

Delaunay Triangulation

Markus Pawellek

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Abstract

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Contents

Contents	i
1 Introduction	1
2 Background	3
2.1 Graph Theory	3
2.2 Geometry	3
2.2.1 Circumcircle	3
2.2.2 Circumsphere	4
2.2.3 Circumscribed n -Sphere	5
2.3 Delaunay Triangulation and Tessellation	5
3 Algorithms and Data Structures	7
3.1 Hash Map	7
3.2 Triangle Mesh	7
3.3 Quad-Edge Data Structure	7
3.4 Radix Sort	7
3.5 Linear Morton Sort	7
3.6 Linear Floating-Point Quad-Tree	7
3.7 Bowyer-Watson Algorithm	7
3.8 Guibas-Stolfi Incremental Algorithm	7
3.9 Guibas-Stolfi Divide-and-Conquer Algorithm	7
3.10 Dwyer Algorithm	7
4 Design and Implementation	9
4.1 API	9
4.2 Robustness	9
4.3 Incremental Algorithm	9
4.4 Divide-and-Conquer Algorithm	9
4.5 Multidimensional Tessellation	9
5 Tests and Testscenes	11
5.1 Uniform Rectangular Distribution	11
5.2 Gaussian Distribution	11
5.3 Robustness Tests	11
6 Benchmarks	13
7 Examples	15
7.1 Image Mosaic and Tessellation	15
7.2 Fluids for Ray Tracing	15
7.3 Finite Element Method	15
7.4 Pareto Frontiers	15
8 Results	17

CONTENTS

9 Conclusions	19
References	21

1 Introduction

2 Background

2.1 Graph Theory

2.2 Geometry

2.2.1 Circumcircle

$$\|a - m\| = r$$

$$\|b - m\| = r$$

$$\|c - m\| = r$$

$$\|a - m\|^2 = r^2$$

$$\|b - m\|^2 = r^2$$

$$\|c - m\|^2 = r^2$$

$$\|\hat{m}\|^2 = r^2$$

$$\|\hat{b} - \hat{m}\| = r^2$$

$$\|\hat{c} - \hat{m}\| = r^2$$

$$\|\hat{b}\|^2 + \|\hat{m}\|^2 - 2\langle \hat{b} | \hat{m} \rangle = r^2$$

$$\|\hat{c}\|^2 + \|\hat{m}\|^2 - 2\langle \hat{c} | \hat{m} \rangle = r^2$$

$$\|\hat{b}\|^2 - 2\langle \hat{b} | \hat{m} \rangle = 0$$

$$\|\hat{c}\|^2 - 2\langle \hat{c} | \hat{m} \rangle = 0$$

$$\|\hat{b}\|^2 - 2\langle \hat{b} | \hat{m} \rangle = 0$$

$$\|\hat{c}\|^2 - 2\langle \hat{c} | \hat{m} \rangle = 0$$

$$\langle \hat{b} | \hat{m} \rangle = \frac{1}{2} \|\hat{b}\|^2$$

$$\langle \hat{c} | \hat{m} \rangle = \frac{1}{2} \|\hat{c}\|^2$$

$$\begin{pmatrix} \hat{b} & \hat{c} \end{pmatrix}^T \hat{m} = \frac{1}{2} \begin{pmatrix} \|\hat{b}\|^2 \\ \|\hat{c}\|^2 \end{pmatrix}$$

$$\hat{m} = \frac{1}{2 \det \begin{pmatrix} \hat{b} & \hat{c} \end{pmatrix}} \text{adj} \begin{pmatrix} \hat{b} & \hat{c} \end{pmatrix}^T \begin{pmatrix} \|\hat{b}\|^2 \\ \|\hat{c}\|^2 \end{pmatrix}$$

$$\hat{m} = \frac{1}{2(u_x v_y - u_y v_x)} \begin{pmatrix} v_y(u_x^2 + u_y^2) - v_x(v_x^2 + v_y^2) \\ u_x(v_x^2 + v_y^2) - u_y(u_x^2 + u_y^2) \end{pmatrix}$$

