# Electoral Engineering through Simulation

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## Electoral System Properties

- Fairness/Proportionality
- Decisiveness/Governability

• Trade-offs are not linear.

#### Measures

- Proportionality
  - Loosemorehanby Index
  - Gallagher Index
- Governability
  - Effective number of parties
  - Sum of squares of Shapley Shubik Power Index
  - Entropy of Shapley Shubik Power Index

- The Electoral Sweet Spot paper claims district magnitude of 2 or 3 is optimal.
- Uses data from 609 real elections in 81 different countries.
- We aim to validate this theoretically by simulation of artificial societies.

#### Models for Artificial Societies

- Polya Eggenberger (Urn model)
- Spatial Model

• We assume districts are independent.

#### Urn Process (Intuition)

- Urn starts with one ball of each colour.
- We select a ball randomly from the Urn, record it, and place the ball and  $\alpha$  copies of the ball back into the Urn.

- $\alpha$  represents the homogeneity of preferences.
- Each colour represents a preference order.

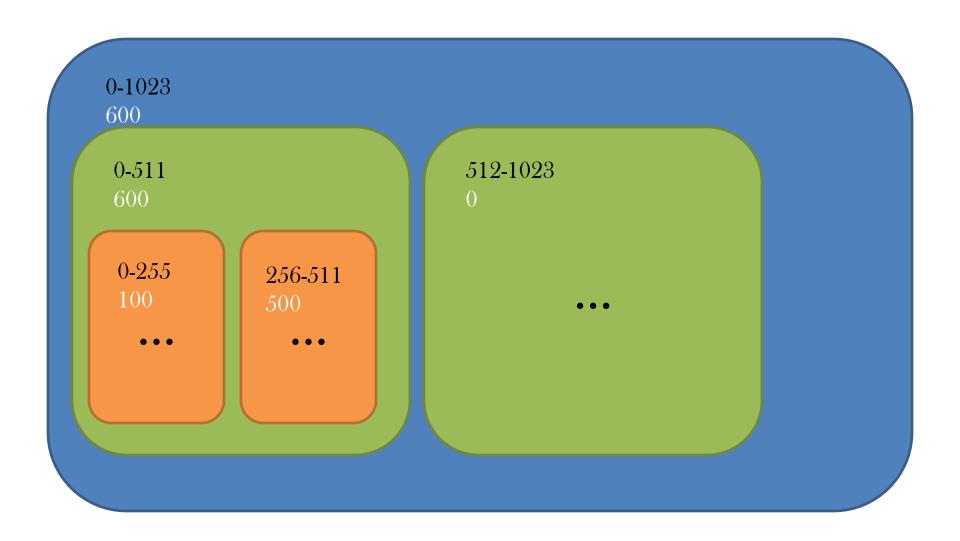
### Sampling from the Urn process

- m = number of parties.
- n = number of voters.
- M = number of preference orders.
- Naive approach would be to simply store the balls in a list and sample from the list.
- M = m!, very large for even relatively small m.  $M \approx 4,000,000$  for m = 10
- Typically n < M
- $n \approx 30,000$  in the case of New Zealand.

## Efficiently Sampling from the Urn

- We take advantage of the fact that almost all preference orders have a weight of 1.
- Time complexity of order  $n \cdot \log(m!)$
- We select **x** uniformly at random between 0 and 1.
- We select the smallest node with the cumulative probability distribution  $\leq x$ .

#### Tree Structure



- Each node in the tree represents a range. Starting from the root, we recursively select a child range that contains x.
- Each node has a parameter w representing the total weight of the range.
- We can compute whether a node contains x in constant time.
- We never actually use most of the nodes.
  Nodes are lazily created.

- We create parties with coordinates in an arbitrary dimensional space (e.g. 2D) called the space of issues.
- Each voter has random coordinates and selects a preference order by distance from itself to each party. The lower the distance, the more a voter prefers a party.
- We vary the mean for each coordinate on the district level.

- Selecting the number of dimensions and variance in each.
- Placement of parties.
  - Normal distribution
  - Normal distribution + traditional major "left" and "right" parties extremely close to the origin.
- We may be able to obtain typical values for these from real world data. (e.g. PCA on votes for each party, by district)

- Variance of distribution used for selecting means for larger district should be lower than that of a smaller district.
- We create a larger district by two methods
  - by combining multiple independent districts
  - the case where the variance is unchanged for the larger district.
- In reality the variance of the distribution for larger districts should be between these two figures. (e.g. it's not very plausible that a far left district and far right district would be combined)

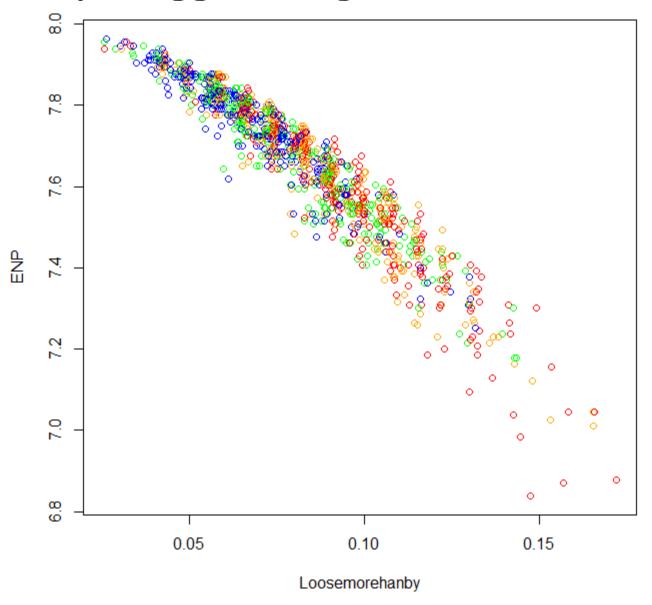
## Preliminary Results

• These results are for STV with different values of district magnitude.

