AAA528: Computational Logic

Lecture 2 — SAT/SMT Solvers

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The Z3 SMT Solver

A popular SMT solver from Microsoft Research:

https://github.com/Z3Prover/z3

- Supported theories:
 - Propositional Logic
 - Theory of Equality
 - Uninterpreted Functions
 - Arithmetic
 - Arrays
 - Bit-vectors, ...
- References
 - Z3 Guide https://rise4fun.com/z3/tutorialcontent/guide
 - ➤ Z3 API in Python
 http://ericpony.github.io/z3py-tutorial/guide-examples.htm

Propositional Logic

```
1 p = Bool('p')
2 q = Bool('q')
3 r = Bool('r')
4 solve(Implies(p, q), r == Not(q), Or(Not(p), r))
```

```
[q = False, p = False, r = True]
```

Arithmetic

```
1 from z3 import *
2
3 x = Int('x')
4 y = Int('y')
5 solve (x > 2, y < 10, x + 2*y == 7)
6
7 x = Real('x')
8 y = Real('y')
9 solve(x**2 + y**2 > 3, x**3 + y < 5)</pre>
```

```
$ python test.py
[y = 0, x = 7]
[x = 1/8, y = 2]
```

BitVectors

no solution [x = 6]

```
1 x = BitVec('x', 32)

2 y = BitVec('y', 32)

3

4 solve(x & y = ~y)

5 solve(x >> 2 = 3)

6 solve(x << 2 = 3)

7 solve(x << 2 = 24)

[y = 4294967295, x = 0]

[x = 12]
```

Uninterpreted Functions

```
1 \times = Int('x')
y = Int('y')
f = Function('f', IntSort(), IntSort())
s = Solver()
6 \text{ s.add}(f(f(x)) = x, f(x) = y, x != y)
8 print s.check()
10 m = s.model()
11 print m
print "f(f(x)) =", m.evaluate(f(f(x)))
print "f(x) = ", m.evaluate(f(x))
 sat
 [x = 0, y = 1, f = [0 \rightarrow 1, 1 \rightarrow 0, else \rightarrow 1]]
 f(f(x)) = 0
```

f(x) = 1

Constraint Generation with Python List

```
1 X = [ Int('x%s' % i) for i in range(5) ]
2 Y = [ Int('y%s' % i) for i in range(5) ]
3 print X, Y
4 X_plus_Y = [ X[i] + Y[i] for i in range(5) ]
5 X_gt_Y = [ X[i] > Y[i] for i in range(5) ]
6 print X_plus_Y
7 print X_gt_Y
8 a = And(X_gt_Y)
9 print a
10 solve(a)
```

```
[x0, x1, x2, x3, x4] [y0, y1, y2, y3, y4]

[x0 + y0, x1 + y1, x2 + y2, x3 + y3, x4 + y4]

[x0 > y0, x1 > y1, x2 > y2, x3 > y3, x4 > y4]

And(x0 > y0, x1 > y1, x2 > y2, x3 > y3, x4 > y4)

[y4 = 0, x4 = 1, y3 = 0, x3 = 1, y2 = 0, x2 = 1, y1 = 0, x1 = 1, y0 = 0, x0 = 1]
```

Problem 1: Program Equivalence

Consider the two code fragments.

```
if (!a&&!b) then h
else if (!a) then g else f
if (a) then f
else if (b) then g else h
```

The latter might have been generated from an optimizing compiler. We would like to prove that the two programs are equivalent.

Encoding in Propositional Logic

The if-then-else construct can be replaced by a PL formula as follows:

if
$$x$$
 then y else $z \equiv (x \wedge y) \vee (\neg x \wedge z)$

The problem of checking the equivalence is to check the validity of the formula:

$$F: (\neg a \wedge \neg b) \wedge h \vee \neg (\neg a \wedge \neg b) \wedge (\neg a \wedge g \vee a \wedge f) \\ \iff a \wedge f \vee \neg a \wedge (b \wedge g \vee \neg b \wedge h)$$

If $\neg F$ is unsatisfiable, the two expressions are equivalent. Write a Python program that checks the validity of the formula F.

