

# 양자 프로그램 자동 합성

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02/06/2023@컴퓨터시스템소사이어티 동계학술대회

# 소프트웨어 분석 연구실@Korea Univ.

- **Members:** 10 PhD and 5 MS students
- **Research areas:** programming languages (PL), software engineering (SE), software security
  - program analysis and testing
  - program synthesis and repair
- **Publication:** top-venues in PL, SE, and Security:
  - **PL:** POPL('22), PLDI('20,'14,'12), OOPSLA('15,'17a,'17b,'18a,'18b,'19,'20,'23)
  - **SE:** ICSE('17,'18,'19,'20,'21'22a,'22b,'23a,'23b,'23c), FSE('18,'19,'20,'21), ASE('18)
  - **Security:** IEEE S&P('17,'20), USENIX Security('21)



<http://prl.korea.ac.kr>

# 연구 목표

- SW 오류 = 사회 모든 영역에서 발생



금융거래SW(2012)



자율주행SW(2017)



의료SW(2018)



블록체인SW(2020)

- SW 오류 = 사회경제적 비용 1.7조 달러/년



**606**  
software fails



**\$1.7**  
trillion



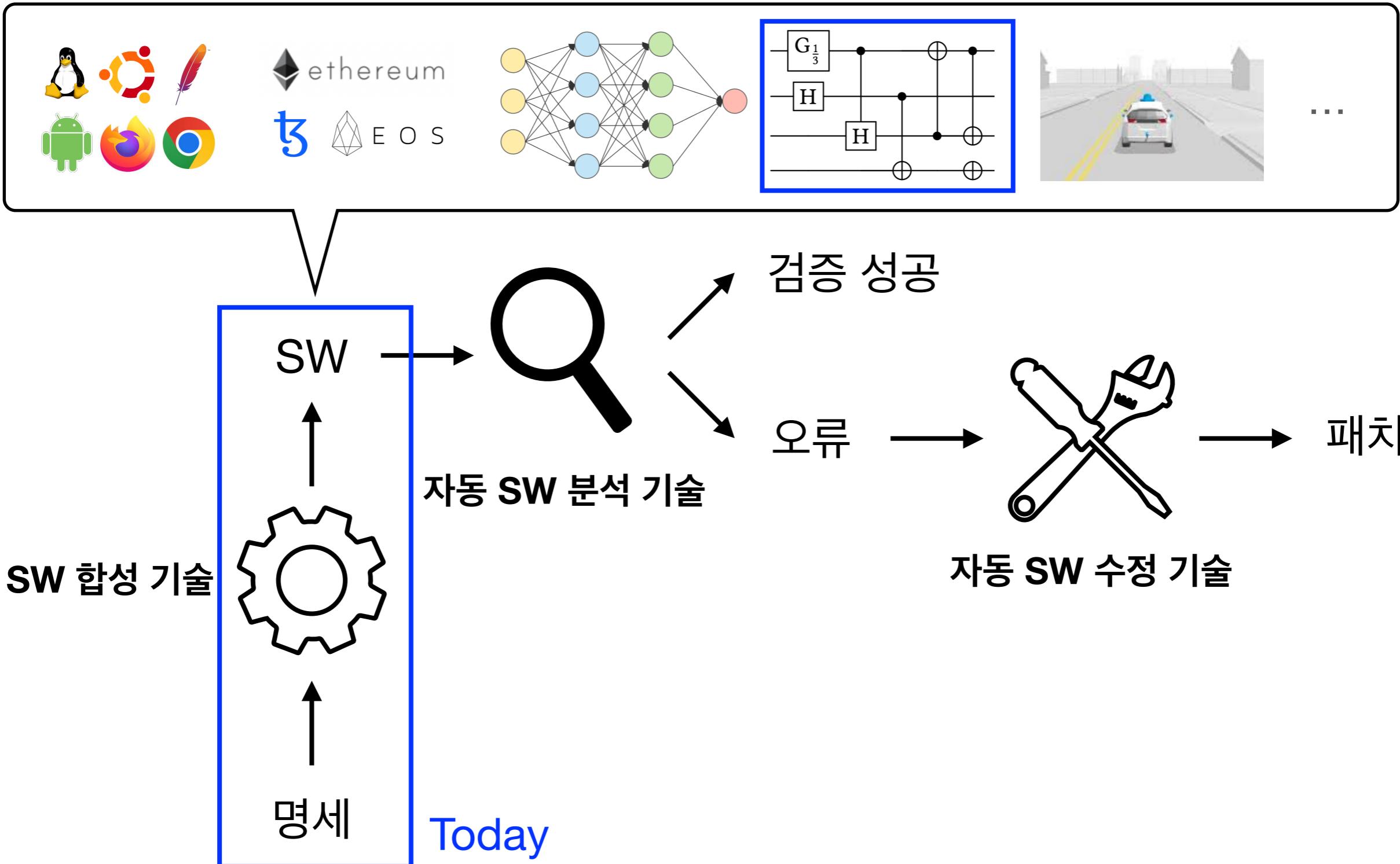
**3.6 billion**  
affected users



**268 years**  
in downtime

Software fail watch (5th edition). 2017

# Our Research



# 프로그램 합성 (Program Synthesis)

- 명세로부터 프로그램 코드를 자동 생성

$\text{reverse}(12) = 21, \text{reverse}(123) = 321$

```
reverse(n) {  
    r := 0;  
    while ( [ ] ) {  
        [ ]  
    };  
    return r;  
}
```



2.5s

```
reverse(n) {  
    r := 0;  
    while ( [ n > 0 ] ) {  
        x := n % 10;  
        r := r * 10;  
        r := r + x;  
        n := n / 10;  
    };  
    return r;  
}
```

# 프로그램 합성 (Program Synthesis)

- 명세로부터 프로그램 코드를 자동 생성

```
# Write a python function to toggle all even bits of a given number.
# Your code should pass these tests:
#
# assert even_bit_toggle_number(10) == 0
# assert even_bit_toggle_number(20) == 30
# assert even_bit_toggle_number(30) == 20

def even_bit_toggle_number(n):
    count = 0
    res = 0
    temp = n
    while temp > 0:
        if count % 2 == 1:
            res |= 1 << count
        count += 1
        temp >>= 1
    return n ^ res
```

