

Wire Images and Path Signals

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measurement and architecture for a middleboxed internet

measurement

architecture

experimentation



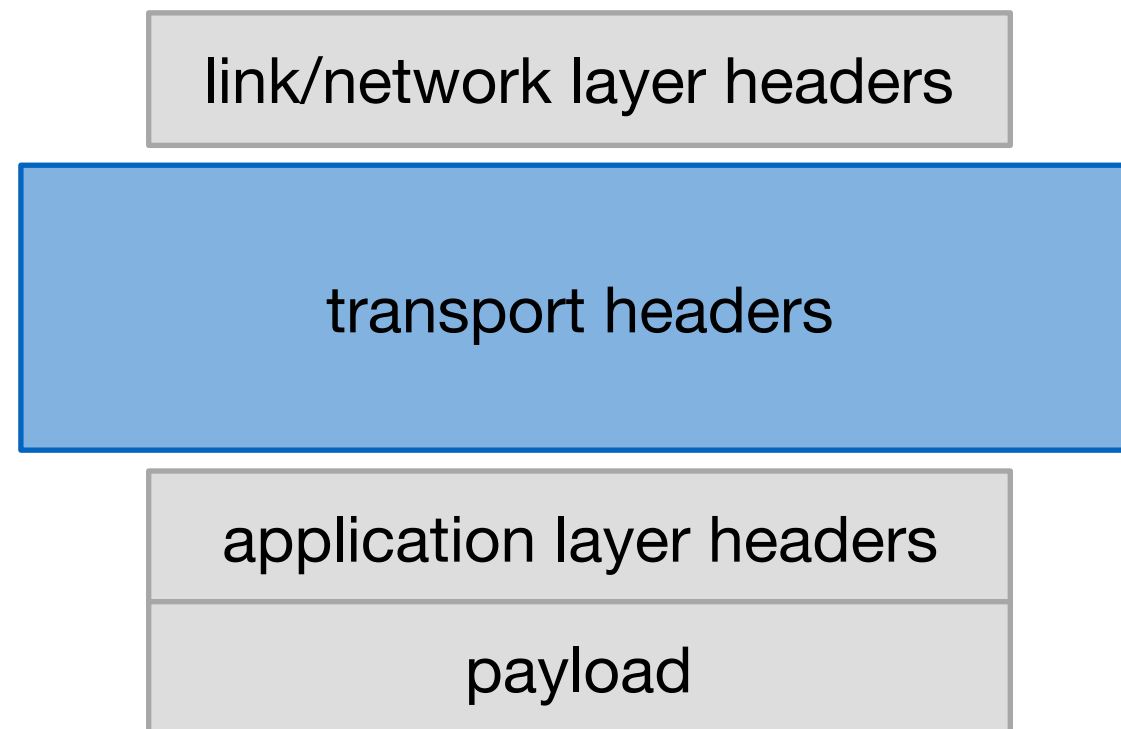
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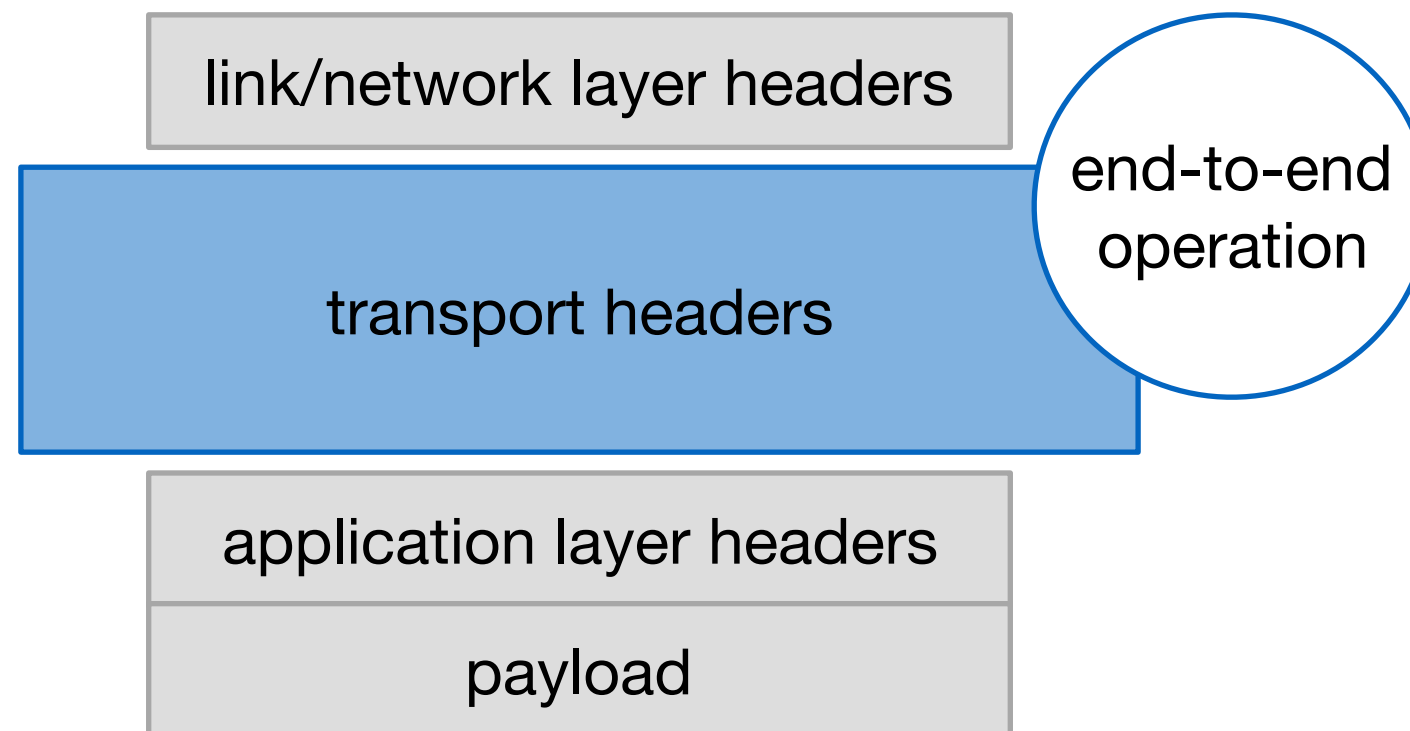


