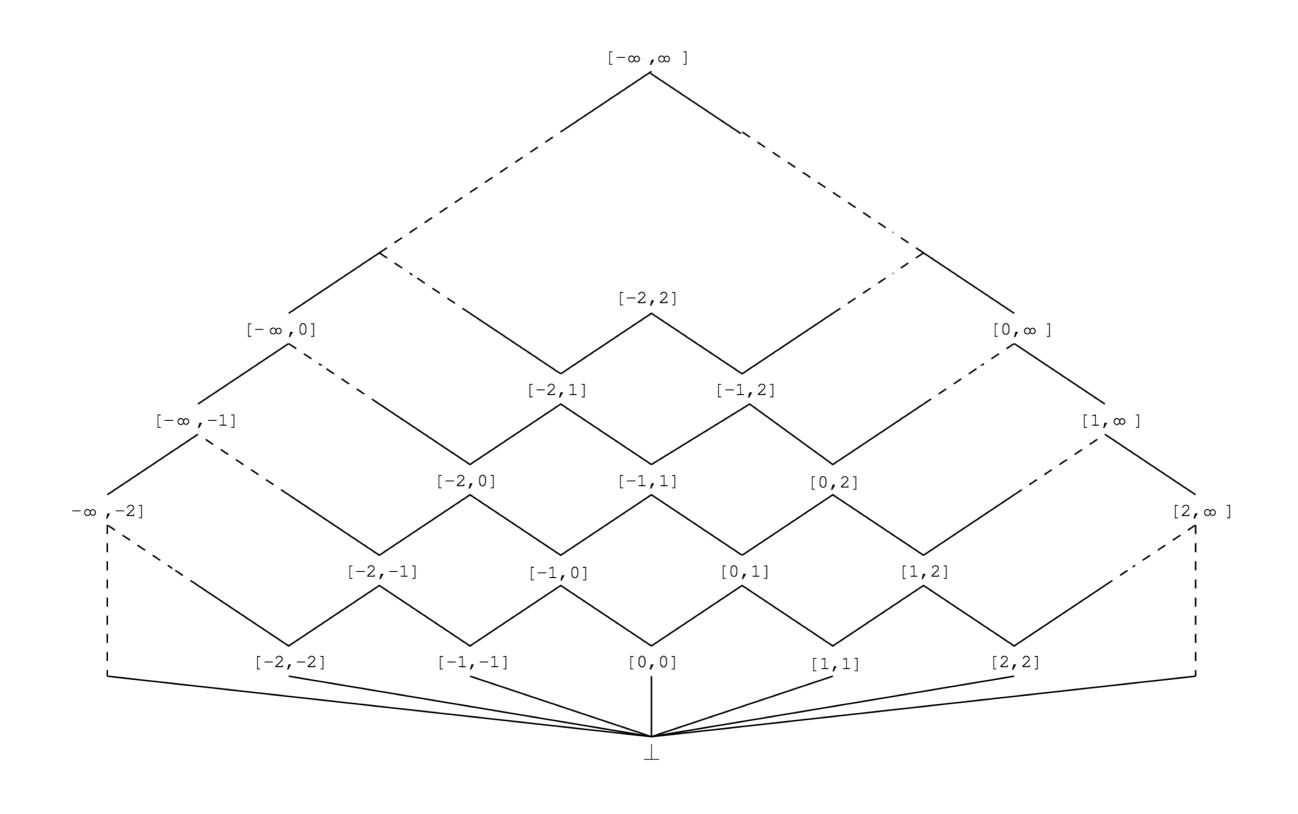
#### **AAA528: Computational Logic**

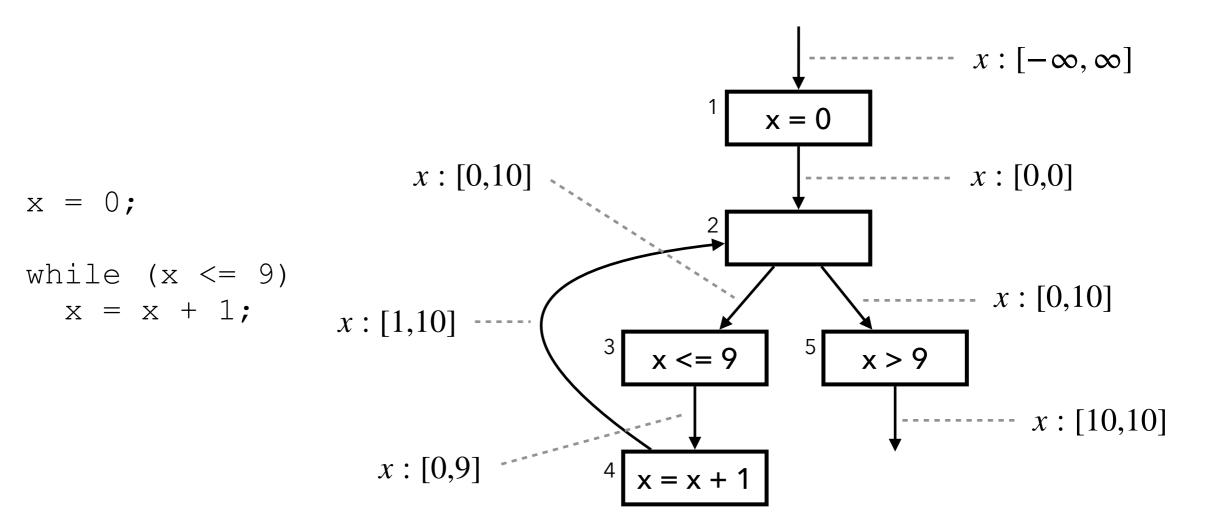
#### Lecture 11 – Static Analysis Examples

