

Formalization of "Knight's Tour Revisited"

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Abstract

This is a formalization of [1]. In [1] the existence of a Knight's path is proved for arbitrary $n \times m$ -boards with $\min(n, m) \geq 5$. If $n \cdot m$ is even, then there exists a Knight's circuit.

A Knight's Path is a sequence of moves of a Knight on a chessboard s.t. the Knight visits every square of a chessboard exactly once. Finding a Knight's path is an instance of the Hamiltonian path problem.

During the formalization two mistakes in the original proof in [1] were discovered. These mistakes are corrected in this formalization.

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theory KnightsTour
  imports Main
begin

```

1 Introduction and Definitions

A Knight's path is a sequence of moves on a chessboard s.t. every step in sequence is a valid move for a Knight and that the Knight visits every square on the boards exactly once. A Knight is a chess figure that is only able to move two squares vertically and one square horizontally or two squares horizontally and one square vertically. Finding a Knight's path is an instance of the Hamiltonian Path Problem. A Knight's circuit is a Knight's path, where additionally the Knight can move from the last square to the first square of the path, forming a loop.

[1] proves the existence of a Knight's path on a $n \times m$ -board for sufficiently large n and m . The main idea for the proof is to inductively construct a Knight's path for the $n \times m$ -board from a few pre-computed Knight's paths for small boards, i.e. 5×5 , 8×6 , ..., 8×9 . The paths for small boards are transformed (i.e. transpose, mirror, translate) and concatenated to create paths for larger boards.

While formalizing the proofs I discovered two mistakes in the original proof in [1]: (i) the pre-computed path for the 6×6 -board that ends in the upper-left (in Figure 2) and (ii) the pre-computed path for the 8×8 -board that ends in the upper-left (in Figure 5) are incorrect. I.e. on the 6×6 -board the Knight cannot step from square 26 to square 27; in the 8×8 -board the Knight cannot step from square 27 to square 28. In this formalization I have replaced the two incorrect paths with correct paths.

A square on a board is identified by its coordinates.

type-synonym *square* = *int* \times *int*

A board is represented as a set of squares. Note, that this allows boards to have an arbitrary shape and do not necessarily need to be rectangular.

type-synonym *board* = *square set*

A (rectangular) $(n \times m)$ -board is the set of all squares (i, j) where $1 \leq i \leq n$

and $1 \leq j \leq m$. $(1,1)$ is the lower-left corner, and (n,m) is the upper-right corner.

definition $board :: nat \Rightarrow nat \Rightarrow board$ **where**

$board\ n\ m = \{(i,j) \mid i\ j.\ 1 \leq i \wedge i \leq int\ n \wedge 1 \leq j \wedge j \leq int\ m\}$

A path is a sequence of steps on a board. A path is represented by the list of visited squares on the board. Each square on the $(n \times m)$ -board is identified by its coordinates (i,j) .

type-synonym $path = square\ list$

A Knight can only move two squares vertically and one square horizontally or two squares horizontally and one square vertically. Thus, a knight at position (i,j) can only move to $(i \pm 1, j \pm 2)$ or $(i \pm 2, j \pm 1)$.

definition $valid-step :: square \Rightarrow square \Rightarrow bool$ **where**

$valid-step\ s_i\ s_j \equiv (case\ s_i\ of\ (i,j) \Rightarrow s_j \in \{(i+1,j+2),(i-1,j+2),(i+1,j-2),(i-1,j-2),$
 $(i+2,j+1),(i-2,j+1),(i+2,j-1),(i-2,j-1)\})$

Now we define an inductive predicate that characterizes a Knight's path. A square s_i can be pre-pended to a current Knight's path $s_j \# ps$ if (i) there is a valid step from the square s_i to the first square s_j of the current path and (ii) the square s_i has not been visited yet.

inductive $knight\text{-}path :: board \Rightarrow path \Rightarrow bool$ **where**

$knight\text{-}path\ \{s_i\}\ [s_i]$
 $\mid s_i \notin b \Longrightarrow valid-step\ s_i\ s_j \Longrightarrow knight\text{-}path\ b\ (s_j \# ps) \Longrightarrow knight\text{-}path\ (b \cup \{s_i\})$
 $(s_i \# s_j \# ps)$

code-pred $knight\text{-}path$.

A sequence is a Knight's circuit iff the sequence is a Knight's path and there is a valid step from the last square to the first square.

definition $knight\text{-}circuit\ b\ ps \equiv (knight\text{-}path\ b\ ps \wedge valid-step\ (last\ ps)\ (hd\ ps))$

2 Executable Checker for a Knight's Path

This section gives the implementation and correctness-proof for an executable checker for a knight's-path wrt. the definition $knight\text{-}path$.

2.1 Implementation of an Executable Checker

fun $row-exec :: nat \Rightarrow int\ set$ **where**

$row-exec\ 0 = \{\}$
 $\mid row-exec\ m = insert\ (int\ m)\ (row-exec\ (m-1))$

fun $board-exec-aux :: nat \Rightarrow int\ set \Rightarrow board$ **where**

$board-exec-aux\ 0\ M = \{\}$

| *board-exec-aux* *k M* = $\{(int\ k, j) \mid j. j \in M\} \cup board-exec-aux\ (k-1)\ M$

Compute a board.

fun *board-exec* :: *nat* \Rightarrow *nat* \Rightarrow *board* **where**
board-exec *n m* = *board-exec-aux* *n* (*row-exec* *m*)

fun *step-checker* :: *square* \Rightarrow *square* \Rightarrow *bool* **where**
step-checker (*i,j*) (*i',j'*) =
 $((i+1,j+2) = (i',j') \vee (i-1,j+2) = (i',j') \vee (i+1,j-2) = (i',j') \vee (i-1,j-2) = (i',j')$
 $\vee (i+2,j+1) = (i',j') \vee (i-2,j+1) = (i',j') \vee (i+2,j-1) = (i',j') \vee (i-2,j-1) = (i',j')$

fun *path-checker* :: *board* \Rightarrow *path* \Rightarrow *bool* **where**
path-checker *b* [] = *False*
| *path-checker* *b* [*s_i*] = ($\{s_i\} = b$)
| *path-checker* *b* (*s_i*#*s_j*#*ps*) = ($s_i \in b \wedge step-checker\ s_i\ s_j \wedge path-checker\ (b - \{s_i\})\ (s_j\#ps)$)

fun *circuit-checker* :: *board* \Rightarrow *path* \Rightarrow *bool* **where**
circuit-checker *b ps* = (*path-checker* *b ps* \wedge *step-checker* (*last ps*) (*hd ps*))

2.2 Correctness Proof of the Executable Checker

lemma *row-exec-leq*: $j \in row-exec\ m \longleftrightarrow 1 \leq j \wedge j \leq int\ m$
by (*induction m*) *auto*

lemma *board-exec-aux-leq-mem*: $(i,j) \in board-exec-aux\ k\ M \longleftrightarrow 1 \leq i \wedge i \leq int\ k \wedge j \in M$
by (*induction k M rule: board-exec-aux.induct*) *auto*

lemma *board-exec-leq*: $(i,j) \in board-exec\ n\ m \longleftrightarrow 1 \leq i \wedge i \leq int\ n \wedge 1 \leq j \wedge j \leq int\ m$
using *board-exec-aux-leq-mem row-exec-leq* **by** *auto*

lemma *board-exec-correct*: *board* *n m* = *board-exec* *n m*
unfolding *board-def* **using** *board-exec-leq* **by** *auto*

lemma *step-checker-correct*: *step-checker* *s_i* *s_j* \longleftrightarrow *valid-step* *s_i* *s_j*
proof

assume *step-checker* *s_i* *s_j*
then show *valid-step* *s_i* *s_j*
unfolding *valid-step-def*
apply (*cases s_i*)
apply (*cases s_j*)
apply *auto*
done
next
assume *assms*: *valid-step* *s_i* *s_j*

```

    then show step-checker  $s_i s_j$ 
      unfolding valid-step-def by auto
qed

```

```

lemma step-checker-rev: step-checker  $(i,j) (i',j') \implies$  step-checker  $(i',j') (i,j)$ 
  apply (simp only: step-checker.simps)
  by (elim disjE) auto

```

```

lemma knights-path-intro-rev:
  assumes  $s_i \in b$  valid-step  $s_i s_j$  knights-path  $(b - \{s_i\}) (s_j \# ps)$ 
  shows knights-path  $b (s_i \# s_j \# ps)$ 
  using assms
proof -
  assume assms:  $s_i \in b$  valid-step  $s_i s_j$  knights-path  $(b - \{s_i\}) (s_j \# ps)$ 
  then have  $s_i \notin (b - \{s_i\})$   $b - \{s_i\} \cup \{s_i\} = b$ 
    by auto
  then show ?thesis
    using assms knights-path.intros(2)[of  $s_i$   $b - \{s_i\}$ ] by auto
qed

```

Final correctness corollary for the executable checker *path-checker*.

```

lemma path-checker-correct: path-checker  $b ps \longleftrightarrow$  knights-path  $b ps$ 
proof
  assume path-checker  $b ps$ 
  then show knights-path  $b ps$ 
  proof (induction rule: path-checker.induct)
    case (3  $s_i s_j xs b$ )
    then show ?case using step-checker-correct knights-path-intro-rev by auto
  qed (auto intro: knights-path.intros)
next
  assume knights-path  $b ps$ 
  then show path-checker  $b ps$ 
  using step-checker-correct
  by (induction rule: knights-path.induct) (auto elim: knights-path.cases)
qed

```

```

corollary knights-path-exec-simp: knights-path  $(board\ n\ m) ps \longleftrightarrow$  path-checker
   $(board-exec\ n\ m) ps$ 
  using board-exec-correct path-checker-correct[symmetric] by simp

```

```

lemma circuit-checker-correct: circuit-checker  $b ps \longleftrightarrow$  knights-circuit  $b ps$ 
  unfolding knights-circuit-def using path-checker-correct step-checker-correct by
  auto

```

```

corollary knights-circuit-exec-simp:
  knights-circuit  $(board\ n\ m) ps \longleftrightarrow$  circuit-checker  $(board-exec\ n\ m) ps$ 
  using board-exec-correct circuit-checker-correct[symmetric] by simp

```

3 Basic Properties of *knights-path* and *knights-circuit*

lemma *board-leq-subset*: $n_1 \leq n_2 \wedge m_1 \leq m_2 \implies \text{board } n_1 \ m_1 \subseteq \text{board } n_2 \ m_2$
unfolding *board-def* **by** *auto*

lemma *finite-row-exec*: *finite* (*row-exec* *m*)
by (*induction* *m*) *auto*

lemma *finite-board-exec-aux*: *finite* *M* \implies *finite* (*board-exec-aux* *n* *M*)
by (*induction* *n*) *auto*

lemma *board-finite*: *finite* (*board* *n* *m*)
using *finite-board-exec-aux* *finite-row-exec* **by** (*simp* *only*: *board-exec-correct*) *auto*

lemma *card-row-exec*: *card* (*row-exec* *m*) = *m*

proof (*induction* *m*)

case (*Suc* *m*)

have *int* (*Suc* *m*) \notin *row-exec* *m*

using *row-exec-leq* **by** *auto*

then have *card* (*insert* (*int* (*Suc* *m*)) (*row-exec* *m*)) = 1 + *card* (*row-exec* *m*)

using *card-Suc-eq* **by** (*metis* *Suc* *plus-1-eq-Suc* *row-exec.simps*(1))

then have *card* (*row-exec* (*Suc* *m*)) = 1 + *card* (*row-exec* *m*)

by *auto*

then show ?*case* **using** *Suc.IH* **by** *auto*

qed *auto*

lemma *set-comp-ins*:

$\{(k,j) \mid j. j \in \text{insert } x \ M\} = \text{insert } (k,x) \ \{(k,j) \mid j. j \in M\}$ (**is** ?*Mi* = ?*iM*)

proof

show ?*Mi* \subseteq ?*iM*

proof

fix *y* **assume** *y* \in ?*Mi*

then obtain *j* **where** [*simp*]: *y* = (*k,j*) **and** *j* \in *insert* *x* *M* **by** *blast*

then have *j* = *x* \vee *j* \in *M* **by** *auto*

then show *y* \in ?*iM* **by** (*elim* *disjE*) *auto*

qed

next

show ?*iM* \subseteq ?*Mi*

proof

fix *y* **assume** *y* \in ?*iM*

then obtain *j* **where** [*simp*]: *y* = (*k,j*) **and** *j* \in *insert* *x* *M* **by** *blast*

then have *j* = *x* \vee *j* \in *M* **by** *auto*

then show *y* \in ?*Mi* **by** (*elim* *disjE*) *auto*

qed

qed

lemma *finite-card-set-comp*: *finite* *M* \implies *card* $\{(k,j) \mid j. j \in M\}$ = *card* *M*

proof (*induction* *M* *rule*: *finite-induct*)

case (*insert* *x* *M*)

then show *?case* **using** *set-comp-ins*[*of k x M*] **by** *auto*
qed *auto*

lemma *card-board-exec-aux*: *finite M* \implies *card (board-exec-aux k M) = k * card M*
proof (*induction k*)

case (*Suc k*)
let *?M' = {(int (Suc k), j) | j. j ∈ M}*
let *?rec-k = board-exec-aux k M*

have *finite*: *finite ?M' finite ?rec-k*
using *Suc finite-board-exec-aux* **by** *auto*
then have *card-Un-simp*: *card (?M' ∪ ?rec-k) = card ?M' + card ?rec-k*
using *board-exec-aux-leq-mem card-Un-Int*[*of ?M' ?rec-k*] **by** *auto*

have *card-M*: *card ?M' = card M*
using *Suc finite-card-set-comp* **by** *auto*

have *card (board-exec-aux (Suc k) M) = card ?M' + card ?rec-k*
using *card-Un-simp* **by** *auto*
also have *... = card M + k * card M*
using *Suc card-M* **by** *auto*
also have *... = (Suc k) * card M*
by *auto*
finally show *?case* .

qed *auto*

lemma *card-board*: *card (board n m) = n * m*
proof –

have *card (board n m) = card (board-exec-aux n (row-exec m))*
using *board-exec-correct* **by** *auto*
also have *... = n * m*
using *card-row-exec card-board-exec-aux finite-row-exec* **by** *auto*
finally show *?thesis* .

qed

lemma *knight-path-board-non-empty*: *knight-path b ps* \implies *b ≠ {}*
by (*induction arbitrary: ps rule: knight-path.induct*) *auto*

lemma *knight-path-board-m-n-geq-1*: *knight-path (board n m) ps* \implies *min n m ≥ 1*
unfolding *board-def* **using** *knight-path-board-non-empty* **by** *fastforce*

lemma *knight-path-non-nil*: *knight-path b ps* \implies *ps ≠ []*
by (*induction arbitrary: b rule: knight-path.induct*) *auto*

lemma *knight-path-set-eq*: *knight-path b ps* \implies *set ps = b*
by (*induction rule: knight-path.induct*) *auto*

lemma *knight-path-subset*:

$\text{knights-path } b_1 \text{ } ps_1 \implies \text{knights-path } b_2 \text{ } ps_2 \implies \text{set } ps_1 \subseteq \text{set } ps_2 \iff b_1 \subseteq b_2$
using *knights-path-set-eq* **by** *auto*

lemma *knights-path-board-unique*: $\text{knights-path } b_1 \text{ } ps \implies \text{knights-path } b_2 \text{ } ps \implies b_1 = b_2$
using *knights-path-set-eq* **by** *auto*

lemma *valid-step-neq*: $\text{valid-step } s_i \text{ } s_j \implies s_i \neq s_j$
unfolding *valid-step-def* **by** *auto*

lemma *valid-step-non-transitive*: $\text{valid-step } s_i \text{ } s_j \implies \text{valid-step } s_j \text{ } s_k \implies \neg \text{valid-step } s_i \text{ } s_k$
proof –
assume *assms*: $\text{valid-step } s_i \text{ } s_j \text{ } \text{valid-step } s_j \text{ } s_k$
obtain $i_i \ j_i \ i_j \ j_j \ i_k \ j_k$ **where** $[simp]: s_i = (i_i, j_i) \ s_j = (i_j, j_j) \ s_k = (i_k, j_k)$ **by** *force*
then have $\text{step-checker } (i_i, j_i) \ (i_j, j_j) \ \text{step-checker } (i_j, j_j) \ (i_k, j_k)$
using *assms step-checker-correct* **by** *auto*
then show $\neg \text{valid-step } s_i \text{ } s_k$
apply (*simp add: step-checker-correct[symmetric]*)
apply (*elim disjE*)
apply *auto*
done
qed

lemma *knights-path-distinct*: $\text{knights-path } b \text{ } ps \implies \text{distinct } ps$
proof (*induction rule: knights-path.induct*)
case ($2 \ s_i \ b \ s_j \ ps$)
then have $s_i \notin \text{set } (s_j \# ps)$
using *knights-path-set-eq valid-step-neq* **by** *blast*
then show $?case$ **using** 2 **by** *auto*
qed *auto*

lemma *knights-path-length*: $\text{knights-path } b \text{ } ps \implies \text{length } ps = \text{card } b$
using *knights-path-set-eq knights-path-distinct* **by** (*metis distinct-card*)

lemma *knights-path-take*:
assumes $\text{knights-path } b \text{ } ps \ 0 < k \ k < \text{length } ps$
shows $\text{knights-path } (\text{set } (\text{take } k \ ps)) \ (\text{take } k \ ps)$
using *assms*
proof (*induction arbitrary: k rule: knights-path.induct*)
case ($2 \ s_i \ b \ s_j \ ps$)
then have $k = 1 \vee k = 2 \vee 2 < k$ **by** *force*
then show $?case$
using 2
proof (*elim disjE*)
assume $k = 2$
then have $\text{take } k \ (s_i \# s_j \# ps) = [s_i, s_j] \ s_i \notin \{s_j\}$ **using** 2 *valid-step-neq* **by** *auto*


```

    then show ?thesis using 2 knights-path.intros by auto
next
  assume 2 < k
  then have k-simps: k-2 = k-1-1 0 < k-2 k-2 < length ps and
    take-simp1: take k (s_i#s_j#ps) = s_i#take (k-1) (s_j#ps) and
    take-simp2: take k (s_i#s_j#ps) = s_i#s_j#take (k-1-1) ps
  using assms 2 take-Cons'[of k s_i s_j#ps] take-Cons'[of k-1 s_j ps] by auto
  then have knights-path (set (take (k-1) (s_j#ps))) (take (k-1) (s_j#ps))
  using 2 k-simps by auto
  then have kp: knights-path (set (take (k-1) (s_j#ps))) (s_j#take (k-2) ps)
  using take-Cons'[of k-1 s_j ps] by (auto simp: k-simps elim: knights-path.cases)

  have no-mem: s_i ∉ set (take (k-1) (s_j#ps))
  using 2 set-take-subset[of k-1 s_j#ps] knights-path-set-eq by blast
  have knights-path (set (take (k-1) (s_j#ps)) ∪ {s_i}) (s_i#s_j#take (k-2) ps)
  using knights-path.intros(2)[OF no-mem ⟨valid-step s_i s_j⟩ kp] by auto
  then show ?thesis using k-simps take-simp2 knights-path-set-eq by metis
qed (auto intro: knights-path.intros)
qed auto

lemma knights-path-drop:
  assumes knights-path b ps 0 < k k < length ps
  shows knights-path (set (drop k ps)) (drop k ps)
  using assms
proof (induction arbitrary: k rule: knights-path.induct)
  case (2 s_i b s_j ps)
  then have (k = 1 ∧ ps = []) ∨ (k = 1 ∧ ps ≠ []) ∨ 1 < k by force
  then show ?case
  using 2
proof (elim disjE)
  assume k = 1 ∧ ps ≠ []
  then show ?thesis using 2 knights-path-set-eq by force
next
  assume 1 < k
  then have 0 < k-1 k-1 < length (s_j#ps) drop k (s_i#s_j#ps) = drop (k-1)
(s_j#ps)
  using assms 2 drop-Cons'[of k s_i s_j#ps] by auto
  then show ?thesis
  using 2 by auto
qed (auto intro: knights-path.intros)
qed auto

```

A Knight's path can be split to form two new disjoint Knight's paths.

corollary *knights-path-split*:

```

  assumes knights-path b ps 0 < k k < length ps
  shows
    ∃ b1 b2. knights-path b1 (take k ps) ∧ knights-path b2 (drop k ps) ∧ b1 ∪ b2 = b
    ∧ b1 ∩ b2 = {}
  using assms

```

proof –
 let $?b_1 = \text{set } (\text{take } k \text{ } ps)$
 let $?b_2 = \text{set } (\text{drop } k \text{ } ps)$
 have $kp1: \text{knight-path } ?b_1 \text{ (take } k \text{ } ps)$ and $kp2: \text{knight-path } ?b_2 \text{ (drop } k \text{ } ps)$
 using *assms knight-path-take knight-path-drop* by *auto*
 have $\text{union: } ?b_1 \cup ?b_2 = b$
 using *assms knight-path-set-eq* by (*metis append-take-drop-id set-append*)
 have $\text{inter: } ?b_1 \cap ?b_2 = \{\}$
 using *assms knight-path-distinct* by (*metis append-take-drop-id distinct-append*)
 show *?thesis* using $kp1 \text{ } kp2 \text{ } \text{union} \text{ } \text{inter}$ by *auto*
qed

Append two disjoint Knight's paths.

corollary *knight-path-append*:

assumes *knight-path* $b_1 \text{ } ps_1$ *knight-path* $b_2 \text{ } ps_2$ $b_1 \cap b_2 = \{\}$ *valid-step* (*last* ps_1) (*hd* ps_2)
 shows *knight-path* $(b_1 \cup b_2)$ $(ps_1 @ ps_2)$
 using *assms*
proof (*induction arbitrary: ps_2 b_2 rule: knight-path.induct*)
 case ($1 \text{ } s_i$)
 then have $s_i \notin b_2$ $ps_2 \neq []$ *valid-step* s_i (*hd* ps_2) *knight-path* b_2 (*hd* $ps_2 \# \text{tl } ps_2$)

using *knight-path-non-nil* by *auto*
 then have *knight-path* $(b_2 \cup \{s_i\})$ $(s_i \# \text{hd } ps_2 \# \text{tl } ps_2)$
 using *knight-path.intros* by *blast*
 then show *?case* using $\langle ps_2 \neq [] \rangle$ by *auto*

next

case ($2 \text{ } s_i \text{ } b_1 \text{ } s_j \text{ } ps_1$)
 then have $s_i \notin b_1 \cup b_2$ *valid-step* $s_i \text{ } s_j$ *knight-path* $(b_1 \cup b_2)$ $(s_j \# ps_1 @ ps_2)$ by *auto*
 then have *knight-path* $(b_1 \cup b_2 \cup \{s_i\})$ $(s_i \# s_j \# ps_1 @ ps_2)$
 using *knight-path.intros* by *auto*
 then show *?case* by *auto*

qed

lemma *valid-step-rev*: *valid-step* $s_i \text{ } s_j \implies \text{valid-step } s_j \text{ } s_i$
 using *step-checker-correct step-checker-rev* by (*metis prod.exhaust-sel*)

Reverse a Knight's path.

corollary *knight-path-rev*:

assumes *knight-path* $b \text{ } ps$
 shows *knight-path* b (*rev* ps)
 using *assms*
proof (*induction rule: knight-path.induct*)
 case ($2 \text{ } s_i \text{ } b \text{ } s_j \text{ } ps$)
 then have *knight-path* $\{s_i\} [s_i] \text{ } b \cap \{s_i\} = \{\}$ *valid-step* (*last* (*rev* $(s_j \# ps)$)) (*hd* $[s_i]$)
 using *valid-step-rev* by (*auto intro: knight-path.intros*)
 then have *knight-path* $(b \cup \{s_i\}) ((\text{rev } (s_j \# ps)) @ [s_i])$

```

    using 2 knights-path-append by blast
  then show ?case by auto
qed (auto intro: knights-path.intros)

```

Reverse a Knight's circuit.