Transport protocol design, 1990s style



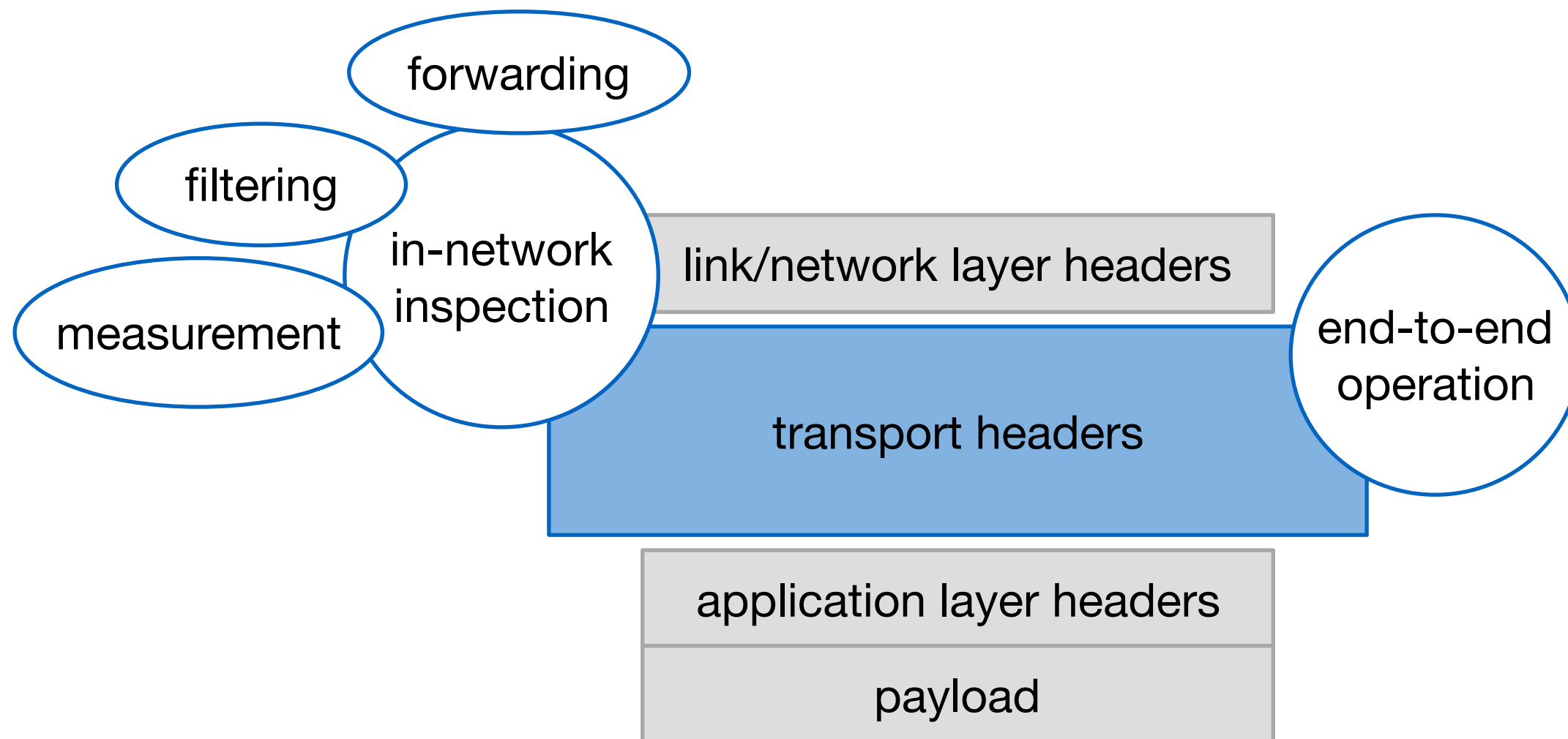


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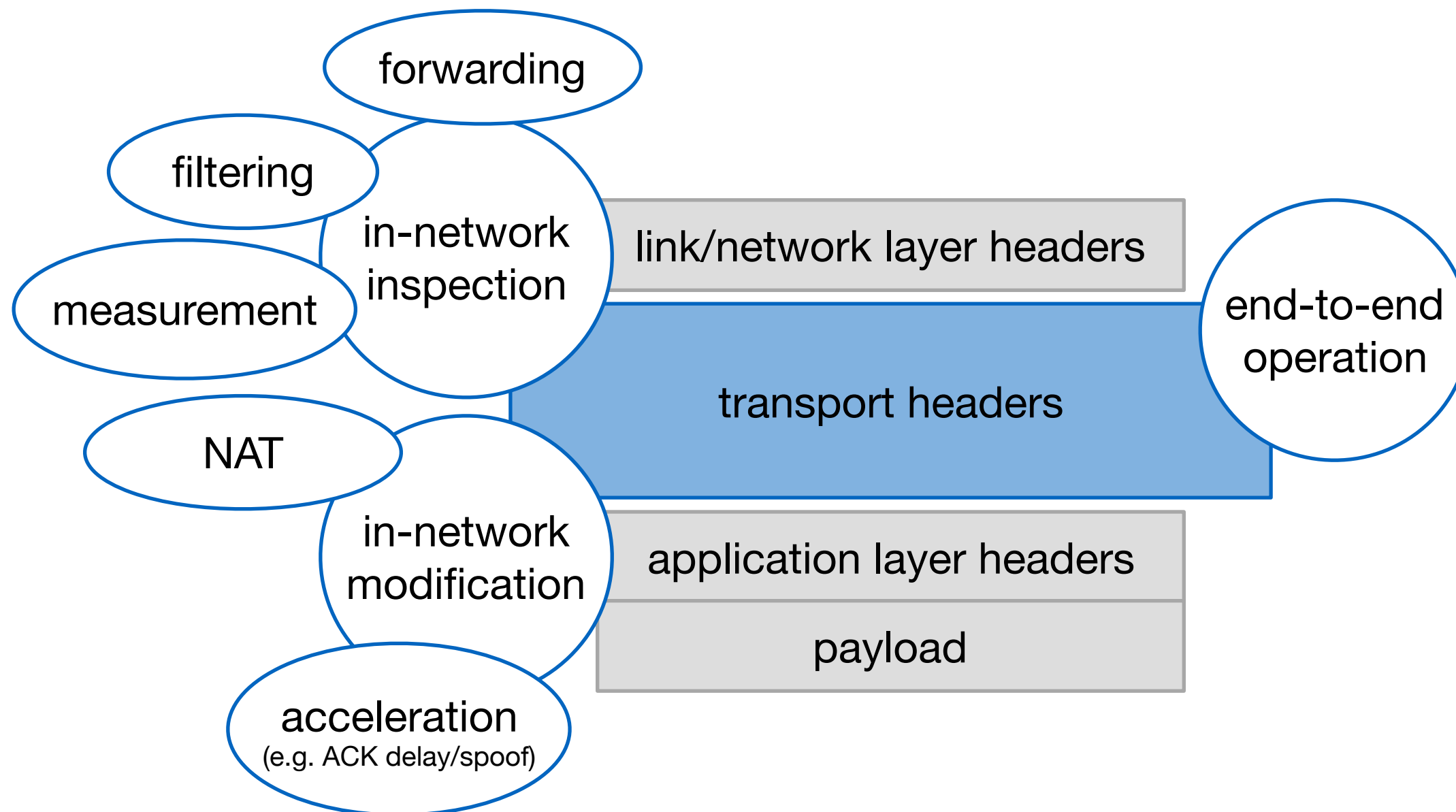


