AMS 250: An Introduction to High Performance Computing

MPI Primer



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Outline

- Introduction
- MPI Implementations and Compilers
- Environment Management Routines
- Point to Point Communication Routines
- Collective Communication Routines
- Derived Data Types
- Group and Communicator Management Routines
- Virtual Topologies

What is MPI?

- MPI = Message Passing Interface
- MPI is a *specification* for the developers and users of message passing libraries.
- MPI primarily addresses the message-passing parallel programming model:

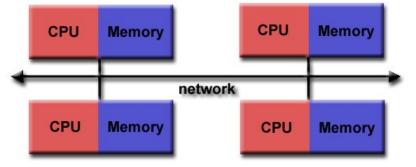
data is moved from the address space of one *process* to that of another *process* through cooperative operations on each *process*

- The goal of MPI is to provide a widely used standard for writing message passing programs. The interface attempts to be:
 - Practical
 - Portable
 - Efficient
 - Flexible

MPI Programming Model

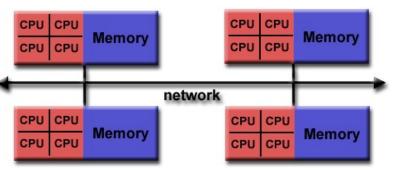
Originally, MPI was designed for distributed memory architectures,
 which were becoming increasingly popular at that time (1980s - early

1990s).



 As architecture trends changed, shared memory SMPs were combined over networks creating hybrid distributed memory / shared

memory systems.



Programming Model (cont'd)

- MPI implementers adapted their libraries to handle both types of underlying memory architectures seamlessly. They also adapted/developed ways of handling different interconnects and protocols.
- Today, MPI runs on virtually any hardware platform:
 - Distributed Memory
 - Shared Memory
 - Hybrid
- The programming model clearly remains a *distributed memory model*, regardless of the underlying physical architecture of the machine.
- All parallelism is *explicit*: the programmer is responsible for correctly identifying parallelism and implementing parallel algorithms using MPI constructs.

Reasons for Using MPI

Standardization:

MPI is the only message passing library which can be considered a standard.
 Practically, it has replaced all previous message passing libraries.

Portability:

- There is little or no need to modify your source code when you port your application to a different platform that supports (and is compliant with) the MPI standard.

Availability:

A variety of implementations are available, both vendor and public domain.

Performance Opportunities:

 Vendor implementations should be able to exploit native hardware features to optimize performance. Any implementation is free to develop optimized algorithms.

Functionality:

 There are over 430 routines defined in MPI-3, which includes the majority of those in MPI-2 and MPI-1.

History

Month/Year	Version
May 1994	MPI-1.0
Jun 1995	MPI-1.1
Jul 1997	MPI-1.2
Jul 2008	MPI-1.3
Jul 1997	MPI-2.0
Sep 2008	MPI-2.1
Sep 2009	MPI-2.2
Sep 2012	MPI-3.0
Jun 2015	MPI-3.1

http://www.mpi-forum.org/docs/

References

- Documentation for all versions of the MPI standard is available at: http://www.mpi-forum.org/docs/
- MPI page at ANL:

http://www.mcs.anl.gov/research/projects/mpi/index.htm
which points to a lot of stuff about MPI, including papers, talks, and tutorials

Books

- Using MPI: Portable Parallel Programming with the Message-Passing Interface, 3rd Ed., by Gropp, Lusk, and Skjellum, MIT Press, 2014. http://ieeexplore.ieee.org/xpl/bkabstractplus.jsp?bkn=6981847
- Using Advanced MPI: Modern Features of the Message-Passing Interface, by Gropp, Hoefler, Thakur, and Lusk, MIT Press, 2014. http://ieeexplore.ieee.org/xpl/bkabstractplus.jsp?bkn=6981848
- MPI: The Complete Reference Vol. 1 The MPI Core, by Snir, Otto, Huss-Lederman, Walker, and Dongarra, MIT Press, 1998.
- MPI: The Complete Reference Vol. 2 The MPI-2 Extensions, by Gropp, Huss-Lederman, Lumsdaine, Lusk, Nitzberg, Saphir, and Snir, MIT Press, 1998.
- Designing and Building Parallel Programs, by Ian Foster, Addison-Wesley, 1995.

http://www.mcs.anl.gov/~itf/dbpp/text/book.html

Tutorials

- MPI tutorial by Blaise Barney at LLNL: https://computing.llnl.gov/tutorials/mpi/
- MPI tutorials at ANL: http://www.mcs.anl.gov/research/projects/mpi/tutorial/index.html
- MPI tutorial by Rolf Rabenseifner at University of Stuttgart:
 https://fs.hlrs.de/projects/par/par prog ws/2004B/03 mpi 1 rab.pdf

Confession: this presentation is built out of those tutorials!

MPI Implementations

- Many open and commercial implementations of the MPI standard are available, targeting C, C++, and Fortran programmers:
 - MPICH
 - Open MPI
 - MVAPICH
 - Intel MPI
- Bindings are available for many other languages, including Perl, Python, R, Ruby, Java and Matlab.

MPICH

- http://www.mpich.org/
- MPICH is a high-performance and widely portable implementation of the MPI standard (MPI-1, MPI-2 and MPI-3).
- Distributed under a BSD-like license.
- Latest stable release: mpich-3.2
- The **CH** part of the name was derived from "Chameleon", which was a portable parallel programming library developed by William Gropp.
- MPICH is the foundation for the vast majority of MPI implementations, including MVAPICH, Intel MPI, Cray MPI, and many others.
- MPICH channels:
 - Nemesis the default channel which supports multiple communication methods
 - Sock a simple channel based on standard Unix sockets

Open MPI

- https://www.open-mpi.org/
- Open MPI is an open source MPI implementation, the merger of 3 well-known MPI implementations:
 - FT-MPI from the University of Tennessee
 - LA-MPI from Los Alamos National Laboratory
 - LAM/MPI from Indiana University
- Latest stable release: v1.10.2, which is fully compliant with MPI-3.1
- thread-safe
- includes Java bindings since v1.7
- Supports various interconnects, including:
 - Infiniband / RoCE / iWARP verbs
 - Shared memory
 - TCP sockets

Open MPI Architecture:

http://www.aosabook.org/en/openmpi.html

MVAPICH

- http://mvapich.cse.ohio-state.edu/
- MPI over InfiniBand, iWARP and RoCE
 - InfiniBand natively supports RDMA (Remote Direct Memory Access)
 - iWARP (internet Wide Area RDMA Protocol) implements RDMA over an IP network
 - Roce (RDMA over Converged Network) allows RDMA over an Ethernet network
- Latest stable release: MVAPICH2 2.1, which supports MPI 3.0
- Other flavors:
 - MVAPICH2-X: with support for PGAS and OpenSHMEM
 - MVAPICH2-GDR, with support for InfiniBand and Nvidia CUDA GPUs
 - MVAPICH2-MIC, with support for InfiniBand and Intel MIC
 - MVAPICH2-Virt, with support for InfiniBand and SR-IOV

Intel MPI

- Intel MPI is a commercial MPI implementation based on MPICH and MVAPICH (http://pleiades.ucsc.edu/doc/intel/impi/4.1.0.024/)
- Supports various network fabrics, including:
 - **shm**: shared-memory
 - dapl: DAPL-capable network fabrics, such as InfiniBand, iWarp, Dolphin & XPMEM
 - tcp: TCP/IP-capable network fabrics, such as Ethernet and InfiniBand (through IPoIB)
 - tmi: Network fabrics with tag matching capabilities through the Tag Matching Interface (TMI), such as Qlogic and Myrinet
 - ofa: Network fabrics supporting OFED (Open Fabrics Enterprise Distribution) verbs, such as InfiniBand
- Default on Hyades
 - Intel MPI 4.1 implements MPI-2.2
 - Intel MPI 5.0 implements MPI-3.0

MPI "Hello, world!" in Fortran 77

```
implicit none
include 'mpif.h'
integer ierr, rank, size

call MPI_INIT(ierr)
call MPI_COMM_RANK(MPI_COMM_WORLD, rank, ierr)
call MPI_COMM_SIZE(MPI_COMM_WORLD, size, ierr)
print *, "Hello, world! I am process ", rank, " of ", size
call MPI_FINALIZE(ierr)
END
```

MPI "Hello, world!" in Fortran 90

```
program hello

use mpi
implicit none
integer :: ierr, rank, size

call MPI_INIT(ierr)
call MPI_COMM_RANK(MPI_COMM_WORLD, rank, ierr)
call MPI_COMM_SIZE(MPI_COMM_WORLD, size, ierr)
print *, "Hello, world! I am process ", rank, " of ", size
call MPI_FINALIZE(ierr)

end program hello
```

MPI "Hello, world!" in C

```
#include "mpi.h"
#include <stdio.h>
int main(int argc, char* argv[])
  int rank, size;
 MPI_Init(&argc, &argv);
  MPI_Comm_rank(MPI_COMM_WORLD, &rank);
  MPI_Comm_size(MPI_COMM_WORLD, &size);
  printf("Hello, world! I am process %d of %d\n", rank, size);
 MPI_Finalize();
  return 0;
```

MPI "Hello, world!" in C++

```
#include "mpi.h"
#include <iostream>
using namespace std;
int main(int argc, char* argv[])
  int rank, size;
  MPI::Init(argc, argv);
  rank = MPI::COMM_WORLD.Get_rank();
  size = MPI::COMM_WORLD.Get_size();
  cout << "Hello, world! I am process " << rank << " of " << size << endl;</pre>
  MPI::Finalize();
  return 0;
```

Quick Note on C++ Bindings

- Only Fortran 77 and ANSI C bindings were included in MPI-1.x
- Fortran 90 and C++ bindings were introduced in MPI-2.0
- C++ bindings were deprecated in MPI-2.2, for several reasons:
 - The MPI Forum didn't have enough C++ expertise to continue to maintain the C++ bindings
 - The C++ bindings really didn't offer anything useful to developers over the C bindings
 - Very few developers actually used the C++ bindings in *real-world applications*
 - The C++ bindings were (intentionally) a more-or-less 1:1 mapping to the C bindings; so C++ MPI applications can be easily converted to use C bindings.
- C++ bindings were removed in MPI-3.0
- We won't cover C++ bindings in this presentation

http://blogs.cisco.com/performance/the-mpi-c-bindings-what-happened-and-why

What if you are a C++ developer?

- Use the C bindings
- Or use the deprecated C++ bindings in MPI-2.2, which are usually provided by the actual MPI implementations
- Or use the popular MPI C++ class library boost.mpi, which is implemented on the MPI C bindings

http://www.boost.org/doc/libs/1 58 0/doc/html/mpi.html

http://blogs.cisco.com/performance/the-mpi-c-bindings-are-gone-what-does-it-mean-to-you

boost.mpi "Hello, world!"

```
#include <boost/mpi.hpp>
#include <iostream>
namespace mpi = boost::mpi;
int main()
  mpi::environment env;
  mpi::communicator world;
  int rank, size;
  rank = world.rank();
  size = world.size();
  std::cout << "Hello, world! I am process "</pre>
            << rank << " of " << size << endl;
  return 0;
```

Intel MPI Compilers

MPI Compiler Command	Default Compiler	Supported Language(s)
mpiicc	icc	С
mpiicpc	icpc	C++
mpiifort	ifort	Fortran 77/90
mpicc	gcc	С
mpicxx	g++	C++
mpifc	gfortran	Fortran 77/90
mpigcc	gcc	С
mpigxx	g++	C++
mpif77	g77	Fortran 77
mpif90	gfortran	Fortran 90

MPI Compiler Wrapper Scripts

The **mpicmds** in the previous table are just wrapper scripts used to simplify the compilation of MPI programs. They automatically include the appropriate MPI include files and link to the necessary MPI libraries and pass switches to the underlying compilers.

```
$ mpiicc -show

icc -I/opt/intel/impi/4.1.3.045/intel64/include \
    -L/opt/intel/impi/4.1.3.045/intel64/lib \
    -xlinker --enable-new-dtags -xlinker -rpath \
    -xlinker /opt/intel/impi/4.1.3.045/intel64/lib \
    -xlinker -rpath -xlinker /opt/intel/mpi-rt/4.1 \
    -lmpigf -lmpi -lmpigi -ldl -lrt -lpthread
```

Compiling & Running MPI Programs on Hyades

```
Compile MPI programs using the default Intel compilers:
```

```
mpiifort [other options] mpi_hello.f -o mpi_hello.x
mpiifort [other options] mpi_hello.f90 -o mpi_hello.x
mpiicc [other options] mpi_hello.c -o mpi_hello.x
mpiicpc [other options] mpi_hello.cpp -o mpi_hello.x
```

(Test) run MPI programs on the master node (32 processes):

```
mpirun -n 32 ./mpi_hello.x
```

For production runs, submit your MPI jobs to the batch scheduler.

mpirun

- MPI start mechanism is implementation dependent
 - Most implementations, including Intel MPI, provide mpirun
 - MPI-2 standard defines mpiexec
- Learn more about **mpirun**:

```
mpirun -help
```

What if you run a MPI program without using mpirun?

```
$ mpirun -n 2 ./hello.x
Hello, world! I am process 1 of 2
Hello, world! I am process 0 of 2
$ ./hello.x
Hello, world! I am process 0 of 1
```

 Quiz: If there are only 2 CPU cores in your laptop, can you run a 4processs (or 8-process) MPI job?