```

corollary knights-circuit-rev:
  assumes knights-circuit b ps
  shows knights-circuit b (rev ps)
  using assms knights-path-rev valid-step-rev
  unfolding knights-circuit-def by (auto simp: hd-rev last-rev)

```

```

lemma knights-circuit-rotate1:
  assumes knights-circuit b (s_i#ps)
  shows knights-circuit b (ps@[s_i])
proof (cases ps = [])
  case True
  then show ?thesis using assms by auto
next
  case False
  have kp1: knights-path b (s_i#ps) valid-step (last (s_i#ps)) (hd (s_i#ps))
    using assms unfolding knights-circuit-def by auto
  then have kp-elim: s_i ∉ (b - {s_i}) valid-step s_i (hd ps) knights-path (b - {s_i})
  ps
    using ⟨ps ≠ []⟩ by (auto elim: knights-path.cases)
  then have vs': valid-step (last (ps@[s_i])) (hd (ps@[s_i]))
    using ⟨ps ≠ []⟩ valid-step-rev by auto

  have kp2: knights-path {s_i} [s_i] (b - {s_i}) ∩ {s_i} = {}
    by (auto intro: knights-path.intros)

  have vs: valid-step (last ps) (hd [s_i])
    using ⟨ps ≠ []⟩ ⟨valid-step (last (s_i#ps)) (hd (s_i#ps))⟩ by auto

  have (b - {s_i}) ∪ {s_i} = b
    using kp1 kp-elim knights-path-set-eq by force
  then show ?thesis
    unfolding knights-circuit-def
    using vs knights-path-append[OF ⟨knights-path (b - {s_i}) ps⟩ kp2] vs' by auto
qed

```

A Knight's circuit can be rotated to start at any square on the board.

```

lemma knights-circuit-rotate-to:
  assumes knights-circuit b ps hd (drop k ps) = s_i k < length ps
  shows ∃ ps'. knights-circuit b ps' ∧ hd ps' = s_i
  using assms
proof (induction k arbitrary: b ps)

```

```

case (Suc k)
let ?sj=hd ps
let ?ps'=tl ps
show ?case
proof (cases si = ?sj)
  case True
  then show ?thesis using Suc by auto
next
  case False
  then have ?ps' ≠ []
  using Suc by (metis drop-Nil drop-Suc drop-eq-Nil2 le-antisym nat-less-le)
  then have knight-circuit b (?sj#?ps')
  using Suc by (metis list.exhaust-sel tl-Nil)
  then have knight-circuit b (?ps'@[?sj]) hd (drop k (?ps'@[?sj])) = si
  using Suc knight-circuit-rotate1 by (auto simp: drop-Suc)
  then show ?thesis using Suc by auto
qed
qed auto

```

For positive boards (1,1) can only have (2,3) and (3,2) as a neighbour.

```

lemma valid-step-1-1:
  assumes valid-step (1,1) (i,j) i > 0 j > 0
  shows (i,j) = (2,3) ∨ (i,j) = (3,2)
  using assms unfolding valid-step-def by auto

```

```

lemma list-len-g-1-split: length xs > 1 ⟹ ∃ x1 x2 xs'. xs = x1#x2#xs'
proof (induction xs)
  case (Cons x xs)
  then have length xs > 0 by auto
  then have length xs ≥ 1 by presburger
  then have length xs = 1 ∨ length xs > 1 by auto
  then show ?case
  proof (elim disjE)
    assume length xs = 1
    then obtain x1 where [simp]: xs = [x1]
    using length-Suc-conv[of xs 0] by auto
    then show ?thesis by auto
  next
    assume 1 < length xs
    then show ?thesis using Cons by auto
  qed
qed auto

```

```

lemma list-len-g-3-split: length xs > 3 ⟹ ∃ x1 x2 xs' x3. xs = x1#x2#xs'@[x3]
proof (induction xs)
  case (Cons x xs)
  then have length xs = 3 ∨ length xs > 3 by auto
  then show ?case
  proof (elim disjE)

```

```

assume length xs = 3
then obtain  $x_1\ xs_1$  where  $[simp]: xs = x_1 \# xs_1$  length xs1 = 2
  using length-Suc-conv[of xs 2] by auto
then obtain  $x_2\ xs_2$  where  $[simp]: xs_1 = x_2 \# xs_2$  length xs2 = 1
  using length-Suc-conv[of xs1 1] by auto
then obtain  $x_3$  where  $[simp]: xs_2 = [x_3]$ 
  using length-Suc-conv[of xs2 0] by auto
then show ?thesis by auto
next
  assume length xs > 3
  then show ?thesis using Cons by auto
qed
qed auto

```

Any Knight's circuit on a positive board can be rotated to start with (1,1) and end with (3,2).

corollary *rotate-knights-circuit*:

```

assumes knights-circuit (board n m) ps min n m ≥ 5
shows  $\exists ps. \text{knights-circuit (board } n\ m) ps \wedge \text{hd } ps = (1,1) \wedge \text{last } ps = (3,2)$ 
using assms
proof –
  let ?b=board n m
  have knights-path ?b ps
    using assms unfolding knights-circuit-def by auto
  then have  $(1,1) \in \text{set } ps$ 
    using assms knights-path-set-eq by  $(\text{auto simp: board-def})$ 
  then obtain  $k$  where  $\text{hd (drop } k\ ps) = (1,1) \ k < \text{length } ps$ 
    by  $(\text{metis hd-drop-conv-nth in-set-conv-nth})$ 
  then obtain  $ps_r$  where psr-prems: knights-circuit ?b psr hd psr = (1,1)
    using assms knights-circuit-rotate-to by blast
  then have  $kp: \text{knights-path ?b } ps_r$  and valid-step (last psr) (1,1)
    unfolding knights-circuit-def by auto

  have  $(1,1) \in ?b\ (1,2) \in ?b\ (1,3) \in ?b$ 
    using assms unfolding board-def by auto
  then have  $(1,1) \in \text{set } ps_r\ (1,2) \in \text{set } ps_r\ (1,3) \in \text{set } ps_r$ 
    using  $kp$  knights-path-set-eq by auto

  have  $3 < \text{card } ?b$ 
    using assms board-leq-subset card-board[of 5 5]
      card-mono[OF board-finite[of n m], of board 5 5] by auto
  then have  $3 < \text{length } ps_r$ 
    using knights-path-length kp by auto
  then obtain  $s_j\ ps'\ s_k$  where  $[simp]: ps_r = (1,1) \# s_j \# ps' @ [s_k]$ 
    using  $\langle \text{hd } ps_r = (1,1) \rangle \text{ list-len-g-3-split[of } ps_r]$  by auto
  have  $s_j \neq s_k$ 
    using  $kp$  knights-path-distinct by force

  have vs-sk: valid-step sk (1,1)

```

```

using ⟨valid-step (last psr) (1,1)⟩ by simp

have vs-sj: valid-step (1,1) sj and kp': knights-path (?b - {(1,1)}) (sj#ps'@[sk])
  using kp by (auto elim: knights-path.cases)

have sj ∈ set psr sk ∈ set psr by auto
then have sj ∈ ?b sk ∈ ?b
  using kp knights-path-set-eq by blast+
then have 0 < fst sj ∧ 0 < snd sj 0 < fst sk ∧ 0 < snd sk
  unfolding board-def by auto
then have sk = (2,3) ∨ sk = (3,2) sj = (2,3) ∨ sj = (3,2)
  using vs-sk vs-sj valid-step-1-1 valid-step-rev by (metis prod.collapse)+
then have sk = (3,2) ∨ sj = (3,2)
  using ⟨sj ≠ sk⟩ by auto
then show ?thesis
proof (elim disjE)
  assume sk = (3,2)
  then have last psr = (3,2) by auto
  then show ?thesis using psr-prems by auto
next
  assume sj = (3,2)
  then have vs: valid-step (last ((1,1)#rev (sj#ps'@[sk]))) (hd ((1,1)#rev
(sj#ps'@[sk])))
    unfolding valid-step-def by auto

  have rev-simp: rev (sj#ps'@[sk]) = sk#(rev ps')@[sj] by auto

  have knights-path (?b - {(1,1)}) (rev (sj#ps'@[sk]))
    using knights-path-rev[OF kp'] by auto
  then have (1,1) ∉ (?b - {(1,1)}) valid-step (1,1) sk
    knights-path (?b - {(1,1)}) (sk#(rev ps')@[sj])
    using assms vs-sk valid-step-rev by (auto simp: rev-simp)
  then have knights-path (?b - {(1,1)} ∪ {(1,1)}) ((1,1)#sk#(rev ps')@[sj])
    using knights-path.intros(2)[of (1,1) ?b - {(1,1)} sk (rev ps')@[sj]] by auto
  then have knights-path ?b ((1,1)#rev (sj#ps'@[sk]))
    using assms by (simp add: board-def insert-absorb rev-simp)
  then have knights-circuit ?b ((1,1)#rev (sj#ps'@[sk]))
    unfolding knights-circuit-def using vs by auto
  then show ?thesis
    using ⟨sj = (3,2)⟩ by auto
qed
qed

```

4 Transposing Paths and Boards

4.1 Implementation of Path and Board Transposition

definition *transpose-square* $s_i = (\text{case } s_i \text{ of } (i,j) \Rightarrow (j,i))$

```

fun transpose :: path  $\Rightarrow$  path where
  transpose [] = []
| transpose (si#ps) = (transpose-square si)#transpose ps

```

```

definition transpose-board :: board  $\Rightarrow$  board where
  transpose-board b  $\equiv$  {(j,i) | i j. (i,j)  $\in$  b}

```

4.2 Correctness of Path and Board Transposition

```

lemma transpose2: transpose-square (transpose-square si) = si
  unfolding transpose-square-def by (auto split: prod.splits)

```

```

lemma transpose-nil: ps = []  $\longleftrightarrow$  transpose ps = []
  using transpose.elims by blast

```

```

lemma transpose-length: length ps = length (transpose ps)
  by (induction ps) auto

```

```

lemma hd-transpose: ps  $\neq$  []  $\implies$  hd (transpose ps) = transpose-square (hd ps)
  by (induction ps) (auto simp: transpose-square-def)

```

```

lemma last-transpose: ps  $\neq$  []  $\implies$  last (transpose ps) = transpose-square (last ps)

```

```

proof (induction ps)
  case (Cons si ps)
  then show ?case
  proof (cases ps = [])
    case True
    then show ?thesis using Cons by (auto simp: transpose-square-def)
  next
    case False
    then show ?thesis using Cons transpose-nil by auto
  qed
qed auto

```

```

lemma take-transpose:

```

```

  shows take k (transpose ps) = transpose (take k ps)
proof (induction ps arbitrary: k)
  case Nil
  then show ?case by auto
next
  case (Cons si ps)
  then obtain i j where si = (i,j) by force
  then have k = 0  $\vee$  k > 0 by auto
  then show ?case
  proof (elim disjE)
    assume k > 0
    then show ?thesis using Cons.IH by (auto simp: <si = (i,j)> take-Cons')
  qed auto
qed

```

```

lemma drop-transpose:
  shows drop k (transpose ps) = transpose (drop k ps)
proof (induction ps arbitrary: k)
  case Nil
  then show ?case by auto
next
  case (Cons si ps)
  then obtain i j where si = (i,j) by force
  then have k = 0 ∨ k > 0 by auto
  then show ?case
  proof (elim disjE)
    assume k > 0
    then show ?thesis using Cons.IH by (auto simp: ⟨si = (i,j)⟩ drop-Cons')
  qed auto
qed

lemma transpose-board-correct: si ∈ b ⟷ (transpose-square si) ∈ transpose-board b
  unfolding transpose-board-def transpose-square-def by (auto split: prod.splits)

lemma transpose-board: transpose-board (board n m) = board m n
  unfolding board-def using transpose-board-correct by (auto simp: transpose-square-def)

lemma insert-transpose-board:
  insert (transpose-square si) (transpose-board b) = transpose-board (insert si b)
  unfolding transpose-board-def transpose-square-def by (auto split: prod.splits)

lemma transpose-board2: transpose-board (transpose-board b) = b
  unfolding transpose-board-def by auto

lemma transpose-union: transpose-board (b1 ∪ b2) = transpose-board b1 ∪ transpose-board b2
  unfolding transpose-board-def by auto

lemma transpose-valid-step:
  valid-step si sj ⟷ valid-step (transpose-square si) (transpose-square sj)
  unfolding valid-step-def transpose-square-def by (auto split: prod.splits)

lemma transpose-knights-path':
  assumes knights-path b ps
  shows knights-path (transpose-board b) (transpose ps)
  using assms
proof (induction rule: knights-path.induct)
  case (1 si)
  then have transpose-board {si} = {transpose-square si} transpose [si] = [transpose-square si]
  using transpose-board-correct by (auto simp: transpose-square-def split: prod.splits)
  then show ?case by (auto intro: knights-path.intros)

```



```

next
case (2 si b sj ps)
then have prems: transpose-square si ∉ transpose-board b
      valid-step (transpose-square si) (transpose-square sj)
      and transpose (sj#ps) = transpose-square sj#transpose ps
      using 2 transpose-board-correct transpose-valid-step by auto
then show ?case
      using 2 knights-path.intros(2)[OF prems] insert-transpose-board by auto
qed

```

corollary *transpose-knights-path*:
 assumes *knights-path* (board n m) ps
 shows *knights-path* (board m n) (transpose ps)
 using *assms transpose-knights-path*[of board n m ps] by (auto simp: transpose-board)

corollary *transpose-knights-circuit*:
 assumes *knights-circuit* (board n m) ps
 shows *knights-circuit* (board m n) (transpose ps)
 using *assms*
proof –
 have *knights-path* (board n m) ps and *vs*: *valid-step* (last ps) (hd ps)
 using *assms unfolding knights-circuit-def* by auto
 then have *kp-t*: *knights-path* (board m n) (transpose ps) and ps ≠ []
 using *transpose-knights-path knights-path-non-nil* by auto
 then have *valid-step* (last (transpose ps)) (hd (transpose ps))
 using *vs hd-transpose last-transpose transpose-valid-step* by auto
 then show ?thesis using *kp-t* by (auto simp: *knights-circuit-def*)
 qed

5 Mirroring Paths and Boards

5.1 Implementation of Path and Board Mirroring

abbreviation *min1* ps ≡ *Min* ((fst) ‘ set ps)
abbreviation *max1* ps ≡ *Max* ((fst) ‘ set ps)
abbreviation *min2* ps ≡ *Min* ((snd) ‘ set ps)
abbreviation *max2* ps ≡ *Max* ((snd) ‘ set ps)

definition *mirror1-square* :: int ⇒ square ⇒ square **where**
mirror1-square n s_i = (case s_i of (i,j) ⇒ (n-i,j))

fun *mirror1-aux* :: int ⇒ path ⇒ path **where**
mirror1-aux n [] = []
| *mirror1-aux* n (s_i#ps) = (*mirror1-square* n s_i)#*mirror1-aux* n ps

definition *mirror1* ps = *mirror1-aux* (*max1* ps + *min1* ps) ps

definition *mirror1-board* :: int ⇒ board ⇒ board **where**

$mirror1-board\ n\ b \equiv \{mirror1-square\ n\ s_i \mid s_i. s_i \in b\}$

definition $mirror2-square :: int \Rightarrow square \Rightarrow square$ **where**
 $mirror2-square\ m\ s_i = (case\ s_i\ of\ (i,j) \Rightarrow (i,m-j))$

fun $mirror2-aux :: int \Rightarrow path \Rightarrow path$ **where**
 $mirror2-aux\ m\ [] = []$
 $| mirror2-aux\ m\ (s_i \# ps) = (mirror2-square\ m\ s_i) \# mirror2-aux\ m\ ps$

definition $mirror2\ ps = mirror2-aux\ (max2\ ps + min2\ ps)\ ps$

definition $mirror2-board :: int \Rightarrow board \Rightarrow board$ **where**
 $mirror2-board\ m\ b \equiv \{mirror2-square\ m\ s_i \mid s_i. s_i \in b\}$

5.2 Correctness of Path and Board Mirroring

lemma $mirror1-board-id: mirror1-board\ (int\ n+1)\ (board\ n\ m) = board\ n\ m$ **(is -**
 $=\ ?b)$

proof

show $mirror1-board\ (int\ n+1)\ ?b \subseteq ?b$

proof

fix s_i'

assume $assms: s_i' \in mirror1-board\ (int\ n+1)\ ?b$

then obtain $i'\ j'$ **where** $[simp]: s_i' = (i',j')$ **by force**

then have $(i',j') \in mirror1-board\ (int\ n+1)\ ?b$

using $assms$ **by auto**

then obtain $i\ j$ **where** $(i,j) \in ?b\ mirror1-square\ (int\ n+1)\ (i,j) = (i',j')$

unfolding $mirror1-board-def$ **by auto**

then have $1 \leq i \wedge i \leq int\ n\ 1 \leq j \wedge j \leq int\ m\ i'=(int\ n+1)-i\ j'=j$

unfolding $board-def\ mirror1-square-def$ **by auto**

then have $1 \leq i' \wedge i' \leq int\ n\ 1 \leq j' \wedge j' \leq int\ m$

by auto

then show $s_i' \in ?b$

unfolding $board-def$ **by auto**

qed

next

show $?b \subseteq mirror1-board\ (int\ n+1)\ ?b$

proof

fix s_i

assume $assms: s_i \in ?b$

then obtain $i\ j$ **where** $[simp]: s_i = (i,j)$ **by force**

then have $(i,j) \in ?b$

using $assms$ **by auto**

then have $1 \leq i \wedge i \leq int\ n\ 1 \leq j \wedge j \leq int\ m$

unfolding $board-def$ **by auto**

then obtain $i'\ j'$ **where** $i'=(int\ n+1)-i\ j'=j$ **by auto**

then have $(i',j') \in ?b\ mirror1-square\ (int\ n+1)\ (i',j') = (i,j)$

using $\langle 1 \leq i \wedge i \leq int\ n \rangle\ \langle 1 \leq j \wedge j \leq int\ m \rangle$

unfolding $mirror1-square-def$ **by** $(auto\ simp: board-def)$