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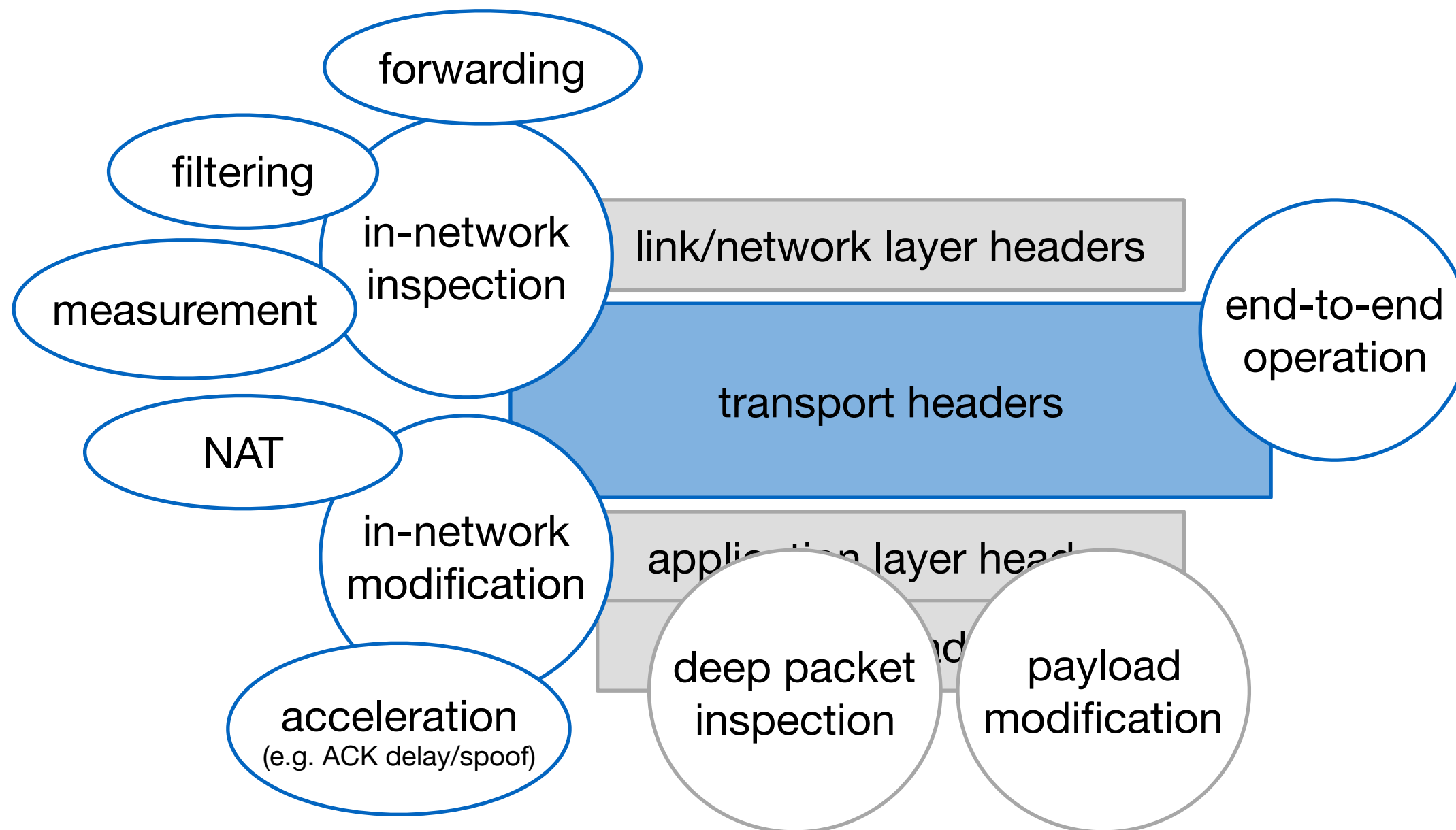


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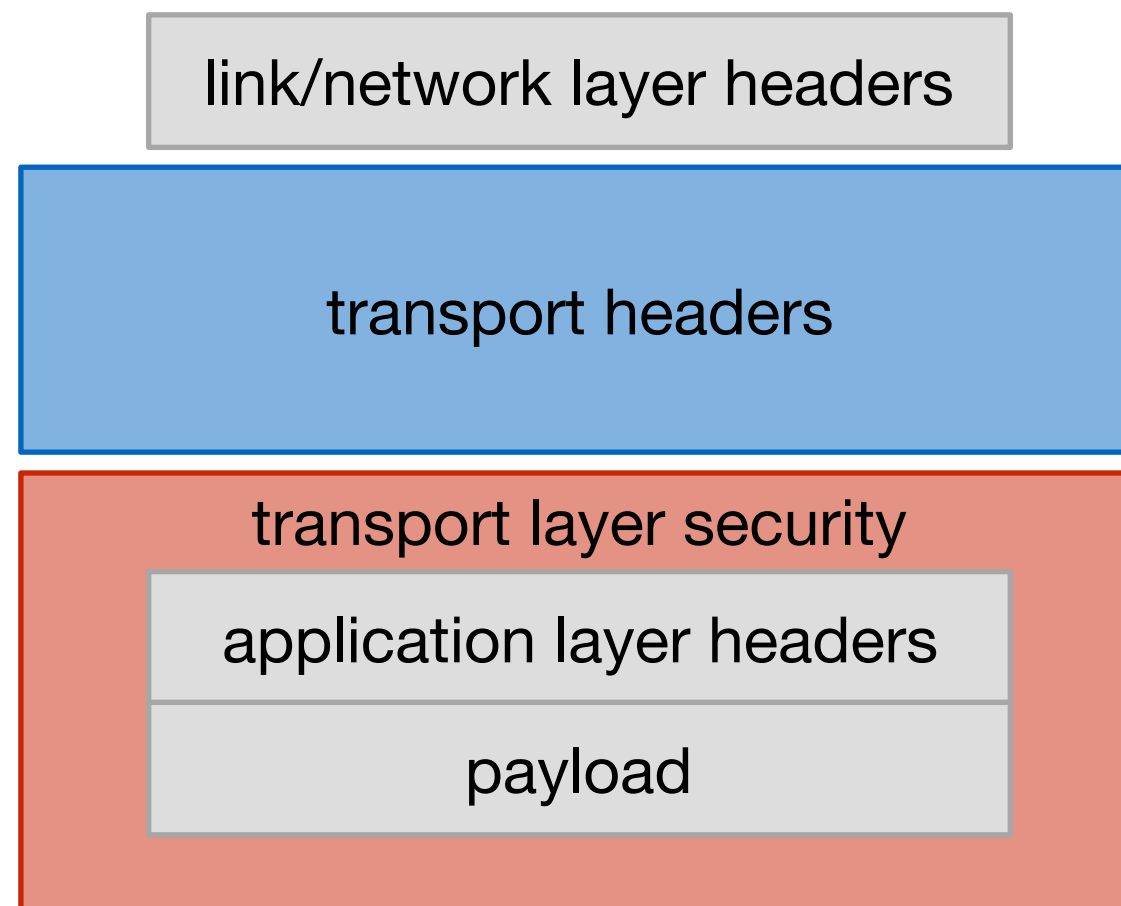