Sample Batch Script for MPI Jobs

Batch script impi.pbs:

```
#!/bin/bash

#PBS -N impi
#PBS -q normal
#PBS -l nodes=4:ppn=16
#PBS -l walltime=4:00:00
#PBS -M shaw@ucsc.edu
#PBS -m abe
#PBS -j oe

cd $PBS_O_WORkDIR
mpirun -genv I_MPI_FABRICS shm:ofa -n 64 ./mpi_hello.x
```

Comments:

```
### your favorite shell

### job name
### job queue
### request 4 node (16 cores per node)
### and 4 hours walltime
### ask Torque to send emails
### when jobs aborts, starts and ends
### merge standard error with standard output

### go to the directory where you submit the job
### run your MPI executable
```

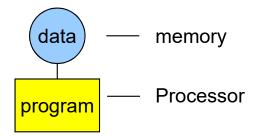
intra-node fabric: **shm** (shared-memory)

inter-node fabric: ofa (InfiniBand OFED verbs)

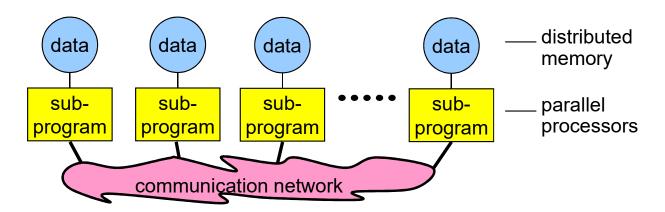
To submit the job: qsub impi.pbs

Programming Paradigms

Sequential Programming Paradigm

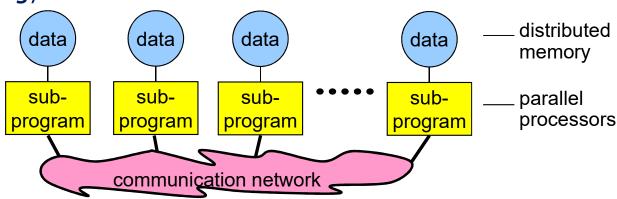


Message-Passing Programming Paradigm



Message-Passing Programming Paradigm

- Each processor in a message passing program runs a sub-program:
 - written in a conventional sequential language, e.g., C/C++ or Fortran
 - typically the same on each processor (SPMD Single Program Multiple Data):
 - the variables of each sub-program have the same names
 - but different locations (distributed memory) and different data!
 - i.e., all variables are private
 - but different sub-programs on different processors are allowed (MPMD)
 - sub-programs communicate via special send & receive routines (message passing)



SPMD vs. MPMD

- Both SPMD and MPMD are subcategories of MIMD (Multiple Instructions Multiple Data) in Flynn's taxonomy
- SPMD Single Program Multiple Data, e.g.:
 mpirun –n 32 ./mpi_hello.x
- MPMD Multiple Program Multiple Data, e.g.: mpirun -n 1 pwd : -n 1 hostname
- MPI allows MPMD
 - But some implementations may only support SPMD, not MPMD
 - MPMD can be emulated with SPMD

Emulation of MPMD with SPMD

C:

```
int main(int argc, char **argv) {
   if (myrank < ... /* process should run the ocean model */) {
      ocean( /* arguments */ );
   }
   else {
      weather( /* arguments */ );
   }
}</pre>
```

Fortran:

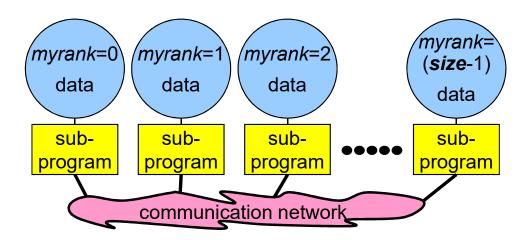
```
PROGRAM
    IF (myrank < ...) THEN! process should run the ocean model
        CALL ocean ( some arguments )
    ELSE
        CALL weather ( some arguments )
    ENDIF
END</pre>
```

Data and Work Distribution

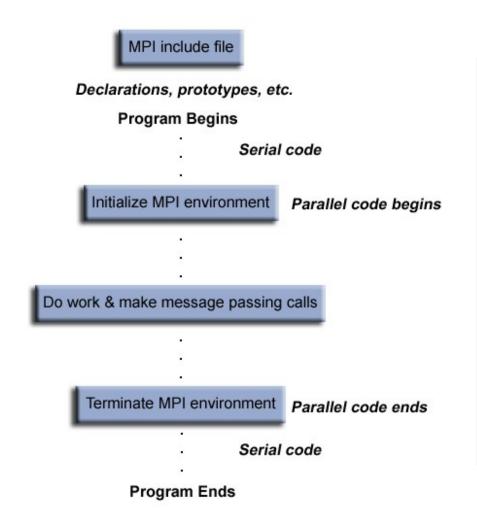
- The value of *myrank* is returned by a special library routine, e.g.,
 MPI_Comm_rank
- The value of *size* (# of processes to be started) is specified as argument(s) to a special MPI initialization program (*mpirun* or *mpiexec*), e.g:

```
mpirun -n 32 ./mpi_hello.x
```

 All distribution decisions are based on myrank, e.g., which process works on which data



General MPI Program Structure



```
#include "mpi.h"
#include <stdio.h>
int main(int argc, char* argv[])
  int rank, size:
 MPI_Init(&argc, &argv);
  MPI_Comm_rank(MPI_COMM_WORLD, &rank);
  MPI_Comm_size(MPI_COMM_WORLD, &size);
  /* do some work */
 MPI_Finalize();
  return 0;
```

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Header File

• Required for all programs that make MPI library calls.

C include file	Fortran include file	Fortran 90 module
#include "mpi.h"	include 'mpif.h'	USE mpi

• With MPI-3 Fortran, USE mpi_f08 is preferred.

Format of MPI Calls

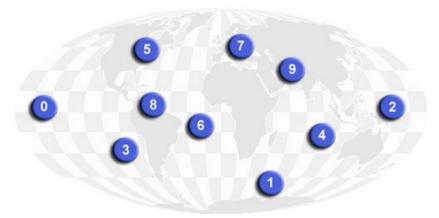
- C names are case sensitive; Fortran names are not.
- Programs must not declare variables or functions with names beginning with the prefix MPI_ or PMPI_ (profiling interface).

C Binding		
Format:	<pre>rc = MPI_Xxxxx(parameter,)</pre>	
Example:	<pre>rc = MPI_Bsend(&buf,count,type,dest,tag,comm)</pre>	
Error code:	Returned as "rc". MPI_SUCCESS if successful	
Fortran Binding		
Format:	CALL MPI_XXXXX(parameter,, ierr)	
Example:	CALL MPI_BSEND(buf,count,type,dest,tag,comm,ierr)	
Error code:	Returned as "ierr" parameter. MPI_SUCCESS if successful	

Communicators and Groups

- MPI uses objects called communicators and groups to define which collection of processes may communicate with each other.
- Most MPI routines require you to specify a communicator as an argument.
- Communicators and groups will be covered in more detail later. For now, simply use MPI_COMM_WORLD whenever a communicator is required it is the predefined communicator that includes all of your MPI processes.

 MPI_COMM_WORLD



Rank

- Within a communicator, every process has its own unique, integer identifier assigned by the system when the process initializes. A rank is sometimes also called a "task ID". Ranks are contiguous and begin at zero.
- Used by the programmer to specify the source and destination of messages. Often used conditionally by the application to control program execution (if rank=0 do this / if rank=1 do that).

Error Handling

- Most MPI routines include a return/error code parameter.
- However, according to the MPI standard, the default behavior of an MPI call is to abort if there is an error. This means you will probably not be able to capture a return/error code other than MPI_SUCCESS (zero).
- The standard does provide a means to override this default error handler. A discussion on how to do this is available at: https://computing.llnl.gov/tutorials/mpi/errorHandlers.pdf
- The types of errors displayed to the user are implementation dependent.

Environment Management Routines

This group of routines is used for interrogating and setting the MPI execution environment, and covers an assortment of purposes, such as initializing and terminating the MPI environment, querying a rank's identity, querying the MPI library's version, etc. The commonly used ones are:

- MPI_Init
- MPI_Comm_size
- MPI_Comm_rank
- MPI_Abort
- MPI_Get_processor_name

- MPI_Get_version
- MPI_Initialized
- MPI_Wtime
- MPI_Wtick
- MPI_Finalize

MPI_Init

Initializes the MPI execution environment. This function must be called in every MPI program, must be called before any other MPI functions and must be called only once in an MPI program. For C programs, MPI_Init may be used to pass the command line arguments to all processes, although this is not required by the standard and is implementation dependent.

```
#include <mpi.h>
int main(int argc, char **argv)
{
    MPI_Init(&argc, &argv);
    ....
```

```
program xxxxx
implicit none
include 'mpif.h'
integer ierror
call MPI_Init(ierror)
....
```

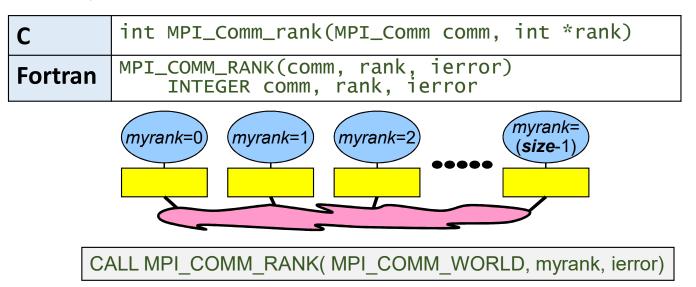
MPI_Comm_size

Returns the total number of MPI processes in the specified communicator, such as MPI_COMM_WORLD. If the communicator is MPI_COMM_WORLD, then it represents the number of MPI tasks available to your application.

С	<pre>int MPI_Comm_size(MPI_Comm comm, int *size)</pre>
Fortran	MPI_COMM_SIZE(comm, size, ierror) INTEGER comm, size, ierror

MPI_Comm_rank

Returns the rank of the calling MPI process within the specified communicator. Initially, each process will be assigned a unique integer rank between 0 and *number of tasks - 1 (size -1)* within the communicator **MPI_COMM_WORLD**. This rank is often referred to as a task ID. If a process becomes associated with other communicators, it will have a unique rank within each of these as well.



MPI_Get_processor_name

Returns the processor name. Also returns the length of the name. The buffer for "name" must be at least MPI_MAX_PROCESSOR_NAME characters in size. What is returned into "name" is implementation dependent - may not be the same as the output of the "hostname" or "host" shell commands.

С	<pre>int MPI_Get_processor_name(char *name, int *resultlen)</pre>
Fortran	MPI_GET_PROCESSOR_NAME(NAME, RESULTLEN, IERROR) CHARACTER*(*) NAME INTEGER RESULTLEN, IERROR

MPI_Get_version

Returns the version and subversion of the MPI standard that's implemented by the library.

С	<pre>int MPI_Get_version(int *version, int *subversion)</pre>
Fortran	MPI_GET_VERSION(VERSION, SUBVERSION, IERROR) INTEGER VERSION, SUBVERSION, IERROR

MPI_Initialized

Indicates whether MPI_Init has been called - returns flag as either logical true (1) or false(0). MPI requires that MPI_Init be called once and only once by each process. This may pose a problem for modules that want to use MPI and are prepared to call MPI_Init if necessary. MPI_Initialized solves this problem.

С	<pre>int MPI_Initialized(int *flag)</pre>
Fortran	MPI_INITIALIZED(FLAG, IERROR) LOGICAL FLAG INTEGER IERROR

MPI_Wtime

Returns an elapsed wall clock time in seconds (double precision) on the calling processor.

С	double MPI_Wtime()
Fortran	DOUBLE PRECISION MPI_WTIME()

MPI_Wtick

Returns the resolution in seconds (double precision) of MPI_Wtime.

С	<pre>double MPI_Wtick()</pre>
Fortran	DOUBLE PRECISION MPI_WTICK()

MPI_Finalize

Terminates the MPI execution environment. This function should be the last MPI routine called in every MPI program - no other MPI routines may be called after it.

С	int MPI_Finalize()
Fortran	MPI_FINALIZE(IERROR) INTEGER IERROR

Handles

- Handles identify MPI objects
- For the programmer, handles are:
 - predefined constants in mpi.h or mpif.h
 - example: MPI_COMM_WORLD
 - predefined values exist only after MPI_Init was called
 - values returned by some MPI routines, to be stored in variables,
 that are defined as
 - o in Fortran: INTEGER
 - in C: special MPI typedefs, e.g., MPI_Comm
- Handles refer to internal MPI data structures

Environment Management Routines Example in C

```
#include "mpi.h"
#include <stdio.h>
int main(int argc, char *argv[]) {
    int size, rank, len, rc;
    char hostname[MPI_MAX_PROCESSOR_NAME];
    rc = MPI_Init(&argc, &argv);
    if (rc != MPI_SUCCESS) {
       printf ("Error starting MPI program. Terminating.\n");
       MPI_Abort(MPI_COMM_WORLD, rc);
    MPI_Comm_size(MPI_COMM_WORLD, &size);
    MPI_Comm_rank(MPI_COMM_WORLD, &rank);
    MPI_Get_processor_name(hostname, &len);
    printf ("Number of tasks= %d My rank= %d Running on %s\n", size, rank,
             hostname):
    /***** do some work *****/
    MPI_Finalize():
                                                                        50
```

Environment Management Routines Example in Fortran

```
program simple
     include 'mpif.h'
     integer size, rank, len, ierr
      character(MPI_MAX_PROCESSOR_NAME) hostname
     call MPI_INIT(ierr)
     if (ierr .ne. MPI_SUCCESS) then
         print *, 'Error starting MPI program. Terminating.'
         call MPI_ABORT(MPI_COMM_WORLD, rc, ierr)
     end if
     call MPI_COMM_RANK(MPI_COMM_WORLD, rank, ierr)
     call MPI_COMM_SIZE(MPI_COMM_WORLD, size, ierr)
     call MPI_GET_PROCESSOR_NAME(hostname, len, ierr)
     print *, 'Number of tasks=', size, ' My rank=', rank,
               ' Running on=', hostname
C ***** do some work *****
     call MPI_FINALIZE(ierr)
      end
                                                                        51
```

Point-to-Point Communication

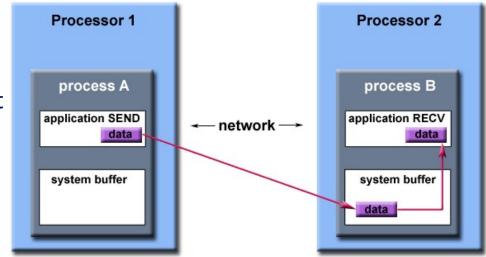
- MPI point-to-point operations typically involve message passing between two, and only two, different MPI tasks/processes. One task is performing a send operation and the other task is performing a matching receive operation.
- There are different types of send and receive routines used for different purposes. For example:
 - Synchronous send
 - Blocking send / blocking receive
 - Non-blocking send / non-blocking receive
 - o Buffered send
 - Combined send/receive
 - o "Ready" send
- Any type of send routine can be paired with any type of receive routine.

