

Formalization of "Knight's Tour Revisited"

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Abstract

This is a formalization of [1]. In [1] the existence of a Knight's path is proved for arbitrary $n \times m$ -boards with $\min(n, m) \geq 5$. If $n \cdot m$ is even, then there exists a Knight's circuit.

A Knight's Path is a sequence of moves of a Knight on a chessboard s.t. the Knight visits every square of a chessboard exactly once. Finding a Knight's path is an instance of the Hamiltonian path problem.

During the formalization two mistakes in the original proof in [1] were discovered. These mistakes are corrected in this formalization.

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theory KnightsTour
  imports Main
begin

```

1 Introduction and Definitions

A Knight's path is a sequence of moves on a chessboard s.t. every step in sequence is a valid move for a Knight and that the Knight visits every square on the boards exactly once. A Knight is a chess figure that is only able to move two squares vertically and one square horizontally or two squares horizontally and one square vertically. Finding a Knight's path is an instance of the Hamiltonian Path Problem. A Knight's circuit is a Knight's path, where additionally the Knight can move from the last square to the first square of the path, forming a loop.

[1] proves the existence of a Knight's path on a $n \times m$ -board for sufficiently large n and m . The main idea for the proof is to inductively construct a Knight's path for the $n \times m$ -board from a few pre-computed Knight's paths for small boards, i.e. 5×5 , 8×6 , ..., 8×9 . The paths for small boards are transformed (i.e. transpose, mirror, translate) and concatenated to create paths for larger boards.

While formalizing the proofs I discovered two mistakes in the original proof in [1]: (i) the pre-computed path for the 6×6 -board that ends in the upper-left (in Figure 2) and (ii) the pre-computed path for the 8×8 -board that ends in the upper-left (in Figure 5) are incorrect. I.e. on the 6×6 -board the Knight cannot step from square 26 to square 27; in the 8×8 -board the Knight cannot step from square 27 to square 28. In this formalization I have replaced the two incorrect paths with correct paths.

A square on a board is identified by its coordinates.

type-synonym $square = int \times int$

A board is represented as a set of squares. Note, that this allows boards to have an arbitrary shape and do not necessarily need to be rectangular.

type-synonym $board = square\ set$

A (rectangular) $(n \times m)$ -board is the set of all squares (i, j) where $1 \leq i \leq n$

and $1 \leq j \leq m$. $(1,1)$ is the lower-left corner, and (n,m) is the upper-right corner.

definition $board :: nat \Rightarrow nat \Rightarrow board$ **where**

$board\ n\ m = \{(i,j) \mid i\ j.\ 1 \leq i \wedge i \leq int\ n \wedge 1 \leq j \wedge j \leq int\ m\}$

A path is a sequence of steps on a board. A path is represented by the list of visited squares on the board. Each square on the $(n \times m)$ -board is identified by its coordinates (i,j) .

type-synonym $path = square\ list$

A Knight can only move two squares vertically and one square horizontally or two squares horizontally and one square vertically. Thus, a knight at position (i,j) can only move to $(i \pm 1, j \pm 2)$ or $(i \pm 2, j \pm 1)$.

definition $valid-step :: square \Rightarrow square \Rightarrow bool$ **where**

$valid-step\ s_i\ s_j \equiv (case\ s_i\ of\ (i,j) \Rightarrow s_j \in \{(i+1,j+2),(i-1,j+2),(i+1,j-2),(i-1,j-2),$
 $(i+2,j+1),(i-2,j+1),(i+2,j-1),(i-2,j-1)\})$

Now we define an inductive predicate that characterizes a Knight's path. A square s_i can be pre-pended to a current Knight's path $s_j \# ps$ if (i) there is a valid step from the square s_i to the first square s_j of the current path and (ii) the square s_i has not been visited yet.

inductive $knight\text{-}path :: board \Rightarrow path \Rightarrow bool$ **where**

$knight\text{-}path\ \{s_i\}\ [s_i]$
 $\mid s_i \notin b \Longrightarrow valid-step\ s_i\ s_j \Longrightarrow knight\text{-}path\ b\ (s_j \# ps) \Longrightarrow knight\text{-}path\ (b \cup \{s_i\})$
 $(s_i \# s_j \# ps)$

code-pred $knight\text{-}path\ \langle proof \rangle$

A sequence is a Knight's circuit iff the sequence is a Knight's path and there is a valid step from the last square to the first square.

definition $knight\text{-}circuit\ b\ ps \equiv (knight\text{-}path\ b\ ps \wedge valid-step\ (last\ ps)\ (hd\ ps))$

2 Executable Checker for a Knight's Path

This section gives the implementation and correctness-proof for an executable checker for a knight's-path wrt. the definition $knight\text{-}path$.

2.1 Implementation of an Executable Checker

fun $row-exec :: nat \Rightarrow int\ set$ **where**

$row-exec\ 0 = \{\}$
 $\mid row-exec\ m = insert\ (int\ m)\ (row-exec\ (m-1))$

fun $board-exec-aux :: nat \Rightarrow int\ set \Rightarrow board$ **where**

$board-exec-aux\ 0\ M = \{\}$

| *board-exec-aux* $k\ M = \{(int\ k, j) \mid j. j \in M\} \cup \text{board-exec-aux } (k-1)\ M$

Compute a board.

fun *board-exec* :: *nat* \Rightarrow *nat* \Rightarrow *board* **where**
board-exec $n\ m = \text{board-exec-aux } n\ (\text{row-exec } m)$

fun *step-checker* :: *square* \Rightarrow *square* \Rightarrow *bool* **where**
step-checker $(i, j)\ (i', j') =$
 $((i+1, j+2) = (i', j') \vee (i-1, j+2) = (i', j') \vee (i+1, j-2) = (i', j') \vee (i-1, j-2)$
 $= (i', j')$
 $\vee (i+2, j+1) = (i', j') \vee (i-2, j+1) = (i', j') \vee (i+2, j-1) = (i', j') \vee (i-2, j-1)$
 $= (i', j'))$

fun *path-checker* :: *board* \Rightarrow *path* \Rightarrow *bool* **where**
path-checker $b\ [] = \text{False}$
| *path-checker* $b\ [s_i] = (\{s_i\} = b)$
| *path-checker* $b\ (s_i \# s_j \# ps) = (s_i \in b \wedge \text{step-checker } s_i\ s_j \wedge \text{path-checker } (b - \{s_i\})\ (s_j \# ps))$

fun *circuit-checker* :: *board* \Rightarrow *path* \Rightarrow *bool* **where**
circuit-checker $b\ ps = (\text{path-checker } b\ ps \wedge \text{step-checker } (\text{last } ps)\ (\text{hd } ps))$

2.2 Correctness Proof of the Executable Checker

lemma *row-exec-leq*: $j \in \text{row-exec } m \longleftrightarrow 1 \leq j \wedge j \leq \text{int } m$
 $\langle \text{proof} \rangle$

lemma *board-exec-aux-leq-mem*: $(i, j) \in \text{board-exec-aux } k\ M \longleftrightarrow 1 \leq i \wedge i \leq \text{int } k \wedge j \in M$
 $\langle \text{proof} \rangle$

lemma *board-exec-leq*: $(i, j) \in \text{board-exec } n\ m \longleftrightarrow 1 \leq i \wedge i \leq \text{int } n \wedge 1 \leq j \wedge j \leq \text{int } m$
 $\langle \text{proof} \rangle$

lemma *board-exec-correct*: $\text{board } n\ m = \text{board-exec } n\ m$
 $\langle \text{proof} \rangle$

lemma *step-checker-correct*: $\text{step-checker } s_i\ s_j \longleftrightarrow \text{valid-step } s_i\ s_j$
 $\langle \text{proof} \rangle$

lemma *step-checker-rev*: $\text{step-checker } (i, j)\ (i', j') \implies \text{step-checker } (i', j')\ (i, j)$
 $\langle \text{proof} \rangle$

lemma *knights-path-intro-rev*:
assumes $s_i \in b\ \text{valid-step } s_i\ s_j\ \text{knights-path } (b - \{s_i\})\ (s_j \# ps)$
shows $\text{knights-path } b\ (s_i \# s_j \# ps)$
 $\langle \text{proof} \rangle$

Final correctness corollary for the executable checker *path-checker*.

lemma *path-checker-correct*: $\text{path-checker } b \text{ } ps \longleftrightarrow \text{knight-path } b \text{ } ps$
 <proof>

corollary *knight-path-exec-simp*: $\text{knight-path } (\text{board } n \text{ } m) \text{ } ps \longleftrightarrow \text{path-checker } (\text{board-exec } n \text{ } m) \text{ } ps$
 <proof>

lemma *circuit-checker-correct*: $\text{circuit-checker } b \text{ } ps \longleftrightarrow \text{knight-circuit } b \text{ } ps$
 <proof>

corollary *knight-circuit-exec-simp*:
 $\text{knight-circuit } (\text{board } n \text{ } m) \text{ } ps \longleftrightarrow \text{circuit-checker } (\text{board-exec } n \text{ } m) \text{ } ps$
 <proof>

3 Basic Properties of *knight-path* and *knight-circuit*

lemma *board-leq-subset*: $n_1 \leq n_2 \wedge m_1 \leq m_2 \implies \text{board } n_1 \text{ } m_1 \subseteq \text{board } n_2 \text{ } m_2$
 <proof>

lemma *finite-row-exec*: $\text{finite } (\text{row-exec } m)$
 <proof>

lemma *finite-board-exec-aux*: $\text{finite } M \implies \text{finite } (\text{board-exec-aux } n \text{ } M)$
 <proof>

lemma *board-finite*: $\text{finite } (\text{board } n \text{ } m)$
 <proof>

lemma *card-row-exec*: $\text{card } (\text{row-exec } m) = m$
 <proof>

lemma *set-comp-ins*:
 $\{(k,j) \mid j. j \in \text{insert } x \text{ } M\} = \text{insert } (k,x) \{(k,j) \mid j. j \in M\} \text{ (is } ?Mi = ?iM)$
 <proof>

lemma *finite-card-set-comp*: $\text{finite } M \implies \text{card } \{(k,j) \mid j. j \in M\} = \text{card } M$
 <proof>

lemma *card-board-exec-aux*: $\text{finite } M \implies \text{card } (\text{board-exec-aux } k \text{ } M) = k * \text{card } M$
 <proof>

lemma *card-board*: $\text{card } (\text{board } n \text{ } m) = n * m$
 <proof>

lemma *knight-path-board-non-empty*: $\text{knight-path } b \text{ } ps \implies b \neq \{\}$
 <proof>

lemma *knight-path-board-m-n-geq-1*: $\text{knight-path } (\text{board } n \text{ } m) \text{ } ps \implies \min n \text{ } m \geq 1$

$\langle proof \rangle$

lemma *knight-path-non-nil*: $knight-path\ b\ ps \implies ps \neq []$
 $\langle proof \rangle$

lemma *knight-path-set-eq*: $knight-path\ b\ ps \implies set\ ps = b$
 $\langle proof \rangle$

lemma *knight-path-subset*:
 $knight-path\ b_1\ ps_1 \implies knight-path\ b_2\ ps_2 \implies set\ ps_1 \subseteq set\ ps_2 \iff b_1 \subseteq b_2$
 $\langle proof \rangle$

lemma *knight-path-board-unique*: $knight-path\ b_1\ ps \implies knight-path\ b_2\ ps \implies b_1 = b_2$
 $\langle proof \rangle$

lemma *valid-step-neg*: $valid-step\ s_i\ s_j \implies s_i \neq s_j$
 $\langle proof \rangle$

lemma *valid-step-non-transitive*: $valid-step\ s_i\ s_j \implies valid-step\ s_j\ s_k \implies \neg valid-step\ s_i\ s_k$
 $\langle proof \rangle$

lemma *knight-path-distinct*: $knight-path\ b\ ps \implies distinct\ ps$
 $\langle proof \rangle$

lemma *knight-path-length*: $knight-path\ b\ ps \implies length\ ps = card\ b$
 $\langle proof \rangle$

lemma *knight-path-take*:
 assumes $knight-path\ b\ ps\ 0 < k\ k < length\ ps$
 shows $knight-path\ (set\ (take\ k\ ps))\ (take\ k\ ps)$
 $\langle proof \rangle$

lemma *knight-path-drop*:
 assumes $knight-path\ b\ ps\ 0 < k\ k < length\ ps$
 shows $knight-path\ (set\ (drop\ k\ ps))\ (drop\ k\ ps)$
 $\langle proof \rangle$

A Knight's path can be split to form two new disjoint Knight's paths.

