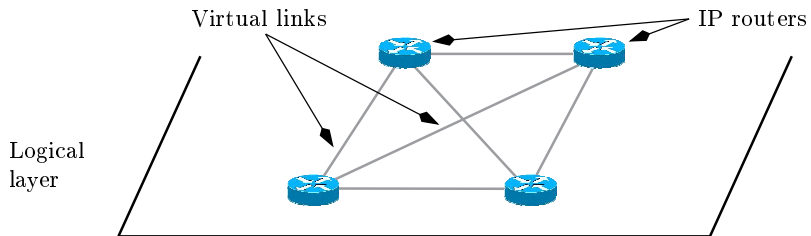


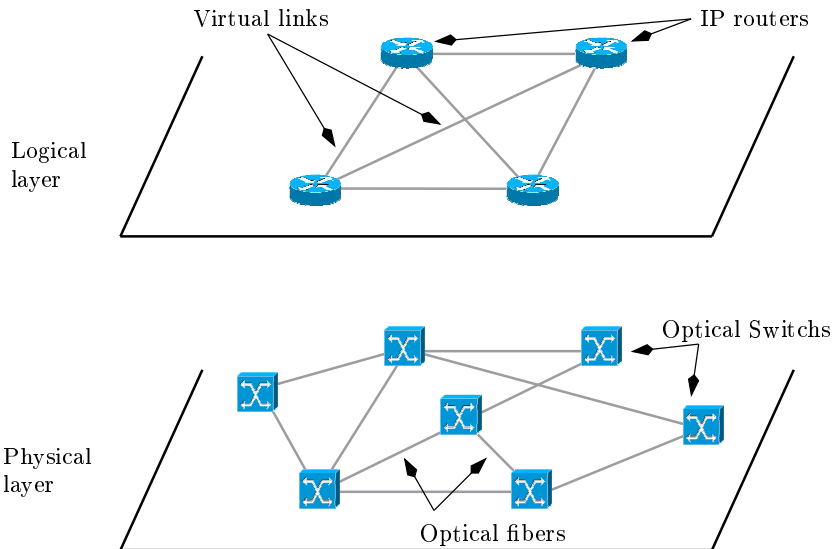
# Multilayer Survivable Optical Network Design

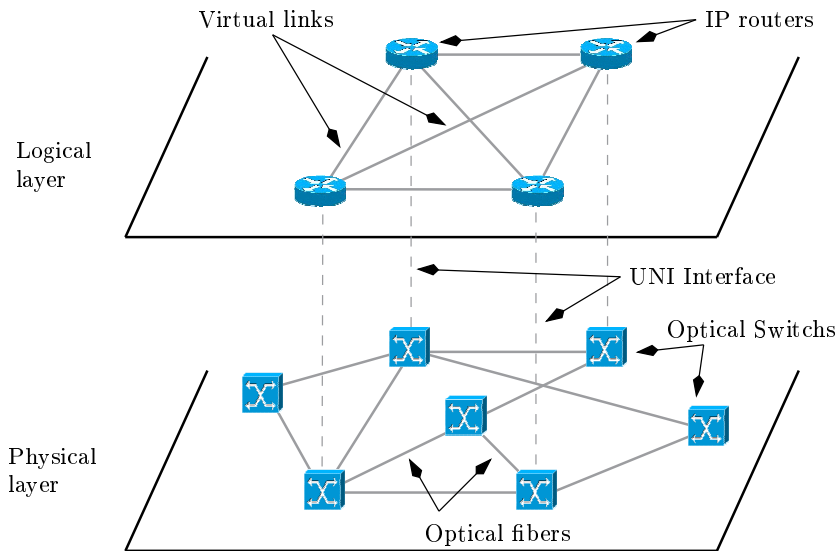
**S. Borne<sup>1</sup>, V. Gabrel<sup>2</sup>, A.R. Mahjoub<sup>2</sup>, R. Taktak<sup>2</sup>**

(1) LIPN, Paris-13 University

(2) LAMSADE, Paris Dauphine University







# MSOND

## Data

- a bilayer network ( $G_1 = (V_1, E_1)$  and  $G_2 = (V_2, E_2)$ );
- to every router  $v_i \in V_1$  corresponds an optical switch  $w_i \in V_2$ ;
- a set  $K$  of demands between pairs  $(O_k, D_k)$ ;
- for each demand, we know two routing paths  $L_k^1$  et  $L_k^2$  **node-disjoint** in the higher layer;
- for each physical edge  $e \in E_2$  in the lower layer  $\mapsto$  an installation cost  $c_e$ ;
- $G_2$  complete, infinite capacities on the edges.