# 양자 회로 자동 합성

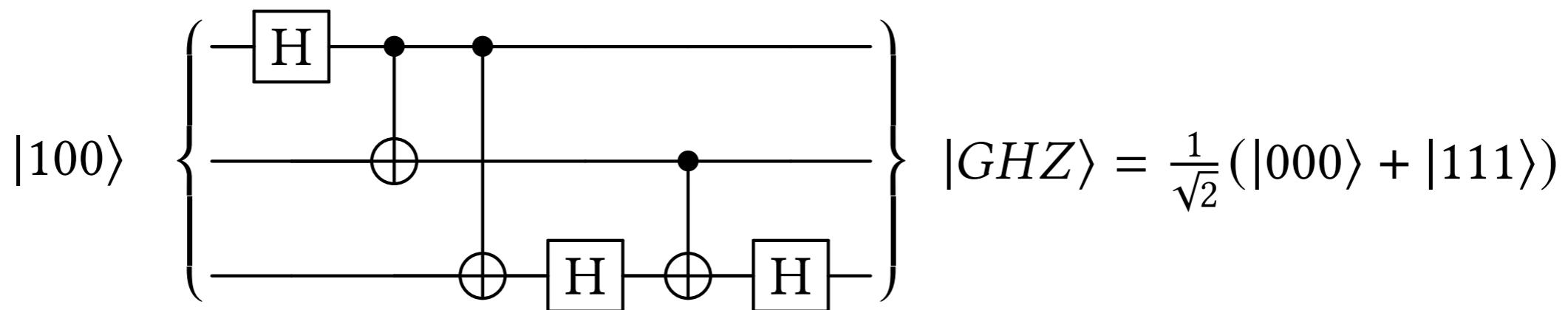
입출력 명세

$$|100\rangle \mapsto \frac{1}{\sqrt{2}}(|000\rangle + |111\rangle)$$

컴포넌트 게이트

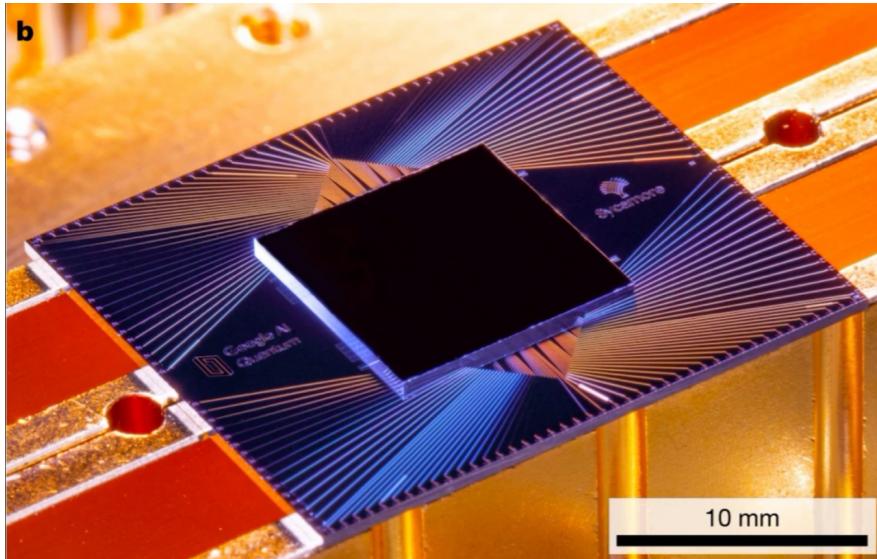
$$H, CNOT$$

양자 회로 합성기



# 양자 컴퓨터

- 하드웨어

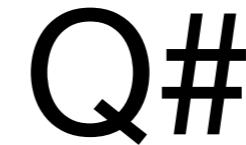


Google Sycamore (2019)



IBM Hummingbird, Eagle, Osprey (2020-2022)

