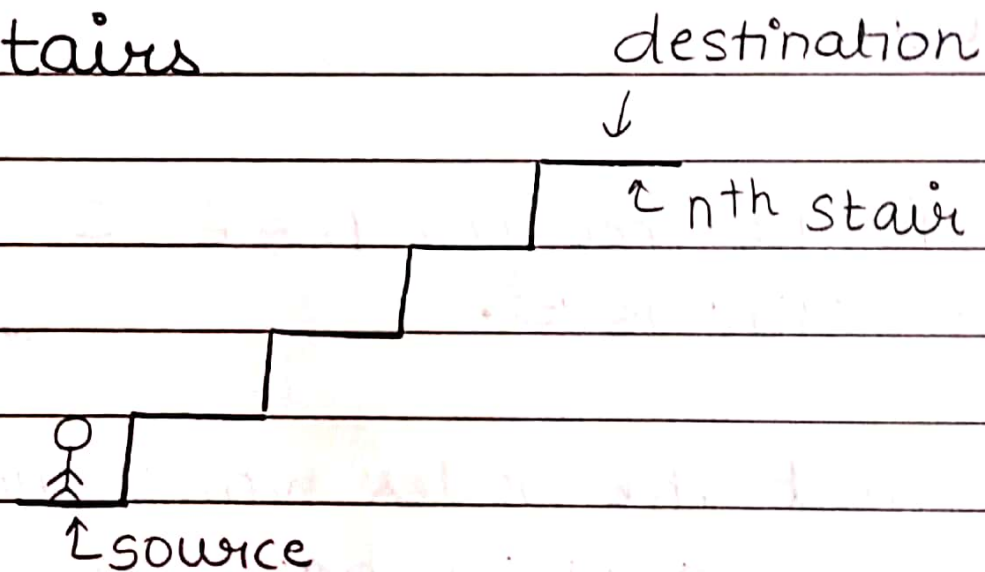


11/03/2023

Questions on Recursion

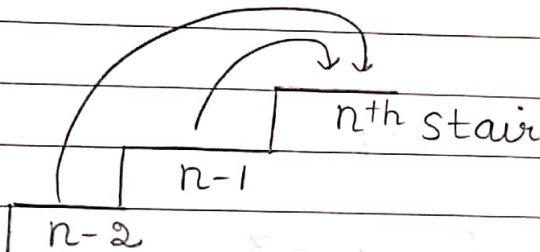
Q1 Climb Stairs



Steps allowed \Rightarrow move 1 stair at a time
move 2 stair at a time

Find the minimum number of steps to reach n th stair.

We are currently on the 0th stair. We just have to solve one case & rest recursion will handle.



To reach the n^{th} stair, we can reach from $(n-1)^{\text{th}}$ stair or $(n-2)^{\text{th}}$ stair.

$$\begin{aligned} n^{\text{th}} \text{ stair} &= (n-1)^{\text{th}} \text{ stair} + (n-2)^{\text{th}} \text{ stair} \\ f(n) &= f(n-1) + f(n-2) \end{aligned}$$

↳ As we have to find total ways

The recursive relation is same as that of the fibonacci series.

Code

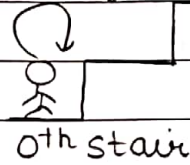
```
int climbStairs (int n) {
    // Base Case
    if (n == 0 || n == 1)
        return 1;
    // Recursive relation
    return climbStairs(n-1) + climbStairs(n-2);
}
```

3

Understanding base case

(only for $n=0$ case)

1 way



0th stair

$n=0$, 1 way

1 way



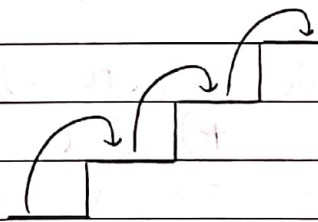
1st stair

$n=1$, 1 way

Dry run for $n=3$

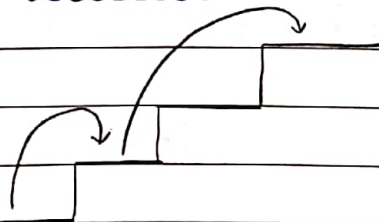
The answer for no. of ways to reach the 3rd stair is 3

1st solution



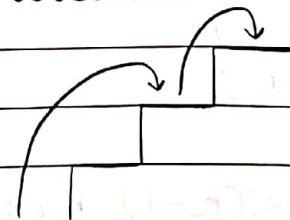
1st way
(1, 1, 1)

2nd solution



2nd way
(1, 2)

3rd solution



3rd way
(2, 1)

The most confusing thing in the question is the base case for 0th stair, that is only 1 way is there to reach 0th stair.