Hakjoo Oh 2025 Spring

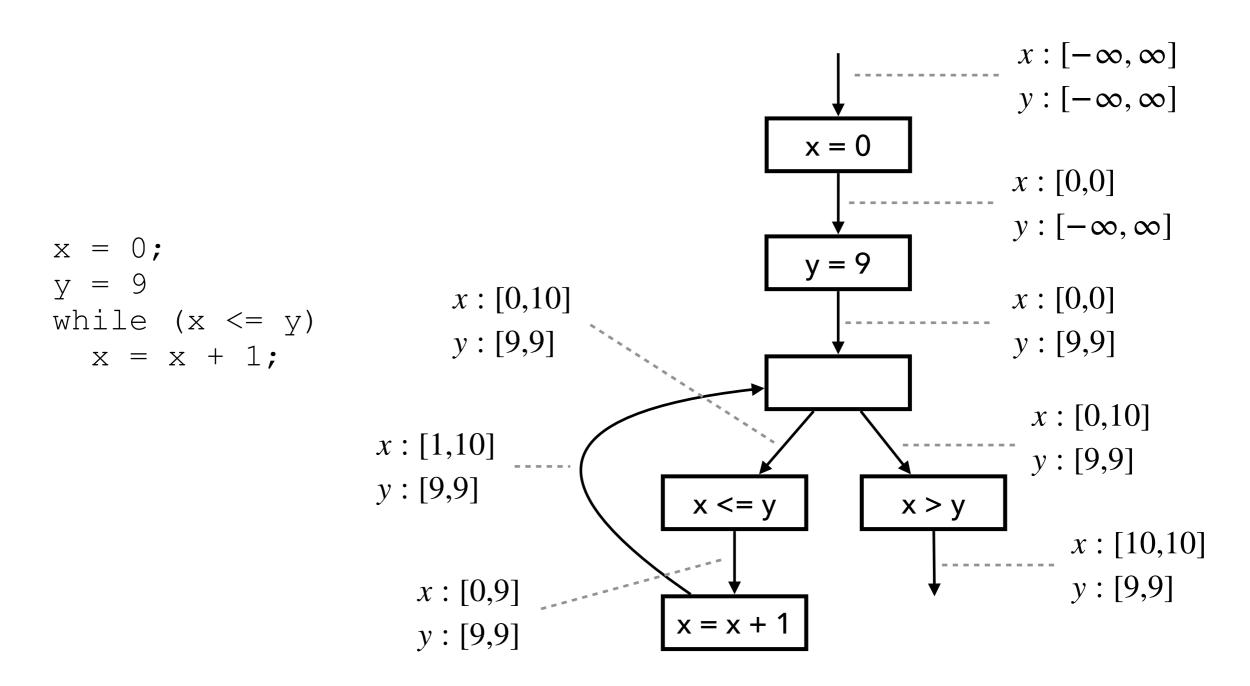
#### The Interval Domain

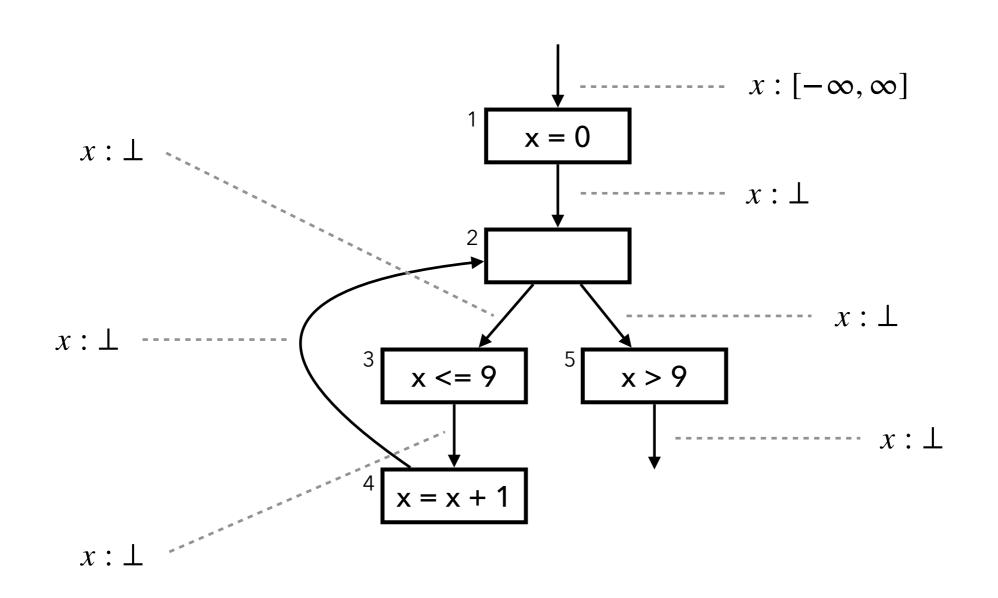


# **Example Program**

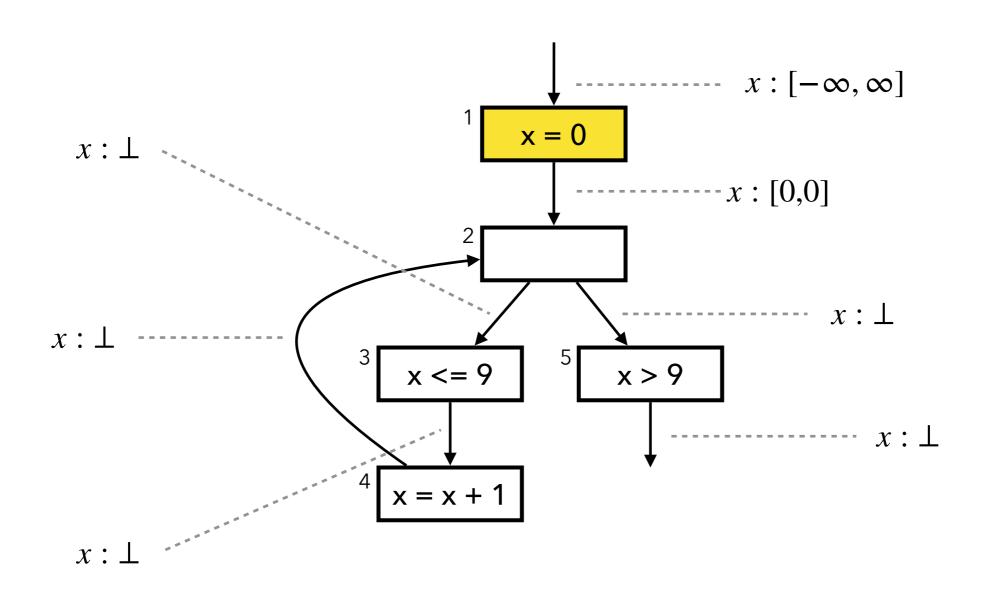


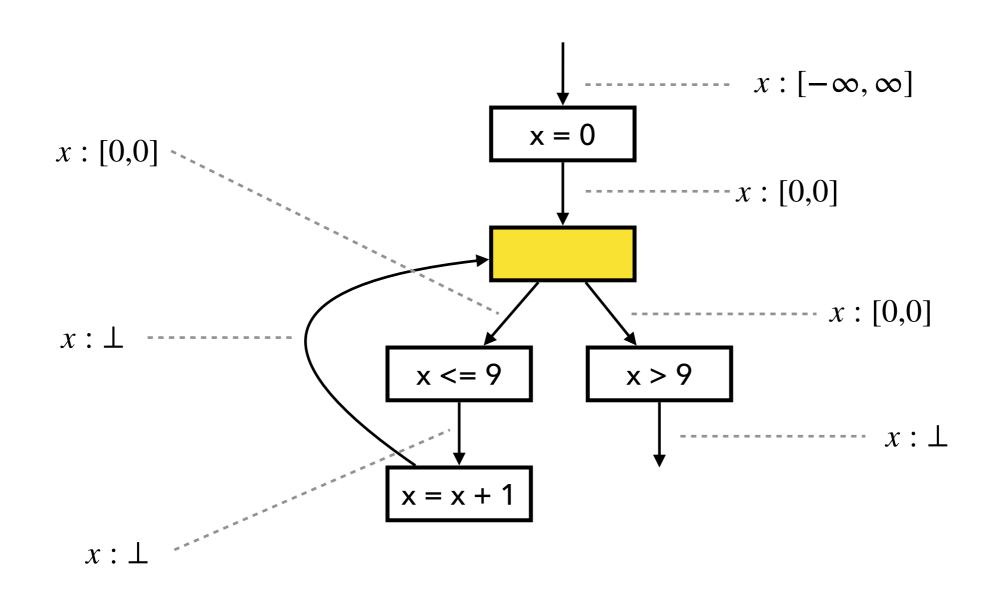
### cf. Multiple Variables



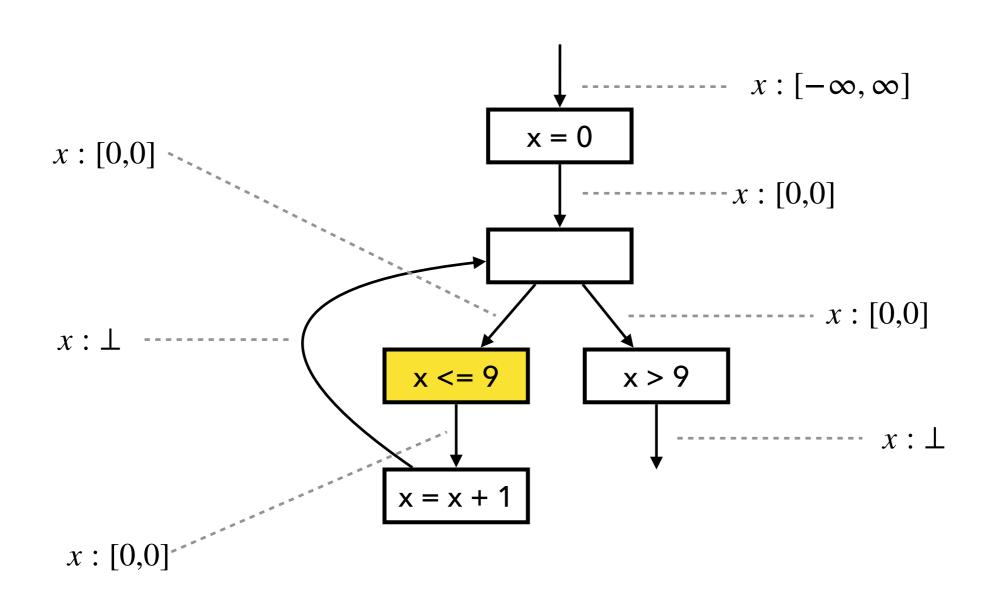


Initial states

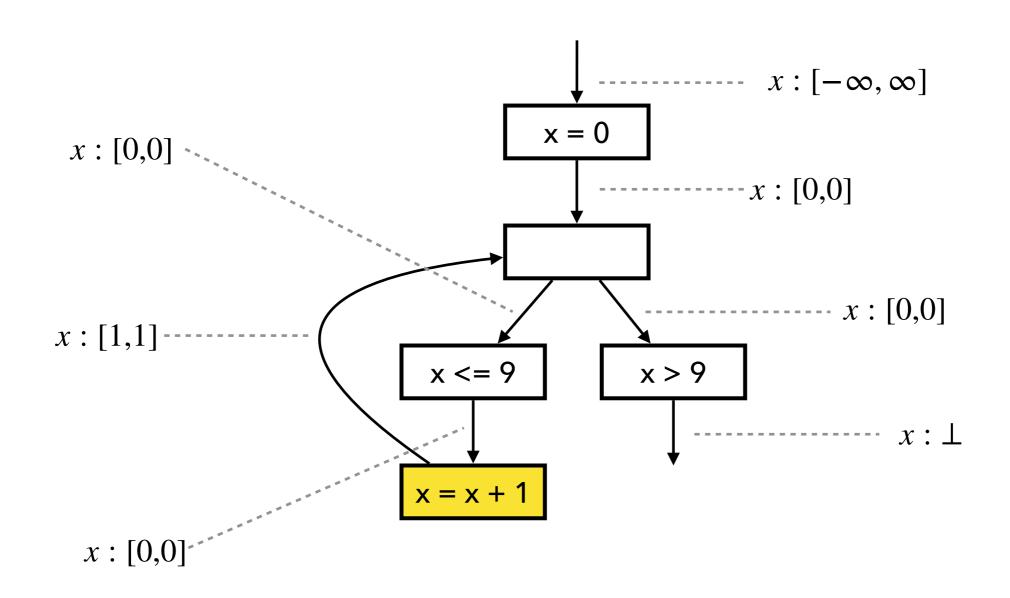


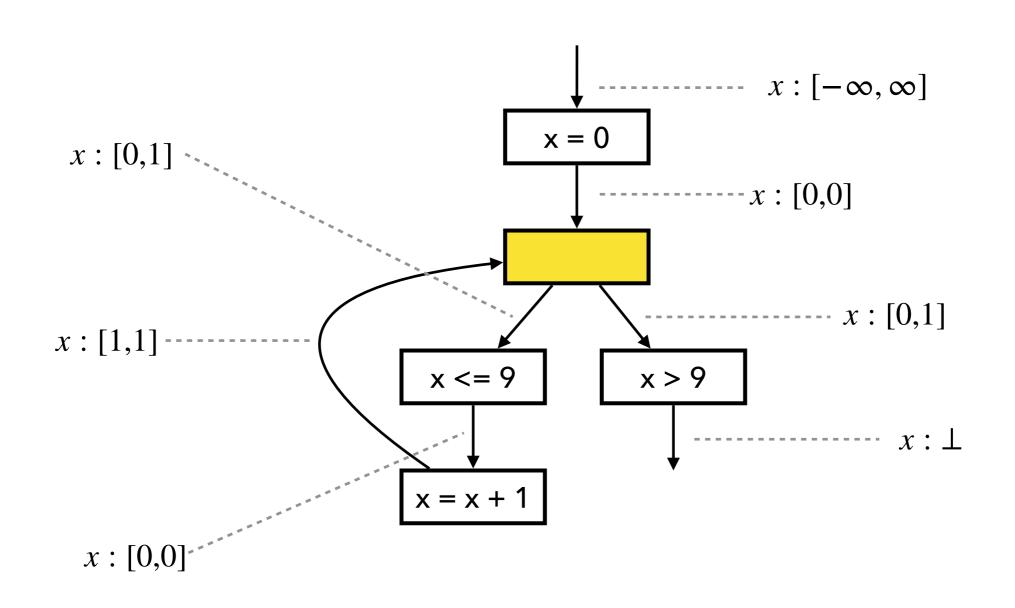


Input state:  $[0,0] \sqcup \bot = [0,0]$ 

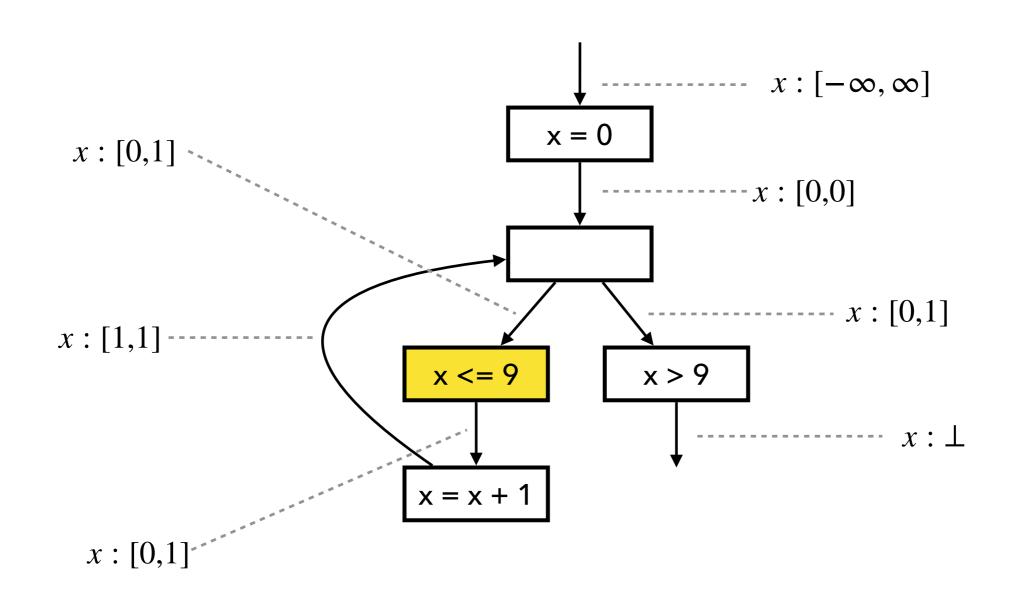


$$[0,0] \sqcap [-\infty,9] = [0,0]$$

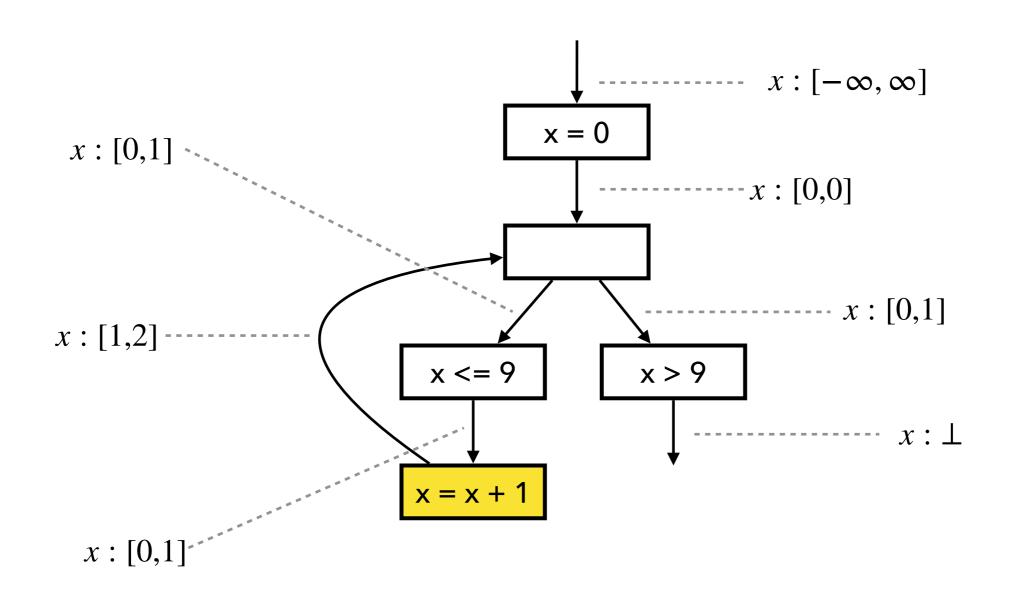


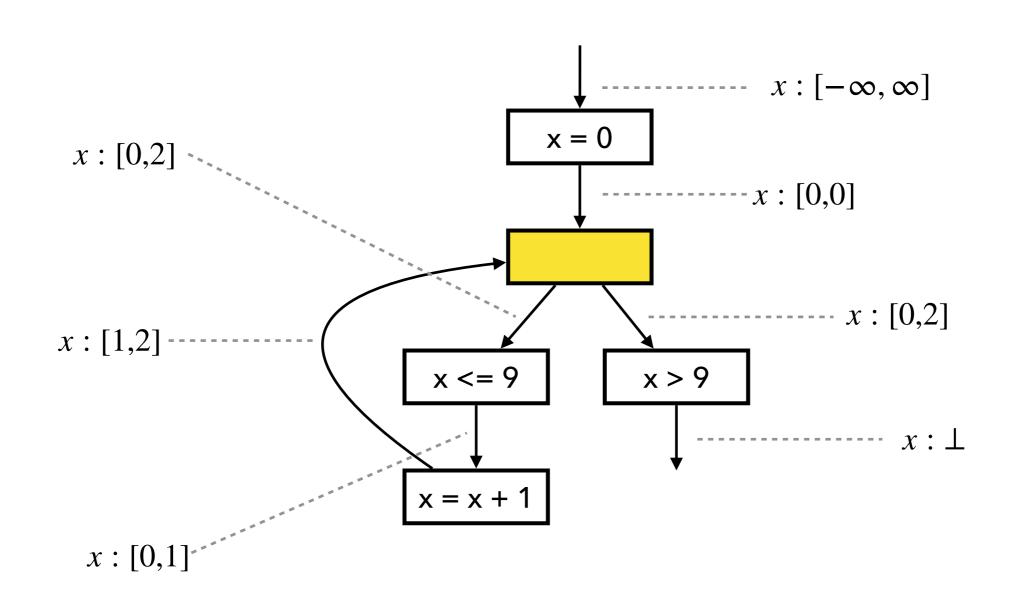


Input state:  $[0,0] \sqcup [1,1] = [0,1]$ (1st iteration of loop)

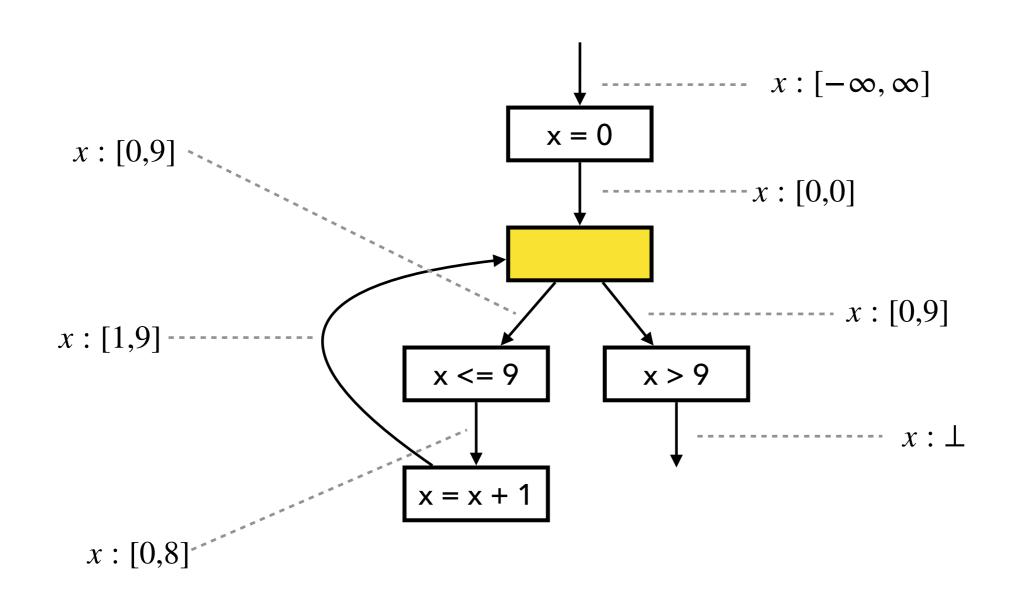


$$[0,1] \sqcap [-\infty,9] = [0,1]$$

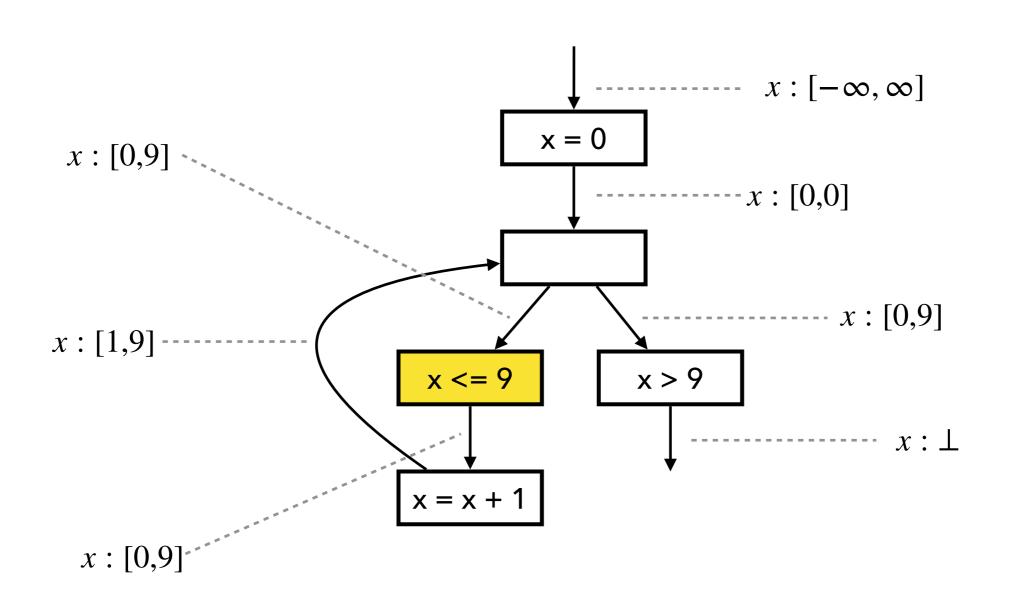




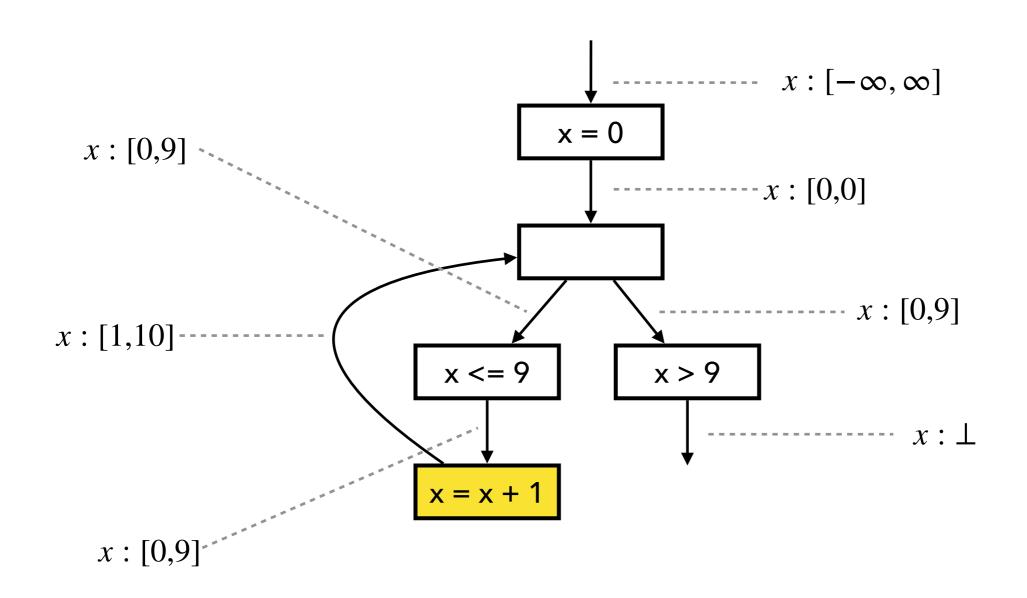
Input state:  $[0,0] \sqcup [1,2] = [0,2]$  (2nd iteration of loop)

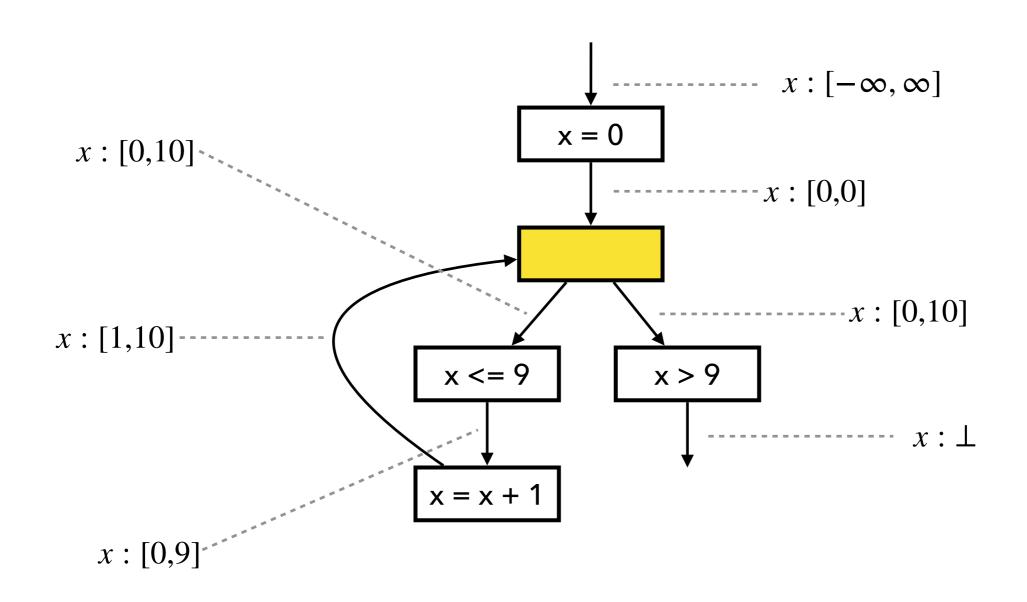


Input state:  $[0,0] \sqcup [1,9] = [0,9]$ (9th iteration of loop)

