Conceptual Database Design

1. Conceptual Database Design

Appl E-R Relational DB

Database Design

Entity-Relationship Model

- The Entity-Relationship (ER) model is a high-level conceptual data model (Chen in 1966).
- ER is used mainly as a design tool.

Entity-Relationship Model(cont)

- Entity type: Group of object with the same properties
- *Entity*: member of an entity type analogous to an object.
- Attribute: a property of object
- Relationship: among objects
 - ER can model "n-way" relationship,
 - ER models a relationship and its inverse by a single relationship.

1.1 Entity and Attributes

- Entities represent things in the real word.
- Attributes describe properties of entities.
- Attributes may be
 - simple(atomic) e.g. sex = 'Female', or
 - composite e.g. name consists of title (Dr), Initials(C.C.), family name (Chen).

- Each entity has values for each attribute.
- Attributes may be
 - single-valued e.g. student number, name, or
 - multivalued e.g. keywords = neural networks,computer graphics, databases.

- Each simple attribute has a *value set (domain)*: the set of possible values for that attribute.
- In a composite attribute $A = (A_1, \ldots, A_n)$, suppose that V_1, \ldots, V_n are the domains of A_1, \ldots, A_n .
- The domain V of A is $V_1 \times \ldots \times V_n$.
- Mathematically, an attribute A of an entity type E is a function

$$A: E \to \wp(V)$$
.

- where V is the domain of A, and $\wp(V)$ is the power set of V
- For single-valued attributes, A(e) must be a singleton.

- An attribute can have a null value if, for example:
 - there is no suitable value e.g. a student may have no interests: keywords = NULL
 - the true value is not known e.g. the marriage date of a person is not known: marriage date = NULL.
- A derived attribute is one whose value can be derived from other attributes and entities. e.g. number of students.

- An *entity* type is a set of entities with the same attributes.
- It is described by an *entity* schema: a name and a list of attributes.
- The set of individual entity *instances* at a particular moment in time is called an extension of the entity type.

Schema (Intension)	RESEACHER Name, Payroll_no, No_of_students, Keywords	DEPARTMENT Name
Instances (Extension)	(Dr C.C. Chen, 230-0013, 3, Neural Networks) (Dr R. Wilkinson, 231-0091, 1, Databases)	Computer Science Psychology Management

- An entity type usually has a *key*: a set of attributes that uniquely identifies an entity. For example:
 - {payroll number} is a key of RESEARCHER,
 - {name} is a key of DEPARTMENT.
- There may be more than one possible key.
- An important constraint is the key constraint: in any
 extension of the entity type, there cannot be two entities
 having the same values for their key attributes.

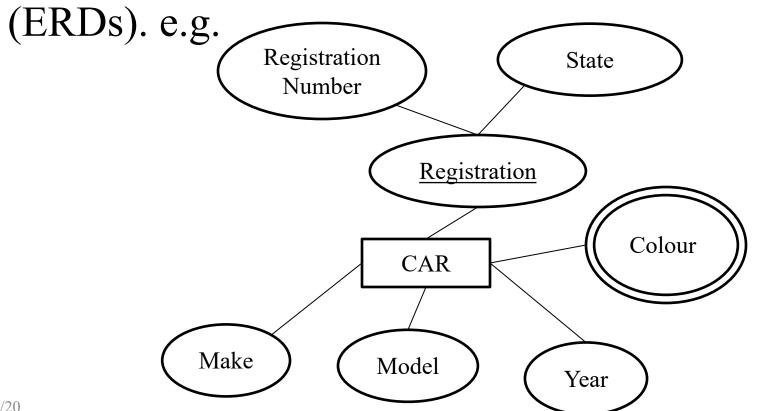
• We can describe schemata with composite attributes using ()'s and with multi-valued attributes using {}'s. e.g.

