BW6 over BLS12-381

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Following the recent work on FFT over non-smooth fields (ECFFT

), it might be viable to instantiate a SNARK with an elliptic curve of non-smooth subgroup order. Particularly, one layer SNARK composition can be achieved with less constraints for the choice of curves. In this note we propose a 2-chain based on the widely used BL12-381 curve. The outer curve is a BW6 with a subgroup order 2-adicity equal to one.

ECFFT

For smooth finite fields \mathbb{F_q}

(i.e., when q-1

factors into small primes) the Fast Fourier Transform (FFT

) leads to the fastest known algebraic algorithms for many basic polynomial operations, such as multiplication, division, interpolation and multi-point evaluation. However, the same operations over fields with no smooth order root of unity suffer from an asymptotic slowdown. [In a recent work by Ben-Sasson, Carmon, Kopparty and Levit, called [ECFFT]

](https://arxiv.org/abs/2107.08473), the authors proposed a new approach to fast algorithms for polynomial operations over all large finite fields.] The key idea is to replace the group of roots of unity with a set of points L \subset F

suitably related to a well-chosen elliptic curve group (the set L

itself is not a group). The key advantage of this approach is that elliptic curve groups can be of any size in the Hasse-Weil interval [q+1\pm2\sqrt{q}]

and thus can have subgroups of large, smooth order, which an FFT-like divide and conquer algorithm can exploit.

(From the paper's abstract

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2-chains

In HG20, the authors proposed a method to construct a fast pairing-friendly elliptic curve over BLS12-377. That is, a curve E_2

with a subgroup order r_2

equal to field size q 1

of BLS12-377. This method can be generalized to any inner curve and especially to BLS12 curves. The choice of BLS12-377 was interesting because q_1-1 \equiv 0 \mod 2^{46}

yielding efficient radix-2 FFT needed to generate SNARK proofs over the outer BW6 curve. With the ECFFT

breakthrough, one can swap BLS12-377 for the more widely used BLS12-381 and construct a new BW6 outer curve. This latter curve would have a subgroup order r_2

s.t. r_2-1 \equiv 0 \mod 2

, but ECFFT

approach might lead to acceptably efficient polynomial operations over \mathbb{F}_{r_2}

.

Inner curve: BLS12-381

 $BLS12-381 is a pairing-friendly \ elliptic \ curve from the family \ Barreto-Lynn-Scott \ with \ an \ embedding \ degree \ k=12$

(Construction 6.6 in FST06, case k\equiv 0 \mod6

). It is defined over a 381-bit field \mathbf{F}_q

and is defined by the equation y^2=x^3+4

. The curve parameters are parametrized by polynomials that are evaluated at the seed u=-15132376222941642752

Parameter

Paramete

Polynomial

value

2-adicity field size q

(x-1)^2/3\cdot r(x)+

1

subgroup order r

x^4-x^2+1

0x73eda753299d7d483339d80809a1d80553bda402fffe5bfefffffff00000001

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Besides the sizes of q

, r

and the the low hamming-weight of \boldsymbol{u}

, which makes the pairing-friendly curve efficient and secure at almost 128-bit level, the high 2-adicity of r-1

makes it a fit for SNARK applications. However, the low 2-adicity of q-1 $\,$

makes it inefficient to construct an outer curve with subgroup order $r_2=q$

.

