

Engine Knock Detection Using Spectral Analysis Techniques With a TMS320 DSP

Application Report



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ABSTRACT

An efficient method of detecting engine knock using spectral analysis is presented. Multiple single-point DFTs are used to condition the measured knock signal. Using multiple frequencies in the detection algorithm provides a better signature of the combustion process and enhances the ability to detect low-level knock across the entire operating range of the engine. The detection strategy compares the DFT outputs to a variable reference to determine a knock intensity metric. Unlike currently used techniques, the algorithm adapts the reference (no-knock condition) to varying engine speeds and loads. An overview of the knock detection problem and current technology is presented. Implementation examples are included to aid in developing system-specific hardware and software. Two systems are presented. The first is based on the TMS320C30 floating-point DSP with software written in C. This system could be used for engine study and algorithm development. The second, which is more production oriented, is based on a TMS320C25 fixed-point DSP with software written in assembly language.

INTRODUCTION

What Is Engine Knock?

Modern engine control systems are designed to minimize exhaust emissions while maximizing power and fuel economy. The ability to maximize power and fuel economy by optimizing spark timing for a given air/fuel ratio is limited by engine knock. Detecting knock and controlling ignition timing to allow an engine to run at the knock threshold provides the best power and fuel economy. Normal combustion occurs when a gaseous mixture of air and fuel is ignited by the spark plug and burns smoothly from the point of ignition to the cylinder walls. Engine knock, or detonation, occurs when the temperature or pressure in the unburned air/fuel mixture (end gases) exceeds a critical level, causing autoignition of the end gases. This produces a shock wave that generates a rapid increase in cylinder pressure. The impulse caused by the shock wave excites a resonance in the cylinder at a characteristic frequency that is dependent primarily on cylinder bore diameter and combustion chamber temperature. Damage to pistons, rings, and exhaust valves can result if sustained heavy knock occurs. Additionally, most automotive customers find the sound of heavy engine knock objectionable.

Knock Sensors

Implementing a knock detection/control strategy requires sensors to monitor the combustion process and provide feedback to the engine controller. Knock sensors can be classified in two broad categories: direct and remote measurements.

Direct Measurements

Pressure sensors measure the pressure inside the combustion chamber of a running engine. This direct measurement of the combustion process provides the best signal to analyze to detect engine knock. However, each cylinder requires its own sensor, and individual sensor costs are still relatively high. As a result, pressure sensors are used primarily in research settings. Currently, Toyota is the only manufacturer that installs pressure sensors in production engines. Pressure sensor usage will increase in the future as sensor costs are reduced and automotive companies develop more sophisticated engine control strategies that monitor the combustion process.

Remote Measurements

Remote measurement sensors use vibrations transmitted through the structure of the engine to detect knock in the combustion chamber. The signal received by remote sensors can be contaminated by sources other

than engine knock, which increases the difficulty of signal detection. This is especially true at higher engine speeds in which background mechanical vibrations are much higher, effectively reducing the signal-to-noise ratio. One advantage of using remote sensors is that, with careful placement, only one or two sensors are required to monitor all cylinders. In addition, the sensors are less expensive, primarily due to a less harsh operating environment.

Two types of remote sensor are being used today: tuned and broadband. Tuned or resonant sensors are used in many low-end knock detection systems. Either mechanically or electronically, the sensor amplifies the magnitude of the signal in the frequency range of the knock-excited resonance (sometimes called the fundamental frequency). A limitation to this approach is that a different sensor can be required for each engine type, due to variations in the characteristic frequency. The resulting part number proliferation increases overall system costs for the manufacturer. To eliminate the cost penalty, sensor bandwidth can be made wide enough to encompass all expected variations in the fundamental frequency. However, doing so can possibly decrease system performance.

Broadband sensors have no resonant peaks below the 20-kHz operating range of the knock-detection system. One sensor works equally well for any engine configuration. Some type of postprocessing is required to identify the characteristic frequency, placing an additional burden on the signal conditioning part of the system. Since variations in the fundamental frequency can be expected for different engine configurations, a programmable solution provides the flexibility to easily modify the frequency range being monitored with minimal impact on system cost.

Knock Detection Overview

Spectral Signature

When engine knock occurs, a shock wave is generated inside the combustion chamber. The shock wave excites a characteristic frequency in the engine, which is typically in the 5 kHz–7 kHz range. Cylinder bore diameter and combustion chamber temperature are the main variables that affect this fundamental frequency. Variations in the fundamental frequency for a given engine configuration can be as much as \pm 400 Hz. Larger diameters and/or lower temperatures result in a lower fundamental frequency.

Signals received by a remote sensor contain additional vibrational modes, which are structural resonances in the engine excited by the shock wave as it hits the cylinder wall. Typically, two to four additional frequency peaks are evident between the fundamental frequency and 20 kHz. Each engine structure can have different higher vibrational modes. Sensor mounting location can affect which modes are detectable and the amplitude of each with respect to the background mechanical noise.

Adaptation Requirements

An engine-knock detection algorithm must be able to adapt to a number of variables to enable the controller to generate optimum spark timing so that the engine can run at the knock threshold. As mentioned previously, the structural design of an engine and the mounting location of the knock sensor(s) affect which frequency modes are detectable by the sensor. <u>Usually, the transfer function between the cylinder and the sensor is different for each cylinder.</u> This causes both the relative and absolute magnitudes of the vibrational modes to be different for each cylinder. A good detection scheme should allow different calibrations for each cylinder.

Another variable that must be accounted for is changes in nonknocking (reference) signal amplitude due to the mechanical vibration of the engine at different RPMs. As the engine speed increases, the background vibration level increases. When a fixed reference is used, a compromise in performance must be made

because signal magnitudes that would indicate knock at lower engine RPMs are equal to or less than the background level at higher engine RPMs. The reference must be set low enough that knock can be detected at lower RPMs, which limits the algorithm's ability to function at higher speeds. For this reason, some knock detection systems are shut off above 4000 RPM, and very conservative spark timings are used to guarantee that knock will not occur. A good detection strategy should adapt to varying levels of background vibration levels to allow trace knock to be detected at all engine speeds.

Finally, an engine's operating characteristics change with time. As an engine wears, tolerances between components change, which could change the magnitudes of the vibrational modes detected by a remote sensor. Normal background vibrational levels could be higher for a given engine speed. The signal-to-noise ratio could decrease at higher engine speeds. A good detection strategy should adjust to changes in daily operating characteristics to allow reliable identification of trace knock without false triggers.

Signal Conditioning

Knock detection systems must perform some type of signal conditioning prior to executing the detection strategy. Information about the signal strength in the frequency range(s) excited by knock must be extracted from the measurement. If a tuned sensor with a very narrow resonant peak about the fundamental frequency is being used, no further signal conditioning is required. In all other situations, either a filtering technique (analog or digital) or a spectral estimation technique must be used.

Analog filtering is the predominant method used today, due to its low cost, ease of implementation, and lack of computational power of the engine controller CPU. The output of a simple analog filter tuned to the fundamental knock frequency is integrated and sent to the engine control unit (ECU) to execute the detection strategy. However, now that higher precision and/or multiple frequency ranges are desired, an analog implementation is becoming cost prohibitive.

Digital filters are starting to become practical as the computational performance of the ECU increases. Programmability allows the same hardware to be shared across a number of engine configurations. Reduction in part numbers can provide big cost savings to manufacturers in all steps in a product's life cycle. Enhancing filter performance or adding additional frequency ranges can be readily accomplished as long as the computational limits of the CPU are not exceeded.

Another digital signal conditioning technique is spectral analysis; for example, Fast Fourier Transform (FFT). An FFT provides a higher level of frequency resolving power than a digital filter. In addition, multiple frequency ranges are available as the basic output of the FFT. Limited computational throughput of the ECU and unfamiliarity with the technique have limited its use to research and development. No current production system uses spectral analysis techniques. The advent of cost-effective digital signal processors like TI's TMS320 fixed-point family is making the computational power available to bring spectral analysis techniques to prominence.

Detection Strategies

Knock detection strategies use the output of a signal conditioning stage to compare with a reference to determine the presence or absence of knock. Most systems today use windowing to isolate periods during the cylinder's firing cycle for analysis when knock is possible. There is a window from approximately 10° to 70° after top dead center (ATDC) of the piston's cycle when detonation is most likely to occur for the firing cylinder. The detection algorithm is run only during this window for that cylinder. By eliminating possible false trigger sources, such as valve closing, the detection algorithm is more robust. The time this window is active varies with engine speed from 20 ms at 500 RPM down to 1.25 ms at 8000 RPM. Tracking changing engine speed to calculate this time variation requires hardware or software overhead for implementation.

An indication of signal strength during the active window period is typically used by the detection algorithm. Integrating the output of a filter is a common method. The magnitude of the signal strength is compared with a reference level, which, in nonadaptive systems, must be predetermined during system development. If the reference level cannot adapt to changes in engine RPM, a compromise must be made. The reference level must be low enough to prevent sustained knock at low speed, yet high enough to prevent false triggers at higher engine speed. Today, at engine speeds above approximately 4000 RPM, a combination of very conservative spark timing maps and shutting off the control strategy is used to guarantee knock-free operation. This results in less than optimal performance and fuel efficiency at higher engine speeds, particularly for systems using only fundamental frequency detection. Even at lower engine speeds, some compromise is required to guarantee that knock is likely occur only during transient operation.

The tradeoffs between using only the fundamental frequency or the combination of fundamental and vibration mode frequencies concern the issues of false triggers vs. complexity, available CPU time, and cost. When multiple frequencies are used, a better signature is available to determine if knock is present. This effectively increases the signal-to-noise ratio of the system. As a result, either the RPM range for reliable detection can be extended or the baseline spark timing at lower engine speeds can be advanced to allow the engine to continuously run closer to the knock threshold.

Control Strategies

Knock control strategies today adjust spark timing to let an engine run at the knock threshold. Look-up tables are used to obtain a baseline setting for a given speed, load, and temperature. Based on the level of knock detected, timing can be advanced (no knock) or retarded (knock). The rate of advance or retardation can also be modified based on knock magnitude and/or offset from the baseline spark timing setting.

The strategies fall into two categories. The simplest strategy, global control, retards the spark timing of all cylinders by the same amount when any knock is detected. This approach has the advantage that only one knock value has to be tracked and only one timing control loop needs to be executed. The computational burden on the CPU is minimized, as are the memory requirements for both data and program. However, engine performance can be compromised if only one or two cylinders are more likely to knock at a given operating condition. Since all cylinders' spark timing is retarded equally, the cylinders not at their knock threshold are not providing optimal power and/or fuel efficiency.

A more sophisticated strategy is to control each cylinder individually so that all cylinders are running at the knock threshold and provide the best power and fuel efficiency. To implement this type of control strategy, knock computations and control updates must be performed individually for each cylinder. The computational burden and memory requirements are higher than with the global control strategy.

Both control strategies are used today. As automotive manufacturers attempt to improve the emissions, power, and fuel efficiency of their engines to meet the competition or government regulations, the individual cylinder control strategy will become predominant. Advanced engine control strategies are being developed that use torque and/or combustion feedback to optimize the operation of each cylinder independently for best spark timing, fuel delivery, and possibly valve timing. When these strategies are implemented, global control will no longer be a viable alternative.

DFT-BASED DETECTION METHOD

An individual cylinder knock detection method has been developed that offers superior performance compared with analog bandpass filtering, digital bandpass filtering, or FFT techniques. This new method uses multiple single-point Discrete Fourier Transforms (DFTs) for the signal conditioning step to monitor

the fundamental frequency plus the vibrational modes of an engine. The DFT algorithm provides better frequency discrimination than low-cost analog filters, provides better frequency discrimination and/or less computational burden on the CPU than digital filters, and is less computationally intensive than an FFT. In addition, because the computation of the DFT algorithm can be spread over all the samples in a block of data, it leaves more time to run the detection algorithm before the next cylinder's data must be processed.

The DFT algorithm can be cost effectively implemented on a TI fixed-point DSP, and that system can replace currently used lower-performance knock detection systems. As a result of the frequency discrimination capability of the DFT and the use of multiple frequency ranges, the effective signal-to-noise ratio is enhanced. In addition, the computational throughput of the DSP allows the reference signal to adapt in real time to changing operating conditions, such as engine speed. At lower speeds, spark timing can be advanced closer to the knock threshold because the reference is adjusted to the lower background vibration levels. It is also possible to resolve a knock signature at higher engine speeds, due to both the multiple frequency ranges, which give a more distinctive signal, and the adaptive reference, which reduces false triggers. This ability allows the variable spark timing of the engine to run closer to the knock threshold to obtain improved power and fuel efficiency.

The following sections present the details of the DFT algorithm, detection strategies, and an adaptation algorithm for engine speed. Two hardware and software implementation methods for TI TMS320 DSPs are discussed. One is a high-level-language version for the TMS320C30 written completely in C, which could be used for research and development. The other is an assembly-language version written for the TMS320C25, which could be used in a production system. Finally, an integration strategy is described that shows a road map, starting with the replacement of an existing analog knock detection system interfaced to an engine ECU with a DSP-based system. The strategy leads to a single-processor solution with increased features and performance.