- 소프트웨어

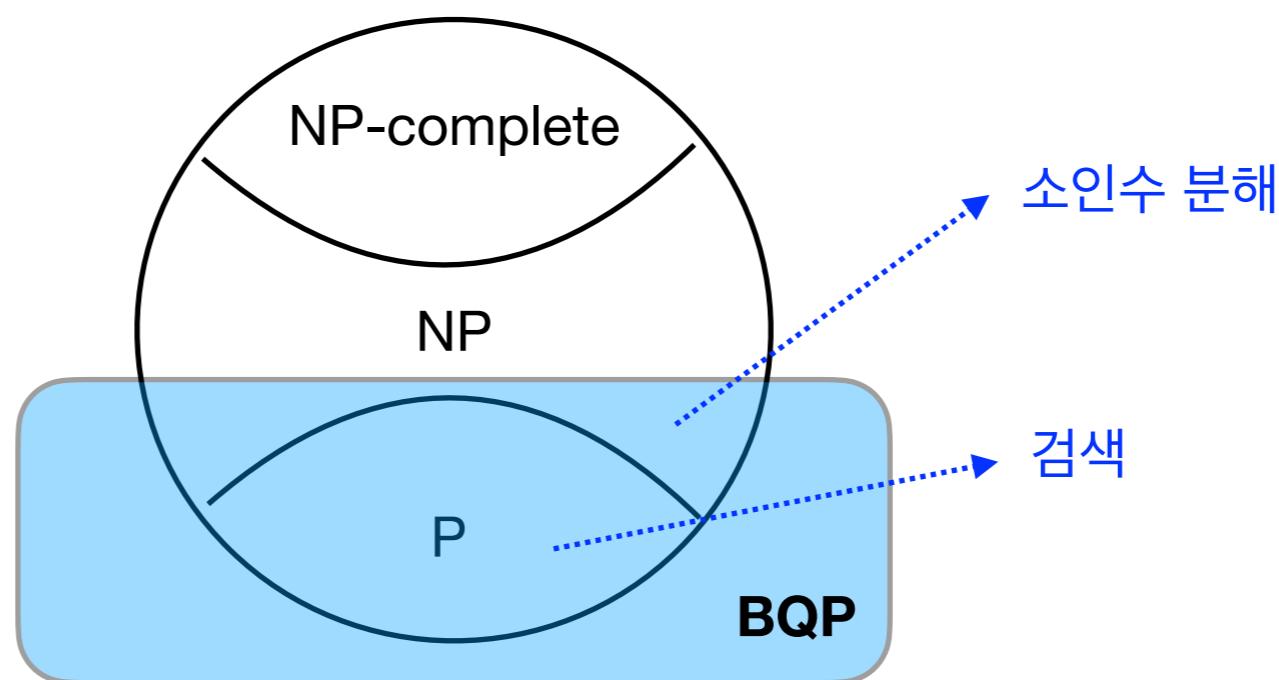


# 양자 컴퓨터

- 고전 컴퓨터의 일반화

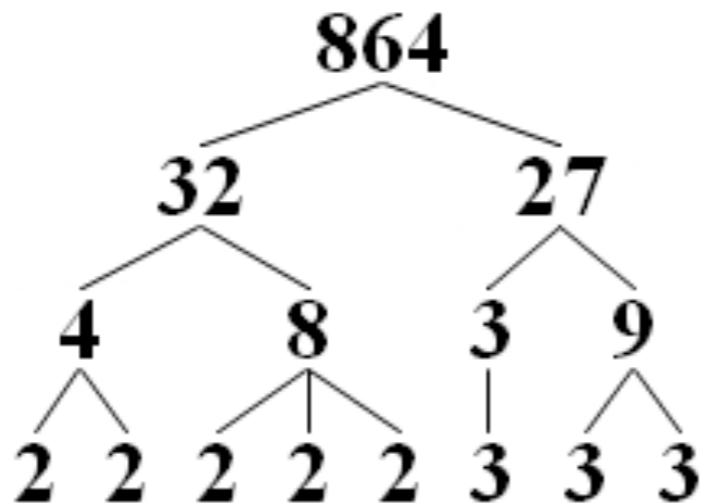


- 가능성 & 한계



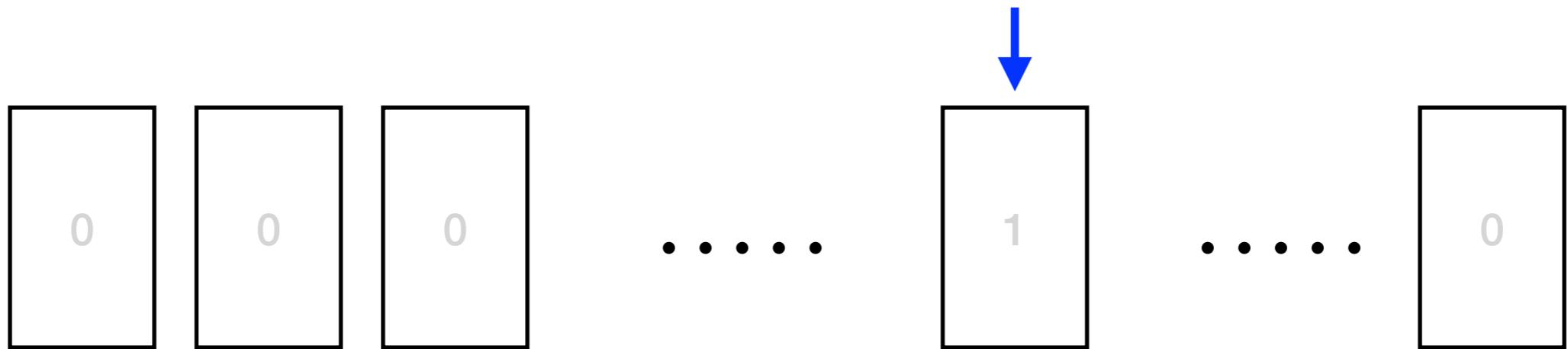
# 적용 사례

(1) 소인수분해 [Shor 1995]



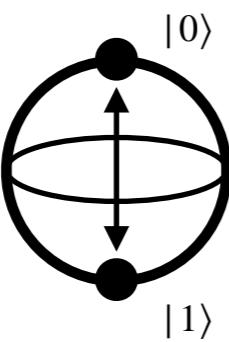
$O(2^N)$  vs  $O(N^3)$

(2) 검색 [Grover 1996]



$O(N)$  vs  $O(\sqrt{N})$

# 큐비트



- 비트의 일반화 (0과 1의 중첩 – 선형 결합)

$$|\psi\rangle = \alpha_0 |0\rangle + \alpha_1 |1\rangle$$

- $\alpha_1, \alpha_2 \in \mathbb{C}$ : 확률 진폭 (probability amplitude)

$$|\alpha_1|^2 + |\alpha_2|^2 = 1$$

- $|0\rangle$
- $|1\rangle$
- $\frac{1}{\sqrt{2}}|0\rangle + \frac{1}{\sqrt{2}}|1\rangle$
- $\frac{1}{\sqrt{2}}|0\rangle - \frac{1}{\sqrt{2}}|1\rangle$
- $\sqrt{\frac{3}{4}}|0\rangle + \frac{i}{\sqrt{4}}|1\rangle$

# N개 큐비트

- 고전 비트 2개로 만들어지는 상태

00, 01, 10, 11

- 큐비트 2개 = 고전 2-bit 상태들의 중첩

$$|\psi\rangle = \alpha_0|00\rangle + \alpha_1|01\rangle + \alpha_2|10\rangle + \alpha_3|11\rangle \quad \left( \sum_i |\alpha_i|^2 = 1 \right)$$

- 큐비트  $N$ 개 = 고전  $N$ -bit 상태들의 중첩

$$|\psi\rangle = \sum_{x \in \{0,1\}^N} \alpha_x |x\rangle$$

2 $N$ 개 확률 진폭을 “저장”

# 병렬처리: 고전 컴퓨터 vs. 양자 컴퓨터

- $f: \{0,1\} \rightarrow \{0,1\}$

$$\begin{array}{c} f(0) \\ \text{vs.} \\ f(1) \end{array} \qquad f\left(\frac{|0\rangle + |1\rangle}{\sqrt{2}}\right)$$

