

I am looking to create a simple smart contract based algorithm
that provides common random number generation
, provided that t
out of N

participants are honest, where $t * 2 > N$

The algorithm has the following properties

- at the end, all honest participants know the common random number R
- malicious participants are not able to influence R

or make the algorithm stuck.

Preliminary setup

Each participant j

register her public keys $P[j]$

with a smartcontract DRNGManager

.

Together with the public key $P[j]$

each participant submits to DRNGManager

a ZK proof that $P[j]$

is valid and that she knows the corresponding private key.

Commit phase

(10 minutes)

Each participant j

generates a random EC polynomial of degree t

$POLY_j$

.

The participant then generates a vector of polynomial evaluations $A[i] = [POLY_j]$

at N

integer points i

.

The participant will then encrypt the evaluations to obtain a vector of encrypted polynomial evaluations $G_{[j]} = [Encrypt(POLY_j)]$

.

It then submits to DRNGManager

- vector $G_{[j]}[i]$
- commitment to $POLY_j$
- a ZK-proof that $G_{[j]}[i]$

were correctly generated from $POLY_j$

.

DRNGManager verifies ZK-proof on receipt

After the commit phase, DRNGManager

will contain j

valid vectors $G[j]$

, where $j \geq t$

.

Reveal phase

. (10 minutes)

Each participant j

will be able to decrypt and reveal to DRNGManager

a vector of points $POLY_j$

. The participants will then submit these vector to DRNGManager

together with a ZK proof that reveal was done correctly.

After the reveal phase, DRNGManager

will include k

reveals, where $k \geq t$

.

RNG computation phase

. For each committed polynomial $POLY[j]$

, each participant is then able calculate random number $R[j]$

$= POLY_j$

. The common random is then XOR of all $R[j]$