	_ Classma	,
(160)	Date Page	1
1	3	

1	
	12/02/2023
	Till now we have studied about ID average.  Today we will be studying about the 2D
	Till now we have studied about the 21
	Today we will be studying about
3	arrays
	a see reply de una -81 distriction de la constante de la const
	2D Arrays
7	72 51 5 6 73
	int our [] = {1,2,5,73;
	This is how ID arrays are stored in the
	memory. But the question is what happens
-	in case of 2D arrays. We can relate 2D arrays with Tic-tac-toe game which have
	9 cells with 11c - tac - tol game which have
	J Cells. W
=	row o
	70W-1
. ,	70W2
j /	colo Coll Col2
Franci	A grid is formed where there are multiple
1111	rows and multiple column. It
×2 [	can be formed by taking 3,1D arrays of
	Size 3 each.
	3,1Darrays

Scarnieu with Cam

	But this is a small axid II II
	But this is a small grid, that's why we created ID average but what if there are
	loso rows & it is not
	1000 ID arrays. Hence 2D arrays make our
	Thame Thorof columns
	UL COOL (1000) 1 10001.
	no of rows
	data type
	Also we can mention different no of rows
	a columns.
	Memory representation of 2D avviays
	int avor [2] [2]; - Total elements = 2x2=4
<u> </u>	over → 1 2 3 4 → 2D averay
	Matrix was just used to visualize the 2D
	arrays but in memory, the ID array
	représentation is used.
	colo col1
	rowo 12 G Just way of
	rows 3 4 J visualizing 2D arrays
	Accessing elements of 2D arrays
	In the above 2D array, say we want to access 3. We can access it by arr [1][0]
	access 3. We can access it by arr [1][0]
	or [0] [0] → 1
	our [0] [1] → 2 rowindex index
	our [1] [1] → 4
/	Formulas for mapping in 1D array representat

Scarined with Cam

	162) Date
	Page
	is C* i + j
	C 7 no · of columns
	i - row-index
	ji column - indez
	and publication of the second
	arr (0)[1] = 2
	Formulal = 2 x0+1 = 0+1 = 1st index in
	1D array.
	0 1 2 3
	THE ROTTE OF THE PROPERTY OF T
	2 is present at 1st index in 1D array
1	aur [1] [1] = 4
	Formulae = 2x1+1=2+1=3rd index, 4 is
(1	Stored in 1D avau.
	Storta III 1 D wordy.
Note	> over [0] [1] + ovr [1], It is just a way
	that in memory ID array is used but the
	actual interface is the matrix only so
ka Jan Pal	use arr [i][j] to access the elements of
	2Darray.
	Pulare Of Ja sansmore Janes
	Initialization of 2D array
	TUBLE STELL WAS ELEVANOR OF THE STELL OF THE
	int aver [2][2] = { {1,23, {3,433;
	row-o row-IL
	$\gamma_{0W-0} \rightarrow 1$ 2
	row-1 → 3 4
Logica	COI-O COI-I
	C01~0 C01-1

<del>o</del>carineu wiiir€am

		Date Page
	Input and Output in 2D Array Input → cin>> aur [i][j]; Output → cout << aur (i)[j];	
	We always have to play within 0 to where n - rows and columns = n.	n-1 index
X0:	e - avr [i][j]  No of rows = n  No of columns = m	
	j∈[0,n-1] ? Never go out j∈[0,m-1] ]	this range
	Accessing elements row-wise int over [3][3] = { {1,2,35, {4,5,63}}	, २७,४,९३};
y	for (int i=0; i<3; i++) {     for (int j=0; j<3; j++) {         cout << aur [i][j] <<	
	cout << endl;	2 3
	3	5 6 8 9 1
	<u>4 5 6</u> 7 8 9	71
	Accessing elements column-wise for (int i=0;i<3; i++){	