Signal Conditioning

A <u>lower-cost broadband knock sensor is used</u> in place of one tuned only to the fundamental frequency to allow the DFT signal-conditioning algorithm to monitor multiple frequencies. The signal-conditioning algorithm starts with the basic definition of the single-point direct-form DFT shown in equation (1).

$$X[k] = \sum_{n=0}^{N-1} x[n] \cdot e^{-j2\pi k[n]/N}$$

where: k = frequency index = sample time ind

n = sample time ind N = sample block length

X[k] = amplitude of sinusoid at frequency index k x[n] = time domain representation of a signal

Three variables determine the effective center frequency and bandwidth (frequency discrimination) of a single-point DFT: <u>sample rate</u>, <u>block size</u>, <u>and indexed frequency</u>. The relationships are:

Nyquist rate : $f_N = f_s/2$ Minimum frequency : $f_{min} = f_N/N$ Indexed frequency : $f[k] = k \cdot f_{min}$ k: 0, ..., N-1

The Nyquist rate (f_N) is one-half the sample rate (f_s) , which determines the highest frequency that can be discriminated without aliasing. For instance, the most significant knock frequencies are typically between

5 kHz-20 kHz; the sample rate for the knock detection system needs to be at least 40 kHz to prevent aliasing. This results in a loop time for the signal-conditioning algorithm of no more than 50 μs. A high computational throughput in the CPU is required to calculate any sophisticated algorithms in such a short period of time.

The minimum frequency (f_{min}) that can be discriminated is determined by the sample rate and the block size. For a given sample rate, the larger the block size, the greater the frequency discrimination of the DFT. Frequency discrimination can be considered similar to the bandwidth of a filter.

The specific frequencies that can be monitored are determined by the indexed frequency (f[k]) once the sample rate and block size are set. One of the advantages of a DFT over an FFT is that the block size can be any length and is not limited to a power of 2. This makes the frequency ranges that can be set using a DFT more flexible than those set using an FFT. A digital filter can be more adjustable than the DFT, but it carries the penalty of higher computational overhead and a larger memory requirement than an equivalent single-point DFT.

The general expression of the DFT allows the input signal x[n] to be complex, but in sampled data systems like this one for knock detection, the input signal is real. This allows a simplification of the final expression. In addition, digital processors cannot easily represent complex exponentials, so a different representation of equation (1) is required. Euler's relationship between complex exponentials and sinusoids is the starting point for developing a set of expressions that can be implemented on a DSP:

$$e^{-j\Theta} = \cos\Theta - j\sin\Theta \tag{2}$$

Substituting (2) into (1) and allowing the input signal to be only a real value yields:

$$X[k] = \sum_{n=0}^{N-1} x[n] \cdot \left(\cos(2\pi [n]k/N) - j\sin(2\pi [n]k/N) \right)$$
 (3)

Expanding the expression into the real and imaginary parts of the computation gives:

$$X[k] = \sum_{n=0}^{N-1} x[n] \cdot \cos(2\pi [n]k/N) - j \sum_{n=0}^{N-1} x[n] \cdot \sin(2\pi [n]k/N)$$
 (4)

Finally, separating the real and imaginary parts of the computation into individual equations gives forms that can be easily programmed on a digital processor:

$$X_{R}[k] = \sum_{n=0}^{N-1} x[n] \cdot \cos(2\pi[n]k/N)$$
 (5a)

$$X_{l}[k] = \sum_{n=0}^{N-1} x[n] \cdot \sin(2\pi[n]k/N)$$
 (5b)

To minimize the computational burden on the DSP, the calculations are spread out over the entire sample block. Each time a sample is received, one step in the summation is computed, which is an advantage of

DFT over FFT implementation. All samples for an FFT must be stored in memory before the calculations can be performed. The DFT can be calculated one sample at a time because there is no linkage between samples as there is with the FFT. Even though the FFT is more efficient in total time required when calculating across all allowable frequency ranges, the DFT makes superior uses of CPU resources when only a few of the possible frequency ranges are being monitored. The actual form of the difference equation implemented on the DSP is:

$$X_R[k] = X_R[k] + x[n] \cdot \cos(2\pi[n]k/N)$$
 (6a)

$$X_{l}[k] = X_{l}[k] + x[n] \cdot \sin(2\pi [n]k/N)$$
 (6b)

These calculations are repeated once every sample period and require two multiplies and two adds per frequency range monitored, regardless of the sample block length used. Note that only a single sample must be stored to implement this algorithm. This is one of the advantages of the DFT over a digital filter, such as a finite impulse response (FIR) filter. Computing one N-tap FIR filter requires that N samples and N coefficients be stored in memory. N multiplies and N accumulates are required each sample period. The number of frequency ranges or the frequency discrimination in each range (the number of filter taps) is much more limited for a given CPU if a filter-based signal-conditioning technique is used, due to the increased computational burden.

The cosine and sine values used in the DFT computation are best implemented as look-up tables for this application. Computing the values requires too much time and the memory required to store the tables is not excessive for the block sizes used for knock detection. To further minimize memory requirements, a single table can be used with two pointers: one for cosine values and another for sine values. The table is constructed to contain N elements that define one period of the lowest frequency (f_{min}) sine wave. Using this single table, any of the allowable frequencies can be monitored. This is another advantage of DFT over digital filter implementation. A digital filter has a unique set of coefficients for each frequency monitored, while the single-point DFT has the "coefficients" as subsets of one N-length table of values for all frequencies.

When the sample block summation is completed, the magnitude of the signal strength for each frequency range must be computed. The standard equation for computing magnitude is:

$$X[k] = \sqrt{X_R[k]^2 + X_I[k]^2}$$
 (7)

The detection algorithm develops the knock metric by comparing the current signal strength to the reference. To simplify the calculations, the square root is not computed, because it does not provide additional information. Using only the sum of the squares amplifies differences in magnitude, because the square root function tends to compress the number range. The equation used to compute the relative magnitude is:

$$X[k]^{2} = X_{R}[k]^{2} + X_{I}[k]^{2}$$
(8)

The DFT signal-conditioning algorithm should run for as much of the variable time period from 10° to 70° ATDC (the knock window) as possible. To make the signal strength calculations consistent, a single block

size is used for all engine speeds. The highest engine speed determines the maximum block size that can be used for a given sample rate. At lower engine speeds, multiple sample blocks are required to span the knock window. When more than one sample block is used, an outer loop around the core DFT algorithm is required to select the sample block with the highest magnitude signals. A high-level pseudocode description of an algorithm to monitor five frequency ranges over multiple sample blocks is shown below:

```
for (i = 0; i < BLOCKS, i ++)
      for(n = 0; n \le N - 1; n ++)
TEST
                if (NO SAMPLE), then go to TEST
                x0 = SAMPLE
               X_R[k1] = X_R[k1] + x0 \cdot \cos(2\pi k1[n]/N)
               X_{I}[k1] = X_{I}[k1] + x0 \cdot \sin(2\pi k1[n]/N)
               X_R[k5] = X_R[k5] + x0 \cdot \cos(2\pi k5[n]/N)
               X_1[k5] = X_1[k5] + x0 \cdot \sin(2\pi k5[n]/N)
           X[k1]^2 = X_R[k1]^2 + X_I[k1]^2
           X[k5]^2 = X_R[k5]^2 + X_I[k5]^2
             X_{sum} = X[k1]^2 + \cdots + X[k5]^2
      if (X_{sum} > X_{max}), then
      {
            X_{sum} = X_{max}
            X[k1]_{max} = X[k1]^2
            X[k5]_{\text{max}} = X[k5]^2
      }
```

When the signal conditioning calculations are complete, the signal magnitudes $X[k]_{max}$ are passed to the detection algorithm.

More complex DFT algorithms can be implemented that use a running magnitude computation technique rather than the sequential block technique shown above. The sliding-mode DFT and the Goertzel algorithm

are two such techniques. Advantages of these algorithms are that they can be used to determine exactly where, during the ignition cycle, the signal reaches the maximum amplitude. This information can be used by the detection strategy to more accurately determine the knock severity. The sooner after spark ignition that engine knock occurs, the more severe the knock and the more aggressive the control response should be. The disadvantages of these techniques are that the computational overhead is slightly higher and the data memory requirements increase, because a sample-block-length number of samples (N) must be retained for the calculations. Only the sequential block technique is discussed in this paper. References [4] and [5] contain details of the more complex methods.

Detection Strategy

The detection strategy uses the output of the DFT-based signal conditioning software module to develop the knock intensity metric that will be used by the control strategy. There are several knock detection strategies that have varying levels of sophistication and computational overhead. A few examples are given here to demonstrate the range of complexity, starting with the simplest and increasing in complexity. The strategy of choice is dependent on the specific engine configurations and control strategy implementations.

The simplest strategy is to use only the DFT-generated frequency magnitudes for the detection algorithm. A more complicated approach is to include the rate of change of the magnitudes, which can be used in predictive algorithms. A more sophisticated approach is to implement a fuzzy logic control strategy. Since the combustion process is highly variable from cycle to cycle, statistical methods can be used to compute trends toward the knock threshold. The magnitude-only approach is presented in more detail.

Magnitude detection strategies can use a number of techniques. One is to count the number of frequencies from the current measurement that exceed the reference, then use the count to indicate knock intensity. This is the simplest technique and is a good starting point when first implementing the DFT-based signal conditioning algorithm. A high-level pseudocode description is shown below:

$$\begin{split} & \text{if } (X[k1]_{max} > X[k1]_{ref}), \text{ then counter } ++\\ & \vdots \\ & \vdots \\ & \text{if } (X[k5]_{max} > X[k5]_{ref}), \text{ then counter } ++\\ & KNOCK = \text{counter} \end{split}$$

Another approach is to weight the count according to which frequency range exceeds the reference. The weighting can also be varied with engine RPM. For instance, the fundamental frequency and one of the vibrational modes could carry a higher weight at lower engine speed and a lower weight at higher speed. In this example, it is assumed that the importance of the frequency ranges is assigned as follows:

Priority	Variable	Frequency	Name
5	range1	7250 Hz	Fundamental
2	range2	9475 Hz	Vibrational mode #1
4	range3	12325 Hz	Vibrational mode #2
3	range4	16000 Hz	Vibrational mode #3
1	range5	18775 Hz	Vibrational mode #4

A simple weighting scheme is shown below in which the fundamental frequency and vibrational mode #2 have weights that vary with engine speed. The pseudocode shows the initialization and the determination of the knock intensity metric.

The amount by which the frequency magnitude exceeds the reference can be used as additional information for determining a knock intensity metric and can be integrated into the weighting factor. This metric can be used to fine-tune the algorithm for detecting trace knock conditions. It requires additional computational resources to implement. The value of $\Delta[k]$ can be different for each frequency being monitored. An example using pseudocode is shown below. Here, the weights vary for both engine RPM and variation between the measurement and the reference value.

```
// Initialization of weights
if (RPM = LO), then
                                                             //06000h
      range1_{LO} = 0.110\,000\,000\,000\,000b
      range1_{HI} = 0\ 111\ 000\ 000\ 000\ 000b
                                                             //07000h
      range3_{LO} = 0\ 000\ 110\ 000\ 000\ 000b
                                                             //00C00h
      range3_{HI} = 0\ 000\ 111\ 000\ 000\ 000b
                                                             //00E00h
}
else
      range1_{LO} = 0.010.000.000.000.000b
                                                             //02000h
      range1_{HI} = 0.011.000.000.000.000b
                                                             //03000h
      range3_{LO} = 0\ 000\ 010\ 000\ 000\ 000b
                                                             //00400h
      range3_{HI} = 0\ 000\ 011\ 000\ 000\ 000b
                                                             //00600h
      range2_{LO} = 0\ 000\ 000\ 000\ 110\ 000b
                                                             //00030h
```

```
//00038h
       range2_{HI} = 0\ 000\ 000\ 000\ 111\ 000b
       range4_{LO} = 0\ 000\ 100\ 000\ 000\ 000b
                                                                      //00800h
       range4_{HI} = 0\ 000\ 101\ 000\ 000\ 000b
                                                                      //00A00h
       range5_{LO} = 0\ 000\ 000\ 000\ 000\ 010b
                                                                      //00002h
       range5_{HI} = 0\ 000\ 000\ 000\ 000\ 011b
                                                                      //00003h
// Compute reference to measurement difference
\Delta[k1] = X[k1]_{max} - X[k1]_{ref}
\Delta[k5] = X[k5]_{max} - X[k5]_{ref}
// Calculate knock intensity metric
counter = 0
if (\Delta[k1] > \Delta[k1]_{HI}), then counter+ = range 1<sub>HI</sub>
else if (\Delta[k1] > \Delta[k1]_{LO}), then counter+ = range 1_{LO}
if (\Delta[k1] > \Delta[k5]_{HI}), then counter+ = range5<sub>HI</sub>
else if (\Delta[k1] > \Delta[k5]_{LO}), then counter+ = range5<sub>LO</sub>
```

In some instances, it may be desirable to predict when knock is about to occur and begin to make adjustments to the control strategy preemptively. The rate of change of the amplitude can be used for this prediction function, and this predictor can be implemented as a simple difference between the current and some previous measurement. The greater the time difference, the less effect cycle-by-cycle variations will have, but the more effect changing the engine speed will have on the result. This technique takes more computational resources and increases memory requirements, because previous values must be retained for the calculation. Normalizing the average rate of change over all the ranges can yield the weighting factor that would be used to adjust the knock intensity metric. The pseudocode for an 8-sample time difference example is shown below. A two-dimensional array is used to store the current and previously measured amplitudes for the frequencies being monitored.

```
// Compute rate of change \Delta_{rate} \ [k1] = \Delta [k1, \, m]_{max} - \Delta [k1, \, m-8]_{max}
```

KNOCK = counter

.

$$\Delta_{\text{rate}} [k5] = \Delta[k5, m]_{\text{max}} - \Delta[k5, m - 8]_{\text{max}}$$

// Normalize the rate of change over frequencies $//(\Delta_{max}$ is largest possible difference)

$$\Delta_{norm} = 1 / \Delta_{max}$$

$$\Delta_{\text{ave}} = 0.2 \cdot \sum_{n=1}^{5} \Delta_{\text{norm}} \cdot \Delta_{\text{rate}} \text{ [kn]}$$

// Adjust knock metric for amplitude rate of change

$$if \ (\Delta_{rate} > 0), \ then \\ KNOCK = \Delta_{ave} \cdot KNOCK$$

Using measurements too close together in time in highly variable systems can cause the control response to become unstable. However, with a carefully designed rate computation, the additional information provided can allow better control during transient operating conditions when speed and/or load are changing rapidly. This technique can even be extended to the rate of change of the knock intensity metric itself.