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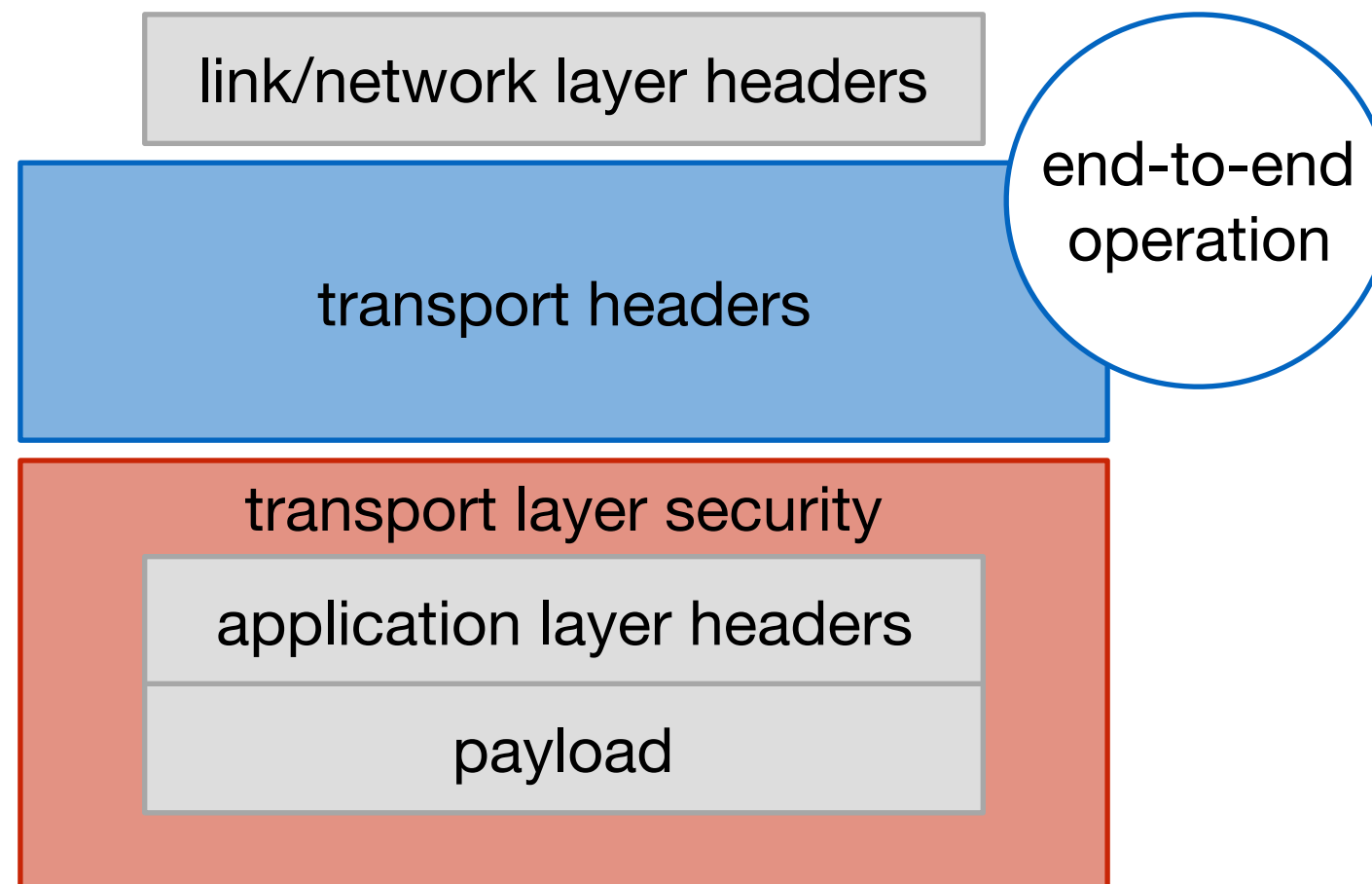


Transport protocol design, now with security!