```

    then show  $s_i \in \text{mirror1-board } (\text{int } n+1) \text{ ?b}$ 
      unfolding mirror1-board-def by force
    qed
  qed

lemma mirror2-board-id:  $\text{mirror2-board } (\text{int } m+1) (\text{board } n \ m) = \text{board } n \ m$  (is -
= ?b)
proof
  show  $\text{mirror2-board } (\text{int } m+1) \text{ ?b} \subseteq \text{?b}$ 
  proof
    fix  $s_i'$ 
    assume assms:  $s_i' \in \text{mirror2-board } (\text{int } m+1) \text{ ?b}$ 
    then obtain  $i' \ j'$  where [simp]:  $s_i' = (i', j')$  by force
    then have  $(i', j') \in \text{mirror2-board } (\text{int } m+1) \text{ ?b}$ 
      using assms by auto
    then obtain  $i \ j$  where  $(i, j) \in \text{?b mirror2-square } (\text{int } m+1) (i, j) = (i', j')$ 
      unfolding mirror2-board-def by auto
    then have  $1 \leq i \wedge i \leq \text{int } n \ 1 \leq j \wedge j \leq \text{int } m \ i' = i \ j' = (\text{int } m+1) - j$ 
      unfolding board-def mirror2-square-def by auto
    then have  $1 \leq i' \wedge i' \leq \text{int } n \ 1 \leq j' \wedge j' \leq \text{int } m$ 
      by auto
    then show  $s_i' \in \text{?b}$ 
      unfolding board-def by auto
    qed
  next
  show  $\text{?b} \subseteq \text{mirror2-board } (\text{int } m+1) \text{ ?b}$ 
  proof
    fix  $s_i$ 
    assume assms:  $s_i \in \text{?b}$ 
    then obtain  $i \ j$  where [simp]:  $s_i = (i, j)$  by force
    then have  $(i, j) \in \text{?b}$ 
      using assms by auto
    then have  $1 \leq i \wedge i \leq \text{int } n \ 1 \leq j \wedge j \leq \text{int } m$ 
      unfolding board-def by auto
    then obtain  $i' \ j'$  where  $i' = i \ j' = (\text{int } m+1) - j$  by auto
    then have  $(i', j') \in \text{?b mirror2-square } (\text{int } m+1) (i', j') = (i, j)$ 
      using  $\langle 1 \leq i \wedge i \leq \text{int } n \rangle \langle 1 \leq j \wedge j \leq \text{int } m \rangle$ 
      unfolding mirror2-square-def by (auto simp: board-def)
    then show  $s_i \in \text{mirror2-board } (\text{int } m+1) \text{ ?b}$ 
      unfolding mirror2-board-def by force
    qed
  qed

lemma knight-path-min1:  $\text{knight-path } (\text{board } n \ m) \ ps \implies \text{min1 } ps = 1$ 
proof -
  assume assms:  $\text{knight-path } (\text{board } n \ m) \ ps$ 
  then have  $\text{min } n \ m \geq 1$ 
    using knight-path-board-m-n-geq-1 by auto
  then have  $(1, 1) \in \text{board } n \ m$  and ge-1:  $\forall (i, j) \in \text{board } n \ m. i \geq 1$ 

```

unfolding board-def by auto
 then have finite: finite ((fst) ‘ board n m) and
 non-empty: (fst) ‘ board n m $\neq \{\}$ and
 mem-1: 1 \in (fst) ‘ board n m
 using board-finite by auto (metis fstI image-eqI)
 then have Min ((fst) ‘ board n m) = 1
 using ge-1 by (auto simp: Min-eq-iff)
 then show ?thesis
 using assms knights-path-set-eq by auto
 qed

lemma knights-path-min2: knights-path (board n m) ps \implies min2 ps = 1
 proof –

assume assms: knights-path (board n m) ps
 then have min n m ≥ 1
 using knights-path-board-m-n-geq-1 by auto
 then have (1,1) \in board n m and ge-1: $\forall (i,j) \in$ board n m. $j \geq 1$
 unfolding board-def by auto
 then have finite: finite ((snd) ‘ board n m) and
 non-empty: (snd) ‘ board n m $\neq \{\}$ and
 mem-1: 1 \in (snd) ‘ board n m
 using board-finite by auto (metis sndI image-eqI)
 then have Min ((snd) ‘ board n m) = 1
 using ge-1 by (auto simp: Min-eq-iff)
 then show ?thesis
 using assms knights-path-set-eq by auto
 qed

lemma knights-path-max1: knights-path (board n m) ps \implies max1 ps = int n
 proof –

assume assms: knights-path (board n m) ps
 then have min n m ≥ 1
 using knights-path-board-m-n-geq-1 by auto
 then have (int n,1) \in board n m and leq-n: $\forall (i,j) \in$ board n m. $i \leq$ int n
 unfolding board-def by auto
 then have finite: finite ((fst) ‘ board n m) and
 non-empty: (fst) ‘ board n m $\neq \{\}$ and
 mem-1: int n \in (fst) ‘ board n m
 using board-finite by auto (metis fstI image-eqI)
 then have Max ((fst) ‘ board n m) = int n
 using leq-n by (auto simp: Max-eq-iff)
 then show ?thesis
 using assms knights-path-set-eq by auto
 qed

lemma knights-path-max2: knights-path (board n m) ps \implies max2 ps = int m
 proof –

assume assms: knights-path (board n m) ps
 then have min n m ≥ 1

using *knight-path-board-m-n-geq-1* by auto
 then have $(1, \text{int } m) \in \text{board } n \ m$ and $\text{leq-}m: \forall (i,j) \in \text{board } n \ m. j \leq \text{int } m$
 unfolding *board-def* by auto
 then have *finite*: $\text{finite } ((\text{snd}) \text{ ` board } n \ m)$ and
 non-empty: $(\text{snd}) \text{ ` board } n \ m \neq \{\}$ and
 mem-1: $\text{int } m \in (\text{snd}) \text{ ` board } n \ m$
 using *board-finite* by auto (*metis sndI image-eqI*)
 then have $\text{Max } ((\text{snd}) \text{ ` board } n \ m) = \text{int } m$
 using *leq-m* by (auto *simp*: *Max-eq-iff*)
 then show *?thesis*
 using *assms knight-path-set-eq* by auto
 qed

lemma *mirror1-aux-nil*: $ps = [] \longleftrightarrow \text{mirror1-aux } m \ ps = []$
 using *mirror1-aux.elims* by blast

lemma *mirror1-nil*: $ps = [] \longleftrightarrow \text{mirror1 } ps = []$
 unfolding *mirror1-def* using *mirror1-aux-nil* by blast

lemma *mirror2-aux-nil*: $ps = [] \longleftrightarrow \text{mirror2-aux } m \ ps = []$
 using *mirror2-aux.elims* by blast

lemma *mirror2-nil*: $ps = [] \longleftrightarrow \text{mirror2 } ps = []$
 unfolding *mirror2-def* using *mirror2-aux-nil* by blast

lemma *length-mirror1-aux*: $\text{length } ps = \text{length } (\text{mirror1-aux } n \ ps)$
 by (*induction ps*) auto

lemma *length-mirror1*: $\text{length } ps = \text{length } (\text{mirror1 } ps)$
 unfolding *mirror1-def* using *length-mirror1-aux* by auto

lemma *length-mirror2-aux*: $\text{length } ps = \text{length } (\text{mirror2-aux } n \ ps)$
 by (*induction ps*) auto

lemma *length-mirror2*: $\text{length } ps = \text{length } (\text{mirror2 } ps)$
 unfolding *mirror2-def* using *length-mirror2-aux* by auto

lemma *mirror1-board-iff*: $s_i \notin b \longleftrightarrow \text{mirror1-square } n \ s_i \notin \text{mirror1-board } n \ b$
 unfolding *mirror1-board-def* *mirror1-square-def* by (auto *split*: *prod.splits*)

lemma *mirror2-board-iff*: $s_i \notin b \longleftrightarrow \text{mirror2-square } n \ s_i \notin \text{mirror2-board } n \ b$
 unfolding *mirror2-board-def* *mirror2-square-def* by (auto *split*: *prod.splits*)

lemma *insert-mirror1-board*:
 $\text{insert } (\text{mirror1-square } n \ s_i) \ (\text{mirror1-board } n \ b) = \text{mirror1-board } n \ (\text{insert } s_i \ b)$
 unfolding *mirror1-board-def* *mirror1-square-def* by (auto *split*: *prod.splits*)

lemma *insert-mirror2-board*:
 $\text{insert } (\text{mirror2-square } n \ s_i) \ (\text{mirror2-board } n \ b) = \text{mirror2-board } n \ (\text{insert } s_i \ b)$

```

unfolding mirror2-board-def mirror2-square-def by (auto split: prod.splits)

lemma (i::int) = i'+1  $\implies$  n-i=n-(i'+1)
  by auto

lemma valid-step-mirror1:
  valid-step  $s_i s_j \iff$  valid-step (mirror1-square n  $s_i$ ) (mirror1-square n  $s_j$ )
proof
  assume assms: valid-step  $s_i s_j$ 
  obtain  $i j i' j'$  where [simp]:  $s_i = (i,j)$   $s_j = (i',j')$  by force
  then have valid-step (n-i,j) (n-i',j')
    using assms unfolding valid-step-def
    apply simp
    apply (elim disjE)
    apply auto
  done
  then show valid-step (mirror1-square n  $s_i$ ) (mirror1-square n  $s_j$ )
    unfolding mirror1-square-def by auto
next
  assume assms: valid-step (mirror1-square n  $s_i$ ) (mirror1-square n  $s_j$ )
  obtain  $i j i' j'$  where [simp]:  $s_i = (i,j)$   $s_j = (i',j')$  by force
  then have valid-step (i,j) (i',j')
    using assms unfolding valid-step-def mirror1-square-def
    apply simp
    apply (elim disjE)
    apply auto
  done
  then show valid-step  $s_i s_j$ 
    unfolding mirror1-square-def by auto
qed

lemma valid-step-mirror2:
  valid-step  $s_i s_j \iff$  valid-step (mirror2-square m  $s_i$ ) (mirror2-square m  $s_j$ )
proof
  assume assms: valid-step  $s_i s_j$ 
  obtain  $i j i' j'$  where [simp]:  $s_i = (i,j)$   $s_j = (i',j')$  by force
  then have valid-step (i,m-j) (i',m-j')
    using assms unfolding valid-step-def
    apply simp
    apply (elim disjE)
    apply auto
  done
  then show valid-step (mirror2-square m  $s_i$ ) (mirror2-square m  $s_j$ )
    unfolding mirror2-square-def by auto
next
  assume assms: valid-step (mirror2-square m  $s_i$ ) (mirror2-square m  $s_j$ )
  obtain  $i j i' j'$  where [simp]:  $s_i = (i,j)$   $s_j = (i',j')$  by force
  then have valid-step (i,j) (i',j')
    using assms unfolding valid-step-def mirror2-square-def

```

```

    apply simp
    apply (elim disjE)
    apply auto
    done
  then show valid-step  $s_i s_j$ 
    unfolding mirror1-square-def by auto
qed

```

```

lemma hd-mirror1:
  assumes knights-path (board  $n m$ )  $ps$   $hd\ ps = (i,j)$ 
  shows  $hd\ (mirror1\ ps) = (int\ n+1-i,j)$ 
  using assms
proof -
  have  $hd\ (mirror1\ ps) = hd\ (mirror1\ aux\ (int\ n+1)\ ps)$ 
    unfolding mirror1-def using assms knights-path-min1 knights-path-max1 by
  auto
  also have  $\dots = hd\ (mirror1\ aux\ (int\ n+1)\ ((hd\ ps)\#(tl\ ps)))$ 
    using assms knights-path-non-nil by (metis list.collapse)
  also have  $\dots = (int\ n+1-i,j)$ 
    using assms by (auto simp: mirror1-square-def)
  finally show ?thesis .
qed

```

```

lemma last-mirror1-aux:
  assumes  $ps \neq []$   $last\ ps = (i,j)$ 
  shows  $last\ (mirror1\ aux\ n\ ps) = (n-i,j)$ 
  using assms
proof (induction  $ps$ )
  case (Cons  $s_i\ ps$ )
  then show ?case
    using mirror1-aux-nil Cons by (cases  $ps = []$ ) (auto simp: mirror1-square-def)
qed auto

```

```

lemma last-mirror1:
  assumes knights-path (board  $n m$ )  $ps$   $last\ ps = (i,j)$ 
  shows  $last\ (mirror1\ ps) = (int\ n+1-i,j)$ 
  unfolding mirror1-def using assms last-mirror1-aux knights-path-non-nil
  by (simp add: knights-path-max1 knights-path-min1)

```

```

lemma hd-mirror2:
  assumes knights-path (board  $n m$ )  $ps$   $hd\ ps = (i,j)$ 
  shows  $hd\ (mirror2\ ps) = (i,int\ m+1-j)$ 
  using assms
proof -
  have  $hd\ (mirror2\ ps) = hd\ (mirror2\ aux\ (int\ m+1)\ ps)$ 
    unfolding mirror2-def using assms knights-path-min2 knights-path-max2 by
  auto
  also have  $\dots = hd\ (mirror2\ aux\ (int\ m+1)\ ((hd\ ps)\#(tl\ ps)))$ 
    using assms knights-path-non-nil by (metis list.collapse)

```

```

    also have ... = (i,int m+1-j)
      using assms by (auto simp: mirror2-square-def)
    finally show ?thesis .
qed

lemma last-mirror2-aux:
  assumes ps ≠ [] last ps = (i,j)
  shows last (mirror2-aux m ps) = (i,m-j)
  using assms
proof (induction ps)
  case (Cons si ps)
  then show ?case
    using mirror2-aux-nil Cons by (cases ps = []) (auto simp: mirror2-square-def)
qed auto

lemma last-mirror2:
  assumes knights-path (board n m) ps last ps = (i,j)
  shows last (mirror2 ps) = (i,int m+1-j)
  unfolding mirror2-def using assms last-mirror2-aux knights-path-non-nil
  by (simp add: knights-path-max2 knights-path-min2)

lemma mirror1-aux-knights-path:
  assumes knights-path b ps
  shows knights-path (mirror1-board n b) (mirror1-aux n ps)
  using assms
proof (induction rule: knights-path.induct)
  case (1 si)
  then have mirror1-board n {si} = {mirror1-square n si}
    unfolding mirror1-board-def by blast
  then show ?case by (auto intro: knights-path.intros)
next
  case (2 si b sj ps)
  then have prems: mirror1-square n si ∉ mirror1-board n b
    valid-step (mirror1-square n si) (mirror1-square n sj)
    and mirror1-aux n (sj#ps) = mirror1-square n sj#mirror1-aux n ps
    using 2 mirror1-board-iff valid-step-mirror1 by auto
  then show ?case
    using 2 knights-path.intros(2)[OF prems] insert-mirror1-board by auto
qed

corollary mirror1-knights-path:
  assumes knights-path (board n m) ps
  shows knights-path (board n m) (mirror1 ps)
  using assms
proof -
  have [simp]: min1 ps = 1 max1 ps = int n
    using assms knights-path-min1 knights-path-max1 by auto
  then have mirror1-board (int n+1) (board n m) = (board n m)
    using mirror1-board-id by auto

```



```

    then have knights-path (board n m) (mirror1-aux (int n+1) ps)
      using assms mirror1-aux-knights-path[of board n m ps int n+1] by auto
    then show ?thesis unfolding mirror1-def by auto
qed

lemma mirror2-aux-knights-path:
  assumes knights-path b ps
  shows knights-path (mirror2-board n b) (mirror2-aux n ps)
  using assms
proof (induction rule: knights-path.induct)
  case (1 si)
  then have mirror2-board n {si} = {mirror2-square n si}
    unfolding mirror2-board-def by blast
  then show ?case by (auto intro: knights-path.intros)
next
  case (2 si b sj ps)
  then have prems: mirror2-square n si ∉ mirror2-board n b
    valid-step (mirror2-square n si) (mirror2-square n sj)
    and mirror2-aux n (sj#ps) = mirror2-square n sj#mirror2-aux n ps
    using 2 mirror2-board-iff valid-step-mirror2 by auto
  then show ?case
    using 2 knights-path.intros(2)[OF prems] insert-mirror2-board by auto
qed

corollary mirror2-knights-path:
  assumes knights-path (board n m) ps
  shows knights-path (board n m) (mirror2 ps)
proof -
  have [simp]: min2 ps = 1 max2 ps = int m
    using assms knights-path-min2 knights-path-max2 by auto
  then have mirror2-board (int m+1) (board n m) = (board n m)
    using mirror2-board-id by auto
  then have knights-path (board n m) (mirror2-aux (int m+1) ps)
    using assms mirror2-aux-knights-path[of board n m ps int m+1] by auto
  then show ?thesis unfolding mirror2-def by auto
qed

```

5.3 Rotate Knight's Paths

Transposing (*transpose*) and mirroring (along first axis *mirror1*) a Knight's path preserves the Knight's path's property. Tranpose+Mirror1 equals a 90deg-clockwise turn.

```

corollary rot90-knights-path:
  assumes knights-path (board n m) ps
  shows knights-path (board m n) (mirror1 (transpose ps))
  using assms transpose-knights-path mirror1-knights-path by auto

```

```

lemma hd-rot90-knights-path:
  assumes knights-path (board n m) ps hd ps = (i,j)

```

```

shows hd (mirror1 (transpose ps)) = (int m+1-j,i)
using assms
proof -
  have hd (transpose ps) = (j,i) knights-path (board m n) (transpose ps)
    using assms knights-path-non-nil hd-transpose transpose-knights-path
    by (auto simp: transpose-square-def)
  then show ?thesis using hd-mirror1 by auto
qed

```

```

lemma last-rot90-knights-path:
  assumes knights-path (board n m) ps last ps = (i,j)
  shows last (mirror1 (transpose ps)) = (int m+1-j,i)
  using assms
proof -
  have last (transpose ps) = (j,i) knights-path (board m n) (transpose ps)
    using assms knights-path-non-nil last-transpose transpose-knights-path
    by (auto simp: transpose-square-def)
  then show ?thesis using last-mirror1 by auto
qed

```

6 Translating Paths and Boards

When constructing knight's paths for larger boards multiple knight's paths for smaller boards are concatenated. To concatenate paths the the coordinates in the path need to be translated. Therefore, simple auxiliary functions are provided.

6.1 Implementation of Path and Board Translation

Translate the coordinates for a path by (k_1, k_2) .

```

fun trans-path :: int × int ⇒ path ⇒ path where
  trans-path (k1,k2) [] = []
| trans-path (k1,k2) ((i,j)#xs) = (i+k1,j+k2)#(trans-path (k1,k2) xs)

```

Translate the coordinates of a board by (k_1, k_2) .

```

definition trans-board :: int × int ⇒ board ⇒ board where
  trans-board t b ≡ (case t of (k1,k2) ⇒ {(i+k1,j+k2) | i j. (i,j) ∈ b})

```

6.2 Correctness of Path and Board Translation

```

lemma trans-path-length: length ps = length (trans-path (k1,k2) ps)
  by (induction ps) auto

```

```

lemma trans-path-non-nil: ps ≠ [] ⇒ trans-path (k1,k2) ps ≠ []
  by (induction ps) auto

```

lemma *trans-path-correct*: $(i,j) \in \text{set } ps \longleftrightarrow (i+k_1, j+k_2) \in \text{set } (\text{trans-path } (k_1, k_2) ps)$
proof (*induction ps*)
 case (*Cons s_i ps*)
 then show ?*case* **by** (*cases s_i*) *auto*
qed *auto*

lemma *trans-path-non-nil-last*:
 $ps \neq [] \implies \text{last } (\text{trans-path } (k_1, k_2) ps) = \text{last } (\text{trans-path } (k_1, k_2) ((i,j)\#ps))$
using *trans-path-non-nil* **by** (*induction ps*) *auto*

lemma *hd-trans-path*:
assumes $ps \neq []$ $\text{hd } ps = (i,j)$
shows $\text{hd } (\text{trans-path } (k_1, k_2) ps) = (i+k_1, j+k_2)$
using *assms* **by** (*induction ps*) *auto*

lemma *last-trans-path*:
assumes $ps \neq []$ $\text{last } ps = (i,j)$
shows $\text{last } (\text{trans-path } (k_1, k_2) ps) = (i+k_1, j+k_2)$
using *assms*
proof (*induction ps*)
 case (*Cons s_i ps*)
 then show ?*case*
 using *trans-path-non-nil-last[symmetric]*
 apply (*cases s_i*)
 apply (*cases ps = []*)
 apply *auto*
 done
qed (*auto*)

lemma *take-trans*:
shows $\text{take } k (\text{trans-path } (k_1, k_2) ps) = \text{trans-path } (k_1, k_2) (\text{take } k ps)$
proof (*induction ps arbitrary: k*)
 case *Nil*
 then show ?*case* **by** *auto*
next
 case (*Cons s_i ps*)
 then obtain *i j* **where** $s_i = (i,j)$ **by** *force*
 then have $k = 0 \vee k > 0$ **by** *auto*
 then show ?*case*
 proof (*elim disjE*)
 assume $k > 0$
 then show ?*thesis* **using** *Cons.IH* **by** (*auto simp: <s_i = (i,j)> take-Cons'*)
 qed *auto*
qed

lemma *drop-trans*:
shows $\text{drop } k (\text{trans-path } (k_1, k_2) ps) = \text{trans-path } (k_1, k_2) (\text{drop } k ps)$
proof (*induction ps arbitrary: k*)

```

    case Nil
    then show ?case by auto
next
case (Cons si ps)
then obtain i j where si = (i,j) by force
then have k = 0 ∨ k > 0 by auto
then show ?case
proof (elim disjE)
  assume k > 0
  then show ?thesis using Cons.IH by (auto simp: ⟨si = (i,j)⟩ drop-Cons')
qed auto
qed

lemma trans-board-correct: (i,j) ∈ b ⟷ (i+k1,j+k2) ∈ trans-board (k1,k2) b
  unfolding trans-board-def by auto

lemma board-subset: n1 ≤ n2 ⟹ m1 ≤ m2 ⟹ board n1 m1 ⊆ board n2 m2
  unfolding board-def by auto

Board concatenation

corollary board-concat:
  shows board n m1 ∪ trans-board (0,int m1) (board n m2) = board n (m1+m2)
  (is ?b1 ∪ ?b2 = ?b)
proof
  show ?b1 ∪ ?b2 ⊆ ?b unfolding board-def trans-board-def by auto
next
  show ?b ⊆ ?b1 ∪ ?b2
  proof
    fix x
    assume x ∈ ?b
    then obtain i j where x-split: x = (i,j) 1 ≤ i ∧ i ≤ int n 1 ≤ j ∧ j ≤ int
      (m1+m2)
    unfolding board-def by auto
    then have j ≤ int m1 ∨ (int m1 < j ∧ j ≤ int (m1+m2)) by auto
    then show x ∈ ?b1 ∪ ?b2
    proof
      assume j ≤ int m1
      then show x ∈ ?b1 ∪ ?b2 using x-split unfolding board-def by auto
    next
      assume asm: int m1 < j ∧ j ≤ int (m1+m2)
      then have (i,j-int m1) ∈ board n m2 using x-split unfolding board-def by
        auto
      then show x ∈ ?b1 ∪ ?b2
      using x-split asm trans-board-correct[of i j-int m1 board n m2 0 int m1] by
        auto
    qed
  qed
qed

```

lemma *transpose-trans-board*:

transpose-board (*trans-board* (k_1, k_2) *b*) = *trans-board* (k_2, k_1) (*transpose-board* *b*)

unfolding *transpose-board-def trans-board-def* **by** *blast*

corollary *board-concatT*:

shows *board* n_1 *m* \cup *trans-board* (*int* $n_1, 0$) (*board* n_2 *m*) = *board* ($n_1 + n_2$) *m* (**is** $?b_1 \cup ?b_2 = ?b$)

proof –

let $?b_1 T = \text{board } m \ n_1$

let $?b_2 T = \text{trans-board } (0, \text{int } n_1) (\text{board } m \ n_2)$

have $?b_1 \cup ?b_2 = \text{transpose-board } (?b_1 T \cup ?b_2 T)$

using *transpose-board2 transpose-union transpose-board transpose-trans-board*

by *auto*

also have ... = *transpose-board* (*board* *m* ($n_1 + n_2$))

using *board-concat* **by** *auto*

also have ... = *board* ($n_1 + n_2$) *m*

using *transpose-board* **by** *auto*

finally show *?thesis* .

qed

lemma *trans-valid-step*:

valid-step (*i, j*) (*i', j'*) \implies *valid-step* ($i + k_1, j + k_2$) ($i' + k_1, j' + k_2$)

unfolding *valid-step-def* **by** *auto*

Translating a path and a boards preserves the validity.

lemma *trans-knights-path*:

assumes *knights-path* *b ps*

shows *knights-path* (*trans-board* (k_1, k_2) *b*) (*trans-path* (k_1, k_2) *ps*)

using *assms*

proof (*induction rule: knights-path.induct*)

case ($2 \ s_i \ b \ s_j \ xs$)

then obtain $i \ j \ i' \ j'$ **where** *split*: $s_i = (i, j)$ $s_j = (i', j')$ **by** *force*

let $?s_i = (i + k_1, j + k_2)$

let $?s_j = (i' + k_1, j' + k_2)$

let $?xs = \text{trans-path } (k_1, k_2) \ xs$

let $?b = \text{trans-board } (k_1, k_2) \ b$

have *simps*: *trans-path* (k_1, k_2) ($s_i \# s_j \# xs$) = $?s_i \# ?s_j \# ?xs$

$?b \cup \{?s_i\} = \text{trans-board } (k_1, k_2) (b \cup \{s_i\})$

unfolding *trans-board-def* **using** *split* **by** *auto*

have $?s_i \notin ?b$ *valid-step* $?s_i \ ?s_j$ *knights-path* $?b$ ($?s_j \# ?xs$)

using $2 \ \text{split}$ *trans-valid-step* **by** (*auto simp: trans-board-def*)

then have *knights-path* ($?b \cup \{?s_i\}$) ($?s_i \# ?s_j \# ?xs$)

using *knights-path.intros* **by** *auto*

then show *?case* **using** *simps* **by** *auto*

qed (*auto simp: trans-board-def intro: knights-path.intros*)

Predicate that indicates if two squares s_i and s_j are adjacent in *ps*.

definition *step-in* :: *path* \Rightarrow *square* \Rightarrow *square* \Rightarrow *bool* **where**

step-in *ps* $s_i \ s_j \equiv (\exists k. \ 0 < k \wedge k < \text{length } ps \wedge \text{last } (\text{take } k \ ps) = s_i \wedge \text{hd } (\text{drop } k \ ps) = s_j)$

$k \text{ ps}) = s_j)$

lemma *step-in-Cons*: $\text{step-in } ps \ s_i \ s_j \implies \text{step-in } (s_k \# ps) \ s_i \ s_j$

proof –

assume *step-in* $ps \ s_i \ s_j$

then obtain k **where** $0 < k \wedge k < \text{length } ps$ **last** $(\text{take } k \ ps) = s_i \ \text{hd} \ (\text{drop } k \ ps) = s_j$

unfolding *step-in-def* **by** *auto*

then have $0 < k+1 \wedge k+1 < \text{length } (s_k \# ps)$

$\text{last } (\text{take } (k+1) \ (s_k \# ps)) = s_i \ \text{hd} \ (\text{drop } (k+1) \ (s_k \# ps)) = s_j$

by *auto*

then show *?thesis*

by $(\text{auto simp: step-in-def})$

qed

lemma *step-in-append*: $\text{step-in } ps \ s_i \ s_j \implies \text{step-in } (ps @ ps') \ s_i \ s_j$

proof –

assume *step-in* $ps \ s_i \ s_j$

then obtain k **where** $0 < k \wedge k < \text{length } ps$ **last** $(\text{take } k \ ps) = s_i \ \text{hd} \ (\text{drop } k \ ps) = s_j$

unfolding *step-in-def* **by** *auto*

then have $0 < k \wedge k < \text{length } (ps @ ps')$

$\text{last } (\text{take } k \ (ps @ ps')) = s_i \ \text{hd} \ (\text{drop } k \ (ps @ ps')) = s_j$

by *auto*

then show *?thesis*

by $(\text{auto simp: step-in-def})$

qed

lemma *step-in-prepend*: $\text{step-in } ps \ s_i \ s_j \implies \text{step-in } (ps' @ ps) \ s_i \ s_j$

using *step-in-Cons* **by** $(\text{induction } ps' \text{ arbitrary: } ps) \ \text{auto}$

lemma *step-in-valid-step*: $\text{knight-path } b \ ps \implies \text{step-in } ps \ s_i \ s_j \implies \text{valid-step } s_i \ s_j$

proof –

assume *assms*: *knight-path* $b \ ps$ *step-in* $ps \ s_i \ s_j$

then obtain k **where** *k-prems*: $0 < k \wedge k < \text{length } ps$ **last** $(\text{take } k \ ps) = s_i \ \text{hd} \ (\text{drop } k \ ps) = s_j$

unfolding *step-in-def* **by** *auto*

then have $k = 1 \vee k > 1$ **by** *auto*

then show *?thesis*

proof $(\text{elim } \text{disjE})$

assume $k = 1$

then obtain ps' **where** $ps = s_i \# s_j \# ps'$

using *k-prems* *list-len-g-1-split* **by** *fastforce*

then show *?thesis*

using *assms* **by** $(\text{auto elim: knight-path.cases})$

next

assume $k > 1$

then have $0 < k-1 \wedge k-1 < \text{length } ps$