Adding sophistication to the detection strategy increases the computational burden on the CPU and can provide only a marginal increase in algorithm performance. Increasing sophistication also requires that the system designer have a high level of understanding of the knock process in the specific engine family on which the knock detection system is being implemented. The implementation examples given in a following section demonstrate one of the simple techniques that make it widely applicable.

Adaptation Strategy

The ability to adapt to changing engine operating conditions, primarily speed and load, greatly improves the performance of a knock detection system. Less compromise is required in the baseline spark timing, because the system is able to more accurately determine low-level knock at all engine speeds. Using a TI DSP provides the computational throughput required to execute both the knock detection and adaptation calculations in real time over the entire speed range of today's passenger car engines.

The adaptation technique used in this application report generates the reference signal from a sliding weighted-average computation. A weighted mean is computed individually for each cylinder and is updated for each firing cycle. The general form of the weighted mean is:

$$Y[kn] \, = \, \sum_{m=0}^{N-1} \, \, \frac{x[k(n\, - \, m)]}{N}$$

$$R[kn] = K \cdot Y[kn]$$

Where: Y[kn] = the mean value of a time history of inputs

x[k(n - m)] = the time history of inputs, where each input is a spectral magnitude

= the number of elements in the time history

R[kn] = the weighted mean value = the weighting coefficient

However, this method for computing Y[kn] requires N multiply/accumulates for each average value. A more effective computational technique is to add the newest input to the running average and subtract the

oldest value. This method requires two multiply/accumulates to compute a new average regardless of the number of values in the time history. The memory storage requirement increases by 1 (to N+1 values) due to the need to store the previous average value. This computational method has the form:

$$Y[kn] \, = \, Y[k(n\, -\, 1)] \, + \frac{x[kn]}{N} \, - \frac{x[k(n\, -\, N)]}{N}$$

The weighting factor (K) is used in the computation of R[kn] to adjust the reference value so that cycle-to-cycle variation does not falsely trigger the detection algorithm. Empirical studies put the value for K at 1.5–2. These values can be different for each cylinder on the same engine due to the effects of sensor placement and structural dynamics. In addition, the weighting coefficient could be modified in real time to track changes in engine speed and/or load if cycle-to-cycle variations change. However, at the current level of control strategy sophistication, the value of investing the additional computational time is probably marginal. Future systems may require statistical analysis of the combustion signature to adjust for cycle-to-cycle variation, but the system implementation described in this report does not use them.

A number of conditions must be taken into account when choosing the averaging period. The averaging period must be long enough that the effects of cycle-to-cycle variations are minimized, but short enough that changing engine speed or load are not filtered out. Values in the range of 8–16 previous samples have been used in development systems with considerable success.

Another consideration when implementing an adaptation strategy is how to handle reference updates in the presence of engine knock. If the reference is updated when knock is occurring, the average value is contaminated by the higher magnitude signals generated by the knock. Assuming that corrective action by the ECU eliminates knock within a small number of firing cycles, a simple method of handling adaptation in the presence of knock can be used: do not update the reference value when knock is detected. This method is practical, because knock is eliminated and the reference updates can restart before any significant changes develop in the operating condition of the engine. The technique described in this section uses this assumption.

For this adaptation strategy to work, there must be a period of time during initialization when nonknocking operation can be guaranteed. Normally, the strategy initializes the adaptation algorithm during the short period of time following engine start when the likelihood of knock is very low. Reinitializing the adaptation algorithm each time the engine is started allows the reference to reflect slowly changing operating conditions.

Implementation Examples

In this section, two implementation examples are presented. The first is a floating-point solution based on the TMS320C30 with all the software written in C. This could be used as a research and development system because of the high computational bandwidth and ease of software development. Due to cost pressures, a production system most likely needs to be a fixed-point solution. The second implementation example is based on the fixed-point TMS320C25 with the software written in assembly language.

While presenting the implementation examples, this report explains the rationale for decisions made on specific structures or variables. This rationale shows how modifications can be made to adapt the concept to new systems. As background for the examples, a discussion on available computational time, sample rate, sample block size, and ADC resolution is presented first, followed by details on the specific target DSPs.

Timing Considerations

The amount of time available for the various stages of the algorithm can affect implementation decisions. The general form of the real-time code for the knock detection software shows the time slots required for the given modules as a function of crankshaft angle. Figure 1 shows the breakdown.

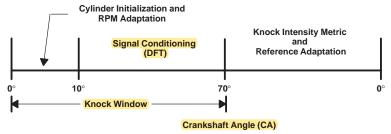


Figure 1. Algorithm Stages Associated With the Crankshaft Angle

A study of the number of cylinders, engine RPM range, sample rate, and sample data block size gives the information required to start making implementation decisions.

Engines today can have operating speeds from 450 RPM at idle up to almost 8000 RPM. Using an execution window comprising 60° of the crankshaft angle (10°–70° CA ATDC) for the monitoring period, the minimum amount of time available to perform the DFT-based signal-conditioning algorithm can be computed. The minimum period occurs when the engine speed reaches its maximum, or 8000 RPM.

$$\frac{min}{8000 \ revs} \cdot \frac{rev}{360^{\circ} \ CA} \cdot \frac{60 \ s}{min} = 20.8\overline{3} \cdot \frac{\mu s}{CA}$$

$$\frac{60^{\circ} \ CA}{knock \ window} \cdot \frac{20.8\overline{3} \ \mu s}{^{\circ}CA} = 1.25 \ \frac{ms}{knock \ window}$$

Even at high engine speeds, there is a significant amount of time available for the signal-conditioning calculations using a TMS320 DSP. The 40-MHz TMS320C30 has time for 25,000 instruction cycles, and the 50-MHz TMS320C25 is able to complete 15,625 cycles.

Next, the time available for the detection and adaptation tasks must be determined. This is the amount of time between 70° ATDC on the monitored cylinder and 0° ATDC on the next cylinder in the firing order, evaluated at maximum engine RPM. This time period is directly affected by the number of cylinders in the engine. Today, the majority of engines use a 4-cycle design, which goes through a complete firing sequence every 720° of crankshaft angle. Most engines have 4, 6, or 8 cylinders, though 3-, 5-, 10-, and 12-cylinder versions exist. Evaluating a number of engine configurations for the time available for the detection and reference adaptation algorithms is shown in the following equations:

$$\left[\frac{720 \circ CA}{firing \ cycle} \times \frac{firing \ cycle}{12 \ cylinders} \right] - \frac{70 \circ CA}{knock \ window} \right] \times \frac{20.8\overline{3} \ \mu s}{\circ CA} = -208.3 \ \mu s$$

$$\left[\left[\frac{720 \circ CA}{firing \ cycle} \times \frac{firing \ cycle}{10 \ cylinders} \right] - \frac{70 \circ CA}{knock \ window} \right] \times \frac{20.8\overline{3} \ \mu s}{\circ CA} = 41.2 \ \mu s$$

$$\left[\left[\frac{720 \circ CA}{firing \ cycle} \times \frac{firing \ cycle}{8 \ cylinders} \right] - \frac{70 \circ CA}{knock \ window} \right] \times \frac{20.8\overline{3} \ \mu s}{\circ CA} = 416.7 \ \mu s$$

$$\left[\left[\frac{720 \circ CA}{firing \ cycle} \times \frac{firing \ cycle}{6 \ cylinders} \right] - \frac{70 \circ CA}{knock \ window} \right] \times \frac{20.8\overline{3} \ \mu s}{\circ CA} = 1.04 \ ms$$

$$\left[\left[\frac{720 \circ CA}{firing \ cycle} \times \frac{firing \ cycle}{6 \ cylinders} \right] - \frac{70 \circ CA}{knock \ window} \right] \times \frac{20.8\overline{3} \ \mu s}{\circ CA} = 2.29 \ ms$$

These calculations show that for 12-cylinder and probably 10-cylinder engines there is computational overlap between the current cylinder's detection and adaptation tasks and the next cylinder's initialization. A method of interleaving the calculations must be implemented for these engine configurations, but such a method exceeds the scope of this report.

To simplify the example, the assumption is made that the knock detection algorithm is being implemented on a 6-cylinder engine. This assures enough time to complete all calculations before the next cylinder's monitoring period begins. The time budget is now established. There are at least 208.3 μ s for the cylinder initialization and RPM adaptation, 1.25 ms for the DFT-based signal-conditioning routine, and 1.04 ms for the detection and reference adaptation routines.

Frequency Selection

Next, the sample rate and block size must be selected. These two parameters determine which center frequencies and bandwidths are possible using the individual single-point DFTs. After defining the fundamental and primary vibrational mode frequencies, the sample rate and block size are varied until a good match is obtained.

A simple C program that searches over a specified range of sample rates and block sizes finds the best matches for a given set of center frequencies. Best match is defined as the combination of sample rate (f_s) and block size (N) that produces the lowest average and lowest difference in standard deviation between the desired and attainable center frequencies. The program listing is contained in Appendix A. When the program runs, it queries the user to define the search parameters and then generates a list that shows the ten best matches.

Tradeoffs on block size, sample rate, and frequency match can then be made, taking into account computation time (the higher the f_s , the less time available between samples) and frequency selectivity (the larger the N, the narrower the effective bandwidth). The effective bandwidth (BW) can be estimated by:

$$BW \approx \frac{f_s}{4 \cdot N}$$

Once the sample rate and block size have been determined, the frequency index numbers that most closely match the desired center frequencies must be determined. The following C code can be used to list the available center frequencies versus index number:

$$\begin{split} f_{min} &= \frac{f_s}{2 \cdot N}; \\ for(k = 0; \ k < N; \ k++) \\ f(k) &= k \cdot f_{min}; \end{split}$$

Choose the values of k that give the closest matches between f(k) and the desired center frequencies.

Analog-to-Digital Converter (ADC) Resolution

The dynamic range of the external sensor must cover several orders of magnitude. Over the operating range of an engine, values can vary from less than 40 mV at low RPM without knock to full scale of the sensor at high RPM and high knock. Full scale ranges typically available are $\pm 10 \text{ V}$, $\pm 5 \text{ V}$, and 0 V - 5 V. To cover this dynamic range, an ADC input system with an effective resolution of 8 bits ($\pm 5 \text{ V}$ and 0 V - 5 V) to 10 bits ($\pm 10 \text{ V}$) is required. Higher resolution is beneficial, because it provides a guard band at either extreme, which can be used for system diagnostics. However, the added resolution increases system cost. Each system must be evaluated to determine the required resolution for that particular implementation.

An 8-bit ADC can be used, even in systems that require higher resolution, by employing external hardware gain switching combined with software control and compensation. A software routine can be incorporated into the input signal-conditioning algorithm to monitor the input signal magnitude and adjust an external gain stage at the input to the ADC. It also must adjust the signal gain to compensate for the external gain changes prior to use in the DFT calculations. The internal compensation is inversely proportional to the external gain adjustment to make the effective input sample dynamic range constant. The implementations presented here assume that the ADC has adequate dynamic range for the sensor that is used, so no gain adjustments are required.

TMS320C30 Implementation

A 'C30-based system is ideal for research and development to develop algorithms and determine analog input system requirements. It is also an excellent platform to use as the sole processor in a high-performance engine control system. The 'C30 is a 32-bit floating-point processor with two on-chip timers and two serial ports. The timers are used to determine engine speed and to trigger execution of the signal-conditioning algorithm. Because of the architectural features and the efficiency of the optimizing C compiler, all the application software can be written in C.

Hardware Description

A high-level block diagram of a 'C30-based system is shown in Figure 2. The system consists of an analog input system (containing analog signal conditioning, a multiplexer, and an A/D converter), the TMS320C30, EPROM, SRAM, processor interface circuitry, and an engine controller (with the assumption that initial tests simply replace an existing knock detection subsystem).

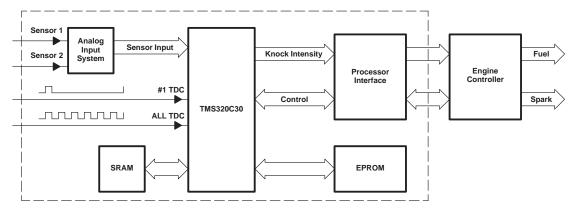


Figure 2. TMS320C30-Based Knock Detection System Block Diagram

Four input signals are required for this configuration: two for knock measurements and two for timing. The two sensor signals are from the engine-mounted accelerometers that monitor the combustion process. Two sensors are used so that one sensor can be put on each bank of a V-configured engine. An in-line engine would require only a single knock sensor. This is standard practice today for the purpose of obtaining a satisfactory signal-to-noise ratio. Single-sensor systems for V-bank engines can also be investigated with this system, but this is not done in this report.

The timing signals are used to compute engine RPM (#1 TDC) and to control when the algorithm runs (ALL TDC) for each cylinder. The #1 TDC signal occurs every 720° of crankshaft rotation. In addition to being used to compute engine RPM, this signal is also used to synchronize the algorithm to the cylinder firing order. The ALL TDC signal occurs every 120° (in the 6-cylinder example) and triggers the algorithm to run for the next cylinder. It is possible to perform both functions using only the #1 TDC signal, but doing so requires more computational overhead.

The memory configuration is set up to allow easy changes to the knock detection system software. Slow EPROM (or EEPROM/Flash) memory is programmed with the current version of the software to be run. During initialization, the real-time code and variables are transferred from EPROM to zero-wait-state SRAM or on-chip RAM.