corollary *knight-path-split*:
 assumes $knight-path\ b\ ps\ 0 < k\ k < length\ ps$
 shows
 $\exists b_1\ b_2. knight-path\ b_1\ (take\ k\ ps) \wedge knight-path\ b_2\ (drop\ k\ ps) \wedge b_1 \cup b_2 = b$
 $\wedge b_1 \cap b_2 = \{\}$
 $\langle proof \rangle$

Append two disjoint Knight's paths.

corollary *knight-path-append*:

assumes *knight-path* b_1 ps_1 *knight-path* b_2 ps_2 $b_1 \cap b_2 = \{\}$ *valid-step* (*last* ps_1) (*hd* ps_2)
shows *knight-path* $(b_1 \cup b_2)$ $(ps_1 @ ps_2)$
<proof>

lemma *valid-step-rev*: *valid-step* s_i $s_j \implies$ *valid-step* s_j s_i
<proof>

Reverse a Knight's path.

corollary *knight-path-rev*:
assumes *knight-path* b ps
shows *knight-path* b (*rev* ps)
<proof>

Reverse a Knight's circuit.

corollary *knight-circuit-rev*:
assumes *knight-circuit* b ps
shows *knight-circuit* b (*rev* ps)
<proof>

lemma *knight-circuit-rotate1*:
assumes *knight-circuit* b $(s_i \# ps)$
shows *knight-circuit* b $(ps @ [s_i])$
<proof>

A Knight's circuit can be rotated to start at any square on the board.

lemma *knight-circuit-rotate-to*:
assumes *knight-circuit* b ps *hd* $(\text{drop } k \text{ } ps) = s_i$ $k < \text{length } ps$
shows $\exists ps'. \text{ knight-circuit } b \text{ } ps' \wedge \text{hd } ps' = s_i$
<proof>

For positive boards (1,1) can only have (2,3) and (3,2) as a neighbour.

lemma *valid-step-1-1*:
assumes *valid-step* $(1,1)$ (i,j) $i > 0$ $j > 0$
shows $(i,j) = (2,3) \vee (i,j) = (3,2)$
<proof>

lemma *list-len-g-1-split*: *length* $xs > 1 \implies \exists x_1 \ x_2 \ xs'. xs = x_1 \# x_2 \# xs'$
<proof>

lemma *list-len-g-3-split*: *length* $xs > 3 \implies \exists x_1 \ x_2 \ xs' \ x_3. xs = x_1 \# x_2 \# xs' @ [x_3]$
<proof>

Any Knight's circuit on a positive board can be rotated to start with (1,1) and end with (3,2).

corollary *rotate-knights-circuit*:

assumes *knights-circuit* (board n m) ps $\min n\ m \geq 5$

shows $\exists ps. \text{knights-circuit (board } n\ m)\ ps \wedge \text{hd } ps = (1,1) \wedge \text{last } ps = (3,2)$

$\langle \text{proof} \rangle$

4 Transposing Paths and Boards

4.1 Implementation of Path and Board Transposition

definition *transpose-square* $s_i = (\text{case } s_i \text{ of } (i,j) \Rightarrow (j,i))$

fun *transpose* :: *path* \Rightarrow *path* **where**

transpose [] = []

| *transpose* ($s_i \# ps$) = (*transpose-square* s_i) # *transpose* ps

definition *transpose-board* :: *board* \Rightarrow *board* **where**

transpose-board $b \equiv \{(j,i) \mid i\ j. (i,j) \in b\}$

4.2 Correctness of Path and Board Transposition

lemma *transpose2*: *transpose-square* (*transpose-square* s_i) = s_i

$\langle \text{proof} \rangle$

lemma *transpose-nil*: $ps = [] \iff \text{transpose } ps = []$

$\langle \text{proof} \rangle$

lemma *transpose-length*: $\text{length } ps = \text{length } (\text{transpose } ps)$

$\langle \text{proof} \rangle$

lemma *hd-transpose*: $ps \neq [] \implies \text{hd } (\text{transpose } ps) = \text{transpose-square } (\text{hd } ps)$

$\langle \text{proof} \rangle$

lemma *last-transpose*: $ps \neq [] \implies \text{last } (\text{transpose } ps) = \text{transpose-square } (\text{last } ps)$

$\langle \text{proof} \rangle$

lemma *take-transpose*:

shows $\text{take } k\ (\text{transpose } ps) = \text{transpose } (\text{take } k\ ps)$

$\langle \text{proof} \rangle$

lemma *drop-transpose*:

shows $\text{drop } k\ (\text{transpose } ps) = \text{transpose } (\text{drop } k\ ps)$

$\langle \text{proof} \rangle$

lemma *transpose-board-correct*: $s_i \in b \iff (\text{transpose-square } s_i) \in \text{transpose-board } b$

$\langle \text{proof} \rangle$

lemma *transpose-board*: *transpose-board* (board n m) = board m n

$\langle \text{proof} \rangle$

lemma *insert-transpose-board*:

insert (transpose-square s_i) (transpose-board b) = transpose-board (insert s_i b)
 ⟨proof⟩

lemma *transpose-board2*: transpose-board (transpose-board b) = b
 ⟨proof⟩

lemma *transpose-union*: transpose-board ($b_1 \cup b_2$) = transpose-board $b_1 \cup$ transpose-board b_2
 ⟨proof⟩

lemma *transpose-valid-step*:

valid-step s_i $s_j \iff$ *valid-step* (transpose-square s_i) (transpose-square s_j)
 ⟨proof⟩

lemma *transpose-knights-path'*:

assumes *knights-path* b ps
shows *knights-path* (transpose-board b) (transpose ps)
 ⟨proof⟩

corollary *transpose-knights-path*:

assumes *knights-path* (board n m) ps
shows *knights-path* (board m n) (transpose ps)
 ⟨proof⟩

corollary *transpose-knights-circuit*:

assumes *knights-circuit* (board n m) ps
shows *knights-circuit* (board m n) (transpose ps)
 ⟨proof⟩

5 Mirroring Paths and Boards

5.1 Implementation of Path and Board Mirroring

abbreviation *min1* $ps \equiv$ *Min* ((fst) ' set ps)

abbreviation *max1* $ps \equiv$ *Max* ((fst) ' set ps)

abbreviation *min2* $ps \equiv$ *Min* ((snd) ' set ps)

abbreviation *max2* $ps \equiv$ *Max* ((snd) ' set ps)

definition *mirror1-square* :: $int \Rightarrow square \Rightarrow square$ **where**

mirror1-square n $s_i =$ (case s_i of $(i,j) \Rightarrow (n-i,j)$)

fun *mirror1-aux* :: $int \Rightarrow path \Rightarrow path$ **where**

mirror1-aux n [] = []

| *mirror1-aux* n ($s_i \# ps$) = (*mirror1-square* n s_i) # *mirror1-aux* n ps

definition *mirror1* $ps =$ *mirror1-aux* (*max1* ps + *min1* ps) ps

definition *mirror1-board* :: *int* \Rightarrow *board* \Rightarrow *board* **where**
mirror1-board *n* *b* $\equiv \{ \text{mirror1-square } n \ s_i \mid s_i. \ s_i \in b \}$

definition *mirror2-square* :: *int* \Rightarrow *square* \Rightarrow *square* **where**
mirror2-square *m* *s_i* = (case *s_i* of (*i*,*j*) \Rightarrow (*i*,*m-j*))

fun *mirror2-aux* :: *int* \Rightarrow *path* \Rightarrow *path* **where**
mirror2-aux *m* [] = []
| *mirror2-aux* *m* (*s_i*#*ps*) = (*mirror2-square* *m* *s_i*)#*mirror2-aux* *m* *ps*

definition *mirror2* *ps* = *mirror2-aux* (*max2* *ps* + *min2* *ps*) *ps*

definition *mirror2-board* :: *int* \Rightarrow *board* \Rightarrow *board* **where**
mirror2-board *m* *b* $\equiv \{ \text{mirror2-square } m \ s_i \mid s_i. \ s_i \in b \}$

5.2 Correctness of Path and Board Mirroring

lemma *mirror1-board-id*: *mirror1-board* (*int* *n+1*) (*board* *n* *m*) = *board* *n* *m* (**is** -
= ?*b*)
⟨*proof*⟩

lemma *mirror2-board-id*: *mirror2-board* (*int* *m+1*) (*board* *n* *m*) = *board* *n* *m* (**is** -
= ?*b*)
⟨*proof*⟩

lemma *knights-path-min1*: *knights-path* (*board* *n* *m*) *ps* \implies *min1* *ps* = 1
⟨*proof*⟩

lemma *knights-path-min2*: *knights-path* (*board* *n* *m*) *ps* \implies *min2* *ps* = 1
⟨*proof*⟩

lemma *knights-path-max1*: *knights-path* (*board* *n* *m*) *ps* \implies *max1* *ps* = *int* *n*
⟨*proof*⟩

lemma *knights-path-max2*: *knights-path* (*board* *n* *m*) *ps* \implies *max2* *ps* = *int* *m*
⟨*proof*⟩

lemma *mirror1-aux-nil*: *ps* = [] \longleftrightarrow *mirror1-aux* *m* *ps* = []
⟨*proof*⟩

lemma *mirror1-nil*: *ps* = [] \longleftrightarrow *mirror1* *ps* = []
⟨*proof*⟩

lemma *mirror2-aux-nil*: *ps* = [] \longleftrightarrow *mirror2-aux* *m* *ps* = []
⟨*proof*⟩

lemma *mirror2-nil*: *ps* = [] \longleftrightarrow *mirror2* *ps* = []
⟨*proof*⟩

lemma *length-mirror1-aux*: $\text{length } ps = \text{length } (\text{mirror1-aux } n \ ps)$
 ⟨proof⟩

lemma *length-mirror1*: $\text{length } ps = \text{length } (\text{mirror1 } ps)$
 ⟨proof⟩

lemma *length-mirror2-aux*: $\text{length } ps = \text{length } (\text{mirror2-aux } n \ ps)$
 ⟨proof⟩

lemma *length-mirror2*: $\text{length } ps = \text{length } (\text{mirror2 } ps)$
 ⟨proof⟩

lemma *mirror1-board-iff*: $s_i \notin b \iff \text{mirror1-square } n \ s_i \notin \text{mirror1-board } n \ b$
 ⟨proof⟩

lemma *mirror2-board-iff*: $s_i \notin b \iff \text{mirror2-square } n \ s_i \notin \text{mirror2-board } n \ b$
 ⟨proof⟩

lemma *insert-mirror1-board*:
 $\text{insert } (\text{mirror1-square } n \ s_i) \ (\text{mirror1-board } n \ b) = \text{mirror1-board } n \ (\text{insert } s_i \ b)$
 ⟨proof⟩

lemma *insert-mirror2-board*:
 $\text{insert } (\text{mirror2-square } n \ s_i) \ (\text{mirror2-board } n \ b) = \text{mirror2-board } n \ (\text{insert } s_i \ b)$
 ⟨proof⟩

lemma $(i::\text{int}) = i' + 1 \implies n - i = n - (i' + 1)$
 ⟨proof⟩

lemma *valid-step-mirror1*:
 $\text{valid-step } s_i \ s_j \iff \text{valid-step } (\text{mirror1-square } n \ s_i) \ (\text{mirror1-square } n \ s_j)$
 ⟨proof⟩

lemma *valid-step-mirror2*:
 $\text{valid-step } s_i \ s_j \iff \text{valid-step } (\text{mirror2-square } m \ s_i) \ (\text{mirror2-square } m \ s_j)$
 ⟨proof⟩

lemma *hd-mirror1*:
assumes *knights-path* (board $n \ m$) ps $\text{hd } ps = (i, j)$
shows $\text{hd } (\text{mirror1 } ps) = (\text{int } n + 1 - i, j)$
 ⟨proof⟩

lemma *last-mirror1-aux*:
assumes $ps \neq []$ $\text{last } ps = (i, j)$
shows $\text{last } (\text{mirror1-aux } n \ ps) = (n - i, j)$
 ⟨proof⟩

lemma *last-mirror1*:
assumes *knights-path* (board $n \ m$) ps $\text{last } ps = (i, j)$

shows $last (mirror1\ ps) = (int\ n+1-i, j)$
 $\langle proof \rangle$

lemma *hd-mirror2*:

assumes *knight's-path* (board $n\ m$) $ps\ hd\ ps = (i, j)$
shows $hd (mirror2\ ps) = (i, int\ m+1-j)$
 $\langle proof \rangle$

lemma *last-mirror2-aux*:

assumes $ps \neq []$ $last\ ps = (i, j)$
shows $last (mirror2-aux\ m\ ps) = (i, m-j)$
 $\langle proof \rangle$

lemma *last-mirror2*:

assumes *knight's-path* (board $n\ m$) $ps\ last\ ps = (i, j)$
shows $last (mirror2\ ps) = (i, int\ m+1-j)$
 $\langle proof \rangle$

lemma *mirror1-aux-knight's-path*:

assumes *knight's-path* $b\ ps$
shows *knight's-path* (mirror1-board $n\ b$) (mirror1-aux $n\ ps$)
 $\langle proof \rangle$

corollary *mirror1-knight's-path*:

assumes *knight's-path* (board $n\ m$) ps
shows *knight's-path* (board $n\ m$) (mirror1 ps)
 $\langle proof \rangle$

lemma *mirror2-aux-knight's-path*:

assumes *knight's-path* $b\ ps$
shows *knight's-path* (mirror2-board $n\ b$) (mirror2-aux $n\ ps$)
 $\langle proof \rangle$

corollary *mirror2-knight's-path*:

assumes *knight's-path* (board $n\ m$) ps
shows *knight's-path* (board $n\ m$) (mirror2 ps)
 $\langle proof \rangle$

5.3 Rotate Knight's Paths

Transposing (*transpose*) and mirroring (along first axis *mirror1*) a Knight's path preserves the Knight's path's property. Tranpose+Mirror1 equals a 90deg-clockwise turn.