```

    using k-prems by auto
  then obtain b where knights-path b (drop (k - 1) ps)
    using assms knights-path-split by blast

  obtain ps' where drop (k - 1) ps = si # sj # ps'
    using k-prems  $\langle 0 < k - 1 \wedge k - 1 < \text{length } ps \rangle$ 
  by (metis Cons-nth-drop-Suc Suc-diff-1 hd-drop-conv-nth last-snoc take-hd-drop)
  then show ?thesis
    using  $\langle \text{knights-path } b \text{ (drop (k-1) ps)} \rangle$  by (auto elim: knights-path.cases)
qed
qed

```

lemma *trans-step-in*:

```

 $\text{step-in } ps \text{ (i,j) (i',j')} \implies \text{step-in (trans-path (k}_1, k_2) ps) (i+k_1, j+k_2) (i'+k_1, j'+k_2)$ 
proof -
  let ?ps' = trans-path (k1, k2) ps
  assume step-in ps (i, j) (i', j')
  then obtain k where  $0 < k \wedge k < \text{length } ps$  last (take k ps) = (i, j) hd (drop k
ps) = (i', j')
    unfolding step-in-def by auto
  then have take k ps  $\neq []$  drop k ps  $\neq []$  by fastforce +
  then have  $0 < k \wedge k < \text{length } ?ps'$ 
    last (take k ?ps') = (i + k1, j + k2) hd (drop k ?ps') = (i' + k1, j' + k2)
  using trans-path-length
    last-trans-path[OF  $\langle \text{take } k \text{ } ps \neq [] \rangle$   $\langle \text{last (take } k \text{ } ps) = (i, j) \rangle$ ] take-trans
    hd-trans-path[OF  $\langle \text{drop } k \text{ } ps \neq [] \rangle$   $\langle \text{hd (drop } k \text{ } ps) = (i', j') \rangle$ ] drop-trans
  by auto
  then show ?thesis
    by (auto simp: step-in-def)
qed

```

lemma *transpose-step-in*:

```

 $\text{step-in } ps \text{ s}_i \text{ s}_j \implies \text{step-in (transpose } ps) (\text{transpose-square } s_i) (\text{transpose-square } s_j)$ 
(is -  $\implies \text{step-in } ?psT \text{ ?s}_i T \text{ ?s}_j T$ )
proof -
  assume step-in ps si sj
  then obtain k where
    k-prems:  $0 < k \wedge k < \text{length } ps$  last (take k ps) = si hd (drop k ps) = sj
    unfolding step-in-def by auto
  then have non-nil: take k ps  $\neq []$  drop k ps  $\neq []$  by fastforce +
  have take k ?psT = transpose (take k ps) drop k ?psT = transpose (drop k ps)
    using take-transpose drop-transpose by auto
  then have last (take k ?psT) = ?si T hd (drop k ?psT) = ?sj T
    using non-nil k-prems hd-transpose last-transpose by auto
  then show step-in ?psT ?si T ?sj T
    unfolding step-in-def using k-prems transpose-length by auto
qed

```

lemma *hd-take*: $0 < k \implies \text{hd } xs = \text{hd } (\text{take } k \text{ } xs)$
by (*induction xs*) *auto*

lemma *last-drop*: $k < \text{length } xs \implies \text{last } xs = \text{last } (\text{drop } k \text{ } xs)$
by (*induction xs*) *auto*

6.3 Concatenate Knight's Paths and Circuits

Concatenate two knight's path on a $n \times m$ -board along the 2nd axis if the first path contains the step $s_i \rightarrow s_j$ and there are valid steps $s_i \rightarrow \text{hd } ps_2'$ and $s_j \rightarrow \text{last } ps_2'$, where ps_2' is ps_2 is translated by m_1 . An arbitrary step in ps_2 is preserved.

corollary *knights-path-split-concat-si-prev*:

assumes *knights-path* (*board n m₁*) ps_1 *knights-path* (*board n m₂*) ps_2
 $\text{step-in } ps_1 \ s_i \ s_j \ \text{hd } ps_2 = (i_h, j_h) \ \text{last } ps_2 = (i_l, j_l) \ \text{step-in } ps_2 \ (i, j) \ (i', j')$
 $\text{valid-step } s_i \ (i_h, \text{int } m_1 + j_h) \ \text{valid-step } (i_l, \text{int } m_1 + j_l) \ s_j$
shows $\exists ps. \text{knights-path } (\text{board } n \ (m_1 + m_2)) \ ps \wedge \text{hd } ps = \text{hd } ps_1$
 $\wedge \text{last } ps = \text{last } ps_1 \wedge \text{step-in } ps \ (i, \text{int } m_1 + j) \ (i', \text{int } m_1 + j')$
using *assms*

proof –

let $?b_1 = \text{board } n \ m_1$
let $?b_2 = \text{board } n \ m_2$
let $?ps_2' = \text{trans-path } (0, \text{int } m_1) \ ps_2$
let $?b' = \text{trans-board } (0, \text{int } m_1) \ ?b_2$
have $kp2'$: *knights-path* $?b' \ ?ps_2'$ **using** *assms trans-knights-path* **by** *auto*
then have $?ps_2' \neq []$ **using** *knights-path-non-nil* **by** *auto*

obtain k **where** *k-prems*:

$0 < k \wedge k < \text{length } ps_1 \ \text{last } (\text{take } k \ ps_1) = s_i \ \text{hd } (\text{drop } k \ ps_1) = s_j$
using *assms unfolding step-in-def* **by** *auto*
let $?ps = (\text{take } k \ ps_1) @ ?ps_2' @ (\text{drop } k \ ps_1)$
obtain $b_1 \ b_2$ **where** *b-prems*: *knights-path* $b_1 \ (\text{take } k \ ps_1) \ \text{knights-path } b_2 \ (\text{drop } k \ ps_1)$
 $b_1 \cup b_2 = ?b_1 \ b_1 \cap b_2 = \{\}$
using *assms* $\langle 0 < k \rangle \langle k < \text{length } ps_1 \rangle$ *knights-path-split* **by** *blast*

have $\text{hd } ?ps_2' = (i_h, \text{int } m_1 + j_h) \ \text{last } ?ps_2' = (i_l, \text{int } m_1 + j_l)$

using *assms knights-path-non-nil hd-trans-path last-trans-path* **by** *auto*

then have $\text{hd } ?ps_2' = (i_h, \text{int } m_1 + j_h) \ \text{last } ((\text{take } k \ ps_1) @ ?ps_2') = (i_l, \text{int } m_1 + j_l)$

using $\langle ?ps_2' \neq [] \rangle$ **by** *auto*

then have *vs*: *valid-step* $(\text{last } (\text{take } k \ ps_1)) \ (\text{hd } ?ps_2')$

valid-step $(\text{last } ((\text{take } k \ ps_1) @ ?ps_2')) \ (\text{hd } (\text{drop } k \ ps_1))$

using *assms k-prems* **by** *auto*

have $?b_1 \cap ?b' = \{\}$ **unfolding** *board-def trans-board-def* **by** *auto*

then have $b_1 \cap ?b' = \{\} \wedge (b_1 \cup ?b') \cap b_2 = \{\}$ **using** *b-prems* **by** *blast*

then have *inter-empty*: $b_1 \cap ?b' = \{\} \wedge (b_1 \cup ?b') \cap b_2 = \{\}$ **by** *auto*

have *knights-path* ($b_1 \cup ?b'$) ((*take* k ps_1) @ $?ps_2'$)
using *kp2'* *b-prems* *inter-empty* *vs* *knights-path-append* **by** *auto*
then have *knights-path* ($b_1 \cup ?b' \cup b_2$) $?ps$
using *b-prems* *inter-empty* *vs* *knights-path-append* [**where** $ps_1 = (\text{take } k \text{ } ps_1) @ ?ps_2'$] **by** *auto*
then have *knights-path* ($?b_1 \cup ?b'$) $?ps$
using *b-prems* *Un-commute* *Un-assoc* **by** *metis*
then have *kp*: *knights-path* (*board* n ($m_1 + m_2$)) $?ps$
using *board-concat* [*of* n m_1 m_2] **by** *auto*

have *hd*: *hd* $?ps = \text{hd } ps_1$
using *assms* $\langle 0 < k \rangle$ *knights-path-non-nil* *hd-take* **by** *auto*

have *last*: *last* $?ps = \text{last } ps_1$
using *assms* $\langle k < \text{length } ps_1 \rangle$ *knights-path-non-nil* *last-drop* **by** *auto*

have *m-simps*: $j + \text{int } m_1 = \text{int } m_1 + j$ $j' + \text{int } m_1 = \text{int } m_1 + j'$ **by** *auto*
have *si*: *step-in* $?ps$ ($i, \text{int } m_1 + j$) ($i', \text{int } m_1 + j'$)
using *assms* *step-in-append* [*OF* *step-in-prepend* [*OF* *trans-step-in*],
of ps_2 i j i' j' *take* k ps_1 0 *int* m_1 *drop* k ps_1]
by (*auto simp: m-simps*)

show *?thesis* **using** *kp* *hd* *last* *si* **by** *auto*
qed

lemma *len1-hd-last*: $\text{length } xs = 1 \implies \text{hd } xs = \text{last } xs$
by (*induction* xs) *auto*

Weaker version of $\llbracket \text{knights-path } (\text{board } ?n \text{ } ?m_1) \text{ } ?ps_1; \text{knights-path } (\text{board } ?n \text{ } ?m_2) \text{ } ?ps_2; \text{step-in } ?ps_1 \text{ } ?s_i \text{ } ?s_j; \text{hd } ?ps_2 = (?i_h, ?j_h); \text{last } ?ps_2 = (?i_l, ?j_l); \text{step-in } ?ps_2 \text{ } (?i, ?j) (?i', ?j'); \text{valid-step } ?s_i (?i_h, \text{int } ?m_1 + ?j_h); \text{valid-step } (?i_l, \text{int } ?m_1 + ?j_l) ?s_j \rrbracket \implies \exists ps. \text{knights-path } (\text{board } ?n \text{ } (?m_1 + ?m_2)) \text{ } ps \wedge \text{hd } ps = \text{hd } ?ps_1 \wedge \text{last } ps = \text{last } ?ps_1 \wedge \text{step-in } ps \text{ } (?i, \text{int } ?m_1 + ?j) (?i', \text{int } ?m_1 + ?j')$.

corollary *knights-path-split-concat*:

assumes *knights-path* (*board* n m_1) ps_1 *knights-path* (*board* n m_2) ps_2
 $\text{step-in } ps_1 \text{ } s_i \text{ } s_j \text{ } \text{hd } ps_2 = (i_h, j_h) \text{ } \text{last } ps_2 = (i_l, j_l)$
 $\text{valid-step } s_i (i_h, \text{int } m_1 + j_h) \text{ } \text{valid-step } (i_l, \text{int } m_1 + j_l) \text{ } s_j$
shows $\exists ps. \text{knights-path } (\text{board } n \text{ } (m_1 + m_2)) \text{ } ps \wedge \text{hd } ps = \text{hd } ps_1 \wedge \text{last } ps = \text{last } ps_1$

proof –

have $\text{length } ps_2 = 1 \vee \text{length } ps_2 > 1$
using *assms* *knights-path-non-nil* **by** (*meson* *length-0-conv* *less-one* *linorder-neqE-nat*)
then show *?thesis*
proof (*elim disjE*)
let $?s_k = (i_h, \text{int } m_1 + j_h)$
assume $\text{length } ps_2 = 1$

then have $(i_h, j_h) = (i_l, j_l)$

```

    using assms len1-hd-last by metis
  then have valid-step  $s_i$   $?s_k$  valid-step  $?s_k$   $s_j$  valid-step  $s_i$   $s_j$ 
    using assms step-in-valid-step by auto
  then show  $?thesis$ 
    using valid-step-non-transitive by blast
next
  assume length  $ps_2 > 1$ 
  then obtain  $i_1$   $j_1$   $i_2$   $j_2$   $ps_2'$  where  $ps_2 = (i_1, j_1) \# (i_2, j_2) \# ps_2'$ 
    using list-len-g-1-split by fastforce
  then have last (take 1  $ps_2$ ) =  $(i_1, j_1)$  hd (drop 1  $ps_2$ ) =  $(i_2, j_2)$  by auto
  then have step-in  $ps_2$   $(i_1, j_1)$   $(i_2, j_2)$  using  $\langle \text{length } ps_2 > 1 \rangle$  by (auto simp:
step-in-def)
  then show  $?thesis$ 
    using assms knights-path-split-concat-si-prev by blast
qed
qed

```

Concatenate two knight's path on a $n \times m$ -board along the 1st axis.

corollary *knights-path-split-concatT*:

```

  assumes knights-path (board  $n_1$   $m$ )  $ps_1$  knights-path (board  $n_2$   $m$ )  $ps_2$ 
    step-in  $ps_1$   $s_i$   $s_j$  hd  $ps_2 = (i_h, j_h)$  last  $ps_2 = (i_l, j_l)$ 
    valid-step  $s_i$   $(\text{int } n_1 + i_h, j_h)$  valid-step  $(\text{int } n_1 + i_l, j_l)$   $s_j$ 
  shows  $\exists ps. \text{knights-path (board } (n_1 + n_2) \text{ } m) \text{ } ps \wedge \text{hd } ps = \text{hd } ps_1 \wedge \text{last } ps =$ 
last  $ps_1$ 
    using assms
  proof -
    let  $?ps_1 T = \text{transpose } ps_1$ 
    let  $?ps_2 T = \text{transpose } ps_2$ 
    have kps: knights-path (board  $m$   $n_1$ )  $?ps_1 T$  knights-path (board  $m$   $n_2$ )  $?ps_2 T$ 
      using assms transpose-knights-path by auto

    let  $?s_i T = \text{transpose-square } s_i$ 
    let  $?s_j T = \text{transpose-square } s_j$ 
    have si: step-in  $?ps_1 T$   $?s_i T$   $?s_j T$ 
      using assms transpose-step-in by auto

    have  $ps_1 \neq []$   $ps_2 \neq []$ 
      using assms knights-path-non-nil by auto
    then have hd-last2: hd  $?ps_2 T = (j_h, i_h)$  last  $?ps_2 T = (j_l, i_l)$ 
      using assms hd-transpose last-transpose by (auto simp: transpose-square-def)

    have vs: valid-step  $?s_i T$   $(j_h, \text{int } n_1 + i_h)$  valid-step  $(j_l, \text{int } n_1 + i_l)$   $?s_j T$ 
      using assms transpose-valid-step by (auto simp: transpose-square-def split:
prod.splits)

    then obtain  $ps$  where
      ps-prems: knights-path (board  $m$   $(n_1 + n_2)$ )  $ps$  hd  $ps = \text{hd } ?ps_1 T$  last  $ps = \text{last } ?ps_1 T$ 
      using knights-path-split-concat[OF kps si hd-last2 vs] by auto
  
```

```

then have  $ps \neq []$  using knight's-path-non-nil by auto
let  $?psT = \text{transpose } ps$ 
have knight's-path (board ( $n_1 + n_2$ ) m)  $?psT \text{ hd } ?psT = \text{hd } ps_1 \text{ last } ?psT = \text{last}$ 
 $ps_1$ 
using  $\langle ps_1 \neq [] \rangle \langle ps \neq [] \rangle$  ps-prems transpose-knight's-path hd-transpose last-transpose

by (auto simp: transpose2)
then show ?thesis by auto
qed

```

Concatenate two Knight's path along the 2nd axis. There is a valid step from the last square in the first Knight's path ps_1 to the first square in the second Knight's path ps_2 .

corollary *knight's-path-concat*:

```

assumes knight's-path (board n  $m_1$ )  $ps_1$  knight's-path (board n  $m_2$ )  $ps_2$ 
 $\text{hd } ps_2 = (i_h, j_h) \text{ valid-step } (\text{last } ps_1) (i_h, \text{int } m_1 + j_h)$ 
shows knight's-path (board n ( $m_1 + m_2$ )) ( $ps_1 @ (\text{trans-path } (0, \text{int } m_1) ps_2)$ )
proof –
let  $?ps_2' = \text{trans-path } (0, \text{int } m_1) ps_2$ 
let  $?b = \text{trans-board } (0, \text{int } m_1) (\text{board } n m_2)$ 
have inter-empty:  $\text{board } n m_1 \cap ?b = \{\}$ 
unfolding board-def trans-board-def by auto
have  $\text{hd } ?ps_2' = (i_h, \text{int } m_1 + j_h)$ 
using assms knight's-path-non-nil hd-trans-path by auto
then have kp: knight's-path (board n  $m_1$ )  $ps_1$  knight's-path  $?b ?ps_2'$  and
 $vs$ : valid-step ( $\text{last } ps_1$ ) ( $\text{hd } ?ps_2'$ )
using assms trans-knight's-path by auto
then show knight's-path (board n ( $m_1 + m_2$ )) ( $ps_1 @ ?ps_2'$ )
using knight's-path-append[OF kp inter-empty vs] board-concat by auto
qed

```

Concatenate two Knight's path along the 2nd axis. The first Knight's path end in $(2, m_1 - 1)$ (lower-right) and the second Knight's paths start in $(1, 1)$ (lower-left).

corollary *knight's-path-lr-concat*:

```

assumes knight's-path (board n  $m_1$ )  $ps_1$  knight's-path (board n  $m_2$ )  $ps_2$ 
 $\text{last } ps_1 = (2, \text{int } m_1 - 1) \text{ hd } ps_2 = (1, 1)$ 
shows knight's-path (board n ( $m_1 + m_2$ )) ( $ps_1 @ (\text{trans-path } (0, \text{int } m_1) ps_2)$ )
proof –
have valid-step ( $\text{last } ps_1$ )  $(1, \text{int } m_1 + 1)$ 
using assms unfolding valid-step-def by auto
then show ?thesis
using assms knight's-path-concat by auto
qed

```

Concatenate two Knight's circuits along the 2nd axis. In the first Knight's path the squares $(2, m_1 - 1)$ and $(4, m_1)$ are adjacent and the second Knight's circuit starts in $(1, 1)$ (lower-left) and end in $(3, 2)$.

corollary *knights-circuit-lr-concat*:

```

assumes knights-circuit (board n m1) ps1 knights-circuit (board n m2) ps2
          step-in ps1 (2,int m1-1) (4,int m1)
          hd ps2 = (1,1) last ps2 = (3,2) step-in ps2 (2,int m2-1) (4,int m2)
shows  $\exists ps. \textit{knights-circuit} \text{ (board } n \text{ (} m_1+m_2 \text{)) } ps \wedge \textit{step-in} \text{ } ps \text{ (} 2, \textit{int} \text{ (} m_1+m_2 \text{)}-1 \text{)}$ 
          (4,int (m1+m2))
proof -
  have kp1: knights-path (board n m1) ps1 and kp2: knights-path (board n m2) ps2

  and vs: valid-step (last ps1) (hd ps1)
  using assms unfolding knights-circuit-def by auto

  have m-simps: int m1 + (int m2-1) = int (m1+m2)-1 int m1 + int m2 = int
    (m1+m2) by auto

  have valid-step (2,int m1-1) (1,int m1+1) valid-step (3,int m1+2) (4,int m1)
  unfolding valid-step-def by auto
  then obtain ps where knights-path (board n (m1+m2)) ps hd ps = hd ps1 last
    ps = last ps1 and
    si: step-in ps (2,int (m1+m2)-1) (4,int (m1+m2))
  using assms kp1 kp2
    knights-path-split-concat-si-prev[of n m1 ps1 m2 ps2 (2,int m1-1)
      (4,int m1) 1 1 3 2 2 int m2-1 4 int m2]
  by (auto simp only: m-simps)
  then have knights-circuit (board n (m1+m2)) ps
  using vs by (auto simp: knights-circuit-def)
  then show ?thesis
  using si by auto
qed

```

7 Parsing Paths

In this section functions are implemented to parse and construct paths. The parser converts the matrix representation ((*nat list*) *list*) used in [1] to a path (*path*).

for debugging

```

fun test-path :: path  $\Rightarrow$  bool where
  test-path (si#sj#xs) = (step-checker si sj  $\wedge$  test-path (sj#xs))
| test-path - = True

```

```

fun f-opt :: ('a  $\Rightarrow$  'a)  $\Rightarrow$  'a option  $\Rightarrow$  'a option where
  f-opt - None = None
| f-opt f (Some a) = Some (f a)

```

```

fun add-opt-fst-sq :: int  $\Rightarrow$  square option  $\Rightarrow$  square option where
  add-opt-fst-sq - None = None
| add-opt-fst-sq k (Some (i,j)) = Some (k+i,j)

```

```

fun find-k-in-col :: nat ⇒ nat list ⇒ int option where
  find-k-in-col k [] = None
| find-k-in-col k (c#cs) = (if c = k then Some 1 else f-opt ((+) 1) (find-k-in-col k
cs))

```

```

fun find-k-sqr :: nat ⇒ (nat list) list ⇒ square option where
  find-k-sqr k [] = None
| find-k-sqr k (r#rs) = (case find-k-in-col k r of
  None ⇒ f-opt (λ(i,j). (i+1,j)) (find-k-sqr k rs)
| Some j ⇒ Some (1,j))

```

Auxiliary function to easily parse pre-computed boards from paper.

