

(*

A farmer is on one shore of a river and has with him a fox, a chicken, and a sack of grain. He has a boat that fits one object besides himself.

In the presence of the farmer nothing gets eaten, but if left without the farmer, the fox will eat the chicken, and the chicken will eat the grain. How can the farmer get all three possessions across the river safely?

*)

```

----- MODULE boat -----
EXTENDS Integers, FiniteSets

CONSTANTS Farmer, Fox, Chicken, Grain

CREATURES  $\triangleq$  {Farmer, Fox, Chicken, Grain}

alone(animals, side)  $\triangleq$  (animals  $\in$  SUBSET side)  $\wedge$   $\neg$ Farmer  $\in$  side

somebodyGetsEaten(l, r)  $\triangleq$   $\vee$  alone({Fox, Chicken}, l)
                         $\vee$  alone({Fox, Chicken}, r)
                         $\vee$  alone({Chicken, Grain}, l)
                         $\vee$  alone({Chicken, Grain}, r)

safe(l, r)  $\triangleq$   $\neg$ somebodyGetsEaten(l, r)

safeBoats(from, to)  $\triangleq$  {boat  $\in$  SUBSET from :  $\wedge$  Farmer  $\in$  boat
                         $\wedge$  Cardinality(boat)  $\leq$  2
                         $\wedge$  safe(from  $\setminus$  boat, to  $\cup$  boat)}

*****

--algorithm RiverCrossing{
variables left = CREATURES ; right = {} ;

process ( LeftToRight = 0 )
{ l: while ( left  $\neq$  {} )
  { await (Farmer  $\in$  left) ;
    with ( boat  $\in$  safeBoats(left, right) )
    {
      left := left  $\setminus$  boat ;
      right := right  $\cup$  boat
    }
  }
}

process ( RightToLeft = 1 )
{ r: while ( right  $\neq$  {} )
  { await (Farmer  $\in$  right) ;
```

```

    with ( boat  $\in$  safeBoats(right, left) )
    {
        left := left  $\cup$  boat ;
        right := right  $\setminus$  boat
    }
}

```

BEGIN TRANSLATION

VARIABLES $left, right, pc$

$vars \triangleq \langle left, right, pc \rangle$

$ProcSet \triangleq \{0\} \cup \{1\}$

$Init \triangleq$ Global variables
 $\wedge left = CREATURES$
 $\wedge right = \{\}$
 $\wedge pc = [self \in ProcSet \mapsto \text{CASE } self = 0 \rightarrow \text{"l"} \\ \square \quad self = 1 \rightarrow \text{"r"}]$

$l \triangleq$ $\wedge pc[0] = \text{"l"}$
 $\wedge \text{IF } left \neq \{\}$
 THEN $\wedge (Farmer \in left)$
 $\wedge \exists boat \in \text{safeBoats}(left, right) :$
 $\wedge left' = left \setminus boat$
 $\wedge right' = (right \cup boat)$
 $\wedge pc' = [pc \text{ EXCEPT } ![0] = \text{"l"}]$
 ELSE $\wedge pc' = [pc \text{ EXCEPT } ![0] = \text{"Done"}]$
 $\wedge \text{UNCHANGED } \langle left, right \rangle$

$LeftToRight \triangleq l$

$r \triangleq$ $\wedge pc[1] = \text{"r"}$
 $\wedge \text{IF } left \neq \{\}$
 THEN $\wedge (Farmer \in right)$
 $\wedge \exists boat \in \text{safeBoats}(right, left) :$
 $\wedge left' = (left \cup boat)$
 $\wedge right' = right \setminus boat$
 $\wedge pc' = [pc \text{ EXCEPT } ![1] = \text{"r"}]$
 ELSE $\wedge pc' = [pc \text{ EXCEPT } ![1] = \text{"Done"}]$
 $\wedge \text{UNCHANGED } \langle left, right \rangle$

$RightToLeft \triangleq r$

$Next \triangleq LeftToRight \vee RightToLeft$

$$\vee \text{ Disjunct to prevent deadlock on termination}$$

$$((\forall self \in ProcSet : pc[self] = \text{"Done"}) \wedge \text{UNCHANGED } vars)$$

$$Spec \triangleq Init \wedge \Box[Next]_{vars}$$

$$Termination \triangleq \Diamond(\forall self \in ProcSet : pc[self] = \text{"Done"})$$

END TRANSLATION

\ * Modification History
\ * Last modified *Wed Jun 04 21:54:45 EDT 2014* by *lorinhochstein*
\ * Created *Mon Jun 02 20:41:25 EDT 2014* by *lorinhochstein*