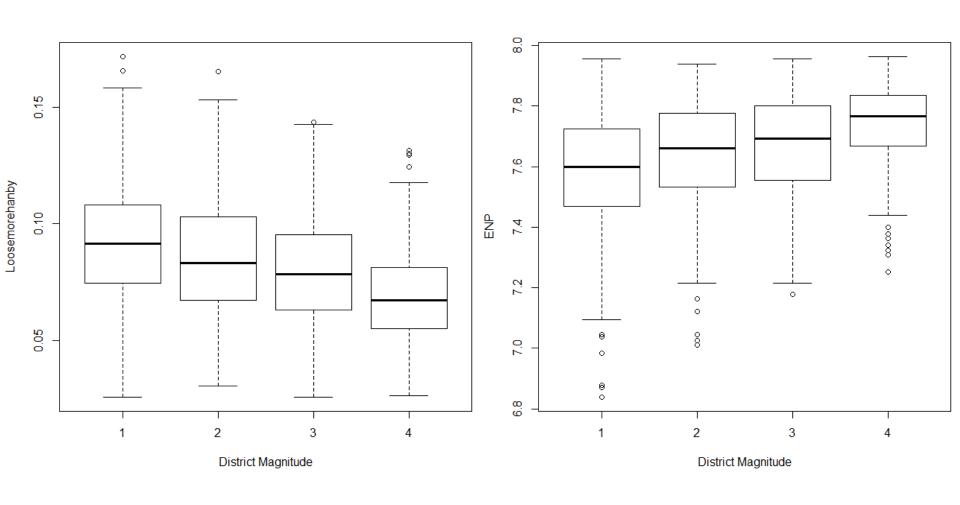
Colour	DM
Red	1
Orange	2
Green	3
Blue	4

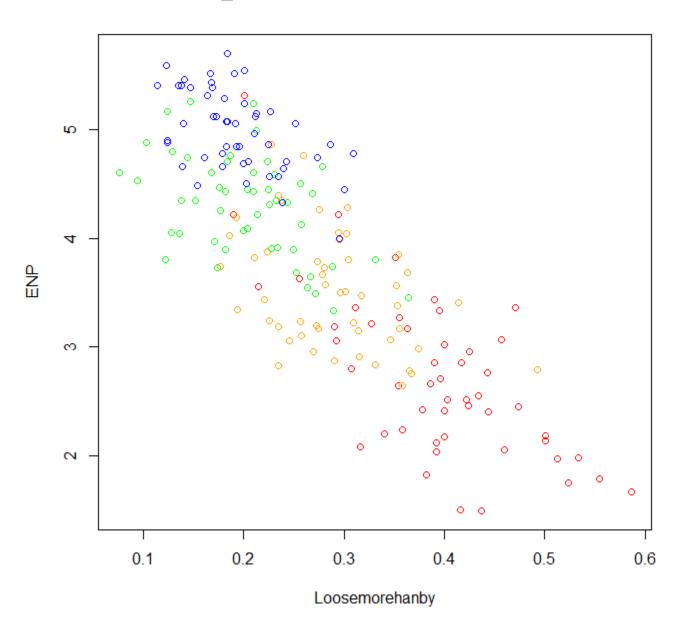
• Points near the bottom left corner are better.

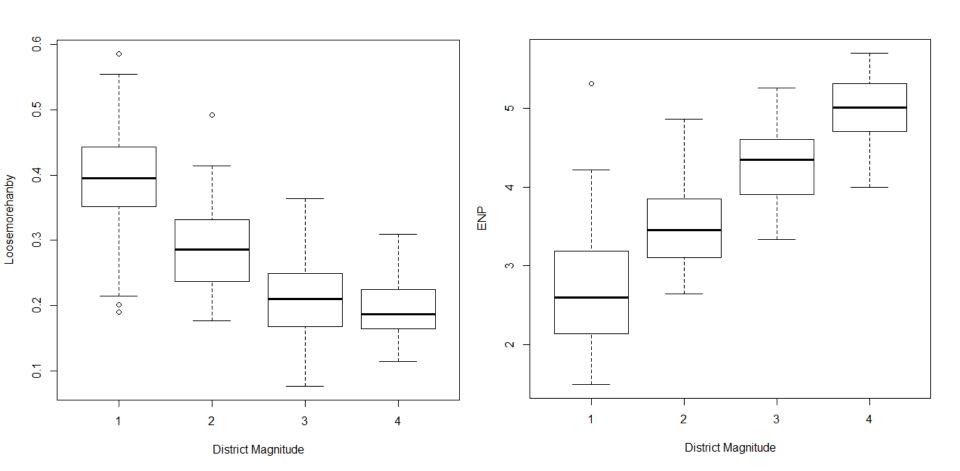
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#### References

- John Carey and Simon Hix (2011) 'The Electoral Sweet Spot: Low-Magnitude Proportional Electoral Systems', *American Journal of Political Science* 55(2) 383-339.
- New Zealand Election Study 2008.
  <a href="http://www.nzes.org/">http://www.nzes.org/</a>