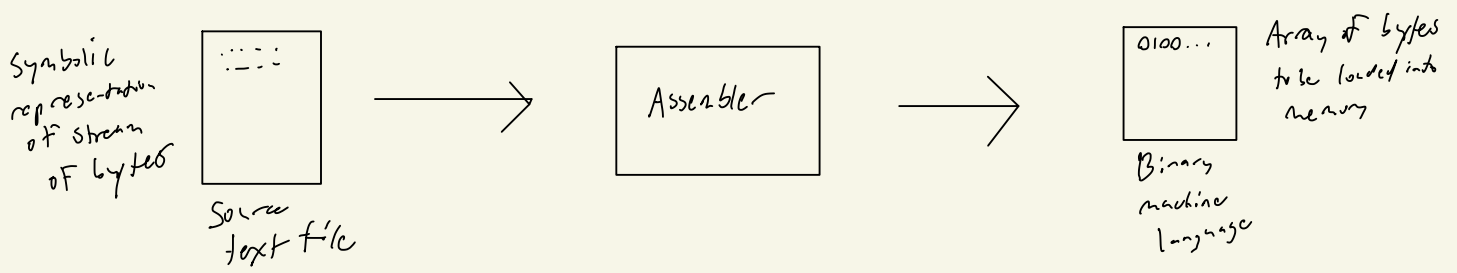


Ch 10



Assembly Language



* This course uses VASM Assembly Language

↳ Main elements

- ↳ comments
- ↳ Symbols → constants
- ↳ Labels → symbols for addresses
- ↳ Macros → instructions (expands into seq. of bytes)

Assembling

- ① Load predefined symbols into a symbol table
- ② Read input line-by-line
 - ↳ expand macros
 - ↳ evaluate expressions

Pseudoinstructions

- convenience macros that expand to one or more real instructions
- extend set of operations without adding to the ISA

USAM Expressions and Layout

- values can be written as expressions
- ↳ assembler evaluates expressions, they are not translated to instructions to compute the value

Summary of Assembly Language

- low-level language, symbolic rep. of seq. of bytes
- ↳ Abstracts: - bit level rep. of instructions.
- address

Universality

- **Universal**: a set of Boolean gates is universal if we can implement any Boolean function using only gates from that set

Turing Machines

- FSM + infinite digital tape that could be read and written at each step
- ↳ encode input as symbols on tape
- ↳ FSM reads tape / writes symbols / changes state until it halts
- ↳ answer encoded on tape
- TMs as integer functions: $y = TM_z[x]$

Church's Thesis

- Every discrete function computable by any realizable machine is computable by some Turing Machine:
- $$f(x) \text{ computable} \iff \text{for some } U, \text{ all } x \quad f(x) = T_U[x]$$

Interpretation

- manipulate coded representations of computing machines, rather than the machines themselves

Turing Universality

- can emulate every other Turing machine
- to show that your computer is Turing Universal, some known UTM demonstrate it can emulate

* Composition is Key

Uncomputability

- no algorithm can compute $f_U(x)$ for arbitrary x
- in finite steps
- e.g.: Halting function:

$$f_U(U, i) = \begin{cases} 1 & \text{if } T_U[i] \text{ halts;} \\ 0 & \text{otherwise} \end{cases}$$

↳ determines whether the U^{th} TM halts when given a tape containing i .