2.2.2 Circumsphere

$$\|a - m\| = r$$

$$\|b - m\| = r$$

$$\|c - m\| = r$$

$$\|d - m\| = r$$

$$\|a - m\|^2 = r^2$$

$$\|b - m\|^2 = r^2$$

$$\|c - m\|^2 = r^2$$

$$\|d - m\|^2 = r^2$$

$$\|\hat{m}\|^2 = r^2$$

$$\|\hat{b} - \hat{m}\|^2 = r^2$$

$$\|\hat{c} - \hat{m}\|^2 = r^2$$

$$\|\hat{d} - \hat{m}\|^2 = r^2$$

$$\|\hat{b}\|^2 + \|\hat{m}\|^2 - 2 \langle \hat{b} | \hat{m} \rangle = r^2$$

$$\|\hat{c}\|^2 + \|\hat{m}\|^2 - 2 \langle \hat{c} | \hat{m} \rangle = r^2$$

$$\|\hat{d}\|^2 + \|\hat{m}\|^2 - 2 \langle \hat{d} | \hat{m} \rangle = r^2$$

$$\|\hat{b}\|^2 - 2 \langle \hat{b} | \hat{m} \rangle = 0$$

$$\|\hat{c}\|^2 - 2 \langle \hat{c} | \hat{m} \rangle = 0$$

$$\|\hat{d}\|^2 - 2 \langle \hat{d} | \hat{m} \rangle = 0$$

$$\|\hat{b}\|^2 - 2 \langle \hat{b} | \hat{m} \rangle = 0$$

$$\begin{aligned}\|\hat{c}\|^2 - 2\langle\hat{c}|\hat{m}\rangle &= 0 \\ \|\hat{d}\|^2 - 2\langle\hat{d}|\hat{m}\rangle &= 0\end{aligned}$$

$$\begin{aligned}\langle\hat{b}|\hat{m}\rangle &= \frac{1}{2}\|\hat{b}\|^2 \\ \langle\hat{c}|\hat{m}\rangle &= \frac{1}{2}\|\hat{c}\|^2 \\ \langle\hat{d}|\hat{m}\rangle &= \frac{1}{2}\|\hat{d}\|^2\end{aligned}$$

$$\begin{pmatrix}\hat{b} & \hat{c} & \hat{d}\end{pmatrix}^T \hat{m} = \frac{1}{2} \begin{pmatrix}\|\hat{b}\|^2 \\ \|\hat{c}\|^2 \\ \|\hat{d}\|^2\end{pmatrix}$$

$$\hat{m} = \frac{1}{2 \det \begin{pmatrix}\hat{b} & \hat{c} & \hat{d}\end{pmatrix}} \text{adj} \begin{pmatrix}\hat{b} & \hat{c} & \hat{d}\end{pmatrix}^T \begin{pmatrix}\|\hat{b}\|^2 \\ \|\hat{c}\|^2 \\ \|\hat{d}\|^2\end{pmatrix}$$

2.2.3 Circumscribed n -Sphere

$$\begin{aligned}\|x_0 - m\| &= r \\ \|x_i - m\| &= r \\ \|x_i - m\|^2 &= r^2 \\ \|\hat{x}_i - \hat{m}\|^2 &= r^2 \\ \|\hat{x}_i\|^2 + \|\hat{m}\|^2 - 2\langle\hat{x}_i|\hat{m}\rangle &= r^2 \\ \|\hat{x}_i\|^2 - 2\langle\hat{x}_i|\hat{m}\rangle &= 0 \\ \langle\hat{x}_i|\hat{m}\rangle &= \frac{1}{2}\|\hat{x}_i\|^2\end{aligned}$$

$$\begin{pmatrix}\hat{x}_1 & \dots & \hat{x}_n\end{pmatrix}^T \hat{m} = \frac{1}{2} \begin{pmatrix}\|\hat{x}_1\|^2 \\ \vdots \\ \|\hat{x}_n\|^2\end{pmatrix}$$

Here, we should not try to build the adjoint. For higher dimensions, it is much more efficient to use a simple matrix solver, like LU decomposition or Cholesky.

2.3 Delaunay Triangulation and Tessellation

Triangulation, Tessellation, Simplicialization, Subdivision, Mesh Generation Two and Three Dimensions

3 Algorithms and Data Structures

3.1 Hash Map

3.2 Triangle Mesh

There are different schemes for storing a triangle. Information about neighbors seem to be important to improve efficiency but not necessary some algorithms.

3.3 Quad-Edge Data Structure

3.4 Radix Sort

3.5 Linear Morton Sort

3.6 Linear Floating-Point Quad-Tree

3.7 Bowyer-Watson Algorithm

Advantages: Easy structures. Low memory consumption. Possible in every dimension. Incremental construction at arbitrary point in domain. Disadvantages: Complexity is $\Theta(n^2)$. Algorithm needs bounding super triangle. Bad parallelization.

3.8 Guibas-Stolfi Incremental Algorithm

3.9 Guibas-Stolfi Divide-and-Conquer Algorithm

3.10 Dwyer Algorithm

4 Design and Implementation

4.1 API

4.2 Robustness

4.3 Incremental Algorithm

4.4 Divide-and-Conquer Algorithm

4.5 Multidimensional Tessellation

5 Tests and Testscenes

5.1 Uniform Rectangular Distribution

5.2 Gaussian Distribution

5.3 Robustness Tests

6 Benchmarks

7 Examples

7.1 Image Mosaic and Tessellation

SVG output is generated. We need the image as PNG or JPG. A seed will be generated automatically or given on command line. Hence, it should be returned as well. Also the triangle count can be defined.

7.2 Fluids for Ray Tracing

7.3 Finite Element Method

7.4 Pareto Frontiers

8 Results

9 Conclusions

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