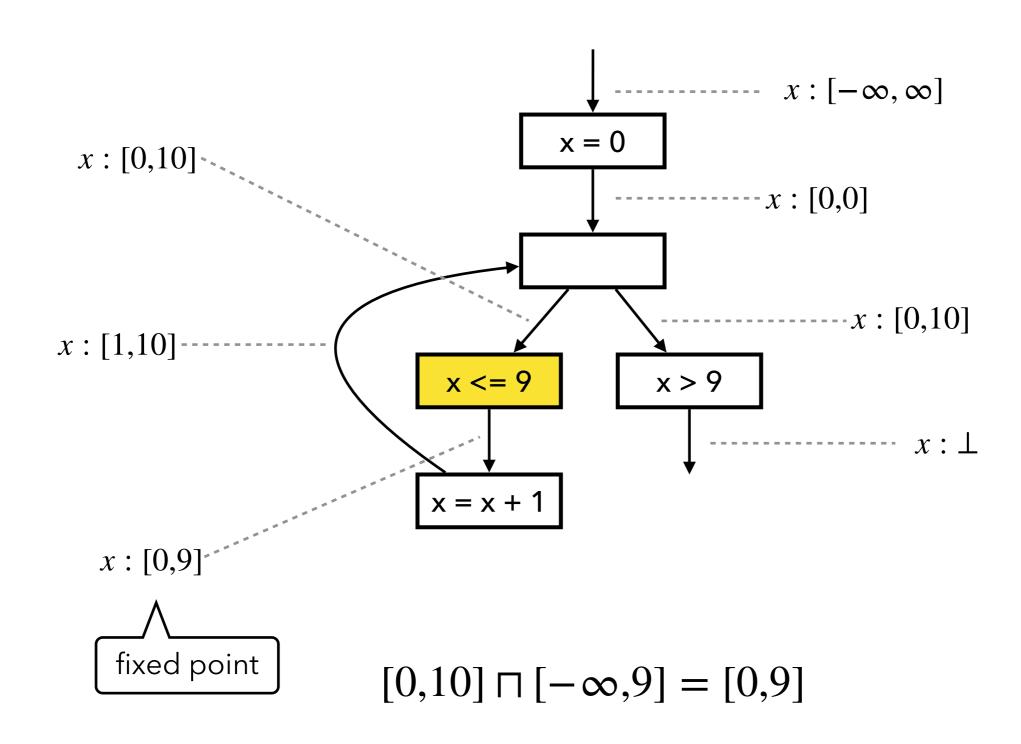


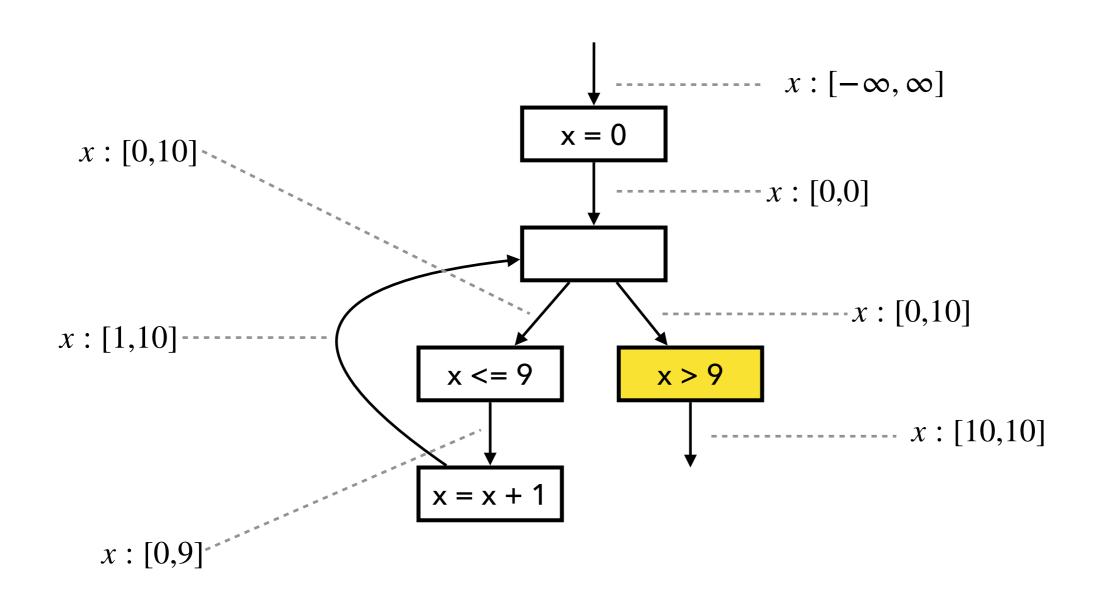
$$[0,9] \sqcap [-\infty,9] = [0,9]$$



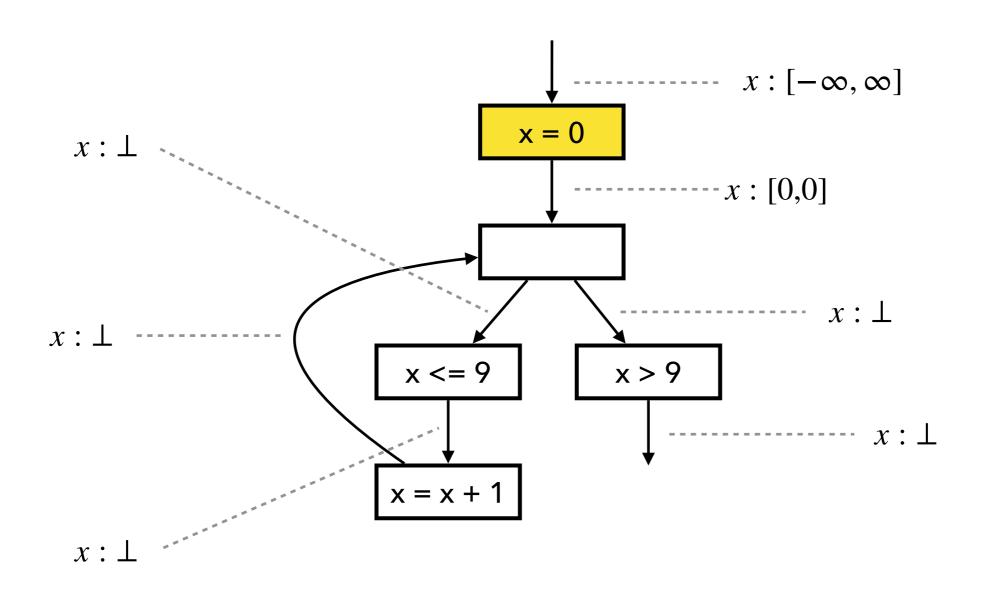


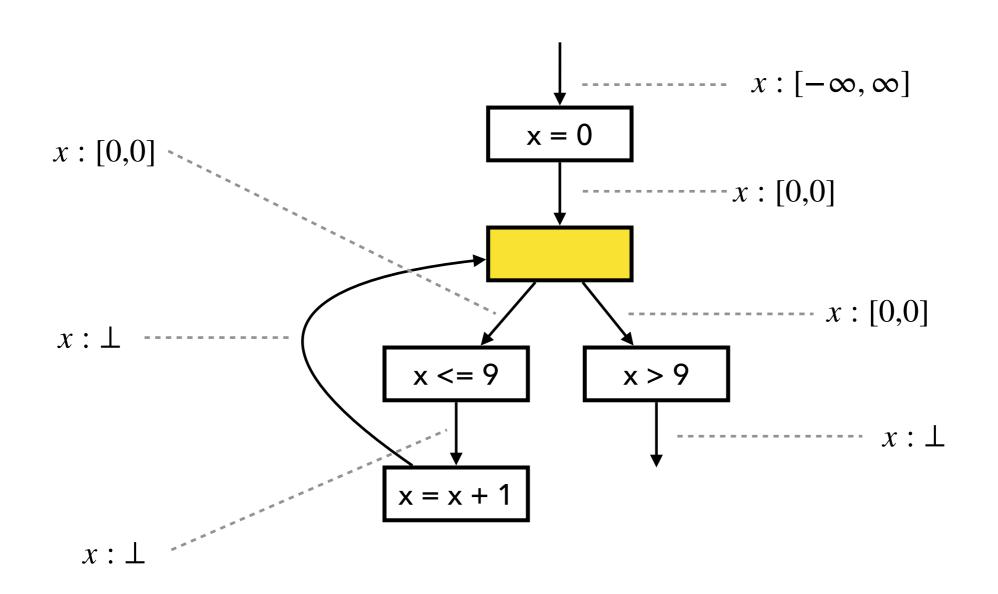
Input state:  $[0,0] \sqcup [1,10] = [0,10]$  (10th iteration of loop)



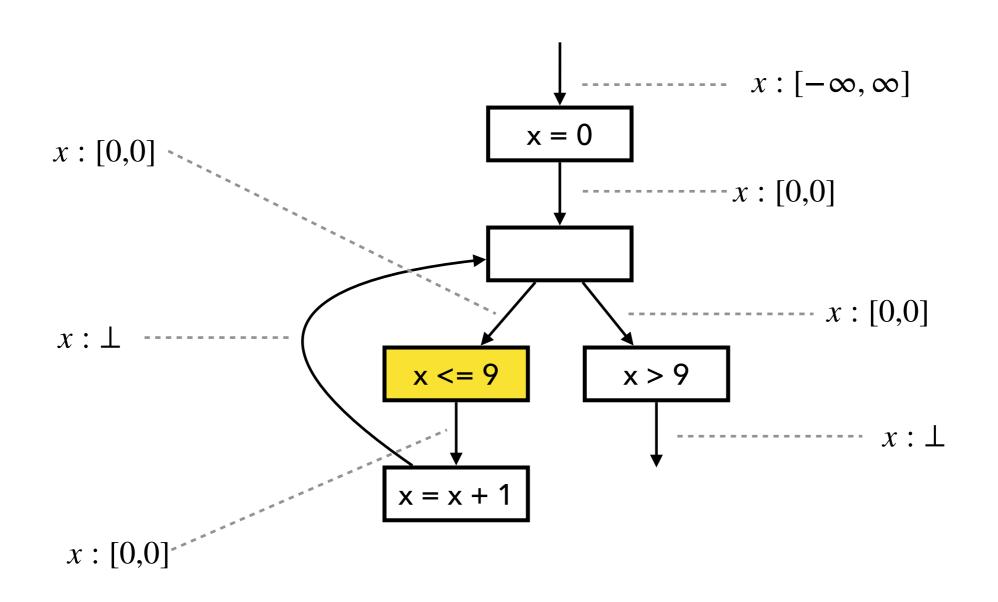


$$[0,10] \sqcap [10,\infty] = [10,10]$$

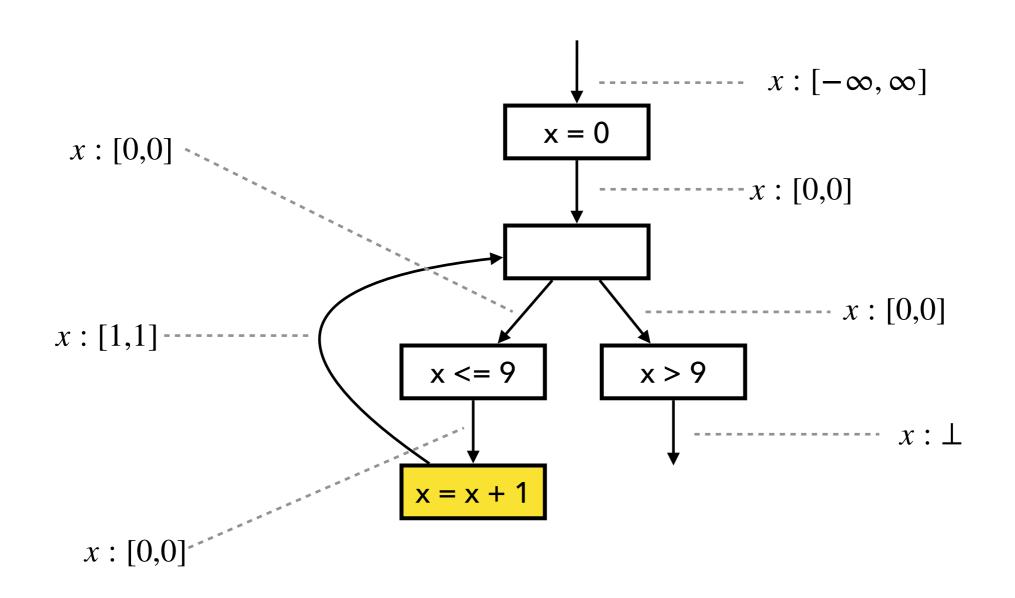




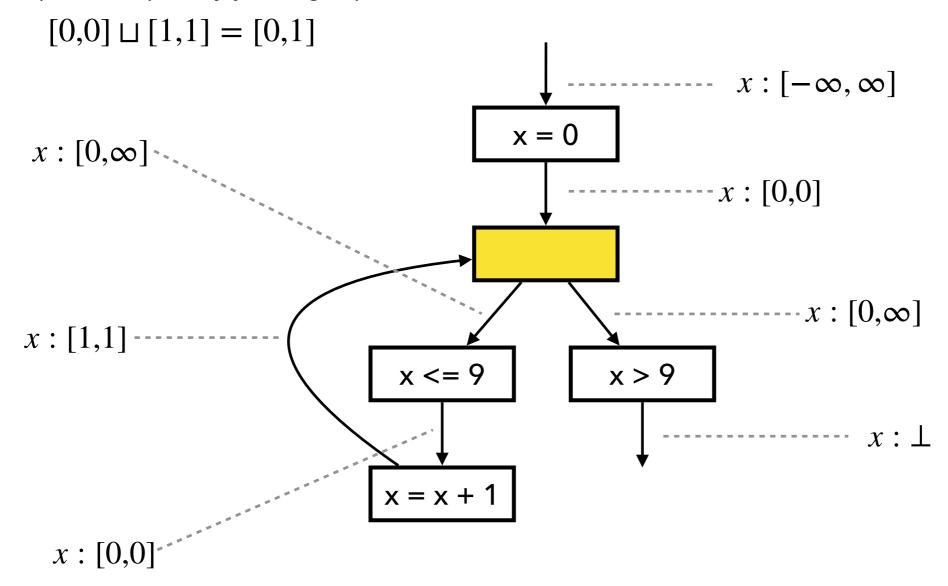
Input state:  $[0,0] \sqcup \bot = [0,0]$ 



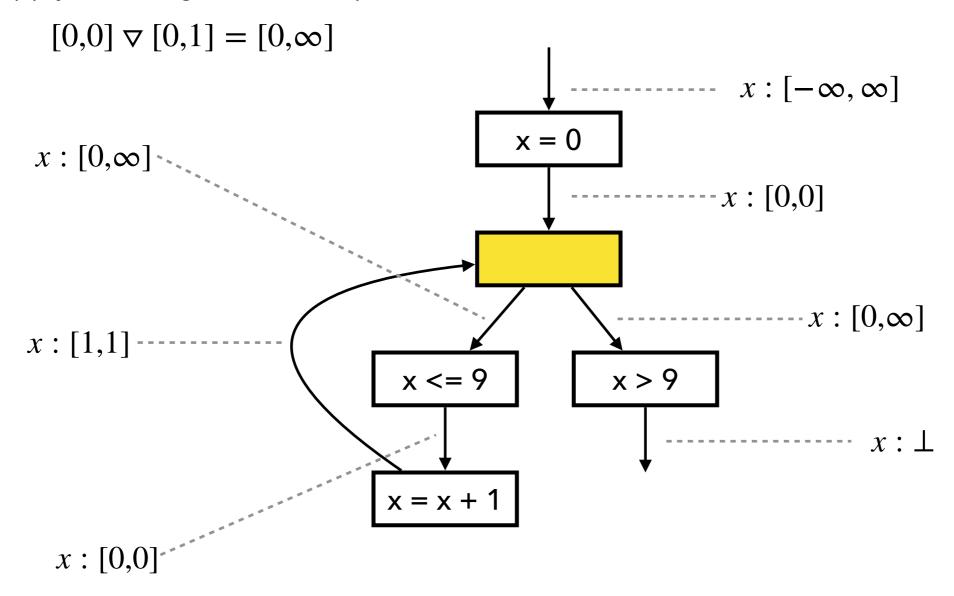
$$[0,0] \sqcap [-\infty,9] = [0,0]$$



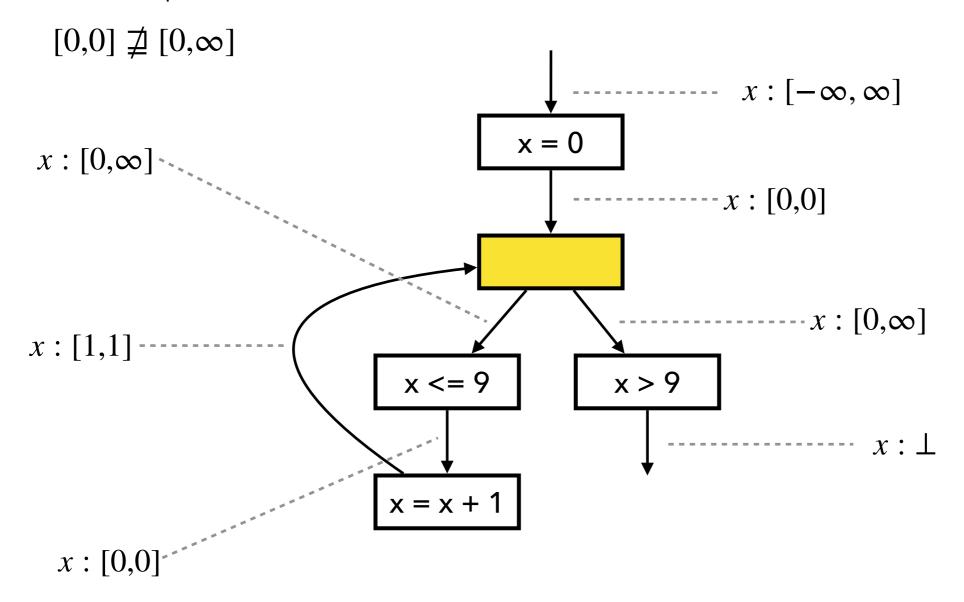
1. Compute output by joining inputs:

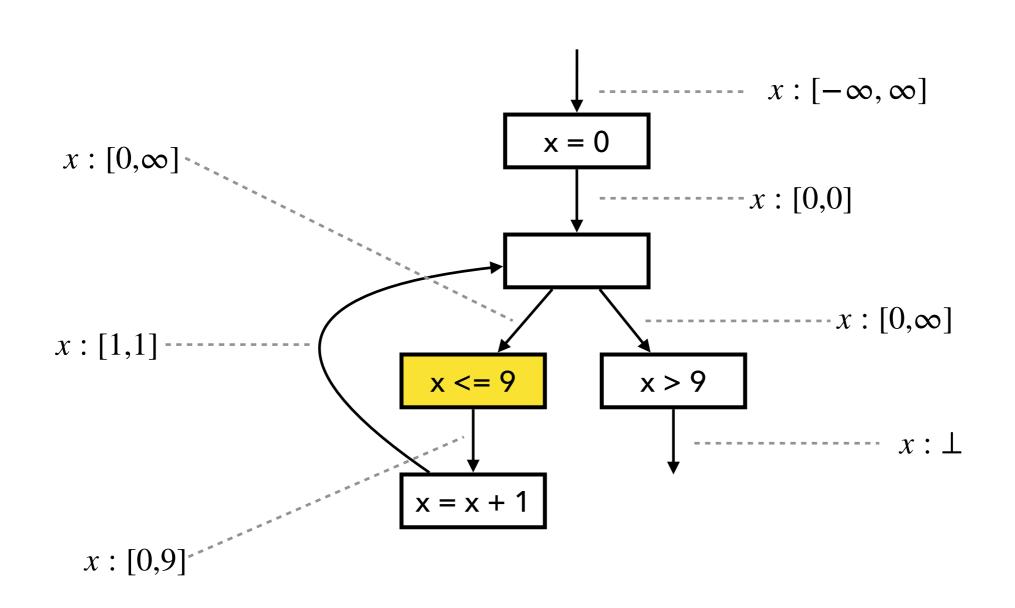


2. Apply widening with old output:

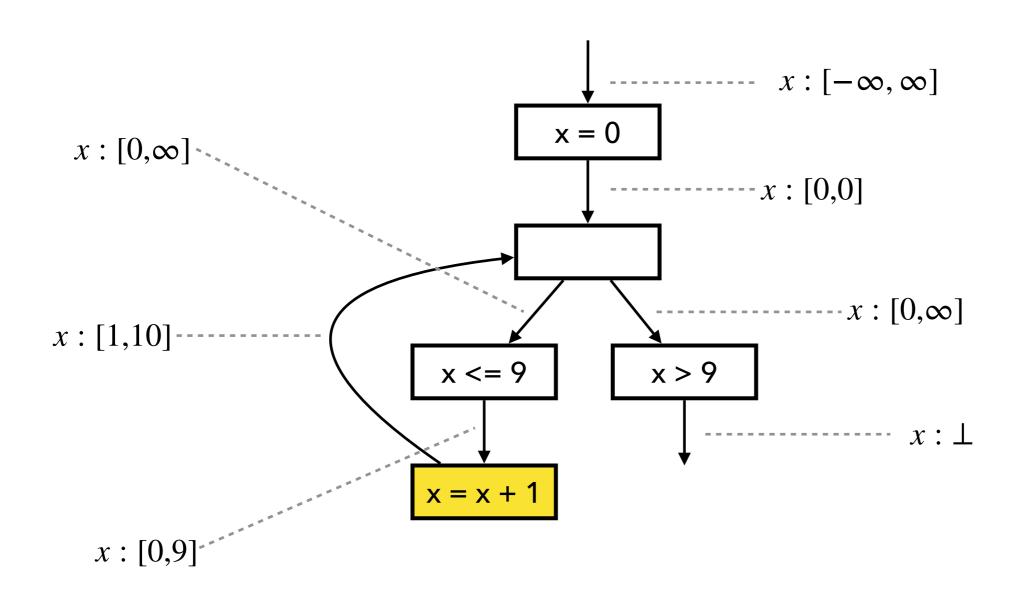


3. Check if fixed point is reached

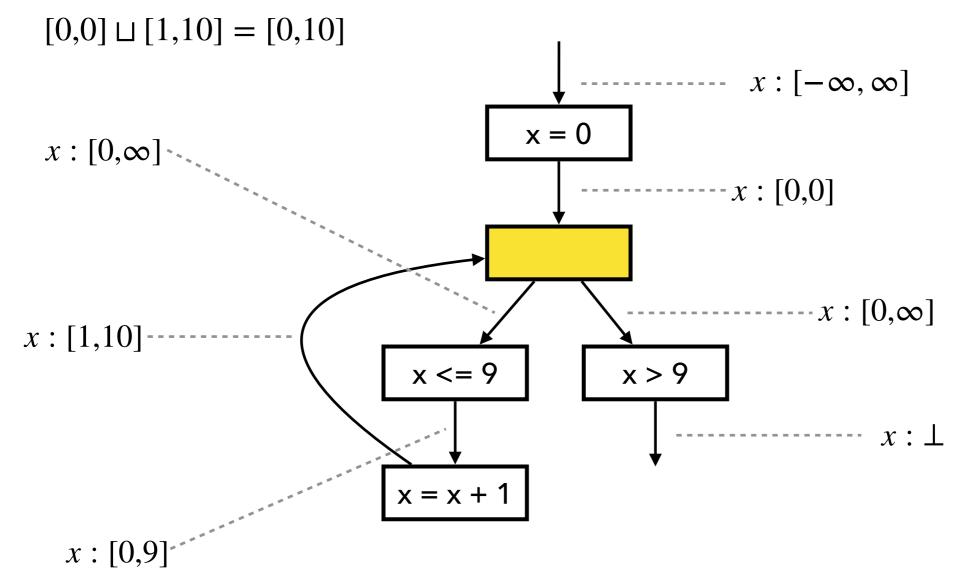




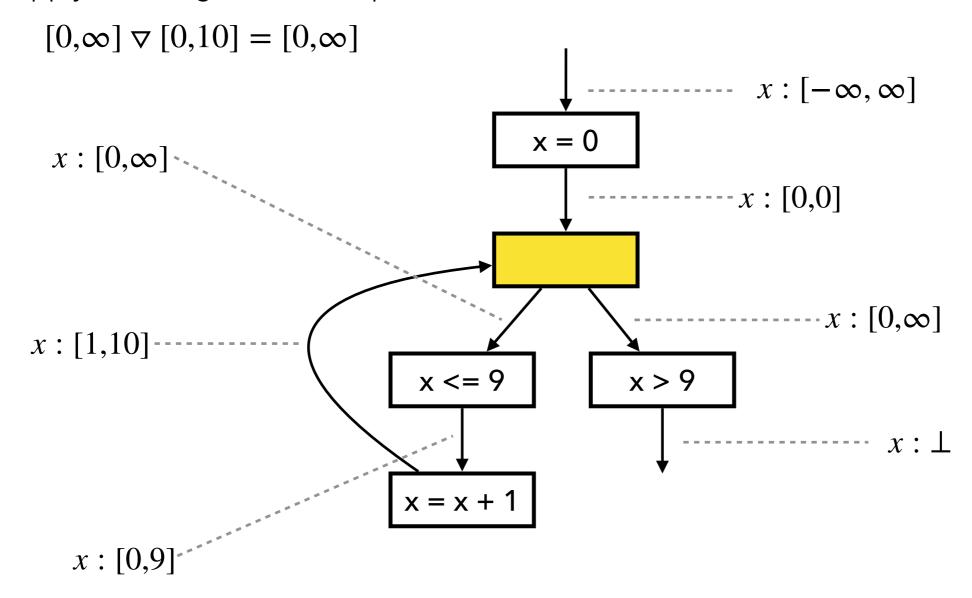
$$[0,\infty] \sqcap [-\infty,9] = [0,9]$$



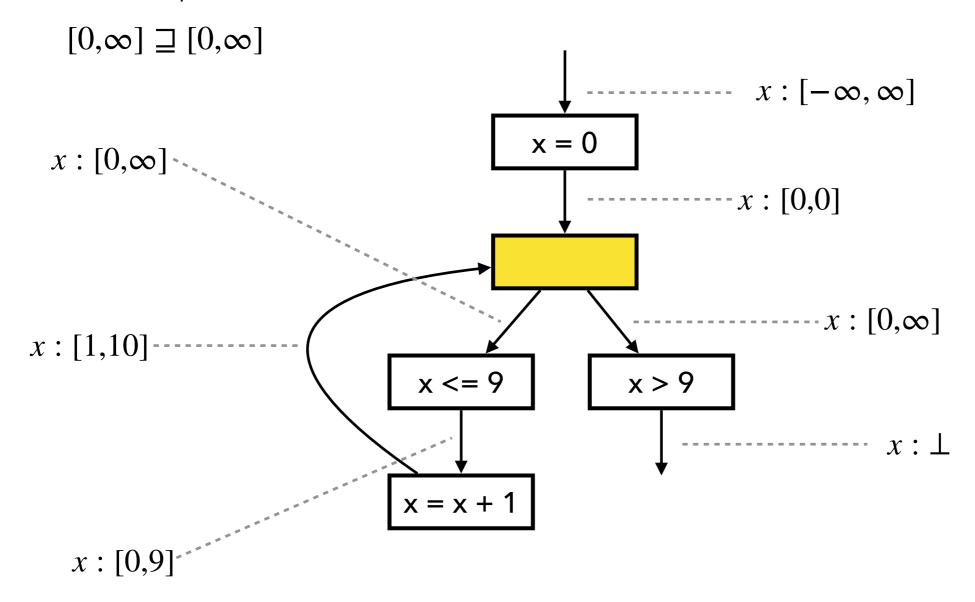
1. Compute output by joining inputs:

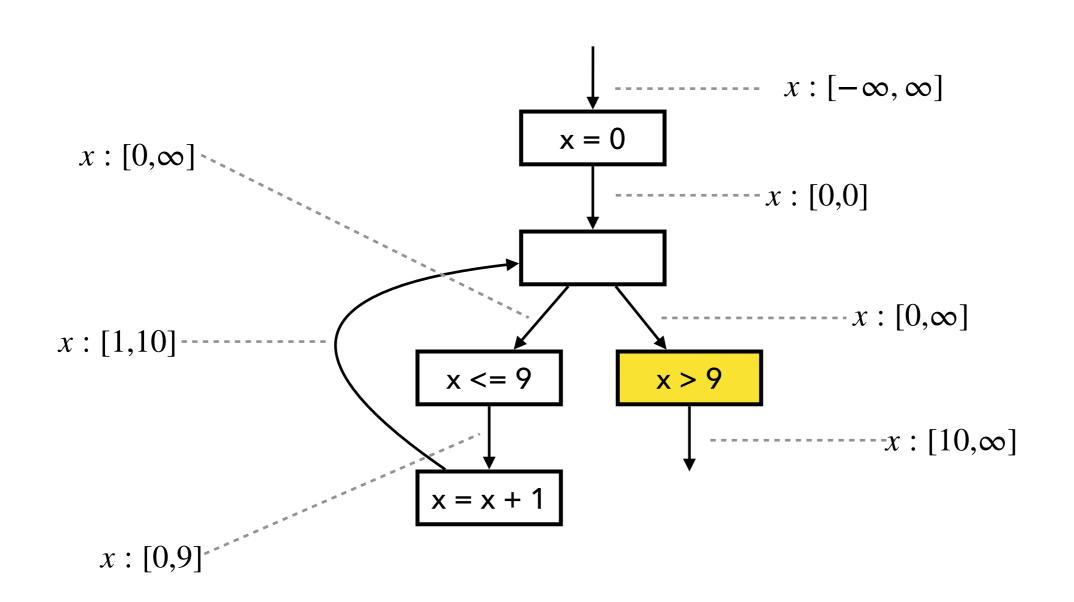


2. Apply widening with old output:



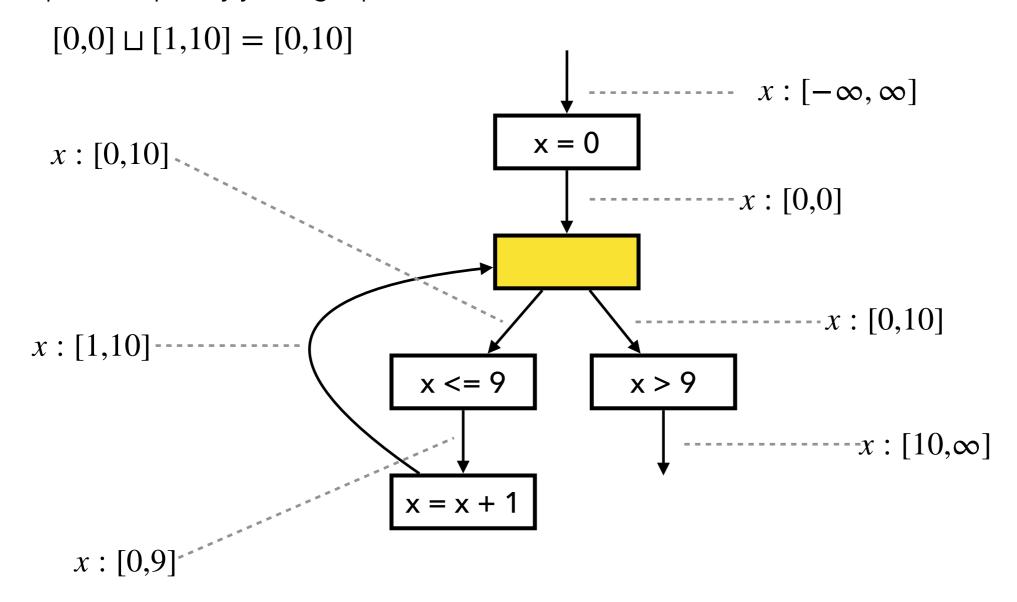
3. Check if fixed point is reached



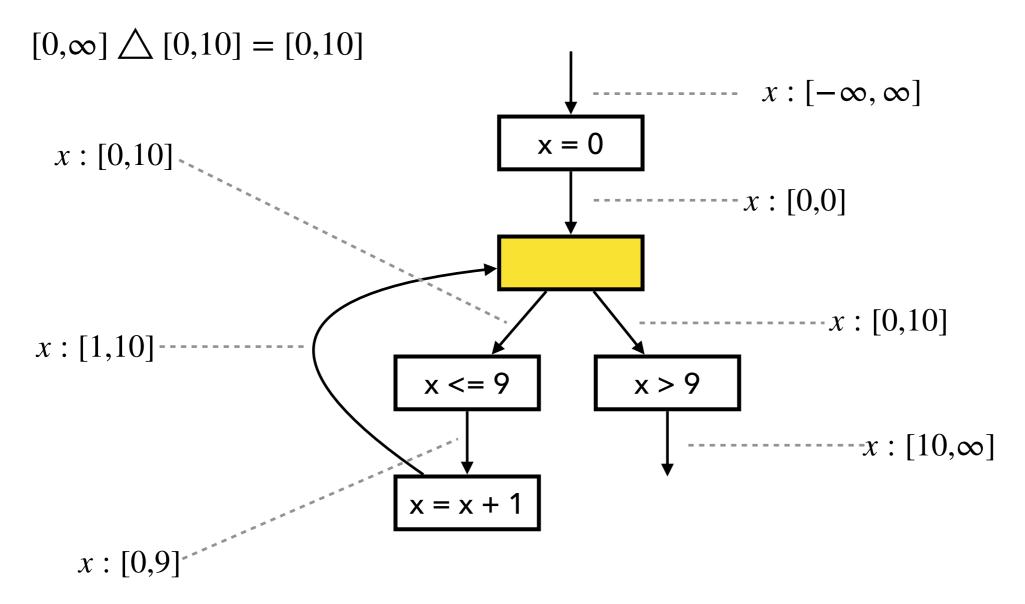


$$[0,\infty] \sqcap [10,\infty] = [10,\infty]$$

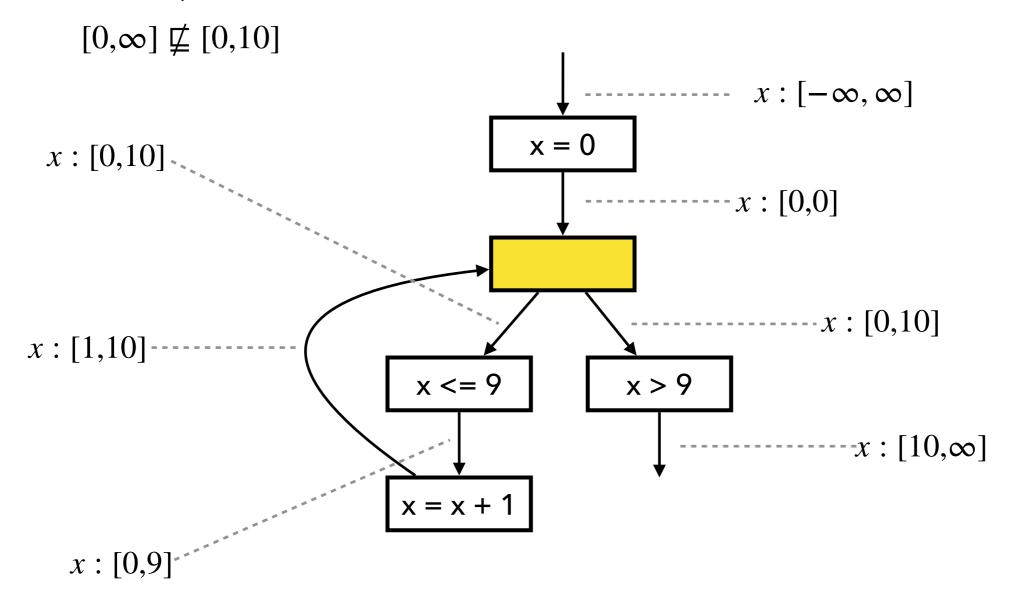
1. Compute output by joining inputs:

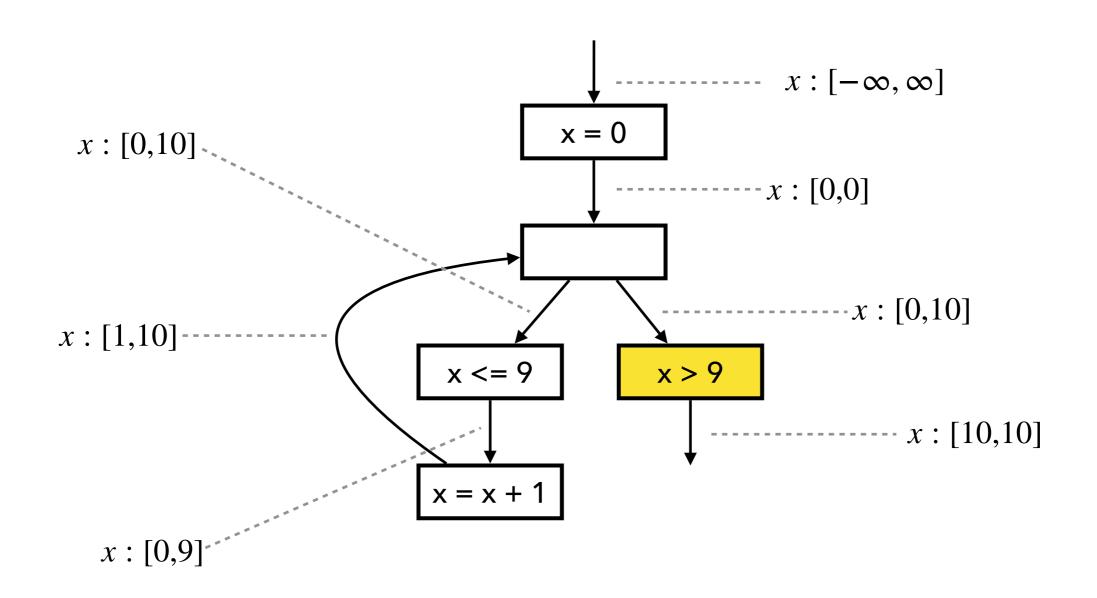


2. Apply narrowing with old output:



3. Check if fixed point is reached:





#### The Interval Domain

The set of intervals:

$$\hat{\mathbb{Z}} = \{ \perp \} \cup \{ [l, u] \mid l, u \in \mathbb{Z} \cup \{ -\infty, \infty \}, l \le u \}$$

Partial order:

$$\bot \sqsubseteq \hat{z}$$
 (for any  $\hat{z} \in \hat{\mathbb{Z}}$ )  $[l_1, u_1] \sqsubseteq [l_2, u_2] \iff l_2 \le l_1 \land u_1 \le u_2$ 

• Join:

$$\perp \sqcup \hat{z} = \hat{z}$$
  $\hat{z} \sqcup \perp = \hat{z}$   $[l_1, u_1] \sqcup [l_2, u_2] = [\min(l_1, l_2), \max(u_1, u_2)]$ 

Meet:

$$[l_1, u_1] \sqcap [l_2, u_2] = [l_2, u_1]$$
 (if  $l_1 \le l_2 \land l_2 \le u_1$ )  
 $[l_1, u_1] \sqcap [l_2, u_2] = [l_1, u_2]$  (if  $l_2 \le l_1 \land l_1 \le u_2$ )  
 $\hat{z}_1 \sqcap \hat{z}_2 = \bot$  (otherwise)

#### The Interval Domain

Widening:

Narrowing:

$$\bot \triangle \hat{z} = \bot$$

$$\hat{z} \triangle \bot = \bot$$

$$[l_1, u_1] \triangle [l_2, u_2] = [l_1 = -\infty?l_2: l_1, u_1 = +\infty?u_2: u_1]$$

#### The Interval Domain

Addition / Subtraction / Multiplication:

$$\begin{split} &[l_1,u_1] + [l_2,u_2] = [l_1 + l_2,u_1 + u_2] \\ &[l_1,u_1] - [l_2,u_2] = [l_1 - u_2,u_1 - l_2] \\ &[l_1,u_1] \times [l_2,u_2] = [\min(l_1l_2,l_1u_2,u_1l_2,u_1u_2),\max(l_1l_2,l_1u_2,u_1l_2,u_1u_2)] \end{split}$$