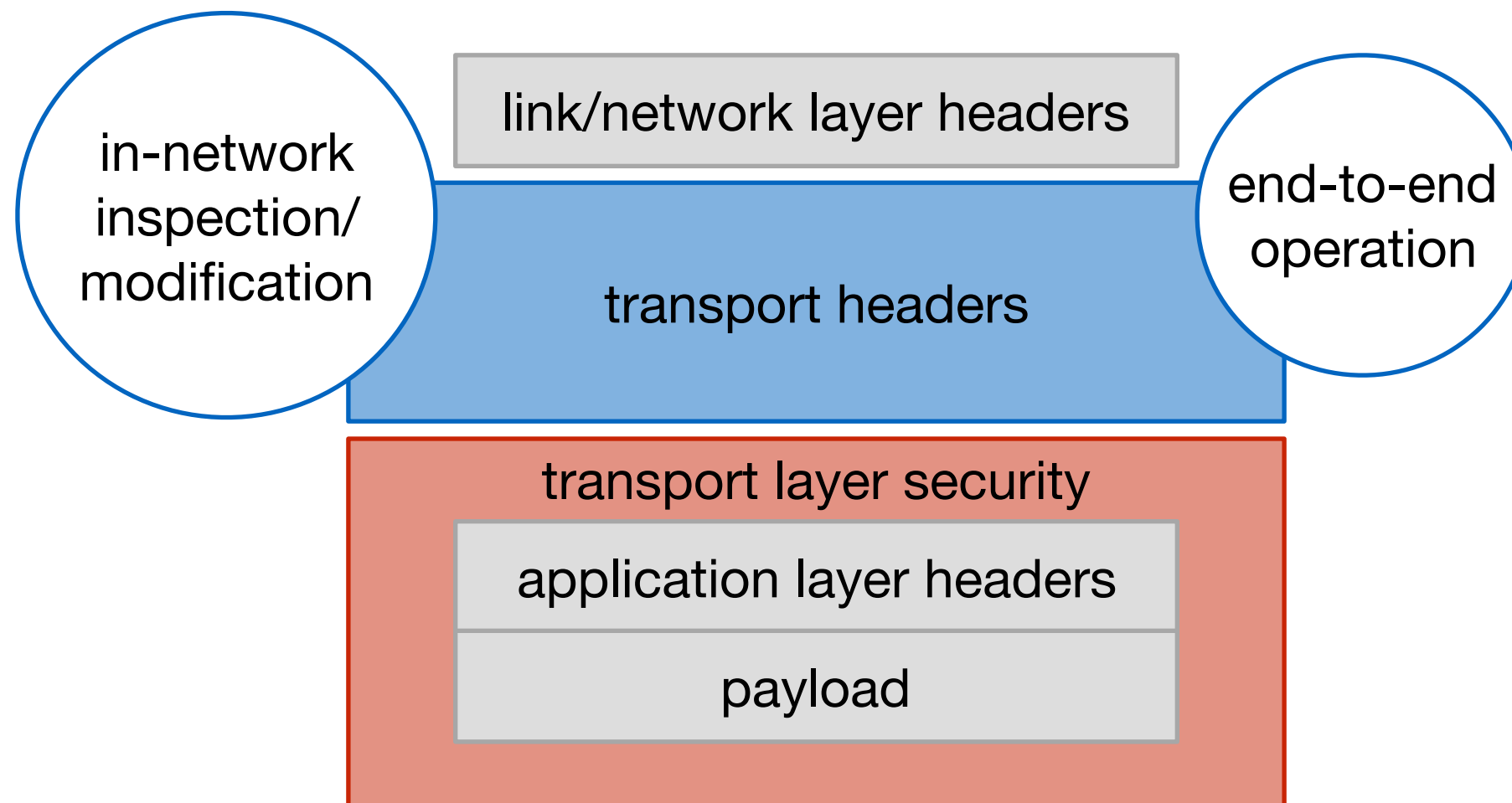


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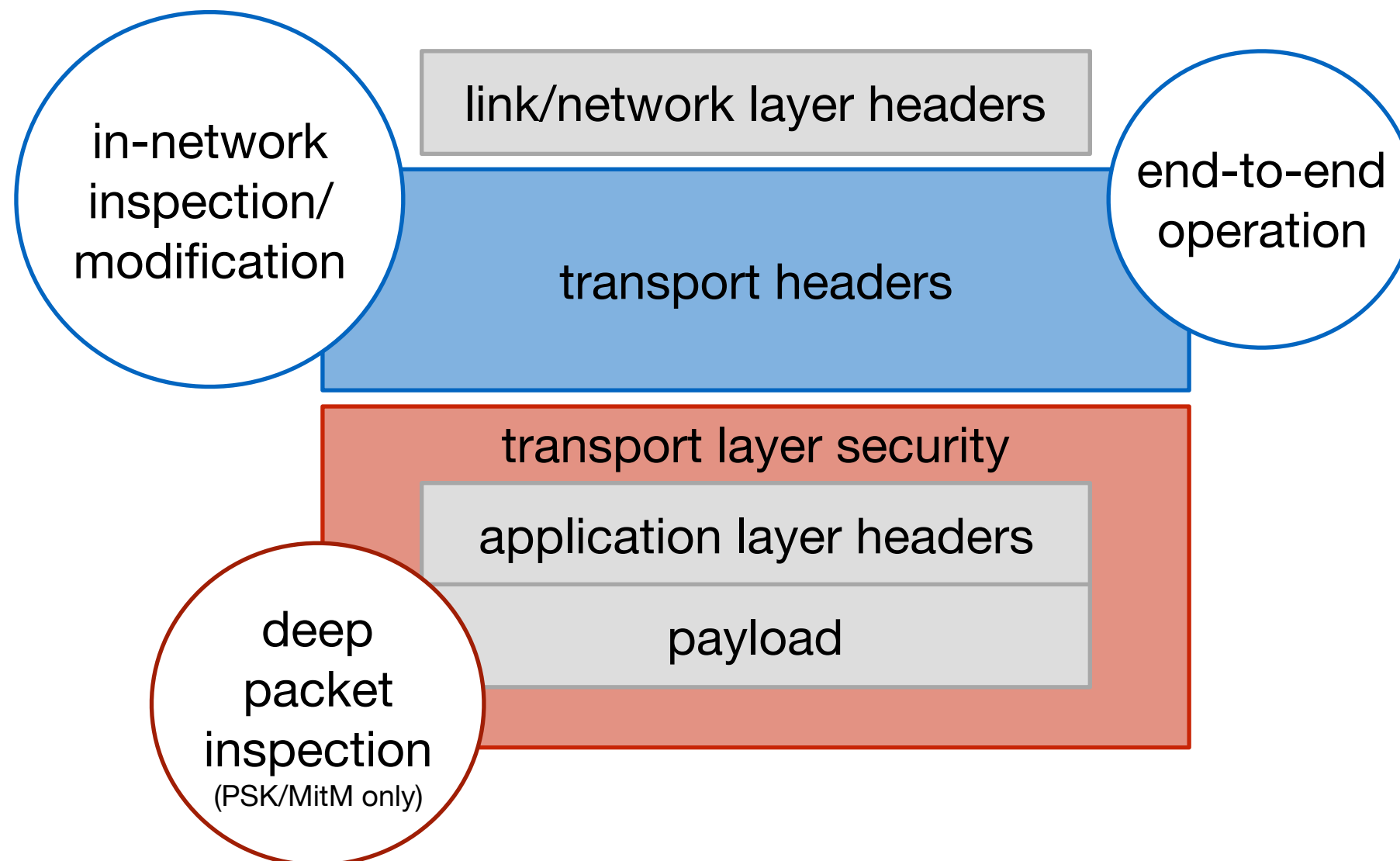


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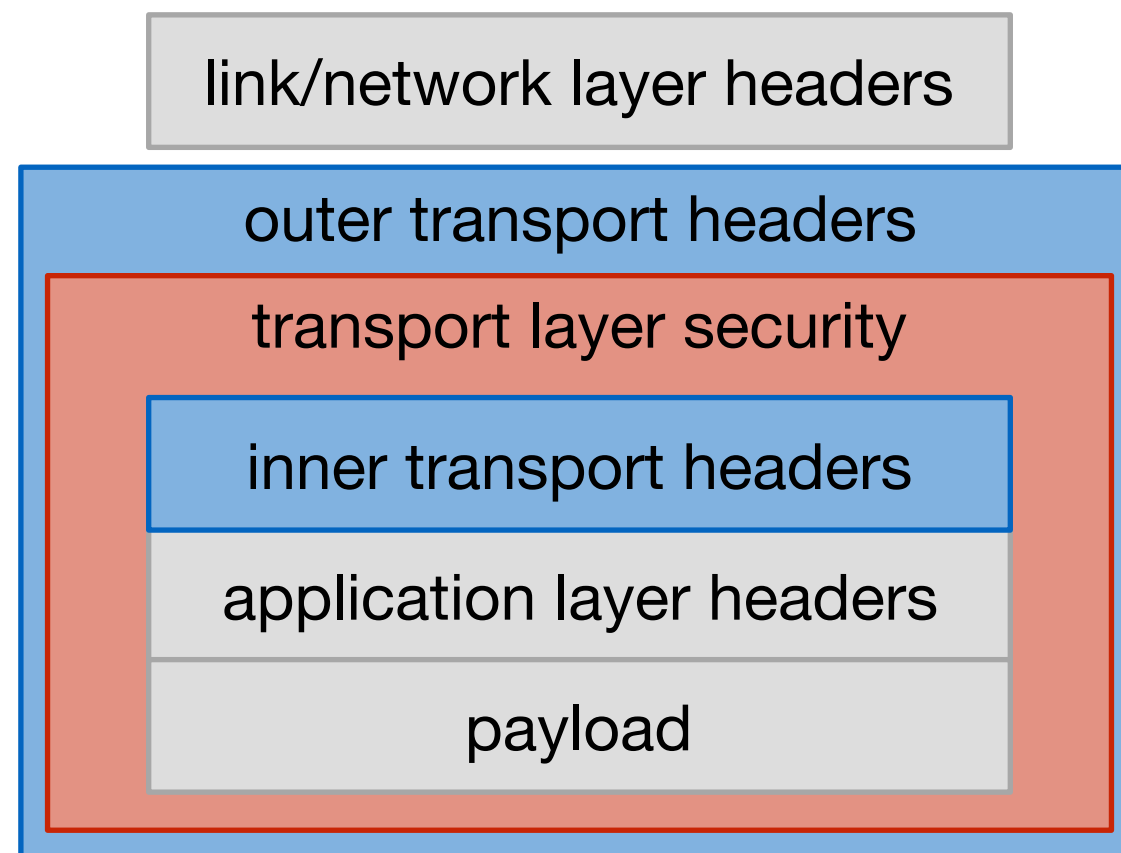