Buffering

- In a perfect world, every send operation would be perfectly synchronized with its matching receive. This is rarely the case. The MPI implementation must be able to deal with storing data when the two tasks are out of sync.
- Consider the following two cases:
 - O A send operation occurs 5 seconds before the receive is ready where is the message while the receive is pending?
 - O Multiple sends arrive at the same receiving task which can only accept one send at a time - what happens to the messages that are "backing up"?
- The MPI implementation (not the MPI standard) decides what happens to data in these types of cases. Typically, a **system buffer** area is reserved to hold data in transit.

Buffering (cont'd)

- System buffer space is:
 - Opaque to the programmer and managed entirely by the MPI library
 - A finite resource that can be easy to exhaust
 - Often mysterious and not well documented
 - Able to exist on the sending side, the receiving side, or both
 - Something that may improve program performance because it allows send - receive operations to be asynchronous.



Path of a message buffered at the receiving process

 User managed address space (i.e. your program variables) is called the application buffer. MPI also provides for a user managed send buffer.

Blocking vs. Non-blocking

• Most of the MPI point-to-point routines can be used in either blocking or non-blocking mode.

Blocking

- A blocking send routine will only "return" after it is safe to modify the *application buffer* (your send data) for reuse. Safe means that modifications will not affect the data intended for the receive task. Safe does not imply that the data was actually received it may very well be sitting in a *system buffer*.
- A blocking send can be synchronous which means there is handshaking occurring with the receive task to confirm a safe send.
- A blocking send can be *asynchronous* if a system buffer is used to hold the data for eventual delivery to the receive.
- A blocking receive only "returns" after the data has arrived and is ready for use by the program.

Non-blocking

- Non-blocking send and receive routines behave similarly they will return almost immediately. They do not wait for any communication events to complete, such as message copying from user memory to system buffer space or the actual arrival of message.
- Non-blocking operations simply "request" the MPI library to perform the operation when it is able. The user can not predict when that will happen.
- It is unsafe to modify the *application buffer* (your variable space) until you know for a fact the requested non-blocking operation was actually performed by the library. There are "wait" routines used to do this.
- Non-blocking communications are primarily used to overlap computation with communication and exploit possible performance gains.

Order and Fairness

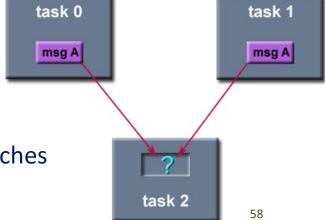
• Order:

- MPI guarantees that messages will not overtake each other.
- If a sender sends two messages (Message 1 and Message 2) in succession to the same destination, and both match the same receive, the receive operation will receive Message 1 before Message 2.
- o If a receiver posts two receives (Receive 1 and Receive 2), in succession, and both are looking for the same message, Receive 1 will receive the message before Receive 2.
- Order rules do not apply if there are multiple threads participating in the communication operations.

• Fairness:

 MPI does not guarantee fairness - it's up to the programmer to prevent "operation starvation".

Example: task 0 sends a message to task 2.
 However, task 1 sends a competing message that matches task 2's receive. Only one of the sends will complete.



MPI Message Passing Routine Arguments

MPI point-to-point communication routines generally have an argument list that takes one of the following formats:

Blocking send	<pre>MPI_Send(buffer,count,type,dest,tag,comm)</pre>
Non-blocking send	<pre>MPI_Isend(buffer,count,type,dest,tag,comm,request)</pre>
Blocking receive	<pre>MPI_Recv(buffer,count,type,source,tag,comm,status)</pre>
Non-blocking receive	<pre>MPI_Irecv(buffer,count,type,source,tag,comm,request)</pre>

Buffer

Program (application) address space that references the data that is to be sent or received. In most cases, this is simply the variable name that is be sent/received. For **C** programs, this argument is passed by reference and usually must be prepended with an ampersand: &var1

Data Count

Indicates the number of data elements of a particular type to be sent.

MPI Message Passing Routine Arguments (cont'd)

Data Type

For reasons of portability, MPI predefines its elementary data types. Programmers may also create their own data types.

○ C Data Types:

MPI_CHAR, MPI_WCHAR, MPI_SHORT, MPI_INT, MPI_LONG, MPI_LONG_LONG_INT, MPI_LONG_LONG, MPI_SIGNED_CHAR, MPI_UNSIGNED_CHAR, MPI_UNSIGNED_SHORT, MPI_UNSIGNED, MPI_UNSIGNED_LONG, MPI_UNSIGNED_LONG, MPI_FLOAT, MPI_DOUBLE, MPI_LONG_DOUBLE, MPI_C_COMPLEX, MPI_C_FLOAT_COMPLEX, MPI_C_DOUBLE_COMPLEX, MPI_C_BOOL, MPI_INT8_T, MPI_INT16_T, MPI_INT32_T, MPI_INT64_T, MPI_BYTE, MPI_PACKED

o Fortran Data Types:

MPI_CHARACTER, MPI_INTEGER, MPI_INTEGER1, MPI_INTEGER2, MPI_INTEGER4, MPI_REAL, MPI_REAL2, MPI_REAL4, MPI_REAL8, MPI_DOUBLE_PRECISION, MPI_COMPLEX, MPI_DOUBLE_COMPLEX, MPI_LOGICAL, MPI_BYTE, MPI_PACKED

MPI Message Passing Routine Arguments (cont'd)

Destination

An argument to send routines that indicates the process where a message should be delivered. Specified as the rank of the receiving process.

Source

An argument to receive routines that indicates the originating process of the message. Specified as the rank of the sending process. This may be set to the wild card MPI_ANY_SOURCE to receive a message from any task.

Tag

Arbitrary non-negative integer assigned by the programmer to uniquely identify a message. Send and receive operations should match message tags. For a receive operation, the wild card **MPI_ANY_TAG** can be used to receive any message regardless of its tag. The MPI standard guarantees that integers 0-32767 can be used as tags, but most implementations allow a much larger range than this.

MPI Message Passing Routine Arguments (cont'd)

Communicator

Indicates the communication context, or set of processes for which the source or destination fields are valid. Unless the programmer is explicitly creating new communicators, the predefined communicator **MPI_COMM_WORLD** is usually used.

Status

For a receive operation, indicates the source of the message and the tag of the message. In C, this argument is a *pointer* to a predefined structure **MPI_Status**. In Fortran, it is an integer array of size **MPI_STATUS_SIZE**. Additionally, the actual number of bytes received is obtainable from *Status* via the **MPI_Get_count** routine.

Request

Used by non-blocking send and receive operations. Since non-blocking operations may return before the requested system buffer space is obtained, the system issues a unique "request number". The programmer uses this system assigned "handle" later (in a *WAIT* type routine) to determine completion of the non-blocking operation. In C, this argument is a pointer to a predefined structure **MPI_Request**. In Fortran, it is an integer.

Blocking Point-to-Point Communication Routines

The more commonly used MPI blocking message passing routines are:

- MPI_Send
- MPI_Recv
- MPI_Ssend
- MPI_Bsend
- MPI_Buffer_attach
- MPI_Buffer_detach

- MPI_Rsend
- MPI_Sendrecv
- MPI_Wait
- MPI_Waitany
- MPI_Waitall
- MPI_Waitsome
- MPI_Probe

Open MPI documentation:

https://www.open-mpi.org/doc/v1.10/

MPI_Send

Basic blocking send operation. Routine returns only after the application buffer in the sending task is free for reuse. Note that this routine may be implemented differently on different systems. The MPI standard permits the use of a system buffer but does not require it. Some implementations may actually use a synchronous send (discussed shortly) to implement the basic blocking send.

С	<pre>int MPI_Send(const void *buf, int count, MPI_Datatype datatype, int dest, int tag, MPI_Comm comm)</pre>
Fortran	MPI_SEND(BUF, COUNT, DATATYPE, DEST, TAG, COMM, IERROR) <type> BUF(*) INTEGER COUNT, DATATYPE, DEST, TAG, COMM, IERROR</type>

MPI_Recv

Receive a message and block until the requested data is available in the application buffer in the receiving task.

С	<pre>int MPI_Recv(void *buf, int count, MPI_Datatype datatype, int source, int tag, MPI_Comm comm, MPI_Status *status)</pre>
Fortran	MPI_RECV(BUF, COUNT, DATATYPE, SOURCE, TAG, COMM, STATUS, IERROR) <pre></pre>

MPI_Ssend

Synchronous blocking send: Send a message and block until the application buffer in the sending task is free for reuse and the destination process has started to receive the message.

	<pre>int MPI_Ssend(const void *buf, int count, MPI_Datatype datatype,</pre>
Fortran	MPI_SSEND(BUF, COUNT, DATATYPE, DEST, TAG, COMM, IERROR) <type> BUF(*) INTEGER COUNT, DATATYPE, DEST, TAG, COMM, IERROR</type>

MPI_Bsend

Buffered blocking send: permits the programmer to allocate the required amount of buffer space into which data can be copied until it is delivered. Insulates against the problems associated with insufficient system buffer space. Routine returns after the data has been copied from application buffer space to the allocated send buffer. Must be used with the MPI_Buffer_attach routine.

```
MPI_Buffer_attach( b, n*sizeof(double) + MPI_BSEND_OVERHEAD );
for (i=0; i<m; i++) {
    MPI_Bsend( buf, n, MPI_DOUBLE, ... );
}</pre>
```

MPI_Buffer_attach & MPI_Buffer_detach

Used by programmer to allocate/deallocate message buffer space to be used by the MPI_Bsend routine. The size argument is specified in actual data bytes - not a count of data elements. Only one buffer can be attached to a process at a time.

С	<pre>int MPI_Buffer_attach(void *buf, int size)</pre>
	<pre>int MPI_Buffer_detach(void *buf, int *size)</pre>
Fortran	MPI_BUFFER_ATTACH(BUF, SIZE, IERROR) <type> BUF(*) INTEGER SIZE, IERROR</type>
	MPI_BUFFER_DETACH(BUF, SIZE, IERROR) <type> BUF(*) INTEGER SIZE, IERROR</type>

MPI_Rsend

Blocking ready send. Should only be used if the programmer is certain that the matching receive has already been posted. Often simply implemented as an **MPI_Send** routine.

С	<pre>int MPI_Rsend(const void *buf, int count, MPI_Datatype datatype, int dest, int tag, MPI_Comm comm)</pre>
Fortran	MPI_RSEND(BUF, COUNT, DATATYPE, DEST, TAG, COMM, IERROR) <type> BUF(*) INTEGER COUNT, DATATYPE, DEST, TAG, COMM, IERROR</type>

MPI_Sendrecv

Send a message and post a receive before blocking. Will block until the sending application buffer is free for reuse and until the receiving application buffer contains the received message.

С	<pre>int MPI_Sendrecv(const void *sendbuf, int sendcount,</pre>
Fortran	MPI_SENDRECV(SENDBUF, SENDCOUNT, SENDTYPE, DEST, SENDTAG, RECVBUF, RECVCOUNT, RECVTYPE, SOURCE, RECVTAG, COMM, STATUS, IERROR) <type> SENDBUF(*), RECVBUF(*) INTEGER SENDCOUNT, SENDTYPE, DEST, SENDTAG INTEGER RECVCOUNT, RECVTYPE, SOURCE, RECVTAG, COMM INTEGER STATUS(MPI_STATUS_SIZE), IERROR</type>

MPI_Wait

MPI_Wait blocks until a specified *non-blocking* send or receive operation has completed. For multiple *non-blocking* operations, the programmer can specify any, all or some completions.

```
int MPI_Wait(MPI_Request *request, MPI_Status *status)
int MPI_Waitany(int count, MPI_Request array_of_requests[],
    int *index, MPI_Status *status)
int MPI_Waitall(int count, MPI_Request array_of_requests[],
        MPI_Status *array_of_statuses)
int MPI_Waitsome(int incount, MPI_Request array_of_requests[],
        int *outcount, int array_of_indices[],
        MPI_Status array_of_statuses[])

Fortran

MPI_WAIT(REQUEST, STATUS, IERROR)
MPI_WAITANY(COUNT, ARRAY_OF_REQUESTS, INDEX, STATUS, IERROR)
MPI_WAITSOME(INCOUNT, ARRAY_OF_REQUESTS, OUTCOUNT,
        ARRAY_OF_INDICES, ARRAY_OF_STATUSES, IERROR)
```

MPI_Probe

Performs a blocking test for a message. The "wildcards" MPI_ANY_SOURCE and MPI_ANY_TAG may be used to test for a message from any source or with any tag. For the C routine, the actual source and tag will be returned in the *status* structure as *status.MPI_SOURCE* and *status.MPI_TAG*. For the Fortran routine, they will be returned in the integer array *status(MPI_SOURCE)* and *status(MPI_TAG)*.

