

### **Problem A: Magic Trick**

**Setter:** Md. Samiul Alam

**Solution:**

**Just print “3”.**

### **Problem B: Cunning Comedian**

**Setter:** Md. Mushfiquur Rahman Sanim

**Solution:**

1. Standard 0-1 BFS
2. If Mr. X takes minimum  $n$  minutes to go to any adjacent cell of the robot then he needs  $2*n$  time to leave the house. So, if  $2*n$  is less than  $D$  then he can successfully take the robots otherwise Mr. Y will keep them.

### **Problem C: K- String**

**Setter:** Simanta Deb Turja

**Category:** Greedy/Binary Search

**Solution:**

1. If  $k$  is greater than the number of distinct characters then answer =  $k - \text{number of distinct characters}$ .
2. If  $k$  is equal to the distinct number of characters then answer = 0
3. If  $k$  is smaller than the number of distinct of characters then answer = sum first (number of distinct characters –  $k$ ) after sorting the characters according to their frequencies.

### **Problem D: Bad Cook**

**Setter:** Mir Imtiaz Mostafiz Naved

**Category:** Math

**Solution:**

1. If  $N$  is odd, answer =  $((N / 2)! * (N / 2 + 1)!) - 1$
2. If  $N$  is even, answer =  $2 * ((N / 2)!)^2 + (((N / 2)!)^2 * (N / 2 - 1)) - 1;$

### **Problem E: Digit Printing**

**Setter:** Sifat Siddiqi Shishir

**Category:** Implementation

**Solution:** Just print what the problem ask for.

## **Problem F: Convert String to Palindrome**

**Setter:** Rajon Bardhan

**Category:** Dynamic Programming

**Solution:**

1. 1. If the last character of the string is same as the first character, no deletion is needed and we recur for the remaining substring  $X[i + 1, j - 1]$
2. 2. If last character of string is different from the first character, then we return one plus maximum of the two values we get by
  - deleting the last character and recursively try for the remaining substring  $X[i, j - 1]$
  - deleting the first character and recursively try for the remaining substring  $X[i + 1, j]$

## **Problem G: Buildings!**

**Setter:** Sifat Rabbi

**Category:** Binary Search/Greedy

**Solution:** Just binary search for the height of the building.

1. If you can build a building of height  $H$ , try with a height bigger the  $H$  else try with a height smaller than  $H$ .

## **Problem H: Tour**

**Setter:** Shakil Ahmed

**Category:** Bitmask DP

**Solution:** Since the number of songs is not so big, try all the subset of songs and check whether all student can perform at least one song. Then take the minimum number of songs from the subset.

## **Problem I: Help**

**Setter:** Fahim Ferdous Neerjhor

**Category:** Geometry

**Solution:**

1. If the tree cannot be trapped with a triangle, then it must be on the intersection point of two diagonals of a quadrilateral.
2. As, there can only be at most  $n/2$  lines that contain the tree, we can check all pairs of lines that contain the tree and check the quadrilateral formed from those pair of lines and take the one with smallest perimeter.

**Problem J: Interesting Pile Game**

**Setter:** Ibnul Tahsin Bhuiyan

**Category:** Game Theory, DP

**Solution:**

Let,  $F(m) = 1$  if current player can win with  $m$  stones and

$F(m) = 0$  if he can't win with  $m$  stones (If both of them play optimally).

To calculate  $F(m)$  we should check if any valid move after taking  $i$  stones makes the next move returning 0, If so then current player can take  $i$  stones and put other player in that losing state in next move.

Otherwise current player can't find such a move and loses.

**Problem K: Phi Numbers in Range**

**Setter:** G. M. Shahariar Shibli

**Category:** Number Theory, Data Structures(Mo's Algorithm, SQRT Decomposition)

**Solution:**

1. Pre-generate Phi numbers up to  $10^5$
2. Co-ordinate compression on the values to map array values in range:  $[1, N]$ .
3. Divide the array into blocks using sqrt decomposition. For each block maintain a bit set. The bitset has  $N$  bits and marks whether the  $i$ -th value (after coordinate compression) is present or not.
4. For all pairs of block numbers  $(i, j)$  find the distinct elements. For any  $(L, R)$ , it is simply the bitwise OR of all bit sets in  $[L, R]$ .

**Alternative Approach:** Mo's Algorithm + Segment Tree