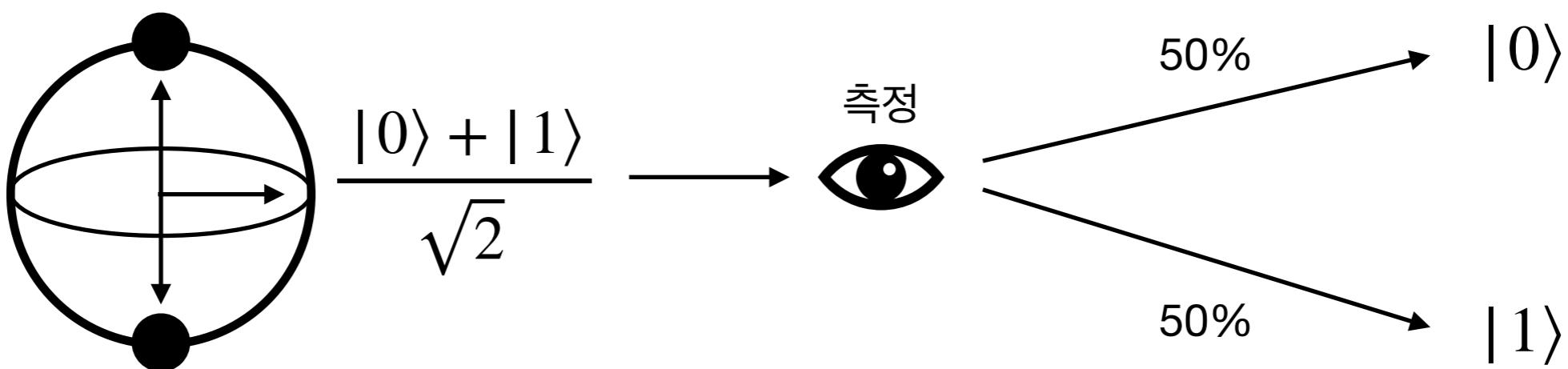
- $f: \{0,1\}^N \rightarrow \{0,1\}$

$$\begin{array}{c} f(0) \\ f(1) \\ \vdots \\ f(2^N - 1) \end{array} \qquad \text{vs.} \qquad f\left(\frac{|0\rangle + |1\rangle + \dots + |2^N - 1\rangle}{\sqrt{2^N}}\right)$$

# 측정과 붕괴

$$|\psi\rangle = \alpha_0 |0\rangle + \alpha_1 |1\rangle$$

- 큐비트의 내부 상태( $\alpha_1, \alpha_2$ )는 알 수 없고, **고전 정보만 관측 가능**
- 큐비트  $|\psi\rangle$ 를 측정하면,
  - $|\alpha_0|^2$ 의 확률로 0이 관측되고  $|\psi\rangle = |0\rangle$ 로 붕괴
  - $|\alpha_1|^2$ 의 확률로 1이 관측되고  $|\psi\rangle = |1\rangle$ 로 붕괴



# N개 큐비트 측정

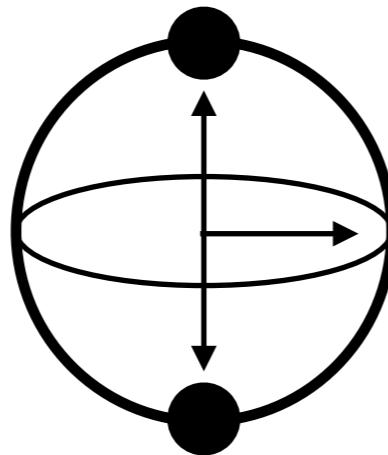
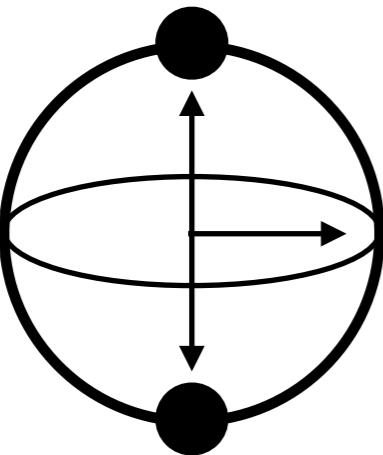
$$|\psi\rangle = \alpha_0|00\rangle + \alpha_1|01\rangle + \alpha_2|10\rangle + \alpha_3|11\rangle$$

- 상태  $|\psi\rangle$ 를 측정하면,
  - $|\alpha_0|^2$ 의 확률로 00이 관측되고  $|\psi\rangle = |00\rangle$ 로 붕괴
  - $|\alpha_1|^2$ 의 확률로 01이 관측되고  $|\psi\rangle = |01\rangle$ 로 붕괴
  - $|\alpha_2|^2$ 의 확률로 10이 관측되고  $|\psi\rangle = |10\rangle$ 로 붕괴
  - $|\alpha_3|^2$ 의 확률로 11이 관측되고  $|\psi\rangle = |11\rangle$ 로 붕괴
- 예:

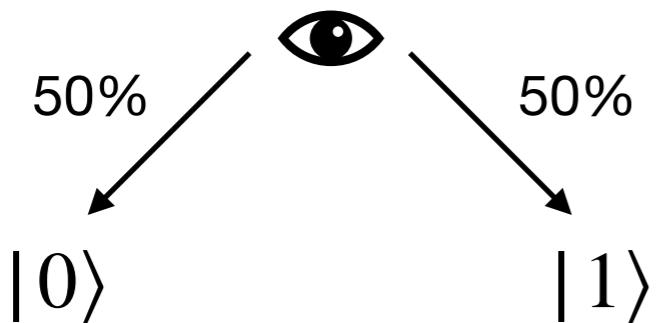
$$|\psi\rangle = \frac{|00\rangle + |01\rangle + |10\rangle + |11\rangle}{\sqrt{4}}$$

# 얽힘 (Entanglement)

- 얽히지 않은 상태 = 개별 큐비트들이 단순 결합된 상태



$$\frac{|0\rangle + |1\rangle}{\sqrt{2}} \times \frac{|0\rangle + |1\rangle}{\sqrt{2}} = \frac{|00\rangle + |01\rangle + |10\rangle + |11\rangle}{2}$$