corollary *rot90-knight's-path*:

assumes *knight's-path* (board $n\ m$) ps
shows *knight's-path* (board $m\ n$) (mirror1 (transpose ps))
 $\langle proof \rangle$

lemma *hd-rot90-knight's-path*:

assumes *knight's-path* (board *n m*) *ps* *hd ps* = (*i,j*)
shows *hd* (*mirror1* (*transpose ps*)) = (*int m+1-j,i*)
 ⟨*proof*⟩

lemma *last-rot90-knight's-path*:
assumes *knight's-path* (board *n m*) *ps* *last ps* = (*i,j*)
shows *last* (*mirror1* (*transpose ps*)) = (*int m+1-j,i*)
 ⟨*proof*⟩

6 Translating Paths and Boards

When constructing knight's paths for larger boards multiple knight's paths for smaller boards are concatenated. To concatenate paths the the coordinates in the path need to be translated. Therefore, simple auxiliary functions are provided.

6.1 Implementation of Path and Board Translation

Translate the coordinates for a path by (*k₁,k₂*).

fun *trans-path* :: *int* × *int* ⇒ *path* ⇒ *path* **where**
 trans-path (*k₁,k₂*) [] = []
 | *trans-path* (*k₁,k₂*) ((*i,j*)#*xs*) = (*i+k₁,j+k₂*)#(*trans-path* (*k₁,k₂*) *xs*)

Translate the coordinates of a board by (*k₁,k₂*).

definition *trans-board* :: *int* × *int* ⇒ *board* ⇒ *board* **where**
 trans-board *t b* ≡ (*case t of* (*k₁,k₂*) ⇒ {(*i+k₁,j+k₂*)|*i j. (i,j) ∈ b*})

6.2 Correctness of Path and Board Translation

lemma *trans-path-length*: *length ps* = *length* (*trans-path* (*k₁,k₂*) *ps*)
 ⟨*proof*⟩

lemma *trans-path-non-nil*: *ps* ≠ [] ⇒ *trans-path* (*k₁,k₂*) *ps* ≠ []
 ⟨*proof*⟩

lemma *trans-path-correct*: (*i,j*) ∈ *set ps* ⇔ (*i+k₁,j+k₂*) ∈ *set* (*trans-path* (*k₁,k₂*) *ps*)
 ⟨*proof*⟩

lemma *trans-path-non-nil-last*:
 ps ≠ [] ⇒ *last* (*trans-path* (*k₁,k₂*) *ps*) = *last* (*trans-path* (*k₁,k₂*) ((*i,j*)#*ps*))
 ⟨*proof*⟩

lemma *hd-trans-path*:
assumes *ps* ≠ [] *hd ps* = (*i,j*)
shows *hd* (*trans-path* (*k₁,k₂*) *ps*) = (*i+k₁,j+k₂*)
 ⟨*proof*⟩

lemma *last-trans-path*:

assumes $ps \neq []$ $last\ ps = (i,j)$

shows $last\ (trans\text{-}path\ (k_1,k_2)\ ps) = (i+k_1,j+k_2)$

$\langle proof \rangle$

lemma *take-trans*:

shows $take\ k\ (trans\text{-}path\ (k_1,k_2)\ ps) = trans\text{-}path\ (k_1,k_2)\ (take\ k\ ps)$

$\langle proof \rangle$

lemma *drop-trans*:

shows $drop\ k\ (trans\text{-}path\ (k_1,k_2)\ ps) = trans\text{-}path\ (k_1,k_2)\ (drop\ k\ ps)$

$\langle proof \rangle$

lemma *trans-board-correct*: $(i,j) \in b \longleftrightarrow (i+k_1,j+k_2) \in trans\text{-}board\ (k_1,k_2)\ b$

$\langle proof \rangle$

lemma *board-subset*: $n_1 \leq n_2 \implies m_1 \leq m_2 \implies board\ n_1\ m_1 \subseteq board\ n_2\ m_2$

$\langle proof \rangle$

Board concatenation

corollary *board-concat*:

shows $board\ n\ m_1 \cup trans\text{-}board\ (0,int\ m_1)\ (board\ n\ m_2) = board\ n\ (m_1+m_2)$

(**is** $?b1 \cup ?b2 = ?b$)

$\langle proof \rangle$

lemma *transpose-trans-board*:

$transpose\text{-}board\ (trans\text{-}board\ (k_1,k_2)\ b) = trans\text{-}board\ (k_2,k_1)\ (transpose\text{-}board\ b)$

$\langle proof \rangle$

corollary *board-concatT*:

shows $board\ n_1\ m \cup trans\text{-}board\ (int\ n_1,0)\ (board\ n_2\ m) = board\ (n_1+n_2)\ m$ (**is**

$?b_1 \cup ?b_2 = ?b$)

$\langle proof \rangle$

lemma *trans-valid-step*:

$valid\text{-}step\ (i,j)\ (i',j') \implies valid\text{-}step\ (i+k_1,j+k_2)\ (i'+k_1,j'+k_2)$

$\langle proof \rangle$

Translating a path and a boards preserves the validity.

lemma *trans-knights-path*:

assumes $knights\text{-}path\ b\ ps$

shows $knights\text{-}path\ (trans\text{-}board\ (k_1,k_2)\ b)\ (trans\text{-}path\ (k_1,k_2)\ ps)$

$\langle proof \rangle$

Predicate that indicates if two squares s_i and s_j are adjacent in ps .

definition *step-in* :: $path \Rightarrow square \Rightarrow square \Rightarrow bool$ **where**

$step\text{-}in\ ps\ s_i\ s_j \equiv (\exists k. 0 < k \wedge k < length\ ps \wedge last\ (take\ k\ ps) = s_i \wedge hd\ (drop\ k\ ps) = s_j)$

lemma *step-in-Cons*: $\text{step-in } ps \ s_i \ s_j \implies \text{step-in } (s_k \# ps) \ s_i \ s_j$
 $\langle \text{proof} \rangle$

lemma *step-in-append*: $\text{step-in } ps \ s_i \ s_j \implies \text{step-in } (ps @ ps') \ s_i \ s_j$
 $\langle \text{proof} \rangle$

lemma *step-in-prepend*: $\text{step-in } ps \ s_i \ s_j \implies \text{step-in } (ps' @ ps) \ s_i \ s_j$
 $\langle \text{proof} \rangle$

lemma *step-in-valid-step*: $\text{knight-path } b \ ps \implies \text{step-in } ps \ s_i \ s_j \implies \text{valid-step } s_i \ s_j$
 $\langle \text{proof} \rangle$

lemma *trans-step-in*:
 $\text{step-in } ps \ (i, j) \ (i', j') \implies \text{step-in } (\text{trans-path } (k_1, k_2) \ ps) \ (i + k_1, j + k_2) \ (i' + k_1, j' + k_2)$
 $\langle \text{proof} \rangle$

lemma *transpose-step-in*:
 $\text{step-in } ps \ s_i \ s_j \implies \text{step-in } (\text{transpose } ps) \ (\text{transpose-square } s_i) \ (\text{transpose-square } s_j)$
 $(\text{is } - \implies \text{step-in } ?psT \ ?s_iT \ ?s_jT)$
 $\langle \text{proof} \rangle$

lemma *hd-take*: $0 < k \implies \text{hd } xs = \text{hd } (\text{take } k \ xs)$
 $\langle \text{proof} \rangle$

lemma *last-drop*: $k < \text{length } xs \implies \text{last } xs = \text{last } (\text{drop } k \ xs)$
 $\langle \text{proof} \rangle$

6.3 Concatenate Knight's Paths and Circuits

Concatenate two knight's path on a $n \times m$ -board along the 2nd axis if the first path contains the step $s_i \rightarrow s_j$ and there are valid steps $s_i \rightarrow \text{hd } ps_2'$ and $s_j \rightarrow \text{last } ps_2'$, where ps_2' is ps_2 is translated by m_1 . An arbitrary step in ps_2 is preserved.

corollary *knight-path-split-concat-si-prev*:

assumes $\text{knight-path } (\text{board } n \ m_1) \ ps_1 \ \text{knight-path } (\text{board } n \ m_2) \ ps_2$
 $\text{step-in } ps_1 \ s_i \ s_j \ \text{hd } ps_2 = (i_h, j_h) \ \text{last } ps_2 = (i_l, j_l) \ \text{step-in } ps_2 \ (i, j) \ (i', j')$
 $\text{valid-step } s_i \ (i_h, \text{int } m_1 + j_h) \ \text{valid-step } (i_l, \text{int } m_1 + j_l) \ s_j$
shows $\exists ps. \text{knight-path } (\text{board } n \ (m_1 + m_2)) \ ps \wedge \text{hd } ps = \text{hd } ps_1$
 $\wedge \text{last } ps = \text{last } ps_1 \wedge \text{step-in } ps \ (i, \text{int } m_1 + j) \ (i', \text{int } m_1 + j')$
 $\langle \text{proof} \rangle$

lemma *len1-hd-last*: $\text{length } xs = 1 \implies \text{hd } xs = \text{last } xs$
 $\langle \text{proof} \rangle$

Weaker version of $\llbracket \text{knight-path } (\text{board } ?n \ ?m_1) \ ?ps_1; \text{knight-path } (\text{board } ?n \ ?m_2) \ ?ps_2; \text{step-in } ?ps_1 \ ?s_i \ ?s_j; \text{hd } ?ps_2 = (?i_h, ?j_h); \text{last } ?ps_2 = (?i_l, ?j_l);$

$step-in\ ?ps_2\ (?i, ?j)\ (?i', ?j')$; $valid-step\ ?s_i\ (?i_h, int\ ?m_1 + ?j_h)$; $valid-step\ (?i_l, int\ ?m_1 + ?j_l)\ ?s_j\] \Longrightarrow \exists ps. knights-path\ (board\ ?n\ (?m_1 + ?m_2))\ ps \wedge hd\ ps = hd\ ?ps_1 \wedge last\ ps = last\ ?ps_1 \wedge step-in\ ps\ (?i, int\ ?m_1 + ?j)\ (?i', int\ ?m_1 + ?j')$.