```

fun to-sqrs :: nat ⇒ (nat list) list ⇒ path option where
  to-sqrs 0 rs = Some []
| to-sqrs k rs = (case find-k-sqr k rs of
  None ⇒ None
| Some si ⇒ f-opt (λps. ps@[si]) (to-sqrs (k-1) rs))

```

```

fun num-elems :: (nat list) list ⇒ nat where
  num-elems (r#rs) = length r * length (r#rs)

```

Convert a matrix $(nat\ list\ list)$ to a path $(path)$. With this function we implicitly define the lower-left corner to be $(1,1)$ and the upper-right corner to be (n,m) .

definition $to-path\ rs \equiv to-sqrs\ (num-elems\ rs)\ (rev\ rs)$

Example

```

value to-path
  [[3,22,13,16,5],
  [12,17,4,21,14],
  [23,2,15,6,9],
  [18,11,8,25,20],
  [1,24,19,10,7::nat]]

```

8 Knight's Paths for $5 \times m$ -Boards

Given here are knight's paths, $kp5xmlr$ and $kp5xmur$, for the $(5 \times m)$ -board that start in the lower-left corner for $m \in \{5,6,7,8,9\}$. The path $kp5xmlr$ ends in the lower-right corner, whereas the path $kp5xmur$ ends in the upper-right corner. The tables show the visited squares numbered in ascending order.

abbreviation $b5x5 \equiv board\ 5\ 5$

A Knight's path for the (5×5) -board that starts in the lower-left and ends in the lower-right.

3	22	13	16	5
12	17	4	21	14
23	2	15	6	9
18	11	8	25	20
1	24	19	10	7

abbreviation $kp5x5lr \equiv$ the (to-path

[[3,22,13,16,5],
 [12,17,4,21,14],
 [23,2,15,6,9],
 [18,11,8,25,20],
 [1,24,19,10,7]])

lemma $kp-5x5-lr$: knights-path b5x5 $kp5x5lr$

by (simp only: knights-path-exec-simp) eval

lemma $kp-5x5-lr-hd$: hd $kp5x5lr = (1,1)$ **by** eval

lemma $kp-5x5-lr-last$: last $kp5x5lr = (2,4)$ **by** eval

lemma $kp-5x5-lr-non-nil$: $kp5x5lr \neq []$ **by** eval

A Knight's path for the (5×5) -board that starts in the lower-left and ends in the upper-right.

7	12	15	20	5
16	21	6	25	14
11	8	13	4	19
22	17	2	9	24
1	10	23	18	3

abbreviation $kp5x5ur \equiv$ the (to-path

[[7,12,15,20,5],
 [16,21,6,25,14],
 [11,8,13,4,19],
 [22,17,2,9,24],
 [1,10,23,18,3]])

lemma $kp-5x5-ur$: knights-path b5x5 $kp5x5ur$

by (simp only: knights-path-exec-simp) eval

lemma $kp-5x5-ur-hd$: hd $kp5x5ur = (1,1)$ **by** eval

lemma $kp-5x5-ur-last$: last $kp5x5ur = (4,4)$ **by** eval

lemma $kp-5x5-ur-non-nil$: $kp5x5ur \neq []$ **by** eval

abbreviation $b5x6 \equiv$ board 5 6

A Knight's path for the (5×6) -board that starts in the lower-left and ends in the lower-right.

7	14	21	28	5	12
22	27	6	13	20	29
15	8	17	24	11	4
26	23	2	9	30	19
1	16	25	18	3	10

abbreviation $kp5x6lr \equiv$ the (to-path

$[[7,14,21,28,5,12],$
 $[22,27,6,13,20,29],$
 $[15,8,17,24,11,4],$
 $[26,23,2,9,30,19],$
 $[1,16,25,18,3,10]]$)

lemma $kp\text{-}5x6\text{-}lr$: *knight's-path b5x6 kp5x6lr*

by (*simp only: knights-path-exec-simp*) *eval*

lemma $kp\text{-}5x6\text{-}lr\text{-}hd$: $hd\ kp5x6lr = (1,1)$ **by** *eval*

lemma $kp\text{-}5x6\text{-}lr\text{-}last$: $last\ kp5x6lr = (2,5)$ **by** *eval*

lemma $kp\text{-}5x6\text{-}lr\text{-}non\text{-}nil$: $kp5x6lr \neq []$ **by** *eval*

A Knight's path for the (5×6) -board that starts in the lower-left and ends in the upper-right.

3	10	29	20	5	12
28	19	4	11	30	21
9	2	17	24	13	6
18	27	8	15	22	25
1	16	23	26	7	14

abbreviation $kp5x6ur \equiv$ the (to-path

$[[3,10,29,20,5,12],$
 $[28,19,4,11,30,21],$
 $[9,2,17,24,13,6],$
 $[18,27,8,15,22,25],$
 $[1,16,23,26,7,14]]$)

lemma $kp\text{-}5x6\text{-}ur$: *knight's-path b5x6 kp5x6ur*

by (*simp only: knights-path-exec-simp*) *eval*

lemma $kp\text{-}5x6\text{-}ur\text{-}hd$: $hd\ kp5x6ur = (1,1)$ **by** *eval*

lemma $kp\text{-}5x6\text{-}ur\text{-}last$: $last\ kp5x6ur = (4,5)$ **by** *eval*

lemma $kp\text{-}5x6\text{-}ur\text{-}non\text{-}nil$: $kp5x6ur \neq []$ **by** *eval*

abbreviation $b5x7 \equiv \text{board } 5 \ 7$

A Knight's path for the (5×7) -board that starts in the lower-left and ends in the lower-right.

3	12	21	30	5	14	23
20	29	4	13	22	31	6
11	2	19	32	7	24	15
28	33	10	17	26	35	8
1	18	27	34	9	16	25

abbreviation $kp5x7lr \equiv \text{the (to-path}$

$[[3,12,21,30,5,14,23],$
 $[20,29,4,13,22,31,6],$
 $[11,2,19,32,7,24,15],$
 $[28,33,10,17,26,35,8],$
 $[1,18,27,34,9,16,25]])$

lemma $kp\text{-}5x7\text{-}lr$: *knight's-path* $b5x7$ $kp5x7lr$

by (*simp only: knight's-path-exec-simp*) *eval*

lemma $kp\text{-}5x7\text{-}lr\text{-}hd$: $hd\ kp5x7lr = (1,1)$ **by** *eval*

lemma $kp\text{-}5x7\text{-}lr\text{-}last$: $last\ kp5x7lr = (2,6)$ **by** *eval*

lemma $kp\text{-}5x7\text{-}lr\text{-}non\text{-}nil$: $kp5x7lr \neq []$ **by** *eval*

A Knight's path for the (5×7) -board that starts in the lower-left and ends in the upper-right.

3	32	11	34	5	26	13
10	19	4	25	12	35	6
31	2	33	20	23	14	27
18	9	24	29	16	7	22
1	30	17	8	21	28	15

abbreviation $kp5x7ur \equiv \text{the (to-path}$

$[3,32,11,34,5,26,13],$
 $[10,19,4,25,12,35,6],$
 $[31,2,33,20,23,14,27],$
 $[18,9,24,29,16,7,22],$
 $[1,30,17,8,21,28,15]])$

lemma $kp\text{-}5x7\text{-}ur$: *knight's-path* $b5x7$ $kp5x7ur$

by (*simp only: knight's-path-exec-simp*) *eval*

lemma $kp\text{-}5x7\text{-}ur\text{-}hd$: $hd\ kp5x7ur = (1,1)$ **by** *eval*

lemma *kp-5x7-ur-last*: *last kp5x7ur = (4,6)* **by** *eval*

lemma *kp-5x7-ur-non-nil*: *kp5x7ur ≠ []* **by** *eval*

abbreviation *b5x8* \equiv *board 5 8*

A Knight's path for the (5×8) -board that starts in the lower-left and ends in the lower-right.

3	12	37	26	5	14	17	28
34	23	4	13	36	27	6	15
11	2	35	38	25	16	29	18
22	33	24	9	20	31	40	7
1	10	21	32	39	8	19	30

abbreviation *kp5x8lr* \equiv *the (to-path*

[[3,12,37,26,5,14,17,28],
[34,23,4,13,36,27,6,15],
[11,2,35,38,25,16,29,18],
[22,33,24,9,20,31,40,7],
[1,10,21,32,39,8,19,30]])

lemma *kp-5x8-lr*: *knight's-path b5x8 kp5x8lr*
by *(simp only: knights-path-exec-simp)* *eval*

lemma *kp-5x8-lr-hd*: *hd kp5x8lr = (1,1)* **by** *eval*

lemma *kp-5x8-lr-last*: *last kp5x8lr = (2,7)* **by** *eval*

lemma *kp-5x8-lr-non-nil*: *kp5x8lr ≠ []* **by** *eval*

A Knight's path for the (5×8) -board that starts in the lower-left and ends in the upper-right.

33	8	17	38	35	6	15	24
18	37	34	7	16	25	40	5
9	32	29	36	39	14	23	26
30	19	2	11	28	21	4	13
1	10	31	20	3	12	27	22

abbreviation *kp5x8ur* \equiv *the (to-path*

[[33,8,17,38,35,6,15,24],
[18,37,34,7,16,25,40,5],
[9,32,29,36,39,14,23,26],
[30,19,2,11,28,21,4,13],
[1,10,31,20,3,12,27,22]])

lemma *kp-5x8-ur*: *knight's-path b5x8 kp5x8ur*
by *(simp only: knights-path-exec-simp)* *eval*

lemma *kp-5x8-ur-hd*: *hd kp5x8ur = (1,1) by eval*

lemma *kp-5x8-ur-last*: *last kp5x8ur = (4,7) by eval*

lemma *kp-5x8-ur-non-nil*: *kp5x8ur ≠ [] by eval*

abbreviation *b5x9* \equiv *board 5 9*

A Knight's path for the (5×9) -board that starts in the lower-left and ends in the lower-right.

9	4	11	16	23	42	33	36	25
12	17	8	3	32	37	24	41	34
5	10	15	20	43	22	35	26	29
18	13	2	7	38	31	28	45	40
1	6	19	14	21	44	39	30	27

abbreviation *kp5x9lr* \equiv *the (to-path*

[[9,4,11,16,23,42,33,36,25],
[12,17,8,3,32,37,24,41,34],
[5,10,15,20,43,22,35,26,29],
[18,13,2,7,38,31,28,45,40],
[1,6,19,14,21,44,39,30,27]]]

lemma *kp-5x9-lr*: *knight's-path b5x9 kp5x9lr*

by *(simp only: knight's-path-exec-simp) eval*

lemma *kp-5x9-lr-hd*: *hd kp5x9lr = (1,1) by eval*

lemma *kp-5x9-lr-last*: *last kp5x9lr = (2,8) by eval*

lemma *kp-5x9-lr-non-nil*: *kp5x9lr ≠ [] by eval*

A Knight's path for the (5×9) -board that starts in the lower-left and ends in the upper-right.

9	4	11	16	27	32	35	40	25
12	17	8	3	36	41	26	45	34
5	10	15	20	31	28	33	24	39
18	13	2	7	42	37	22	29	44
1	6	19	14	21	30	43	38	23

abbreviation *kp5x9ur* \equiv *the (to-path*

[[9,4,11,16,27,32,35,40,25],
[12,17,8,3,36,41,26,45,34],
[5,10,15,20,31,28,33,24,39],
[18,13,2,7,42,37,22,29,44],

$[1,6,19,14,21,30,43,38,23]]$
lemma *kp-5x9-ur: knights-path b5x9 kp5x9ur*
by (*simp only: knights-path-exec-simp*) *eval*
lemma *kp-5x9-ur-hd: hd kp5x9ur = (1,1)* **by** *eval*
lemma *kp-5x9-ur-last: last kp5x9ur = (4,8)* **by** *eval*
lemma *kp-5x9-ur-non-nil: kp5x9ur ≠ []* **by** *eval*
lemmas *kp-5xm-lr =*
kp-5x5-lr kp-5x5-lr-hd kp-5x5-lr-last kp-5x5-lr-non-nil
kp-5x6-lr kp-5x6-lr-hd kp-5x6-lr-last kp-5x6-lr-non-nil
kp-5x7-lr kp-5x7-lr-hd kp-5x7-lr-last kp-5x7-lr-non-nil
kp-5x8-lr kp-5x8-lr-hd kp-5x8-lr-last kp-5x8-lr-non-nil
kp-5x9-lr kp-5x9-lr-hd kp-5x9-lr-last kp-5x9-lr-non-nil

lemmas *kp-5xm-ur =*
kp-5x5-ur kp-5x5-ur-hd kp-5x5-ur-last kp-5x5-ur-non-nil
kp-5x6-ur kp-5x6-ur-hd kp-5x6-ur-last kp-5x6-ur-non-nil
kp-5x7-ur kp-5x7-ur-hd kp-5x7-ur-last kp-5x7-ur-non-nil
kp-5x8-ur kp-5x8-ur-hd kp-5x8-ur-last kp-5x8-ur-non-nil
kp-5x9-ur kp-5x9-ur-hd kp-5x9-ur-last kp-5x9-ur-non-nil

For every $5 \times m$ -board with $m \geq 5$ there exists a knight's path that starts in $(1,1)$ (bottom-left) and ends in $(2,m-1)$ (bottom-right).

lemma *knights-path-5xm-lr-exists:*
assumes $m \geq 5$
shows $\exists ps. \text{knights-path (board 5 } m) ps \wedge \text{hd } ps = (1,1) \wedge \text{last } ps = (2, \text{int } m-1)$
using *assms*
proof (*induction m rule: less-induct*)
case (*less m*)
then have $m \in \{5,6,7,8,9\} \vee 5 \leq m-5$ **by** *auto*
then show *?case*
proof (*elim disjE*)
assume $m \in \{5,6,7,8,9\}$
then show *?thesis* **using** *kp-5xm-lr* **by** *fastforce*
next
assume $m-ge: 5 \leq m-5$
then obtain ps_1 **where** $ps_1\text{-IH: knights-path (board 5 (m-5)) } ps_1 \text{ hd } ps_1 =$
 $(1,1)$
 $\text{last } ps_1 = (2, \text{int } (m-5)-1) \text{ } ps_1 \neq []$
using *less.IH[of m-5] knights-path-non-nil* **by** *auto*
let $?ps_2 = kp5x5lr$
let $?ps_2' = ps_1 @ \text{trans-path (0, int (m-5)) } ?ps_2$
have *knights-path b5x5 ?ps_2* **hd** $?ps_2 = (1, 1)$ $?ps_2 \neq []$ **last** $?ps_2 = (2,4)$
using *kp-5xm-lr* **by** *auto*
then have $1: \text{knights-path (board 5 } m) ?ps_2'$

```

    using m-ge ps1-IH knights-path-lr-concat[of 5 m-5 ps1 5 ?ps2] by auto

    have 2: hd ?ps2' = (1,1) using ps1-IH by auto

    have last (trans-path (0,int (m-5)) ?ps2) = (2,int m-1)
      using m-ge last-trans-path[OF ‹?ps2 ≠ []› ‹last ?ps2 = (2,4)›] by auto
    then have 3: last ?ps2' = (2,int m-1)
      using last-appendR[OF trans-path-non-nil[OF ‹?ps2 ≠ []›],symmetric] by
metis

    show ?thesis using 1 2 3 by auto
  qed
qed

For every 5×m-board with  $m \geq 5$  there exists a knight's path that starts
in (1,1) (bottom-left) and ends in (4,m-1) (top-right).

lemma knights-path-5xm-ur-exists:
  assumes m ≥ 5
  shows ∃ ps. knights-path (board 5 m) ps ∧ hd ps = (1,1) ∧ last ps = (4,int m-1)
  using assms
proof -
  have m ∈ {5,6,7,8,9} ∨ 5 ≤ m-5 using assms by auto
  then show ?thesis
  proof (elim disjE)
    assume m ∈ {5,6,7,8,9}
    then show ?thesis using kp-5xm-ur by fastforce
  next
    assume m-ge: 5 ≤ m-5
    then obtain ps1 where ps-prems: knights-path (board 5 (m-5)) ps1 hd ps1 =
(1,1)
      last ps1 = (2,int (m-5)-1) ps1 ≠ []
      using knights-path-5xm-lr-exists[of (m-5)] knights-path-non-nil by auto
    let ?ps2=kp5x5ur
    let ?ps'=ps1 @ trans-path (0,int (m-5)) ?ps2
    have knights-path b5x5 ?ps2 hd ?ps2 = (1, 1) ?ps2 ≠ []
      last ?ps2 = (4,4)
      using kp-5xm-ur by auto
    then have 1: knights-path (board 5 m) ?ps'
      using m-ge ps-prems knights-path-lr-concat[of 5 m-5 ps1 5 ?ps2] by auto

    have 2: hd ?ps' = (1,1) using ps-prems by auto

    have last (trans-path (0,int (m-5)) ?ps2) = (4,int m-1)
      using m-ge last-trans-path[OF ‹?ps2 ≠ []› ‹last ?ps2 = (4,4)›] by auto
    then have 3: last ?ps' = (4,int m-1)
      using last-appendR[OF trans-path-non-nil[OF ‹?ps2 ≠ []›],symmetric] by
metis