• Equality (=) produces T except for the cases:

$$[l_1, u_1] \triangleq [l_2, u_2] = true$$
 (if  $l_1 = u_1 = l_2 = u_2$ )  
 $[l_1, u_1] \triangleq [l_2, u_2] = false$  (no overlap)

• "Less than" (<) produces T except for the cases:

$$[l_1, u_1] \ \hat{<} \ [l_2, u_2] = true \quad (if \ u_1 < l_2)$$
  
 $[l_1, u_1] \ \hat{<} \ [l_2, u_2] = false \quad (if \ l_1 > u_2)$ 

#### **Abstract Memory**

$$\hat{\mathbb{M}} = \mathbf{Var} \to \hat{\mathbb{Z}}$$

$$m_1 \sqsubseteq m_2 \iff \forall x \in \mathbf{Var} . \ m_1(x) \sqsubseteq m_2(x)$$
 
$$m_1 \sqcup m_2 = \lambda x . \ m_1(x) \sqcup m_2(x)$$
 
$$m_1 \sqcap m_2 = \lambda x . \ m_1(x) \sqcap m_2(x)$$
 
$$m_1 \bigvee m_2 = \lambda x . \ m_1(x) \bigvee m_2(x)$$
 
$$m_1 \bigtriangleup m_2 = \lambda x . \ m_1(x) \bigtriangleup m_2(x)$$

## Worklist Algorithm

Fixpoint comp. with widening

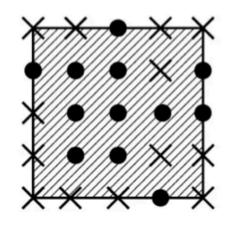
```
W := \mathsf{Node}
T:=\lambda n . \perp_{\hat{\mathbb{M}}}
while W \neq \emptyset
   n := choose(W)
   W := W \setminus \{n\}
   in := input o f(n, T)
   out := analyze(n, in)
   if out \not\sqsubseteq T(n)
      if widening is needed
         T(n) := T(n) \nabla out
     else
         T(n) := T(n) \sqcup out
     W := W \cup succ(n)
```

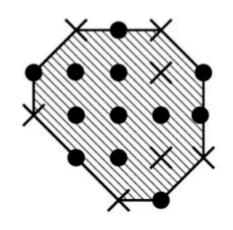
Fixpoint comp. with narrowing

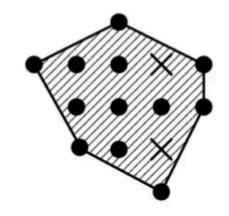
$$W := \mathbf{Node}$$
 $while \ W \neq \emptyset$ 
 $n := choose(W)$ 
 $W := W \setminus \{n\}$ 
 $in := inputof(n, T)$ 
 $out := analyze(n, in)$ 
 $if \ T(n) \not\sqsubseteq out$ 
 $T(n) := T(n) \triangle out$ 
 $W := W \cup succ(n)$ 

#### **Relational Abstract Domains**

• Intervals vs. Octagons vs. Polyhedra







Focus: Core idea of the Octagon domain\*

int a[10];
x = 0; y = 0;

while (x < 9) {
 x++; y++;
}
a[y] = 0;</pre>
Octagon analysis

y:[9,9] x - y:[0,0]x + y:[18,18]

x : [9,9]

x: [9,9] $y: [0,\infty]$ 

#### Difference Bound Matrix (DBM)

•  $(N+1) \times (N+1)$  matrix (N: the number of variables): e.g.,

Example

$$\begin{bmatrix} 0 & 10 & 10 \\ 0 & 0 & 0 \\ 0 & 0 & 0 \end{bmatrix} \iff \begin{matrix} 0 \le x \le 10 \\ 0 \le y \le 10 \\ y - x \le 0 \\ x - y \le 0 \end{matrix} \qquad \begin{bmatrix} 0 & 10 & +\infty \\ -1 & 0 & -1 \\ 0 & 1 & 0 \end{bmatrix} \iff \begin{matrix} 1 \le x \le 10 \\ 0 \le y \\ y - x \le -1 \\ x - y \le 1 \end{matrix}$$

#### Difference Bound Matrix (DBM)

A DBM represents a set of program states (N-dim points)

$$\gamma \left( \begin{bmatrix} 0 & 10 & +\infty \\ -1 & 0 & -1 \\ 0 & 1 & 0 \end{bmatrix} \right) = \{ (x, y) \mid 1 \le x \le 10, 0 \le y, y - x \le -1, x - y \le 1 \}$$

A DBM can also be represented by a directed graph

#### Difference Bound Matrix (DBM)

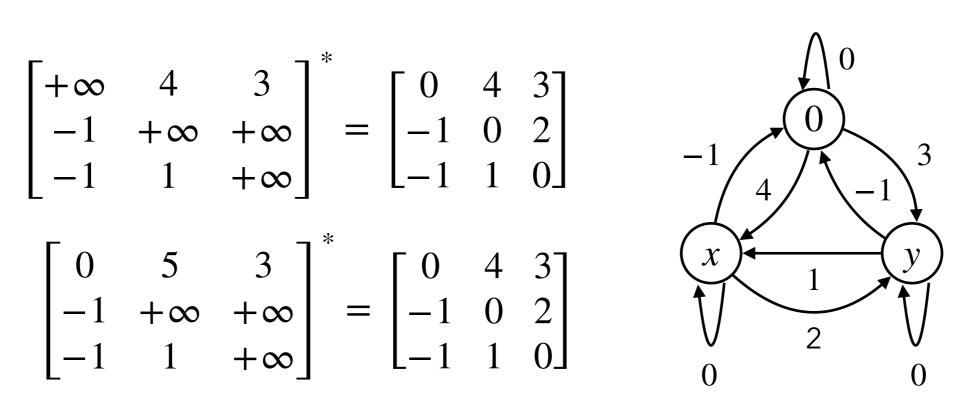
Two different DBMs can represent the same set of points

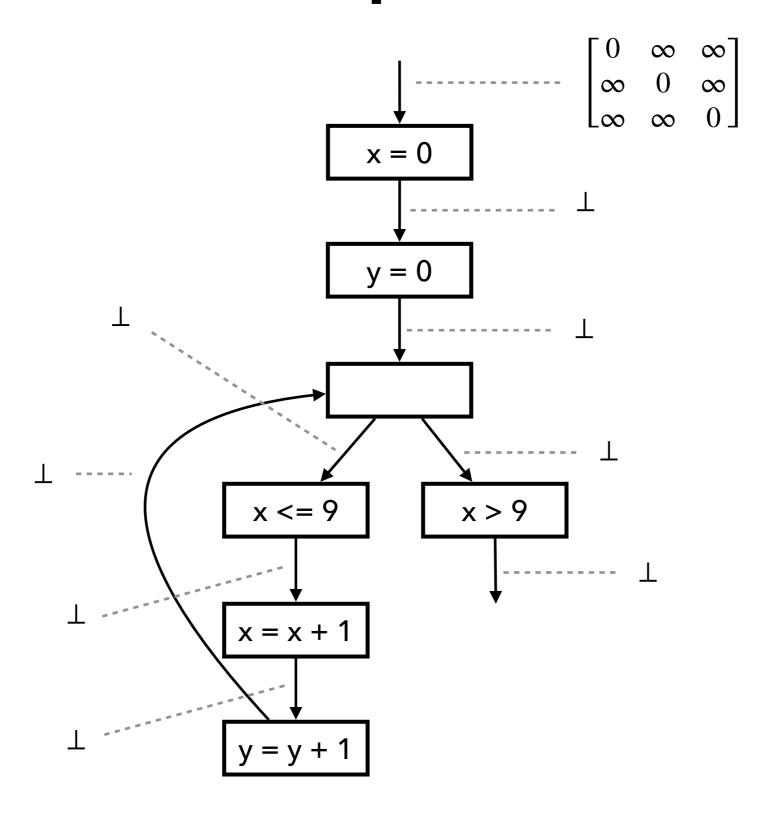
$$\gamma \left[ \begin{bmatrix} +\infty & 4 & 3 \\ -1 & +\infty & +\infty \\ -1 & 1 & +\infty \end{bmatrix} \right] = \gamma \left[ \begin{bmatrix} 0 & 5 & 3 \\ -1 & +\infty & +\infty \\ -1 & 1 & +\infty \end{bmatrix} \right]$$

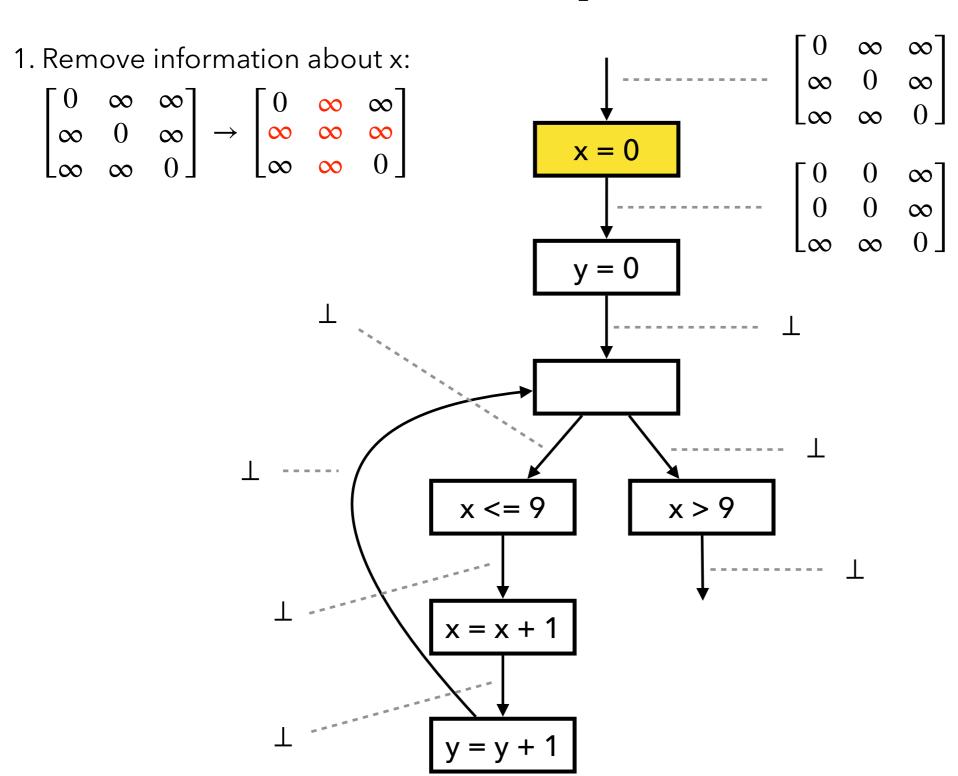
Closure (normalization) via the Floyd-Warshall algorithm

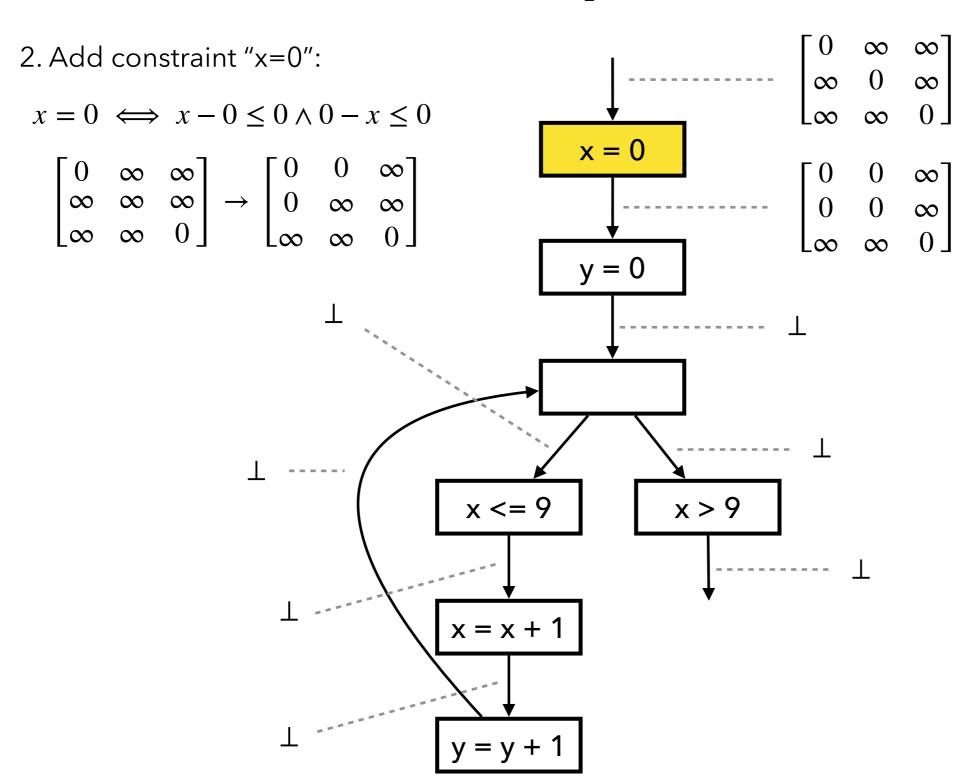
$$\begin{bmatrix} +\infty & 4 & 3 \\ -1 & +\infty & +\infty \\ -1 & 1 & +\infty \end{bmatrix}^* = \begin{bmatrix} 0 & 4 & 3 \\ -1 & 0 & 2 \\ -1 & 1 & 0 \end{bmatrix}$$

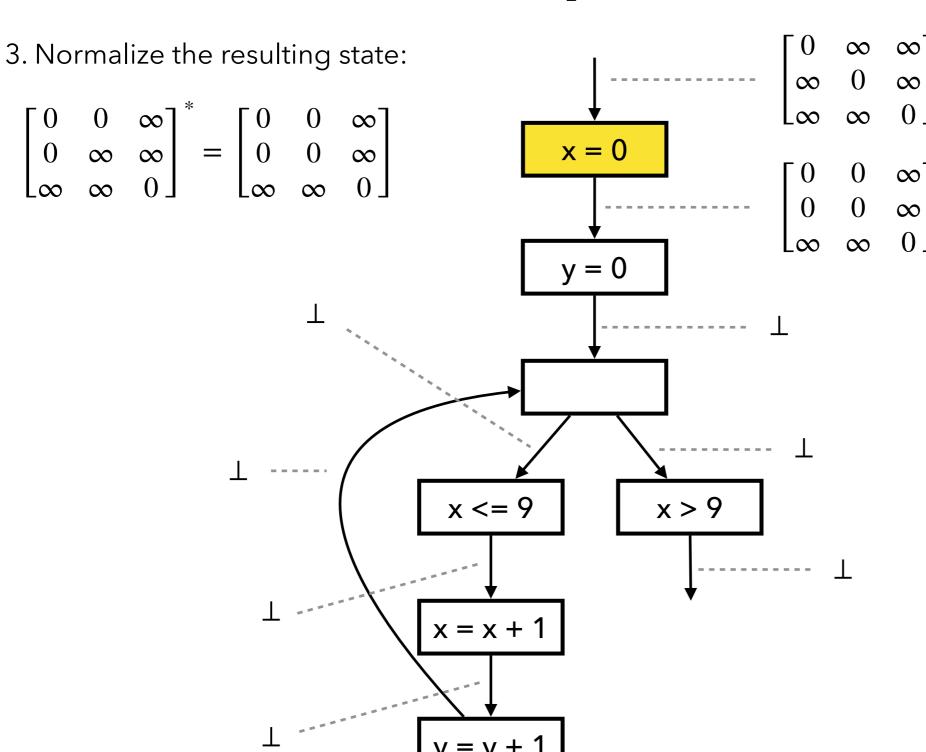
$$\begin{bmatrix} 0 & 5 & 3 \\ -1 & +\infty & +\infty \\ -1 & 1 & +\infty \end{bmatrix}^{*} = \begin{bmatrix} 0 & 4 & 3 \\ -1 & 0 & 2 \\ -1 & 1 & 0 \end{bmatrix}$$

