# Math 325K - Lecture 5 Section 3.1 & 3.2 Predicates and quantified statements

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## Outline

- Predicates and quantifiers.
- Universal and existential quantified statements.
- Other forms of quantified statements.

## Motivation

Now we can do reasoning by valid argument forms. But there are a lot more complicated arguments. For example: all humans are mortal; Socrates is a human; : Socrates is mortal.

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Now we can do reasoning by valid argument forms. But there are a lot more complicated arguments. For example: all humans are mortal; Socrates is a human; : Socrates is mortal.

This should be a valid argument form, but it does not fit any patterns we introduced last time. In fact, our tools to deal with statements are not enough. And as a result, we need to introduce some new concepts and tools to study these arguments, which is called predicate logic.

## **Predicates**

In grammar, the word predicate refers to the part of a sentence that gives information about the subject.

#### Definition

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#### Remark

Usually we can interpret a predicate as a function whose co-domain consists of statements.

# Examples of predicates as functions

## Example

Let S(x) be the predicate  $x^2 > x$  with domain  $\mathbb{R}$ . Rewrite the following statements in sentences:

- (a) S(2);
- (b)  $\sim S(1) \wedge S(0)$ .

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#### Solution

(a) is simply  $2^2 > 2$  and it is true. (b) is '(not  $1^2 > 1$ ) and  $0^2 > 0$ ', which is equivalently ' $1^2 < 1$  and  $0^2 > 0$ ', so it is false.

## Truth set

Suppose we have a predicate P(x) with domain D, then for each  $y \in D$ , P(y) is a specific statement and it is either true or false. Then we would like to know that for what elements  $y \in D$ , P(y) is true?

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#### Definition

If P(x) is a predicate with domain D, the **truth set** of P(x) is the set of all elements of D that make P(x) true when they are substituted for x. The truth set of P(x) is denoted

$$\{x \in D \mid P(x) \text{ is true } \}.$$

# Quantifiers

In a sentence, even if we fix the subject and the predicate, there is still a twist: the number of subjects referred to? For example, the following sentences have very different meanings. As a result, we must consider the quantifiers too.

- all humans are mortal;
- some humans are mortal;
- one human is mortal;
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- one human is mortal;
- no human is mortal.

#### Definition

**Quantifiers** are words that refer to quantities such as 'some' or 'all' and tell for how many elements a given predicate is true.



# The quantifiers $\forall$ and $\exists$

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#### Definition

The universal quantifier, written as  $\forall$  and read 'for all', refers to all elements in the domain.

#### Definition

The existential quantifier, written as ∃ and read 'there exists/there exist', refers to at least one element in the domain.

## Universal statements

#### Definition

Let Q(x) be a predicate and D the domain of x. A universal statement is a statement of the form ' $\forall x \in D, Q(x)$ '. It is defined to be true if and only if Q(x) is true for every x in D. It is defined to be false if and only if Q(x) is false for at least one x in D. A value for x for which Q(x) is false is called a counterexample to the universal statement.

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#### Remark

One needs to specify the domain of universal statements. Sometimes this set may be implicitly given.

# Examples of universal statements

## Example

Let R(x) be the predicate ' $\sqrt{x}$  is a rational number' with domain  $\mathbb{N}$  (why not  $\mathbb{Z}$ ?), and L(x) be the predicate 'x < 2x' with domain  $\mathbb{N}$ . Rewrite the following universal statements in sentences and figure out whether they are true or false.

- (a)  $\forall x \in \mathbb{N}, R(x)$ ;
- (b)  $\forall x \in \mathbb{N}, L(x).$

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#### Solution

(a) is 'the square root of all positive integers are rational numbers'. It is false as 2 is a counterexample.

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- (a)  $\forall x \in \mathbb{N}, R(x)$ ;
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#### Solution

(a) is 'the square root of all positive integers are rational numbers'. It is false as 2 is a counterexample. (b) is 'for all positive integers x, x is less than 2x', which is true.