Software Description

Details of the hardware implementation affect how the initialization and input/output code is written. Therefore, the software in this example is not a complete knock detection system implementation, but is limited to the core algorithm. The assumption is made that no hardware or software input scaling is required since such scaling is also implementation specific. Also, a simple indication of knock intensity is passed to the engine controller in a single word. The algorithm is structured to monitor five frequencies and uses the simple comparison method (number of monitored frequencies exceeding the reference) for detection. Tradeoff decisions between speed and code size in this example favor speed to maximize computational throughput.

Both 'C30 timers are required for this implementation. One is used to determine the engine RPM, which is required to adjust the delay period from top dead center (TDC) to 10° ATDC and the number of sample blocks analyzed in the 10°–70° ATDC knock window. Execution of the code to compute engine RPM computation occurs once per combustion cycle (Cylinder #1 TDC). During the ignition cycles of the remaining cylinders, this code is bypassed. The other timer is used to control when the DFT algorithm begins execution at 10° ATDC.

Figure 3 shows the structure of the program used with the 'C30 floating-point DSP. The first three modules are written in assembler. They are used to define the interrupt vectors, initialize the C environment, and transfer real-time code from EPROM to SRAM or on-chip RAM memory. All remaining modules that pertain to the algorithm are written in C. Except for the initialization module, all functions reside in interrupt service routines (ISRs). An external interrupt triggered by the ALL TDC signal drives the cylinder initialization ISR. This ISR also contains the RPM adaptation code. One of the onboard timers is used to measure the time to rotate through 720° CA, which is directly proportional to RPM. The other timer is used to trigger the signal-conditioning and knock detection ISR to run after the adjustable delay time between 0° to 10°ATDC has elapsed. The comments in the program listing in Appendix B discuss more details of the program's functionality.

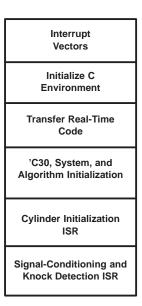


Figure 3. TMS320C30 Floating-Point DSP Knock Detection Program Structure

Test Results

Limited hardware testing has been performed on a 'C30-based platform. Assuming a V6 engine with a maximum speed of 8000 RPM, a 64-sample block size, and 16 frequencies to monitor, initial indications are that sample rates of 75 kHz are achievable. Sample rates below 60 kHz will probably be used when the desired effective bandwidths of the DFTs are taken into consideration. This indicates that a 'C30 with software written in a high-level language like C is an excellent platform for algorithm development and system testing.

TMS320C25 Implementation

To put a DSP-based knock detection system into production, a fixed-point DSP would most likely be used to keep system costs down. The TMS320C25 is the processor chosen for this example. It is the midpoint of TI's fixed-point DSP family. This device is a 16-bit fixed-point processor with a single onboard timer. Depending on the sample rate and the number of frequencies to be monitored in a specific implementation, a member of the TMS320C1x family might be used. To maximize system efficiency, the software is written in assembler, which allows the user to best take advantage of the Harvard architecture used in TI's fixed-point DSP family.

Hardware Description

A high-level block diagram of a possible 'C25-based system is shown in Figure 4. The system consists of an analog input system (containing analog signal conditioning, a multiplexer, and an A/D converter), the 'C25, the engine RPM timer, processor interface circuitry, and an engine controller (with the assumption that the initial implementation retains the existing microprocessor).

The primary differences between the 'C30-based system and the 'C25-based system are twofold. First, no external memory is used, which reduces system cost. Second, a 22–24-bit timer is required to compute engine RPM (down to 100 RPM). The onboard timer on the 'C25 is used to control the DFT algorithm execution (see *Software Description*, page 17, for details). The functionality in the 'C25-based knock detection hardware subsystem can be integrated into a cDSP (customized DSP) device if the specifics of a given implementation warrant doing so.

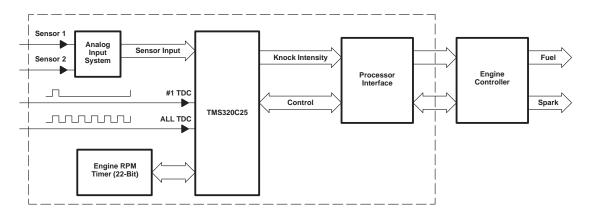


Figure 4. TMS320C25-Based Knock Detection System Block Diagram

Software Description

As in the 'C30 example, the software presented in this example is not a complete knock detection system implementation but is limited to the core algorithm code. The assumptions made for this implementation are the same as those made for the 'C30 example — that no hardware or software input scaling is required, that a simple indication of knock intensity is passed to the engine controller in a single word, that the algorithm monitors five frequencies, and that a simple frequency comparison method is used for detection.

In the case of the 'C25, tradeoff decisions between speed and code size favor code size to reduce memory usage. Specifically, the cosine and sine coefficient tables used to compute the DFTs are not precomputed for the entire sample block. Instead, a single sine wave period one sample block long is used. This saves $(2 \times \text{number of frequencies } -1) \times N$ memory locations, but costs five instruction cycles per frequency monitored.

Figure 5 shows the structure of the program used with the 'C25 fixed-point DSP. All modules shown are written in assembler. An external interrupt triggered by the ALL TDC signal drives the cylinder initialization ISR. The onboard timer triggers the signal-conditioning and knock detection ISR to run after the adjustable delay time from 0° to 10° ATDC has elapsed. The comments in the program listing in Appendix C contain more details of the program's functionality.

Interrupt
Vectors

'C25, System, and
Algorithm Initialization

Cylinder Initialization
ISR

Signal-Conditioning and
Knock
Detection ISR

Figure 5. TMS320C25 Fixed-Point DSP Knock Detection Program Structure

Test Results

The fixed-point implementation is still being developed. Preliminary testing to date has been restricted to simulation. Again, assuming a V6 engine with a maximum speed of 8000 RPM, a 64-sample block size, and 16 frequencies to monitor, a sample rate of approximately 50 kHz should be achievable. Reducing the number of frequencies to ten or fewer should increase the sample rate to over 60 kHz.

The most significant differences in performance between the 'C30 and the 'C25 result from the need to eliminate external memory on the fixed-point DSP system. Therefore, the large interleaved cosine and sine coefficient tables constructed in the 'C30 example are not used. Instead of precalculating all coefficients during initialization, the fixed-point implementation must carry the overhead of coefficient-table pointer manipulation in the real-time code. This adds at least five instruction cycles for each frequency monitored to the real-time code that runs between each sample input. If only three frequencies were monitored, the coefficient-table pointers could be kept permanently in six of the eight auxiliary registers. This would eliminate four of the five additional instruction cycles of overhead penalty for the 'C25, boosting performance significantly.

Integration Road Map

The knock detection system based on a TI DSP can be used as a building block to incorporate advanced capabilities in an engine control system. An implementation road map is shown in Figure 6. The DSP-based knock detection system is initially used as a coprocessor to the current engine control microprocessor. This allows new features to be added cost effectively to the ECU with minimal change to existing hardware. The microprocessor can be retained in the system to perform spark, fuel, idle speed, and evaporative recovery system actuation while the DSP incorporates new, advanced analysis and control strategy computations. If additional cost savings are desired, the microprocessor functions can be integrated into the DSP and the microprocessor can be eliminated from the system.

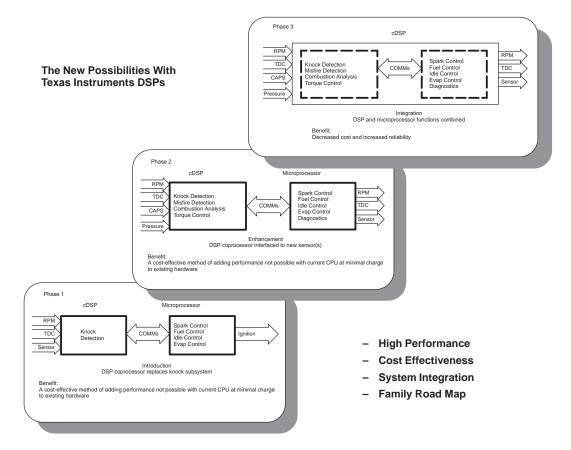


Figure 6. DSP Road Map From Knock Detection Subsystem to Full Engine Control

Texas Instruments has the capability to achieve system integration of digital signal processors, microcontrollers, memory, digital logic, and analog functions on a single chip. The TI integration road map offers customers the benefits of reduced board space, fewer packages, lower power consumption, and reduced total cost of ownership for their unique application.

MISFIRE DETECTION

Misfire detection may be possible using a DFT-based algorithm, which would require only minimal change to the detection part of the system. Misfire is defined as the lack of combustion. When the air/fuel mixture in the firing cylinder does not ignite, a reduction in excitation of the engine structure occurs, and the signal received by an external sensor shows a reduction in signal strength in the fundamental or vibrational mode frequencies. Comparing the signature of a nonfiring cylinder to the computed reference signal may indicate misfire.

The signal-conditioning algorithm using the DFT is the same for knock or misfire detection; only the detection part of the algorithm is different. When misfire occurs, the amplitude of one or more of the monitored frequency ranges is reduced relative to the reference. The change in signal strength when the cylinder misfires must be determined experimentally. Once this is done, the detection algorithm for knock just needs to be modified to trigger when the current signal strength is less than, rather than greater than, the reference magnitude.

Unlike knock detection in which corrective action is taken immediately, California's Onboard Diagnostics (OBDII) regulations are written to detect the occurrence of misfire as a percentage of 200 crankshaft revolutions (for system damage) or a percentage of 1000 revolutions (for emission control performance). This places the additional burden on the system of tracking misfire over a relatively long period before any action is taken.

Using a signal-strength analysis technique like the DFT algorithm described in the preceding sections requires only a small additional computational resource to implement both knock and misfire detection. Two other techniques are widely discussed in the literature. One is accurately measuring crankshaft speed fluctuation during a cylinder's ignition cycle and then using a software model to compute engine torque. Significant computational effort is required to compute the torque using an engine model to correct for torsional and vibrational dynamics. The other method is to obtain direct torque measurements using a special sensor (not yet commercially available), which significantly adds to system cost.

Incorporating a system that is able to detect both knock and misfire provides a cost-effective means of improving performance and fuel economy (knock detection) while complying with government emissions system monitoring regulations (misfire detection).