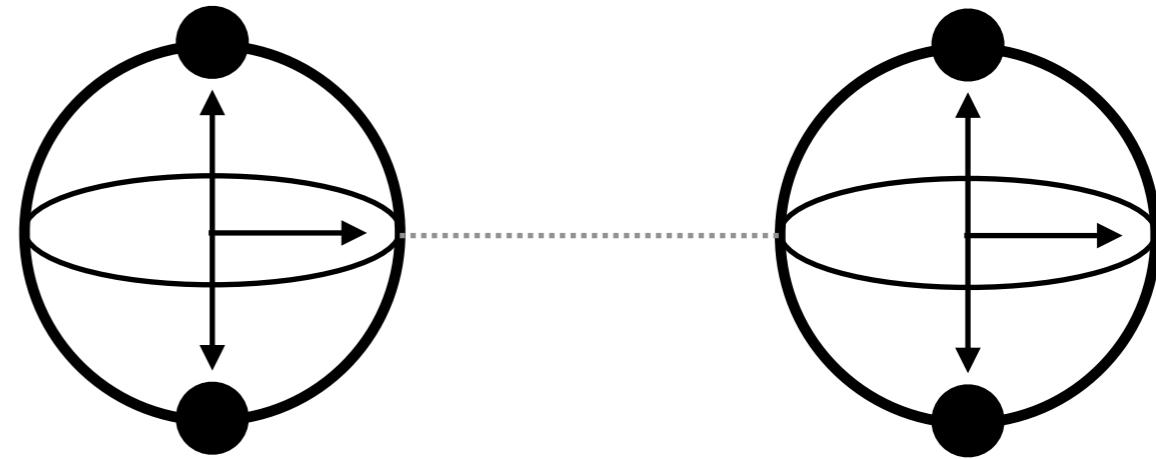


- 하나의 큐비트 관찰 결과가 다른 큐비트 상태에 영향을 주지 않음

# 얽힘 (Entanglement)

- 얽힌 상태 = 개별 큐비트 상태들로 분리 불가능

$$\frac{|00\rangle + |11\rangle}{\sqrt{2}}$$

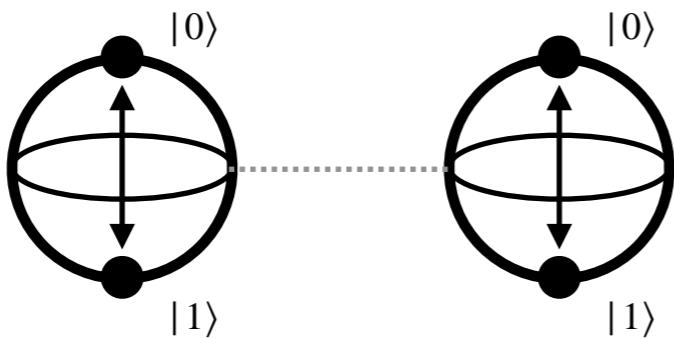


$$\neq (\alpha_0 |0\rangle + \alpha_1 |1\rangle) \times (\beta_0 |0\rangle + \beta_1 |1\rangle)$$

- 하나의 큐비트 관찰이 다른 큐비트 상태에 영향을 줌
  - 하나의 큐비트에서 0을 관측하면 다른 큐비트 상태는  $|0\rangle$ 로 결정
  - 하나의 큐비트에서 1을 관측하면 다른 큐비트 상태는  $|1\rangle$ 로 결정

# EPR 패러독스

1. 두 큐비트



를  $\frac{|00\rangle + |11\rangle}{\sqrt{2}}$  상태로 준비

2. 두 큐비트를 서로 멀리 떨어지도록 이동



$|0\rangle$   
 $|1\rangle$

$$\frac{|00\rangle + |11\rangle}{\sqrt{2}}$$

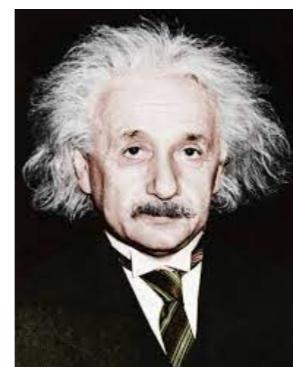


$|0\rangle$   
 $|1\rangle$



3. 측정

“Spooky action at a distance”



# 양자 게이트

$$\alpha_0 |0\rangle + \alpha_1 |1\rangle \xrightarrow{\quad X \quad} \alpha_1 |0\rangle + \alpha_0 |1\rangle$$

- 큐비트 상태 = 벡터

$$|\psi\rangle = \alpha_0 |0\rangle + \alpha_1 |1\rangle = \alpha_0 \begin{bmatrix} 1 \\ 0 \end{bmatrix} + \alpha_1 \begin{bmatrix} 0 \\ 1 \end{bmatrix} = \begin{bmatrix} \alpha_0 \\ \alpha_1 \end{bmatrix}$$

- 게이트 연산 = 행렬 곱

$$\begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \times \begin{bmatrix} \alpha_0 \\ \alpha_1 \end{bmatrix} = \begin{bmatrix} \alpha_1 \\ \alpha_0 \end{bmatrix}$$

# 양자 게이트

$$\alpha_0 |0\rangle + \alpha_1 |1\rangle \xrightarrow{\quad} \boxed{Z} \xrightarrow{\quad} \alpha_0 |0\rangle - \alpha_1 |1\rangle$$

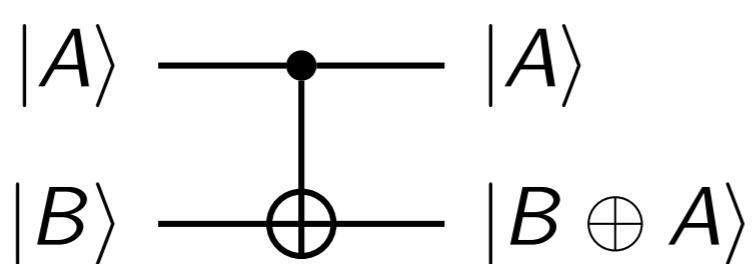
$$\begin{bmatrix} 1 & 0 \\ 0 & -1 \end{bmatrix}$$

$$\alpha_0 |0\rangle + \alpha_1 |1\rangle \xrightarrow{\quad} \boxed{H} \xrightarrow{\quad} \alpha_0 \left( \frac{|0\rangle + |1\rangle}{\sqrt{2}} \right) + \alpha_1 \left( \frac{|0\rangle - |1\rangle}{\sqrt{2}} \right)$$