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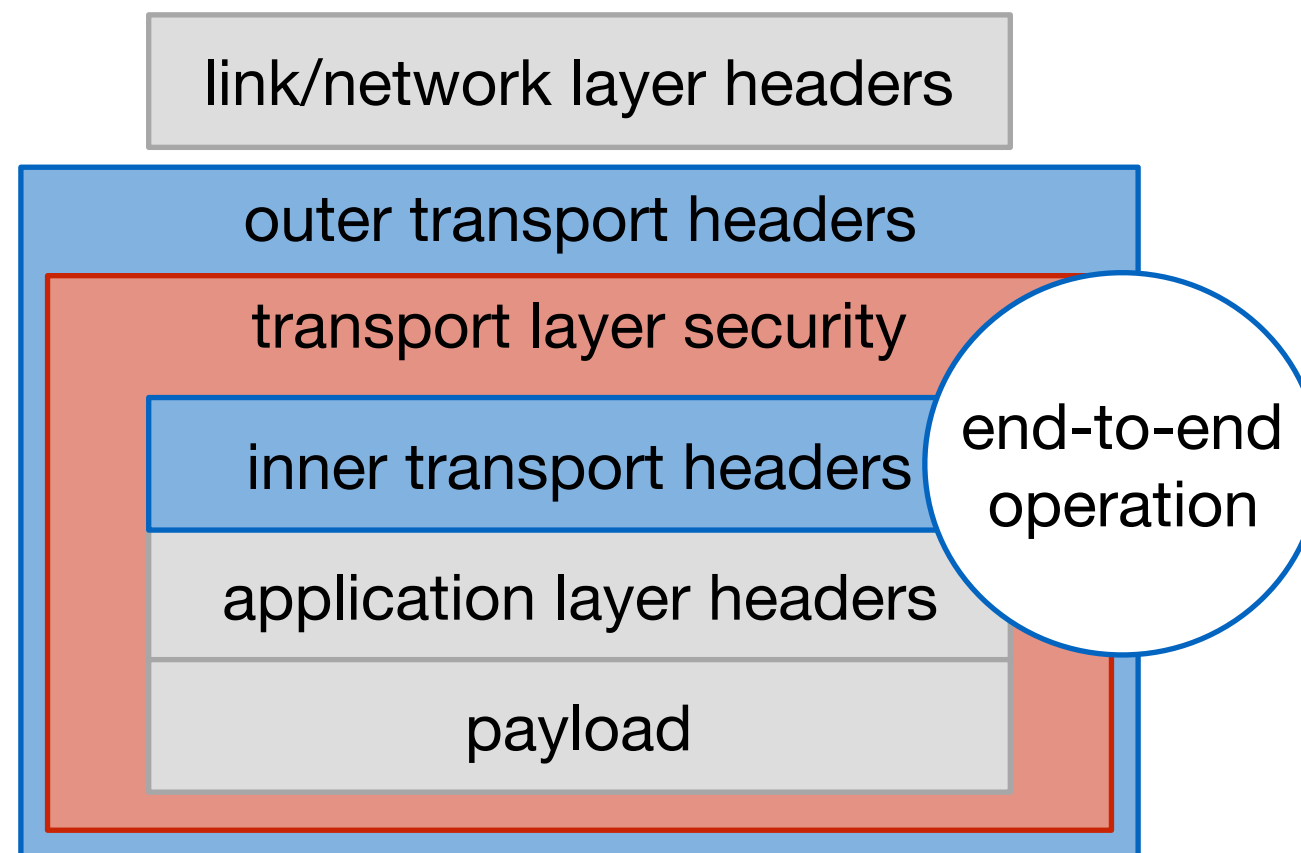


Encrypted transport protocol design



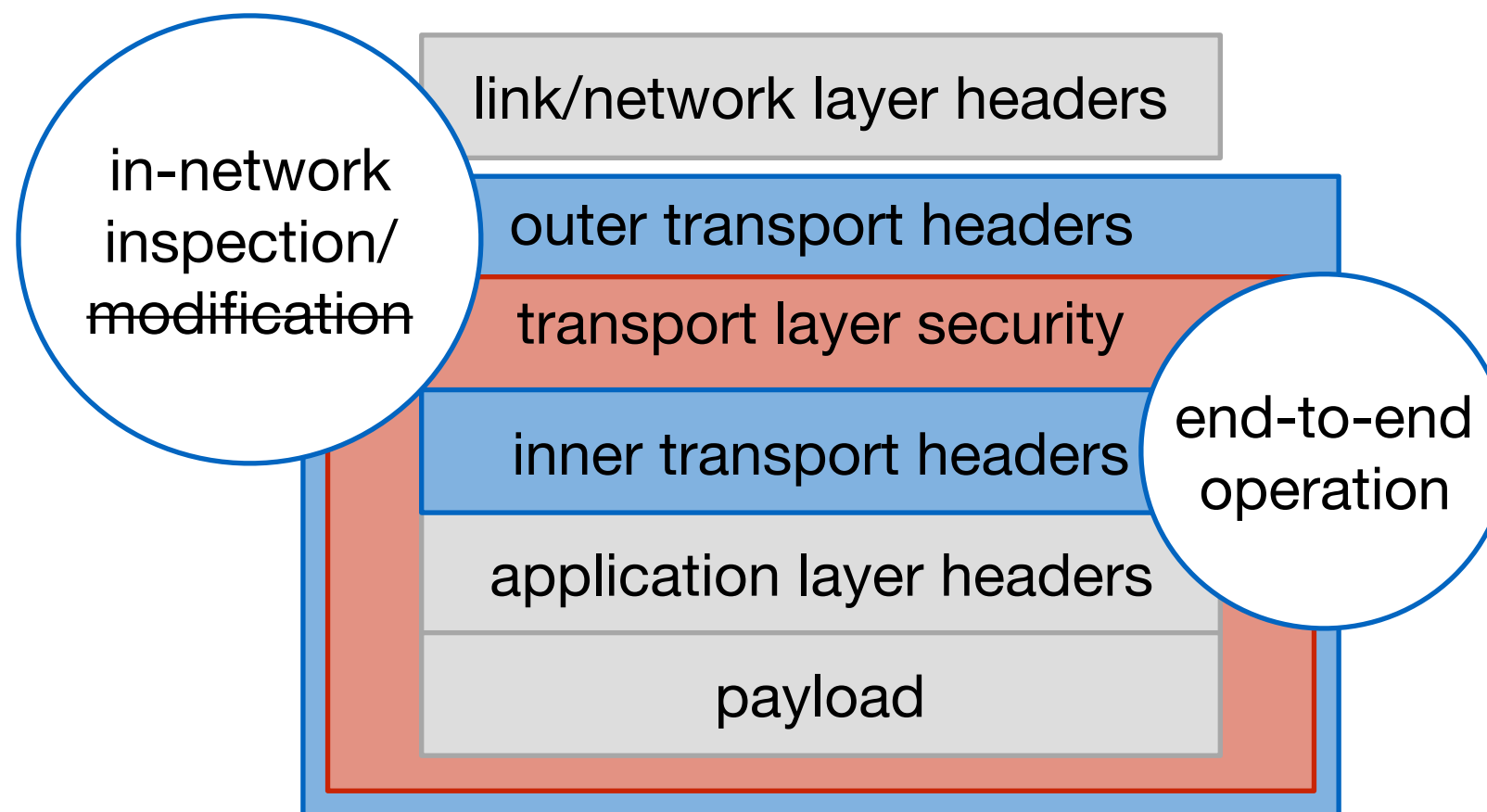


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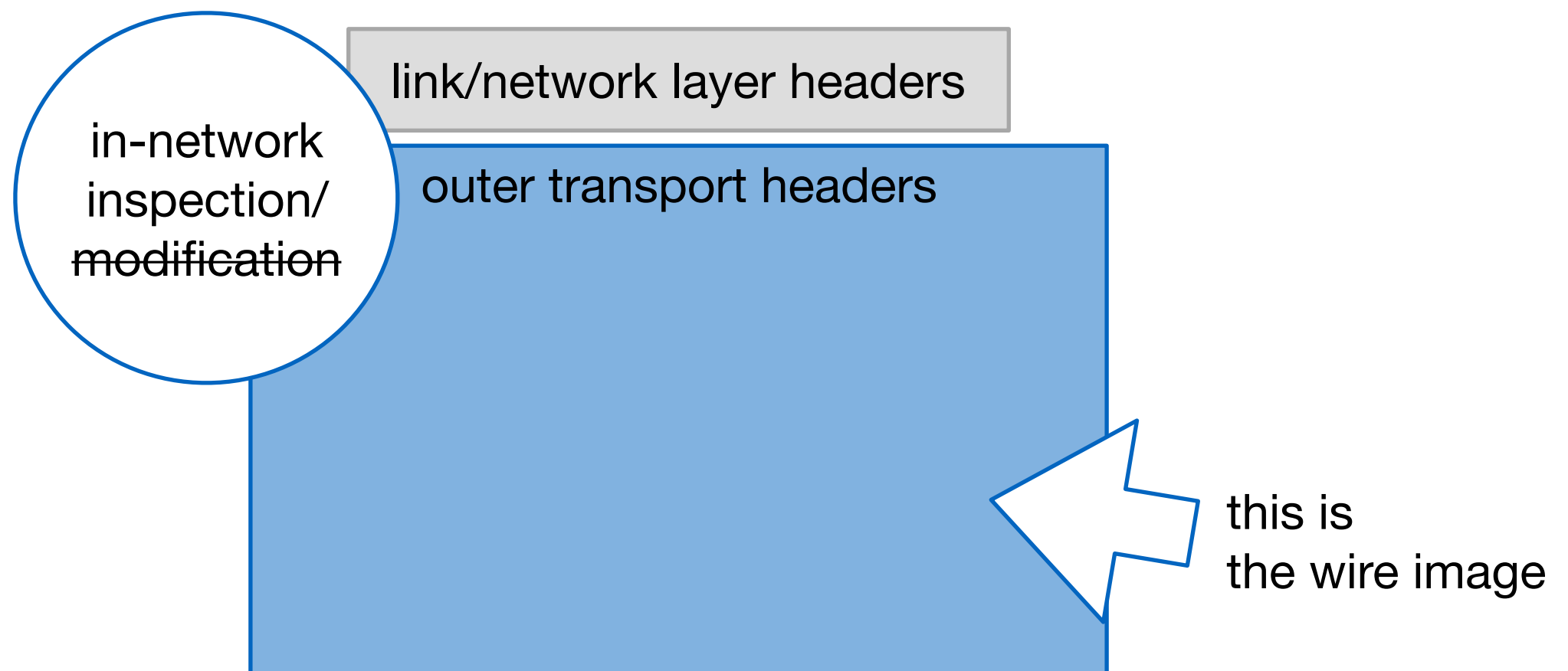