corollary *knights-path-split-concat*:

assumes $knights-path\ (board\ n\ m_1)\ ps_1\ knights-path\ (board\ n\ m_2)\ ps_2$
 $step-in\ ps_1\ s_i\ s_j\ hd\ ps_2 = (i_h, j_h)\ last\ ps_2 = (i_l, j_l)$
 $valid-step\ s_i\ (i_h, int\ m_1 + j_h)\ valid-step\ (i_l, int\ m_1 + j_l)\ s_j$
shows $\exists ps. knights-path\ (board\ n\ (m_1 + m_2))\ ps \wedge hd\ ps = hd\ ps_1 \wedge last\ ps = last\ ps_1$
 $\langle proof \rangle$

Concatenate two knight's path on a $n \times m$ -board along the 1st axis.

corollary *knights-path-split-concatT*:

assumes $knights-path\ (board\ n_1\ m)\ ps_1\ knights-path\ (board\ n_2\ m)\ ps_2$
 $step-in\ ps_1\ s_i\ s_j\ hd\ ps_2 = (i_h, j_h)\ last\ ps_2 = (i_l, j_l)$
 $valid-step\ s_i\ (int\ n_1 + i_h, j_h)\ valid-step\ (int\ n_1 + i_l, j_l)\ s_j$
shows $\exists ps. knights-path\ (board\ (n_1 + n_2)\ m)\ ps \wedge hd\ ps = hd\ ps_1 \wedge last\ ps = last\ ps_1$
 $\langle proof \rangle$

Concatenate two Knight's path along the 2nd axis. There is a valid step from the last square in the first Knight's path ps_1 to the first square in the second Knight's path ps_2 .

corollary *knights-path-concat*:

assumes $knights-path\ (board\ n\ m_1)\ ps_1\ knights-path\ (board\ n\ m_2)\ ps_2$
 $hd\ ps_2 = (i_h, j_h)\ valid-step\ (last\ ps_1)\ (i_h, int\ m_1 + j_h)$
shows $knights-path\ (board\ n\ (m_1 + m_2))\ (ps_1 @ (trans-path\ (0, int\ m_1)\ ps_2))$
 $\langle proof \rangle$

Concatenate two Knight's path along the 2nd axis. The first Knight's path end in $(2, m_1 - 1)$ (lower-right) and the second Knight's paths start in $(1, 1)$ (lower-left).

corollary *knights-path-lr-concat*:

assumes $knights-path\ (board\ n\ m_1)\ ps_1\ knights-path\ (board\ n\ m_2)\ ps_2$
 $last\ ps_1 = (2, int\ m_1 - 1)\ hd\ ps_2 = (1, 1)$
shows $knights-path\ (board\ n\ (m_1 + m_2))\ (ps_1 @ (trans-path\ (0, int\ m_1)\ ps_2))$
 $\langle proof \rangle$

Concatenate two Knight's circuits along the 2nd axis. In the first Knight's path the squares $(2, m_1 - 1)$ and $(4, m_1)$ are adjacent and the second Knight's circuit starts in $(1, 1)$ (lower-left) and end in $(3, 2)$.

corollary *knights-circuit-lr-concat*:

assumes $knights-circuit\ (board\ n\ m_1)\ ps_1\ knights-circuit\ (board\ n\ m_2)\ ps_2$
 $step-in\ ps_1\ (2, int\ m_1 - 1)\ (4, int\ m_1)$
 $hd\ ps_2 = (1, 1)\ last\ ps_2 = (3, 2)\ step-in\ ps_2\ (2, int\ m_2 - 1)\ (4, int\ m_2)$

shows $\exists ps. \text{knights-circuit } (\text{board } n \ (m_1+m_2)) \ ps \wedge \text{step-in } ps \ (2, \text{int } (m_1+m_2)-1)$
 $(4, \text{int } (m_1+m_2))$
 $\langle \text{proof} \rangle$

7 Parsing Paths

In this section functions are implemented to parse and construct paths. The parser converts the matrix representation $((\text{nat list}) \text{ list})$ used in [1] to a path (path) .

for debugging

fun *test-path* :: *path* \Rightarrow *bool* **where**
test-path (*s*_i#*s*_j#*xs*) = (*step-checker* *s*_i *s*_j \wedge *test-path* (*s*_j#*xs*))
| *test-path* - = *True*

fun *f-opt* :: (*'a* \Rightarrow *'a*) \Rightarrow *'a option* \Rightarrow *'a option* **where**
f-opt - *None* = *None*
| *f-opt* *f* (*Some a*) = *Some* (*f a*)

fun *add-opt-fst-sq* :: *int* \Rightarrow *square option* \Rightarrow *square option* **where**
add-opt-fst-sq - *None* = *None*
| *add-opt-fst-sq* *k* (*Some* (*i,j*)) = *Some* (*k+i,j*)

fun *find-k-in-col* :: *nat* \Rightarrow *nat list* \Rightarrow *int option* **where**
find-k-in-col *k* [] = *None*
| *find-k-in-col* *k* (*c*#*cs*) = (*if* *c* = *k* *then Some 1* *else f-opt* ((+) 1) (*find-k-in-col* *k* *cs*))

fun *find-k-sqr* :: *nat* \Rightarrow (*nat list*) *list* \Rightarrow *square option* **where**
find-k-sqr *k* [] = *None*
| *find-k-sqr* *k* (*r*#*rs*) = (*case find-k-in-col* *k* *r* *of*
None \Rightarrow *f-opt* ($\lambda(i,j). (i+1,j)$) (*find-k-sqr* *k* *rs*)
| *Some j* \Rightarrow *Some* (*1,j*))

Auxiliary function to easily parse pre-computed boards from paper.

fun *to-sqrs* :: *nat* \Rightarrow (*nat list*) *list* \Rightarrow *path option* **where**
to-sqrs 0 *rs* = *Some* []
| *to-sqrs* *k* *rs* = (*case find-k-sqr* *k* *rs* *of*
None \Rightarrow *None*
| *Some* *s*_i \Rightarrow *f-opt* ($\lambda ps. ps@[s_i]$) (*to-sqrs* (*k-1*) *rs*))

fun *num-elems* :: (*nat list*) *list* \Rightarrow *nat* **where**
num-elems (*r*#*rs*) = *length* *r* * *length* (*r*#*rs*)

Convert a matrix (nat list list) to a path (path) . With this function we implicitly define the lower-left corner to be $(1,1)$ and the upper-right corner to be (n,m) .

definition $to\text{-}path\ rs \equiv to\text{-}sqrs\ (num\text{-}elems\ rs)\ (rev\ rs)$

Example

value $to\text{-}path$
 $[[3,22,13,16,5],$
 $[12,17,4,21,14],$
 $[23,2,15,6,9],$
 $[18,11,8,25,20],$
 $[1,24,19,10,7::nat]]$

8 Knight's Paths for $5 \times m$ -Boards

Given here are knight's paths, $kp5xmlr$ and $kp5xmur$, for the $(5 \times m)$ -board that start in the lower-left corner for $m \in \{5, 6, 7, 8, 9\}$. The path $kp5xmlr$ ends in the lower-right corner, whereas the path $kp5xmur$ ends in the upper-right corner. The tables show the visited squares numbered in ascending order.

abbreviation $b5x5 \equiv board\ 5\ 5$

A Knight's path for the (5×5) -board that starts in the lower-left and ends in the lower-right.

3	22	13	16	5
12	17	4	21	14
23	2	15	6	9
18	11	8	25	20
1	24	19	10	7

abbreviation $kp5x5lr \equiv the\ (to\text{-}path$

$[[3,22,13,16,5],$
 $[12,17,4,21,14],$
 $[23,2,15,6,9],$
 $[18,11,8,25,20],$
 $[1,24,19,10,7]])$

lemma $kp\text{-}5x5\text{-}lr$: $knight\text{-}path\ b5x5\ kp5x5lr$
 $\langle proof \rangle$

lemma $kp\text{-}5x5\text{-}lr\text{-}hd$: $hd\ kp5x5lr = (1,1)\ \langle proof \rangle$

lemma $kp\text{-}5x5\text{-}lr\text{-}last$: $last\ kp5x5lr = (2,4)\ \langle proof \rangle$

lemma $kp\text{-}5x5\text{-}lr\text{-}non\text{-}nil$: $kp5x5lr \neq []\ \langle proof \rangle$

A Knight's path for the (5×5) -board that starts in the lower-left and ends in the upper-right.

7	12	15	20	5
16	21	6	25	14
11	8	13	4	19
22	17	2	9	24
1	10	23	18	3

abbreviation $kp5x5ur \equiv$ the (to-path

[[7,12,15,20,5],
 [16,21,6,25,14],
 [11,8,13,4,19],
 [22,17,2,9,24],
 [1,10,23,18,3]])

lemma $kp-5x5-ur$: knights-path $b5x5$ $kp5x5ur$
 $\langle proof \rangle$

lemma $kp-5x5-ur-hd$: $hd\ kp5x5ur = (1,1)$ $\langle proof \rangle$

lemma $kp-5x5-ur-last$: $last\ kp5x5ur = (4,4)$ $\langle proof \rangle$

lemma $kp-5x5-ur-non-nil$: $kp5x5ur \neq []$ $\langle proof \rangle$

abbreviation $b5x6 \equiv$ board 5 6

A Knight's path for the (5×6) -board that starts in the lower-left and ends in the lower-right.

7	14	21	28	5	12
22	27	6	13	20	29
15	8	17	24	11	4
26	23	2	9	30	19
1	16	25	18	3	10

abbreviation $kp5x6lr \equiv$ the (to-path

[[7,14,21,28,5,12],
 [22,27,6,13,20,29],
 [15,8,17,24,11,4],
 [26,23,2,9,30,19],
 [1,16,25,18,3,10]])

lemma $kp-5x6-lr$: knights-path $b5x6$ $kp5x6lr$
 $\langle proof \rangle$

lemma $kp-5x6-lr-hd$: $hd\ kp5x6lr = (1,1)$ $\langle proof \rangle$

lemma $kp-5x6-lr-last$: $last\ kp5x6lr = (2,5)$ $\langle proof \rangle$

lemma $kp-5x6-lr-non-nil$: $kp5x6lr \neq []$ $\langle proof \rangle$

A Knight's path for the (5×6) -board that starts in the lower-left and ends in the upper-right.

3	10	29	20	5	12
28	19	4	11	30	21
9	2	17	24	13	6
18	27	8	15	22	25
1	16	23	26	7	14

abbreviation $kp5x6ur \equiv$ the (to-path

$[[3,10,29,20,5,12],$
 $[28,19,4,11,30,21],$
 $[9,2,17,24,13,6],$
 $[18,27,8,15,22,25],$
 $[1,16,23,26,7,14]]$)

lemma $kp-5x6-ur$: knights-path $b5x6$ $kp5x6ur$
 $\langle proof \rangle$

lemma $kp-5x6-ur-hd$: $hd\ kp5x6ur = (1,1)$ $\langle proof \rangle$

lemma $kp-5x6-ur-last$: $last\ kp5x6ur = (4,5)$ $\langle proof \rangle$

lemma $kp-5x6-ur-non-nil$: $kp5x6ur \neq []$ $\langle proof \rangle$

abbreviation $b5x7 \equiv$ board 5 7

A Knight's path for the (5×7) -board that starts in the lower-left and ends in the lower-right.

