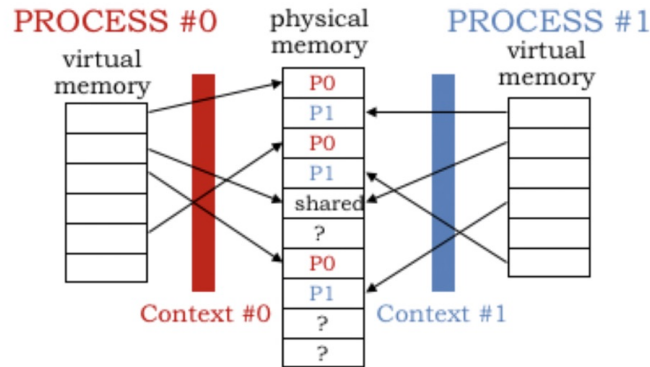


CL 17



Building a Virtual Machine (VM)



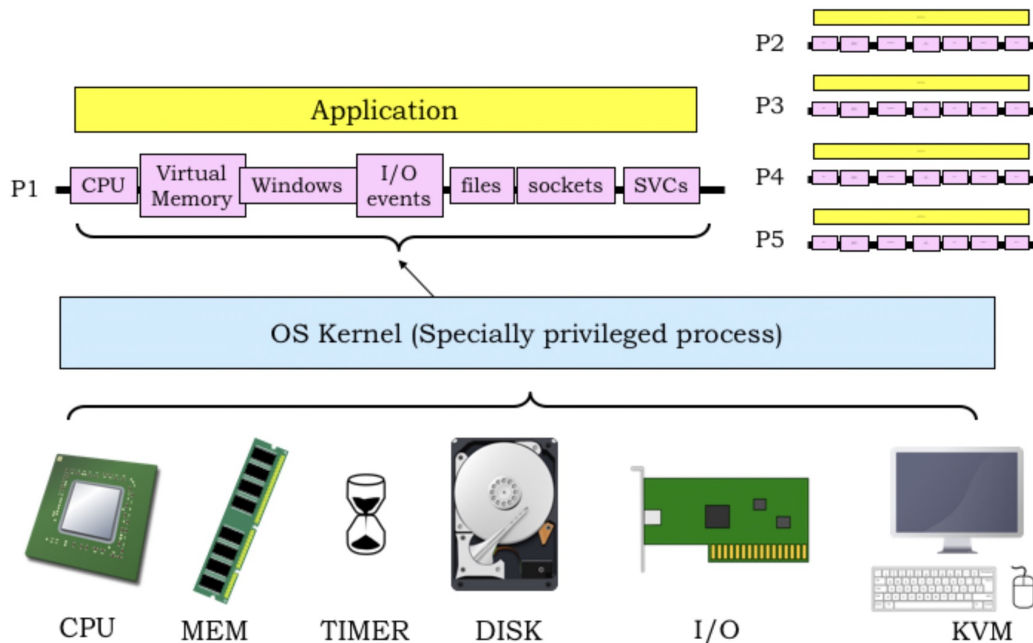
Goal: give each program its own “VIRTUAL MACHINE”;
programs don’t “know” about each other...

New abstraction: a process which has its own

- machine state: R0, ..., R30
- context (virtual address space)
- PC, stack
- program (w/ shared code)
- virtual I/O devices

“OS Kernel” is a special, privileged process running in its own context. It manages the execution of other processes and handles real I/O devices, emulating virtual I/O devices for each process.

One VM For Each Process



Communicating with the OS

User-mode programs need to communicate with OS code:

Access virtual I/O devices

Communicate with other processes

...



*But if OS Kernel is in
another context (ie, not in
user-mode address space)
how do we get to it?*

Solution:

Abstraction: a supervisor call (SVC) with args in registers –
result in R0 or maybe user-mode memory

Implementation:

use *illegal instructions* to cause an exception --
OS code will recognize these particular illegal
instructions as a user-mode SVCs

*Okay...
show me
how it
works!*

