

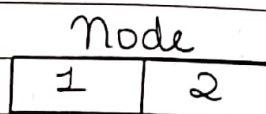
8/04/2023

## linked list

- 1) Suppose that in our OS we need to make an array of 10 MB size but the memory available is 10 MB but is not contiguous & hence we can't use arrays here but linked list can work here & hence it can work on non-contiguous memory locations.
- 2) At run-time, we can do creation & deletion of node at run-time and hence there would be no wastage of memory which was there in case of arrays.
- 3) In arrays, insertion of element takes  $O(n)$  time complexity but in linked list insertion can be done in  $O(1)$  time complexity provided the pointer is at that position only.
- 4) There is no concept of indexing in case of linked list but concept of address is used here.

### Definition

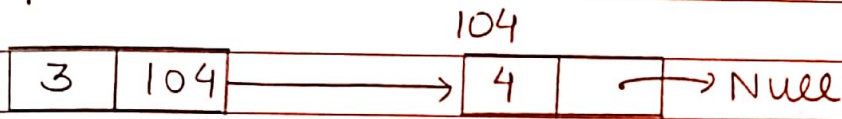
linked list can be defined as collection of nodes.



1 → data

2 → address of next node

## Simple linked list



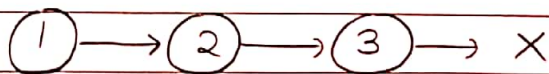
```

class Node {
    int data ;
    Node* next ;
}
  
```

As we create an integer pointer by `int *`, we create pointer to Node by `Node *`.

## Types of linked list

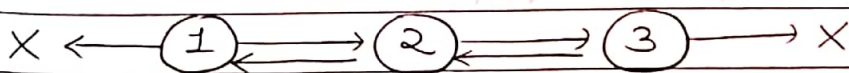
### 1) Singly linked list



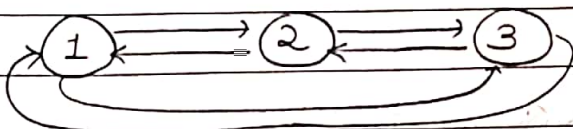
### 2) Circular linked list



### 3) Doubly linked list (prev & next pointers)



### 4) Doubly circular linked list (prev & next pointers)



Note - Linked list is Hindi & this is a magical line but don't tell this to any interviewer. Also linked list is a linear data structure as only one descendant is there.



Creation of Linked List (Better method is also there)

```
class Node {  
    public :  
        int data ;  
        Node* next ;  
        Node() {  
            this->data = 0 ;  
            this->next = NULL ;  
        }  
        Node (int data) {  
            this->data = data ;  
            this->next = NULL ;  
        }  
};
```

```
main() {  
    Node* first = new Node(1);  
    Node* second = new Node(2);  
    Node* third = new Node(3);  
  
    first->next = second;  
    second->next = third;  
}
```

① → ② → ③ → NULL

Printing a linked list

- 1) Print the data of current node.
- 2) Move pointer forward.
- 3) Stop when we reach null.

```
void print (Node* & head) {
```

```
    // Temporary pointer
```

```
    Node* temp = head;
```

```
    while (temp != NULL) { // Step-3
```

```
        // Step-1
```

```
        cout << temp->data << " ";
```

```
    // Step-2    temp = temp->next;
```

```
    }
```

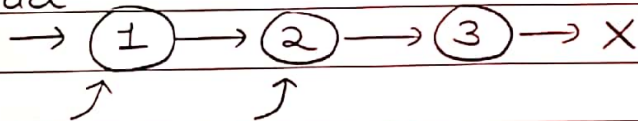
```
}
```

Output

1      2      3

Meaning of  $temp = temp \rightarrow next$

head



temp      temp → next

```
temp = temp → next
      = head → next
```

Hence we are moving forward in the linked list.

Why we created temp?

It is a good practice not to change the head and we created temp which is pointing to head & hence we move temp forward leaving head at same place.