→ Time Limit Exceeded

Note → There will many TLEs in case of recursion which will further be optimized.

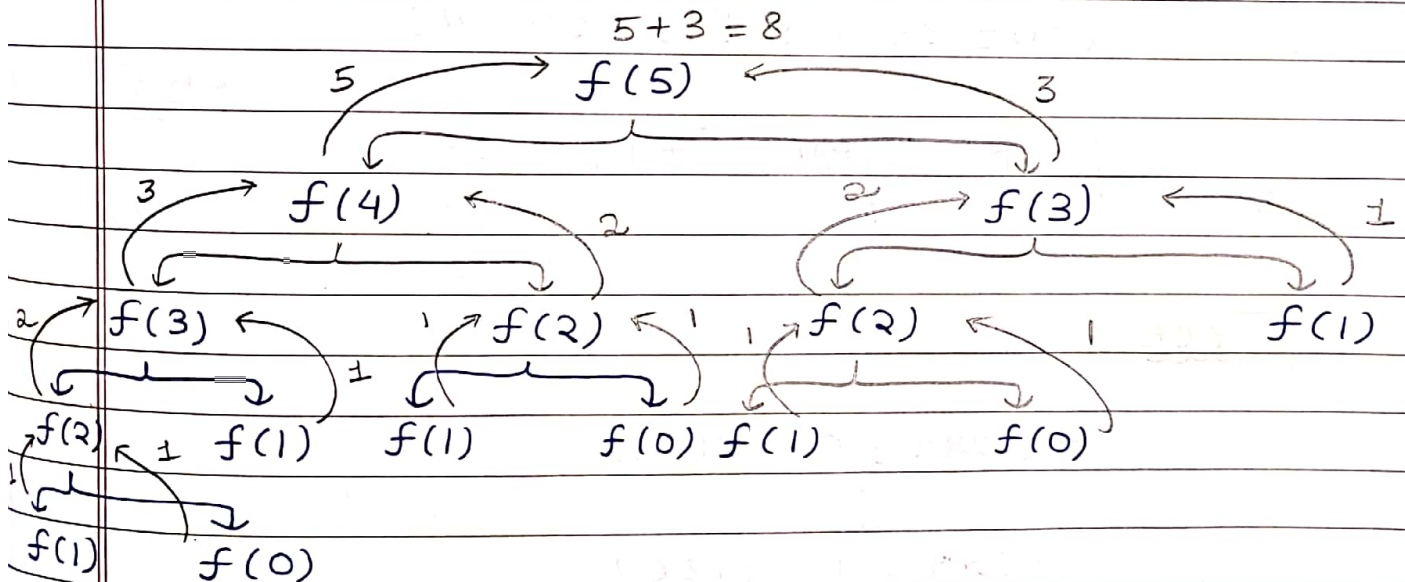
$$f(5) = f(4) + f(3)$$

$$\hookrightarrow f(3) + f(2)$$

$$\hookrightarrow f(2) + f(1)$$

$$\hookrightarrow f(1) + f(0)$$

Similarity for $f(3)$ in 1st line, calls will be sent. See the recursion tree below.



Hence to reach the 5th stair, there are 8 possible ways.

Note → $f(0) = 0$ is not selected as it means that we can't reach the 0th step/stair but we are already on 0th stair & hence we take it as 1.

Q2 Print the elements of the array.
We will simply take the base case as if index becomes equal to size of array, we need to return. After that simply print

the element & rest recursion will handle.

Code

```
void print (int arr [], int i, int size) {
    // Base Case
    if (i == size)
        return i;
    // Processing
    cout << arr [i] << " ";
    // Recursive Relation
    print (arr, i+1, size);
}
```

↳ move to next index

Tree

```
f (arr, 0, size)
    ↓ cout << arr [0];
f (arr, 1, size)
    ↓ cout << arr [1];
f (arr, 2, size)
    ↓ cout << arr [2];
    |
    | } goes on until base
    | } condition is achieved
```

Alternative code

```
void print (int arr [], int size) {
    // Base Case
    if (size == 0)
        return;
```


// Processing

cout << arr[0] << " ";

print (arr + 1, n - 1); // Recursive

3

Relation

Q3 Find maximum element in the array.

