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$S = mid \Rightarrow S = 3$ $O = 1$	[17]	1 10 10 10 5 (Dim) 100 18 (D) 212 12 12 100 2 (D) 11 1 1
0 1 2 3 4 5 6 3 4 5 6 7 1 2 3 4 5 6 7 1 2 3 4 5 6 7 1 2 3 4 5 6 7 1 2 3 4 5 6 7 1 2 3 4 5 6 7 1 2 3 4 5 6 7 1 2 3 4 5 6 7 1 2 3 4 5 6 7 1 2 3 4 5 6 7 1 2 3 4 5 6 7 1 2 3 4 5 6 7 1 1 2 3 4 5 6 7 1 1 2 1 1 1 1 1 1 1 1 1 1		$Steb-1 \qquad 6 > = 2$
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Step-2 $mid = 3+6 = 4$ $avr [mid] = 7$		3 4 5 6 7
Step-2 $mid = 3+6 = 4$ $wur (mid) = 7$		[him] reso < (1-bim) reso & & o = < 1-bim) di
mid = 3+6 = 4 $avr(mid) = 7$		Steb-2
our [mid] = 7		mid = 3+6 = 4
		our [mid] = 7

```
> aur [mid] > aur [mid+1] 4 Page Page thence return mid i.e 4th index.
7>=2 3 True, then search in the
right averay by doing s=mid
Step-3 (Extra step)
arr [mid] > arr [mid+1] => False
arr [mid-1] > arr [mid] + True & hence
return mid-1 i.e 4th index & hence
the pivot element is 7.
         int bivot Element (vector <int) aver).
    int s = o i
    inte = arr. size () -1;
    int mid = s+ (e-s)/2;
while (S < e) {
  if (mid+1 < ary. size() && aver [mid) > arr (mid)
    return mids
 if (mid-1) = 0 & & avr [mid-1] > avr [mid]
    return mid-1;
```