3	12	21	30	5	14	23
20	29	4	13	22	31	6
11	2	19	32	7	24	15
28	33	10	17	26	35	8
1	18	27	34	9	16	25

abbreviation $kp5x7lr \equiv$ the (to-path

$[[3,12,21,30,5,14,23],$
 $[20,29,4,13,22,31,6],$
 $[11,2,19,32,7,24,15],$
 $[28,33,10,17,26,35,8],$
 $[1,18,27,34,9,16,25]]$)

lemma $kp-5x7-lr$: knights-path $b5x7$ $kp5x7lr$
 $\langle proof \rangle$

lemma $kp-5x7-lr-hd$: $hd\ kp5x7lr = (1,1)$ $\langle proof \rangle$

lemma $kp-5x7-lr-last$: $last\ kp5x7lr = (2,6)$ $\langle proof \rangle$

lemma *kp-5x7-lr-non-nil*: $kp5x7lr \neq []$ *<proof>*

A Knight's path for the (5×7) -board that starts in the lower-left and ends in the upper-right.

3	32	11	34	5	26	13
10	19	4	25	12	35	6
31	2	33	20	23	14	27
18	9	24	29	16	7	22
1	30	17	8	21	28	15

abbreviation *kp5x7ur* \equiv the (to-path

$[[3,32,11,34,5,26,13],$
 $[10,19,4,25,12,35,6],$
 $[31,2,33,20,23,14,27],$
 $[18,9,24,29,16,7,22],$
 $[1,30,17,8,21,28,15]])$

lemma *kp-5x7-ur*: *knight's-path b5x7 kp5x7ur*
<proof>

lemma *kp-5x7-ur-hd*: $hd\ kp5x7ur = (1,1)$ *<proof>*

lemma *kp-5x7-ur-last*: $last\ kp5x7ur = (4,6)$ *<proof>*

lemma *kp-5x7-ur-non-nil*: $kp5x7ur \neq []$ *<proof>*

abbreviation *b5x8* \equiv board 5 8

A Knight's path for the (5×8) -board that starts in the lower-left and ends in the lower-right.

3	12	37	26	5	14	17	28
34	23	4	13	36	27	6	15
11	2	35	38	25	16	29	18
22	33	24	9	20	31	40	7
1	10	21	32	39	8	19	30

abbreviation *kp5x8lr* \equiv the (to-path

$[3,12,37,26,5,14,17,28],$
 $[34,23,4,13,36,27,6,15],$
 $[11,2,35,38,25,16,29,18],$
 $[22,33,24,9,20,31,40,7],$
 $[1,10,21,32,39,8,19,30]])$

lemma *kp-5x8-lr*: *knight's-path b5x8 kp5x8lr*
<proof>

lemma *kp-5x8-lr-hd*: $hd\ kp5x8lr = (1,1)$ $\langle proof \rangle$

lemma *kp-5x8-lr-last*: $last\ kp5x8lr = (2,7)$ $\langle proof \rangle$

lemma *kp-5x8-lr-non-nil*: $kp5x8lr \neq []$ $\langle proof \rangle$

A Knight's path for the (5×8) -board that starts in the lower-left and ends in the upper-right.

33	8	17	38	35	6	15	24
18	37	34	7	16	25	40	5
9	32	29	36	39	14	23	26
30	19	2	11	28	21	4	13
1	10	31	20	3	12	27	22

abbreviation *kp5x8ur* \equiv the (to-path

$[[33,8,17,38,35,6,15,24],$
 $[18,37,34,7,16,25,40,5],$
 $[9,32,29,36,39,14,23,26],$
 $[30,19,2,11,28,21,4,13],$
 $[1,10,31,20,3,12,27,22]]$)

lemma *kp-5x8-ur*: *knight's-path* *b5x8* *kp5x8ur*
 $\langle proof \rangle$

lemma *kp-5x8-ur-hd*: $hd\ kp5x8ur = (1,1)$ $\langle proof \rangle$

lemma *kp-5x8-ur-last*: $last\ kp5x8ur = (4,7)$ $\langle proof \rangle$

lemma *kp-5x8-ur-non-nil*: $kp5x8ur \neq []$ $\langle proof \rangle$

abbreviation *b5x9* \equiv board 5 9

A Knight's path for the (5×9) -board that starts in the lower-left and ends in the lower-right.

9	4	11	16	23	42	33	36	25
12	17	8	3	32	37	24	41	34
5	10	15	20	43	22	35	26	29
18	13	2	7	38	31	28	45	40
1	6	19	14	21	44	39	30	27

abbreviation *kp5x9lr* \equiv the (to-path

$[9,4,11,16,23,42,33,36,25],$
 $[12,17,8,3,32,37,24,41,34],$
 $[5,10,15,20,43,22,35,26,29],$
 $[18,13,2,7,38,31,28,45,40],$
 $[1,6,19,14,21,44,39,30,27]]$)

lemma *kp-5x9-lr: knights-path b5x9 kp5x9lr*
 $\langle proof \rangle$

lemma *kp-5x9-lr-hd: hd kp5x9lr = (1,1) $\langle proof \rangle$*

lemma *kp-5x9-lr-last: last kp5x9lr = (2,8) $\langle proof \rangle$*

lemma *kp-5x9-lr-non-nil: kp5x9lr $\neq []$ $\langle proof \rangle$*

A Knight's path for the (5×9) -board that starts in the lower-left and ends in the upper-right.

9	4	11	16	27	32	35	40	25
12	17	8	3	36	41	26	45	34
5	10	15	20	31	28	33	24	39
18	13	2	7	42	37	22	29	44
1	6	19	14	21	30	43	38	23

abbreviation *kp5x9ur \equiv the (to-path*

[[9,4,11,16,27,32,35,40,25],
[12,17,8,3,36,41,26,45,34],
[5,10,15,20,31,28,33,24,39],
[18,13,2,7,42,37,22,29,44],
[1,6,19,14,21,30,43,38,23]])

lemma *kp-5x9-ur: knights-path b5x9 kp5x9ur*
 $\langle proof \rangle$

lemma *kp-5x9-ur-hd: hd kp5x9ur = (1,1) $\langle proof \rangle$*

lemma *kp-5x9-ur-last: last kp5x9ur = (4,8) $\langle proof \rangle$*

lemma *kp-5x9-ur-non-nil: kp5x9ur $\neq []$ $\langle proof \rangle$*

lemmas *kp-5xm-lr =*

kp-5x5-lr kp-5x5-lr-hd kp-5x5-lr-last kp-5x5-lr-non-nil
kp-5x6-lr kp-5x6-lr-hd kp-5x6-lr-last kp-5x6-lr-non-nil
kp-5x7-lr kp-5x7-lr-hd kp-5x7-lr-last kp-5x7-lr-non-nil
kp-5x8-lr kp-5x8-lr-hd kp-5x8-lr-last kp-5x8-lr-non-nil
kp-5x9-lr kp-5x9-lr-hd kp-5x9-lr-last kp-5x9-lr-non-nil

lemmas *kp-5xm-ur =*

kp-5x5-ur kp-5x5-ur-hd kp-5x5-ur-last kp-5x5-ur-non-nil
kp-5x6-ur kp-5x6-ur-hd kp-5x6-ur-last kp-5x6-ur-non-nil
kp-5x7-ur kp-5x7-ur-hd kp-5x7-ur-last kp-5x7-ur-non-nil
kp-5x8-ur kp-5x8-ur-hd kp-5x8-ur-last kp-5x8-ur-non-nil
kp-5x9-ur kp-5x9-ur-hd kp-5x9-ur-last kp-5x9-ur-non-nil

For every $5 \times m$ -board with $m \geq 5$ there exists a knight's path that starts in $(1,1)$ (bottom-left) and ends in $(2,m-1)$ (bottom-right).

lemma *knight's-path-5xm-lr-exists*:

assumes $m \geq 5$

shows $\exists ps. \text{knight's-path } (\text{board } 5 \ m) \ ps \wedge \text{hd } ps = (1,1) \wedge \text{last } ps = (2, \text{int } m-1)$
 $\langle \text{proof} \rangle$

For every $5 \times m$ -board with $m \geq 5$ there exists a knight's path that starts in $(1,1)$ (bottom-left) and ends in $(4, m-1)$ (top-right).

lemma *knight's-path-5xm-ur-exists*:

assumes $m \geq 5$

shows $\exists ps. \text{knight's-path } (\text{board } 5 \ m) \ ps \wedge \text{hd } ps = (1,1) \wedge \text{last } ps = (4, \text{int } m-1)$
 $\langle \text{proof} \rangle$

$5 \leq ?m \implies \exists ps. \text{knight's-path } (\text{board } 5 \ ?m) \ ps \wedge \text{hd } ps = (1, 1) \wedge \text{last } ps = (2, \text{int } ?m - 1)$ and $5 \leq ?m \implies \exists ps. \text{knight's-path } (\text{board } 5 \ ?m) \ ps \wedge \text{hd } ps = (1, 1) \wedge \text{last } ps = (2, \text{int } ?m - 1)$ formalize Lemma 1 from [1].

lemmas *knight's-path-5xm-exists* = *knight's-path-5xm-lr-exists knight's-path-5xm-ur-exists*

9 Knight's Paths and Circuits for $6 \times m$ -Boards

abbreviation $b6x5 \equiv \text{board } 6 \ 5$

A Knight's path for the (6×5) -board that starts in the lower-left and ends in the upper-left.

10	19	4	29	12
3	30	11	20	5
18	9	24	13	28
25	2	17	6	21
16	23	8	27	14
1	26	15	22	7

abbreviation $kp6x5ul \equiv \text{the } (to\text{-path}$

$[[10,19,4,29,12],$

$[3,30,11,20,5],$

$[18,9,24,13,28],$

$[25,2,17,6,21],$

$[16,23,8,27,14],$

$[1,26,15,22,7]])$

lemma *kp-6x5-ul*: *knight's-path* $b6x5 \ kp6x5ul$

$\langle \text{proof} \rangle$

lemma *kp-6x5-ul-hd*: $\text{hd } kp6x5ul = (1,1)$ $\langle \text{proof} \rangle$

lemma *kp-6x5-ul-last*: $\text{last } kp6x5ul = (5,2)$ $\langle \text{proof} \rangle$

lemma *kp-6x5-ul-non-nil*: $kp6x5ul \neq []$ $\langle \text{proof} \rangle$

A Knight's circuit for the (6×5) -board.

16	9	6	27	18
7	26	17	14	5
10	15	8	19	28
25	30	23	4	13
22	11	2	29	20
1	24	21	12	3

abbreviation $kc6x5 \equiv$ the (to-path

$[[16,9,6,27,18],$
 $[7,26,17,14,5],$
 $[10,15,8,19,28],$
 $[25,30,23,4,13],$
 $[22,11,2,29,20],$
 $[1,24,21,12,3]]$)

lemma $kc-6x5$: knights-circuit $b6x5$ $kc6x5$
 $\langle proof \rangle$

lemma $kc-6x5-hd$: $hd\ kc6x5 = (1,1)$ $\langle proof \rangle$

lemma $kc-6x5-non-nil$: $kc6x5 \neq []$ $\langle proof \rangle$

abbreviation $b6x6 \equiv$ board 6 6

The path given for the 6×6 -board that ends in the upper-left is wrong. The Knight cannot move from square 26 to square 27.

14	23	6	28	12	21
7	36	13	22	5	27
24	15	29	35	20	11
30	8	17	26	34	4
16	25	2	32	10	19
1	31	9	18	3	33

abbreviation $kp6x6ul-false \equiv$ the (to-path

$[[14,23,6,28,12,21],$
 $[7,36,13,22,5,27],$
 $[24,15,29,35,20,11],$
 $[30,8,17,26,34,4],$
 $[16,25,2,32,10,19],$
 $[1,31,9,18,3,33]]$)

lemma \neg knights-path $b6x6$ $kp6x6ul-false$
 $\langle proof \rangle$

I have computed a correct Knight's path for the 6×6 -board that ends in the

upper-left. A Knight's path for the (6×6) -board that starts in the lower-left and ends in the upper-left.