Encrypted transport protocol design





Encrypted transport protocol design





Three (and a half) levels of accessibility

- Unprotected (as TCP today): all bits can be seen at all points along the path, and can be modified without endpoint knowledge.
 - Supports manipulation of transport internals.
- Integrity-protected (e.g. QUIC outer header): all bits can be seen at all points along the path, but modification is endpoint detectable (and leads to packet rejection).
 - "Scratch space" (e.g. PLUS extended header): bits intended for modification on the path; content not protected, length/meaning integrity-protected.
- Encrypted (e.g. QUIC frame headers): no bits can be interpreted, modification leads to corruption and rejection.



The Wire Image, more formally (draft-trammell-wire-image)

- The sequence of messages sent by each participant in the protocol
 - each expressed as a sequence of bits,
 - with an associated time at which each was sent.
- Only unencrypted bits have assignable semantics
 - encrypted size → upper bound on information content
- Separating transport machinery from the understandable parts of the wire image is new, presents both problems and opportunities.



Signaling in the wire image (draft-hardie-path-signals)

- Encrypting transport mechanics limits the availability of *implicit signaling* in transport headers to in-network functions.
- We have four options to deal with this:
 - Do nothing.
 - Replace transport information with network-layer signaling.
 - Add *explicit signals* on a per-transport basis
 - Add *explicit signals* on a common wire image shared by multiple transports



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Transport-layer headers were never meant to be used by the network. It's the network's fault for abusing them, and solely the problem of in-network function developers and users to solve.



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 - Do nothing.
 - Replace transport information with network-layer signaling.
 - Add *explicit signaling*. Talking to the network is a network-layer function. Therefore any explicit signaling should happen in the network layer.
 - Add *explicit signaling* to multiple transports.



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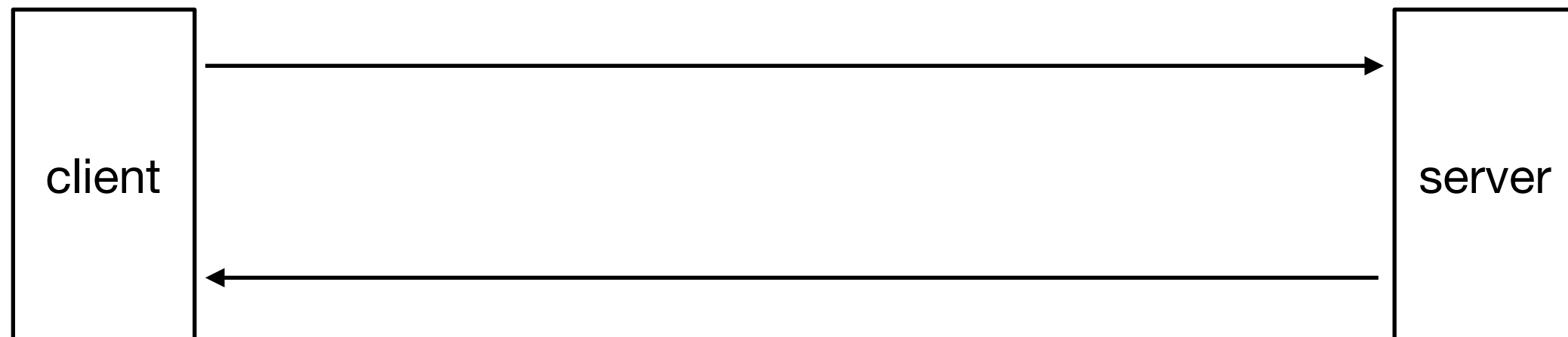
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An example explicit signal: draft-trammell-quic-spin

- Proposed bit in QUIC short header, visible/integrity-protected.
- Algorithm ensures bit changes $0 \rightarrow 1$ or $1 \rightarrow 0$ once per RTT.
- Allows on-path determination of RTT on a per-flow basis.
 - Replaces TCP seq/ack and tsval/tescr for this purpose.