    show ?thesis using 1 2 3 by auto
  end
end

```

qed
qed

$5 \leq ?m \implies \exists ps. \text{knights-path } (\text{board } 5 \ ?m) \ ps \wedge \text{hd } ps = (1, 1) \wedge \text{last } ps = (2, \text{int } ?m - 1)$ and $5 \leq ?m \implies \exists ps. \text{knights-path } (\text{board } 5 \ ?m) \ ps \wedge \text{hd } ps = (1, 1) \wedge \text{last } ps = (2, \text{int } ?m - 1)$ formalize Lemma 1 from [1].

lemmas *knights-path-5xm-exists = knights-path-5xm-lr-exists knights-path-5xm-ur-exists*

9 Knight's Paths and Circuits for $6 \times m$ -Boards

abbreviation *b6x5* \equiv *board 6 5*

A Knight's path for the (6×5) -board that starts in the lower-left and ends in the upper-left.

10	19	4	29	12
3	30	11	20	5
18	9	24	13	28
25	2	17	6	21
16	23	8	27	14
1	26	15	22	7

abbreviation *kp6x5ul* \equiv *the (to-path*

[[10,19,4,29,12],
[3,30,11,20,5],
[18,9,24,13,28],
[25,2,17,6,21],
[16,23,8,27,14],
[1,26,15,22,7]])

lemma *kp-6x5-ul: knights-path b6x5 kp6x5ul*
by *(simp only: knights-path-exec-simp) eval*

lemma *kp-6x5-ul-hd: hd kp6x5ul = (1,1) by eval*

lemma *kp-6x5-ul-last: last kp6x5ul = (5,2) by eval*

lemma *kp-6x5-ul-non-nil: kp6x5ul $\neq []$ by eval*

A Knight's circuit for the (6×5) -board.

16	9	6	27	18
7	26	17	14	5
10	15	8	19	28
25	30	23	4	13
22	11	2	29	20
1	24	21	12	3

abbreviation $kc6x5 \equiv$ the (to-path

[[16,9,6,27,18],
 [7,26,17,14,5],
 [10,15,8,19,28],
 [25,30,23,4,13],
 [22,11,2,29,20],
 [1,24,21,12,3]])

lemma $kc-6x5$: knights-circuit $b6x5$ $kc6x5$
 by (simp only: knights-circuit-exec-simp) eval

lemma $kc-6x5$ -hd: hd $kc6x5 = (1,1)$ by eval

lemma $kc-6x5$ -non-nil: $kc6x5 \neq []$ by eval

abbreviation $b6x6 \equiv$ board 6 6

The path given for the 6×6 -board that ends in the upper-left is wrong. The Knight cannot move from square 26 to square 27.

14	23	6	28	12	21
7	36	13	22	5	27
24	15	29	35	20	11
30	8	17	26	34	4
16	25	2	32	10	19
1	31	9	18	3	33

abbreviation $kp6x6ul$ -false \equiv the (to-path

[[14,23,6,28,12,21],
 [7,36,13,22,5,27],
 [24,15,29,35,20,11],
 [30,8,17,26,34,4],
 [16,25,2,32,10,19],
 [1,31,9,18,3,33]])

lemma \neg knights-path $b6x6$ $kp6x6ul$ -false
 by (simp only: knights-path-exec-simp) eval

I have computed a correct Knight's path for the 6×6 -board that ends in the upper-left. A Knight's path for the (6×6) -board that starts in the lower-left and ends in the upper-left.

8	25	10	21	6	23
11	36	7	24	33	20
26	9	34	3	22	5
35	12	15	30	19	32
14	27	2	17	4	29
1	16	13	28	31	18

abbreviation $kp6x6ul \equiv$ the (to-path

$[[8,25,10,21,6,23],$
 $[11,36,7,24,33,20],$
 $[26,9,34,3,22,5],$
 $[35,12,15,30,19,32],$
 $[14,27,2,17,4,29],$
 $[1,16,13,28,31,18]])$

lemma $kp-6x6-ul$: knights-path $b6x6$ $kp6x6ul$

by (simp only: knights-path-exec-simp) eval

lemma $kp-6x6-ul-hd$: $hd\ kp6x6ul = (1,1)$ **by** eval

lemma $kp-6x6-ul-last$: $last\ kp6x6ul = (5,2)$ **by** eval

lemma $kp-6x6-ul-non-nil$: $kp6x6ul \neq []$ **by** eval

A Knight's circuit for the (6×6) -board.

4	25	34	15	18	7
35	14	5	8	33	16
24	3	26	17	6	19
13	36	23	30	9	32
22	27	2	11	20	29
1	12	21	28	31	10

abbreviation $kc6x6 \equiv$ the (to-path

$[[4,25,34,15,18,7],$
 $[35,14,5,8,33,16],$
 $[24,3,26,17,6,19],$
 $[13,36,23,30,9,32],$
 $[22,27,2,11,20,29],$
 $[1,12,21,28,31,10]])$

lemma $kc-6x6$: knights-circuit $b6x6$ $kc6x6$

by (simp only: knights-circuit-exec-simp) eval

lemma $kc-6x6-hd$: $hd\ kc6x6 = (1,1)$ **by** eval

lemma $kc-6x6-non-nil$: $kc6x6 \neq []$ **by** eval

abbreviation $b6x7 \equiv$ board 6 7

A Knight's path for the (6×7) -board that starts in the lower-left and ends in the upper-left.

18	23	8	39	16	25	6
9	42	17	24	7	40	15
22	19	32	41	38	5	26
33	10	21	28	31	14	37
20	29	2	35	12	27	4
1	34	11	30	3	36	13

abbreviation $kp6x7ul \equiv$ the (to-path

[[18,23,8,39,16,25,6],
 [9,42,17,24,7,40,15],
 [22,19,32,41,38,5,26],
 [33,10,21,28,31,14,37],
 [20,29,2,35,12,27,4],
 [1,34,11,30,3,36,13]])

lemma $kp\text{-}6x7\text{-}ul$: *knight's-path* $b6x7\text{ }kp6x7ul$
by (*simp only: knight's-path-exec-simp*) *eval*

lemma $kp\text{-}6x7\text{-}ul\text{-}hd$: $hd\text{ }kp6x7ul = (1,1)$ **by** *eval*

lemma $kp\text{-}6x7\text{-}ul\text{-}last$: $last\text{ }kp6x7ul = (5,2)$ **by** *eval*

lemma $kp\text{-}6x7\text{-}ul\text{-}non\text{-}nil$: $kp6x7ul \neq []$ **by** *eval*

A Knight's circuit for the (6×7) -board.

26	37	8	17	28	31	6
9	18	27	36	7	16	29
38	25	10	19	30	5	32
11	42	23	40	35	20	15
24	39	2	13	22	33	4
1	12	41	34	3	14	21

abbreviation $kc6x7 \equiv$ the (to-path

[[26,37,8,17,28,31,6],
 [9,18,27,36,7,16,29],
 [38,25,10,19,30,5,32],
 [11,42,23,40,35,20,15],
 [24,39,2,13,22,33,4],
 [1,12,41,34,3,14,21]])

lemma $kc\text{-}6x7$: *knight's-circuit* $b6x7\text{ }kc6x7$
by (*simp only: knight's-circuit-exec-simp*) *eval*

lemma $kc\text{-}6x7\text{-}hd$: $hd\text{ }kc6x7 = (1,1)$ **by** *eval*

lemma $kc\text{-}6x7\text{-}non\text{-}nil$: $kc6x7 \neq []$ **by** *eval*

abbreviation $b6x8 \equiv$ board 6 8

A Knight's path for the (6×8) -board that starts in the lower-left and ends in the upper-left.

18	31	8	35	16	33	6	45
9	48	17	32	7	46	15	26
30	19	36	47	34	27	44	5
37	10	21	28	43	40	25	14
20	29	2	39	12	23	4	41
1	38	11	22	3	42	13	24

abbreviation $kp6x8ul \equiv$ the (to-path

$[[18,31,8,35,16,33,6,45],$
 $[9,48,17,32,7,46,15,26],$
 $[30,19,36,47,34,27,44,5],$
 $[37,10,21,28,43,40,25,14],$
 $[20,29,2,39,12,23,4,41],$
 $[1,38,11,22,3,42,13,24]]$)

lemma $kp\text{-}6x8\text{-}ul$: *knight's-path* $b6x8$ $kp6x8ul$

by (*simp only: knights-path-exec-simp*) *eval*

lemma $kp\text{-}6x8\text{-}ul\text{-}hd$: $hd\ kp6x8ul = (1,1)$ **by** *eval*

lemma $kp\text{-}6x8\text{-}ul\text{-}last$: $last\ kp6x8ul = (5,2)$ **by** *eval*

lemma $kp\text{-}6x8\text{-}ul\text{-}non\text{-}nil$: $kp6x8ul \neq []$ **by** *eval*

A Knight's circuit for the (6×8) -board.

30	35	8	15	28	39	6	13
9	16	29	36	7	14	27	38
34	31	10	23	40	37	12	5
17	48	33	46	11	22	41	26
32	45	2	19	24	43	4	21
1	18	47	44	3	20	25	42

abbreviation $kc6x8 \equiv$ the (to-path

$[[30,35,8,15,28,39,6,13],$
 $[9,16,29,36,7,14,27,38],$
 $[34,31,10,23,40,37,12,5],$
 $[17,48,33,46,11,22,41,26],$
 $[32,45,2,19,24,43,4,21],$
 $[1,18,47,44,3,20,25,42]]$)

lemma $kc\text{-}6x8$: *knight's-circuit* $b6x8$ $kc6x8$

by (*simp only: knights-circuit-exec-simp*) *eval*

lemma $kc\text{-}6x8\text{-}hd$: $hd\ kc6x8 = (1,1)$ **by** *eval*

lemma *kc-6x8-non-nil*: $kc6x8 \neq []$ **by** *eval*

abbreviation *b6x9* \equiv *board 6 9*

A Knight's path for the (6×9) -board that starts in the lower-left and ends in the upper-left.

22	45	10	53	20	47	8	35	18
11	54	21	46	9	36	19	48	7
44	23	42	37	52	49	32	17	34
41	12	25	50	27	38	29	6	31
24	43	2	39	14	51	4	33	16
1	40	13	26	3	28	15	30	5

abbreviation *kp6x9ul* \equiv *the (to-path*

[[22,45,10,53,20,47,8,35,18],
[11,54,21,46,9,36,19,48,7],
[44,23,42,37,52,49,32,17,34],
[41,12,25,50,27,38,29,6,31],
[24,43,2,39,14,51,4,33,16],
[1,40,13,26,3,28,15,30,5]])

lemma *kp-6x9-ul*: *knight's-path b6x9 kp6x9ul*
by (*simp only: knight's-path-exec-simp*) *eval*

lemma *kp-6x9-ul-hd*: *hd kp6x9ul = (1,1)* **by** *eval*

lemma *kp-6x9-ul-last*: *last kp6x9ul = (5,2)* **by** *eval*

lemma *kp-6x9-ul-non-nil*: $kp6x9ul \neq []$ **by** *eval*

A Knight's circuit for the (6×9) -board.

14	49	4	51	24	39	6	29	22
3	52	13	40	5	32	23	42	7
48	15	50	25	38	41	28	21	30
53	2	37	12	33	26	31	8	43
16	47	54	35	18	45	10	27	20
1	36	17	46	11	34	19	44	9

abbreviation *kc6x9* \equiv *the (to-path*

[[14,49,4,51,24,39,6,29,22],
[3,52,13,40,5,32,23,42,7],
[48,15,50,25,38,41,28,21,30],
[53,2,37,12,33,26,31,8,43],
[16,47,54,35,18,45,10,27,20],
[1,36,17,46,11,34,19,44,9]])

lemma *kc-6x9*: *knight's-circuit b6x9 kc6x9*

by (*simp only: knights-circuit-exec-simp*) *eval*

lemma *kc-6x9-hd*: *hd kc6x9 = (1,1)* **by** *eval*

lemma *kc-6x9-non-nil*: *kc6x9 ≠ []* **by** *eval*

lemmas *kp-6xm-ul =*

kp-6x5-ul kp-6x5-ul-hd kp-6x5-ul-last kp-6x5-ul-non-nil
kp-6x6-ul kp-6x6-ul-hd kp-6x6-ul-last kp-6x6-ul-non-nil
kp-6x7-ul kp-6x7-ul-hd kp-6x7-ul-last kp-6x7-ul-non-nil
kp-6x8-ul kp-6x8-ul-hd kp-6x8-ul-last kp-6x8-ul-non-nil
kp-6x9-ul kp-6x9-ul-hd kp-6x9-ul-last kp-6x9-ul-non-nil

lemmas *kc-6xm =*

kc-6x5 kc-6x5-hd kc-6x5-non-nil
kc-6x6 kc-6x6-hd kc-6x6-non-nil
kc-6x7 kc-6x7-hd kc-6x7-non-nil
kc-6x8 kc-6x8-hd kc-6x8-non-nil
kc-6x9 kc-6x9-hd kc-6x9-non-nil

For every $6 \times m$ -board with $m \geq 5$ there exists a knight's path that starts in $(1,1)$ (bottom-left) and ends in $(5,2)$ (top-left).

lemma *knights-path-6xm-ul-exists*:

assumes $m \geq 5$

shows $\exists ps. \text{knights-path (board } 6 \text{ } m) \text{ } ps \wedge \text{hd } ps = (1,1) \wedge \text{last } ps = (5,2)$

using *assms*

proof (*induction m rule: less-induct*)

case (*less m*)

then have $m \in \{5,6,7,8,9\} \vee 5 \leq m-5$ **by** *auto*

then show *?case*

proof (*elim disjE*)

assume $m \in \{5,6,7,8,9\}$

then show *?thesis* **using** *kp-6xm-ul* **by** *fastforce*

next

let *?ps₁=kp6x5ul*

let *?b₁=board 6 5*

have *ps₁-prems*: *knights-path ?b₁ ?ps₁ hd ?ps₁ = (1,1) last ?ps₁ = (5,2)*

using *kp-6xm-ul* **by** *auto*

assume *m-ge*: $5 \leq m-5$

then obtain *ps₂* **where** *ps₂-IH*: *knights-path (board 6 (m-5)) ps₂ hd ps₂ = (1,1)*

last ps₂ = (5,2)

using *less.IH[of m-5]* *knights-path-non-nil* **by** *auto*

have $27 < \text{length } ?ps_1 \text{ last (take } 27 \text{ } ?ps_1) = (2,4) \text{ hd (drop } 27 \text{ } ?ps_1) = (4,5)$

by *eval+*

then have *step-in ?ps₁ (2,4) (4,5)*

unfolding *step-in-def* **using** *zero-less-numeral* **by** *blast*

then have *step-in ?ps₁ (2,4) (4,5)*

```

      valid-step (2,4) (1,int 5+1)
      valid-step (5,int 5+2) (4,5)
    unfolding valid-step-def by auto
  then show ?thesis
    using ⟨5 ≤ m-5⟩ ps1-prems ps2-IH knights-path-split-concat[of 6 5 ?ps1 m-5
ps2] by auto
  qed
qed

```

For every $6 \times m$ -board with $m \geq 5$ there exists a knight's circuit.

lemma *knights-circuit-6xm-exists*:

```

  assumes m ≥ 5
  shows ∃ ps. knights-circuit (board 6 m) ps
  using assms
proof -
  have m ∈ {5,6,7,8,9} ∨ 5 ≤ m-5 using assms by auto
  then show ?thesis
  proof (elim disjE)
    assume m ∈ {5,6,7,8,9}
    then show ?thesis using kc-6xm by fastforce
  next
    let ?ps1 = rev kc6x5
    have knights-circuit b6x5 ?ps1 last ?ps1 = (1,1)
      using kc-6xm knights-circuit-rev by (auto simp: last-rev)
    then have ps1-prems: knights-path b6x5 ?ps1 valid-step (last ?ps1) (hd ?ps1)
      unfolding knights-circuit-def using valid-step-rev by auto
    assume m-ge: 5 ≤ m-5
    then obtain ps2 where ps2-prems: knights-path (board 6 (m-5)) ps2 hd ps2
      = (1,1)
      last ps2 = (5,2)
      using knights-path-6xm-ul-exists[of (m-5)] knights-path-non-nil by auto

    have 2 < length ?ps1 last (take 2 ?ps1) = (2,4) hd (drop 2 ?ps1) = (4,5) by
eval+
    then have step-in ?ps1 (2,4) (4,5)
      unfolding step-in-def using zero-less-numeral by blast
    then have step-in ?ps1 (2,4) (4,5)
      valid-step (2,4) (1,int 5+1)
      valid-step (5,int 5+2) (4,5)
      unfolding valid-step-def by auto
    then have ∃ ps. knights-path (board 6 m) ps ∧ hd ps = hd ?ps1 ∧ last ps = last
      ?ps1
      using m-ge ps1-prems ps2-prems knights-path-split-concat[of 6 5 ?ps1 m-5
ps2] by auto
    then show ?thesis using ps1-prems by (auto simp: knights-circuit-def)
  qed
qed

```

$5 \leq ?m \implies \exists ps. \text{knights-path (board 6 ?m) } ps \wedge \text{hd } ps = (1, 1) \wedge \text{last } ps$

$= (5, 2)$ and $5 \leq ?m \implies \exists ps. \text{ knights-circuit } (\text{board } 6 \ ?m) \ ps$ formalize Lemma 2 from [1].

lemmas *knights-path-6xm-exists = knights-path-6xm-ul-exists knights-circuit-6xm-exists*

10 Knight's Paths and Circuits for $8 \times m$ -Boards

abbreviation *b8x5* \equiv *board 8 5*

A Knight's path for the (8×5) -board that starts in the lower-left and ends in the upper-left.

28	7	22	39	26
23	40	27	6	21
8	29	38	25	14
37	24	15	20	5
16	9	30	13	34
31	36	33	4	19
10	17	2	35	12
1	32	11	18	3

abbreviation *kp8x5ul* \equiv *the (to-path*

[[28,7,22,39,26],
[23,40,27,6,21],
[8,29,38,25,14],
[37,24,15,20,5],
[16,9,30,13,34],
[31,36,33,4,19],
[10,17,2,35,12],
[1,32,11,18,3]])

lemma *kp-8x5-ul: knights-path b8x5 kp8x5ul*

by (*simp only: knights-path-exec-simp*) *eval*

lemma *kp-8x5-ul-hd: hd kp8x5ul = (1,1)* **by** *eval*

lemma *kp-8x5-ul-last: last kp8x5ul = (7,2)* **by** *eval*

lemma *kp-8x5-ul-non-nil: kp8x5ul $\neq []$* **by** *eval*

A Knight's circuit for the (8×5) -board.

26	7	28	15	24
31	16	25	6	29
8	27	30	23	14
17	32	39	34	5
38	9	18	13	22
19	40	33	4	35
10	37	2	21	12
1	20	11	36	3

abbreviation $kc8x5 \equiv$ the (to-path

[[26,7,28,15,24],
 [31,16,25,6,29],
 [8,27,30,23,14],
 [17,32,39,34,5],
 [38,9,18,13,22],
 [19,40,33,4,35],
 [10,37,2,21,12],
 [1,20,11,36,3]])

lemma $kc\text{-}8x5$: *knight's-circuit* $b8x5\ kc8x5$
 by (simp only: *knight's-circuit-exec-simp*) eval

lemma $kc\text{-}8x5\text{-}hd$: $hd\ kc8x5 = (1,1)$ by eval

lemma $kc\text{-}8x5\text{-}last$: $last\ kc8x5 = (3,2)$ by eval

lemma $kc\text{-}8x5\text{-}non\text{-}nil$: $kc8x5 \neq []$ by eval

lemma $kc\text{-}8x5\text{-}si$: *step-in* $kc8x5\ (2,4)\ (4,5)$ (is *step-in* ?ps - -)

proof -

have $0 < (21::nat)\ 21 < length\ ?ps\ last\ (take\ 21\ ?ps) = (2,4)\ hd\ (drop\ 21\ ?ps)$
 $= (4,5)$

by eval+

then show ?thesis unfolding *step-in-def* by blast

qed

abbreviation $b8x6 \equiv$ board 8 6

A Knight's path for the (8×6) -board that starts in the lower-left and ends in the upper-left.

42	11	26	9	34	13
25	48	43	12	27	8
44	41	10	33	14	35
47	24	45	20	7	28
40	19	32	3	36	15
23	46	21	6	29	4
18	39	2	31	16	37
1	22	17	38	5	30

abbreviation $kp8x6ul \equiv$ the (to-path

[[42,11,26,9,34,13],
 [25,48,43,12,27,8],
 [44,41,10,33,14,35],
 [47,24,45,20,7,28],
 [40,19,32,3,36,15],
 [23,46,21,6,29,4],
 [18,39,2,31,16,37],
 [1,22,17,38,5,30]])

lemma $kp\text{-}8x6\text{-}ul$: *knight's-path* $b8x6$ $kp8x6ul$
 by (simp only: *knight's-path-exec-simp*) eval

lemma $kp\text{-}8x6\text{-}ul\text{-}hd$: $hd\ kp8x6ul = (1,1)$ by eval

lemma $kp\text{-}8x6\text{-}ul\text{-}last$: $last\ kp8x6ul = (7,2)$ by eval

lemma $kp\text{-}8x6\text{-}ul\text{-}non\text{-}nil$: $kp8x6ul \neq []$ by eval

A Knight's circuit for the (8×6) -board. I have reversed circuit s.t. the circuit steps from $(2,5)$ to $(4,6)$ and not the other way around. This makes the proofs easier.

8	29	24	45	12	37
25	46	9	38	23	44
30	7	28	13	36	11
47	26	39	10	43	22
6	31	4	27	14	35
3	48	17	40	21	42
32	5	2	19	34	15
1	18	33	16	41	20

abbreviation $kc8x6 \equiv$ the (to-path

[[8,29,24,45,12,37],
 [25,46,9,38,23,44],
 [30,7,28,13,36,11],
 [47,26,39,10,43,22],
 [6,31,4,27,14,35],

```

[3,48,17,40,21,42],
[32,5,2,19,34,15],
[1,18,33,16,41,20]])
lemma kc-8x6: knights-circuit b8x6 kc8x6
  by (simp only: knights-circuit-exec-simp) eval

lemma kc-8x6-hd: hd kc8x6 = (1,1) by eval

lemma kc-8x6-non-nil: kc8x6 ≠ [] by eval

lemma kc-8x6-si: step-in kc8x6 (2,5) (4,6) (is step-in ?ps -)
proof -
  have  $0 < (34::nat)$   $34 < \text{length } ?ps$ 
     $\text{last } (\text{take } 34 \text{ } ?ps) = (2,5)$   $\text{hd } (\text{drop } 34 \text{ } ?ps) = (4,6)$  by eval+
  then show ?thesis unfolding step-in-def by blast
qed

```

abbreviation $b8x7 \equiv \text{board } 8 \ 7$

A Knight's path for the (8×7) -board that starts in the lower-left and ends in the upper-left.

38	19	6	55	46	21	8
5	56	39	20	7	54	45
18	37	4	47	34	9	22
3	48	35	40	53	44	33
36	17	52	49	32	23	10
51	2	29	14	41	26	43
16	13	50	31	28	11	24
1	30	15	12	25	42	27

abbreviation $kp8x7ul \equiv \text{the } (to\text{-path}$
 $[[38,19,6,55,46,21,8],$
 $[5,56,39,20,7,54,45],$
 $[18,37,4,47,34,9,22],$
 $[3,48,35,40,53,44,33],$
 $[36,17,52,49,32,23,10],$
 $[51,2,29,14,41,26,43],$
 $[16,13,50,31,28,11,24],$
 $[1,30,15,12,25,42,27]])$
lemma *kp-8x7-ul: knights-path b8x7 kp8x7ul*
by (*simp only: knights-path-exec-simp*) *eval*

lemma *kp-8x7-ul-hd: hd kp8x7ul = (1,1)* **by** *eval*

lemma *kp-8x7-ul-last: last kp8x7ul = (7,2)* **by** *eval*

lemma *kp-8x7-ul-non-nil: kp8x7ul ≠ []* **by** *eval*

A Knight's circuit for the (8×7) -board. I have reversed circuit s.t. the circuit steps from $(2,6)$ to $(4,7)$ and not the other way around. This makes the proofs easier.

36	31	18	53	20	29	44
17	54	35	30	45	52	21
32	37	46	19	8	43	28
55	16	7	34	27	22	51
38	33	26	47	6	9	42
3	56	15	12	25	50	23
14	39	2	5	48	41	10
1	4	13	40	11	24	49

abbreviation $kc8x7 \equiv$ the (to-path