8	25	10	21	6	23
11	36	7	24	33	20
26	9	34	3	22	5
35	12	15	30	19	32
14	27	2	17	4	29
1	16	13	28	31	18

abbreviation $kp6x6ul \equiv$ the (to-path

$[[8,25,10,21,6,23],$
 $[11,36,7,24,33,20],$
 $[26,9,34,3,22,5],$
 $[35,12,15,30,19,32],$
 $[14,27,2,17,4,29],$
 $[1,16,13,28,31,18]]$)

lemma $kp\text{-}6x6\text{-}ul$: knights-path $b6x6$ $kp6x6ul$
 $\langle proof \rangle$

lemma $kp\text{-}6x6\text{-}ul\text{-}hd$: $hd\ kp6x6ul = (1,1)$ $\langle proof \rangle$

lemma $kp\text{-}6x6\text{-}ul\text{-}last$: $last\ kp6x6ul = (5,2)$ $\langle proof \rangle$

lemma $kp\text{-}6x6\text{-}ul\text{-}non\text{-}nil$: $kp6x6ul \neq []$ $\langle proof \rangle$

A Knight's circuit for the (6×6) -board.

4	25	34	15	18	7
35	14	5	8	33	16
24	3	26	17	6	19
13	36	23	30	9	32
22	27	2	11	20	29
1	12	21	28	31	10

abbreviation $kc6x6 \equiv$ the (to-path

$[[4,25,34,15,18,7],$
 $[35,14,5,8,33,16],$
 $[24,3,26,17,6,19],$
 $[13,36,23,30,9,32],$
 $[22,27,2,11,20,29],$
 $[1,12,21,28,31,10]]$)

lemma $kc\text{-}6x6$: knights-circuit $b6x6$ $kc6x6$
 $\langle proof \rangle$

lemma $kc\text{-}6x6\text{-}hd$: $hd\ kc6x6 = (1,1)$ $\langle proof \rangle$

lemma *kc-6x6-non-nil*: $kc6x6 \neq []$ *<proof>*

abbreviation *b6x7* \equiv *board 6 7*

A Knight's path for the (6×7) -board that starts in the lower-left and ends in the upper-left.

18	23	8	39	16	25	6
9	42	17	24	7	40	15
22	19	32	41	38	5	26
33	10	21	28	31	14	37
20	29	2	35	12	27	4
1	34	11	30	3	36	13

abbreviation *kp6x7ul* \equiv *the (to-path*

[[18,23,8,39,16,25,6],
[9,42,17,24,7,40,15],
[22,19,32,41,38,5,26],
[33,10,21,28,31,14,37],
[20,29,2,35,12,27,4],
[1,34,11,30,3,36,13]])

lemma *kp-6x7-ul*: *knight's-path b6x7 kp6x7ul*
<proof>

lemma *kp-6x7-ul-hd*: *hd kp6x7ul = (1,1)* *<proof>*

lemma *kp-6x7-ul-last*: *last kp6x7ul = (5,2)* *<proof>*

lemma *kp-6x7-ul-non-nil*: $kp6x7ul \neq []$ *<proof>*

A Knight's circuit for the (6×7) -board.

26	37	8	17	28	31	6
9	18	27	36	7	16	29
38	25	10	19	30	5	32
11	42	23	40	35	20	15
24	39	2	13	22	33	4
1	12	41	34	3	14	21

abbreviation *kc6x7* \equiv *the (to-path*

[[26,37,8,17,28,31,6],
[9,18,27,36,7,16,29],
[38,25,10,19,30,5,32],
[11,42,23,40,35,20,15],
[24,39,2,13,22,33,4],
[1,12,41,34,3,14,21]])

lemma *kc-6x7*: *knight's-circuit b6x7 kc6x7*

$\langle proof \rangle$

lemma *kc-6x7-hd*: $hd\ kc6x7 = (1,1)$ $\langle proof \rangle$

lemma *kc-6x7-non-nil*: $kc6x7 \neq []$ $\langle proof \rangle$

abbreviation *b6x8* \equiv *board 6 8*

A Knight's path for the (6×8) -board that starts in the lower-left and ends in the upper-left.

18	31	8	35	16	33	6	45
9	48	17	32	7	46	15	26
30	19	36	47	34	27	44	5
37	10	21	28	43	40	25	14
20	29	2	39	12	23	4	41
1	38	11	22	3	42	13	24

abbreviation *kp6x8ul* \equiv *the (to-path*

$[[18,31,8,35,16,33,6,45],$
 $[9,48,17,32,7,46,15,26],$
 $[30,19,36,47,34,27,44,5],$
 $[37,10,21,28,43,40,25,14],$
 $[20,29,2,39,12,23,4,41],$
 $[1,38,11,22,3,42,13,24]]$

lemma *kp-6x8-ul*: *knight's-path b6x8 kp6x8ul*
 $\langle proof \rangle$

lemma *kp-6x8-ul-hd*: $hd\ kp6x8ul = (1,1)$ $\langle proof \rangle$

lemma *kp-6x8-ul-last*: $last\ kp6x8ul = (5,2)$ $\langle proof \rangle$

lemma *kp-6x8-ul-non-nil*: $kp6x8ul \neq []$ $\langle proof \rangle$

A Knight's circuit for the (6×8) -board.

30	35	8	15	28	39	6	13
9	16	29	36	7	14	27	38
34	31	10	23	40	37	12	5
17	48	33	46	11	22	41	26
32	45	2	19	24	43	4	21
1	18	47	44	3	20	25	42

abbreviation *kc6x8* \equiv *the (to-path*

$[30,35,8,15,28,39,6,13],$
 $[9,16,29,36,7,14,27,38],$
 $[34,31,10,23,40,37,12,5],$

[17,48,33,46,11,22,41,26],
 [32,45,2,19,24,43,4,21],
 [1,18,47,44,3,20,25,42]])

lemma *kc-6x8: knights-circuit b6x8 kc6x8*
 ⟨proof⟩

lemma *kc-6x8-hd: hd kc6x8 = (1,1) ⟨proof⟩*

lemma *kc-6x8-non-nil: kc6x8 ≠ [] ⟨proof⟩*

abbreviation *b6x9 ≡ board 6 9*

A Knight's path for the (6×9) -board that starts in the lower-left and ends in the upper-left.

22	45	10	53	20	47	8	35	18
11	54	21	46	9	36	19	48	7
44	23	42	37	52	49	32	17	34
41	12	25	50	27	38	29	6	31
24	43	2	39	14	51	4	33	16
1	40	13	26	3	28	15	30	5

abbreviation *kp6x9ul ≡ the (to-path*

[[22,45,10,53,20,47,8,35,18],
 [11,54,21,46,9,36,19,48,7],
 [44,23,42,37,52,49,32,17,34],
 [41,12,25,50,27,38,29,6,31],
 [24,43,2,39,14,51,4,33,16],
 [1,40,13,26,3,28,15,30,5]])

lemma *kp-6x9-ul: knights-path b6x9 kp6x9ul*
 ⟨proof⟩

lemma *kp-6x9-ul-hd: hd kp6x9ul = (1,1) ⟨proof⟩*

lemma *kp-6x9-ul-last: last kp6x9ul = (5,2) ⟨proof⟩*

lemma *kp-6x9-ul-non-nil: kp6x9ul ≠ [] ⟨proof⟩*

A Knight's circuit for the (6×9) -board.

14	49	4	51	24	39	6	29	22
3	52	13	40	5	32	23	42	7
48	15	50	25	38	41	28	21	30
53	2	37	12	33	26	31	8	43
16	47	54	35	18	45	10	27	20
1	36	17	46	11	34	19	44	9

abbreviation *kc6x9 ≡ the (to-path*

$[[14,49,4,51,24,39,6,29,22],$
 $[3,52,13,40,5,32,23,42,7],$
 $[48,15,50,25,38,41,28,21,30],$
 $[53,2,37,12,33,26,31,8,43],$
 $[16,47,54,35,18,45,10,27,20],$
 $[1,36,17,46,11,34,19,44,9]]]$

lemma *kc-6x9: knights-circuit b6x9 kc6x9*
 $\langle \text{proof} \rangle$

lemma *kc-6x9-hd: hd kc6x9 = (1,1) $\langle \text{proof} \rangle$*

lemma *kc-6x9-non-nil: kc6x9 $\neq []$ $\langle \text{proof} \rangle$*

lemmas *kp-6xm-ul =*
kp-6x5-ul kp-6x5-ul-hd kp-6x5-ul-last kp-6x5-ul-non-nil
kp-6x6-ul kp-6x6-ul-hd kp-6x6-ul-last kp-6x6-ul-non-nil
kp-6x7-ul kp-6x7-ul-hd kp-6x7-ul-last kp-6x7-ul-non-nil
kp-6x8-ul kp-6x8-ul-hd kp-6x8-ul-last kp-6x8-ul-non-nil
kp-6x9-ul kp-6x9-ul-hd kp-6x9-ul-last kp-6x9-ul-non-nil

lemmas *kc-6xm =*
kc-6x5 kc-6x5-hd kc-6x5-non-nil
kc-6x6 kc-6x6-hd kc-6x6-non-nil
kc-6x7 kc-6x7-hd kc-6x7-non-nil
kc-6x8 kc-6x8-hd kc-6x8-non-nil
kc-6x9 kc-6x9-hd kc-6x9-non-nil

For every $6 \times m$ -board with $m \geq 5$ there exists a knight's path that starts in $(1,1)$ (bottom-left) and ends in $(5,2)$ (top-left).

lemma *knights-path-6xm-ul-exists:*
assumes $m \geq 5$
shows $\exists ps. \text{knights-path (board } 6 \text{ } m) \text{ } ps \wedge \text{hd } ps = (1,1) \wedge \text{last } ps = (5,2)$
 $\langle \text{proof} \rangle$

For every $6 \times m$ -board with $m \geq 5$ there exists a knight's circuit.

lemma *knights-circuit-6xm-exists:*
assumes $m \geq 5$
shows $\exists ps. \text{knights-circuit (board } 6 \text{ } m) \text{ } ps$
 $\langle \text{proof} \rangle$

$5 \leq ?m \implies \exists ps. \text{knights-path (board } 6 \text{ } ?m) \text{ } ps \wedge \text{hd } ps = (1, 1) \wedge \text{last } ps = (5, 2)$ and $5 \leq ?m \implies \exists ps. \text{knights-circuit (board } 6 \text{ } ?m) \text{ } ps$ formalize Lemma 2 from [1].

lemmas *knights-path-6xm-exists = knights-path-6xm-ul-exists knights-circuit-6xm-exists*

10 Knight's Paths and Circuits for $8 \times m$ -Boards

abbreviation *b8x5 \equiv board 8 5*

A Knight's path for the (8×5) -board that starts in the lower-left and ends in the upper-left.

28	7	22	39	26
23	40	27	6	21
8	29	38	25	14
37	24	15	20	5
16	9	30	13	34
31	36	33	4	19
10	17	2	35	12
1	32	11	18	3

abbreviation $kp8x5ul \equiv$ the (to-path

[[28,7,22,39,26],
 [23,40,27,6,21],
 [8,29,38,25,14],
 [37,24,15,20,5],
 [16,9,30,13,34],
 [31,36,33,4,19],
 [10,17,2,35,12],
 [1,32,11,18,3]])

lemma $kp-8x5-ul$: knights-path $b8x5$ $kp8x5ul$
 ⟨proof⟩

lemma $kp-8x5-ul-hd$: hd $kp8x5ul = (1,1)$ ⟨proof⟩

lemma $kp-8x5-ul-last$: $last$ $kp8x5ul = (7,2)$ ⟨proof⟩

lemma $kp-8x5-ul-non-nil$: $kp8x5ul \neq []$ ⟨proof⟩

A Knight's circuit for the (8×5) -board.