First create a variable maxi & initialize it with INT_MIN. Now as we start traversing the array, we will check that if the maxi is less than arr[i], if true then update maxi.

Code

```
void findMax (int arr[], int i, int size,
              int & maxi) {
```

```
    // Base Case  $\rightarrow$  initialized to INT_MIN
```

```
    if (i == size)
```

```
        return;
```

```
    // Processing
```

```
    if (arr[i] > maxi)
```

```
        maxi = arr[i];
```

```
    // Recursive Relation
```

```
    findMax (arr, i + 1, size, maxi);
```

3

The dry run of the above code is similar but instead of printing we will be comparing the maxi & arr[i].

Q4 Find minimum element in the array.

```
void findMin (int arr[], int i, int size,
              int & mini) {
```

```
    // Base case
```

```
    if (i == size)
```

```
         $\rightarrow$  initialized to INT_MAX
```

```

return i;
// Processing
if (mini > arr[i])
    mini = arr[i];
// Recursive Relation
findMin(arr, i+1, size, mini);
3

```

The approach is same as that of Q3 but the if condition has changed.

Q5 Find whether key is present in string or not.

i/p → lovebabbar , Key = 'r'
o/p → Yes

Will be done by recursion
 (lovebabbar)
 → Checking this only (1st case solve)

Base condition ⇒ If we have traversed whole string, then simply return false. If while traversing, we found the key then simply return true otherwise we move on to the next index.

Code

```

bool checkKey(string str, int i, int n,
               int key) {
    // Base condition
    if (i == n)
        return false;

```


// 1 case solve

if (str[i] == key)

return true;

// Recursive relation

return checkkey (str, i+1, n, key);

} not mandatory

↳ move to next index.

Note → By doing &, TLE can get eliminated. So here string can be passed as reference.

Here we can return index of the key in string. Just change bool type to int, false to -1 and true to i. If there are multiple occurrences of the key in the string, then we can store those indexes in vector.

Q6 Find the digits of the number.

i/p → 647

o/p → digits

6

4

7

$647 \% 10 = 7$ } 1st digit

$647 / 10 = 64$

$64 \% 10 = 4$ } 2nd digit

$64 / 10 = 6$

$6 \% 10 = 6$ } 3rd digit

$6 / 10 = 0 \rightarrow \text{stop}$

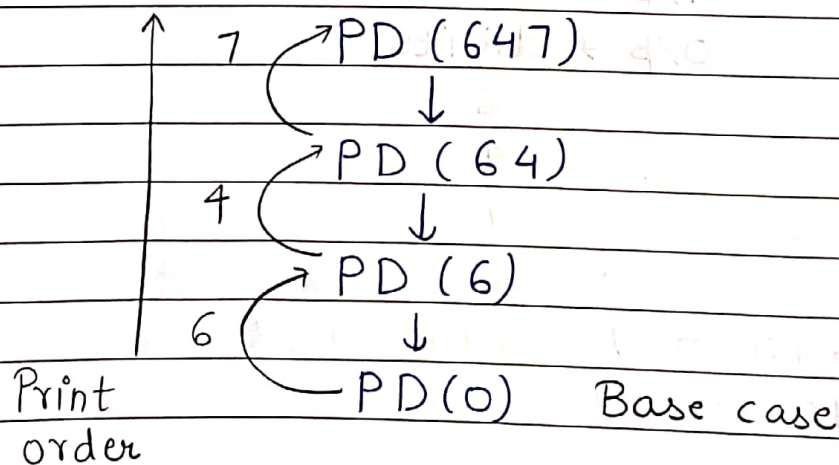
Base condition \Rightarrow When n becomes 0, we need to stop.

Code

```
void printDigits (int n) {
    // Base case
    if (n == 0)
        return;
    // Recursive Relation
    printDigits (n/10);
    // Processing
    int digit = n % 10;
    cout << digit << " ";
}
```

3

Tree



6 4 7 will be printed which are the digits of number 647.