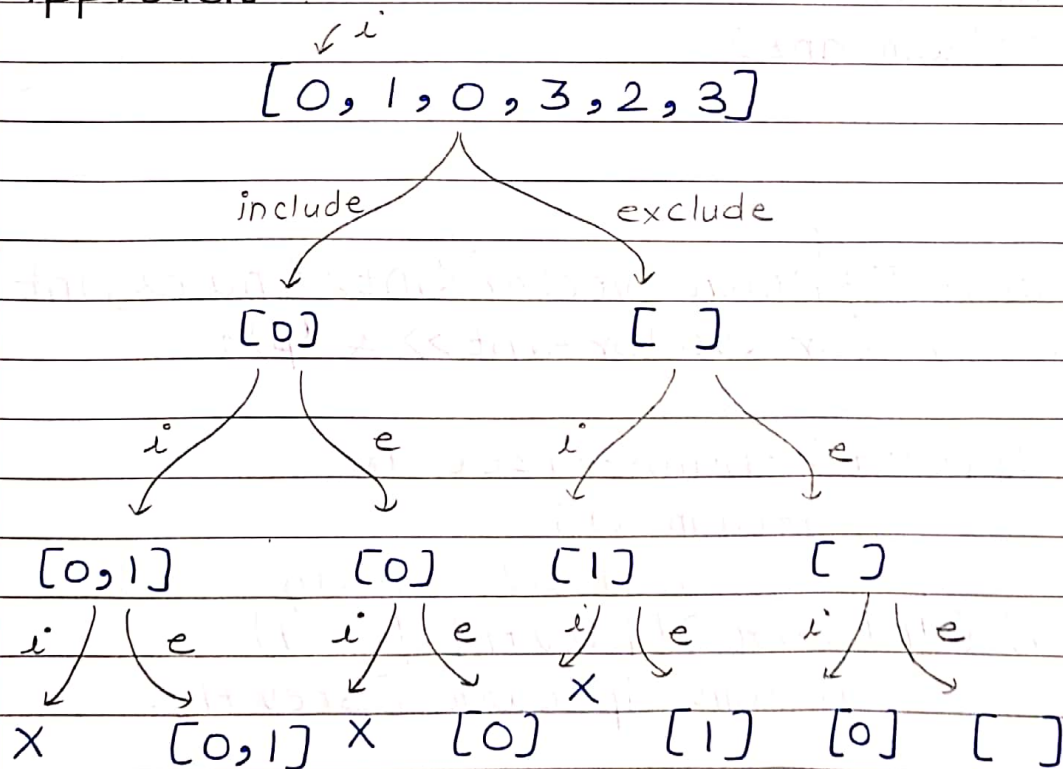


16/06/2023Q1 Longest increasing subsequencei/p  $\rightarrow \{5, 8, 3, 2, 1, 9, 7\}$ 

Total subsequences possible =  $2^n$ . Filter out the increasing subsequence and we have to return the longest length subsequence.

Increasing subsequence  $\rightarrow [5, 8, 9] \rightarrow$  example of increasing subsequence

Approach

Note  $\rightarrow$  We need to include a particular element only when it is bigger than previous element.

Code

```

int solveRec (vector<int> &nums, int curr, int prev)
    // Base case
    if (curr >= nums.size())
        return 0;
    // include call
    int include = 0;
    // 1st element or previous element smaller
    if (prev == -1 || nums[curr] > nums[prev]){
        include = 1 + solveRec (nums, curr+1, curr);
    }
    // exclude call
    int exclude = 0 + solveRec (nums, curr+1, prev);
    int ans = max(include, exclude);
    return ans;

```

3

remains  
some

// Top down approach

```

int solveTopDown (vector<int> &nums, int curr,
int prev, vector<vector<int>> &dp) {

```

// Base case

```

    if (curr >= nums.size())

```

```

        return 0;

```

// Step 3: Check if ans already exists

```

    if (dp[curr][prev+1] != -1)

```

```

        return dp[curr][prev+1];

```

// include call

```

    int include = 0;

```

```

    if (prev != -1 || nums[curr] > nums[prev])

```

```

        include = 1 + solveTopDown (nums,
curr+1, curr, dp);

```

3

// exclude



```
int exclude = 0 + solveTopDown(nums, curr+1,
prev, dp);
```

// Step 2 : Store ans in dp array

// prev+1 as prev = -1 can be there and prev+1 = 0.

```
dp[curr][prev+1] = ans;
return dp[curr][prev+1];
```

valid index

```
}
```

// Bottom up approach

```
int solveTab (vector <int> &nums) {
```

```
    int n = nums.size();
```