Ex →

```

    graph LR
      temp --> 1((1))
      1 --> 2((2))
      2 --> 3((3))
      3 --> X[X]
  
```

temp → next → data = 2

temp → next → next → data = 3



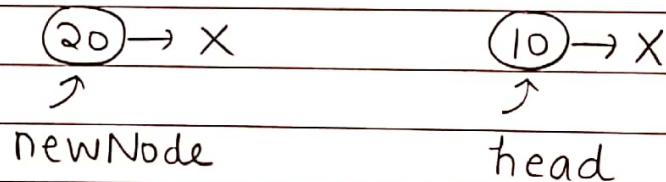
## Better way of creating linked list

### (i) Insertion at head

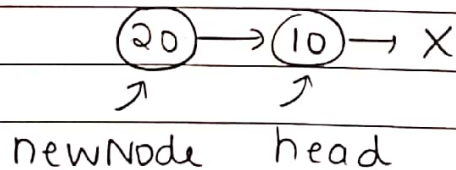
Initial scenario  $\rightarrow (10) \rightarrow X$

Inserting 20 at the head

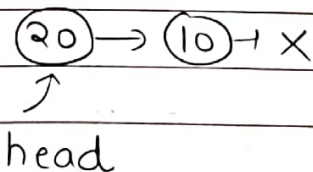
#### 1) Creating node with value 20



#### 2) newNode to be connected to head



#### 3) Update head

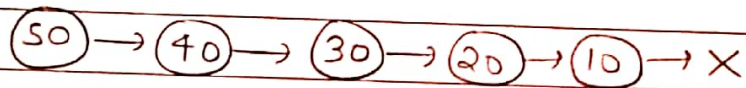


```

void insert (Node* &head, int data) {
    //Step-1
    Node* newNode = new Node(data);
    //Step-2
    newNode->next = head;
    //Step-3
    head = newNode;
}
  
```

Suppose now we inserted 20, 30, 40 & 50

using the function, then the linked list would become.



Printing linked list will give the output as 50 40 30 20 10

Note  $\rightarrow$  head  $\Rightarrow$  starting point of linked list  
tail  $\Rightarrow$  ending point of linked list

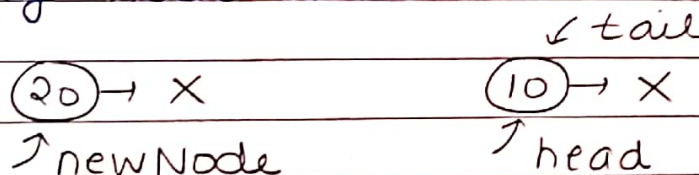
point = node here.

## (ii) Insertion at tail

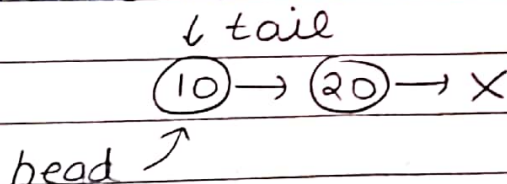
Initial scenario  $\rightarrow (10) \rightarrow X$

Inserting 20 at tail

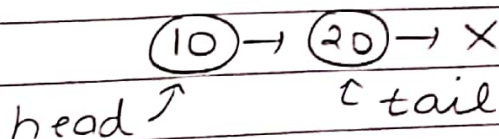
1) Creating node with value 20



2) newNode to be connected to the tail



3) Update tail



```
void insert (Node* &head, Node* &tail,
             int data) {
```



// Step-1

Node\* newNode = new Node(data);

// Step-2

tail->next = newNode;

// Step-3

tail = newNode;

}

Note → If initially head & tail was initialized to NULL, then we need to handle this explicitly as when we are creating the new node, initialize the head & tail with that node & this would be when we are creating the 1st node.

if (head == NULL) {

head = newNode;

tail = newNode;

}

} 1st node creation

Better way & easy to understand for handling empty linked list case.

if (head == NULL) {

Node\* newNode = new Node(Data);

head = newNode;

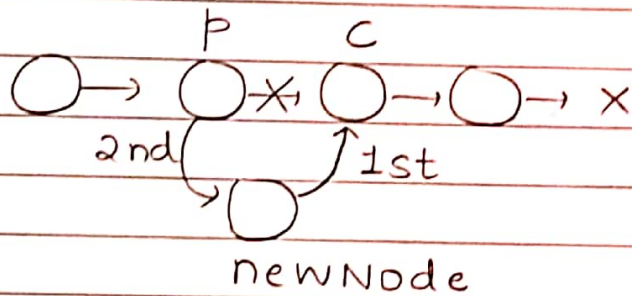
tail = newNode;

return; // As node has been created

}

(iii) Inserting at specified position.

- 1) Check for empty linked list case.
- 2) If non-empty then
  - (i) Traverse to that position.
  - (ii) Create a node



- (iii)  $\text{newNode} \rightarrow \text{next} = \text{C};$
- (iv)  $\text{p} \rightarrow \text{next} = \text{newNode};$

What if we do  $\text{p} \rightarrow \text{next} = \text{newNode}$  first and then update  $\text{newNode} \rightarrow \text{next}$  to  $\text{curr}$  but by this, we will be losing the track of further linked list.