Problem 2: Seat Assignment

Consider three persons A, B, and C who need to be seated in a row. There are three constraints:

- A does not want to sit next to C
- A does not want to sit in the leftmost chair
- B does not want to sit to the right of C

We would like to check if there is a seat assignment for the three persons that satisfies the above constraints.

Encoding in Propositional Logic

To encode the problem, let X_{ij} be boolean variables such that

$$X_{ij} \iff \mathsf{person}\; i \; \mathsf{seats} \; \mathsf{in} \; \mathsf{chair}\; j$$

We need to encode two types of constraints.

- Valid assignments:
 - Every person is seated

$$igwedge_i igwedge_j X_{ij}$$

Every seat is occupied

$$\bigwedge_{j}\bigvee_{i}X_{ij}$$

One person per seat

$$\bigwedge_{i,j}(X_{ij}\implies \bigwedge_{k\neq j}\neg X_{ik})$$

Encoding in Propositional Logic

- Problem constraints:
 - A does not want to sit next to C:

$$(X_{00} \implies \neg X_{21}) \wedge (X_{01} \implies (\neg X_{20} \wedge \neg X_{22})) \wedge (X_{02} \implies \neg X_{21})$$

A does not want to sit in the leftmost chair

$$\neg X_{00}$$

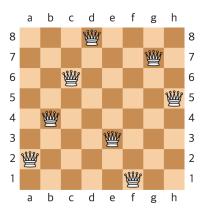
B does not want to sit to the right of C

$$(X_{20} \implies \neg X_{11}) \land (X_{21} \implies \neg X_{12})$$

Write a Python program that solves the problem.

Problem 3: Eight Queens

The eight queens puzzle is the problem of placing eight chess queens on an 8x8 chessboard so that no two queens attack each other. Thus, a solution requires that no two queens share the same row, column, or diagonal.



Encoding

Define boolean variables Q_i as follows:

 Q_i : the column position of the queen in row i

• Each queen is in a column $\{1,\ldots,8\}$:

$$igwedge_{i=1}^{8} 1 \leq Q_i \wedge Q_i \leq 8$$

• No queens share the same column:

$$\bigwedge_{i=1}^{8} \bigwedge_{j=1}^{8} (i \neq j \implies Q_i \neq Q_j)$$

• No queens share the same diagonal:

$$igwedge^8 igwedge^i_{i=1} igwedge^i_{j=1} (i
eq j \implies Q_i - Q_j
eq i - j \wedge Q_i - Q_j
eq j - i)$$

In Python

```
1 from z3 import *
3 def print_board (r):
   for i in range(8):
4
        for j in range(8):
            if r[i] = j+1:
6
                 print 1.
             else:
8
                 print 0,
        print ""
10
Q = [Int ("Q-\%i" \% (i+1)) for i in range(8)]
14 val_c = [ And (1 <= Q[i], Q[i] <= 8) for i in range(8) ]
col_c = [Implies (i \Leftrightarrow j, Q[i] \Leftrightarrow Q[j]) for i in range(8)
      for j in range(8) ]
diag_c = [Implies (i \Leftrightarrow j, And (Q[i]-Q[j] != i-j, Q[i]-Q[j])]
     != i-i) for i in range(8) for j in range(i)
```

In Python

```
s = Solver()
s.add (val_c + col_c + diag_c)
res = s.check()
if res == sat:
    m = s.model ()
    r = [ m.evaluate (Q[i]) for i in range(8) ]
    print_board (r)
    print ""
```

Exercise: Finding All Solutions

There are multiple solutions to the eight queens problem. For example, the following can also be a solution:

How many different solutions can you find? Write a Python program that finds all solutions of the problem.