26	7	28	15	24
31	16	25	6	29
8	27	30	23	14
17	32	39	34	5
38	9	18	13	22
19	40	33	4	35
10	37	2	21	12
1	20	11	36	3

abbreviation $kc8x5 \equiv$ the (to-path

[[26,7,28,15,24],
 [31,16,25,6,29],
 [8,27,30,23,14],
 [17,32,39,34,5],

[38,9,18,13,22],
 [19,40,33,4,35],
 [10,37,2,21,12],
 [1,20,11,36,3]])

lemma *kc-8x5: knights-circuit b8x5 kc8x5*
 ⟨proof⟩

lemma *kc-8x5-hd: hd kc8x5 = (1,1) ⟨proof⟩*

lemma *kc-8x5-last: last kc8x5 = (3,2) ⟨proof⟩*

lemma *kc-8x5-non-nil: kc8x5 ≠ [] ⟨proof⟩*

lemma *kc-8x5-si: step-in kc8x5 (2,4) (4,5) (is step-in ?ps - -)*
 ⟨proof⟩

abbreviation *b8x6 ≡ board 8 6*

A Knight's path for the (8×6) -board that starts in the lower-left and ends in the upper-left.

42	11	26	9	34	13
25	48	43	12	27	8
44	41	10	33	14	35
47	24	45	20	7	28
40	19	32	3	36	15
23	46	21	6	29	4
18	39	2	31	16	37
1	22	17	38	5	30

abbreviation *kp8x6ul ≡ the (to-path*

[[42,11,26,9,34,13],
 [25,48,43,12,27,8],
 [44,41,10,33,14,35],
 [47,24,45,20,7,28],
 [40,19,32,3,36,15],
 [23,46,21,6,29,4],
 [18,39,2,31,16,37],
 [1,22,17,38,5,30]])

lemma *kp-8x6-ul: knights-path b8x6 kp8x6ul*
 ⟨proof⟩

lemma *kp-8x6-ul-hd: hd kp8x6ul = (1,1) ⟨proof⟩*

lemma *kp-8x6-ul-last: last kp8x6ul = (7,2) ⟨proof⟩*

lemma *kp-8x6-ul-non-nil: kp8x6ul ≠ [] ⟨proof⟩*

A Knight's circuit for the (8×6) -board. I have reversed circuit s.t. the circuit steps from $(2,5)$ to $(4,6)$ and not the other way around. This makes the proofs easier.

8	29	24	45	12	37
25	46	9	38	23	44
30	7	28	13	36	11
47	26	39	10	43	22
6	31	4	27	14	35
3	48	17	40	21	42
32	5	2	19	34	15
1	18	33	16	41	20

abbreviation $kc8x6 \equiv$ the (to-path

$[[8,29,24,45,12,37],$
 $[25,46,9,38,23,44],$
 $[30,7,28,13,36,11],$
 $[47,26,39,10,43,22],$
 $[6,31,4,27,14,35],$
 $[3,48,17,40,21,42],$
 $[32,5,2,19,34,15],$
 $[1,18,33,16,41,20]]$

lemma $kc-8x6$: knights-circuit $b8x6$ $kc8x6$
 $\langle proof \rangle$

lemma $kc-8x6-hd$: $hd\ kc8x6 = (1,1)$ $\langle proof \rangle$

lemma $kc-8x6-non-nil$: $kc8x6 \neq []$ $\langle proof \rangle$

lemma $kc-8x6-si$: step-in $kc8x6$ $(2,5)$ $(4,6)$ (is step-in ?ps - -)
 $\langle proof \rangle$

abbreviation $b8x7 \equiv$ board 8 7

A Knight's path for the (8×7) -board that starts in the lower-left and ends in the upper-left.

38	19	6	55	46	21	8
5	56	39	20	7	54	45
18	37	4	47	34	9	22
3	48	35	40	53	44	33
36	17	52	49	32	23	10
51	2	29	14	41	26	43
16	13	50	31	28	11	24
1	30	15	12	25	42	27

abbreviation $kp8x7ul \equiv$ the (to-path

$[[38,19,6,55,46,21,8],$
 $[5,56,39,20,7,54,45],$
 $[18,37,4,47,34,9,22],$
 $[3,48,35,40,53,44,33],$
 $[36,17,52,49,32,23,10],$
 $[51,2,29,14,41,26,43],$
 $[16,13,50,31,28,11,24],$
 $[1,30,15,12,25,42,27]]])$

lemma *kp-8x7-ul: knights-path b8x7 kp8x7ul*
 $\langle proof \rangle$

lemma *kp-8x7-ul-hd: hd kp8x7ul = (1,1) $\langle proof \rangle$*

lemma *kp-8x7-ul-last: last kp8x7ul = (7,2) $\langle proof \rangle$*

lemma *kp-8x7-ul-non-nil: kp8x7ul $\neq []$ $\langle proof \rangle$*

A Knight's circuit for the (8×7) -board. I have reversed circuit s.t. the circuit steps from $(2,6)$ to $(4,7)$ and not the other way around. This makes the proofs easier.

36	31	18	53	20	29	44
17	54	35	30	45	52	21
32	37	46	19	8	43	28
55	16	7	34	27	22	51
38	33	26	47	6	9	42
3	56	15	12	25	50	23
14	39	2	5	48	41	10
1	4	13	40	11	24	49

abbreviation *kc8x7 \equiv the (to-path*

$[[36,31,18,53,20,29,44],$
 $[17,54,35,30,45,52,21],$
 $[32,37,46,19,8,43,28],$
 $[55,16,7,34,27,22,51],$
 $[38,33,26,47,6,9,42],$
 $[3,56,15,12,25,50,23],$
 $[14,39,2,5,48,41,10],$
 $[1,4,13,40,11,24,49]]])$

lemma *kc-8x7: knights-circuit b8x7 kc8x7*
 $\langle proof \rangle$

lemma *kc-8x7-hd: hd kc8x7 = (1,1) $\langle proof \rangle$*

lemma *kc-8x7-non-nil: kc8x7 $\neq []$ $\langle proof \rangle$*

lemma *kc-8x7-si: step-in kc8x7 (2,6) (4,7) (is step-in ?ps - -)*
 $\langle proof \rangle$

abbreviation $b8x8 \equiv \text{board } 8 \ 8$

The path given for the 8×8 -board that ends in the upper-left is wrong. The Knight cannot move from square 27 to square 28.

24	11	37	9	26	21	39	7
36	64	24	22	38	8	27	20
12	23	10	53	58	49	6	28
63	35	61	50	55	52	19	40
46	13	54	57	48	59	29	5
34	62	47	60	51	56	41	18
14	45	2	32	16	43	4	30
1	33	15	44	3	31	17	42

abbreviation $kp8x8ul\text{-}false \equiv \text{the (to-path}$

$[[24,11,37,9,26,21,39,7],$
 $[36,64,25,22,38,8,27,20],$
 $[12,23,10,53,58,49,6,28],$
 $[63,35,61,50,55,52,19,40],$
 $[46,13,54,57,48,59,29,5],$
 $[34,62,47,60,51,56,41,18],$
 $[14,45,2,32,16,43,4,30],$
 $[1,33,15,44,3,31,17,42]]$)

lemma $\neg \text{knight's-path } b8x8 \text{ } kp8x8ul\text{-}false$
 $\langle \text{proof} \rangle$

I have computed a correct Knight's path for the 8×8 -board that ends in the upper-left.

38	41	36	27	32	43	20	25
35	64	39	42	21	26	29	44
40	37	6	33	28	31	24	19
5	34	63	14	7	22	45	30
62	13	4	9	58	49	18	23
3	10	61	52	15	8	57	46
12	53	2	59	48	55	50	17
1	60	11	54	51	16	47	56

abbreviation $kp8x8ul \equiv \text{the (to-path}$

$[[38,41,36,27,32,43,20,25],$
 $[35,64,39,42,21,26,29,44],$
 $[40,37,6,33,28,31,24,19],$
 $[5,34,63,14,7,22,45,30],$
 $[62,13,4,9,58,49,18,23],$

[3,10,61,52,15,8,57,46],
 [12,53,2,59,48,55,50,17],
 [1,60,11,54,51,16,47,56]])

lemma *kp-8x8-ul: knights-path b8x8 kp8x8ul*
 ⟨proof⟩

lemma *kp-8x8-ul-hd: hd kp8x8ul = (1,1) ⟨proof⟩*

lemma *kp-8x8-ul-last: last kp8x8ul = (7,2) ⟨proof⟩*

lemma *kp-8x8-ul-non-nil: kp8x8ul ≠ [] ⟨proof⟩*

A Knight's circuit for the (8×8) -board.

48	13	30	9	56	45	28	7
31	10	47	50	29	8	57	44
14	49	12	55	46	59	6	27
11	32	37	60	51	54	43	58
36	15	52	63	38	61	26	5
33	64	35	18	53	40	23	42
16	19	2	39	62	21	4	25
1	34	17	20	3	24	41	22

abbreviation *kc8x8 ≡ the (to-path*

[[48,13,30,9,56,45,28,7],
 [31,10,47,50,29,8,57,44],
 [14,49,12,55,46,59,6,27],
 [11,32,37,60,51,54,43,58],
 [36,15,52,63,38,61,26,5],
 [33,64,35,18,53,40,23,42],
 [16,19,2,39,62,21,4,25],
 [1,34,17,20,3,24,41,22]])

lemma *kc-8x8: knights-circuit b8x8 kc8x8*
 ⟨proof⟩

lemma *kc-8x8-hd: hd kc8x8 = (1,1) ⟨proof⟩*

lemma *kc-8x8-non-nil: kc8x8 ≠ [] ⟨proof⟩*

lemma *kc-8x8-si: step-in kc8x8 (2,7) (4,8) (is step-in ?ps - -)*
 ⟨proof⟩

abbreviation *b8x9 ≡ board 8 9*

A Knight's path for the (8×9) -board that starts in the lower-left and ends in the upper-left.

32	47	6	71	30	45	8	43	26
5	72	31	46	7	70	27	22	9
48	33	4	29	64	23	44	25	42
3	60	35	62	69	28	41	10	21
34	49	68	65	36	63	24	55	40
59	2	61	16	67	56	37	20	11
50	15	66	57	52	13	18	39	54
1	58	51	14	17	38	53	12	19

abbreviation $kp8x9ul \equiv$ the (to-path

$[[32,47,6,71,30,45,8,43,26],$
 $[5,72,31,46,7,70,27,22,9],$
 $[48,33,4,29,64,23,44,25,42],$
 $[3,60,35,62,69,28,41,10,21],$
 $[34,49,68,65,36,63,24,55,40],$
 $[59,2,61,16,67,56,37,20,11],$
 $[50,15,66,57,52,13,18,39,54],$
 $[1,58,51,14,17,38,53,12,19]]$)

lemma $kp-8x9-ul$: knights-path $b8x9$ $kp8x9ul$
 $\langle proof \rangle$

lemma $kp-8x9-ul-hd$: $hd\ kp8x9ul = (1,1)$ $\langle proof \rangle$

lemma $kp-8x9-ul-last$: $last\ kp8x9ul = (7,2)$ $\langle proof \rangle$

lemma $kp-8x9-ul-non-nil$: $kp8x9ul \neq []$ $\langle proof \rangle$

A Knight's circuit for the (8×9) -board.