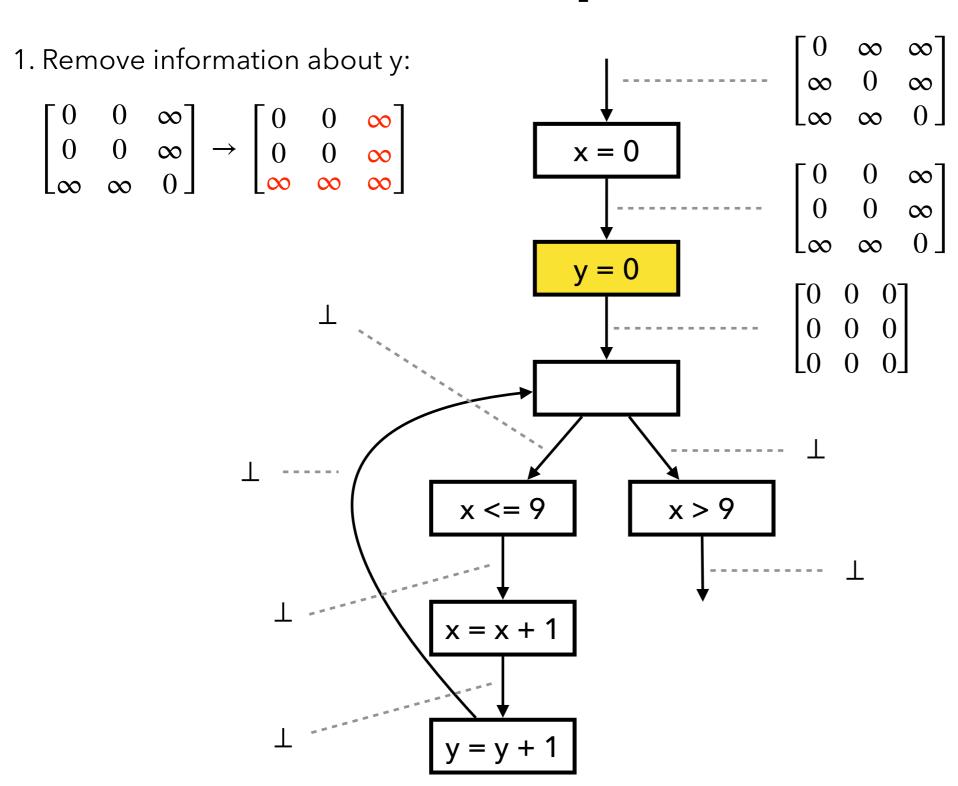


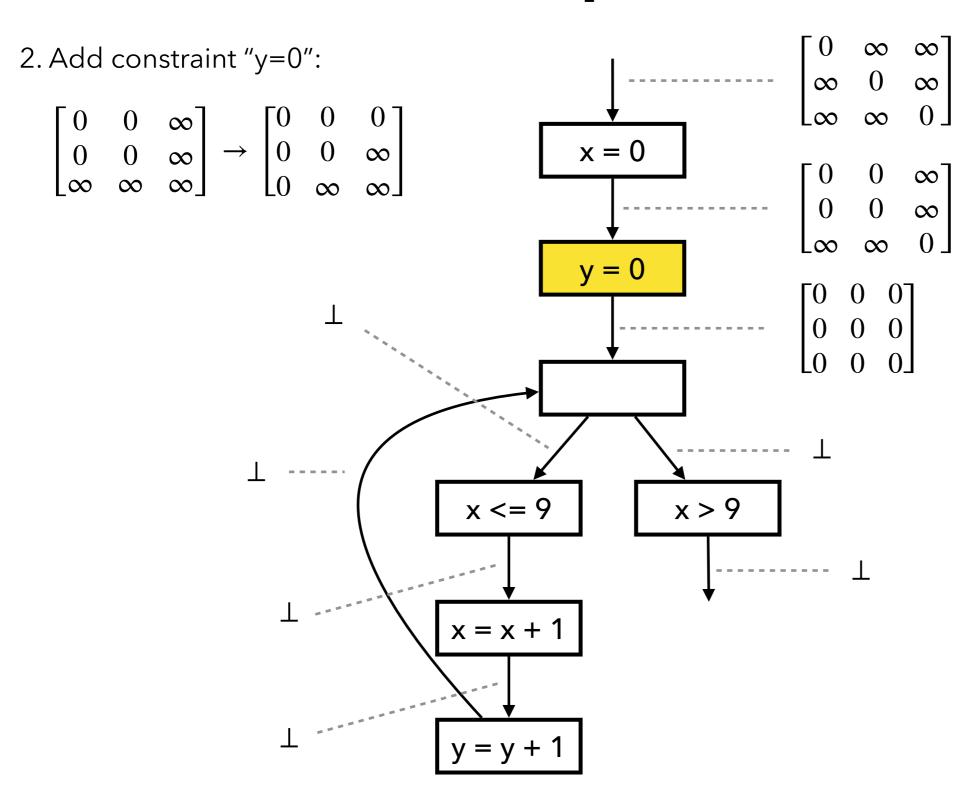


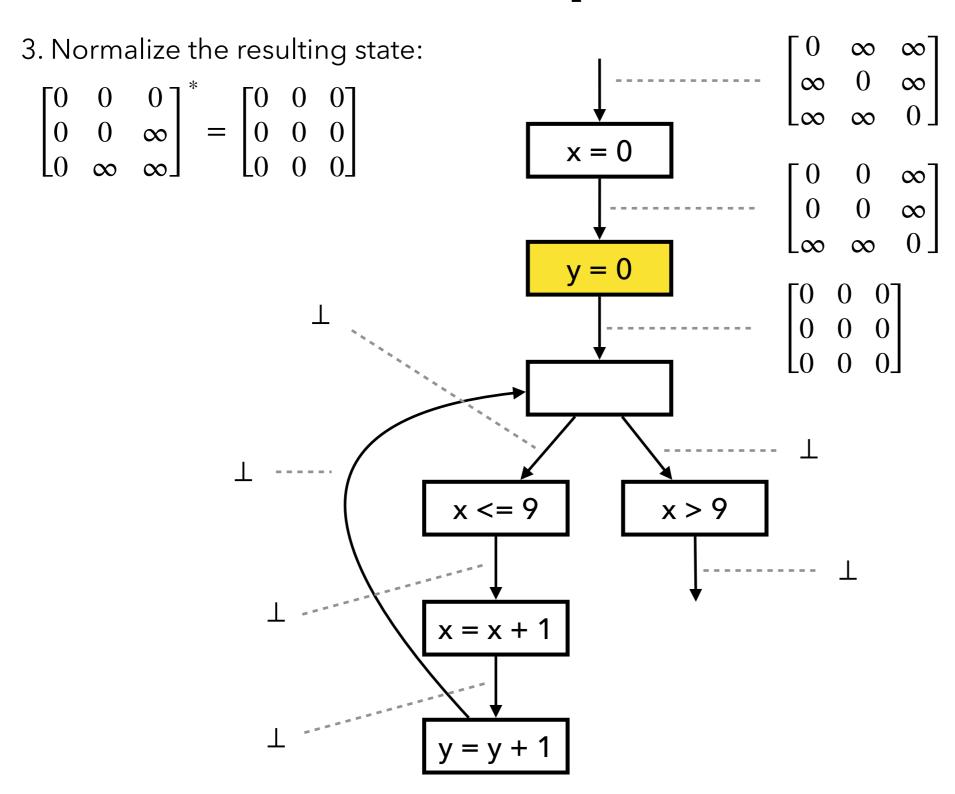


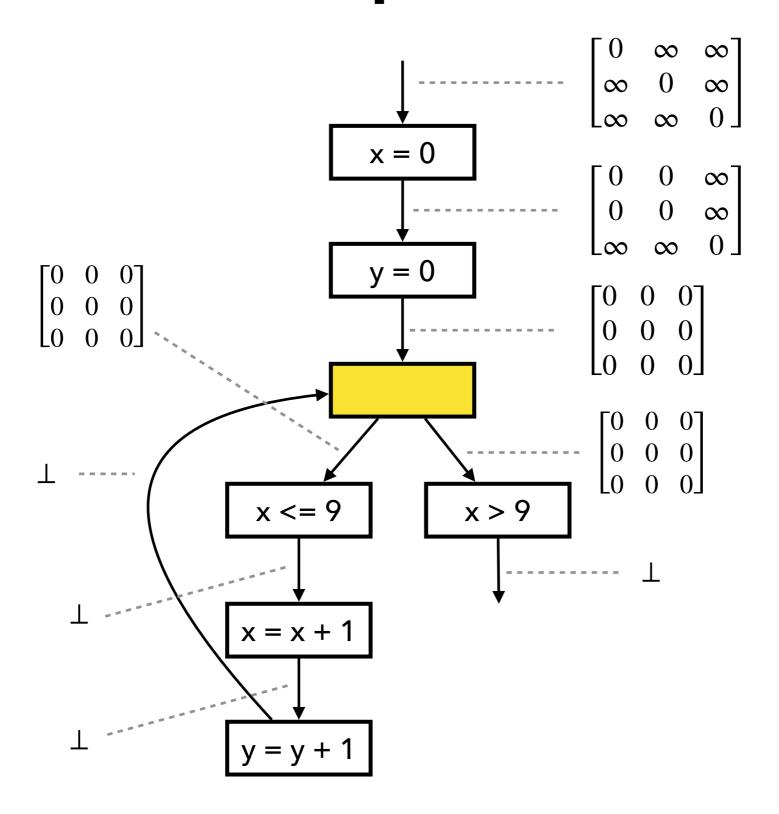


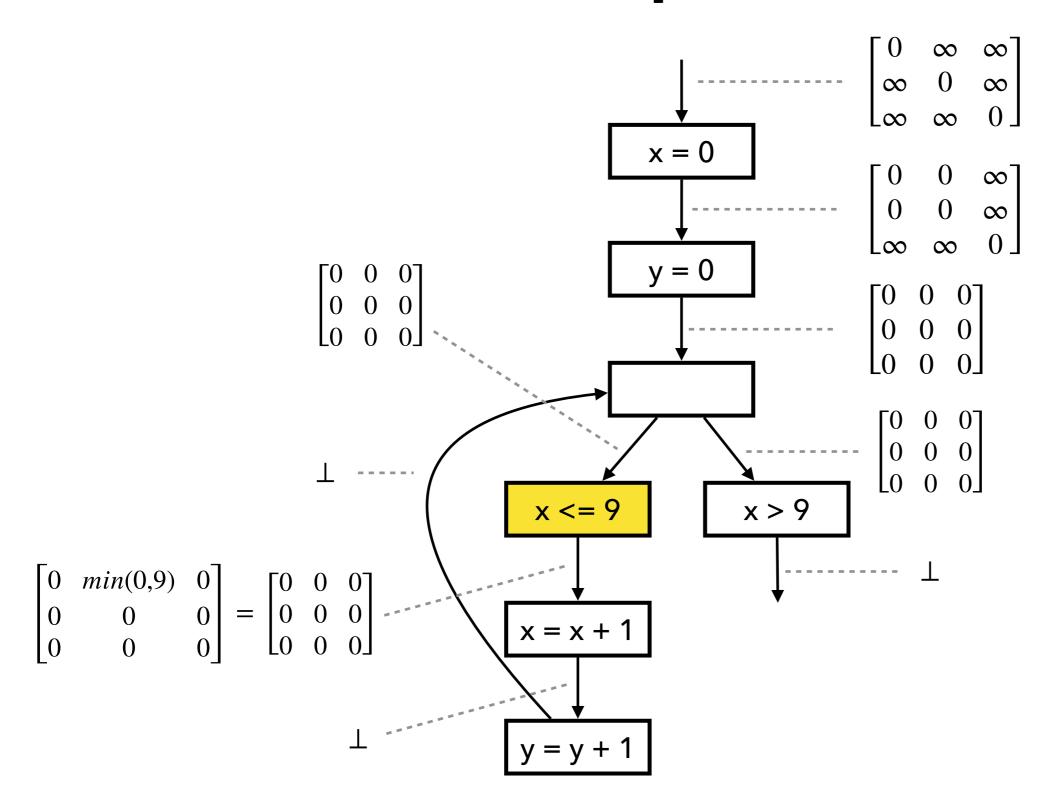


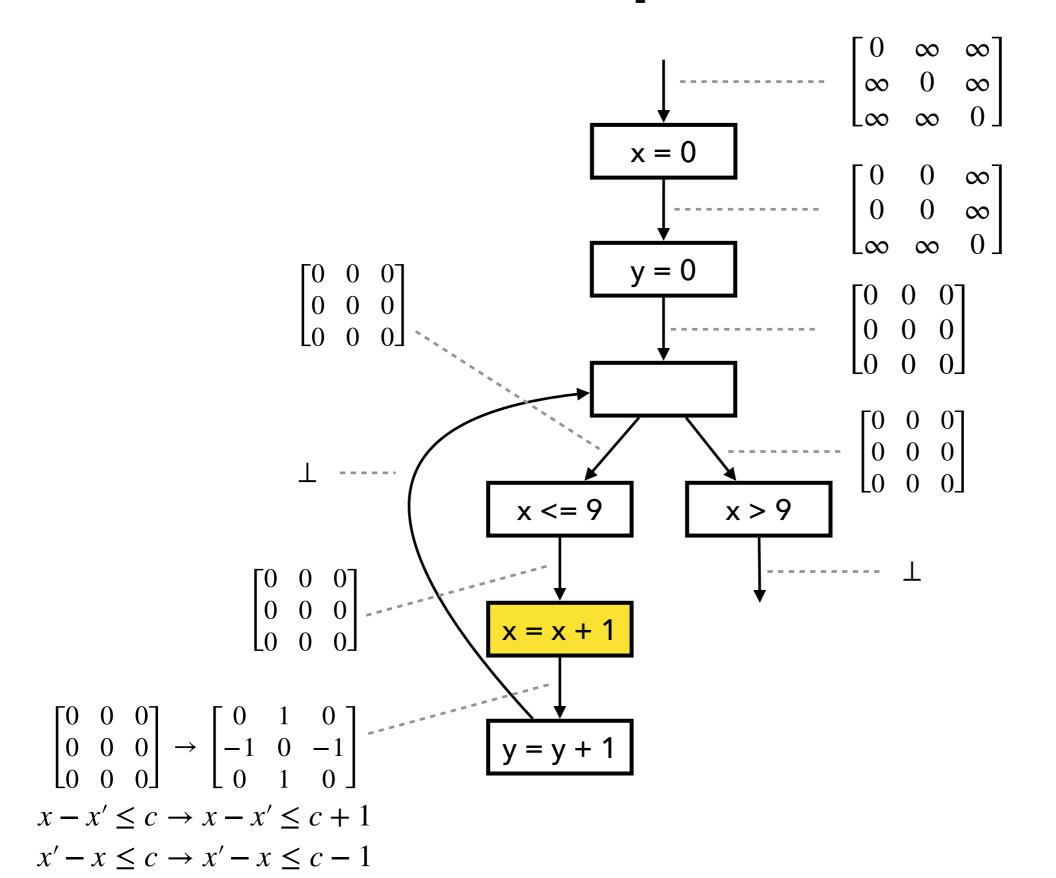


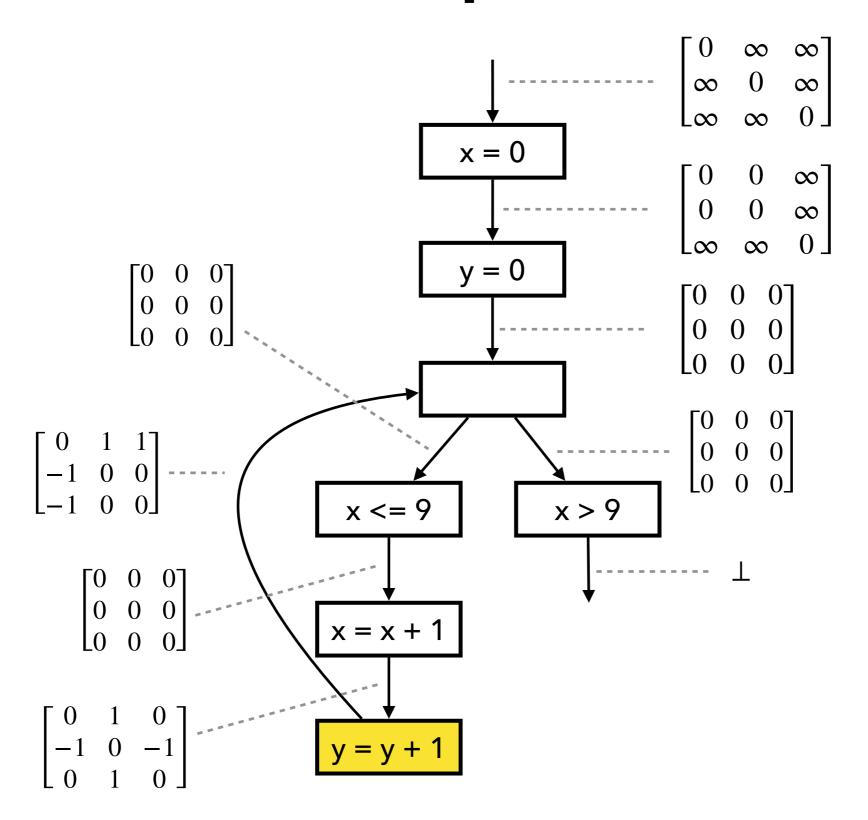


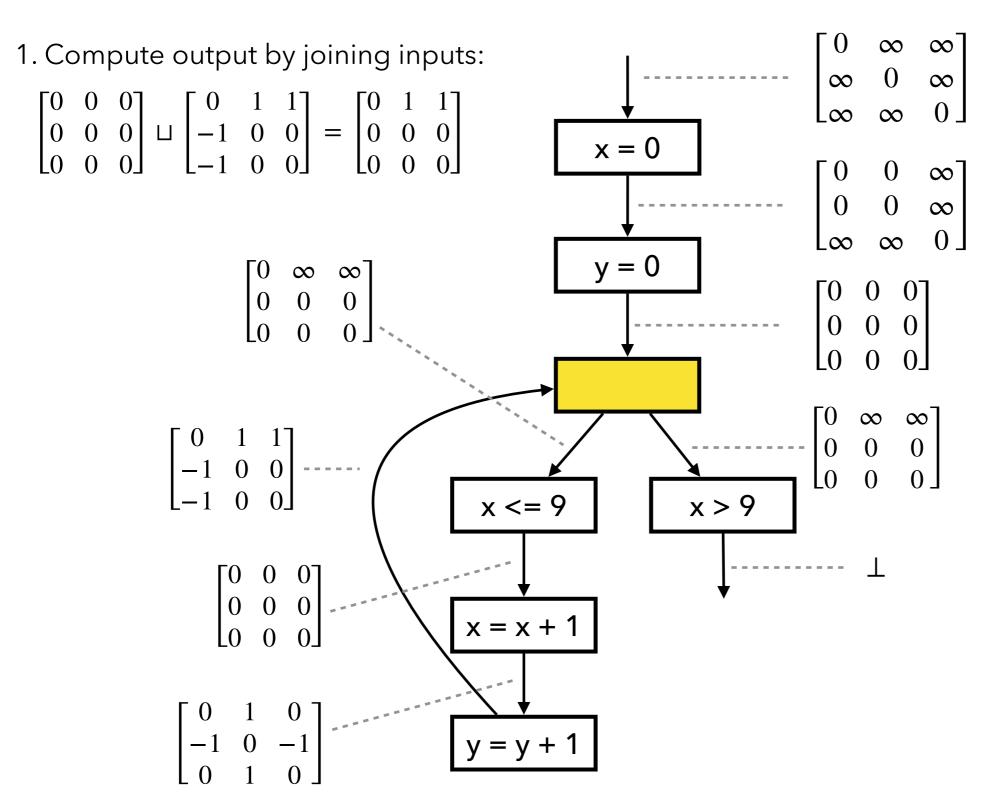


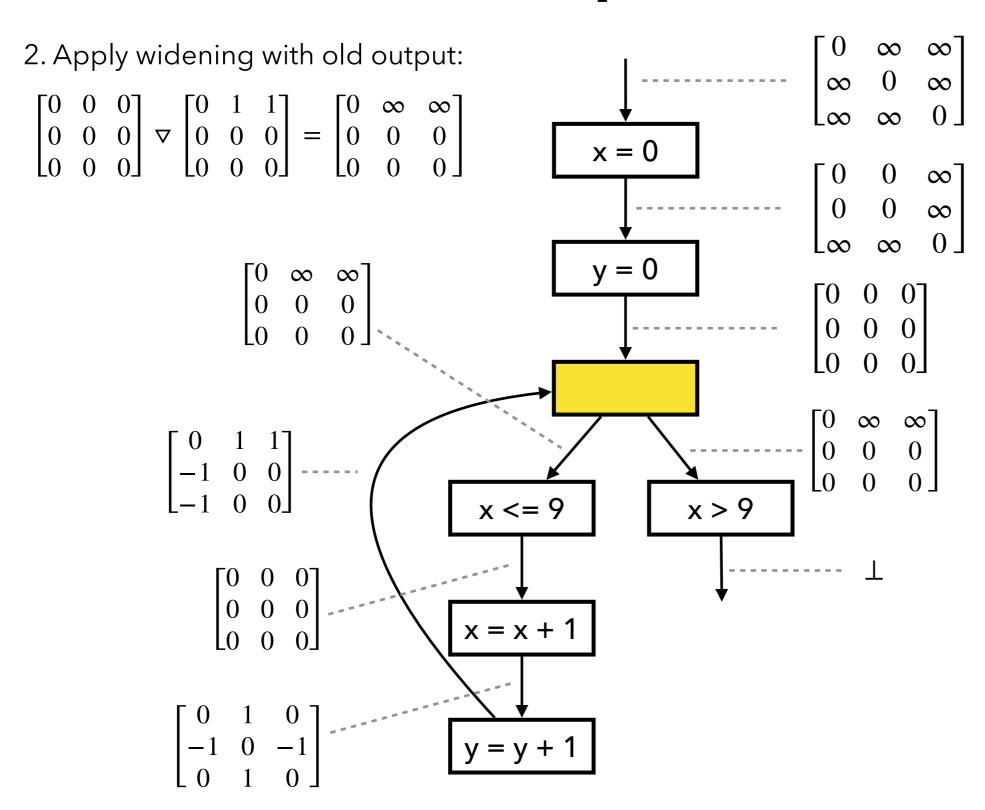