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Appendix A: Sample-Rate and Block-Size Search Software

```
/* Sample Rate and Block Size Search Routine
      File: FINDFBIN.C
                                         Rev: 1.0
     Start Date: 5/7/93
Written by: Thomas G. Horner
                                         Last Change: 6/9/93
                    Texas Instruments
#include <stdlib.h>
#include <stdio.h>
#include <math.h>
#include <conio.h>
#include <time.h>
           PROTOTYPES
/*
void
       define(void);
       file_open(void);
void
void
       make_space(void);
void
       search(void);
           DEFINITIONS
/*
#define
           list_length 10
         GLOBAL VARIABLES
/*
FILE
            *fptr;
float *delta;
float *fbin;
float *freq_mean;
float *freq_stddev;
float *fs_list;
float *N_list;
float freq[16];
float dif;
float fmin;
float freqn;
float fs;
float fs_max;
float fs_min;
float fs_range;
float fs_step;
float mean;
float std_dev;
float sum;
float var_sum;
int fs_i;
int nfreq;
int N;
int Ni;
int Nmax;
int Nmin;
int Nover2;
int Nrange;
int Nstep;
           MAIN FUNCTION
/*
/* main() is used for overall system control.   
*/   
/* Functions are called which do all the work.   
*/
```

```
void main(void)
   file_open();
   define();
   make_space();
   search();
        FUNCTION DECLARATIONS
/*
       Search Parameter Definition
/* define() is used to define the search
/* parameters by questioning the user.
void define(void)
   int i;
   clrscr();
/**** Define search by questioning user. ****/
   printf("Please answer the following questions to ");
printf("define search parameters.\n");
     /* Define the number of frequencies to
     /* search
   printf("\nHow many frequencies do you want to monitor (1-16)? ");
   scanf("%d",&nfreq);
     /* Define the frequencies
   printf("\nNow enter the frequencies to be \n");
   printf("monitored one at a time.\n");
   printf("Enter a number between 2000-20000.\n");
   for(i=0 ; i<nfreq ;
                         i++)
     printf("\nEnter frequency No.%2d ",i+1);
     scanf("%f",&freq[i]);
     /* Define the block size min/max
   printf("\n\next enter the sample block size \n");
   printf("minimum and maximum values. \n");
   printf("Use a number from 32 to 128.\n");
printf("\nEnter the minimum block size. ");
   scanf("%d", &Nmin);
   printf("\nEnter the maximum block size. ");
   scanf("%d",&Nmax);
     /* Define block size step size
   printf("\n\nNow enter the step size to be used ");
   printf("to go from the minimum");
   printf("\nto the maximum value of block size. ");
   scanf("%d",&Nstep);
     /*Define the sample rate min/max
   printf("\n\nNext you are asked to enter the sample rate ");
   printf("minimum and maximum values.");
   printf("\n Use a number between 40000 and 60000.\n");
   printf("\nEnter the minimum sample rate.
   scanf("%f",&fs_min);
   printf("\nEnter the maximum sample rate. ");
   scanf("%f",&fs_max);
```

```
/* Define sample rate step size
   printf("\n\nNow enter the step size to be used ");
   printf("for the sample rate ");
   printf("sweep\nfrom the minimum to the maximum value. ");
   scanf("%f",&fs_step);
/**** Initialize parameters for search
     /* Compute block size range
   Nrange = Nmax-Nmin;
     /* Compute block size search repeats
   Ni = Nrange/Nstep + 1;
     /* Compute sample rate range
   fs_range = fs_max-fs_min;
     /* Compute sample rate search repeats
   fs_i = (int)(fs_range/fs_step) + 1;
     /* Write search definition to output file */
   fprintf(fptr,"\nNmax = %d \n",Nmax);
fprintf(fptr,"Nmin = %d \n",Nmin);
   fprintf(fptr, "Nrange = %d \n", Nrange);
   fprintf(fptr,"Nstep = %d \n",Nstep);
fprintf(fptr,"Ni = %d \n",Ni);
   fprintf(fptr,"\nfs_max = %5.0f \n",fs_max);
   fprintf(fptr,"fs_min = %5.0f \n",fs_min);
fprintf(fptr,"fs_range = %5.0f \n",fs_range);
fprintf(fptr,"fs_step = %5.0f \n",fs_step);
   fprintf(fptr,"fs_i = %d \n\n",fs_i);
/*
   File Allocation
                         * /
/* file_open() is used to open any files which */
/* might needed during program execution.
void file_open(void)
   if ((fptr = fopen("findfbin.dat", "w")) == NULL)
     printf("error in opening output file #1.");
     exit(0);
   Array Allocation */
/* make_space() is used to allocate memory
/* space for arrays whose length is determined */
/* at runtime. Checks are made to verify that */
/* there is enough memory available. Some of */
/* the arrays are initialized to start values. */
void
       make_space(void)
   int i;
   fbin = (float*)malloc(sizeof(float)*(Nmax/2+1));
   if (fbin==NULL)
     printf("fbin[] allocation error \n");
     exit(0);
   delta = (float*)malloc(sizeof(float)*nfreq);
   if (delta==NULL)
```

```
printf("delta[] allocation error \n");
     exit(0);
   fs_list = (float*)malloc(sizeof(float)*list_length);
   if (fs_list==NULL)
     printf("fs_list[] allocation error \n");
     exit(0);
   for(i=0;i<list_length;i++)</pre>
     fs_list[i]=0.;
   N_list = (float*)malloc(sizeof(float)*list_length);
   if (N_list==NULL)
     printf("N_list[] allocation error \n");
     exit(0);
   for(i=0;i<list_length;i++)</pre>
     N_list[i]=0;
   freq_mean = (float*)malloc(sizeof(float)*list_length);
   if (freq_mean==NULL)
     printf("freq_mean[] allocation error \n");
     exit(0);
   for(i=0;i<list_length;i++)</pre>
     freq_mean[i]=1e+04;
   freq_stddev = (float*)malloc(sizeof(float)*list_length);
   if (freq_stddev==NULL)
     printf("freq_stddev[] allocation error \n");
     exit(0);
   for(i=0;i<list_length;i++)</pre>
     freq_stddev[i]=1e+04;
}
/* Frequency Bin Search */
/* search() carries out a search to find the
/* sample rate and block size that gives
/* center frequencies closest to the user's
/* specifications. The top 10 matches are
/* listed. The average error and deviation
/* are used to determine the best match.
void search(void)
   int i,j,k,l,m,n;
   int TESTING;
/* The frequency bin search strategy is to have */
/* an outer loop for sampling frequency and an
/* inner loop for block size. For a given set
/* of fs and N:
/* 1. Compute the frequency bins.
/* 2. Compare against desired frequencies and
    find deltas for closest matches.
/* 3. Compute delta mean and std. deviation.
/* 4. Compare against top 10 matches and
    update match array if necessary.
```

```
1 = 0;
/* Search over frequency range and
/* block size range.
for(i=0;i<fs_i;i++)</pre>
    fs = (float)(fs_min+i*fs_step);
    for(j=0;j<Ni;j++)</pre>
       1+=1;
       N = Nmin+j*Nstep;
       Nover2 = N/2;
        for(m=0;m<nfreq;m++)</pre>
       delta[m]=1e+10;
    /* Compute min frequency for current
    /* sample rate and block size and
    /* compute frequency bins values.
        fmin = (float)(fs/N);
        for(k=0;k\leq Nover2;k++)
           fbin[k] = k*fmin;
           /\,^{\star} Check if new frequency range less /\,^{\star} than previous smallest difference
           for(m=0;m<nfreq;m++)</pre>
               dif=abs(fbin[k]-freq[m]);
               if(dif<delta[m])</pre>
                   delta[m]=dif;
        /* Compute average and std deviation
                                                     * /
    sum=0.;
    var_sum=0.;
    freqn = (float)nfreq;
    for(m=0;m<nfreq;m++)</pre>
       sum+=delta[m];
    mean=sum/freqn;
    for(m=0;m<nfreq;m++)</pre>
       var_sum+=(delta[m]-mean)*(delta[m]-mean);
    std_dev=sqrt(var_sum/freqn);
        /* Update results if better matches
    for(n=0,TESTING=1;(n<list_length)&&TESTING;n++)</pre>
        if(mean<freq_mean[n] && std_dev<freq_stddev[n])</pre>
           TESTING=0;
            for(m=list_length-2;m>=n;m--)
               fs_list[m+1]=fs_list[m];
               N_list[m+1]=N_list[m];
               freq_mean[m+1]=freq_mean[m];
               freq_stddev[m+1]=freq_stddev[m];
           fs_list[n]=fs;
           N_list[n]=N;
           freq_mean[n]=mean;
           freq_stddev[n]=std_dev;
    }
```

```
/* Write results to output file */
fprintf(fptr,"\n");
for(n=0;n<nfreq;n++)
{
    fprintf(fptr,"freq%2d=%5.0f \n",n+1,freq[n]);
}
fprintf(fptr,"\n");
for(n=0;n<list_length;n++)
{
    fprintf(fptr,"N[%2d]=%4d \t",n,(int)N_list[n]);
    fprintf(fptr,"fs[%2d]=%6.0f \t",n,fs_list[n]);
    fprintf(fptr,"mean[%2d]=%6.0f \t",n,freq_mean[n]);
    fprintf(fptr,"stddev[%2d]=%6.0f \n",n,freq_stddev[n]);
}</pre>
```

Appendix B: TMS320C30 Implementation Software

This appendix contains software to implement the core elements of the knock detection algorithm for a TI TMS320C30 DSP. It is not intended to be a complete description of the code required to implement a knock detection system.

Time-critical code is loaded to on-chip memory block B0. To minimize pipeline conflicts, variables accessed during the same instruction are located in different memory spaces. Variables required for the real-time code are split between the heap (dynamically allocated), which is placed in on-chip memory block B1, and external zero-wait-state SRAM.

The layout of the code is grouped into four sections:

- Definitions
 - Function prototypes
 - Definitions of constants
 - Global variables/pointers defined
- Mair
- Interrupt Service Routines (ISR)
 - External interrupt #0: c_int0()
 - Timer interrupt: c int09()
- Function Declarations
 - Global variable initialization: define()
 - Dynamic memory allocation: make_space()

Description

Dynamic Memory Initialization The function make_space() contains the code to dynamically allocate

structures and arrays.

RPM Adaptation The ISR c_int01() contains the RPM adaptation code.

Cylinder Initialization The ISR c_int01() contains the cylinder init code.

DFT Signal Conditioning The ISR c_int09() contains the DFT algorithm.

Knock Intensity Metric The ISR c_int09() contains the knock intensity metric calculation code.

Reference Adaptation The ISR c_int09() contains the reference adaptation code.

Listing

```
FUNCTION PROTOTYPES
                                 /* External Interrupt #1
/* Timer #1 Interrupt
      c_int01(void);
void
      c_int09(void);
C LANGUAGE MACROS */
DEFINITIONS
#define pi 3.14159265358979323846
#define cnt2sec clk2cnt/clkin /* Timer cnt period (sec)
#define inst2sec clk2isnt/clkin /* Inst cycle time (sec)
#define rev2delay tdc2knock/720 /* DFT delay per CC
DFT frequency bin numbers for fs=47 kHz and N=45. These numbers are used to
generate the sine/cosine tables needed for the DFT calculations. The bin
numbers are different for different block sizes, sample rates, and desired
frequencies. If any of these values change, then new bin numbers must be
determined and entered here.
#define nfreq1 6 /* f1 = 6.14 KHz

#define nfreq2 8 /* f2 = 8.18 KHz

#define nfreq3 11 /* f3 = 11.25 KHz

#define nfreq4 14 /* f4 = 14.32 KHz

#define nfreq5 19 /* f5 = 19.43 KHz
```

```
/*
The structure defined here contains the cylinder-specific information which
is used by the time-critical part of the DFT calculations. The definition of
the elements of the structure are:
   MAG[] = computed DFT magnitude history array
   REF[] b = scaled/averaged DFT array
   knock[] = knock metric histogram (diagnostics)
typedef struct cylinder
          float MAG[nrange*(history+1)];
          float REF[nrange];
          float knock[knock_levels];
         } cyl_type, *cyl_pntr;
typedef unsigned short UINT; /* Shorthand definition
           GLOBAL VARIABLES
                                                          * /
/*
                                             * /
/* Create array of structures (1/cyl)
cyl_pntr cyl_list[cyl_number];
/* Create pointers to elements of structure */
cyl_pntr cp;
float *Mag_p;
float
          *Mag_p_src;
float
          *Mag_p_dst;
float
          *REFp;
         *kp;
float
float *real_p;
                              /* DFT[real] sum
float *imag_p;
                              /* DFT[imag] sum
/* Allocate arrays for external SRAM memory
float real_exp[N]; /* Base cosine table array
                                 /* Base sine table array
float imag_exp[N];
                               /* Interleaved cosine table array
/* Interleaved sine table array
float real_table[(nfreq)*N];
float imag_table[(nfreq)*N];
                                 /* Gain array used for REF[] calc
float K[cyl_number][nrange];
/* Allocate global variables*/
float delay;
                                 / \, ^{\star} Delay period in Timer counts
float fs;
                                 /* Sample rate (Hz)
float k_nblocks;
                                 /* Conversion of Timer cnts to
                                 /* sample blocks / cylinder
/* CC degrees / knock degrees
float rev2knock;
                                 /* Number Timer cnts / DFT calc loop */
float DFT_loop;
float DFT_calc;
                                 /* Number Timer cnts / max DFT
                                  /* update loop
Offset pntr into base sine/cosine tables, which is used to construct the
interleaved sine/cosine tables. Each frequency being monitored requires
separate pointer into base tables.
int
      offset1;
```

```
int
     offset2;
    offset3;
int
      offset4;
int
      offset5;
int
                       /* Counter of DFT>REF for cyl */
/* Number RPM from RESET */
/* Current cylinder pointer */
int level_cntr;
int nrpm;
int
    ncyl;
    MAIN FUNCTION
Function main() is used for overall system control. It calls the functions
that do the initialization and setup tasks.
void main(void)
     :
  comp_exp(void);
  define();
  makespace();
  : -
  ; /* Infinite WAIT loop. All real-time */
/* code in interrupt routines. */
/* INTERRUPT SERVICE ROUTINES
/* INTO ISR
Function c_int01() is the INTO ISR, which is triggered by the ALL TDC signal.
The software always initializes Timer 0 Interrupt after the engine RPM calc
has stabilized (delay period needs to be known before DFT calcs started). If
CYL=1 (ncly=0), then delay period and nblocks are updated, otherwise those
blocks of code are bypassed.
void c_int01(void)
   UINT *mem_pntr;
                            /* Pntr to memory */
                             /* Pntr to |DFT| array */
   float *sum_p;
                              /* Counter
   int i;
Determine if CYL #1 by checking #1 TDC signal (LOW = yes). If yes, then read
720° CA period and run RPM adaptation routine.
   mem_pntr = (UINT *) READ_#1_TDC;
   temp = *mem_pntr;
   if (temp == 0)
   {
      mem_pntr = (UINT *) T1_CNTR;
                                         /* Read 720° period */
      Tlcnts = *mem_pntr;
      mem_pntr = (UINT *) T1_CNTL; /* Reset Timer 1 & */
```

```
if(nblocks==0) nblocks = 1;
                                              /* nblocks ≠ zero */
/* Update delay and */
       delay = T1cnts*rev2delay;
                                              /* Timer 0 PRD
       mem_pntr = (UINT *) T0_PRD;
                                               /* (for next cyl)
       *mem_pntr = delay;
                                               /* Reset cyl cntr
       ncyl = 0;
/* Cylinder initialization */
   cp = cyl_list[ncyl];
Mag_p = cp -> MAG;
                                       /* Create pointers to
/* structure for current
                                       /* cylinder
   REFp = cp -> REF;
   kp = cp -> knock;
                                       /* Init final |DFT| array */
/* to zero */
   sum_p = DFT_p;
   for (i=0 ; i<nfreq ; i++)
        *sum_p++ = 0;
                          TIMER 0 ISR
Function c_int09() is the TIMER 0 ISR which is triggered at the end of Timer
0 count down period (10 ^{\circ}CA). All DFT algorithm routines are run in this ISR.
Communications of knock intensity to ECU also occurs in this ISR.
void
         c_int09(void)
                           /* Pntr to C30 memory
/* Pntr to cos table
   UINT *mem_pntr;
   float *rtable_p;
                           /* Pntr to sine table
   float *itable_p;
                           /* Pntr to sine table
/* Pntr to DFT real sum
/* Pntr to DFT imag sum
/* Pntr to intermediate |DFT|
/* Pntr to final |DFT|
/* Storage for ADC input
/* Counters
   float *rvar_p;
   float *ivar_p;
float *tsum_p;
   float *sum_p;
   float temp;
   int i,j;
/* Init ISR */
   reset_T0_ie();
```

```
each sample block that is computed. It initializes inner loop variables, and
computes intermediate and final |DFT|s.
   for (j=0 ; j<nblocks ; j++)
                                   /* Zero DFT real/imag
/* sums. There is one
       rvar_p = real_p;
       ivar_p = imag_p;
       for (i=0; i<nfreq; i++) /* real/imag pair for
                                   /* each frequency being
           *rvar_p++ = 0;
                                       /* monitored.
          *ivar_p++ = 0;
       rtable_p = real_table;
                                   /* Init pointers to
                                   /* interleaved sine and
       itable_p = imag_table;
                                   /* cosine tables.
                                   /* Init pointers to
       tsum_p = tDFT_p;
                                   /* intermediate and
       sum_p = DFT_p;
                                   /* final |DFT| arrays.
                                   /* Clear INT2 bit in IF
       clr_int2();
                                   /* to init ADC polling
Start of inner loop of DFT real-time calculations. This loop is run once per
sample and computes the real/imaginary DFT running sums. It is reinitialized
for every sample block. The ADC sampling is polled rather than
interrupt-driven to save interrupt latency time.
       for (i=0 ; i< N ; i++)
          wait_int2();
                                       /* Wait for INT2 flag
          ADC_in = *mem_pntr;
                                       /* Read ADC
                                       /* Clear INT2 IF
          clr_int2();
                                       /* Convert to float
          temp = (float) ADC_in;
                                       /* Init pntrs to DFT
          rvar_p = real_p;
                                       /* real/imag running
          ivar_p = imag_p;
                                       /* sums.
This section computes the real and imaginary parts of the DFT. Inline code is
used to increase execution speed. Perform real/imag 2 mults and 2 adds for
each frequency being monitored.
           *rvar_p++ += temp * *rtable_p++; /* Freq #1 */
          *ivar_p++ += temp * *itable_p++;
          *rvar_p++ += temp * *rtable_p++; /* Freq #2 */
           *ivar_p++ += temp * *itable_p++;
           *rvar_p++ += temp * *rtable_p++; /* Freq #3 */
          *ivar_p++ += temp * *itable_p++;
           *rvar_p++ += temp * *rtable_p++; /* Freq #4 */
           *ivar_p++ += temp * *itable_p++;
           *rvar_p += temp * *rtable_p++;
                                              /* Freq #5 */
          *ivar_p += temp * *itable_p++;
/* Completed inner loop
This section computes the intermediate | DFT | ^2 for the current sample block
and determines final |\mathrm{DFT}|^2 for all sample blocks for current cylinder firing. The |\mathrm{DFT}|^2 is used instead of |\mathrm{DFT}| to save CPU processing time. The
tests made to determine knock intensity are looking for changes, so do not
require mathematical definition of |DFT| to function correctly.
       rvar_p = real_p;
                           /* Reinit pointers to */
                         /* real/imag DFT sums. */
       ivar_p = imag_p;
/* Intermediate | DFT | ^2 */
```