// Step 1 : Create dp array

```
    vector <vector <int>> dp(n+1, vector <int> (n+1, 0))
```

// Step 2 : Observe base case of top-down → Handled

// Step 3 : Reverse flow of top-down

```
    for (int curr = n-1; curr >= 0; curr--) {
```

```
        for (int prev = curr-1; prev >= -1; prev--) {
```

// include call

```
            int include = 0;
```

```
            if (prev == -1 || nums[curr] > nums[prev]) {
```

```
                include = 1 + dp[curr+1][curr+1];
```

```
            }
```

// exclude

```
            int exclude = 0 + dp[curr+1][prev+1];
```

```
            dp[curr][prev+1] = max(include, exclude);
```

```
        }
```

```
    }
```

```
    return dp[0][0];
```

```
}
```

\* Space optimization possible or not ?

Yes it is possible.

- 1) Create 2, 1D arrays.  
vector<int> currRow (n+1, 0);  
vector<int> nextRow (n+1, 0);
- 2) Replace  $dp[curr+1]--$  with  $nextRow[--]$   
and replace  $dp[curr]--$  with  $currRow[--]$
- 3) Shift  $\rightarrow$  nextRow = currRow  $\rightarrow$  going upwards

DP with binary search

arr  $\rightarrow$  [10, 9, 2, 5, 3, 7, 101, 18]

$i=0 \rightarrow$  [10]

$i=1 \rightarrow$  [9] (Overwrite)

$i=2 \rightarrow$  [2] (Overwrite)

$i=3 \rightarrow$  [2, 5] (Include)

$i=4 \rightarrow$  [2, 3] (Overwrite)

$i=5 \rightarrow$  [2, 3, 7] (Include)

$i=6 \rightarrow$  [2, 3, 7, 101] (Include)

$i=7 \rightarrow$  [2, 3, 7, 18]  $\rightarrow$  length = 4  
(Overwrite)

Note  $\rightarrow$  Overwrite and include is to be done.

Code

```
int solveOptimal (vector<int> &arr) {
```

```
    if (arr.size() == 0)
```

```
        return 0;
```

```
    vector<int> ans;
```

```
    ans.push-back(arr[0]);
```



```

for (int i = 1; i < arr.size(); i++) {
    if (arr[i] > ans.back())
        ans.push_back(arr[i]);
    else { // overwrite
        // find index of just bigger element
        int index = lower_bound(ans.begin(),
                                ans.end(), arr[i]) - ans.begin();
        ans[index] = arr[i];
    }
}
return ans.size();
}

```

TC =  $O(n \log n)$

\* Russian doll approach

arr → { [5, 4], [6, 4], [6, 7], [2, 3] }

↑ height

width ↑

Arrange them in increasing order of width. If width is same, then take higher height. Simply apply LIS solution on the width now.

[2, 3]

[5, 4]

[6, 7]

[6, 4]

Now apply LIS on [2, 5, 6, 6] and hence here 3 is the answer.

Q2 Max height by stacking cuboid.

i/p → [50, 45, 20]

[95, 37, 53]

[45, 23, 12]

1) First sort. (We need maximum height.)

w l h

[20, 45, 50]

[37, 53, 95]

[12, 23, 45]

2) Now sort according to the width.

w l h

[12, 23, 45]

[20, 45, 50]

[37, 53, 95]

3) Now simply apply LIS.

Code

```
int maxHeight (vector <vector <int>> & cuboids) {
    // Sort every array
    for (auto &a : cuboids)
        sort (a.begin(), a.end());
    // Sort 2D array
    sort (cuboids.begin(), cuboids.end());
    // Apply LIS logic
    int ans = SpaceOpt (cuboids);
    return ans;
}
```

3

```
bool check (vector <int> a, vector <int> b) {  
    // Compare length, width & height  
    if (b[0] <= a[0] && b[1] <= a[1] &&  
        b[2] <= a[2])  
        return true;  
    else  
        return false;  
}
```

Changes in LIS code (Space optimization)

majorly changed. ↵

```
if (prev == -1 && check (nums[curr], nums[prev])  
{ // height to be added instead of 1  
    include = nums[curr][2] + nextRow[curr+1];  
}
```