С	<pre>int MPI_Probe(int source, int tag, MPI_Comm comm, MPI_Status *status)</pre>
Fortran	MPI_PROBE(SOURCE, TAG, COMM, STATUS, IERROR) INTEGER SOURCE, TAG, COMM, STATUS(MPI_STATUS_SIZE), IERROR

Blocking Message Passing Routines Example in C

```
#include "mpi.h"
#include <stdio.h>
int main(int argc, char *argv[]) {
  int size, rank, dest, source, rc, count, tag=1;
  char inmsq, outmsq='x';
  MPI_Status Stat:
  MPI_Init(&argc,&argv);
  MPI_Comm_size(MPI_COMM_WORLD, &size);
  MPI_Comm_rank(MPI_COMM_WORLD, &rank);
  if (rank == 0) {
    dest = 1:
    source = 1;
    rc = MPI_Send(&outmsg, 1, MPI_CHAR, dest, tag, MPI_COMM_WORLD);
    rc = MPI_Recv(&inmsg, 1, MPI_CHAR, source, tag, MPI_COMM_WORLD, &Stat);
```

Try it out

```
$ mpiicc ping.c -o ping.x
$ mpirun -n 2 ./ping.x
Task 0: Received 1 char(s) from task 1 with tag 1
Task 1: Received 1 char(s) from task 0 with tag 1
```

Blocking Message Passing Routines Example in Fortran

```
program ping
 include 'mpif.h'
 integer size, rank, dest, source, count, tag, ierr
 integer stat(MPI_STATUS_SIZE)
 character inmsg, outmsg
 outmsq = 'x'
 tag = 1
 call MPI_INIT(ierr)
 call MPI_COMM_RANK(MPI_COMM_WORLD, rank, ierr)
 call MPI_COMM_SIZE(MPI_COMM_WORLD, size, ierr)
 if (rank .eq. 0) then
   dest = 1
   source = 1
   call MPI_SEND(outmsg, 1, MPI_CHARACTER, dest, tag,
            MPI_COMM_WORLD, ierr)
&
   call MPI_RECV(inmsg, 1, MPI_CHARACTER, source, tag,
            MPI_COMM_WORLD, stat, ierr)
&
```

```
else if (rank .eq. 1) then
   dest = 0
   source = 0
   call MPI_RECV(inmsg, 1, MPI_CHARACTER, source, tag,
&
            MPI_COMM_WORLD, stat, err)
   call MPI_SEND(outmsg, 1, MPI_CHARACTER, dest, tag,
&
            MPI_COMM_WORLD, err)
 endif
 call MPI_GET_COUNT(stat, MPI_CHARACTER, count, ierr)
 print *, 'Task ',rank,': Received', count, 'char(s) from task',
          stat(MPI_SOURCE), 'with tag', stat(MPI_TAG)
 call MPI_FINALIZE(ierr)
 end
```

Try it out

```
$ mpiifort ping.f -o ping.x
$ mpirun -n 2 ./ping.x
Task 0: Received 1 char(s) from task 1 with tag 1
Task 1: Received 1 char(s) from task 0 with tag 1
```

Non-Blocking Point-to-Point Communication Routines

The more commonly used MPI non-blocking message passing routines are:

- MPI_Isend
- MPI_Irecv
- MPI_Issend
- MPI_Ibsend
- MPI_Irsend

- MPI_Test
- MPI_Testany
- MPI_Testall
- MPI_Testsome
- MPI_Iprobe

Open MPI documentation:

https://www.open-mpi.org/doc/v1.10/

MPI_Isend

Identifies an area in memory to serve as a send buffer. Processing continues *immediately* without waiting for the message to be copied out from the application buffer. A communication *request handle* is returned for handling the pending message status. The program should not modify the application buffer until subsequent calls to **MPI_Wait** or **MPI_Test** indicate that the non-blocking send has completed.

С		<pre>int MPI_Isend(const void *buf, int count, MPI_Datatype datatype, int dest, int tag, MPI_Comm comm, MPI_Request *request)</pre>
For	tran	MPI_ISEND(BUF, COUNT, DATATYPE, DEST, TAG, COMM, REQUEST, IERROR) <type> BUF(*) INTEGER COUNT, DATATYPE, DEST, TAG, COMM, REQUEST, IERROR</type>

MPI_Irecv

Identifies an area in memory to serve as a receive buffer. Processing continues *immediately* without actually waiting for the message to be received and copied into the the application buffer. A communication *request handle* is returned for handling the pending message status. The program must use calls to **MPI_Wait** or **MPI_Test** to determine when the non-blocking receive operation completes and the requested message is available in the application buffer.

С	<pre>int MPI_Irecv(void *buf, int count, MPI_Datatype datatype, int source, int tag, MPI_Comm comm, MPI_Request *request)</pre>
Fortran	<pre>MPI_ISEND(BUF, COUNT, DATATYPE, DEST, TAG, COMM, REQUEST, IERROR)</pre>

MPI_Issend

Non-blocking synchronous send. Similar to MPI_Isend, except MPI_Wait or MPI_Test indicates when the destination process has received the message.

С		<pre>int MPI_Issend(const void *buf, int count, MPI_Datatype datatype, int dest, int tag, MPI_Comm comm, MPI_Request *request)</pre>
Fo	rtran	<pre>MPI_ISSEND(BUF, COUNT, DATATYPE, DEST, TAG, COMM, REQUEST, IERROR)</pre>

MPI_Ibsend

Non-blocking buffered send. Similar to MPI_Bsend except MPI_Wait or MPI_Test indicates when the destination process has received the message. Must be used with the MPI_Buffer_attach routine.

С	<pre>int MPI_Ibsend(const void *buf, int count, MPI_Datatype datatype, int dest, int tag, MPI_Comm comm, MPI_Request *request)</pre>
Fortran	MPI_IBSEND(BUF, COUNT, DATATYPE, DEST, TAG, COMM, REQUEST, IERROR) <type> BUF(*) INTEGER COUNT, DATATYPE, DEST, TAG, COMM, REQUEST, IERROR</type>

MPI_Irsend

Non-blocking ready send. Similar to MPI_Rsend except MPI_Wait or MPI_Test indicates when the destination process has received the message. Should only be used if the programmer is certain that the matching receive has already been posted.

С		<pre>int MPI_Irsend(const void *buf, int count, MPI_Datatype datatype, int dest, int tag, MPI_Comm comm, MPI_Request *request)</pre>
Fc	ortran	<pre>MPI_IRSEND(BUF, COUNT, DATATYPE, DEST, TAG, COMM, REQUEST, IERROR)</pre>

MPI_Test

MPI_Test checks the status of a specified non-blocking send or receive operation. The "flag" parameter is returned logical true (1) if the operation has completed, and logical false (0) if not. For multiple non-blocking operations, the programmer can specify any, all or some completions.

C	<pre>int MPI_Test(MPI_Request *request, int *flag, MPI_Status *status) int MPI_Testany(int count, MPI_Request array_of_requests[], int *index, int *flag, MPI_Status *status) int MPI_Testall(int count, MPI_Request array_of_requests[], int *flag, MPI_Status array_of_statuses[]) int MPI_Testsome(int incount, MPI_Request array_of_requests[], int *outcount, int array_of_indices[], MPI_Status array_of_statuses[])</pre>
Fortran	MPI_TEST(REQUEST, FLAG, STATUS, IERROR) MPI_TESTANY(COUNT, ARRAY_OF_REQUESTS, INDEX, FLAG, STATUS, IERROR) MPI_WAITALL(COUNT, ARRAY_OF_REQUESTS, ARRAY_OF_STATUSES, IERROR) MPI_TESTALL(COUNT, ARRAY_OF_REQUESTS, FLAG, ARRAY_OF_STATUSES, IERROR) MPI_TESTSOME(INCOUNT, ARRAY_OF_REQUESTS, OUTCOUNT,

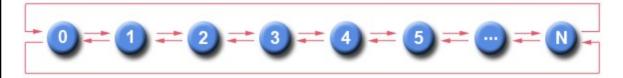
MPI_Iprobe

Performs a non-blocking test for a message. The "wildcards" MPI_ANY_SOURCE and MPI_ANY_TAG may be used to test for a message from any source or with any tag. The integer "flag" parameter is returned logical true (1) if a message has arrived, and logical false (0) if not. For the C routine, the actual source and tag will be returned in the status structure as <code>status.MPI_SOURCE</code> and <code>status.MPI_TAG</code>. For the Fortran routine, they will be returned in the integer array <code>status(MPI_SOURCE)</code> and <code>status(MPI_TAG)</code>.

С	<pre>int MPI_Iprobe(int source, int tag, MPI_Comm comm, int *flag,</pre>
Fortran	MPI_IPROBE(SOURCE, TAG, COMM, FLAG, STATUS, IERROR) LOGICAL FLAG
	<pre>INTEGER SOURCE, TAG, COMM, STATUS(MPI_STATUS_SIZE), IERROR</pre>

Non-Blocking Message Passing Routines Example in C

Nearest neighbor exchange in a ring topology



```
#include "mpi.h"
#include <stdio.h>

int main(int argc, char *argv[]) {
   int size, rank, next, prev, buf[2], tag1=1, tag2=2;
   MPI_Request reqs[4];
   MPI_Status stats[4];

   MPI_Init(&argc,&argv);
   MPI_Comm_size(MPI_COMM_WORLD, &size);
   MPI_Comm_rank(MPI_COMM_WORLD, &rank);
```

```
prev = rank-1;
next = rank+1:
if (rank == 0) prev = size - 1;
if (rank == (size - 1)) next = 0;
MPI_Irecv(&buf[0], 1, MPI_INT, prev, tag1, MPI_COMM_WORLD, &regs[0]);
MPI_Irecv(&buf[1], 1, MPI_INT, next, tag2, MPI_COMM_WORLD, &reqs[1]);
MPI_Isend(&rank, 1, MPI_INT, prev, tag2, MPI_COMM_WORLD, &regs[2]);
MPI_Isend(&rank, 1, MPI_INT, next, tag1, MPI_COMM_WORLD, &reqs[3]);
/* do some work */
MPI_Waitall(4, regs, stats);
MPI_Finalize();
```

Non-Blocking Message Passing Routines Example in Fortran

Nearest neighbor exchange in a ring topology



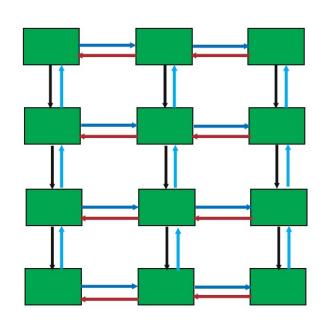
```
program ringtopo
include 'mpif.h'

integer size, rank, next, prev, buf(2), tag1, tag2, ierr
integer stats(MPI_STATUS_SIZE,4), reqs(4)
tag1 = 1
tag2 = 2

call MPI_INIT(ierr)
call MPI_COMM_RANK(MPI_COMM_WORLD, rank, ierr)
call MPI_COMM_SIZE(MPI_COMM_WORLD, size, ierr)
prev = rank - 1
next = rank + 1
```

```
if (rank .eq. 0) then
   prev = size - 1
 endif
 if (rank .eq. size - 1) then
  next = 0
 endif
 call MPI_IRECV(buf(1), 1, MPI_INTEGER, prev, tag1,
&
          MPI_COMM_WORLD, regs(1), ierr)
 call MPI_IRECV(buf(2), 1, MPI_INTEGER, next, tag2,
          MPI_COMM_WORLD, regs(2), ierr)
 call MPI_ISEND(rank, 1, MPI_INTEGER, prev, tag2,
          MPI_COMM_WORLD, reqs(3), ierr)
&
 call MPI_ISEND(rank, 1, MPI_INTEGER, next, tag1,
&
          MPI_COMM_WORLD, reqs(4), ierr)
do some work
 call MPI_WAITALL(4, reqs, stats, ierr);
 call MPI_FINALIZE(ierr)
 end
                                                                   90
```

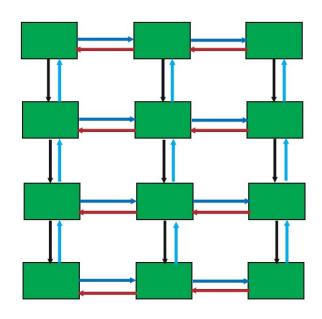
Mesh Exchange



Deadlocks!

Fix 1: Sequentialization

```
if ( odd patch ) then
  call MPI_Send( ... down ... )
  call MPI_Recv( ... down ... )
endif
if ( even patch ) then
  call MPI_Recv( ... up ... )
  call MPI_Send( ... up ... )
endif
...
```



Fix 2: Using Irecv

This fix does not perform well in practice. Why?

Fix 3: Using Irecv and Isend

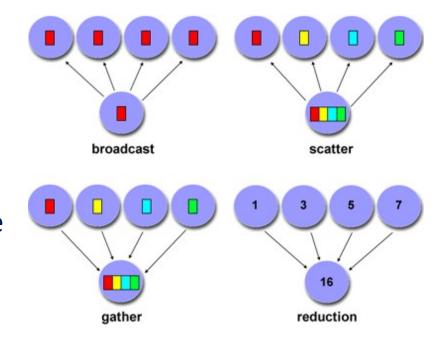
 All of the sends may block, waiting for a matching receive (will for large enough messages) – call Irecv first

Collective Communication Routines

- Collective communication routines must involve **all** processes within the scope of a communicator.
 - All processes are, by default, members in the communicator MPI_COMM_WORLD.
 - Additional communicators can be defined by the programmer.
- Unexpected behavior, including program failure, can occur if even one task in the communicator doesn't participate.
- It is the programmer's responsibility to ensure that all processes within a communicator participate in any collective operations.

Types of Collective Operations

- **Synchronization** processes wait until all members of the group have reached the synchronization point.
- Data Movement broadcast, scatter/gather, all to all.
- Collective Computation (reductions) one member of the group collects data from the other members and performs an operation (min, max, add, multiply, etc.) on that data.



Programming Considerations and Restrictions

- Collective communication routines do not take message tag arguments.
- Collective operations within subsets of processes are accomplished by first partitioning the subsets into new groups and then attaching the new groups to new communicators.
- Can only be used with MPI predefined datatypes not with MPI Derived Data Types.
- MPI-2 extended most collective operations to allow data movement between intercommunicators.
- With MPI-3, collective operations can be blocking or non-blocking.

MPI_Barrier

Synchronization operation. Creates a barrier synchronization in a group. Each task, when reaching the MPI_Barrier call, blocks until all tasks in the group reach the same MPI_Barrier call. Then all tasks are free to proceed.

С	<pre>int MPI_Barrier(MPI_Comm comm)</pre>
Fortran	MPI_BARRIER(COMM, IERROR) INTEGER COMM, IERROR

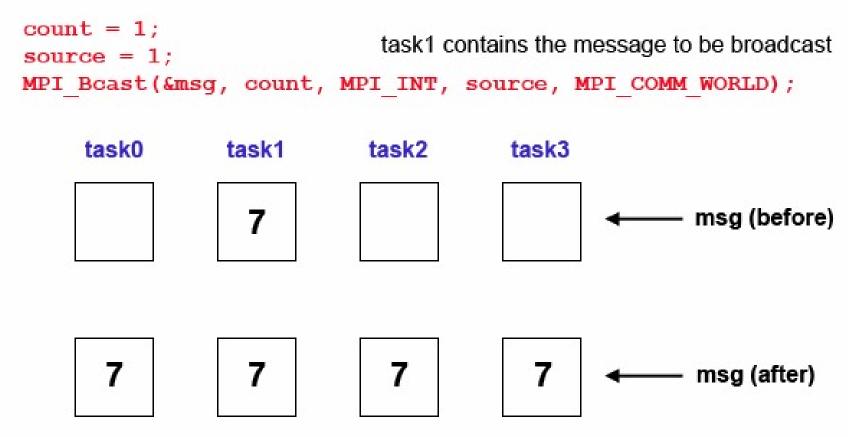
MPI_Bcast

Data movement operation. Broadcasts (sends) a message from the process with rank "root" to all other processes in the group.