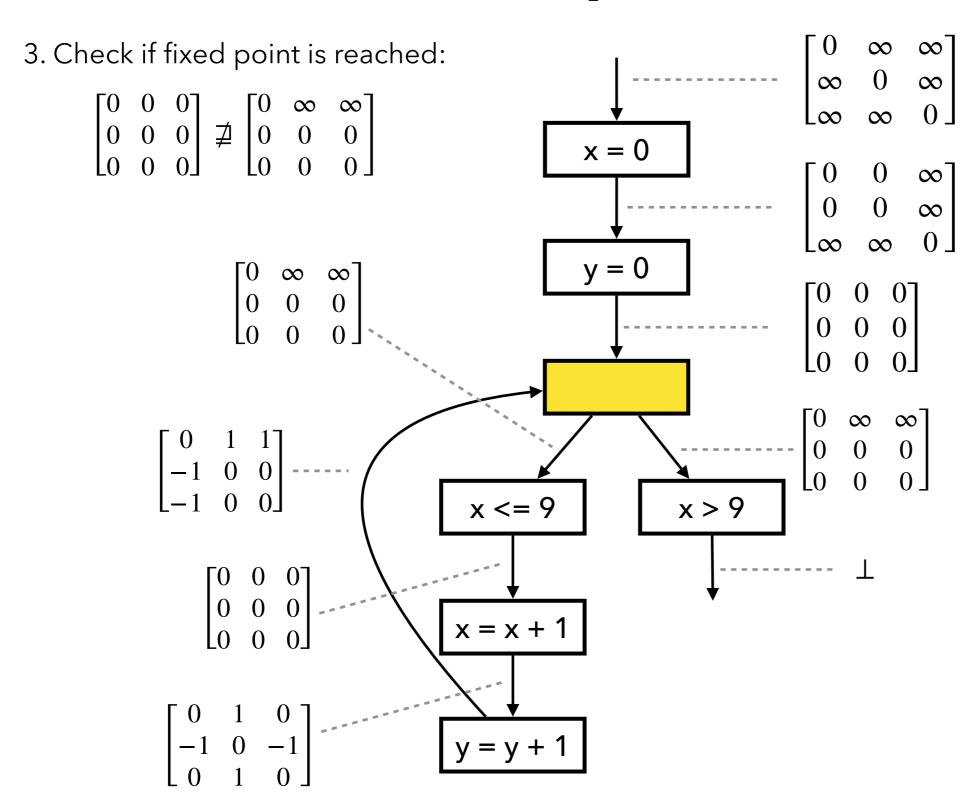


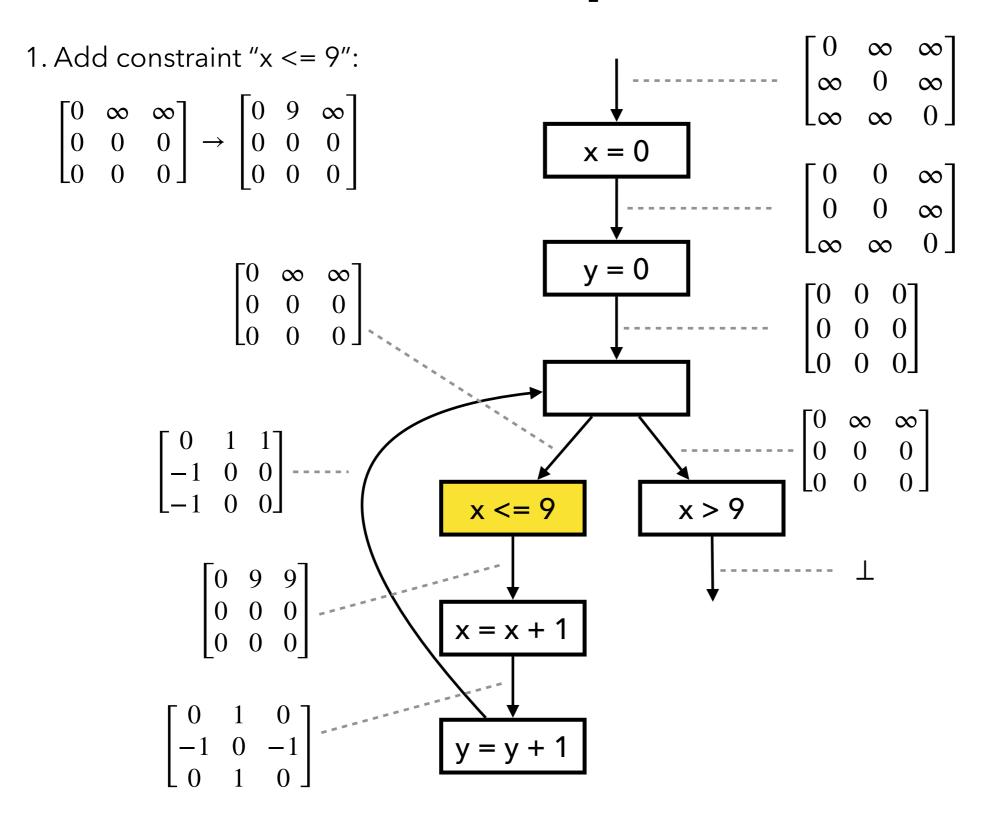


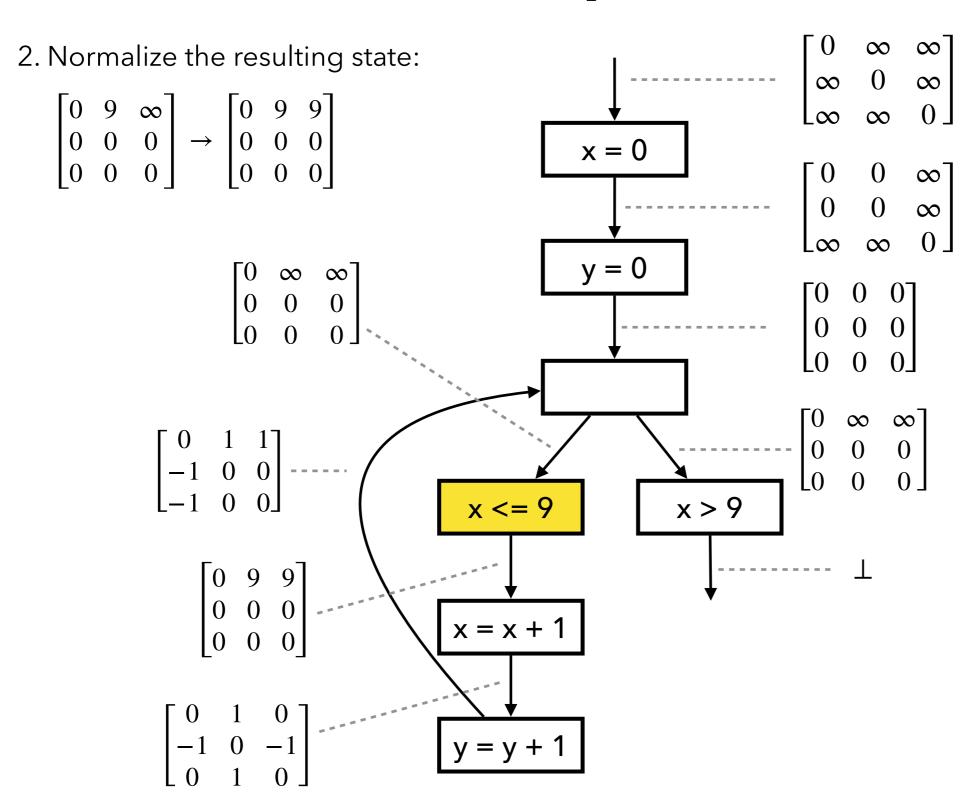


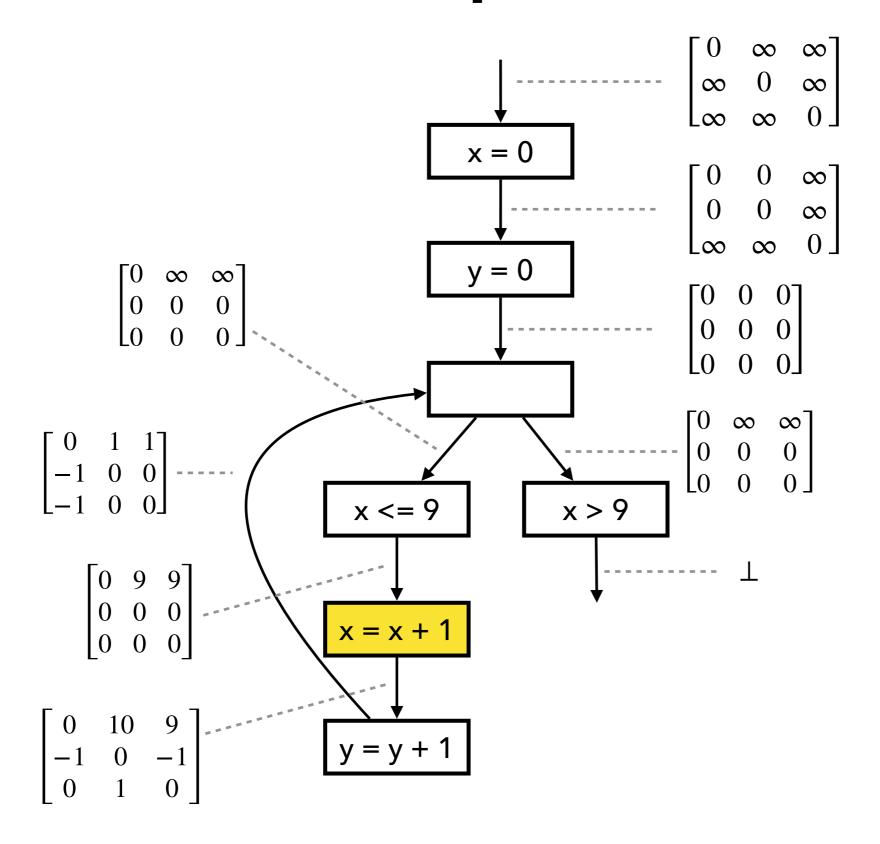


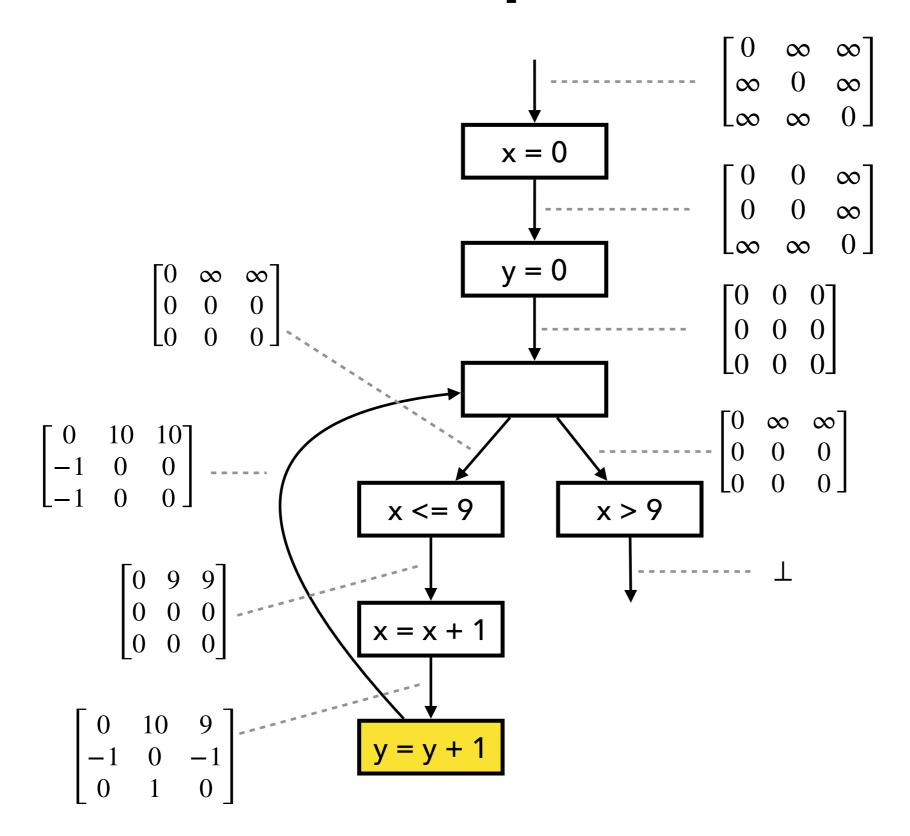


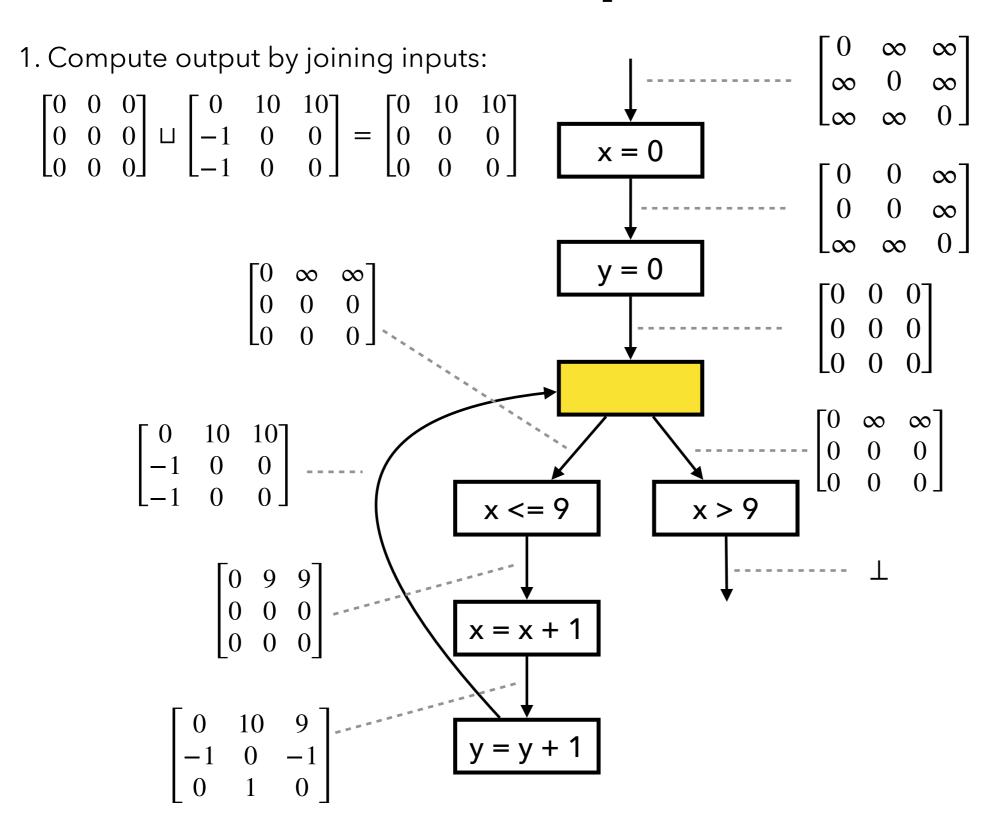


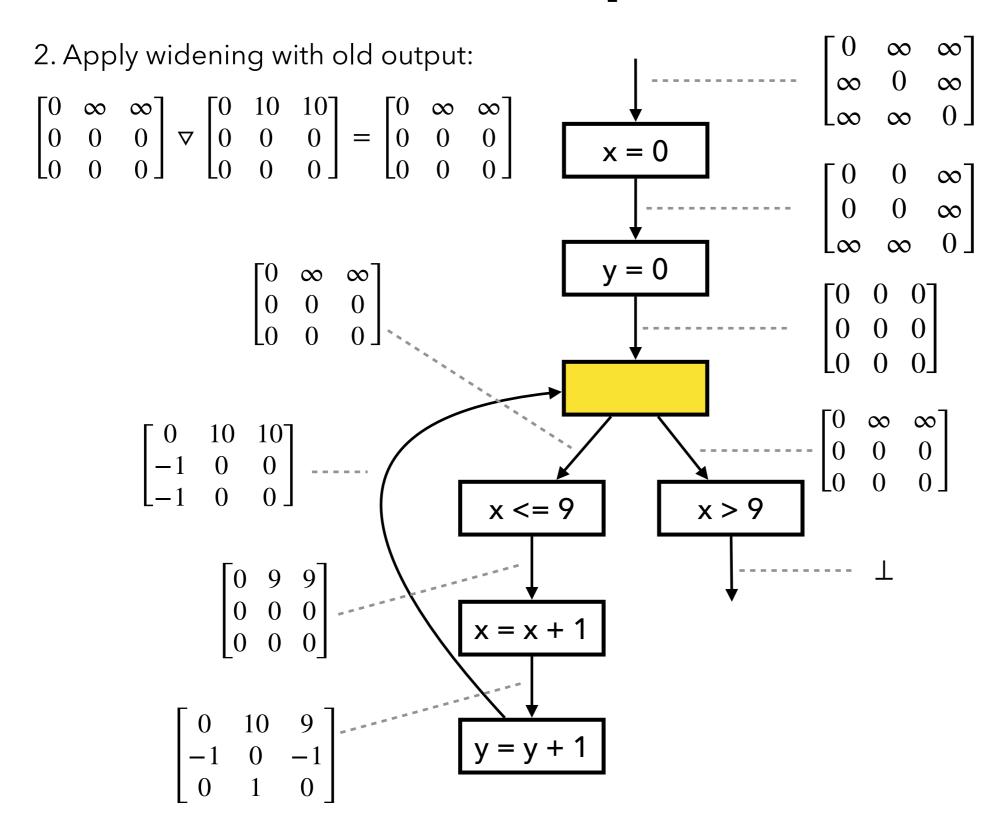


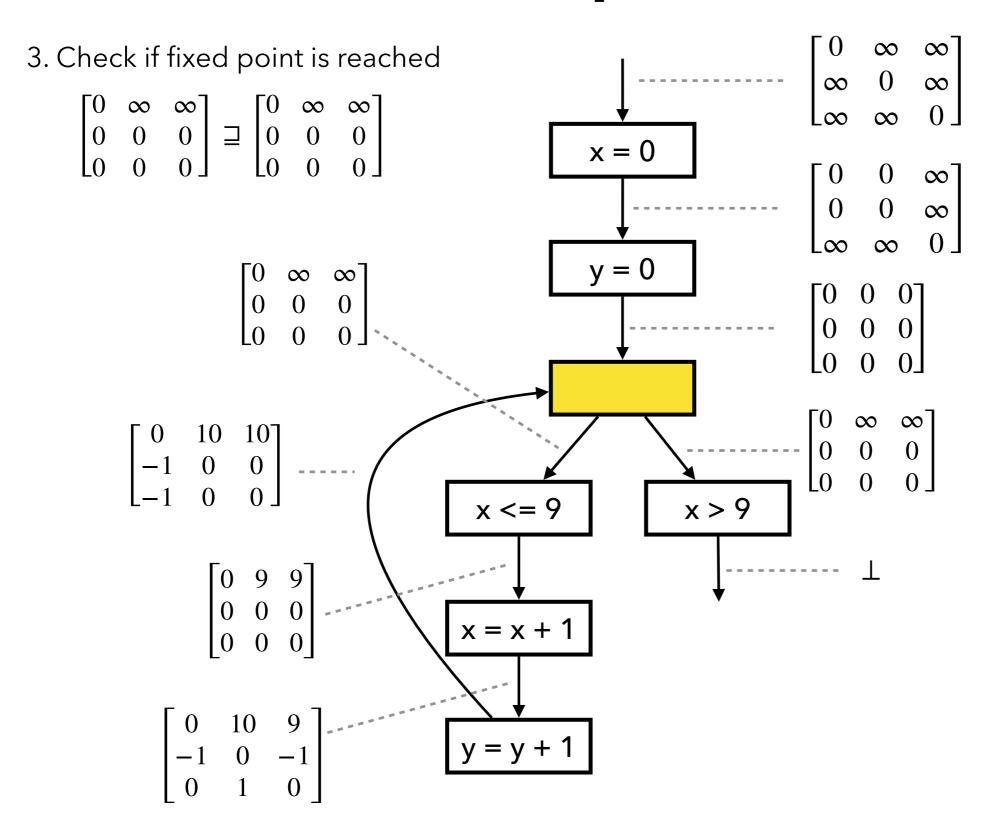


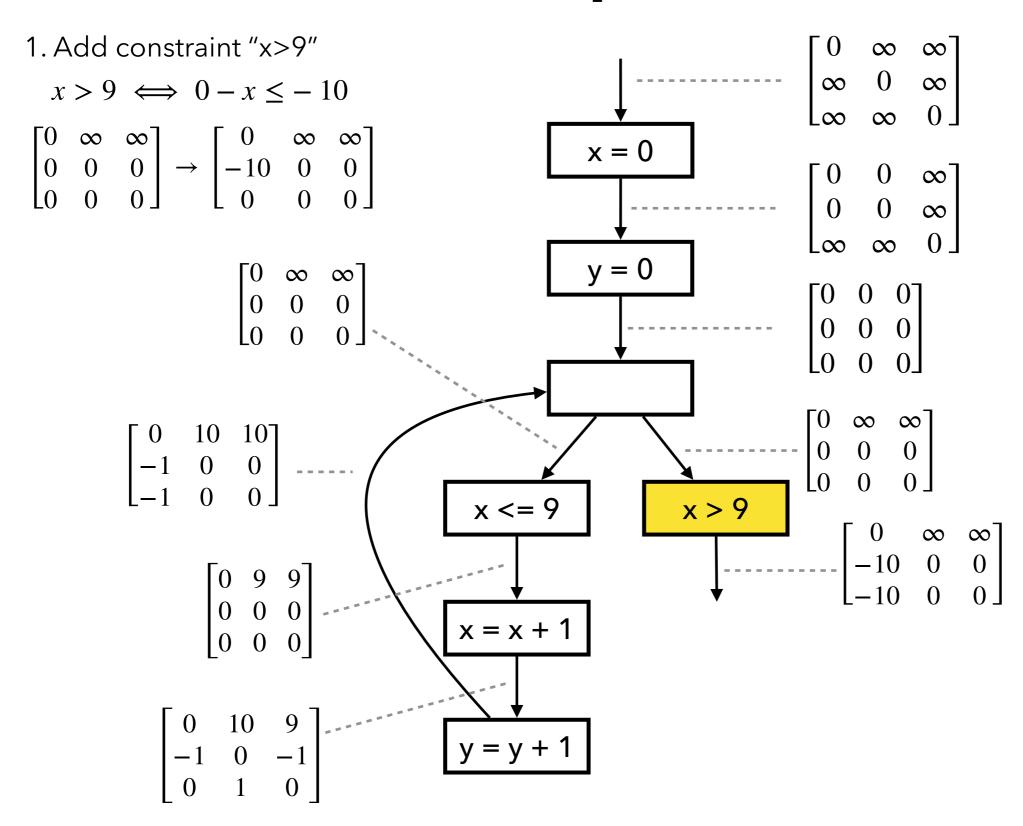


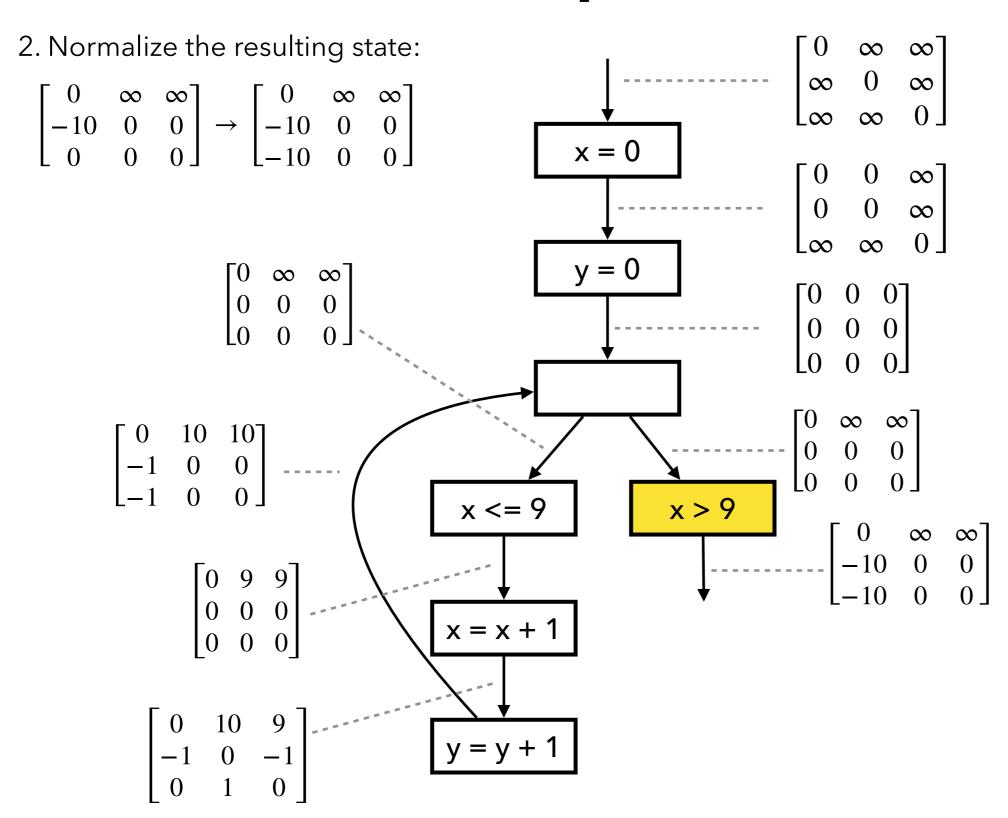