```

[[36,31,18,53,20,29,44],
 [17,54,35,30,45,52,21],
 [32,37,46,19,8,43,28],
 [55,16,7,34,27,22,51],
 [38,33,26,47,6,9,42],
 [3,56,15,12,25,50,23],
 [14,39,2,5,48,41,10],
 [1,4,13,40,11,24,49]]

```

lemma $kc\text{-}8x7$: knights-circuit $b8x7\ kc8x7$

by (simp only: knights-circuit-exec-simp) eval

lemma $kc\text{-}8x7\text{-}hd$: $hd\ kc8x7 = (1,1)$ **by** eval

lemma $kc\text{-}8x7\text{-}non\text{-}nil$: $kc8x7 \neq []$ **by** eval

lemma $kc\text{-}8x7\text{-}si$: step-in $kc8x7\ (2,6)\ (4,7)$ (is step-in ?ps - -)

proof -

have $0 < (41::nat)\ 41 < length\ ?ps$

$last\ (take\ 41\ ?ps) = (2,6)\ hd\ (drop\ 41\ ?ps) = (4,7)$ **by** eval+

then show ?thesis **unfolding** step-in-def **by** blast

qed

abbreviation $b8x8 \equiv board\ 8\ 8$

The path given for the 8×8 -board that ends in the upper-left is wrong. The Knight cannot move from square 27 to square 28.

24	11	37	9	26	21	39	7
36	64	24	22	38	8	27	20
12	23	10	53	58	49	6	28
63	35	61	50	55	52	19	40
46	13	54	57	48	59	29	5
34	62	47	60	51	56	41	18
14	45	2	32	16	43	4	30
1	33	15	44	3	31	17	42

abbreviation $kp8x8ul\text{-}false \equiv$ the (to-path

[[24,11,37,9,26,21,39,7],
 [36,64,25,22,38,8,27,20],
 [12,23,10,53,58,49,6,28],
 [63,35,61,50,55,52,19,40],
 [46,13,54,57,48,59,29,5],
 [34,62,47,60,51,56,41,18],
 [14,45,2,32,16,43,4,30],
 [1,33,15,44,3,31,17,42]])

lemma $\neg\text{knights-path } b8x8 \text{ } kp8x8ul\text{-}false$
 by (simp only: knights-path-exec-simp) eval

I have computed a correct Knight's path for the 8×8 -board that ends in the upper-left.

38	41	36	27	32	43	20	25
35	64	39	42	21	26	29	44
40	37	6	33	28	31	24	19
5	34	63	14	7	22	45	30
62	13	4	9	58	49	18	23
3	10	61	52	15	8	57	46
12	53	2	59	48	55	50	17
1	60	11	54	51	16	47	56

abbreviation $kp8x8ul \equiv$ the (to-path

[[38,41,36,27,32,43,20,25],
 [35,64,39,42,21,26,29,44],
 [40,37,6,33,28,31,24,19],
 [5,34,63,14,7,22,45,30],
 [62,13,4,9,58,49,18,23],
 [3,10,61,52,15,8,57,46],
 [12,53,2,59,48,55,50,17],
 [1,60,11,54,51,16,47,56]])

lemma $kp\text{-}8x8\text{-}ul$: knights-path $b8x8 \text{ } kp8x8ul$
 by (simp only: knights-path-exec-simp) eval

lemma *kp-8x8-ul-hd*: *hd kp8x8ul = (1,1)* **by** *eval*

lemma *kp-8x8-ul-last*: *last kp8x8ul = (7,2)* **by** *eval*

lemma *kp-8x8-ul-non-nil*: *kp8x8ul ≠ []* **by** *eval*

A Knight's circuit for the (8×8) -board.

48	13	30	9	56	45	28	7
31	10	47	50	29	8	57	44
14	49	12	55	46	59	6	27
11	32	37	60	51	54	43	58
36	15	52	63	38	61	26	5
33	64	35	18	53	40	23	42
16	19	2	39	62	21	4	25
1	34	17	20	3	24	41	22

abbreviation *kc8x8* \equiv *the (to-path*

[[48,13,30,9,56,45,28,7],
[31,10,47,50,29,8,57,44],
[14,49,12,55,46,59,6,27],
[11,32,37,60,51,54,43,58],
[36,15,52,63,38,61,26,5],
[33,64,35,18,53,40,23,42],
[16,19,2,39,62,21,4,25],
[1,34,17,20,3,24,41,22]])

lemma *kc-8x8*: *knight-circuit b8x8 kc8x8*
by (*simp only: knight-circuit-exec-simp*) *eval*

lemma *kc-8x8-hd*: *hd kc8x8 = (1,1)* **by** *eval*

lemma *kc-8x8-non-nil*: *kc8x8 ≠ []* **by** *eval*

lemma *kc-8x8-si*: *step-in kc8x8 (2,7) (4,8) (is step-in ?ps -)*

proof –

have $0 < (4 :: nat)$ $4 < \text{length } ?ps$

last (take 4 ?ps) = (2,7) *hd (drop 4 ?ps) = (4,8)* **by** *eval+*

then show *?thesis* **unfolding** *step-in-def* **by** *blast*

qed

abbreviation *b8x9* \equiv *board 8 9*

A Knight's path for the (8×9) -board that starts in the lower-left and ends in the upper-left.

32	47	6	71	30	45	8	43	26
5	72	31	46	7	70	27	22	9
48	33	4	29	64	23	44	25	42
3	60	35	62	69	28	41	10	21
34	49	68	65	36	63	24	55	40
59	2	61	16	67	56	37	20	11
50	15	66	57	52	13	18	39	54
1	58	51	14	17	38	53	12	19

abbreviation $kp8x9ul \equiv$ the (to-path

$[[32,47,6,71,30,45,8,43,26],$
 $[5,72,31,46,7,70,27,22,9],$
 $[48,33,4,29,64,23,44,25,42],$
 $[3,60,35,62,69,28,41,10,21],$
 $[34,49,68,65,36,63,24,55,40],$
 $[59,2,61,16,67,56,37,20,11],$
 $[50,15,66,57,52,13,18,39,54],$
 $[1,58,51,14,17,38,53,12,19]]$)

lemma $kp-8x9-ul$: *knights-path* $b8x9$ $kp8x9ul$
by (*simp only: knights-path-exec-simp*) *eval*

lemma $kp-8x9-ul-hd$: $hd\ kp8x9ul = (1,1)$ **by** *eval*

lemma $kp-8x9-ul-last$: $last\ kp8x9ul = (7,2)$ **by** *eval*

lemma $kp-8x9-ul-non-nil$: $kp8x9ul \neq []$ **by** *eval*

A Knight's circuit for the (8×9) -board.

42	19	38	5	36	21	34	7	60
39	4	41	20	63	6	59	22	33
18	43	70	37	58	35	68	61	8
3	40	49	64	69	62	57	32	23
50	17	44	71	48	67	54	9	56
45	2	65	14	27	12	29	24	31
16	51	72	47	66	53	26	55	10
1	46	15	52	13	28	11	30	25

abbreviation $kc8x9 \equiv$ the (to-path

$[[42,19,38,5,36,21,34,7,60],$
 $[39,4,41,20,63,6,59,22,33],$
 $[18,43,70,37,58,35,68,61,8],$
 $[3,40,49,64,69,62,57,32,23],$
 $[50,17,44,71,48,67,54,9,56],$
 $[45,2,65,14,27,12,29,24,31],$
 $[16,51,72,47,66,53,26,55,10],$

```

[1,46,15,52,13,28,11,30,25]])
lemma kc-8x9: knights-circuit b8x9 kc8x9
  by (simp only: knights-circuit-exec-simp) eval

lemma kc-8x9-hd: hd kc8x9 = (1,1) by eval

lemma kc-8x9-non-nil: kc8x9 ≠ [] by eval

lemma kc-8x9-si: step-in kc8x9 (2,8) (4,9) (is step-in ?ps -)
proof -
  have  $0 < (55::nat)$   $55 < \text{length } ?ps$ 
     $\text{last } (\text{take } 55 ?ps) = (2,8)$   $\text{hd } (\text{drop } 55 ?ps) = (4,9)$  by eval+
  then show ?thesis unfolding step-in-def by blast
qed

lemmas kp-8xm-ul =
  kp-8x5-ul kp-8x5-ul-hd kp-8x5-ul-last kp-8x5-ul-non-nil
  kp-8x6-ul kp-8x6-ul-hd kp-8x6-ul-last kp-8x6-ul-non-nil
  kp-8x7-ul kp-8x7-ul-hd kp-8x7-ul-last kp-8x7-ul-non-nil
  kp-8x8-ul kp-8x8-ul-hd kp-8x8-ul-last kp-8x8-ul-non-nil
  kp-8x9-ul kp-8x9-ul-hd kp-8x9-ul-last kp-8x9-ul-non-nil

lemmas kc-8xm =
  kc-8x5 kc-8x5-hd kc-8x5-last kc-8x5-non-nil kc-8x5-si
  kc-8x6 kc-8x6-hd kc-8x6-non-nil kc-8x6-si
  kc-8x7 kc-8x7-hd kc-8x7-non-nil kc-8x7-si
  kc-8x8 kc-8x8-hd kc-8x8-non-nil kc-8x8-si
  kc-8x9 kc-8x9-hd kc-8x9-non-nil kc-8x9-si

For every  $8 \times m$ -board with  $m \geq 5$  there exists a knight's circuit.

lemma knights-circuit-8xm-exists:
  assumes  $m \geq 5$ 
  shows  $\exists ps. \text{knights-circuit } (\text{board } 8 \ m) \ ps \wedge \text{step-in } ps \ (2, \text{int } m-1) \ (4, \text{int } m)$ 
  using assms
proof (induction m rule: less-induct)
  case (less m)
  then have  $m \in \{5,6,7,8,9\} \vee 5 \leq m-5$  by auto
  then show ?case
  proof (elim disjE)
    assume  $m \in \{5,6,7,8,9\}$ 
    then show ?thesis using kc-8xm by fastforce
  next
    let  $?ps_2 = \text{kc8x5}$ 
    let  $?b_2 = \text{board } 8 \ 5$ 
    have  $ps_2\text{-prems: knights-circuit } ?b_2 \ ?ps_2 \ \text{hd } ?ps_2 = (1,1) \ \text{last } ?ps_2 = (3,2)$ 
    using kc-8xm by auto
    have  $21 < \text{length } ?ps_2$   $\text{last } (\text{take } 21 ?ps_2) = (2, \text{int } 5-1) \ \text{hd } (\text{drop } 21 ?ps_2) =$ 
     $(4, \text{int } 5)$ 
    by eval+

```

```

then have si: step-in ?ps2 (2,int 5-1) (4,int 5)
  unfolding step-in-def using zero-less-numeral by blast
assume m-ge: 5 ≤ m-5
then obtain ps1 where ps1-IH: knights-circuit (board 8 (m-5)) ps1
  step-in ps1 (2,int (m-5)-1) (4,int (m-5))
  using less.IH[of m-5] knights-path-non-nil by auto
then show ?thesis
  using m-ge ps2-prems si knights-circuit-lr-concat[of 8 m-5 ps1 5 ?ps2] by
auto
qed
qed

```

For every $8 \times m$ -board with $m \geq 5$ there exists a knight's path that starts in $(1,1)$ (bottom-left) and ends in $(7,2)$ (top-left).

lemma *knights-path-8xm-ul-exists*:

```

assumes m ≥ 5
shows ∃ ps. knights-path (board 8 m) ps ∧ hd ps = (1,1) ∧ last ps = (7,2)
using assms
proof -
  have m ∈ {5,6,7,8,9} ∨ 5 ≤ m-5 using assms by auto
  then show ?thesis
  proof (elim disjE)
    assume m ∈ {5,6,7,8,9}
    then show ?thesis using kp-8xm-ul by fastforce
  next
    let ?ps1=kp8x5ul
    have ps1-prems: knights-path b8x5 ?ps1 hd ?ps1 = (1,1) last ?ps1 = (7,2)
      using kp-8xm-ul by auto
    assume m-ge: 5 ≤ m-5
    then have b-prems: 5 ≤ min 8 (m-5)
      unfolding board-def by auto

    obtain ps2 where knights-circuit (board 8 (m-5)) ps2
      using m-ge knights-circuit-8xm-exists[of (m-5)] knights-path-non-nil by auto
    then obtain ps2' where ps2'-prems': knights-circuit (board 8 (m-5)) ps2'
      hd ps2' = (1,1) last ps2' = (3,2)
      using b-prems ⟨5 ≤ min 8 (m-5)⟩ rotate-knights-circuit by blast
    then have ps2'-path: knights-path (board 8 (m-5)) (rev ps2')
      valid-step (last ps2') (hd ps2') hd (rev ps2') = (3,2) last (rev ps2') = (1,1)
      unfolding knights-circuit-def using knights-path-rev by (auto simp: hd-rev
last-rev)

    have 34 < length ?ps1 last (take 34 ?ps1) = (4,5) hd (drop 34 ?ps1) = (2,4)
  by eval+
  then have step-in ?ps1 (4,5) (2,4)
    unfolding step-in-def using zero-less-numeral by blast
  then have step-in ?ps1 (4,5) (2,4)
    valid-step (4,5) (3,int 5+2)
    valid-step (1,int 5+1) (2,4)

```

```

unfolding valid-step-def by auto
then have  $\exists ps. \text{knights-path } (\text{board } 8 \ m) \ ps \wedge \text{hd } ps = \text{hd } ?ps_1 \wedge \text{last } ps = \text{last } ?ps_1$ 
using m-ge ps1-prems ps2'-prems' ps2'-path
knights-path-split-concat[of 8 5 ?ps1 m-5 rev ps2] by auto
then show ?thesis using ps1-prems by auto
qed
qed

```

$5 \leq ?m \implies \exists ps. \text{knights-circuit } (\text{board } 8 \ ?m) \ ps \wedge \text{step-in } ps \ (2, \text{int } ?m - 1) \ (4, \text{int } ?m)$ and $5 \leq ?m \implies \exists ps. \text{knights-path } (\text{board } 8 \ ?m) \ ps \wedge \text{hd } ps = (1, 1) \wedge \text{last } ps = (7, 2)$ formalize Lemma 3 from [1].

lemmas *knights-path-8xm-exists = knights-circuit-8xm-exists knights-path-8xm-ul-exists*

11 Knight's Paths and Circuits for $n \times m$ -Boards

In this section the desired theorems are proved. The proof uses the previous lemmas to construct paths and circuits for arbitrary $n \times m$ -boards.

A Knight's path for the (5×5) -board that starts in the lower-left and ends in the upper-left.

7	20	9	14	5
10	25	6	21	16
19	8	15	4	13
24	11	2	17	22
1	18	23	12	3

abbreviation *kp5x5ul* \equiv *the (to-path*

[[7,20,9,14,5],
[10,25,6,21,16],
[19,8,15,4,13],
[24,11,2,17,22],
[1,18,23,12,3]])

lemma *kp-5x5-ul: knights-path b5x5 kp5x5ul*

by (*simp only: knights-path-exec-simp*) *eval*

A Knight's path for the (5×7) -board that starts in the lower-left and ends in the upper-left.

17	14	25	6	19	8	29
26	35	18	15	28	5	20
13	16	27	24	7	30	9
34	23	2	11	32	21	4
1	12	33	22	3	10	31

abbreviation $kp5x7ul \equiv$ the (to-path

$[[17,14,25,6,19,8,29],$
 $[26,35,18,15,28,5,20],$
 $[13,16,27,24,7,30,9],$
 $[34,23,2,11,32,21,4],$
 $[1,12,33,22,3,10,31]])$

lemma $kp-5x7-ul$: knights-path $b5x7$ $kp5x7ul$

by (simp only: knights-path-exec-simp) eval

A Knight's path for the (5×9) -board that starts in the lower-left and ends in the upper-left.

7	12	37	42	5	18	23	32	27
38	45	6	11	36	31	26	19	24
13	8	43	4	41	22	17	28	33
44	39	2	15	10	35	30	25	20
1	14	9	40	3	16	21	34	29

abbreviation $kp5x9ul \equiv$ the (to-path

$[[7,12,37,42,5,18,23,32,27],$
 $[38,45,6,11,36,31,26,19,24],$
 $[13,8,43,4,41,22,17,28,33],$
 $[44,39,2,15,10,35,30,25,20],$
 $[1,14,9,40,3,16,21,34,29]])$

lemma $kp-5x9-ul$: knights-path $b5x9$ $kp5x9ul$

by (simp only: knights-path-exec-simp) eval

abbreviation $b7x7 \equiv$ board $7\ 7$

A Knight's path for the (7×7) -board that starts in the lower-left and ends in the upper-left.

9	30	19	42	7	32	17
20	49	8	31	18	43	6
29	10	41	36	39	16	33
48	21	38	27	34	5	44
11	28	35	40	37	26	15
22	47	2	13	24	45	4
1	12	23	46	3	14	25

abbreviation $kp7x7ul \equiv$ the (to-path

$[[9,30,19,42,7,32,17],$
 $[20,49,8,31,18,43,6],$
 $[29,10,41,36,39,16,33],$
 $[48,21,38,27,34,5,44],$
 $[11,28,35,40,37,26,15],$
 $[22,47,2,13,24,45,4],$

[1,12,23,46,3,14,25]])
lemma *kp-7x7-ul: knights-path b7x7 kp7x7ul*
by (*simp only: knights-path-exec-simp*) *eval*

abbreviation *b7x9* \equiv *board 7 9*

A Knight's path for the (7×9) -board that starts in the lower-left and ends in the upper-left.

59	4	17	50	37	6	19	30	39
16	63	58	5	18	51	38	7	20
3	60	49	36	57	42	29	40	31
48	15	62	43	52	35	56	21	8
61	2	13	26	45	28	41	32	55
14	47	44	11	24	53	34	9	22
1	12	25	46	27	10	23	54	33

abbreviation *kp7x9ul* \equiv *the (to-path*
[[59,4,17,50,37,6,19,30,39],
[16,63,58,5,18,51,38,7,20],
[3,60,49,36,57,42,29,40,31],
[48,15,62,43,52,35,56,21,8],
[61,2,13,26,45,28,41,32,55],
[14,47,44,11,24,53,34,9,22],
[1,12,25,46,27,10,23,54,33]])

lemma *kp-7x9-ul: knights-path b7x9 kp7x9ul*
by (*simp only: knights-path-exec-simp*) *eval*

abbreviation *b9x7* \equiv *board 9 7*

A Knight's path for the (9×7) -board that starts in the lower-left and ends in the upper-left.

5	20	53	48	7	22	31
52	63	6	21	32	55	8
19	4	49	54	47	30	23
62	51	46	33	56	9	58
3	18	61	50	59	24	29
14	43	34	45	28	57	10
17	2	15	60	35	38	25
42	13	44	27	40	11	36
1	16	41	12	37	26	39

abbreviation *kp9x7ul* \equiv *the (to-path*
[[5,20,53,48,7,22,31],
[52,63,6,21,32,55,8],
[19,4,49,54,47,30,23],
[62,51,46,33,56,9,58],
[3,18,61,50,59,24,29],
[14,43,34,45,28,57,10],
[17,2,15,60,35,38,25],
[42,13,44,27,40,11,36],
[1,16,41,12,37,26,39]])

```

[19,4,49,54,47,30,23],
[62,51,46,33,56,9,58],
[3,18,61,50,59,24,29],
[14,43,34,45,28,57,10],
[17,2,15,60,35,38,25],
[42,13,44,27,40,11,36],
[1,16,41,12,37,26,39]])
lemma kp-9x7-ul: knights-path b9x7 kp9x7ul
  by (simp only: knights-path-exec-simp) eval

```

abbreviation *b9x9* \equiv *board 9 9*

A Knight's path for the (9×9) -board that starts in the lower-left and ends in the upper-left.

13	26	39	52	11	24	37	50	9
40	81	12	25	38	51	10	23	36
27	14	53	58	63	68	73	8	49
80	41	64	67	72	57	62	35	22
15	28	59	54	65	74	69	48	7
42	79	66	71	76	61	56	21	34
29	16	77	60	55	70	75	6	47
78	43	2	31	18	45	4	33	20
1	30	17	44	3	32	19	46	5

```

abbreviation kp9x9ul  $\equiv$  the (to-path
[[13,26,39,52,11,24,37,50,9],
[40,81,12,25,38,51,10,23,36],
[27,14,53,58,63,68,73,8,49],
[80,41,64,67,72,57,62,35,22],
[15,28,59,54,65,74,69,48,7],
[42,79,66,71,76,61,56,21,34],
[29,16,77,60,55,70,75,6,47],
[78,43,2,31,18,45,4,33,20],
[1,30,17,44,3,32,19,46,5]])
lemma kp-9x9-ul: knights-path b9x9 kp9x9ul
  by (simp only: knights-path-exec-simp) eval

```

The following lemma is a sub-proof used in Lemma 4 in [1]. I moved the sub-proof out to a separate lemma.