## Existential statements

#### Definition

Let Q(x) be a predicate and D the domain of x. An existential statement is a statement of the form  $\exists x \in D$  such that Q(x). It is defined to be true if and only if Q(x) is true for at least one x in D. It is defined to be false if and only if Q(x) is false for all x in D.

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#### Remark

Existential statements are somehow dual to universal statements, and we will introduce their connections in a while.

# Examples of existential statements

## Example

Rewrite the following existential statement in symbols and figure out whether it is true or false: there exists an integer that is prime and greater than 10.

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Rewrite the following existential statement in symbols and figure out whether it is true or false: there exists an integer that is prime and greater than 10.

#### Solution

The answer is not unique:

- $\exists n \in \mathbb{Z}$  such that n is prime and n > 10;
- $\exists n \in \{x \in \mathbb{Z} \mid x > 10\}$  such that n is prime;
- $\exists n \in \text{the set of prime numbers such that } n > 10.$

Since 11 is prime and greater than 10, this existential statement is true.

# Universal conditional statements

#### Definition

A universal conditional statement is of the form

$$\forall x \in D$$
, if  $P(x)$  then  $Q(x)$ ,

where P(x), Q(x) are predicates with domain D.

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where P(x), Q(x) are predicates with domain D.

#### Remark

Let U be the truth set of P(x), then the above universal conditional statement is logically equivalent to

$$\forall x \in U, Q(x).$$

# Implicit quantification

As we have already seen, some statements are written in a way without quantifiers. However, we can equivalently rewrite them with  $\forall$  or  $\exists$ .

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## Example

Write the following statement using symbols: if a person is a member of UT, then he/she has a UT EID. You may need to define predicates by yourselves.

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## Example

Write the following statement using symbols: if a person is a member of UT, then he/she has a UT EID. You may need to define predicates by yourselves.

#### Solution

Let H be the set of humans, M(x) be the predicate  ${}'x$  is a member of UT' with domain H and E(x) be the predicate  ${}'x$  has a UT EID ${}'$  with domain H. The statement becomes

$$\forall x \in H, M(x) \to E(x).$$

# Equivalent notations

#### Definition

Let P(x) and Q(x) be predicates with domain D. For convenience we write  $P(x) \Rightarrow Q(x)$  for

$$\forall x \in D, P(x) \to Q(x).$$

And we write  $P(x) \Leftrightarrow Q(x)$  for

$$\forall x \in D, P(x) \leftrightarrow Q(x).$$

# Finding the truth value of quantified statements

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If the domain is finite, we can simply substitute the variable by each element in the domain. This approach is called **exhaustion** and it is guaranteed to work in this case.

# Finding the truth value of quantified statements

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If the domain is finite, we can simply substitute the variable by each element in the domain. This approach is called **exhaustion** and it is guaranteed to work in this case.

But in mathematics, many sets are infinite, which means we need to consider more efficient way. There is one shortcut: one counterexample is enough to show a universal statement being false and one example is enough to show a existential statement being true.

# Negation of universal statements

## Proposition

The negation of a statement of the form

$$\forall x \in D, Q(x)$$

is logically equivalent to a statement of the form

$$\exists x \in D \text{ such that } \sim Q(x).$$

Symbolically,  $\neg(\forall x \in D, Q(x)) \equiv (\exists x \in D, \text{ such that } \neg Q(x)).$ 

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#### Remark

As a corollary, universal statements and existential statements can be defined by each other plus negation.

# Negation of existential statements

### Similarly we have

#### Proposition

The negation of a statement of the form

$$\exists x \in D \text{ such that } Q(x)$$

is logically equivalent to a statement of the form

$$\forall x \in D, \sim Q(x).$$

Symbolically,  $\sim (\exists x \in D \text{ such that } Q(x)) \equiv (\forall x \in D, \sim Q(x)).$ 

# Negation of universal conditional statements

### Proposition

The negation of a statement of the form

$$\forall x \in D, P(x) \to Q(x)$$

is logically equivalent to a statement of the form

$$\exists x \in D \text{ such that } \sim (P(x) \to Q(x)).$$

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 $\equiv \exists x \in D \text{ such that } \sim (P(x) \to Q(x)).$ 

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### Example

Write the formal (with symbols) negations of the following quantified statements:

- (a) All computer programs are finite.
- (b) There is a computer program in the programming language Lisp.
- (c) If a computer program has more than 100,000 lines, then it contains a bug.