С		<pre>int MPI_Bcast(void *buffer, int count, MPI_Datatype datatype, int root, MPI_Comm comm)</pre>
Foi	rtran	<pre>MPI_BCAST(BUFFER, COUNT, DATATYPE, ROOT, COMM, IERROR)</pre>

MPI_Bcast

Broadcasts a message from one task to all other tasks in communicator



MPI_Scatter

Data movement operation. Distributes distinct messages from a single source task to each task in the group.

С	<pre>int MPI_Scatter(const void *sendbuf, int sendcount, MPI_Datatype sendtype, void *recvbuf, int recvcount, MPI_Datatype recvtype, int root, MPI_Comm comm)</pre>
Fortran	MPI_SCATTER(SENDBUF, SENDCOUNT, SENDTYPE, RECVBUF, RECVCOUNT, RECVTYPE, ROOT, COMM, IERROR) <type> SENDBUF(*), RECVBUF(*) INTEGER SENDCOUNT, SENDTYPE, RECVCOUNT, RECVTYPE, ROOT INTEGER COMM, IERROR</type>

MPI_Scatter

Sends data from one task to all other tasks in communicator

```
sendont = 1;
recvent = 1;
                                task1 contains the data to be scattered
src = 1;
MPI Scatter(sendbuf, sendcnt, MPI INT
             recybuf, recycnt, MPI INT
             src, MPI COMM WORLD);
task0
            task1
                       task2
                                  task3
             2
                                                 - sendbuf (before)
              3
             4
                                                 recvbuf (after)
             2
                         3
```

MPI_Gather

Data movement operation. Gathers distinct messages from each task in the group to a single destination task. This routine is the reverse operation of MPI_Scatter.

С	<pre>int MPI_Gather(const void *sendbuf, int sendcount, MPI_Datatype sendtype, void *recvbuf, int recvcount, MPI_Datatype recvtype, int root, MPI_Comm comm)</pre>
Fortran	MPI_GATHER(SENDBUF, SENDCOUNT, SENDTYPE, RECVBUF, RECVCOUNT, RECVTYPE, ROOT, COMM, IERROR) <type> SENDBUF(*), RECVBUF(*) INTEGER SENDCOUNT, SENDTYPE, RECVCOUNT, RECVTYPE, ROOT INTEGER COMM, IERROR</type>

MPI_Gather

Gathers data from all tasks in communicator to a single task

```
sendcnt = 1;
recvcnt = 1;
                                  message will be gathered into task1
src = 1;
MPI Gather(sendbuf, sendont, MPI INT
            recybuf, recycnt, MPI INT
            src, MPI COMM WORLD);
task0
           task1
                                  task3
                       task2
             2
                         3
  1
                                   4
                                                - sendbuf (before)
             2
                                                 recvbuf (after)
             3
             4
```

MPI_Allgather

Data movement operation. Concatenation of data to all tasks in a group. Each task in the group, in effect, performs a one-to-all broadcasting operation within the group.

С	<pre>int MPI_Allgather(const void *sendbuf, int sendcount, MPI_Datatype sendtype, void *recvbuf, int recvcount, MPI_Datatype recvtype, MPI_Comm comm)</pre>
Fortran	MPI_ALLGATHER(SENDBUF, SENDCOUNT, SENDTYPE, RECVBUF, RECVCOUNT, RECVTYPE, COMM, IERROR) <type> SENDBUF (*), RECVBUF (*) INTEGER SENDCOUNT, SENDTYPE, RECVCOUNT, RECVTYPE, COMM, INTEGER IERROR</type>

MPI_Allgather

Gathers data from all tasks and then distributes to all tasks in communicator

```
sendont = 1;
recvent = 1;
MPI Allgather (sendbuf, sendont, MPI INT
               recvbuf, recvcnt, MPI_INT
               MPI COMM WORLD);
     task0
                task1
                           task2
                                      task3
                             3
                                                   sendbuf (before)
                                        1
                             2
                                                   recvbuf (after)
      3
                             3
      4
                                        4
                 4
                             4
```

MPI_Reduce

Collective computation operation. Applies a reduction operation on all tasks in the group and places the result in one task.

С	<pre>int MPI_Reduce(const void *sendbuf, void *recvbuf, int count,</pre>
Fortran	MPI_REDUCE(SENDBUF, RECVBUF, COUNT, DATATYPE, OP, ROOT, COMM, IERROR) <type> SENDBUF(*), RECVBUF(*) INTEGER COUNT, DATATYPE, OP, ROOT, COMM, IERROR</type>

MPI_Reduce

Perform reduction across all tasks in communicator and store result in 1 task

```
count = 1;
                                                task1 will contain result
dest = 1;
MPI Reduce (sendbuf, recvbuf, count, MPI INT,
            MPI SUM, dest, MPI COMM WORLD);
     task0
                 task1
                             task2
                                         task3

    sendbuf (before)

                                           4
                  10
                                                        recvbuf (after)
```

Predefined MPI Reduction Operations

MPI Red	uction Operation	C Data Types	Fortran Data Types
MPI_MAX	maximum	integer, float	integer, real, complex
MPI_MIN	minimum	integer, float	integer, real, complex
MPI_SUM	sum	integer, float	integer, real, complex
MPI_PROD	product	integer, float	integer, real, complex
MPI_LAND	logical AND	integer	logical
MPI_BAND	bit-wise AND	integer, MPI_BYTE	integer, MPI_BYTE
MPI_LOR	logical OR	integer	logical
MPI_BOR	bit-wise OR	integer, MPI_BYTE	integer, MPI_BYTE
MPI_LXOR	logical XOR	integer	logical
MPI_BXOR	bit-wise XOR	integer, MPI_BYTE	integer, MPI_BYTE
MPI_MAXLOC	max value and location	float, double and long double	real, complex, double precision
MPI_MINLOC	min value and location	float, double and long double	real, complex, double precision

Users can also define their own reduction functions by using the MPI_Op_create routine.

MPI_Allreduce

Collective computation operation + data movement. Applies a reduction operation and places the result in all tasks in the group. This is equivalent to an MPI_Reduce followed by an MPI_Bcast.

С	<pre>int MPI_Allreduce(const void *sendbuf, void *recvbuf, int count,</pre>
Fortran	MPI_ALLREDUCE(SENDBUF, RECVBUF, COUNT, DATATYPE, OP, COMM, IERROR) <type> SENDBUF(*), RECVBUF(*) INTEGER COUNT, DATATYPE, OP, COMM, IERROR</type>

MPI_Allreduce

Perform reduction and store result across all tasks in communicator

```
count = 1;
MPI Allreduce (sendbuf, recvbuf, count, MPI_INT,
               MPI SUM, MPI COMM WORLD);
    task0
               task1
                          task2
                                      task3
                                                  - sendbuf (before)
                 2
                            3
                                        4
     10
                 10
                            10
                                       10
                                                 recvbuf (after)
```

MPI_Reduce_scatter

Collective computation operation + data movement. First does an element-wise reduction on a vector across all tasks in the group. Next, the result vector is split into disjoint segments and distributed across the tasks. This is equivalent to an MPI_Reduce followed by an MPI_Scatter operation.

С	<pre>int MPI_Reduce_scatter(const void *sendbuf, void *recvbuf, const int recvcounts[], MPI_Datatype datatype, MPI_Op op, MPI_Comm comm)</pre>
Fortran	MPI_REDUCE_SCATTER(SENDBUF, RECVBUF, RECVCOUNTS, DATATYPE, OP, COMM, IERROR) <type> SENDBUF(*), RECVBUF(*) INTEGER RECVCOUNTS(*), DATATYPE, OP, COMM, IERROR</type>

MPI_Reduce_scatter

Perform reduction on vector elements and distribute segments of result vector across all tasks in communicator

```
recvent = 1;
MPI Reduce scatter(sendbuf, recvbuf, recvcount,
                    MPI INT, MPI SUM, MPI COMM WORLD);
task0
            task1
                       task2
                                  task3
                         1
                         2
                                    2
                                                - sendbuf (before)
  3
                         3
                                    3
                         4
                                    4
  4
             4

    recvbuf (after)

  4
             8
                        12
                                   16
```

MPI_Alltoall

Data movement operation. Each task in a group performs a scatter operation, sending a distinct message to all the tasks in the group in order by index.

С	<pre>int MPI_Alltoall(const void *sendbuf, int sendcount, MPI_Datatype sendtype, void *recvbuf, int recvcount, MPI_Datatype recvtype, MPI_Comm comm)</pre>
Fortran	MPI_ALLTOALL(SENDBUF, SENDCOUNT, SENDTYPE, RECVBUF, RECVCOUNT, RECVTYPE, COMM, IERROR) <type> SENDBUF(*), RECVBUF(*) INTEGER SENDCOUNT, SENDTYPE, RECVCOUNT, RECVTYPE INTEGER COMM, IERROR</type>

MPI_Alltoall

Scatter data from all tasks to all tasks in communicator

```
sendont = 1;
recvent = 1;
MPI Alltoall(sendbuf, sendont, MPI INT
           recvbuf, recvcnt, MPI_INT
           MPI COMM WORLD);
 task0
           task1
                     task2
                               task3
            5
                                13
                       9
  2
            6
                      10
                                14
                                          - sendbuf (before)
                      11
                                14
  4
            8
                      12
                                16
            2
                       3
                                 4
  1
  5
                       7
                                 8
            6
                                            recvbuf (after)
  9
            10
                      11
                                12
  13
            14
                      15
                                16
```

MPI_Scan

Performs a scan operation with respect to a reduction operation across a task group.

С	<pre>int MPI_Scan(const void *sendbuf, void *recvbuf, int count,</pre>
Fortran	MPI_SCAN(SENDBUF, RECVBUF, COUNT, DATATYPE, OP, COMM, IERROR) <pre></pre>

MPI_Scan

Computes the scan (partial reductions) across all tasks in communicator

Collective Communications Example in C

Perform a scatter operation on the rows of an array

```
#include "mpi.h"
#include <stdio.h>
#define SIZE 4
int main(int argc, char *argv[]) {
  int numtasks, rank, sendcount, recvcount, source;
  float sendbuf[SIZE][SIZE] = {
    \{1.0, 2.0, 3.0, 4.0\}.
    {5.0, 6.0, 7.0, 8.0},
    \{9.0, 10.0, 11.0, 12.0\},\
    {13.0, 14.0, 15.0, 16.0} };
  float recvbuf[SIZE];
  MPI_Init(&argc,&argv);
  MPI_Comm_size(MPI_COMM_WORLD, &numtasks);
  MPI_Comm_rank(MPI_COMM_WORLD, &rank);
```

```
if (numtasks == SIZE) {
  source = 1;
  sendcount = SIZE;
  recvcount = SIZE:
  MPI_Scatter(sendbuf, sendcount, MPI_FLOAT, recvbuf, recvcount,
              MPI_FLOAT, source, MPI_COMM_WORLD);
  printf("rank= %d Results: %f %f %f %f\n", rank, recvbuf[0],
         recvbuf[1], recvbuf[2], recvbuf[3]);
else
  printf("Must specify %d processors. Terminating.\n", SIZE);
MPI_Finalize();
```

Try it out

```
$ mpiicc scatter.c -o scatter.x

$ mpirun -n 4 ./scatter.x

rank= 2 Results: 9.000000 10.000000 11.000000 12.000000

rank= 3 Results: 13.000000 14.000000 15.000000 16.000000

rank= 1 Results: 5.000000 6.000000 7.000000 8.000000

rank= 0 Results: 1.000000 2.000000 3.000000 4.000000
```

Collective Communications Example in Fortran

Perform a scatter operation on the rows of an array

```
program scatter
   include 'mpif.h'
   integer SIZE
   parameter(SIZE=4)
   integer numtasks, rank, sendcount, recvcount, source, ierr
   real*4 sendbuf(SIZE, SIZE), recvbuf(SIZE)
Fortran stores this array in column major order, so the
scatter will actually scatter columns, not rows.
   data sendbuf /1.0, 2.0, 3.0, 4.0,
            5.0, 6.0, 7.0, 8.0,
  &
  &
            9.0, 10.0, 11.0, 12.0,
            13.0, 14.0, 15.0, 16.0 /
   call MPI_INIT(ierr)
   call MPI_COMM_RANK(MPI_COMM_WORLD, rank, ierr)
   call MPI_COMM_SIZE(MPI_COMM_WORLD, numtasks, ierr)
                                                                    121
```

```
if (rank .eq. 0) then
  prev = size - 1
 endif
 if (numtasks .eq. SIZE) then
   source = 1
   sendcount = SIZE
  recvcount = SIZE
   call MPI_SCATTER(sendbuf, sendcount, MPI_REAL, recvbuf,
&
            recvcount, MPI_REAL, source, MPI_COMM_WORLD, ierr)
   print *, 'rank= ',rank,' Results: ', recvbuf
 else
   print *, 'Must specify', SIZE,' processors. Terminating.'
 endif
 call MPI_FINALIZE(ierr)
 end
```

Try it out

```
$ mpiifort scatter.f -o scatter.x

$ mpirun -n 4 ./scatter.x

rank= 2 Results: 9.000000 10.000000 11.000000 12.000000

rank= 3 Results: 13.000000 14.000000 15.000000 16.000000

rank= 1 Results: 5.000000 6.000000 7.000000 8.000000

rank= 0 Results: 1.000000 2.000000 3.000000 4.000000
```

Derived Data Types

- As previously mentioned, MPI predefines its primitive data types.
- MPI also provides facilities for you to define your own data structures based upon sequences of the MPI primitive data types. Such user defined structures are called derived data types.
- Primitive data types are contiguous. Derived data types allow you to specify non-contiguous data in a convenient manner and to treat it as though it was contiguous.
- MPI provides several methods for constructing derived data types:
 - Contiguous
 - Vector
 - Indexed
 - Struct

Derived Data Type Routines

- MPI_Type_contiguous
- MPI_Type_vector
- MPI_Type_hvector
- MPI_Type_indexed
- MPI_Type_hindexed
- MPI_Type_struct
- MPI_Type_extent
- MPI_Type_commit
- MPI_Type_free

Open MPI documentation:

https://www.open-mpi.org/doc/v1.10/

MPI_Type_contiguous

The simplest constructor. Produces a new data type by making count copies of an existing data type.