Outer curve: BW6-767

Following HG20, we're looking for a curve E 2

with a subgroup order r_2

equal to BLS12-381 q

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. The same observations apply: we need q 2
 to be less than 768 bits and D=-3
  . We use the modified Brezing-Weng method with tailored Frobenius trace t(x)=t 0(x)+h t \cdot r(x)
 and CM parameter y(x)=y_0(x)+h_y \cdot cdot r(x)
 (from CM equation 4q=t^2+Dy^2
 ). The lifting cofactors h t
and h y
 must be chosen carefully according to the desired bit lengths and the CM discriminant -D
 . Given these constraints, we found 4 curves, which boils down basically to two pair of curves (a curve and its twist over the base field \mathbb{F}_q
 Our script, which will be soon open-sourced under MIT license, outputs the following curves:
test_vector_BLS12_381_BW6_768 = [ {'u':-0xd20100000010000, 'D:3, 'ht':-4, 'hy':-6, 'b':1, 'pnbits':767, 'rmbits':381, 'px':[31,65,37,115,112,-100,56,29,-191,79,142,-127,31], 'px_denom':9, 'rx': [1,1,0,2,0,-2,1], 'rx_denom':3, 'cx';52_16_21,41,12,-65,31], 'cx_denom':3, 'yx':[6,7,0,9,3,-13,6], 'yx_denom':3, 'tx':[-4,-1,0,-17,9,5,-4], 'tx_denom':3, 'betax':[72,62,-16,258,-86,-152,234,-97,-190,300,-158,31], 'betax_denom':3, 'label':"BLS12_BW6_767"], {'u':-0xd201000000010000, 'D':3, 'ht':5, 'hy':9, 'b':9, 'pnbits':768, 'rmbits':381, 'px':[67,128,64,286,238,-262,209,38,-515,304,241,-262,67], 'px_denom':9, 'rx':[1,1,0,2,0,-2,1], 'rx_denom':3, 'tx':[5,8,0,1,9,-13,5], 'tx_denom':3, 'bx:denom':3, 'px:denom':3, 'tx':[5,8,0,1,9,-13,5], 'tx_denom':3, 'bx:denom':3, 'px:denom':3, 'tx':[5,8,0,1,9,-13,5], 'tx_denom':3, 'batel':"BLS12_BW6_768"], {'u':-0xd201000000010000, 'D':3, 'ht'.7, 'hy':5, 'b':3, 'pnbits':767, 'rmbits':381, 'px::[31,65,37,115,112,-10,65,29,-191,79,142,-127,31], 'px_denom':9, 'rx':[1,1,0,2,0,-2,1], 'rx_denom':3, 'cx:[19,6,2,14,112,-65,31], 'cx_denom':3, 'yx:[5,4,0,13,-3,-9,5], 'yx_denom':3, 'tx':[7,10,0,59,-17,7], 'tx_denom':3, 'batax':[72,62,-16,258,-86,-152,234,-97,-190,300,-158,31], 'batax_denom':3, 'px::[31,65,37,115,112,-10,65,29], 'px:denom':3, 'px::[31,65,37,115,112,-10,65,29], 'px
 All these curves benefit from endomorphism-based optimizations as D=-3
 and from fast Miller loop as the seed u
 has low hamming-weight (same as BLS12-381)
 That is said, the best curve is the first one: y^2=x^3+1
 defined over a 767-bit field with D=-3
 and h_y=-6
       • The field \mathbb{F} o
 can be implemented in twelve 64-bit limbs, with one bit to spare for carries or for compressed y-coordinate flags.
       . The cofactors h t,h y
 give the fastest final exponentiation in the pairing computation as an exponentiation by h_t^2+3h_y^2
 is needed.

