# Lecture Notes

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Math 151A
Lecture Notes

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### 1 Lecture 1

The goal of this class is to solve mathematical problems with the help of computers.

### 1.1 Chapter Summaries

- 1. Computations in computers
  - How to store (real) numbers in the computer

Note. If the number has finitely many digits, then it is simple. What about numbers with infinitely many digits, i.e.  $\frac{1}{3}$ ? We have to truncate or round, and store an approximation with finitely many digits.

• How to perform computations

Note. From regular math, we know that  $\frac{1}{3} + \frac{1}{3} = \frac{2}{3}$ . However, in the computer, due to errors, we have  $\frac{1}{3} \oplus \frac{1}{3} = ? \oplus ? = ??$ .

- Errors
- 2. Find roots of f(x) = 0 using bisection, Newton's method, ...
  - Convergence
  - Convergence order (how fast it converges)
- 3. Polynomial interpolation
  - Approximate a function f(x) by a polynomial P(x), where  $f(x_i) = P(x_i)$  for finitely many  $x_i$
  - Accuracy of the polynomial approximations
- 4. Numerical differentiations and numerical integrations
  - Using the approximations from chapter 3, we can approximate using

$$f'(x^*) \approx P'(x^*) = \sum_{i=0}^{k} f(x_k)c_k,$$

and

$$\int_{a}^{b} f(x) dx \approx \int_{a}^{b} P(x) dx = \sum_{i=1}^{k} f(\overline{x_k}) \overline{c_k}.$$

- Error analysis
- 6.7. Solving linear systems of equations
  - Direct methods: Gaussian elimination (computationally expensive)
  - Iterative methods: (faster and cheaper)
  - Solution stability

#### 1.2 Round-off errors and computer arithmetics

There are three kinds of errors:

- Modeling Error: Occurs when we convert a problem from the real world into the mathematical world.
- Method Error: Occurs when we try to solve the mathematical problem numerically.
- Round-off error: Occurs when the computer gives an incorrect result with the correct algorithm (comes from storage and computation).

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#### 1.2.1 Storage

Infinite digit real numbers are *stored* as finite digit numbers, using the normalized decimal form of real numbers.

**Definition.** Normalized decimal form of a real number

For any  $y \in \mathbb{R}$ , we may write

$$y = \pm 0.d_1 d_2 d_3 \dots d_k d_{k+1} \dots \cdot 10^n$$

where  $0 < d_1 \le 9$ ,  $0 \le d_i \le 9$ , n are integers. For the particular case where y = 0, we write  $y = 0.0 \cdot 10^0$ .

**Definition.** Normalized machine numbers (Floating-point form)

Any machine number y can be written as

$$y = \pm 0.d_1 d_2 \cdots d_k \cdot 10^n,$$

where  $0 < d_1 \le 9$ ,  $0 \le d_i \le 9$ , n are integers.

We can think of the storage process as mapping normalized real numbers to normalized machine numbers. We do this via rounding or truncating.

Consider some  $y \in \mathbb{R} \setminus \{0\}$ .

• Truncating (k-digit truncation of  $y = \pm 0.d_1d_2d_3...d_kd_{k+1}d_{k+2}...\cdot 10^n$ ) Simply omit the digits from  $d_{k+1}$  and onwards, in other words

$$f_{\ell}(\pm 0.d_1d_2d_3\dots d_kd_{k+1}d_{k+2}\dots \cdot 10^n) = \pm 0.d_1d_2\dots d_k\cdot 10^n.$$

Thus we have  $y \approx f_{\ell}(y)$ .

• Rounding (k-digit rounding of  $y = \pm 0.d_1d_2d_3...d_kd_{k+1}d_{k+2}...\cdot 10^n$ )

If  $d_{k+1} < 5$ , then we drop  $d_{k+1}d_{k+2}...$  (same with truncating)

If  $d_{k+1} \geq 5$ , then add 1 to  $d_k$  and drop  $d_{k+1}d_{k+2}...$ 

$$f_{\ell}(\pm 0.d_1d_2d_3\dots d_kd_{k+1}d_{k+2}\dots \cdot 10^n) = \pm \delta_1\delta_2\dots \delta_k \cdot 10^m.$$

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## 2 Lecture 2