$$\frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix}$$

# 양자 게이트

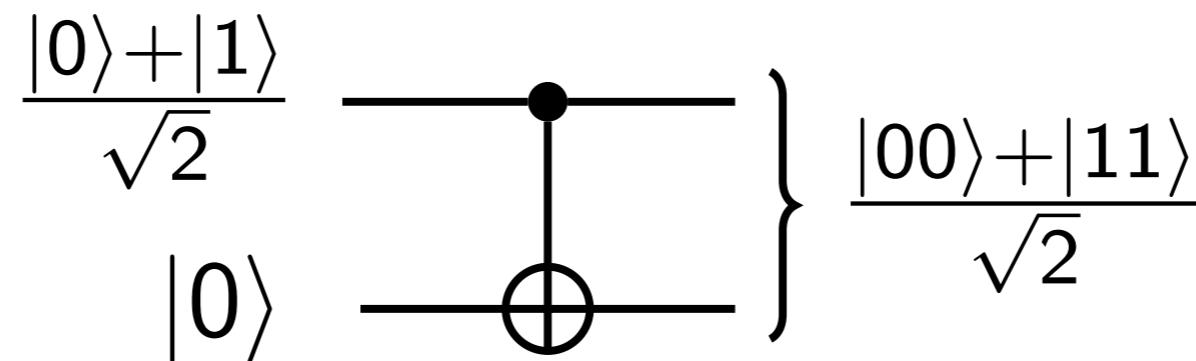
- CNOT (Controlled-Not) 게이트



$$\text{CNOT} = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{bmatrix}$$

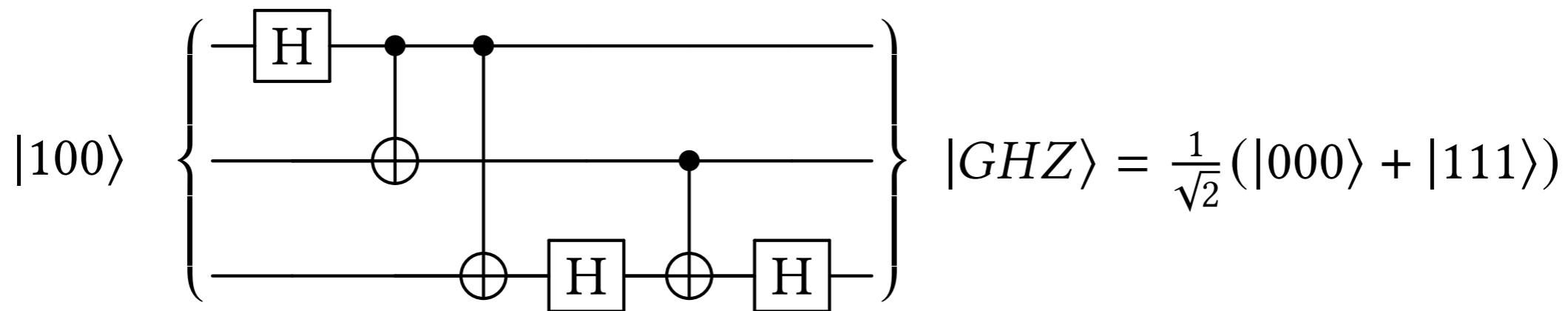
$$\begin{aligned} \text{CNOT}(\alpha_{00} |00\rangle + \alpha_{01} |01\rangle + \alpha_{10} |10\rangle + \alpha_{11} |11\rangle) \\ = \alpha_{00} |00\rangle + \alpha_{01} |01\rangle + \alpha_{10} |11\rangle + \alpha_{11} |10\rangle \end{aligned}$$

- CNOT: 얹힌 상태를 만드는데 사용



# 양자 회로

- 양자 게이트들의 조합



```
circ = QuantumCircuit(3)
circ.h(0)
circ.cx(0, 1)
circ.cx(0, 2)
circ.h(2)
circ.cx(1, 2)
circ.h(2)
```



# 양자 프로그래밍의 어려움

- 고전 프로그래밍(e.g. Python, C, Java, ...)과는 전혀 다른 방식

- 양자 프로그래밍은 비직관적

- 데이터 값 : 벡터

$$\begin{array}{c} \text{---} \\ \boxed{H} \\ \text{---} \end{array} = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix}$$

- 프로그램 연산자 : 행렬

$$\begin{array}{c} \text{---} \\ \bullet \\ \text{---} \\ \text{---} \\ \oplus \\ \text{---} \end{array} = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix}$$