# MSOND

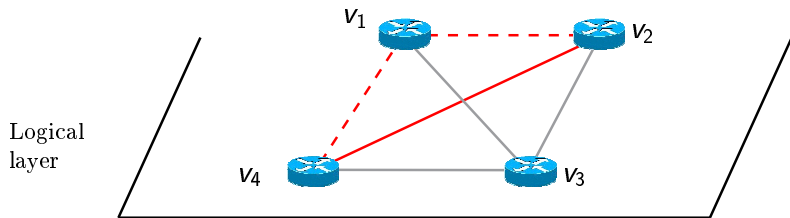
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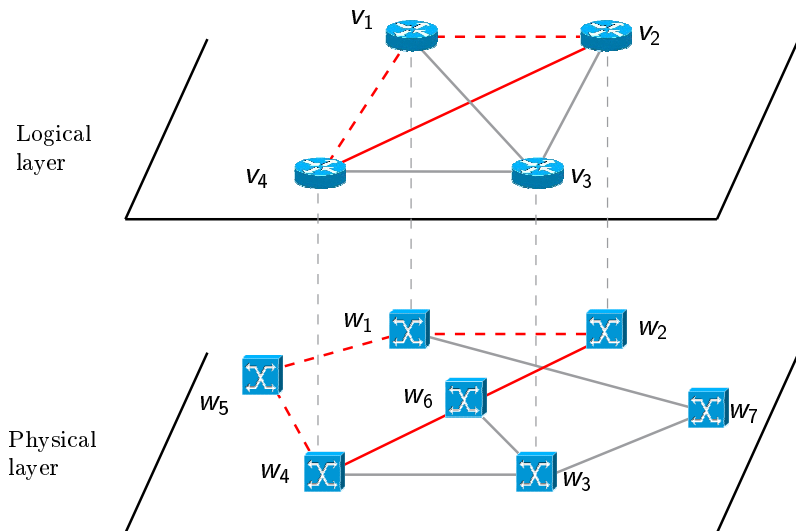
## Objective

find for each demand two **elementary** physical **node-disjoint** paths **respecting the routers order** in  $L_k^1$  and  $L_k^2$  such as installation's **total cost is minimum**.

## Example(1)

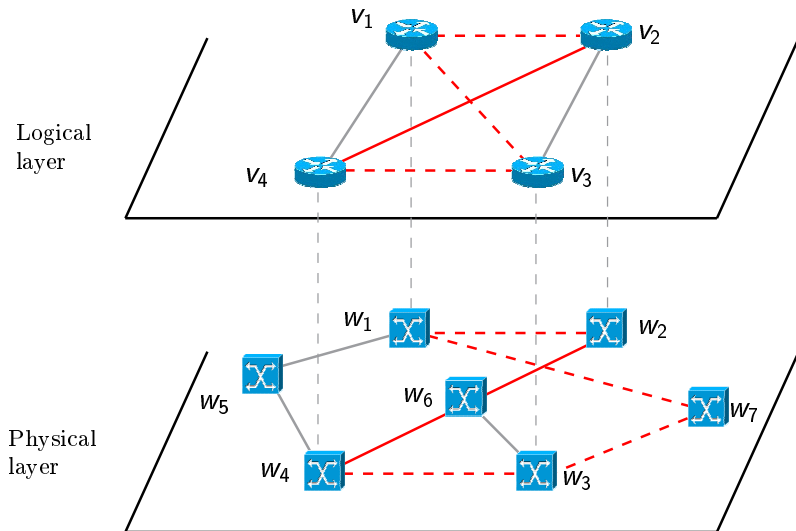


# Example(1)





## Example(2)



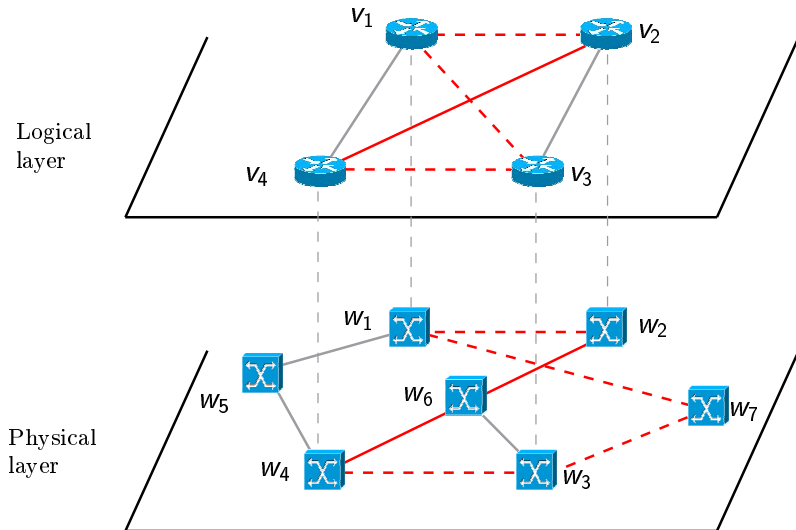
# Motivations

- 1 importance of survivability in the telecommunication context ;