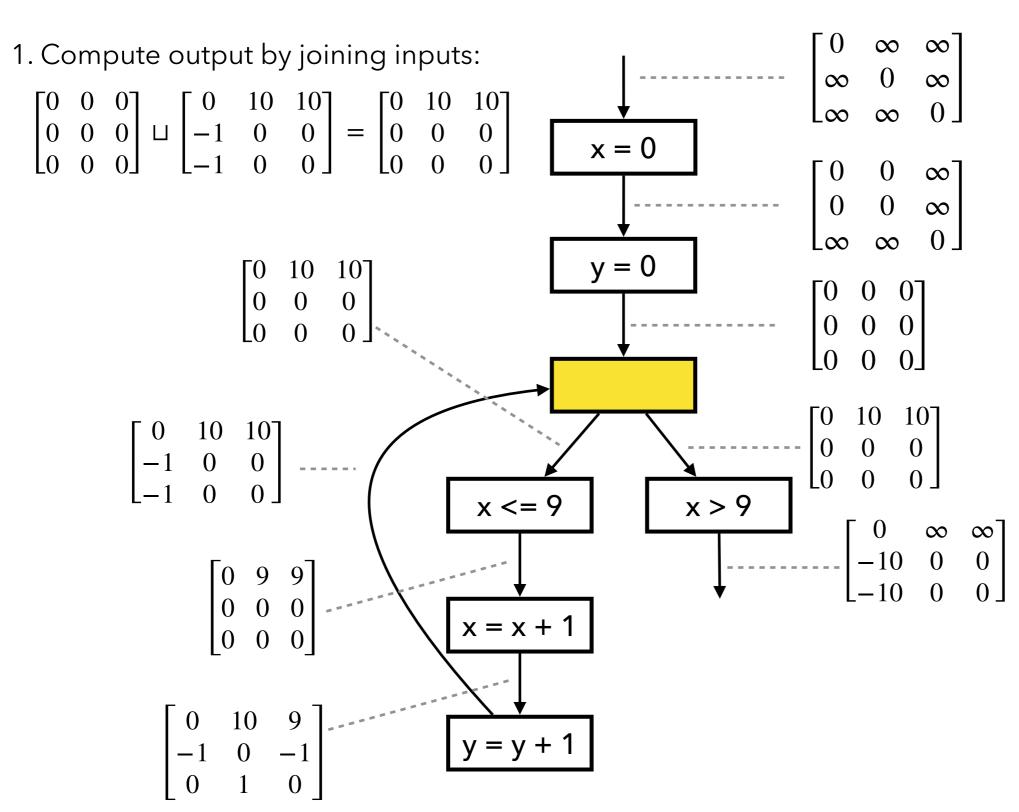


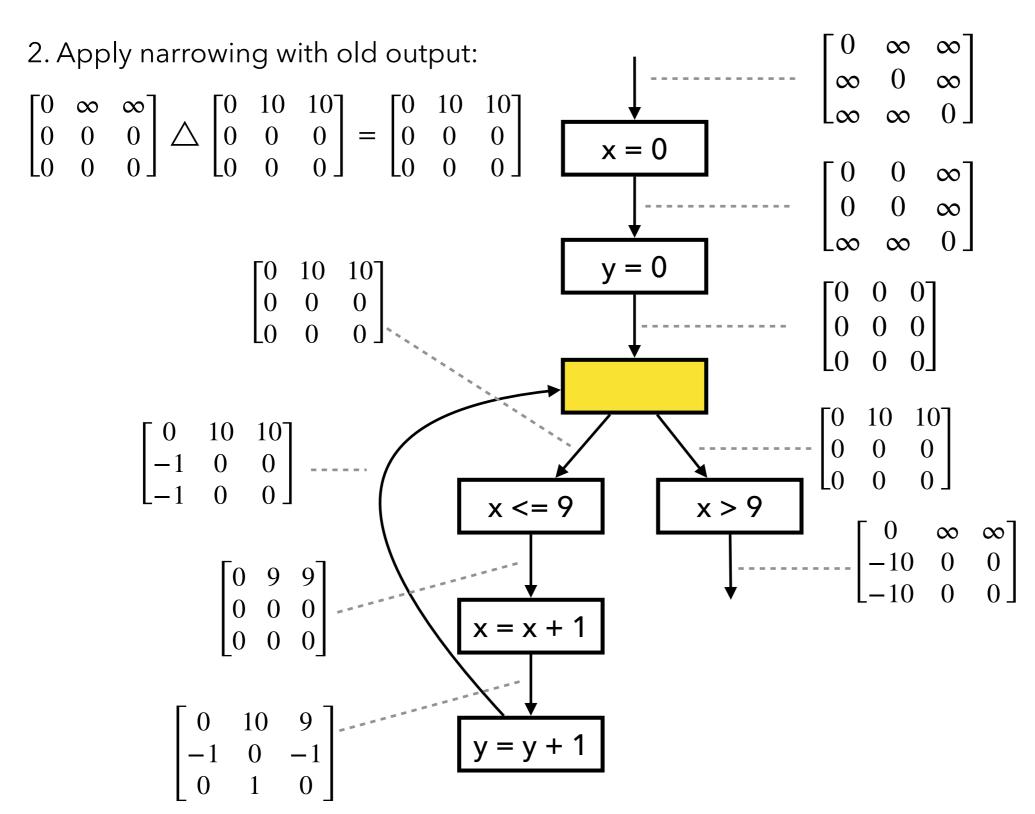


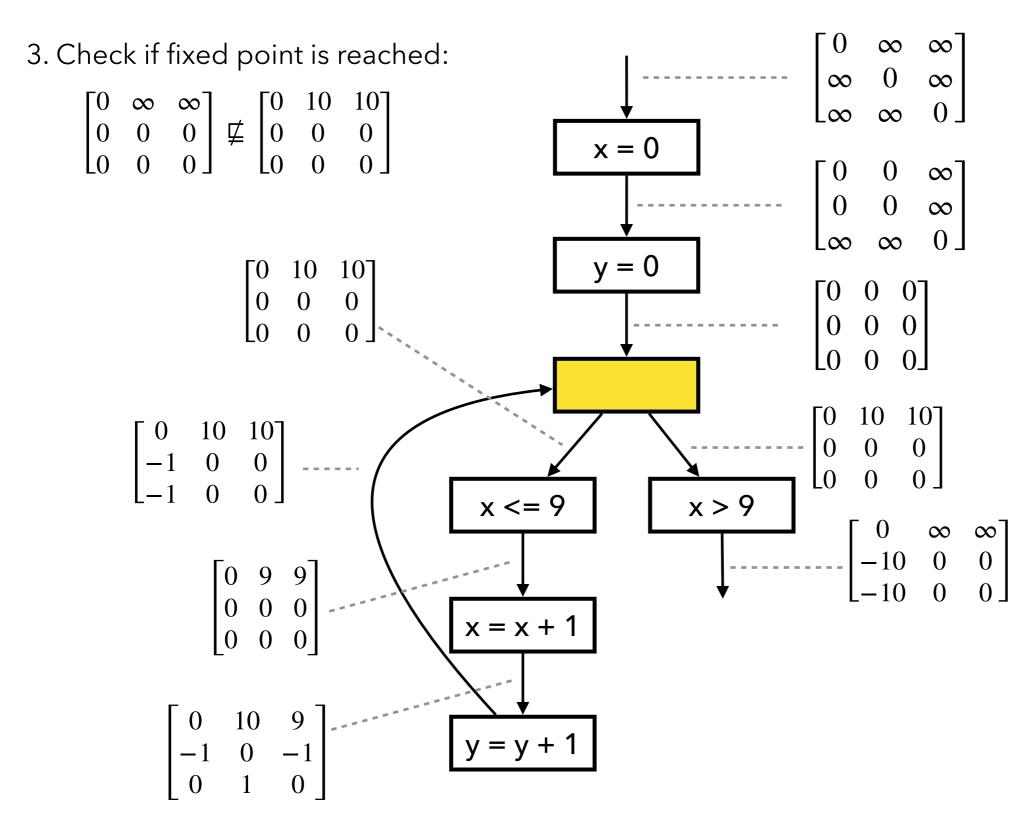


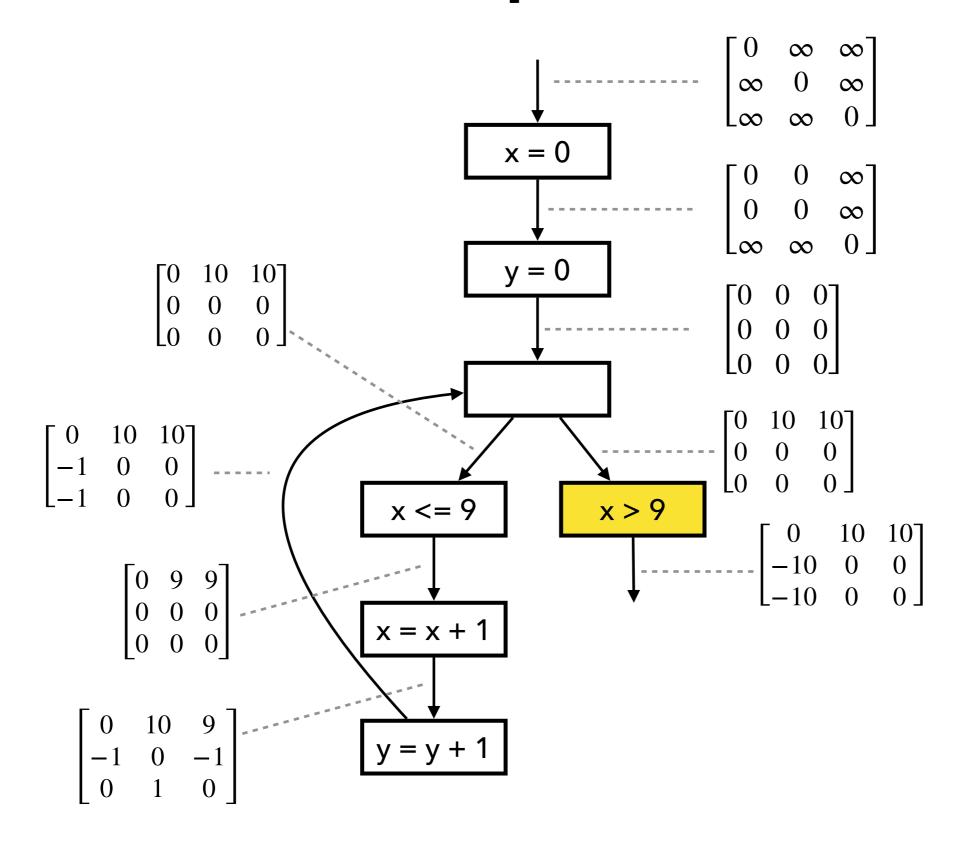












#### Example

Describe how the zone analysis works for the following example.

```
// a >= 0, b >= 0
q = 0;
r = a;
while (r >= b) {
   r = r - b;
   q = q + 1;
}
assert(q >= 0);
assert(r >= 0);
```

