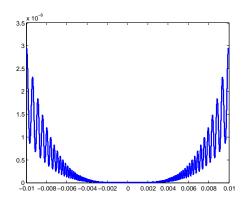
## **Solutions to Part A of Problem Sheet 1**

## Solution (1.1)

- (a) The function  $f(x) = x^4$  has a strict minimum at x = 0, but the second derivative satisfies f''(0) = 0.
- (b) We construct a function that has a strict minimizer  $x^*$ , but such that every open neighourhood U of  $x^*$  contains other local minimizers. One such function is

$$f(x) = \begin{cases} x^4(\cos(1/x) + 2) & \text{for } x \neq 0 \\ 0 & \text{for } x = 0. \end{cases}$$



We explain the construction of this function:

- 1. Start out with  $g(x) = \cos(1/x) + 2$  for  $x \neq 0$  and g(0) = 1. This function has minimizers  $x_0 = 0$  and  $x_k = 1/(\pi k)$  for k > 0, with values  $g(x_k) = 1$  at all minimizers. Therefore, any open interval around 0 contains (infinitely many) local minimizers  $x_k$  other than  $x_0 = 0$ .
- 2. Multipy  $x^4$  to the function:  $f(x) = x^4 g(x)$ . This ensures that f(0) = 0 and f(x) > 0 for  $x \neq 0$ . There are still local minima in every neighbourhood of 0. To see this, compute the derivative

$$f'(x) = x^2 (4x\cos(1/x) + \sin(1/x) + 8x). \tag{1}$$

Set  $z_m=1/(\pi/2+m\pi)$  for m>0. Since  $\sin(1/z_m)=\sin(\pi/2+m\pi)=1$  for m even and -1 for m odd, for m sufficiently large the derivative (1) changes signs between successive  $z_m$ . Since f'(x) is continuous, it has roots between any  $z_m$  and  $z_{m+1}$  for large enough m, and these correspond to maxima and minima of f.

The function is in  $C^2(\mathbb{R})$ . For  $x \neq 0$  this is clear, and to verify this at x = 0, one shows that the right and left limits as  $x \to 0$  of f'(x) and f''(x) coincide (they are in fact 0).

Note the subtle point that one minimizer  $x^*$  can have local minimizers that are arbitrary close: while each open interval I surrounding  $x^*$  has another local minimizer  $\tilde{x}$ , every such  $\tilde{x}$  has an interval  $\tilde{I}$  surrounding it where this  $\tilde{x}$  is the only minimizer!

## **Solution (1.2)** We want to show that the function

$$f(\boldsymbol{x}) = \frac{1}{2} \boldsymbol{x}^{\top} \boldsymbol{A} \boldsymbol{x} + \boldsymbol{b}^{\top} \boldsymbol{x} + c$$
 (2)

is convex if and only A is positive semidefinite. To see this, we compute the partial derivatives and the Hessian of f. The parts  $b^{\top}x$  and c disappear when computing second derivatives. The function  $x^{\top}Ax$  can be written as

$$oldsymbol{x}^{ op} oldsymbol{A} oldsymbol{x} = \sum_{i=1}^m \sum_{j=1}^n a_{ij} x_i x_j,$$

so that the first derivative is

$$\frac{\partial f}{\partial x_i} = \frac{1}{2} \sum_{j \neq i} (a_{ij} + a_{ji}) x_j + a_{ii} x_i + b_i = \sum_{j=1}^n a_{ij} x_j + b_i,$$

where we used the symmetry of A (i.e.,  $a_{ij} = a_{ji}$ ). The gradient and Hessian are therefore just given by

$$\nabla f(\boldsymbol{x}) = \boldsymbol{A}\boldsymbol{x} + \boldsymbol{b}, \quad \nabla^2 f(\boldsymbol{x}) = \boldsymbol{A}.$$

An interesting special case is when the (2) arises in the form

$$f(x) = \frac{1}{2} ||Ax - b||_2^2.$$
 (3)

The quadratic system then has the form

$$\|Ax - b\|_{2}^{2} = (Ax - b)^{\top}(Ax - b) = x^{\top}A^{\top}Ax - 2b^{\top}Ax + \|b\|_{2}^{2}.$$
 (4)

The matrix  $A^{\top}A$  is always symmetric and positive semidefinite:

$$(\boldsymbol{A}^{\top}\boldsymbol{A})^{\top} = \boldsymbol{A}^{\top}(\boldsymbol{A}^{\top})^{\top} = \boldsymbol{A}^{\top}\boldsymbol{A}, \quad \text{ and } \quad \boldsymbol{x}^{\top}\boldsymbol{A}^{\top}\boldsymbol{A}\boldsymbol{x} = \|\boldsymbol{A}\boldsymbol{x}\|_{2}^{2} > 0 \text{ if and only if } \boldsymbol{x} \neq \boldsymbol{0},$$

so that the function (3) is convex. From (4) we also see that the derivative of (3) is

$$A^{\top}(Ax-b)$$
.

## Solution (1.3)

(a) The unit circle with respect to the  $\infty$ -norm is the square with corners  $(\pm 1, \pm 1)^{\top}$ .

(b) The trick in transforming the unconstrained problem

minimize 
$$\|Ax - b\|_{\infty}$$
 (5)

into a constrained linear programming problem is to characterise the  $\infty$ -norm as the solution of a minimization problem. In fact, for any set of numbers  $x_1, \ldots, x_n$ ,

$$\max_{1 \le i \le n} |x_i| = \min_{\forall i \colon |x_i| \le t} t.$$

Put simply, the *maximum* of a set of non-negative numbers is the *smallest* upper bound on these numbers. We can further replace the condition  $|x_i| \leq t$  by  $-t \leq x_i \leq t$ , so that the problem (5) becomes

minimize 
$$t$$
 subject to  $-t \le \boldsymbol{a}_1^{\top} \boldsymbol{x} \le t$   $\cdots$   $-t \le \boldsymbol{a}_m^{\top} \boldsymbol{x} \le t$ , (6)

where  $a_i^{\top}$  are the rows of the matrix A. This problem can be brought into standard form by replacing each condition with the pair of conditions

$$\mathbf{a}_i^{\top} \mathbf{x} - t \le 0$$
$$-\mathbf{a}_i^{\top} \mathbf{x} - t \le 0.$$

The solution x of Problem (5) can be read off the solution (x, t) of Problem (6).