42	19	38	5	36	21	34	7	60
39	4	41	20	63	6	59	22	33
18	43	70	37	58	35	68	61	8
3	40	49	64	69	62	57	32	23
50	17	44	71	48	67	54	9	56
45	2	65	14	27	12	29	24	31
16	51	72	47	66	53	26	55	10
1	46	15	52	13	28	11	30	25

abbreviation $kc8x9 \equiv$ the (to-path

$[[42,19,38,5,36,21,34,7,60],$
 $[39,4,41,20,63,6,59,22,33],$
 $[18,43,70,37,58,35,68,61,8],$
 $[3,40,49,64,69,62,57,32,23],$
 $[50,17,44,71,48,67,54,9,56],$
 $[45,2,65,14,27,12,29,24,31],$
 $[16,51,72,47,66,53,26,55,10],$

$[1,46,15,52,13,28,11,30,25]]$)
lemma *kc-8x9: knights-circuit b8x9 kc8x9*
 $\langle \text{proof} \rangle$

lemma *kc-8x9-hd: hd kc8x9 = (1,1) $\langle \text{proof} \rangle$*

lemma *kc-8x9-non-nil: kc8x9 $\neq []$ $\langle \text{proof} \rangle$*

lemma *kc-8x9-si: step-in kc8x9 (2,8) (4,9) (is step-in ?ps - -)*
 $\langle \text{proof} \rangle$

lemmas *kp-8xm-ul =*
kp-8x5-ul kp-8x5-ul-hd kp-8x5-ul-last kp-8x5-ul-non-nil
kp-8x6-ul kp-8x6-ul-hd kp-8x6-ul-last kp-8x6-ul-non-nil
kp-8x7-ul kp-8x7-ul-hd kp-8x7-ul-last kp-8x7-ul-non-nil
kp-8x8-ul kp-8x8-ul-hd kp-8x8-ul-last kp-8x8-ul-non-nil
kp-8x9-ul kp-8x9-ul-hd kp-8x9-ul-last kp-8x9-ul-non-nil

lemmas *kc-8xm =*
kc-8x5 kc-8x5-hd kc-8x5-last kc-8x5-non-nil kc-8x5-si
kc-8x6 kc-8x6-hd kc-8x6-non-nil kc-8x6-si
kc-8x7 kc-8x7-hd kc-8x7-non-nil kc-8x7-si
kc-8x8 kc-8x8-hd kc-8x8-non-nil kc-8x8-si
kc-8x9 kc-8x9-hd kc-8x9-non-nil kc-8x9-si

For every $8 \times m$ -board with $m \geq 5$ there exists a knight's circuit.

lemma *knights-circuit-8xm-exists:*
assumes $m \geq 5$
shows $\exists ps. \text{knights-circuit } (\text{board } 8 \ m) \ ps \wedge \text{step-in } ps \ (2, \text{int } m-1) \ (4, \text{int } m)$
 $\langle \text{proof} \rangle$

For every $8 \times m$ -board with $m \geq 5$ there exists a knight's path that starts in $(1,1)$ (bottom-left) and ends in $(7,2)$ (top-left).

lemma *knights-path-8xm-ul-exists:*
assumes $m \geq 5$
shows $\exists ps. \text{knights-path } (\text{board } 8 \ m) \ ps \wedge \text{hd } ps = (1,1) \wedge \text{last } ps = (7,2)$
 $\langle \text{proof} \rangle$

$5 \leq ?m \implies \exists ps. \text{knights-circuit } (\text{board } 8 \ ?m) \ ps \wedge \text{step-in } ps \ (2, \text{int } ?m - 1) \ (4, \text{int } ?m)$ and $5 \leq ?m \implies \exists ps. \text{knights-path } (\text{board } 8 \ ?m) \ ps \wedge \text{hd } ps = (1, 1) \wedge \text{last } ps = (7, 2)$ formalize Lemma 3 from [1].

lemmas *knights-path-8xm-exists = knights-circuit-8xm-exists knights-path-8xm-ul-exists*

11 Knight's Paths and Circuits for $n \times m$ -Boards

In this section the desired theorems are proved. The proof uses the previous lemmas to construct paths and circuits for arbitrary $n \times m$ -boards.

A Knight's path for the (5×5) -board that starts in the lower-left and ends in the upper-left.

7	20	9	14	5
10	25	6	21	16
19	8	15	4	13
24	11	2	17	22
1	18	23	12	3

abbreviation $kp5x5ul \equiv$ the (to-path

$[[7,20,9,14,5],$
 $[10,25,6,21,16],$
 $[19,8,15,4,13],$
 $[24,11,2,17,22],$
 $[1,18,23,12,3]]$)

lemma $kp-5x5-ul$: knights-path $b5x5$ $kp5x5ul$
 $\langle proof \rangle$

A Knight's path for the (5×7) -board that starts in the lower-left and ends in the upper-left.

17	14	25	6	19	8	29
26	35	18	15	28	5	20
13	16	27	24	7	30	9
34	23	2	11	32	21	4
1	12	33	22	3	10	31

abbreviation $kp5x7ul \equiv$ the (to-path

$[[17,14,25,6,19,8,29],$
 $[26,35,18,15,28,5,20],$
 $[13,16,27,24,7,30,9],$
 $[34,23,2,11,32,21,4],$
 $[1,12,33,22,3,10,31]]$)

lemma $kp-5x7-ul$: knights-path $b5x7$ $kp5x7ul$
 $\langle proof \rangle$

A Knight's path for the (5×9) -board that starts in the lower-left and ends in the upper-left.

7	12	37	42	5	18	23	32	27
38	45	6	11	36	31	26	19	24
13	8	43	4	41	22	17	28	33
44	39	2	15	10	35	30	25	20
1	14	9	40	3	16	21	34	29

abbreviation $kp5x9ul \equiv$ the (to-path

$[[7,12,37,42,5,18,23,32,27],$
 $[38,45,6,11,36,31,26,19,24],$
 $[13,8,43,4,41,22,17,28,33],$
 $[44,39,2,15,10,35,30,25,20],$
 $[1,14,9,40,3,16,21,34,29]]$

lemma *kp-5x9-ul: knights-path b5x9 kp5x9ul*
<proof>

abbreviation *b7x7* \equiv *board 7 7*

A Knight's path for the (7×7) -board that starts in the lower-left and ends in the upper-left.

9	30	19	42	7	32	17
20	49	8	31	18	43	6
29	10	41	36	39	16	33
48	21	38	27	34	5	44
11	28	35	40	37	26	15
22	47	2	13	24	45	4
1	12	23	46	3	14	25

abbreviation *kp7x7ul* \equiv *the (to-path*

$[[9,30,19,42,7,32,17],$
 $[20,49,8,31,18,43,6],$
 $[29,10,41,36,39,16,33],$
 $[48,21,38,27,34,5,44],$
 $[11,28,35,40,37,26,15],$
 $[22,47,2,13,24,45,4],$
 $[1,12,23,46,3,14,25]]$

lemma *kp-7x7-ul: knights-path b7x7 kp7x7ul*
<proof>

abbreviation *b7x9* \equiv *board 7 9*

A Knight's path for the (7×9) -board that starts in the lower-left and ends in the upper-left.

59	4	17	50	37	6	19	30	39
16	63	58	5	18	51	38	7	20
3	60	49	36	57	42	29	40	31
48	15	62	43	52	35	56	21	8
61	2	13	26	45	28	41	32	55
14	47	44	11	24	53	34	9	22
1	12	25	46	27	10	23	54	33

abbreviation *kp7x9ul* \equiv *the (to-path*

$[[59,4,17,50,37,6,19,30,39],$
 $[16,63,58,5,18,51,38,7,20],$
 $[3,60,49,36,57,42,29,40,31],$
 $[48,15,62,43,52,35,56,21,8],$
 $[61,2,13,26,45,28,41,32,55],$
 $[14,47,44,11,24,53,34,9,22],$
 $[1,12,25,46,27,10,23,54,33]]$

lemma *kp-7x9-ul: knights-path b7x9 kp7x9ul*
<proof>

abbreviation *b9x7* \equiv *board 9 7*

A Knight's path for the (9×7) -board that starts in the lower-left and ends in the upper-left.

5	20	53	48	7	22	31
52	63	6	21	32	55	8
19	4	49	54	47	30	23
62	51	46	33	56	9	58
3	18	61	50	59	24	29
14	43	34	45	28	57	10
17	2	15	60	35	38	25
42	13	44	27	40	11	36
1	16	41	12	37	26	39

abbreviation *kp9x7ul* \equiv *the (to-path*

$[[5,20,53,48,7,22,31],$
 $[52,63,6,21,32,55,8],$
 $[19,4,49,54,47,30,23],$
 $[62,51,46,33,56,9,58],$
 $[3,18,61,50,59,24,29],$
 $[14,43,34,45,28,57,10],$
 $[17,2,15,60,35,38,25],$
 $[42,13,44,27,40,11,36],$
 $[1,16,41,12,37,26,39]]$

lemma *kp-9x7-ul: knights-path b9x7 kp9x7ul*
<proof>

abbreviation *b9x9* \equiv *board 9 9*

A Knight's path for the (9×9) -board that starts in the lower-left and ends in the upper-left.

13	26	39	52	11	24	37	50	9
40	81	12	25	38	51	10	23	36
27	14	53	58	63	68	73	8	49
80	41	64	67	72	57	62	35	22
15	28	59	54	65	74	69	48	7
42	79	66	71	76	61	56	21	34
29	16	77	60	55	70	75	6	47
78	43	2	31	18	45	4	33	20
1	30	17	44	3	32	19	46	5

abbreviation $kp9x9ul \equiv$ the (to-path

$[[13,26,39,52,11,24,37,50,9],$
 $[40,81,12,25,38,51,10,23,36],$
 $[27,14,53,58,63,68,73,8,49],$
 $[80,41,64,67,72,57,62,35,22],$
 $[15,28,59,54,65,74,69,48,7],$
 $[42,79,66,71,76,61,56,21,34],$
 $[29,16,77,60,55,70,75,6,47],$
 $[78,43,2,31,18,45,4,33,20],$
 $[1,30,17,44,3,32,19,46,5]]$)

lemma $kp-9x9-ul$: knights-path $b9x9$ $kp9x9ul$
 $\langle proof \rangle$

The following lemma is a sub-proof used in Lemma 4 in [1]. I moved the sub-proof out to a separate lemma.

lemma $knights-circuit-exists-even-n-gr10$:

assumes $even\ n\ n \geq 10\ m \geq 5$

$\exists ps. knights-path\ (board\ (n-5)\ m)\ ps \wedge hd\ ps = (int\ (n-5),1)$
 $\wedge last\ ps = (int\ (n-5)-1, int\ m-1)$

shows $\exists ps. knights-circuit\ (board\ m\ n)\ ps$
 $\langle proof \rangle$

For every $n \times m$ -board with $min\ n\ m \geq 5$ and odd n there exists a Knight's path that starts in $(n,1)$ (top-left) and ends in $(n-1, m-1)$ (top-right).

This lemma formalizes Lemma 4 from [1]. Formalizing the proof of this lemma was quite challenging as a lot of details on how to exactly combine the boards are left out in the original proof in [1].

lemma $knights-path-odd-n-exists$:

assumes $odd\ n\ min\ n\ m \geq 5$

shows $\exists ps. knights-path\ (board\ n\ m)\ ps \wedge hd\ ps = (int\ n,1) \wedge last\ ps = (int\ n-1, int\ m-1)$

$\langle proof \rangle$

Auxiliary lemma that constructs a Knight's circuit if $m \geq 5$ and $n \geq 10 \wedge even\ n$.

lemma $knights-circuit-exists-n-even-gr-10$:

assumes $n \geq 10 \wedge \text{even } n \ m \geq 5$
shows $\exists ps. \text{ knights-circuit } (\text{board } n \ m) \ ps$
 $\langle \text{proof} \rangle$

Final Theorem 1: For every $n \times m$ -board with $\min n \ m \geq 5$ and $n*m$ even there exists a Knight's circuit.

theorem *knights-circuit-exists*:
assumes $\min n \ m \geq 5 \text{ even } (n*m)$
shows $\exists ps. \text{ knights-circuit } (\text{board } n \ m) \ ps$
 $\langle \text{proof} \rangle$

Final Theorem 2: for every $n \times m$ -board with $\min n \ m \geq 5$ there exists a Knight's path.

theorem *knights-path-exists*:
assumes $\min n \ m \geq 5$
shows $\exists ps. \text{ knights-path } (\text{board } n \ m) \ ps$
 $\langle \text{proof} \rangle$

THE END

end

References

- [1] P. Cull and J. D. Curtins. Knight's tour revisited. *Fibonacci Quarterly*, 16:276–285, 1978.