Start of outer loop of DFT real-time calculations. This loop is run once for

```
*tsum_p++ = *rvar_p * *rvar_p++ + *ivar_p * *ivar_p++;
       *tsum_p++ = *rvar_p * *rvar_p++ + *ivar_p * *ivar_p++;
       *tsum_p++ = *rvar_p * *rvar_p++ + *ivar_p * *ivar_p++;
       *tsum_p++ = *rvar_p * *rvar_p++ + *ivar_p * *ivar_p++;
               = *rvar_p * *rvar_p +
                                        *ivar_p * *ivar_p;
       *tsum_p
/* Final |DFT|^2 */
                                  /* Init pntr to int |DFT| */
       tsum_p = tDFT_p;
       *sum_p++ = MAX(*sum_p,tsum_p[0]);
       *sum_p++ = MAX(*sum_p,tsum_p[1]);
       *sum_p++ = MAX(*sum_p,tsum_p[2]);
       *sum_p++ = MAX(*sum_p,tsum_p[3]);
       *sum_p = MAX(*sum_p,tsum_p[4]);
/* Completed outer loop */
Compute knock metric. This is simple direct comparison method. Each monitored
frequency range is compared against the adaptive reference. The knock metric
is the number of frequency ranges that the measured magnitude is higher than
the reference magnitude.
   level_cntr = 0;
                                         /* Init cntr for knock tests
   if(DFT_p[0] > REFp[0]) level_cntr++; /* Test DFTn>REFn for all ranges */
if(DFT_p[1] > REFp[1]) level_cntr++;
   if(DFT_p[2] > REFp[2]) level_cntr++;
   if(DFT_p[3] > REFp[3]) level_cntr++;
   if(DFT_p[4] > REFp[4]) level_cntr++;
   kp = level_cntr;
If revs > rev_count then, the knock intensity metric is valid so output to
ECU is enabled.
   if(knock_out_flag)
     mem_pntr = (UINT *) WRITE_PA1;
     *mem_pntr = kp;
This section tests level_cntr to determine if knock has occurred. If no, then
new scaled average magnitudes (REF) are computed and the arrays in structure
holding past values (MAG) are updated.
   if (level_cntr == 0)
/* Copy measured magnitudes into history array in structure */
     Mag_p[0] = DFT_p[0];
     Mag_p[1] = DFT_p[1];
     Mag_p[2] = DFT_p[2];
     Mag_p[3] = DFT_p[3];
     Mag_p[4] = DFT_p[4];
/* Move Mag[n]'s back one time step in array */
     for (i=0; i<history; i++)
        Mag_p_dst = Mag_p + ((history-i)*nrange);
        Mag_p_src = Mag_p_dst - nrange;
        for (j=0 ; j<nrange ; j++)
          Mag_p_dst[j] = Mag_p_src[j];
Compute new scaled average magnitude. The reference is the scaled average of
the last 8 non-knocking events. A unique scaling factor is used for each
frequency range of each cylinder.
```

```
for (i=0 ; i<nrange ; i++) /* Zero REF array</pre>
                                                             * /
      REFp[i] = 0;
                                /* Compute new REF for all
   for (i=0 ; i<nrange ; i++)</pre>
                                /* the frequency ranges.
      for (j=1; j<=history; j++)</pre>
         REFp[i] += Mag_p[(j*nrange)+i];
      REFp[i] = 0.125 * REFp[i];
      REFp[i] = K[ncyl][i] * REFp[i];
  ncyl++;
                                /* Increment cylinder counter */
    _____*/
              FUNCTION DECLARATIONS
             Define Complex Exponential
Function comp_exp() computes the complex cos (real) and sin (imag) table
arrays and then constructs the interleaved tables used for the DFT
calculations. Block size determines the table length.
void
     comp_exp(void)
                         * /
  int j,k; /* Counters
/* This section computes the base sine/cosine arrays which are N long. */
  for (j=0 ; j<N ; j++)
    real_exp[j]=cos(2*pi*j/N);
    imag_exp[j]=sin(2*pi*j/N);
/*
This section constructs the interleaved sine/cosine arrays from the base
arrays. This was done to save the time required to manage numerous circular
buffers during the real-time code execution. It is feasible in an
experimental system because of the extra memory available on the C30 board.
The interleaved arrays are each nfreq*N long. The method used to manage the
pointer to the base arrays works with any block length N.
   for (j=0 ; j<N ; j++)
    k=j*nfreq;
                       /* Init offset pointer */
    real_table[k] = real_exp[offset1];
    imag_table[k] = imag_exp[offset1];
    offset1 += nfreq1;
    if (offset1 >= N)
       offset1 -= N;
    k ++;
    real_table[k] = real_exp[offset2];
    imag_table[k] = imag_exp[offset2];
    offset2 += nfreq2;
    if (offset2 >= N)
       offset2 -= N;
    k ++;
    real_table[k] = real_exp[offset3];
    imag_table[k] = imag_exp[offset3];
```

```
offset3 += nfreq3;
     if (offset3 >= N)
        offset3 -= N;
     k ++;
     real_table[k] = real_exp[offset4];
imag_table[k] = imag_exp[offset4];
     offset4 += nfreq4;
     if (offset4 >= N)
         offset4 -= N;
     k ++;
     real_table[k] = real_exp[offset5];
     imag_table[k] = imag_exp[offset5];
     offset5 += nfreq5;
     if (offset5 >= N)
         offset5 -= N;
                 Global Variable Initialization
/*
Function define() is used to define the initial values of global variables.
void
      define(void)
   ADC_in = 0;
   knock_out_flag = 0;
   level_cntr = 0;
   ncyl = 0;
   nrpm = 0;
NOTE: This section initializes the Kn's for each cylinder separately. If the number of cylinders for the engine is not 6 (current value of cyl_number),
then the number of the Kn sets below has to be modified to match.
   K[0][0] = 2.;
                       /* [Cylinder #1] [Range #1]
   K[0][1] = 2.;
                       /*
                                            [Range #2]
                                                           * /
                       /*
   K[0][2] = 2.;

K[0][3] = 2.;
                                                           */
                                            [Range #3]
                                                           * /
                       /*
                                            [Range #4]
                       /*
   K[0][4] = 2.;
                                            [Range #5]
                                                           * /
                                                           * /
                       /* Cylinder #2
   K[1][0] = 2.;
   K[1][1] = 2.;
   K[1][2] = 2.;
   K[1][3] = 2.;
   K[1][4] = 2.;
   K[2][0] = 2.;

K[2][1] = 2.;
                       /* Cylinder #3
   K[2][2] = 2.;
   K[2][3] = 2.;
   K[2][4] = 2.;
   K[3][0] = 2.;
                       /* Cylinder #4
                                                           * /
   K[3][1] = 2.;
   K[3][2] = 2.;
   K[3][3] = 2.;
   K[3][4] = 2.;
   K[4][0] = 2.;
                       /* Cylinder #5 */
   K[4][1] = 2.;
```

```
K[4][2] = 2.;
   K[4][3] = 2.;
   K[4][4] = 2.;
   K[5][0] = 2.;
                 /* Cylinder #6 */
   K[5][1] = 2.;
   K[5][2] = 2.;
   K[5][3] = 2.;
   K[5][4] = 2.;
/* This section computes a coefficient "k_nblocks" which is multiplied
/* the variable "Tlcnts" (720 crank angle degrees period) to determine
/* the number of sample blocks to use for the current engine RPM. The
/* calculation is defined below:
   K_NBLOCKS = REV2KNOCK * -----
                            (DFT_LOOP + DFT_CALC)
/* WHERE:
/* REV2KNOCK: Fraction of Combustion Cycle (720 deg) that knock window
     exists for a single cylinder.
   revprd 60 deg
   720 deg knockprd
/* CNT2SEC: Number of seconds per Timer count.
   4 clkins/Tlcnt
   32M clkins/sec
  DFT_LOOP: Number of seconds required for inner loop for given sample
       rate and sample block size.
               1 sec
   N samples
    1 block
             fs samples
/* DFT_CALC: Number of seconds required to compute |DFT|^2 and select
            the largest value. This number is approximate and is
            determined experimentally. If the technique to calculate
            magnitude or select max value is changed the time required ^{*}/
            for this section should be checked and modified.
    340 inst 2 clkins/inst
```

```
block
              32M clkins/sec
  rev2knock = revprd2deg*knockdeg;
  DFT_loop = N/fs
  DFT_calc = DFT_calc_cnts*inst2sec;
  k_nblocks = rev2knock*(cnt2sec/(DFT_loop+DFT_calc));
                        Array Allocation
Function make_space() is used to allocate memory space for arrays whose
length is determined at runtime. The values of the variable arrays in the structure are initialized also. The dynamic memory allocation stores the
variables on the heap (ON-CHIP RAM BLK1) and accesses them using pointers.
This method is used to speed execution time of knock intensity metric.
void make_space(void)
    int i,j,k;
/* Allocate structure */
  for (i=0 ; i<cyl_number ; i++)</pre>
        cyl_list[i] = (cyl_pntr)malloc(sizeof(cyl_type));
/* Initialize structure */
  for (i=0; i<cyl_number; i++)</pre>
        cp = cyl_list[i]; /* Create pointers
                                                               * /
       Mag_p = cp -> D;
       REFp = cp -> REF;
       kp = cp -> knock;
        for (j=0; j<knock_levels; j++)
    kp[j] = 0.;    /* Init knock histogram</pre>
        for (j=0; j<nrange; j++)
            /* Init Mag's & REF's large pos. number.
            for (k=0 ; k<=history ; k++)
               Mag_p[(k*nrange)+j] = 1.0e+37;
            REFp[j] = 2.0*Mag_p[j];
/* Allocate dynamic arrays */
    tDFT_p = (float*)malloc(sizeof(float)*nfreq);
    DFT_p = (float*)malloc(sizeof(float)*nfreq);
real_p = (float*)malloc(sizeof(float)*nfreq);
    imag_p = (float*)malloc(sizeof(float)*nfreq);
```

Appendix C: TMS320C25 Implementation Software

This appendix contains software to implement the core elements of the knock detection algorithm for a TI TMS320C25 DSP. It is not intended to be a complete description of the code required to implement a knock detection system.

The fixed-point implementation is intended to present code that is more appropriate for a production environment than the code in the 'C30 example. It is not as easily configurable because it is assumed that most constants are finalized and don't need to be changed as often as in a development system. Many of the constants and variables computed in the 'C30 example must be defined by the user in this one.

The layout of the code is grouped into three sections:

- Definitions
 - Definitions of constants
 - Global variables/pointers defined
- Main
 - Global variable initialization
- Interrupt Service Routines (ISR)
 - External interrupt #0: int0
 - Timer interrupt: tint

Description

Algorithm Initialization	The main routine contains the algorithm initialization code.
RPM Adaptation	The INTO ISR contains the RPM adaptation code.
Cylinder Initialization	The INT0 ISR contains the cylinder init code.
DFT Signal Conditioning	The TIMER ISR contains the DFT algorithm
Knock Intensity Metric	The TIMER ISR contains the knock intensity metric calculation code.
Reference Adaptation	The TIMER ISR contains the reference adaptation code.