	CLASSA
(164)	Date Page
P	rage

for (int j = 0; j < 3; j + +) {
Cout << arc 13 c  3  Cout << endl;  Coumn wise)  3  Output  1 4 7  2 5 8  3 6 9  Similarly we can take row-wise input & column - wise input.  Problem Solving  1) Row sum print  1/p 7 1 2 3 4 5 6
Cout << endl; (Swapped in case of coumn wise)  3  Output  1 4 7  2 5 8  3 6 9  Similarly we can take row-wise input & column - wise input:  Problem Solving  1) Row sum print  1/p 7 1 2 3  4 5 6
Output  1 4 7  2 5 8  3 6 9  Similarly we can take row-wise input & column - wise input:  Problem Solving  1) Row sum print  1/p 7 1 2 3  4 5 6
Output  1 4 7 2 5 8 3 6 9  Similarly we can take row-wise input & column - wise input:  Problem Solving  1) Row sum print  1/p -> 1/p -
J 4 7  2 5 8  3 6 9  Similarly we can take row-wise input & column - wise input:  Problem Solving  1) Row sum print  1/p 7   1 2 3   4 5 6
Similarly we can take row-wise input & column - wise input.  Problem Solving  1) Row sum print  1/p -> 1 2 3 4 5 6
Similarly we can take row-wise input & column - wise input.  Problem Solving  1) Row sum print  1/p -> 1 2 3 4 5 6
Problem Solving  1) Row sum print  1/p -> 1 2 3  4 5 6
Problem Solving  1) Row sum print  1/p -> 1 2 3  4 5 6
1) Row sum print  1/p -> 1 2 3  4 5 6
1) Row sum print  1/p -> 1 2 3  4 5 6
1/p → 1 2 3 4 5 6
4 5 6
- Let IL 2 port 28 many
0/6-16
24
then print We have to simply do your
Wise traversal.
Code
Void print Row Sum (int aver [][3], int rows
int cols) {

Just cols) {

	Page
	for (int i=0) i < rows; 1++){
	int sum = 0; - After each row, sum = 0
	for (int j = 0 ; j < cols ; j + +) {
	Sum = sum + any [i][i]i
	3
	cout << "Row" << i << "sum is"
	< <sum <<="" endl;<="" th=""></sum>
	3
	^
Not	e → While passing 2D avorage in function,
	We need to mention the bound of
	column. It is a rule Formulae of
	Column. It is a rule Formulae of C* i + j is the reason behind this
2)	Column sum print
	en net en
	1/p-1 Same array as that of Ques-1
	0/6-1 12
	15 Delvi. Arador 1
	18
	We just have to do column - wise traversal:
	& print the sum just like we did in row
	wise sum print
	(2.9E 12 ) 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
	In the above code just change
	sym = sym + our [i][j] to
	sum = sum + aur [j](i)
·	
_3)	Linear search in 2D array
	1/p→ Same array as that of ques-1
	element = 3
	0/p+ True

1.11

<del>Scarmed with C</del>am

	Class	
(166)	Date	
	Page	
3		

	rice traversal a
	We can do either row-wise traversal on
	We can do either row-Wise Traversal & comparu each column -wise traversal & that is given to
	column - wise traversal & composition to element with the element is matched
	Af element is matched
	element with the element is matched us as in i/p. If element is matched
	us as in i/p. If element is more if not retwin true else retwin false if not found even after traversing full away.
	tound even after.
	Code
	bool 52537 int voint c
	find key (int over [][3], int y, int c,
	for (int i=0; i< x; i++) {
	C = C(x+1=0) + C(x+1) + C(x+
	if ( our [I][] = = R)?
	return true;
	Taronto de dont a granda sono la la la
	2
	3 return false;
	ξ
4	
4)	Maximum and minimum in 2D averay.
	Maximum and minimum at a b working.
	i/b > Cama ayyay as that al ayya
	1/p→ Same array as that of ques-1
	$O/p \rightarrow Maximum = 9$
	Minimum = 1
	Initialize maxi° = INT_MIN & mini = INT-H
	and compare maxi & mini with each element
S 1/2	the state of the s
<b>/</b> /	maxi (arr [1]) + update maxi
1	maxi (arr[i][j] - update maxi mini > arr[i][j] - update mini
11	Scarneu with G

L\
, p
\h.
-1
••,
-7/
,,,
,,
- 4,,
·· <sub>\</sub>
-· <sub>V</sub> .
·\

	14. x01/5 into
	We just have to change the rows into the column we need to per form the swap on
	We just have I ber form the swap on
	Column · We need To poiss
	Swap (aur [i](j], aur (j)(i));
	Shoop (any [i] (j), worth
	But the inner for loop won't run fully as some elements will be swapped twice?
	would reach to the original position and
	hence the matrix would be same.
	Code
	$\frac{1}{1}$
	void transpose (int avu [][3], int r, intc){-
	for (int i=0; i< x; i++){
	for (int $j = 0$ ; $j < l$ ; $j + +$ ) {
	swap (ave [i][j], ave [j][i]
	24 24 24 24 24 24 24 24 24 24 24 24 24 2
	3
	3
	J
	Better way to use 2D arrays
	Better way to use 2D arrays  ID array - vector <int> array :  2D array - We will be using vector of vector concept</int>
	2D array - We will be using
	vector concept vector of
	→ Vecha:
	> vector of vector having
	3 cells
Ø	
	And the second of the second o
-17/4	
13	

Scarnieu with Cam