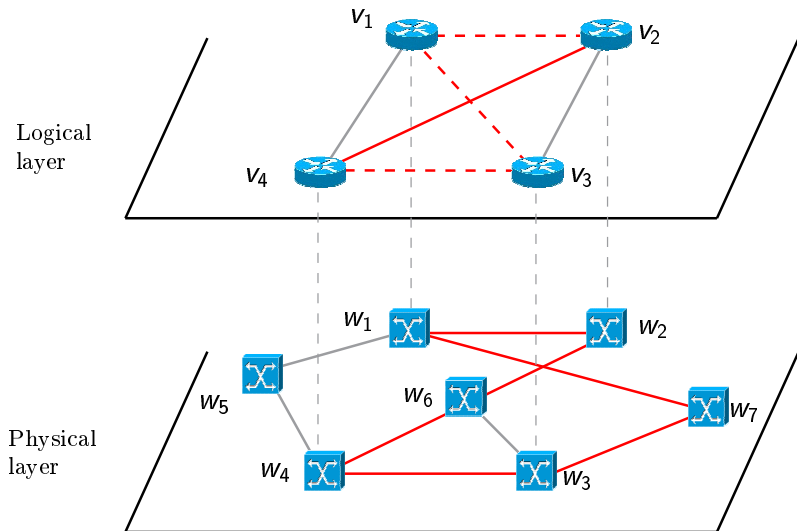
# Motivations

- ① importance of survivability in the telecommunication context ;
- ② tight relationship with some classical problems and in particular the TSP problem ;

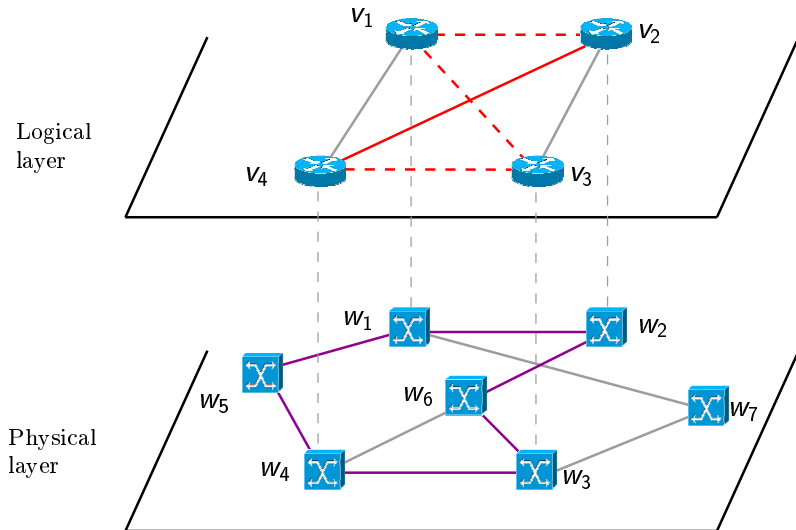
## Example(2)



## Example(2)



## Example(3)



## Particularities

- 1 Elementarity ;
- 2 Steiner and terminal nodes ;
- 3 Precedence constraint.

# Classical Related Problems (1)

Problem	elementarity	Steiner	Order	Complexity
Shortest Path Problem with specified nodes		x		polynomial (special cases) <a href="#">Bajaj (1971)</a> , <a href="#">Ibaraki (1973)</a> , <a href="#">Laporte (1984)</a> , <a href="#">Volgenant (1987)</a>
Steiner cycle		x		<ul style="list-style-type: none"> <li>- polynomial (series-parallel graphs, graphical case) <a href="#">Cornuejols and al. (1985)</a></li> <li>- complete polyhedral description (series-parallel graphs) <a href="#">Baiou and Mahjoub (2002)</a></li> <li>- NP-hard <a href="#">Salazar-Gonzalez (2003)</a></li> <li>- approximation results <a href="#">Steinová (2009)</a></li> </ul>
Steiner TSP	x	x		
TSP with precedence constraints	x		x	NP-hard <a href="#">Balas and al. (1993)</a> , <a href="#">Ruland (1997)</a> , <a href="#">Gouveia and pesneau (2006)</a> , <a href="#">Dumitrescu and al. (2005, 2008)</a> , <a href="#">Acheuer and al. (1995, 2000)</a>
Sequential Ordering Problem (Hamiltonian Path with PC)				



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TSP with precedence constraints	x		x	NP-hard <a href="#">Balas and al. (1993)</a> , <a href="#">Ruland (1997)</a> , <a href="#">Gouveia and pesneau (2006)</a> , <a href="#">Dumitrescu and al. (2005, 2008)</a> , <a href="#">Acheuer and al. (1995, 2000)</a>
Sequential Ordering Problem (Hamiltonian Path with PC)				
SC-MSOND	x	x	x	complexity ?

# Motivations

- ① importance of survivability in the telecommunication context ;
- ② tight relationship with classical problems and in particular the TSP problem ;
- ③ NP-hard even for a single commodity (except for some special cases).