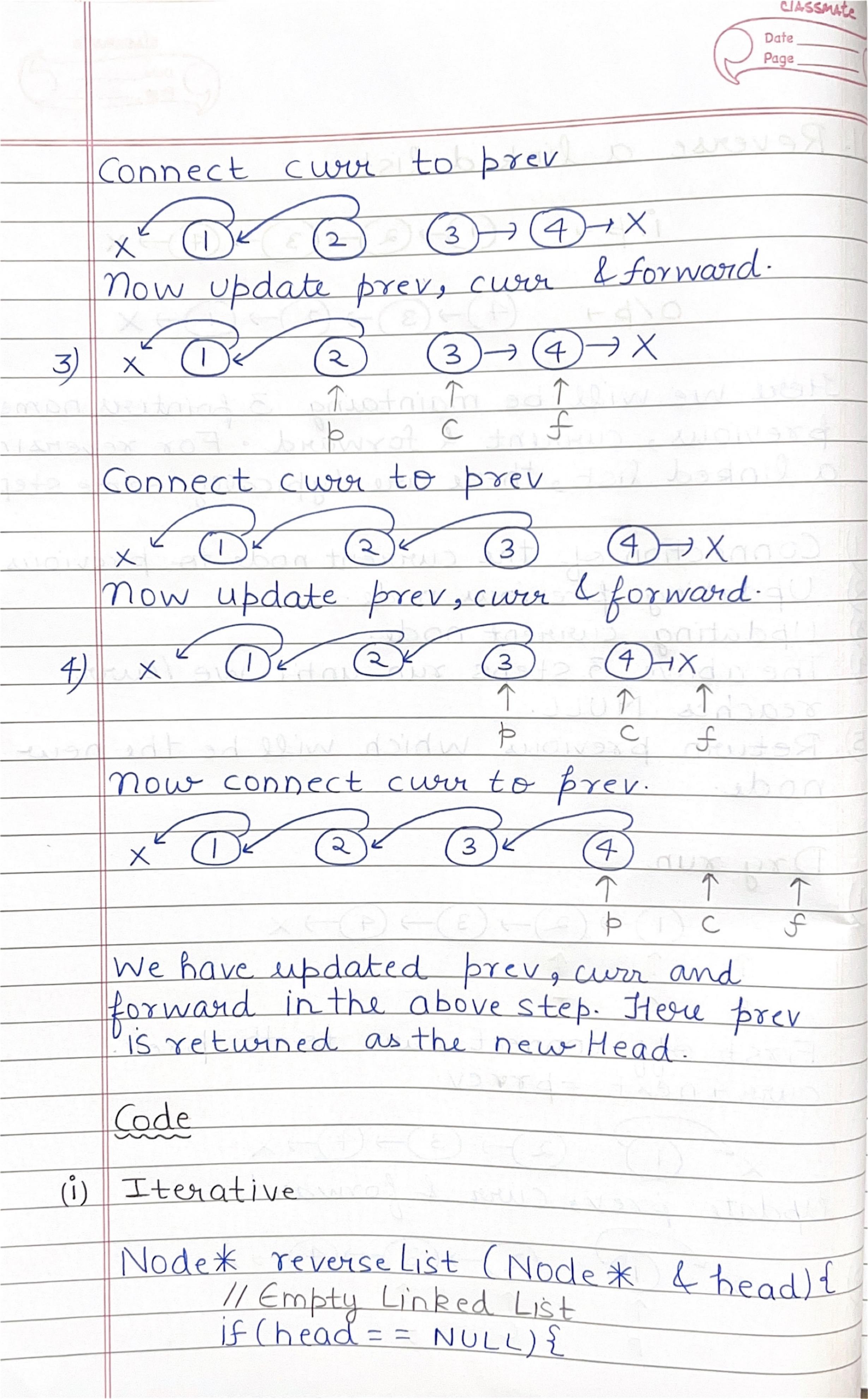
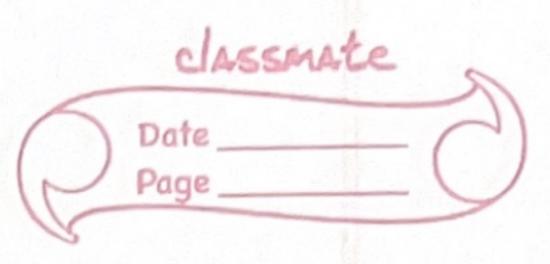


	Page
Q1.	Reverse a linked list.
	$i'/p \rightarrow (D \rightarrow Q) \rightarrow (3) \rightarrow (4) \rightarrow X$
	now whole body a formand.
	$0/b \rightarrow (4) \rightarrow (3) \rightarrow (2) \rightarrow (1) \rightarrow X$
	$X \leftarrow (A) \leftarrow (E) $
	Here we will be maintaing 3 pointers namely
	previous, current & forward. For reversing
	a linked list, there are typically some steps
1)	Connection of the current node to previous.
2)	Opdating previous node.
3)	Opdating current node.
4)	The above 3 steps run until we / curr
	reaches NULL.
5)	Return previous which will be the new
	node. Nord at mens tradad till
1)	Dry run
	$(1) \rightarrow (2) \rightarrow (3) \rightarrow (4) \rightarrow \times$
	P CT F
	Fixe + All connect curry to the breve
	First off connect curr to the prev. current = pr ev.
	χ^{2} (1) (2) (3) (4) (4)
	Update prev, curr & forward
2)	$\sqrt{(1)} (2) \rightarrow (3) \rightarrow (4) \rightarrow x$
	P C III F - Total Di





return head // Nothing to reverse Maintaing frev & cour pointers
Node* prev = NULL;
Node* cour = head; // Run loop until cur 1 = NULL While (com 1 = NULL) { 1/3 rd Bointer Node * forward = cwr + next i cwr + next = prev : // Connection prev = curi // Updating prev curr = forward : // updating curr getuin previ //new Head is prev. (ii) Iterative - V, Recursive implementation Node* reverselist (Node* & Brev, Node* 4 curr) { // Base case if (cur = = NULL) { return previ //new head 1/maintaing forward (I case solve) Node * forward = curr + next; curry-next = previ 1/ Recursive relation return reverse List (cover, forward) Time complexity = 0(n) Space complexity = O(n) -1 Recursive else O(1)