#### Solution

Let P be the set of all computer programs. F(x) be the predicate x is finite; L(x) be the predicate x is in Lisp; Lisp; Lisp; Lisp; Lisp; Lisp; Lisp; Lisp; Lisp; and Li

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(a) is  $\forall x \in P, F(x)$ . So the negation is  $\exists x \in P$  such that  $\sim F(x)$ .

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- (a) is  $\forall x \in P, F(x)$ . So the negation is  $\exists x \in P$  such that  $\sim F(x)$ .
- (b) is  $\exists x \in P$  such that L(x). So its negation is  $\forall x \in P, \neg L(x)$ .

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- (b) is  $\exists x \in P$  such that L(x). So its negation is  $\forall x \in P, \neg L(x)$ .
- (c) is  $\forall x \in P, T(x) \to B(x)$ . So its negation is  $\exists x \in P$  such that  $\sim (T(x) \to B(x))$ , equivalently

 $\exists x \in P \text{ such that } (T(x) \land \sim B(x)).$ 

## Connection with $\land$ and $\lor$

Let P(x) be a predicate with a finite domain  $D = \{x_1, x_2, \dots, x_n\}$ . Then we have

#### Proposition

The universal statement  $\forall x \in D, P(x)$  is logically equivalent to

$$P(x_1) \wedge P(x_2) \wedge \cdots \wedge P(x_n)$$
.

And the existential statement  $\exists x \in D$  such that P(x) is logically equivalent to

$$P(x_1) \vee P(x_2) \vee \cdots \vee P(x_n).$$

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#### Remark

We don't have a similar result when D is infinite.

## Vacuous truth of universal statements

### Example

Is the following universal conditional statement true?
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#### Solution

It is of the form  $\forall x \in D, P(x) \to Q(x)$  while for all  $x \in D, P(x)$  is false. Then for each  $x \in D$ , the conditional statement  $P(x) \to Q(x)$  is by default true, hence the universal conditional statement is also true. And this is called the vacuous truth of them

## Other forms of universal conditional statements

#### Definition

Let P(x), Q(x) be predicates with domain D and s be the statement

$$\forall x \in D, P(x) \to Q(x)$$

The contrapositive of s is

$$\forall x \in D, \sim Q(x) \rightarrow \sim P(x).$$

The converse of s is

$$\forall x \in D, Q(x) \to P(x).$$

And the inverse of s is

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## Warning: $\forall$ and $\exists$ are not commutative

As we will discuss in detail next time, the quantifiers are not commutative in general.

### Example

Let G(x,y) be the binary predicate x < y with domain  $\mathbb{N} \times \mathbb{N}$ . Are the following statements true or false?

- (a)  $\forall x \in \mathbb{N}, \exists y \in \mathbb{N} \text{ such that } G(x,y).$
- (b)  $\exists y \in \mathbb{N}$  such that  $\forall x \in \mathbb{N}, G(x, y)$ .

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#### Solution

(a) is true, as given  $x \in N$ , we can always choose y = x + 1 to justify the existential statement.

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#### Solution

- (a) is true, as given  $x \in N$ , we can always choose y = x + 1 to justify the existential statement.
- (b) is false, as given  $y \in N$ , x = y + 1 is a counterexample to the universal statement, so for all y it is false. So is (b).

## HW #2 - these sections

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Section 3.1 Exercise 6, 17(b), 23(b), 28(a)(c).
Section 3.2 Exercise 2, 14, 23, 46.
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