С	<pre>int MPI_Type_contiguous(int count, MPI_Datatype oldtype,</pre>	
Fortran	MPI_TYPE_CONTIGUOUS(COUNT, OLDTYPE, NEWTYPE, IERROR) INTEGER COUNT, OLDTYPE, NEWTYPE, IERROR	

MPI_Type_vector & MPI_Type_hvector

Similar to contiguous, but allows for regular gaps (stride) in the displacements. **MPI_Type_hvector** is identical to **MPI_Type_vector** except that stride is specified in bytes.

С	<pre>int MPI_Type_vector(int count, int blocklength, int stride,</pre>
	<pre>int MPI_Type_hvector(int count, int blocklength, MPI_Aint stride,</pre>
Fortran	MPI_TYPE_VECTOR(COUNT, BLOCKLENGTH, STRIDE, OLDTYPE, NEWTYPE, IERROR) INTEGER COUNT, BLOCKLENGTH, STRIDE, OLDTYPE INTEGER NEWTYPE, IERROR
	MPI_TYPE_HVECTOR(COUNT, BLOCKLENGTH, STRIDE, OLDTYPE, NEWTYPE, IERROR) INTEGER COUNT, BLOCKLENGTH, STRIDE, OLDTYPE INTEGER NEWTYPE, IERROR

MPI_Type_indexed & MPI_Type_hindexed

An array of displacements of the input data type is provided as the map for the new data type. MPI_Type_hindexed is identical to MPI_Type_indexed except that offsets are specified in bytes.

C	<pre>int MPI_Type_indexed(int count, const int array_of_blocklengths[],</pre>
Fortran	MPI_TYPE_INDEXED(COUNT, ARRAY_OF_BLOCKLENGTHS,

MPI_Type_struct

The new data type is formed according to completely defined map of the component data types.

NOTE: This function is deprecated in MPI-2.0 and replaced by **MPI_Type_create_struct** in MPI-3.0

С	<pre>int MPI_Type_struct(int count, int *array_of_blocklengths,</pre>
Fortran	MPI_TYPE_STRUCT(COUNT, ARRAY_OF_BLOCKLENGTHS,

MPI_Type_extent

Returns the size in bytes of the specified data type. Useful for the MPI subroutines that require specification of offsets in bytes.

NOTE: This function is deprecated in MPI-2.0 and replaced by MPI_Type_get_extent in MPI-3.0

С	<pre>int MPI_Type_extent(MPI_Datatype datatype, MPI_Aint *extent)</pre>
Fortran	MPI_TYPE_EXTENT(DATATYPE, EXTENT, IERROR) INTEGER DATATYPE, EXTENT, IERROR

MPI_Type_commit

Commits new datatype to the system. Required for all user constructed (derived) datatypes.

С	<pre>int MPI_Type_commit(MPI_Datatype *datatype)</pre>
Fortran	MPI_TYPE_COMMIT(DATATYPE, IERROR) INTEGER DATATYPE, IERROR

MPI_Type_free

Deallocates the specified datatype object. Use of this routine is especially important to prevent memory exhaustion if many datatype objects are created, as in a loop.

С	<pre>int MPI_Type_free(MPI_Datatype *datatype)</pre>
Fortran	MPI_TYPE_FREE(DATATYPE, IERROR) INTEGER DATATYPE, IERROR

Contiguous Derived Data Type Example

MPI_Type_contiguous

Create a data type representing a column of an array and distribute a different row to all processes.

count = 4;
MPI_Type_contiguous(count, MPI_FLOAT, &rowtype);

1.0	2.0	3.0	4.0
5.0	6.0	7.0	8.0
9.0	10.0	11.0	12.0
13.0	14.0	15.0	16.0

a[4][4]

MPI_Send(&a[2][0], 1, rowtype, dest, tag, comm);

9.0	10.0	11.0	120

1 element of rowtype

Contiguous Derived Data Type Example in C

```
#include "mpi.h"
#include <stdio.h>
#define SIZE 4
int main(int argc, char *argv[]) {
  int numtasks, rank, source=0, dest, tag=1, i;
  float a[SIZE][SIZE] =
    \{1.0, 2.0, 3.0, 4.0,
     5.0. 6.0. 7.0. 8.0.
     9.0, 10.0, 11.0, 12.0,
     13.0, 14.0, 15.0, 16.0}:
  float b[SIZE];
  MPI_Status stat;
  MPI_Datatype rowtype;
  MPI_Init(&argc,&argv);
  MPI_Comm_rank(MPI_COMM_WORLD, &rank);
  MPI_Comm_size(MPI_COMM_WORLD, &numtasks);
```

```
MPI_Type_contiguous(SIZE, MPI_FLOAT, &rowtype);
MPI_Type_commit(&rowtype);
if (numtasks == SIZE) {
  if (rank == 0) {
    for (i=0: i<numtasks: i++)
      MPI_Send(&a[i][0], 1, rowtype, i, tag, MPI_COMM_WORLD);
 MPI_Recv(b, SIZE, MPI_FLOAT, source, tag, MPI_COMM_WORLD, &stat);
  printf("rank= %d b= %3.1f %3.1f %3.1f \n",
         rank, b[0], b[1], b[2], b[3]);
else
  printf("Must specify %d processors. Terminating.\n",SIZE);
MPI_Type_free(&rowtype);
MPI_Finalize();
```

Try it out

```
$ mpiicc contiguous.c -o contiguous.x

$ mpirun -n 4 ./contiguous.x
Fatal error in MPI_Send: Other MPI error, error stack:
MPI_Send(186): MPI_Send(buf=0x7fff5160f580, count=1,
dtype=USER<contig>, dest=0, tag=1, MPI_COMM_WORLD) failed
MPID_Send(53): DEADLOCK: attempting to send a message to the
local process without a prior matching receive
```

DEADLOCK in the example codes in LLNL MPI tutorial: https://computing.llnl.gov/tutorials/mpi/#Derived Data Types

Fixed Contiguous Derived Data Type Example in C

```
#include "mpi.h"
#include <stdio.h>
#define SIZE 4
int main(int argc, char *argv[]) {
  int numtasks, rank, source=0, dest, tag=1, i;
  float a[SIZE][SIZE] =
    \{1.0, 2.0, 3.0, 4.0,
     5.0. 6.0. 7.0. 8.0.
     9.0, 10.0, 11.0, 12.0,
     13.0, 14.0, 15.0, 16.0};
  float b[SIZE];
  MPI_Request req; /* fix */
  MPI_Status stat:
  MPI_Datatype rowtype;
  MPI_Init(&argc,&argv);
  MPI_Comm_rank(MPI_COMM_WORLD, &rank);
  MPI_Comm_size(MPI_COMM_WORLD, &numtasks);
```

```
MPI_Type_contiquous(SIZE, MPI_FLOAT, &rowtype);
MPI_Type_commit(&rowtype);
if (numtasks == SIZE) {
 /* fix */
  MPI_Irecv(b, SIZE, MPI_FLOAT, source, tag, MPI_COMM_WORLD, &req);
  if (rank == 0) {
    for (i=0; i<numtasks; i++)
      MPI_Send(&a[i][0], 1, rowtype, i, tag, MPI_COMM_WORLD);
  MPI_Wait(&req, &stat); /* fix */
  /* MPI_Recv(b, SIZE, MPI_FLOAT, source, tag, MPI_COMM_WORLD, &stat); */
  printf("rank= %d b= %3.1f %3.1f %3.1f %3.1f\n",
         rank, b[0], b[1], b[2], b[3]);
else
  printf("Must specify %d processors. Terminating.\n",SIZE);
MPI_Type_free(&rowtype);
MPI_Finalize();
                                                                      138
```

Try it out

```
$ mpiicc contiguous_fixed.c -o contiguous_fixed.x

$ mpirun -n 4 ./contiguous_fixed.x
rank= 0 b= 1.0 2.0 3.0 4.0
rank= 1 b= 5.0 6.0 7.0 8.0
rank= 2 b= 9.0 10.0 11.0 12.0
rank= 3 b= 13.0 14.0 15.0 16.0
```

Contiguous Derived Data Type Example in Fortran

```
program contiquous
   include 'mpif.h'
   integer SIZE
   parameter(SIZE=4)
   integer numtasks, rank, source, dest, tag, i, ierr
   real*4 a(0:SIZE-1,0:SIZE-1), b(0:SIZE-1)
   integer stat(MPI_STATUS_SIZE), columntype
Fortran stores this array in column major order
   data a /1.0, 2.0, 3.0, 4.0,
            5.0, 6.0, 7.0, 8.0,
  &
            9.0, 10.0, 11.0, 12.0,
  &
            13.0, 14.0, 15.0, 16.0 /
   call MPI_INIT(ierr)
   call MPI_COMM_RANK(MPI_COMM_WORLD, rank, ierr)
   call MPI_COMM_SIZE(MPI_COMM_WORLD, numtasks, ierr)
```

```
call MPI_TYPE_CONTIGUOUS(SIZE, MPI_REAL, columntype, ierr)
 call MPI_TYPE_COMMIT(columntype, ierr)
 tag = 1
 if (numtasks .eq. SIZE) then
   if (rank .eq. 0) then
     do i=0, numtasks-1
       call MPI_SEND(a(0,i), 1, columntype, i, tag,
&
                     MPI_COMM_WORLD, ierr)
     enddo
   endif
   source = 0
   call MPI_RECV(b, SIZE, MPI_REAL, source, tag,
                 MPI_COMM_WORLD, stat, ierr)
&
   print *, 'rank= ',rank,' b= ',b
 else
   print *, 'Must specify', SIZE,' processors. Terminating.'
 endif
 call MPI_TYPE_FREE(columntype, ierr)
 call MPI_FINALIZE(ierr)
 end
```

Try it out

```
$ mpiifort contiguous.f -o contiguous.x
$ mpirun -n 4 ./contiguous.x
Fatal error in MPI_Send: Other MPI error, error stack:
MPI_Send(186): MPI_Send(buf=0x69eee0, count=1,
dtype=USER<contig>, dest=0, tag=1, MPI_COMM_WORLD) failed
MPID_Send(53): DEADLOCK: attempting to send a message to the
local process without a prior matching receive
```

DEADLOCK in the example codes in LLNL MPI tutorial: https://computing.llnl.gov/tutorials/mpi/#Derived Data Types

Fixed Contiguous Derived Data Type Example in Fortran

```
program contiquous
   include 'mpif.h'
   integer SIZE
   parameter(SIZE=4)
   integer numtasks, rank, source, dest, tag, i, ierr
   real*4 a(0:SIZE-1,0:SIZE-1), b(0:SIZE-1)
   integer req, stat(MPI_STATUS_SIZE), columntype
Fortran stores this array in column major order
   data a /1.0, 2.0, 3.0, 4.0,
            5.0, 6.0, 7.0, 8.0,
  &
            9.0, 10.0, 11.0, 12.0,
  &
            13.0, 14.0, 15.0, 16.0 /
   call MPI_INIT(ierr)
   call MPI_COMM_RANK(MPI_COMM_WORLD, rank, ierr)
   call MPI_COMM_SIZE(MPI_COMM_WORLD, numtasks, ierr)
```

```
call MPI_TYPE_CONTIGUOUS(SIZE, MPI_REAL, columntype, ierr)
 call MPI_TYPE_COMMIT(columntype, ierr)
 tag = 1
 if (numtasks .eq. SIZE) then
   source = 0
   call MPI_IRECV(b, SIZE, MPI_REAL, source, tag,
&
                 MPI_COMM_WORLD, req, ierr)
   if (rank .eq. 0) then
     do i=0, numtasks-1
       call MPI_SEND(a(0,i), 1, columntype, i, tag,
&
                     MPI_COMM_WORLD. ierr)
     enddo
   endif
   call MPI_WAIT(req, stat, ierr)
   print *, 'rank= ',rank,' b= ',b
 else
   print *, 'Must specify', SIZE,' processors. Terminating.'
 endif
 call MPI_TYPE_FREE(columntype, ierr)
 call MPI_FINALIZE(ierr)
 end
                                                                   144
```

```
$ mpiifort contiguous_fixed.f -o contiguous_fixed.x

$ mpirun -n 4 ./contiguous_fixed.x
  rank= 0 b= 1.000000 2.000000 3.000000 4.000000
  rank= 1 b= 5.000000 6.000000 7.000000 8.000000
  rank= 2 b= 9.000000 10.00000 11.00000 12.00000
  rank= 3 b= 13.00000 14.00000 15.00000 16.00000
```

Vector Derived Data Type Example

Create a data type representing a column of an array and distribute different columns to all processes.

MPI_Type_vector

1.0	2.0	3.0	4.0
5.0	6.0	7.0	8.0
9.0	10.0	11.0	12.0
13.0	14.0	15.0	16.0

a[4][4]

MPI_Send(&a[0][1], 1, columntype, dest, tag, comm);

2.0 6.0 10.0 14.0

1 element of columnty pe

Vector Derived Data Type Example in C

```
#include "mpi.h"
#include <stdio.h>
#define SIZE 4
int main(int argc, char *argv[]) {
  int numtasks, rank, source=0, dest, tag=1, i;
  float a[SIZE][SIZE] =
    \{1.0, 2.0, 3.0, 4.0,
     5.0, 6.0, 7.0, 8.0,
     9.0, 10.0, 11.0, 12.0,
     13.0, 14.0, 15.0, 16.0};
  float b[SIZE];
  MPI_Request req;
  MPI_Status stat:
  MPI_Datatype columntype;
  MPI_Init(&argc,&argv);
  MPI_Comm_rank(MPI_COMM_WORLD, &rank);
  MPI_Comm_size(MPI_COMM_WORLD, &numtasks);
                                                                         147
```

```
MPI_Type_vector(SIZE, 1, SIZE, MPI_FLOAT, &columntype);
MPI_Type_commit(&columntype);
if (numtasks == SIZE) {
  MPI_Irecv(b, SIZE, MPI_FLOAT, source, tag, MPI_COMM_WORLD, &req);
  if (rank == 0) {
    for (i=0; i<numtasks; i++)
      MPI_Send(&a[0][i], 1, columntype, i, tag, MPI_COMM_WORLD);
  MPI_Wait(&req, &stat);
  printf("rank= %d b= \%3.1f \%3.1f \%3.1f \%3.1f n",
         rank, b[0], b[1], b[2], b[3]);
else
  printf("Must specify %d processors. Terminating.\n",SIZE);
MPI_Type_free(&columntype);
MPI_Finalize();
```

```
$ mpiicc vector.c -o vector.x
$ mpirun -n 4 ./vector.x
rank= 1 b= 2.0 6.0 10.0 14.0
rank= 2 b= 3.0 7.0 11.0 15.0
rank= 3 b= 4.0 8.0 12.0 16.0
rank= 0 b= 1.0 5.0 9.0 13.0
```

Note: same **DEADLOCK** bug in all "Derived Data Types" examples in the LLNL MPI tutorial:

https://computing.llnl.gov/tutorials/mpi/#Derived Data Types

I've applied the simple **IRECV** fix in this presentation!

