r}}_{\bar{S}_{I}}, \ \mathbf{\bar{y}}_{\bar{S}_{IJ}} \leq \mathbf{\bar{r}}_{\bar{S}_{IJ}} \end{array}$$

- Pricing domain keeps the m-s equality rows corresponding to \mathbf{V}^{\perp} and $\mathbf{\bar{y}}_{\bar{5}} \geq \mathbf{0}, \ \mathbf{\bar{y}}_{\bar{5}} \geq \mathbf{0}, \ \mathbf{\bar{y}}_{\bar{5}_{ij}} = \mathbf{0}, \ \mathbf{\bar{y}}_{\bar{5}_{ij}} = \mathbf{0}.$
- Master domain keeps the s equality rows corresponding to V, and all remaining lower and upper bounds of variables.

Special case $\#4: F \subseteq S \subseteq B$

Dynamic Constraint Aggregation for Set Partitioning (Elhallaoui et al. 2005)

The partition of the row-set is derived from the f groups of identical rows of A_F .

Properties

 $F \subseteq S \subseteq B$ for fractional solutions but S = F (or $S = \emptyset$) for binary solutions. Degenerate pivots may occur, $\rho \ge 0$, if r > f.

A_{S} : A SUBSET OF s COLUMNS OF THE SIMPLEX BASIS A_{B}