# 양자 회로 자동 합성 (OOPSLA 2023)

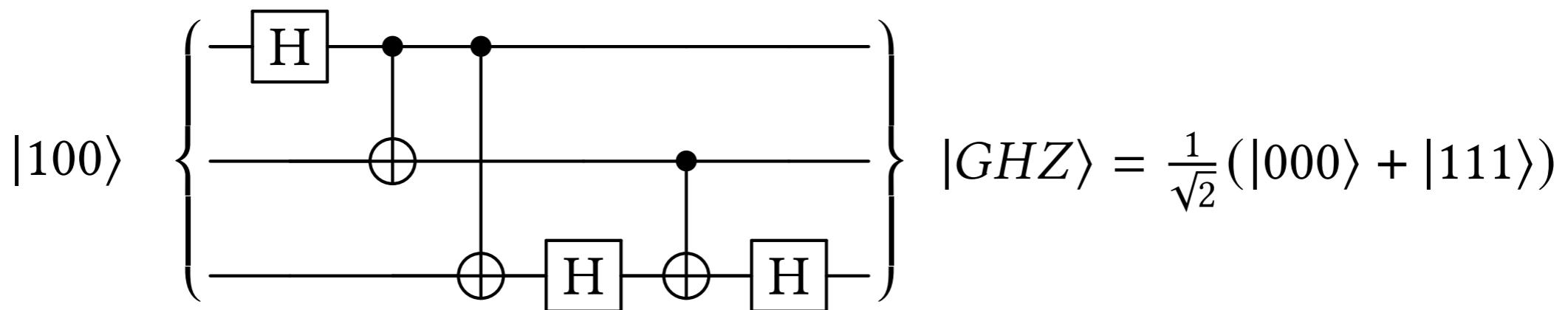
입출력 명세

$$|100\rangle \mapsto \frac{1}{\sqrt{2}}(|000\rangle + |111\rangle)$$

컴포넌트 게이트

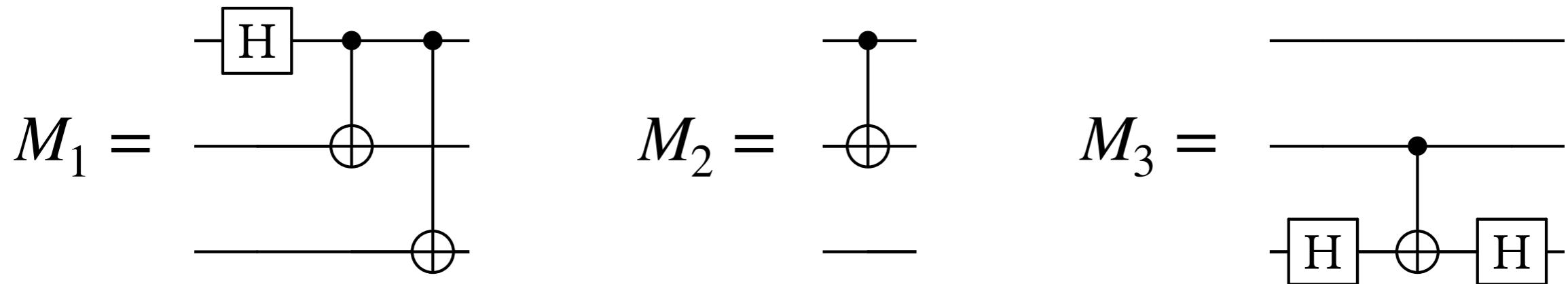
$$H, CNOT$$

양자 회로 합성기

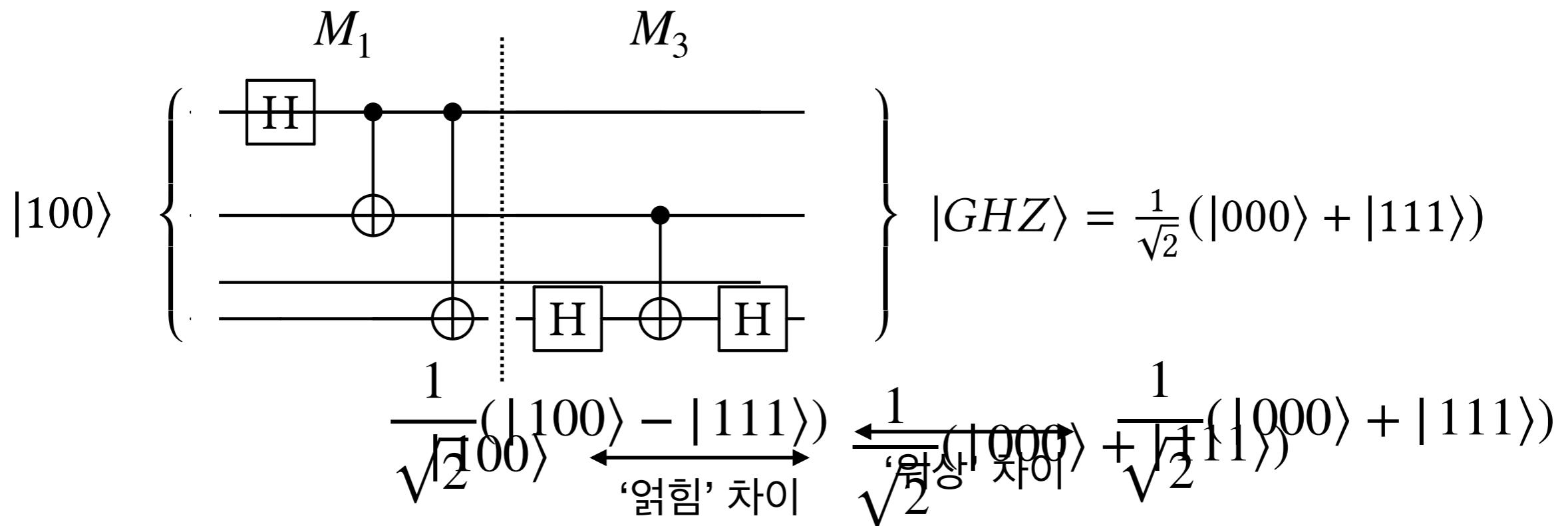


# 아이디어: 모듈 레벨 양자 회로 합성

1. 모듈 = 게이트 시퀀스 (up to some fixed length)



2. 명세와의 특성 차이를 줄이는 방향으로 모듈 쌓기



# 합성 알고리즘의 안전성 (Soundness)

“특정 가정하에서 모듈 기반 합성 알고리즘이 항상 정답을 찾아냄”

THEOREM 4.17. Let  $E = \{(|in\rangle, |out\rangle)\}$  be an example and  $C^* = M_1; \dots; M_k$  ( $M_i \in \mathcal{M}$  and attribute of each  $M_i$  is not IDENTITY) be the solution circuit to be synthesized such that  $C^*(|in\rangle) = |out\rangle$ . Suppose  $C^*$  is monotonically decreasing (by input  $|in\rangle$ ). Then, for any prefix  $C = M_1; \dots; M_{l-1}$  ( $l \leq k$ ) of  $C^*$ ,