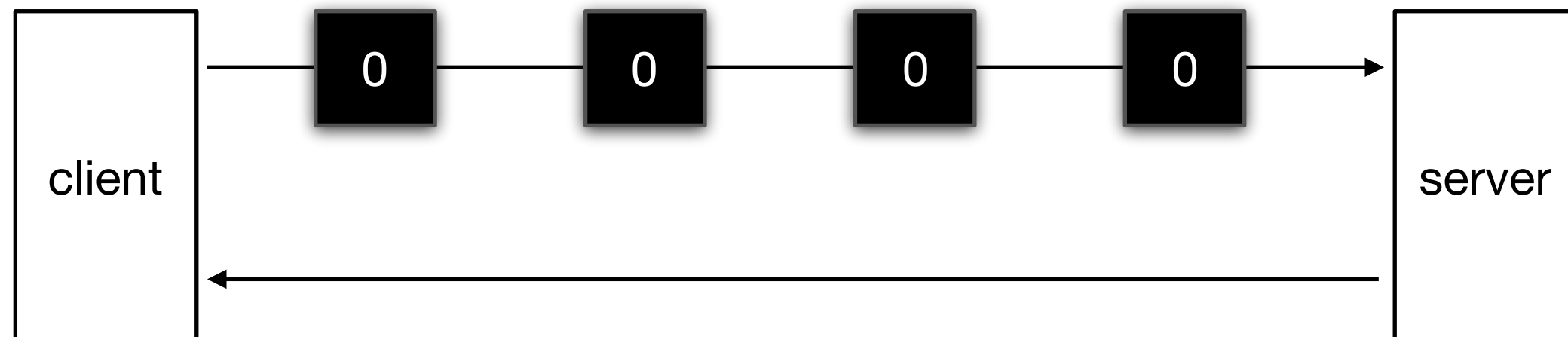


- Discussion in QUIC WG, 09:30-12:00 Thursday



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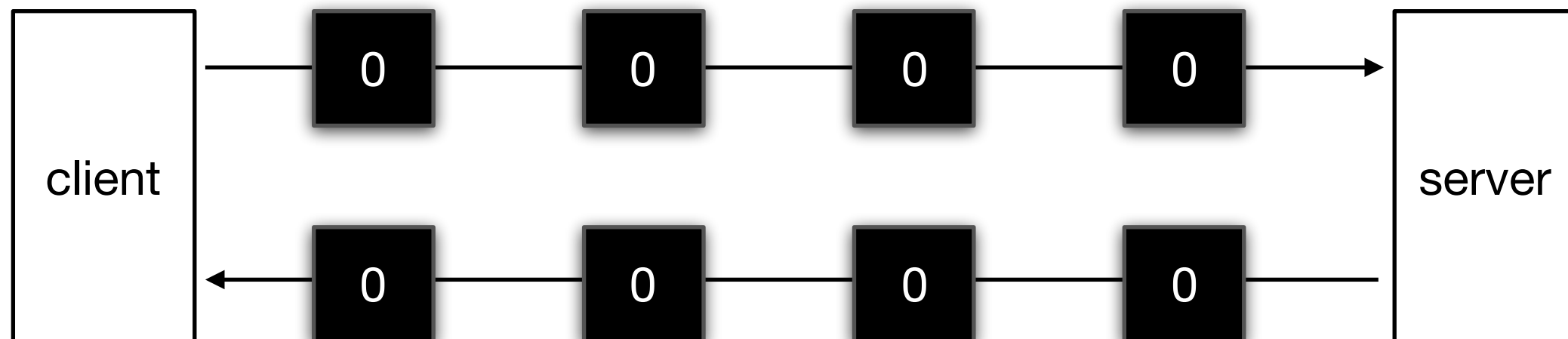


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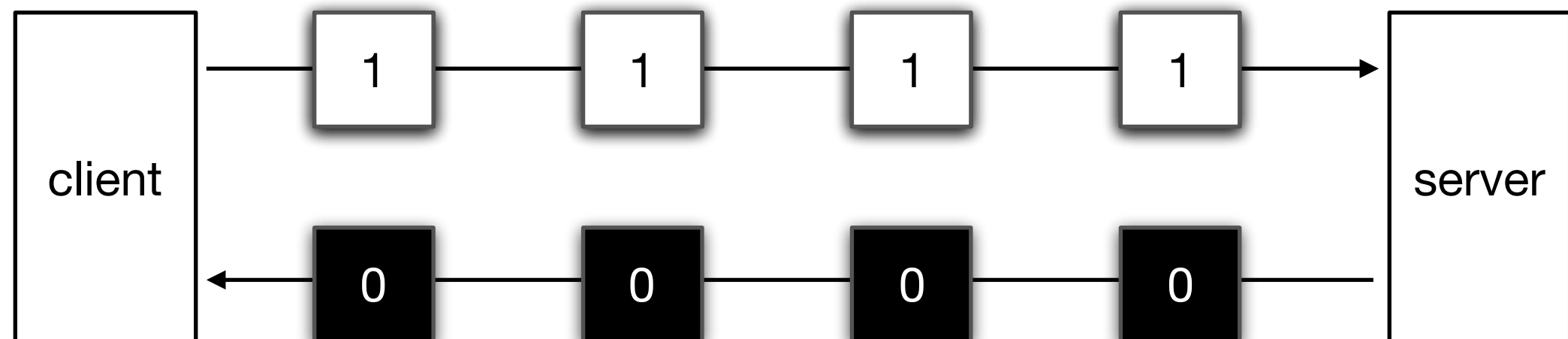


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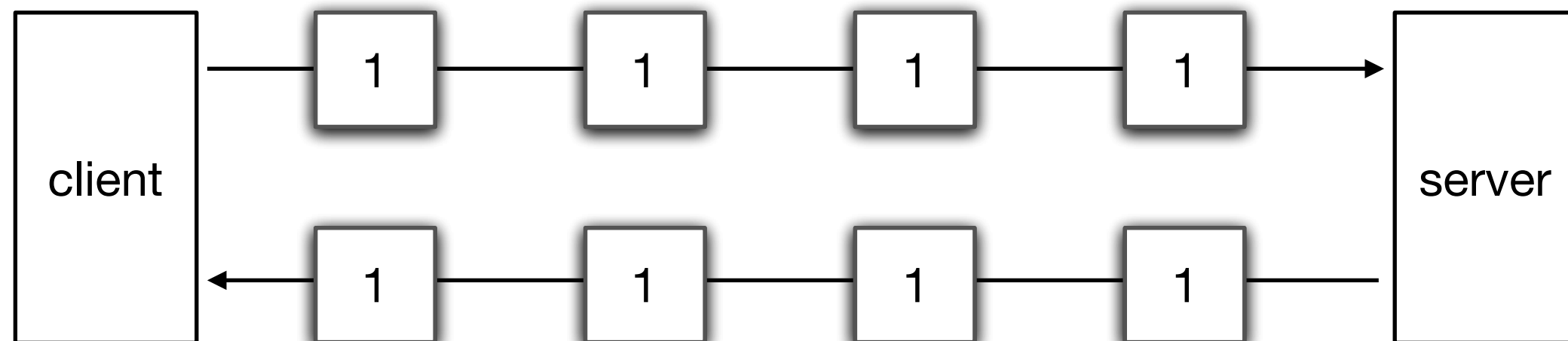


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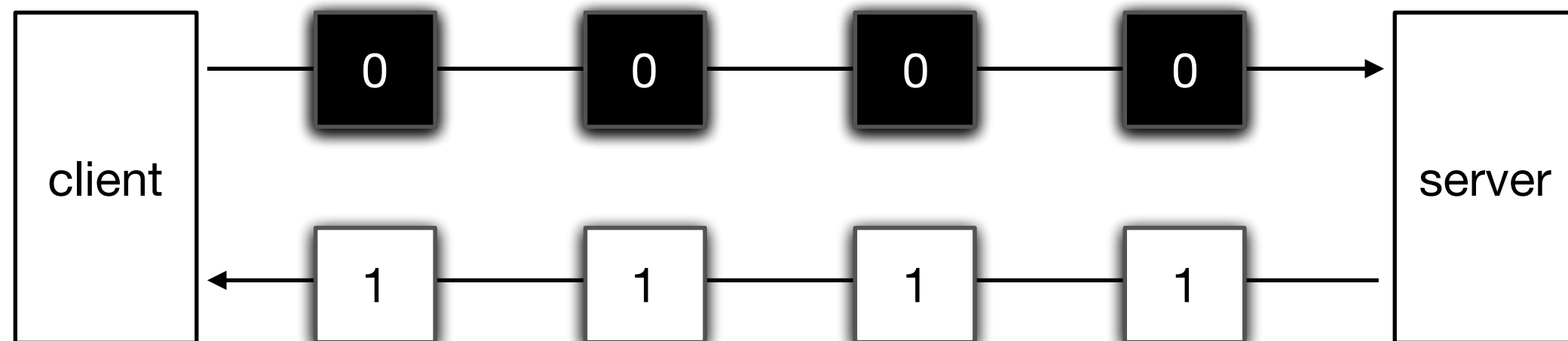


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The Way Forward

- Header encryption is necessary for transport layer evolution.
 - Every visible bit of information in a protocol's wire image will eventually be used → we now know "everyone can see everything" is the wrong approach.
- Finding a balance between security, privacy, evolvability, and maintenance of current in-network functions is crucial.
 - Explicit signaling to the path must be done on a per-protocol, per-function basis.
 - Signal design must consider *who* needs to see *what* and *why*, and provide the *minimum necessary information* to drive the function.