 The coefficient b=1

 is convenient for checking that a point is on the curve.
       • There is a 2-isogeny from a curve with j \ne 0
 or 1728
 , allowing use of the "simplified SWU" method for hashing to an elliptic curve.
 The constant c=3
 is a sextic non-residue in \mathbb{F} q
 and so the twisted curve is a M-twist that corresponds to the third output curve: y^2=x^3+3
 over the same field \mathbb{F}_q
 with D=-3
 , h t=7
 and h_y=5
 . The extension fields can be constructed as \mathbb{F}_{q^3}[u]=\mathbb{F}_{q^3}[u]=\mathbb{F}_{q^3}[u]
 and \mathbb{F}_{q^6}[v]=\mathbb{F}_{q^3}/v^2-u
 \rightarrow
 Below are the polynomial forms of the curves' parameters:
 Curve
 subgroup order r(x)
 (same as BLS12-381 q(x)
 field size q(x)
 cofactor c(x)
 Frobenius trace t(x)
 CM y(x)
 GLV \lambda(x)
 GLV \beta(x)
 y^2=x^3+1
 (x-1)^2/3\cdot (x^4-x^2+1)+x
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(31x^{12} - 127x^{11} + 142x^{10} + 79x^9 - 191x^8 + 29x^7 + 56x^6 - 100x^5 + 112x^4 + 115x^3 + 37x^2 + 65x + 31)/9
 (93x^6 - 195x^5 + 36x^4 + 123x^3 + 63x^2 + 48x + 156)/9
 (-4x^6 + 5x^5 + 9x^4 - 17x^3 - x - 4)/3
 (6x^6 - 13x^5 + 3x^4 + 9x^3 + 7x + 6)/3
 1-x+3x^3-3x^4+x^5
 (72+62x-16x^2+258x^3-86x^4-152x^5+234x^6-97x^7-190x^8+300x^9-158x^{\{10\}}+31x^{\{11\}})/33
y^2=x^3+3
 (twist)
(x-1)^2/3\cdot (x^4-x^2+1)+x
 (31x^{12} - 127x^{11} + 142x^{10} + 79x^9 - 191x^8 + 29x^7 + 56x^6 - 100x^5 + 112x^4 + 115x^3 + 37x^2 + 65x + 31)/9 + 12x^4 + 115x^3 + 37x^2 + 65x + 31)/9 + 12x^4 + 115x^3 + 37x^2 + 65x + 31)/9 + 12x^4 + 115x^3 + 37x^2 + 65x + 31)/9 + 12x^4 + 115x^3 + 37x^2 + 65x + 31)/9 + 12x^4 + 115x^3 + 37x^2 + 65x + 31)/9 + 12x^4 + 115x^3 + 37x^2 + 65x + 31)/9 + 12x^4 + 115x^3 + 37x^2 + 65x + 31)/9 + 12x^4 + 115x^3 + 37x^2 + 65x + 31)/9 + 12x^4 + 115x^3 + 37x^2 + 65x + 31)/9 + 12x^4 + 115x^3 + 37x^2 + 65x + 31)/9 + 12x^4 + 115x^3 + 37x^2 + 65x + 31)/9 + 12x^4 + 115x^3 + 37x^2 + 65x + 31)/9 + 12x^4 + 115x^3 + 37x^2 + 65x + 31)/9 + 12x^4 + 115x^3 + 37x^2 + 65x + 31)/9 + 12x^2 + 12x^
 (93x^6 - 195x^5 + 36x^4 + 123x^3 + 63x^2 + 48x + 57)/9
 (7x^6 - 17x^5 + 9x^4 + 5x^3 + 10x + 7)/3
(5x^6 - 9x^5 - 3x^4 + 13x^3 + 4x + 5)/3
 1-x+3x^3-3x^4+x^5
 (72+62x-16x^2+258x^3-86x^4-152x^5+234x^6-97x^7-190x^8+300x^9-158x^{\{10\}}+31x^{\{11\}})/33
where \lambda(x)
 is the endomorphism eigenvalue which when multiplied by a point P=(x_P,y_P)
on the curve acts as (x_P,y_P) \pmod (\beta(x)\cdot (x_P,y_P)
 \rightarrow
 Below are the hexadecimal values of the shared curves' parameters, when evaluated at the seed u=-015132376222941642752
 subgroup order r
(same as BLS12-381 q
 field size a
GLV \beta
 GLV \lambda
 y^2=x^3+1
and v^2=x^3+3
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0x51e2bcf25fa8992238259ea59a063294c36dc4098befce4230f8d18f41e3fc19665e4360b872007d3dd5a1b865cbe8dadc2ce0c034926d18fe0ef8c1c63df7d97cbc118805598e5c31732000974254c83a38t
0x4a7108dc56f65caaa3748933617f8fb542e1d6e9ef88f166a011051b5f65b0b728456c117f8e93508e38b5b4682991ce3dae358b1b36f832f0bb174392eb6c806d9cd8a33550dbe01e63601ff328f27c5a594f
0x1a0111ea397fe699ec02408663d4de85aa0d857d89759ad4897d29650fb85f9b409427eb4f49fffd8bfd00000000aaac

Implementation

A first version of an optimized implementation of the curve arithmetic and the underlying finite fields, alongside optimal ate pairing computation (with fast final exp) was written in Golang undernark-crypto library (zkTeam @ConsenSys

). It can be found in this fork's branch:gnark-crypto/ecc/bw6-767 at feat/bw6_on_bls12-381 · yelhousni/gnark-crypto · GitHub

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Some more optimizations are WIP:

- [] fast cofactor clearing
- [] fast subgroups check
- [] hash-to-curve

Conclusion

For interoperability of blockchain projects, it would be great to use the same elliptic curve. BLS12-381 is a secure, optimized and widely used pairing-friendly elliptic curve in many patforms. For projects that need one layer proof composition, it might be viable to stick to BLS12-381 and use this BW6-767 alongside ECFFT for composing proofs.

A detailed paper will be published soon that proposes a generic framework for constructing optimized 2-chains for SNARK composition with open-sourced code.