

# DIVIDE AND CONQUER PARADIGM

ANGEL NOÉ MARTÍNEZ GONZÁLEZ

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# THE PARADIGM I

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This paradigm solves a problem **recursively** applying three steps at each level of the recursion

- ▶ **DIVIDE** the problem into a smaller sub-problems.
- ▶ **CONQUER** via recursive calls.
- ▶ **COMBINE** solutions of sub-problems into one of the original problem.

Consider also a **base case** for small enough sub-problems.

# THE PARADIGM II

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## INSIGHTS AND HINTS

- ▶ Sub-problems can be any size:  $1/2$ ,  $1/3$ , etc.
- ▶ Generally, third step is the key to achieve good performance.
- ▶ Base case is often too ingnue.

# THE PARADIGM III

## MERGE SORT (RETAKE)

MERGE( $A, p, q, r$ )

MERGE-SORT( $A, p, r$ )

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1: if  $p < r$  then
2:    $q = \lfloor (p + r) / 2 \rfloor$ 
3:   MERGE-SORT( $A, p, q$ )
4:   MERGE-SORT( $A, q, r$ )
5:   MERGE( $A, p, q, r$ )
6: end if

```

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1:  $B = 1^{st}$  part of array.
2:  $C = 2^{nd}$  part of array.
3:  $i = 1, j = 1$ 
4: for  $k = 1$  to  $n$  do
5:   if  $B[i] < C[j]$  then
6:      $A[k] = B[i++]$ 
7:   else
8:      $A[k] = C[j++]$ 
9:   end if
10: end for

```

# COUNTING INVERSIONS I

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# THE MAXIMUM SUBARRAY I

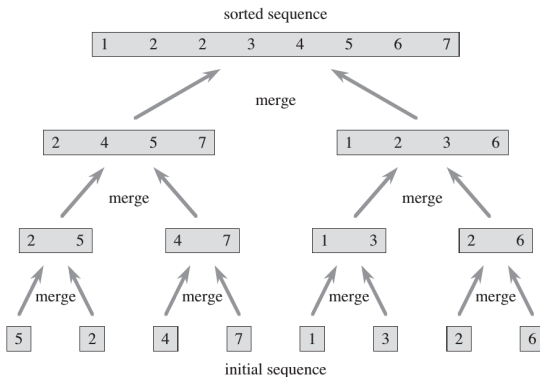
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# BINARY SEARCH I

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# RECURRENT TREE

## MERGE SORT RECURSION TREE



At each level  $j = 0, 1, \dots, \log_2(n)$ , there are  $2^j$  subproblems of size  $n/2^j$ .



# THE MASTER METHOD I

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## A BLACK BOX