Listing

```
DEFINITIONS
*
ENABLE_int0 .set 01h
DISABLE_int0 .set 0fffeh
get 08h
                                ;OR value w/ IMR to enable ext interrupt #0
                                ;AND value w/ IMR to disable ext interrupt #0
ENABLE_tint .set 08h
DISABLE_tint .set 0fff7h
                                ;OR value w/ IMR to enable timer interrupt
                                ;AND value w/ IMR to disable timer
              .set 44
                                ;Sample block size
             .set N/4
c_offset
                                ;Offset to start of cosine wave
             .set 5
.set 6
nfregs
                                ; Number of monitored frequencies
cyl_num
                                ; Number of cylinders
             .set 8
history
                               ;DFT average filter length
power
                               ;2^power = history
             .set 16-power
shift
                               ;scaling for ave mag calculation
```

```
* The constant array_size defines the amount of memory required to
* store the individual cylinder data. The total amount of space
 available on-chip is 384 words. Each cylinder contains the
* following arrays:
  K[nfreqs]
  REF[nfreqs]
  MAG[nfreqs*history]
* The constant cyl_size defines the number of memory locations for
^{\star} each cylinder. It is used to compute the constant array_size and
* to generate pointers into the array.
\mbox{\ensuremath{^{\star}}} The total number of memory locations required for this example is
* shown here:
  array_size = 6*(2*5+5*8) = 300
              .set 2*nfreqs+nfreqs*history
cyl size
array_size
              .set cyl_num*cyl_size
* Gain factors for K[] arrays. This example assumes the gain for each
* frequency is the same for each cylinder across the entire operating
* range. Might not be true in all cases.
ΚO
       .set 6
к1
       .set 2
K2
       .set 4
K3
       .set 5
К4
       .set 4
* DFT frequency bin numbers for fs=47 KHz and N=45. These numbers
* are used to modify the sine/cosine table pointers needed for the * DFT calculations. The bin numbers are different for various
* block sizes, sample rates, and desired frequencies. If any of
* these values change, then new bin numbers must be determined and
* entered here.
                         ; f1 = 6.14 \text{ KHz}
nfreq0
           .set 6
           .set 8
                         i f2 = 8.18 \text{ KHz}
nfreq1
nfreq2
           .set 11
                         ; f3 = 11.25 \text{ KHz}
                         ; f4 = 14.32 \text{ KHz}
nfreq3
           .set 14
           .set 19
                         ; f5 = 19.43 KHz
nfreq4
* This section shows how to compute the coefficient "k_nblocks" which is
* multiplied by the 24 bit variable "Timer" (time for \overline{7}20^{\circ} crank angle)
 to determine the number of sample blocks to use for the current engine
* RPM. The calculation is defined below using example numbers:
                                     CNT2SEC
    K NBLOCKS = REV2KNOCK *
                              (DFT_LOOP + DFT_CALC)
* WHERE:
 REV2KNOCK: Fraction of Combustion Cycle (720 deg) that knock window
              exists for a single cylinder.
   revprd
               60 deg
               ----- = 8.33*10
   720 deg
             knockprd
* CNT2SEC: Number of seconds per internal Timer count.
```

```
-7
   4 clkins/Tcnt
    ----- = 1*10
  40M clkins/sec
* DFT_LOOP: Number of seconds required for inner loop for given sample
           rate and sample block size.
  44 samples 1 sec ----- * ----- = 9.362*10
   1 block 47K samples
* DFT_CALC: Number of seconds required to compute |DFT|^2 and select
            the largest value. This number is approximate and is
            determined experimentally. If the technique to calculate
           magnitude or select max value is changed the time required for this section should be checked and modified.
   18 inst 5 freq 4 clkins/inst ------ * ------ = 9*10
     freq block 40M clkins/sec
* WHICH GIVES:
  -4 -6
                              9.362*10 + 9*10
KNBLOCK
             .set 148
* This section shows how to compute the coefficient "delay_frac" which is
* multiplied by the 24 bit variable "Timer" to determine the number of
* internal timer counts required for the specified delay period between
* TDC and the beginning of the knock window. The example below is for a
* desired delay period of 9°:
   9 deg 15
DELAY_FRACT = ----- * 2 = 409.6
                   720 deg
DELAYFRACT
             .set 410
     DATA TABLE
* The sine table is a multiple of four long. When this
* configuration is used, the values at 0°, 90°, 180°, and 270°
* are included which gives more accuracy in the calculations
sintabl
  .word
          00h
   .word
          01238h
  .word
         02410h
  .word 0352ch
   .word
         04534h
```

```
.word
            060bch
            06baeh
   .word
            0746eh
   .word
   .word
            07ad0h
   .word
            07eb2h
   .word
            07fffh
   .word
            07eb2h
            07ad0h
   .word
   .word
            0746eh
            06baeh
   .word
   .word
            060bch
            053d2h
   .word
            04534h
   .word
   .word
            0352ch
            02410h
   .word
   .word
            01238h
   .word
            01h
           0edc8h
   .word
            0dbf0h
   .word
            0cad4h
   .word
   .word
            0bacch
   .word
            0ac2eh
            09f44h
   .word
   .word
            09452h
   .word
            08b92h
   .word
            08530h
            0814eh
   .word
   .word
            08000h
   .word
            0814eh
   .word
           08530h
   .word
            08b92h
            09452h
   .word
   .word
           09f44h
            0ac2eh
   .word
   .word
            0bacch
   .word
            0cad4h
            0dbf0h
   .word
   .word
            0edc8h
endtable
         GLOBAL VARIABLES
; DFT algorithm variables
; Knock detection algorithm variables
; Cylinder history variables
; BLOCK B2
; These variables are used by the DFT algorithm and apply to all cylinders.
delay_frac .usect "dft_vars", 1
delay
              .usect "dft_vars", 1
              .usect "dft_vars", 1
flags
             .usect "dft_vars", 1
k_nblocks
             .usect "dft_vars", 1
.usect "dft_vars", 1
.usect "dft_vars", 1
metric
nblocks
ncyl
             .usect "dft_vars", 1
tablend
             .usect "dft_vars", 1
.usect "dft_vars", 1
tdc#1
temp
             .usect "dft_vars", 1
timer_hi
timer_lo
             .usect "dft_vars", 1
```