Vector Derived Data Type Example in Fortran

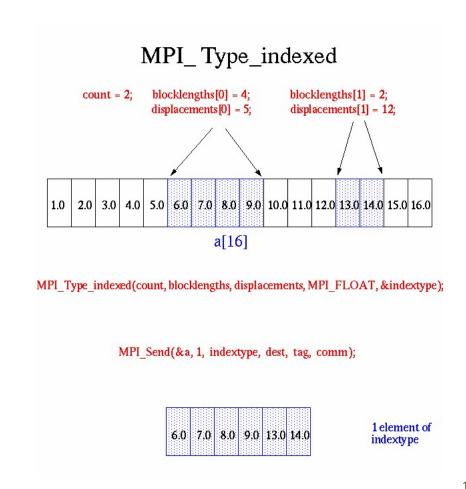
```
program vector
   include 'mpif.h'
   integer SIZE
   parameter(SIZE=4)
   integer numtasks, rank, source, dest, tag, i, ierr
   real*4 a(0:SIZE-1,0:SIZE-1), b(0:SIZE-1)
   integer req, stat(MPI_STATUS_SIZE), rowtype
Fortran stores this array in column major order
   data a /1.0, 2.0, 3.0, 4.0,
            5.0, 6.0, 7.0, 8.0,
  &
            9.0, 10.0, 11.0, 12.0,
  &
            13.0, 14.0, 15.0, 16.0 /
   call MPI_INIT(ierr)
   call MPI_COMM_RANK(MPI_COMM_WORLD, rank, ierr)
   call MPI_COMM_SIZE(MPI_COMM_WORLD, numtasks, ierr)
```

```
call MPI_TYPE_VECTOR(SIZE, 1, SIZE, MPI_REAL, rowtype, ierr)
 call MPI_TYPE_COMMIT(rowtype, ierr)
 tag = 1
 if (numtasks .eq. SIZE) then
   if (rank .eq. 0) then
     do i=0, numtasks-1
       call MPI_SEND(a(i,0), 1, rowtype, i, tag,
&
                     MPI_COMM_WORLD, ierr)
     enddo
   endif
   source = 0
   call MPI_RECV(b, SIZE, MPI_REAL, source, tag,
                 MPI_COMM_WORLD, stat, ierr)
&
   print *, 'rank= ',rank,' b= ',b
 else
   print *, 'Must specify',SIZE,' processors. Terminating.'
 endif
 call MPI_TYPE_FREE(rowtype, ierr)
 call MPI_FINALIZE(ierr)
 end
```

```
$ mpiifort vector.f -o vector.x
$ mpirun -n 4 ./vector.x
rank= 1 b= 2.000000 6.000000 10.00000 14.00000
rank= 2 b= 3.000000 7.000000 11.00000 15.00000
rank= 3 b= 4.000000 8.000000 12.00000 16.00000
rank= 0 b= 1.000000 5.000000 9.000000 13.00000
```

Indexed Derived Data Type Example

Create a datatype by extracting variable portions of an array and distribute to all tasks.



Indexed Derived Data Type Example in C

```
#include "mpi.h"
#include <stdio.h>
#define NELEMENTS 6
int main(int argc, char *argv[]) {
  int numtasks, rank, source=0, dest, tag=1, i;
  int blocklengths[2], displacements[2];
  float a[16] =
    \{1.0, 2.0, 3.0, 4.0, 5.0, 6.0, 7.0, 8.0,
     9.0, 10.0, 11.0, 12.0, 13.0, 14.0, 15.0, 16.0};
  float b[NELEMENTS]:
  MPI_Request req;
  MPI_Status stat:
  MPI_Datatype indextype;
  MPI_Init(&argc,&argv);
  MPI_Comm_rank(MPI_COMM_WORLD, &rank);
  MPI_Comm_size(MPI_COMM_WORLD, &numtasks);
```

```
blocklengths[0] = 4;
blocklengths[1] = 2;
displacements[0] = 5;
displacements[1] = 12;
MPI_Type_indexed(2, blocklengths, displacements, MPI_FLOAT, &indextype);
MPI_Type_commit(&indextype);
MPI_Irecv(b, NELEMENTS, MPI_FLOAT, source, tag, MPI_COMM_WORLD, &req);
if (rank == 0) {
  for (i=0; i<numtasks; i++)
    MPI_Send(a, 1, indextype, i, tag, MPI_COMM_WORLD);
MPI_Wait(&req, &stat);
printf("rank= %d b= %3.1f %3.1f %3.1f %3.1f %3.1f \%3.1f\n",
       rank, b[0], b[1], b[2], b[3], b[4], b[5]);
MPI_Type_free(&indextype);
MPI_Finalize();
                                                                      155
```

```
$ mpiicc indexed.c -o indexed.x
$ mpirun -n 4 ./indexed.x
rank= 0 b= 6.0 7.0 8.0 9.0 13.0 14.0
rank= 1 b= 6.0 7.0 8.0 9.0 13.0 14.0
rank= 2 b= 6.0 7.0 8.0 9.0 13.0 14.0
rank= 3 b= 6.0 7.0 8.0 9.0 13.0 14.0
```

Indexed Derived Data Type Example in Fortran

```
program indexed
include 'mpif.h'
integer NELEMENTS
parameter(NELEMENTS=6)
integer numtasks, rank, source, dest, tag, i, ierr
integer blocklengths(0:1), displacements(0:1)
real*4 a(0:15), b(0:NELEMENTS-1)
integer req, stat(MPI_STATUS_SIZE), indextype
data a /1.0, 2.0, 3.0, 4.0, 5.0, 6.0, 7.0, 8.0,
         9.0, 10.0, 11.0, 12.0, 13.0, 14.0, 15.0, 16.0 /
call MPI_INIT(ierr)
call MPI_COMM_RANK(MPI_COMM_WORLD, rank, ierr)
call MPI_COMM_SIZE(MPI_COMM_WORLD, numtasks, ierr)
```

```
blocklengths(0) = 4
 blocklengths(1) = 2
 displacements(0) = 5
 displacements(1) = 12
 call MPI_TYPE_INDEXED(2, blocklengths, displacements, MPI_REAL,
                       indextype, ierr)
 call MPI_TYPE_COMMIT(indextype, ierr)
 tag = 1
 source = 0
 call MPI_IRECV(b, NELEMENTS, MPI_REAL, source, tag,
&
                MPI_COMM_WORLD, req, ierr)
 if (rank .eq. 0) then
   do i=0, numtasks-1
     call MPI_SEND(a, 1, indextype, i, tag, MPI_COMM_WORLD, ierr)
   enddo
 endif
 call MPI_WAIT(req, stat, ierr)
 print *, 'rank= ',rank,' b= ',b
 call MPI_TYPE_FREE(indextype, ierr)
 call MPI_FINALIZE(ierr)
 end
                                                                   158
```

```
$ mpiifort indexed.f -o indexed.x
$ mpirun -n 4 ./indexed.x
rank= 1 b= 6.000000 7.000000 8.000000 9.000000 13.00000 14.00000
rank= 2 b= 6.000000 7.000000 8.000000 9.000000 13.00000 14.00000
rank= 3 b= 6.000000 7.000000 8.000000 9.000000 13.00000 14.00000
rank= 0 b= 6.000000 7.000000 8.000000 9.000000 13.00000 14.00000
```

Struct Derived Data Type Example

Create a data type that represents a particle and distribute an array of such particles to all processes.

MPI_ Type_struct

typedef struct { float x, y, z, velocity; int n, type; } Particle;
Particle particles[NELEM];

MPI_Type_extent(MPI_FLOAT, &extent);

count = 2; oldtypes[0] = MPI_FLOAT; oldtypes[1] = MPI_INT offsets[0] = 0; offsets[1] = 4 * extent; blockcounts[1] = 2;

particles[NELEM]

MPI_Type_struct(count, blockcounts, offsets, oldtypes, &particletype);

MPI_Send(particles, NELEM, particletype, dest, tag, comm);

Sends entire (NELEM) array of particles, each particle being comprised four floats and two integers.

Struct Derived Data Type Example in C

```
#include "mpi.h"
#include <stdio.h>
#define NELEM 25
int main(int argc, char *argv[]) {
  int numtasks, rank, source=0, dest, tag=1, i;
  typedef struct {
    float x, y, z;
    float velocity;
    int n, type;
     Particle:
  Particle p[NELEM], particles[NELEM];
  MPI_Datatype particletype, oldtypes[2];
  int
              blockcounts[2];
  /* MPI_Aint type used to be consistent with syntax of
    MPI_Type_extent routine */
  MPI_Aint offsets[2], extent;
  MPI_Request req;
  MPI_Status stat;
```

```
MPI_Init(&argc,&argv);
MPI_Comm_rank(MPI_COMM_WORLD, &rank);
MPI_Comm_size(MPI_COMM_WORLD, &numtasks);
/* Setup description of the 4 MPI_FLOAT fields x, y, z, velocity */
offsets[0] = 0;
oldtypes[0] = MPI_FLOAT;
blockcounts[0] = 4;
/* Setup description of the 2 MPI_INT fields n, type
   Need to first figure offset by getting size of MPI_FLOAT */
MPI_Type_extent(MPI_FLOAT, &extent);
offsets[1] = 4 * extent;
oldtypes[1] = MPI_INT;
blockcounts[1] = 2;
/* Now define structured type and commit it */
MPI_Type_struct(2, blockcounts, offsets, oldtypes, &particletype);
MPI_Type_commit(&particletype);
```

```
MPI_Irecv(p, NELEM, particletype, source, tag, MPI_COMM_WORLD, &req);
/* Initialize the particle array and then send it to each task */
if (rank == 0) {
  for (i=0; i<NELEM; i++) {
    particles[i].x = i * 1.0;
    particles[i].y = i * -1.0;
    particles[i].z = i * 1.0;
    particles[i].velocity = 0.25:
    particles[i].n = i;
    particles[i].type = i % 2;
  for (i=0; i<numtasks; i++)
    MPI_Send(particles, NELEM, particletype, i, tag, MPI_COMM_WORLD);
MPI_Wait(&req, &stat);
/* Print a sample of what was received */
printf("rank= %d  %3.2f %3.2f %3.2f %3.2f %d %d\n", rank, p[3].x,
       p[3].y, p[3].z, p[3].velocity, p[3].n, p[3].type);
MPI_Type_free(&particletype);
MPI_Finalize();
                                                                      163
```

```
$ mpiicc struct.c -o struct.x
$ mpirun -n 4 ./struct.x
rank= 0    3.00 -3.00 3.00 0.25 3 1
rank= 1    3.00 -3.00 3.00 0.25 3 1
rank= 2    3.00 -3.00 3.00 0.25 3 1
rank= 3    3.00 -3.00 3.00 0.25 3 1
```

Struct Derived Data Type Example in Fortran

```
program struct
include 'mpif.h'
integer NELEM
parameter(NELEM=25)
integer numtasks, rank, source, dest, tag, i, ierr
integer req, stat(MPI_STATUS_SIZE)
type Particle
  sequence
  real*4 x, y, z, velocity
  integer n, type
end type Particle
type (Particle) p(NELEM), particles(NELEM)
integer particletype, oldtypes(0:1), blockcounts(0:1),
        offsets(0:1), extent
```

```
call MPI_INIT(ierr)
   call MPI_COMM_RANK(MPI_COMM_WORLD, rank, ierr)
   call MPI_COMM_SIZE(MPI_COMM_WORLD, numtasks, ierr)
Setup description of the 4 MPI_REAL fields x, y, z, velocity
   offsets(0) = 0
   oldtypes(0) = MPI_REAL
   blockcounts(0) = 4
Setup description of the 2 MPI_INTEGER fields n, type
Need to first figure offset by getting size of MPI_REAL
   call MPI_TYPE_EXTENT(MPI_REAL, extent, ierr)
   offsets(1) = 4 * extent
   oldtypes(1) = MPI_INTEGER
   blockcounts(1) = 2
Now define structured type and commit it
   call MPI_TYPE_STRUCT(2, blockcounts, offsets, oldtypes,
                         particletype, ierr)
   call MPI_TYPE_COMMIT(particletype, ierr)
```

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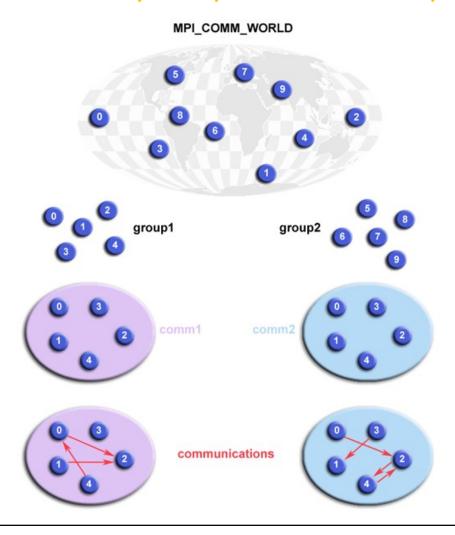
```
Initialize the particle array and then send it to each task
   tag = 1
   source = 0
   call MPI_IRECV(p, NELEM, particletype, source, tag,
  &
                  MPI_COMM_WORLD, req, ierr)
   if (rank .eq. 0) then
     do i=0, NELEM-1
       particles(i) = Particle (1.0*i, -1.0*i, 1.0*i,
  &
                                 0.25, i, mod(i,2))
     enddo
     do i=0, numtasks-1
       call MPI_SEND(particles, NELEM, particletype, i, tag,
  &
                     MPI_COMM_WORLD, ierr)
     enddo
   endif
   call MPI_WAIT(reg, stat, ierr)
   print *, 'rank= ',rank,' p(3)= ',p(3)
   call MPI_TYPE_FREE(particletype, ierr)
   call MPI_FINALIZE(ierr)
   end
```

```
$ mpiifort struct.f -o struct.x
$ mpirun -n 4 ./struct.x
rank= 0 p(3)= 3.000000 -3.000000 3.000000 0.2500000 3 1
rank= 1 p(3)= 3.000000 -3.000000 3.000000 0.2500000 3 1
rank= 2 p(3)= 3.000000 -3.000000 3.000000 0.2500000 3 1
rank= 3 p(3)= 3.000000 -3.000000 3.000000 0.2500000 3 1
```

Groups and Communicators

- A group is an ordered set of processes. Each process in a group is associated with a unique integer rank. Rank values start at zero and go to N-1, where N is the number of processes in the group. In MPI, a group is represented within system memory as an object. It is accessible to the programmer only by a "handle". A group is always associated with a communicator object.
- A communicator encompasses a group of processes that may communicate with each other. All MPI messages must specify a communicator. In the simplest sense, the communicator is an extra "tag" that must be included with MPI calls. Like groups, communicators are represented within system memory as objects and are accessible to the programmer only by "handles". For example, the handle for the communicator that comprises all tasks is MPI_COMM_WORLD.
- From the programmer's perspective, a group and a communicator are one. The group routines are primarily used to specify which processes should be used to construct a communicator.