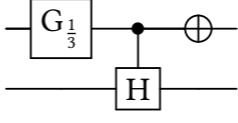
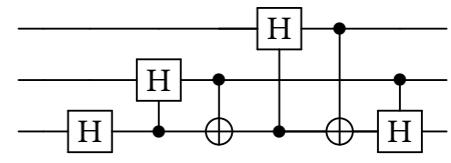
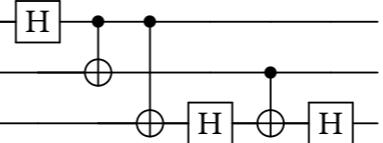
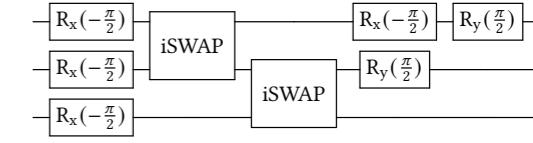
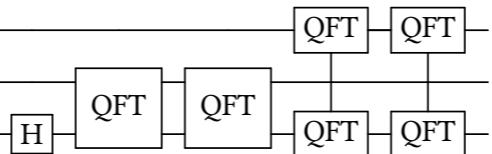
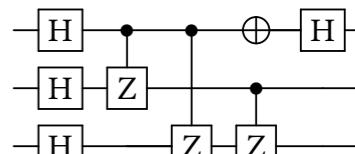
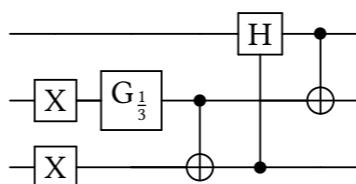
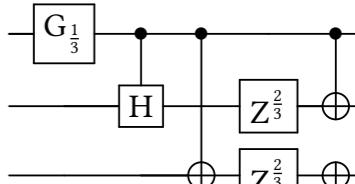
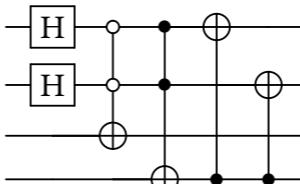
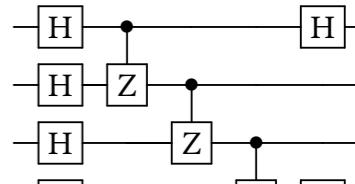
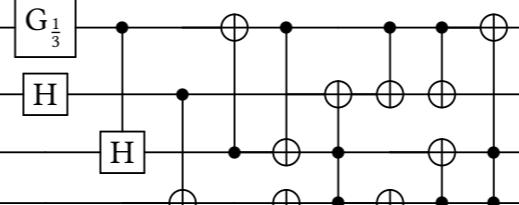
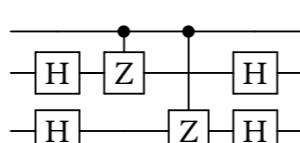
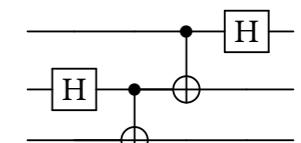
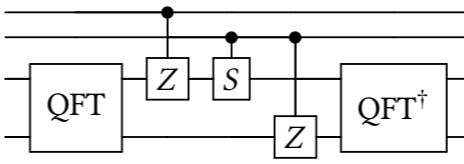
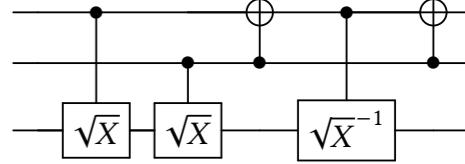
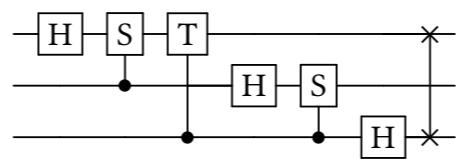
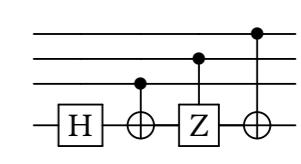
$$\text{is\_gap\_filled}(C, M_l, (|in\rangle, |out\rangle)) = \text{True}.$$

PROOF. Let  $|\psi\rangle = C|in\rangle = M_{l-1} \cdots M_1|in\rangle$  be the output vector of  $C$ . By definition  $M_k \cdots M_l|\psi\rangle = |out\rangle$  and thus

$$\text{Att}_{|\psi\rangle}(M_k M_{k-1} \cdots M_l) = |out\rangle \ominus |\psi\rangle = \text{gap\_att}_{|in\rangle, |out\rangle}(C).$$

Since  $C^*$  (and so  $C$ ) is decreasing, by Lemma 4.16  $\text{Att}_{|\psi\rangle}(M_k M_{k-1} \cdots M_l) = \text{Att}_{|\psi\rangle}(M_l)$ . Therefore,  $\text{gap\_att}_{|in\rangle, |out\rangle}(C) = \text{Att}_{|\psi\rangle}(M_l)$ , which is satisfying the criterion.  $\square$

# 벤치마크

Type	ID	Circuit	ID	Circuit
State Preparation	three_superpose		M_valued	
	GHZ_from_100		GHZ_by_iSWAP	
	GHZ_by_QFT		GHZ_Game	
	W_orthog		W_phased	
	W_four		cluster	
	bit_measure			
	flip		teleportation	
	Multi IO			
	draper		toffoli_by_sqrt_X	
	QFT		indexed_bell	

## 게이트 레벨 합성 알고리즘

## 모듈 레벨 합성 알고리즘

ID	Base <sub>no_prune</sub>	Base	Ours <sub>no_prune</sub>	Ours	Spd-up
three_superpose	0.14	0.12	0.12	0.09	1x
M_valued	1764.75	1126.79	666.25	3.89	290x
GHZ_from_100	106.81	48.16	—	0.47	102x
GHZ_by_iSWAP	—	—	690.67	2.19	-
GHZ_by_QFT	116.90	101.65	101.17	39.26	3x
GHZ_Game	—	2305.71	4.51	0.57	4058x
W_orthog	2927.20	2075.06	248.23	2.43	854x
W_phased	—	—	258.56	5.43	-
W_four	—	—	2851.10	254.88	-
cluster	—	—	3560.38	8.91	-
bit_measure	—	—	—	—	-
flip	18.56	3.95	0.76	0.83	5x
teleportation	2.02	1.30	1.35	1.35	1x
indexed_bell	14.67	11.66	1.60	1.52	8x
toffoli_by_√X	956.29	716.28	306.10	264.66	3x
QFT	—	—	—	220.87	-
draper	—	—	933.47	737.99	-
Avg. (excluding —)	656.37	639.07	687.45	96.58	20x

# Summary

- 양자 컴퓨팅 및 프로그래밍 소개
- 양자 회로 자동 합성 소개
- 양자 SW + X
  - 양자 SW 합성, 최적화, 분석, 검증, 수정, 언어 디자인

감사합니다!