.word

053d2h

```
ADC_in
                                                                     .usect "dft_vars", 1
                                                                      .usect "dft_vars", 1
                                                                       .usect "dft_vars", 1
.usect "dft_vars", 1
MAG
REF
 ; BLOCK B0 (DP=4)
 ; sinewave table
                   .bss sinewave, N
 ; Pointers for magnitude and detection algorithms
                 .bss c_pntr, nfreqs
                   .bss
                                                s_pntr,
                                                                                                         nfreqs
                                                r_sum,
                   .bss
                                                                                                          nfreqs
                                                                                                       nfregs
                  .bss
                   .bss
                                                tDFT,
                                                                                                         nfreqs
                   .bss
                                                   DFT,
                                                                                                          nfreqs
; BLOCK B0 (DP=5) and BLOCK B1 (DP=6/7)
; This memory section starts at 0280h and contains arrays % \left( 1\right) =\left( 1\right) \left( 1\right) \left(
 ; for individual cylinder variables. Each cylinder requires
 ; the following arrays:
              K[nfreqs]
; REF[nfreqs]
            MAG[nfreqs*history]
; The magnitude array is organized as:
             +----
                                              history -->
                    nfreqs
;
cylinder .usect "dft_array", array_size
                                             MAIN FUNCTION
                .mmregs
Reset
 ; Disable all interrupts
                  ldpk 0 lack 0
                                                                               ;DP = mmregs data page
                                                                                          ; Reset IMR bits to disable interrupts
                  sacl
 ; Disable OVERFLOW mode (ACC won't saturate)
                  rovm
 ; Setup P register shift mode
                                             1
                                                                                      ;Shift left 1 on PREG ==> ACC
 ; Enable sign extention mode
                  ssxm
 ; Init variables
                                                                        delay_frac
                  ldpk
                  lalk
                                                                        DELAYFRACT
                                                                        delay_frac
                  sacl
                  lalk
                                                                        KNBLOCK
                  sacl
                                                                       k_nblock
                  zac
                  sacl
                                                                      ncyl
```

```
; Init sine table and tablend pointer
   ldpk
              sinewave
   larp
              AR1
              AR1, sinewave
   lrlk
   rptk
              endtable-sintabl-1
              sintabl, *+
   blkp
   lalk
              sinewave
   addk
              N-1
              tablend
   ldpk
              tablend
   sacl
; Init K[] array. This example uses same freq gain for each cylinder.
; In actual implementation all gains might be different.
           ARO, nfreq-1
                             ;Init counter
          AR1, cylinder
   lrlk
                             ;Init cylinder pointer
repeat
   larp
           AR1
   lack
          K0
                             ;Init K[0]
   sacl
           *+
   lack
          К1
                             ;Init K[1]
   sacl
   lack
         K2
                             ;Init K[2]
   sacl
          K3
   lack
                             ;Init K[3]
   sacl
         K4
   lack
                             ;Init K[4]
   sacl
         cyl_size-nfreqs,ARO ;Point to K[] for next cylinder
   adrk
   banz repeat, AR1
Wait
   b
                            ; Endless loop
          Wait
          INTERRUPT SERVICE ROUTINES
    INTO ISR
* ISR INTO is triggered by the ALL TDC signal. The software
* always initializes Timer 0 Interrupt after the engine
* RPM calc has stabilized (delay period needs to be known * before DFT calcs started. If CYL=1 (ncly=0), then delay
* period and nblocks are updated, otherwise those blocks of
* code are bypassed.
; Determine if cylinder \#1 by checking \#1 TDC signal (LOW = yes).
; If yes, then read external timer to determine 720^{\circ} period and
; run RPM adaptation routine.
          tdc#1, PA2
tdc#1, Ofh
                            ; Read #1 TDC signal
   in
   bit
                             ; Test for cylinder #1
                             ; If no, skip delay/nblock calc
   bbnz
          not_#1
; Read 24 bit external timer to obtain combustion cycle period.
          timer_hi, PA3 ; Read upper byte of timer
timer_lo, PA4 ; Read lower word of timer
   in
   in
; Delay count calculation
   lt
          delay_frac
   mpyu
          timer_lo
```

```
timer_hi
   mpya
   sach
          temp
   pac
   add
          temp
   sacl
          delay
; NBLOCKS calculation
   spm
          0
   lt
          k_nblock
          timer_lo
timer_hi
   mpyu
   mpya
   sach
          temp
   pac
   add
          temp
          temp
   sacl
   lac
          temp, 8
          nblocks
   sach
   spm
                            ;Set ncyl=0
   zac
   ldpk
          ncyl
   sacl
          ncyl
not_#1
; Initialize pointers to cylinder arrays K[], REF[], and MAG[]
                            ;ACC = cylinder address
          cylinder
   lalk
   spm
                            ;Setup for Q27 multiply
          temp
   ldpk
          nfreqs
   lack
   sacl
          temp
                            ;Calculate offset for
   lt
          ncyl
   mpyk
          cyl_size
                            ; current cylinder
                            ;Compute K[] address
   apac
   sach
          K
   addh
          temp
                            ;Compute REF[] address
   sach
          REF
                            ;Compute MAG[] address
   addh
          temp
   sach
          MAG
                            ; Restore PM to Q30 multiply
   spm
          1
; Zero DFT array for use in knock detection algorithm
   lrlk
          ARO, DFT
   larp
          AR0
   zac
   rptk
          nfreq-1
   sacl
    TIMER ISR
The TIMER ISR is triggered at the end of the on-chip timer count down period
(10 ^{\circ}\text{CA}). All DFT algorithm routines are run in this ISR. Communications of
knock intensity to ECU also occurs here.
tint
; disable Timer interrupt
   ldpk
   lac
          IMR
          DISABLE_tint
   andk
   sacl
          IMR
          ADC_in, PA0
                                ;Init ADC conversions
;START BLOCK PROCESSING
```

```
ldpk
          nblocks
   lar
          AR4, nblocks
                               ; Init block counter
blocks
;Zero real and imag sums
   lrlk
                               ;AR1 = real sum array
         AR1, r_sum
   lrlk
          AR2, i_sum
                               ;AR2 = imag sum array
   lark
                               ;AR3 = real/imag sum array counter
          AR3, nfreqs-1
   larp
          AR1
   zac
zero
          *+, AR2
                            ;Zero real sum element
   sacl
          *+, AR3
                            ; Zero imag sum element
   sacl
                            ;Check if done
   banz
          zero, AR1
; Init cosine/sine table pointers
   lrlk
          AR1, c_pntr
                            ;AR1 = cosine pointer array
          AR2, s_pntr
AR3, nfreqs-1
   lrlk
                            ;AR2 = sine pointer array
   lark
                            ;AR3 = cos/sin pntr array counter
   larp
          AR2
setup
   lalk
          sinewave
                            ;ACC = start of sine table
          *+, AR1
                            ; Init sine pointer element
   sacl
          c_offset
                            ;ACC = start of cosine table
   addk
                            ; Init cosine pointer element
   sacl
          *+, AR3
                            ;Check if done
   banz
          setup, AR2
   lark
          AR3, N-1
                            ; Init sample counter
                            ;ARP = AR1 (cos table pointer)
   larp
          AR1
;START DFT PROCESSING
; The BIO pin is polled to determine when ADC conversion is complete.
; This can be used to eliminate interrupt latency. If knock detection
; is not implemented on dedicated processor, an interrupt could be used
; for this section of the calculations.
; backward
                            ; If ADC complete continue
          forward
   bioz
   b
          backward
                            ; else wait
forward
; Read latest measurement from Port 0. ADC input is left justified.
   ldpk
          ADC_in
          ADC_in, PA0
   in
   lt
          ADC_in
                            ;Load T reg with latest measurement
; Compute running DFT calculation. A different index is loaded into ARO
; for each frequency being monitored.
; Frequency 0
   lark
          ARO, nfreq0
                            ;AR0 = freq0 index
   lar
                            ;AR1 = freq0 cos table pntr
          AR1, c_pntr
          AR2, s_pntr
*0+, AR2
   lar
                            ;AR2 = freq0 sin table pntr
   mpy
                            ;P = x*cos(freq0)
   zalh
                            ;ACC = real sum0
          r_sum
   mpya
          *0+, AR1
                            ;P = x*sin(freq0) & accum real product
          r_sum
                            ;store updated real sum0
   sach
                            ;ACC = imag sum0
   zalh
          i_sum
          AR1, c_pntr
   sar
                            ;store freq0 cos pntr
          AR2, s_pntr
                            ;store freq0 sin pntr
   sar
; Frequency 1
                            ;AR0 = freq1 index
   lark
          ARO, nfreq1
   lar
          AR1, c_pntr+1
                            ;AR1 = freq1 cos table pntr
   lar
          AR2, s_pntr+1
                            ;AR2 = freq1 sin table pntr
```

```
*0+, AR2
                            ;P = x*cos(freq1) & accum imag product
   mpya
   sach
          i_sum+0
                            ;store updated imag sum0
          r_sum+1
   zalh
                            ;ACC = real sum1
   mpya
          *0+, AR1
                            ;P = x*sin(freq1) & accum real product
   sach
          r_sum+1
                            ;store updated real sum1
          i_sum+1
                            ;ACC = imag sum1
   zalh
          AR1, c_pntr+1
   Sar
                            ;store freq1 cos pntr
                            ;store freq1 sin pntr
          AR2, s_pntr+1
   sar
; Frequency 2
   lark
          ARO, nfreq2
                            ;AR0 = freq2 index
                            ;AR1 = freq2 cos table pntr
   lar
          AR1, c_pntr+2
   lar
          AR2, s_pntr+2
                            ;AR2 = freq2 sin table pntr
                            ;P = x*cos(freq2) & accum imag product
   mpya
          *0+, AR2
                            ;store updated imag sum1
   sach
          i_sum+1
   zalh
          r_sum+2
                            ;ACC = real sum2
          *0+, AR1
                            ;P = x*sin(freq2) & accum real product
   mpya
   sach
          r_sum+2
                            ;store updated real sum2
          i_sum+2
                            ;ACC = imag sum2
   zalh
          AR1, c_pntr+2
                            ;store freq2 cos pntr
   sar
   sar
          AR2, s_pntr+2
                            ;store freq2 sin pntr
; Frequency 3
                            ;AR0 = freq3 index
   lark
          ARO, nfreq3
   lar
          AR1, c_pntr+3
                            ;AR1 = freq3 cos table pntr
                            ;AR2 = freq3 sin table pntr
   lar
          AR2, s_pntr+3
          *0+, AR2
                            ;P = x*cos(freq3) & accum imag product
   mpya
                            ;store updated imag sum2
   sach
          i_sum+2
                            ;ACC = real sum3
   zalh
          r_sum+3
   mpya
          *0+, AR1
                            ;P = x*sin(freq3) & accum real product
          r_sum+3
                            ;store updated real sum3
   sach
                            ;ACC = imag sum3
   zalh
          i sum+3
          AR1, c_pntr+3
                            ;store freq3 cos pntr
   sar
   sar
          AR2, s_pntr+3
                            ;store freq3 sin pntr
; Frequency 4
   lark
          ARO, nfreq4
                            ;AR0 = freq4 index
          AR1, c_pntr+4
                            ;AR1 = freq4 cos table pntr
   lar
   lar
          AR2, s_pntr+4
                            ;AR2 = freq4 sin table pntr
          *0+, AR2
                            ;P = x*cos(freq4) & accum imag product
   mpya
          i\_sum+3
                            ;store updated imag sum3
   sach
   zalh
          r_sum+4
                            ;ACC = real sum4
          *0+, AR1
                            ;P = x*sin(freq4) & accum real product
   mpya
   sach
          r_sum+4
                            ;store updated real sum4
   zalh
          i_sum+4
                            ; ACC = imag sum4
                            ;acc imag product
   apac
   sach
          i_sum+4
                            ;store updated imag sum4
          AR1, c_pntr+4
   sar
                            ;store freq4 cos pntr
   sar
          AR2, s_pntr+4
                            ;store freq4 sin pntr
; Check for cosine/sine table pointer overrun
          c_pntr
                            ; Check freq0 cosine pointer
   lac
          tablend
   sub
   blez
          spntr0
   adlk
                            ; Adjust pointer if overrun occurs
          sinewave
   sacl
          c_pntr
spntr0
   lac
          s_pntr
                            ; Check freq0 sine pointer
   sub
          tablend
   blez
          rpntr1
   adlk
          sinewave
                            ; Adjust pointer if overrun occurs
   sacl
          s_pntr
rpntr1
   lac
          c_pntr+1
                            ; Check freq1 cosine pointer
   sub
          tablend
```

```
blez
          spntr1
                            ; Adjust pointer if overrun occurs
   adlk
          sinewave
   sacl
          c_pntr+1
spntr1
                            ; Check freq1 sine pointer
   lac
          s_pntr+1
   sub
          tablend
   blez
          rpntr2
                            ; Adjust pointer if overrun occurs
   adlk
          sinewave
   sacl
          s_pntr+1
rpntr2
                            ; Check freq2 cosine pointer
   lac
          c_pntr+2
   sub
          tablend
   blez
          spntr2
   adlk
          sinewave
                            ; Adjust pointer if overrun occurs
          c_pntr+2
   sacl
spntr2
   lac
          s_pntr+2
                            ; Check freq2 sine pointer
          tablend
   sub
   blez
          rpntr3
   adlk
          sinewave
                            ; Adjust pointer if overrun occurs
   sacl
          s_pntr+2
rpntr3
   lac
          c_pntr+3
                            ; Check freq3 cosine pointer
          tablend
   sub
   blez
          spntr3
   adlk
          sinewave
                            ; Adjust pointer if overrun occurs
          c_pntr+3
   sacl
spntr3
          s_pntr+3
                            ; Check freq3 sine pointer
   lac
          tablend
   sub
   blez
          rpntr4
                            ; Adjust pointer if overrun occurs
   adlk
          sinewave
   sacl
          s_pntr+3
rpntr4
                            ; Check freq4 cosine pointer
   lac
          c_pntr+4
          tablend
   sub
   blez
          spntr4
                            ; Adjust pointer if overrun occurs
   adlk
          sinewave
   sacl
          c_pntr+4
spntr4
                            ; Check freq4 sine pointer
   lac
          s_pntr+4
   sub
          tablend
   blez
          pntr_end
                            ; Adjust pointer if overrun occurs
   adlk
          sinewave
   sacl
          s_pntr+4
pntr_end
          dft,AR1
                            ; Go back if more samples
   banz
; end inner loop
;Compute magnitudes: tDFT
                            ;Clear ACC
   zac
   mpyk
          Ω
                            ;Clear P reg
                            ;P = real0^2
   sqra
          r_sum
                            ;P = imag0^2 / ACC = real0^2
          i_sum
   sqra
                            ;P = real1^2 / ACC = real0^2 + imag0^2
   sqra
          r_sum+1
                            ;Store |DFT0|
          tDFT
   sach
                            ;Clear ACC
   zac
                            ;P = imag1^2 / ACC = real1^2
   sqra
          i_sum+1
                            P = real2^2 / ACC = real1^2 + imag1^2
   sqra
          r_sum+2
          tDFT+1
   sach
                            ;Store |DFT1|
```

```
;Clear ACC
   zac
                            ;P = imag2^2 / ACC = real2^2
   sqra
          i_sum+2
   sqra
          r_sum+3
                            P = real3^2 / ACC = real2^2 + imag2^2
          tDFT+2
                            ;Store | DFT2 |
   sach
   zac
                            ;Clear ACC
   sqra
          i_sum+3
                            ;P = imag3^2 / ACC = real3^2
   sqra
          r sum+4
                            P = real4^2 / ACC = real3^2 + imag3^2
   sach
          tDFT+3
                            ;Store | DFT3 |
                            ;Clear ACC
   zac
                            P = imag4^2 / ACC = real4^2
          i_sum+4
   sqra
                            ; ACC = real4^2 + imag4^2
   apac
          tDFT+4
   sach
                            ;Store |DFT4|
;Determine max magnitude: MAX(DFT,tDFT)
   ldpk
         DFT
   lac
          DFT
                            ; ACC = DFT0
                            ; ACC = DFTO - tDFTO
   sub
          tDFT
   bqz
                            ; If DFT > tDFT, no change
          max1
          tDFT
                            ;Else DFT = tDFT
   lac
   sacl
          DFT
max1
          DFT+1
                            ; ACC = DFT1
   lac
          tDFT+1
                            ; ACC = DFT1 - tDFT1
   sub
          max2
                            ; If DFT > tDFT, no change
   bgz
   lac
          tDFT+1
                            ;Else DFT = tDFT
   sacl
          DFT+1
max2
   lac
          DFT+2
                            ; ACC = DFT2
          tDFT+2
                            ; ACC = DFT2 - tDFT2
   sub
   bgz
          max3
                            ; If DFT > tDFT, no change
                            ;Else DFT = tDFT
   lac
          tDFT+2
   sacl
          DFT+2
max3
          DFT+3
                            ; ACC = DFT3
   lac
   sub
          tDFT+3
                            ; ACC = DFT3 - tDFT3
                            ; If DFT > tDFT, no change
   bgz
          max4
                            ;Else DFT = tDFT
   lac
          tDFT+3
   sacl
          DFT+3
max4
                            ; ACC = DFT4
   lac
          DFT+4
          tDFT+4
                            ; ACC = DFT4 - tDFT4
   sub
   bgz
          max_end
                            ; If DFT > tDFT, no change
   lac
          tDFT+4
                            ;Else DFT = tDFT
          DFT+4
   sacl
max_end
   larp
          AR4
                            ;Set ARP
          blocks
                            ;Go back if more blocks
   banz
; end outer loop
;Knock detection algorithm
   lark
          AR0, 0
                            ; Init counter
   ldpk
          DFT
                            ; Point to magnitude array
   lar
          AR1, DFT
   ldpk
          REF
   lar
          AR2, REF
                            ; Point to reference array
   larp
          AR1
                            ; Activate AR1
   lac
          *+, AR2
                            ; ACC = DFT
   sub
          *+, AR1
                            ; ACC = DFT-REF
          dif1
                            ; Is DFT>REF?
   blez
   larp
          AR0
                            ; Yes - increment counter
   mar
          *+, AR1
dif1
                             ; No - test next frequency range
          *+, AR2
                            ; ACC = DFT
   lac
```

```
sub
          *+. AR1
                           ; ACC = DFT-REF
   blez
          dif2
                           ; Is DFT>REF?
   larp
          AR0
                           ; Yes - increment counter
   mar
          *+, AR1
dif2
                           ; No - test next frequency range
          *+, AR2
                           ; ACC = DFT
   lac
          *+, AR1
   sub
                           ; ACC = DFT-REF
          dif3
   blez
                           ; Is DFT>REF?
          AR0
                           ; Yes - increment counter
   larp
          *+, AR1
   mar
dif3
                           ; No - test next frequency range
   lac
          *+, AR2
                           ; ACC = DFT
          *+, AR1
   sub
                           ; ACC = DFT-REF
          dif4
                           ; Is DFT>REF?
   blez
   larp
          AR0
                           ; Yes - increment counter
          *+, AR1
   mar
dif4
                           ; No - test next frequency range
                           ; ACC = DFT
   lac
          *+, AR2
          *+, AR1
   sub
                           ; ACC = DFT-REF
   blez
          dif_end
                           ; Is DFT>REF?
                           ; Yes - increment counter
          AR0
   larp
          *+, AR1
   mar
dif_end
                           ; No - save test results
          ARO, metric
   sar
; If knock_flag set (revs > rev_count) then, the knock intensity metric
; is valid so output to ECU is enabled.
                           ; Test knock_flag
   bit
          flags, Ofh
   bbz
          update
                           ; Skip output if flag=0
          PA1, metric
                           ; Output to ECU
   out
; Reference adaptation if no knock occurred
update
   lac
          metric
                           ; Test for knock
   bez
          done
                           ; Skip reference update if knock occurred
; Copy measured magnitudes into history array in structure
   larp
          AR1
   lark
          ARO, history
   lrlk
          AR1, MAG
   rptk
          nfreqs-1
   blkd
          DFT, *0+
; Move Mag[] elements back one time step in array for each frequency
; (must be onchip to use DMOV)
   lark
         ARO, nfreq-1
          AR1, MAG+(nfreqs*history-2)
   lrlk
move
          history-1
                           ; Move current frequency history
   rptk
   dmov
          *_
   mar
                           ;Setup for next frequency
          *_
   mar
          AR0
   larp
   banz
          move, AR1
                           ;Check if done
; Compute new average magnitude (use power of two length: 4,8,16)
                        ; Point to start of magnitude array
          AR1, MAG
   lrlk
   lrlk
          AR2, REF
                           ; Point to start of reference array
          AR3, nfreqs-1
   lark
                           ; Init frequency counter
   larp
          AR1
average
                           ; Zero ACC
   zac
          history-1
                           ; Sum history of magnitudes
   rptk
                           ; w/ 1/history scaling for average
   add
          *+, shift, AR2
          *+, AR3
   sach
                           ; Store average to REF array
   banz
          average, AR1
```

```
; Scale average magnitudes to create new REF array. This technique
; assumes that K*REF<32767. If this condition is not met, then additional
; scaling must be done.
   ldpk
          K
           AR1, K
                             ; Point to start of scale factor array
   lar
           AR2, REF
AR3, nfreqs-1
                             ; Point to start of reference array ; Init frequency counter
   lar
   lark
scale
           *+, AR2
                             ; T = K (scale factor)
   lt
                             ; P = K*REF (scaled ave)
; Store back to REF array
   mpy
   spl
           *+, AR3
   banz
           scale, AR1
                              ; Check for more
done
   ldpk
           ncyl
   lac
                             ; Increment cylinder counter
           ncyl
   addk
           1
   sacl
          nycl
   eint
                              ; and exit
   ret
```

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