CAR

Registration(Registration No, State), Make, Model, Year, {Colour}

((ARQ) 595, Vic), Datsun, 120Y, 1972, {green}) ((8HR) 696, WA), Mazda, 929, 1979, {grey, black})

• Entities and their attributes can also be described with Entity-Relationship Diagrams

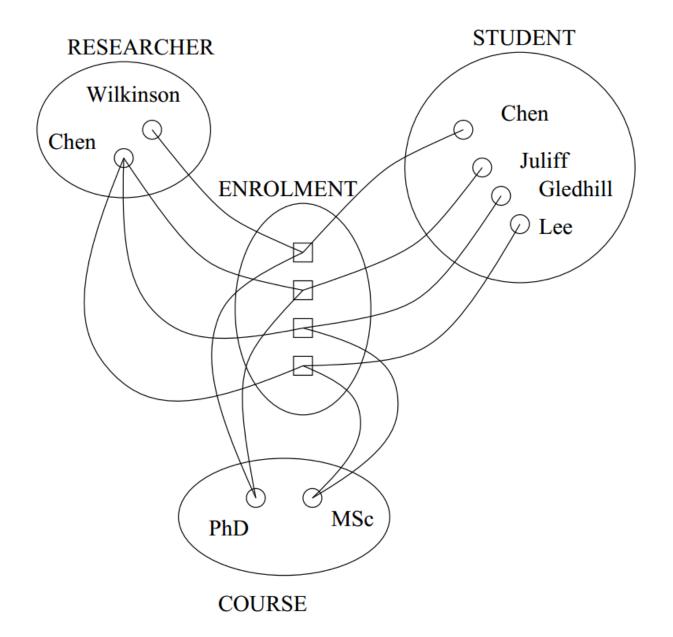


1.2 Relationships

- A relationship represents an association between things.
- A relationship type R among n entity types E_1, \ldots, E_n is a set of associations among entities from these types.
- Mathematically, a relationship type R among entity types E_1, \ldots, E_n is a subset of $E_1 \times \ldots \times E_n$.
- Each instance $r = (e_1, \ldots, e_n)$ in R is a relationship.

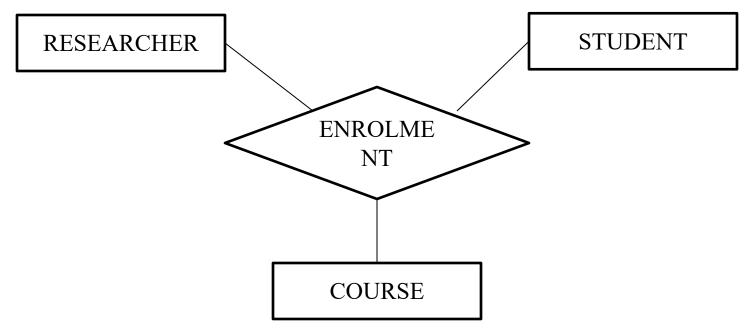
1.2 Relationships(cont)

- We say that E_1, \ldots, E_n participate in R.
- Similarly if $r = (e_1, \ldots, e_n)$ is an instance of R, we say that each e_i participates in r.
- The *degree* of *R* is the number of participating entity types. For example,
 - ENROLMENT could be a ternary (degree 3) relationship between RESEARCHER, STUDENT and COURSE.
- We can illustrate this using an occurrence diagram:



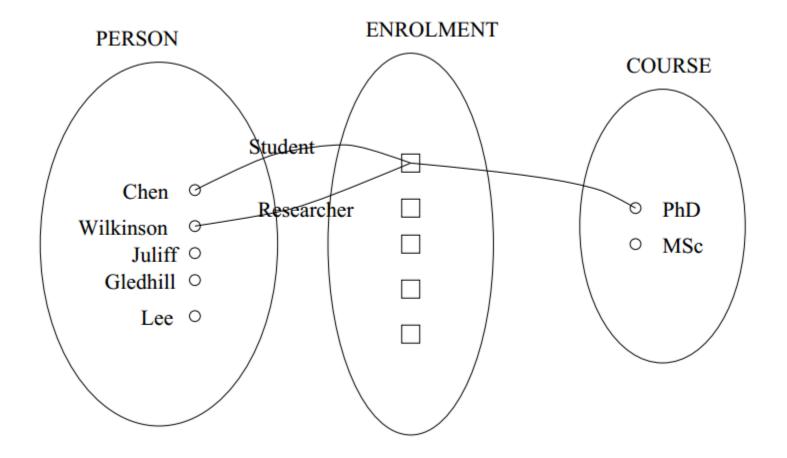
1.2 Relationships(cont)

• Entities and their relationships can also be represented using Entity-Relationship diagrams:



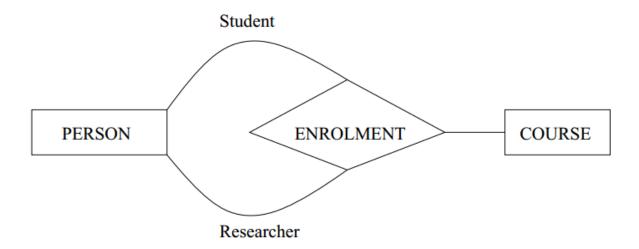
1.2 Relationships(cont)

- Each entity type that participates in a relationship plays a particular *role* in the relationship.
- An entity type can play
 - different roles in different relationships, or
 - more than one role in a relationship.
- A role name can be used to distinguish these.
- For example, ENROLMENT could be a relationship between PERSON(as researcher), PERSON(as student) and COURSE as in the diagram below:



1.2 Relationships(cont)

• Or, using an ERD:



• This is called a recursive relationship.

1.3 Weak entity types

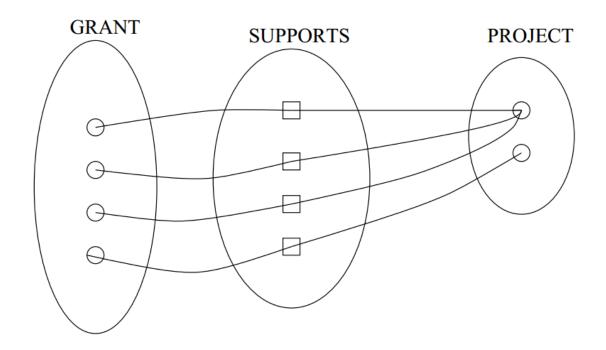
- Some entity types do not have a key of their own.
- Such entity types are called weak entity types.
- Entities of a weak entity type can be identified by a partial key and by being related to another entity type *owner*.
- The relationship type between a weak entity type to its owner is the *identifying relationship* of the weak entity type.

1.3 Weak entity types(cont)

- For example, a TAX PAYER entity may be related to several DEPENDENT, identified by their names.
- In this example, DEPENDENT is called a weak entity, {Name} is a partial key for it. The identifying relationship between DEPENDENT and TAX PAYER is IS DEPENDENT OF. TAX PAYER is said to *own* DEPENDENT.

- Relationship types usually have certain constraints that limit the possible combinations of entities participating in relationship instances.
- They should reflect the correct factors
- *Cardinality ratio constraint*: specifies the number of relationship instances an entity can participate in.
- Example: A research grant supports only one research project, but a research project may be supported by many grants. PROJECT:GRANT is a 1: N relationship.

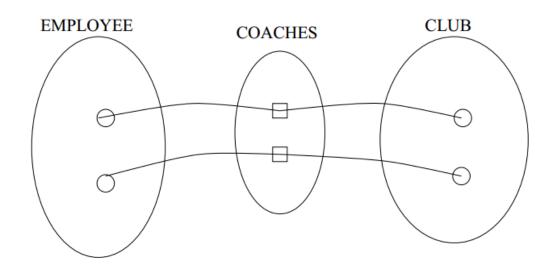
• This is illustrated in the occurrence diagram below:



• We can also show this in an ERD:



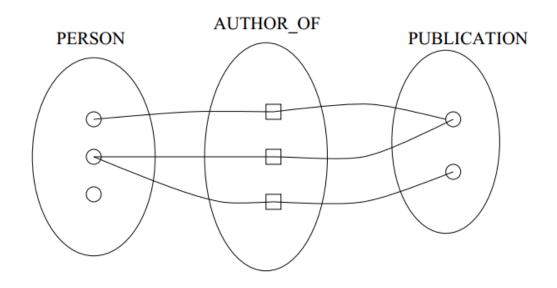
• Example: Consider a database of AFL (here substitute your favourite team sport) statistics. The relationship of head coaches to clubs is an example of a 1 : 1 relationship.



• With an ERD:



• Example: An example of an N : M relationship is authorship of publications:



• The equivalent ERD:



1.2.4 Constraints on relationship types_(cont)

- Another kind of constraint that can be represented using the ER model is a
 - Participation constraint: participation of an entity in a relationship can be:
 - *total*: every entity must participate e.g. every publication has an author.
 - partial: not necessarily total. e.g. not every person has publications.