Some edge cases

- 1)  $\text{pos} == 0 \rightarrow \text{insert At Head}$
  - 2)  $\text{pos} == \text{len} \rightarrow \text{insert At Tail}$
- } Reusing func<sup>n</sup>.

Length of linked list

```
int findLength (Node* &head) {
    int len = 1;
    Node* temp = head;
    while (temp->next != NULL) {
        temp = temp->next;
        len++;
    }
    return len;
}
```

}



## Deletion operation in Linked list

### (i) Deleting head of linked list

$(10) \rightarrow (20) \rightarrow (30) \rightarrow X$        $temp = head$   
 $h$

- 1)  $head = head \rightarrow next$
- 2)  $temp \rightarrow next = null$  ;
- 3) delete temp

$(20) \rightarrow (30) \rightarrow X$

### (ii) Deleting tail in linked list

$(10) \rightarrow (20) \rightarrow (30) \rightarrow X$        $temp = tail$   
 $\downarrow prev$

- 1) Find prev
- 2)  $prev \rightarrow next = null$
- 3)  $tail = prev$
- 4) delete temp

$(10) \rightarrow (20) \rightarrow X$

```
void deleteQ(int pos, Node* & head,
             Node* & tail) {
```

```
    // Empty linked list
```

```
    if (head == NULL) {
```

```
        return ;
```

```
    }
```

```
    // Delete head
```

```
    if (pos == 1)
```

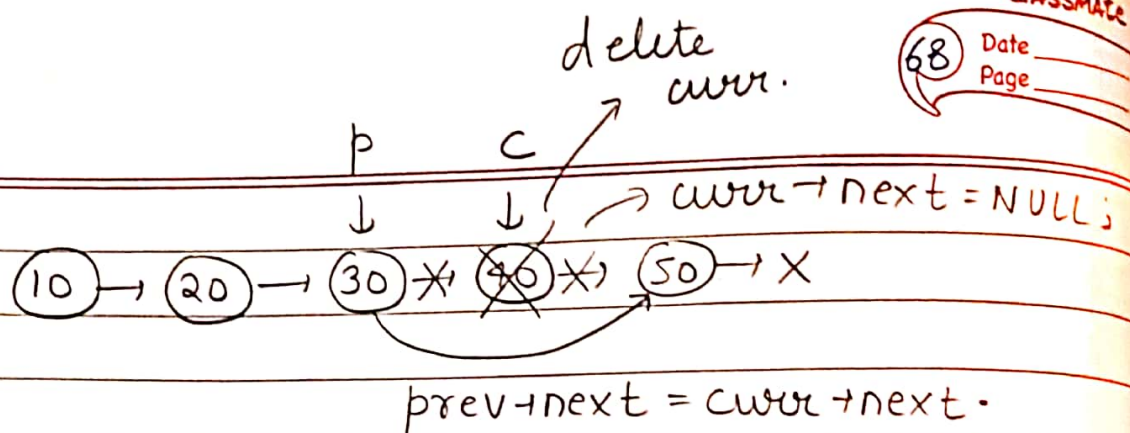
```
{  
    Node* temp = head;  
    head = head->next; //S-1  
    temp->next = NULL; //S-2  
    delete temp; //S-3  
}
```

```
// Deleting tail  
int len = findLength(head);  
if (pos == len) {  
    // Find prev (S-1)  
    int i = 1;  
    Node* prev = head;  
  
    while (i < pos - 1) {  
        prev = prev->next;  
        i++;  
    }  
  
    //S-2  
    prev->next = NULL;  
    //S-3  
    tail = prev;  
    //S-4  
    delete temp;  
}
```

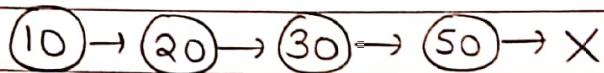
(iii) Deleting middle node (specified position)

- 1) Find prev & curr.
- 2) Do prev->next = curr->next;
- 3) curr->next = NULL;
- 4) delete curr.





Updated linked list



```
int i = 1;
```

```
Node* prev = head;
```

```
// Finding prev
```

```
while (i < pos - 1) {
```

```
    prev = prev → next;
```

```
    i++;
```

```
}
```

```
// Finding curr
```

```
Node* curr = prev → next;
```

```
// Step-2
```

```
prev → next = curr → next;
```

```
// Step-3
```

```
curr → next = NULL;
```

```
// Step-4
```

```
delete curr;
```

```
}
```