

The problem can be attacked in multiple front. One of the main point is to observe that the maximum cost has to be under 500(the query weight). So we can do DFS, BFS or modified Dijkstra with only checking for weight to not cross the query weight **W**.

Judge solution was treating the pair of node and cost to reach that node as a composite node. So we have $100 * 500$ composite nodes. Then use a DFS from source **U** and cost 0 as starting point in the DFS. Then do the DFS without exceeding the cost from **W**. After the DFS, for finishing node **V** start checking from **V** and cost **W** to **W-1**, **W-2** until you find a composite node which was visited in the DFS. That cost is the result. If no cost found then it is impossible to go to **V** without cost exceeding **W**.