```

lemma knights-circuit-exists-even-n-gr10:
  assumes even n n ≥ 10 m ≥ 5
     $\exists ps. \text{knights-path (board (n-5) m) ps} \wedge \text{hd ps} = (\text{int (n-5)}, 1)$ 
     $\wedge \text{last ps} = (\text{int (n-5)}-1, \text{int m}-1)$ 
  shows  $\exists ps. \text{knights-circuit (board m n) ps}$ 
  using assms
proof –

```

```

let ?b2=board (n-5) m
assume n ≥ 10
then obtain ps2 where ps2-prems: knights-path ?b2 ps2 hd ps2 = (int (n-5),1)

    last ps2 = (int (n-5)-1,int m-1)
    using assms by auto
let ?ps2-m2=mirror2 ps2
have ps2-m2-prems: knights-path ?b2 ?ps2-m2 hd ?ps2-m2 = (int (n-5),int m)
    last ?ps2-m2 = (int (n-5)-1,2)
    using ps2-prems mirror2-knights-path hd-mirror2 last-mirror2 by auto

obtain ps1 where ps1-prems: knights-path (board 5 m) ps1 hd ps1 = (1,1)last
ps1 = (2,int m-1)
    using assms knights-path-5xm-exists by auto
let ?ps1'=trans-path (int (n-5),0) ps1
let ?b1'=trans-board (int (n-5),0) (board 5 m)
have ps1'-prems: knights-path ?b1' ?ps1' hd ?ps1' = (int (n-5)+1,1)
    last ?ps1' = (int (n-5)+2,int m-1)
    using ps1-prems trans-knights-path knights-path-non-nil hd-trans-path last-trans-path
by auto

let ?ps=?ps1'@?ps2-m2
let ?psT=transpose ?ps

have n-5 ≥ 5 using ⟨n ≥ 10⟩ by auto
have inter: ?b1' ∩ ?b2 = {}
    unfolding trans-board-def board-def using ⟨n-5 ≥ 5⟩ by auto
have union: ?b1' ∪ ?b2 = board n m
    using ⟨n-5 ≥ 5⟩ board-concatT[of n-5 m 5] by auto

have vs: valid-step (last ?ps1') (hd ?ps2-m2) and valid-step (last ?ps2-m2) (hd
?ps1')
    unfolding valid-step-def using ps1'-prems ps2-m2-prems by auto
then have vs-c: valid-step (last ?ps) (hd ?ps)
    using ps1'-prems ps2-m2-prems knights-path-non-nil by auto

have knights-path (board n m) ?ps
    using ps1'-prems ps2-m2-prems inter vs union knights-path-append[of ?b1' ?ps1'
?b2 ?ps2-m2]
    by auto
then have knights-circuit (board n m) ?ps
    unfolding knights-circuit-def using vs-c by auto
then show ?thesis using transpose-knights-circuit by auto
qed

```

For every $n \times m$ -board with $\min n m \geq 5$ and odd n there exists a Knight's path that starts in $(n,1)$ (top-left) and ends in $(n-1,m-1)$ (top-right).

This lemma formalizes Lemma 4 from [1]. Formalizing the proof of this lemma was quite challenging as a lot of details on how to exactly combine

the boards are left out in the original proof in [1].

lemma *knight-path-odd-n-exists*:

assumes *odd n min n m* ≥ 5

shows $\exists ps. \text{knight-path } (\text{board } n \ m) \ ps \wedge \text{hd } ps = (\text{int } n, 1) \wedge \text{last } ps = (\text{int } n-1, \text{int } m-1)$

using *assms*

proof –

obtain *x* **where** $x = n + m$ **by** *auto*

then show *?thesis*

using *assms*

proof (*induction x arbitrary: n m rule: less-induct*)

case (*less x*)

then have $m = 5 \vee m = 6 \vee m = 7 \vee m = 8 \vee m = 9 \vee m \geq 10$ **by** *auto*

then show *?case*

proof (*elim disjE*)

assume [*simp*]: $m = 5$

have *odd n n* ≥ 5 **using** *less* **by** *auto*

then have $n = 5 \vee n = 7 \vee n = 9 \vee n-5 \geq 5$ **by** *presburger*

then show *?thesis*

proof (*elim disjE*)

assume [*simp*]: $n = 5$

let *?ps=mirror1* (*transpose kp5x5ul*)

have *kp: knight-path (board n m) ?ps*

using *kp-5x5-ul rot90-knight-path* **by** *auto*

have *hd ?ps = (int n, 1) last ?ps = (int n-1, int m-1)*

by (*simp only: <m = 5> <n = 5> | eval*)**+**

then show *?thesis* **using** *kp* **by** *auto*

next

assume [*simp*]: $n = 7$

let *?ps=mirror1* (*transpose kp5x7ul*)

have *kp: knight-path (board n m) ?ps*

using *kp-5x7-ul rot90-knight-path* **by** *auto*

have *hd ?ps = (int n, 1) last ?ps = (int n-1, int m-1)*

by (*simp only: <m = 5> <n = 7> | eval*)**+**

then show *?thesis* **using** *kp* **by** *auto*

next

assume [*simp*]: $n = 9$

let *?ps=mirror1* (*transpose kp5x9ul*)

have *kp: knight-path (board n m) ?ps*

using *kp-5x9-ul rot90-knight-path* **by** *auto*

have *hd ?ps = (int n, 1) last ?ps = (int n-1, int m-1)*

by (*simp only: <m = 5> <n = 9> | eval*)**+**

then show *?thesis* **using** *kp* **by** *auto*

next

let *?b₂=board m (n-5)*

assume $n-5 \geq 5$

then have $\exists ps. \text{knight-circuit } ?b_2 \ ps$

proof –

have $n-5 = 6 \vee n-5 = 8 \vee n-5 \geq 10$

```

    using  $\langle n-5 \geq 5 \rangle$  less by presburger
  then show ?thesis
proof (elim disjE)
  assume  $n-5 = 6$ 
  then obtain ps where knights-circuit (board (n-5) m) ps
    using knights-path-6xm-exists[of m] by auto
  then show ?thesis
    using transpose-knights-circuit by auto
next
  assume  $n-5 = 8$ 
  then obtain ps where knights-circuit (board (n-5) m) ps
    using knights-path-8xm-exists[of m] by auto
  then show ?thesis
    using transpose-knights-circuit by auto
next
  assume  $n-5 \geq 10$ 
  then show ?thesis
    using less less.IH[of n-10+m n-10 m]
      knights-circuit-exists-even-n-gr10[of n-5 m] by auto
qed
qed
then obtain ps2 where knights-circuit ?b2 ps2 hd ps2 = (1,1) last ps2 =
(3,2)
  using  $\langle n-5 \geq 5 \rangle$  rotate-knights-circuit[of m n-5] by auto
then have rev-ps2-prems: knights-path ?b2 (rev ps2) valid-step (last ps2) (hd
ps2)
  hd (rev ps2) = (3,2) last (rev ps2) = (1,1)
  unfolding knights-circuit-def using knights-path-rev by (auto simp: hd-rev
last-rev)

let ?ps1=kp5x5ul
have ps1-prems: knights-path (board 5 5) ?ps1 hd ?ps1 = (1,1) last ?ps1 =
(4,2)
  using kp-5x5-ul by simp eval+

  have 16 < length ?ps1 last (take 16 ?ps1) = (4,5) hd (drop 16 ?ps1) =
(2,4) by eval+
  then have si: step-in ?ps1 (4,5) (2,4)
    unfolding step-in-def using zero-less-numeral by blast

  have vs: valid-step (4,5) (3,int 5+2) valid-step (1,int 5+1) (2,4)
    unfolding valid-step-def by auto

obtain ps where knights-path (board m n) ps hd ps = (1,1) last ps = (4,2)
  using  $\langle n-5 \geq 5 \rangle$  ps1-prems rev-ps2-prems si vs
    knights-path-split-concat[of 5 5 ?ps1 n-5 rev ps2 (4,5) (2,4)] by auto
then show ?thesis
  using rot90-knights-path hd-rot90-knights-path last-rot90-knights-path by
fastforce

```

```

qed
next
  assume [simp]: m = 6
  then obtain ps where
    ps-prems: knights-path (board m n) ps hd ps = (1,1) last ps = (int m-1,2)
    using less knights-path-6xm-exists[of n] by auto
  let ?ps'=mirror1 (transpose ps)
  have knights-path (board n m) ?ps' hd ?ps' = (int n,1) last ?ps' = (int n-1,int
m-1)
    using ps-prems rot90-knights-path hd-rot90-knights-path last-rot90-knights-path
  by auto
  then show ?thesis by auto
next
  assume [simp]: m = 7
  have odd n n ≥ 5 using less by auto
  then have n = 5 ∨ n = 7 ∨ n = 9 ∨ n-5 ≥ 5 by presburger
  then show ?thesis
  proof (elim disjE)
    assume [simp]: n = 5
    let ?ps=mirror1 kp5x7lr
    have kp: knights-path (board n m) ?ps
      using kp-5x7-lr mirror1-knights-path by auto
    have hd ?ps = (int n,1) last ?ps = (int n-1,int m-1)
      by (simp only: ⟨m = 7⟩ ⟨n = 5⟩ | eval)+
    then show ?thesis using kp by auto
  next
    assume [simp]: n = 7
    let ?ps=mirror1 (transpose kp7x7ul)
    have kp: knights-path (board n m) ?ps
      using kp-7x7-ul rot90-knights-path by auto
    have hd ?ps = (int n,1) last ?ps = (int n-1,int m-1)
      by (simp only: ⟨m = 7⟩ ⟨n = 7⟩ | eval)+
    then show ?thesis using kp by auto
  next
    assume [simp]: n = 9
    let ?ps=mirror1 (transpose kp7x9ul)
    have kp: knights-path (board n m) ?ps
      using kp-7x9-ul rot90-knights-path by auto
    have hd ?ps = (int n,1) last ?ps = (int n-1,int m-1)
      by (simp only: ⟨m = 7⟩ ⟨n = 9⟩ | eval)+
    then show ?thesis using kp by auto
  next
    let ?b2=board m (n-5)
    let ?b2T=board (n-5) m
    assume n-5 ≥ 5
    then have ∃ ps. knights-circuit ?b2 ps
    proof -
      have n-5 = 6 ∨ n-5 = 8 ∨ n-5 ≥ 10
        using ⟨n-5 ≥ 5⟩ less by presburger

```

```

then show ?thesis
proof (elim disjE)
  assume  $n-5 = 6$ 
  then obtain  $ps$  where  $knights-circuit$  (board  $(n-5)$   $m$ )  $ps$ 
    using  $knights-path-6xm-exists[of\ m]$  by auto
  then show ?thesis
    using  $transpose-knights-circuit$  by auto
next
  assume  $n-5 = 8$ 
  then obtain  $ps$  where  $knights-circuit$  (board  $(n-5)$   $m$ )  $ps$ 
    using  $knights-path-8xm-exists[of\ m]$  by auto
  then show ?thesis
    using  $transpose-knights-circuit$  by auto
next
  assume  $n-5 \geq 10$ 
  then show ?thesis
    using  $less\ less.IH[of\ n-10+m\ n-10\ m]$ 
       $knights-circuit-exists-even-n-gr10[of\ n-5\ m]$  by auto
qed
qed
then obtain  $ps_2$  where  $ps_2-prems: knights-circuit\ ?b_2\ ps_2\ hd\ ps_2 = (1,1)$ 
   $last\ ps_2 = (3,2)$ 
  using  $\langle n-5 \geq 5 \rangle\ rotate-knights-circuit[of\ m\ n-5]$  by auto
let  $?ps_2T = transpose\ ps_2$ 
have  $ps_2T-prems: knights-path\ ?b_2T\ ?ps_2T\ hd\ ?ps_2T = (1,1)\ last\ ?ps_2T =$ 
 $(2,3)$ 
  using  $ps_2-prems\ transpose-knights-path\ knights-path-non-nil\ hd-transpose$ 
 $last-transpose$ 
  unfolding  $knights-circuit-def\ transpose-square-def$  by auto

let  $?ps_1 = kp5x7lr$ 
have  $ps_1-prems: knights-path\ b5x7\ ?ps_1\ hd\ ?ps_1 = (1,1)\ last\ ?ps_1 = (2,6)$ 
  using  $kp-5x7-lr$  by simp eval+

have  $29 < length\ ?ps_1\ last\ (take\ 29\ ?ps_1) = (4,2)\ hd\ (drop\ 29\ ?ps_1) =$ 
 $(5,4)$  by eval+
then have  $si: step-in\ ?ps_1\ (4,2)\ (5,4)$ 
  unfolding  $step-in-def$  using  $zero-less-numeral$  by blast

have  $vs: valid-step\ (4,2)\ (int\ 5+1,1)\ valid-step\ (int\ 5+2,3)\ (5,4)$ 
  unfolding  $valid-step-def$  by auto

obtain  $ps$  where  $knights-path$  (board  $n$   $m$ )  $ps\ hd\ ps = (1,1)\ last\ ps = (2,6)$ 
  using  $\langle n-5 \geq 5 \rangle\ ps_1-prems\ ps_2T-prems\ si\ vs$ 
   $knights-path-split-concatT[of\ 5\ m\ ?ps_1\ n-5\ ?ps_2T\ (4,2)\ (5,4)]$  by auto
then show ?thesis
  using  $mirror1-knights-path\ hd-mirror1\ last-mirror1$  by fastforce
qed
next

```

```

    assume [simp]: m = 8
    then obtain ps where ps-prems: knights-path (board m n) ps hd ps = (1,1)
      last ps = (int m-1,2)
      using less knights-path-8xm-exists[of n] by auto
    let ?ps'=mirror1 (transpose ps)
    have knights-path (board n m) ?ps' hd ?ps' = (int n,1) last ?ps' = (int n-1,int
m-1)
      using ps-prems rot90-knights-path hd-rot90-knights-path last-rot90-knights-path
by auto
    then show ?thesis by auto
  next
    assume [simp]: m = 9
    have odd n n ≥ 5 using less by auto
    then have n = 5 ∨ n = 7 ∨ n = 9 ∨ n-5 ≥ 5 by presburger
    then show ?thesis
  proof (elim disjE)
    assume [simp]: n = 5
    let ?ps=mirror1 kp5x9lr
    have kp: knights-path (board n m) ?ps
      using kp-5x9-lr mirror1-knights-path by auto
    have hd ?ps = (int n,1) last ?ps = (int n-1,int m-1)
      by (simp only: ⟨m = 9⟩ ⟨n = 5⟩ | eval)+
    then show ?thesis using kp by auto
  next
    assume [simp]: n = 7
    let ?ps=mirror1 (transpose kp9x7ul)
    have kp: knights-path (board n m) ?ps
      using kp-9x7-ul rot90-knights-path by auto
    have hd ?ps = (int n,1) last ?ps = (int n-1,int m-1)
      by (simp only: ⟨m = 9⟩ ⟨n = 7⟩ | eval)+
    then show ?thesis using kp by auto
  next
    assume [simp]: n = 9
    let ?ps=mirror1 (transpose kp9x9ul)
    have kp: knights-path (board n m) ?ps
      using kp-9x9-ul rot90-knights-path by auto
    have hd ?ps = (int n,1) last ?ps = (int n-1,int m-1)
      by (simp only: ⟨m = 9⟩ ⟨n = 9⟩ | eval)+
    then show ?thesis using kp by auto
  next
    let ?b2=board m (n-5)
    let ?b2T=board (n-5) m
    assume n-5 ≥ 5
    then have ∃ ps. knights-circuit ?b2 ps
    proof -
      have n-5 = 6 ∨ n-5 = 8 ∨ n-5 ≥ 10
        using ⟨n-5 ≥ 5⟩ less by presburger
      then show ?thesis
    proof (elim disjE)

```



```

    assume  $n-5 = 6$ 
    then obtain  $ps$  where  $knights-circuit$  (board  $(n-5)$   $m$ )  $ps$ 
      using  $knights-path-6xm-exists[of\ m]$  by auto
    then show  $?thesis$ 
      using  $transpose-knights-circuit$  by auto
  next
    assume  $n-5 = 8$ 
    then obtain  $ps$  where  $knights-circuit$  (board  $(n-5)$   $m$ )  $ps$ 
      using  $knights-path-8xm-exists[of\ m]$  by auto
    then show  $?thesis$ 
      using  $transpose-knights-circuit$  by auto
  next
    assume  $n-5 \geq 10$ 
    then show  $?thesis$ 
      using  $less\ less.IH[of\ n-10+m\ n-10\ m]$ 
         $knights-circuit-exists-even-n-gr10[of\ n-5\ m]$  by auto
  qed
qed
then obtain  $ps_2$  where  $ps_2-prems: knights-circuit\ ?b_2\ ps_2\ hd\ ps_2 = (1,1)$ 
   $last\ ps_2 = (3,2)$ 
  using  $\langle n-5 \geq 5 \rangle rotate-knights-circuit[of\ m\ n-5]$  by auto
let  $?ps_2T = transpose\ (rev\ ps_2)$ 
have  $ps_2T-prems: knights-path\ ?b_2T\ ?ps_2T\ hd\ ?ps_2T = (2,3)\ last\ ?ps_2T =$ 
 $(1,1)$ 
  using  $ps_2-prems\ knights-path-rev\ transpose-knights-path\ knights-path-non-nil$ 
   $hd-transpose\ last-transpose$ 
  unfolding  $knights-circuit-def\ transpose-square-def$  by (auto simp:  $hd-rev$ 
 $last-rev$ )

let  $?ps_1 = kp5x9lr$ 
have  $ps_1-prems: knights-path\ b5x9\ ?ps_1\ hd\ ?ps_1 = (1,1)\ last\ ?ps_1 = (2,8)$ 
  using  $kp-5x9-lr$  by simp eval+

have  $16 < length\ ?ps_1\ last\ (take\ 16\ ?ps_1) = (5,4)\ hd\ (drop\ 16\ ?ps_1) =$ 
 $(4,2)$  by eval+
then have  $si: step-in\ ?ps_1\ (5,4)\ (4,2)$ 
  unfolding  $step-in-def$  using  $zero-less-numeral$  by blast

have  $vs: valid-step\ (5,4)\ (int\ 5+2,3)\ valid-step\ (int\ 5+1,1)\ (4,2)$ 
  unfolding  $valid-step-def$  by auto

obtain  $ps$  where  $knights-path$  (board  $n\ m$ )  $ps\ hd\ ps = (1,1)\ last\ ps = (2,8)$ 
  using  $\langle n-5 \geq 5 \rangle ps_1-prems\ ps_2T-prems\ si\ vs$ 
   $knights-path-split-concatT[of\ 5\ m\ ?ps_1\ n-5\ ?ps_2T\ (5,4)\ (4,2)]$  by auto
then show  $?thesis$ 
  using  $mirror1-knights-path\ hd-mirror1\ last-mirror1$  by fastforce
qed
next

```

```

let ?b1=board n 5
let ?b2=board n (m-5)
assume m ≥ 10
then have n+5 < x 5 ≤ min n 5 n+(m-5) < x 5 ≤ min n (m-5)
  using less by auto
then obtain ps1 ps2 where kp-prems:
  knights-path ?b1 ps1 hd ps1 = (int n,1) last ps1 = (int n-1,4)
  knights-path (board n (m-5)) ps2 hd ps2 = (int n,1) last ps2 = (int n-1,int
(m-5)-1)
  using less.prem less.IH[of n+5 n 5] less.IH[of n+(m-5) n m-5] by auto
let ?ps=ps1@trans-path (0,int 5) ps2
have valid-step (last ps1) (int n,int 5+1)
  unfolding valid-step-def using kp-prems by auto
then have knights-path (board n m) ?ps hd ?ps = (int n,1) last ?ps = (int
n-1,int m-1)
  using ⟨m ≥ 10⟩ kp-prems knights-path-concat[of n 5 ps1 m-5 ps2]
  knights-path-non-nil trans-path-non-nil last-trans-path by auto
then show ?thesis by auto
qed
qed
qed

```

Auxiliary lemma that constructs a Knight's circuit if $m \geq 5$ and $n \geq 10 \wedge$ even n .

```

lemma knights-circuit-exists-n-even-gr-10:
  assumes n ≥ 10 ∧ even n m ≥ 5
  shows ∃ps. knights-circuit (board n m) ps
  using assms
proof –
  obtain ps1 where ps1-prems: knights-path (board 5 m) ps1 hd ps1 = (1,1)
    last ps1 = (2,int m-1)
  using assms knights-path-5xm-exists by auto
  let ?ps1'=trans-path (int (n-5),0) ps1
  let ?b5xm'=trans-board (int (n-5),0) (board 5 m)
  have ps1'-prems: knights-path ?b5xm' ?ps1' hd ?ps1' = (int (n-5)+1,1)
    last ?ps1' = (int (n-5)+2,int m-1)
  using ps1-prems trans-knights-path knights-path-non-nil hd-trans-path last-trans-path
by auto

  assume n ≥ 10 ∧ even n
  then have odd (n-5) min (n-5) m ≥ 5 using assms by auto
  then obtain ps2 where ps2-prems: knights-path (board (n-5) m) ps2 hd ps2 =
(int (n-5),1)
    last ps2 = (int (n-5)-1,int m-1)
  using knights-path-odd-n-exists[of n-5 m] by auto
  let ?ps2'=mirror2 ps2
  have ps2'-prems: knights-path (board (n-5) m) ?ps2' hd ?ps2' = (int (n-5),int
m)
    last ?ps2' = (int (n-5)-1,2)

```

```

using ps2-prems mirror2-knights-path hd-mirror2 last-mirror2 by auto

have inter:  $?b5xm' \cap \text{board } (n-5) \ m = \{\}$ 
unfolding trans-board-def board-def by auto

have union:  $\text{board } n \ m = ?b5xm' \cup \text{board } (n-5) \ m$ 
using  $\langle n \geq 10 \ \wedge \ \text{even } n \rangle$  board-concatT[of n-5 m 5] by auto

have vs: valid-step (last  $?ps_1'$ ) (hd  $?ps_2'$ ) valid-step (last  $?ps_2'$ ) (hd  $?ps_1'$ )
using ps1'-prems ps2'-prems unfolding valid-step-def by auto

let  $?ps = ?ps_1' @ ?ps_2'$ 
have last  $?ps = \text{last } ?ps_2'$  hd  $?ps = \text{hd } ?ps_1'$ 
using ps1'-prems ps2'-prems knights-path-non-nil by auto
then have vs-c: valid-step (last  $?ps$ ) (hd  $?ps$ )
using vs by auto

have knights-path (board  $n \ m$ )  $?ps$ 
using ps1'-prems ps2'-prems inter union vs knights-path-append by auto
then show ?thesis
using vs-c unfolding knights-circuit-def by blast
qed

```

Final Theorem 1: For every $n \times m$ -board with $\min n \ m \geq 5$ and $n*m$ even there exists a Knight's circuit.

```

theorem knights-circuit-exists:
  assumes  $\min n \ m \geq 5 \ \text{even } (n*m)$ 
  shows  $\exists ps. \text{knights-circuit } (\text{board } n \ m) \ ps$ 
  using assms
proof –
  have  $n = 6 \vee m = 6 \vee n = 8 \vee m = 8 \vee (n \geq 10 \ \wedge \ \text{even } n) \vee (m \geq 10 \ \wedge \ \text{even } m)$ 
  using assms by auto
  then show ?thesis
  proof (elim disjE)
    assume  $n = 6$ 
    then show ?thesis
    using assms knights-path-6xm-exists by auto
  next
    assume  $m = 6$ 
    then obtain ps where knights-circuit (board  $m \ n$ ) ps
    using assms knights-path-6xm-exists by auto
    then show ?thesis
    using transpose-knights-circuit by auto
  next
    assume  $n = 8$ 
    then show ?thesis
    using assms knights-path-8xm-exists by auto
  next

```

```

    assume  $m = 8$ 
    then obtain  $ps$  where  $knights-circuit$  (board  $m$   $n$ )  $ps$ 
      using  $assms$   $knights-path-8xm-exists$  by  $auto$ 
    then show  $?thesis$ 
      using  $transpose-knights-circuit$  by  $auto$ 
  next
    assume  $n \geq 10 \wedge even\ n$ 
    then show  $?thesis$ 
      using  $assms$   $knights-circuit-exists-n-even-gr-10$  by  $auto$ 
  next
    assume  $m \geq 10 \wedge even\ m$ 
    then obtain  $ps$  where  $knights-circuit$  (board  $m$   $n$ )  $ps$ 
      using  $assms$   $knights-circuit-exists-n-even-gr-10$  by  $auto$ 
    then show  $?thesis$ 
      using  $transpose-knights-circuit$  by  $auto$ 
  qed
qed

```

Final Theorem 2: for every $n \times m$ -board with $min\ n\ m \geq 5$ there exists a Knight's path.

```

theorem  $knights-path-exists$ :
  assumes  $min\ n\ m \geq 5$ 
  shows  $\exists ps. knights-path$  (board  $n$   $m$ )  $ps$ 
  using  $assms$ 
proof -
  have  $odd\ n \vee odd\ m \vee even\ (n*m)$  by  $simp$ 
  then show  $?thesis$ 
  proof (elim  $disjE$ )
    assume  $odd\ n$ 
    then show  $?thesis$ 
      using  $assms$   $knights-path-odd-n-exists$  by  $auto$ 
  next
    assume  $odd\ m$ 
    then obtain  $ps$  where  $knights-path$  (board  $m$   $n$ )  $ps$ 
      using  $assms$   $knights-path-odd-n-exists$  by  $auto$ 
    then show  $?thesis$ 
      using  $transpose-knights-path$  by  $auto$ 
  next
    assume  $even\ (n*m)$ 
    then show  $?thesis$ 
      using  $assms$   $knights-circuit-exists$  by (auto simp:  $knights-circuit-def$ )
  qed
qed

```

THE END

end

References

- [1] P. Cull and J. D. Curtins. Knight's tour revisited. *Fibonacci Quarterly*, 16:276–285, 1978.