- ▶ Def: cycle(s) =  $\{s_i...s_{|s|}s_1...s_{i-1} \mid 1 \le i \le |s|\}$
- Ex:  $cycle(123) = \{123, 231, 312\}.$
- ▶ Def: cycle(L) =  $\cup_{s \in L}$  cycle(s)
- ▶ Prove or disprove: if *L* is regular, then cycle(*L*) is also regular.

```
Def: merge(x,[]) = {x}
Def: merge([], y) = {y}
Def: merge(x : xs, y : ys) =
    {x ∘ s | s ∈ merge(xs, y : ys)} ∪
    {y ∘ s | s ∈ merge(x : xs, ys)}
Ex: merge(ab, 12) = {ab12, a1b2, a12b, 1ab2, 1a2b, 12ab}
Def: merge(L<sub>1</sub>, L<sub>2</sub>) = ∪<sub>s1∈L1,s2∈L2</sub>merge(s<sub>1</sub>, s<sub>2</sub>)
Prove or disprove: if L<sub>1</sub> and L<sub>2</sub> are regular, then merge(L<sub>1</sub>, L<sub>2</sub>) is also regular.
```

- ▶ Def: permute is the permutation function.
- ▶ Prove or disprove: if *L* is regular, then permute(*L*) is also regular.

- ▶ Let  $\Sigma = \{0, 1\}$
- ▶ Let n(s) = number of 1's in s
- ▶ Prove or disprove: if for every predicate p,  $L_1 = \{1^n | p(n) = \text{true}\}$  is regular iff  $L_2 = \{s | p(n(s)) = \text{true}\}$  is regular

- ▶ Let  $\Sigma = \{0, 1\}$
- Let bin(s) = value of s read as binary number
- Prove or disprove: if for every predicate p,  $L_1 = \{1^n | p(n) = \text{true}\}$  is regular iff  $L_2 = \{s | p(bin(s)) = \text{true}\}$  is regular

- ▶ A read-twice DFA, on deciding input w, is fed  $w \circ w$ . (So it gets to read the input twice).
- Ex: to decide the string "Hello", it would be fed "HelloHello"
- Prove or disprove: a read-twice DFA can only recognize regular languages.

- ▶ A read-twice-reverse DFA, on deciding input w, is fed  $w \circ w^r$ , where  $w^r$  is w reversed.
- ► Ex: to decide the string "Hello", it would be fed "HelloolleH"
- Prove or disprove: a read-twice-reverse DFA can only recognize regular languages.