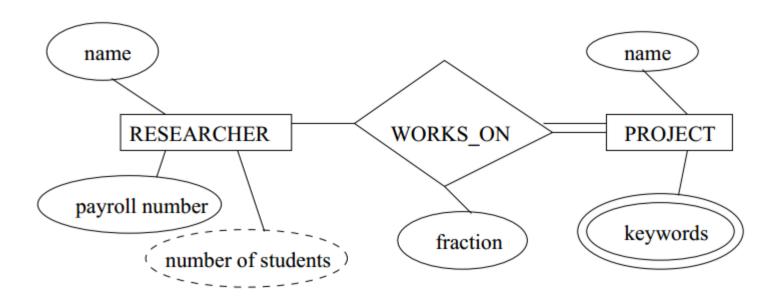
• This can be shown with an ERD like the one below:



1.5 Attributes of relationship types

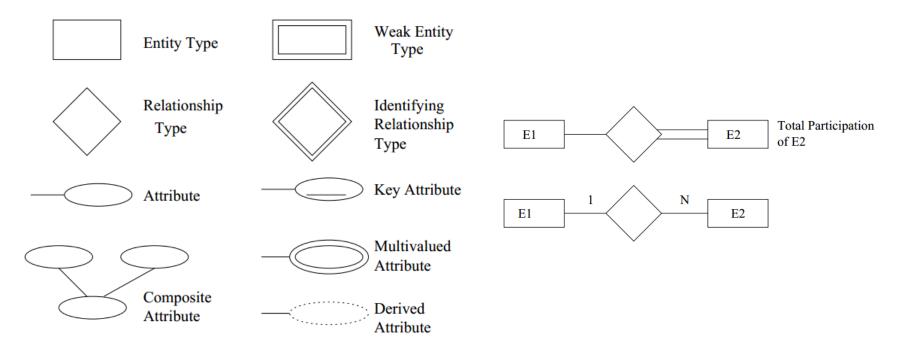
- Relationship types can have attributes for example,
 - a researcher may work on several projects. The fraction of her time devoted to a particular project could be an attribute of the WORKS ON relationship type.
- This can be shown in an ERD as below:

1.5 Attributes of relationship types(cont)



1.5 Attributes of relationship types(cont)

• The notation used for ERDs is summarised in Elmasre/Navathe Figure 3.15.



1.6 Enhanced ER (EER) model

Designers must use additionally modelling concepts

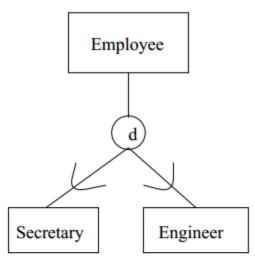
 represent the requirements from applications as accurately and explicitly as possible.

1.6 Enhanced ER (EER) model(cont)

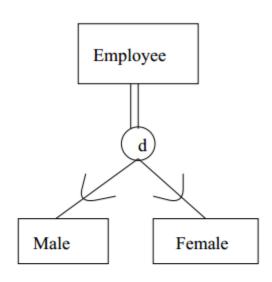
- There are many extensions to the ER model. We will look at one:
 - Specialisation: the process of defining a set of subclasses of an entity type; this entity type is called the superclass of the specialization.
 - Generalisation: a reverse process of specialisation.
- A subclass inherits all the attributes of the superclasses.

1.6 Enhanced ER (EER) model(cont)

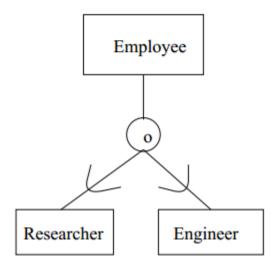
- A specialisation involves the following aspects:
 - Define a set of subclasses of an entity type.
 - Associate additional specific attributes with each subclass.
 - Establish additional specific relationship types between each subclass and other entity types, or other subclasses.
- A subclass may have multiple superclasses.
- A specialisation:
 - may be either total or partial; and
 - may be either disjoint or overlapping.

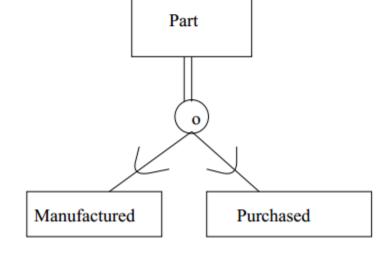


Partial disjoint



Total disjoint





2/14/2 **Partial Overlapping**

Total Ovelapping

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1.7 Design Principles

- Faithfulness: reflect reality.
- Avoid redundancy.
- Picking the right kind of element.