Primary Purposes of Group and Communicator Objects



- 1. Allow you to organize tasks, based upon function, into task groups.
- 2. Enable Collective Communications operations across a subset of related tasks.
- Provide basis for implementing user defined virtual topologies
- 4. Provide for safe communications

Programming Considerations and Restrictions

- Groups/communicators are dynamic they can be created and destroyed during program execution.
- Processes may be in more than one group/communicator. They will have a unique rank within each group/communicator.
- MPI provides over 40 routines related to groups, communicators, and virtual topologies.
- Typical usage:
 - 1. Extract handle of global group from MPI_COMM_WORLD using MPI_Comm_group
 - 2. Form new group as a subset of global group using MPI_Group_incl
 - 3. Create new communicator for new group using MPI_Comm_create
 - 4. Determine new rank in new communicator using MPI_Comm_rank
 - 5. Conduct communications using any MPI message passing routine
 - 6. When finished, free up new communicator and group (optional) using MPI_Comm_free and MPI_Group_free

Group and Communicator Routines Example in C

```
#include "mpi.h"
#include <stdio.h>
#define NPROCS 8
int main(int argc, char *argv[]) {
  int rank, new_rank, sendbuf, recvbuf, numtasks;
  int ranks1[4]=\{0,1,2,3\}, ranks2[4]=\{4,5,6,7\};
  MPI_Group orig_group, new_group;
  MPI Comm
             new_comm;
  MPI_Init(&argc,&argv);
  MPI_Comm_rank(MPI_COMM_WORLD, &rank);
  MPI_Comm_size(MPI_COMM_WORLD, &numtasks);
  if (numtasks != NPROCS) {
    printf("Must specify MP_PROCS= %d. Terminating.\n", NPROCS);
    MPI_Finalize();
    exit(0);
```

```
sendbuf = rank;
/* Extract the original group handle */
MPI_Comm_group(MPI_COMM_WORLD, &orig_group);
/* Divide tasks into two distinct groups based upon rank */
if (rank < NPROCS/2) {
  MPI_Group_incl(orig_group, NPROCS/2, ranks1, &new_group);
else {
  MPI_Group_incl(orig_group, NPROCS/2, ranks2, &new_group);
/* Create new communicator and then perform collective communications */
MPI_Comm_create(MPI_COMM_WORLD, new_group, &new_comm);
MPI_Allreduce(&sendbuf, &recvbuf, 1, MPI_INT, MPI_SUM, new_comm);
MPI_Group_rank(new_group, &new_rank);
printf("rank= %d newrank= %d recvbuf= %d\n",rank, new_rank, recvbuf);
MPI_Finalize();
```

```
$ mpiicc group.c -o group.x
$ mpirun -n 8 ./group.x
rank= 1 newrank= 1 recvbuf= 6
rank= 2 newrank= 2 recvbuf= 6
rank= 3 newrank= 3 recvbuf= 6
rank= 4 newrank= 0 recvbuf= 22
rank= 6 newrank= 2 recvbuf= 22
rank= 0 newrank= 0 recvbuf= 6
rank= 5 newrank= 1 recvbuf= 22
rank= 7 newrank= 3 recvbuf= 22
```

Group and Communicator Routines Example in Fortran

```
program group
include 'mpif.h'
integer NPROCS
parameter(NPROCS=8)
integer rank, new_rank, sendbuf, recvbuf, numtasks
integer ranks1(4), ranks2(4), ierr
integer orig_group, new_group, new_comm
data ranks1 /0, 1, 2, 3/, ranks2 /4, 5, 6, 7/
call MPI_INIT(ierr)
call MPI_COMM_RANK(MPI_COMM_WORLD, rank, ierr)
call MPI_COMM_SIZE(MPI_COMM_WORLD, numtasks, ierr)
if (numtasks .ne. NPROCS) then
  print *, 'Must specify NPROCS= ', NPROCS, ' Terminating.'
  call MPI_FINALIZE(ierr)
  stop
endif
```

```
sendbuf = rank
Extract the original group handle
   call MPI_COMM_GROUP(MPI_COMM_WORLD, orig_group, ierr)
Divide tasks into two distinct groups based upon rank
   if (rank .lt. NPROCS/2) then
     call MPI_GROUP_INCL(orig_group, NPROCS/2, ranks1,
  &
                          new_group, ierr)
   else
     call MPI_GROUP_INCL(orig_group, NPROCS/2, ranks2,
  &
                          new_group, ierr)
   endif
   call MPI_COMM_CREATE(MPI_COMM_WORLD, new_group,
  &
                         new_comm, ierr)
   call MPI_ALLREDUCE(sendbuf, recvbuf, 1, MPI_INTEGER,
                      MPI_SUM, new_comm, ierr)
   call MPI_GROUP_RANK(new_group, new_rank, ierr)
   print *, 'rank= ', rank, ' newrank= ', new_rank, ' recvbuf= ',
            recybuf
   call MPI_FINALIZE(ierr)
   end
```

```
$ mpiifort group.f -o group.x
$ mpirun -n 8 ./group.x
rank= 0 newrank= 0 recvbuf= 6
rank= 1 newrank= 1 recvbuf= 6
rank= 2 newrank= 2 recvbuf= 6
rank= 3 newrank= 3 recvbuf= 6
rank= 4 newrank= 0 recvbuf= 22
rank= 5 newrank= 1 recvbuf= 22
rank= 6 newrank= 2 recvbuf= 22
rank= 7 newrank= 3 recvbuf= 22
```

Virtual Topologies

- A virtual topology describes a mapping/ordering of MPI processes into a geometric "shape".
- The two main types of topologies supported by MPI are Cartesian (grid) and Graph.
- MPI topologies are virtual there may be no relation between the physical structure of the parallel machine and the process topology.
- Virtual topologies are built upon MPI communicators and groups.
- Must be "programmed" by the application developer.

Why use Virtual Topologies?

Convenience

- Virtual topologies may be useful for applications with specific communication patterns - patterns that match an MPI topology structure.
- For example, a Cartesian topology might prove convenient for an application that requires 4-way nearest neighbor communications for grid based data.

Communication Efficiency

- Some hardware architectures may impose penalties for communications between successively distant "nodes".
- A particular implementation may optimize process mapping based upon the physical characteristics of a given parallel machine.
- The mapping of processes into an MPI virtual topology is dependent upon the MPI implementation, and may be totally ignored.

Virtual Topology Example

Mapping of processes into a Cartesian virtual topology

0	1	2	3
(0,0)	(0,1)	(0,2)	(0,3)
4	5	6	7
(1,0)	(1,1)	(1,2)	(1,3)
8	9	10	11
(2,0)	(2,1)	(2,2)	(2,3)
(3,0)	13	14	15
	(3,1)	(3,2)	(3,3)

Cartesian Virtual Topology Example in C

```
#include "mpi.h"
#include <stdio.h>
#define SIZE 16
#define UP
#define DOWN 1
#define LFFT 2
#define RIGHT 3
int main(int argc, char *argv[]) {
  int numtasks, rank, source, dest, outbuf, i, tag=1,
    inbuf[4]={MPI_PROC_NULL,MPI_PROC_NULL,MPI_PROC_NULL,MPI_PROC_NULL,},
    nbrs[4], dims[2]={4,4},
    periods[2]={0,0}, reorder=0, coords[2];
  MPI_Request reqs[8];
  MPI_Status stats[8];
  MPI_Comm cartcomm;
  MPI_Init(&argc,&argv);
  MPI_Comm_size(MPI_COMM_WORLD, &numtasks);
```

```
if (numtasks == SIZE) {
   MPI_Cart_create(MPI_COMM_WORLD, 2, dims, periods, reorder, &cartcomm);
   MPI_Comm_rank(cartcomm, &rank);
   MPI_Cart_coords(cartcomm, rank, 2, coords);
   MPI_Cart_shift(cartcomm, 0, 1, &nbrs[UP], &nbrs[DOWN]);
    MPI_Cart_shift(cartcomm, 1, 1, &nbrs[LEFT], &nbrs[RIGHT]);
    printf("rank= %d coords= %d %d neighbors(u,d,1,r)= %d %d %d %d\n",
           rank, coords[0], coords[1], nbrs[UP], nbrs[DOWN], nbrs[LEFT],
           nbrs[RIGHT]):
    outbuf = rank:
    for (i=0; i<4; i++) {
     dest = nbrs[i]:
     source = nbrs[i];
     MPI_Isend(&outbuf, 1, MPI_INT, dest, tag, MPI_COMM_WORLD, &reqs[i]);
     MPI_Irecv(&inbuf[i], 1, MPI_INT, source, tag, MPI_COMM_WORLD, &regs[i+4]);
   MPI_Waitall(8, reqs, stats);
    printf("rank= %d
                                      inbuf(u,d,l,r) = %d %d %d %d n",
         rank,inbuf[UP], inbuf[DOWN], inbuf[LEFT], inbuf[RIGHT]); }
 else
    printf("Must specify %d processors. Terminating.\n",SIZE);
 MPI_Finalize();
}
                                                                               182
```

Cartesian Virtual Topology Example in Fortran

```
program cartesian
 include 'mpif.h'
 integer SIZE, UP, DOWN, LEFT, RIGHT
 parameter(SIZE=16)
 parameter(UP=1)
 parameter(DOWN=2)
 parameter(LEFT=3)
 parameter(RIGHT=4)
 integer numtasks, rank, source, dest, outbuf, i, tag, ierr,
         inbuf(4), nbrs(4), dims(2), coords(2),
&
         stats(MPI_STATUS_SIZE, 8), reqs(8), cartcomm,
&
         periods(2), reorder
 data inbuf /MPI_PROC_NULL, MPI_PROC_NULL, MPI_PROC_NULL,
      MPI_PROC_NULL/, dims /4,4/, tag /1/,
&
      periods /0,0/, reorder /0/
&
 call MPI_INIT(ierr)
 call MPI_COMM_SIZE(MPI_COMM_WORLD, numtasks, ierr)
                                                                   184
```

```
if (numtasks .eq. SIZE) then
   call MPI_CART_CREATE(MPI_COMM_WORLD, 2, dims, periods, reorder,
&
                        cartcomm, ierr)
   call MPI_COMM_RANK(cartcomm, rank, ierr)
   call MPI_CART_COORDS(cartcomm, rank, 2, coords, ierr)
   call MPI_CART_SHIFT(cartcomm, 0, 1, nbrs(UP), nbrs(DOWN), ierr)
   call MPI_CART_SHIFT(cartcomm, 1, 1, nbrs(LEFT), nbrs(RIGHT),
&
                       ierr)
  write(*,*) rank, coords(1), coords(2), nbrs(UP), nbrs(DOWN),
              nbrs(LEFT), nbrs(RIGHT)
&
   outbuf = rank
   do i=1.4
     dest = nbrs(i)
     source = nbrs(i)
     call MPI_ISEND(outbuf, 1, MPI_INTEGER, dest, tag,
&
                    MPI_COMM_WORLD, reqs(i), ierr)
     call MPI_IRECV(inbuf(i), 1, MPI_INTEGER, source, tag,
                    MPI_COMM_WORLD, regs(i+4), ierr)
&
   enddo
   call MPI_WAITALL(8, regs, stats, ierr)
   write(*,*) rank, inbuf
 else
   print *, 'Must specify', SIZE,' processors. Terminating.'
 endif
 call MPI FINALIZE(ierr)
 end
```

```
$ mpiifort cartesian.f -o cartesian.x
$ mpirun -n 16 ./cartesian.x
0     0     0     -1     4     -1     1
1     0     1     -1     5     0     2
2     0     2     -1     6     1     3
8     2     0     4     12     -1     9
...
1     -1     5     0     2
8     4     12     -1     9
2     -1     6     1     3
0     -1     4     -1     1
...
```

A Brief Word on MPI-2

- Intentionally, the MPI-1 specification did not address several "difficult" issues. For reasons of expediency, these issues were deferred to a second specification, called MPI-2 in 1997.
- MPI-2 was a major revision to MPI-1 adding new functionality and corrections.
- Key areas of new functionality in MPI-2:
 - O **Dynamic Processes** extensions that remove the static process model of MPI. Provides routines to create new processes after job startup.
 - One-Sided Communications provides routines for one directional communications.
 Include shared memory operations (put/get) and remote accumulate operations.
 - Extended Collective Operations allows for the application of collective operations to inter-communicators
 - External Interfaces defines routines that allow developers to layer on top of MPI, such as for debuggers and profilers.
 - Additional Language Bindings describes C++ bindings and discusses Fortran-90 issues.
 - Parallel I/O describes MPI support for parallel I/O.

A Brief Word on MPI-3

- The MPI-3 standard was adopted in 2012, and contains significant extensions to MPI-1 and MPI-2 functionality including:
 - Non-blocking Collective Operations permits tasks in a collective to perform operations without blocking, possibly offering performance improvements.
 - New One-sided Communication Operations to better handle different memory models.
 - Neighborhood Collectives Extends the distributed graph and Cartesian process topologies with additional communication power.
 - Fortran 2008 Bindings expanded from Fortran90 bindings
 - MPIT Tool Interface This new tool interface allows the MPI implementation to expose certain internal variables, counters, and other states to the user (most likely performance tools).
 - Matched Probe Fixes an old bug in MPI-2 where one could not probe for messages in a multi-threaded environment.